

3.

a) i. Repeatedly multiplying by  $x$ . Each  $x$  has an error  $\epsilon_x$ . Since  $n$  just counts how many operations we do, it doesn't necessarily have any error.

However, if it's a parameter we input, then in principle it could.

Thankfully, error propagation in multiplication is additive. So,

$$\begin{aligned} [x(1+\epsilon_x)]^{n(1+\epsilon_n)} &\approx [x^n(1+n\epsilon_x)]^{(1+\epsilon_n)} \\ &\approx [x^n(1+n\epsilon_x)](1+\epsilon_n) \\ &\approx x^n(1+\epsilon_n+n\epsilon_x) \end{aligned}$$

So the relative error is:

$$\boxed{\epsilon_{x^n} = \epsilon_n + n\epsilon_x}$$

ii. What's the error in  $\log x$ ?

$$\begin{aligned} \text{Well, } \log x(1+\epsilon_x) &= \log x + \log(1+\epsilon_x) \\ &\approx \log x + \epsilon_x \end{aligned}$$

multiply by  $n(1+\epsilon_n)$ .

$$n(1+\epsilon_n)\log x(1+\epsilon_x) = n\log x + n\epsilon_x + n\epsilon_n\log x$$

We needed to get an  $\epsilon$  in  $\log x$ .

$$\text{So } \log x(1+\epsilon) = [\log x](1+\epsilon) + \epsilon_x = \log x + \epsilon\log x + \epsilon_x$$

Then we have:

$$f(n\log x) = n\log x + n\epsilon_x + n\epsilon_n\log x + n\epsilon\log x$$

$$\text{Now, } f(e^{n\log x}) = (1+\epsilon) e^{n\log x + n(\epsilon_n + \epsilon)\log x + n\epsilon_x}$$

$$= e^{n \log x} (1 + \epsilon) (1 + n\epsilon_x + n(\epsilon + \epsilon_n) \log x)$$

$$= e^{n \log x} (1 + \epsilon + n\epsilon_x + n(\epsilon + \epsilon_n) \log x)$$

$$\Rightarrow \Sigma_{\text{tot}} = \epsilon + n\epsilon_x + n(\epsilon + \epsilon_n) \log x$$

How does this compare to  $\epsilon_n + n\epsilon_x$ ?  
 If all errors are of the same magnitude, it's  
 Surely worse.

$$\epsilon + n\epsilon < \epsilon + n\epsilon + 2n\epsilon \log x$$

Therefore, (i) is preferable to (ii).

As  $\epsilon_x n$  comes into each side, it all hinges on  $\epsilon_n$  and  $\log x$ . When  $\log x$  is small, then  $\Sigma_{\text{tot}} - n\epsilon_x \approx \epsilon$ , while

$$\epsilon_x n - n\epsilon_x \approx \epsilon_n. \quad \epsilon_n \text{ may often be zero.}$$

If it is indeed zero, then  $e^{n \log x}$  is better. Otherwise they're similar. Now, if  $\log x$  is fairly large, then  $e^{n \log x}$  is much worse than  $x^n$ . Assume  $\Sigma_{\text{tot}} \approx n(\epsilon + \epsilon_n) \log x \gg \epsilon_n + n\epsilon_x$

We need quite a large  $x$  for this to happen.

b) the advantage that  $e^{a \log x} = x^a$  has is that it works for non-integers. Let's assume  $a$  is not necessarily an integer.

$$\begin{aligned} \text{i) } f(x^a) &= [e^{a \log x}]^{(1 + \epsilon_a + \epsilon)} \quad \rightarrow \text{from the log.} \\ &= e^{a \log x (1 + (\epsilon_a + \epsilon))} \\ \Rightarrow \quad \boxed{\epsilon_{\text{tot}} &= (\epsilon_a + \epsilon) a \log x} \end{aligned}$$

$$\begin{aligned} \text{(ii) } f(x^a) &= e^{a(1+\epsilon) \log x (1+\epsilon_x)} \\ &= e^{a \log x + \epsilon \log x + a \epsilon_x} \\ &= e^{a \log x} e^{\epsilon \log x + a \epsilon_x} \\ &\approx e^{a \log x} (1 + \epsilon \log x + a \epsilon_x) \end{aligned}$$

$$\Rightarrow \quad \boxed{\epsilon_{\text{tot}} = a \epsilon_x + \epsilon \log x}$$

We technically must add  $\epsilon$  to each of these for the error in exp, but that won't change our comparison.

In case (i), if  $a \gtrsim 1/(\epsilon_a + \epsilon)$  then the error could be substantial (provided  $\log x \gtrsim 1$ )

In case (ii), if  $a \gtrsim 1/\epsilon_x$ , we could have substantial error.  $\log x$  will never be huge, so  $a$  is the thing to worry about.