FUNCTIONAL PROGRAMMING MT2018

Sheet 5

9.1 Suppose a type of natural numbers is defined by

```
> data Nat = Zero | Succ Nat
```

Use recursion to define functions $int :: Nat \rightarrow Int$ and $nat :: Int \rightarrow Nat$ which embed the natural numbers in Int in the obvious way.

Use recursion (on the second argument) to define functions

$$add, mul, pow, tet :: Nat \rightarrow Nat \rightarrow Nat$$

which implement addition, multiplication, exponentiation, and what Goodstein calls tetration.

(x ‘tet‘ $n = x$ ^ x ^ $\cdots$ ^ x where there are n copies of x.)

9.2 What property characterises $foldNat$, the fold for Nat? Define $foldNat$.

What are the deconstructors for Nat, and what characterises the unfold $unfoldNat$? Define $unfoldNat$.

Express each of int and nat as either $foldNat$ or $unfoldNat$.

Finally express add, mul, pow and tet as folds.

10.1 Prove directly by induction that

$$fold\ c\ n\ (xs +\!\!+ ys) \quad = \quad fold\ c\ (fold\ c\ n\ ys)\ xs$$

for all lists xs and ys (whether partial, finite or infinite).

10.2 Use fold fusion to show that the section $(+\!\!+ bs)$ is a fold.

Deduce without resort to induction that

$$fold\ c\ n\ (xs +\!\!+ ys) \quad = \quad fold\ c\ (fold\ c\ n\ ys)\ xs$$

10.3 Use fold fusion to show that $filter$ is a fold.

Deduce that

$$filter\ p\ (xs +\!\!+ ys) \quad = \quad filter\ p\ xs +\!\!+ filter\ p\ ys$$

10.4 A data type very like that of lists might be defined by

```
> data Liste a = Snoc (Liste a) a | Lin
```

There will be elements of *Liste* α and of $[\alpha]$ corresponding to finite lists, for example *Snoc* (*Snoc* (*Snoc Lin* 1) 2) 3 correspond to $1 : (2 : (3 : [\,]))$, that is $[1, 2, 3]$.

Write a recursive definition of a function $cat :: Liste\ \alpha \to Liste\ \alpha \to Liste\ \alpha$ which concatenates two elements of *Liste*.

Define a function *folde* which is the natural fold for *Liste* α.

Express *cat* in terms of *folde*.

Define (as folds) functions $list :: Liste\ \alpha \to [\alpha]$ and $liste :: [\alpha] \to Liste\ \alpha$ which express the identification of finite lists represented as elements of *Liste* α and of $[\alpha]$. (That is, they should be mutually inverse on finite lists.)

What does *liste* return when applied to an infinite list? What are the infinite objects of type *Liste* α?

Find equivalent definitions of *list* and *liste* as instances of *tailfold*.

10.5 Recall that the unfold function for $[\alpha]$

```
> unfold n h t x | n x       = []
>                | otherwise = h x : unfold n h t (t x)
```

yields the identity function *unfold null head tail* when applied to the deconstructors for $[\alpha]$.

Using the same property for the identity of *Liste* α define the unfold function *unfolde* for *Liste* α.

Write *list* and *liste* as the appropriate unfolds.

You may want to know that there are predefined functions $init :: [\alpha] \to [\alpha]$ and $last :: [\alpha] \to \alpha$ for which $xs = init\ xs +\!\!+ [last\ xs]$ for all non-null xs. It might help to work out first how these might be defined by recursion.

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