QUESTION 3

(a)

> (4) :: [a] -> [a] -> [a]

> [] # 1/s = 1/s > (x:xs) + ys = x: (xs + ys)

The time-complexity of (++) is O(N), where N=lungth xs

> neverse :: [a] -> [a]

) reverse [] = []

> neverse (x:xs) = neverse xs # [x]

(b) Let T(N) be the number of steps needed for neverse xs, where length xs = N Then, we have $T(N) \leq T(N-1) + (N-1) \Rightarrow T(N) \leq 1 + 2 + \dots + (N-1) = \frac{(N-1)N}{3} \Rightarrow T(N) = O(N^2)$

(c) > data Catlist a = Nil | One a | Cat (Catlist a) (Catlist a)

· [] can be represented as Nil, Cat Nil Nil, Cat (Cat Nil Nil) Nil

· [x] can be represented as One x, Cat (One x) Nil, Cat Nil (One x)

(q)

> nep :: [a] -> Catlist a > nep []=Nil

> nep (x:xs) = Cat (One x) (nep xs)

(e)

> abs :: Catlist a -> [a] > abs' = flatCat []

> flatCot :: [a] -> CatList a -> [a]

> flatCat ys Nil = ys

> flatcat ys (One x) = x: ys

> flatcat ys (cat l n) = flatcat (flatcat ys n) l

The time complexity needed is O(N) as we only made use of (:), and not (#), where N is the size of the list of type Catlist.

(f) The "abs" function is not in general linear in the length of the returned list, as for example (f) line returned list could be made of N time Nil and the element, where N is each element the returned list, so it makes it run graduation the in N

of the neturned list, so it makes it run graduatically in N.

(g)

> nev :: Catlist a -> Catlist a

> nev Nil = Nil

> nev (One x) = One x

> nev (Catl n) = Cat (nev n) (nev l)

This nums in O(N), where N is the size of the imput.

(h)

> never = abs'. nev. nep

We first transform the list into a Catlist using "nep" im O(N) time, then neverse it using "nev" im
O(N) time and then transform it back to a normal list using "abs', using O(N) time.

Therefore, the new definition of "neverse" is linear in the length of the argument.