Question 3

- (a) > interleave :: [a] -> [a] -> [a]
 - > interleave (x:xs) ys = x: interleave ys xs
 - > interleave [] ys = ys
- (b) The first 6 elements of the result of interleave [2,2] [1...] are 2,1,2,2,3,4 as we first alternate the first 2 elements of each list and then we print the rest of the infinite list.
 - (c) > interleave List :: [[a]] → [a]
 - > interleave List = folds interleave []

As the function interleave ensures us that every element of an input list can be found after a finite number of steps, then the interleave list function, which is built based on multiple applications of the interleave function, is ensured to have the same property.

Comment: Here, I feel that my explication lacks some more specific justification and it is more of an intuitive argumentation. How can I express myself more rigorously?

- (d) The first 8 elements of the result of interleave List [[1,2], [3,4], [5,6], [7,8], [9,10]...] are 1,3,2,5,4,7,6,9. Let's say that we got to the point where we have the elements of the accumulator x and y and we now interleave the accumulator with [9,10]. Therefore, the first 4 elements of the resulting list will be [9, x, 10, y], according to the definition of interleave. Then, after interleaving the list with [7,8], the first 6 elements will be [7,9,8,x,10,y], then [5,7,6,9,8,x,10,y], then [3,5,4,7,6,9,8,x,10,y], and finally [1,3,2,5,4,7,6,9,8,x,10,y], so the first 8 elements are 1,3,2,5,4,7,6,9.
 - (e) > allpains :: [a] -> [b] -> [(a,b)]

> allpains xs ys = [(x,y) | x <- xs, y <- ys]

The problem is that if the ystis infinite, then there will be pairs which won't be reached in a finite number of steps, therefore we don't get a correct result. Let's imagine that is a list with at least 2 dements, the first being a and the second, b. Then, with the current definition of allpairs, the pairs of the form (a,y) with y from ys are returned first, and there fore mo pair of the form (b,y) with y from ys will be reached in a finite number of steps, as there are infinitely many pairs before.

(f) > allpains 2 :: [a] -> [b] -> [(a,b)]
> allpains 2 xs ys = interleave List [[(x,y) | y <- ys] | x <- xs]

By interleaving (with interlegist) a finite or infinite number of lists (finite or infinite that together form the Cartesian product of the two imput lists, we are ensured (as we proved at (c)) that every pair can be found in a finite number of steps.

(g) > data Tree = Nil | Fork Tree Tree

> all trees :: [Tree]

> all thees = Nil: nest

> where nest = map (uncurry Fork) (allpains 2 alltrees all trees)

This way, we create all the finite elements of the datatype Tree. We apply the (uncurry Fork) function to every pair of the Cartesian product of two alltrees arguments (therefore this definition is recursive) to obtain new trees which have the left subtree the first element of a specific pair from (allpairs 2 alltrees alltrees) and the right subtree the second one. Recall that we used uncurry defined as:

> um curry :: (a -> b -> c) -> (a,b) -> c

> uncourry f (x,y) = f x y

FUNCTIONAL PROGRAMMING - 2012

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Question 4
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```
(a) > type Event = String
     > type Country = String
    > data Medal = Gold | Silver | Bronze deriving (Eg, Show)
     > type Winners= [(Event, Country, Medal)]
```

- (b) > countraedals :: Wimmers -> Medal -> Country -> Int > countredals winners medal country
 - > | null winners = 0
 - > | (medal == m) && (country == c) = 1 + countredals (tail winners) medal country
 - > 1 otherwise = countraedals (tail winners) medal country
 - > where (e, c, m) = head winners
- (c) > score :: Winners -> Country -> Int
 - > score winners country = 3 x countraedals winners Gold country +
 - 2 x countmedals winners Silver country +
 - countredals winners Bronze country
- (d) > namking :: Winners -> Country -> Int
 - > nanking winners country = 1+ length (filter better winners)

where better (e, c, m) = scone winners c > scone winners country

This function calculates the namking of a particular country by comparing its "score" with the "scores" of all the other countries from the list.

- > namk :: Wimmers -> [Country] -> [int]
- > namk wimmers []=[]
- > namk wimmers (C:CS) = namking wimmers C: namk wimmers cs
- (e) > ljustify :: Country -> int -> String
 - > ljustify country m = country ++ concat (take (m-lm) (nepeat " "))
 - where In=length country

We justify to the left the mame of the countries, which appear in the first column

- > njustify : String -> Int -> String
- > njustify number m = concat (take (m-ln) (nepeat " ")) ++ number > where lm = lungth number We justify to the night the data we obtain from the results (winners)

```
> table :: Winners -> [Country] -> String

> table winners [] = (ljustify "Country" 10) ++ " Gold Silver Bronze Ramk m"

> table winners (c: cs) =

> table winners cs ++ (ljustify c 10) ++ " " ++ (njustify (show gold) 4) ++ " "

++ (njustify (show silver) 6) ++ " " ++ (njustify (show bronze) 6) ++ " "

> ++ (njustify (show pos) 4) ++ " m"

> where gold = countraedals winners Gold c

> silver = countraedals winners Silver c

> bronze = countraedals winners Bronze c

> pos = nanking winners c

This function creates the string we want to print to obtain the "medaltable". We justify
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This function creates the string we want to print to obtain the "medaltable". We justify the countries to the left with 10 as we are told that is the maximum lungth of the mame of a country and the numbers according to their column. We observe that we use the list of countries in the reversed order for us to be able to write the headings, so we'll print the string with the list of countries revused:

> medaltable :: Winners -> [Country] -> 10()

> medaltable wimmers cs = put Str (table wimmers (neverse cs))