

FUNCTIONAL PROGRAMMING MT2018

Sheet 2

3.1 Suppose that the class of ordered types is declared by something like

```
class Eq a => Ord a where
  (<), (<=), (>), (>=) :: a -> a -> Bool
  x < y = not (x >= y)
  x > y = not (x <= y)
  x >= y = x == y || x > y
```

(It includes a couple of other things, and these are not quite the default definitions.) Lists are lexicographically ordered, like the words of a dictionary. Write an instance declaration for `Ord [a]`.

3.2 Suppose $h\ x\ y = f\ (g\ x\ y)$. Which of the following are true, which are false, and (in each case) why?

1. $h = f \cdot g$
2. $h\ x = f \cdot g\ x$
3. $h\ x\ y = (f \cdot g)\ x\ y$

3.3 Give most general types for the following, where possible

```
> subst f g x = (f x) (g x)
> fix f = f (fix f)
> twice f = f . f
> selfie f = f f
```

You should try to work out what the most general type is by hand, but you can check that you are right by using an interpreter; and if you are wrong, check that you understand why.

3.4 Which of these equations are badly typed? For the others, what can you say about the type of xs , and whether and when the equation holds?

$a) [] : xs = xs$	$e) [] : xs = [], xs$	$i) xs : [] = xs$
$b) xs : [] = [xs]$	$f) xs : xs = [xs, xs]$	$j) xs : [xs] = [xs, xs]$
$c) [[]] ++ xs = xs$	$g) [[]] ++ xs = [xs]$	$k) [[]] ++ xs = [], xs$
$d) [[]] ++ [xs] = [], xs$	$h) [xs] ++ [] = [xs]$	$l) [xs] ++ [xs] = [xs, xs]$

4.1 Show that if f and g are strict, so is the composition $f \cdot g$.

Is the converse true: that if $f \cdot g$ is strict, so must f and g be?

- 4.2 If we count $\perp :: \text{Bool}$ as well as the proper values, there are three values of type *Bool*. So how many functions are there of type $\text{Bool} \rightarrow \text{Bool}$? How many of these are computable? Are all the computable ones definable in Haskell?
- 4.3 By evaluating expressions like `False && undefined` and other expressions using combinations of *True*, *False*, and *undefined* find exactly which function `(&&)` is implemented in the standard prelude. Give a definition which would produce that behaviour.

There is exactly one computable function, `(&&&)` say, which satisfies

$$\begin{aligned} \text{False} \ \&\&\& \ y &= \text{False} \\ x \ \&\&\& \ \text{False} &= \text{False} \\ \text{True} \ \&\&\& \ \text{True} &= \text{True} \end{aligned}$$

for all x and y , including \perp . Given that it is computable, explain what this function does in other cases, and so why there is only one such function. Is it definable in Haskell?

- 4.4 Time (as Richard Bird might say) for a song:

```
One man went to mow
Went to mow a meadow
One man and his dog
Went to mow a meadow
```

```
Two men went to mow
Went to mow a meadow
Two men, one man and his dog
Went to mow a meadow
```

```
Three men went to mow
Went to mow a meadow
Three men, two men, one man and his dog
Went to mow a meadow
```

Write a Haskell function $\text{song} :: \text{Int} \rightarrow \text{String}$ so that $\text{song } n$ is the song when there are n men (and a dog). Assume $n \leq 10$. To print the song, type for example: `putStr (song 5)`. You may want to start from

```
> song 0 = ""
> song n = song (n-1) ++ "\n" ++ verse n
> verse n = line1 n ++ line ++ line3 n ++ line
```

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