The AC is a lie - ICPC Library

ont	tent	ts
	ont	ontent

1	Stri	ng Algorithms	1
	1.1	String Alignment	
	1.2	KMP	
	1.3	Trie	
	1.4	Aho-Corasick	
	1.5	Algoritmo de Z	
	1.6	Suffix Array	
		·	
2	Dat	a Structures	9
_	2.1	BIT - Binary Indexed Tree	٠
	2.2	BIT 2D	
	2.3	BIT 2D Comprimida	
	$\frac{2.5}{2.4}$	Iterative Segment Tree	
	2.5	Iterative Segment Tree with Interval Updates	
	2.6	Recursive Segment Tree	
	2.7	Segment Tree with Lazy Propagation	
	2.8	Persistent Segment Tree	
	2.9	Color Updates Structure	
	2.10	Policy Based Structures	
	2.11	Heavy Light Decomposition	
	2.12	Centroid Decomposition	
	2.13	Sparse Table	
	2.14	Li Chao Tree	
3	Cra	ph Algorithms	ć
J	3.1	Dinic Max Flow	. 9
	3.2	Euler Path and Circuit	
	3.3	Articulation Points/Bridges/Biconnected Components	. :
	3.4	SCC - Strongly Connected Components / 2SAT	
	$3.4 \\ 3.5$	LCA - Lowest Common Ancestor	
	3.6	Sack	
	3.7	Min Cost Max Flow	
	3.8	Hungarian Algorithm - Maximum Cost Matching	
	3 .0	Trungarian Angorrania - Maximum Cost Materinig	. 11
4	Mat	հ	12
4			
	4.1	Discrete Logarithm	
	$\frac{4.2}{4.3}$	GCD - Greatest Common Divisor	
		D. 4 1 . 1 D 1; 1	
		Extended Euclides	. 13
	4.4	Fast Exponentiation	. 13
	$\frac{4.4}{4.5}$	Fast Exponentiation	. 13 . 13
	$4.4 \\ 4.5 \\ 4.6$	Fast Exponentiation	. 13 . 13 . 13
	4.4 4.5 4.6 4.7	Fast Exponentiation	. 13 . 13 . 13
	4.4 4.5 4.6 4.7 4.8	Fast Exponentiation	. 13 . 13 . 13 . 14
	4.4 4.5 4.6 4.7	Fast Exponentiation	. 13 . 13 . 13 . 14
F	4.4 4.5 4.6 4.7 4.8 4.9	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod	. 13 . 13 . 13 . 14 . 14
5	4.4 4.5 4.6 4.7 4.8 4.9	Fast Exponentiation	. 1; . 1; . 1; . 1; . 1 ² . 1;
5	4.4 4.5 4.6 4.7 4.8 4.9 Geo	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry	. 13 . 13 . 13 . 14 . 14 . 15
5	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull	. 1; . 1; . 1; . 1; . 1; . 12 . 1; . 16 . 16
5	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair	. 13 . 13 . 13 . 14 . 14 . 14 . 16 . 16 . 18
5	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points	. 13 . 13 . 14 . 15 . 14 . 14 . 15 . 16 . 16 . 18
5	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair	. 13 . 13 . 14 . 15 . 14 . 14 . 15 . 16 . 16 . 18
	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library	113 113 114 115 115 115 115 115 115 115 115 115
5	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points	. 13 . 13 . 14 . 15 . 14 . 14 . 15 . 16 . 16 . 18
	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library	13 13 14 15 16 18 18 18 18 18 18 18 18 18
	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyn	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library mamic Programming	13 13 14 15 16 18 18 18 18 18 18 18 18 18
	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyn 6.1	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library mamic Programming	13 13 14 15 16 18 18 18 18 18 18 18 18 18
6	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyn 6.1	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library mamic Programming Convex Hull Trick cellaneous	. 13 . 13 . 13 . 14 . 14 . 16 . 18 . 18 . 18 . 18 . 18
6	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyn 6.1	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library mamic Programming Convex Hull Trick cellaneous LIS - Longest Increasing Subsequence	. 13 . 15 . 15 . 16 . 16 . 16 . 18 . 18 . 18 . 18 . 18
6	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyr 6.1 Mis 7.1 7.2	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library mamic Programming Convex Hull Trick cellaneous LIS - Longest Increasing Subsequence Binary Search	. 13 . 13 . 13 . 14 . 14 . 14 . 18 . 16 . 18 . 18 . 18 . 18 . 18 . 20 . 20 . 20 . 20
6	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyn 6.1 Mis 7.1 7.2 7.3	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library manic Programming Convex Hull Trick cellaneous LIS - Longest Increasing Subsequence Binary Search Ternary Search	. 13 . 13 . 14 . 14 . 14 . 18 . 16 . 18 . 18 . 18 . 19 . 20 . 20 . 20 . 20 . 20 . 20 . 20 . 20
6	4.4 4.5 4.6 4.7 4.8 4.9 Geo 5.1 5.2 5.3 5.4 5.5 Dyr 6.1 Mis 7.1 7.2	Fast Exponentiation Matrix Fast Exponentiation FFT - Fast Fourier Transform NTT - Number Theoretic Transform Miller and Rho Determinant using Mod metry Geometry Convex Hull Closest Pair Intersection Points Java Geometry Library mamic Programming Convex Hull Trick cellaneous LIS - Longest Increasing Subsequence Binary Search	. 13 . 13 . 13 . 14 . 14 . 14 . 16 . 18 . 18 . 18 . 18 . 19 . 19 . 19 . 19 . 19 . 19 . 19 . 19

```
      8 Teoremas e formulas uteis
      22

      8.1 Grafos
      2

      8.2 Math
      2

      8.3 Geometry
      2

      8.4 Mersenne's Primes
      2
```

1 String Algorithms

1.1 String Alignment

```
int pd[ms][ms];
int edit_distance(string &a, string &b) {
    int n = a.size(), m = b.size();
    for(int i = 0; i <= n; i++) pd[i][0] = i;
    for(int j = 0; j <= m; j++) pd[0][j] = j;
    for(int i = 1; i <= n; i++) {
        for(int j = 1; j <= m; j++) {
            int del = pd[i][j-1] + 1;
            int ins = pd[i-1][j] + 1;
            int mod = pd[i-1][j-1] + (a[i-1] != b[j-1]);
            pd[i][j] = min(del, min(ins, mod));
        }
    }
    return pd[n][m];
}</pre>
```

1.2 KMP

```
string p, t;
int b[ms], n, m;
void kmpPreprocess() {
  int i = 0, j = -1;
 b[0] = -1;
  while (i < m) {
    while (j \ge 0 \&\& p[i] != p[j]) j = b[j];
    b[++i] = ++j;
void kmpSearch() {
  int i = 0, j = 0, ans = 0;
  while (i < n) {
    while(j \ge 0 \&\& t[i] != p[j]) j = b[j];
    <u>i</u>++; j++;
    if(j == m) {
      //ocorrencia agui comecando em i - j
      ans++;
      j = b[j];
  return ans;
```

1.3 Trie

```
int trie[ms][sigma], terminal[ms], z;
void init() {
 memset(trie[0], -1, sizeof trie[0]);
 z = 1;
int get_id(char c) {
  return c - 'a';
void insert(string &p) {
  int cur = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int id = get_id(p[i]);
    if(trie[cur][id] == -1)
      memset(trie[z], -1, sizeof trie[z]);
      trie[cur][id] = z++;
    cur = trie[cur][id];
  terminal[cur]++;
int count(string &p) {
  int cur = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int id = get_id(p[i]);
    if(trie[cur][id] == -1) {
      return false;
    cur = trie[cur][id];
  return terminal[cur];
```

} } int search(string &txt) { int node = 0; int ans = 0; for(int i = 0; i < txt.length(); i++) { int pos = get_id(txt[i]); node = trie[node][pos]; // processar informacoes no no atual ans += terminal[node]; } return ans; }</pre>

1.5 Algoritmo de Z

```
string s;
int fz[ms], n;

void zfunc() {
  fz[0] = n;
  for(int i = 1, l = 0, r = 0; i < n; i++) {
    if(l && i + fz[i-l] < r)
        fz[i] = fz[i-l];
    else {
        int j = min(l ? fz[i-l] : 0, i > r ? 0 : r - i);
        while(s[i+j] == s[j] && ++j);
        fz[i] = j; l = i; r = i + j;
    }
}
```

1.4 Aho-Corasick

```
// Construa a Trie do seu dicionario com o codigo acima
int fail[ms];
queue<int> q;
void buildFailure() {
  q.push(0);
  while(!q.empty()) {
    int node = q.front();
    q.pop();
    for(int pos = 0; pos < sigma; pos++) {</pre>
      int &v = trie[node][pos];
      int f = node == 0 ? 0 : trie[fail[node]][pos];
      if(v == -1) {
       v = f:
      } else {
       fail[v] = f;
        q.push(v);
        // juntar as informacoes da borda para o V ja q um match em V
            implica um match na borda
        terminal[v] += terminal[f];
```

1.6 Suffix Array

```
namespace SA {
  typedef pair<int, int> ii;
  vector<int> buildSA(string s) {
    int n = (int) s.size();
    vector<int> ids(n), pos(n);
    vector<ii>> pairs(n);
    for (int i = 0; i < n; i++) {
      ids[i] = i;
      pairs[i] = ii(s[i], -1);
    sort(ids.begin(), ids.end(), [&](int a, int b) -> bool {
      return pairs[a] < pairs[b];</pre>
    });
    int on = 0:
    for(int i = 0; i < n; i++) {</pre>
      if (i && pairs[ids[i - 1]] != pairs[ids[i]]) on++;
      pos[ids[i]] = on;
    for(int offset = 1; offset < n; offset <<= 1) {</pre>
      //ja tao ordenados pelos primeiros offset caracteres
      for(int i = 0; i < n; i++) {
```

```
pairs[i].first = pos[i];
        if (i + offset < n) {
          pairs[i].second = pos[i + offset];
          pairs[i].second = -1;
      sort(ids.begin(), ids.end(), [&](int a, int b) -> bool {
        return pairs[a] < pairs[b];</pre>
      });
      int on = 0;
      for (int i = 0; i < n; i++) {
        if (i && pairs[ids[i - 1]] != pairs[ids[i]]) on++;
        pos[ids[i]] = on;
   return ids;
  vector<int> buildLCP(string s, vector<int> sa) {
   int n = (int) s.size();
   vector<int> pos(n), lcp(n, 0);
    for (int i = 0; i < n; i++) {
      pos[sa[i]] = i;
   int k = 0;
    for(int i = 0; i < n; i++) {</pre>
      if (pos[i] + 1 == n) {
       \mathbf{k} = 0;
        continue;
      int j = sa[pos[i] + 1];
      while(i + k < n && j + k < n && s[i + k] == s[j + k]) k++;
      lcp[pos[i]] = k;
      k = \max(k - 1, 0);
   return lcp;
};
```

2 Data Structures

2.1 BIT - Binary Indexed Tree

```
int bit[ms], n;

void update(int v, int idx) {
  while(idx <= n) {
    bit[idx] += v;
    idx += idx & -idx;
  }
}

int query(int idx) {
  int r = 0;
  while(idx > 0) {
    r += bit[idx];
    idx -= idx & -idx;
}
```

```
}
return r;
```

2.2 BIT 2D

```
int bit[ms][ms], n, m;

void update(int v, int x, int y) {
    while(x <= n) {
        bit[x][y] += v;
        y += y&-y;
        }
        x += x&-x;
    }
}

int query(int x, int y) {
    int r = 0;
    while(x > 0) {
        while(y > 0) {
            r += bit[x][y];
            y -= y&-y;
        }
        x -= x&-x;
    }

    return r;
}
```

2.3 BIT 2D Comprimida

```
// by TFG
#include <vector>
#include <utility>
#include <algorithm>
typedef std::pair<int, int> ii;
struct Bit2D {
public:
  Bit2D(std::vector<ii> pts) {
    std::sort(pts.begin(), pts.end());
    for(auto a : pts) {
      if(ord.empty() || a.first != ord.back())
        ord.push_back(a.first);
    fw.resize(ord.size() + 1);
    coord.resize(fw.size());
    for(auto &a : pts)
      std::swap(a.first, a.second);
    std::sort(pts.begin(), pts.end());
    for(auto &a : pts) {
      std::swap(a.first, a.second);
      for(int on = std::upper_bound(ord.begin(), ord.end(), a.first) -
           ord.begin(); on < fw.size(); on += on & -on) {
        if(coord[on].empty() || coord[on].back() != a.second);
          coord[on].push_back(a.second);
```

```
for(int i = 0; i < fw.size(); i++) {</pre>
      fw[i].assign(coord[i].size() + 1, 0);
  void upd(int x, int y, int v) {
    for(int xx = std::upper bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx < fw.size(); xx += xx & -xx) {
      for(int yy = std::upper_bound(coord[xx].begin(), coord[xx].end()
          , y) - coord[xx].begin(); yy < fw[xx].size(); yy += yy & -yy
          ) {
        fw[xx][yy] += v;
  int gry(int x, int y) {
    int ans = 0:
    for(int xx = std::upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx > 0; xx -= xx & -xx) {
      for(int yy = std::upper_bound(coord[xx].begin(), coord[xx].end()
         , y) - coord[xx].begin(); yy > 0; yy -= yy & -yy) {
        ans += fw[xx][yy];
    return ans;
private:
  std::vector<int> ord;
  std::vector<std::vector<int>> fw, coord;
};
```

2.4 Iterative Segment Tree

```
int n, t[2 * ms];
void build() {
  for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1];
void update(int p, int value) { // set value at position p
  for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1];
int query(int 1, int r) {
  int res = 0;
  for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if(1&1) res += t[1++];
    if(r\&1) res += t[--r];
  return res:
// If is non-commutative
S query(int 1, int r) {
  S resl, resr;
  for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
  if (l\&1) resl = combine(resl, t[l++]);
```

```
if (r&1) resr = combine(t[--r], resr);
}
return combine(resl, resr);
```

2.5 Iterative Segment Tree with Interval Updates

```
int n, t[2 * ms];
void build() {
  for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1];
void update(int 1, int r, int value) {
  for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1)
    if(1&1) t[1++] += value;
    if(r&1) t[--r] += value;
int query(int p) {
  int res = 0;
  for (p += n; p > 0; p >>= 1) res += t[p];
  return res:
void push() { // push modifications to leafs
 for (int i = 1; i < n; i++) {
   t[i<<1] += t[i];
   t[i<<1|1] += t[i];
    t[i] = 0;
```

2.6 Recursive Segment Tree

```
int arr[4 * ms], seg[4 * ms], n;
void build(int idx = 0, int l = 0, int r = n - 1) {
  int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
  if(1 == r) {
    seg[idx] = arr[l];
    return;
 build(left, 1, mid); build(right, mid + 1, r);
  seq[idx] = seq[left] + seq[right];
int query(int L, int R, int idx = 0, int l = 0, int r = n - 1) {
 int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
 if(R < 1 \mid | L > r) return 0;
  if(L <= 1 && r <= R) return seg[idx];</pre>
  return query(L, R, left, l, mid) + query(L, R, right, mid + 1, r);
void update(int V, int I, int idx = 0, int l = 0, int r = n - 1) {
  int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
  if(1 > I \mid | r < I) return;
```

```
if(1 == r) {
    arr[I] = V;
    seg[idx] = V;
    return;
}
update(V, I, left, 1, mid); update(V, I, right, mid + 1, r);
seg[idx] = seg[left] + seg[right];
}
```

2.7 Segment Tree with Lazy Propagation

```
int arr[4 * ms], seq[4 * ms], lazy[4 * ms], n;
void build(int idx = 0, int l = 0, int r = n - 1) {
  int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
 lazy[idx] = 0;
 if(1 == r) {
   seg[idx] = arr[l];
   return;
 build(left, 1, mid); build(right, mid + 1, r);
  seg[idx] = seg[left] + seg[right];
void propagate(int idx, int 1, int r) {
  int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
  if(lazv[idx]) {
   seg[idx] += lazy[idx] * (r - l + 1);
   if(1 < r) {
     lazv[2*idx+1] += lazv[idx];
      lazy[2*idx+2] += lazy[idx];
    lazy[idx] = 0;
int query(int L, int R, int idx = 0, int l = 0, int r = n - 1) {
 int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
 propagate(idx, l, r);
 if(R < 1 \mid | L > r) return 0:
 if(L <= 1 && r <= R) return seg[idx];</pre>
 return query(L, R, left, 1, mid) + query(L, R, right, mid + 1, r);
void update(int V, int L, int R, int idx = 0, int 1 = 0, int r = n -1)
  int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
 propagate(idx, l, r);
  if(1 > R | | r < L) return;
 if(L \le 1 \&\& r \le R) {
   lazv[idx] += V;
   propagate(idx, l, r);
  update(V, L, R, left, l, mid); update(V, L, R, right, mid + 1, r);
  seq[idx] = seq[left] + seq[right];
```

2.8 Persistent Segment Tree

```
struct PSEGTREE{
 private:
   int z, t, sz, *tree, *L, *R, head[112345];
   void _build(int 1, int r, int on, vector<int> &v) {
      if(1 == r){
       tree[on] = v[1];
        return:
      L[on] = ++z;
      int mid = (1+r) >> 1;
      _build(l, mid, L[on], v);
     R[on] = ++z;
      _build(mid+1, r, R[on], v);
      tree[on] = tree[L[on]] + tree[R[on]];
    int _upd(int ql, int qr, int val, int l, int r, int on) {
      if(1 > qr \mid | r < ql) return on;
      int curr = ++z;
      if(1 >= ql && r <= qr) {
       tree[curr] = tree[on] + val;
       return curr:
      int mid = (1+r) >> 1;
      L[curr] = \_upd(ql, qr, val, l, mid, L[on]);
      R[curr] = \_upd(ql, qr, val, mid+1, r, R[on]);
      tree[curr] = tree[L[curr]] + tree[R[curr]];
      return curr;
   int _query(int ql, int qr, int l, int r, int on){
     if(1 > qr || r < ql) return 0;
      if(1 >= q1 && r <= qr) {
        return tree[on];
      int mid = (1+r) >> 1;
      return _query(ql, qr, l, mid, L[on]) + _query(ql, qr, mid+1, r,
          R[on]);
 public:
   PSEGTREE (vector<int> &v) {
     tree = new int[1123456];
      L = new int[1123456];
      R = new int[1123456];
     build(v):
    void build(vector<int> &v) {
     t = 0, z = 0;
      sz = v.size():
     head[0] = 0;
      _{\text{build}}(0, \text{sz-1}, 0, \text{v});
    void upd(int pos, int val, int idx){
     head[++t] = \_upd(pos, pos, val, 0, sz-1, head[idx]);
```

```
int query(int 1, int r, int idx) {
    return _query(1, r, 0, sz-1, head[idx]);
    }
};
```

2.9 Color Updates Structure

```
struct range {
  int 1. r:
  int v;
  range(int l = 0, int r = 0, int v = 0) : l(1), r(r), v(v) {}
 bool operator < (const range &a) const {</pre>
    return 1 < a.1;</pre>
};
set<range> ranges;
vector<range> update(int 1, int r, int v) { // [1, r)
  vector<range> ans;
  if(l >= r) return ans;
  auto it = ranges.lower_bound(1);
  if(it != ranges.begin()) {
    it--;
    if(it->r > 1) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, 1, cur.v));
      ranges.insert(range(l, cur.r, cur.v));
  it = ranges.lower_bound(r);
  if(it != ranges.begin()) {
    it--;
    if(it->r>r) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, r, cur.v));
      ranges.insert(range(r, cur.r, cur.v));
  for(it = ranges.lower_bound(l); it != ranges.end() && it->l < r; it</pre>
      ++) {
    ans.push_back(*it);
  ranges.erase(ranges.lower_bound(l), ranges.lower_bound(r));
  ranges.insert(range(l, r, v));
  return ans;
int query(int v) { // Substituir -1 por flag para quando nao houver
  auto it = ranges.upper_bound(v);
  if(it == ranges.begin()) {
    return -1;
  it--;
```

```
return it->r >= v ? it->v : -1;
```

2.10 Policy Based Structures

2.11 Heavy Light Decomposition

```
// HLD + Euler Tour by adamant
// OTREE3
#include <bits/stdc++.h>
using namespace std:
const int N = 112345, inf = 0x3f3f3f3f3f;
vector<int> q[N];
void add(int a, int b) {
        g[a].push_back(b), g[b].push_back(a);
int sz[N], par[N], h[N];
void dfs_sz(int v = 0, int p = -1) {
        sz[v] = 1;
        for(auto &u: g[v]){
                if(u == p) continue;
                h[u] = 1+h[v], par[u] = v;
                dfs_sz(u, v);
                sz[v] += sz[u];
                if(g[v][0] == p || sz[u] > sz[g[v][0]]) {
                        swap(u, g[v][0]);
int t, in[N], out[N], rin[N], nxt[N];
void dfs_hld(int v = 0, int p = -1) {
        in[v] = t++;
        rin[in[v]] = v;
        for(auto &u : g[v]){
                if(u == p) continue;
                nxt[u] = (u == g[v][0] ? nxt[v] : u);
                dfs_hld(u, v);
```

```
out[v] = t;
vector<int> tree;
void upd(int p, int value) {
        for(tree[p += sz[0]] = value; p > 1; p >>= 1) tree[p>>1] = min
            (tree[p], tree[p^1]);
int rmg(int 1, int r) {
        int res = inf:
        for (1 + sz[0], r + sz[0]; 1 < r; 1 >>= 1, r >>= 1) {
                if(1&1) res = min(res, tree[1++]);
                if(r\&1) res = min(res, tree[--r]);
        return res;
int up(int v) {
        return (nxt[v] != v) ? nxt[v] : (~par[v] ? par[v] : v);
int getLCA(int a, int b) {
        while(nxt[a] != nxt[b]){
                if(h[a] == 0 || h[up(a)] < h[up(b)]) swap(a, b);
                a = up(a);
        return h[a] < h[b] ? a : b;
int queryUp(int a, int p = 0){
        int ans = inf;
        while(nxt[a] != nxt[p]){
                ans = min(ans, rmq(in[nxt[a]], in[a]+1));
                a = par[nxt[a]];
        ans = min(ans, rmq(in[p], in[a]+1));
        return (ans == inf) ? -1 : rin[ans]+1;
int main(){
        scanf("%d %d", &n, &q);
        for (int i = 0; i < n-1; i++) {
                int a, b;
                scanf("%d %d", &a, &b);
                a--, b--;
                add(a, b);
        dfs_sz();
        t = 0;
        dfs sz();
        tree.assign(2*sz[0], inf);
        dfs_hld();
        for (int i = 0; i < q; i++) {
                int o, v;
                scanf("%d %d", &o, &v);
```

2.12 Centroid Decomposition

```
//Centroid decomposition1
void dfsSize(int v, int pa) {
  sz[v] = 1;
  for(int u : adj[v]) {
   if (u == pa || rem[u]) continue;
   dfsSize(u, v);
    sz[v] += sz[u];
int getCentroid(int v, int pa, int tam) {
  for(int u : adj[v]) {
   if (u == pa || rem[u]) continue;
   if (2 * sz[u] > tam) return getCentroid(u, v, tam);
  return v:
void decompose (int v, int pa = -1) {
 //cout << v << ' ' << pa << '\n';
 dfsSize(v, pa);
  int c = getCentroid(v, pa, sz[v]);
  //cout << c << '\n';
  par[c] = pa;
  rem[c] = 1;
  for(int u : adj[c]) {
   if (!rem[u] && u != pa) decompose(u, c);
  adj[c].clear();
//Centroid decomposition2
void dfsSize(int v, int par) {
  sz[v] = 1;
 for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   dfsSize(u, v);
   sz[v] += sz[u];
int getCentroid(int v, int par, int tam) {
 for(int u : adj[v]) {
   if (u == par || removed[u]) continue;
   if (2 * sz[u] > tam) return getCentroid(u, v, tam);
```

```
return v;
}

void setDis(int v, int par, int nv, int d) {
    dis[v][nv] = d;
    for(int u : adj[v]) {
        if (u == par || removed[u]) continue;
        setDis(u, v, nv, d + 1);
    }
}

void decompose(int v, int par, int nv) {
    dfsSize(v, par);
    int c = getCentroid(v, par, sz[v]);
    ct[c] = par;
    removed[c] = 1;
    setDis(c, par, nv, 0);
    for(int u : adj[c]) {
        if (!removed[u]) {
            decompose(u, c, nv + 1);
        }
    }
}
```

2.13 Sparse Table

```
template < class Info_t>
class SparseTable {
private:
  vector<int> log2;
  vector<vector<Info t>> table;
  Info_t merge(Info_t &a, Info_t &b) {
  }
public:
  SparseTable(int n, vector<Info_t> v) {
    log2.resize(n + 1);
    log2[1] = 0;
    for (int i = 2; i <= n; i++) {
      log2[i] = log2[i >> 1] + 1;
    table.resize(n + 1);
    for (int i = 0; i < n; i++) {
      table[i].resize(log2[n] + 1);
    for (int i = 0; i < n; i++) {
      table[i][0] = v[i];
    for (int i = 0; i < log2[n]; i++) {
      for (int j = 0; j < n; j++) {
        if (i + (1 << i)) >= n) break:
        table[j][i + 1] = merge(table[j][i], table[j + (1 << i)][i]);
  int get(int 1, int r) {
    int k = log2[r - 1 + 1];
```

```
return merge(table[l][k], table[r - (1 << k) + 1][k]);
};</pre>
```

2.14 Li Chao Tree

```
// by luucasv
typedef long long T;
const T INF = 1e18, EPS = 1;
const int BUFFER SIZE = 1e4;
struct Line {
 T m, b;
 Line (T m = 0, T b = INF) : m(m), b(b) {}
 T apply(T x) { return x * m + b; }
};
struct Node {
 Node *left, *right;
 Line line;
  Node(): left(NULL), right(NULL) {}
};
struct LiChaoTree {
 Node *root, buffer[BUFFER_SIZE];
  T min value, max value:
  int buffer_pointer;
  LiChaoTree(T min_value, T max_value): min_value(min_value),
      max_value(max_value + 1) { clear(); }
  void clear() { buffer_pointer = 0; root = newNode(); }
  void insert_line(T m, T b) { update(root, min_value, max_value, Line
      (m, b)); }
  T eval(T x) { return query(root, min_value, max_value, x); }
  void update(Node *cur, T l, T r, Line line) {
    T m = 1 + (r - 1) / 2;
    bool left = line.apply(1) < cur->line.apply(1);
    bool mid = line.apply(m) < cur->line.apply(m);
    bool right = line.apply(r) < cur->line.apply(r);
    if (mid) {
      swap(cur->line, line);
    if (r - 1 <= EPS) return;</pre>
    if (left == right) return;
    if (mid != left) {
      if (cur->left == NULL) cur->left = newNode();
      update(cur->left, 1, m, line);
    } else {
      if (cur->right == NULL) cur->right = newNode();
      update(cur->right, m, r, line);
  T query(Node *cur, T l, T r, T x) {
    if (cur == NULL) return INF;
    if (r - 1 <= EPS) {
      return cur->line.apply(x);
    T m = 1 + (r - 1) / 2;
    T ans:
```

```
if (x < m) {
    ans = query(cur->left, l, m, x);
} else {
    ans = query(cur->right, m, r, x);
}
    return min(ans, cur->line.apply(x));
}
Node* newNode() {
    buffer[buffer_pointer] = Node();
    return &buffer[buffer_pointer++];
}
};
```

3 Graph Algorithms

3.1 Dinic Max Flow

```
const int ms = 1e3; // Quantidade maxima de vertices
const int me = 1e5; // Quantidade maxima de arestas
int adj[ms], to[me], ant[me], wt[me], z, n;
int copy_adj[ms], fila[ms], level[ms];
void clear() {
 memset(adj, -1, sizeof adj);
  z = 0;
int add(int u, int v, int k) {
 to[z] = v;
  ant[z] = adj[u];
 wt[z] = k;
  adj[u] = z++;
  swap(u, v);
  to[z] = v;
  ant[z] = adj[u];
  wt[z] = 0; // Lembrar de colocar = 0
  adi[u] = z++;
int bfs(int source, int sink) {
 memset(level, -1, sizeof level);
  level[source] = 0;
  int front = 0, size = 0, v;
  fila[size++] = source;
  while(front < size)</pre>
  v = fila[front++];
  for(int i = adj[v]; i != -1; i = ant[i]) {
    if(wt[i] && level[to[i]] == -1) {
    level[to[i]] = level[v] + 1;
    fila[size++] = to[i];
  return level[sink] != -1;
int dfs(int v, int sink, int flow) {
```

```
if(v == sink) return flow;
int f;
for(int &i = copy_adj[v]; i != -1; i = ant[i]) {
   if(wt[i] && level[to[i]] == level[v] + 1 &&
        (f = dfs(to[i], sink, min(flow, wt[i])))) {
        wt[i] -= f;
        wt[i ^ 1] += f;
        return f;
   }
} return 0;
}
int maxflow(int source, int sink) {
   int ret = 0, flow;
   while(bfs(source, sink)) {
        memcpy(copy_adj, adj, sizeof adj);
        while((flow = dfs(source, sink, 1 << 30))) {
        ret += flow;
   }
} return ret;
}</pre>
```

3.2 Euler Path and Circuit

```
int pathV[me], szV, del[me], pathE, szE;
int adj[ms], to[me], ant[me], wt[me], z, n;

// Funcao de add e clear no dinic

void eulerPath(int u) {
  for(int i = adj[u]; ~i; i = ant[u]) if(!del[i]) {
    del[i] = del[i^1] = 1;
    eulerPath(to[i]);
    pathE[szE++] = i;
  }
  pathV[szV++] = u;
}
```

3.3 Articulation Points/Bridges/Biconnected Components

```
const int ms = 1e3; // Quantidade maxima de vertices
const int me = 1e5; // Quantidade maxima de arestas

int adj[ms], to[me], ant[me], z, n;
int idx[ms], bc[me], ind, nbc, child, st[me], top;

// Funcao de add e clear no dinic

void generateBc(int edge) {
   while(st[--top] != edge) {
     bc[st[top]] = nbc;
   }
   bc[edge] = nbc++;
}

int dfs(int v, int par = -1) {
```

```
int low = idx[v] = ind++;
  for(int i = adj[v]; i > -1; i = ant[i]) {
    if(idx[to[i]] == -1) {
      if (par == -1) child++;
      st[top++] = i;
      int temp = dfs(to[i], v);
      if(par == -1 \&\& child > 1 || par \&\& temp >= idx[v]) generateBc(
      if(temp >= idx[v]) art[v] = true;
      if(temp > idx[v]) bridge[i] = true;
      low = min(low, temp);
    } else if(to[i] != par && idx[to[i]] < low) {</pre>
      low = idx[to[i]];
      st[top++] = i;
  return low:
void biconnected() {
  ind = 0;
 nbc = 0;
 top = -1;
 memset(idx, -1, sizeof idx);
 memset(art, 0, sizeof art);
  memset(bridge, 0, sizeof bridge);
  for(int i = 0; i < n; i++) if(idx[i] == -1) {</pre>
    child = 0;
    dfs(i);
```

3.4 SCC - Strongly Connected Components / 2SAT

```
vector<int> g[ms];
int idx[ms], low[ms], z, comp[ms], ncomp, st[ms], top;
int dfs(int u) {
  if(~idx[u]) return idx[u] ? idx[u] : z;
  low[u] = idx[u] = z++;
  st.push(u);
  for(int v : g[u]) {
   low[u] = min(low[u], dfs(v));
  if(low[u] == idx[u]) {
    idx[st.top()] = 0;
    st.pop();
    while(st.top() != u) {
      int v = st.top();
      st.pop();
      idx[v] = 0;
      low[v] = low[u];
      comp[v] = ncomp;
    comp[u] = ncomp++;
  return low[u];
bool solveSat() {
```

```
memset(idx, -1, sizeof idx);
ind = 1; top = -1;
for(int i = 0; i < n; i++) dfs(i);
for(int i = 0; i < n; i++) if(comp[i] == comp[i^1]) return false;
return true;
}

// Operacoes comuns de 2-sat
// ~v = "nao v"
#define trad(v) (v<0?((~v)*2)^1:v*2)
void addImp(int a, int b) { g[trad(a)].push(trad(b)); }
void addOr(int a, int b) { addImp(~a, b); addImp(~b, a); }
void addDiff(int a, int b) { addOr(a, ~b); addOr(~a, b); }
void addDiff(int a, int b) { addEqual(a, ~b); }
// valoracao: value[v] = comp[trad(v)] < comp[trad(~v)]</pre>
```

3.5 LCA - Lowest Common Ancestor

```
int par[ms][mlg+1], lvl[ms];
vector<int> q[ms];
void dfs(int v, int p, int l = 0) {
 lv1[v] = 1;
  par[v][0] = p;
  for(int u : q[v]) {
    if (u != p) dfs(u, v, l + 1);
void processAncestors(int root = 0) {
  dfs(root, root);
  for(int k = 1; k <= mlg; k++) {</pre>
    for (int i = 0; i < n; i++) {
      par[i][k] = par[par[i][k-1]][k-1];
int lca(int a, int b) {
 if(lvl[b] > lvl[a]) swap(a, b);
  for(int i = mlg; i >= 0; i--) {
    if(lvl[a] - (1 << i) >= lvl[b]) a = par[a][i];
  if(a == b) return a;
  for (int i = mlg; i >= 0; i--) {
    if(par[a][i] != par[b][i]) a = par[a][i], b = par[b][i];
  return par[a][0];
```

3.6 Sack

```
void solve(int a, int p, bool f) {
  int big = -1;
  for(auto &b : adj[a]) {
    if(b != p && (big == -1 || en[b]-st[b] > en[big]-st[big])) {
      big = b;
    }
```

```
for(auto &b : adj[a]){
 if(b == p || b == big) continue;
 solve(b, a, 0);
if(~big) solve(big, a, 1);
add(cnt[v[a]], -1);
cnt[v[a]]++;
add(cnt[v[a]], +1);
for(auto &b : adj[a]){
 if(b == p || b == big) continue;
 for(int i = st[b]; i < en[b]; i++) {</pre>
    add(cnt[ett[i]], -1);
    cnt[ett[i]]++;
    add(cnt[ett[i]], +1);
for(auto &q : Q[a]){
 ans[q.first] = query(mx-1)-query(q.second-1);
if(!f){
  for (int i = st[a]; i < en[a]; i++) {
    add(cnt[ett[i]], -1);
    cnt[ett[i]]--;
    add(cnt[ett[i]], +1);
```

3.7 Min Cost Max Flow

```
template <class flow_t, class cost_t>
class MinCostMaxFlow {
private:
  typedef pair<cost_t, int> ii;
  struct Edge {
    int to;
    flow_t cap;
    cost t cost;
    Edge(int to, flow_t cap, cost_t cost) : to(to), cap(cap), cost(
        cost) {}
  };
  int n;
  vector<vector<int>> adi;
  vector<Edge> edges:
  vector<cost_t> dis;
  vector<int> prev, id_prev;
        vector<int> q:
        vector<bool> ing;
  pair<flow t, cost t> spfa(int src, int sink) {
    fill(dis.begin(), dis.end(), int(1e9)); //cost_t inf
    fill(prev.begin(), prev.end(), -1);
    fill(ing.begin(), ing.end(), false);
    q.clear();
    q.push_back(src);
    inq[src] = true;
    dis[src] = 0;
```

```
for(int on = 0; on < (int) q.size(); on++) {
        int cur = q[on];
        ing[cur] = false;
        for(auto id : adj[cur]) {
                if (edges[id].cap == 0) continue;
                int to = edges[id].to;
                if (dis[to] > dis[cur] + edges[id].cost) {
                        prev[to] = cur;
                        id_prev[to] = id;
                        dis[to] = dis[cur] + edges[id].cost;
                        if (!ing[to]) {
                                q.push_back(to);
                                ing[to] = true;
    flow_t mn = flow_t(1e9);
    for(int cur = sink; prev[cur] != -1; cur = prev[cur]) {
      int id = id_prev[cur];
      mn = min(mn, edges[id].cap);
    if (mn == flow_t(1e9) || mn == 0) return make_pair(0, 0);
    pair<flow t, cost t> ans(mn, 0);
    for(int cur = sink; prev[cur] != -1; cur = prev[cur]) {
      int id = id_prev[cur];
      edges[id].cap -= mn;
      edges[id ^ 1].cap += mn;
      ans.second += mn * edges[id].cost;
    return ans;
public:
 MinCostMaxFlow(int a = 0) {
    n = a:
    adj.resize(n + 2);
    edges.clear():
    dis.resize(n + 2);
    prev.resize(n + 2);
    id prev.resize(n + 2):
    inq.resize(n + 2);
  void init(int a) {
    n = a:
    adj.resize(n + 2);
    edges.clear();
    dis.resize(n + 2);
    prev.resize(n + 2);
    id prev.resize(n + 2);
    inq.resize(n + 2);
  void add(int from, int to, flow_t cap, cost_t cost) {
    adj[from].push_back(int(edges.size()));
                edges.push back(Edge(to, cap, cost));
                adj[to].push_back(int(edges.size()));
                edges.push_back(Edge(from, 0, -cost));
  pair<flow_t, cost_t> maxflow(int src, int sink) {
    pair<flow_t, cost_t> ans(0, 0), got;
    while((got = spfa(src, sink)).first > 0) {
      ans.first += got.first;
```

```
ans.second += got.second;
}
return ans;
};
```

const int inf = 0x3f3f3f3f;

3.8 Hungarian Algorithm - Maximum Cost Matching

```
int n, w[ms][ms], maxm;
int lx[ms], ly[ms], xy[ms], yx[ms];
int slack[ms], slackx[ms], prev[ms];
bool S[ms], T[ms];
void init labels() {
 memset(lx, 0, sizeof lx); memset(ly, 0, sizeof ly);
  for (int x = 0; x < n; x++) for (int y = 0; y < n; y++) {
    lx[x] = max(lx[x], cos[x][y]);
  }
void updateLabels() {
  int delta = inf;
  for (int y = 0; y < n; y++) if (!T[y]) delta = min(delta, slack[y]);
  for (int x = 0; x < n; x++) if (S[x]) lx[x] -= delta;
  for (int y = 0; y < n; y++) if (T[y]) |y[y]| += delta;
  for (int y = 0; y < n; y++) if (!T[y]) slack[y] -= delta;
void addTree(int x, int prevx) {
 S[x] = 1; prev[x] = prevx;
  for (int y = 0; y < n; y++) if (lx[x] + ly[y] - w[x][y] < slack[y]) {
    slack[y] = lx[x] + ly[y] - cost[x][y];
    slackx[y] = x;
void augment() {
  if(maxm == n) return;
  int x, y, root;
  int q[ms], wr = 0, rd = 0;
  memset(S, 0, sizeof S); memset(T, 0, sizeof T);
  memset (prev, -1, sizeof prev);
  for (int x = 0; x < n; x++) if (xy[x] == -1) {
    q[wr++] = root = x;
    prev[x] = -2;
    S[x] = 1;
    break;
  for (int y = 0; y < n; y++) {
    slack[y] = lx[root] + ly[y] - w[root][y];
    slackx[y] = root;
  while(true) {
    while(rd < wr) {</pre>
      for(y = 0; y < n; y++) if(w[x][y] == lx[x] + ly[y] && !T[y]) {
        if(yx[y] == -1) break;
        T[y] = 1;
```

```
q[wr++] = yx[y];
        addTree(yx[y], x);
      if(y < n) break;</pre>
    if(v < n) break;</pre>
    updateLabels();
    wr = rd = 0;
    for(y = 0; y < n; y++) if(!T[y] && !slack[y]) {
      if(yx[y] == -1) {
        x = slackx[y];
        break;
      } else {
        T[y] = true;
        if(!S[yx[y]])
          q[wr++] = yx[y];
          addTree(yx[y], slackx[y]);
    if(y < n) break;</pre>
  if(y < n) {
    for(int cx = x, cy = y, ty; cx != -2; cx = prev[cx], cy = ty) {
      ty = xy[cx];
      yx[cy] = cx;
      xy[cx] = cy;
    augment();
int hungarian() {
  int ans = 0; maxm = 0;
  memset (xy, -1, sizeof xy); memset (yx, -1, sizeof yx);
  initLabels(); augment();
  for (int x = 0; x < n; x++) ans += w[x][xy[x]];
  return ans;
```

4 Math

4.1 Discrete Logarithm

```
1l discreteLog(ll a, ll b, ll m) {
    // a^ans == b mod m
    // ou -1 se nao existir
    ll cur = a, on = 1;
    for(int i = 0; i < 100; i++) {
        cur = cur * a % m;
    }
    while(on * on <= m) {
        cur = cur * a % m;
        on++;
    }
    map<1l, ll> position;
    for(ll i = 0, x = 1; i * i <= m; i++) {</pre>
```

```
position[x] = i * on;
    x = x * cur % m;
}
for(ll i = 0; i <= on + 20; i++) {
    if(position.count(b)) {
        return position[b] - i;
    }
    b = b * a % m;
}
return -1;
}</pre>
```

4.2 GCD - Greatest Common Divisor

```
11 gcd(11 a, 11 b) {
  while(b) a %= b, swap(a, b);
  return a;
}
```

4.3 Extended Euclides

```
// euclides estendido: acha u e v da equacao:
// u * x + v * y = gcd(x, y);
// u eh inverso modular de x no modulo y
// v eh inverso modular de y no modulo x

pair<ll, ll> euclides(ll a, ll b) {
    ll u = 0, oldu = 1, v = 1, oldv = 0;
    while(b) {
        ll q = a / b;
        oldv = oldv - v * q;
        oldu = oldu - u * q;
        a = a - b * q;
        swap(a, b);
        swap(u, oldu);
        swap(v, oldv);
    }
    return make_pair(oldu, oldv);
}
```

4.4 Fast Exponentiation

```
const l1 mod = 1e9+7;

l1 fExp(l1 a, l1 b) {
    l1 ans = 1;
    while(b) {
        if(b & 1) ans = ans * a % mod;
        a = a * a % mod;
        b >>= 1;
    }
    return ans;
}
```

4.5 Matrix Fast Exponentiation

```
const 11 \mod = 1e9+7;
const int m = 2; // size of matrix
struct Matrix {
  11 mat[m][m];
 Matrix operator * (const Matrix &p) {
    Matrix ans;
    for (int i = 0; i < m; i++)
      for (int j = 0; j < m; j++)
        for (int k = ans.mat[i][j] = 0; k < m; k++)
          ans.mat[i][j] = (ans.mat[i][j] + mat[i][k] * p.mat[k][j]) %
    return ans;
};
Matrix fExp(Matrix a, ll b) {
  Matrix ans:
  for (int i = 0; i < m; i++) for (int j = 0; j < m; j++)
    ans.mat[i][j] = i == j;
  while(b) {
   if(b \& 1) ans = ans * a;
    a = a * a;
    b >>= 1;
  return ans;
```

4.6 FFT - Fast Fourier Transform

```
typedef complex<double> Complex;
typedef long double ld;
typedef long long 11;
const int ms = maiortamanhoderesposta * 2;
const ld pi = acosl(-1.0);
int rbit[1 << 23];</pre>
Complex a[ms], b[ms];
void calcReversedBits(int n) {
  int lq = 0;
  while (1 << (lq + 1) < n) {
    lq++;
  for(int i = 1; i < n; i++) {</pre>
    rbit[i] = (rbit[i >> 1] >> 1) | ((i & 1) << lq);
void fft(Complex a[], int n, bool inv = false) {
  for(int i = 0; i < n; i++) {
    if(rbit[i] > i) swap(a[i], a[rbit[i]]);
  double ang = inv ? -pi : pi;
```

```
for (int m = 1; m < n; m += m) {
    Complex d(cosl(ang/m), sinl(ang/m));
    for (int i = 0; i < n; i += m+m) {
      Complex cur = 1;
      for (int j = 0; j < m; j++) {
        Complex u = a[i + j], v = a[i + j + m] * cur;
        a[i + j] = u + v;
        a[i + j + m] = u - v;
        cur *= d;
  if(inv) {
    for(int i = 0; i < n; i++) a[i] /= n;</pre>
void multiply(11 x[], 11 y[], 11 ans[], int nx, int ny, int &n) {
 n = 1:
  while (n < nx+ny) n <<= 1;
  calcReversedBits(n);
  for (int i = 0; i < n; i++) {
   a[i] = Complex(x[i]);
   b[i] = Complex(v[i]);
  fft(a, n); fft(b, n);
  for(int i = 0; i < n; i++) {</pre>
   a[i] = a[i] * b[i];
  fft(a, n, true);
  for (int i = 0; i < n; i++) {
    ans[i] = 11(a[i].real() + 0.5);
 n = nx + ny;
```

4.7 NTT - Number Theoretic Transform

```
long long int mod = (11911 << 23) + 1, c_root = 3;

namespace NTT {
    typedef long long int 11;

    11 fexp(11 base, 11 e) {
        11 ans = 1;
        while(e > 0) {
            if (e & 1) ans = ans * base % mod;
            base = base * base % mod;
            e >>= 1;
        }
        return ans;
}

ll inv_mod(11 base) {
        return fexp(base, mod - 2);
}

void ntt(vector<11>& a, bool inv) {
        int n = (int) a.size();
        if (n == 1) return;
}
```

```
for(int i = 0, j = 0; i < n; i++) {
      if (i > j) {
        swap(a[i], a[j]);
      for (int 1 = n / 2; († \hat{} = 1) < 1; 1 >>= 1);
    for(int sz = 1; sz < n; sz <<= 1) {</pre>
      11 delta = fexp(c_root, (mod - 1) / (2 * sz)); //delta = w_2sz
      if (inv) {
        delta = inv_mod(delta);
      for (int i = 0; i < n; i += 2 * sz) {
        11 w = 1;
        for (int j = 0; j < sz; j++) {
          ll u = a[i + j], v = w * a[i + j + sz] % mod;
          a[i + j] = (u + v + mod) % mod;
          a[i + i] = (a[i + i] + mod) % mod;
          a[i + j + sz] = (u - v + mod) % mod;
          a[i + j + sz] = (a[i + j + sz] + mod) % mod;
          w = w * delta % mod;
    if (inv) {
      ll inv n = inv mod(n);
      for(int i = 0; i < n; i++) {</pre>
        a[i] = a[i] * inv n % mod;
    for (int i = 0; i < n; i++) {
      a[i] %= mod;
      a[i] = (a[i] + mod) % mod;
  void multiply(vector<11> &a, vector<11> &b, vector<11> &ans) {
    int lim = (int) max(a.size(), b.size());
    int n = 1:
    while (n < lim) n <<= 1;
    n <<= 1;
    a.resize(n);
    b.resize(n);
    ans.resize(n);
    ntt(a, false);
    ntt(b, false);
    for (int i = 0; i < n; i++) {
      ans[i] = a[i] * b[i] % mod;
    ntt(ans, true);
};
```

4.8 Miller and Rho

```
typedef long long int 11;
bool overflow(ll a, 11 b) {
  return b && (a >= (111 << 62) / b);</pre>
```

```
11 add(ll a, ll b, ll md) {
  return (a + b) % md;
11 mul(ll a, ll b, ll md) {
  if (!overflow(a, b)) return (a * b) % md;
 11 \text{ ans} = 0;
  while(b) {
   if (b & 1) ans = add(ans, a, md);
   a = add(a, a, md);
   b >>= 1;
  return ans;
ll fexp(ll a, ll e, ll md) {
 11 \text{ ans} = 1;
 while(e) {
   if (e & 1) ans = mul(ans, a, md);
   a = mul(a, a, md);
   e >>= 1;
  return ans;
11 my_rand() {
 ll ans = rand();
 ans = (ans \ll 31) \mid rand();
  return ans;
11 gcd(ll a, ll b) {
  while(b) {
   11 t = a % b;
   a = b:
   b = t;
  return a;
bool miller(ll p, int iteracao) {
 if(p < 2) return 0;
  if(p % 2 == 0) return (p == 2);
 11 s = p - 1;
  while (s % 2 == 0) s >>= 1;
  for(int i = 0; i < iteracao; i++) {</pre>
    11 a = rand() % (p - 1) + 1, temp = s;
   11 \mod = fexp(a, temp, p);
    while (temp != p - 1 && mod != 1 && mod != p - 1) {
      mod = mul(mod, mod, p);
      temp <<= 1;
    if(mod != p - 1 && temp % 2 == 0) return 0;
  return 1;
ll rho(ll n) {
  if (n == 1 || miller(n, 10)) return n;
```

```
if (n % 2 == 0) return 2;
while(1) {
    ll x = my_rand() % (n - 2) + 2, y = x;
    ll c = 0, cur = 1;
    while(c == 0) {
        c = my_rand() % (n - 2) + 1;
    }
    while(cur == 1) {
        x = add(mul(x, x, n), c, n);
        y = add(mul(y, y, n), c, n);
        y = add(mul(y, y, n), c, n);
        cur = gcd((x >= y ? x - y : y - x), n);
    }
    if (cur != n) return cur;
}
```

4.9 Determinant using Mod

```
// by zchao1995
// Determinante com coordenadas inteiras usando Mod
11 mat[ms][ms];
11 det (int n) {
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < n; j++) {
      mat[i][i] %= mod;
  11 \text{ res} = 1:
 for (int i = 0; i < n; i++) {</pre>
    if (!mat[i][i])
     bool flag = false;
      for (int j = i + 1; j < n; j++) {
        if (mat[j][i]) {
          flag = true;
          for (int k = i; k < n; k++) {
            swap (mat[i][k], mat[j][k]);
          res = -res;
          break;
      if (!flag) {
        return 0;
    for (int j = i + 1; j < n; j++) {
      while (mat[j][i]) {
        11 t = mat[i][i] / mat[j][i];
        for (int k = i; k < n; k++) {
          mat[i][k] = (mat[i][k] - t * mat[j][k]) % mod;
          swap (mat[i][k], mat[j][k]);
        res = -res;
    res = (res * mat[i][i]) % mod;
```

```
return (res + mod) % mod;
}
```

5 Geometry

5.1 Geometry

```
const double inf = 1e100, eps = 1e-9;
struct PT {
        double x, y;
        PT (double x = 0, double y = 0) : x(x), y(y) {}
        PT operator + (const PT &p) { return PT(x + p.x, y + p.y); }
        PT operator - (const PT &p) { return PT(x - p.x, y - p.y); }
        PT operator * (double c) { return PT(x * c, y * c); }
        PT operator / (double c) { return PT(x / c, y / c); }
        bool operator <(const PT &p) const {</pre>
                if(fabs(x - p.x) >= eps) return x < p.x;</pre>
                return fabs(y - p.y) >= eps && y < p.y;</pre>
        bool operator == (const PT &p) const {
                return fabs (x - p.x) < eps && fabs (y - p.y) < eps;
};
double dot(PT p, PT q) { return p.x * q.x + p.y * q.y; }
double dist2(PT p, PT q) { return dot(p - q, p - q); }
double dist(PT p, PT q) {return hypot(p.x-q.x, p.y-q.y); }
double cross(PT p, PT q) { return p.x * q.y - p.y * q.x; }
// Rotaciona o ponto CCW ou CW ao redor da origem
PT rotateCCW90(PT p) { return PT(-p.y, p.x); }
PT rotateCW90(PT p) { return PT(p.y, -p.x); }
PT rotateCCW(PT p, double d) {
    return PT(p.x * cos(t) - p.y * sin(t), p.x * sin(t) + p.y * cos(t)
// Projeta ponto c na linha a - b assumindo a != b
PT projectPointLine(PT a, PT b, PT c) {
    return a + (b - a) * dot(c - a, b - a) / dot(b - a, b - a);
// Projeta ponto c no segmento a - b
PT projectPointSegment(PT a, PT b, PT c) {
    double r = dot(b - a, b - a);
    if(abs(r) < eps) return a;</pre>
    r = dot(c - a, b - a) / r;
    if(r < 0) return a;</pre>
    if(r > 1) return b;
    return a + (b - a) * r;
// Calcula distancia entre o ponto c e o segmento a - b
double distancePointSegment(PT a, PT b, PT c) {
    return sqrt(dist2(c, projectPointSegment(a, b, c)));
// Determina se o ponto c esta em um segmento a - b
```

```
bool ptInSegment(PT a, PT b, PT c) {
  bool x = min(a.x, b.x) \le c.x \le c.x \le max(a.x, b.x);
  bool y = min(a.y, b.y) \le c.y \le c.y \le max(a.y, b.y);
  return x && y && (cross((b-a),(c-a)) == 0); // testar com eps se for
       double
// Calcula distancia entre o ponto (x, y, z) e o plano ax + by + cz =
double distancePointPlane(double x, double y, double z, double a,
    double b, double c, double d) {
    return abs(a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c
        );
// Determina se as linhas a - b e c - d sao paralelas ou colineares
bool linesParallel(PT a, PT b, PT c, PT d) {
   return abs(cross(b - a, c - d)) < eps;</pre>
bool linesCollinear(PT a, PT b, PT c, PT d) {
    return linesParallel(a, b, c, d) && abs(cross(a - b, a - c)) < eps</pre>
         && abs(cross(c - d, c - a)) < eps;
// Determina se o segmento a - b intersecta com o segmento c - d
bool segmentsIntersect(PT a, PT b, PT c, PT d) {
   if(linesCollinear(a, b, c, d)) {
       if(dist2(a, c) < eps || dist2(a, d) < eps || dist2(b, c) < eps</pre>
             || dist2(b, d) < eps) return true;</pre>
        if(dot(c - a, c - b) > 0 \&\& dot(d - a, d - b) > 0 \&\& dot(c - b)
            d - b > 0 return false;
       return true:
    if(cross(d - a, b - a) * cross(c - a, b - a) > 0) return false:
    if(cross(a - c, d - c) * cross(b - c, d - c) > 0) return false;
    return true;
// Calcula a intersecao entre as retas a - b e c - d assumindo que uma
     unica intersecao existe
// Para intersecao de segmentos, cheque primeiro se os segmentos se
    intersectam e que nao paralelos
PT computeLineIntersection(PT a, PT b, PT c, PT d) {
   b = b - a; d = c - d; c = c - a;
   assert(cross(b, d) != 0); // garante que as retas nao sao
        paralelas, remover pra evitar tle
   return a + b * cross(c, d) / cross(b, d);
// Calcula centro do circulo dado tres pontos
PT computeCircleCenter(PT a, PT b, PT c) {
   b = (a + b) / 2;
   c = (a + c) / 2;
   return computeLineIntersection(b, b + rotateCW90(a - b), c, c +
        rotateCW90(a - c));
// Determina se o ponto p esta dentro do triangulo (a, b, c)
bool ptInsideTriangle(PT p, PT a, PT b, PT c) {
 if(cross(b-a, c-b) < 0) swap(a, b);
```

```
11 x = cross(b-a, p-b);
  11 y = cross(c-b, p-c);
  11 z = cross(a-c, p-a);
  if (x > 0 \& \& y > 0 \& \& z > 0) return true;
  if(!x) return ptInSegment(a,b,p);
  if(!v) return ptInSegment(b,c,p);
  if(!z) return ptInSegment(c,a,p);
  return false;
// Determina se o ponto esta num poligono convexo em O(lqn)
bool pointInConvexPolygon(const vector<PT> &p, PT q) {
 PT pivot = p[0];
  int x = 1, y = p.size();
 while (y-x != 1)
    int z = (x+y)/2;
    PT diagonal = pivot - p[z];
    if(cross(p[x] - pivot, q - pivot) * cross(q-pivot, p[z] - pivot)
        >= 0) v = z;
    else x = z;
  return ptInsideTriangle(q, p[x], p[y], pivot);
// Determina se o ponto esta num poligono possivelmente nao-convexo
// Retorna 1 para pontos estritamente dentro, 0 para pontos
    estritamente fora do poligno
// e 0 ou 1 para os pontos restantes
// Eh possivel converter num teste exato usando inteiros e tomando
    cuidado com a divisao
// e entao usar testes exatos para checar se esta na borda do poligno
bool pointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i + 1) % p.size();
    if((p[i].y \le q.y \&\& q.y < p[j].y || p[j].y \le q.y \&\& q.y < p[i].y
      q.x < p[i].x + (p[i].x - p[i].x) * (q.y - p[i].y) / (p[i].y - p[i].y)
          i].y))
      c = !c;
  return c;
// Determina se o ponto esta na borda do poligno
bool pointOnPolygon(const vector<PT> &p, PT q) {
  for(int i = 0; i < p.size(); i++)</pre>
    if(dist2(projectPointSegment(p[i], p[(i + 1) % p.size()], q), q) <</pre>
      return true;
    return false;
// Calcula intersecao da linha a - b com o circulo centrado em c com
    raio r > 0
vector<PT> circleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ans;
 b = b - a:
  a = a - c;
  double x = dot(b, b);
  double y = dot(a, b);
```

```
double z = dot(a, a) - r * r;
  double w = y * y - x * z;
  if (w < -eps) return ans;</pre>
  ans.push_back(c + a + b \star (-y + sqrt(w + eps)) / x);
  if (w > eps)
    ans.push_back(c + a + b \star (-y - sqrt(w)) / x);
  return ans;
// Calcula intersecao do circulo centrado em a com raio r e o centrado
     em b com raio R
vector<PT> circleCircleIntersection(PT a, PT b, double r, double R) {
 vector<PT> ans;
  double d = sqrt(dist2(a, b));
  if (d > r + R \mid \mid d + min(r, R) < max(r, R)) return ans;
  double x = (d * d - R * R + r * r)/(2 * d);
  double y = sqrt(r * r - x * x);
  PT v = (b - a) / d;
  ans.push_back(a + v * x + rotateCCW90(v) * y);
  if (y > 0)
    ans.push_back(a + v * x - RotateCCW90(v) * y);
  return ans;
// Calcula a area ou o centroide de um poligono (possivelmente nao-
    convexo)
// assumindo que as coordenadas estao listada em ordem horaria ou anti
    -horaria
// O centroide eh equivalente a o centro de massa ou centro de
    gravidade
double computeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i + 1) % p.size();
    area += p[i].x * p[j].y - p[j].x * p[i].y;
  return area / 2.0:
double computeArea(const vector<PT> &p) {
  return abs(computeSignedArea(p));
PT computeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for(int i = 0; i < p.size(); i++){</pre>
   int j = (i + 1) % p.size();
    c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
  return c / scale;
// Testa se o poligno listada em ordem CW ou CCW eh simples (nenhuma
    linha se intersecta)
bool isSimple(const vector<PT> &p) {
  for(int i = 0; i < p.size(); i++) {</pre>
    for(int k = i + 1; k < p.size(); k++) {</pre>
      int j = (i + 1) % p.size();
      int 1 = (k + 1) % p.size();
      if (i == 1 \mid | j == k) continue;
```

```
if (segmentsIntersect(p[i], p[j], p[k], p[l]))
    return false;
}
return true;
}
```

5.2 Convex Hull

```
vector<PT> convexHull(vector<PT> p)) {
  int n = p.size(), k = 0;
  vector<PT> h(2 * n);
  sort(p.begin(), p.end());
  for(int i = 0; i < n; i++) {
    while(k >= 2 && cross(h[k - 1] - h[k - 2], p[i] - h[k - 2]) <= 0)
        k--;
    h[k++] = p[i];
  }
  for(int i = n - 2, t = k + 1; i >= 0; i--) {
    while(k >= t && cross(h[k - 1] - h[k - 2], p[i] - h[k - 2]) <= 0)
        k--;
    h[k++] = p[i];
  }
  h.resize(k);
  return h;
}</pre>
```

5.3 Closest Pair

```
double closestPair(vector<PT> p) {
  int n = p.size(), k = 0;
  sort(p.begin(), p.end());
  double d = inf;
  set<PT> ptsInv;
  for(int i = 0; i < n; i++) {
    while(k < i && p[k].x < p[i].x - d) {
      ptsInv.erase(swapCoord(p[k++]));
    }
  for(auto it = ptsInv.lower_bound(PT(p[i].y - d, p[i].x - d));
    it != ptsInv.end() && it->x <= p[i].y + d; it++) {
      d = min(d, !(p[i] - swapCoord(*it)));
    }
    ptsInv.insert(swapCoord(p[i]));
}
return d;
}</pre>
```

5.4 Intersection Points

```
// Utiliza uma seg tree
int intersectionPoints(vector<pair<PT, PT>> v) {
  int n = v.size();
  vector<pair<int, int>> events, vertInt;
  for(int i = 0; i < n; i++) {
    if(v.first.x == v.second.x) { // Segmento Vertical}</pre>
```

```
int y0 = min(v.first.y, v.second.y), y1 = max(v.first.y, v.
        second.v);
    events.push_back({v.first.x, vertInt.size()}); // Tipo = Indice
        no array
    vertInt.push_back({y0, y1});
  } else { // Segmento Horizontal
    int x0 = min(v.first.x, v.second.x), x1 = max(v.first.x, v.
        second.x);
    events.push_back({x0, -1}); // Inicio de Segmento
    events.push_back({x1, inf}); // Final de Segmento
sort(events.begin(), events.end());
int ans = 0;
for(int i = 0; i < events.size(); i++) {</pre>
  int t = events[i].second;
  if(t == -1) {
    segUpdate(events[i].first, 1);
  } else if(t == inf) {
    segUpdate(events[i].first, 0);
    ans += segQuery(vertInt[t].first, vertInt[t].second);
return ans;
```

5.5 Java Geometry Library

```
import java.util.*;
import java.io.*;
import java.awt.geom.*;
import java.lang.*;
//Lazy Geometry
class AWT{
  static Area makeArea(double[] pts){
    Path2D.Double p = new Path2D.Double();
    p.moveTo(pts[0], pts[1]);
    for(int i = 2; i < pts.length; i+=2) {</pre>
      p.lineTo(pts[i], pts[i+1]);
    p.closePath();
    return new Area(p);
  static double computePolygonArea(ArrayList<Point2D.Double> points) {
    Point2D.Double[] pts = points.toArray(new Point2D.Double[points.
        size()]);
    double area = 0;
    for (int i = 0; i < pts.length; i++) {</pre>
      int j = (i+1) % pts.length;
      area += pts[i].x * pts[j].y - pts[j].x * pts[i].y;
    return Math.abs(area)/2:
  static double computeArea(Area area) {
    double totArea = 0;
    PathIterator iter = area.getPathIterator(null);
    ArrayList<Point2D.Double> points = new ArrayList<Point2D.Double>()
    while (!iter.isDone()) {
```

```
double[] buffer = new double[6];
switch (iter.currentSegment (buffer)) {
    case PathIterator.SEG_MOVETO:
    case PathIterator.SEG_LINETO:
        points.add(new Point2D.Double(buffer[0], buffer[1]));
        break;
    case PathIterator.SEG_CLOSE:
        totArea += computePolygonArea(points);
        points.clear();
        break;
    }
    iter.next();
}
return totArea;
}
```

6 Dynamic Programming

6.1 Convex Hull Trick

```
typedef long double double_t;
typedef long long int 11;
class HullDynamic {
  const double t inf = 1e9;
  struct Line {
    11 m, b;
    double_t start;
    bool is_query;
    Line() {}
    Line(ll _m, ll _b, double_t _start, bool _is_query) : m(_m), b(_b)
        , start(_start), is_query(_is_query) {}
    11 eval(11 x) {
      return m * x + b;
    double t intersect (const Line& 1) const {
      return (double_t) (1.b - b) / (m - 1.m);
    bool operator< (const Line& 1) const {</pre>
      if (is_query == 0) return m > 1.m;
      return (start < 1.start);</pre>
  };
  typedef set<Line>::iterator iterator_t;
  bool has_prev(iterator_t it) {
    return (it != hull.begin());
```

```
return (++it != hull.end());
bool irrelevant(iterator_t it) {
 if (!has prev(it) || !has next(it)) return 0;
  iterator_t prev = it, next = it;
 prev--;
 next++;
  return next->intersect(*prev) <= it->intersect(*prev);
void update_left(iterator_t it) {
 if (it == hull.begin()) return;
 iterator_t pos = it;
  --it:
  vector<Line> rem:
  while(has_prev(it)) {
   iterator_t prev = it;
    if (prev->intersect(*pos) <= prev->intersect(*it)) {
      rem.push back(*it);
    } else {
     break;
    --it:
  double_t start = pos->intersect(*it);
  Line f = *pos;
  for (Line r : rem) hull.erase(r);
 hull.erase(f);
  f.start = start;
 hull.insert(f);
void update_right(iterator_t it) {
  if (!has next(it)) return;
  iterator_t pos = it;
  ++it;
  vector<Line> rem:
  while(has_next(it)) {
    iterator_t next = it;
    ++next;
    if (next->intersect(*pos) <= pos->intersect(*it)) {
     rem.push_back(*it);
    } else {
     break;
  double_t start = pos->intersect(*it);
  Line f = *it;
  for (Line r : rem) hull.erase(r);
 hull.erase(f);
 f.start = start;
 hull.insert(f);
void insert_line(ll m, ll b) {
 Line f(m, b, -inf, 0);
 iterator_t it = hull.lower_bound(f);
```

bool has_next(iterator_t it) {

```
if (it != hull.end() && it->m == f.m) {
      if (it->b <= f.b) {
        return;
      } else if (it->b > f.b) {
        hull.erase(it);
    hull.insert(f);
    it = hull.lower bound(f);
    if (irrelevant(it)) {
     hull.erase(it);
      return:
    update_left(it);
    it = hull.lower_bound(f);
    update_right(it);
  11 get(ll x) {
    Line f(0, 0, x, 1);
    iterator_t it = hull.upper_bound(f);
    assert(it != hull.begin());
   --it;
    return it->m * x + it->b;
private:
  set<Line> hull;
};
```

7 Miscellaneous

7.1 LIS - Longest Increasing Subsequence

```
int arr[ms], lisArr[ms], n;
// int bef[ms], pos[ms];
int lis() {
  int len = 1;
  lisArr[0] = arr[0];
  // bef[0] = -1;
  for (int i = 1; i < n; i++) {
    // upper bound se non-decreasing
    int x = lower_bound(lisArr, lisArr + len, arr[i]) - lisArr;
    len = max(len, x + 1);
    lisArr[x] = arr[i];
    // pos[x] = i;
    // bef[i] = x ? pos[x-1] : -1;
  return len;
vi getLis() {
  int len = lis();
  for(int i = pos[lisArr[len - 1]]; i >= 0; i = bef[i]) {
    ans.push_back(arr[i]);
```

```
reverse(ans.begin(), ans.end());
return ans;
```

7.2 Binary Search

```
int smallestSolution() {
   int x = -1;
   for(int b = z; b >= 1; b /= 2) {
     while(!ok(x+b)) x += b;
   }
   return x + 1;
}

int maximumValue() {
   int x = -1;
   for(int b = z; b >= 1; b /= 2) {
     while(f(x+b) < f(x+b+1)) x += b;
   }
   return x + 1;
}</pre>
```

7.3 Ternary Search

```
for (int i = 0; i < LOG; i++) {
  long double m1 = (A * 2 + B) / 3.0;
  long double m2 = (A + 2 * B) / 3.0;
  if(f(m1) > f(m2))
   A = m1;
  else
    B = m2:
ans = f(A);
//Z
while (B - A > 4) {
  int m1 = (A + B) / 2;
  int m2 = (A + B) / 2 + 1;
  if(f(m1) > f(m2))
   A = m1;
  else
    B = m2;
ans = inf;
for (int i = A; i \le B; i++) ans = min(ans, f(i));
```

7.4 Random Number Generator

```
// mt19937_64 se LL
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
// Random_Shuffle
shuffle(v.begin(), v.end(), rng);
// Random number in interval
int randomInt = uniform_int_distribution(0, i)(rng);
```

7.5 Submask Enumeration

```
for (int s=m; ; s=(s-1)&m) {
    ... you can use s ...
    if (s==0) break;
}
```

7.6 Java Fast I/O

```
import java.util.*;
import java.io.*;
// https://www.spoj.com/problems/INTEST/
 public static void main(String[] args) throws Exception{
    Reader s = new Reader();
   PrintWriter out = new PrintWriter(new BufferedOutputStream(System.
        out));
   int n = s.nextInt();
   int k = s.nextInt();
   int count=0;
   while (s.hasNext())
     int x = s.nextInt();
     if (x%k == 0)
     count++:
   out.printf("%d\n", count);
   out.close():
    s.close();
  // fast io
  static class Reader {
    final private int BUFFER_SIZE = 1 << 16;</pre>
   private DataInputStream din;
   private byte[] buffer;
   private int bufferPointer, bytesRead;
   public Reader() {
      din = new DataInputStream(System.in);
      buffer = new byte[BUFFER_SIZE];
      bufferPointer = bytesRead = 0;
    public Reader(String file_name) throws IOException {
      din = new DataInputStream(new FileInputStream(file_name));
     buffer = new byte[BUFFER_SIZE];
      bufferPointer = bvtesRead = 0:
    public String readLine() throws IOException {
      byte[] buf = new byte[64]; // line length
      int cnt = 0, c;
      while ((c = read()) != -1) {
       if (c == '\n') break;
```

```
buf[cnt++] = (byte) c;
  return new String(buf, 0, cnt);
public int nextInt() throws IOException {
  int ret = 0;
  byte c = read();
  while (c <= ' ') c = read();</pre>
  boolean neg = (c == '-');
  if (neg) c = read();
  do {
   ret = ret * 10 + c - '0';
  \} while ((c = read()) >= '0' && c <= '9');
  if (neg) return -ret;
  return ret;
public long nextLong() throws IOException {
  long ret = 0;
 byte c = read();
  while (c <= ' ') c = read();</pre>
 boolean neg = (c == '-');
  if (neg) c = read();
  do {
    ret = ret * 10 + c - '0';
  \} while ((c = read()) >= '0' && c <= '9');
  if (neg) return -ret;
  return ret;
public double nextDouble() throws IOException {
  double ret = 0, div = 1;
 byte c = read();
 while (c <= ' ')
  c = read():
 boolean neg = (c == '-');
  if (neg) c = read();
  do {
   ret = ret * 10 + c - '0';
  \} while ((c = read()) >= '0' && c <= '9');
  if (c == '.') {
   while ((c = read()) >= '0' && c <= '9') 
      ret += (c - '0') / (div *= 10);
  if (neg) return -ret;
  return ret;
private void fillBuffer() throws IOException {
 bytesRead = din.read(buffer, bufferPointer = 0, BUFFER_SIZE);
  if (bytesRead == -1) buffer[0] = -1;
public boolean hasNext() throws IOException {
```

```
|W| \le |vizinhosW| onde |W| eh quantos vertices tem em W
8.2 Math
   Goldbach's: todo numero par n > 2 pode ser representado com n = a + b
       onde a e b sao primos
   Twin prime: existem infinitos pares p, p + 2 onde ambos sao primos
   Legendre's: sempre tem um primo entre n^2 e (n+1)^2
   Lagrange's: todo numero inteiro pode ser inscrito como a soma de 4
       quadrados
   Zeckendorf's: todo numero pode ser representado pela soma de dois
       numeros de fibonnacis diferentes e nao consecutivos
   Euclid's: toda tripla de pitagoras primitiva pode ser gerada com
       (n^2 - m^2, 2nm, n^2+m^2) onde n, m sao coprimos e um deles eh par
   Wilson's: n \in primo quando (n-1)! \mod n = n - 1
   Mcnugget: Para dois coprimos x, y o maior inteiro que nao pode ser
       escrito como ax + by eh (x-1)(y-1)/2
   Fermat: Se p eh primo entao a(p-1) % p = 1
   Se x e m tambem forem coprimos entao x^k % m = x^k (k \mod (m-1)) % m
   Euler's theorem: x^(phi(m)) mod m = 1 onde phi(m) eh o totiente de
       euler
   Chinese remainder theorem:
   Para equacoes no formato x = a1 \mod m1, ..., x = an \mod mn onde todos
        os pares m1, \dots, mn sao coprimos
   Deixe Xk = m1 * m2 * ... * mn/mk e Xk^-1 mod mk = inverso de Xk mod mk, entao
   x = somatorio com k de 1 ate n de ak*Xk*(Xk,mk^-1 mod mk)
   Para achar outra solucao so somar m1*m2*..*mn a solucao existente
   Catalan number: exemplo expressoes de parenteses bem formadas
   C0 = 1, Cn = somatorio de <math>i=0 \rightarrow n-1 de Ci \star C(n-1+1)
   outra forma: Cn = (2n \text{ escolhe } n)/(n+1)
   Bertrand's ballot theorem: p votos tipo A e q votos tipo B com p>q,
       prob de em todo ponto ter mais As do que Bs antes dele = (p-q)/(p+
   Se puder empates entao prob = (p+1-q)/(p+1), para achar quantidade de
       possibilidades nos dois casos basta multiplicar por (p + q escolhe
```

houver aresta i-j ou 0 caso contrario, remove uma linha e uma coluna qualquer e o numero de spanning trees nesse grafo eh o det da matriz

Formula de Euler: V - E + F = 2 (para grafo planar)

Handshaking: Numero par de vertices tem grau impar

if (bufferPointer < bytesRead) return true;</pre>

if (bufferPointer == bytesRead) fillBuffer();

if(buffer[0] == -1) return false;

private byte read() throws IOException {

public void close() throws IOException {

return buffer[bufferPointer++];

Teoremas e formulas uteis

if (din == null) return;

fillBuffer();

return true;

din.close();

Grafo contem caminho hamiltoniano se:
Dirac's theorem: Se o grau de cada vertice for pelo menos n/2
Ore's theorem: Se a soma dos graus que cada par nao-adjacente de
 vertices for pelo menos n

Kirchhoff's Theorem: Monta matriz onde Mi, i = Grau[i] e Mi, i = -1 se

Trees:

8.1 Grafos

Tem Catalan(N) Binary trees de N vertices
Tem Catalan(N-1) Arvores enraizadas com N vertices
Caley Formula: n^(n-2) arvores em N vertices com label
Prufer code: Cada etapa voce remove a folha com menor label e o label
 do vizinho eh adicionado ao codigo ate ter 2 vertices

Flow:

Max Edge-disjoint paths: Max flow com arestas com peso 1
Max Node-disjoint paths: Faz a mesma coisa mas separa cada vertice em
um com as arestas de chegadas e um com as arestas de saida e uma
aresta de peso 1 conectando o vertice com aresta de chegada com
ele mesmo com arestas de saida

Konig's Theorem: minimum node cover = maximum matching se o grafo for bipartido, complemento eh o maximum independent set

Min Node disjoint path cover: formar grafo bipartido de vertices duplicados, onde aresta sai do vertice tipo A e chega em tipo B, entao o path cover eh N - matching

Min General path cover: Mesma coisa mas colocando arestas de A pra B sempre que houver caminho de A pra B

Dilworth's Theorem: Min General Path cover = Max Antichain (set de vertices tal que nao existe caminho no grafo entre vertices desse set)

Hall's marriage: um grafo tem um matching completo do lado X se para cada subconjunto W de X,

```
Distribuicao uniforme a,a+1, ..., b Expected[X] = (a+b)/2 Distribuicao binomial com n tentativas de probabilidade p, X = sucessos:
```

Hockey-stick: Somatorio de $i = r \rightarrow n$ de (i escolhe r) = (n + 1)

Vandermonde: $(m+n \ escolhe \ r) = somatorio \ de \ k = 0 -> r \ de \ (m \ escolhe \ k$

Burnside lemma: colares diferentes nao contando rotacoes quando m =

 $\label{eq:problem} \begin{array}{lll} P\left(X=x\right) = p^{n}x \;\star\; (1-p)^{n}(n-x) \;\star\; (n \; escolhe \; x) \; e \; E[X] = p \star n \\ \text{Distribuicao geometrica onde continuamos ate ter sucesso, } X = \\ & \text{tentativas:} \end{array}$

 $P(X = x) = (1-p)^(x-1) * p \in E[X] = 1/p$

 $(m^n + somatorio i = 1 - > n-1 de m^gcd(i, n))/n$

escolhe r + 1)

) \star (n escolhe r - k)

cores e n = comprimento

Linearity of expectation: Tendo duas variaveis X e Y e constantes a e
b, o valor esperado de aX + bY = a*E[X] + b*E[X]

8.3 Geometry

Formula de Euler: V - E + F = 2

Pick Theorem: Para achar pontos em coords inteiras num poligono Area = i + b/2 - 1 onde i eh o o numero de pontos dentro do poligono e b de pontos no perimetro do poligono

Two ears theorem: Todo poligono simples com mais de 3 vertices tem pelo menos 2 orelhas, vertices que podem ser removidos sem criar um crossing, remover orelhas repetidamente triangula o poligono

Incentro triangulo: (a(Xa, Ya) + b(Xb, Yb) + c(Xc, Yc))/(a+b+c) onde a = lado oposto ao vertice a, incentro eh onde cruzam as bissetrizes, eh o centro da circunferencia inscrita e eh equidistante aos lados

8.4 Mersenne's Primes

Primos de Mersenne 2^n - 1
Lista de Ns que resultam nos primeiros 41 primos de Mersenne:
2; 3; 5; 7; 13; 17; 19; 31; 61; 89; 107; 127; 521; 607; 1.279; 2.203;
2.281; 3.217; 4.253; 4.423; 9.689; 9.941; 11.213; 19.937; 21.701;
23.209; 44.497; 86.243; 110.503; 132.049; 216.091; 756.839;
859.433; 1.257.787; 1.398.269; 2.976.221; 3.021.377; 6.972.593;
13.466.917; 20.996.011; 24.036.583;