

# PROBABILITY PROBLEM SHEET 5

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1. Let  $p_m$  be the probability that the frog is on ① after  $m$  jumps. In the same manner, let  $q_m$  and  $s_m$  be the probabilities that the frog is on ② or on ③, respectively, after  $m$  jumps.

Obviously,  $p_m + q_m + s_m = 1$ , for all  $m \in \mathbb{N}$ .

We want to prove that  $q_m = s_m$  for all  $m \in \mathbb{N}$  by induction on  $m$ .

The base case  $Q(0)$ :  $q_0 = s_0$ .

At first, the frog is on ①, therefore  $p_0 = 1$ ,  $q_0 = 0$  and  $s_0 = 0 \Rightarrow q_0 = s_0$ .

The inductive step

IH: We know that  $Q(k)$ :  $q_k = s_k$  and we want to prove  $Q(k+1)$ :  $q_{k+1} = s_{k+1}$ .

From the fact that the frog jumps with probability  $\frac{1}{2}$  from a vertex to either of its adjacent vertices, we know that  $q_{k+1} = \frac{1}{2}(p_k + s_k)$  and  $s_{k+1} = \frac{1}{2}(p_k + q_k)$ . However, these two are equal because  $s_k = q_k$  from IH. So,  $q_{k+1} = s_{k+1}$ .

$$\text{Now, we have } p_m = \frac{1}{2}(q_{m-1} + s_{m-1}) \stackrel{Q(m-1)}{=} q_{m-1} = \frac{p_{m-2} + s_{m-2}}{2} \stackrel{Q(m-2)}{=} \frac{p_{m-2} + q_{m-2}}{2} = \frac{p_{m-2} + \frac{1}{2}(q_{m-3} + s_{m-3})}{2} = \frac{p_{m-2} + p_{m-3}}{2} \Rightarrow 2p_m - p_{m-1} - p_{m-2} = 0 \text{ for all } m \geq 2$$

$$\text{Let } p_m = A \lambda^m \Rightarrow 2A \lambda^m - A \lambda^{m-1} - A \lambda^{m-2} = 0 \quad | : A \lambda^{m-2} \neq 0$$

$$\begin{matrix} (\lambda \neq 0) \\ (A \neq 0) \end{matrix}$$

$$2\lambda^2 - \lambda - 1 = 0 \Rightarrow \lambda_1 = 1, \lambda_2 = -\frac{1}{2} \Rightarrow p_m = A_1 \cdot 1^m + A_2 \cdot \left(-\frac{1}{2}\right)^m$$

$$\text{From } p_0 = 1 \text{ and } p_1 = 0 \text{ we get } \begin{cases} A_1 + A_2 = 1 \\ A_1 - \frac{1}{2}A_2 = 0 \end{cases} \Rightarrow \begin{cases} A_1 = \frac{1}{3} \\ A_2 = \frac{2}{3} \end{cases} \Rightarrow p_m = \frac{1}{3} + \frac{2}{3} \left(-\frac{1}{2}\right)^m \text{ for all } m \in \mathbb{N}.$$

$$\text{When } m \rightarrow \infty \text{ we have } \lim_{m \rightarrow \infty} p_m = \lim_{m \rightarrow \infty} \frac{1}{3} + \frac{2}{3} \cdot \left(-\frac{1}{2}\right)^m = \frac{1}{3} + \frac{2}{3} \cdot 0 = \frac{1}{3} \Rightarrow \boxed{\lim_{m \rightarrow \infty} p_m = \frac{1}{3}}$$

2. The floor plan of the house can be represented by a graph with 6 nodes (each representing a room), where the edges represent the fact that the mouse can get from one room to another directly.

Let  $X$  be the number of minutes until the mouse gets in the 6th room. We write  $e_i$  for the expectation of  $X$  when the mouse starts from the  $i$ th room, with  $i \in \{1, 2, 3, 4, 5, 6\}$ . We know the mouse walks randomly, so we can say that if he has  $m$  choices for the next step, then each of them has a probability of  $\frac{1}{m}$  to be followed.

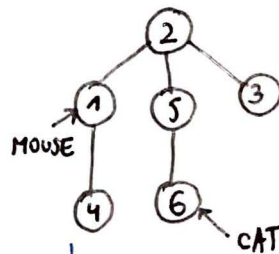
We want to find  $e_1$ .

$$\text{Now, } e_1 = \frac{1}{2}(e_4 + 1) + \frac{1}{2}(e_2 + 1)$$

$$e_2 = \frac{1}{3}(e_1 + 1) + \frac{1}{3}(e_5 + 1) + \frac{1}{3}(e_3 + 1)$$

$$e_3 = e_2 + 1$$

$$e_4 = e_1 + 1$$



$$e_5 = \frac{1}{2}(e_2+1) + \frac{1}{2}(e_6+1)$$

$e_6 = 0$  (the mouse already is in the 6<sup>th</sup> room)

By solving the linear system we get  $e_2 = 16, e_3 = 17, e_4 = 20, e_5 = 9, e_6 = 0$  and

$$\boxed{e_1 = 19}$$

3. (Gambler's ruin, symmetric case)

The gambler starts with £ $m$ , where  $m \in \{1, 2, \dots, M-1\}$ , he wins £1 with probability  $\frac{1}{2}$  and loses £1 with the same probability. The game ends when he reaches £0 or £ $M$ .

(a) Let  $X$  denote the amount of money the gambler ends up with. We write  $e_m$  for the expectation of  $X$  when he starts with £ $m$ . Then

$$e_m = \frac{1}{2} \mathbb{E}(X | \text{first game win}) + \frac{1}{2} \mathbb{E}(X | \text{first game lose})$$

$$e_m = \frac{1}{2}(e_{m+1}+1) + \frac{1}{2}(e_{m-1}-1)$$

$$e_m = \frac{1}{2}e_{m+1} + \frac{1}{2}e_{m-1} \Rightarrow e_{m+1} - 2e_m + e_{m-1} = 0 \quad \left| \Rightarrow A\lambda^{m+1} - 2A\lambda^m + A\lambda^{m-1} = 0 : A\lambda^{m-1} \neq 0 \right.$$

$$\text{Let } e_m = A\lambda^m, A \neq 0, \lambda \neq 0 \quad \left| \begin{array}{l} \lambda^2 - 2\lambda + 1 = 0 \\ \lambda_{1/2} = 1 \Rightarrow \lambda = 1 \end{array} \right.$$

$$\lambda_{1/2} = 1 \Rightarrow \lambda = 1$$

$$\text{Then, } e_m = (A+Bm)\lambda^m$$

$$\text{From } e_0 = 0 \text{ and } e_M = M \text{ we have } (A+B \cdot 0)\lambda^0 = 0 \Rightarrow A = 0 \text{ and } (A+B \cdot M)\lambda^M = M \Rightarrow$$

$$\Rightarrow B \cdot M \cdot 1^M = M \Rightarrow B = 1.$$

$$\text{Therefore, } \boxed{e_m = m}$$

(b) We want to calculate the conditional probability of the fact that he won £1 on his first game, given the fact that he ends the game with £ $M$ . ← these two are not independent

$$\begin{aligned} P(\text{win on 1st game} | \text{ends with } \pounds M) &= \frac{P(\{\text{win on 1st game}\} \cap \{\text{ends with } \pounds M\})}{P(\{\text{ends with } \pounds M\})} \quad \leftarrow \text{from lectures, this is } \frac{m}{M} \\ &= \frac{P(\{\text{win on 1st game}\} \cap \{\text{ends the game with } \pounds M, \text{ starting from } \pounds(m+1)\})}{P(\{\text{ends the game with } \pounds M, \text{ starting from } \pounds m\})} \quad \leftarrow \text{these 2 are independent} \\ &= \frac{P(\text{win on 1st game}) P(\text{ends with } \pounds M, \text{ starting from } \pounds(m+1))}{P(\text{ends with } \pounds M, \text{ starting from } \pounds m)} = \frac{\frac{1}{2} \cdot \frac{m+1}{M}}{\frac{m}{M}} = \frac{m+1}{2m} \Rightarrow \end{aligned}$$

$$\Rightarrow \boxed{P(\text{win } \pounds 1 \text{ on first game} | \text{he ends the game with } \pounds M, \text{ starting from } \pounds m) = \frac{m+1}{2m}}$$



(c) Now  $X$  denotes the number of steps (length) until the game ends. We now write  $e_m$  for the expectation of  $X$  when the game starts from  $\&m$ . Then

$$e_m = \frac{1}{2} \mathbb{E}(X | \text{first round is a win}) + \frac{1}{2} \mathbb{E}(X | \text{first round is a lose})$$

$$e_m = \frac{1}{2} (e_{m+1} + 1) + \frac{1}{2} (e_{m-1} + 1)$$

$$e_m = \frac{1}{2} e_{m+1} + \frac{1}{2} e_{m-1} + 1 \Rightarrow \frac{1}{2} e_{m+1} - e_m + \frac{1}{2} e_{m-1} = -1 \quad | \cdot 2$$

$$e_{m+1} - 2e_m + e_{m-1} = -2$$

From (a) we know that  $\boxed{W_m = m}$ .

Now we'll try to find the particular solution  $v_m$ . Trying  $v_m = Cm^2 + Dm$  gets us to:

$$C(m+1)^2 + D(m+1) - 2Cm^2 - 2Dm + C(m-1)^2 + D(m-1) = -2$$

$$Cm^2 + 2Cm + C + Dm + D - 2Cm^2 - 2Dm + Cm^2 - 2Cm + C + Dm - D = -2$$

$$2C = -2 \Rightarrow \boxed{C = -1}$$

Therefore, we have  $\boxed{v_m = -m^2 + Dm} \Rightarrow e_m = -m^2 + Dm + m$ .

By using the boundary conditions to find  $D$ , we get  $e_0 = 0$ , which happens for all  $D \in \mathbb{R}$ , and for  $e_M = 0$ , we obtain  $-M^2 + DM + M = 0 \quad | : M \neq 0$

$$-M + D + 1 = 0 \Rightarrow D = M - 1 \Rightarrow e_m = -m^2 + (M-1)m + m \Rightarrow$$

$$\Rightarrow e_m = -m^2 + Mm \Rightarrow \boxed{e_m = m(M-m)} \text{ for each } m \in \{1, 2, \dots, M-1\} \text{ (it works for } m=0 \text{ and } m=M, \text{ too).}$$

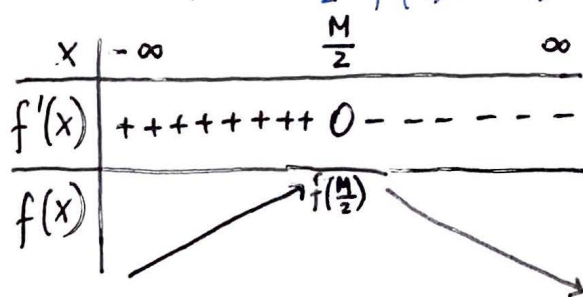
We want to know for which value of  $m$ ,  $e_m$  is the largest.

For that we'll consider  $f: \mathbb{R} \rightarrow \mathbb{R}$ ,  $f(x) = xM - x^2$ , which is a continuous function. We have

$f'(x) = M - 2x$ , so for  $x \in (-\infty, \frac{M}{2})$   $f'(x) > 0 \Rightarrow f(x)$  is strictly increasing on  $(-\infty, \frac{M}{2})$

for  $x \in (\frac{M}{2}, \infty)$   $f'(x) < 0 \Rightarrow f(x)$  is strictly decreasing on  $(\frac{M}{2}, \infty)$

for  $x = \frac{M}{2}$   $f'(x) = 0 \Rightarrow x = \frac{M}{2}$  is a critical point



Therefore,  $\max(f(x)) = f(\frac{M}{2})$ .

Now, coming back to the problem, we can deduce that if  $M$  is even, then  $\max(e_m) = e_{\frac{M}{2}} \Rightarrow$

$$\Rightarrow \boxed{\max(e_m) = \frac{M}{2} \cdot \frac{M}{2}} \text{ and if } M \text{ is odd, then we can deduce that } \max(e_m) = e_{\lfloor \frac{M}{2} \rfloor} = e_{\lceil \frac{M}{2} \rceil} \Rightarrow$$

$$\Rightarrow \boxed{\max(e_m) = \lfloor \frac{M}{2} \rfloor \cdot \lceil \frac{M}{2} \rceil}$$

⊛ Proof: Let  $M = 2p+1 \Rightarrow \lfloor \frac{M}{2} \rfloor = p$ ,

$$\lceil \frac{M}{2} \rceil = p+1. e_{\lfloor \frac{M}{2} \rfloor} = e_p = p(2p+1-p) \Rightarrow$$

$$\Rightarrow e_{\lfloor \frac{M}{2} \rfloor} = p(p+1) \text{ and } e_{\lceil \frac{M}{2} \rceil} = e_{p+1} \Rightarrow$$

$$\Rightarrow e_{p+1} = (p+1)(2p+1-p-1) = (p+1)p \Rightarrow e_p = e_{p+1}$$

$$\Rightarrow \boxed{e_{\lfloor \frac{M}{2} \rfloor} = e_{\lceil \frac{M}{2} \rceil}}$$

easy to prove ⊛

In conclusion,  $\boxed{\max(e_m) = \lfloor \frac{M}{2} \rfloor^2}$  for all  $M \in \mathbb{N}$  (for  $M$  even  $\lfloor \frac{M}{2} \rfloor = \frac{M}{2}$ ) and it

happens when  $\boxed{m = \lfloor \frac{M}{2} \rfloor}$ .

(a) Let  $X \sim \text{Geom}(p) \Rightarrow P_k = p(1-p)^{k-1}$ ,  $k=1, 2, \dots$

We have  $G_X(s) = \sum_{k=1}^{\infty} s^k p(1-p)^{k-1} = \sum_{k=1}^{\infty} (ps) \underset{a=k-1}{(s(1-p))^{k-1}} = (ps) \sum_{a=0}^{\infty} (s(1-p))^a = (ps) \cdot \frac{1}{1-s(1-p)}$

Therefore,  $G_X(s) = \frac{ps}{1-(1-p)s}$ , provided that  $|s| < \frac{1}{1-p}$ .

this has to be less than 1 (the absolute value), so  $|s| < \frac{1}{1-p}$ .

(b)  $E(X) = G'_X(1)$

$$G'_X(s) = \frac{p(1-s+ps) - ps(-(1-p))}{(1-s+ps)^2} = \frac{p - p/s + p^2/s + p/s - p^2/s}{(1-s+ps)^2} = \frac{p}{(1-s+ps)^2} \Rightarrow$$

$$\Rightarrow E(X) = \frac{p}{(1-1+p)^2} = \frac{p}{p^2} \Rightarrow \boxed{E(X) = \frac{1}{p}}$$

$$G''_X(s) = \left( \frac{p}{(1-s+ps)^2} \right)' = \frac{-p \cdot 2(1-s+ps) \cdot (-1+p)}{(1-s+ps)^4} = \frac{(-2p)(p-1)}{(1-s+ps)^3} = \frac{2p(1-p)}{(1-s+ps)^3}$$

$$\text{Var}(X) = G''_X(1) + G'_X(1) - (G'_X(1))^2$$

$$\text{Var}(X) = \frac{2p(1-p)}{p^3} + \frac{p^2}{p} - \frac{p}{p^2} = \frac{2p - 2p^2 + p^2 - p}{p^3} = \frac{p - p^2}{p^3} \Rightarrow \boxed{\text{Var}(X) = \frac{1-p}{p^2}}$$

5. (a) A fair coin is tossed  $n$  times  $\Rightarrow p(H) = p(T) = \frac{1}{2}$ . Let  $r_n$  be the probability that the sequence of tosses never has a head followed by a head.

So,  $r_n = P(\text{n-sequence ok}) = P(\text{n-sequence ok} | T \text{ last}) \cdot P(T \text{ last}) + P(\text{n-sequence ok} | TH \text{ last 2}) \cdot P(TH \text{ last 2}) + P(\text{n-sequence ok} | HH \text{ last 2}) \cdot P(HH \text{ last 2})$

$$r_n = P((n-1)\text{-sequence ok}) \cdot \frac{1}{2} + P((n-2)\text{-sequence ok}) \cdot \frac{1}{4} + 0 \cdot \frac{1}{4}$$

$$\boxed{r_n = \frac{1}{2} r_{n-1} + \frac{1}{4} r_{n-2}} \text{ for } n \geq 2 \text{ and } r_0 = r_1 = 1.$$

$$r_2 = \frac{1}{2} + \frac{1}{4} = \frac{3}{4} : \text{favourable cases} = 3 \text{ (HT, TH, TT)} / \text{total cases} = 4 \text{ (HH, HT, TH, TT)}$$

(b) Let  $X$  denote the number of coin tosses needed until we first get two heads in a row. We want to calculate  $P(X=k)$ , for  $k \in \mathbb{N}, k \geq 2$ . So, in order to first get two heads in a row after exactly  $k$  coin tosses we need that we didn't get any HH for the first  $k-3$  tosses and the last 3 results must be T, H, H. Therefore,

$$\boxed{P(X=k) = r_{k-3} \cdot \frac{1}{8}} \text{, for } k \geq 3, \text{ and for } k=2 \text{ it is } P(X=2) = \frac{1}{4} \text{ (HT)}$$

$$(a) \quad r_n - \frac{1}{2} r_{n-1} - \frac{1}{4} r_{n-2} = 0 \Rightarrow A \lambda^n - \frac{1}{2} A \lambda^{n-1} - \frac{1}{4} A \lambda^{n-2} = 0 \mid : A \lambda^{n-2} \neq 0$$

$$\text{Let } r_n = A \lambda^n, A \neq 0, \lambda \neq 0$$

$$\lambda^2 - \frac{1}{2} \lambda - \frac{1}{4} = 0 \mid \cdot 4$$

$$4\lambda^2 - 2\lambda - 1 = 0$$

$$(\Delta = 4 + 16 = 20) \Rightarrow \lambda_1 = \frac{1+\sqrt{5}}{4}, \lambda_2 = \frac{1-\sqrt{5}}{4}$$

Therefore,  $r_n = A \lambda_1^n + B \lambda_2^n$  and by using the boundary conditions we get:



$$\begin{cases} A+B=1 \Rightarrow B=1-A \\ A \cdot \frac{1+\sqrt{5}}{4} + B \cdot \frac{1-\sqrt{5}}{4} = 1 \end{cases} \Rightarrow A \cdot \frac{1+\sqrt{5}}{4} + \frac{1-\sqrt{5}}{4} - A \cdot \frac{1-\sqrt{5}}{4} = 1$$

$$A \cdot \left( \frac{1+\sqrt{5}}{4} - \frac{1-\sqrt{5}}{4} \right) = 1 - \frac{1-\sqrt{5}}{4}$$

$$A \cdot \frac{2\sqrt{5}}{4} = \frac{3+\sqrt{5}}{4} \Rightarrow A = \frac{3+\sqrt{5}}{2\sqrt{5}} \Rightarrow \boxed{A = \frac{5+3\sqrt{5}}{10}} \Rightarrow \boxed{B = \frac{5-3\sqrt{5}}{10}}$$

$$\text{So, } r_m = \frac{5+3\sqrt{5}}{10} \cdot \left( \frac{1+\sqrt{5}}{4} \right)^m + \frac{5-3\sqrt{5}}{10} \cdot \left( \frac{1-\sqrt{5}}{4} \right)^m$$

$$r_2 = \frac{5+3\sqrt{5}}{10} \cdot \frac{6+2\sqrt{5}}{16} + \frac{5-3\sqrt{5}}{10} \cdot \frac{6-2\sqrt{5}}{16} = \frac{30+28\sqrt{5}+30}{160} + \frac{30-28\sqrt{5}+30}{160} = \frac{120}{160} = \frac{3}{4}, \text{ so it's correct!}$$

$$(c) P(X=k) = \frac{r_{k-3}}{8} \text{ from (b), for } k \geq 3$$

$$\text{For } k=0 \text{ we have } P(X=0)=0, \text{ for } k=1 \text{ } P(X=1)=0 \text{ and for } P(X=2)=\frac{1}{4}.$$

Now,

$$G_X(s) = \sum_{k=0}^{\infty} s^k P(X=k) = 0 \cdot s^0 + 0 \cdot s^1 + \frac{1}{4} \cdot s^2 + \sum_{k=3}^{\infty} s^k \cdot \frac{1}{8} r_{k-3} = \frac{1}{4} s^2 + \frac{1}{8} \sum_{k=3}^{\infty} s^k \cdot r_{k-3}$$

if we replace  $k-3$  with  $a$  we get:

$$G_X(s) = \frac{1}{4} s^2 + \frac{1}{8} \sum_{a=0}^{\infty} s^{a+3} \cdot r_a = \frac{1}{4} s^2 + \frac{1}{8} s^3 \cdot \sum_{a=0}^{\infty} s^a \cdot r_a$$

$$\text{But } r_a = A \cdot \lambda_1^a + B \cdot \lambda_2^a, \text{ so}$$

$$G_X(s) = \frac{1}{4} s^2 + \frac{1}{8} s^3 \cdot \left( A \sum_{a=0}^{\infty} (s\lambda_1)^a + B \sum_{a=0}^{\infty} (s\lambda_2)^a \right)$$

The two sums converge only if  $|s\lambda_1| < 1$  and  $|s\lambda_2| < 1$  and that implies  $|s| < \frac{4}{1+\sqrt{5}}$ . We will need the function value (the derivative) for  $s=1$ , so we can calculate it.

$$G_X(s) = \frac{1}{4} s^2 + \frac{1}{8} s^3 \cdot \left( \frac{A}{1-s\lambda_1} + \frac{B}{1-s\lambda_2} \right) = \frac{1}{4} s^2 + \frac{1}{8} s^3 \cdot \frac{A(1-s\lambda_2) + B(1-s\lambda_1)}{(1-s\lambda_1)(1-s\lambda_2)}$$

Now we'll make use of these results:

$$A+B=1, \lambda_1+\lambda_2=\frac{1}{2}, \lambda_1\lambda_2=-\frac{1}{4}, A\lambda_2+B\lambda_1=-\frac{1}{2}, \text{ so}$$

$$G_X(s) = \frac{1}{4} s^2 + \frac{1}{8} s^3 \cdot \frac{(A+B) - s(A\lambda_2 + B\lambda_1)}{1 - s(\lambda_1 + \lambda_2) + (\lambda_1\lambda_2)s^2} = \frac{1}{4} s^2 + \frac{1}{8} s^3 \cdot \frac{1 + \frac{1}{2}s}{1 - \frac{1}{2}s - \frac{1}{4}s^2} = \frac{s^2}{4} + \frac{s^3}{8} \cdot \frac{4+2s}{4-2s-s^2}$$

$$G_X(s) = \frac{s^2}{4} + \frac{2s^4 + 4s^3}{-8s^2 - 16s + 32} = \frac{s^2}{4} + \frac{s^4 + 2s^3}{-4s^2 - 8s + 16} = \frac{s^2(-s^2 - 2s + 4) + s^3 + 2s^3}{-4s^2 - 8s + 16}$$

$$G_X(s) = \frac{-s^4 - 2s^3 + 4s^2 + 2s^3}{-4s^2 - 8s + 16} \Rightarrow \boxed{G_X(s) = \frac{s^2}{-s^2 - 2s + 4}}, \text{ provided that } |s| < \frac{4}{1+\sqrt{5}}$$

If we want to calculate the mean of  $X$ , we need

$$G'_X(1), \text{ so we'll calculate } G'_X(s) = \left( \frac{s^2}{-s^2 - 2s + 4} \right)' = \frac{2s(-s^2 - 2s + 4) - s^2(-2s - 2)}{(s^2 + 2s - 4)^2}$$

$$G'_X(s) = \frac{-2s^3 - 4s^2 + 8s + 2s^3 + 2s^2}{(s^2 + 2s - 4)^2} = \frac{-2s^2 + 8s}{(s^2 + 2s - 4)^2} \Rightarrow \boxed{G'_X(s) = \frac{2s(4-s)}{(s^2 + 2s - 4)^2}}$$

The answer I got at Q6 on Sheet 3 was  $\frac{p+1}{p^2}$ , where  $p$  was the probability to get heads.  
 here,  $p = \frac{1}{2} \Rightarrow \mathbb{E}(X) = \frac{\frac{3}{2}}{\frac{1}{4}} = 6$ .

Calculating  $G_X'(1) = \frac{2 \cdot 1 \cdot 3}{(1+2-1)^2} = \frac{6}{(-1)^2} = 6 \Rightarrow \boxed{\mathbb{E}(X) = 6}$  (checks with the answer!)

(d) Let  $Y$  denote the number of coin tosses needed until we first see a TH.

Proceeding the same way as before, let  $t_n$  be the probability that we didn't obtain any TH in a sequence of  $n$  tosses. Then,  $t_0 = t_1 = 1$  and

$$t_n = P(\text{n-sequence OK}) = P(\text{n-sequence OK} \mid \text{last T})P(\text{last T}) + P(\text{n-sequence OK} \mid \text{last 2 TH})P(\text{TH}) + P(\text{n-sequence OK} \mid \text{last 2 HH})$$

$$t_n = t_{n-1} \cdot \frac{1}{2} + 0 \cdot \frac{1}{4} + \frac{1}{2^n} \Rightarrow \boxed{t_n = \frac{1}{2} t_{n-1} + \frac{1}{2^n}}, \text{ with } n \geq 1$$

↑ if we have last 2 HH, all tosses must be H

Solving it, we obtain  $\boxed{t_n = \frac{n+1}{2^n}}$ , for all  $n \in \mathbb{N}$  (can be proven by induction)

$P(Y > n)$  is the probability to get the first TH in a sequence with more than  $n$  tosses  $\Rightarrow$  it is the probability to not get a TH in a sequence of  $n$  tosses  $\Rightarrow$

$$\Rightarrow P(Y > n) = t_n \Rightarrow \boxed{P(Y > n) = \frac{n+1}{2^n}}$$

In the same manner,  $P(X > n) = r_n \Rightarrow P(X > n) = \frac{5+3\sqrt{5}}{10} \left(\frac{1+\sqrt{5}}{4}\right)^n + \frac{5-3\sqrt{5}}{10} \left(\frac{1-\sqrt{5}}{4}\right)^n$ .

We will calculate  $\frac{r_n}{r_{n-1}}$ , as because of the <sup>fact</sup> that when  $n \rightarrow \infty$  we get close to the point

where  $\frac{r_n}{r_{n-1}} \approx \frac{r_{n-1}}{r_{n-2}} = a$ . We have  $a = \frac{r_n}{r_{n-1}} = \frac{\frac{1}{2} r_{n-1} + \frac{1}{4} r_{n-2}}{r_{n-1}} = \frac{1}{2} + \frac{1}{4} \frac{r_{n-2}}{r_{n-1}} = \frac{1}{2} + \frac{1}{4} \cdot \frac{1}{a}$

$$a = \frac{1}{2} + \frac{1}{4a} \Rightarrow 4a^2 = 2a + 1$$

$$4a^2 - 2a - 1 = 0 \Rightarrow a_{1/2} = \frac{2 \pm 2\sqrt{5}}{8} = \frac{1 \pm \sqrt{5}}{4}$$

But  $a > 0 \Rightarrow a = \frac{1+\sqrt{5}}{4}$ , so  $r_n$  decreases with a factor of  $\frac{4}{1+\sqrt{5}}$

in general.

Now,  $\frac{t_n}{t_{n-1}} = \frac{n+1}{2^n} \cdot \frac{2^{n-1}}{n} = \frac{n+1}{2n} \xrightarrow{n \rightarrow \infty} \frac{1}{2}$ , so  $t_n$  decreases with a factor of 2.

From the fact that  $2 > \frac{4}{1+\sqrt{5}}$  ( $2+2\sqrt{5} > 4 \Leftrightarrow 2\sqrt{5} > 2 \Leftrightarrow \sqrt{5} > 1$  YES) we deduce that

$t_n < r_n \Rightarrow \boxed{P(Y > n) < P(X < n)}$ , which is not surprising because it is definitely more plausible to get a HH later than a TH in a configuration of tosses.



6. We consider a symmetric random walk on a cycle with  $N$  sites, labelled  $0, 1, 2, \dots, N-1$ . We start at 0 and from  $i$  we can go to  $(i+1) \bmod N$  or  $(i-1) \bmod N$  with  $p = \frac{1}{2}$  (for each) independently.

(a) We want to find the expected number of steps until every site has been visited. Let  $X_m$  denote the number of steps until we have visited  $m$  different sites (the last visit being a new site). Therefore, let  $e_m$  be the expected number of steps until  $m$  different sites have been visited.

Notice that because from one site we can only go to one of its neighbours, the set of visited sites is always of the form  $\{N-j, N-j+1, \dots, N-2, N-1, 0, 1, 2, \dots, m-j-1\}$  when we have visited  $m$  different sites, with  $j \in \{0, 1, 2, \dots, m+1\}$ . We can renumber them for  $\{1, 2, \dots, m\}$ . We now are at 1 or  $m$  (after  $e_m$  number of steps). We want to find the expected number of steps until we either get to 0 or to  $(m+1)$ . We can view this problem similarly to the Gambler's Ruin problem. You start from 1 or  $m$ , the ground level is 0, the top level is  $m+1$ .

From 3.c), the expected number of steps is  $1(m+1-x) = m(x+1-x) = m$ , therefore we have  $e_{m+1} = e_m + m$ , for  $m \in \mathbb{N}_+$  and  $e_1 = 0$ . So,  $e_m = \frac{(m-1)m}{2}$ , and as we want to see the number of steps needed for us to visit all the  $N$  different sites, then the result is  $e_N = \frac{(N-1)N}{2}$ .

(b) Now we want to calculate the probability that  $k$  is the last site to be visited, for each  $k = 1, 2, \dots, N-1$ .

$$P(\text{last site was } k) = P(\text{last site was } k \mid \text{before, we visited } k-1) P(\text{before, we visited } k-1) + P(\text{last site was } k \mid \text{before, we visited } k+1) P(\text{before, we visited } k+1).$$

$$\textcircled{1} P(\text{last site was } k \mid \text{before, we visited } k-1) = ?$$

So, now we are in  $k-1$  and we want to get to  $k$ , but without going through the way with  $k+1 \rightarrow k$ . For an easier reasoning, we'll renumber the sites (we'll add  $k$  to each, and then mod  $m$ ). So, instead of having  $0, 1, 2, \dots, N-1$ , we'll have  $k, k+1, \dots, 0, \dots, k-1$ . We'll also use the fact that  $P(\text{last visited was } k) = P(\text{last visited was } N-k)$ , as the directions in which we go through the cycle are left or right and they equal probability.

Now, we basically want to get to the site  $(N-k+k)$  from site  $k$ , by going through 1 (its right). We can see this as a gambler's ruin game, where the bottom is 0, we are at 1, the top is  $N$  (we want to win  $N$ ). So the probability to win here is  $\frac{1}{N}$ .

$$\textcircled{2} P(\text{before, we visited } k-1) = ?$$

In our reasoning, we want to first get to  $N-k+1+k$ , which is 1. So, the bottom is now 1, we are at  $k$ , the top is  $N$ , so basically we can say that we want to lose a game where the bottom is 0, the position is  $k-1$  and the top is  $N-1$ , where the probability is  $(1 - \frac{k-1}{N-1})$ .

③  $P(\text{last site was } k | \text{before, we visited } k+1) = ?$

So, now we are at  $N-1$  and we want to get to 0 from its left (or to  $N$  if we replace it with  $N$ , as  $N \bmod N = 0$  anyways). So, we want to lose the game where the ground is 0, our position is  $N-1$  and the top is  $N$  (we don't want to get to  $N$ ). Therefore, the probability is  $1 - \frac{N-1}{N} = \frac{1}{N}$ .

④  $P(\text{before, we visited } k+1) = ?$

In our reordering, we first want to get to  $N-1$ , starting from  $k$ , without going through 0. So, we want to win the game where the ground is 0, our position is  $k$  and the top is  $N-1$ , so the probability is  $\frac{k}{N-1}$ .

Therefore, our probability is:

$$P(\text{last site was } k) = \frac{1}{N} \left(1 - \frac{k-1}{N-1}\right) + \frac{1}{N} \cdot \frac{k}{N-1} = \frac{1}{N} \left(1 - \frac{k-1}{N-1} + \frac{k}{N-1}\right) = \frac{1}{N} \cdot \frac{N}{N-1} = \frac{1}{N-1}$$

As  $k$  takes values from 1 to  $N-1$ , we conclude that there's an equal probability for any site to be visited last (starting from site 0).

The reordering transformed the sequence:

$k-2, k-1, \boxed{k}, k+1, k+2, \dots, N-1, 0, 1, \dots, k-1, \boxed{k}, k+1$

$\leftarrow \text{FINISH}$                        $\leftarrow \text{START}$                        $\leftarrow \text{FINISH}$   
 $\uparrow$                        $\uparrow$                        $\uparrow$   
 Case 2                      Case 1

into:

$N-2, N-1, \boxed{0}, 1, 2, \dots, k-1, k, k+1, \dots, N-1, \boxed{0}, 1$

$\leftarrow \text{FINISH}$                        $\leftarrow \text{START}$                        $\leftarrow \text{FINISH}$   
 $\uparrow$                        $\uparrow$                        $\uparrow$   
 Case 1                      Case 2