

$$1. P[|\epsilon_{in}(g) - \epsilon_{out}(g)| > \varepsilon] \leq 2M e^{-2\varepsilon^2 N}$$

$$\varepsilon = 0.05 \rightarrow \text{want } 2M e^{-2(0.05)^2 \cdot N} \leq 0.03$$

$$M=1 : e^{-2(0.05)^2 N} \leq \frac{0.03}{2}$$

$$-2 \cdot 0.05^2 N \leq \ln(0.015)$$

$$N \geq -\frac{\ln(0.015)}{2 \cdot 0.05^2}$$

$$N \geq 839.94$$

The least number of examples among the choices is 1000 [b]

$$2 \cdot \text{for } M=10: 2(10)e^{-2(0.05)^2 \cdot N} \leq 0.03$$

$$e^{-2(0.05)^2 \cdot N} \leq \frac{0.03}{20}$$

$$-2(0.05)^2 N \leq \ln(0.03/20)$$

$$N \geq -\frac{\ln(0.03/20)}{2(0.05)^2}$$

$$N \geq 1300.5$$

the least number of examples required of  
the given choices is  $1500 \Rightarrow [c]$

3. for  $M=100$ : (same eq. as 1 and 2)

$$N \geq -\frac{\ln(0.03/200)}{2(0.05)^2}$$

$$N \geq 1761.0$$

the least number of examples required if  
the given choices is 2000  $\Rightarrow [d]$

4. The break point in  $\mathbb{R}^2$  is 4, because there is always a line that separates any 3 points. In  $\mathbb{R}^3$ , this means that there is always a plane that separates any 4 points, because a plane in  $\mathbb{R}^3$  is similar to a line in  $\mathbb{R}^2$ . So, adding one more point does not guarantee that a plane will separate the 5 points. For example, if there are 4 points on a plane and a 5th one that is not on that plane, then there is no plane that can separate all possible labelings of the 5 points.

answer: [b]

So from lecture 6, we know that  $m_H(N)$  is polynomial in  $N$ , and from lecture 5, we know that  $m_H(N) \leq 2^N$  and  $m_H(N)$  is not polynomial in  $N$  only when  $m_H(N)$  does not have a break point. So, the growth formulas that are polynomial in  $N$  or equal  $2^N$  are valid.

i)  $1+N$  ✓ this is possible because it is linear in  $N$  and  $\leq 2^N$ . In fact, we know that this is the growth function for positive rays.

$$\text{ii) } 1+N + \binom{N}{2} = 1+N + \frac{N!}{(N-2)! \cdot 2!} = 1+N + \frac{(N)(N-1)}{2}$$

✓ this is a possible growth function too because it is polynomial in  $N$ .

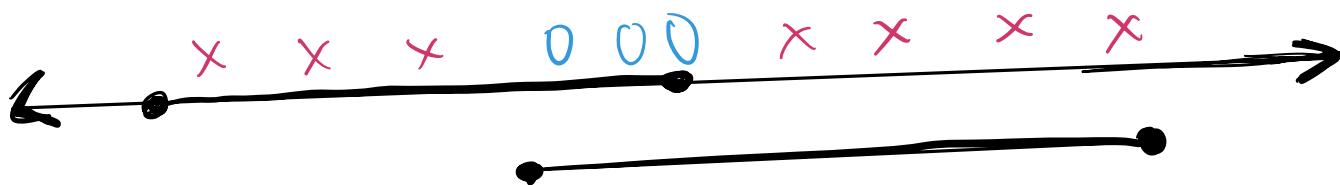
iii)  $\sum_{i=1}^{\lfloor \sqrt{N} \rfloor} \binom{N}{i} \rightarrow$  This is not polynomial in  $N$  and not equal to  $2^N$ .

This is not a valid growth function because it is not polynomial in  $N$  and not equal to  $2^N$  (for a  $H$  with no break points).  $m_H(N) \leq \sum_{i=0}^{\lfloor \sqrt{N} \rfloor} \binom{N}{i}$  this inequality will not hold for large  $N$  and small  $K$ .

- iv)  $2^{\lfloor N/2 \rfloor}$  this growth function is not polynomial in  $N$ , so it is not valid. X
- v)  $2^{\sim}$  this is the growth function for a hypothesis with no break point, so this is a valid growth function. ✓
- i, ii, and v are valid  $\Rightarrow$  [b]

$$6. h: \mathbb{R} \rightarrow \{-1, +1\}$$

$h(x) = 1$  if it is within two randomly chosen intervals  
 $h(x) = -1$  otherwise



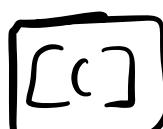
we want to know the smallest number of points where there are no two intervals that can shatter any labeling of the points

try 3:  $\begin{matrix} X & X & O & \checkmark \\ X & O & X & \checkmark \end{matrix}$  4 points also works because all arrangements will be similar to 3 points other than:

$\begin{matrix} X & O & O & \checkmark \\ O & O & X & \checkmark \end{matrix}$  In which case each  $O$  can correspond to one interval.

$\begin{matrix} O & X & O & \checkmark \\ X & X & X & - \end{matrix}$  Any arrangement at 5 points cannot necessarily be shattered by any

intervals outside of points  $\rightarrow$  At least one interval covers all points

two intervals because:  $\textcircled{0} \times \textcircled{0} \times \textcircled{0}$   
requires that each  $\textcircled{0}$  corresponds to an interval, since the intervals are continuous,  
there cannot be a  $\times$  in between two  $\textcircled{0}$ 's  
that are within the same interval. Thus,  
5 is the smallest break point for the  
2-intervals learning model.  $\Rightarrow$  

7.  $m_n(N)$  for "2-intervals"

Based on slide 11 of lecture 6,

$$m_n(N) \leq \sum_{i=0}^{k-1} \binom{N}{i} \text{ for } m_n(N) \text{ w/ break point } k$$

We have 3 cases for the 2-intervals:

We know from problem 6 that 4 is the maximum number of points that can be shattered by 2 intervals, so the '—' are the possible spots

— 1 — 2 — 3 — 4 — for the endpoints of the intervals and the numbers are the points.

Case 1: intervals are not overlapping at all

$\Rightarrow 4$  endpoints for total of 2 intervals :

$$\binom{N+1}{4} \xrightarrow{\text{N+1 spots for}} N \text{ points}$$

Case 2: intervals overlap, so there are

exactly 2 end points :  $\binom{N+1}{2}$

Case 3: all the points are outside both intervals, which is just one case (all points are -1)

Final answer:  $\boxed{\binom{N+1}{u} + \binom{N+1}{2} + 1 \Rightarrow [c]}$

8. M-intervals case where point falls in any of M randomly chosen intervals

For the two intervals case, we found that the break point is 5 because if all 5 points were alternating in terms of classification (starting with +1), then it required 3 intervals.

With M intervals, if we have  $2M+1$  points that are alternating in classification (starting with +1), then  $M+1$  intervals are required to shatter these points. Again, since intervals are continuous, two o's with a  $\times$  in between must be in two distinct intervals. Thus, the smallest number of break points for M intervals is  $2M+1 \Rightarrow [d]$

9. The largest number of points that can be shattered by this hypothesis set where +1 is inside a randomly chosen triangle.

1, 3, 5, 7, 9 are options for max that can be shattered

2 points → The triangle can always enclose 0, 1, or 2 of the points because one side of the triangle can divide any two of the points (a line can always split 2 points) or it can be made large or small enough to have both/no points.

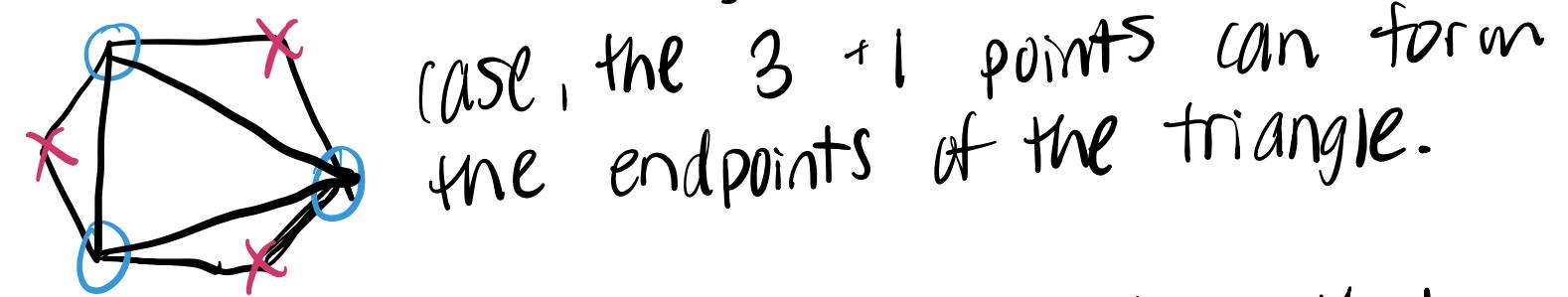
So, 1 is not the maximum.

4 points → A triangle can also cut across any two sides of the quadrilateral formed by the 4 points to omit any arrangement of points.

This means that 3 is also not the maximum. Worst case scenario for 4 points if the ones that are supposed to be inside are diametrically opposite.



6 points → Again, we can show that the worst-case scenario for 6 points where the labels are alternating around a hexagon can still be shattered by a triangle. In this

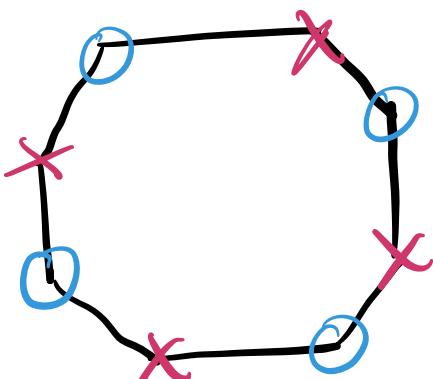


8 points → Now, we attempt to show that alternating labels around an octagon can also be shattered by a triangle. There is no triangle that can enclose the  $O$ 's here. In the worst-case scenario, where the

8 points are equidistant around a circle and alternatively labeled  $+1$  and  $-1$ , the convex properties of a triangle do not allow for a valid triangle to only enclose the  $O$ 's, it must enclose an adjacent  $X$  to do so, which means that our hypothesis set cannot shatter the hypothesis.

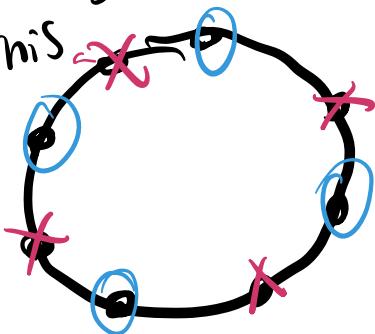
So, 8 points cannot be shattered, but 6 points can, so we know the answer is

$[d] \Rightarrow 7$



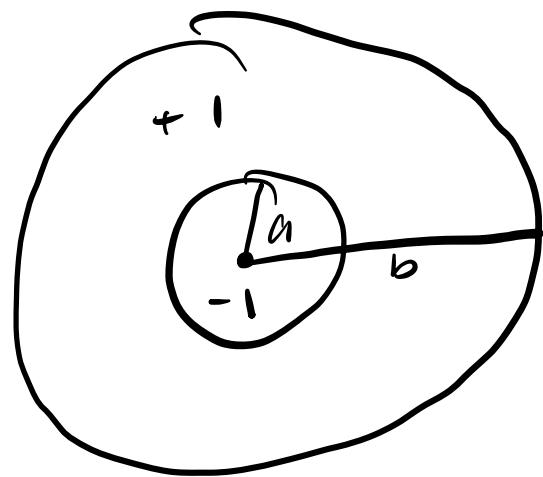
$x^2 + y^2 = r^2$

all 8 points  
satisfy this



$$IV. \quad a^2 \leq x_1^2 + x_2^2 \leq b^2$$

This is essentially labeling all the points inside the larger circle but outside the smaller circle as +1, and -1 otherwise. This is the same as labeling points inside the interval  $[a, b]$  as +1, and outside the interval as -1.



The selection of  $a$  and  $b$  are the only constraints that matter, because this selected interval is just rotated at one point to create a circle. Thus, we can calculate the growth function in a very similar manner to that of problem 7. In this case, there are only two scenarios for the interval: the two end points are amongst the  $N+1$  gaps between  $N$  points ( $\binom{N+1}{2}$ ) or all the points are outside the interval (1).

Final answer:  $\boxed{\binom{N+1}{2} + 1 \Rightarrow [b]}$