```
#include <map>
#include <vector>
#include <string.h>
#include <signal.h>
#include <set>
using namespace std;
// Definition of type script_function, that masks a void* (*)(void*
typedef void* (*script_function)(void*);
// Definition of type script_vector, a std::vector of type script_function
typedef std::vector<script_function> script_vector;
// structure: _loop_spec_thread, used to encapsulate a pthread,
             which will run a speculative thread for loop speculation,
             keeping as well some related data.
struct _loop_spec_thread{
        pthread_mutex_t _thread_mutex;
        pthread_t _thread;
        set <void*> written data; //will be used on thread-cancelation,
        //to undo any changes it has made.
        set <void*> _read_data; //will be used on threcad-cancelation,
//to delete it's logs as reader.
        void* (* thread_instructions) (void*);
        void* _const_args;
        bool _commit;//helps to determine if the loop iteration
        //that the loop spec thread represents has committed or not.
        loop spec thread(){
               pthread mutex init (& thread mutex, NULL);
                _commit=false;
        };
// structure: _data_access_log, used to encapsulate the use history of a shared data,
            item: it's writers, readers and previous values. It is through this
             structure that data dependence violation is tracked.
struct _data_access_log{
        unsigned int size;//size of data
        set <int> _readers;
        vector <int> _writers;
        vector <void*> _previous_values; //this will be used to solve anti-dependence
                                         //and output dependence violations (WAR and WAW)
                                         //restoring the data to a value that is safe.
        pthread_mutex_t _data_mutex;
        _data_access_log(){
                pthread_mutex_init (&_data_mutex, NULL);
        template <class T>
        bool _output_violations_are_false(T* data_to_write, int pos){ //this function determines if the output dependence
        //violations are false, that is, if the value to write is the same that is already written by later iterations.
        //Note that in this version this check to avoid unnecessary thread-resetting is only possible for WAW violations, since
        //there is book-keeping of the values that are being written, but not of the values that are being read.
                void* prev_value;
                int first_reader=-2;;
                for (int i=0; i< writers.size(); i++){
                        if (_writers[i]==-1){
                                prev_value=_previous_values[i];
                        else if (_writers[i]>pos){
                                if (data_to_write!=(T*)_previous_values[i]){
                                         if (first_reader==-2){
                                                 if (_previous_values[i]!=prev_value){
                                                        return false;
                                                 first_reader=_writers[i];
                                         else if (first reader== writers[i]) {
                                                 if ( previous values[i]!=prev value){
```

#include <pthread.h>

```
return false;
                                 else{
                                         return false:
                        }
                1
        if (first_reader==-2){
                _writers.push_back(pos);
                __previous_values.push_back((void*)data_to_write);
                return true;
        for (int i=0; i<_writers.size(); i++){</pre>
                if (_writers[i]==first_reader){
                        _previous_values[i]=(void*)data_to_write;
        }
        writers.push back (pos);
        _previous_values.push_back(prev_value);
        return true;
};
vector <int> _cancel_higher_readers(int pos){
//if used before a write_data and the return vector is not empty, a true-dependence violation has ocurred (RAW).
//and the higher readers have to be restarted.
        vector <int> to_cancel;
        set <int>:: iterator it=_readers.begin();
        while ((!_readers.empty()) && it!=_readers.end()){
                if ((*it)>pos){
                        to_cancel.push_back(*it);
                        _readers.erase(*it);
                        it= readers.begin();
                else {
                        it++;
        if ((!to cancel.empty())&&(! writers.empty())){
                for (unsigned int i=0; i<to_cancel.size(); i++){
                        unsigned int j=0;
                         while((! writers.empty())&& j< writers.size()){</pre>
                                 if (to_cancel[i]==_writers[j]){
                                         _writers.erase(_writers.begin()+j); //all the writings of the higher readers on this data
                                         __previous_values.erase(_previous_values.begin()+j);
                                         j=0;
                                 1
                                 else
                                         j++;
                1
        return to_cancel;
};
void* _get_previous_value(int pos){
        void* value;
        int currPos=-2;
        //The logic to find the previous value is to get the previous value from the lesser of the higher writers,
        //if that fails, then it is necessary to find the previous value of the latest of the lower writers.
        for (int i=(static_cast<int>(_writers.size())-1); i>-1; i--){ //this order matters because if a thread has written several
                                                   //it's first should be restored
                if (static_cast<int>(_writers[i])>pos){
                        if (currPos==-2) {
                                 value=_previous_values[i];
                                 currPos=_writers[i];
                        else if (static_cast<int>(_writers[i])<=currPos){</pre>
                                 value= previous values[i];
                                 currPos= writers[i];
                        }
                }
        if (currPos==-2) {
                for (int i=(static cast<int>( writers.size())-1); i>-1; i--){
                        if (static_cast<int>(_writers[i])<=pos){
    if (currPos==-2){</pre>
                                         value=_previous_values[i];
                                         currPos= writers[i];
                                 else if (static cast<int>( writers[i])>currPos){
                                         value=_previous_values[i];
```

```
currPos= writers[i];
                        }
                }
                if (currPos!=-2) {
                        return (void*) value;
                1
                return (void*)-1:
        1:
        vector <int> cancel higher writers(int pos){
        //if used on a write function and the return vector is not empty, an output dependence violation has ocurred (WAW) and the
        //value has to be restored to it's state before pos, so it can be written over.
        //if used on a read function and the return vector is not empty, an anti-dependence violation has ocurred (WAR) and the
        //value has to be restored to it's state before pos, so it can be read.
                set <int> to_cancel_set;
                unsigned int i=0;
                while ((!_writers.empty())&& (i<_writers.size())){</pre>
                        if (_writers[i]>pos){
                                to cancel set.insert( writers[i]);
                                _writers.erase(_writers.begin()+i);
                                 previous values.erase( previous values.begin()+i);
                                i=0;
                        else
                vector <int> to cancel;
                to_cancel.insert(to_cancel.end(), to_cancel_set.begin(), to_cancel_set.end());
                if ((!to_cancel.empty()) &&(!_readers.empty())) {
                        for (unsigned int i=0; i<to cancel.size(); i++){</pre>
                                _readers.erase(to_cancel[i]); //all the readings of the higher writers on this data item have to be unlogg
                }
                return to_cancel;
        1:
//! \class loop_speculator
//! \brief implements a class that takes n ordered interations from a loop and
//!
            executes them speculatively in parallel, while keeping the sequential
//!
//!
//!
//!
//!
//!
            consistency. In order to do this the class relies on write_data and
            read data functions, that check for data dependence violations, as well
            as a commit function that helps to verify that an iteration has ended it's
            execution.
            Additional functionality is provided with functions like append and
            valid_til, that allow the pre-loop section to dynamically add new
            iterations or substract iterations to an on-going speculative execution.
//!
            This allows the model to mimic do-while and do-until behaviour.
//!
//! Limitations:
//! * If a speculative thread is to be canceled, it cannot use functions
//! that involve system mutexes, such as printf, etc. In this case, it
//! is possible that the thread can be canceled while holding such a mutex.
//! and the application con go into deadlock. In order to prevent this the
//! user has to surround this "dangerous" code with:
        "pthread\_setcancelstate(PTHREAD\_CANCEL\ DISABLE,\ NULL);"\ and
//!
        "pthread setcancelstate(PTHREAD CANCEL ENABLE, NULL);".
//!
//! * Sequential consistency is mostly guaranteed, save for exception
//! behaviour.
class loop_speculator {
private:
        //bool values to set if the speculator is active, or has control
        //of the pre-loop section
        bool is active, has pre loop;
        //mutex for synchronized use of the previous group of variables
        pthread_mutex_t _is_active_mutex;
        //thread to manage the pre-loop section, boolean value to check if
        //it has ended it's execution & related mutex
        pthread_t _pl;
```

```
bool pl commit;
pthread_mutex_t _pl_mutex;
//the speculative threads and it's data
vector <_loop_spec_thread> _spec_threads;
//red-black tree, used as a thread index to relate a pthread_id with
//it's model id, i.e. it's position in \_spec\_threads
map <pthread_t, int> _thread_index;
//mutex for a synchronized access to both of the former
pthread_mutex_t _spec_threads_mutex;
//shared_data data between the pre-loop section & the speculative threads
void*& shared data;
//auxiliary data to reset _shared_data.
int _null_data;
//variable indicating until what iteration the pre-loop section has validated
//the execution (-1 if not used) & related mutex
int valid til;
pthread_mutex_t _valid_til_mutex;
//red-black tree that relates the reference of a data element from
//shared_data data with a \_data\_access\_log used to keep track of all the //possible data dependence violations
map <void*, _data_access_log> _access_log;
//mutex related to the _access_log as a whole
pthread_mutex_t _global_access_log_mutex;
//! private method that allows to cancel speculative threads, whose
//! positions are passed as an argument. It is used if there is a
//! error on starting the speculation, and also for resetting the object.
//! TO BE NOTED:
//! * this method takes all class mutexes. On invocation
     _spec_threads_mutex and _global_access_log_mutex should
//!
      be on hold and no other class mutex, save those related
//!
     to a thread that is not about to be canceled or restarted.
//! * if called with an empty vector, it will reset all logs and
//! wait for the threads in deferred cancel to finish.
void _cancel_spec_threads_and_reset_logs (vector <int> threads_to_delete) {
        bool valid_arguments=true;
        if (!threads_to_delete.empty()){
                for (i=0; i<threads to delete.size(); i++){</pre>
                        if (threads to delete[i]<0 || threads to delete[i]> static cast<int>( spec threads.size())){
                                valid arguments=false; //one thread has an invalid id.
                                i=threads_to_delete.size();
                        }
                }
        else{
                valid arguments=false; //no threads to delete
                /*This option will be used to initialize the object, reseting it's inner arrays*/
                pthread mutex lock(& is active mutex);
                pthread mutex lock(& valid til mutex);
                if ( has pre loop)
                        pthread mutex lock(& pl mutex);
                if (!_spec_threads.empty()){
                        _spec_threads.clear();
                if (!_thread_index.empty()){
                        /*in order to clear the thread index, it will be necessary to guarantee that all the threads in deferred
                        cancel have finished, and thus will not use the object*/
```

```
if ( has pre loop)
                         pthread mutex unlock(& pl mutex);
                 pthread mutex unlock(& valid til mutex);
                 pthread_mutex_unlock(&_is_active_mutex);
                 pthread_mutex_unlock(&_spec_threads_mutex);
                 pthread_mutex_unlock(&_global_access_log_mutex);
                 map <pthread_t, int>::iterator it;
for (it=_thread_index.begin(); it!=_thread_index.end(); it++){
                         if (it->second==-1) {
                                  int a;
                                  do {
                                           a= pthread_kill(it->first, 0);
                                  } while (a==0);
                 pthread_mutex_lock(&_global_access_log_mutex);
                 pthread_mutex_lock(&_spec_threads_mutex);
                 _thread_index.clear();
                 if (! access log.empty()){
                         _access_log.clear();
        else {
                 if (! access log.empty()){
                         _access_log.clear();
                 if (_has_pre_loop)
                         pthread mutex unlock(& pl mutex);
                 pthread_mutex_unlock(&_valid_til_mutex);
                 pthread_mutex_unlock(&_is_active_mutex);
}
if (valid_arguments) {
        for (unsigned int i=threads_to_delete.size()-1; i>=0; i--){
                 pthread_mutex_lock(&_spec_threads[threads_to_delete[i]]._thread_mutex);
        pthread_t thrd_id=pthread_self();
        pthread_mutex_lock(&_is_active_mutex);
        pthread_mutex_lock(&_valid_til_mutex);
        if (_has_pre_loop)
                 pthread mutex lock(& pl mutex);
        for (unsigned int i=0; i<threads_to_delete.size(); i++){
     if (pthread_equal(_spec_threads[threads_to_delete[i]]._thread, thrd_id)==0){</pre>
                         if (pthread_kill(_spec_threads[threads_to_delete[i]]._thread, 0)==0){
                                  pthread_cancel(_spec_threads[threads_to_delete[i]]._thread);
                         }
                          _thread_index[_spec_threads[threads_to_delete[i]]._thread]= -1;
                         pthread_mutex_destroy(&_spec_threads[threads_to_delete[i]]._thread_mutex);
                 }
        if (_has_pre_loop)
                 pthread_mutex_unlock(&_pl_mutex);
        pthread_mutex_unlock(&_valid_til_mutex);
        pthread_mutex_unlock(&_is_active_mutex);
```

```
//! private method that allows to cancel speculative threads higher than the
//! one passed as an argument, it is used in the valid_til function.
//!
//! TO BE NOTED:
//! * this method takes all class mutexes. On invocation
//! _spec_threads_mutex and _global_access_log_mutex should
//!
      be on hold and no other class mutex, save those related
//! to a thread that is not about to be canceled or restarted.
//!
//!
void _cancel_higher_spec_threads (int top_thread_pos) {
        bool valid arguments=false;
        if (top_thread_pos>=-1 && top_thread_pos< static_cast<int>(_spec_threads.size())){
                set <int> threads to cancel;
                set <void*> data_to_check;
                int aux_value=top_thread_pos;
                if (aux_value==-1) {
                        aux value=0;
                for (unsigned int i=aux_value; i<_spec_threads.size(); i++){</pre>
                        threads_to_cancel.insert(i);
                }
                set <int>::iterator it;
                for (it=threads_to_cancel.begin(); it!=threads_to_cancel.end(); it++){
                        if (*it!=top_thread_pos){
                                pthread mutex lock(& spec threads[*it]. thread mutex);
                }
                set<void*>::iterator it2;
                for (it=threads to cancel.begin(); it!=threads to cancel.end(); it++){
                        for (it2= spec threads[*it]. read data.begin(); it2!= spec threads[*it]. read data.end(); it2++){
                                 if (data to check.find(*it2)==data_to_check.end()){
                                         data to check.insert(*it2);
                        if (*it!=top_thread_pos){
                                 _spec_threads[*it]._read_data.clear();
                        1
                        for (it2=_spec_threads[*it]._written_data.begin(); it2!=_spec_threads[*it]._written_data.end(); it2++){
                                 if (data_to_check.find(*it2)==data_to_check.end()){
                                        data_to_check.insert(*it2);
                        }
                        if (*it!=top_thread_pos){
                                 _spec_threads[*it]._written_data.clear();
                        }
                pthread_mutex_lock(&_is_active_mutex);
                pthread_mutex_lock(&_valid_til_mutex);
                if (_has_pre_loop)
                        pthread mutex lock(& pl mutex);
                for (it=threads to cancel.begin(); it!=threads to cancel.end(); it++){
                        if (*it!=top_thread_pos){
                                 _spec_threads[*it]._commit=false;
if (pthread_kill(_spec_threads[*it]._thread, 0)==0){
                                        pthread_cancel(_spec_threads[*it]._thread);
                                 _thread_index[_spec_threads[*it]._thread]= -1;
                                 \verb|pthread_mutex_destroy(&\_spec_threads[*it]._thread_mutex)|;\\
                        }
                for (it2=data_to_check.begin(); it2!=data_to_check.end(); it2++){
```

```
void* new_val= _access_log.find(*it2)->second. get_previous_value(top_thread_pos);
vector <int> writers_to_cancel=_access_log.find(*it2)->second._cancel_higher_writers(top_thread_pos);
                         vector <int> readers_to_cancel=_access_log.find(*it2)->second._cancel_higher_readers(top_thread_pos);
                         if (!writers to cancel.emptv()){
                                 memcpy((void*&)const_cast<void*&>(*it2), (void*)&new_val, _access_log.find(*it2)->second._size);
                if (top_thread_pos==-1){
                         _spec_threads.clear();
                1
                else {
                         _spec_threads.resize(top_thread_pos+1);
                1
                if (_has_pre_loop)
                         pthread_mutex_unlock(&_pl_mutex);
                pthread_mutex_unlock(&_valid_til_mutex);
                pthread mutex unlock (& is active mutex);
/\!/! private method that allows to reset speculative threads, which has to be called
//! from a thread whose position is passed as argument. This method is used for reseting
//! threads that committed some form of data dependence violation.
//! TO BE NOTED:
//! * this method takes all class mutexes. On invocation
//! _spec_threads_mutex and _global_access_log_mutex should
      be on hold and no other class mutex, save those related
//!
    to a thread that is not about to be canceled or restarted.
//!
void _reset_spec_threads (int current_thread, void* data_currently_held, vector <int> threads_to_reset){
        bool valid arguments=true;
        if (!threads_to_reset.empty()){
                 for (i=0; i<threads_to_reset.size(); i++){</pre>
                         if (threads_to_reset[i]<0 || threads_to_reset[i]> static_cast<int>(_spec_threads.size())){
                                 valid arguments=false; //one thread has an invalid id.
                                 i=threads_to_reset.size();
                         }
                }
        else{
                valid arguments=false; //no threads to delete
        if (current_thread<-1) {</pre>
                valid_arguments=false;
        1
        if (valid arguments) {
                 set <int> threads_currently_held;
                threads_currently_held.insert(current_thread);
                set <int> set_of_threads_to_reset;
                set of threads to reset.insert(threads to reset.begin(), threads to reset.end());
                set<int>:: iterator it;
                it=set of threads to reset.begin();
                set <void*> data already checked;
                data already checked.insert(data currently held);
                while (it!=set_of_threads_to_reset.end()){
                         unsigned int previous size=set of threads to reset.size();
                         if (threads currently held.find(*it)==threads currently held.end()){
                                 pthread_mutex_lock(&_spec_threads[*it]._thread_mutex);
                                 threads_currently_held.insert(*it);
                                 set<void*>::iterator it2;
                                 for (it2=_spec_threads[*it]._read_data.begin(); it2!=_spec_threads[*it]._read_data.end(); it2++){
                                          _access_log.find(*it2)->second._readers.erase(*it);
```

```
for (it2= spec threads[*it]. written data.beqin(); it2!= spec threads[*it]. written data.end(); it
                                             if (data already checked.find(const cast<void*&>(*it2)) == data already checked.end()) {
                                                            data already checked.insert(const cast<void*&>(*it2));
                                                            void* value to restore= access log.find(*it2)->second. get previous value(current
                                                            \verb|vector < int>| writers_to_cancel=| access_log.find(*it2)-> second._cancel_higher| write| | to cancel=| access_log.find(*it2)-> second._cancel_higher| write| | to cancel=| access_log.find(*it2)-> second._cancel=| access_log.find(*it2)-> second._cancel_higher| write| | to cancel=| access_log.find(*it2)-> second._cancel_higher| | to cancel=| access_log.find(*it2)-> second._cancel=| access_log.find(*it2)-> 
                                                            vector <int> readers_to_cancel=_access_log.find(*it2)->second._cancel_higher_reade
                                                            if (!writers_to_cancel.empty()){//implied or secondary WAW and WAR violations have
                                                                           memcpy((void*&)const_cast<void*&>(*it2), (void*)&value_to_restore, _access
                                                            if (!readers to cancel.empty()) {//implied or secondary RAW violations have ocurred
                                                                           writers to cancel.insert(writers to cancel.end(), readers to cancel.begin()
                                                            if (!writers_to_cancel.empty()){
                                                                           for (unsigned int i=0; i<writers to cancel.size(); i++){
                                                                                          set_of_threads_to_reset.insert(writers_to_cancel[i]);
                              _spec_threads[*it]._written_data.clear();
                              if ((*it!=current_thread)&&(pthread_equal(_spec_threads[*it]._thread, pthread_self())==0)){
                                             pthread_mutex_lock(&_is_active_mutex);
                                             pthread_mutex_lock(&_valid_til_mutex);
                                             if (_has_pre_loop)
                                                            pthread mutex lock(& pl mutex);
                                             _spec_threads[*it]._commit=false;
                                             if (pthread_kill(_spec_threads[*it]._thread, 0)==0){
                                                            pthread_cancel(_spec_threads[*it]._thread);
                                             _thread_index[_spec_threads[*it]._thread]= -1;
                                             if (_has_pre_loop)
                                                           pthread mutex unlock(& pl mutex);
                                             pthread_mutex_unlock(&_valid_til_mutex);
                                             pthread_mutex_unlock(&_is_active_mutex);
                              }
               if (set of threads to reset.size()>previous size) {
                              it=set of threads to reset.begin();
               else {
                              it++;
pthread_mutex_lock(&_is_active_mutex);
pthread mutex lock(&_valid_til_mutex);
if (_has_pre_loop)
              pthread_mutex_lock(&_pl_mutex);
pthread_attr_t attr;
pthread_attr_init (&attr);
pthread attr setschedpolicy(&attr, SCHED RR);
for (it=set of threads to reset.begin(); it!=set of threads to reset.end(); it++){
               if (*it!=current thread) {
                              if (_valid_til!=-1){
                                             if (*it<=_valid_til) {</pre>
                                                            pthread create (& spec threads[*it]. thread, &attr, spec threads[*it]. thread inst
                                                            _thread_index[_spec_threads[*it]._thread]= *it;
                                             pthread create(& spec threads[*it]. thread, &attr, spec threads[*it]. thread instructions
                                             thread index[ spec threads[*it]. thread] = *it;
```

spec threads[*it]. read data.clear();

1

```
if (_has_pre loop)
                                  pthread mutex unlock(& pl mutex);
                         pthread mutex unlock(& valid til mutex);
                         pthread_mutex_unlock(&_is_active_mutex);
                         for (it=threads_currently_held.begin(); it!=threads_currently_held.end(); it++){
                                  if (*it!=current_thread)
                                           pthread_mutex_unlock(&_spec_threads[*it]._thread_mutex);
                         1
                }
        };
public:
        //! default constructor
        loop_speculator():_shared_data((void*&)_null_data){
                 _is_active=false;
                 _pl_commit=false;
                 _has_pre_loop=false;
                 _valid_til=-1;//the valid til funcionality is not used
                 null data=-1;
                pthread_mutex_init (&_is_active_mutex, NULL);
pthread_mutex_init (&_pl_mutex, NULL);
                pthread_mutex_init (& spec_threads_mutex, NULL);
pthread_mutex_init (& valid_til_mutex, NULL);
                 pthread_mutex_init (&_global_access_log_mutex, NULL);
        };
        //! function providing access to the shared_data as a whole
        void*& get_shared_data (){
                 pthread_mutex_lock(&_is_active_mutex);
                 bool manages_pre_loop=_has_pre_loop;
                 if (!_is_active) {
                         pthread_mutex_unlock(&_is_active_mutex);
                         return (void*&) null data; //the object is inactive
                 pthread_mutex_unlock(&_is_active_mutex);
                 pthread_mutex_lock(&_spec_threads_mutex);
                 if ( thread index.empty() || thread index.find(pthread self())!= thread index.end()){
                         pthread_mutex_unlock(&_spec_threads_mutex);
                         return _shared_data; //valid thread or thread in deferred cancelation
                 pthread_mutex_unlock(&_spec_threads_mutex);
                 if (manages_pre_loop) {
                         pthread mutex lock(& pl mutex);
                         if (pthread_self() == _pl) {
                                  pthread_mutex_unlock(&_pl_mutex);
                                  return shared data;
                         }
                         else {
                                  pthread_mutex_unlock(&_pl_mutex);
                                  return (void*&)_null_data; //invalid caller
                         }
                 }
```

```
return shared data; //unmanaged pre-loop or thread in deferred cancelation
};
//! required function that signals that a given thread has ended
//! it's execution
//!
void commit (){
        pthread_mutex_lock(&_is_active_mutex);
        if (_is_active){
                bool manages pre loop= has pre loop;
                pthread mutex unlock(& is active mutex);
                pthread mutex lock(& spec threads mutex);
                if (!_thread_index.empty()){ //thread in deferred cancel.
                        pthread_t thread_id=pthread_self();
                        if (_thread_index.find(thread_id)!=_thread_index.end()){
                                int pos=_thread_index.find(thread_id)->second;
                                if (pos>=0){ //valid thread
                                        pthread_mutex_lock(&_spec_threads[pos]._thread_mutex);
                                         _spec_threads[pos]._commit=true;
                                        pthread_mutex_unlock(&_spec_threads[pos]._thread_mutex);
                        else if (manages_pre_loop){
                                if (thread id== pl) {
                                        pthread_mutex_lock(&_pl_mutex);
                                        _pl_commit=true;
pthread_mutex_unlock(&_pl_mutex);
                        else {
                                //unmanaged pre-loop
                                pthread_mutex_lock(&_pl_mutex);
                                _pl_commit=true;
                                pthread_mutex_unlock(&_pl_mutex);
                pthread_mutex_unlock(&_spec_threads_mutex);
        else{
                pthread_mutex_unlock(&_is_active_mutex);
        }
};
//!
//! function to be called from the speculative threads
//! and the concurrent pre-loop section, in order to read the shared_data
//! while keeping the expected sequential data consistency
//!
template <class T>
T read_data(T*& data_to_be_read){
        pthread_mutex_lock(&_is_active_mutex);
        if (_is_active){
                bool manages_pre_loop=_has_pre_loop;
                pthread_mutex_unlock(&_is_active_mutex);
                int current_thread=-2;
                pthread_mutex_lock(&_spec_threads_mutex);
                pthread_t thread_id=pthread_self();
                if ( thread index.empty()){//thread in deferred cancel.
                        current_thread=-3;
                else if ( thread index.find(thread id)!= thread index.end()){
                        current thread= thread index.find(thread id)->second;
                        if (current thread==-1) //thread in deferred cancel.
                                current thread=-3;
                else {
                        if (manages_pre_loop) {
                                pthread_mutex_lock(&_pl_mutex);
```

```
if (thread id== pl){
                                                                   current thread=-1;
                                             }
                                             else {
                                                                   current thread=-3; //invalid caller
                                             pthread mutex unlock(& pl mutex);
                      else {
                                             current_thread=-1;
                      1
if (current thread>=-1) { //is a valid read.
                      T retval;
                      pthread_mutex_lock(&_valid_til_mutex);
                      if (_valid_til!=-1 && _valid_til<current_thread) {</pre>
                                             retval=*(T*)data to be read;
                                             pthread mutex unlock (& valid til mutex);
                                             pthread mutex unlock(& spec threads mutex);
return retval; //thread that is no longer valid.
                      \verb|pthread_mutex_unlock(&_valid_til_mutex)|;|\\
                      pthread_mutex_lock(&_global_access_log_mutex);
                      if (_access_log.find((void*)data_to_be_read)==_access_log.end()){
                      //if the data has never been accesed, it's log should be created
                                             pthread_mutex_t temp;
                                             pthread mutex init (&temp, NULL);
                                             _data_access_log temp2;
                                             //since it's the data's first use, it's assumed to be the value in the pre-loop
                                              //so the data is initialized on that premise
                                             temp2._size=sizeof(T);
                                             temp2. writers.push back((int)-1);
                                             temp2._previous_values.push_back((void*)*(T*) data_to_be_read);
                                             if (current thread!=-1) {
                                                                   pthread_mutex_lock(&_spec_threads[current_thread]._thread_mutex);
                                                                     spec threads[current thread]. read data.insert((void*)data to be read);
                                                                   temp2. readers.insert(current_thread);//the current read is logged
                                             ì
                                             pthread_mutex_unlock(&_spec_threads_mutex);
                                             _access_log.insert (pair<void*, _data_access_log>((void*)data_to_be_read, temp2));
                                             retval=*(T*)data_to_be_read;
                                             if (current_thread!=-1) {
                                                                    \verb|pthread_mutex_unlock(&\_spec_threads[current_thread]._thread_mutex|);|
                                             pthread mutex unlock(& global access log mutex);
                                             return retval;
                       _data_access_log* ptr=&_access_log.find((void*)data_to_be_read)->second;
                      pthread_mutex_lock(&ptr->_data_mutex);
                      pthread_mutex_unlock(&_global_access_log_mutex);
                      retval= (T) ptr->_get_previous_value(current_thread);
                      vector <int> threads to cancel= ptr-> cancel higher writers (current thread);
                      // {\it if} the log exists, then anti-dependence violation is checked (WAR)
                      if (current thread!=-1) {
                                             pthread_mutex_lock(&_spec_threads[current_thread]._thread_mutex);
                      if (!threads to cancel.empty()){//an anti-dependence violation has occurred
                                             \verb|memcpy((void*)data_to_be_read, (void*)&retval, sizeof(T)); // the previous value is restored.
                                             _reset_spec_threads(current_thread, (void*)data_to_be_read, threads_to_cancel);
                      else{
                                             retval=*(T*) data_to_be_read;//no anti-dependence violation has ocurred, the data can be read on i
                      pthread_mutex_unlock(&_spec_threads_mutex);
                      //now that the reading is done, it has to be logged % \left( 1\right) =\left( 1\right) \left( 1\right) \left
                      if (current thread!=-1) {
                                             ptr->_readers.insert(current thread);
                                             _spec_threads[current_thread]._read_data.insert((void*)data_to_be_read);
```

```
pthread mutex unlock (& spec threads [current thread]. thread mutex);
                        pthread mutex unlock(&ptr-> data mutex);
                        return retval;
                pthread mutex unlock (& spec threads mutex);
        else (
                pthread mutex unlock (& is active mutex);
        1
        return *(T*) data_to_be_read; //invalid caller, inactive object or thread in deferred cancel.
                                     //it's calling data is returned instead of NULL, to prevent segmentation faults.
};
//!
//! function to be called from the speculative threads
//! and concurrent pre-loop section, in order to write the shared_data
//! while keeping the expected sequential data consistency
//!
template <class T>
int write_data(T*& data_to_be_written_upon, T* data_to_write){
        pthread_mutex_lock(&_is_active_mutex);
        if (_is_active){
                bool manages_pre_loop=_has_pre_loop;
                {\tt pthread\_mutex\_unlock(\&\_is\_active\_mutex);}
                int current_thread=-2;
                pthread_mutex_lock(&_spec_threads_mutex);
                pthread t thread id=pthread self();
                if ( thread index.empty()){//thread in deferred cancel.
                        current_thread=-3;
                else if ( thread index.find(thread id)!= thread index.end()){
                        current thread= thread index.find(thread id)->second;
                        if (current thread==-1) //thread in deferred cancel.
                                current thread=-3;
                else {
                        if (manages_pre_loop) {
                                {\tt pthread\_mutex\_lock(\&\_pl\_mutex);}
                                if (thread_id==_pl){
                                        current_thread=-1;
                                }
                                else {
                                        current_thread=-3; //invalid caller
                                pthread_mutex_unlock(&_pl_mutex);
                        else {
                                current_thread=-1;
                if (current thread>=-1){
                        pthread mutex lock(& valid til mutex);
                        pthread_mutex_unlock(&_spec_threads_mutex);
                                return -1;
                        pthread_mutex_unlock(&_valid_til_mutex);
                        {\tt pthread\_mutex\_lock} \, ( \&\_{\tt global\_access\_log\_mutex} ) \, ; \\
                        if (_access_log.find((void*)data_to_be_written_upon) == _access_log.end()) {
                        //if data has no log entry, then it's entry should be created.
                                pthread mutex t temp;
                                pthread_mutex_init (&temp, NULL);
                                data access log temp2;
                                //since it's the data's first use, it's assumed to be the value in the pre-loop
                                //so the data is initialized on that premise
                                temp2._size=sizeof(T);
                                temp2._writers.push_back((int)-1);
                                temp2. previous values.push back((void*)*(T*)data to be written upon);
                                if (current_thread!=-1) {
                                        pthread mutex lock(& spec threads[current thread]. thread mutex);
                                        _spec_threads[current_thread]._written_data.insert((void*)data_to_be_written_upon);
                                        temp2. writers.push back(current thread);
                                        temp2. previous values.push back((void*)*(T*)data to be written upon);
```

```
else{
                        temp2. writers.push back(-1);
                        temp2._previous_values.push_back((void*)data to write);
                pthread mutex unlock(& spec threads mutex);
                access log.insert (pair<void*, data access log>((void*)data to be written upon, temp2));
                //and finally the writting is done
                memcpy ((void*)data_to_be_written_upon, (void*)&data_to_write, sizeof(T));
                if (current_thread!=-1) {
                       pthread mutex unlock(& spec threads[current thread]. thread mutex);
                pthread mutex unlock(& global access log mutex);
                return 0;
//if the log exists, then possible true dependence and output dependence violations are checked (RAW and WAW)
        _data_access_log* ptr=&_access_log.find((void*)data_to_be_written_upon)->second;
        pthread_mutex_lock(&ptr->_data_mutex);
        pthread mutex unlock(& global access log mutex);
        if (current thread!=-1) {
                pthread mutex lock(& spec threads[current thread]. thread mutex);
        T value:
        vector <int> writers to cancel;
        bool no output violation=false;
        if ((void*)*(T*)data_to_be_written_upon==(void*)data_to_write){
        /\!/\!Since\ restarting\ threads\ has\ such\ a\ significant\ cost,\ it\ is\ better\ to\ check\ if\ it\ is\ really\ necessary\ tc
        //So, we will evaluate false positives in output dependence violations.
                if (!ptr->_output_violations_are_false(data_to_write, current_thread)) {
                value= (T) ptr->_get_previous_value(current_thread);
                writers to cancel=ptr-> cancel higher writers(current thread);
                }
                else {
                        no output violation=true;
        elsef
                value= (T) ptr-> get previous value(current thread);
                writers to cancel=ptr-> cancel higher writers (current thread);
        vector <int> readers to cancel=ptr-> cancel higher readers(current thread);
        if (no_output_violation && readers_to_cancel.empty()){
                if (current thread!=-1) {
                        _spec_threads[current_thread]._written_data.insert((void*)data_to_be_written_upon);
                        pthread_mutex_unlock(&_spec_threads[current_thread]._thread_mutex);
                pthread_mutex_unlock(&ptr->_data_mutex);
                pthread_mutex_unlock(&_spec_threads_mutex);
                return 0;
        }
        if (!writers_to_cancel.empty()){
        //an output dependence violation has ocurred, it's previous value has to be restored before writting over,
        //those writers have to be restarted
                ptr->_writers.push_back(current_thread);
                ptr->_previous_values.push_back((void*)value);
        /*This branch has been turned into a comment, since the logging for all cases where there are false output
        is done on output violations are false(...). However, should a unforseeable case arise where this does not
        (for instance, a particular thread reset pattern), then it should be un-commented.*/
        else if (!no_output_violation){
                //no output dependence violation, it's current value is valid and logged as the previous value of
                //thread write..
                ptr-> writers.push back(current thread);
                ptr->_previous_values.push_back((void*)*(T*)data_to_be_written upon);
        if (!readers_to_cancel.empty()){//a true dependence violation has ocurred, those readers should be restart
                writers_to_cancel.insert(writers_to_cancel.end(), readers_to_cancel.begin(), readers_to_cancel.end
        if (!writers to cancel.empty()){
                //the violating threads are restarted
```

```
reset spec threads (current thread, (void*) data to be written upon, writers to cancel);
                        pthread_mutex_unlock(&_spec_threads_mutex);
                         //now some more logging of the write
                         if (current thread==-1) {
                                 ptr->_writers.push_back(-1);
                                 ptr-> previous values.push back((void*)data to write);
                        else {
                                 spec threads[current thread]. written data.insert((void*)data to be written upon);
                        }
                         //and finally the write is made
                        memcpy ((void*)data_to_be_written_upon, (void*)&data_to_write, sizeof(T));
                        if (current thread!=-1){
                                 pthread_mutex_unlock(&_spec_threads[current_thread]._thread_mutex);
                        pthread_mutex_unlock(&ptr->_data_mutex);
                         return 0;
                pthread mutex unlock(& spec threads mutex);
        else {
                pthread mutex unlock(& is active mutex);
        return -1; //invalid caller, thread in deferred cancel or inactive object.
};
//! function that permits the pre_loop section to dynamically append
//! a new speculative thread at the end of the array in an on-going
//! speculative execution.
//!
int append(void* (f)(void*), void* consts){
        pthread_mutex_lock(&_is_active_mutex);
        if (!_is_active){
                pthread mutex unlock(& is active mutex);
                return -1; //the object is not active
        bool manages_pre_loop=_has_pre_loop;
        pthread_mutex_unlock(&_is_active_mutex);
        bool is_pre_loop=false;
        pthread mutex lock(& spec threads mutex);
        if (manages_pre_loop) {
                pthread_mutex_lock(&_pl_mutex);
                if (pthread_self() == _pl) {
                        is_pre_loop=true;
                pthread_mutex_unlock(&_pl_mutex);
        if (!is_pre_loop){
                pthread_t thrd_id=pthread_self();
                if ((!_thread_index.empty()) && (_thread_index.find(thrd_id)==_thread_index.end())){
                        if (!manages pre loop){
                                 is_pre_loop=true;
        if (is pre loop) {
                pthread_attr_t attr;
pthread_attr_init (&attr);
                pthread_attr_setschedpolicy(&attr, SCHED_RR);
                int success;
                int current size= spec threads.size();
                pthread mutex lock(& valid til mutex);
                if ( valid_til!=-1 && (current_size-1)>=_valid_til){
                        pthread_mutex_unlock(& valid_til_mutex);
pthread_mutex_unlock(& spec_threads_mutex);
                        return -1;//the thread to append does not fit with the valid_til condition
                pthread_mutex_unlock(&_valid_til_mutex);
                 _spec_threads.resize(_spec_threads.size()+1);
                pthread_mutex_lock(&_spec_threads[_spec_threads.size()-1]._thread_mutex);
                int i=_spec_threads.size()-1;
                _spec_threads[i]._commit=false;
                _spec_threads[i]._thread_instructions=f;
                _spec_threads[i]._const_args=consts;
```

```
success=pthread create(& spec threads[i]. thread, &attr, f, consts);
                if (success!=0){
                        pthread mutex unlock(& spec threads[i]. thread mutex);
                        _ spec_threads.resize(_spec_threads.size()-1);
pthread_mutex_unlock(& spec_threads_mutex);
                         return -1;//the thread could not be created, thus is not appended
                1
                _thread_index.insert(pair<pthread_t, int>(_spec_threads[_spec_threads.size()-1]._thread, _spec_threads.size()-1));
                pthread_mutex_unlock(&_spec_threads[i]._thread_mutex);
                pthread_mutex_unlock(&_spec_threads_mutex);
                return 0;
        pthread mutex unlock(& spec threads mutex);
        return -1; //invalid caller
};
//! function that allows the pre loop section to dynamically cancel
//! a number of speculative threads at the end of the array in an on-going
//! speculative execution, this is done by indicating until what position
//! the execution remains valid.
//! TO BE NOTED: It can only be called once during all the execution.
//!
int valid til (int pos) {
        if (pos<-1)
                return -1;//invalid argument
        pthread mutex lock(& is active mutex);
        if (!_is_active) {
                pthread_mutex_unlock(&_is_active_mutex);
                return -1; //the object is not active
        bool manages pre loop= has pre loop;
        pthread_mutex_unlock(&_is_active_mutex);
        pthread mutex lock(& valid til mutex);
        if ( valid til!=-1) {
                pthread mutex unlock(& valid til mutex);
                return -1; //the valid til function has already been called
        pthread mutex unlock (& valid til mutex);
        bool is_pre_loop=false;
        {\tt pthread\_mutex\_lock\,(\&\_spec\_threads\_mutex);}
        if (manages_pre_loop){
                pthread_mutex_lock(&_pl_mutex);
                if (pthread_self() == _pl) {
                        is pre loop=true;
                pthread_mutex_unlock(&_pl_mutex);
        if (!is_pre_loop){
                pthread_t thrd_id=pthread_self();
                 if ((!_thread_index.empty()) & (_thread_index.find(thrd_id) == _thread_index.end())){
                        if (!manages_pre_loop){
                                 is_pre_loop=true;
        if (is_pre_loop){
                pthread mutex lock(& valid til mutex);
                 valid til=pos;
                pthread_mutex_unlock(&_valid_til_mutex);
                if ( (static cast<int>( spec threads.size())-1)>pos){
                        _cancel_higher_spec_threads (pos);
                {\tt pthread\_mutex\_unlock\,(\&\_spec\_threads\_mutex);}
                {\tt return} 0; //no thread was canceled, the restriction will be
                          //imposed through the append function
        pthread_mutex_unlock(&_spec_threads_mutex);
        return -1;//invalid caller
1:
//!
//! speculate: a complete function that takes orderly the instructions and arguments
//! of the pre-loop section, as well as those of the loop iterations; and starts their
//! speculatively parallel execution, while maintaining the sequential consistency.
```

```
//! This function returns in the &shared data, the results of the computation.
//! TO BE NOTED:
       * In this version of the function, the object is in control of the pre-
//!
//!
       loop section. This means that on invocation, the caller blocks until the pre-
//!
      loop and all the iterations up to the end or valid til have committed.
//!
       * If the pre-loop section does not explicitly commit through the commit function,
//!
      then the whole results of the speculation will not be communicated to the
//!
       shared data and the speculation will be canceled.
//!
       * If any iteration does not commit, then only the results up to it will be taken
//!
       into account, the remaining iterations will be canceled.
//!
int speculate (void*& shared_data, void* (fpl)(void*), void* const_args_pl, script_vector thread_instructions, vector <void*> cons
        if (thread instructions.size()!=const args.size()){
                return -1; //invalid parameters
        pthread_mutex_lock(&_is_active_mutex);
        if (_is_active){
                pthread mutex unlock(& is active mutex);
                return -1; //the object is active
        }
         is active=true:
        {\tt pthread\_mutex\_unlock(\&\_is\_active\_mutex);}
        pthread_mutex_lock(&_spec_threads_mutex);
        pthread_mutex_lock(&_global_access_log_mutex);
        vector <int> threads to cancel;//a selective reset is performed, waiting for the threads in
        _cancel_spec_threads_and_reset_logs(threads_to_cancel);//deferred cancel to finish, and reseting the
        //arrays that the speculation uses.
        pthread mutex_lock(&_is_active_mutex);
        has pre loop=true;
        pthread_mutex_unlock(&_is_active_mutex);
        {\tt pthread\_mutex\_lock\,(\&\_valid\_til\_mutex);}
         valid til=-1;
        \verb|pthread_mutex_unlock(&_valid_til mutex)|;|\\
        pthread_mutex_lock(&_pl_mutex);
        _pl_commit=false;
        pthread_mutex_unlock(&_pl_mutex);
        shared data=shared data;
        pthread_attr_t attr;
        pthread attr init (&attr);
        pthread_attr_setschedpolicy(&attr, SCHED_RR);
        //the threads are created
        int success;
        int valid til thread=-1;
         spec threads.resize(thread instructions.size());
        pthread_mutex_lock(&_pl_mutex);
        success=pthread_create (&_pl, &attr, fpl, const_args_pl);
        pthread_mutex_unlock(&_pl_mutex);
        if (success!=0) {
                _spec_threads.clear();
                pthread_mutex_unlock(&_global_access_log_mutex);
                pthread_mutex_unlock(&_spec_threads_mutex);
                pthread_mutex_lock(&_is_active_mutex);
                 _is_active=false;
                pthread_mutex_unlock(&_is_active_mutex);
                return -1; //the pre-loop section could not be created
        1
        for (unsigned int i=0; i<thread instructions.size(); i++){</pre>
                pthread mutex lock(& spec threads[i]. thread mutex);
        for (unsigned int i=0; i<thread instructions.size(); i++){</pre>
                _spec_threads[i]._thread_instructions=thread_instructions[i];
                 _spec_threads[i]._const_args=const_args[i];
                success = pthread\_create \ (\textbf{\&\_spec\_threads[i].\_thread}, \ \textbf{\&} attr, \ \_spec\_threads[i].\_thread\_instructions, \ \_spec\_threads[i].
                if (success==0) {
                         _thread_index.insert(pair<pthread_t, int>(_spec_threads[i]._thread, i));
                1
                pthread_mutex_unlock(&_spec_threads[i]._thread_mutex);
                         for (unsigned int j=i+1; j<thread instructions.size(); j++){</pre>
```

```
pthread_mutex_unlock(&_spec_threads[j]._thread_mutex);
                 for (unsigned int j=0; j<i; j++){
                         threads_to_cancel.push_back(j);
                 _cancel_spec_threads_and_reset_logs(threads_to_cancel);
                 pthread_mutex_unlock(&_global_access_log_mutex);
                 {\tt pthread\_mutex\_unlock\,(\&\_spec\_threads\_mutex);}
                 {\tt pthread\_mutex\_lock(\&\_is\_active\_mutex);}
                 _is_active=false;
                 pthread_mutex_unlock(&_is_active_mutex);
                 return -1;// a thread could not be created.
pthread_mutex_unlock(&_global_access_log_mutex);
pthread_mutex_unlock(&_spec_threads_mutex);
bool commit made=false;
int joined_but_didnt_commit=0;
do{
        success=pthread_join (_pl, NULL);
        if (success!=0){
                pthread mutex lock(& pl mutex);
                 if (_pl_commit){
                         commit_made=true;
                 }
                 else{
                         if (joined_but_didnt_commit<10){</pre>
                                  success=0;
                                  joined_but_didnt_commit++;
                         1
                 }
                 /*just in case the commit is not logged in a timely fasion, it is tried
                 10 times to see if it get's made, if not then the object asumes there was no commitment in the pre-loop.
                 Although slightly un-elegant, the validity of this solution has not been disproved by empirical
                 pthread mutex unlock(& pl mutex);
} while (success==0);
pthread mutex lock(& spec threads mutex);
if (!commit made) {
                 _cancel_higher_spec_threads(-1);
                 pthread_mutex_unlock(&_spec_threads_mutex);
                 pthread_mutex_lock(&_is_active_mutex);
                 _is_active=false;
                 {\tt pthread\_mutex\_unlock\,(\&\_is\_active\_mutex);}
                return -1; //the pre-loop did not commit.
}
int i=0;
while (_spec_threads.size()>i){
        pthread_mutex_unlock(&_spec_threads_mutex);
        commit made=false;
        joined_but_didnt_commit=0;
                 success=pthread join ( spec threads[i]. thread, NULL);
                 if (success!=0){
                         pthread_mutex_lock(&_spec_threads[i]._thread_mutex);
if (_spec_threads[i]._commit)(
                                 commit_made=true;
                         }
                         else(
                                  if (joined_but_didnt_commit<100){</pre>
                                          success=0;
                                          joined but didnt commit++;
```

```
pthread mutex unlock(& spec threads[i]. thread mutex);
                                          /*just in case the commit is not logged in a timely fasion, it is tried
                                         10 times to see if it get's made, if not then the object asumes there was no commitment in the iteration.
                                         Although slightly un-elegant, the validity of this solution has not been disproved by empirical
                                         tests. */
                            } while (success==0);
                            if (!commit made){
                                         pthread mutex lock(& spec threads mutex);
                                          cancel higher spec threads(i); //This is perhaps a bit excesive, but for consistency purposes it is manda
                                         pthread mutex unlock(& spec threads mutex);
                                         pthread_mutex_lock(&_is_active_mutex);
                                           is active=false;
                                         pthread mutex unlock(& is active mutex);
                                         return -1; //the pre-loop did not commit.
                            }
                           i++:
                           pthread mutex lock(& spec threads mutex);
                            pthread_mutex_lock(&_valid_til_mutex);
                            valid til thread= valid til;
                            pthread_mutex_unlock(&_valid_til_mutex);
                            if (valid til thread!=-1 && i>valid til thread) {
                                         i=_spec_threads.size();
              pthread mutex unlock(& spec threads mutex);
              pthread_mutex_lock(&_is_active_mutex);
              is active=false;
              pthread mutex unlock(& is active mutex);
              return 0:
1:
//! speculate: a complete function that takes orderly the instructions and arguments
//! of the loop iterations; and starts their speculatively parallel execution, while
//! maintaining the sequential consistency with an un-managed pre-loop section.
//! This function will return in the &shared data, the results of the computation, when
//! get_results is called.
//!
//! TO BE NOTED:
//!
           * In this version of the function, the object is not in control of the pre-
            loop section. This means that on invocation, the caller only blocks for the creation
          of the threads, and can resume it's execution as a possible pre-loop, until
//!
           calling get results(), when the caller blocks until all valid loop iterations/threads return.
//!
int speculate (void*& shared_data, script_vector thread_instructions, vector <void*> const_args){
             if (thread instructions.size()!=const_args.size()){
                           return -1; //invalid parameters
              pthread_mutex_lock(&_is_active_mutex);
              if (_is_active){
                           pthread mutex unlock(& is active mutex);
                           return -1; //the object is active
               is active=true;
              pthread mutex unlock(& is active mutex);
              pthread mutex lock(& spec threads mutex);
             pthread_mutex_lock(&_global_access_log_mutex);
             \verb|vector < int>| threads_to_cancel|; //a | selective | reset | is | performed, | waiting | for | the | threads | in | in | threads | t
              \verb|_cancel_spec_threads_and_reset_logs(threads_to_cancel); // \textit{deferred cancel to finish, and reseting the logs}|
              //arrays that the speculation uses.
             pthread mutex lock(& is active mutex);
              has pre loop=false;
              pthread_mutex_unlock(&_is_active_mutex);
             pthread mutex lock(& valid til mutex);
              valid til=-1;
              pthread_mutex_unlock(&_valid_til_mutex);
```

```
pthread mutex lock(& pl mutex);
        pl commit=false;
        pthread_mutex_unlock(&_pl_mutex);
        shared data=shared data;
        pthread_attr_t attr;
        pthread_attr_init (&attr);
        pthread_attr_setschedpolicy(&attr, SCHED_FIFO);
        //the threads are created
        int success;
        _spec_threads.resize(thread_instructions.size());
        for (unsigned int i=0; i<thread instructions.size(); i++){</pre>
                pthread mutex lock(& spec threads[i]. thread mutex);
        for (unsigned int i=0; i<thread instructions.size(); i++){</pre>
                _spec_threads[i]._thread_instructions=thread_instructions[i];
                 _spec_threads[i]._const_args=const_args[i];
                success = pthread\_create \ \ \textbf{(\&\_spec\_threads[i].\_thread, \&attr, \_spec\_threads[i].\_thread\_instructions, \ spec\_threads[i].}
                if (success==0) {
                         _thread_index.insert(pair<pthread_t, int>(_spec_threads[i]._thread, i));
                pthread_mutex_unlock(&_spec_threads[i]._thread_mutex);
                if (success!=0){
                         for (unsigned int j=i+1; j<thread_instructions.size(); j++){</pre>
                                 pthread_mutex_unlock(&_spec_threads[j]._thread_mutex);
                         for (unsigned int j=0; j<i; j++){</pre>
                                 threads_to_cancel.push_back(j);
                         _cancel_spec_threads_and_reset_logs(threads_to_cancel);
                         pthread_mutex_unlock(&_global_access_log_mutex);
                         pthread mutex unlock (& spec threads mutex);
                         pthread_mutex_lock(&_is_active_mutex);
                          is active=false;
                         pthread_mutex_unlock(&_is_active_mutex);
                         return -1;// a thread could not be created.
        pthread_mutex_unlock(&_global_access_log_mutex);
pthread_mutex_unlock(&_spec_threads_mutex);
        return 0:
};
//! function to get the results of the speculation once the un-managed pre-
//! loop section has ended it's execution.
//!
//! TO BE NOTED:
//!
       * If the pre-loop section does not explicitly commit through the commit function,
      then the whole results of the speculation will not be communicated to the
//!
//!
      shared data and the speculation will be canceled.
//!
      * If any iteration does not commit, then the object will be in an infinite loop wa-
//! iting for it to commit.
int get_results(){
        pthread mutex lock(& is active mutex);
        if (!_is_active || has_pre_loop){
                pthread_mutex_unlock(&_is_active_mutex);
                return -1;//the object is inactive or manages it's pre-loop section, thus this function cannot be run.
        pthread mutex unlock(& is active mutex);
        bool commit_made=false;
        unsigned int joined_but_didnt_commit=0;
        int success;
        do f
                success=-1;
                 if (success!=0){
                         pthread_mutex_lock(&_pl_mutex);
                         if (_pl_commit){
                                 commit made=true;
                         else{
```

```
if (joined but didnt commit<10){</pre>
                                 success=0;
                                 joined but didnt commit++;
                pthread_mutex_unlock(&_pl_mutex);
} while (success==0);
int i=0:
pthread_mutex_lock(&_spec_threads_mutex);
if (!commit_made) {
                  cancel_higher_spec_threads(-1);
                pthread_mutex_unlock(&_spec_threads_mutex);
                {\tt pthread\_mutex\_lock(\&\_is\_active\_mutex);}
                 _is_active=false;
                pthread_mutex_unlock(&_is_active_mutex);
                return -1; //the pre-loop did not commit.
while ( spec threads.size()>i){
        pthread_t thread_to_join=_spec_threads[i]._thread;
        pthread_mutex_unlock(&_spec_threads_mutex);
        commit made=false;
        joined_but_didnt_commit=0;
        do (
                success=pthread_join (thread_to_join, NULL);
                if (success!=0) {
                         pthread_mutex_lock(&_spec_threads[i]._thread_mutex);
                         if (_spec_threads[i]._commit){
                                 commit_made=true;
                         else{
                                 if (joined_but_didnt_commit<100){</pre>
                                          success=0;
                                          joined_but_didnt_commit++;
                         pthread_mutex_unlock(&_spec_threads[i]._thread_mutex);
                }
        / * \verb"just" in case the commit is not logged in a timely fasion, it is tried
        10 times to see if it get's made, if not then the object asumes there was no commitment in the iteration.
        Although slightly un-elegant, the validity of this solution has not been disproved by empirical
        tests. */
        } while (success==0);
        if (!commit_made) {
                pthread_mutex_lock(&_spec_threads_mutex);
                 _cancel_higher_spec_threads(i); //This is perhaps a bit excesive, but for consistency purposes it is manda
                pthread_mutex_unlock(&_spec_threads_mutex);
                pthread_mutex_lock(&_is_active_mutex);
                  is_active=false;
                pthread_mutex_unlock(&_is_active_mutex);
                return -1; //some thread did not commit.
        }
        i++;
        pthread_mutex_lock(&_spec_threads_mutex);
        pthread_mutex_lock(&_valid_til_mutex);
int valid_til_thread=_valid_til;
        pthread_mutex_unlock(&_valid_til_mutex);
        if (valid_til_thread!=-1 && i>valid_til_thread) {
                i=_spec_threads.size();
```

```
pthread_mutex_unlock(&_spec_threads_mutex);
pthread_mutex_lock(&_is_active_mutex);
_is_active=false;
pthread_mutex_unlock(&_is_active_mutex);
return 0;
};
```

};