## Section 2 Overview

## Agenda

- Inductive Proof
  - Inductive Proof Steps
  - o Recurrence Relation
- Dynamic Programming
- Optional Problem
  - o Recursive Python
  - o Pascal's Triangle

### **Inductive Proof**

### **Inductive Proof Steps**

Steps of Induction:

- 1. Problem Statement (kind of optional but nice)
- 2. Base Case
- 3. Inductive Hypothesis
- 4. Inductive Leap of Faith (Inductive Step)
- 5. Conclusion

### **Inductive Proof**

# 1 Money

## 1.1 Problem

Prove that  $\sum_{i=1}^{n} [(i) \times (i+1)] = \frac{(n)(n+1)(n+2)}{3}$  where  $n \geq 1$  using a beautiful inductive proof.

### 1.2 Solution

Base Case:

Let 
$$n = 1$$
.  

$$1 \times 2 = 2$$

$$\frac{1 \times 2 \times 3}{3} = 2$$

$$2 = 2$$

Inductive Hypothesis:

Assume, for n=k, that 
$$\sum_{i=1}^{k} \left[ (i) \times (i+1) \right] = \frac{(k)(k+1)(k+2)}{3}$$

Inductive Step:

By the Inductive Hypothesis,

$$\sum_{i=1}^{k} [(i) \times (i+1)] = \frac{(k)(k+1)(k+2)}{3}$$

Add  $((k+1)\times(k+2))$  to both sides of the equation.

$$\sum_{i=1}^{k} \left[ (i) \times (i+1) \right] + ((k+1) \times (k+2)) = \frac{(k)(k+1)(k+2)}{3} + ((k+1) \times (k+2))$$

Simplify the left side of the equation.

$$\sum_{i=1}^{k+1} \left[ (i) \times (i+1) \right] = \frac{(k)(k+1)(k+2)}{3} + ((k+1) \times (k+2))$$

Simplify the right side of the equation.

$$\sum_{i=1}^{k+1} \left[ (i) \times (i+1) \right] = \frac{(k)(k+1)(k+2)}{3} + \frac{3((k+1)(k+2))}{3}$$

$$\sum_{i=1}^{k+1} [(i) \times (i+1)] = \frac{(k+1)(k+2)(k+3)}{3}$$

$$\sum_{i=1}^{k+1} [(i) \times (i+1)] = \frac{(k+1)((k+1)+1)((k+1)+2))}{3} \quad \Box$$

#### **Recurrence Relation Solution**

#### RECURRENCE

The game of Hanoi Tower is to play with a set of disks of graduated size with holes in their centers and a playing board having three spokes for holding the disks.

The object of the game is to transfer all the disks from spoke A to spoke C by moving one disk at a time without placing a larger disk on top of a smaller one. The minimal number of moves required to solve the problem with n disks can be modeled by the following recurrence relation:

$$a_n = 2a_{n-1} + 1, n \ge 1$$
  
 $a_1 = 1$ 

Plug and Chug Solution:

$$a_1 = 1$$

$$a_2 = 2a_{2-1} + 1 = 2*1 + 1 = 3$$

$$a_3 = 2a_{3-1} + 1 = 2*3 + 1 = 7$$

$$a_4 = 2a_{4-1} + 1 = 2*7 + 1 = 15$$

$$a_n = 2a_{n-1} + 1 = 2^n - 1, n \ge 1$$

## **Dynamic Programming**

Convert some amount of money M into a given a list of denominations (decreasing order), using the smallest possible number of coins. Return the smallest number of coins (not the denominations used)

Greedy approach:

**Input:** An amount of money, and an array (denoms) of d denominations = (c1, c2, ..., cd), in decreasing order (c1>c2>...cd).

Output: Min number of coins to make amt
greedy change(amt, denoms):t

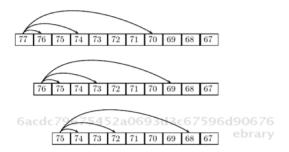
```
remainder = amt //remaining amount to make change for
pieces = [ ] //output array of denominations (number of each denomination)
for k = 0 to denom.length: //for each denomination (starting with largest)
    pieces[k] = remainder/denoms[k]
    remainder = remainder % denoms[k] //practice mod!
return sum(pieces)
```

## **Dynamic Approach:**

Looks at the minimum of previous choices then adds the denomination needed.

For example, if you were making 77 cents from 1, 3, and 7 cents. 77 cents would depend on:

- 1. The best combination for 77-1 = 76 cents, plus 1 cent coin
- 2. The best combination for 77-3 = 74 cents, plus a 3 cent coin
- 3. The best combination for 77-7 = 70 cents, plus a 7 cent coin



**Figure 6.1** The relationships between optimal solutions in the Change problem. The smallest number of coins for 77 cents depends on the smallest number of coins for 76, 74, and 70 cents; the smallest number of coins for 76 cents depends on the smallest number of coins for 75, 73, and 69 cents, and so on.

## **Optional Problems**

### Sum All

Given a positive integer x, recursively find the sum of all numbers from 1 to x.

```
def sumAll(x):
    '''Function that recursively sums positive integers less
    than or equal to x'''
    if x==0:
        return 0
    else:
        return sumAll(x-1)+x
```

### Pascal's Triangle

Write an algorithm that takes an integer n and returns the nth row of Pascal's Triangle (see here for an explanation: <a href="https://en.wikipedia.org/wiki/Pascal%27s">https://en.wikipedia.org/wiki/Pascal%27s</a> triangle)

```
def pascal(n):
   11 11 11
    input: an int, n, the line of pascal's triangle that you
want to return
    output: a list representing the nth line of pascal's
triangle. If n is less than 2, then the function should return
[1]
    .....
    if n < 2:
        return [1]
    prev line = pascal(n-1)
    line = [prev line[i]+prev line[i+1] for i in
range(len(prev line)-1)]
    line.insert(0,1)
    line.append(1)
    return line
```