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fw[i].assign(coord[i].size() + 1, 0);

```
void upd(T x, T y, T v) {
    for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin();
          xx < fw.size(); xx += xx & -xx) {
      for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y)
           - \operatorname{coord}[xx].\operatorname{begin}(); yy < \operatorname{fw}[xx].\operatorname{size}(); yy += yy & -yy) {
         fw[xx][yy] += v;
  T qry(T x, T y) {
    T ans = 0;
    for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin();
          xx > 0; xx -= xx & -xx) {
      for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y)
           - coord[xx].begin(); yy > 0; yy -= yy & -yy) {
        ans += fw[xx][vv]:
    return ans;
  T qry(T x1, T y1, T x2, T y2) {
    return qry(x^2, y^2) - qry(x^2, y^2 - 1) - qry(x^2 - 1, y^2) + qry(x^2 - 1)
        1, y1 - 1);
  void upd(T x1, T y1, T x2, T y2, T v) { // !insert these points
    upd(x1, y1, v);
    upd(x1, y2 + 1, -v);
    upd(x2 + 1, y1, -v);
    upd(x2 + 1, y2 + 1, v);
private:
  vector<T> ord:
  vector<vector<T>> fw, coord;
};
```

## 1.2 Iterative Segment Tree

```
int n, t[2 * ms];

void build() {
   for(int i = n - 1; i > 0; --i) t[i] = t[i<<1] + t[i<<1|1]; // Merge
}

void update(int p, int value) { // set value at position p
   for(t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1]; //
        Merge
}

int query(int l, int r) {
   int res = 0;
   for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1) {
      if(1&1) res += t[1++]; // Merge
      if(r&1) res += t[--r]; // Merge
   }
   return res;
```

```
// If is non-commutative
S query(int 1, int r) {
S resl, resr;
for (1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1) {
if (1&1) resl = combine(resl, t[1++]);
if (r&1) resr = combine(t[--r], resr);
}
return combine(resl, resr);
}
```

## 1.3 Iterative Segment Tree with Lazy Propagation

```
struct LazyContext {
  LazyContext() { }
  void reset() { }
  void operator += (LazyContext o) { }
  // atributes
};
struct Node {
  Node() {
    // neutral element
  Node() {
    // init
  Node (Node 1, Node r) {
    // merge
  bool canBreak(LazyContext lazy) {
    // return true if can break without applying lazy
  bool canApply(LazyContext lazy) {
    // returns true if can apply lazy
  void apply(LazyContext &lazy) {
    // changes lazy if needed
  // atributes
};
template <class i_t, class e_t, class lazy_cont>
class SegmentTree {
public:
  void init(std::vector<e_t> base) {
    n = base.size();
    \mathbf{h} = 0;
    while((1 << h) < n) h++;
    tree.resize(2 * n);
    dirty.assign(n, false);
    lazy.resize(n);
    for (int i = 0; i < n; i++) {
      tree[i + n] = i_t(base[i]);
    for (int i = n - 1; i > 0; i--) {
      tree[i] = i_t(tree[i + i], tree[i + i + 1]);
      lazv[i].reset();
```

```
i_t qry(int 1, int r) {
    if(l >= r) return i_t();
    1 += n, r += n;
    push(1);
    push(r - 1);
    i_t lp, rp;
    for (; 1 < r; 1 /= 2, r /= 2) {
      if(l & 1) lp = i t(lp, tree[l++]);
      if(r & 1) rp = i_t(tree[--r], rp);
    return i_t(lp, rp);
  void upd(int 1, int r, lazy_cont lc) {
    if(1 >= r) return;
    1 += n, r += n;
    push(1);
    push(r - 1):
    int 10 = 1, r0 = r;
    for (; 1 < r; 1 /= 2, r /= 2) {
     if(1 & 1) downUpd(l++, lc);
      if(r & 1) downUpd(--r, lc);
    build(10);
    build(r0 - 1);
  void upd(int pos, e_t v) {
    pos += n;
    push (pos);
    tree[pos] = i_t(v);
    build(pos);
private:
  int n, h;
  std::vector<bool> dirty:
  std::vector<i_t> tree;
  std::vector<lazy_cont> lazy;
  void apply(int p, lazy_cont lc) {
    tree[p].apply(lc);
    if(p < n) {
      dirty[p] = true;
      lazy[p] += lc;
  void pushSingle(int p) {
    if(dirty[p]) {
      downUpd(p + p, lazy[p]);
      downUpd(p + p + 1, lazy[p]);
      lazy[p].reset();
      dirty[p] = false;
  void push(int p) {
    for(int s = h; s > 0; s--) {
      pushSingle(p >> s);
```

```
void downUpd(int p, lazy_cont lc) {
    if(tree[p].canBreak(lc)) {
      return:
    } else if(tree[p].canApply(lc)) {
      apply(p, lc);
    } else {
      pushSingle(p);
      downUpd(p + p, lc);
      downUpd(p + p + 1, lc);
      tree[p] = i_t(tree[p + p], tree[p + p + 1]);
  void build(int p) {
    for (p /= 2; p > 0; p /= 2) {
      tree[p] = i_t(tree[p + p], tree[p + p + 1]);
      if(dirtv[p]) {
        tree[p].apply(lazy[p]);
};
```

#### 1.4 Segment Tree with Lazy Propagation

```
int arr[ms], seq[4 * ms], lazy[4 * ms], n;
void build(int idx = 0, int l = 0, int r = n-1) {
  int mid = (1+r)/2;
 lazy[idx] = 0;
 if(1 == r) {
    seq[idx] = arr[l];
    return;
  build(2*idx+1, 1, mid); build(2*idx+2, mid+1, r);
  seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
void apply(int idx, int 1, int r) {
  if(lazy[idx] && !canBreak) { // if not beats canBreak = false
   if(1 < r) {
     lazy[2*idx+1] += lazy[idx]; // Merge de lazy
      lazv[2*idx+2] += lazv[idx]; // Merge de lazv
    if(canApply) { // if not beats canApply = true
      seg[idx] += lazy[idx] * (r - 1 + 1); // Aplicar lazy no seg
      apply (2*idx+1, 1, mid); apply (2*idx+2, mid+1, r);
      seq[idx] = seq[2*idx+1] + seq[2*idx+2]; // Merge
  lazy[idx] = 0; // Limpar a lazy
int query(int L, int R, int idx = 0, int 1 = 0, int r = n-1) {
  int mid = (1+r)/2;
  apply(idx, l, r);
  if(l > R | | r < L) return 0; // Valor que nao atrapalhe</pre>
```

```
if(L <= 1 && r <= R) return seg[idx];
return query(L, R, 2*idx+1, 1, mid) + query(L, R, 2*idx+2, mid+1, r)
    ; // Merge
}

void update(int L, int R, int V, int idx = 0, int 1 = 0, int r = n-1)
    {
    int mid = (1+r)/2;
    apply(idx, 1, r);
    if(1 > R || r < L) return;
    if(L <= 1 && r <= R) {
        lazy[idx] = V;
        apply(idx, 1, r);
        return;
    }
    update(L, R, V, 2*idx+1, 1, mid); update(L, R, V, 2*idx+2, mid+1, r)
    ;
    seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
}</pre>
```

## 1.5 Treap

```
mt19937 rng ((int) chrono::steady_clock::now().time_since_epoch().
    count());
typedef int Value:
typedef struct item * pitem;
struct item {
 item () {}
  item (Value v) { // add key if not implicit
   prio = uniform_int_distribution<int>() (rng);
   cnt = 1:
   rev = 0;
   1 = r = 0:
  pitem 1, r;
 Value value;
  int prio, cnt;
 bool rev;
};
int cnt (pitem it) {
  return it ? it->cnt : 0;
void fix (pitem it) {
 if (it)
    it->cnt = cnt(it->1) + cnt(it->r) + 1;
void pushLazv (pitem it) {
  if (it && it->rev) {
   it->rev = false;
   swap(it->1, it->r);
   if (it->1) it->1->rev ^= true;
   if (it->r) it->r->rev ^= true;
```

```
void merge (pitem & t, pitem 1, pitem r) {
  pushLazy (1); pushLazy (r);
  if (!1 || !r) t = 1 ? 1 : r;
  else if (l->prio > r->prio)
   merge (1->r, 1->r, r), t = 1;
   merge (r->1, 1, r->1), t = r;
  fix (t);
void split (pitem t, pitem & l, pitem & r, int key) {
  if (!t) return void( l = r = 0 );
  pushLazy (t);
  int cur_key = cnt(t->1); // t->key if not implicit
  if (kev <= cur kev)</pre>
   split (t->1, 1, t->1, key), r = t;
    split (t->r, t->r, r, key - (1 + cnt(t->1))), 1 = t;
  fix (t);
void reverse (pitem t, int l, int r) {
 pitem t1, t2, t3;
  split (t, t1, t2, 1);
  split (t2, t2, t3, r-l+1);
  t2->rev ^= true;
 merge (t, t1, t2);
 merge (t, t, t3);
void unite (pitem & t, pitem l, pitem r) {
 if (!1 || !r) return void ( t = 1 ? 1 : r );
  if (1->prio < r->prio) swap (1, r);
 pitem lt, rt;
 split (r, lt, rt, l->key);
 unite (1->1, 1->1, 1t);
 unite (1-> r, 1->r, rt);
  t = 1;
```

## 1.6 Persistent Treap

```
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count())
;

typedef int Key;
struct Treap {
   Treap(){}
   Treap(char k) {
      key = 1;
      size = 1;
      l = r = NULL;
      val = k;
   }

   Treap *1, *r;
   Key key;
   char val;
   int size;
```

```
};
typedef Treap * PTreap;
bool leftSide(PTreap 1, PTreap r) {
  return (int) (rng() % (1->size + r->size)) < 1->size;
void fix(PTreap t) {
  if (t == NULL) {
    return;
  t \rightarrow size = 1:
  t->key = 1;
  if (t->1) {
   t\rightarrow size += t\rightarrow l\rightarrow size:
   t->key += t->l->size;
  if (t->r) {
    t->size += t->r->size;
void split(PTreap t, Key key, PTreap &1, PTreap &r) {
  if (t == NULL) {
    1 = r = NULL;
  } else if (t->key <= key) {</pre>
    1 = new Treap();
    *1 = *t;
    split(t->r, key - t->key, l->r, r);
    fix(1);
  } else {
    r = new Treap();
    *r = *t;
    split(t->1, key, l, r->l);
    fix(r);
void merge(PTreap &t, PTreap 1, PTreap r) {
  if (!l || !r) {
    t = 1 ? 1 : r;
    return;
  t = new Treap();
  if (leftSide(l, r)) {
    *t = *1;
    merge(t->r, 1->r, r);
  } else {
    *t = *r;
    merge(t->1, 1, r->1);
  fix(t);
vector<PTreap> ver = {NULL};
PTreap build(int 1, int r, string& s) {
 if (1 >= r) return NULL;
  int mid = (1 + r) >> 1;
  auto ans = new Treap(s[mid]);
```

```
ans->1 = build(1, mid, s);
  ans->r = build(mid + 1, r, s);
  fix(ans);
  return ans;
int last = 0;
void go(PTreap t, int f) {
  if (!t) return;
  qo(t->1, f);
  cout << t->val;
  last += (t->val <math>== 'c') * f;
  qo(t->r, f);
void insert(PTreap t, int pos, string& s) {
 PTreap 1, r;
  split(t, pos + 1, l, r);
 PTreap mid = build(0, s.size(), s);
 merge(mid, 1, mid);
 merge(mid, mid, r);
  ver.push_back(mid);
void erase(PTreap t, int L, int R) {
 PTreap 1, mid, r;
  split(t, L, l, mid);
  split(mid, R - L + 1, mid, r);
 merge(1, 1, r);
  ver.push_back(1);
```

#### 1.7 KD-Tree

```
int d:
long long getValue(const PT &a) {return (d & 1) == 0 ? a.x : a.y; }
bool comp (const PT &a, const PT &b) {
  if((d & 1) == 0) { return a.x < b.x; }</pre>
  else { return a.v < b.v; }</pre>
long long sqrDist(PT a, PT b) { return (a - b) * (a - b); }
class KD_Tree {
public:
  struct Node {
   PT point;
    Node *left, *right;
  void init(std::vector<PT> pts) {
    if(pts.size() == 0) {
      return;
    int n = 0;
    tree.resize(2 * pts.size());
    build(pts.begin(), pts.end(), n);
    //assert(n <= (int) tree.size());</pre>
```

```
long long nearestNeighbor(PT point) {
    // assert(tree.size() > 0);
    long long ans = (long long) le18;
    nearestNeighbor(&tree[0], point, 0, ans);
    return ans;
private:
  std::vector<Node> tree;
  Node* build(std::vector<PT>::iterator 1, std::vector<PT>::iterator r
      , int &n, int h = 0) {
    int id = n++;
    if(r - 1 == 1) {
      tree[id].left = tree[id].right = NULL;
      tree[id].point = *1;
    \} else if (r - 1 > 1) {
      std::vectorPT>::iterator mid = 1 + ((r - 1) / 2);
      std::nth_element(l, mid - 1, r, comp);
      tree[id].point = \star (mid - 1);
      // BE CAREFUL!
      // DO EVERYTHING BEFORE BUILDING THE LOWER PART!
      tree[id].left = build(l, mid, n, h^1);
      tree[id].right = build(mid, r, n, h^1);
    return &tree[id];
  void nearestNeighbor(Node* node, PT point, int h, long long &ans) {
    if(!node) {
      return;
    if(point != node->point) {
      // THIS WAS FOR A PROBLEM
      // THAT YOU DON'T CONSIDER THE DISTANCE TO ITSELF!
      ans = std::min(ans, sqrDist(point, node->point));
    d = h;
    long long delta = getValue(point) - getValue(node->point);
    if(delta <= 0) {
      nearestNeighbor(node->left, point, h^1, ans);
      if(ans > delta * delta) {
        nearestNeighbor(node->right, point, h^1, ans);
    } else {
      nearestNeighbor(node->right, point, h^1, ans);
      if(ans > delta * delta) {
        nearestNeighbor(node->left, point, h^1, ans);
};
```

#### 1.8 Link Cut Tree

```
* Description: Represents a forest of unrooted trees. You can add and
      remove
 * edges (as long as the result is still a forest), and check whether
 * two nodes are in the same tree.
 * Time: All operations take amortized O(\log N).
 * Status: Fuzz-tested a bit for N <= 20
#pragma once
struct Node { // Splay tree. Root's pp contains tree's parent.
 Node *p = 0, *pp = 0, *c[2];
  bool flip = 0;
  Node() { c[0] = c[1] = 0; fix(); }
 void fix() {
   if (c[0]) c[0]->p = this;
    if (c[1]) c[1]->p = this;
    // (+ update sum of subtree elements etc. if wanted)
  void push flip() {
    if (!flip) return;
    flip = 0; swap(c[0], c[1]);
    if (c[0]) c[0]->flip ^= 1;
    if (c[1]) c[1]->flip ^= 1;
  int up() { return p ? p \rightarrow c[1] == this : -1; }
  void rot(int i, int b) {
    int h = i \hat{b};
    Node *x = c[i], *y = b == 2 ? x : x -> c[h], *z = b ? y : x;
    if ((y->p = p)) p->c[up()] = y;
    c[i] = z -> c[i ^ 1];
    if (b < 2) {
      x - c[h] = y - c[h^1];
      z - c[h ^1] = b ? x : this;
    y - > c[i ^1] = b ? this : x;
    fix(); x->fix(); y->fix();
    if (p) p->fix();
    swap (pp, y->pp);
  void splay() { /// Splay this up to the root. Always finishes
      without flip set.
    for (push_flip(); p; ) {
     if (p->p) p->p->push_flip();
      p->push_flip(); push_flip();
      int c1 = up(), c2 = p->up();
      if (c2 == -1) p->rot(c1, 2);
      else p->p->rot(c2, c1 != c2);
  Node* first() { /// Return the min element of the subtree rooted at
      this, splayed to the top.
    push_flip();
    return c[0] ? c[0]->first() : (splay(), this);
};
struct LinkCut {
 vector<Node> node;
 LinkCut(int N) : node(N) {}
  void link(int u, int v) { // add an edge (u, v)
```

```
assert(!connected(u, v));
    make root(&node[u]);
    node[u].pp = &node[v];
  void cut(int u, int v) { // remove an edge (u, v)
    Node *x = &node[u], *top = &node[v];
    make_root(top); x->splay();
    assert(top == (x->pp ?: x->c[0]));
    if (x->pp) x->pp = 0;
    else {
      x->c[0] = top->p = 0;
      x->fix();
  bool connected(int u, int v) { // are u, v in the same tree?
    Node* nu = access(&node[u]) ->first();
    return nu == access(&node[v])->first();
  void make_root(Node* u) { /// Move u to root of represented tree.
    access(u);
    u->splay();
    if(u->c[0]) {
      u - c[0] - p = 0;
      u \rightarrow c[0] \rightarrow flip ^= 1;
      u - c[0] - pp = u;
      u -> c[0] = 0;
      u->fix();
  Node* access(Node* u) { /// Move u to root aux tree. Return the root
       of the root aux tree.
    u->splay();
    while (Node* pp = u->pp) {
      pp \rightarrow splay(); u \rightarrow pp = 0;
      if (pp->c[1]) {
        pp->c[1]->p = 0; pp->c[1]->pp = pp; }
      pp - c[1] = u; pp - fix(); u = pp;
    return u;
};
```

## 1.9 Sparse Table

```
vector<vector<int>> table;
vector<int> lg2;

void build(int n, vector<int> v) {
    lg2.resize(n + 1);
    lg2[1] = 0;
    for (int i = 2; i <= n; i++) {
        lg2[i] = lg2[i >> 1] + 1;
    }
    table.resize(lg2[n] + 1);
    for (int i = 0; i < lg2[n] + 1; i++) {
        table[i].resize(n + 1);
    }
    for (int i = 0; i < n; i++) {
        table[0][i] = v[i];
    }
}</pre>
```

```
for (int i = 0; i < lg2[n]; i++) {
    for (int j = 0; j < n; j++) {
        if (j + (1 << i) >= n) break;
        table[i + 1][j] = min(table[i][j], table[i][j + (1 << i)]);
     }
}
int get(int l, int r) {
    int k = lg2[r - 1 + 1];
    return min(table[k][l], table[k][r - (1 << k) + 1]);
}</pre>
```

## 1.10 Max Queue

```
// src: tfq50
template <class T, class C = std::less<T>>
struct MaxQueue {
 MaxOueue() {
    clear();
  void clear() {
    id = 0:
    q.clear();
  void push(T x) {
    std::pair<int, T> nxt(1, x);
    while(q.size() > id && cmp(q.back().second, x)) {
      nxt.first += q.back().first;
      q.pop_back();
    q.push_back(nxt);
  T qry() {
    return q[id].second;
  void pop() {
    q[id].first--;
    if(q[id].first == 0) {
      id++;
private:
  std::vector<std::pair<int, T>> q;
  int id:
  C cmp;
};
```

## 1.11 Policy Based Structures

```
using namespace __gnu_pbds;

typedef tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;

ordered_set X;
X.insert(1);
X.find_by_order(0);
X.order_of_key(-5);
end(X), begin(X);
```

## 1.12 Color Updates Structure

```
struct range {
  int 1, r;
  int v;
  range(int l = 0, int r = 0, int v = 0) : l(l), r(r), v(v) {}
 bool operator < (const range &a) const {</pre>
    return 1 < a.1;</pre>
};
set<range> ranges;
vector<range> update(int 1, int r, int v) { // [1, r)
  vector<range> ans:
  if(1 >= r) return ans;
  auto it = ranges.lower_bound(1);
  if(it != ranges.begin()) {
    it--;
    if(it->r>1) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, 1, cur.v));
      ranges.insert(range(l, cur.r, cur.v));
  it = ranges.lower bound(r);
  if(it != ranges.begin()) {
    it--;
    if(it->r>r) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, r, cur.v));
      ranges.insert(range(r, cur.r, cur.v));
  for(it = ranges.lower_bound(1); it != ranges.end() && it->1 < r; it</pre>
      ++) {
    ans.push_back(*it);
  ranges.erase(ranges.lower_bound(1), ranges.lower_bound(r));
  ranges.insert(range(l, r, v));
  return ans;
int query(int v) { // Substituir -1 por flag para quando nao houver
    resposta
```

```
auto it = ranges.upper_bound(v);
if(it == ranges.begin()) {
    return -1;
}
it--;
return it->r >= v ? it->v : -1;
}
```

# 2 Graph Algorithms

## 2.1 Simple Disjoint Set

```
struct dsu {
 vector<int> hist, par, sz;
 vector<ii> changes;
 int n;
 dsu (int n) : n(n) {
   hist.assign(n, 1e9);
   par.resize(n);
   iota(par.begin(), par.end(), 0);
   sz.assign(n, 1);
 int root (int x, int t) {
   if(hist[x] > t) return x;
   return root(par[x], t);
 void join (int a, int b, int t) {
   a = root(a, t);
   b = root(b, t);
   if (a == b) { changes.emplace_back(-1, -1); return; }
   if (sz[a] > sz[b]) swap(a, b);
   par[a] = b;
   sz[b] += sz[a];
   hist[a] = t;
   changes.emplace_back(a, b);
 bool same (int a, int b, int t) {
    return root(a, t) == root(b, t);
 void undo () {
   int a, b;
   tie(a, b) = changes.back();
   changes.pop_back();
   if (a == -1) return;
   sz[b] = sz[a];
   par[a] = a;
   hist[a] = 1e9;
   n++;
 int when (int a, int b) {
   while (1) {
     if (hist[a] > hist[b]) swap(a, b);
```

```
if (par[a] == b) return hist[a];
  if (hist[a] == 1e9) return 1e9;
    a = par[a];
  }
};
```

#### 2.2 Boruvka

```
struct edge {
  int u, v;
 int w;
 int id;
 edge () {};
 edge (int u, int v, int w = 0, int id = 0) : u(u), v(v), w(w), id(id)
     ) {};
 bool operator < (edge &other) const { return w < other.w; };</pre>
vector<edge> boruvka (vector<edge> &edges, int n) {
 vector<edge> mst;
 vector<edge> best(n);
 initDSU(n):
 bool f = 1;
 while (f) {
   f = 0:
   for (int i = 0; i < n; i++) best[i] = edge(i, i, inf);</pre>
   for (auto e : edges) {
      int pu = root(e.u), pv = root(e.v);
      if (pu == pv) continue;
      if (e < best[pu]) best[pu] = e;</pre>
      if (e < best[pv]) best[pv] = e;</pre>
    for (int i = 0; i < n; i++) {
      edge e = best[root(i)];
      if (e.w != inf) {
        join(e.u, e.v);
        mst.push_back(e);
        f = 1;
  return mst;
```

## 2.3 Dinic Max Flow

```
const int ms = 1e3; // Quantidade maxima de vertices
const int me = 1e5; // Quantidade maxima de arestas

int adj[ms], to[me], ant[me], wt[me], z, n;
int copy_adj[ms], fila[ms], level[ms];

void clear() { // Lembrar de chamar no main
   memset(adj, -1, sizeof adj);
   z = 0;
}
```

```
void add(int u, int v, int k) {
  to[z] = v;
  ant[z] = adj[u];
  wt[z] = k;
  adj[u] = z++;
  swap(u, v);
  to[z] = v;
  ant[z] = adj[u];
 wt[z] = 0; // Lembrar de colocar = 0
  adj[u] = z++;
int bfs(int source, int sink) {
 memset(level, -1, sizeof level);
  level[source] = 0;
  int front = 0, size = 0, v;
  fila[size++] = source;
 while(front < size) {</pre>
   v = fila[front++];
    for(int i = adj[v]; i != -1; i = ant[i]) {
      if(wt[i] && level[to[i]] == -1) {
       level[to[i]] = level[v] + 1;
       fila[size++] = to[i];
    }
  return level[sink] != -1;
int dfs(int v, int sink, int flow) {
 if(v == sink) return flow;
  int f;
  for(int &i = copy_adj[v]; i != -1; i = ant[i]) {
    if(wt[i] && level[to[i]] == level[v] + 1 &&
      (f = dfs(to[i], sink, min(flow, wt[i])))) {
      wt[i] -= f;
      wt[i ^ 1] += f;
      return f;
  return 0;
int maxflow(int source, int sink) {
  int ret = 0, flow;
 while(bfs(source, sink)) {
    memcpy(copy_adj, adj, sizeof adj);
    while((flow = dfs(source, sink, 1 << 30))) {</pre>
      ret += flow;
    }
  return ret;
```

## 2.4 Minimum Vertex Cover

```
// + Dinic
vector<int> coverU, U, coverV, V; // ITA - Parti o U LEFT,
    parti o V RIGHT, 0 indexed
bool Zu[mx], Zv[mx];
```

```
int pairU[mx], pairV[mx];
void getreach(int u) {
  if (u == -1 \mid | Zu[u]) return;
  Zu[u] = true;
  for (int i = adj[u]; ~i; i = ant[i]) {
    int v = to[i];
    if (v == SOURCE || v == pairU[u]) continue;
    Zv[v] = true;
    getreach(pairV[v]);
void minimumcover () {
 memset(pairU, -1, sizeof pairU);
 memset(pairV, -1, sizeof pairV);
  for (auto i : U) {
    for (int j = adj[i]; ~j; j = ant[j]) {
      if (!(j&1) && !wt[j]) {
        pairU[i] = to[j], pairV[to[j]] = i;
  memset (Zu, 0, sizeof Zu);
  memset (Zv, 0, sizeof Zv);
  for (auto u : U) {
   if (pairU[u] == -1) getreach(u);
  coverU.clear(), coverV.clear();
  for (auto u : U) {
    if (!Zu[u]) coverU.push_back(u);
  for (auto v : V) {
    if (Zv[v]) coverV.push_back(v);
```

#### 2.5 Min Cost Max Flow

```
template <class T = int>
class MCMF {
public:
  struct Edge {
    Edge(int a, T b, T c) : to(a), cap(b), cost(c) {}
   int to;
    T cap, cost;
  };
  MCMF(int size) {
    n = size:
    edges.resize(n);
    pot.assign(n, 0);
    dist.resize(n);
    visit.assign(n, false);
  std::pair<T, T> mcmf(int src, int sink) {
    std::pair<T, T> ans(0, 0);
    if(!SPFA(src, sink)) return ans;
    fixPot();
    // can use dijkstra to speed up depending on the graph
```

```
while(SPFA(src, sink)) {
      auto flow = augment(src, sink);
      ans.first += flow.first;
      ans.second += flow.first * flow.second;
      fixPot();
    return ans;
  void addEdge(int from, int to, T cap, T cost) {
    edges[from].push_back(list.size());
    list.push_back(Edge(to, cap, cost));
    edges[to].push_back(list.size());
    list.push_back(Edge(from, 0, -cost));
private:
  int n:
  std::vector<std::vector<int>> edges;
  std::vector<Edge> list:
  std::vector<int> from;
  std::vector<T> dist, pot;
  std::vector<bool> visit;
  /*bool dij(int src, int sink) {
    T INF = std::numeric_limits<T>::max();
    dist.assign(n, INF);
    from.assign(n, -1);
    visit.assign(n, false);
    dist[src] = 0;
    for(int i = 0; i < n; i++) {
      int best = -1;
      for (int j = 0; j < n; j++) {
       if(visit[j]) continue;
        if(best == -1 \mid | dist[best] > dist[j]) best = j;
      if(dist[best] >= INF) break;
      visit[best] = true;
      for(auto e : edges[best]) {
        auto ed = list[e];
        if (ed.cap == 0) continue;
        T toDist = dist[best] + ed.cost + pot[best] - pot[ed.to];
        assert(toDist >= dist[best]);
        if(toDist < dist[ed.to]) {</pre>
          dist[ed.to] = toDist;
          from[ed.to] = e;
    return dist[sink] < INF;
  1 */
  std::pair<T, T> augment(int src, int sink) {
    std::pair<T, T> flow = {list[from[sink]].cap, 0};
    for(int v = sink; v != src; v = list[from[v]^1].to) {
      flow.first = std::min(flow.first, list[from[v]].cap);
      flow.second += list[from[v]].cost;
    for(int v = sink; v != src; v = list[from[v]^1].to) {
     list[from[v]].cap -= flow.first;
      list[from[v]^1].cap += flow.first;
```

```
return flow;
  std::queue<int> q;
 bool SPFA(int src, int sink) {
   T INF = std::numeric limits<T>::max();
   dist.assign(n, INF);
   from.assign(n, -1);
   q.push(src);
   dist[src] = 0;
   while(!q.empty()) {
      int on = q.front();
      q.pop();
      visit[on] = false;
      for(auto e : edges[on]) {
        auto ed = list[e];
        if(ed.cap == 0) continue;
        T toDist = dist[on] + ed.cost + pot[on] - pot[ed.to];
        if(toDist < dist[ed.to])</pre>
          dist[ed.to] = toDist;
          from[ed.to] = e;
          if(!visit[ed.to]) {
            visit[ed.to] = true;
            q.push(ed.to);
    return dist[sink] < INF;</pre>
 void fixPot() {
   T INF = std::numeric_limits<T>::max();
   for(int i = 0; i < n; i++) {
      if(dist[i] < INF) pot[i] += dist[i];</pre>
};
```

#### 2.6 Euler Path and Circuit

```
int pathV[me], szV, del[me], pathE, szE;
int adj[ms], to[me], ant[me], wt[me], z, n;

// Funcao de add e clear no dinic

void eulerPath(int u) {
  for(int i = adj[u]; ~i; i = ant[u]) if(!del[i]) {
    del[i] = del[i^1] = 1;
    eulerPath(to[i]);
    pathE[szE++] = i;
  }
  pathV[szV++] = u;
}
```

## 2.7 Articulation Points/Bridges/Biconnected Components

```
int adj[ms], to[me], ant[me], z;
```

```
int num[ms], low[ms], timer;
int art[ms], bridge[me], rch;
int bc[ms], nbc;
stack<int> st;
bool f[me];
void clear() { // Lembrar de chamar no main
  memset(adj, -1, sizeof adj);
void add(int u, int v) {
  to[z] = v;
  ant[z] = adj[u];
 adj[u] = z++;
void generateBc (int v) {
  while (!st.empty()) {
    int u = st.top();
    st.pop();
   bc[u] = nbc;
    if (v == u) break;
  ++nbc;
void dfs (int v, int p) {
  st.push(v);
  low[v] = num[v] = ++timer;
  for (int i = adj[v]; i != -1; i = ant[i]) {
    if (f[i] || f[i^1]) continue;
    f[i] = 1;
    int u = to[i];
    if (num[u] == -1) {
      dfs(u, v):
      if (low[u] > num[v]) bridge[i] = bridge[i^1] = 1;
      art[v] |= p != -1 \&\& low[u] >= num[v];
      if (p == -1 \&\& rch > 1) art[v] = 1;
      else rch ++;
      low[v] = min(low[v], low[u]);
    } else {
      low[v] = min(low[v], num[u]);
  if (low[v] == num[v]) generateBc(v);
void biCon (int n)
  nbc = 0, timer = 0;
  memset(num, -1, sizeof num);
  memset (bc, -1, sizeof bc);
  memset (bridge, 0, sizeof bridge);
  memset(art, 0, sizeof art);
  memset(f, 0, sizeof f);
  for (int i = 0; i < n; i++) {
    if (num[i] == -1) {
      rch = 0;
      dfs(i, 0);
```

## 2.8 SCC - Strongly Connected Components / 2SAT

```
vector<int> q[ms];
int idx[ms], low[ms], z, comp[ms], ncomp;
stack<int> st;
int dfs(int u) {
  if(~idx[u]) return idx[u] ? idx[u] : z;
  low[u] = idx[u] = z++;
  st.push(u):
  for(int v : q[u]) {
   low[u] = min(low[u], dfs(v));
  if(low[u] == idx[u]) {
    while(st.top() != u) {
      int v = st.top();
      idx[v] = 0;
      low[v] = low[u];
      comp[v] = ncomp;
      st.pop();
    idx[st.top()] = 0;
    st.pop();
    comp[u] = ncomp++;
  return low[u];
bool solveSat() {
 memset(idx, -1, sizeof idx);
  z = 1; ncomp = 0;
  for (int i = 0; i < n; i++) dfs(i);
  for(int i = 0; i < n; i++) if(comp[i] == comp[i^1]) return false;</pre>
  return true;
// Operacoes comuns de 2-sat
// v = "nao v"
#define trad(v) (v<0?((~v)*2)^1:v*2)
void addImp(int a, int b) { g[trad(a)].push(trad(b)); }
void addOr(int a, int b) { addImp(~a, b); addImp(~b, a); }
void addEqual(int a, int b) { addOr(a, b); addOr(a, b); }
void addDiff(int a, int b) { addEqual(a, ~b); }
// valoracao: value[v] = comp[trad(v)] < comp[trad(~v)]</pre>
```

#### 2.9 LCA - Lowest Common Ancestor

```
int par[ms][mlg+1], lvl[ms];
vector<int> g[ms];

void dfs(int v, int p, int 1 = 0) { // chamar como dfs(root, root)
  lvl[v] = 1;
  par[v][0] = p;
  for(int k = 1; k <= mlg; k++) {
    par[v][k] = par[par[v][k-1]][k-1];</pre>
```

```
for(int u : g[v]) {
    if(u != p) dfs(u, v, l + 1);
}

int lca(int a, int b) {
    if(lvl[b] > lvl[a]) swap(a, b);
    for(int i = mlg; i >= 0; i--) {
        if(lvl[a] - (1 << i) >= lvl[b]) a = par[a][i];
    }
    if(a == b) return a;
    for(int i = mlg; i >= 0; i--) {
        if(par[a][i] != par[b][i]) a = par[a][i], b = par[b][i];
    }
    return par[a][0];
}
```

## 2.10 Heavy Light Decomposition

```
// src: tfq
class HLD {
public:
  void init(int n) {
   // this doesn't delete edges!
    sz.resize(n):
    in.resize(n);
    out.resize(n):
    rin.resize(n);
    p.resize(n);
    edges.resize(n);
    nxt.resize(n);
    h.resize(n);
  void addEdge(int u, int v) {
    edges[u].push_back(v);
    edges[v].push_back(u);
  void setRoot(int n) {
   t = 0;
    p[n] = n;
    h[n] = 0;
    prep(n, n);
    nxt[n] = n;
    hld(n);
  int getLCA(int u, int v) {
    while(!inSubtree(nxt[u], v)) {
      u = p[nxt[u]];
    while(!inSubtree(nxt[v], u)) {
      v = p[nxt[v]];
    return in[u] < in[v] ? u : v;</pre>
  bool inSubtree(int u, int v) {
```

```
// is v in the subtree of u?
    return in[u] <= in[v] && in[v] < out[u];</pre>
  vector<pair<int, int>> getPathtoAncestor(int u, int anc) {
    // returns ranges [1, r) that the path has
    vector<pair<int, int>> ans;
    //assert(inSubtree(anc, u));
    while(nxt[u] != nxt[anc]) {
      ans.emplace_back(in[nxt[u]], in[u] + 1);
      u = p[nxt[u]];
    // this includes the ancestor!
    ans.emplace_back(in[anc], in[u] + 1);
    return ans;
private:
  vector<int> in, out, p, rin, sz, nxt, h;
  vector<vector<int>> edges:
  int t;
  void prep(int on, int par) {
    sz[on] = 1;
    p[on] = par;
    for(int i = 0; i < (int) edges[on].size(); i++) {</pre>
      int &u = edges[on][i];
      if(u == par) {
        swap(u, edges[on].back());
        edges[on].pop_back();
        i--;
      } else {
        h[u] = 1 + h[on];
        prep(u, on);
        sz[on] += sz[u];
        if(sz[u] > sz[edges[on][0]]) {
          swap(edges[on][0], u);
  void hld(int on) {
    in[on] = t++;
    rin[in[on]] = on;
    for(auto u : edges[on]) {
      nxt[u] = (u == edges[on][0] ? nxt[on] : u);
     hld(u);
    out[on] = t;
};
```

## 2.11 Centroid Decomposition

```
//Centroid decomposition1
void dfsSize(int v, int pa) {
   sz[v] = 1;
   for(int u : adj[v]) {
      if (u == pa || rem[u]) continue;
}
```

```
dfsSize(u, v);
    sz[v] += sz[u];
int getCentroid(int v, int pa, int tam) {
  for(int u : adj[v]) {
   if (u == pa || rem[u]) continue;
   if (2 * sz[u] > tam) return getCentroid(u, v, tam);
  return v;
void decompose (int v, int pa = -1) {
  //cout << v << ' ' << pa << '\n';
  dfsSize(v, pa);
 int c = getCentroid(v, pa, sz[v]);
 //cout << c << '\n';
 par[c] = pa:
 rem[c] = 1;
  for(int u : adj[c]) {
   if (!rem[u] && u != pa) decompose(u, c);
  adi[c].clear();
//Centroid decomposition2
void dfsSize(int v, int par) {
  sz[v] = 1;
  for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   dfsSize(u, v);
   sz[v] += sz[u];
int getCentroid(int v, int par, int tam) {
 for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   if (2 * sz[u] > tam) return getCentroid(u, v, tam);
  return v;
void setDis(int v, int par, int nv, int d) {
 dis[v][nv] = d;
  for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   setDis(u, v, nv, d + 1);
void decompose(int v, int par, int nv) {
  dfsSize(v, par);
 int c = getCentroid(v, par, sz[v]);
  ct[c] = par;
 removed[c] = 1;
  setDis(c, par, nv, 0);
  for(int u : adj[c]) {
   if (!removed[u]) {
```

```
decompose(u, c, nv + 1);
}
}
```

### 2.12 Sack

```
void dfs(int v, int par = -1, bool keep = 0) {
    int big = -1;
    for (int u : adj[v]) {
        if (u == par) continue;
        if (big == -1 || sz[u] > sz[big]) {
            biq = u;
    for (int u : adj[v]) {
        if (u == par || u == big) {
           continue;
        dfs(u, v, 0);
   if (big !=-1) {
        dfs(big, v, 1);
   for (int u : adj[v]) {
        if (u == par || u == big) {
           continue;
        put(u, v);
   if (!keep) {
```

## 2.13 Hungarian Algorithm - Maximum Cost Matching

```
//input: matrix n x m, n <= m</pre>
//return vector p of size n, where p[i] is the match for i
// and minimum cost
// time complexity: O(n^2 * m)
int u[ms], v[ms], p[ms], way[ms], minv[ms];
bool used[ms];
pair<vector<int>, int> solve(const vector<vector<int>> &matrix) {
  int n = matrix.size();
  if(n == 0) return {vector<int>(), 0};
  int m = matrix[0].size();
  assert(n <= m);</pre>
  memset(u, 0, (n+1)*sizeof(int));
  memset(v, 0, (m+1)*sizeof(int));
  memset(p, 0, (m+1)*sizeof(int));
  for(int i = 1; i <= n; i++) {</pre>
    memset(minv, 0x3f, (m+1)*sizeof(int));
    memset(way, 0, (m+1) *sizeof(int));
    for (int j = 0; j \le m; j++) used [j] = 0;
    p[0] = i;
```

```
int k0 = 0;
  do {
    used[k0] = 1;
    int i0 = p[k0], delta = inf, k1;
    for (int j = 1; j \le m; j++) {
      if(!used[i]) {
        int cur = matrix[i0-1][j-1] - u[i0] - v[j];
        if (cur < minv[j]) {</pre>
          minv[i] = cur;
          way[j] = k0;
        if(minv[j] < delta) {</pre>
          delta = minv[j];
          k1 = j;
    for(int j = 0; j \le m; j++) {
      if(used[i]) {
        u[p[j]] += delta;
        v[j] -= delta;
      } else {
        minv[j] -= delta;
    k0 = k1;
  } while(p[k0]);
    int k1 = way[k0];
    p[k0] = p[k1];
    k0 = k1;
 } while(k0);
vector<int> ans(n, -1);
for (int j = 1; j <= m; j++) {
 if(!p[j]) continue;
 ans[p[j] - 1] = j - 1;
return {ans, -v[0]};
```

#### 2.14 Burunduk

```
struct edge {
   int a, b, l, r;
};

typedef vector <edge> List;

int cnt[N + 1], ans[N], u[N], color[N], deg[N];
vi g[N];

void add (int a, int b) {
   g[a].pb(b), g[b].pb(a);
}

void dfs (int v, int value) {
   u[v] = 1, color[v] = value;
   forn(i, sz(g[v]))
      if (!u[g[v][i])
```

```
dfs(g[v][i], value);
int compress (List &v1, int vn, int &add_vn) {
  int vn1 = 0;
  forn (i, vn) u[i] = 0;
  forn (i, vn) {
   if (!u[i]) deg[vn1] = 0, dfs(i, vn1++);
  forn (i, sz(v1)) {
   v1[i].a = color[v1[i].a];
    v1[i].b = color[v1[i].b];
   if (v1[i].a != v1[i].b)
      deg[v1[i].a]++, deg[v1[i].b]++;
  vn = vn1, vn1 = 0;
  forn (i, vn) {
   u[i] = vn1, vn1 += (deg[i] > 0), add_vn += !deg[i];
  forn (i, sz(v1)) {
   v1[i].a = u[v1[i].a];
   v1[i].b = u[v1[i].b];
  return vn1;
void go (int 1, int r, const List &v, int vn, int add_vn) {
  if (cnt[l] == cnt[r]) return;
  if (!sz(v)){
    while (1 < r)
      ans[l++] = vn + add_vn;
    return:
  List v1:
  forn (i, vn) {
   g[i].clear();
  forn (i, sz(v)) {
    if (v[i].a != v[i].b) {
      if (v[i].l \le l \&\& v[i].r >= r)
        add(v[i].a, v[i].b);
      else if (1 < v[i].r \&\& r > v[i].1)
        v1.pb(v[i]);
  int vn1 = compress(v1, vn, add_vn);
  int m = (1 + r) / 2;
  go(1, m, v1, vn1, add vn);
  go(m, r, v1, vn1, add_vn);
```

#### 2.15 Minimum Arborescence

```
// uncommented O(V^2) arborescence
struct Edges {
   //set<pair<long long, int>> cost; O(Elog^2)
   long long cost[ms];

// possible optimization, use vector of size n
```

```
// instead of ms
  long long sum = 0;
  Edges() {
   memset(cost, 0x3f, sizeof cost);
  void addEdge(int u, long long c) {
    // cost.insert({c - sum, u}); O(Elog^2)
    cost[u] = min(cost[u], c - sum);
  pair<long long, int> getMin() {
    //return *cost.begin(); O(E*log^2)
    pair<long long, int> ans(cost[0], 0);
    // in this loop can change ms to n to make it faster for many
        cases
    for(int i = 1; i < ms; i++) {</pre>
      if(cost[i] < ans.first) {</pre>
        ans = pair<long long, int>(cost[i], i);
    return ans;
  void unite(Edges &e) {
    O(E*log^2E)
    if(e.cost.size() > cost.size()) {
      cost.swap(e.cost);
      swap(sum, e.sum);
    for(auto i : e.cost) {
      addEdge(i.second, i.first + e.sum);
    e.cost.clear():
    */
    // O(V^2)
    // can change ms to n
    for (int i = 0; i < ms; i++) {
      cost[i] = min(cost[i], e.cost[i] + e.sum - sum);
  }
};
typedef vector<vector<pair<long long, int>>> Graph;
Edges ed[ms];
int par[ms];
long long best[ms];
int col[ms];
int getPar(int x) { return par[x] < 0 ? x : par[x] = getPar(par[x]); }</pre>
void makeUnion(int a, int b) {
 a = getPar(a);
 b = getPar(b);
  if(a == b) return;
 ed[a].unite(ed[b]);
  par[b] = a;
```

```
long long arborescence(Graph edges) {
  // root is 0
 // edges has transposed adjacency list (cost, from)
 // edge from i to j cost c is
  // edge[i].emplace back(c, i)
  int n = (int) edges.size();
  long long ans = 0;
  for (int i = 0; i < n; i++) {
   ed[i] = Edges();
   par[i] = -1;
   for(auto j : edges[i]) {
     ed[i].addEdge(j.second, j.first);
   col[i] = 0;
  // to change the root you can simply change this next line to
  // col[root] = 2;
  col[0] = 2;
  for (int i = 0; i < n; i++) {
   if(col[getPar(i)] == 2) {
      continue;
   int on = getPar(i);
   vector<int> st;
   while(col[on] != 2) {
     assert(getPar(on) == on);
      if(col[on] == 1) {
        // found cycle
        int v = on;
        vector<int> cycle;
        //cout << "found cycle\n";</pre>
        while(st.back() != v) {
         //cout << st.back() << endl;</pre>
          cycle.push_back(st.back());
         st.pop_back();
        // compress cycle
        for(auto u : cycle) {
         makeUnion(v, u);
        v = getPar(v);
        col[v] = 0;
        on = v;
      } else {
        // still no cycle
        // while best is in compressed cycle, remove
        THIS IS TO MAKE O(E*log^2) ALGORITHM!!
        while (!ed[on].cost.empty() && getPar(on) == getPar(ed[on].
            getMin().second)) {
          ed[on].cost.erase(ed[on].cost.begin());
        // O(V^2)
        for (int x = 0; x < n; x++) {
         if(on == getPar(x)) {
            ed[on].cost[x] = 0x3f3f3f3f3f3f3f3f1LL;
```

```
// best edge
        auto e = ed[on].getMin();
        // O(E*log^2) assert(!ed[on].cost.empty()) if every vertex
            appears in the arborescence
        // O(V^2)
        assert(e.first < 0x3f3f3f3f3f3f3f3f3f1LL);</pre>
        int v = getPar(e.second);
        //cout << "found not cycle to " << v << " of cost " << e.first
             + ed[on].sum << '\n';
        assert (v != on);
        best[on] = e.first + ed[on].sum;
        ans += best[on];
        // compress edges
        ed[on].sum = -(e.first);
        st.push_back(on);
        col[on] = 1;
        on = v:
    // make everything 2
    for(auto u : st) {
      assert(getPar(u) == u);
      col[u] = 2;
  return ans;
int main() {
  cin.tie(NULL);
  ios_base::sync_with_stdio(NULL);
  // https://open.kattis.com/problems/fastestspeedrun
  int n;
  cin >> n;
  Graph edges(n+1);
  for(int i = 1; i <= n; i++) {</pre>
      int x, s;
      cin >> x >> s;
      edges[i].emplace_back(s, x);
    for (int j = 0; j \le n; j++) {
      int x;
      cin >> x;
      edges[i].emplace_back(x, j);
  cout << arborescence(edges) << '\n';</pre>
  /*int n;
  cin >> n;
  vector<int> a(n), b(n);
  for(int i = 0; i < n; i++) {
    cin >> a[i];
  for(int i = 0; i < n; i++) {
   cin >> b[i]:
  Graph edges(n+1);
  for (int i = 0; i < n; i++) {
```

```
edges[i+1].emplace_back(a[i] ^ b[i], 0);
for(int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    if(i == j) continue;
    edges[i+1].emplace\_back(a[i] ^ b[j], j+1);
long long cost = arborescence(edges);
cout << cost << '\n';</pre>
vector<bool> got(n, false);
long long cur = 0;
for (int i = 0; i < n; i++) {
  int j = 0;
  while(1) {
    while (got[i]) {
      //cout << "skipping " << j << '\n';
    //cout << "testing " << j << endl;
    for (auto &e : edges) e.clear();
    int mn = a[j] ^ b[j];
    for(int k = 0; k < n; k++)  {
      if(got[k] \mid \mid k == j)  {
        mn = min(mn, a[j] ^ b[k]);
      } else {
        int mine = a[k] \hat{b}[k];
        for (int x = 0; x < n; x++) {
          if(got[x] | | x == j | | x == k)  {
            mine = min(mine, a[k] \hat{b}[x]);
             edges[k+1].emplace_back(a[k] \hat{b}[x], x+1);
        edges[k+1].emplace_back(mine, 0);
    //cout << "got here!" << endl;
    long long gott = arborescence(edges);
    //cout << "!" << gott + cur + mn << "\n";
    if(gott + cur + mn == cost) {
      cout << j + 1 << (i + 1 == n ? ' \n' : ' ');
      cur += mn;
      //cout << endl;</pre>
      got[j] = true;
      break;
    j++;
} */
```

# 3 Dynamic Programming

### 3.1 Line Container

```
typedef long long int 11;
```

```
bool Q;
struct Line {
  mutable 11 k, m, p;
  bool operator<(const Line& o) const {</pre>
    return Q ? p < o.p : k < o.k;
};
struct LineContainer : multiset<Line> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b)
  const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = inf; return false; }
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
  void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p)
      isect(x, erase(y));
  11 query(11 x) {
    assert(!empty());
    Q = 1; auto 1 = *lower_bound({0,0,x}); Q = 0;
    return l.k * x + l.m;
};
```

#### 3.2 Li Chao Tree

```
// by luucasv
typedef long long T;
const T INF = 1e18, EPS = 1;
const int BUFFER_SIZE = 1e4;
struct Line {
 T m, b;
 Line (T m = 0, T b = INF): m(m), b(b) {}
 T apply(T x) { return x * m + b; }
};
struct Node {
 Node *left, *right;
 Line line;
  Node(): left(NULL), right(NULL) {}
} ;
struct LiChaoTree {
 Node *root, buffer[BUFFER_SIZE];
 T min_value, max_value;
  int buffer pointer;
  LiChaoTree (T min_value, T max_value): min_value (min_value),
      max_value(max_value + 1) { clear(); }
```

```
void clear() { buffer_pointer = 0; root = newNode(); }
  void insert_line(T m, T b) { update(root, min_value, max_value, Line
      (m, b)); }
  T eval(T x) { return query(root, min_value, max_value, x); }
  void update(Node *cur, T 1, T r, Line line) {
   T m = 1 + (r - 1) / 2;
   bool left = line.apply(1) < cur->line.apply(1);
   bool mid = line.apply(m) < cur->line.apply(m);
   bool right = line.apply(r) < cur->line.apply(r);
   if (mid) {
     swap(cur->line, line);
   if (r - 1 <= EPS) return;</pre>
   if (left == right) return;
   if (mid != left) {
      if (cur->left == NULL) cur->left = newNode();
      update(cur->left, 1, m, line);
      if (cur->right == NULL) cur->right = newNode();
      update(cur->right, m, r, line);
  T query(Node *cur, T l, T r, T x) {
   if (cur == NULL) return INF;
   if (r - 1 <= EPS) {
      return cur->line.apply(x);
   T m = 1 + (r - 1) / 2;
   T ans:
   if (x < m) {
      ans = query(cur->left, 1, m, x);
     ans = query(cur->right, m, r, x);
   return min(ans, cur->line.apply(x));
 Node* newNode() {
     buffer[buffer_pointer] = Node();
      return &buffer[buffer_pointer++];
};
```

## 3.3 Divide and Conquer Optimization

```
int n, k;
ll dpold[ms], dp[ms], c[ms][ms]; // c(i, j) pode ser funcao

void compute(int l, int r, int optl, int optr) {
    if(l>r) return;
    int mid = (l+r)/2;
    pair<ll, int> best = {inf, -1}; // long long inf
    for(int k = optl; k <= min(mid, optr); k++) {
        best = min(best, {dpold[k-1] + c[k][mid], k});
    }
    dp[mid] = best.first;
    int opt = best.second;
    compute(l, mid-1, optl, opt);
    compute(mid+1, r, opt, optr);
}</pre>
```

```
11 solve() {
    dp[0] = 0;
    for(int i = 1; i <= n; i++) dp[i] = inf; // initialize row 0 of
        the dp
    for(int i = 1; i <= k; i++) {
        swap(dpold, dp);
        compute(0, n, 0, n); // solve row i of the dp
    }
    return dp[n]; // return dp[k][n]
}</pre>
```

## 3.4 Knuth Optimization

## 4 Math

## 4.1 Chinese Remainder Theorem

```
#include<bits/stdc++.h>
using namespace std;
const long long N = 20;

long long GCD(long long a, long long b) {
  return (b == 0) ? a : GCD(b, a % b);
}
inline long long get_LCM(long long a, long long b) {
  return a / GCD(a, b) * b;
}
inline long long normalize(long long x, long long mod) {
  x %= mod;
  if (x < 0) x += mod;
  return x;
}

struct GCD_type {
  long long x, y, d;
};</pre>
```

```
GCD_type ex_GCD(long long a, long long b) {
  if (b == 0) return {1, 0, a};
  GCD_type pom = ex_GCD(b, a % b);
  return {pom.y, pom.x - a / b * pom.y, pom.d};
long long testCases;
long long t;
long long a[N], n[N], ans, LCM;
int main() {
  ios_base::sync_with_stdio(0);
  cin.tie(0);
 t = 2;
 long long T;
  cin >> T:
  while (T--)
    for (long long i = 1; i \le t; i++) {
      cin >> a[i] >> n[i];
      normalize(a[i], n[i]);
    ans = a[1];
   LCM = n[1];
    bool impossible = false;
    for (long long i = 2; i \le t; i++) {
      auto pom = ex_GCD(LCM, n[i]);
      long long x1 = pom.x;
      long long d = pom.d;
      if((a[i] - ans) % d != 0) {
        impossible = true;
      ans = normalize(ans + x1 * (a[i] - ans) / d % (n[i] / d) * LCM,
          LCM * n[i] / d);
      LCM = get_LCM(LCM, n[i]);
    if (impossible) cout << "no solution\n";</pre>
    else cout << ans << " " << LCM << endl;</pre>
  return 0;
```

## 4.2 Diophantine Equations

```
int gcd_ext(int a, int b, int& x, int &y) {
    if (b == 0) {
        x = 1, y = 0;
        return a;
    }
    int nx, ny;
    int gc = gcd_ext(b, a % b, nx, ny);
    x = ny;
    y = nx - (a / b) * ny;
    return gc;
}

vector<int> diophantine(int D, vector<int> 1) {
    int n = 1.size();
    vector<int> gc(n), ans(n);
    gc[n - 1] = 1[n - 1];
    for (int i = n - 2; i >= 0; i--) {
```

```
int x, y;
  gc[i] = gcd_ext(l[i], gc[i + 1], x, y);
if (D % gc[0] != 0) {
 return vector<int>();
for (int i = 0; i < n; i++) {
 if (i == n - 1) {
    ans[i] = D / l[i];
    D = l[i] * ans[i];
    continue;
  int x, y;
  gcd_ext(l[i] / gc[i], gc[i + 1] / gc[i], x, y);
 ans[i] = (long long int) D / gc[i] * x % (gc[i + 1] / gc[i]);
 if (D < 0 \&\& ans[i] > 0) {
    ans[i] -= (qc[i + 1] / qc[i]);
  if (D > 0 \&\& ans[i] < 0) {
    ans[i] += (qc[i + 1] / qc[i]);
 D = 1[i] * ans[i];
return ans;
```

## 4.3 Discrete Logarithm

```
11 discreteLog (11 a, 11 b, 11 m) {
    a %= m; b %= m;
    11 n = (11) sqrt (m + .0) + 1, an = 1;
    for (11 i = 0; i < n; i++) {
        an = (an * a) % m;
    }
    map<11, 11> vals;
    for (11 i = 1, cur = an; i <= n; i++) {
        if (!vals.count(cur)) vals[cur] = i;
        cur = (cur * an) % m;
    }
    11 ans = 1e18; //inf
    for (11 i = 0, cur = b; i <= n; i++) {
        if (vals.count(cur)) {
            ans = min(ans, vals[cur] * n - i);
        }
        cur = (cur * a) % m;
    }
    return ans;
}</pre>
```

#### 4.4 Discrete Root

```
//x^k = a % mod

ll discreteRoot(ll k, ll a, ll mod) {
    ll g = primitiveRoot(mod);
    ll y = discreteLog(fexp(g, k, mod), a, mod);
    if (y == -1) {
        return y;
    }
}
```

```
return fexp(g, y, mod);
}
```

#### 4.5 Primitive Root

```
int primitiveRoot(int p) {
  vector<int> fact;
  int phi = p - 1, n = phi;
  for (int i = 2; i * i <= n; i++) {</pre>
   if (n % i == 0) {
      fact.push back(i):
      while (n \% i == 0) {
       n /= i;
  if (n > 1) {
   fact.push_back(n);
  for (int res = 2; res <= p; res++) {</pre>
   bool ok = true;
    for (auto it : fact) {
      ok &= fexp(res, phi / it, p) != 1;
      if (!ok) {
       break;
      }
   if (ok) {
      return res;
  return -1;
```

#### 4.6 Extended Euclides

```
// euclides estendido: acha u e v da equacao:
// u * x + v * y = gcd(x, y);
// u eh inverso modular de x no modulo y
// v eh inverso modular de y no modulo x

pair<11, 11> euclides(11 a, 11 b) {
  11 u = 0, oldu = 1, v = 1, oldv = 0;
  while(b) {
   11 q = a / b;
   oldv = oldv - v * q;
   oldu = oldu - u * q;
   a = a - b * q;
   swap(a, b);
   swap(u, oldu);
   swap(v, oldv);
}

return make_pair(oldu, oldv);
}
```

## 4.7 Matrix Fast Exponentiation

```
const 11 \mod = 1e9+7;
const int m = 2; // size of matrix
struct Matrix {
  11 mat[m][m];
  Matrix operator * (const Matrix &p) {
    Matrix ans;
    for (int i = 0; i < m; i++)
      for (int j = 0; j < m; j++)
        for (int k = ans.mat[i][j] = 0; k < m; k++)
          ans.mat[i][j] = (ans.mat[i][j] + mat[i][k] * p.mat[k][j]) %
    return ans;
} ;
Matrix fExp(Matrix a, ll b) {
 Matrix ans;
  for (int i = 0; i < m; i++) for (int j = 0; j < m; j++)
    ans.mat[i][j] = i == j;
  while(b) {
    if(b \& 1) ans = ans * a;
    a = a * a;
    b >>= 1;
  return ans;
```

#### 4.8 FFT - Fast Fourier Transform

```
typedef double 1d;
const 1d PI = acos(-1);
struct Complex {
  ld real, imag;
  Complex conj() { return Complex(real, -imag); }
  Complex (1d = 0, 1d = 0) : real(a), imag(b) {}
  Complex operator + (const Complex &o) const { return Complex(real +
      o.real, imag + o.imag); }
  Complex operator - (const Complex &o) const { return Complex(real -
      o.real, imag - o.imag); }
  Complex operator * (const Complex &o) const { return Complex(real *
      o.real - imag * o.imag, real * o.imag + imag * o.real); }
  Complex operator / (ld o) const { return Complex(real / o, imag / o)
  void operator *= (Complex o) { *this = *this * o; }
  void operator /= (ld o) { real /= o, imag /= o; }
typedef std::vector<Complex> CVector;
const int ms = 1 << 22:
int bits[ms];
Complex root[ms];
void initFFT() {
  root[1] = Complex(1);
  for(int len = 2; len < ms; len += len) {</pre>
```

```
Complex z(cos(PI / len), sin(PI / len));
    for(int i = len / 2; i < len; i++) {</pre>
      root[2 * i] = root[i];
      root[2 * i + 1] = root[i] * z;
void pre(int n) {
  int LOG = 0;
  while (1 << (LOG + 1) < n) {
    LOG++;
  for (int i = 1; i < n; i++) {
    bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << LOG);
CVector fft(CVector a, bool inv = false) {
  int n = a.size();
  pre(n);
  if(inv) {
    std::reverse(a.begin() + 1, a.end());
  for (int i = 0; i < n; i++) {
    int to = bits[i];
    if(to > i) {
      std::swap(a[to], a[i]);
  for(int len = 1; len < n; len *= 2) {</pre>
    for(int i = 0; i < n; i += 2 * len) {
      for(int j = 0; j < len; j++) {
        Complex u = a[i + j], v = a[i + j + len] * root[len + j];
        a[i + j] = u + v;
        a[i + j + len] = u - v;
    }
  if(inv) {
    for (int i = 0; i < n; i++)
     a[i] /= n;
  return a;
void fft2in1(CVector &a, CVector &b) {
  int n = (int) a.size();
  for (int i = 0; i < n; i++) {
    a[i] = Complex(a[i].real, b[i].real);
  auto c = fft(a);
  for (int i = 0; i < n; i++) {
    a[i] = (c[i] + c[(n-i) % n].conj()) * Complex(0.5, 0);
    b[i] = (c[i] - c[(n-i) % n].conj()) * Complex(0, -0.5);
void ifft2in1(CVector &a, CVector &b) {
  int n = (int) a.size();
```

```
for(int i = 0; i < n; i++) {
    a[i] = a[i] + b[i] * Complex(0, 1);
 a = fft(a, true);
  for(int i = 0; i < n; i++) {
   b[i] = Complex(a[i].imag, 0);
    a[i] = Complex(a[i].real, 0);
std::vector<long long> mod_mul(const std::vector<long long> &a, const
    std::vector<long long> &b, long long cut = 1 << 15) {</pre>
  // TODO cut memory here by /2
  int n = (int) a.size();
 CVector C[4];
  for(int i = 0; i < 4; i++) {
    C[i].resize(n);
  for(int i = 0; i < n; i++) {
   C[0][i] = a[i] % cut;
    C[1][i] = a[i] / cut;
    C[2][i] = b[i] % cut;
    C[3][i] = b[i] / cut;
  fft2in1(C[0], C[1]);
  fft2in1(C[2], C[3]);
  for(int i = 0; i < n; i++) {
    // 00, 01, 10, 11
    Complex cur[4];
    for(int j = 0; j < 4; j++) cur[j] = C[j/2+2][i] * C[j % 2][i];
    for (int j = 0; j < 4; j++) C[j][i] = cur[j];
  ifft2in1(C[0], C[1]);
  ifft2in1(C[2], C[3]);
  std::vector<long long> ans(n, 0);
  for(int i = 0; i < n; i++) {</pre>
    // if there are negative values, care with rounding
    ans[i] += (long long) (C[0][i].real + 0.5);
    ans[i] += (long long) (C[1][i].real + C[2][i].real + 0.5) * cut;
    ans[i] += (long long) (C[3][i].real + 0.5) * cut * cut;
  return ans;
std::vector<int> mul(const std::vector<int> &a, const std::vector<int>
     (d3
  int n = 1;
  while (n - 1 < (int) \ a.size() + (int) \ b.size() - 2) \ n += n;
  CVector poly(n);
  for(int i = 0; i < n; i++) {</pre>
    if(i < (int) a.size()) {
      poly[i].real = a[i];
   if(i < (int) b.size()) {
      poly[i].imag = b[i];
  poly = fft(poly);
  for(int i = 0; i < n; i++) {</pre>
    poly[i] *= poly[i];
```

```
}
poly = fft(poly, true);
std::vector<int> c(n, 0);
for(int i = 0; i < n; i++) {
   c[i] = (int) (poly[i].imag / 2 + 0.5);
}
while (c.size() > 0 && c.back() == 0) c.pop_back();
return c;
```

#### 4.9 NTT - Number Theoretic Transform

```
long long int mod = (11911 << 23) + 1, c root = 3;
namespace NTT {
  typedef long long int 11;
 11 fexp(ll base, ll e) {
   11 \text{ ans} = 1;
    while(e > 0) {
      if (e & 1) ans = ans * base % mod;
      base = base * base % mod;
      e >>= 1:
    return ans:
  11 inv mod(ll base) {
    return fexp(base, mod - 2);
  void ntt(vector<ll>& a, bool inv) {
    int n = (int) a.size();
    if (n == 1) return;
    for (int i = 0, j = 0; i < n; i++) {
      if (i > j) {
        swap(a[i], a[j]);
      for (int 1 = n / 2; († \hat{} = 1) < 1; 1 >>= 1);
    for(int sz = 1; sz < n; sz <<= 1) {</pre>
      11 delta = fexp(c_root, (mod - 1) / (2 * sz)); //delta = w_2sz
      if (inv) {
        delta = inv_mod(delta);
      for (int i = 0; i < n; i += 2 * sz) {
        11 w = 1:
        for(int j = 0; j < sz; j++) {
          11 \ u = a[i + j], \ v = w * a[i + j + sz] % mod;
          a[i + j] = (u + v + mod) % mod;
          a[i + j] = (a[i + j] + mod) % mod;
          a[i + j + sz] = (u - v + mod) % mod;
          a[i + j + sz] = (a[i + j + sz] + mod) % mod;
          w = w * delta % mod;
    if (inv) {
```

```
ll inv_n = inv_mod(n);
      for(int i = 0; i < n; i++) {
        a[i] = a[i] * inv n % mod;
    for(int i = 0; i < n; i++) {
      a[i] %= mod;
      a[i] = (a[i] + mod) % mod;
  void multiply(vector<ll> &a, vector<ll> &b, vector<ll> &ans) {
    int lim = (int) max(a.size(), b.size());
    int n = 1;
    while (n < lim) n <<= 1;
    n <<= 1;
    a.resize(n):
    b.resize(n);
    ans.resize(n):
    ntt(a, false);
    ntt(b, false);
    for(int i = 0; i < n; i++) {
     ans[i] = a[i] * b[i] % mod;
   ntt(ans, true);
};
```

## 4.10 Fast Walsh Hadamard Transform

```
vector<ll> FWHT(char oper, vector<ll> a, const bool inv = false) {
 int n = (int) a.size();
  for(int len = 1; len < n; len += len) {</pre>
    for(int i = 0; i < n; i += 2 * len) {
      for(int j = 0; j < len; j++) {
        auto u = a[i + j] % mod, v = a[i + j + len] % mod;
        if(oper == '^') {
          a[i + j] = (u + v) % mod;
          a[i + j + len] = (u - v + mod) % mod;
        if(oper == '|') {
          if(!inv) {
            a[i + j + len] = (u + v) % mod;
            a[i + j + len] = (v - u + mod) % mod;
        if(oper == '&') {
          if(!inv) {
            a[i + j] = (u + v) \% mod;
          } else {
            a[i + j] = (u - v + mod) % mod;
  if(oper == '^' && inv) {
    11 \text{ rev} = \text{fexp(n, mod - 2);}
    for (int i = 0; i < n; i++) {
```

```
a[i] = a[i] * rev % mod;
  return a;
vector<ll> multiply(char oper, vector<ll> a, vector<ll> b) {
  int n = 1;
  while (n < (int) max(a.size(), b.size())) {</pre>
   n <<= 1:
  vector<ll> ans(n);
  while (a.size() < ans.size()) a.push_back(0);</pre>
  while (b.size() < ans.size()) b.push_back(0);</pre>
  a = FWHT(oper, a);
 b = FWHT (oper, b);
  for (int i = 0; i < n; i++) {</pre>
   ans[i] = a[i] * b[i] % mod;
  ans = FWHT (oper, ans, true);
  return ans;
const int mxlog = 17;
vector<ll> subset_multiply(vector<ll> a, vector<ll> b) {
  while (n < (int) max(a.size(), b.size())) {</pre>
   n <<= 1;
  vector<ll> ans(n);
  while (a.size() < ans.size()) a.push back(0);</pre>
  while (b.size() < ans.size()) b.push_back(0);</pre>
  vector<vector<ll>> A(mxlog + 1, vector<ll>(a.size())), B(mxlog + 1,
      vector<ll>(b.size()));
  for (int i = 0; i < n; i++)</pre>
    A[__builtin_popcount(i)][i] = a[i];
    B[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= mxlog; i++) {</pre>
    A[i] = FWHT('|', A[i]);
    B[i] = FWHT('|', B[i]);
  for (int i = 0; i <= mxlog; i++) {</pre>
    vector<ll> C(n);
    for (int x = 0; x <= i; x++) {
      int v = i - x;
      for (int j = 0; j < n; j++) {
        C[\dot{\uparrow}] = (C[\dot{\uparrow}] + A[x][\dot{\uparrow}] * B[y][\dot{\uparrow}] % mod) % mod;
    C = FWHT('|', C, true);
    for (int j = 0; j < n; j++) {
      if ( builtin popcount(j) == i) {
        ans[j] = (ans[j] + C[j]) % mod;
  return ans:
```

#### 4.11 Miller and Rho

```
typedef long long int 11;
bool overflow(ll a, ll b) {
  return b && (a >= (111 << 62) / b);
11 add(11 a, 11 b, 11 md) {
  return (a + b) % md;
11 mul(ll a, ll b, ll md) {
  if (!overflow(a, b)) return (a * b) % md;
  11 \text{ ans} = 0;
  while(b) {
    if (b & 1) ans = add(ans, a, md);
    a = add(a, a, md);
    b >>= 1;
  return ans;
11 fexp(ll a, ll e, ll md) {
 11 \text{ ans} = 1;
  while(e) {
    if (e & 1) ans = mul(ans, a, md);
    a = mul(a, a, md):
    e >>= 1;
  return ans:
11 my_rand() {
 ll ans = rand();
  ans = (ans \ll 31) \mid rand();
  return ans;
11 gcd(ll a, ll b) {
  while(b) {
    11 t = a % b;
    a = b;
    b = t;
  return a;
bool miller(ll p, int iteracao) {
  if(p < 2) return 0:
  if(p % 2 == 0) return (p == 2);
 11 s = p - 1;
  while(s \% 2 == 0) s >>= 1;
  for(int i = 0; i < iteracao; i++) {</pre>
   11 a = rand() % (p - 1) + 1, temp = s;
    11 \mod = fexp(a, temp, p);
    while (temp != p - 1 && mod != 1 && mod != p - 1) {
      mod = mul(mod, mod, p);
      temp <<= 1;
```

```
if(mod != p - 1 && temp % 2 == 0) return 0;
 return 1;
11 rho(ll n) {
 if (n == 1 || miller(n, 10)) return n;
 if (n % 2 == 0) return 2;
  while(1) {
   11 x = my_rand() % (n - 2) + 2, y = x;
   11 c = 0, cur = 1;
   while(c == 0)
      c = my_rand() % (n - 2) + 1;
   while(cur == 1) {
      x = add(mul(x, x, n), c, n);
     y = add(mul(y, y, n), c, n);
     y = add(mul(y, y, n), c, n);
     cur = gcd((x >= y ? x - y : y - x), n);
   if (cur != n) return cur;
```

### 4.12 Determinant using Mod

// by zchao1995

```
// Determinante com coordenadas inteiras usando Mod
11 mat[ms][ms];
11 det (int n) {
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      mat[i][j] %= mod;
  11 \text{ res} = 1;
  for (int i = 0; i < n; i++) {
    if (!mat[i][i]) {
      bool flag = false;
      for (int j = i + 1; j < n; j++) {
        if (mat[j][i]) {
          flag = true;
          for (int k = i; k < n; k++) {
            swap (mat[i][k], mat[j][k]);
          res = -res;
          break;
      if (!flag) {
        return 0;
    for (int j = i + 1; j < n; j++) {
      while (mat[j][i]) {
        11 t = mat[i][i] / mat[j][i];
        for (int k = i; k < n; k++) {
          mat[i][k] = (mat[i][k] - t * mat[j][k]) % mod;
```

```
swap (mat[i][k], mat[j][k]);

    res = -res;
    }
} res = (res * mat[i][i]) % mod;
}
return (res + mod) % mod;
}
```

## 4.13 Lagrange Interpolation

```
class LagrangePoly {
public:
  LagrangePoly(std::vector<long long> _a) {
    //f(i) = \underline{a}[i]
    //interpola o vetor em um polinomio de grau y.size() - 1
    den.resize(y.size());
    int n = (int) y.size();
    for (int i = 0; i < n; i++) {
      y[i] = (y[i] % MOD + MOD) % MOD;
      den[i] = ifat[n - i - 1] * ifat[i] % MOD;
      if((n - i - 1) % 2 == 1) {
        den[i] = (MOD - den[i]) % MOD;
  long long getVal(long long x) {
    int n = (int) y.size();
    x \% = MOD;
    if(x < n) {
      //return y[(int) x];
    std::vector<long long> 1, r;
    l.resize(n);
    1[0] = 1;
    for (int i = 1; i < n; i++) {
     l[i] = l[i - 1] * (x - (i - 1) + MOD) % MOD;
    r.resize(n);
    r[n - 1] = 1;
    for (int i = n - 2; i >= 0; i--) {
      r[i] = r[i + 1] * (x - (i + 1) + MOD) % MOD;
    long long ans = 0;
    for (int i = 0; i < n; i++) {
      long long coef = l[i] * r[i] % MOD;
      ans = (ans + coef * y[i] % MOD * den[i]) % MOD;
    return ans;
  std::vector<long long> y, den;
int main(){
 fat[0] = ifat[0] = 1;
```

```
for(int i = 1; i < ms; i++) {
    fat[i] = fat[i - 1] * i % MOD;
    ifat[i] = fexp(fat[i], MOD - 2);
}

// Codeforces 622F
int x, k;
std::cin >> x >> k;
std::vector<long long> a;
a.push_back(0);
for(long long i = 1; i <= k + 1; i++) {
    a.push_back((a.back() + fexp(i, k)) % MOD);
}
LagrangePoly f(a);
std::cout << f.getVal(x) << '\n';</pre>
```

#### 4.14 Count integer points inside triangle

```
//gcd(p, q) == 1
ll get(ll p, ll q, ll n, bool floor = true) {
   if (n == 0) {
      return 0;
   }
   if (p % q == 0) {
      return n * (n + 1) / 2 * (p / q);
   }
   if (p > q) {
      return n * (n + 1) / 2 * (p / q) + get(p % q, q, n, floor);
   }
   ll new_n = p * n / q;
   ll ans = (n + 1) * new_n - get(q, p, new_n, false);
   if (!floor) {
      ans += n - n / q;
   }
   return ans;
}
```

## 5 Geometry

## 5.1 Geometry

```
const double inf = le100, eps = le-9;
const double PI = acos(-1.0L);

int cmp (double a, double b = 0) {
   if (abs(a-b) < eps) return 0;
   return (a < b) ? -1 : +1;
}

struct PT {
   double x, y;
   PT (double x = 0, double y = 0) : x(x), y(y) {}
   PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
   PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
   PT operator * (double c) const { return PT(x*c, y*c); }
   PT operator / (double c) const { return PT(x*c, y*c); }</pre>
```

```
bool operator < (const PT &p) const {
    if(cmp(x, p.x) != 0) return x < p.x;
    return cmp(y, p.y) < 0;
 bool operator ==(const PT &p) const {
    return !cmp(x, p.x) && !cmp(y, p.y);
  bool operator != (const PT &p) const {
    return ! (p == *this);
double dot (PT p, PT q) { return p.x * q.x + p.y*q.y; }
double cross (PT p, PT q) { return p.x * q.y - p.y*q.x; }
double dist2 (PT p, PT q = PT(0, 0)) { return dot(p-q, p-q); }
double dist (PT p, PT q) { return hypot(p.x-q.x, p.y-q.y); }
double norm (PT p) { return hypot(p.x, p.y); }
PT normalize (PT p) { return p/hypot(p.x, p.y); }
double angle (PT p, PT q) { return atan2(cross(p, q), dot(p, q)); }
double angle (PT p) { return atan2(p.y, p.x); }
double polarAngle (PT p) {
  double a = atan2(p.y,p.x);
  return a < 0 ? a + 2*PI : a;
// - p.y*sen(+90), p.x*sen(+90)
PT rotateCCW90 (PT p) { return PT(-p.y, p.x); }
// - p.y*sen(-90), p.x*sen(-90)
PT rotateCW90 (PT p) { return PT(p.y, -p.x); }
PT rotateCCW (PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// !!! PT (int, int)
typedef pair<PT, int> Line;
PT getDir (PT a, PT b) {
 if (a.x == b.x) return PT(0, 1);
  if (a.y == b.y) return PT(1, 0);
  int dx = b.x-a.x;
  int dy = b.y-a.y;
  int g = __gcd(abs(dx), abs(dy));
  if (dx < 0) q = -q;
  return PT(dx/q, dy/q);
Line getLine (PT a, PT b) {
 PT dir = getDir(a, b);
  return {dir, cross(dir, a)};
// Projeta ponto c na linha a - b assumindo a != b
// a.b = |a| cost * |b|
PT projectPointLine (PT a, PT b, PT c) {
  return a + (b-a) * dot (b-a, c-a)/dot (b-a, b-a);
PT reflectPointLine (PT a, PT b, PT c) {
 PT p = projectPointLine(a, b, c);
  return p*2 - c;
// Projeta ponto c no segmento a - b
```

```
PT projectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (cmp(r) == 0) return a;
  r = dot(b-a, c-a)/r;
  if (cmp(r, 0) < 0) return a;
  if (cmp(r, 1) > 0) return b;
  return a + (b - a) * r;
// Calcula distancia entre o ponto c e o segmento a - b
double distancePointSegment (PT a, PT b, PT c) {
  return dist(c, projectPointSegment(a, b, c));
// Parallel and opposite directions
// Determina se o ponto c esta em um segmento a - b
bool ptInSegment (PT a, PT b, PT c) {
 if (a == b) return a == c;
 a = a-c, b = b-c;
 return cmp(cross(a, b)) == 0 && cmp(dot(a, b)) <= 0;
// Determina se as linhas a - b e c - d sao paralelas ou colineares
bool parallel (PT a, PT b, PT c, PT d) {
 return cmp (cross (b - a, c - d)) == 0;
bool collinear (PT a, PT b, PT c, PT d) {
  return parallel(a, b, c, d) && cmp(cross(a - b, a - c)) == 0 && cmp(
      cross(c - d, c - a)) == 0;
// Calcula distancia entre o ponto (x, y, z) e o plano ax + by + cz =
double distancePointPlane(double x, double y, double z, double a,
    double b, double c, double d) {
    return abs(a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c
        );
// Determina se o segmento a - b intersecta com o segmento c - d
bool segmentsIntersect (PT a, PT b, PT c, PT d) {
 if (collinear(a, b, c, d)) {
   if (cmp(dist(a, c)) == 0 \mid | cmp(dist(a, d)) == 0 \mid | cmp(dist(b, c))
        ) == 0 || cmp(dist(b, d)) == 0) return true;
   if (cmp(dot(c - a, c - b)) > 0 \&\& cmp(dot(d - a, d - b)) > 0 \&\&
        cmp(dot(c - b, d - b)) > 0) return false;
   return true;
  if (cmp(cross(d - a, b - a) * cross(c - a, b - a)) > 0) return false
  if (cmp(cross(a - c, d - c) * cross(b - c, d - c)) > 0) return false
  return true;
// Calcula a intersecao entre as retas a - b e c - d assumindo que uma
     unica intersecao existe
// Para intersecao de segmentos, cheque primeiro se os segmentos se
    intersectam e que nao sao paralelos
// r = a1 + t*d1, (r - a2) x d2 = 0
```

```
PT computeLineIntersection (PT a, PT b, PT c, PT d) {
 b = b - a; d = c - d; c = c - a;
 assert (cmp(cross(b, d)) != 0);
  return a + b * cross(c, d) / cross(b, d);
// Calcula centro do circulo dado tres pontos
PT computeCircleCenter (PT a, PT b, PT c) {
 b = (a + b) / 2; // bissector
 c = (a + c) / 2; // bissector
 return computeLineIntersection(b, b + rotateCW90(a - b), c, c +
      rotateCW90(a - c));
vector<PT> circle2PtsRad (PT p1, PT p2, double r) {
 vector<PT> ret;
  double d2 = dist2(p1, p2);
  double det = r * r / d2 - 0.25;
  if (det < 0.0) return ret;</pre>
  double h = sgrt(det);
  for (int i = 0; i < 2; i++) {
    double x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
    double y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;
   ret.push back(PT(x, v));
   swap(p1, p2);
  return ret;
// Calcula intersecao da linha a - b com o circulo centrado em c com
    raio r > 0
bool circleLineIntersection(PT a, PT b, PT c, double r) {
    return cmp(dist(c, projectPointLine(a, b, c)), r) <= 0;</pre>
vector<PT> circleLine (PT a, PT b, PT c, double r) {
  vector<PT> ret;
  PT p = projectPointLine(a, b, c), p1;
  double h = norm(c-p);
  if (cmp(h,r) == 0) {
    ret.push_back(p);
  else if (cmp(h,r) < 0) {
    double k = sqrt(r*r - h*h);
    p1 = p + (b-a) / (norm(b-a)) *k;
    ret.push_back(p1);
   p1 = p - (b-a)/(norm(b-a)) *k;
    ret.push_back(p1);
  return ret;
bool ptInsideTriangle(PT p, PT a, PT b, PT c) {
  if(cross(b-a, c-b) < 0) swap(a, b);
  if(ptInSegment(a,b,p)) return 1;
  if(ptInSegment(b,c,p)) return 1;
  if(ptInSegment(c,a,p)) return 1;
  bool x = cross(b-a, p-b) < 0;
  bool y = cross(c-b, p-c) < 0;
  bool z = cross(a-c, p-a) < 0;
  return x == y && y == z;
```

```
// Determina se o ponto esta num poligono convexo em O(lgn)
bool pointInConvexPolygon(const vector<PT> &p, PT g) {
  if (p.size() == 1) return p.front() == q;
  int 1 = 1, r = p.size()-1;
  while (abs(r-1) > 1) {
    int m = (r+1)/2;
    if(cross(p[m]-p[0], q-p[0]) < 0) r = m;
    else 1 = m;
  return ptInsideTriangle(q, p[0], p[1], p[r]);
// Determina se o ponto esta num poligono possivelmente nao-convexo
// Retorna 1 para pontos estritamente dentro, 0 para pontos
    estritamente fora do poligno
// e 0 ou 1 para os pontos restantes
// Eh possivel converter num teste exato usando inteiros e tomando
    cuidado com a divisao
// e entao usar testes exatos para checar se esta na borda do poligno
bool pointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for(int i = 0; i < p.size(); i++){</pre>
    int j = (i + 1) % p.size();
    if((p[i].y \le q.y \& q.y < p[j].y || p[j].y \le q.y \& q.y < p[i].y
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y)
         i].y))
      c = !c;
  return c;
// Determina se o ponto esta na borda do poligno
bool pointOnPolygon(const vector<PT> &p, PT q) {
  for(int i = 0; i < p.size(); i++)</pre>
    if(cmp(dist2(projectPointSegment(p[i], p[(i + 1) % p.size()], q),
        (q)) < 0)
      return true;
    return false:
// area / semiperimeter
double rIncircle (PT a, PT b, PT c) {
  double ab = norm(a-b), bc = norm(b-c), ca = norm(c-a);
  return abs(cross(b-a, c-a)/(ab+bc+ca));
// Calcula intersecao do circulo centrado em a com raio r e o centrado
     em b com raio R
vector<PT> circleCircle (PT a, double r, PT b, double R) {
  vector<PT> ret;
  double d = norm(a-b);
  if (d > r + R \mid \mid d + min(r, R) < max(r, R)) return ret;
  double x = (d*d - R*R + r*r) / (2*d); // x = r*cos(R opposite angle)
  double y = sqrt(r*r - x*x);
  PT v = (b - a)/d;
  ret.push_back(a + v*x + rotateCCW90(v)*y);
  if (cmp(y) > 0)
    ret.push_back(a + v*x - rotateCCW90(v)*y);
  return ret;
```

```
double circularSegArea (double r, double R, double d) {
  double ang = 2 * acos((d*d - R*R + r*r) / (2*d*r)); // cos(R)
      opposite angle) = x/r
  double tri = sin(ang) * r * r;
  double sector = ang * r * r;
  return (sector - tri) / 2;
// Calcula a area ou o centroide de um poligono (possivelmente nao-
    convexo)
// assumindo que as coordenadas estao listada em ordem horaria ou anti
    -horaria
// O centroide eh equivalente a o centro de massa ou centro de
    gravidade
double computeSignedArea (const vector<PT> &p) {
  double area = 0:
  for (int i = 0; i < p.size(); i++) {</pre>
    int i = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area/2.0;
double computeArea(const vector<PT> &p) {
  return abs(computeSignedArea(p));
PT computeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for(int i = 0; i < p.size(); i++) {</pre>
   int j = (i + 1) % p.size();
    c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
  return c / scale;
// Testa se o poligno listada em ordem CW ou CCW eh simples (nenhuma
    linha se intersecta)
bool isSimple(const vector<PT> &p) {
  for(int i = 0; i < p.size(); i++) {</pre>
    for (int k = i + 1; k < p.size(); k++) {
      int j = (i + 1) % p.size();
      int 1 = (k + 1) % p.size();
      if (i == 1 \mid | i == k) continue;
      if (segmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
vector< pair<PT, PT> > getTangentSegs (PT c1, double r1, PT c2, double
     r2) {
  if (r1 < r2) swap(c1, c2), swap(r1, r2);
  vector<pair<PT, PT> > ans;
  double d = dist(c1, c2);
 if (cmp(d) <= 0) return ans;</pre>
  double dr = abs(r1 - r2), sr = r1 + r2;
  if (cmp(dr, d) >= 0) return ans;
```

```
double u = acos(dr / d);
PT dc1 = normalize(c2 - c1)*r1;
PT dc2 = normalize(c2 - c1)*r2;
ans.push_back(make_pair(c1 + rotateCCW(dc1, +u), c2 + rotateCCW(dc2, +u)));
ans.push_back(make_pair(c1 + rotateCCW(dc1, -u), c2 + rotateCCW(dc2, -u)));
if (cmp(sr, d) >= 0) return ans;
double v = acos(sr / d);
dc2 = normalize(c1 - c2)*r2;
ans.push_back({c1 + rotateCCW(dc1, +v), c2 + rotateCCW(dc2, +v)});
ans.push_back({c1 + rotateCCW(dc1, -v), c2 + rotateCCW(dc2, -v)});
return ans;
```

#### 5.2 Convex Hull

```
vector<PT> convexHull(vector<PT> p, bool needs = 1) {
  if(needs) sort(p.begin(), p.end());
  p.erase(unique(p.begin(), p.end()), p.end());
  int n = p.size(), k = 0;
  if(n <= 1) return p;</pre>
  vector<PT> h(n + 2);
  for (int i = 0; i < n; i++) {
    while (k \ge 2 \&\& cross(h[k-1] - h[k-2], p[i] - h[k-2]) \le 0)
        k--:
   h[k++] = p[i];
  for (int i = n - 2, t = k + 1; i >= 0; i--) {
    while (k > = t \&\& cross(h[k - 1] - h[k - 2], p[i] - h[k - 2]) \le 0)
        k--:
   h[k++] = p[i];
  h.resize(k); // n+1 points where the first is equal to the last
  return h;
void sortByAngle (vector<PT>::iterator first, vector<PT>::iterator
    last, const PT o) {
  first = partition(first, last, [&o] (const PT &a) { return a == o;
      });
  auto pivot = partition(first, last, [&o] (const PT &a) {
    return ! (a < 0 | | a == 0); // PT(a.v, a.x) < PT(o.v, o.x)
  });
  auto acmp = [&o] (const PT &a, const PT &b) { // C++11 only
    if (cmp(cross(a-o, b-o)) != 0) return cross(a-o, b-o) > 0;
    else return cmp(norm(a-o), norm(b-o)) < 0;</pre>
  };
  sort(first, pivot, acmp);
  sort(pivot, last, acmp);
vector<PT> graham (vector<PT> v) {
  sort(v.begin(), v.end());
  sortByAngle(v.begin(), v.end(), v[0]);
 vector<PT> u (v.size());
  int top = 0;
  for (int i = 0; i < v.size(); i++) {</pre>
    while (top > 1 \&\& cmp(cross(u[top-1] - u[top-2], v[i]-u[top-2]))
        <= 0) top--;
```

```
u[top++] = v[i];
  u.resize(top);
  return u;
vector<PT> splitHull(const vector<PT> &hull) {
  vector<PT> ans(hull.size());
  for (int i = 0, j = (int) hull.size()-1, k = 0; k < (int) hull.size()
      ; k++) {
    if(hull[i] < hull[j]) {</pre>
      ans[k] = hull[i++];
    } else {
      ans[k] = hull[j--];
  return ans;
vector<PT> ConvexHull(const vector<PT> &a, const vector<PT> &b) {
  auto A = splitHull(a);
  auto B = splitHull(b);
  vector<PT> C(A.size() + B.size());
  merge(A.begin(), A.end(), B.begin(), B.end(), C.begin());
  return ConvexHull(C, false);
int maximizeScalarProduct(const vector<PT> &hull, PT vec) {
  // this code assumes that there are no 3 colinear points
  int ans = 0;
  int n = hull.size();
  if(n < 20) {
    for (int i = 0; i < n; i++) {
      if (dot (hull[i], vec) > dot (hull[ans], vec)) {
        ans = i:
  } else {
    if(dot(hull[1], vec) > dot(hull[ans], vec)) {
      ans = 1:
    for(int rep = 0; rep < 2; rep++) {</pre>
      int 1 = 2, r = n - 1;
      while (1 != r) {
        int mid = (1 + r + 1) / 2;
        bool flag = dot(hull[mid], vec) >= dot(hull[mid-1], vec);
        if(rep == 0) { flag = flag && dot(hull[mid], vec) >= dot(hull
            [0], vec); }
        else { flag = flag || dot(hull[mid-1], vec) < dot(hull[0], vec</pre>
            ); }
        if(flag) {
         1 = mid;
        } else {
          r = mid - 1;
      if(dot(hull[ans], vec) < dot(hull[1], vec)) {</pre>
        ans = 1:
```

```
return ans;
```

## 5.3 Cut Polygon

```
struct Segment {
  typedef long double T;
  PT p1, p2;
 T a, b, c;
  Segment() {}
  Segment (PT st, PT en) {
    p1 = st, p2 = en;
   a = -(st.y - en.y);
   b = st.x - en.x;
    c = a * en.x + b * en.y;
  T plug(T x, T y) {
    // plug >= 0 is to the right
    return a * x + b * y - c;
  T plug(PT p) {
    return plug(p.x, p.y);
 bool inLine(PT p) { return cross((p - p1), (p2 - p1)) == 0; }
 bool inSegment(PT p) {
    return inLine(p) && dot((p1 - p2), (p - p2)) >= 0 && dot((p2 - p1)
        (p - p1)) >= 0;
  PT lineIntersection(Segment s) {
    long double A = a, B = b, C = c;
    long double D = s.a, E = s.b, F = s.c;
    long double x = (long double) C * E - (long double) B * F;
    long double y = (long double) A * F - (long double) C * D;
    long double tmp = (long double) A \star E - (long double) B \star D;
    x /= tmp;
    y /= tmp;
    return PT(x, y);
  bool polygonIntersection(const vector<PT> &poly) {
    long double 1 = -1e18, r = 1e18;
    for(auto p : poly) {
      long double z = plug(p);
     1 = \max(1, z);
      r = min(r, z);
    return 1 - r > eps:
};
vector<PT> cutPolygon(vector<PT> poly, Segment seg) {
  int n = (int) poly.size();
  vector<PT> ans;
  for (int i = 0; i < n; i++) {
```

```
double z = seg.plug(poly[i]);
  if(z > -eps) {
    ans.push_back(poly[i]);
  }
  double z2 = seg.plug(poly[(i + 1) % n]);
  if((z > eps && z2 < -eps) || (z < -eps && z2 > eps)) {
    ans.push_back(seg.lineIntersection(Segment(poly[i], poly[(i + 1) % n])));
    }
  }
  return ans;
}
```

## 5.4 Smallest Enclosing Circle

```
typedef pair<PT, double> circle;
bool inCircle (circle c, PT p) {
  return cmp(dist(c.first, p), c.second) <= 0;</pre>
PT circumcenter (PT p, PT q, PT r) {
 PT a = p-r, b = q-r;
  PT c = PT(dot(a, p+r)/2, dot(b, q+r)/2);
  return PT(cross(c, PT(a.y,b.y)), cross(PT(a.x,b.x), c)) / cross(a, b
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
circle spanningCircle (vector<PT> &v) {
  int n = v.size();
  shuffle(v.begin(), v.end(), rng);
  circle C(PT(), -1);
  for (int i = 0; i < n; i++) if (!inCircle(C, v[i])) {</pre>
    C = circle(v[i], 0);
    for (int j = 0; j < i; j++) if (!inCircle(C, v[j])) {
      C = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
      for (int k = 0; k < j; k++) if (!inCircle(C, v[k])) {
        PT o = circumcenter(v[i], v[j], v[k]);
        C = circle(o, dist(o, v[k]));
  return C;
```

## 5.5 Minkowski

```
bool comp(PT a, PT b) {
  int hp1 = (a.x < 0 || (a.x==0 && a.y<0));
  int hp2 = (b.x < 0 || (b.x==0 && b.y<0));
  if(hp1 != hp2) return hp1 < hp2;
  long long R = cross(a, b);
  if(R) return R > 0;
  return dot(a, a) < dot(b, b);
}

// This code assumes points are ordered in ccw and the first points
  in both vectors is the min lexicographically</pre>
```

```
vector<PT> minkowskiSum(const vector<PT> &a, const vector<PT> &b) {
  if(a.empty() || b.empty()) return vector<PT>(0);
  vector<PT> ret;
  int n1 = a.size(), n2 = b.size();
  if(min(n1, n2) < 2){
   for (int i = 0; i < n1; i++) {
      for (int j = 0; j < n2; j++) {
        ret.push_back(a[i]+b[j]);
   return ret;
  PT v1, v2, p = a[0]+b[0];
  ret.push_back(p);
  for (int i = 0, j = 0; i + j + 1 < n1+n2; ) {
   v1 = a[(i+1)%n1]-a[i];
   v2 = b[(j+1) n2] - b[j];
   if(j == n2 \mid | (i < n1 \&\& comp(v1, v2))) p = p + v1, i++;
   else p = p + v2, j++;
   while(ret.size() >= 2 && cmp(cross(p-ret.back(), p-ret[(int))ret.
        size()-2])) == 0) {
      // removing colinear points
     // needs the scalar product stuff it the result is a line
      ret.pop back();
    ret.push_back(p);
  return ret;
```

#### 5.6 Half Plane Intersection

```
struct L {
   PT a, b;
    L(){}
    L(PT a, PT b) : a(a), b(b) {}
};
double angle (L la) { return atan2(-(la.a.y - la.b.y), la.b.x - la.a.x
    ); }
bool comp (L la, L lb) {
    if (cmp(angle(la), angle(lb)) == 0) return cross((lb.b - lb.a), (
        la.b - lb.a)) > eps;
    return cmp(angle(la), angle(lb)) < 0;</pre>
PT computeLineIntersection (L la, L lb) {
    return computeLineIntersection(la.a, la.b, lb.a, lb.b);
bool check (L la, L lb, L lc) {
    PT p = computeLineIntersection(lb, lc):
    double det = cross((la.b - la.a), (p - la.a));
    return cmp(det) < 0;</pre>
vector<PT> hpi (vector<L> line) { // salvar (i, j) CCW, (j, i) CW
    sort(line.begin(), line.end(), comp);
    vector<L> pl(1, line[0]);
```

```
for (int i = 0; i < (int)line.size(); ++i) if (cmp(angle(line[i]),</pre>
     angle(pl.back())) != 0) pl.push_back(line[i]);
deque<int> dq;
dq.push_back(0);
dq.push_back(1);
for (int i = 2; i < (int)pl.size(); ++i) {</pre>
    while ((int)dq.size() > 1 && check(pl[i], pl[dq.back()], pl[dq
        [dq.size() - 2]])) dq.pop_back();
    while ((int)dq.size() > 1 && check(pl[i], pl[dq[0]], pl[dq
        [1]])) dq.pop_front();
    dq.push_back(i);
while ((int)dq.size() > 1 && check(pl[dq[0]], pl[dq.back()], pl[dq
    [dq.size() - 2]])) dq.pop_back();
while ((int)dq.size() > 1 && check(pl[dq.back()], pl[dq[0]], pl[dq
    [1]])) dq.pop_front();
vector<PT> res;
for (int i = 0; i < (int)dq.size(); ++i){</pre>
    res.emplace_back(computeLineIntersection(pl[dq[i]], pl[dq[(i +
         1) % dq.size()]]));
return res;
```

#### 5.7 Closest Pair

```
double closestPair(vector<PT> p) {
  int n = p.size(), k = 0;
  sort(p.begin(), p.end());
  double d = inf;
  set<PT> ptsInv;
  for(int i = 0; i < n; i++) {
    while(k < i && p[k].x < p[i].x - d) {
        ptsInv.erase(swapCoord(p[k++]));
    }
  for(auto it = ptsInv.lower_bound(PT(p[i].y - d, p[i].x - d));
        it != ptsInv.end() && it->x <= p[i].y + d; it++) {
        d = min(d, dist(p[i] - swapCoord(*it), PT(0, 0)));
    }
    ptsInv.insert(swapCoord(p[i]));
}
return d;
}</pre>
```

## 5.8 Delaunay Triangulation

```
bool ge(const ll& a, const ll& b) { return a >= b; }
bool le(const ll& a, const ll& b) { return a <= b; }
bool eq(const ll& a, const ll& b) { return a == b; }
bool gt(const ll& a, const ll& b) { return a == b; }
bool lt(const ll& a, const ll& b) { return a > b; }
bool lt(const ll& a, const ll& b) { return a < b; }
int sgn(const ll& a) { return a >= 0 ? a ? l : 0 : -1; }

struct pt {
    ll x, y;
    pt() { }
    pt(ll _x, ll _y) : x(_x), y(_y) { }
    pt operator-(const pt& p) const {
```

```
return pt(x - p.x, y - p.y);
    11 cross(const pt& p) const {
        return x * p.y - y * p.x;
   ll cross(const pt& a, const pt& b) const {
        return (a - *this).cross(b - *this);
    11 dot(const pt& p) const {
        return x * p.x + y * p.y;
   11 dot (const pt& a, const pt& b) const {
        return (a - *this).dot(b - *this);
    11 sgrLength() const {
        return this->dot(*this);
   bool operator == (const pt& p) const {
        return eq(x, p.x) && eq(v, p.v);
};
const pt inf_pt = pt(1e18, 1e18);
struct QuadEdge {
   pt origin;
    OuadEdge* rot = nullptr;
   QuadEdge* onext = nullptr;
   bool used = false;
    OuadEdge* rev() const {
        return rot->rot;
   OuadEdge* lnext() const {
        return rot->rev()->onext->rot;
   QuadEdge* oprev() const {
        return rot->onext->rot;
    pt dest() const {
        return rev()->origin;
};
QuadEdge* make_edge(pt from, pt to) {
    OuadEdge* e1 = new OuadEdge;
    OuadEdge* e2 = new OuadEdge;
    OuadEdge* e3 = new OuadEdge;
    QuadEdge* e4 = new QuadEdge;
   e1->origin = from;
   e2->origin = to;
   e3->origin = e4->origin = inf pt;
   e1->rot = e3;
   e2 \rightarrow rot = e4;
   e3 \rightarrow rot = e2;
   e4->rot = e1;
   e1->onext = e1:
   e2 \rightarrow onext = e2;
    e3 - > onext = e4;
   e4->onext = e3:
    return e1:
```

```
void splice(OuadEdge* a, OuadEdge* b) {
    swap(a->onext->rot->onext, b->onext->rot->onext);
    swap(a->onext, b->onext);
void delete_edge(QuadEdge* e) {
    splice(e, e->oprev());
    splice(e->rev(), e->rev()->oprev());
    delete e->rot:
    delete e->rev()->rot;
    delete e:
    delete e->rev();
OuadEdge* connect(OuadEdge* a, OuadEdge* b) {
    QuadEdge* e = make_edge(a->dest(), b->origin);
    splice(e, a->lnext());
    splice(e->rev(), b);
    return e:
bool left_of(pt p, QuadEdge* e) {
    return gt(p.cross(e->origin, e->dest()), 0);
bool right_of(pt p, QuadEdge* e) {
    return lt(p.cross(e->origin, e->dest()), 0);
template <class T>
T det3(T a1, T a2, T a3, T b1, T b2, T b3, T c1, T c2, T c3) {
    return a1 * (b2 * c3 - c2 * b3) - a2 * (b1 * c3 - c1 * b3) +
           a3 * (b1 * c2 - c1 * b2);
bool in_circle(pt a, pt b, pt c, pt d) {
// If there is __int128, calculate directly.
// Otherwise, calculate angles.
#if defined( LP64 ) | | defined(WIN64)
    int128 det = -det3<__int128>(b.x, b.y, b.sqrLength(), c.x, c.y,
                                    c.sqrLength(), d.x, d.y, d.
                                        sqrLength());
    det += det3 < __int128 > (a.x, a.y, a.sqrLength(), c.x, c.y, c.
        sqrLength(), d.x,
                          d.v, d.sqrLength());
    det -= det3 < __int128 > (a.x, a.y, a.sqrLength(), b.x, b.y, b.
        sqrLength(), d.x,
                          d.v, d.sqrLength());
    det += det3 < __int128 > (a.x, a.y, a.sqrLength(), b.x, b.y, b.
        sqrLength(), c.x,
                          c.y, c.sqrLength());
    return det > 0:
#else
    auto ang = [](pt l, pt mid, pt r) {
        11 \times = mid.dot(1, r):
        11 y = mid.cross(1, r);
        long double res = atan2((long double)x, (long double)y);
        return res:
    long double kek = ang(a, b, c) + ang(c, d, a) - ang(b, c, d) - ang
```

```
(d, a, b);
    if (kek > 1e-8)
        return true;
    else
        return false;
#endif
pair<OuadEdge*, OuadEdge*> build tr(int 1, int r, vector<pt>& p) {
   if (r - 1 + 1 == 2) {
        OuadEdge* res = make_edge(p[1], p[r]);
        return make_pair(res, res->rev());
   if (r - 1 + 1 == 3) {
        QuadEdge \star a = make\_edge(p[1], p[1 + 1]), \star b = make\_edge(p[1 +
        splice(a->rev(), b);
        int sg = sgn(p[1].cross(p[1 + 1], p[r]));
        if (sq == 0)
            return make_pair(a, b->rev());
        QuadEdge* c = connect(b, a);
        if (sq == 1)
            return make_pair(a, b->rev());
            return make_pair(c->rev(), c);
    int mid = (1 + r) / 2;
    QuadEdge *ldo, *ldi, *rdo, *rdi;
    tie(ldo, ldi) = build_tr(l, mid, p);
   tie(rdi, rdo) = build_tr(mid + 1, r, p);
   while (true) {
        if (left_of(rdi->origin, ldi)) {
            ldi = ldi->lnext();
            continue;
        if (right_of(ldi->origin, rdi)) {
            rdi = rdi->rev()->onext;
            continue;
        break:
    QuadEdge* basel = connect(rdi->rev(), ldi);
    auto valid = [&basel](QuadEdge* e) { return right_of(e->dest(),
        basel); };
    if (ldi->origin == ldo->origin)
        ldo = basel->rev();
   if (rdi->origin == rdo->origin)
        rdo = basel;
    while (true) {
        QuadEdge* lcand = basel->rev()->onext;
        if (valid(lcand)) {
            while (in_circle(basel->dest(), basel->origin, lcand->dest
                (),
                             lcand->onext->dest())) {
                QuadEdge* t = lcand->onext;
                delete_edge(lcand);
                lcand = t;
        QuadEdge* rcand = basel->oprev();
        if (valid(rcand)) {
```

```
while (in_circle(basel->dest(), basel->origin, rcand->dest
                             rcand->oprev()->dest())) {
                QuadEdge* t = rcand->oprev();
                delete_edge(rcand);
                rcand = t;
        if (!valid(lcand) && !valid(rcand))
            break:
        if (!valid(lcand) ||
            (valid(rcand) && in_circle(lcand->dest(), lcand->origin,
                                       rcand->origin, rcand->dest())))
            basel = connect(rcand, basel->rev());
            basel = connect(basel->rev(), lcand->rev());
    return make_pair(ldo, rdo);
vector<tuple<pt, pt, pt>> delaunay(vector<pt> p) {
    sort(p.begin(), p.end(), [](const pt& a, const pt& b) {
        return lt(a.x, b.x) || (eq(a.x, b.x) && lt(a.y, b.y));
    auto res = build_tr(0, (int)p.size() - 1, p);
    QuadEdge* e = res.first;
    vector<OuadEdge*> edges = {e};
    while (lt(e->onext->dest().cross(e->dest(), e->origin), 0))
    auto add = [&p, &e, &edges]() {
        QuadEdge* curr = e;
        do {
            curr->used = true;
            p.push_back(curr->origin);
            edges.push_back(curr->rev());
            curr = curr->lnext();
        } while (curr != e);
    };
    add();
    p.clear();
    int kek = 0;
    while (kek < (int)edges.size()) {</pre>
        if (!(e = edges[kek++])->used)
            add();
    vector<tuple<pt, pt, pt>> ans;
    for (int i = 0; i < (int)p.size(); i += 3) {</pre>
        ans.push_back(make_tuple(p[i], p[i + 1], p[i + 2]));
    return ans;
```

## 5.9 Java Geometry Library

```
import java.util.*;
import java.io.*;
import java.awt.geom.*;
import java.lang.*;
//Lazy Geometry
class AWT{
```

```
static Area makeArea(double[] pts){
 Path2D.Double p = new Path2D.Double();
 p.moveTo(pts[0], pts[1]);
  for(int i = 2; i < pts.length; i+=2) {</pre>
    p.lineTo(pts[i], pts[i+1]);
 p.closePath();
 return new Area(p);
static double computePolygonArea(ArrayList<Point2D.Double> points) {
 Point2D.Double[] pts = points.toArray(new Point2D.Double[points.
      size()]);
 double area = 0;
 for (int i = 0; i < pts.length; i++) {</pre>
    int j = (i+1) % pts.length;
    area += pts[i].x * pts[j].y - pts[j].x * pts[i].y;
 return Math.abs(area)/2;
static double computeArea(Area area) {
 double totArea = 0;
 PathIterator iter = area.getPathIterator(null);
 ArrayList<Point2D.Double> points = new ArrayList<Point2D.Double>()
 while (!iter.isDone()) {
    double[] buffer = new double[6];
    switch (iter.currentSegment(buffer)) {
      case PathIterator.SEG_MOVETO:
      case PathIterator.SEG LINETO:
       points.add(new Point2D.Double(buffer[0], buffer[1]));
      case PathIterator.SEG CLOSE:
       totArea += computePolygonArea(points);
       points.clear();
       break;
    iter.next():
 return totArea;
```

# 6 String Algorithms

## 6.1 KMP

```
string p, t;
int b[ms], n, m;

void kmpPreprocess() {
  int i = 0, j = -1;
  b[0] = -1;
  while(i < m) {
    while(j >= 0 && p[i] != p[j]) j = b[j];
    b[++i] = ++j;
  }
}
```

```
void kmpSearch() {
  int i = 0, j = 0, ans = 0;
  while(i < n) {
    while(j >= 0 && t[i] != p[j]) j = b[j];
    i++; j++;
    if(j == m) {
        //ocorrencia aqui comecando em i - j
        ans++;
        j = b[j];
    }
} return ans;
}
```

## 6.2 KMP Automaton

```
const int limit =
vector<vector<int>> build automaton(string s) {
    s += '#'; //tem que ser diferente de todos os caracteres
    int n = (int) s.size();
    vector<vector<int>> ans(n, vector<int>(limit));
    vector<int> fail(n);
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < limit; j++) {</pre>
            if (i == 0) {
                if (s[i] == j + 'a') {
                    ans[i][j] = i + 1;
                     ans[i][j] = 0;
            } else {
                if (s[i] == j + 'a') {
                    ans[i][j] = i + 1;
                } else {
                    ans[i][j] = ans[fail[i - 1]][j];
        if (i == 0) {
            continue;
        int j = fail[i - 1];
        while (j > 0 \&\& s[i] != s[j]) {
            j = fail[j - 1];
        fail[i] = j + (s[i] == s[j]);
    return ans;
```

#### **6.3** Trie

```
int trie[ms][sigma], terminal[ms], z;

void init() {
  memset(trie[0], -1, sizeof trie[0]);
  z = 1;
```

```
int get_id(char c) {
  return c - 'a';
void insert(string &p) {
  int cur = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int id = get_id(p[i]);
    if(trie[cur][id] == -1) {
      memset(trie[z], -1, sizeof trie[z]);
      trie[cur][id] = z++;
    cur = trie[cur][id];
  terminal[cur]++;
int count(string &p) {
  int cur = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int id = get_id(p[i]);
    if(trie[cur][id] == -1) {
      return false;
    cur = trie[cur][id];
  return terminal[cur];
```

## 6.4 Aho-Corasick

```
// Construa a Trie do seu dicionario com o codigo acima
int fail[ms];
queue<int> q;
void buildFailure() {
  q.push(0);
  while(!q.empty()) {
    int node = q.front();
    q.pop();
    for(int pos = 0; pos < sigma; pos++) {</pre>
      int &v = trie[node][pos];
      int f = node == 0 ? 0 : trie[fail[node]][pos];
      if(v == -1) {
        v = f;
      } else {
        fail[v] = f;
        q.push(v);
        // juntar as informacoes da borda para o V ja q um match em V
            implica um match na borda
        terminal[v] += terminal[f];
int search(string &txt) {
```

```
int node = 0;
int ans = 0;
for(int i = 0; i < txt.length(); i++) {
   int pos = get_id(txt[i]);
   node = trie[node][pos];
   // processar informacoes no no atual
   ans += terminal[node];
}
return ans;
}</pre>
```

## 6.5 Algoritmo de Z

```
string s;
int fz[ms], n;

void zfunc() {
  fz[0] = n;
  for(int i = 1, l = 0, r = 0; i < n; i++) {
    fz[i] = max(0, min(r-i, fz[i-l]));
    while(s[fz[i]] == s[i+fz[i]]) ++fz[i];
    if(i + fz[i] > r) {
      l = i;
      r = i + fz[i];
    }
}
```

## 6.6 Suffix Array

```
vector<int> buildSa(const string& in) {
  int n = in.size(), c = 0;
  vector<int> temp(n), posBucket(n), bucket(n), bpos(n), out(n);
  for (int i = 0; i < n; i++) out[i] = i;</pre>
  sort(out.begin(), out.end(), [&](int a, int b) { return in[a] < in[b</pre>
      1; });
  for (int i = 0; i < n; i++) {
   bucket[i] = c;
    if (i + 1 == n || in[out[i]] != in[out[i + 1]]) c++;
  for (int h = 1; h < n && c < n; h <<= 1) {
    for (int i = 0; i < n; i++) posBucket[out[i]] = bucket[i];</pre>
    for (int i = n - 1; i >= 0; i--) bpos[bucket[i]] = i;
    for (int i = 0; i < n; i++) {
      if (out[i] >= n - h) temp[bpos[bucket[i]]++] = out[i];
    for (int i = 0; i < n; i++) {
      if (out[i] >= h) temp[bpos[posBucket[out[i] - h]]++] = out[i] -
          h;
    for (int i = 0; i + 1 < n; i++) {
        int a = (bucket[i] != bucket[i + 1]) || (temp[i] >= n - h)
            || (posBucket[temp[i + 1] + h] != posBucket[temp[i] + h]);
       bucket[i] = c;
        c += a;
    bucket[n-1] = c++;
```

```
temp.swap(out);
 return out;
vector<int> buildLcp(string s, vector<int> sa) {
 int n = (int) s.size();
  vector<int> pos(n), lcp(n, 0);
  for (int i = 0; i < n; i++) {
   pos[sa[i]] = i;
 int k = 0;
  for (int i = 0; i < n; i++) {
   if (pos[i] + 1 == n) {
      k = 0;
      continue;
   int j = sa[pos[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
   lcp[pos[i]] = k;
   k = max(k - 1, 0);
  return lcp;
```

#### 6.7 Suffix Tree

```
//by adamant
#define fpos adla
const int inf = 1e9;
const int maxn = 1e4;
char s[maxn];
map<int, int> to[maxn];
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
int sz = 1, n = 0;
int make_node(int _pos, int _len)
    fpos[sz] = _pos;
    len [sz] = _len;
    return sz++;
void go_edge()
    while(pos > len[to[node][s[n - pos]]])
        node = to[node][s[n - pos]];
        pos -= len[node];
void add_letter(int c)
    s[n++] = c;
    pos++;
    int last = 0;
    while(pos > 0)
```

```
go_edge();
        int edge = s[n - pos];
        int &v = to[node][edge];
        int t = s[fpos[v] + pos - 1];
        if(v == 0)
            v = make_node(n - pos, inf);
            link[last] = node;
            last = 0;
        else if(t == c)
            link[last] = node;
            return;
        else
            int u = make_node(fpos[v], pos - 1);
            to[u][c] = make_node(n - 1, inf);
            to[u][t] = v;
            fpos[v] += pos - 1;
            len [v] -= pos - 1;
            v = u;
            link[last] = u;
            last = u;
        if(node == 0)
            pos--;
        else
            node = link[node];
//len[0] = inf
```

#### 6.8 Suffix Automaton

```
int len[ms*2], link[ms*2], nxt[ms*2][sigma];
int sz, last;
void build(string &s) {
  len[0] = 0; link[0] = -1;
  sz = 1; last = 0;
  memset (nxt[0], -1, sizeof nxt[0]);
  for(char ch : s) {
    int c = ch-'a', cur = sz++;
    len[cur] = len[last]+1;
    memset(nxt[cur], -1, sizeof nxt[cur]);
    int p = last;
    while (p != -1 \&\& nxt[p][c] == -1) {
      nxt[p][c] = cur; p = link[p];
    if(p == -1) {
     link[cur] = 0;
    } else {
      int q = nxt[p][c];
      if(len[p] + 1 == len[q]) {
       link[cur] = q;
      } else {
```

```
len[sz] = len[p]+1; link[sz] = link[q];
    memcpy(nxt[sz], nxt[q], sizeof nxt[q]);
    while (p != -1 && nxt[p][c] == q) {
        nxt[p][c] = sz; p = link[p];
        }
        link[q] = link[cur] = sz++;
    }
    }
    last = cur;
}
```

## 7 Miscellaneous

## 7.1 LIS - Longest Increasing Subsequence

```
int arr[ms], lisArr[ms], n;
// int bef[ms], pos[ms];
int lis() {
  int len = 1;
  lisArr[0] = arr[0];
  // bef[0] = -1;
  for(int i = 1; i < n; i++) {
    // upper_bound se non-decreasing
    int x = lower_bound(lisArr, lisArr + len, arr[i]) - lisArr;
    len = max(len, x + 1);
    lisArr[x] = arr[i];
    // pos[x] = i;
    // bef[i] = x ? pos[x-1] : -1;
  return len;
vi getLis() {
 int len = lis();
  for(int i = pos[lisArr[len - 1]]; i >= 0; i = bef[i]) {
    ans.push_back(arr[i]);
  reverse(ans.begin(), ans.end());
  return ans;
```

## 7.2 Ternary Search

```
// R
for(int i = 0; i < LOG; i++) {
  long double m1 = (A * 2 + B) / 3.0;
  long double m2 = (A + 2 * B) / 3.0;

if(f(m1) > f(m2))
    A = m1;
  else
    B = m2;
}
ans = f(A);
```

```
// Z
while (B - A > 4) {
  int m1 = (A + B) / 2;
  int m2 = (A + B) / 2 + 1;
  if (f (m1) > f (m2))
    A = m1;
  else
    B = m2;
}
ans = inf;
for (int i = A; i <= B; i++) ans = min (ans , f(i));</pre>
```

#### 7.3 Count Sort

```
int H[(1 << 15) +1], to [mx], b[mx];
void sort(int m, int a[]) {
 memset(H, 0, sizeof H);
  for (int i = 1; i <= m; i++) {</pre>
   H[a[i] % (1 << 15)] ++;
  for (int i = 1; i < 1<<15; i++) {
   H[i] += H[i-1];
  for (int i = m; i; i--) {
    to[i] = H[a[i] % (1 << 15)] --;
  for (int i = 1; i <= m; i++) {
    b[to[i]] = a[i];
  memset(H, 0, sizeof H);
  for (int i = 1; i <= m; i++) {
   H[b[i] >> 15] ++;
  for (int i = 1; i < 1<<15; i++) {
   H[i] += H[i-1];
  for (int i = m; i ; i--) {
    to[i] = H[b[i] >> 15] --;
  for (int i = 1; i <= m; i++) {
    a[to[i]] = b[i];
```

## 7.4 Random Number Generator

#### 7.5 Rectangle Hash

```
namespace {
  struct safe_hash {
    static uint64_t splitmix64(uint64_t x) {
      // http://xorshift.di.unimi.it/splitmix64.c
      x += 0x9e3779b97f4a7c15;
      x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
      x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
      return x ^ (x >> 31);
    size_t operator()(uint64_t x) const {
      static const uint64_t FIXED_RANDOM = std::chrono::steady_clock::
          now().time_since_epoch().count();
      return splitmix64(x + FIXED_RANDOM);
  };
struct rect {
  int x1, y1, x2, y2; // x1 < x2, y1 < y2
  rect () {};
  rect (int x1, int y1, int x2, int y2) : x1(x1), x2(x2), y1(y1), y2(
      y2) {};
  rect inter (rect other) {
    int x3 = max(x1, other.x1);
    int y3 = max(y1, other.y1);
    int x4 = min(x2, other.x2);
    int y4 = min(y2, other.y2);
    return rect (x3, y3, x4, y4);
  uint64_t get_hash() {
    safe_hash sh;
    uint64_t ret = sh(x1);
    ret ^= sh(ret ^ y1);
    ret ^= sh(ret ^ x2);
    ret ^= sh(ret ^ y2);
    return ret;
};
```

## 7.6 Unordered Map Tricks

```
// pair<int, int> hash function
struct HASH{
    size_t operator() (const pair<int,int>&x) const{
        return (size_t) x.first * 37U + (size_t) x.second;
    }
};
unordered_map<int,int>mp;
mp.reserve(1024);
mp.max_load_factor(0.25);
```

#### 7.7 Submask Enumeration

```
for (int s=m; ; s=(s-1)&m) {
   ... you can use s ...
   if (s==0) break;
}
```

#### 7.8 Sum over Subsets DP

```
// F[i] = Sum of all A[j] where j is a submask of i
for(int i = 0; i < (1 < N); ++i)
  F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1 < N); ++mask) {
  if(mask & (1 < i))
    F[mask] += F[mask^(1 < i)];
}</pre>
```

## 7.9 Java Fast I/O

```
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.PrintWriter;
import java.util.Arrays;
import java.util.Random;
import java.io.IOException;
import java.io.InputStreamReader;
import java.util.StringTokenizer;
import java.io.BufferedReader;
import java.io.InputStream;
import java.util.*;
import java.io.*;
// src petr
public class Main {
 public static void main(String[] args) {
    InputStream inputStream = System.in;
    OutputStream outputStream = System.out;
    InputReader in = new InputReader(inputStream);
    PrintWriter out = new PrintWriter(outputStream);
    TaskA solver = new TaskA();
    solver.solve(1, in, out);
    out.close();
  static class TaskA {
    public void solve(int testNumber, InputReader in, PrintWriter out)
  static class InputReader {
    public BufferedReader reader;
    public StringTokenizer tokenizer;
    public InputReader(InputStream stream) {
```

#### **7.10** Dates

```
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"
    };
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
  return
   1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
   3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
   d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
  n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x = 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
  x = j / 11;
 m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
```

```
import java.util.*;
import java.util.regex.*;

public class Main {
   public static String BuildRegex () {
      return "^" + sentence + "$";
   }

   public static void main (String args[]) {
      String regex = BuildRegex();
      // check pattern documentation
      Pattern pattern = Pattern.compile (regex);
      Scanner s = new Scanner(System.in);
      String sentence = s.nextLine().trim();
      boolean found = pattern.matcher(sentence).find()
   }
}
```

## 7.12 Lat Long

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
struct 11
  double r, lat, lon;
struct rect
  double x, y, z;
11 convert(rect& P)
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
  return 0;
rect convert(ll& 0)
 rect P;
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.z = Q.r*sin(Q.lat*M_PI/180);
  return P;
```

## 8 Teoremas e formulas uteis

## 8.1 Grafos

```
Formula de Euler: V - E + F = 2 (para grafo planar)
Handshaking: Numero par de vertices tem grau impar
Kirchhoff's Theorem: Monta matriz onde Mi,i = Grau[i] e Mi,j = -1 se
   houver aresta i-j ou 0 caso contrario, remove uma linha e uma
   coluna qualquer e o numero de spanning trees nesse grafo eh o det
   da matriz

Grafo contem caminho hamiltoniano se:
Dirac's theorem: Se o grau de cada vertice for pelo menos n/2
Ore's theorem: Se a soma dos graus que cada par nao-adjacente de
   vertices for pelo menos n
Trees:
Tem Catalan(N) Binary trees de N vertices
```

Tem Catalan(N) Binary trees de N Vertices

Tem Catalan(N-1) Arvores enraizadas com N vertices

Caley Formula: n^(n-2) arvores em N vertices com label

Prufer code: Cada etapa voce remove a folha com menor label e o label

do vizinho eh adicionado ao codigo ate ter 2 vertices

#### Flow:

Max Edge-disjoint paths: Max flow com arestas com peso 1
Max Node-disjoint paths: Faz a mesma coisa mas separa cada vertice em
um com as arestas de chegadas e um com as arestas de saida e uma
aresta de peso 1 conectando o vertice com aresta de chegada com
ele mesmo com arestas de saida

Konig's Theorem: minimum node cover = maximum matching se o grafo for bipartido, complemento eh o maximum independent set

Min Node disjoint path cover: formar grafo bipartido de vertices duplicados, onde aresta sai do vertice tipo A e chega em tipo B, entao o path cover eh N - matching

Min General path cover: Mesma coisa mas colocando arestas de A pra B sempre que houver caminho de A pra B

Dilworth's Theorem: Min General Path cover = Max Antichain (set de vertices tal que nao existe caminho no grafo entre vertices desse set)

Hall's marriage: um grafo tem um matching completo do lado X se para cada subconjunto W de X,

|W| <= |vizinhosW| onde |W| eh quantos vertices tem em W

### 8.2 Math

Goldbach's: todo numero par n > 2 pode ser representado com n = a + b onde a e b sao primos

Twin prime: existem infinitos pares p, p + 2 onde ambos sao primos

Legendre's: sempre tem um primo entre n^2 e (n+1)^2

Lagrange's: todo numero inteiro pode ser inscrito como a soma de 4 quadrados

Zeckendorf's: todo numero pode ser representado pela soma de dois numeros de fibonnacis diferentes e nao consecutivos

Euclid's: toda tripla de pitagoras primitiva pode ser gerada com (n^2 - m^2, 2nm, n^2+m^2) onde n, m sao coprimos e um deles eh par

Wilson's: n eh primo quando (n-1)! mod n = n - 1

Mcnugget: Para dois coprimos x, y o maior inteiro que nao pode ser escrito como ax + by eh (x-1)(y-1)/2

Fermat: Se p eh primo entao a^(p-1) % p = 1 Se x e m tambem forem coprimos entao x^k % m = x^(k mod(m-1)) % m Euler's theorem: x^(phi(m)) mod m = 1 onde phi(m) eh o totiente de euler

```
Para equacoes no formato x = a1 mod m1, ..., x = an mod mn onde todos
      os pares m1, ..., mn sao coprimos
Deixe Xk = m1 * m2 * ... * mn/mk e Xk^-1 mod mk = inverso de Xk mod mk, entao
x = somatorio com k de 1 ate n de ak*Xk*(Xk,mk^-1 mod mk)
Para achar outra solucao so somar m1*m2*..*mn a solucao existente
Catalan number: exemplo expressoes de parenteses bem formadas
C0 = 1, Cn = somatorio de <math>i=0 \rightarrow n-1 de Ci \times C(n-1+1)
outra forma: Cn = (2n \text{ escolhe } n)/(n+1)
Bertrand's ballot theorem: p votos tipo A e q votos tipo B com p>q,
    prob de em todo ponto ter mais As do que Bs antes dele = (p-q)/(p+
Se puder empates entao prob = (p+1-q)/(p+1), para achar quantidade de
    possibilidades nos dois casos basta multiplicar por (p + q escolhe
Propriedades de Coeficientes Binomiais:
Somatorio de k = 0 -> m de (-1)^k * (n escolhe k) = (-1)^m * (n -1)
    escolhe m)
(N \text{ escolhe } K) = (N \text{ escolhe } N-K)
(N \text{ escolhe } K) = N/K * (n-1 \text{ escolhe } k-1)
Somatorio de k = 0 \rightarrow n de (n escolhe k) = 2^n
Somatorio de m = 0 \rightarrow n de (m = scolhe k) = (n+1 = scolhe k + 1)
Somatorio de k = 0 \rightarrow m de (n+k \text{ escolhe } k) = (n+m+1 \text{ escolhe } m)
Somatorio de k = 0 \rightarrow n de (n \text{ escolhe } k)^2 = (2n \text{ escolhe } n)
Somatorio de k = 0 ou 1 \rightarrow n de k \times (n \text{ escolhe } k) = n \times 2^n (n-1)
Somatorio de k = 0 \rightarrow n de (n-k \text{ escolhe } k) = \text{Fib}(n+1)
Hockey-stick: Somatorio de i = r \rightarrow n de (i escolhe r) = (n + 1)
    escolhe r + 1)
Vandermonde: (m+n \text{ escolhe r}) = \text{somatorio de } k = 0 \rightarrow r \text{ de } (m \text{ escolhe } k
    ) \star (n escolhe r - k)
Burnside lemma: colares diferentes nao contando rotacoes quando m =
    cores e n = comprimento
(m^n + somatorio i = 1 - > n-1 de m^qcd(i, n))/n
Distribuicao uniforme a,a+1, ..., b Expected[X] = (a+b)/2
Distribuicao binomial com n tentativas de probabilidade p, X =
    sucessos:
    P(X = x) = p^x * (1-p)^(n-x) * (n escolhe x) e E[X] = p*n
Distribuicao geometrica onde continuamos ate ter sucesso, X =
    tentativas:
    P(X = x) = (1-p)^(x-1) * p e E[X] = 1/p
Linearity of expectation: Tendo duas variaveis X e Y e constantes a e
    b, o valor esperado de aX + bY = a*E[X] + b*E[X]
V(X) = E((X-u)^2)
V(X) = E(X^2) - E(X^2)
PG: a1 * (q^n - 1)/(q - 1)
```

## 8.3 Geometry

Chinese remainder theorem:

Formula de Euler: V - E + F = 2Pick Theorem: Para achar pontos em coords inteiras num poligono Area = i + b/2 - 1 onde i eh o o numero de pontos dentro do poligono e b de pontos no perimetro do poligono Two ears theorem: Todo poligono simples com mais de 3 vertices tem

pelo menos 2 orelhas, vertices que podem ser removidos sem criar

um crossing, remover orelhas repetidamente triangula o poligono Incentro triangulo: (a(Xa, Ya) + b(Xb, Yb) + c(Xc, Yc))/(a+b+c) onde a = lado oposto ao vertice a, incentro eh onde cruzam as bissetrizes, eh o centro da circunferencia inscrita e eh equidistante aos lados

Delaunay Triangulation: Triangulacao onde nenhum ponto esta dentro de nenhum circulo circunscrito nos triangulos

Eh uma triangulacao que maximiza o menor angulo e a MST euclidiana de um conjunto de pontos eh um subconjunto da triangulacao

Brahmagupta s formula: Area cyclic quadrilateral s = (a+b+c+d)/2 area = sqrt((s-a)\*(s-b)\*(s-c)\*(s-d))  $d = 0 \Rightarrow area = sqrt((s-a)*(s-b)*(s-c)*s)$ 

#### 8.4 Mersenne's Primes

Primos de Mersenne 2^n - 1
Lista de Ns que resultam nos primeiros 41 primos de Mersenne:
2; 3; 5; 7; 13; 17; 19; 31; 61; 89; 107; 127; 521; 607; 1.279; 2.203;
2.281; 3.217; 4.253; 4.423; 9.689; 9.941; 11.213; 19.937; 21.701;
23.209; 44.497; 86.243; 110.503; 132.049; 216.091; 756.839;
859.433; 1.257.787; 1.398.269; 2.976.221; 3.021.377; 6.972.593;
13.466.917; 20.996.011; 24.036.583;