## Linguagem L1 com listas e exceções

Arthur Zachow

Gabriel Haggstrom

Gabriel Pakulski

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## Sintaxe abstrata

```
\in
                      Expr
      ::= n
                     e_1 op e_2
                     \mathtt{not}\ e_1
                     if e_1 then e_2 else e_3
                     e_1 e_2
                     \mathtt{fn}\ x\!:\!T\Rightarrow e
                     \mathtt{fn}\; x \Rightarrow e
                     \mathtt{let}\ x \colon\! T = e_1\ \mathtt{in}\ e_2
                     \mathbf{let}\ x = e_1\ \mathbf{in}\ e_2
                     let rec f:T_1 \to T_2 = (\operatorname{fn} y:T_1 \Rightarrow e_1) in e_2
                     let rec f = (\operatorname{fn} y \Rightarrow e_1) in e_2
                     nil
                     e_1 :: e_2
                     \mathtt{isempty}\; e
                     {\tt hd}\ e
                     {\tt tl}\; e
                     raise
                     {	t try} \ e_1 \ {	t with} \ e_2
                    Types
T \hspace{0.1in} ::= \hspace{0.1in} X \hspace{0.1in} | \hspace{0.1in} \mathsf{int} \hspace{0.1in} | \hspace{0.1in} \mathsf{bool} \hspace{0.1in} | \hspace{0.1in} T_1 \rightarrow \hspace{0.1in} T_2 \hspace{0.1in} | \hspace{0.1in} T \hspace{0.1in} \mathsf{list}
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onde

```
n \in conjunto de numerais inteiros

b \in \{true, false\}

x \in Ident

op \in \{+, -, *, div, ==, and, or\}
```

Semântica operacional big-step

$$v \in Values \\ v ::= n \mid b \mid nil \mid v_1 :: v_2 \mid \langle x, e, \text{env} \rangle \mid \langle f, x, e, \text{env} \rangle \\ env \vdash n \downarrow n \\ env \vdash n \downarrow n \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow n_2 \\ env \vdash e_1 = e_2 \downarrow raise \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_1 = e_2 \downarrow raise \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow raise \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_2 \downarrow raise \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_1 \downarrow raise \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_1 \downarrow raise \\ env \vdash e_1 \downarrow raise \\ env \vdash e_1 \downarrow n_1 \\ env \vdash e_1 \downarrow raise \\ env \vdash e_1 \downarrow rai$$

$$\frac{\operatorname{env} \vdash c_1 \Downarrow \operatorname{true} \quad \operatorname{env} \vdash c_2 \Downarrow v}{\operatorname{env} \vdash if \; c_1 \operatorname{then} \; c_2 \operatorname{else} \; c_3 \Downarrow v} \text{ (BS-IFRS)}$$

$$\frac{\operatorname{env} \vdash e_1 \Downarrow \operatorname{raise} \; \operatorname{env} \vdash e_1 \Downarrow v' \quad \{x \mapsto v'\} + \operatorname{env} \vdash e_2 \Downarrow v \; \operatorname{env} \vdash e_1 \Downarrow v' \quad \{BS-IFRS2\}$$

$$\frac{\operatorname{env} \vdash e_1 \Downarrow \operatorname{raise} \; \operatorname{env} \vdash e_2 \Downarrow v' \quad \{x \mapsto v'\} + \operatorname{env} \vdash e_2 \Downarrow v \; \operatorname{env} \vdash e_1 \Downarrow v' \quad \{BS-IFRS2\} \\ \operatorname{env} \vdash \operatorname{env} \vdash \operatorname{env} \vdash e_1 \Downarrow v' \quad \{x \mapsto v'\} + \operatorname{env} \vdash e_2 \Downarrow v \; \operatorname{env} \vdash \operatorname{env} \vdash$$

$$\begin{array}{c} \operatorname{env} \vdash \operatorname{nil} \Downarrow \operatorname{nil} & (\operatorname{BS-NIL}) & \frac{\operatorname{env} \vdash e_1 \Downarrow v_1 \quad \operatorname{env} \vdash e_2 \Downarrow v_2}{\operatorname{env} \vdash e_1 :: e_2 \Downarrow v_1 :: v_2} \ (\operatorname{BS-CONS}) \\ \\ \frac{\operatorname{env} \vdash e_1 \Downarrow \operatorname{raise}}{\operatorname{env} \vdash e_1 :: e_2 \Downarrow \operatorname{raise}} & (\operatorname{BS-CONSRS1}) & \frac{\operatorname{env} \vdash e_1 \Downarrow v_1 \quad \operatorname{env} \vdash e_2 \Downarrow \operatorname{raise}}{\operatorname{env} \vdash e_1 :: e_2 \Downarrow \operatorname{raise}} \ (\operatorname{BS-CONSRS2}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{nil}}{\operatorname{env} \vdash \operatorname{isempty}} e \Downarrow \operatorname{true} & (\operatorname{BS-ISEMPTYNIL}) & \frac{\operatorname{env} \vdash e \Downarrow v_1 :: v_2}{\operatorname{env} \vdash \operatorname{isempty}} e \Downarrow \operatorname{false} \ (\operatorname{BS-ISEMPTYCONS}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{hd} e \Downarrow \operatorname{raise}} & (\operatorname{BS-ISEMPTYRS}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{hd} e \Downarrow \operatorname{raise}} & (\operatorname{BS-HDCONS}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{hd} e \Downarrow \operatorname{raise}} & (\operatorname{BS-HDCONS}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{tl} e \Downarrow \operatorname{raise}} & (\operatorname{BS-HDCONS}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{tl} e \Downarrow \operatorname{raise}} & (\operatorname{BS-TLCONS}) \\ \\ \frac{\operatorname{env} \vdash e \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{tl} e \Downarrow \operatorname{raise}} & (\operatorname{BS-TLCONS}) \\ \\ \frac{\operatorname{env} \vdash e \sqcup \operatorname{raise}}{\operatorname{env} \vdash \operatorname{tl} e \Downarrow \operatorname{raise}} & (\operatorname{BS-TLRS}) \\ \\ \frac{\operatorname{env} \vdash e_1 \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{tl} e \Downarrow \operatorname{tr} e_2 \Downarrow \operatorname{raise}} & (\operatorname{BS-TRYRS2}) \\ \\ \frac{\operatorname{env} \vdash e_1 \Downarrow \operatorname{raise}}{\operatorname{env} \vdash \operatorname{tr} e_1 \Downarrow \operatorname{tr} \operatorname{tith} e_2 \Downarrow \operatorname{raise}} & (\operatorname{BS-TRYRS2}) \\ \\ \frac{\operatorname{env} \vdash \operatorname{tr} e_1 \Downarrow \operatorname{tr} \operatorname{taise}}{\operatorname{env} \vdash \operatorname{tr} e_2 \Downarrow \operatorname{raise}} & (\operatorname{BS-TRYRS2}) \\ \\ \frac{\operatorname{env} \vdash \operatorname{tr} e_1 \Downarrow \operatorname{tr} \operatorname{taise}}{\operatorname{env} \vdash \operatorname{tr} e_2 \Downarrow \operatorname{tr} \operatorname{taise}} & (\operatorname{BS-TRYRS2}) \\ \\ \frac{\operatorname{env} \vdash \operatorname{tr} \operatorname{tr}$$