**Esercizio 1.** Let  $L = \{<\}$  and let N be a  $\omega_1$ -saturated extension of  $\mathbb{Q}$ . Prove that there is an embedding  $f : \mathbb{R} \to N$ . Is it elementary? Can it be an isomorphism?

**Soluzione 1.** From Theorem 9.6 N  $\omega_1$ -saturated implies that N is  $\omega_1$ -rich in the category  $\mathbb{M}$  with elementary maps as arrows. As N is an extension of  $\mathbb{Q}$  there is an elementary embedding  $\mathbb{Q} \hookrightarrow N$ ; this can be regarded as an elementary map  $\mathbb{R} \to N$  of cardinality  $< \omega_1$ . Now using that N is  $\omega_1$  rich together with the finite character of morphisms (c2 from Definition 7.1) we can extend that map to obtain an embedding f of  $\mathbb{R}$  in N that is elementary.

Finally suppose that f is an isomorphism and consider the type  $p(x) = \{n < x : n \in \mathbb{N}\}$  with parameters in  $\mathbb{N}$  of cardinality  $< \omega_1$ . This type is realized in N by  $\omega_1$ -saturation. Now, since f is an isomorphism, there must be a  $x^* \in \mathbb{R}$  that realizes the type in  $\mathbb{R}$  but this is absurd.

**Esercizio 2.** Let M and N be elementarily homogeneous structures of the same cardinality  $\lambda$ . Suppose that  $M \models \exists x \, p(x) \Leftrightarrow N \models \exists x \, p(x)$  for every  $p(x) \subseteq L$  such that  $|x| < \lambda$ . Prove that the two structures are isomorphic. (Hint: see Theorem 7.8)

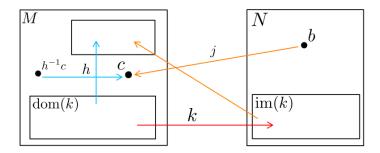
**Soluzione 2.** We would like to do a back-and-forth construction but in order to do so we need a "morphism extension lemma".

Let  $k: M \to N$  be an elementary map such that  $|\operatorname{dom}(k)| < \lambda$ . Now pick  $b \in N$  (that we can assume to be outside  $\operatorname{im}(k)$ ) and consider the type

$$p(x) = \operatorname{tp}_N(\operatorname{im} k, b)$$

where we consider  $\operatorname{im}(k)$  as a tuple. Since  $|\operatorname{im}(k)| < \lambda$  by assumption p(x) is realized by some  $a \in M^{|\operatorname{im}(k)|}$  and  $c \in M$ . Now let  $j : N \to M$  be the map with  $\operatorname{dom}(j) = \operatorname{im}(k) \cup \{b\}$  that sends the tuple  $\operatorname{rm}(k)$  to the tuple a and the element b to the element c. This map is elementary because given  $\varphi(x) \in L$  and  $a \in N^{|x|}$  such that  $N \models \varphi(a)$  we have  $\varphi(x) \in p(x)$  and thus  $M \models \varphi(ja)$ . The composition of elementary maps is elementary so  $j \circ k : M \to M$  is elementary and, by  $\lambda$ -homogeneity, extends to an automorphism  $h : M \to M$ .

Now consider the map  $k' : M \to N$  with  $dom(k') = dom(k) \cup \{h^{-1}c\}$  defined as  $k' = j^{-1} \circ h$ . This is an elmentary map and an extension of k such that  $b \in im(k')$ .



We can now obtain an isomorphism by a back-and-forth construction using the above "morphism extension lemma" and familiar induction techniques.

**Esercizio 3.** Let  $A \subseteq N \models T_{\text{acf}}$  what is the cardinality of  $S_x(A)$ , where |x| = 1? Recall that  $S_x(A)$  is the set of complete types  $p(x) \subseteq L(A)$ , finitely consistent in N.

Answer the same question for  $A \subseteq N \models T_{rg}$ .

## Soluzione 3.

**Algebraically closed fields.** Consider  $\overline{A}$  made of all  $a \in \mathcal{U}$  (with  $\mathcal{U}$  some monster model) that are non-zero and such that there is a polinomial equation  $\varphi_a(x) \in L(A)$  such that  $\varphi_a(a)$  is true. Since  $T_{\rm acf}$  has quantifier elimination we can assume all formulas to be polynomial (un)equations eventually combined with connectives.

Consider  $p(x) \in S_x(A)$  and let a be a realization of p. By completeness unless p is the type of all polynomial unequations there is some polynomial equation  $\varphi_a(x) \in p$  such that  $\varphi_a(a)$  is true and so  $a \in \overline{A}$ . Now if  $b \in \overline{A}$  is another realization of p then  $\varphi_a(b)$  must be true as well. This tells us that the type p can only have a finite number of realizations because if this were not the case then the polynomial given by the equation  $\varphi_a(x)$  would have arbitrarily large degree. This tells us that  $|S_x(A)| \leq |\overline{A}|$ .

Now consider  $a \in \overline{A}$  and let  $p(x) \in S_x(A)$  be the unique (by completeness) type with realization a. By remembering that p(x) has only a finite number of realizations we have  $|\overline{A}| \le |S_x(A)|$ . So, in the end,  $|S_x(A)| = |\overline{A}|$ .

Finally if *A* is infinite then  $|A| = |\overline{A}|$  and if *A* is finite then  $|\overline{A}| = \aleph_0$ .

**Random graphs.** We know that  $T_{rg}$  has quantifier elimination so we can assume that all formulas in our types are of the kind r(x, a) or  $\neg r(x, a)$  for some  $a \in A$ , eventually combined with the binary connectives.

By completeness every type  $p(x) \in S_x(A)$  is completely determined by a binary choice for every  $a \in A$  and thus  $S_x(A)$  is in bijection with the set  $2^A$ . So we conclude  $|S_x(A)| = |2^A|$ .