# The AC is a lie - ICPC Library

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# 1 Data Structures

# 1.1 BIT 2D Comprimida

```
// src: tfg50
template<class T = int>
struct Bit2D {
public:
  Bit2D(vector<pair<T, T>> pts) {
    sort(pts.begin(), pts.end());
    for(auto a : pts) {
      if(ord.empty() || a.first != ord.back()) {
        ord.push_back(a.first);
    fw.resize(ord.size() + 1);
    coord.resize(fw.size());
    for(auto &a : pts) {
      swap(a.first, a.second);
    sort(pts.begin(), pts.end());
    for(auto &a : pts) {
      swap(a.first, a.second);
      for(int on = upper_bound(ord.begin(), ord.end(), a.first) - ord.
          begin(); on < \text{fw.size}(); on += \text{on } \& -\text{on}) {
        if(coord[on].empty() || coord[on].back() != a.second) {
          coord[on].push_back(a.second);
    for(int i = 0; i < fw.size(); i++) {</pre>
      fw[i].assign(coord[i].size() + 1, 0);
  void upd(T x, T y, T v) {
    for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin();
         xx < fw.size(); xx += xx & -xx) {
```

```
for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y)
           - \operatorname{coord}[xx].\operatorname{begin}(); yy < \operatorname{fw}[xx].\operatorname{size}(); yy += yy & -yy) {
         fw[xx][yy] += v;
  T qry(T x, T y) {
    T ans = 0;
    for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin();
          xx > 0; xx -= xx & -xx) {
       for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y)
           - coord[xx].begin(); yy > 0; yy -= yy & -yy) {
         ans += fw[xx][yy];
      }
    return ans;
  T \operatorname{qry}(T \times 1, T \times 1, T \times 2, T \times 2) {
    return qry(x2, y2) - qry(x2, y1 - 1) - qry(x1 - 1, y2) + qry(x1 -
         1, y1 - 1);
  void upd(T x1, T y1, T x2, T y2, T v) { // !insert these points
    upd(x1, y1, v);
    upd(x1, y2 + 1, -v);
    upd(x2 + 1, y1, -v);
    upd(x2 + 1, y2 + 1, v);
private:
  vector<T> ord;
  vector<vector<T>> fw, coord;
```

# 1.2 Iterative Segment Tree

```
for (l += n, r += n+1; l < r; l >>= 1, r >>= 1) {
   if (l&1) resl = combine(resl, t[l++]);
   if (r&1) resr = combine(t[--r], resr);
}
return combine(resl, resr);
```

# 1.3 Iterative Segment Tree with Lazy Propagation

```
struct LazyContext {
  int v;
 LazyContext(int v = 0) : v(v) { }
  void reset() {
  v = 0;
  void operator += (LazyContext o) {
 v += o.v;
};
struct Node
  int sz, v;
 Node() { // neutral element
        v = 0; sz = 0;
  Node(int i) { // init
        v = i; sz = 1;
  Node (Node &1, Node &r) { // merge
        sz = 1.sz + r.sz;
        v = 1.v + r.v;
  void apply(LazyContext lazy) {
  v += lazy.v * sz;
};
Node tree[2*ms];
LazyContext lazy[ms];
bool dirty[ms];
int n, h, a[ms];
void init() {
    h = 0;
    while ((1 << h) < n) h++;
    for (int i = 0; i < n; i++) {
        tree[i + n] = Node(a[i]);
    for (int i = n - 1; i > 0; i--) {
        tree[i] = Node(tree[i + i], tree[i + i + 1]);
        lazv[i].reset();
        dirty[i] = 0;
```

```
void apply(int p, LazyContext &lc) {
    tree[p].apply(lc);
    if(p < n) {
        dirty[p] = true;
        lazy[p] += lc;
void push(int p) {
    for (int s = h; s > 0; s--) {
        int i = p \gg s;
        if(dirty[i]) {
            apply(i + i, lazy[i]);
            apply(i + i + 1, lazy[i]);
            lazy[i].reset();
            dirty[i] = false;
void build(int p) {
    for (p /= 2; p > 0; p /= 2) {
        tree[p] = Node(tree[p + p], tree[p + p + 1]);
        if(dirty[p]) {
            tree(p).apply(lazy(p));
Node query(int 1, int r) {
    if(1 > r) return Node();
    1 += n, r += n+1;
    push(1);
    push(r - 1);
    Node lp, rp:
    for(; 1 < r; 1 /= 2, r /= 2) {
        if(1 & 1) lp = Node(lp, tree[l++]);
        if(r \& 1) rp = Node(tree[--r], rp);
    return Node(lp, rp);
void update(int 1, int r, LazyContext lc) {
    if(l > r) return;
    1 += n, r += n+1;
    push(1);
    push(r - 1);
    int 10 = 1, r0 = r;
    for (; 1 < r; 1 /= 2, r /= 2) {
        if(l & 1) apply(l++, lc);
        if(r & 1) apply(--r, lc);
    build(10);
    build(r0 - 1);
```

```
int arr[ms], seg[4 * ms], lazy[4 * ms], n;
void build (int idx = 0, int l = 0, int r = n-1) {
  int mid = (1+r)/2;
  lazy[idx] = 0;
  if(1 == r) {
    seg[idx] = arr[l];
    return;
  build(2*idx+1, 1, mid); build(2*idx+2, mid+1, r);
  seq[idx] = seq[2*idx+1] + seq[2*idx+2]; // Merge
void apply(int idx, int 1, int r) {
  int mid = (1+r)/2;
 if(lazy[idx] && !canBreak) { // if not beats canBreak = false
    if(1 < r) {
      lazy[2*idx+1] += lazy[idx]; // Merge de lazy
      lazv[2*idx+2] += lazv[idx]; // Merge de lazv
    if(canApply) { // if not beats canApply = true
      seq[idx] += lazv[idx] * (r - 1 + 1); // Aplicar lazv no seq
    } else {
      apply (2*idx+1, 1, mid); apply (2*idx+2, mid+1, r);
      seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
  lazy[idx] = 0; // Limpar a lazy
int query(int L, int R, int idx = 0, int l = 0, int r = n-1) {
  int mid = (1+r)/2;
  apply(idx, l, r);
  if(l > R | | r < L) return 0; // Valor que nao atrapalhe</pre>
  if(L <= 1 && r <= R) return seq[idx];</pre>
  return query(L, R, 2*idx+1, 1, mid) + query(L, R, 2*idx+2, mid+1, r)
      : // Merge
void update(int L, int R, int V, int idx = 0, int 1 = 0, int r = n-1)
  int mid = (1+r)/2;
  apply(idx, l, r);
  if(1 > R | | r < L) return;
  if(L <= 1 && r <= R) {
   lazy[idx] = V;
    apply(idx, l, r);
    return;
  update(L, R, V, 2*idx+1, 1, mid); update(L, R, V, 2*idx+2, mid+1, r)
  seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
```

# 1.5 Persistent Segment Tree

```
struct PSEGTREE{
  private:
   int z, t, sz, *tree, *L, *R, head[112345];
```

```
void _build(int 1, int r, int on, vector<int> &v) {
      if(| == r){
        tree[on] = v[1];
        return;
      L[on] = ++z;
      int mid = (1+r) >> 1;
      _build(l, mid, L[on], v);
      R[on] = ++z;
      _build(mid+1, r, R[on], v);
      tree[on] = tree[L[on]] + tree[R[on]];
    int _upd(int ql, int qr, int val, int l, int r, int on) {
      if(1 > qr \mid | r < ql) return on;
      int curr = ++z:
      if(1 >= ql && r <= qr) {
       tree[curr] = tree[on] + val;
        return curr;
      int mid = (1+r) >> 1;
      L[curr] = \_upd(ql, qr, val, l, mid, L[on]);
      R[curr] = _upd(ql, qr, val, mid+1, r, R[on]);
      tree[curr] = tree[L[curr]] + tree[R[curr]];
      return curr;
    int _query(int ql, int qr, int l, int r, int on) {
      if(1 > qr \mid | r < ql) return 0;
      if(1 >= q1 \&\& r <= qr) {
       return tree[on];
      int mid = (1+r) >> 1;
      return _query(ql, qr, l, mid, L[on]) + _query(ql, qr, mid+1, r,
          R[on]);
 public:
   PSEGTREE (vector<int> &v) {
      tree = new int[1123456]:
      L = new int[1123456];
      R = new int[1123456];
     build(v):
   void build(vector<int> &v) {
     t = 0, z = 0;
      sz = v.size();
     head[0] = 0;
      build(0, sz-1, 0, v);
   void upd(int pos, int val, int idx){
      head[++t] = _upd(pos, pos, val, 0, sz-1, head[idx]);
   int query(int 1, int r, int idx){
      return _query(1, r, 0, sz-1, head[idx]);
};
```

## 1.6 Treap

```
mt19937 rng ((int) chrono::steady_clock::now().time_since_epoch().
    count());
typedef int Value;
typedef struct item * pitem;
struct item {
 item () {}
 item (Value v) { // add key if not implicit
   value = v:
   prio = uniform int distribution<int>() (rng);
   cnt = 1;
   rev = 0;
   1 = r = 0;
  pitem 1, r;
 Value value;
  int prio, cnt;
 bool rev;
int cnt (pitem it) {
 return it ? it->cnt : 0;
void fix (pitem it) {
 if (it)
    it->cnt = cnt(it->1) + cnt(it->r) + 1;
void pushLazy (pitem it) {
  if (it && it->rev) {
   it->rev = false;
   swap(it->1, it->r):
   if (it->1) it->1->rev ^= true;
   if (it->r) it->r->rev ^= true;
void merge (pitem & t, pitem l, pitem r) {
  pushLazy (1); pushLazy (r);
  if (!1 || !r) t = 1 ? 1 : r;
  else if (l->prio > r->prio)
   merge (1->r, 1->r, r), t = 1;
   merge (r->1, 1, r->1), t = r;
  fix (t):
void split (pitem t, pitem & l, pitem & r, int key) {
  if (!t) return void( l = r = 0 );
  pushLazy (t);
  int cur_key = cnt(t->1); // t->key if not implicit
  if (key <= cur_key)</pre>
   split (t->1, 1, t->1, key), r = t;
    split (t->r, t->r, r, key - (1 + cnt(t->1))), 1 = t;
  fix (t);
```

## 1.7 Persistent Treap

```
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count())
typedef int Key;
struct Treap {
  Treap(){}
  Treap(char k) {
    key = 1;
    size = 1;
    1 = r = NULL;
    val = k;
  Treap *1, *r;
  Key key;
  char val;
  int size;
};
typedef Treap * PTreap;
bool leftSide(PTreap 1, PTreap r) {
  return (int) (rnq() % (l->size + r->size)) < l->size;
void fix(PTreap t) {
  if (t == NULL) {
    return;
  t \rightarrow size = 1:
  t->key = 1;
  if (t->1) {
    t->size += t->l->size;
    t->kev += t->l->size;
  if (t->r) {
    t \rightarrow size += t \rightarrow r \rightarrow size;
```

```
void split(PTreap t, Key key, PTreap &1, PTreap &r) {
  if (t == NULL) {
   1 = r = NULL;
  } else if (t->key <= key) {</pre>
   1 = new Treap();
    *1 = *t;
    split(t->r, key - t->key, l->r, r);
    fix(1);
  } else {
   r = new Treap();
    *r = *t;
    split(t->1, key, l, r->l);
    fix(r);
void merge(PTreap &t, PTreap 1, PTreap r) {
  if (!l || !r) {
   t = 1 ? 1 : r;
    return;
  t = new Treap();
  if (leftSide(l, r)) {
    *t = *1;
    merge(t->r, l->r, r);
    *t = *r;
    merge(t->1, 1, r->1);
  fix(t);
vector<PTreap> ver = {NULL};
PTreap build(int 1, int r, string& s) {
 if (1 >= r) return NULL;
  int mid = (1 + r) >> 1;
  auto ans = new Treap(s[mid]);
  ans->1 = build(1, mid, s);
 ans->r = build(mid + 1, r, s);
  fix(ans);
  return ans;
int last = 0;
void go(PTreap t, int f) {
 if (!t) return;
 qo(t->1, f);
  cout << t->val;
  last += (t->val <math>== 'c') * f;
  go(t->r, f);
void insert(PTreap t, int pos, string& s) {
 PTreap 1, r;
 split(t, pos + 1, l, r);
 PTreap mid = build(0, s.size(), s);
```

```
merge(mid, l, mid);
merge(mid, mid, r);
ver.push_back(mid);
}

void erase(PTreap t, int L, int R) {
  PTreap l, mid, r;
  split(t, L, l, mid);
  split(mid, R - L + l, mid, r);
  merge(l, l, r);
  ver.push_back(l);
}
```

#### 1.8 KD-Tree

```
int d;
long long getValue(const PT &a) {return (d & 1) == 0 ? a.x : a.y; }
bool comp (const PT &a, const PT &b) {
  if((d & 1) == 0) { return a.x < b.x; }</pre>
  else { return a.y < b.y; }</pre>
long long sqrDist(PT a, PT b) { return (a - b) * (a - b); }
class KD_Tree {
public:
  struct Node {
    PT point:
    Node *left, *right;
  void init(std::vector<PT> pts) {
    if(pts.size() == 0) {
      return;
    int n = 0;
   tree.resize(2 * pts.size());
    build(pts.begin(), pts.end(), n);
    //assert(n <= (int) tree.size());</pre>
  long long nearestNeighbor(PT point) {
    // assert(tree.size() > 0);
    long long ans = (long long) le18;
    nearestNeighbor(&tree[0], point, 0, ans);
    return ans;
private:
  std::vector<Node> tree;
  Node* build(std::vector<PT>::iterator l, std::vector<PT>::iterator r
      , int &n, int h = 0) {
    int id = n++;
    if(r - 1 == 1) {
      tree[id].left = tree[id].right = NULL;
      tree[id].point = *1;
    \} else if (r - 1 > 1) {
      std::vectorPT>::iterator mid = 1 + ((r - 1) / 2);
      std::nth_element(l, mid - 1, r, comp);
      tree[id].point = *(mid - 1);
```

```
// BE CAREFUL!
      // DO EVERYTHING BEFORE BUILDING THE LOWER PART!
      tree[id].left = build(l, mid, n, h^1);
      tree[id].right = build(mid, r, n, h^1);
    return &tree[id];
  void nearestNeighbor(Node* node, PT point, int h, long long &ans) {
    if(!node) {
      return;
    if(point != node->point) {
      // THIS WAS FOR A PROBLEM
      // THAT YOU DON'T CONSIDER THE DISTANCE TO ITSELF!
      ans = std::min(ans, sqrDist(point, node->point));
    d = h;
    long long delta = getValue(point) - getValue(node->point);
    if(delta <= 0) {
      nearestNeighbor(node->left, point, h^1, ans);
      if(ans > delta * delta) {
        nearestNeighbor(node->right, point, h^1, ans);
    } else {
      nearestNeighbor(node->right, point, h^1, ans);
      if(ans > delta * delta)
        nearestNeighbor(node->left, point, h^1, ans);
};
```

## 1.9 Sparse Table

```
template < class Info_t>
class SparseTable {
private:
  vector<int> log2;
  vector<vector<Info t>> table;
  Info_t merge(Info_t &a, Info_t &b) {
  }
  SparseTable(int n, vector<Info_t> v) {
    log2.resize(n + 1);
    log2[1] = 0;
    for (int i = 2; i <= n; i++) {
      log2[i] = log2[i >> 1] + 1;
    table.resize(n + 1):
    for (int i = 0; i < n; i++) {
      table[i].resize(log2[n] + 1);
    for (int i = 0; i < n; i++) {
      table[i][0] = v[i];
    for (int i = 0; i < log2[n]; i++) {
```

```
for (int j = 0; j < n; j++) {
    if (j + (1 << i) >= n) break;
    table[j][i + 1] = merge(table[j][i], table[j + (1 << i)][i]);
    }
}
int get(int l, int r) {
    int k = log2[r - 1 + 1];
    return merge(table[l][k], table[r - (1 << k) + 1][k]);
}
};</pre>
```

## 1.10 Max Queue

```
// src: tfq50
template <class T, class C = std::less<T>>
struct MaxQueue {
 MaxOueue() {
    clear();
  void clear() {
   id = 0:
    q.clear();
  void push(T x) {
    std::pair<int, T> nxt(1, x);
    while(q.size() > id && cmp(q.back().second, x)) {
      nxt.first += q.back().first;
      q.pop_back();
    q.push_back(nxt);
  T qry() {
    return q[id].second;
  void pop() {
    q[id].first--;
    if(q[id].first == 0) {
      id++:
private:
  std::vector<std::pair<int, T>> q;
 int id;
 C cmp;
};
```

# 1.11 Policy Based Structures

```
using namespace __gnu_pbds;

typedef tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;

ordered_set X;
X.insert(1);
X.find_by_order(0);
X.order_of_key(-5);
end(X), begin(X);
```

## 1.12 Color Updates Structure

```
struct range {
  int 1, r;
  int v;
  range(int l = 0, int r = 0, int v = 0) : l(l), r(r), v(v) {}
  bool operator < (const range &a) const {
    return 1 < a.1;</pre>
};
set<range> ranges;
vector<range> update(int 1, int r, int v) { // [1, r)
 vector<range> ans:
  if(1 >= r) return ans;
  auto it = ranges.lower_bound(1);
  if(it != ranges.begin()) {
    it--;
    if(it->r>1) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, 1, cur.v));
      ranges.insert(range(l, cur.r, cur.v));
  it = ranges.lower bound(r);
  if(it != ranges.begin()) {
    it--;
    if(it->r>r) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.l, r, cur.v));
      ranges.insert(range(r, cur.r, cur.v));
  for(it = ranges.lower_bound(1); it != ranges.end() && it->1 < r; it</pre>
      ++) {
    ans.push_back(*it);
  ranges.erase(ranges.lower_bound(1), ranges.lower_bound(r));
  ranges.insert(range(l, r, v));
  return ans:
int query(int v) { // Substituir -1 por flag para quando nao houver
    resposta
```

```
auto it = ranges.upper_bound(v);
if(it == ranges.begin()) {
    return -1;
}
it--;
return it->r >= v ? it->v : -1;
}
```

# 2 Dynamic Programming

## 2.1 Dynamic Hull

```
typedef long double double_t;
typedef long long int 11;
class HullDynamic {
public:
  const double t inf = 1e9;
  struct Line {
    11 m, b;
    double_t start;
    bool is_query;
    Line() {}
    Line(ll _m, ll _b, double_t _start, bool _is_query) : m(_m), b(_b)
        , start(_start), is_query(_is_query) {}
    11 eval(11 x) {
      return m * x + b;
    double t intersect(const Line& 1) const {
      return (double_t) (1.b - b) / (m - 1.m);
    bool operator< (const Line& 1) const {</pre>
      if (is_query == 0) return m > 1.m;
      return (start < 1.start);</pre>
  };
  typedef set<Line>::iterator iterator_t;
  bool has_prev(iterator_t it) {
    return (it != hull.begin());
  bool has_next(iterator_t it) {
    return (++it != hull.end());
  bool irrelevant(iterator_t it) {
    if (!has_prev(it) || !has_next(it)) return 0;
    iterator_t prev = it, next = it;
    prev--;
    next++;
```

```
return next->intersect(*prev) <= it->intersect(*prev);
void update_left(iterator_t it) {
  if (it == hull.begin()) return;
  iterator t pos = it;
  --it;
  vector<Line> rem;
  while(has prev(it)) {
    iterator_t prev = it;
    if (prev->intersect(*pos) <= prev->intersect(*it)) {
      rem.push_back(*it);
    } else {
     break:
    --it:
  double_t start = pos->intersect(*it);
 Line f = *pos;
  for (Line r : rem) hull.erase(r);
 hull.erase(f);
  f.start = start;
 hull.insert(f);
void update_right(iterator_t it) {
  if (!has_next(it)) return;
  iterator_t pos = it;
  ++it;
  vector<Line> rem;
  while (has next(it))
    iterator_t next = it;
    if (next->intersect(*pos) <= pos->intersect(*it)) {
      rem.push_back(*it);
    } else {
     break;
    ++it;
  double_t start = pos->intersect(*it);
  Line f = *it:
  for (Line r : rem) hull.erase(r);
 hull.erase(f);
 f.start = start;
 hull.insert(f);
void insert_line(ll m, ll b) {
 Line f(m, b, -inf, 0);
  iterator_t it = hull.lower_bound(f);
 if (it != hull.end() && it->m == f.m) {
    if (it->b <= f.b) {
      return;
    } else if (it->b > f.b) {
     hull.erase(it);
 hull.insert(f);
  it = hull.lower_bound(f);
```

```
if (irrelevant(it)) {
    hull.erase(it);
    return;
}
update_left(it);
it = hull.lower_bound(f);
update_right(it);
}

ll get(ll x) {
    Line f(0, 0, x, 1);
    iterator_t it = hull.upper_bound(f);
    assert(it != hull.begin());
    --it;
    return it->m * x + it->b;
}

private:
    set<Line> hull;
};
```

#### 2.2 Line Container

```
typedef long long int 11;
bool O;
struct Line {
 mutable 11 k, m, p;
 bool operator<(const Line& o) const {</pre>
    return 0 ? p < o.p : k < o.k;
};
struct LineContainer : multiset<Line> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b)
  const ll inf = LLONG_MAX;
 11 div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b); }
 bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = inf; return false; }
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
  void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p)
     isect(x, erase(v));
  ll query(ll x) {
    assert(!emptv());
    Q = 1; auto 1 = *lower_bound({0,0,x}); Q = 0;
    return 1.k * x + 1.m;
};
```

#### 2.3 Li Chao Tree

```
// by luucasv
typedef long long T;
const T INF = 1e18, EPS = 1;
const int BUFFER_SIZE = 1e4;
struct Line {
  T m, b;
 Line (T m = 0, T b = INF): m(m), b(b) {}
 T apply(T x) { return x * m + b; }
};
struct Node {
 Node *left, *right;
 Line line;
 Node(): left(NULL), right(NULL) {}
};
struct LiChaoTree {
 Node *root, buffer[BUFFER SIZE];
  T min_value, max_value;
  int buffer_pointer;
  LiChaoTree(T min_value, T max_value): min_value(min_value),
      max_value(max_value + 1) { clear(); }
  void clear() { buffer_pointer = 0; root = newNode(); }
  void insert_line(T m, T b) { update(root, min_value, max_value, Line
      (m, b)); }
  T eval(T x) { return query(root, min_value, max_value, x); }
  void update(Node *cur, T l, T r, Line line) {
    T m = 1 + (r - 1) / 2:
    bool left = line.apply(l) < cur->line.apply(l);
    bool mid = line.apply(m) < cur->line.apply(m);
    bool right = line.apply(r) < cur->line.apply(r);
    if (mid) {
      swap(cur->line, line);
    if (r - 1 <= EPS) return;</pre>
    if (left == right) return;
    if (mid != left) {
      if (cur->left == NULL) cur->left = newNode();
      update(cur->left, 1, m, line);
    } else {
      if (cur->right == NULL) cur->right = newNode();
      update(cur->right, m, r, line);
  T query (Node *cur, T l, T r, T x) {
    if (cur == NULL) return INF;
    if (r - 1 <= EPS) {
      return cur->line.apply(x);
    T m = 1 + (r - 1) / 2;
    T ans;
    if (x < m) {
      ans = query(cur->left, 1, m, x);
    } else {
      ans = query(cur->right, m, r, x);
```

```
}
  return min(ans, cur->line.apply(x));
}
Node* newNode() {
  buffer[buffer_pointer] = Node();
  return &buffer[buffer_pointer++];
}
};
```

## 2.4 Divide and Conquer Optimization

```
int n, k;
11 dpold[ms], dp[ms], c[ms][ms]; // c(i, j) pode ser funcao
void compute(int 1, int r, int opt1, int optr) {
    if(l>r) return;
    int mid = (1+r)/2;
    pair<11, int> best = {inf, -1}; // long long inf
    for(int k = optl; k <= min(mid, optr); k++) {</pre>
        best = min(best, \{dpold[k-1] + c[k][mid], k\});
    dp[mid] = best.first;
    int opt = best.second;
    compute(l, mid-1, optl, opt);
    compute(mid+1, r, opt, optr);
11 solve() {
    dp[0] = 0;
    for(int i = 1; i <= n; i++) dp[i] = inf; // initialize row 0 of</pre>
    for(int i = 1; i <= k; i++) {
        swap(dpold, dp);
        compute(0, n, 0, n); // solve row i of the dp
    return dp[n]; // return dp[k][n]
```

# 2.5 Knuth Optimization

# 3 Geometry

## 3.1 Geometry

```
const double inf = 1e100, eps = 1e-9;
const double PI = acos(-1.0L);
int cmp (double a, double b = 0) {
  if (abs(a-b) < eps) return 0;</pre>
  return (a < b) ? -1 : +1;
struct PT {
  double x, v;
  PT (double x = 0, double y = 0) : x(x), y(y) {}
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
  PT operator * (double c) const { return PT(x*c, y*c); }
  PT operator / (double c) const { return PT(x/c, y/c); }
  bool operator < (const PT &p) const {
    if(cmp(x, p.x) != 0) return x < p.x;
    return cmp(y, p.y) < 0;
  bool operator == (const PT &p) const {
    return !cmp(x, p.x) && !cmp(y, p.y);
  bool operator != (const PT &p) const +
    return ! (p == *this);
};
double dot (PT p, PT q) { return p.x * q.x + p.y*q.y; }
double cross (PT p, PT q) { return p.x * q.y - p.y*q.x; }
double dist2 (PT p, PT q = PT(0, 0)) { return dot(p-q, p-q); }
double dist (PT p, PT q) { return hypot(p.x-q.x, p.y-q.y); }
double norm (PT p) { return hypot(p.x, p.y); }
PT normalize (PT p) { return p/hypot(p.x, p.y); }
double angle (PT p, PT q) { return atan2(cross(p, q), dot(p, q)); }
double angle (PT p) { return atan2(p.y, p.x); }
double polarAngle (PT p) {
  double a = atan2(p.v,p.x);
  return a < 0 ? a + 2*PI : a;
// - p.y*sen(+90), p.x*sen(+90)
PT rotateCCW90 (PT p) { return PT(-p.y, p.x); }
// - p.y*sen(-90), p.x*sen(-90)
PT rotateCW90 (PT p) { return PT(p.y, -p.x); }
PT rotateCCW (PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// !!! PT (int, int)
typedef pair<PT, int> Line;
PT getDir (PT a, PT b) {
  if (a.x == b.x) return PT(0, 1);
 if (a.y == b.y) return PT(1, 0);
```

```
int dx = b.x-a.x;
  int dv = b.v-a.v;
  int g = \underline{gcd(abs(dx), abs(dy))};
  if (dx < 0) q = -q;
  return PT(dx/g, dy/g);
Line getLine (PT a, PT b) {
 PT dir = getDir(a, b);
 return {dir, cross(dir, a)};
// Projeta ponto c na linha a - b assumindo a != b
// a.b = |a| cost * |b|
PT projectPointLine (PT a, PT b, PT c) {
 return a + (b-a) * dot (b-a, c-a) /dot (b-a, b-a);
PT reflectPointLine (PT a, PT b, PT c) {
 PT p = projectPointLine(a, b, c):
 return p*2 - c;
// Projeta ponto c no segmento a - b
PT projectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (cmp(r) == 0) return a;
  r = dot(b-a, c-a)/r;
  if (cmp(r, 0) < 0) return a;
  if (cmp(r, 1) > 0) return b;
  return a + (b - a) * r;
// Calcula distancia entre o ponto c e o segmento a - b
double distancePointSegment (PT a, PT b, PT c) {
  return dist(c, projectPointSegment(a, b, c));
// Parallel and opposite directions
// Determina se o ponto c esta em um segmento a - b
bool ptInSegment (PT a, PT b, PT c) {
 if (a == b) return a == c;
 a = a-c, b = b-c;
 return cmp(cross(a, b)) == 0 && cmp(dot(a, b)) <= 0;
// Determina se as linhas a - b e c - d sao paralelas ou colineares
bool parallel (PT a, PT b, PT c, PT d) {
  return cmp(cross(b - a, c - d)) == 0;
bool collinear (PT a, PT b, PT c, PT d) {
  return parallel(a, b, c, d) && cmp(cross(a - b, a - c)) == 0 && cmp(
      cross(c - d, c - a)) == 0;
// Calcula distancia entre o ponto (x, y, z) e o plano ax + by + cz =
double distancePointPlane(double x, double y, double z, double a,
    double b, double c, double d) {
    return abs(a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c
```

```
// Determina se o segmento a - b intersecta com o segmento c - d
bool segmentsIntersect (PT a, PT b, PT c, PT d) {
  if (collinear(a, b, c, d)) {
    if (cmp(dist(a, c)) == 0 \mid | cmp(dist(a, d)) == 0 \mid | cmp(dist(b, c))
        ) == 0 || cmp(dist(b, d)) == 0) return true;
    if (cmp(dot(c - a, c - b)) > 0 \&\& cmp(dot(d - a, d - b)) > 0 \&\&
        cmp(dot(c - b, d - b)) > 0) return false;
    return true;
  if (cmp(cross(d - a, b - a) * cross(c - a, b - a)) > 0) return false
  if (cmp(cross(a - c, d - c) * cross(b - c, d - c)) > 0) return false
  return true:
// Calcula a intersecao entre as retas a - b e c - d assumindo que uma
     unica intersecao existe
// Para intersecao de segmentos, cheque primeiro se os segmentos se
    intersectam e que nao sao paralelos
// r = a1 + t*d1, (r - a2) x d2 = 0
PT computeLineIntersection (PT a, PT b, PT c, PT d) {
 b = b - a; d = c - d; c = c - a;
 assert(cmp(cross(b, d)) != 0);
  return a + b * cross(c, d) / cross(b, d);
// Calcula centro do circulo dado tres pontos
PT computeCircleCenter (PT a, PT b, PT c) {
 b = (a + b) / 2; // bissector
 c = (a + c) / 2; // bissector
  return computeLineIntersection(b, b + rotateCW90(a - b), c, c +
      rotateCW90(a - c));
vector<PT> circle2PtsRad (PT p1, PT p2, double r) {
 vector<PT> ret:
 double d2 = dist2(p1, p2);
  double det = r * r / d2 - 0.25;
  if (det < 0.0) return ret;</pre>
  double h = sgrt(det);
  for (int i = 0; i < 2; i++) {
    double x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
    double y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;
    ret.push_back(PT(x, y));
    swap(p1, p2);
  return ret;
// Calcula intersecao da linha a - b com o circulo centrado em c com
    raio r > 0
bool circleLineIntersection(PT a, PT b, PT c, double r) {
    return cmp(dist(c, projectPointLine(a, b, c)), r) <= 0;</pre>
vector<PT> circleLine (PT a, PT b, PT c, double r) {
 vector<PT> ret;
  PT p = projectPointLine(a, b, c), p1;
```

```
double h = norm(c-p);
  if (cmp(h,r) == 0) {
    ret.push back(p);
  else if (cmp(h,r) < 0) 
    double k = sqrt(r*r - h*h);
    p1 = p + (b-a)/(norm(b-a)) *k;
    ret.push_back(p1);
    p1 = p - (b-a) / (norm(b-a)) *k;
    ret.push back(p1);
  return ret;
bool ptInsideTriangle(PT p, PT a, PT b, PT c) {
  if(cross(b-a, c-b) < 0) swap(a, b);
  long long x = cross(b-a, p-b);
  long long y = cross(c-b, p-c);
  long long z = cross(a-c, p-a);
  if (x > 0 \& \& v > 0 \& \& z > 0) return true;
  if(!x) return ptInSegment(a,b,p);
  if(!y) return ptInSegment(b,c,p);
  if(!z) return ptInSegment(c,a,p);
  return false;
// Determina se o ponto esta num poligono convexo em O(lgn)
bool pointInConvexPolygon(const vector<PT> &hull, PT point) {
  int n = hull.size();
  if(cmp(cross(point - hull[0], hull[1] - hull[0])) || cmp(cross(point
       - hull[0], hull[n-1] - hull[0]))) return false;
  int 1 = 1, r = n - 1;
  while (1 != r) {
    int mid = (1 + r + 1) / 2;
    if(cmp(cross(point - hull[0], hull[mid] - hull[0])) < 0) 1 = mid;</pre>
    else r = mid - 1;
  return cmp(cross(hull[(l+1)%n] - hull[1], point - hull[1])) >= 0;
// Determina se o ponto esta num poligono possivelmente nao-convexo
// Retorna 1 para pontos estritamente dentro, 0 para pontos
    estritamente fora do poligno
// e 0 ou 1 para os pontos restantes
// Eh possivel converter num teste exato usando inteiros e tomando
    cuidado com a divisao
// e entao usar testes exatos para checar se esta na borda do poligno
bool pointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for(int i = 0; i < p.size(); i++){</pre>
    int j = (i + 1) % p.size();
    if((p[i].y \le q.y \& q.y < p[j].y || p[j].y \le q.y \& q.y < p[i].y
        ) & &
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y)
      c = !c;
  return c;
// Determina se o ponto esta na borda do poligno
bool pointOnPolygon(const vector<PT> &p, PT q) {
```

```
for(int i = 0; i < p.size(); i++)</pre>
    if(cmp(dist2(projectPointSegment(p[i], p[(i + 1) % p.size()], q),
        (0)
      return true;
    return false:
// area / semiperimeter
double rIncircle (PT a, PT b, PT c) {
  double ab = norm(a-b), bc = norm(b-c), ca = norm(c-a);
  return abs(cross(b-a, c-a)/(ab+bc+ca));
// Calcula intersecao do circulo centrado em a com raio r e o centrado
     em b com raio R
vector<PT> circleCircle (PT a, double r, PT b, double R) {
 vector<PT> ret;
  double d = norm(a-b);
  if (d > r + R \mid | d + min(r, R) < max(r, R)) return ret;
  double x = (d*d - R*R + r*r) / (2*d); // x = r*cos(R opposite angle)
  double y = sqrt(r*r - x*x);
 PT v = (b - a)/d;
  ret.push back(a + v*x + rotateCCW90(v)*v);
  if (cmp(y) > 0)
    ret.push_back(a + v*x - rotateCCW90(v)*y);
  return ret;
double circularSegArea (double r, double R, double d) {
  double ang = 2 * acos((d*d - R*R + r*r) / (2*d*r)); // cos(R)
      opposite angle) = x/r
  double tri = sin(ang) * r * r;
  double sector = ang * r * r;
  return (sector - tri) / 2;
// Calcula a area ou o centroide de um poligono (possivelmente nao-
    convexo)
// assumindo que as coordenadas estao listada em ordem horaria ou anti
    -horaria
// O centroide eh equivalente a o centro de massa ou centro de
    gravidade
double computeSignedArea (const vector<PT> &p) {
  double area = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area/2.0;
double computeArea(const vector<PT> &p) {
  return abs(computeSignedArea(p));
PT computeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for(int i = 0; i < p.size(); i++){</pre>
   int j = (i + 1) % p.size();
    c = c + (p[i] + p[i]) * (p[i].x * p[i].y - p[i].x * p[i].y);
```

```
return c / scale;
// Testa se o poligno listada em ordem CW ou CCW eh simples (nenhuma
    linha se intersecta)
bool isSimple(const vector<PT> &p) {
  for(int i = 0; i < p.size(); i++) {</pre>
    for (int k = i + 1; k < p.size(); k++) {
      int j = (i + 1) % p.size();
      int 1 = (k + 1) % p.size();
      if (i == 1 \mid | j == k) continue;
      if (segmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
vector< pair<PT, PT> > getTangentSegs (PT c1, double r1, PT c2, double
  if (r1 < r2) swap(c1, c2), swap(r1, r2);</pre>
  vector<pair<PT, PT> > ans;
  double d = dist(c1, c2);
  if (cmp(d) <= 0) return ans;</pre>
  double dr = abs(r1 - r2), sr = r1 + r2;
  if (cmp(dr, d) >= 0) return ans;
  double u = acos(dr / d);
  PT dc1 = normalize(c2 - c1) *r1;
 PT dc2 = normalize(c2 - c1) \star r2;
  ans.push_back(make_pair(c1 + rotateCCW(dc1, +u), c2 + rotateCCW(dc2,
  ans.push_back(make_pair(c1 + rotateCCW(dc1, -u), c2 + rotateCCW(dc2,
       -u)));
  if (cmp(sr, d) >= 0) return ans;
  double v = acos(sr / d);
  dc2 = normalize(c1 - c2) *r2;
  ans.push_back({c1 + rotateCCW(dc1, +v), c2 + rotateCCW(dc2, +v)});
  ans.push_back(\{c1 + rotateCCW(dc1, -v), c2 + rotateCCW(dc2, -v)\});
  return ans;
```

## 3.2 Convex Hull

```
void sortByAngle (vector<PT>::iterator first, vector<PT>::iterator
    last, const PT o) {
  first = partition(first, last, [&o] (const PT &a) { return a == o;
      });
  auto pivot = partition(first, last, [&o] (const PT &a) {
    return ! (a < o | | a == o); // PT(a.y, a.x) < PT(o.y, o.x)
  auto acmp = [&o] (const PT &a, const PT &b) { // C++11 only
    if (cmp(cross(a-o, b-o)) != 0) return cross(a-o, b-o) > 0;
    else return cmp(norm(a-o), norm(b-o)) < 0;</pre>
  sort(first, pivot, acmp);
  sort(pivot, last, acmp);
vector<PT> graham (vector<PT> v) {
  sort(v.begin(), v.end());
  sortBvAngle(v.begin(), v.end(), v[0]);
  vector<PT> u (v.size());
  int top = 0:
  for (int i = 0; i < v.size(); i++) {</pre>
    while (top > 1 \&\& cmp(cross(u[top-1] - u[top-2], v[i]-u[top-2]))
        <= 0) top--;
    u[top++] = v[i];
  u.resize(top);
  return u;
```

## 3.3 Cut Polygon

```
struct Segment {
 typedef long double T;
 PT p1, p2;
 T a, b, c;
 Segment() {}
  Segment (PT st, PT en) {
   p1 = st, p2 = en;
   a = -(st.y - en.y);
   b = st.x - en.x;
   c = a * en.x + b * en.y;
 T plug(T x, T y) {
    // plug >= 0 is to the right
    return a * x + b * y - c;
 T plug(PT p) {
    return plug(p.x, p.y);
 bool inLine(PT p) { return cross((p - p1), (p2 - p1)) == 0; }
 bool inSegment(PT p) {
   return inLine(p) && dot((p1 - p2), (p - p2)) >= 0 && dot((p2 - p1)
        (p - p1)) >= 0;
```

```
PT lineIntersection(Segment s) {
    long double A = a, B = b, C = c;
    long double D = s.a, E = s.b, F = s.c;
    long double x = (long double) C * E - (long double) B * F;
    long double y = (long double) A * F - (long double) C * D;
    long double tmp = (long double) A * E - (long double) B * D;
    x /= tmp;
    v /= tmp;
    return PT(x, y);
 bool polygonIntersection(const vector<PT> &poly) {
    long double 1 = -1e18, r = 1e18;
    for(auto p : poly) {
      long double z = plug(p);
      1 = \max(1, z);
      r = min(r, z);
    return 1 - r > eps;
};
vector<PT> cutPolygon(vector<PT> poly, Segment seg) {
  int n = (int) poly.size();
  vector<PT> ans;
  for (int i = 0; i < n; i++) {
    double z = seg.plug(poly[i]);
    if(z > -eps) {
      ans.push_back(poly[i]);
    double z2 = seg.plug(poly[(i + 1) % n]);
    if((z > eps \&\& z2 < -eps) || (z < -eps \&\& z2 > eps)) {
      ans.push_back(seq.lineIntersection(Segment(poly[i], poly[(i + 1)
           % n])));
  return ans;
```

## 3.4 Smallest Enclosing Circle

```
typedef pair<PT, double> circle;
bool inCircle (circle c, PT p) {
   return cmp(dist(c.first, p), c.second) <= 0;
}

PT circumcenter (PT p, PT q, PT r) {
   PT a = p-r, b = q-r;
   PT c = PT(dot(a, p+r)/2, dot(b, q+r)/2);
   return PT(cross(c, PT(a.y,b.y)), cross(PT(a.x,b.x), c)) / cross(a, b );
}

mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
circle spanningCircle (vector<PT> &v) {
   int n = v.size();
   shuffle(v.begin(), v.end(), rng);
   circle C(PT(), -1);
   for (int i = 0; i < n; i++) if (!inCircle(C, v[i])) {</pre>
```

```
C = circle(v[i], 0);
for (int j = 0; j < i; j++) if (!inCircle(C, v[j])) {
    C = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
    for(int k = 0; k < j; k++) if (!inCircle(C, v[k])) {
        PT o = circumcenter(v[i], v[j], v[k]);
        C = circle(o, dist(o, v[k]));
    }
    }
}
return C;
}</pre>
```

#### 3.5 Minkowski

```
bool comp(PT a, PT b) {
  int hp1 = (a.x < 0 \mid | (a.x==0 \&\& a.y<0));
  int hp2 = (b.x < 0 \mid | (b.x==0 \&\& b.y<0));
  if (hp1 != hp2) return hp1 < hp2;</pre>
  long long R = cross(a, b);
  if(R) return R > 0;
  return dot(a, a) < dot(b, b);</pre>
vector<PT> minkowskiSum(const vector<PT> &a, const vector<PT> &b) {
  if(a.empty() || b.empty()) return vector<PT>(0);
  vector<PT> ret;
  int n1 = a.size(), n2 = b.size();
  if(min(n1, n2) < 2){
    for (int i = 0; i < n1; i++) {
      for (int j = 0; j < n2; j++) {
        ret.push_back(a[i]+b[j]);
    return ret;
  auto insert = [&](PT p) {
    while(ret.size() >= 2 && cmp(cross(p-ret.back(), p-ret[(int))ret.
        size()-2])) == 0) {
      // removing colinear points
      // needs the scalar product stuff it the result is a line
      ret.pop_back();
    ret.push_back(p);
  PT v1, v2, p = a[0]+b[0];
  ret.push_back(p);
  for (int i = 0, j = 0; i + j + 1 < n1+n2; ) {
    v1 = a[(i+1)%n1]-a[i];
    v2 = b[(j+1) n2] - b[j];
    if(j == n2 \mid | (i < n1 && comp(v1, v2))) p = p + v1, i++;
    else p = p + v2, j++;
    insert(p);
  return ret;
```

#### 3.6 Half Plane Intersection

```
struct L {
    PT a, b;
    L(){}
    L(PT a, PT b) : a(a), b(b) {}
};
double angle (L la) { return atan2(-(la.a.y - la.b.y), la.b.x - la.a.x
    ); }
bool comp (L la, L lb) {
    if (cmp(angle(la), angle(lb)) == 0) return cross((lb.b - lb.a), (
        la.b - lb.a)) > eps;
    return cmp(angle(la), angle(lb)) < 0;</pre>
PT computeLineIntersection (L la, L lb) {
    return computeLineIntersection(la.a, la.b, lb.a, lb.b);
bool check (L la, L lb, L lc) {
    PT p = computeLineIntersection(lb, lc);
    double det = cross((la.b - la.a), (p - la.a));
    return cmp(det) < 0;</pre>
vector<PT> hpi (vector<L> line) { // salvar (i, j) CCW, (j, i) CW
    sort(line.begin(), line.end(), comp);
    vector<L> pl(1, line[0]);
    for (int i = 0; i < (int)line.size(); ++i) if (cmp(angle(line[i]),</pre>
         angle(pl.back())) != 0) pl.push_back(line[i]);
    deque<int> dq;
    dq.push_back(0);
    dq.push_back(1);
    for (int i = 2; i < (int)pl.size(); ++i) {</pre>
        while ((int)dq.size() > 1 && check(pl[i], pl[dq.back()], pl[dq
            [dq.size() - 2]])) dq.pop_back();
        while ((int)dq.size() > 1 && check(pl[i], pl[dq[0]], pl[dq
            [1]])) dq.pop_front();
        dq.push_back(i);
    while ((int)dq.size() > 1 && check(pl[dq[0]], pl[dq.back()], pl[dq
        [dq.size() - 2]])) dq.pop_back();
    while ((int)dq.size() > 1 && check(pl[dq.back()], pl[dq[0]], pl[dq
        [1]])) dq.pop_front();
    vector<PT> res;
    for (int i = 0; i < (int)dq.size(); ++i){</pre>
        res.emplace_back(computeLineIntersection(pl[dq[i]], pl[dq[(i +
             1) % dq.size()]]));
    return res;
```

## 3.7 Closest Pair

```
double closestPair(vector<PT> p) {
  int n = p.size(), k = 0;
  sort(p.begin(), p.end());
  double d = inf;
  set<PT> ptsInv;
  for(int i = 0; i < n; i++) {</pre>
```

```
while(k < i && p[k].x < p[i].x - d) {
    ptsInv.erase(swapCoord(p[k++]));
}
for(auto it = ptsInv.lower_bound(PT(p[i].y - d, p[i].x - d));
    it != ptsInv.end() && it->x <= p[i].y + d; it++) {
        d = min(d, dist(p[i] - swapCoord(*it), PT(0, 0)));
    }
    ptsInv.insert(swapCoord(p[i]));
}
return d;
}</pre>
```

# 3.8 Delaunay Triangulation

```
bool ge(const 11& a, const 11& b) { return a >= b; }
bool le(const ll& a, const ll& b) { return a <= b; }
bool eq(const ll& a, const ll& b) { return a == b; }
bool gt(const ll& a, const ll& b) { return a > b; }
bool lt(const ll& a, const ll& b) { return a < b; }</pre>
int sqn(const 11& a) { return a >= 0 ? a ? 1 : 0 : -1; }
struct pt {
    11 x, y;
    pt() { }
    pt(ll _x, ll _y) : x(_x), y(_y) { }
    pt operator-(const pt& p) const {
        return pt(x - p.x, y - p.y);
    11 cross(const pt& p) const {
        return x * p.y - y * p.x;
    11 cross(const pt& a, const pt& b) const {
        return (a - *this).cross(b - *this);
    11 dot(const pt& p) const {
        return x * p.x + y * p.y;
    11 dot(const pt& a, const pt& b) const {
        return (a - *this).dot(b - *this);
    11 sqrLength() const {
        return this->dot(*this);
    bool operator == (const pt& p) const {
        return eq(x, p.x) && eq(y, p.y);
};
const pt inf_pt = pt(1e18, 1e18);
struct QuadEdge {
    pt origin;
    OuadEdge* rot = nullptr:
    QuadEdge* onext = nullptr;
    bool used = false;
    QuadEdge* rev() const {
        return rot->rot;
    QuadEdge* lnext() const {
        return rot->rev()->onext->rot;
```

```
OuadEdge* oprev() const {
        return rot->onext->rot;
    pt dest() const {
        return rev()->origin;
};
QuadEdge* make_edge(pt from, pt to) {
    QuadEdge* e1 = new QuadEdge;
    QuadEdge* e2 = new QuadEdge;
    QuadEdge* e3 = new QuadEdge;
    QuadEdge* e4 = new QuadEdge;
    e1->origin = from;
    e2->origin = to;
    e3->origin = e4->origin = inf_pt;
    e1->rot = e3;
    e2 - > rot = e4:
    e3 \rightarrow rot = e2;
    e4->rot = e1:
    e1->onext = e1;
    e2 - > onext = e2;
    e3 \rightarrow onext = e4;
    e4 \rightarrow onext = e3;
    return e1;
void splice(OuadEdge* a, OuadEdge* b) {
    swap(a->onext->rot->onext, b->onext->rot->onext);
    swap(a->onext, b->onext);
void delete_edge(QuadEdge* e) {
    splice(e, e->oprev());
    splice(e->rev(), e->rev()->oprev());
    delete e->rot;
    delete e->rev()->rot;
    delete e:
    delete e->rev();
QuadEdge* connect (QuadEdge* a, QuadEdge* b) {
    QuadEdge* e = make_edge(a->dest(), b->origin);
    splice(e, a->lnext());
    splice(e->rev(), b);
    return e;
bool left_of(pt p, QuadEdge* e) {
    return gt(p.cross(e->origin, e->dest()), 0);
bool right of(pt p, OuadEdge* e) {
    return lt(p.cross(e->origin, e->dest()), 0);
template <class T>
T det3(T a1, T a2, T a3, T b1, T b2, T b3, T c1, T c2, T c3) {
    return a1 * (b2 * c3 - c2 * b3) - a2 * (b1 * c3 - c1 * b3) +
           a3 * (b1 * c2 - c1 * b2);
```

```
bool in_circle(pt a, pt b, pt c, pt d) {
// If there is int128, calculate directly.
// Otherwise, calculate angles.
#if defined( LP64 ) || defined( WIN64)
    __int128 det = -det3<__int128>(b.x, b.y, b.sqrLength(), c.x, c.y,
                                    c.sqrLength(), d.x, d.y, d.
                                        sqrLength());
    det += det3<__int128>(a.x, a.y, a.sqrLength(), c.x, c.y, c.
        sqrLength(), d.x,
                          d.y, d.sqrLength());
    det -= det3 < __int128 > (a.x, a.y, a.sqrLength(), b.x, b.y, b.
        sqrLength(), d.x,
                          d.y, d.sqrLength());
    det += det3 < __int128 > (a.x, a.y, a.sqrLength(), b.x, b.y, b.
        sqrLength(), c.x,
                          c.y, c.sqrLength());
    return det > 0:
#else
    auto ang = [](pt l, pt mid, pt r) {
        11 x = mid.dot(1, r);
        11 y = mid.cross(l, r);
        long double res = atan2((long double)x, (long double)y);
        return res;
    long double kek = ang(a, b, c) + ang(c, d, a) - ang(b, c, d) - ang
        (d, a, b);
    if (kek > 1e-8)
        return true;
    else
        return false:
#endif
pair<QuadEdge*, QuadEdge*> build_tr(int 1, int r, vector<pt>& p) {
    if (r - 1 + 1 == 2) {
        QuadEdge* res = make_edge(p[l], p[r]);
        return make_pair(res, res->rev());
    if (r - 1 + 1 == 3) {
        QuadEdge *a = make\_edge(p[1], p[1 + 1]), *b = make\_edge(p[1 + 1])
            1], p[r]);
        splice(a->rev(), b);
        int sg = sgn(p[1].cross(p[1 + 1], p[r]));
        if (sq == 0)
            return make_pair(a, b->rev());
        QuadEdge* c = connect(b, a);
        if (sq == 1)
            return make_pair(a, b->rev());
        else
            return make_pair(c->rev(), c);
    int mid = (1 + r) / 2;
    QuadEdge *ldo, *ldi, *rdo, *rdi;
    tie(ldo, ldi) = build_tr(l, mid, p);
    tie(rdi, rdo) = build_tr(mid + 1, r, p);
    while (true) {
        if (left_of(rdi->origin, ldi)) {
            ldi = ldi->lnext();
            continue;
```

```
if (right_of(ldi->origin, rdi)) {
            rdi = rdi->rev()->onext;
            continue;
        break;
    QuadEdge* basel = connect(rdi->rev(), ldi);
    auto valid = [&basel](OuadEdge* e) { return right of(e->dest(),
        basel); };
    if (ldi->origin == ldo->origin)
        ldo = basel->rev();
   if (rdi->origin == rdo->origin)
        rdo = basel;
   while (true) {
        QuadEdge* lcand = basel->rev()->onext;
        if (valid(lcand)) {
            while (in_circle(basel->dest(), basel->origin, lcand->dest
                (),
                             lcand->onext->dest())) {
                QuadEdge* t = lcand->onext;
                delete edge(lcand);
                lcand = t;
        QuadEdge* rcand = basel->oprev();
        if (valid(rcand)) {
            while (in_circle(basel->dest(), basel->origin, rcand->dest
                             rcand->oprev()->dest())) {
                QuadEdge* t = rcand->oprev();
                delete_edge(rcand);
                rcand = t;
        if (!valid(lcand) && !valid(rcand))
            break:
        if (!valid(lcand) ||
            (valid(rcand) && in_circle(lcand->dest(), lcand->origin,
                                       rcand->origin, rcand->dest())))
            basel = connect(rcand, basel->rev());
            basel = connect(basel->rev(), lcand->rev());
    return make_pair(ldo, rdo);
vector<tuple<pt, pt, pt>> delaunay(vector<pt> p) {
    sort(p.begin(), p.end(), [](const pt& a, const pt& b) {
        return lt(a.x, b.x) || (eq(a.x, b.x) && lt(a.y, b.y));
    auto res = build_tr(0, (int)p.size() - 1, p);
    QuadEdge* e = res.first;
   vector<OuadEdge*> edges = {e};
   while (lt(e->onext->dest().cross(e->dest(), e->origin), 0))
        e = e->onext;
    auto add = [&p, &e, &edges]() {
        QuadEdge* curr = e;
        do {
            curr->used = true;
            p.push_back(curr->origin);
```

## 3.9 Java Geometry Library

```
import java.util.*;
import java.io.*;
import java.awt.geom.*;
import java.lang.*;
//Lazy Geometry
class AWT{
  static Area makeArea(double[] pts){
    Path2D.Double p = new Path2D.Double();
    p.moveTo(pts[0], pts[1]);
    for (int i = 2; i < pts.length; i+=2) {
      p.lineTo(pts[i], pts[i+1]);
    p.closePath();
    return new Area(p);
  static double computePolygonArea(ArrayList<Point2D.Double> points) {
    Point2D.Double[] pts = points.toArray(new Point2D.Double[points.
        size()]);
    double area = 0;
    for (int i = 0; i < pts.length; i++) {
      int j = (i+1) % pts.length;
      area += pts[i].x * pts[j].y - pts[j].x * pts[i].y;
    return Math.abs(area)/2;
  static double computeArea(Area area) {
    double totArea = 0;
    PathIterator iter = area.getPathIterator(null);
    ArrayList<Point2D.Double> points = new ArrayList<Point2D.Double>()
    while (!iter.isDone()) {
      double[] buffer = new double[6];
      switch (iter.currentSeament(buffer)) {
        case PathIterator.SEG_MOVETO:
        case PathIterator.SEG_LINETO:
          points.add(new Point2D.Double(buffer[0], buffer[1]));
          break;
        case PathIterator.SEG CLOSE:
          totArea += computePolygonArea(points);
          points.clear();
```

```
break;
    }
    iter.next();
    }
    return totArea;
}
```

# 4 Graph Algorithms

# 4.1 Simple Disjoint Set

struct dsu {

```
vector<int> hist, par, sz;
vector<ii> changes;
int n;
dsu (int n) : n(n) {
 hist.assign(n, 1e9);
 par.resize(n);
 iota(par.begin(), par.end(), 0);
 sz.assign(n, 1);
int root (int x, int t) {
 if(hist[x] > t) return x;
 return root(par[x], t);
void join (int a, int b, int t) {
 a = root(a, t);
 b = root(b, t);
 if (a == b) { changes.emplace_back(-1, -1); return; }
 if (sz[a] > sz[b]) swap(a, b);
 par[a] = b;
 sz[b] += sz[a];
 hist[a] = t;
 changes.emplace_back(a, b);
bool same (int a, int b, int t) {
 return root(a, t) == root(b, t);
void undo () {
 int a, b;
 tie(a, b) = changes.back();
 changes.pop_back();
 if (a == -1) return;
 sz[b] = sz[a];
 par[a] = a;
 hist[a] = 1e9;
 n++;
int when (int a, int b) {
 while (1) {
   if (hist[a] > hist[b]) swap(a, b);
```

```
if (par[a] == b) return hist[a];
   if (hist[a] == le9) return le9;
   a = par[a];
}
}
};
```

#### 4.2 Boruvka

```
struct edge {
  int u, v;
  int w;
  int id;
  edge () {};
  edge (int u, int v, int w = 0, int id = 0) : u(u), v(v), w(w), id(id)
  bool operator < (edge &other) const { return w < other.w; };</pre>
vector<edge> boruvka (vector<edge> &edges, int n) {
  vector<edge> mst;
  vector<edge> best(n);
  initDSU(n);
  bool f = 1:
  while (f) {
    f = 0:
    for (int i = 0; i < n; i++) best[i] = edge(i, i, inf);</pre>
    for (auto e : edges) {
      int pu = root(e.u), pv = root(e.v);
      if (pu == pv) continue;
      if (e < best[pu]) best[pu] = e;</pre>
      if (e < best[pv]) best[pv] = e;</pre>
    for (int i = 0; i < n; i++) {</pre>
      edge e = best[root(i)];
      if (e.w != inf) {
        join(e.u, e.v);
        mst.push_back(e);
        f = 1;
  return mst;
```

#### 4.3 Dinic Max Flow

```
const int ms = 1e3; // Quantidade maxima de vertices
const int me = 1e5; // Quantidade maxima de arestas

int adj[ms], to[me], ant[me], wt[me], z, n;
int copy_adj[ms], fila[ms], level[ms];

void clear() { // Lembrar de chamar no main
   memset(adj, -1, sizeof adj);
   z = 0;
}
```

```
void add(int u, int v, int k) {
  to[z] = v;
  ant[z] = adj[u];
  wt[z] = k;
  adj[u] = z++;
  swap(u, v);
 to[z] = v;
  ant[z] = adj[u];
 wt[z] = 0; // Lembrar de colocar = 0
  adj[u] = z++;
int bfs(int source, int sink) {
 memset(level, -1, sizeof level);
  level[source] = 0;
  int front = 0, size = 0, v;
  fila[size++] = source;
  while(front < size) {</pre>
    v = fila[front++];
    for(int i = adj[v]; i != -1; i = ant[i]) {
      if(wt[i] && level[to[i]] == -1) {
        level[to[i]] = level[v] + 1;
        fila[size++] = to[i];
    }
  return level[sink] != -1;
int dfs(int v, int sink, int flow) {
  if(v == sink) return flow;
  int f;
  for(int &i = copy_adj[v]; i != -1; i = ant[i]) {
    if(wt[i] && level[to[i]] == level[v] + 1 &&
      (f = dfs(to[i], sink, min(flow, wt[i])))) {
      wt[i] -= f;
      wt[i ^ 1] += f;
      return f;
  return 0;
int maxflow(int source, int sink) {
  int ret = 0, flow;
 while(bfs(source, sink)) {
    memcpy(copy_adj, adj, sizeof adj);
    while((flow = dfs(source, sink, 1 << 30))) {</pre>
      ret += flow;
  return ret;
```

#### 4.4 Minimum Vertex Cover

```
// + Dinic
vector<int> coverU, U, coverV, V; // ITA - Parti o U LEFT,
    parti o V RIGHT, 0 indexed
bool Zu[mx], Zv[mx];
```

```
int pairU[mx], pairV[mx];
void getreach(int u) {
 if (u == -1 || Zu[u]) return;
  Zu[u] = true;
  for (int i = adj[u]; ~i; i = ant[i]) {
    int v = to[i];
    if (v == SOURCE || v == pairU[u]) continue;
    Zv[v] = true;
    getreach(pairV[v]);
void minimumcover () {
 memset(pairU, -1, sizeof pairU);
 memset(pairV, -1, sizeof pairV);
  for (auto i : U) {
    for (int j = adj[i]; ~j; j = ant[j]) {
     if (!(j&1) && !wt[j]) {
       pairU[i] = to[j], pairV[to[j]] = i;
  memset (Zu, 0, sizeof Zu);
 memset (Zv, 0, sizeof Zv);
  for (auto u : U) {
    if (pairU[u] == -1) getreach(u);
  coverU.clear(), coverV.clear();
  for (auto u : U) {
    if (!Zu[u]) coverU.push_back(u);
  for (auto v : V) {
    if (Zv[v]) coverV.push_back(v);
```

#### 4.5 Min Cost Max Flow

```
template <class flow_t, class cost_t>
class MinCostMaxFlow {
private:
  typedef pair<cost_t, int> ii;
  struct Edge {
    int to;
    flow_t cap;
    cost_t cost;
    Edge(int to, flow_t cap, cost_t cost) : to(to), cap(cap), cost(
        cost) {}
  };
  int n;
  vector<vector<int>> adi;
  vector<Edge> edges;
 vector<cost_t> dis;
  vector<int> prev, id_prev;
        vector<int> q;
        vector<bool> ing;
  pair<flow_t, cost_t> spfa(int src, int sink) {
```

```
fill(dis.begin(), dis.end(), int(1e9)); //cost_t inf
    fill(prev.begin(), prev.end(), -1);
    fill(ing.begin(), ing.end(), false);
    q.clear();
    q.push_back(src);
    ing[src] = true;
    dis[src] = 0;
    for (int on = 0; on < (int) q.size(); on++) {
        int cur = q[on];
        inq[cur] = false;
        for(auto id : adj[cur]) {
                if (edges[id].cap == 0) continue;
                int to = edges[id].to;
                if (dis[to] > dis[cur] + edges[id].cost) {
                        prev[to] = cur;
                        id_prev[to] = id;
                        dis[to] = dis[cur] + edges[id].cost;
                        if (!inq[to]) {
                                q.push_back(to);
                                inq[to] = true;
    flow_t mn = flow_t(1e9);
    for(int cur = sink; prev[cur] != -1; cur = prev[cur]) {
      int id = id prev[cur];
      mn = min(mn, edges[id].cap);
    if (mn == flow_t(1e9) || mn == 0) return make_pair(0, 0);
    pair<flow_t, cost_t> ans(mn, 0);
    for(int cur = sink; prev[cur] != -1; cur = prev[cur]) {
      int id = id_prev[cur];
      edges[id].cap -= mn;
      edges[id ^ 1].cap += mn;
      ans.second += mn * edges[id].cost;
    return ans;
public:
 MinCostMaxFlow(int a = 0) {
    n = a:
    adj.resize(n + 2);
    edges.clear();
    dis.resize(n + 2);
    prev.resize(n + 2);
    id_prev.resize(n + 2);
    inq.resize(n + 2);
  void init(int a) {
    n = a;
    adj.resize(n + 2);
    edges.clear();
    dis.resize(n + 2);
    prev.resize(n + 2);
    id_prev.resize(n + 2);
    inq.resize(n + 2);
  void add(int from, int to, flow_t cap, cost_t cost) {
    adj[from].push_back(int(edges.size()));
                edges.push_back(Edge(to, cap, cost));
```

#### 4.6 Euler Path and Circuit

```
int pathV[me], szV, del[me], pathE, szE;
int adj[ms], to[me], ant[me], wt[me], z, n;

// Funcao de add e clear no dinic

void eulerPath(int u) {
  for(int i = adj[u]; ~i; i = ant[u]) if(!del[i]) {
    del[i] = del[i^1] = 1;
    eulerPath(to[i]);
    pathE[szE++] = i;
  }
  pathV[szV++] = u;
}
```

# 4.7 Articulation Points/Bridges/Biconnected Components

```
int adj[ms], to[me], ant[me], z;
int num[ms], low[ms], timer;
int art[ms], bridge[me], rch;
int bc[ms], nbc;
stack<int> st;
bool f[me];
void clear() { // Lembrar de chamar no main
 memset(adj, -1, sizeof adj);
  z = 0;
void add(int u, int v) {
  to[z] = v;
  ant[z] = adj[u];
  adj[u] = z++;
void generateBc (int v) {
  while (!st.empty()) {
   int u = st.top();
   st.pop();
   bc[u] = nbc;
    if (v == u) break;
  ++nbc;
```

```
void dfs (int v, int p) {
  st.push(v);
  low[v] = num[v] = ++timer;
  for (int i = adj[v]; i != -1; i = ant[i]) {
    if (f[i] || f[i^1]) continue;
    f[i] = 1;
    int u = to[i];
    if (num[u] == -1) {
      dfs(u, v);
      if (low[u] > num[v]) bridge[i] = bridge[i^1] = 1;
      art[v] \mid = p != -1 \&\& low[u] >= num[v];
      if (p == -1 \&\& rch > 1) art[v] = 1;
      else rch ++;
      low[v] = min(low[v], low[u]);
    } else {
      low[v] = min(low[v], num[u]);
  if (low[v] == num[v]) generateBc(v);
void biCon (int n) {
 nbc = 0, timer = 0;
 memset(num, -1, sizeof num);
 memset(bc, -1, sizeof bc);
 memset(bridge, 0, sizeof bridge);
 memset(art, 0, sizeof art);
 memset(f, 0, sizeof f);
  for (int i = 0; i < n; i++) {
    if (num[i] == -1) {
      rch = 0;
      dfs(i, 0);
```

# 4.8 SCC - Strongly Connected Components / 2SAT

```
vector<int> g[ms];
int idx[ms], low[ms], z, comp[ms], ncomp;
stack<int> st;
int dfs(int u) {
  if(~idx[u]) return idx[u] ? idx[u] : z;
  low[u] = idx[u] = z++;
  st.push(u);
  for(int v : g[u]) {
    low[u] = min(low[u], dfs(v));
  if(low[u] == idx[u]) {
    while(st.top() != u) {
      int v = st.top();
      idx[v] = 0;
      low[v] = low[u];
      comp[v] = ncomp;
      st.pop();
    idx[st.top()] = 0;
```

```
st.pop();
    comp[u] = ncomp++;
  return low[u];
bool solveSat() {
  memset(idx, -1, sizeof idx);
  z = 1; ncomp = 0;
  for(int i = 0; i < n; i++) dfs(i);</pre>
  for(int i = 0; i < n; i++) if(comp[i] == comp[i^1]) return false;</pre>
  return true;
// Operacoes comuns de 2-sat
// ~v = "nao v"
#define trad(v) (v<0?((~v)*2)^1:v*2)
void addImp(int a, int b) { g[trad(a)].push(trad(b));
void addOr(int a, int b) { addImp(~a, b); addImp(~b, a); }
void addEqual(int a, int b) { addOr(a, ~b); addOr(~a, b); }
void addDiff(int a, int b) { addEqual(a, ~b); }
// valoracao: value[v] = comp[trad(v)] < comp[trad(~v)]</pre>
```

#### 4.9 LCA - Lowest Common Ancestor

```
int par[ms][mlg+1], lvl[ms];
vector<int> q[ms];
void dfs(int v, int p, int 1 = 0) { // chamar como dfs(root, root)
  lvl[v] = 1;
  par[v][0] = p;
  for(int k = 1; k <= mlg; k++) {</pre>
    par[v][k] = par[par[v][k-1]][k-1];
  for(int u : g[v]) {
    if (u != p) dfs(u, v, l + 1);
int lca(int a, int b) {
  if(lvl[b] > lvl[a]) swap(a, b);
  for(int i = mlg; i >= 0; i--) {
    if(lvl[a] - (1 << i) >= lvl[b]) a = par[a][i];
  if(a == b) return a;
  for(int i = mlg; i >= 0; i--) {
    if(par[a][i] != par[b][i]) a = par[a][i], b = par[b][i];
  return par[a][0];
```

# 4.10 Heavy Light Decomposition

```
// src: tfg
class HLD {
public:
    void init(int n) {
        // this doesn't delete edges!
```

```
sz.resize(n);
    in.resize(n);
    out.resize(n);
    rin.resize(n);
    p.resize(n);
    edges.resize(n);
    nxt.resize(n);
    h.resize(n);
  void addEdge(int u, int v) {
    edges[u].push_back(v);
    edges[v].push_back(u);
  void setRoot(int n) {
    t = 0;
    p[n] = n;
    h[n] = 0;
    prep(n, n);
    nxt[n] = n;
    hld(n);
  int getLCA(int u, int v) {
    while(!inSubtree(nxt[u], v)) {
      u = p[nxt[u]];
    while(!inSubtree(nxt[v], u)) {
      v = p[nxt[v]];
    return in[u] < in[v] ? u : v;</pre>
  bool inSubtree(int u, int v) {
    // is v in the subtree of u?
    return in[u] <= in[v] && in[v] < out[u];</pre>
  vector<pair<int, int>> getPathtoAncestor(int u, int anc) {
    // returns ranges [l, r) that the path has
    vector<pair<int, int>> ans;
    //assert(inSubtree(anc, u));
    while(nxt[u] != nxt[anc]) {
      ans.emplace_back(in[nxt[u]], in[u] + 1);
      u = p[nxt[u]];
    // this includes the ancestor!
    ans.emplace_back(in[anc], in[u] + 1);
    return ans;
private:
  vector<int> in, out, p, rin, sz, nxt, h;
  vector<vector<int>> edges;
  int t;
  void prep(int on, int par) {
    sz[on] = 1;
    p[on] = par;
    for(int i = 0; i < (int) edges[on].size(); i++) {</pre>
      int &u = edges[on][i];
```

```
if(u == par) {
        swap(u, edges[on].back());
        edges[on].pop_back();
        i--;
      } else {
       h[u] = 1 + h[on];
        prep(u, on);
        sz[on] += sz[u];
        if(sz[u] > sz[edges[on][0]]) {
          swap(edges[on][0], u);
  void hld(int on) {
    in[on] = t++;
    rin[in[on]] = on;
    for(auto u : edges[on]) {
     nxt[u] = (u == edges[on][0] ? nxt[on] : u);
     hld(u);
    out[on] = t;
};
```

## 4.11 Centroid Decomposition

```
//Centroid decomposition1
void dfsSize(int v, int pa) {
  sz[v] = 1;
  for(int u : adj[v]) {
   if (u == pa || rem[u]) continue;
    dfsSize(u, v);
    sz[v] += sz[u];
int getCentroid(int v, int pa, int tam) {
  for(int u : adj[v]) {
    if (u == pa || rem[u]) continue;
    if (2 * sz[u] > tam) return getCentroid(u, v, tam);
  return v;
void decompose (int v, int pa = -1) {
  //cout << v << ' ' << pa << '\n';
 dfsSize(v, pa);
 int c = getCentroid(v, pa, sz[v]);
  //cout << c << '\n';
 par[c] = pa;
  rem[c] = 1;
  for(int u : adj[c]) {
    if (!rem[u] && u != pa) decompose(u, c);
  adi[c].clear();
```

```
//Centroid decomposition2
void dfsSize(int v, int par) {
  sz[v] = 1;
  for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   dfsSize(u, v);
   sz[v] += sz[u];
int getCentroid(int v, int par, int tam) {
  for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   if (2 * sz[u] > tam) return getCentroid(u, v, tam);
  return v;
void setDis(int v, int par, int nv, int d) {
  dis[v][nv] = d;
  for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   setDis(u, v, nv, d + 1);
void decompose(int v, int par, int nv) {
  dfsSize(v, par);
  int c = getCentroid(v, par, sz[v]);
  ct[c] = par;
  removed[c] = 1;
  setDis(c, par, nv, 0);
  for(int u : adj[c]) {
   if (!removed[u]) {
      decompose(u, c, nv + 1);
```

#### 4.12 Sack

```
void solve(int a, int p, bool f) {
  int big = -1;
  for(auto &b : adj[a]) {
    if(b != p && (big == -1 || en[b]-st[b] > en[big]-st[big])) {
      big = b;
    }
  }
  for(auto &b : adj[a]) {
    if(b == p || b == big) continue;
    solve(b, a, 0);
  }
  if(~big) solve(big, a, 1);
  add(cnt[v[a]], -1);
  cnt[v[a]]++;
  add(cnt[v[a]], +1);
  for(auto &b : adj[a]) {
    if(b == p || b == big) continue;
    for(int i = st[b]; i < en[b]; i++) {</pre>
```

```
add(cnt[ett[i]], -1);
    cnt[ett[i]]++;
    add(cnt[ett[i]], +1);
}

for(auto &q: Q[a]) {
    ans[q.first] = query(mx-1)-query(q.second-1);
}

if(!f) {
    for(int i = st[a]; i < en[a]; i++) {
        add(cnt[ett[i]], -1);
        cnt[ett[i]]--;
        add(cnt[ett[i]], +1);
    }
}</pre>
```

# 4.13 Hungarian Algorithm - Maximum Cost Matching

```
const int inf = 0x3f3f3f3f3f;
int n, w[ms][ms], maxm;
int lx[ms], ly[ms], xy[ms], yx[ms];
int slack[ms], slackx[ms], prev[ms];
bool S[ms], T[ms];
void init_labels() {
  memset(lx, 0, sizeof lx); memset(ly, 0, sizeof ly);
  for (int x = 0; x < n; x++) for (int y = 0; y < n; y++) {
    lx[x] = max(lx[x], cos[x][y]);
void updateLabels() {
  int delta = inf;
  for(int y = 0; y < n; y++) if(!T[y]) delta = min(delta, slack[y]);</pre>
  for(int x = 0; x < n; x++) if(S[x]) lx[x] -= delta;
  for(int y = 0; y < n; y++) if(T[y]) ly[y] += delta;
  for(int y = 0; y < n; y++) if(!T[y]) slack[y] -= delta;</pre>
void addTree(int x, int prevx) {
 S[x] = 1; prev[x] = prevx;
  for(int y = 0; y < n; y++) if(lx[x] + ly[y] - w[x][y] < slack[y]) {
    slack[y] = lx[x] + ly[y] - cost[x][y];
    slackx[y] = x;
void augment() {
  if (maxm == n) return;
  int x, y, root;
  int q[ms], wr = 0, rd = 0;
 memset(S, 0, sizeof S); memset(T, 0, sizeof T);
 memset (prev, -1, sizeof prev);
  for (int x = 0; x < n; x++) if (xy[x] == -1) {
    q[wr++] = root = x;
    prev[x] = -2;
    S[x] = 1;
    break;
```

```
for (int y = 0; y < n; y++) {
    slack[y] = lx[root] + ly[y] - w[root][y];
    slackx[y] = root;
  while(true) {
    while(rd < wr) {</pre>
      x = q[rd++];
      for (y = 0; y < n; y++) if (w[x][y] == 1x[x] + 1y[y] && !T[y]) {
        if(yx[y] == -1) break;
        T[y] = 1;
        q[wr++] = yx[y];
        addTree(yx[y], x);
      if(y < n) break;</pre>
    if(y < n) break;</pre>
    updateLabels();
    wr = rd = 0:
    for(y = 0; y < n; y++) if(!T[y] && !slack[y]) {
      if(yx[y] == -1) {
        x = slackx[v];
        break;
      } else {
        T[y] = true;
        if(!S[yx[y]]) {
          q[wr++] = yx[y];
          addTree(yx[y], slackx[y]);
    if(y < n) break;</pre>
  if(y < n) {
    for (int cx = x, cy = y, ty; cx != -2; cx = prev[cx], cy = ty) {
      ty = xy[cx];
      yx[cy] = cx;
      xy[cx] = cy;
    augment();
int hungarian() {
  int ans = 0; maxm = 0;
  memset(xy, -1, sizeof xy); memset(yx, -1, sizeof yx);
  initLabels(); augment();
  for (int x = 0; x < n; x++) ans += w[x][xy[x]];
  return ans;
```

# 5 Math

#### 5.1 Chinese Remainder Theorem

```
#include<bits/stdc++.h>
using namespace std;
const long long N = 20;
long long GCD(long long a, long long b) {
  return (b == 0) ? a : GCD(b, a % b);
inline long long get_LCM(long long a, long long b) {
  return a / GCD(a, b) * b;
inline long long normalize(long long x, long long mod) {
 x %= mod:
  if (x < 0) x += mod;
  return x;
struct GCD_type {
  long long x, y, d;
GCD_type ex_GCD(long long a, long long b) {
  if (b == 0) return {1, 0, a};
  GCD_type pom = ex_GCD(b, a % b);
  return {pom.y, pom.x - a / b * pom.y, pom.d};
long long testCases;
long long t;
long long a[N], n[N], ans, LCM;
int main() {
 ios_base::sync_with_stdio(0);
  cin.tie(0);
  t = 2;
  long long T;
  cin >> T;
  while(T--)
    for(long long i = 1; i <= t; i++) {</pre>
      cin >> a[i] >> n[i];
      normalize(a[i], n[i]);
    ans = a[1];
    LCM = n[1];
    bool impossible = false;
    for (long long i = 2; i \le t; i++) {
      auto pom = ex_GCD(LCM, n[i]);
      long long x1 = pom.x;
      long long d = pom.d;
      if((a[i] - ans) % d != 0) {
        impossible = true;
      ans = normalize(ans + x1 * (a[i] - ans) / d % (n[i] / d) * LCM,
          LCM * n[i] / d);
      LCM = get_LCM(LCM, n[i]);
    if (impossible) cout << "no solution\n";</pre>
    else cout << ans << " " << LCM << endl;
  return 0;
```

## 5.2 Diophantine Equations

```
int gcd_ext(int a, int b, int& x, int &y) {
  if (b == 0) {
   x = 1, y = 0;
   return a:
  int nx, ny;
  int gc = gcd_ext(b, a % b, nx, ny);
 y = nx - (a / b) * ny;
 return gc;
vector<int> diophantine(int D, vector<int> 1) {
  int n = l.size();
  vector<int> gc(n), ans(n);
  qc[n-1] = 1[n-1];
  for (int i = n - 2; i >= 0; i--) {
   int x, y;
   gc[i] = gcd_ext(l[i], gc[i + 1], x, y);
  if (D % qc[0] != 0) {
   return vector<int>();
  for (int i = 0; i < n; i++) {
   if (i == n - 1) {
      ans[i] = D / l[i];
      D = l[i] * ans[i];
      continue;
   int x, y;
   gcd_ext(l[i] / gc[i], gc[i + 1] / gc[i], x, y);
   ans[i] = (long long int) D / gc[i] * x % (gc[i + 1] / gc[i]);
   if (D < 0 \&\& ans[i] > 0) {
     ans[i] -= (gc[i + 1] / gc[i]);
   if (D > 0 \&\& ans[i] < 0) {
      ans[i] += (gc[i + 1] / gc[i]);
   D = l[i] * ans[i];
  return ans;
```

# 5.3 Discrete Logarithm

```
11 discreteLog(ll a, ll b, ll m) {
    // a^ans == b mod m
    // ou -1 se nao existir
    ll cur = a, on = 1;
    for(int i = 0; i < 100; i++) {
        cur = cur * a % m;
    }
    while(on * on <= m) {
        cur = cur * a % m;
        on++;
    }</pre>
```

```
map<11, ll> position;
for(ll i = 0, x = 1; i * i <= m; i++) {
    position[x] = i * on;
    x = x * cur % m;
}
for(ll i = 0; i <= on + 20; i++) {
    if(position.count(b)) {
        return position[b] - i;
    }
    b = b * a % m;
}
return -1;
}</pre>
```

## 5.4 Discrete Root

```
//x^k = a % mod

ll discreteRoot(ll k, ll a, ll mod) {
    ll g = primitiveRoot(mod);
    ll y = discreteLog(fexp(g, k, mod), a, mod);
    if (y == -1) {
        return y;
    }
    return fexp(g, y, mod);
}
```

## 5.5 Primitive Root

```
int primitiveRoot(int p) {
  vector<int> fact;
  int phi = p - 1, n = phi;
  for (int i = 2; i * i <= n; i++) {</pre>
    if (n \% i == 0) {
      fact.push_back(i);
      while (n \% i == 0) {
        n /= i;
  if (n > 1) {
    fact.push back(n);
  for (int res = 2; res <= p; res++) {</pre>
    bool ok = true;
    for (auto it : fact) {
      ok &= fexp(res, phi / it, p) != 1;
      if (!ok) {
        break:
    if (ok) {
      return res;
  return -1;
```

#### 5.6 Extended Euclides

```
// euclides estendido: acha u e v da equacao:
// u * x + v * y = gcd(x, y);
// u eh inverso modular de x no modulo y
// v eh inverso modular de y no modulo x

pair<11, ll> euclides(ll a, ll b) {
    ll u = 0, oldu = 1, v = 1, oldv = 0;
    while(b) {
        ll q = a / b;
        oldv = oldv - v * q;
        oldu = oldu - u * q;
        a = a - b * q;
        swap(a, b);
        swap(u, oldu);
        swap(v, oldv);
    }
    return make_pair(oldu, oldv);
}
```

## 5.7 Matrix Fast Exponentiation

```
const 11 \mod = 1e9+7;
const int m = 2; // size of matrix
struct Matrix {
  11 mat[m][m];
 Matrix operator * (const Matrix &p) {
    Matrix ans:
    for (int i = 0; i < m; i++)
      for (int j = 0; j < m; j++)
        for (int k = ans.mat[i][j] = 0; k < m; k++)
          ans.mat[i][j] = (ans.mat[i][j] + mat[i][k] * p.mat[k][j]) %
              mod:
    return ans;
};
Matrix fExp(Matrix a, ll b) {
 Matrix ans;
  for (int i = 0; i < m; i++) for (int j = 0; j < m; j++)
    ans.mat[i][j] = i == j;
  while(b) {
    if(b \& 1) ans = ans * a;
    a = a * a;
    b >>= 1;
  return ans;
```

## 5.8 FFT - Fast Fourier Transform

```
typedef double ld;
const ld PI = acos(-1);
```

```
struct Complex {
        ld real, imag:
        Complex conj() { return Complex(real, -imag); }
        Complex(1d = 0, 1d = 0): real(a), imag(b) {}
        Complex operator + (const Complex &o) const { return Complex(
            real + o.real, imag + o.imag); }
        Complex operator - (const Complex &o) const { return Complex(
            real - o.real, imag - o.imag); }
        Complex operator * (const Complex &o) const { return Complex(
            real * o.real - imag * o.imag, real * o.imag + imag * o.
        Complex operator / (ld o) const { return Complex(real / o,
            imag / o); }
        void operator *= (Complex o) { *this = *this * o; }
        void operator /= (ld o) { real /= o, imag /= o; }
};
typedef std::vector<Complex> CVector;
const int ms = 1 << 22:
int bits[ms];
Complex root[ms];
void initFFT() {
        root[1] = Complex(1);
        for(int len = 2; len < ms; len += len) {</pre>
                Complex z(cos(PI / len), sin(PI / len));
                for(int i = len / 2; i < len; i++) {</pre>
                        root[2 * i] = root[i];
                        root[2 * i + 1] = root[i] * z;
void pre(int n) {
        int LOG = 0;
        while (1 << (LOG + 1) < n) {
                LOG++;
        for (int i = 1; i < n; i++) {
                bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << LOG);
CVector fft(CVector a, bool inv = false) {
        int n = a.size();
        pre(n);
        if(inv) {
                std::reverse(a.begin() + 1, a.end());
        for(int i = 0; i < n; i++) {
                int to = bits[i];
                if(to > i) {
                        std::swap(a[to], a[i]);
        for (int len = 1; len < n; len \star= 2) {
                for (int i = 0; i < n; i += 2 * len) {
```

```
for(int j = 0; j < len; j++) {
                                Complex u = a[i + j], v = a[i + j +
                                    len] * root[len + j];
                                a[i + j] = u + v;
                                a[i + j + len] = u - v;
        if(inv) {
                for (int i = 0; i < n; i++)
                       a[i] /= n;
        return a;
void fft2in1(CVector &a, CVector &b) {
        int n = (int) a.size();
        for (int i = 0; i < n; i++) {
               a[i] = Complex(a[i].real, b[i].real);
        auto c = fft(a);
        for(int i = 0; i < n; i++) {
                a[i] = (c[i] + c[(n-i) % n].conj()) * Complex(0.5, 0);
                b[i] = (c[i] - c[(n-i) % n].conj()) * Complex(0, -0.5)
                   ;
void ifft2in1(CVector &a, CVector &b) {
        int n = (int) a.size();
        for(int i = 0; i < n; i++) {</pre>
                a[i] = a[i] + b[i] * Complex(0, 1);
        a = fft(a, true);
        for (int i = 0; i < n; i++) {
              b[i] = Complex(a[i].imag, 0);
                a[i] = Complex(a[i].real, 0);
std::vector<long long> mod_mul(const std::vector<long long> &a, const
    std::vector<long long> &b, long long cut = 1 << 15) {</pre>
        // TODO cut memory here by /2
        int n = (int) a.size();
        CVector C[4];
        for (int i = 0; i < 4; i++) {
               C[i].resize(n);
        for (int i = 0; i < n; i++) {
                C[0][i] = a[i] % cut;
                C[1][i] = a[i] / cut;
                C[2][i] = b[i] % cut;
                C[3][i] = b[i] / cut;
        fft2in1(C[0], C[1]);
        fft2in1(C[2], C[3]);
        for (int i = 0; i < n; i++) {
               // 00, 01, 10, 11
                Complex cur[4];
                for(int j = 0; j < 4; j++) cur[j] = C[j/2+2][i] * C[j
```

```
% 2][i];
                for(int j = 0; j < 4; j++) C[j][i] = cur[j];</pre>
        ifft2in1(C[0], C[1]);
       ifft2in1(C[2], C[3]);
        std::vector<long long> ans(n, 0);
        for (int i = 0; i < n; i++) {
                // if there are negative values, care with rounding
                ans[i] += (long long) (C[0][i].real + 0.5);
                ans[i] += (long long) (C[1][i].real + C[2][i].real +
                    0.5) * cut;
                ans[i] += (long long) (C[3][i].real + 0.5) * cut * cut
        return ans;
std::vector<int> mul(const std::vector<int> &a, const std::vector<int>
     } (d&
       int n = 1:
        while (n - 1 < (int) a.size() + (int) b.size() - 2) n += n;
        CVector poly(n);
        for (int i = 0; i < n; i++) {
                if(i < (int) a.size()) {
                        poly[i].real = a[i];
                if(i < (int) b.size()) {
                        poly[i].imag = b[i];
        poly = fft(poly);
        for (int i = 0; i < n; i++) {
                poly[i] *= poly[i];
       poly = fft(poly, true);
        std::vector<int> c(n, 0);
        for(int i = 0; i < n; i++) {
               c[i] = (int) (poly[i].imag / 2 + 0.5);
        while (c.size() > 0 && c.back() == 0) c.pop_back();
        return c;
```

## 5.9 NTT - Number Theoretic Transform

```
long long int mod = (11911 << 23) + 1, c_root = 3;

namespace NTT {
  typedef long long int 11;

ll fexp(11 base, 11 e) {
    11 ans = 1;
    while(e > 0) {
        if (e & 1) ans = ans * base % mod;
        base = base * base % mod;
        e >>= 1;
    }
    return ans;
}
```

```
11 inv_mod(ll base) {
    return fexp(base, mod - 2);
  void ntt(vector<ll>& a, bool inv) {
    int n = (int) a.size();
    if (n == 1) return;
    for (int i = 0, j = 0; i < n; i++) {
      if (i > j) {
        swap(a[i], a[j]);
      for(int 1 = n / 2; (j = 1) < 1; 1 >>= 1);
    for(int sz = 1; sz < n; sz <<= 1) {</pre>
      11 delta = fexp(c_root, (mod - 1) / (2 * sz)); //delta = w_2sz
      if (inv) {
        delta = inv mod(delta);
      for (int i = 0; i < n; i += 2 * sz) {
       11 w = 1;
        for (int j = 0; j < sz; j++) {
         11 u = a[i + j], v = w * a[i + j + sz] % mod;
          a[i + j] = (u + v + mod) % mod;
          a[i + j] = (a[i + j] + mod) % mod;
          a[i + j + sz] = (u - v + mod) % mod;
          a[i + j + sz] = (a[i + j + sz] + mod) % mod;
          w = w * delta % mod;
    if (inv) {
     11 inv_n = inv_mod(n);
     for(int i = 0; i < n; i++) {</pre>
       a[i] = a[i] * inv_n % mod;
    for (int i = 0; i < n; i++) {
      a[i] %= mod:
      a[i] = (a[i] + mod) % mod;
  void multiply(vector<ll> &a, vector<ll> &b, vector<ll> &ans) {
    int lim = (int) max(a.size(), b.size());
    int n = 1;
    while(n < lim) n <<= 1;
    n <<= 1;
    a.resize(n);
    b.resize(n);
    ans.resize(n);
    ntt(a, false);
    ntt(b, false);
    for(int i = 0; i < n; i++) {
     ans[i] = a[i] * b[i] % mod;
    ntt(ans, true);
};
```

#### 5.10 Miller and Rho

```
typedef long long int 11;
bool overflow(ll a, ll b) {
  return b && (a >= (111 << 62) / b);
11 add(11 a, 11 b, 11 md) {
  return (a + b) % md;
11 mul(ll a, ll b, ll md) {
  if (!overflow(a, b)) return (a * b) % md;
  11 \text{ ans} = 0;
  while(b) {
    if (b & 1) ans = add(ans, a, md);
    a = add(a, a, md);
    b >>= 1;
  return ans;
11 fexp(ll a, ll e, ll md) {
 11 \text{ ans} = 1;
  while(e) {
    if (e & 1) ans = mul(ans, a, md);
    a = mul(a, a, md);
    e >>= 1;
  return ans;
11 my_rand() {
 11 ans = rand();
  ans = (ans << 31) \mid rand();
  return ans;
11 gcd(ll a, ll b) {
  while(b) {
    11 t = a % b;
    a = b;
    b = t;
  return a;
bool miller(ll p, int iteracao) {
  if(p < 2) return 0:
  if(p % 2 == 0) return (p == 2);
  11 s = p - 1;
  while(s \% 2 == 0) s >>= 1;
  for(int i = 0; i < iteracao; i++) {</pre>
   11 a = rand() % (p - 1) + 1, temp = s;
    11 mod = fexp(a, temp, p);
    while (temp != p - 1 && mod != 1 && mod != p - 1) {
      mod = mul(mod, mod, p);
      temp <<= 1;
```

```
if(mod != p - 1 && temp % 2 == 0) return 0;
 return 1;
11 rho(ll n) {
 if (n == 1 || miller(n, 10)) return n;
 if (n % 2 == 0) return 2;
  while(1) {
   11 x = my_rand() % (n - 2) + 2, y = x;
   11 c = 0, cur = 1;
   while(c == 0) {
      c = my_rand() % (n - 2) + 1;
   while(cur == 1) {
      x = add(mul(x, x, n), c, n);
     y = add(mul(y, y, n), c, n);
     y = add(mul(y, y, n), c, n);
     cur = gcd((x >= y ? x - y : y - x), n);
   if (cur != n) return cur;
```

## 5.11 Determinant using Mod

// by zchao1995

```
// Determinante com coordenadas inteiras usando Mod
11 mat[ms][ms];
11 det (int n) {
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      mat[i][j] %= mod;
  11 \text{ res} = 1;
  for (int i = 0; i < n; i++) {</pre>
    if (!mat[i][i])
      bool flag = false;
      for (int j = i + 1; j < n; j++) {
        if (mat[j][i]) {
          flag = true;
          for (int k = i; k < n; k++) {
            swap (mat[i][k], mat[j][k]);
          res = -res;
          break;
      if (!flag) {
        return 0;
    for (int j = i + 1; j < n; j++) {
      while (mat[j][i]) {
        11 t = mat[i][i] / mat[j][i];
        for (int k = i; k < n; k++) {
          mat[i][k] = (mat[i][k] - t * mat[j][k]) % mod;
```

```
swap (mat[i][k], mat[j][k]);
    res = -res;
    }
} res = (res * mat[i][i]) % mod;
}
return (res + mod) % mod;
}
```

## 5.12 Lagrange Interpolation

```
class LagrangePoly {
public:
        LagrangePoly(std::vector<long long> _a) {
                //f(i) = \underline{a}[i]
                //interpola o vetor em um polinomio de grau y.size() -
                y = _a;
                den.resize(y.size());
                int n = (int) y.size();
                for (int i = 0; i < n; i++) {
                        y[i] = (y[i] % MOD + MOD) % MOD;
                         den[i] = ifat[n - i - 1] * ifat[i] % MOD;
                        if((n - i - 1) % 2 == 1) {
                                 den[i] = (MOD - den[i]) % MOD;
        long long getVal(long long x) {
                int n = (int) y.size();
                x %= MOD;
                if(x < n) {
                         //return y[(int) x];
                std::vector<long long> 1, r;
                l.resize(n);
                1[0] = 1;
                for(int i = 1; i < n; i++) {
                        l[i] = l[i - 1] * (x - (i - 1) + MOD) % MOD;
                r.resize(n);
                r[n - 1] = 1;
                for (int i = n - 2; i >= 0; i--) {
                        r[i] = r[i + 1] * (x - (i + 1) + MOD) % MOD;
                long long ans = 0;
                for(int i = 0; i < n; i++) {</pre>
                        long long coef = l[i] * r[i] % MOD;
                         ans = (ans + coef * y[i] % MOD * den[i]) % MOD
                return ans;
private:
        std::vector<long long> y, den;
};
```

## 6 Miscellaneous

## 6.1 LIS - Longest Increasing Subsequence

```
int arr[ms], lisArr[ms], n;
// int bef[ms], pos[ms];
int lis() {
  int len = 1;
  lisArr[0] = arr[0];
  // bef[0] = -1;
  for(int i = 1; i < n; i++) {</pre>
    // upper_bound se non-decreasing
    int x = lower_bound(lisArr, lisArr + len, arr[i]) - lisArr;
    len = max(len, x + 1);
    lisArr[x] = arr[i];
    // pos[x] = i;
    // bef[i] = x ? pos[x-1] : -1;
  return len;
vi getLis() {
  int len = lis();
  for (int i = pos[lisArr[len - 1]]; i >= 0; i = bef[i]) {
   ans.push_back(arr[i]);
  reverse(ans.begin(), ans.end());
  return ans;
```

# 6.2 Ternary Search

```
// R
for(int i = 0; i < LOG; i++) {
  long double m1 = (A * 2 + B) / 3.0;
  long double m2 = (A + 2 * B) / 3.0;</pre>
```

```
if(f(m1) > f(m2))
    A = m1;
else
    B = m2;
}
ans = f(A);

// Z
while(B - A > 4) {
    int m1 = (A + B) / 2;
    int m2 = (A + B) / 2 + 1;
    if(f(m1) > f(m2))
        A = m1;
else
        B = m2;
}
ans = inf;
for(int i = A; i <= B; i++) ans = min(ans , f(i));</pre>
```

## 6.3 Count Sort

```
int H[(1 << 15) +1], to[mx], b[mx];
void sort(int m, int a[]) {
  memset(H, 0, sizeof H);
  for (int i = 1; i <= m; i++) {</pre>
    H[a[i] % (1 << 15)] ++;
  for (int i = 1; i < 1<<15; i++) {
    H[i] += H[i-1];
  for (int i = m; i; i--) {
    to[i] = H[a[i] % (1 << 15)] --;
  for (int i = 1; i <= m; i++) {
    b[to[i]] = a[i];
  memset(H, 0, sizeof H);
  for (int i = 1; i <= m; i++) {
   H[b[i]>>15]++;
  for (int i = 1; i < 1 << 15; i++) {
    H[i] += H[i-1];
  for (int i = m; i : i--) {
    to[i] = H[b[i] >> 15] --;
  for (int i = 1; i <= m; i++) {</pre>
    a[to[i]] = b[i];
```

# 6.4 Random Number Generator

```
// mt19937_64 se LL
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
// Random_Shuffle
shuffle(v.begin(), v.end(), rng);
// Random number in interval
```

```
unordered_map<int,int>mp;
mp.reserve(1024);
mp.max_load_factor(0.25);
```

## 6.5 Rectangle Hash

```
namespace {
  struct safe hash {
    static uint64_t splitmix64(uint64_t x) {
      // http://xorshift.di.unimi.it/splitmix64.c
      x += 0x9e3779b97f4a7c15;
      x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
      x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
      return x ^ (x >> 31);
    size_t operator()(uint64_t x) const {
      static const uint64_t FIXED_RANDOM = std::chrono::steady_clock::
          now().time_since_epoch().count();
      return splitmix64(x + FIXED_RANDOM);
  };
struct rect {
  int x1, y1, x2, y2; // x1 < x2, y1 < y2
  rect (int x1, int y1, int x2, int y2) : x1(x1), x2(x2), y1(y1), y2(
      y2) {};
  rect inter (rect other) {
   int x3 = max(x1, other.x1);
   int y3 = max(y1, other.y1);
   int x4 = min(x2, other.x2);
   int y4 = min(y2, other.y2);
    return rect (x3, y3, x4, y4);
  uint64_t get_hash() {
   safe_hash sh;
   uint64_t ret = sh(x1);
   ret ^= sh(ret ^ v1);
   ret ^= sh(ret ^ x2);
    ret ^= sh(ret ^ y2);
   return ret:
};
```

# 6.6 Unordered Map Tricks

```
// pair<int, int> hash function
struct HASH{
    size_t operator() (const pair<int,int>&x) const{
      return (size_t) x.first * 37U + (size_t) x.second;
    }
};
```

### 6.7 Submask Enumeration

```
for (int s=m; ; s=(s-1)&m) {
    ... you can use s ...
    if (s==0) break;
}
```

## 6.8 Sum over Subsets DP

## 6.9 Java Fast I/O

```
import java.io.OutputStream;
import java.io.IOException;
import java.io.InputStream;
import java.io.PrintWriter;
import java.util.Arrays;
import java.util.Random;
import java.io.IOException;
import java.io.InputStreamReader;
import java.util.StringTokenizer;
import java.io.BufferedReader;
import java.io.InputStream;
import java.util.*;
import java.io.*;
// src petr
public class Main {
  public static void main(String[] args) {
    InputStream inputStream = System.in;
    OutputStream outputStream = System.out;
    InputReader in = new InputReader(inputStream);
    PrintWriter out = new PrintWriter(outputStream);
    TaskA solver = new TaskA();
    solver.solve(1, in, out);
    out.close();
  static class TaskA {
    public void solve(int testNumber, InputReader in, PrintWriter out)
```

```
static class InputReader {
 public BufferedReader reader;
 public StringTokenizer tokenizer;
 public InputReader(InputStream stream) {
   reader = new BufferedReader(new InputStreamReader(stream),
        32768);
    tokenizer = null;
 public String next() {
   while (tokenizer == null || !tokenizer.hasMoreTokens()) {
      try {
       tokenizer = new StringTokenizer(reader.readLine());
      } catch (IOException e) {
       throw new RuntimeException(e);
   return tokenizer.nextToken();
 public int nextInt() {
   return Integer.parseInt(next());
```

## **6.10** Dates

```
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
  return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 \star (m - 2 - (m - 14) / 12 \star 12) / 12 -
    3 \star ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = id + 68569;
  n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x = 1461 * i / 4 - 31;
  \dot{j} = 80 * x / 2447;
  d = x - 2447 * j / 80;
 x = \frac{1}{2} / 11;
 m = 1 + 2 - 12 * x:
 y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
```

# 7 String Algorithms

#### 7.1 KMP

```
string p, t;
int b[ms], n, m;
void kmpPreprocess() {
  int i = 0, j = -1;
  b[0] = -1;
  while (i < m) {
    while (j \ge 0 \&\& p[i] != p[j]) j = b[j];
    b[++i] = ++j;
void kmpSearch() {
  int i = 0, j = 0, ans = 0;
  while (i < n) {
    while (j \ge 0 \&\& t[i] != p[j]) j = b[j];
    <u>i</u>++; j++;
    if(i == m) {
      //ocorrencia agui comecando em i - j
      j = b[j];
  return ans;
```

## 7.2 KMP Automaton

```
const int limit =
vector<vector<int>> build_automaton(string s) {
    s += '#'; //tem que ser diferente de todos os caracteres
    int n = (int) s.size();
    vector<vector<int>> ans(n, vector<int>(limit));
    vector<int> fail(n);
    for (int i = 0; i < n; i++) {</pre>
        for (int j = 0; j < limit; j++) {
            if (i == 0) {
                if (s[i] == j + 'a') {
                    ans[i][j] = i + 1;
                } else {
                    ans[i][j] = 0;
            } else {
                if (s[i] == j + 'a') {
                    ans[i][j] = i + 1;
                } else {
                    ans[i][j] = ans[fail[i - 1]][j];
        if (i == 0) {
            continue;
```

```
}
int j = fail[i - 1];
while (j > 0 && s[i] != s[j]) {
    j = fail[j - 1];
}
fail[i] = j + (s[i] == s[j]);
}
return ans;
}
```

#### 7.3 Trie

```
int trie[ms][sigma], terminal[ms], z;
void init() {
 memset(trie[0], -1, sizeof trie[0]);
  z = 1;
int get_id(char c) {
  return c - 'a';
void insert(string &p) {
  int cur = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int id = get_id(p[i]);
    if(trie[cur][id] == -1) {
      memset(trie[z], -1, sizeof trie[z]);
      trie[cur][id] = z++;
    cur = trie[cur][id];
  terminal[cur]++;
int count(string &p) {
  int cur = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int id = get_id(p[i]);
    if(trie[cur][id] == -1) {
      return false:
    cur = trie[cur][id];
  return terminal[cur];
```

## 7.4 Aho-Corasick

```
// Construa a Trie do seu dicionario com o codigo acima
int fail[ms];
queue<int> q;

void buildFailure() {
   q.push(0);
   while(!q.empty()) {
```

```
int node = q.front();
    q.pop();
    for(int pos = 0; pos < sigma; pos++) {</pre>
      int &v = trie[node][pos];
      int f = node == 0 ? 0 : trie[fail[node]][pos];
      if(v == -1) {
        v = f;
      } else {
        fail[v] = f;
        q.push(v);
        // juntar as informacoes da borda para o V ja q um match em V
            implica um match na borda
        terminal[v] += terminal[f];
int search(string &txt) {
 int node = 0;
 int ans = 0;
 for(int i = 0; i < txt.length(); i++) {</pre>
   int pos = get_id(txt[i]);
   node = trie[node][pos];
    // processar informacoes no no atual
    ans += terminal[node];
 return ans;
```

# 7.5 Algoritmo de Z

```
string s;
int fz[ms], n;

void zfunc() {
  fz[0] = n;
  for(int i = 1, l = 0, r = 0; i < n; i++) {
    fz[i] = max(0, min(r-i, fz[i-1]));
    while(s[fz[i]] == s[i+fz[i]]) ++fz[i];
    if(i + fz[i] > r) {
      l = i;
      r = i + fz[i];
    }
}
```

# 7.6 Suffix Array

```
namespace SA {
  typedef pair<int, int> ii;

vector<int> buildSA(string s) {
  int n = (int) s.size();
  vector<int> ids(n), pos(n);
  vector<ii> pairs(n);
  for(int i = 0; i < n; i++) {
  ids[i] = i;</pre>
```

```
pairs[i] = ii(s[i], -1);
    sort(ids.begin(), ids.end(), [&](int a, int b) -> bool {
      return pairs[a] < pairs[b];</pre>
    });
    int on = 0;
    for (int i = 0; i < n; i++) {
      if (i && pairs[ids[i - 1]] != pairs[ids[i]]) on++;
      pos[ids[i]] = on;
    for(int offset = 1; offset < n; offset <<= 1) {</pre>
      //ia tao ordenados pelos primeiros offset caracteres
      for (int i = 0; i < n; i++) {
        pairs[i].first = pos[i];
        if (i + offset < n) {
          pairs[i].second = pos[i + offset];
        } else {
          pairs[i].second = -1;
      sort(ids.begin(), ids.end(), [&](int a, int b) -> bool {
        return pairs[a] < pairs[b];</pre>
      });
      int on = 0;
      for (int i = 0; i < n; i++) {
        if (i && pairs[ids[i - 1]] != pairs[ids[i]]) on++;
        pos[ids[i]] = on;
    return ids;
  vector<int> buildLCP(string s, vector<int> sa) {
    int n = (int) s.size();
    vector<int> pos(n), lcp(n, 0);
    for(int i = 0; i < n; i++) {</pre>
      pos[sa[i]] = i:
    int k = 0:
    for(int i = 0; i < n; i++) {
      if (pos[i] + 1 == n) {
       \mathbf{k} = 0:
        continue;
      int j = sa[pos[i] + 1];
      while(i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
      lcp[pos[i]] = k;
      k = \max(k - 1, 0);
    return lcp;
};
//nlogn
vector<int> suffix_array(const string& in) {
    int n = (int) in.size(), c = 0;
    vector<int> temp(n), pos2bckt(n), bckt(n), bpos(n), out(n);
    for (int i = 0; i < n; i++) out[i] = i;</pre>
    sort(out.begin(), out.end(), [&](int a, int b) { return in[a] < in</pre>
        [b]; });
```

```
for (int i = 0; i < n; i++) {</pre>
   bckt[i] = c;
    if (i + 1 == n || in[out[i]] != in[out[i + 1]]) c++;
/*Start*/
for (int h = 1; h < n \&\& c < n; h <<= 1) {// executes log n times
    for (int i = 0; i < n; i++) pos2bckt[out[i]] = bckt[i];</pre>
    for (int i = n - 1; i >= 0; i--) bpos[bckt[i]] = i;
    for (int i = 0; i < n; i++)
        if (out[i] >= n - h) temp[bpos[bckt[i]]++] = out[i];
    for (int i = 0; i < n; i++)
        if (out[i] >= h) temp[bpos[pos2bckt[out[i] - h]]++] = out[
            il - h:
    c = 0;
    for (int i = 0; i + 1 < n; i++) {
        int a = (bckt[i] != bckt[i + 1]) || (temp[i] >= n - h)
            | | (pos2bckt[temp[i + 1] + h] | = pos2bckt[temp[i] + h]
        bckt[i] = c;
        c += a;
   bckt[n - 1] = c++;
   temp.swap(out);
return out;
```

## 8 Teoremas e formulas uteis

#### 8.1 Grafos

```
Formula de Euler: V - E + F = 2 (para grafo planar)
Handshaking: Numero par de vertices tem grau impar
Kirchhoff's Theorem: Monta matriz onde Mi, i = Grau[i] e Mi, j = -1 se
    houver aresta i-j ou 0 caso contrario, remove uma linha e uma
    coluna qualquer e o numero de spanning trees nesse grafo eh o det
    da matriz
Grafo contem caminho hamiltoniano se:
Dirac's theorem: Se o grau de cada vertice for pelo menos n/2
Ore's theorem: Se a soma dos graus que cada par nao-adjacente de
    vertices for pelo menos n
Tem Catalan(N) Binary trees de N vertices
Tem Catalan (N-1) Arvores enraizadas com N vertices
Caley Formula: n^(n-2) arvores em N vertices com label
Prufer code: Cada etapa voce remove a folha com menor label e o label
    do vizinho eh adicionado ao codigo ate ter 2 vertices
Flow:
Max Edge-disjoint paths: Max flow com arestas com peso 1
Max Node-disjoint paths: Faz a mesma coisa mas separa cada vertice em
    um com as arestas de chegadas e um com as arestas de saida e uma
    aresta de peso 1 conectando o vertice com aresta de chegada com
    ele mesmo com arestas de saida
Konig's Theorem: minimum node cover = maximum matching se o grafo for
    bipartido, complemento eh o maximum independent set
```

- Min Node disjoint path cover: formar grafo bipartido de vertices duplicados, onde aresta sai do vertice tipo A e chega em tipo B, entao o path cover eh N matching
- Min General path cover: Mesma coisa mas colocando arestas de A pra B sempre que houver caminho de A pra B
- Dilworth's Theorem: Min General Path cover = Max Antichain (set de vertices tal que nao existe caminho no grafo entre vertices desse set)
- Hall's marriage: um grafo tem um matching completo do lado X se para cada subconjunto W de X,
  - $|W| \le |vizinhosW|$  onde |W| eh quantos vertices tem em W

## 8.2 Math

```
Goldbach's: todo numero par n > 2 pode ser representado com n = a + b
    onde a e b sao primos
Twin prime: existem infinitos pares p, p + 2 onde ambos sao primos
Legendre's: sempre tem um primo entre n^2 e (n+1)^2
Lagrange's: todo numero inteiro pode ser inscrito como a soma de 4
Zeckendorf's: todo numero pode ser representado pela soma de dois
    numeros de fibonnacis diferentes e nao consecutivos
Euclid's: toda tripla de pitagoras primitiva pode ser gerada com
    (n^2 - m^2, 2nm, n^2+m^2) onde n, m sao coprimos e um deles eh par
Wilson's: n \in P primo quando (n-1)! \mod n = n - 1
Mcnugget: Para dois coprimos x, y o maior inteiro que nao pode ser
    escrito como ax + by eh (x-1)(y-1)/2
Fermat: Se p eh primo entao a(p-1) % p = 1
Se x e m tambem forem coprimos entao x^k % m = x^(k \mod (m-1)) % m
Euler's theorem: x^(phi(m)) mod m = 1 onde phi(m) eh o totiente de
    euler
Chinese remainder theorem:
Para equações no formato x = al \mod ml, ..., x = an \mod mn onde todos
     os pares m1, ..., mn sao coprimos
Deixe Xk = m1*m2*..*mn/mk e Xk^-1 mod mk = inverso de Xk mod mk, entao
x = somatorio com k de 1 ate n de ak*Xk*(Xk,mk^-1 mod mk)
Para achar outra solucao so somar m1*m2*..*mn a solucao existente
Catalan number: exemplo expressoes de parenteses bem formadas
C0 = 1, Cn = somatorio de <math>i=0 \rightarrow n-1 de Ci*C(n-1+1)
outra forma: Cn = (2n \text{ escolhe } n) / (n+1)
Bertrand's ballot theorem: p votos tipo A e q votos tipo B com p>q,
    prob de em todo ponto ter mais As do que Bs antes dele = (p-q)/(p+
Se puder empates entao prob = (p+1-q)/(p+1), para achar quantidade de
    possibilidades nos dois casos basta multiplicar por (p + q escolhe
Propriedades de Coeficientes Binomiais:
Somatorio de k = 0 \rightarrow m de (-1)^k \star (n \text{ escolhe } k) = (-1)^m \star (n -1)
    escolhe m)
(N \text{ escolhe } K) = (N \text{ escolhe } N-K)
(N \text{ escolhe } K) = N/K * (n-1 \text{ escolhe } k-1)
Somatorio de k = 0 \rightarrow n de (n escolhe k) = 2^n
Somatorio de m = 0 \rightarrow n de (m \text{ escolhe } k) = (n+1 \text{ escolhe } k + 1)
Somatorio de k = 0 \rightarrow m de (n+k \text{ escolhe } k) = (n+m+1 \text{ escolhe } m)
Somatorio de k = 0 \rightarrow n de (n \text{ escolhe } k)^2 = (2n \text{ escolhe } n)
```

```
Somatorio de k = 0 ou 1 \rightarrow n de k*(n escolhe k) = n * 2^(n-1)
Somatorio de k = 0 \rightarrow n de (n-k \text{ escolhe } k) = \text{Fib}(n+1)
Hockey-stick: Somatorio de i = r \rightarrow n de (i escolhe r) = (n + 1)
    escolhe r + 1)
Vandermonde: (m+n \text{ escolhe } r) = \text{somatorio de } k = 0 -> r \text{ de } (m \text{ escolhe } k)
    ) \star (n escolhe r - k)
Burnside lemma: colares diferentes nao contando rotacoes quando m =
    cores e n = comprimento
(m^n + somatorio i = 1 - > n-1 de m^qcd(i, n))/n
Distribuicao uniforme a,a+1, ..., b Expected[X] = (a+b)/2
Distribuicao binomial com n tentativas de probabilidade p, X =
    P(X = x) = p^x * (1-p)^(n-x) * (n escolhe x) e E[X] = p*n
Distribuicao geometrica onde continuamos ate ter sucesso, X =
    tentativas:
    P(X = x) = (1-p)^(x-1) * p e E[X] = 1/p
Linearity of expectation: Tendo duas variaveis X e Y e constantes a e
    b, o valor esperado de aX + bY = a*E[X] + b*E[X]
```

## 8.3 Geometry

```
Formula de Euler: V - E + F = 2
Pick Theorem: Para achar pontos em coords inteiras num poligono Area =
     i + b/2 - 1 onde i eh o o numero de pontos dentro do poligono e b
     de pontos no perimetro do poligono
Two ears theorem: Todo poligono simples com mais de 3 vertices tem
    pelo menos 2 orelhas, vertices que podem ser removidos sem criar
    um crossing, remover orelhas repetidamente triangula o poligono
Incentro triangulo: (a(Xa, Ya) + b(Xb, Yb) + c(Xc, Yc))/(a+b+c) onde
    a = lado oposto ao vertice a, incentro eh onde cruzam as
    bissetrizes, eh o centro da circunferencia inscrita e eh
    equidistante aos lados
Delaunay Triangulation: Triangulacao onde nenhum ponto esta dentro de
    nenhum circulo circunscrito nos triangulos
Eh uma triangulacao que maximiza o menor angulo e a MST euclidiana de
    um conjunto de pontos eh um subconjunto da triangulacao
Brahmagupta s formula: Area cyclic quadrilateral
s = (a+b+c+d)/2
area = sqrt((s-a)*(s-b)*(s-c)*(s-d))
d = 0 \Rightarrow area = sqrt((s-a)*(s-b)*(s-c)*s)
```

#### 8.4 Mersenne's Primes

```
Primos de Mersenne 2^n - 1
Lista de Ns que resultam nos primeiros 41 primos de Mersenne:
2; 3; 5; 7; 13; 17; 19; 31; 61; 89; 107; 127; 521; 607; 1.279; 2.203;
2.281; 3.217; 4.253; 4.423; 9.689; 9.941; 11.213; 19.937; 21.701;
23.209; 44.497; 86.243; 110.503; 132.049; 216.091; 756.839;
859.433; 1.257.787; 1.398.269; 2.976.221; 3.021.377; 6.972.593;
13.466.917; 20.996.011; 24.036.583;
```