

# **Computer Networks - Xarxes de Computadors**

#### **Outline**

- Course Syllabus
- Unit 1: Introduction
- Unit 2. Network applications
- Unit 3. IP Networks
- Unit 5. TCP
- Unit 5. LANs



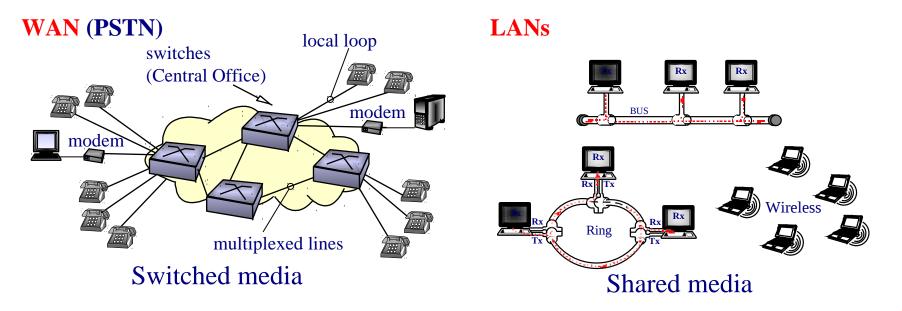
#### **Outline**

- Introduction
- IEEE LAN Architecture
- Ethernet
- Ethernet Switches
- Wireless LANs



#### Unit 5. Local Area Networks, LANs Introduction – WAN and LAN differences

- WANs:
  - Main goal: scalability.
  - Switched network with mesh topology.
- LANs:
  - Multy-access network with shared media.
  - A Medium Access Control (MAC) protocol is needed.





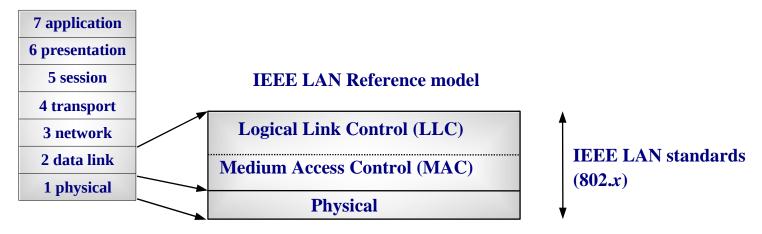
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# Unit 5. Local Area Networks, LANs IEEE LAN Architecture

#### **OSI Reference model:**



- **LLC** sublayer (802.2):
  - Common to all 802.x MAC standards.
  - Define the interface with the upper layer and specifies several services (operational modes):
    - (i) unacknowledged connectionless, (ii) connection oriented, (iii) acknowledged connectionless.
- MAC sublayer:
  - Define the medium access protocol. It is different for each LAN technology.



# Unit 5. Local Area Networks, LANs IEEE LAN Architecture – IEEE 802 standards (some)

802.1: LAN/MAN architecture.

802.2 Logical Link Control (LLC)

802.3 Ethernet

802.4 Token Bus

802.5 Token Ring

802.8 FDDI

802.11 WiFi: Wireless LANs.

802.15 Personal Area Networks or short distance wireless networks (WPAN)

802.15.1 Bluetooth

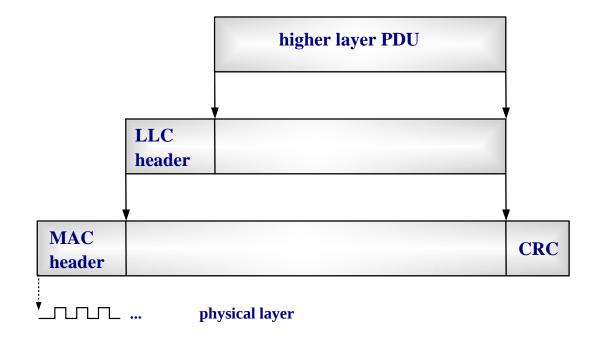
802.15.4 low data rate and low cost sensor devices

802.16 WiMAX: broadband Wireless Metropolitan Area Networks.

See: http://grouper.ieee.org/groups/802/1, 2, ...

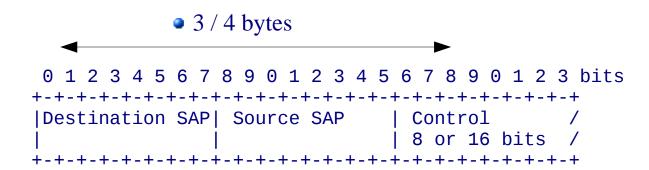


#### **IEEE LAN Architecture – LAN encapsulation**





#### Unit 5. Local Area Networks, LANs IEEE LAN Architecture – LLC header



- Service Access Point (SAP): Identifies the upper layer protocol.
- Control: Identifies the frame type. It can be 8 or 16 bits long, 8 bits for unnumbered frames (used in connectionless modes).

SAP (hex)	Protocol				
06	ARPANET Internet Protocol (IP)				
08	SNA				
42	3IEEE 802.1 Bridge Spanning Tree Protocol				
98	ARPANET Address Resolution Protocol (ARP)				
AA	SubNetwork Access Protocol (SNAP)				
E0	Novell Netware				
F0	IBM NetBIOS				
FF	FF Global LSAP				

Example of some IEEE SAP values.

**SNAP: used in TCP/IP** 



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#### Unit 5. Local Area Networks, LANs Ethernet – Introduction

- Designed by Bob Metcalfe at Xerox in mid-70s.
- Initially was commercialized by Digital, Intel and Xerox consortium (DIX).
- Ethernet was standardized by IEEE (802.3) in 1983.
- Nowadays Ethernet is the leading LAN technology. There are numerous Ethernet standards with different transmission mediums, and line bitrates. There are several active Ethernet working groups inside IEEE 802.3.



# Unit 5. Local Area Networks, LANs Types of MACs

- Token Passing:
  - Only the station having the token can transmit. After transmission the token is passed to another station.
  - Examples: FDDI and Token-Ring
- Random:
  - There is no token. Instead, there is a non null collision probability. In case of collision, the frame is retransmitted after a random *backoff* time.
  - Examples: Ethernet



# Unit 5. Local Area Networks, LANs Carrier Sense Multiple Access/Collision Detection (CSMA/CD)

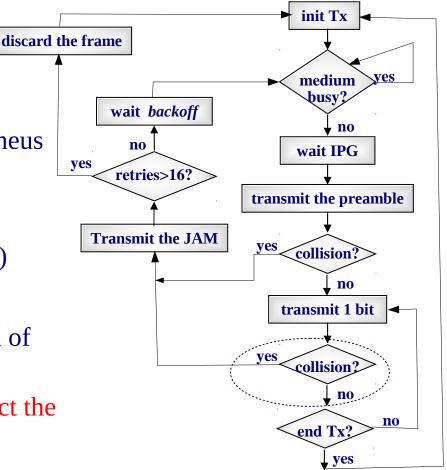
- Is a random MAC where the stations "listen" the medium (*carrier sense*) before transmission.
- When the medium is becomes free the frame is transmitted immediately, and the medium is listened to detect collisions.
- In case of collision, the frame is retransmitted after a random *backoff* time.



# Unit 4. Local Area Networks, LANs Ethernet – CSMA/CD Ethernet protocol (simplified)

#### Legend:

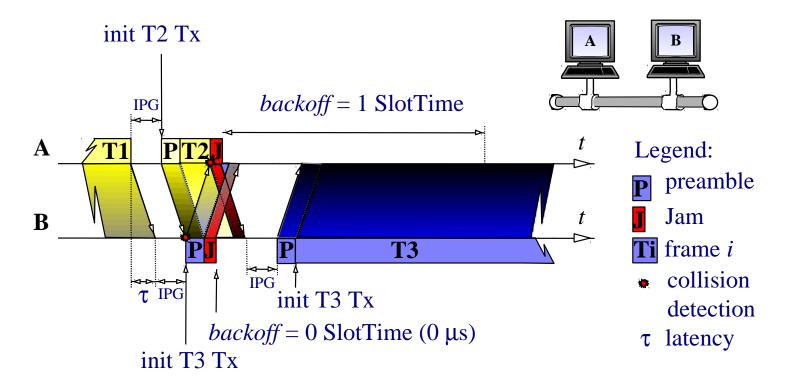
- InterPacket Gap (IPG): 96 bits.
- JAM: 32 bits that produce an erroneus CRC.
- backoff =  $n T_{512}$
- $T_{512}$ : *SlotTime* (51,2 µs at 10 Mbps)
- $n = \text{random}\{0, 2^{\min\{N, 10\}}-1\},$ 
  - *N*: number of retransmission of the same frame (1, 2...)
- The transmitting station must detect the collision (no ack is sent).





# Unit 4. Local Area Networks, LANs Ethernet – Collision example

Stations A y B have frames ready to Tx:



NOTE: The preamble is not interrupted in case of collision, and the JAM is Tx immediately after.



# Unit 5. Local Area Networks, LANs Ethernet – Half Duplex and full-duplex

- Half Duplex: Using CSMA/CD only one NIC can be simultaneously transmitting into the medium.
- Full Duplex: When 2 Ethernet NICs are connected point-to-point, some Ethernet standards allow a full-duplex Tx, .
- Ethernet NICs have an auto-negotiation mechanism to detect the full-duplex availability.
- In full-duplex mode Ethernet NICs deactivate CSMA/CD (no collisions can occur).



#### **Ethernet – Frames**

Ethernet II (DIX):

• IEEE 802.3

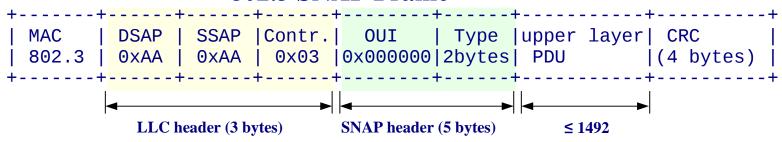
- Preamble: Give time to detect, synchronize and start reception.
- Type: Identifies the upper layer protocol (IP, ARP, etc. RFC 1700, Assigned numbers). This value is always > 1500.
- Length: Payload size (0~1500).



## Unit 5. Local Area Networks, LANs Ethernet – IEEE Sub-Network Access Protocol (SNAP)

- Allows the specification of protocols, and vendor-private identifiers, not supported by the 8-bit 802.2 Service Access Point (SAP) field.
- It is used to encapsulate TCP/IP protocols over IEEE 802.2 with OUI=0x000000 and Type equal to the RFC 1700 (used for DIX).

#### **802.3 SNAP Frame**



• Note: The MSS indicated by TCP would be of 1460 if DIX, and 1452 if IEEE encapsulation is used.



# Unit 5. Local Area Networks, LANs Ethernet – Different Ethernet Standards (some)

Commercial name	bps	Standard y		year Name	Cabling	UTP/OF Pairs	Connector	Codification	segment distance*	
			year						Half duplex	Full duplex
Ethernet	10Mbps	802.3	1983	10Base5	Coax-thick	-	AUI	Manchester	500m	n/a
		802.3a	1985	10Base2	Coax-thin	-	BNC	Manchester	185m	n/a
		802.3i	1990	10BaseT	UTP-cat.3	2	RJ45	Manchester	100m	100m
		802.3j	1993	10BASE-FL	FO	2	SC	on/off Manchester	2000m	>2000m
Fast Ethernet	100Mbps	802.3u	1995	100BaseTX	UTP-cat.5	2	RJ45	4B/5B	100m	100m
		802.3u	1995	100BaseFX	FO	2	SC	4B/5B	412m	2000m
		TIA/EIA-785	1999	100BaseSX	FO/led	2	SC	4B/5B	300m	300m
Gigabit-Eth.	1Gbps	802.3z	1998	1000BaseSX	FO	2	SC	8B/10B	275-316m	275-550m
		802.3z	1998	1000BaseLX	FO	2	SC	8B/10B	316m	550-10000m
		802.3z	1998	1000BaseLH	FO	2	SC	8B/10B	n/a	100km
		802.3ab	1999	1000BaseT	UTP-cat. 5e	4	RJ45	PAM5	100m	100m
10Gigabit- Eth.	10Gbps	802.3ae	2002	10GBASE-CX4	InfiniBand	4	CX4	8B/10B	n/a	15m
		802.3ae	2002	10GBASE-SR	FO	2	SC	64B/66B	n/a	26-300m
		802.3ae	2002	10GBASE-LR	FO	2	SC	64B/66B	n/a	10km
		802.3ae	2002		FO	2	SC		n/a	

<sup>\*</sup>With OF the distance depends on the OF type.

**xBasey** 



#### Unit 5. Local Area Networks, LANs Ethernet – Different Ethernet Standards

#### **Denomination:**

#### Line bitrate:

- 10: 10 Mbps
- 100: 100 Mbps
- 1000: 1000 Mbps (1 Gbps)
- 10G: 10 Gbps

Base band signal.

**Broad**: translated band signal.

#### **Various meanings:**

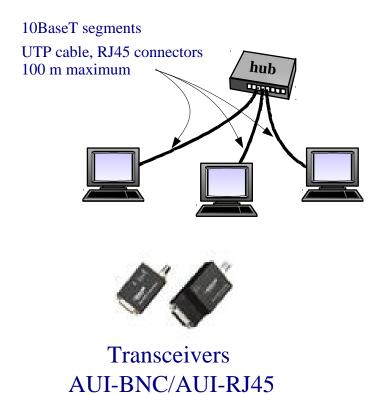
- Number: Maximum segment distant in hundreds of m.
- Reference to the medium type:
  - T: UTP
  - F: Optical Fiber
  - Other:
  - T4: Uses 4 UTP pairs.
  - TX: Full Duplex
  - ...

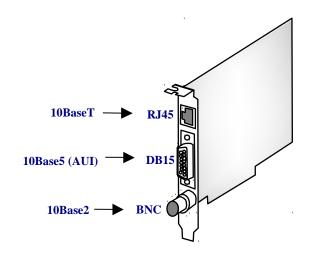


#### Unit 5. Local Area Networks, LANs Ethernet – Different Ethernet Standards: 10BaseT

#### 1990. Cable UTP-cat 3.

- Hub: Is a multi-port repeater (layer 1).
- The signal received in 1 port is retransmitted by all the others.





NIC "combo": Supports 10Base5, 10Base2, 10BaseT



#### Unit 5. Local Area Networks, LANs Ethernet – Different Ethernet Standards: after 10BaseT

#### All standards use UTP o OF (except 10GBaseCX4):

- Fast Ethernet (1995). 100BaseTX: UTP-cat. 5
- Gigabit Ethernet (1998). 1000BaseT: UTP-cat 5e
- 10Gigabit Ethernet (2002). Uses optical fiber. The only copper standard is Infiniband with segment size  $\leq 15$ m.



NIC 10/100 – RJ45 10BaseT-100BaseTX \$11.99



NIC 10/100/1000 - SC 10BaseFL-100BaseFX-1000Base-SX \$151



NIC 10Gbps – CX4 10GBaseCX4 \$795



Infiniband cable with CX4 connectors



#### **Outline**

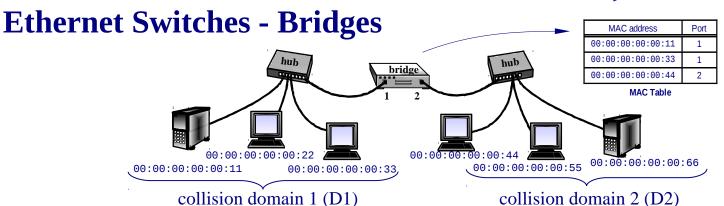
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# Unit 5. Local Area Networks, LANs Ethernet Switches - Introduction

- Hub problem: If many stations are connected, may be inefficient due to collisions.
- Solution: bridges and switches.
- Ethernet bridge:
  - "plug and play" layer 2 device.
  - In each port there is a NIC in "promiscuous" mode: Capturing all frames.
  - The source address is used to "learn" which MAC is present in each port (MAC table). Each entry has the MAC and the port numbers.
  - The destination MAC is used to decide whether the frame needs to be retransmitted by another port.
  - Segments the "collision domain".





#### How the bridge works:

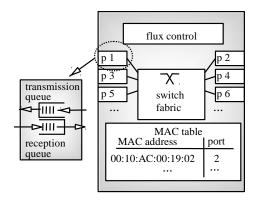
- If a frame is received with a source address on in the MAC table, it is added (*learning bridge*).
- If a frame from D1 is received with a destination address: (i) is in D2, (ii) it is not in the table, (ii) it is broadcast: It is sent into D2 (*flooding*).
- If it is received a frame from D1 addressed to another station from D1, it is discarded (*filtering*).
- The entries have an *aging timer*. Each time an entry is used, it is refreshed. If the aging timer expires, the entry is removed.

#### **Advantages:**

- Segments the collision domain (less collisions).
- Clients in D1 and D2 can simultaneously access their servers.



#### Unit 5. Local Area Networks, LANs Ethernet Switches - Switch Architecture



Switch#show mac-	address-table
Address	Dest Interface
00D0.5868.F583 00E0.1E74.6ADA 00E0.1E74.6AC0 0060.47D5.2770 00D0.5868.F580	FastEthernet 2 FastEthernet 1 FastEthernet 1 FastEthernet 3 FastEthernet 5





MAC Table in a CISCO Switch

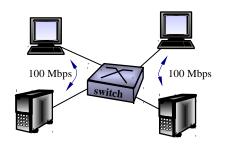
Edge and backbone CISCO switches.

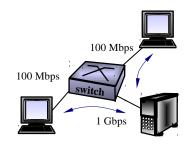
#### How the switch works:

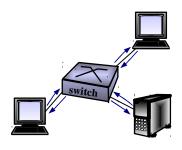
- It is equivalent to a "multiport bridge".
- When a frame is received with a source address not in the table, it is added.
- If a frame is received with a destination address: (i) not in the table, (ii) broadcast or multicast: copy the frame in all transmission buffer of the other ports (flooding).
- If a frame is received with the address from another port: It is switched as fast as possible the transmission buffer of that port.
- If receives a frame addressed to another station from the same port, it is discarded (filtering).



# Unit 5. Local Area Networks, LANs **Ethernet Switches - Switch Capabilities**







Simultaneous Transmissions Ports with Different bitrates

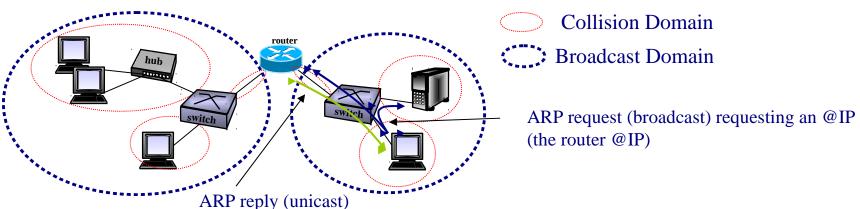
Full Duplex Ports

- Each port is different a collision domain (less collisions).
- Different ports can be simultaneously Tx/Rx.
- Ports can have different bitrates.
- Ports may be full-duplex (usable if only one host is connected).
- There can be ports simultaneouly in half or full duplex mode.
- Security: Stations can only capture the traffic of their collision domain.



#### Unit 5. Local Area Networks, LANs Ethernet Switches - Broadcast and Collision Domains

- Broadcast Domain: Set of stations that will received a broadcast frame sent by any of them.
- Unless Virtual LANs are used, a switch does not segment the broadcast domain.
- A router segment the broadcast domain.
- The broadcast reachability is important because allows reaching stations having one hop connectivity (with ARP).



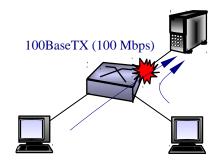
ARP cannot solve an @IP out of the broadcast domain. To leave the broadcast domain a router is required.



#### Unit 5. Local Area Networks, LANs Ethernet Switches – Flux Control

- Switch Flux Control: Consists of adapting the rate at which the switch receives the frames, and the rate at which the switch can send them.
- Examples:





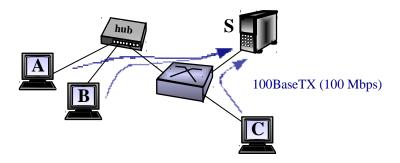
If no flux control is used, frames could be lost by buffer overflow.

- Flux control techniques (back pressure):
  - Jabber signal (half duplex): The switch sends a signal into the port which need to be throttled down, such that CSMA see the medium busy.
  - Pause frames (full duplex): The switch send special *pause frames*. These frames have an integer (2 bytes) indicating the number of slot-times (512 bits) that the NICs receiving the frame must be silent.



# Unit 5. Local Area Networks, LANs Ethernet Switches – Line bitrate sharing

- Hub: If the hub is the bottleneck for all the active ports, the capacity is equally shared between all ports where frames are transmitted.
- Switch: If one congested port is the bottleneck for all ports sending traffic to it, the port bit rate is equally shared between all ports sending traffic to it.
- Example:



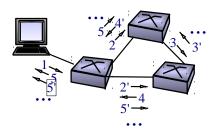
• If A, B and C simultaneously transmit to S: throughput C  $\approx$  100 Mbps / 2 = 50 Mbps throughput A = throughput B  $\approx$  (100 Mbps / 2) / 2 = 25 Mbps



# Unit 5. Local Area Networks, LANs Ethernet Switches – Spanning Tree Protocol (STP)

- The basic principle of the "layer 2 routing" done by Ethernet switches is based on having a unique port to forward the frame towards the destination. Therefore, loops are not allowed.
- In practice loops can appear because:
  - They are introduced by accident.
  - The are desirable to have redundant path (fault tolerance).
- If loops are introduced without protection a *broadcast storm* is produced, and the network blocks:

Frames multiply and remain turning indefinitely in the loop!



#### Other problems:

- Reception of duplicated frames
- MAC Tables instability

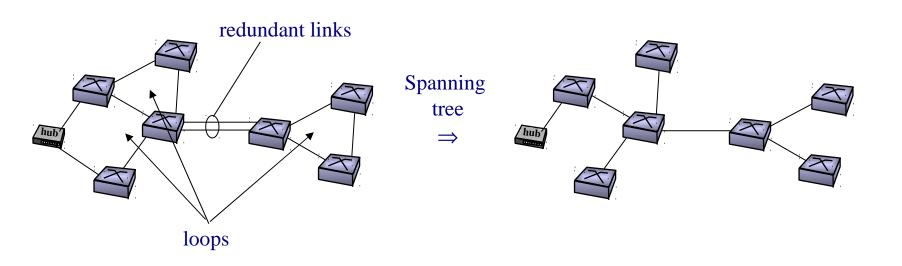


Solution: IEEE 802.1D Spanning Tree Protocol (STP)



# Unit 5. Local Area Networks, LANs Ethernet Switches – Spanning Tree Protocol (STP)

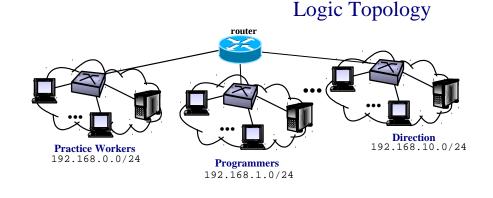
- STP goal: Build a loop free topology (STP-tree) with optimal paths. The ports that do not belong to the STP tree are blocked.
- The switches send 802.1D messages to their neighbors to build up the STP-tree. If the topology changes (e.g. due to a link failure), a new STP-tree is setup.



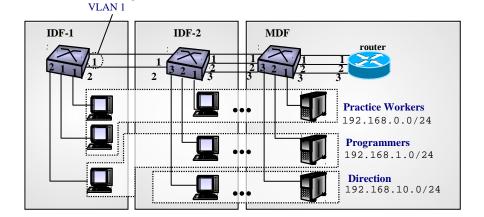


## Unit 5. Local Area Networks, LANs Ethernet Switches – Virtual LANs, VLANs

- Motivation:
- Grouping related servers and hosts in different broadcast domains.
- How VLANs work:
- Each switch port belongs to a VLAN.
- The switch isolates
   different VLANs: The
   switch flooding is done
   only towards the ports of
   the same VLAN. Each
   VLAN is equivalent to a
   different physical switch.
- A router is needed to send traffic to a different VLAN.



Physical Topology



Port configured in



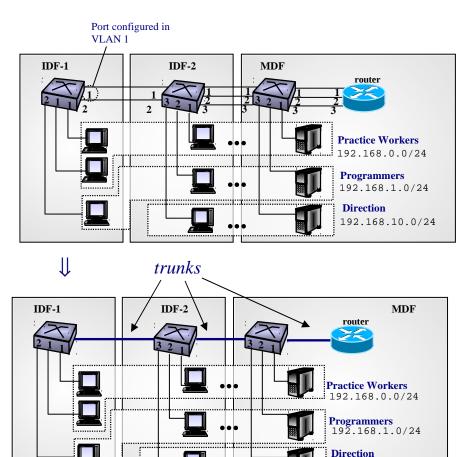
## Unit 5. Local Area Networks, LANs Ethernet Switches – Virtual LANs, VLANs

- Advantages:
- Flexibility of the physical placement of the devices.
- Facilitates the network grow.
- Facilitates the network management: Changing the topology, adding new subnetworks, moving ports from one network to another.
- NOTE: Since each VLAN is a different broadcast domain, usually a different STP instantiation is used for each VLAN. Thus, a different STP-tree is build in each VLAN.



# Unit 5. Local Area Networks, LANs Ethernet Switches – VLAN Trunking

- Problem:
- Why connecting several ports between the same devices?
- Trunking:
- The port configured as trunk belongs to several VLANs (maybe all).
- The traffic sent in one VLAN is also sent to the trunk the VLAN belongs to.
- A tagging mechanism is used in the trunk to discriminate the traffic from different VLANs.



192.168.10.0/24



# Unit 5. Local Area Networks, LANs Ethernet Switches – VLAN Trunking

- Trunking Protocols:
- Inter-Switch Link (ISL). CISCO propietary protocol.
- IEEE-802.1Q.

IEEE-802.3 frame with the 802.1Q tag.

#### Legend:

- *Tag Protocol Identifier* (TPID): Field with the hex. value 8100 for an Ethernet frame.
- *Tag Control Information* (TCI): Contains several fields. The most important is the VLAN ID (12 bits), which identify the VLAN.



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## Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – Brief WLAN History

- 1971: Prof. Norman Abramson develops ALOHANET for the University of Hawaii
- 1990: many companies develop proprietary WLANs products.
- 1996: ETSI approves HIPERLAN/1 and 1997 IEEE approves 802.11
- Late 90 and 2000: Wi-Fi Alliance, tremendous growth of 802.11 products.
- 1999: 802.11a, 802.11b. 2003: 802.11g ...



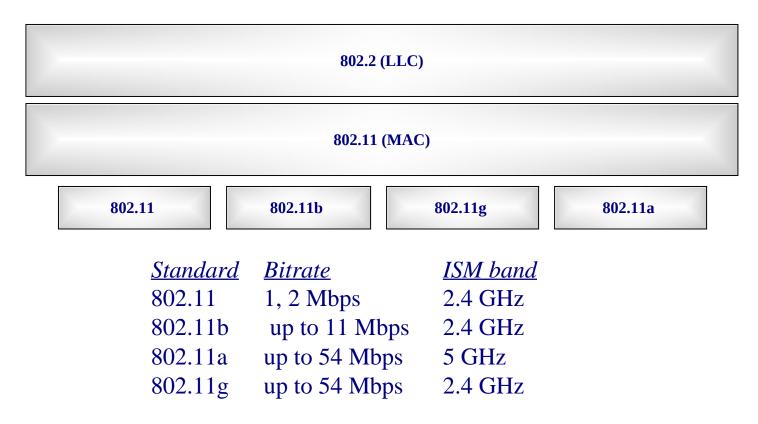




Home made antenna 802.11 Antennas



## Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11

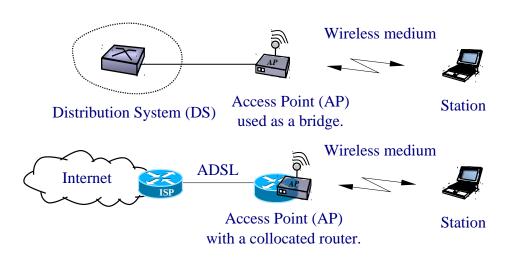


• ISM: Industrial Scientific and Medical. Free band for non commercial usage.



## Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 Components

- Distribution System (DS):
  - Used by APs to exchange frames with one another and with wired networks. (e.g. an ethernet switch).
- Access Point (AP)
  - Simplify communication between stations.
  - All transmissions go through the AP.
  - APs are bridges and may have a collocated router.





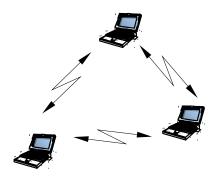
### Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 Components

- Basic Service Set (BSS)
  - Set of stations communicating with each other.
  - Are identified by: (i) a Service Set identifier (SSID), or Network name: String with <32 characters; and (ii) a BSS Identifier (BSSID): 48 bits number.
  - If the network is composed of more than 1 BSS it is called Extended Service Set (ESS).



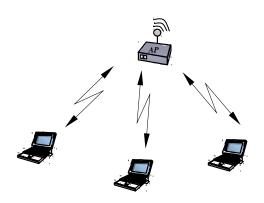
## Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 Components

#### **Independent BSS, IBSS** (ad-hoc mode)

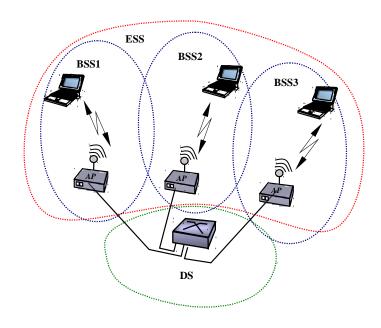


#### **Infrastructure BSS** (*infrastructure mode*)

- An station must associate with an AP.
- All transmissions go through the APs.



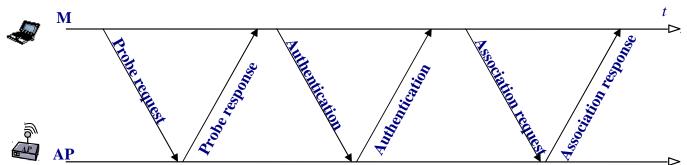
#### **Extended Service Set (ESS)**





# Unit 5. Local Area Networks, LANs 802.11: Protocol description- Components

- Beacons
  - Special frames carrying information related to the BSS (e.g. the BSSID).
  - In infrastructure BSS are sent by the APs, in IBSS there is a contention algorithm for electing the station generating beacons.
  - BSSID are: (i) the MAC@ of the AP in infrastructure BSS, and (ii) the MAC@ of the station generating beacons in IBSS.
- AP Association:
  - Probe
  - Authentication
  - Association





## Unit 5. Local Area Networks, LANs 802.11: Protocol description- Features

- Fragmentation
  - Optional mechanism to reduce the effect of Tx errors. If the frame size is larger than the threshold, it is fragmented into multiple frames.
- Power-saving mechanism
  - Optional mechanism to save battery: The AP sends periodically a TIM (Traffic Information MAP), informing which stations have buffered traffic. The stations wake up at the TIM Tx periods, and request the frames, if any.
- WEP (Wired Equivalent Privacy):
  - Frame payload is encrypted using a 64/128 key.



## Unit 5. Local Area Networks, LANs 802.11: Protocol description- Frames

- Data frames
- Control frames: handle reliable transmission of data frames
  - ACK, RTS, CTS and polling
  - Typical time scales: Frame transmission time (<1ms)
- Management frames: communication between stations and APs
  - Beacons, association, Probe and authentication.
  - Typical time scales: 100 ms, minutes, hours,...



### Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 Addresses

- Designed to be compatible with ethernet.
- 48 bits (6 bytes).
- Use non overlapping ranges with ethernet.
- Broadcast: FF:FF:FF:FF:FF
- The frame may have up to 4 addresses. The meaning of the addresses is specified by the bits to-DS and from-DS of the control.
- The BSSID is always present to identify frames belonging to the BSS.
- When a station is searching for the BSS it uses the broadcast BSSID:
   FF:FF:FF:FF:FF

2	2	6	6	6	2	6	Variable: 0-2312	4	
Frame Contro	Duration	Address 1	Address 2	Address 3	Seq Ctrl	Address 4	Payload	FCS	

**Generic frame format** 

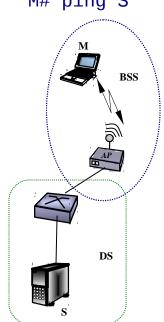


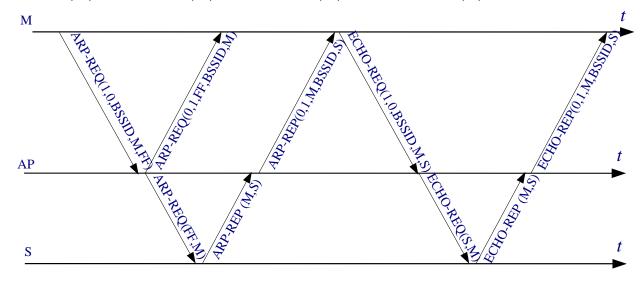
### Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 Addresses

Scenario	Usage	to-DS	from-DS	Address1	Address2	Address3	Address4
STA→STA	Ad-hoc	0	0	DA	SA	BSSID	=
STA→AP	Infrastructure	1	0	BSSID	SA	DA	-
AP → STA	Infrastructure	0	1	DA	BSSID	SA	-
$AP \rightarrow AP$	WDS	1	1	RA	TA	DA	SA

Legend: Destination Address (DA), Source Address (SA), Receiver Address (RA), Transmitter Address (TA)

## Example: M# ping S





#### Legend, frames 802.11:

MESSAGE-TYPE(to-DS, from-DS, Address1, Address2, Address3)

#### **Legend, frames ethernet:**

MESSAGE-TYPE(destination address, source address)

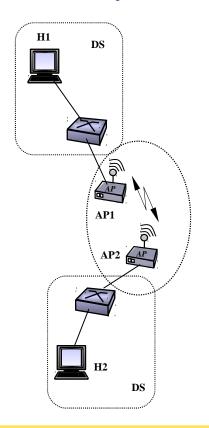
FF is the broadcast address

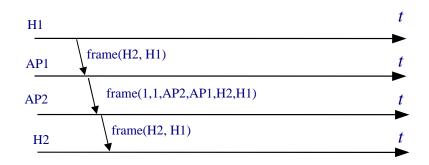


### Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 Addresses

Scenario	Usage	to-DS	from-DS	Address1	Address2	Address3	Address4
STA→STA	Ad-hoc	0	0	DA	SA	BSSID	=
STA→AP	Infrastructure	1	0	BSSID	SA	DA	-
AP → STA	Infrastructure	0	1	DA	BSSID	SA	-
$AP \rightarrow AP$	WDS	1	1	RA	TA	DA	SA

Legend: Destination Address (DA), Source Address (SA), Receiver Address (FA), Transmitter Address (TA)





#### Legend, frames 802.11:

frame(to-DS, from-DS, Address1, Address2, Address3, Address4)

#### **Legend, frames ethernet:**

frame(destination address, source address)



### Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 MAC

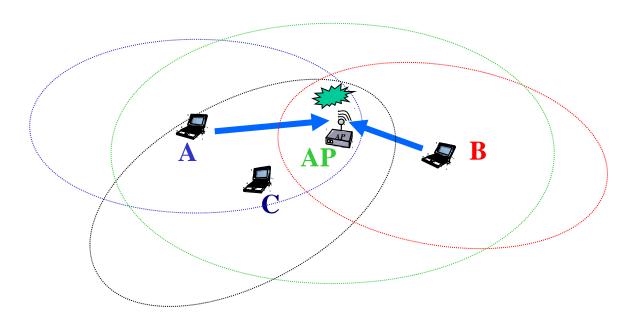
Two Coordination Functions (CF) are defined:

- Distributed CF (DCF):
  - Contention MAC.
  - Best effort service
  - Upon receiving a correct frame, an ack is immediatelly transmitted.
  - Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA): In contrast to CSMA/CD, always wait a random backoff before Tx. Acks are needed to detect whether a transmitted frame collided.
- Optional Point CF (PCF):
  - Contention free MAC built on top of DCF.
  - Centralized polling scheme. The AP poll each PCF station for Tx.
  - A contention free period (CFP) using PCF and a contention period (CP) using DCF follow each beacon.



## Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – Hidden Node Problem

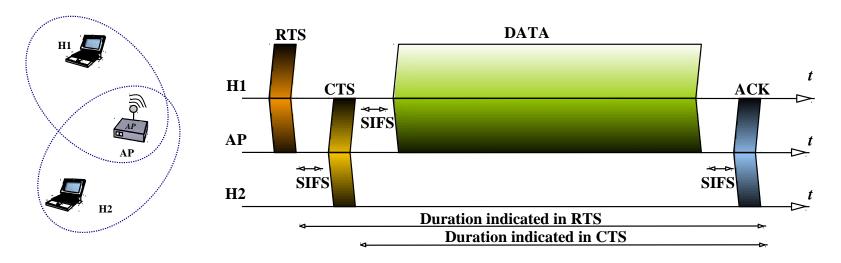
- Node A is in coverage with AP and C
- A and B cannot hear each other
- When A transmits to AP, B cannot detect the transmission using the carrier sense mechanism
- If B transmits, a collision will occur at AP





## Unit 5. Local Area Networks, LANs Wireless LANs (WLANs) – 802.11 RTS/CTS

Optional mechanism to solve the hidden node problem.



- RTS is sent using the basic access mechanism.
- Upon receiving a RTS/CTS, the station set the Network Allocation Vector (NAV) to the indicated duration. While the NAV is non zero, the virtual carrier sensing indicates that the medium is busy.
- RTS/CTS is only used for unicast Tx.
- There is a threshold indicating the minimum frame size for using RTS/CTS.