

Currying

Principles of Functional Programming

Motivation

Look again at the summation functions:

```
def sumInts(a: Int, b: Int) = sum(x => x, a, b)
def sumCubes(a: Int, b: Int) = sum(x => x * x * x, a, b)
def sumFactorials(a: Int, b: Int) = sum(fact, a, b)
```

Question

Note that a and b get passed unchanged from sumInts and sumCubes into sum.

Can we be even shorter by getting rid of these parameters?

Functions Returning Functions

Let's rewrite sum as follows.

```
def sum(f: Int => Int): (Int, Int) => Int = {
  def sumF(a: Int, b: Int): Int =
    if (a > b) 0
    else f(a) + sumF(a + 1, b)
  sumF
}
```

sum is now a function that returns another function.

The returned function sumF applies the given function parameter f and sums the results.

Stepwise Applications

We can then define:

```
def sumInts = sum(x => x)
def sumCubes = sum(x => x * x * x)
def sumFactorials = sum(fact)
```

These functions can in turn be applied like any other function:

```
sumCubes(1, 10) + sumFactorials(10, 20)
```

Consecutive Stepwise Applications

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```

Of course:

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- sum(cube) is therefore equivalent to sumCubes.
- ► This function is next applied to the arguments (1, 10).

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Generally, function application associates to the left:

```
sum(cube)(1, 10) == (sum (cube))(1, 10)
```

Multiple Parameter Lists

The definition of functions that return functions is so useful in functional programming that there is a special syntax for it in Scala.

For example, the following definition of sum is equivalent to the one with the nested sumF function, but shorter:

```
def sum(f: Int => Int)(a: Int, b: Int): Int =
  if (a > b) 0 else f(a) + sum(f)(a + 1, b)
```

Expansion of Multiple Parameter Lists

In general, a definition of a function with multiple parameter lists

$$def f(args_1)...(args_n) = E$$

where n > 1, is equivalent to

$$\mathsf{def}\ \mathsf{f}(\mathsf{args}_1)...(\mathsf{args}_{n-1}) = \{\mathsf{def}\ \mathsf{g}(\mathsf{args}_n) = \mathsf{E}; \mathsf{g}\}$$

where g is a fresh identifier. Or for short:

$$\mathsf{def}\ \mathsf{f}(\mathsf{args}_1)...(\mathsf{args}_{\mathsf{n}-1}) = (\mathsf{args}_\mathsf{n} \Rightarrow \mathsf{E})$$

Expansion of Multiple Parameter Lists (2)

By repeating the process n times

$$def \ f(args_1)...(args_{n-1})(args_n) = E$$

is shown to be equivalent to

$$\mathsf{def}\ \mathsf{f} = (\mathsf{args}_1 \Rightarrow (\mathsf{args}_2 \Rightarrow ... (\mathsf{args}_n \Rightarrow \mathsf{E})...))$$

This style of definition and function application is called *currying*, named for its instigator, Haskell Brooks Curry (1900-1982), a twentieth century logician.

In fact, the idea goes back even further to Schönfinkel and Frege, but the term "currying" has stuck.

More Function Types

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Question: Given,
  def sum(f: Int => Int)(a: Int, b: Int): Int = ...
What is the type of sum?
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  def sum(f: Int => Int)(a: Int, b: Int): Int = ...
What is the type of sum?
Answer:
  (Int => Int) => (Int, Int) => Int
Note that functional types associate to the right. That is to say that
    Int => Int => Int
is equivalent to
    Int => (Int => Int)
```

Exercise

- 1. Write a product function that calculates the product of the values of a function for the points on a given interval.
- 2. Write factorial in terms of product.
- 3. Can you write a more general function, which generalizes both sum and product?