

04/19/23 - Symbolic Dynamics - Prof Wolf

Final:

- informal oral exam
- favorite proof, idea how to prove it, where you might go from there
- how you liked the class

An entire class of Ergodic theory in one lecture

$$X \xrightarrow{\quad} h_{\text{top}} \quad \begin{array}{l} \text{general shift space} \\ \uparrow \\ \text{only computable at } X \text{ where entropy } \geq 0. \end{array}$$

What about more specific shift spaces? e.g. coded shift spaces
if a generating set \mathcal{G} $\rightarrow X(\mathcal{G})$ a coded shift
notion of uniquely representable coded shifts.

Consider $\mathcal{U} = \{G\}$ generating sets: $X(G)$ uniquely representable

$$\text{map } G \mapsto X(G) \rightarrow h_{\text{top}}(X(G)) \text{ for } G \in \mathcal{U}.$$

in general this isn't computable, but there is a classification
of some that are computable.

→ to do this classification, we need a lot of ergodic theory.

Let (X, d) a compact metric space, $f: X \rightarrow X$ continuous. $\mu(X) = 1$.
invariant probability measures: $\mathcal{M} = \{ \mu : \mu \text{ is } f\text{-invariant Borel prob. measure} \}$
 $\mu(f^{-1}(A)) = \mu(A) \forall A \in \mathcal{B}(X)$

endowed w/ weak* topology $\phi \in C(X, \mathbb{R}) \quad \mu \mapsto \int \phi d\mu$

↪ look for smallest topology that continuous functions make all these functionals continuous.

in weak*, things behave a little confusingly

$$\mu_n \rightarrow \mu \text{ iff. } \int \phi d\mu_n \rightarrow \int \phi d\mu \quad \forall \phi \in C(X, \mathbb{R})$$

fact: \mathcal{M} is compact in this topology

does such a measure exist? i.e. is $\mathcal{M} \neq \emptyset$?

Thm: (Krylov-Bogoliubov)
 $\mathcal{M} \neq \emptyset$.

proof: (idea) Pick $x \in X$. let $\{\varphi_m\}$ be dense in $C(X, \mathbb{R})$ (using the sup norm)

$$\text{fix } m \in \mathbb{N}. \Rightarrow \frac{1}{n} \sum_{k=0}^{n-1} \varphi_m(f^k(x))$$

so $\exists n_k$ s.t. $\lim_{k \rightarrow \infty} \frac{1}{n_k} \sum_{k=0}^{n_k-1} \varphi_m(f^k(x))$ converges

(standard existence of a convergent subsequence)

now we need convergence for every m

using a diagonal argument, $\exists n_k$ s.t. $\lim_{k \rightarrow \infty} \frac{1}{n_k} \sum_{t=0}^{n_k-1} \varphi_m(f^t(x))$ for $\forall m \in \mathbb{N}$

$$= J(\varphi_m).$$

since φ_m is dense, and this limit exists if φ_m , we can prove it converges for all φ . (by "standard arguments") $J(\varphi) \exists \forall \varphi \in C(X, \mathbb{R})$. ■

note $J: C(X, \mathbb{R}) \rightarrow \mathbb{R}$ is linear, bounded, positive.

$$\Rightarrow \exists \mu \text{ Borel probability measure s.t. } J(\varphi) = \int \varphi d\mu \quad \forall \varphi$$

} Riesz Representation Theorem

One can now use the construction to show that μ is invariant.

Simplest example of a uniquely ergodic system:

- irrational rotation \hookrightarrow unique measure
(measure is the Lebesgue measure)

def: M_X convex if $M_1, M_2 \in M_X, \alpha \in [0, 1] \Rightarrow \alpha M_1 + (1-\alpha) M_2 \in M_X$

if you have a compact, convex top. space, you can define an extreme point.

def: $M_E \subseteq M_X$ is the subset of extreme points in M_X .

def: $\mu \ni$ ergodic if $A = f^{-1}(A)$ means $\mu(A) = 0$ or $= 1$.

turns out M_E is also the set of all Ergodic measures.
(otherwise you could decompose μ into nontrivial μ_1 and μ_2)

Ergodic decomposition

$\mu \in M_X \Rightarrow \exists \lambda_\mu$ borel prob. measure in M_X s.t.
 $\lambda(M_E) = 1$ and $\int \phi d\lambda = \int \int \phi d\mu d\lambda$ $\forall \phi \in C(X, \mathbb{R})$

(the non-ergodic measures form a set of measure 0 for λ)

Poincaré Recurrence Theorem:

Let $\mu \in M_X$, $A \in \mathcal{B}(X)$ with $\mu(A) > 0$.

$\Rightarrow \mu(\{x \in A : \exists N \text{ s.t. } f^n(x) \notin A \text{ if } n \geq N\}) = 0$

the set of points that only return a finite number of times

equivalently:

$\mu(\{x \in A : f^n(x) \in A \text{ infinitely many times}\}) = \mu(A)$

\hookrightarrow what this lets us do is look at what

Birkhoff Ergodic Theorem

$\mu \in M_X$, let $\phi \in C^1(X, \mathbb{R})$. Then the time average, given by
 $\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k=0}^{n-1} \phi(f^k(x)) = \phi_f(x)$ exists for x "μ almost everywhere"

also $\phi_f \in C(X, \mathbb{R})$. furthermore if $\mu \in M_E$ then ϕ_f is constant a.e.

and $\underbrace{\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k=0}^{n-1} \phi(f^k(x))}_{\text{time average}} = \underbrace{\int \phi d\mu}_{\text{space average}}$ for a.e. $x \in X$

(or: let $B(\mu)$ the basin of $\mu = \{x \in X : \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k=0}^{n-1} S^k \phi(x) = \mu\}$.
If $\mu \in M_E$ then $\mu(B(\mu)) = 1$

def: let $\mu, \nu \in M_X$. Define $W_1(\mu, \nu) = \sup \left| \int \phi d\mu - \int \phi d\nu \right| \phi \in \text{Lip}(X, \mathbb{R})$
the Wasserstein-Kantorovich distance.
 \hookrightarrow this metric induces weak*

Let X be a shift space, $\mu \in M_X$.

def: $\lim_{n \rightarrow \infty} -\frac{1}{n} \sum_{\tau^n t \in \Sigma_n(X)} \mu(\tau^n) \log \mu(\tau^n)$

take all words of length n , take the measure of the cylinder, then $\mu(\cdot) \log \mu(\cdot)$

$\log x$
(has great properties)

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(2)

Thm: (Kakutani-Brown) Let $\mu \in M_x$. Then the limit

$$\lim_{r \rightarrow 0} \lim_{n \rightarrow \infty} -\frac{1}{n} \log \mu(B_n(x, r)) = h_\mu(x) \text{ exists } \mu\text{-a.e. } x.$$

$$\hookrightarrow \text{recall } = \max_{k \in \{0, \dots, n-1\}} d(f^k(x), f^k(r))$$

idea: for this to be positive the measure must live on points that are pushed apart.

also $x \mapsto h_\mu(x)$ is M -measurable and if $\mu \in M_E$, $h_\mu(\cdot)$ is cont. μ -a.e., and $h_\mu(f) = \int h_\mu(x) d\mu(x)$

Connection between metric and topological entropy:

Variational Principle:

$$h_{\text{top}}(f) = \sup_{\mu \in M_x} h_\mu(f)$$

and if $h_\mu(f) = h_{\text{top}}(f)$ then μ is a measure of maximal entropy (mme).

in some systems there is no mme! (bc sup).
but in symbolic systems there is.

shift maps are always expansive.

n.z.e property: f expansive means $\mu \mapsto h_\mu(f)$ is upper semi continuous

Let $X = X(f)$ be a coded shift.

$$h_{\text{seq}}(X(f)) = \sup \{ h_\mu(\sigma) : \mu \in M_x, \mu(X_{\text{seq}}) = 1 \}$$

$$h_{\text{res}}(X(f)) = \sup \{ h_\mu(\sigma) : \mu \in M_x, \mu(X_{\text{res}}) = 1 \}$$

$$M_m \rightarrow M \Rightarrow$$

$$\limsup_{n \rightarrow \infty} h_\mu|f \leq h_\mu(A)$$

(in the limit you can only go up).

Idea: entropy computable if $h_{\text{seq}} > h_{\text{res}}$, and sometimes if $h_{\text{seq}} = h_{\text{res}}$.
Where does the entropy live? Is it on the concatenations? Then a computer can find it.
(It's in the limit, a computer cannot find it.)

works if there's a mme that lives on h_{seq} .