Amigos do Beto - ICPC Library

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1 Data Structures

KMP

1.1 BIT 2D Comprimida

```
template<class T = int>
struct Bit2D {
public:
  // send updated points
  Bit2D(vector<pair<T, T>> pts) {
   sort(pts.begin(), pts.end());
   for(auto a : pts)
     if(ord.empty() || a.first != ord.back()) {
       ord.push_back(a.first);
   fw.resize(ord.size() + 1);
   coord.resize(fw.size());
   for(auto &a : pts) {
     swap(a.first, a.second);
   sort(pts.begin(), pts.end());
   for(auto &a : pts) {
     swap(a.first, a.second);
     for(int on = upper_bound(ord.begin(), ord.end(), a.first) - ord.begin(); on < fw</pre>
          .size(); on += on & -on) {
       if(coord[on].empty() || coord[on].back() != a.second) {
         coord[on].push_back(a.second);
   for(int i = 0; i < fw.size(); i++) {</pre>
     fw[i].assign(coord[i].size() + 1, 0);
  void upd(T x, T y, T v) {
   for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin(); xx < fw.size();</pre>
         xx += xx & -xx) {
      for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y) - coord[xx].
       T qry(T x, T y) {
  T ans = 0;
   for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin(); xx > 0; xx -=
        xx & -xx) {
```

1.2 Seg Tree Lazy

```
int arr[ms], seg[4 * ms], lazy[4 * ms], n;
struct LazyContext {
  LazyContext() { }
 void reset() { }
  void operator += (LazyContext o) { }
struct Node {
 Node() { }
Node() { }
 Node (Node 1, Node r) { }
  bool canBreak(LazyContext lazy) { } // false if non beats
  bool canApply(LazyContext lazy) { } // true if non beats
  void apply(LazyContext &lazy) { }
void build(int idx = 0, int l = 0, int r = n-1) {
 int mid = (1+r)/2;
  lazy[idx] = 0;
 if(\bar{1} == r) {
    seg[idx] = arr[l];
  build(2*idx+1, 1, mid); build(2*idx+2, mid+1, r);
  seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
void apply(int idx, int 1, int r) {
  if(lazy[idx] && !canBreak) { // if not beats canBreak = false
   if(1 < r) {
      lazy[2*idx+1] += lazy[idx]; // Merge de lazy
      lazy[2*idx+2] += lazy[idx]; // Merge de lazy
    if(canApply) { // if not beats canApply = true
      seg[idx] += lazy[idx] * (r - l + 1); // Aplicar lazy no seg
      apply (2*idx+1, 1, mid); apply (2*idx+2, mid+1, r);
      seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
  lazy[idx] = 0; // Limpar a lazy
int query(int L, int R, int idx = 0, int 1 = 0, int r = n-1) {
  int mid = (1+r)/2;
 apply(idx, l, r);
 if(1 > R || r < L) return 0; // Valor que nao atrapalhe</pre>
  if(L <= 1 && r <= R) return seg[idx];</pre>
  return query(L, R, 2*idx+1, 1, mid) + query(L, R, 2*idx+2, mid+1, r); // Merge
void update(int L, int R, int V, int idx = 0, int 1 = 0, int r = n-1) {
  int mid = (1+r)/2;
 apply(idx, l, r);
  if(1 > R \mid \mid r < L) return;
 if(L <= 1 && r <= R) {
   lazy[idx] = V;
```

```
apply(idx, 1, r);
  return;
}
update(L, R, V, 2*idx+1, 1, mid); update(L, R, V, 2*idx+2, mid+1, r);
seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
}
```

1.3 Persistent Segment Tree

```
struct Node {
        int v = 0;
        Node *1 = this, *r = this;
};
int CNT = 1;
Node buffer[ms * 20];
Node* update(Node *root, int 1, int r, int idx, int val){
   Node *node = &buffer[CNT++];
        *node = *root;
        int mid = (1 + r) / 2;
        node->v += val;
        if(1+1 != r){
                 if(idx < mid) node->1 = update(root->1, 1, mid, idx, val);
                 else node->r = update(root->r, mid, r, idx, val);
        return node;
int query(Node *node, int tl, int tr, int l, int r){
        if(1 <= t1 && tr <= r) return node->v;
        if(tr <= 1 || t1 >= r) return 0;
        int mid = (tl+tr) / 2;
        return query(node->1, t1, mid, 1, r) + query(node->r, mid, tr, 1, r);
```

1.4 Treap

```
mt19937 rng ((int) chrono::steady_clock::now().time_since_epoch().count());
typedef int Value;
typedef struct item * pitem;
struct item {
 item () {}
  item (Value v) { // add key if not implicit
   value = v:
    prio = uniform_int_distribution<int>() (rng);
    cnt = 1;
    rev = 0;
    1 = r = 0;
  pitem 1. r:
  Value value;
  int prio, cnt;
  bool rev;
int cnt (pitem it) {
  return it ? it->cnt : 0;
void fix (pitem it) {
  if (it)
    it\rightarrow cnt = cnt(it\rightarrow 1) + cnt(it\rightarrow r) + 1;
void pushLazy (pitem it) {
  if (it && it->rev) {
    it->rev = false;
    swap(it->1, it->r);
    if (it->1) it->1->rev ^= true;
    if (it->r) it->r->rev ^= true;
void insert (pitem & t, pitem it) {
  if (!t)
    t = it;
  else if (it->prio > t->prio)
    split (t, it->key, it->l, it->r), t = it;
  else
    insert (t->key <= it->key ? t->r : t->l, it);
```

```
void merge (pitem & t, pitem 1, pitem r) {
  pushLazy (1); pushLazy (r);
  if (!1 || !r) t = 1 ? 1 : r;
  else if (l->prio > r->prio)
    merge (1->r, 1->r, r), t = 1;
    merge (r->1, 1, r->1), t = r;
  fix (t);
void erase (pitem & t, int key) {
  if (t->key == key) {
    pitem th = t;
    merge (t, t->1, t->r);
    delete th;
  else
    erase (key < t->key ? t->l : t->r, key);
void split (pitem t, pitem & 1, pitem & r, int key) {
  if (!t) return void( 1 = r = 0 );
  pushLazy (t);
  int cur_key = cnt(t->1); // t->key if not implicit
  if (key <= cur_key)</pre>
    split (t->1, 1, t->1, key), r = t;
  else
   split (t->r, t->r, r, key - (1 + cnt(t->1))), 1 = t;
  fix (t);
void reverse (pitem t, int 1, int r) {
  pitem t1, t2, t3;
  split (t, t1, t2, 1);
  split (t2, t2, t3, r-1+1);
  t2->rev ^= true;
  merge (t, t1, t2);
  merge (t, t, t3);
void unite (pitem & t, pitem 1, pitem r) {
  if (!l || !r) return void ( t = l ? l : r );
  if (l->prio < r->prio) swap (l, r);
  pitem lt, rt;
  split (r, lt, rt, l->key);
  unite (1->1, 1->1, 1t);
  unite (1-> r, 1->r, rt);
  t = 1;
pitem kth_element(pitem t, int k) {
        if(!t) return NULL;
        if(t->1)
                if(t->l->size >= k) return kth_element(t->l, k);
                else k = t->l->cnt;
        return (k == 1) ? t : kth_element(t->r, k - 1);
int countLeft(pitem t, int key) {
        if(!t) {
                return 0:
        } else if(t->key < key) {</pre>
                return 1 + (t\rightarrow 1 ? t\rightarrow 1\rightarrow size : 0) + countLeft(t\rightarrow r, key);
        ) else
                return countLeft(t->1, key);
```

1.5 KD-Tree

```
return;
    int n = 0:
    tree.resize(2 * pts.size());
    build(pts.begin(), pts.end(), n);
  long long nearestNeighbor(PT point)
    long long ans = (long long) le18;
    nearestNeighbor(&tree[0], point, 0, ans);
    return ans;
private:
  std::vector<Node> tree;
  Node* build(std::vector<PT>::iterator 1, std::vector<PT>::iterator r, int &n, int h
    int id = n++;
    if(r - 1 == 1) {
      tree[id].left = tree[id].right = NULL;
      tree[id].point = *1;
    \} else if (r - 1 > 1) {
      std::vector\langle PT \rangle::iterator mid = 1 + ((r - 1) / 2);
      d = h;
      std::nth_element(1, mid - 1, r, comp);
      tree[id].point = *(mid - 1);
      // BE CAREFUL!
// DO EVERYTHING BEFORE BUILDING THE LOWER PART!
      tree[id].left = build(l, mid, n, h^1);
      tree[id].right = build(mid, r, n, h^1);
    return &tree[id];
  void nearestNeighbor(Node* node, PT point, int h, long long &ans) {
    if(!node) {
      return;
    if(point != node->point)
      // THIS WAS FOR A PROBLEM
// THAT YOU DON'T CONSIDER THE DISTANCE TO ITSELF!
      ans = std::min(ans, sqrDist(point, node->point));
    long long delta = getValue(point) - getValue(node->point);
    if(delta <= 0) {
      nearestNeighbor(node->left, point, h^1, ans);
      if(ans > delta * delta) {
        nearestNeighbor(node->right, point, h^1, ans);
      nearestNeighbor(node->right, point, h^1, ans);
      if(ans > delta * delta)
        nearestNeighbor(node->left, point, h^1, ans);
};
```

1.6 Sparse Table

```
vector<vector<int>> table;
vector<int> lg2;
void build(int n, vector<int> v) {
  lg2.resize(n + 1);
  lg2[1] = 0;
  for (int i = 2; i <= n; i++) {
    lg2[i] = lg2[i >> 1] + 1;
  table.resize(lg2[n] + 1);
  for (int i = 0; i < lg2[n] + 1; i++) {
    table[i].resize(n + 1);
  for (int i = 0; i < n; i++) {</pre>
    table[0][i] = v[i];
  for (int i = 0; i < lg2[n]; i++) {</pre>
    for (int j = 0; j < n; j++) {
      if (j + (1 \ll i) >= n) break;
      table[i + 1][j] = min(table[i][j], table[i][j + (1 << i)]);
```

```
} 
int get(int 1, int r) {
   int k = lg2[r - 1 + 1];
   return min(table[k][1], table[k][r - (1 << k) + 1]);
}</pre>
```

1.7 Policy Based Structures

```
#include <ext/pb_ds/assoc_container.hpp> // Common file
#include <ext/pb_ds/tree_policy.hpp> // Including tree_order_statistics_node_update
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
ordered_set X;
X.insert(1); X.find_by_order(0);
X.order_of_key(-5); end(X), begin(X);
```

1.8 Color Updates Structure

```
struct range {
 int 1, r;
 range(int 1 = 0, int r = 0, int v = 0) : 1(1), r(r), v(v) {}
 bool operator < (const range &a) const {
   return 1 < a 1;
set<range> ranges;
vector<range> update(int 1, int r, int v) { // [1, r)
 vector<range> ans;
 if(l >= r) return ans;
 auto it = ranges.lower_bound(1);
 if(it != ranges.begin()) {
   it--;
   if(it->r>1) {
     auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, 1, cur.v));
      ranges.insert(range(l, cur.r, cur.v));
 it = ranges.lower_bound(r);
 if(it != ranges.begin()) {
   it--;
   if(it->r>r) {
     auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.l, r, cur.v));
      ranges.insert(range(r, cur.r, cur.v));
 for(it = ranges.lower_bound(1); it != ranges.end() && it->1 < r; it++) {</pre>
   ans.push_back(*it);
 ranges.erase(ranges.lower_bound(1), ranges.lower_bound(r));
 ranges.insert(range(l, r, v));
int query(int v) { // Substituir -1 por flag para quando nao houver resposta
 auto it = ranges.upper_bound(v);
 if(it == ranges.begin()) {
   return -1;
 it--:
 return it->r>=v ? it->v : -1;
```

2 Graph Algorithms

2.1 Blossom

```
#define MAXN 110
#define MAXM MAXN*MAXN
int n, m;
int mate[MAXN], first[MAXN], label[MAXN];
int adj[MAXN] [MAXN], nadj[MAXN], from[MAXM], to[MAXM];
queue<int> q;
#define OUTER(x) (label[x] >= 0)
void L(int x, int y, int nxy) {
  int join, v, r = first[x], s = first[y];
  if (r == s) { return; }
  nxy += n + 1;
  label[r] = label[s] = -nxy;
  while (1) {
    if (s != 0) { swap(r, s); }
    r = first[label[mate[r]]];
    if (label[r] != -nxy) { label[r] = -nxy; }
    else {
      break:
  v = first[x];
  while (v != join) {
    if (!OUTER(v)) { q.push(v); }
    first[v] = join;
    v = first[label[mate[v]]];
  v = first[y];
  while (v != join) {
    if (!OUTER(v)) { q.push(v); }
    label[v] = nxy;
    first[v] = join;
    v = first[label[mate[v]]];
  for (int i = 0; i <= n; i++) {
    if (OUTER(i) && OUTER(first[i])) { first[i] = join; }
void R(int v, int w) {
  int t = mate[v];
  mate[v] = w;
  if (mate[t] != v) { return;
  if (label[v] >= 1 \&\& label[v] <= n) {
    mate[t] = label[v];
R(label[v], t);
  int x = from[label[v] - n - 1], y = to[label[v] - n - 1];
  R(x, y);
R(y, x);
int E() {
  memset(mate, 0, sizeof(mate));
  int r = 0;
  bool e7;
  for (int u = 1; u \le n; u++) {
   memset(label, -1, sizeof(label));
    while (!q.empty()) { q.pop(); }
    if (mate[u]) { continue; }
    label[u] = first[u] = 0;
    q.push(u);
e7 = false;
    while (!q.empty() && !e7) {
      int x = q.front();
      q.pop();
      for (int i = 0; i < nadj[x]; i++) {</pre>
        int y = from[adj[x][i]];
        if (y == x) { y = to[adj[x][i]]; }
        if (!mate[y] && y != u) {
  mate[y] = x;
          R(x, y);
          r++;
e7 = true;
          break;
```

```
} else if (OUTER(y)) { L(x, y, adj[x][i]); }
          int v = mate[y];
          if (!OUTER(v)) {
           label[v] = x;
           first[v] = y;
           q.push(v);
    label[0] = -1;
  return r;
/*Exemplo simples de uso*/
memset(nadj, 0, sizeof nadj);
for (int i = 0; i < m; ++i) { // arestas
  scanf("%d%d", &a, &b);
 a++, b++; // nao utilizar o vertice 0
  adj[a][nadj[a]++] = i;
  adj[b][nadj[b]++] = i;
  from[i] = a;
 to[i] = b;
printf("O emparelhamento tem tamanho %d\n", E());
for (int i = 1; i <= n; i++) {
  if (mate[i] > i) { printf("%d com %d\n", i - 1, mate[i] - 1); }
```

2.2 Dinic Max Flow

```
const int ms = 1e3; // vertices
const int me = 1e5; // arestas
int adj[ms], to[me], ant[me], wt[me], z, n;
int copy_adj[ms], fila[ms], level[ms];
void clear() { // Lembrar de chamar no main
  memset(adj, -1, sizeof adj);
  z = 0;
void add(int u, int v, int k) {
  to[z] = v;
  ant[z] = adj[u];
wt[z] = k;
  adj[u] = z++;
  swap(u, v);
to[z] = v;
 ant[z] = adj[u];
wt[z] = 0; // Lembrar de colocar = 0
  adj[u] = z++;
int bfs(int source, int sink) {
  memset(level, -1, sizeof level);
level[source] = 0;
  int front = 0, size = 0, v;
  fila[size++] = source;
  while(front < size) {</pre>
    v = fila[front++];
    for(int i = adj[v]; i != -1; i = ant[i]) {
      if(wt[i] && level[to[i]] == -1) {
        level[to[i]] = level[v] + 1;
        fila[size++] = to[i];
  return level[sink] != -1;
int dfs(int v, int sink, int flow) {
  if(v == sink) return flow;
  for(int &i = copy_adj[v]; i != -1; i = ant[i]) {
    if(wt[i] && level[to[i]] == level[v] + 1 &&
      (f = dfs(to[i], sink, min(flow, wt[i])))) {
      wt[i] -= f;
      wt[i ^ 1] += f;
return f;
  return 0;
```

```
int maxflow(int source, int sink) {
  int ret = 0, flow;
  while(bfs(source, sink)) {
    memcpy(copy_adj, adj, sizeof adj);
    while((flow = dfs(source, sink, 1 << 30))) {
      ret += flow;
    }
  }
  return ret;</pre>
```

2.3 Min Cost Max Flow

```
template <class T = int>
class MCMF
public:
  struct Edge {
   Edge(int a, T b, T c) : to(a), cap(b), cost(c) {}
   int to;
   T cap, cost;
  }:
  MCMF (int size) {
   n = size;
    edges.resize(n);
   pot.assign(n, 0);
   dist.resize(n);
    visit.assign(n, false);
  pair<T, T> mcmf(int src, int sink) {
   pair<T, T> ans(0, 0);
    if(!SPFA(src, sink)) return ans;
    fixPot();
    // can use dijkstra to speed up depending on the graph
    while(SPFA(src, sink)) {
      auto flow = augment(src, sink);
      ans.first += flow.first;
      ans.second += flow.first * flow.second;
      fixPot();
    return ans;
  void addEdge(int from, int to, T cap, T cost) {
    edges[from].push_back(list.size());
    list.push_back(Edge(to, cap, cost));
    edges[to].push_back(list.size());
    list.push_back(Edge(from, 0, -cost));
private:
  int n;
  vector<vector<int>> edges;
  vector<Edge> list;
  vector<int> from;
  vector<T> dist, pot;
  vector<bool> visit;
  /*bool dij(int src, int sink) {
   T INF = std::numeric_limits<T>::max();
    dist.assign(n, INF);
    from.assign(n, -1);
    visit.assign(n, false);
    dist[src] = 0;
    for(int i = 0; i < n; i++) {
      int best = -1;
      for (int j = 0; j < n; j++) {
        if(visit[j]) continue;
        if(best == -1 || dist[best] > dist[j]) best = j;
      if(dist[best] >= INF) break;
      visit[best] = true;
      for(auto e : edges[best]) {
        auto ed = list[e];
        if (ed.cap == 0) continue;
        T toDist = dist[best] + ed.cost + pot[best] - pot[ed.to];
        assert(toDist >= dist[best]);
        if(toDist < dist[ed.to]) {</pre>
          dist[ed.to] = toDist;
          from[ed.to] = e;
```

```
return dist[sink] < INF;
  pair<T, T> augment(int src, int sink) {
    pair<T, T> flow = {list[from[sink]].cap, 0};
    for(int v = sink; v != src; v = list[from[v]^1].to) {
      flow.first = min(flow.first, list[from[v]].cap);
      flow.second += list[from[v]].cost;
    for(int v = sink; v != src; v = list[from[v]^1].to) {
     list[from[v]].cap -= flow.first;
      list[from[v]^1].cap += flow.first;
    return flow;
 queue<int> q;
  bool SPFA(int src, int sink) {
   T INF = numeric_limits<T>::max();
    dist.assign(n, INF);
    from.assign(n, -1);
    q.push(src);
    dist[src] = 0;
    while(!q.emptv()) {
      int on = q.front();
      q.pop();
      visit[on] = false;
      for(auto e : edges[on]) {
        auto ed = list[e];
        if(ed.cap == 0) continue;
        T toDist = dist[on] + ed.cost + pot[on] - pot[ed.to];
        if(toDist < dist[ed.to]) {</pre>
          dist[ed.to] = toDist;
          from[ed.to] = e;
          if(!visit[ed.to])
            visit[ed.to] = true;
            q.push(ed.to);
    return dist[sink] < INF;</pre>
 void fixPot() {
   T INF = numeric limits<T>::max();
    for(int i = 0; i < n; i++) {</pre>
      if(dist[i] < INF) pot[i] += dist[i];</pre>
};
```

2.4 Euler Path and Circuit

```
int del[me],adj[ms], to[me], ant[me], wt[me], z, n;
vector<int> pathE, pathV;
// Funcao de add e clear no dinic
void eulerPath(int u) {
  for(int &i = adj[u]; ~i; i = ant[i]) if(!del[i]) {
    del[i] = del[i^1] = 1;
    eulerPath(to[i]);
    pathE.emplace_back(i);
  }
  pathV.emplace_back(u);
}
```

2.5 Articulation Points/Bridges/Biconnected Components

```
int adj[ms], to[me], ant[me], z;
int num[ms], low[ms], timer;
bool art[ms], bridge[me], f[me];
int bc[ms], nbc;
stack<int> st, stk;
vector<vector<int>> comps;
```

```
void clear() { // Lembrar de chamar no main
  memset(adj, -1, sizeof adj);
  7 = 0:
void add(int u, int v) {
  to[z] = v;
  ant[z] = adj[u];
  adj[u] = z++;
void generateBc (int v) {
  while (!st.empty()) {
    int u = st.top();
    st.pop();
bc[u] = nbc;
    if (v == u) break;
  ++nbc;
void dfs (int v, int p) {
  st.push(v), stk.push(v);
  low[v] = num[v] = ++timer;
  for (int i = adj[v]; i != -1; i = ant[i]) {
    if (f[i] || f[i^1]) continue;
    f[i] = 1;
    int u = to[i];
    if (num[u] == -1) {
      dfs(u, v);
      low[v] = min(low[v], low[u]);
      if (low[u] > num[v]) bridge[i] = bridge[i^1] = 1;
      if (low[u] >= num[v]) {
        art[v] = (num[v] > 1 || num[u] > 2);
        comps.push_back({v});
        while (comps.back().back() != u)
          comps.back().push_back(stk.top()), stk.pop();
    } else {
      low[v] = min(low[v], num[u]);
  if (low[v] == num[v]) generateBc(v);
void biCon (int n) {
  nbc = 0, timer = 0;
  memset (num, -1, sizeof num);
  memset(bc, -1, sizeof bc);
  memset (bridge, 0, sizeof bridge);
  memset(art, 0, sizeof art);
  memset(f, 0, sizeof f);
  for (int i = 0; i < n; i++) {
    if (num[i] == -1) {
      timer = 0;
      dfs(i, 0);
vector<int> g[ms];
int id[ms];
void buildBlockCut (int n) {
  int z = 0;
  for (int u = 0; u < n; ++u) {
    if (art[u]) id[u] = z++;
  for (auto &comp : comps) {
    int node = z++;
    for (int u : comp)
      if (!art[u]) id[u] = node;
        g[node].push_back(id[u]);
        g[id[u]].push_back(node);
```

2.6 SCC - Strongly Connected Components / 2SAT

```
const int ms = 212345;
vector<int> g[ms];
int idx[ms], low[ms], z, comp[ms], ncomp;
stack<int> st;
int dfs(int u) {
  if(~idx[u]) return idx[u] ? idx[u] : z;
  low[u] = idx[u] = z++;
  st.push(u);
  for(int v : q[u]) {
    low[u] = min(low[u], dfs(v));
  if(low[u] == idx[u]) {
    while(st.top() != u) {
      int v = st.top();
       idx[v] = 0;
      low[v] = low[u];
comp[v] = ncomp;
      st.pop();
    idx[st.top()] = 0;
    st.pop();
comp[u] = ncomp++;
  return low[u];
bool solveSat(int n) {
  memset(idx, -1, sizeof idx);
  z = 1; ncomp = 0;
  for (int i = 0; i < 2*n; i++) dfs(i);
  for(int i = 0; i < 2*n; i++) if(comp[i] == comp[i^1]) return false;</pre>
int trad(int v) { return v < 0 ?(~v) *2^1 : v * 2; }</pre>
void add(int a, int b) { g[trad(a)].push_back(trad(b)); }
void addOr(int a, int b) { add(~a, b); add(~b, a); }
void addImp(int a, int b) { addOr(~a, b); }
void addEqual(int a, int b) { addOr(a, ~b); addOr(~a, b); }
void addDiff(int a, int b) { addEqual(a, ~b); }
// value[i] = comp[trad(i)] < comp[trad(~id)];</pre>
```

2.7 LCA O(1)

```
template<class T>
struct RMQ {
  vector<vector<T>> jmp;
  RMQ(const vector<T>& V) : jmp(1, V) {
    for (int pw = 1, k = 1; pw * 2 <= (int) size(V); pw *= 2, ++k) {
      jmp.emplace_back(size(V) - pw * 2 + 1);
      for (int j = 0; j < (int) size(jmp[k]); ++j)
        jmp[k][j] = min(jmp[k - 1][j], jmp[k - 1][j + pw]);
 T query(int a, int b) {
   assert(a < b); // or return inf if a == b</pre>
    int dep = 31 - __builtin_clz(b - a);
    return min(jmp[dep][a], jmp[dep][b - (1 << dep)]);</pre>
};
struct LCA {
  int T = 0;
  vector<int> time, path, ret;
  RMQ<int> rmq;
  LCA(vector<vector<int>>& C) : time(size(C)), rmq((dfs(C,0,-1), ret)) {}
  void dfs(vector<vector<int>>& C, int v, int par) {
   time[v] = T++;
    for (int y : C[v]) if (y != par) {
      path.push_back(v), ret.push_back(time[v]);
      dfs(C, y, v);
  int lca(int a, int b) {
   if (a == b) return a;
    tie(a, b) = minmax(time[a], time[b]);
```

```
return path[rmq.query(a, b)];
};
```

2.8 Heavy Light Decomposition

```
class HLD {
public:
  void init(int n) { /* resize everything */ }
  void addEdge(int u, int v) {
    edges[u].push_back(v);
    edges[v].push_back(u);
  void setRoot(int r) {
    t = 0;
    p[r] = r;
h[r] = 0;
    prep(r, r);
nxt[r] = r;
    hld(r);
  int getLCA(int u, int v) {
    while(!inSubtree(nxt[u], v)) u = p[nxt[u]];
    while(!inSubtree(nxt[v], u)) v = p[nxt[v]];
    return in[u] < in[v] ? u : v;</pre>
  // is v in the subtree of u?
  bool inSubtree(int u, int v)
    return in[u] <= in[v] && in[v] < out[u];</pre>
  // returns ranges [1, r) that the path has
  vector<pair<int, int>> getPath(int u, int anc) {
    vector<std::pair<int, int>> ans;
    //assert(inSubtree(anc, u));
    while(nxt[u] != nxt[anc]) {
      ans.emplace_back(in[nxt[u]], in[u] + 1);
      u = p[nxt[u]];
    // this includes the ancestor! care
    ans.emplace_back(in[anc], in[u] + 1);
    return ans;
private:
  vector<int> in, out, p, rin, sz, nxt, h;
  vector<vector<int>> edges;
  int t;
  void prep(int on, int par) {
    sz[on] = 1;
p[on] = par;
    for(int i = 0; i < (int) edges[on].size(); i++) {</pre>
      int &u = edges[on][i];
      if(u == par) {
        swap(u, edges[on].back());
        edges[on].pop_back();
      } else {
  h[u] = 1 + h[on];
        prep(u, on);
sz[on] += sz[u];
        if(sz[u] > sz[edges[on][0]]) {
          swap(edges[on][0], u);
  void hld(int on) {
    in[on] = t++;
    rin[in[on]] = on;
    for(auto u : edges[on]) {
  nxt[u] = (u == edges[on][0] ? nxt[on] : u);
      hld(u);
    out[on] = t;
};
```

2.9 Centroid Decomposition

```
vector<int> g[ms];
int dis[ms][30];
int par[ms], sz[ms], rem[ms], h[ms];
void dfsSubtree(int u, int p) {
  sz[u] = 1;
  for(auto v : q[u]) {
   if(v != p && !rem[v]) {
      dfsSubtree(v, u);
      sz[u] += sz[v];
int getCentroid(int u, int p, int size) {
  for(auto v : g[u]) {
   if(v != p \&\& !rem[v] \&\& sz[v] * 2 >= size)
      return getCentroid(v, u, size);
 return u;
void setDis(int u, int p, int nv){
  for (auto v : g[u]) {
   if (v == p || rem[v]) continue;
   dis[v][nv] = dis[u][nv]+1;
    setDis(v, u, nv);
void decompose (int u, int p = -1, int nv = 0) {
  dfsSubtree(u, -1);
  int ctr = getCentroid(u, -1, sz[u]);
  par[ctr] = p;
  h[ctr] = nv;
  rem[ctr] = 1;
  setDis(ctr, p, nv);
 for(auto v : g[ctr]) {
   if(v != p && !rem[v]) {
      decompose(v, ctr, nv+1);
```

2.10 Hungarian Algorithm - Maximum Cost Matching

```
int u[ms], v[ms], p[ms], way[ms], minv[ms];
bool used[ms];
pair<vector<int>, int> solve(const vector<vector<int>> &matrix) {
  int n = matrix.size();
  if(n == 0) return {vector<int>(), 0};
  int m = matrix[0].size();
  assert (n <= m);
  memset(u, 0, (n+1)*sizeof(int));
  memset(v, 0, (m+1)*sizeof(int));
  memset(p, 0, (m+1)*sizeof(int));
  for(int i = 1; i <= n; i++) {
    memset (minv, 0x3f, (m+1)*sizeof(int));
    memset(way, 0, (m+1)*sizeof(int));
    for(int j = 0; j <= m; j++) used[j] = 0;</pre>
    p[0] = i;
    int k0 = 0;
      used[k0] = 1;
      int i0 = p[k0], delta = inf, k1;
      for(int j = 1; j <= m; j++) {</pre>
        if(!used[j]) {
          int cur = matrix[i0-1][j-1] - u[i0] - v[j];
          if (cur < minv[j]) {</pre>
            minv[j] = cur;
            way[j] = k0;
          if(minv[j] < delta) {</pre>
            delta = minv[j];
            k1 = j;
      for (int j = 0; j \le m; j++) {
        if(used[j]) {
```

```
u[p[j]] += delta;
  v[j] -= delta;
  } else {
    minv[j] -= delta;
  }
  k0 = k1;
} while(p[k0]);
do {
  int k1 = way[k0];
  p[k0] = p[k1];
  k0 = k1;
} while(k0);
}
vector<int> ans(n, -1);
for(int j = 1; j <= m; j++) {
  if(!p[j]) continue;
  ans[p[j] - 1] = j - 1;
}
return {ans, -v[0]};</pre>
```

2.11 Minimum Arborescence

```
// uncommented O(V^2) arborescence
struct Edges {
  //set<pair<long long, int>> cost; O(Elog^2)
  long long cost[ms];
  // possible optimization, use vector of size n
  // instead of ms
  long long sum = 0;
  Edges() {
    memset(cost, 0x3f, sizeof cost);
  void addEdge(int u, long long c) {
    // cost.insert({c - sum, u}); O(Elog^2)
    cost[u] = min(cost[u], c - sum);
  pair<long long, int> getMin() {
    //return *cost.begin(); O(E*log^2)
    pair<long long, int> ans(cost[0], 0);
    // in this loop can change ms to n to make it faster for many cases
    for(int i = 1; i < ms; i++) {
      if(cost[i] < ans.first)</pre>
        ans = pair<long long, int>(cost[i], i);
    return ans;
  void unite(Edges &e) {
    O(E*log^2E)
    if(e.cost.size() > cost.size()) {
      cost.swap(e.cost);
      swap(sum, e.sum);
    for(auto i : e.cost) {
      addEdge(i.second, i.first + e.sum);
    e.cost.clear();
    // O(V^2)
    // can change ms to n
    for(int i = 0; i < ms; i++) {</pre>
      cost[i] = min(cost[i], e.cost[i] + e.sum - sum);
};
typedef vector<vector<pair<long long, int>>> Graph;
int par[ms];
long long best[ms];
int col[ms];
int getPar(int x) { return par[x] < 0 ? x : par[x] = getPar(par[x]); }</pre>
void makeUnion(int a, int b) {
 a = getPar(a);
b = getPar(b);
  if(a == b) return;
  ed[a].unite(ed[b]);
  par[b] = a;
```

```
long long arborescence(Graph edges) {
  // root is 0
  // edges has transposed adjacency list (cost, from)
  // edge from i to j cost c is
  // edge[j].emplace_back(c, i)
 int n = (int) edges.size();
 long long ans = 0;
  for (int i = 0; i < n; i++) {
   ed[i] = Edges();
    par[i] = -1;
    for(auto j : edges[i]) {
      ed[i].addEdge(j.second, j.first);
   col[i] = 0;
  // to change the root you can simply change this next line to
  // col[root] = 2;
  col[0] = 2;
  for (int i = 0; i < n; i++) {
   if(col[getPar(i)] == 2) {
      continue;
    int on = getPar(i);
    vector<int> st;
    while(col[on] != 2) {
  assert(getPar(on) == on);
      if(col[on] == 1) {
        int v = on;
        vector<int> cycle;
//cout << "found cycle\n";</pre>
        while(st.back() != v) {
          //cout << st.back() << endl;</pre>
          cycle.push_back(st.back());
          st.pop_back();
        for(auto u : cycle) { // compress cycle
          makeUnion(v, u);
        v = getPar(v);
        col[v] = 0;
      } else {
        // still no cycle
        // while best is in compressed cycle, remove
        // THIS IS TO MAKE O(E*log^2) ALGORITHM!!
        // while(!ed[on].cost.empty() && getPar(on) == getPar(ed[on].getMin().second))
             ed[on].cost.erase(ed[on].cost.begin());
        // O(V^2)
        for (int x = 0; x < n; x++) {
          if(on == getPar(x)) {
            ed[on].cost[x] = 0x3f3f3f3f3f3f3f3f1L;
        // best edge
        auto e = ed[on].getMin();
        // O(E*log^2) assert(!ed[on].cost.empty()) if every vertex appears in the
        assert (e.first < 0x3f3f3f3f3f3f3f3f3f1LL);
        int v = getPar(e.second);
        //cout << "found not cycle to " << v << " of cost " << e.first + ed[on].sum << '\n':
        assert (v != on);
        best[on] = e.first + ed[on].sum;
        ans += best[on];
        // compress edges
        ed[on].sum = -(e.first);
        st.push_back(on);
col[on] = 1;
        on = v;
    // make everything 2
    for(auto u : st)
      assert(getPar(u) == u);
col[u] = 2;
  return ans;
```

2.12 Dominator Tree

```
struct dominator_tree {
  vector<basic_string<int>> g, rg, bucket;
  vector<int> arr, par, rev, sdom, dom, dsu, label;
  dominator_tree(int n) : g(n), rg(n), bucket(n), arr(n, -1),
    par(n), rev(n), sdom(n), dom(n), dsu(n), label(n), n(n), t(0) {}
  void add_edge(int u, int v) { g[u] += v; }
  void dfs(int u) {
    arr[u] = t;
rev[t] = u;
    label[t] = sdom[t] = dsu[t] = t;
    t++;
    for (int w : q[u]) {
      if (arr[w] == -1) {
        dfs(w);
        par[arr[w]] = arr[u];
      rg[arr[w]] += arr[u];
  int find(int u, int x=0) {
    if (u == dsu[u])
     return x ? -1 : u;
    int v = find(dsu[u], x+1);
    if (v < 0)
     return u;
    if (sdom[label[dsu[u]]] < sdom[label[u]])
  label[u] = label[dsu[u]];</pre>
    return x ? v : label[u];
  vector<int> run(int root) {
    dfs(root);
    iota(dom.begin(), dom.end(), 0);
    for (int i=t-1; i>=0; i--) {
      for (int w : rg[i])
        sdom[i] = min(sdom[i], sdom[find(w)]);
        bucket[sdom[i]] += i;
      for (int w : bucket[i]) {
        int v = find(w);
        if (sdom[v] == sdom[w])
          dom[w] = sdom[w];
        else
          dom[w] = v;
      if (i > 1)
        dsu[i] = par[i];
    for (int i=1; i<t; i++) {</pre>
      if (dom[i] != sdom[i])
        dom[i] = dom[dom[i]];
    vector<int> outside_dom(n);
    iota(begin(outside_dom), end(outside_dom), 0);
    for (int i=0; i<n; i++)</pre>
     outside_dom[rev[i]] = rev[dom[i]];
    return outside_dom;
};
```

3 Dynamic Programming

3.1 Line Container

```
struct Line {
  mutable ll k, m, p;
  bool operator<(const Line& o) const { return k < o.k; }
  bool operator<(ll x) const { return p < x; }
};

struct LineContainer : multiset<Line, less<>> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b)
  static const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
```

```
return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) return x \rightarrow p = inf, 0;
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
  void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p)
      isect(x, erase(y));
  11 query(11 x) +
    assert(!empty());
    auto l = *lower_bound(x);
return l.k * x + l.m;
};
```

3.2 Li Chao Tree

```
typedef long long T;
const T INF = 2e18, EPS = 1;
struct Line {
  T m b:
  Line(T m = 0, T b = INF): m(m), b(b){}
  T apply(T x) { return x * m + b; }
struct Node {
  Node *1 = this, *r = this;
  Line line;
};
Node buffer[mx * 31];
const T MIN_VALUE = 0, MAX_VALUE = 1e9;
int CNT = 1;
Node* update(Node *root, Line line, T l = MIN_VALUE, T r = MAX_VALUE+1) {
  Node *node = &buffer[CNT++];
  *node = *root;
  T m = 1 + (r - 1) / 2;
  bool left = line.apply(1) < node->line.apply(1);
  bool mid = line.apply(m) < node->line.apply(m);
  bool right = line.apply(r) < node->line.apply(r);
  if (mid) swap(node->line, line);
  if (r - 1 <= EPS) return node;</pre>
  if (left == right) return node;
  if (mid != left) node->l = update(root->l, line, l, m);
  else node->r = update(root->r, line, m, r);
  return node;
T query (Node *root, T x, T 1 = MIN_VALUE, T r = MAX_VALUE+1) {
  if (!root) return INF;
  if (r - 1 <= EPS) return root->line.apply(x);
  T m = 1 + (r - 1) / 2;
  if (x < m) ans = query(root->1, x, 1, m);
  else ans = query(root->r, x, m, r);
  return min(ans, root->line.apply(x));
```

3.3 Knuth Optimization

```
int n, m, mid[ms][ms];
11 dp[ms][ms];
void knuth() {
  for(int i = n; i >= 0; i--) { // limites entre 0 e n
    dp[i][i+1] = 0; mid[i][i+1] = i; // caso base
    for(int j = i+2; j <= n; j++) {
        dp[i][j] = inf; // long long inf
        for(int k = mid[i][j-1]; k <= mid[i+1][j]; k++) {</pre>
```

```
if(dp[i][j] > dp[i][k] + dp[k][j]) {
    dp[i][j] = dp[i][k] + dp[k][j];
    mid[i][j] = k;
    }
    dp[i][j] += c(i, j); // custo associado ao intervalo
}
}
```

4 Math

4.1 Chinese Remainder Theorem

```
long long modinverse(long long a, long long b, long long s0 = 1, long long s1 = 0) {
  if(!b) return s0;
  else return modinverse(b, a % b, s1, s0 - s1 \star (a / b));
long long gcd(long long a, long long b) {
  if(!b) return a:
  else return gcd(b, a % b);
ll mul(ll a, ll b, ll m) {
 ll q = (long double) a * (long double) b / (long double) m;
ll r = a * b - q * m;
  return (r + 5 * m) % m;
long long safemod(long long a, long long m) {
  return (a % m + m) % m;
struct equation(
  equation(long long a, long long m) {mod = m, ans = a, valid = true;}
  equation() {valid = false;}
  equation (equation a, equation b) {
    if(!a.valid || !b.valid) {
      valid = false;
      return;
    long long g = gcd(a.mod, b.mod);
    if((a.ans - b.ans) % q != 0) {
      valid = false;
      return;
    valid = true;
    mod = a.mod * (b.mod / g);
    ans = a.ans +
     mul(a.mod, modinverse(a.mod, b.mod), mod),
      (b.ans - a.ans) / g
      , mod);
    ans = safemod(ans, mod);
  long long mod, ans;
  bool valid:
  void print()
    if(!valid)
      std::cout << "equation is not valid\n";</pre>
      std::cout << "equation is " << ans << " mod " << mod << '\n';
};
```

4.2 Diophantine Equations

```
int gcd_ext(int a, int b, int& x, int &y) {
   if (b == 0) {
      x = 1, y = 0;
      return a;
   }
   int nx, ny;
```

```
int gc = gcd_ext(b, a % b, nx, ny);
 x = ny;
y = nx - (a / b) * ny;
 return gc;
vector<int> diophantine(int D, vector<int> 1) {
 int n = l.size();
  vector<int> gc(n), ans(n);
  gc[n-1] = 1[n-1];
 for (int i = n - 2; i >= 0; i--) {
   int x, y;
   gc[i] = gcd_ext(l[i], gc[i + 1], x, y);
  if (D % gc[0] != 0) {
   return vector<int>();
  for (int i = 0; i < n; i++) {</pre>
   if (i == n - 1) {
      ans[i] = D / l[i];
      D = 1[i] * ans[i];
      continue;
    gcd_ext(l[i] / gc[i], gc[i + 1] / gc[i], x, y);
    ans[i] = (long long int) D / gc[i] * x % (gc[i + 1] / gc[i]);
    if (D < 0 \&\& ans[i] > 0) {
      ans[i] -= (gc[i + 1] / gc[i]);
    if (D > 0 \&\& ans[i] < 0) {
      ans[i] += (gc[i + 1] / gc[i]);
    D -= 1[i] * ans[i];
  return ans;
```

4.3 Discrete Logarithm

```
ll discreteLog (ll a, ll b, ll m) {
    a %= m; b %= m;
    ll n = (ll) sqrt (m + .0) + 1, an = 1;
    for (ll i = 0; i < n; i++) {
        an = (an * a) % m;
    }
    map<11, ll> vals;
    for (ll i = 1, cur = an; i <= n; i++) {
        if (!vals.count(cur)) vals[cur] = i;
        cur = (cur * an) % m;
    }
    ll ans = le18; //inf
    for (ll i = 0, cur = b; i <= n; i++) {
        if (vals.count(cur)) {
            ans = min(ans, vals[cur] * n - i);
        }
        cur = (cur * a) % m;
    return ans;
}</pre>
```

4.4 Discrete Root

```
//x^k = a % mod

11 discreteRoot(11 k, 11 a, 11 mod) {
    11 g = primitiveRoot(mod);
    11 y = discreteLog(fexp(g, k, mod), a, mod);
    if (y == -1) {
        return y;
    }
    return fexp(g, y, mod);
}
```

4.5 Division Trick

```
for(int l = 1, r; l <= n; l = r + 1) {
    r = n / (n / 1);
    // n / i has the same value for l <= i <= r</pre>
```

4.6 Modular Sum

```
//calcula (sum(0 <= i <= n) P(i)) % mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
  ll calc(ll a, ll b, ll n, ll mod) {
    assert (a&&b);
    if(a >= b){
      11 ret = ((n*(n+1)/2)*mod)*(a/b);
      if(a%b) ret = (ret + calc(a%b,b,n,mod))%mod;
      else ret = (ret+n+1)%mod;
      return ret;
    return ((n+1)*(((n*a)/b+1)*mod) - calc(b,a,(n*a)/b,mod) + mod + n/b + 1)*mod;
  ll solve(ll a, ll n, ll m, ll mod) {
    a = (a%m + m)%m;
    if(!a) return 0;
    11 ret = (n*(n+1)/2)%mod;
    ret = (ret*a)%mod;
    11 g = __gcd(a,m);
ret -= m*(calc(a/g,m/g,n,mod)-n-1);
    return (ret%mod + mod)%mod;
//P(i) = a + r*i mod m
ll solve(ll a, ll r, ll n, ll m, ll mod){
   a = (a%m + m)%m;
    r = (r%m + m)%m;
    if(!r) return (a*(n+1))%mod;
    if(!a) return solve(r, n, m, mod);
    ll g, x, y;
    g = gcdExtended(r, m, x, y);
x = (x%m + m)%m;
    11 d = a - (a/g) *g;
    a -= d:
    x = (x*(a/q))%m;
    return (solve(r, n+x, m, mod) - solve(r, x-1, m, mod) + mod + d*(n+1)) % mod;
};
```

4.7 Primitive Root

```
//is n primitive root of p ?
bool test(long long x, long long p) {
    long long m = p - 1;
    for(int i = 2; i * i <= m; ++i) if(!(m % i)) {
        if(fexp(x, i, p) == 1) return false;
        if(fexp(x, m / i, p) == 1) return false;
    }
    return true;
}
//find the smallest primitive root for p
int search(int p) {
    for(int i = 2; i < p; i++) if(test(i, p)) return i;
    return -1;
}</pre>
```

4.8 Extended Euclides

```
// euclides estendido: acha u e v da equacao:
// u * x + v * y = gcd(x, y);
// u eh inverso modular de x no modulo y
// v eh inverso modular de y no modulo x

pair<ll, ll> euclides(ll a, ll b) {
    ll u = 0, oldu = 1, v = 1, oldv = 0;
    while(b) {
        ll q = a / b;
        oldv = oldv - v * q;
        oldu = oldu - u * q;
    }
}
```

```
a = a - b * q;
swap(a, b);
swap(u, oldu);
swap(v, oldv);
}
return make_pair(oldu, oldv);
```

4.9 Matrix

4.10 FFT - Fast Fourier Transform

```
typedef double ld;
const ld PI = acos(-1);
struct Complex {
  ld real, imag;
  Complex conj() { return Complex(real, -imag); }
  Complex(1d = 0, 1d b = 0) : real(a), imag(b) {}
  Complex operator + (const Complex &o) const { return Complex(real + o.real, imag + o
       imag): }
  Complex operator - (const Complex &o) const { return Complex (real - o.real, imag - o
       .imag); }
  Complex operator * (const Complex &o) const { return Complex(real * o.real - imag *
      o.imag, real * o.imag + imag * o.real); }
  Complex operator / (ld o) const { return Complex(real / o, imag / o); }
  void operator *= (Complex o) { *this = *this * o; }
  void operator /= (ld o) { real /= o, imag /= o; }
typedef std::vector<Complex> CVector;
const int ms = 1 << 22;</pre>
int bits[ms];
Complex root[ms];
void initFFT() {
  root[1] = Complex(1);
  for(int len = 2; len < ms; len += len) {</pre>
    Complex z(cos(PI / len), sin(PI / len));
    for(int i = len / 2; i < len; i++) {</pre>
     root[2 * i] = root[i];
      root[2 * i + 1] = root[i] * z;
  }
void pre(int n) {
  int LOG = 0;
  while (1 << (LOG + 1) < n) {
  for(int i = 1; i < n; i++) {</pre>
    bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << LOG);
CVector fft(CVector a, bool inv = false) {
  int n = a.size();
  pre(n);
  if(inv) {
    std::reverse(a.begin() + 1, a.end());
  for (int i = 0; i < n; i++) {
   int to = bits[i];
    if(to > i) {
```

```
std::swap(a[to], a[i]);
  for (int len = 1; len < n; len \star= 2) {
    for(int i = 0; i < n; i += 2 * len) {
      for(int j = 0; j < len; j++) {</pre>
        Complex u = a[i + j], v = a[i + j + len] * root[len + j];
        a[i + j] = u + v;
        a[i + j + len] = u - v;
  if(inv) {
    for (int i = 0; i < n; i++)
      a[i] /= n;
  return a;
void fft2in1(CVector &a, CVector &b) {
  int n = (int) a.size();
  for (int i = 0; i < n; i++) {
    a[i] = Complex(a[i].real, b[i].real);
  auto c = fft(a);
  for (int i = 0; i < n; i++) {
   a[i] = (c[i] + c[(n-i) % n].conj()) * Complex(0.5, 0);
    b[i] = (c[i] - c[(n-i) % n].conj()) * Complex(0, -0.5);
void ifft2in1(CVector &a, CVector &b) {
  int n = (int) a.size();
  for(int i = 0; i < n; i++) a[i] = a[i] + b[i] * Complex(0, 1);</pre>
  a = fft(a, true);
  for(int i = 0; i < n; i++) {</pre>
   b[i] = Complex(a[i].imag, 0);
    a[i] = Complex(a[i].real, 0);
std::vector<long long> mod_mul(const std::vector<long long> &a, const std::vector<long</pre>
      long> &b, long long cut = 1 << 15) {
  int n = (int) a.size();
  CVector C[4];
  for(int i = 0; i < 4; i++) C[i].resize(n);</pre>
  for(int i = 0; i < n; i++) {
    C[0][i] = a[i] % cut;
    C[1][i] = a[i] / cut;
    C[2][i] = b[i] % cut;
    C[3][i] = b[i] / cut;
  fft2in1(C[0], C[1]);
  fft2in1(C[2], C[3]);
  for(int i = 0; i < n; i++) {</pre>
    // 00, 01, 10,
Complex cur[4];
    for(int j = 0; j < 4; j++) cur[j] = C[j/2+2][i] * C[j % 2][i];
for(int j = 0; j < 4; j++) C[j][i] = cur[j];</pre>
  ifft2in1(C[0], C[1]);
  ifft2in1(C[2], C[3]);
  std::vector<long long> ans(n, 0);
  for (int i = 0; i < n; i++) {
    // if there are negative values, care with rounding
    ans[i] += (long long) (C[0][i].real + 0.5);
    ans[i] += (long long) (C[1][i].real + C[2][i].real + 0.5) * cut;
    ans[i] += (long long) (C[3][i].real + 0.5) * cut * cut;
  return ans:
std::vector<int> mul(const std::vector<int> &a, const std::vector<int> &b) {
  int n = 1:
  while (n-1 < (int) \ a.size() + (int) \ b.size() - 2) \ n += n;
  CVector poly(n);
  for(int i = 0; i < n; i++) {
    if(i < (int) a.size()) {</pre>
      poly[i].real = a[i];
    if(i < (int) b.size()) {
      poly[i].imag = b[i];
  poly = fft(poly);
  for(int i = 0; i < n; i++) {</pre>
```

```
poly[i] *= poly[i];
}
poly = fft(poly, true);
std::vector<int> c(n, 0);
for(int i = 0; i < n; i++) {
    c[i] = (int) (poly[i].imag / 2 + 0.5);
}
while (c.size() > 0 && c.back() == 0) c.pop_back();
return c;
```

4.11 NTT - Number Theoretic Transform

```
const int MOD = 998244353;
const int me = 15;
const int ms = 1 << me;</pre>
#define add(x, y) x+y>=MOD?x+y-MOD:x+y
const int gen = 3; // use search() from PrimitiveRoot.cpp if MOD isn't 998244353
int bits[ms], root[ms];
void initFFT() {
  root[1] = 1;
  for(int len = 2; len < ms; len += len) {</pre>
    int z = fexp(gen, (MOD - 1) / len / 2);
    for(int i = len / 2; i < len; i++) {</pre>
     root[2 * i] = root[i];
      root[2 * i + 1] = (long long) root[i] * z % MOD;
void pre(int n) {
  int I_iOG = 0:
  while (1 << (LOG + 1) < n) {
  for(int i = 1; i < n; i++) {</pre>
    bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << LOG);
vector<int> fft (vector<int> a, bool inv = false) {
  int n = (int) a.size();
  pre(n);
  if(inv) {
    reverse(a.begin() + 1, a.end());
  for(int i = 0; i < n; i++) {</pre>
   int to = bits[i];
    if(i < to)
      swap(a[i], a[to]);
  for(int len = 1; len < n; len *= 2) {</pre>
    for (int i = 0; i < n; i += len * 2) {
      for(int j = 0; j < len; j++) {
        int u = a[i + j], v = (11) a[i + j + len] * root[len + j] % mod;
        a[i + j] = add(u, v);
        a[i + j + len] = add(u, mod - v);
  if(inv) {
    int rev = fexp(n, mod-2, mod);
    for(int i = 0; i < n; i++)
      a[i] = (ll) a[i] * rev % mod;
  return a;
std::vector<int> operator *(std::vector<int> a, std::vector<int> b) {
  while(!a.empty() && a.back() == 0) a.pop_back();
  while(!b.empty() && b.back() == 0) b.pop back();
  if(a.empty() || b.empty()) return std::vector<int>(0, 0);
  int n = 1:
  while (n-1 < (int) a.size() + (int) b.size() - 2) n += n;
  a.resize(n. 0);
  b.resize(n, 0);
a = fft(a, false);
  b = fft(b, false);
  for (int i = 0; i < n; i++) {
```

```
 a[i] = (int) \ ((long \ long) \ a[i] \ * \ b[i] \ % \ MOD);  return fft(a, true);
```

4.12 Fast Walsh Hadamard Transform

```
vector<11> FWHT(char oper, vector<11> a, const bool inv = false)
  int n = (int) a.size();
  for(int len = 1; len < n; len += len) +</pre>
    for (int i = 0; i < n; i += 2 * len)
      for(int j = 0; j < len; j++) {
        auto u = a[i + j] % mod, v = a[i + j + len] % mod;
        if(oper == '^') {
          a[i + j] = (u + v) % mod;
          a[i + j + len] = (u - v + mod) % mod;
        if(oper == '|') {
          if(!inv) {
            a[i + j + len] = (u + v) % mod;
          } else {
            a[i + j + len] = (v - u + mod) % mod;
        if(oper == '&') {
          if(!inv) {
            a[i + j] = (u + v) % mod;
          } else {
            a[i + j] = (u - v + mod) % mod;
  if(oper == '^' && inv) {
    11 \text{ rev} = \text{fexp(n, mod - 2);}
    for(int i = 0; i < n; i++) {
      a[i] = a[i] * rev % mod;
  return a;
vector<ll> multiply(char oper, vector<ll> a, vector<ll> b) {
  int n = 1:
  while (n < (int) max(a.size(), b.size())) {</pre>
    n <<= 1;
  vector<ll> ans(n);
  while (a.size() < ans.size()) a.push_back(0);</pre>
  while (b.size() < ans.size()) b.push_back(0);</pre>
  a = FWHT(oper, a);
b = FWHT(oper, b);
  for (int i = 0; i < n; i++) {
    ans[i] = a[i] * b[i] % mod;
  ans = FWHT (oper, ans, true);
  return ans:
const int mxlog = 17;
vector<ll> subset_multiply(vector<ll> a, vector<ll> b) {
  int n = 1;
  while (n < (int) max(a.size(), b.size())) {</pre>
    n <<= 1;
  vector<11> ans(n);
  while (a.size() < ans.size()) a.push_back(0);</pre>
  while (b.size() < ans.size()) b.push_back(0);</pre>
  vector<vector<ll>> A(mxlog + 1, vector<ll>(a.size())), B(mxlog + 1, vector<ll>(b.
       size()));
  for (int i = 0; i < n; i++) {
    A[__builtin_popcount(i)][i] = a[i];
    B[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= mxlog; i++) {</pre>
    A[i] = FWHT('|', A[i]);
    B[i] = FWHT('|', B[i]);
```

```
for (int i = 0; i <= mxlog; i++) {
    vector<ll> C(n);
    for (int x = 0; x <= i; x++) {
        int y = i - x;
        for (int j = 0; j < n; j++) {
            C[j] = (C[j] + A[x][j] * B[y][j] % mod) % mod;
        }
    }
    C = FWHT('|', C, true);
    for (int j = 0; j < n; j++) {
        if (_builtin_popcount(j) == i) {
            ans[j] = (ans[j] + C[j]) % mod;
    }
    }
}
return ans;</pre>
```

4.13 Miller and Rho

```
//miller_rabin
typedef unsigned long long ull;
typedef long double ld;
ull fmul(ull a, ull b, ull m) { ull q = (ld) \ a * (ld) \ b / (ld) \ m; ull \ r = a * b - q * m; }
  return (r + m) % m;
bool miller(ull p, ull a) {
  ull s = p - 1;
  while(s % 2 == 0) s >>= 1;
  while (a >= p) a >>= 1;
  ull mod = fexp(a, s, p);
  while(s != p - 1 && mod != 1 && mod != p - 1) {
    mod = fmul(mod, mod, p);
    s <<= 1;
  if(mod != p - 1 && s % 2 == 0)return false;
  else return true;
bool prime (ull p) {
  if(p <= 3)
    return true;
  if(p % 2 == 0)
    return false;
  return miller(p, 2) && miller(p, 3)
    && miller(p, 5) && miller(p, 7)
    && miller(p, 11) && miller(p, 13)
    && miller(p, 17) && miller(p, 19)
    && miller(p, 23) && miller(p, 29)
    && miller(p, 31) && miller(p, 37);
//pollard_rho
ull func(ull x, ull c, ull n) {
  return (fmul(x, x, n) + c) % n;
ull gcd(ull a, ull b) {
  if(!b) return a;
  else return gcd(b, a % b);
ull rho(ull n) {
  if(n % 2 == 0) return 2;
  if(prime(n)) return n;
  while(1) {
     ull c;
    do {
      c = rand() % n;
     while(c == 0 || (c + 2) % n == 0);
ull x = 2, y = 2, d = 1;
ull pot = 1, lam = 1;
    do {
       if(pot == lam) {
         x = y;
pot <<= 1;
         lam = 0;
       \dot{y} = func(y, c, n);
       lam++;
       d = gcd(x >= y ? x - y : y - x, n);
     } while(d == 1);
    if(d != n) return d;
```

```
vector<ull> factors(ull n) {
 vector<ull> ans, rest, times;
  if(n == 1) return ans;
  rest push_back(n);
  times.push_back(1);
  while(!rest.empty()) {
    ull x = rho(rest.back());
    if(x == rest.back()) {
      int freq = 0;
      for(int i = 0; i < rest.size(); i++) {</pre>
        int cur_freq = 0;
        while(rest[i] % x == 0) {
          rest[i] /= x;
          cur_freq++;
        freq += cur_freq * times[i];
        if(rest[i] == 1) {
          swap(rest[i], rest.back());
          swap(times[i], times.back());
rest.pop_back();
          times.pop_back();
          i--;
      while(freq--) {
  ans.push_back(x);
      continue;
    ull e = 0;
    while (rest.back() % x == 0) {
      rest.back() /= x;
      e++;
    e *= times.back();
    if(rest.back() == 1)
      rest.pop_back();
      times.pop_back();
    rest push_back(x);
   times.push_back(e);
  return ans;
```

4.14 Determinant using Mod

```
// by zchao1995
// Determinante com coordenadas inteiras usando Mod
11 mat[ms][ms];
11 det (int n) {
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < n; j++) {
      mat[i][j] %= mod;
  11 res = 1;
  for (int i = 0; i < n; i++) {</pre>
    if (!mat[i][i]) {
      bool flag = false;
for (int j = i + 1; j < n; j++) {
        if (mat[j][i]) {
          flag = true;
          for (int k = i; k < n; k++) {
            swap (mat[i][k], mat[j][k]);
           res = -res;
          break;
      if (!flag) {
        return 0:
    for (int j = i + 1; j < n; j++) {
      while (mat[j][i]) {
```

```
11 t = mat[i][i] / mat[j][i];
  for (int k = i; k < n; k++) {
    mat[i][k] = (mat[i][k] - t * mat[j][k]) % mod;
    swap (mat[i][k], mat[j][k]);
  }
  res = -res;
  }
  res = (res * mat[i][i]) % mod;
}
return (res + mod) % mod;</pre>
```

4.15 Gauss

```
const double eps = 1e-9;
int gauss (vector<vector<double>> a, vector<double> & ans) 
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        int sel = row;
        for (int i=row; i<n; ++i) {</pre>
             if (abs (a[i][col]) > abs (a[sel][col]))
                 sel = i;
        if (abs (a[sel][col]) < eps) continue;</pre>
        for (int i=col; i<=m; ++i)</pre>
             swap (a[sel][i], a[row][i]);
        where[col] = row;
        for (int i=0; i<n; ++i) {</pre>
             if (i != row) {
                 double c = a[i][col] / a[row][col];
for (int j=col; j<=m; ++j)</pre>
                     a[i][j] = a[row][j] * c;
         ++row;
    ans.assign (m, 0);
    for (int i=0; i<m; ++i) {</pre>
        if (where[i] != -1)
             ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {</pre>
        double sum = 0;
        for (int j=0; j<m; ++j)
            sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > eps)
             return 0:
    for (int i=0; i<m; ++i) {</pre>
        if (where[i] == -1)
             return INF;
    return 1;
// mod 2 (xor);
int gauss (vector <bitset<ms>> a, int m, bitset<ms> &ans) {
    int n = (int) a.size();
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        for (int i=row; i<n; ++i){</pre>
            if (a[i][col]) {
                 swap (a[i], a[row]);
                 break;
        if (!a[row][col]) continue;
        where[col] = row;
        for (int i=0; i<n; ++i) {</pre>
            if (i != row && a[i][col])
                 a[i] ^= a[row];
        ++row:
    for (int i = 0; i < m; ++i)
```

```
if(where[i] != -1) {
    ans[i] = a[where[i]][m];
}

for(int i = 0; i < n; ++i) {
    int sum = 0;
    for(int j = 0; j < m; ++j) {
        sum ^= (ans[j] & a[i][j]);
    }
    if(sum != a[i][m]) {
        return 0;
    }
}

for(int i = 0; i < m; ++i)
    if(where[i] == -1)
        return 1e9;
return 1;</pre>
```

4.16 Lagrange Interpolation

```
class LagrangePoly {
public:
  LagrangePoly(vector<long long> _a) {
    //f(i) = \underline{a[i]}
    //interpola o vetor em um polinomio de grau y.size() - 1
    y = \underline{a};
    den.resize(y.size());
    int n = (int) y.size();
    for (int i = 0; i < n; i++) {
      y[i] = (y[i] % MOD + MOD) % MOD;
      den[i] = ifat[n - i - 1] * ifat[i] % MOD;
      if((n - i - 1) % 2 == 1) {
        den[i] = (MOD - den[i]) % MOD;
  long long getVal(long long x) {
    int n = (int) y.size();
x %= MOD;
    if(x < n) {
      //return y[(int) x];
    vector<long long> 1, r;
    1.resize(n);
    for(int i = 1; i < n; i++) {</pre>
      l[i] = l[i - 1] * (x - (i - 1) + MOD) % MOD;
    r.resize(n);
    r[n - 1] = 1;
    for (int i = n - 2; i >= 0; i--) {
      r[i] = r[i + 1] * (x - (i + 1) + MOD) % MOD;
    long long ans = 0;
    for (int i = 0; i < n; i++) {
      long long coef = 1[i] * r[i] % MOD;
      ans = (ans + coef * y[i] % MOD * den[i]) % MOD;
    return ans;
vector<long long> y, den;
int main(){
  fat[0] = ifat[0] = 1;
  for(int i = 1; i < ms; i++) {</pre>
    fat[i] = fat[i - 1] * i % MOD;
    ifat[i] = fexp(fat[i], MOD - 2);
  // Codeforces 622F
  int x, k;
  cin >> x >> k:
  vector<long long> a;
  a.push_back(0);
  for (long long i = 1; i \le k + 1; i++) {
```

```
a.push_back((a.back() + fexp(i, k)) % MOD);
}
LagrangePoly f(a);
cout << f.getVal(x) << '\n';
}</pre>
```

4.17 Lagrange extracting polynomial

```
// O(n^2), receve v {x, y} e retorna o polinomio em fracao
vector<pair<int, int>> interpolate(vector<ii> v) {
  int n = v.size();
 vector<int> prod(n+1);
  prod[0] = 1;
  for (auto p : v) {
    for (int i = n; i > 0; i--) {
      prod[i] = prod[i-1] - p.first * prod[i];
   prod[0] = -p.first * prod[0];
  vector<pair<int, int>> ans(n+1);
  for(int i = 0; i <= n; i++) ans[i].second = 1;</pre>
  for (int i = 0; i < n; i++) {
   vector<int> pol(n+1); // (x - v[i].first)
    for (int j = n; j > 0; j--) {
      pol[j-1] = prod[j] + pol[j] * v[i].first;
    for (int j = 0; j < n; j++) {
      pol[j] *= v[i].second;
    int k = 1;
    for (int j = 0; j < n; j++) {
     if(i==j) continue;
      k *= v[i].first - v[j].first;
    if(k < 0) {
        = -k;
      for(auto &p : pol) p = -p;
    for (int i = 0; i < n; i++) {
      ans[i] = {ans[i].first*k + pol[i]*ans[i].second, k*ans[i].second};
      if(ans[i].first == 0) ans[i].second = 1;
      else (
        int gc = __gcd(abs(ans[i].first), ans[i].second);
        ans[i].first /= gc;
        ans[i].second /= qc;
  return ans:
```

4.18 Count integer points inside triangle

```
//gcd(p, q) == 1
ll get(ll p, ll q, ll n, bool floor = true) {
    if (n == 0) {
        return 0;
    }
    if (p % q == 0) {
            return n * (n + 1) / 2 * (p / q);
    }
    if (p > q) {
        return n * (n + 1) / 2 * (p / q) + get(p % q, q, n, floor);
    }
    ll new_n = p * n / q;
    ll ans = (n + 1) * new_n - get(q, p, new_n, false);
    if (!floor) {
        ans += n - n / q;
    }
    return ans;
}
```

4.19 Prime Counting

```
const int ms = 5001000, lim_n = 3e5, lim_p = 1e2;
std::vector<int> primes;
int id[ms];
int memo[lim_n][lim_p];
void pre() {
  std::vector<bool> isPrime(ms, true);
  for(int i = 2; i < ms; i++) {</pre>
    id[i] = (int) primes.size();
    if(!isPrime[i]) continue;
    id[i]++;
    primes.push_back(i);
    for(int j = i+i; j < ms; j += i) isPrime[j] = false;</pre>
  for(int i = 1; i < lim_n; i++) {</pre>
    memo[i][0] = i;
    for(int j = 1; j < \lim_{p} j++) memo[i][j] = memo[i][j-1] - memo[i/primes[j-1]][j
int cbc(long long n) {
  int ans = std::max(0, (int) pow((double) n, 1.0 / 3) - 2);
  while((11) ans * ans * ans < n) ans++;</pre>
  return ans;
long long dp (long long n, int i) {
  if(n == 0) return 0; if(i == 0) return n;
  if(primes[i-1] >= n) return 1;
  if((11) primes[i-1] * primes[i-1] > n && n < ms) return id[n] - (i-1);</pre>
  else if(n < lim_n && i < lim_p) return memo[n][i];</pre>
  else return dp(n, i-1) - dp(n / primes[i-1], i-1);
long long primeFunction(long long n) {
  if(n < ms) return id[(int)n];</pre>
  int i = id[cbc(n)];
  long long ans = dp(n, i) + i - 1;
  while((long long) primes[i] * primes[i] <= n) {</pre>
   ans -= primeFunction(n / primes[i]) - i;
    <u>i</u>++;
  return ans;
```

4.20 Berlekamp Massey

```
vector<int> berlekampMassey(const vector<int> &s) {
    int n = (int) s.size(), l = 0, m = 1;
    vector<int> b(n), c(n);
    int 1d = b[0] = c[0] = 1;
    for (int i=0; i<n; i++, m++) {</pre>
        int d = s[i];
        for (int j=1; j<=1; j++)</pre>
            d = (d + c[j] * s[i-j]) % mod;
        if (d == 0)
            continue;
        vector<int> temp = c;
        int coef = d * fexp(ld, mod-2) % mod;
        for (int j=m; j<n; j++)</pre>
             c[j] = ((c[j] - coef * b[j-m]) % mod + mod) % mod;
        if (2 * 1 <= i) {
             1 = i + 1 - 1;
             b = temp;
             ld = d;
            \mathbf{m} = 0;
    c.resize(1 + 1);
    c.erase(c.begin());
    for (int &x : c)
    return c;
void mull(vector<int> &p, vector<int> &q, vector<int> &h, int m) {
        vector<int> t_(m+m);
        for(int i=0;i<m;++i) if(p[i])</pre>
                 for (int j=0; j<m; ++j)</pre>
                          t_{[i+j]} = (t_{[i+j]} + p[i] * q[j]) * mod;
```

```
for(int i=m+m-1;i>=m;--i) if(t_[i])
                 //miuns t_{[i]}x^{i-m}(x^m-\sum_{j=0}^{m-1}x^{m-j-1}h_{j})
                 for(int j=m-1; ~j; --j)
                         t_{[i-j-1]} = (t_{[i-j-1]} + t_{[i]} * h[j]) %mod;
        for (int i=0; i < m; ++i) p[i] = t_[i];</pre>
// a = caso base, h = recorrencia, m = tamanho da recorrencia
inline int calc(vector<int> &a, vector<int> &h, int K, int m) {
        vector<int> s(m), t(m);
        //init
        s[0]=1; if(m!=1) t[1]=1; else t[0]=h[0];
        //binary-exponentiation
        while(K) {
                if(K&1) mull(s,t,h,m);
                 mull(t,t,h,m); K>>=1;
        int su=0;
        for (int i=0;i<m;++i) su=(su+s[i]*a[i])%mod;</pre>
        return (su%mod+mod)%mod;
```

4.21 Polynomial exp

```
// by ijmg
vector<int> power(vector<int> &a, int k, int limit = -1) {
    while(a.back() == 0) a.pop_back();
    if(a.size() == 0 || limit == 0) return {};
    if(limit == -1) {
        limit = (a.size() - 1) * k;
    }
    vector<int> ans(limit + 1, 0);
    ans[0] = fexp(a[0], k);
    for(int i = 1; i <= limit; ++i) {
        for(int j = 1; j <= min(i, (int) a.size() - 1); ++j) {
            ans[i] += a[j] * ans[i - j] * (k * j - (i - j));
        }
        ans[i] /= i * a[0];
    }
    return ans;
}</pre>
```

5 Geometry

5.1 Geometry

```
const double inf = le100, eps = le-9;
const double PI = acos(-1.0L);
int cmp (double a, double b = 0) {
   if (abs(a-b) < eps) return 0;
    return (a < b) ? -1 : +1;
}
struct PT {
   double x, y;
   PT (double x = 0, double y = 0) : x(x), y(y) {}
   PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
   PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
   PT operator * (double c) const { return PT(x/c, y/c); }
   PT operator / (double c) const { return PT(x/c, y/c); }
   bool operator <(const PT &p) const {
        if(cmp(x, p.x) != 0) return x < p.x;
        return cmp(y, p.y) < 0;
}
   bool operator ==(const PT &p) const {return !cmp(x, p.x) && !cmp(y, p.y); }
   bool operator != (const PT &p) const {return !(p == *this); }
};
ostream &operator<<((ostream &os, const PT &p) {
        return os << "(" << p.x << "," << p.y << ")";
}
double dot (PT p, PT q) { return p.x * q.x + p.y*q.y; }
double cross (PT p, PT q) { return p.x * q.x + p.y*q.x; }
double dist2 (PT p, PT q = PT(0, 0)) { return dot(p-q, p-q); }
double dist (PT p, PT q) = return hypot(p.x-q.x, p.y-q.y); }
double norm (PT p) { return hypot(p.x, p.y); }</pre>
```

```
PT normalize (PT p) { return p/hypot(p.x, p.y); }
double angle (PT p, PT q) { return atan2(cross(p, q), dot(p, q)); }
double angle (PT p) { return atan2(p.y, p.x); }
double polarAngle (PT p) {
  double a = atan2(p.y,p.x);
return a < 0 ? a + 2*PI : a;</pre>
PT rotateCCW90 (PT p) { return PT(-p.y, p.x); }
PT rotateCW90 (PT p) { return PT(p.y, -p.x); }
PT rotateCCW (PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
typedef pair<PT, int> Line;
PT getDir (PT a, PT b) {
  if (a.x == b.x) return PT(0, 1);
  if (a.y == b.y) return PT(1, 0);
  int dx = b.x-a.x;
  int dy = b.y-a.y;
  int g = __gcd(abs(dx), abs(dy));
  if (dx < 0) g = -g;
  return PT(dx/g, dy/g);
Line getLine (PT a, PT b)
  PT dir = getDir(a, b);
  return {dir, cross(dir, a)};
PT projPtLine (PT a, PT b, PT c) { // ponto c na linha a - b, a.b = |a| cost * |b| return a + (b-a) * dot(b-a, c-a)/dot(b-a, b-a);
PT reflectPointLine (PT a, PT b, PT c) {
  PT p = projPtLine(a, b, c);
return p*2 - c;
PT projPtSeg (PT a, PT b, PT c) { // c no segmento a - b double r = dot(b-a, b-a);
  if (cmp(r) == 0) return a;
  r = dot(b-a, c-a)/r;
  if (cmp(r, 0) < 0) return a;
  if (cmp(r, 1) > 0) return b;
return a + (b - a) * r;
double distancePointSegment (PT a, PT b, PT c) { // ponto c e o segmento a - b
  return dist(c, projPtSeg(a, b, c));
bool ptInSegment (PT a, PT b, PT c) { // ponto c esta em um segmento a - b
  if (a == b) return a == c;
  return cmp(cross(a, b)) == 0 \&\& cmp(dot(a, b)) <= 0;
bool parallel (PT a, PT b, PT c, PT d) {
  return cmp(cross(b - a, c - d)) == 0;
bool collinear (PT a, PT b, PT c, PT d) {
  return parallel(a, b, c, d) && cmp(cross(a - b, a - c)) == 0 && cmp(cross(c - d, c -
         a)) == 0;
 // Calcula distancia entre o ponto (x, y, z) e o plano ax + by + cz = d
double distPtPlane(double x, double y, double z, double a, double b, double c, double
    return abs(a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c);
bool segInter (PT a, PT b, PT c, PT d) {
  if (collinear(a, b, c, d)) {
    if (a == c || a == d || b == c || b == d) return true;
    if (cmp(dot(c - a, c - b)) > 0 && cmp(dot(d - a, d - b)) > 0 && cmp(dot(c - b, d - b))) > 0 && cmp(dot(c - b, d - b)))
            b)) > 0) return false;
  if (cmp(cross(d - a, b - a) * cross(c - a, b - a)) > 0) return false;
  if (cmp(cross(a - c, d - c) * cross(b - c, d - c)) > 0) return false;
// Calcula a intersecao entre as retas a - b e c - d assumindo que uma unica
      intersecao existe
// Para intersecao de segmentos, cheque primeiro se os segmentos se intersectam e que
     nao sao paralelos
PT lineLine (PT a, PT b, PT c, PT d) {
b = b - a; d = c - d; c = c - a;
// assert(cmp(cross(b, d)) != 0);
return a + b * cross(c, d) / cross(b, d);
PT circleCenter (PT a, PT b, PT c) {
  b = (a + b) / 2; // bissector
```

c = (a + c) / 2; // bissector

```
return lineLine(b, b + rotateCW90(a - b), c, c + rotateCW90(a - c));
vector<PT> circle2PtsRad (PT p1, PT p2, double r) {
  vector<PT> ret:
  double d2 = dist2(p1, p2);
double det = r * r / d2 - 0.25;
  if (det < 0.0) return ret;</pre>
  double h = sgrt(det);
  for (int i = 0; i < 2; i++) {
    double x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;

double y = (p1.y + p2.y) * 0.5 + (p2.x - p1.x) * h;

ret.push_back(PT(x, y));
    swap(p1, p2);
  return ret;
bool circleLineIntersection(PT a, PT b, PT c, double r) {
    return cmp(dist(c, projPtLine(a, b, c)), r) <= 0;</pre>
vector<PT> circleLine (PT a, PT b, PT c, double r) {
  vector<PT> ret;
  PT p = projPtLine(a, b, c), p1;
  double h = norm(c-p);
 if (cmp(h,r) == 0) {
  ret.push_back(p);
  else if (cmp(h,r) < 0)
    double k = sqrt(r*r - h*h);
    p1 = p + (b-a)/(norm(b-a))*k;
    ret.push_back(p1);
    p1 = p - (b-a)/(norm(b-a))*k;
    ret.push_back(p1);
  return ret:
bool ptInsideTriangle(PT p, PT a, PT b, PT c) {
  if(cross(b-a, c-b) < 0) swap(a, b);
  if(ptInSegment(a,b,p)) return 1;
  if(ptInSegment(b,c,p)) return 1;
  if(ptInSegment(c,a,p)) return 1;
  bool x = cross(b-a, p-b) < 0;
bool y = cross(c-b, p-c) < 0;
bool z = cross(a-c, p-a) < 0;
  return x == y && y == z;
bool pointInConvexPolygon(const vector<PT> &p, PT q) {
  if (p.size() == 1) return p.front() == q;
  int 1 = 1, r = p.size()-1;
  while (abs(r-1) > 1) {
    int m = (r+1)/2;
    if(cross(p[m]-p[0], q-p[0]) < 0) r = m;
  return ptInsideTriangle(q, p[0], p[1], p[r]);
// Determina se o ponto esta num poligono possivelmente nao-convexo
// Retorna 1 para pontos estritamente dentro, 0 para pontos estritamente fora do
     poligno
// e 0 ou 1 para os pontos restantes
bool pointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0:
  for(int i = 0; i < p.size(); i++){</pre>
    int j = (i + 1) % p.size();
    if((p[i].y \le q.y \&\& q.y < p[j].y || p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c;
// area / semiperimeter
double rIncircle (PT a, PT b, PT c) {
  double ab = norm(a-b), bc = norm(b-c), ca = norm(c-a);
  return abs(cross(b-a, c-a)/(ab+bc+ca));
vector<PT> circleCircle (PT a, double r, PT b, double R) {
  vector<PT> ret;
  double d = norm(a-b);
  if (d > r + R \mid \mid d + min(r, R) < max(r, R)) return ret;
  double x = (d*d - R*R + r*r) / (2*d); // x = r*cos(R opposite angle)
  double y = sqrt(r*r - x*x);
  PT v = (b - a)/d;
  ret.push_back(a + v*x + rotateCCW90(v)*y);
  if (cmp(y) > 0)
    ret.push back(a + v*x - rotateCCW90(v)*v);
  return ret;
```

```
double circularSegArea (double r, double R, double d) {
  double ang = 2 \times a\cos((d*d - R*R + r*r) / (2*d*r)); // cos(R opposite angle) = x/r
  double tri = sin(ang) * r * r;
  double sector = ang * r * r;
  return (sector - tri) / 2;
double computeSignedArea (const vector<PT> &p) {
  double area = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area/2.0;
double computeArea(const vector<PT> &p) { return abs(computeSignedArea(p)); }
PT computeCentroid(const vector<PT> &p) {
  double scale = 6.0 * computeSignedArea(p);
  for(int i = 0; i < p.size(); i++){</pre>
    int j = (i + 1) % p.size();
    c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
  return c / scale;
// Testa se o poligno listada em ordem CW ou CCW eh simples (nenhuma linha se
     intersecta)
bool isSimple(const vector<PT> &p) {
  for(int i = 0; i < p.size(); i++) {</pre>
    for(int k = i + 1; k < p.size(); k++) {</pre>
     int j = (i + 1) % p.size();
      int 1 = (k + 1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (segInter(p[i], p[j], p[k], p[l]))
        return false;
  return true;
vector< pair<PT, PT> > getTangentSegs (PT c1, double r1, PT c2, double r2) {
  if (r1 < r2) swap(c1, c2), swap(r1, r2);</pre>
  vector<pair<PT, PT> > ans;
  double d = dist(c1, c2);
  if (cmp(d) <= 0) return ans;</pre>
  double dr = abs(r1 - r2), sr = r1 + r2;
  if (cmp(dr, d) >= 0) return ans;
  double u = acos(dr / d);
  PT dc1 = normalize(c2 - c1) *r1;
  PT dc2 = normalize(c2 - c1) \star r2;
  ans.push_back(make_pair(c1 + rotateCCW(dc1, +u), c2 + rotateCCW(dc2, +u)));
  ans.push_back(make_pair(c1 + rotateCCW(dc1, -u), c2 + rotateCCW(dc2, -u)));
  if (cmp(sr, d) >= 0) return ans;
  double v = acos(sr / d);
  dc2 = normalize(c1 - c2)*r2;
  ans.push_back({c1 + rotateCCW(dc1, +v), c2 + rotateCCW(dc2, +v)});
  ans.push_back({c1 + rotateCCW(dc1, -v), c2 + rotateCCW(dc2, -v)});
  return ans:
```

5.2 Convex Hull

```
vector<PT> convexHull(vector<PT> p, bool needs = 1) {
  if(needs) sort(p.begin(), p.end());
  p.erase(unique(p.begin(), p.end()), p.end());
  int n = p.size(), k = 0;
  if(n <= 1) return p;</pre>
  vector<PT> h(2*n + 5);
  for (int i = 0; i < n; i++) {
    while (k \ge 2 \&\& cross(h[k-1] - h[k-2], p[i] - h[k-2]) \le 0) k--;
   h[k++] = p[i];
  for (int i = n - 2, t = k + 1; i >= 0; i--) {
   while (k \ge t \&\& cross(h[k-1] - h[k-2], p[i] - h[k-2]) \le 0) k--;
    h[k++] = p[i];
  h.resize(k); // n+1 points where the first is equal to the last
  return h;
vector<PT> splitHull(const vector<PT> &hull) {
  vector<PT> ans(hull.size()):
```

```
for (int i = 0, j = (int) hull.size()-1, k = 0; k < (int) hull.size(); k++) {
   if(hull[i] < hull[j]) {</pre>
     ans[k] = hull[i++];
   else
     ans[k] = hull[j--];
 return ans:
// uniao de convex hulls
vector<PT> ConvexHull(const vector<PT> &a, const vector<PT> &b) {
 auto A = splitHull(a);
 auto B = splitHull(b);
 vector<PT> C(A.size() + B.size());
 merge(A.begin(), A.end(), B.begin(), B.end(), C.begin());
 return ConvexHull(C, false);
int maximizeScalarProduct(const vector<PT> &hull, PT vec) {
 // this code assumes that there are no 3 colinear points
 int ans = 0;
 int n = hull.size();
 if(n < 20) {
   for (int i = 0; i < n; i++) {
     if(dot(hull[i], vec) > dot(hull[ans], vec)) {
       ans = i;
 } else {
   if(dot(hull[1], vec) > dot(hull[ans], vec)) {
     ans = 1:
   for(int rep = 0; rep < 2; rep++) {</pre>
     int 1 = 2, r = n - 1;
     while (1 != r) {
       int mid = (1 + r + 1) / 2;
       bool flag = dot(hull[mid], vec) >= dot(hull[mid-1], vec);
       if(rep == 0) { flag = flag && dot(hull[mid], vec) >= dot(hull[0], vec); }
       else { flag = flag || dot(hull[mid-1], vec) < dot(hull[0], vec); }</pre>
       if(flag) {
          1 = mid;
       } else {
         r = mid - 1;
     if(dot(hull[ans], vec) < dot(hull[1], vec)) {</pre>
       ans = 1:
 return ans:
```

5.3 Cut Polygon

```
struct Segment {
    typedef long double T;
    typ pl, p2;
    T a, b, c;

Segment() {}

Segment(PT st, PT en) {
    p1 = st, p2 = en;
    a = -(st.y - en.y);
    b = st.x - en.x;
    c = a * en.x + b * en.y;
}

T plug(T x, T y) {
    // plug >= 0 is to the right
    return a * x + b * y - c;
}

T plug(PT p) {
    return plug(p.x, p.y);
}

bool inLine(PT p) { return cross((p - p1), (p2 - p1)) == 0; }

bool inSegment(PT p) {
    return inLine(p) && dot((p1 - p2), (p - p2)) >= 0 && dot((p2 - p1), (p - p1)) >=
    0;
}
```

```
PT lineIntersection(Segment s) {
    long double A = a, B = b, C = c;
long double D = s.a, E = s.b, F = s.c;
long double x = (long double) C * E - (long double) B * F;
long double y = (long double) A * F - (long double) B * D;
    long double tmp = (long double) A * E - (long double) B * D;
    x /= tmp;
    v /= tmp;
    return PT(x, y);
  bool polygonIntersection(const vector<PT> &poly) {
    long double 1 = -1e18, r = 1e18;
    for(auto p : poly) {
       long double z = plug(p);
       1 = \max(1, z);
       r = min(r, z);
    return 1 - r > eps;
};
vector<PT> cutPolygon(vector<PT> poly, Segment seg) {
  int n = (int) poly.size();
  vector<PT> ans;
  for (int i = 0; i < n; i++)
    double z = seg.plug(poly[i]);
    if(z > -eps) {
       ans.push_back(poly[i]);
    double z2 = seg.plug(poly[(i + 1) % n]);
    if((z > eps && z2 < -eps) || (z < -eps && z2 > eps)) {
       ans.push_back(seg.lineIntersection(Segment(poly[i], poly[(i + 1) % n])));
  return ans;
```

5.4 Smallest Enclosing Circle

```
typedef pair<PT, double> circle;
bool inCircle (circle c, PT p) {
  return cmp(dist(c.first, p), c.second) <= 0;</pre>
PT circumcenter (PT p, PT q, PT r) {
  PT a = p-r, b = q-r;
PT c = PT(dot(a, p+r)/2, dot(b, q+r)/2);
  return PT(cross(c, PT(a.y,b.y)), cross(PT(a.x,b.x), c)) / cross(a, b);
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
circle spanningCircle (vector<PT> &v) {
  int n = v.size();
  shuffle(v.begin(), v.end(), rng);
  circle C(PT(), -1);
  for (int i = 0; i < n; i++) if (!inCircle(C, v[i])) {</pre>
    C = circle(v[i], 0);
    for (int j = 0; j < i; j++) if (!inCircle(C, v[j])) {</pre>
      C = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
      for (int k = 0; k < j; k++) if (!inCircle(C, v[k])) {
        PT o = circumcenter(v[i], v[j], v[k]);
        C = circle(o, dist(o, v[k]));
  return C;
```

5.5 Minkowski

```
bool comp(PT a, PT b) {
  int hp1 = (a.x < 0 || (a.x==0 && a.y<0));
  int hp2 = (b.x < 0 || (b.x==0 && b.y<0));
  if(hp1 != hp2) return hp1 < hp2;</pre>
```

```
long long R = cross(a, b);
 if(R) return R > 0;
 return dot(a, a) < dot(b, b);</pre>
 // This code assumes points are ordered in ccw and the first points in both vectors
     is the min lexicographically
vector<PT> minkowskiSum(const vector<PT> &a, const vector<PT> &b) {
 if(a.empty() || b.empty()) return vector<PT>(0);
 vector<PT> ret;
 int n1 = a.size(), n2 = b.size();
 if(min(n1, n2) < 2){
   for(int i = 0; i < n1; i++) {
      for (int j = 0; j < n2; j++) {
       ret.push_back(a[i]+b[j]);
   return ret;
 PT v1, v2, p = a[0]+b[0];
 ret.push_back(p);
 for (int i = 0, j = 0; i + j + 1 < n1+n2; ) {
   v1 = a[(i+1)*n1]-a[i];
   v2 = b[(j+1) n2] - b[j];
   if(j == n2 || (i < n1 && comp(v1, v2))) p = p + v1, i++;</pre>
   else p = p + v2, j++;
   while(ret.size() >= 2 && cmp(cross(p-ret.back(), p-ret[(int)ret.size()-2])) == 0)
      // removing colinear points
      // needs the scalar product stuff it the result is a line
      ret.pop_back();
   ret .push_back (p);
 return ret;
```

5.6 Half Plane Intersection

```
struct L { // salvar (p[i], p[i + 1]) poligono CCW, (p[i + 1], p[i]) poligono CW
    PT a, b, dir;
    L(){}
    L(PT a, PT b) : a(a), b(b) {
      dir = b - a;
    int quadrant() const {
      if (dir.y > 0 && dir.x >= 0) return 0;
      if (dir.x < 0 && dir.y >= 0) return 1;
      if (dir.y < 0 && dir.x <= 0) return 2;</pre>
      return 3:
    bool operator < (const L &1) const {
      int q1 = quadrant(), q2 = 1.quadrant();
      if (q1 != q2) return q1 < q2;</pre>
      double c = cross(dir, 1.dir);
      if(cmp(c) == 0) {
        return cmp(cross((1.b - 1.a), (b - 1.a))) > 0;
      return cmp(c) > 0;
};
PT computeLineIntersection (L la, L lb) {
    return lineLine(la.a, la.b, lb.a, lb.b);
bool check (L la, L lb, L lc) {
    PT p = computeLineIntersection(lb, lc);
double det = cross((la.b - la.a), (p - la.a));
    return cmp(det) < 0;
vector<PT> hpi (vector<L> line) {
    vector<PT> box = {PT(inf, inf), PT(-inf, inf), PT(-inf, -inf), PT(inf, -inf)};
    for(int i = 0; i < 4; i++)
        line.emplace_back(box[i], box[(i + 1) % 4]);
    sort(line.begin(), line.end());
    vector<L> pl(1, line[0]);
    for (int i = 0; i < (int)line.size(); ++i) if (cmp(cross(line[i].dir, pl.back().</pre>
         dir)) != 0) pl.push_back(line[i]);
    vector<int> dq;
    int start = 0;
    for (int i = 0; i < (int)pl.size(); ++i) {</pre>
```

```
while ((int)dq.size() - start > 1 && check(pl[i], pl[dq.back()], pl[dq[dq.size
         () - 2]])) dq.pop_back();
    while ((int)dq.size() - start > 1 && check(pl[i], pl[dq[start]], pl[dq[start +
          1]])) start++;
    if((int)dq.size() - start > 0 && cmp(cross(pl[i].dir, pl[dq.back()].dir)) ==
      if(cmp(dot(pl[i].dir, pl[dq.back()].dir)) < 0) return vector<PT>();
      if(cmp(cross(pl[i].dir, pl[dq.back()].a - pl[i].a)) < 0) dq.pop_back();
      else continue:
    dq.push back(i);
while ((int)dq.size() - start > 1 && check(pl[dq[start]], pl[dq.back()], pl[dq[dq.
     size() - 2]])) dq.pop_back();
while ((int)dq.size() - start > 1 && check(pl[dq.back()], pl[dq[start]], pl[dq[
     start + 1]])) start++;
vector<PT> res;
if((int)dq.size() - start < 3) return vector<PT>(); // remove this if res can be
     point/line
for (int i = start; i < (int)dq.size(); ++i){</pre>
  res.emplace_back(computeLineIntersection(pl[dq[i]], pl[dq[i + 1 == dq.size() ?
      start : i + 1]]));
return res;
```

5.7 Closest Pair

```
double closestPair(vector<PT> p) {
  int n = p.size(), k = 0;
  sort(p.begin(), p.end());
  double d = inf;
  set<PT> ptsInv;
  for(int i = 0; i < n; i++) {
    while(k < i && p[k].x < p[i].x - d) {
      ptsInv.erase(swapCoord(p[k++]));
    }
  for(auto it = ptsInv.lower_bound(PT(p[i].y - d, p[i].x - d));
      it != ptsInv.end() && it->x <= p[i].y + d; it++) {
      d = min(d, dist(p[i] - swapCoord(*it), PT(0, 0)));
    }
  ptsInv.insert(swapCoord(p[i]));
}
return d;
}</pre>
```

5.8 Voronoi

```
Segment getBisector(PT a, PT b) {
  Segment ans (a, b);
  swap (ans.a, ans.b);
  ans.b \star = -1;
  ans.c = ans.a \star (a.x + b.x) \star 0.5 + ans.b \star (a.y + b.y) \star 0.5;
  return ans;
// BE CAREFUL!
// the first point may be any point
// O(N^3)
vector<PT> getCell(vector<PT> pts, int i) {
  vector<PT> ans;
  ans.emplace_back(0, 0);
  ans.emplace_back(1e6, 0);
  ans.emplace_back(1e6,
  ans.emplace_back(0, 1e6);
  for(int j = 0; j < (int) pts.size(); <math>j++) {
   if(j!= i) {
      ans = cutPolygon(ans, getBisector(pts[i], pts[j]));
  return ans;
// O(N^2) expected time
vector<vector<PT>> getVoronoi(vector<PT> pts) {
  // assert (pts.size() > 0);
  int n = (int) pts.size();
  vector<int> p(n, 0);
```

```
for(int i = 0; i < n; i++) {</pre>
 p[i] = i;
shuffle(p.begin(), p.end(), rng);
vector<vector<PT>> ans(n);
ans[0].emplace_back(0, 0);
ans[0].emplace_back(w, 0);
ans[0].emplace_back(w, h);
ans[0].emplace_back(0, h);
for (int i = 1; i < n; i++) {
 ans[i] = ans[0];
for(auto i : p) {
 for(auto j : p) {
    if(j == i) break;
    auto bi = getBisector(pts[j], pts[i]);
    if(!bi.polygonIntersection(ans[j])) continue;
    ans[j] = cutPolygon(ans[j], getBisector(pts[j], pts[i]));
    ans[i] = cutPolygon(ans[i], getBisector(pts[i], pts[j]));
return ans;
```

6 String Algorithms

6.1 KMP

```
vector<int> getBorder(string str) {
  int n = str.size();
  vector<int> border(n, -1);
  for (int i = 1, j = -1; i < n; i++) {
    while (j \ge 0 \&\& str[i] != str[j+1]) {
      j = border[j];
    if(str[i] == str[j + 1]) {
      j++;
   border[i] = j;
  return border;
int matchPattern(const string &txt, const string &pat, const vector<int> &border) {
  int freq = 0;
  for(int i = 0, j = -1; i < txt.size(); i++) {</pre>
    while (j \ge 0 \&\& txt[i] != pat[j + 1]) {
      j = border[j];
    if(pat[j + 1] == txt[i]) {
     j++;
    if(j + 1 == (int) pat.size()) {
       //found occurence
      frea++;
      j = border[j];
  return freq;
```

6.2 Aho-Corasick

```
const int ms = le6;  // quantidade de caracteres
const int sigma = 26;  // tamanho do alfabeto
int trie[ms][sigma], fail[ms], superfail[ms], terminal[ms], z = 1;

void add(string &s) {
   int node = 0;
   for (char ch : s) {
     int pos = val(ch);  // no caso de alfabeto a-z: val(ch) = ch - 'a'
     if (!trie[node][pos]) {
        terminal[z] = 0;
        trie[node][pos] = z++;
   }
```

```
node = trie[node][pos];
  ++terminal[node]; // trocar pela info que quiser
void buildFailure() {
 memset(fail, 0, sizeof(int) * z), memset(superfail, 0, sizeof(int) * z);
  queue<int> Q;
  O.push(0);
  while (Q.size()) {
   int node = Q.front();
    for (int pos = 0; pos < sigma; ++pos) {</pre>
     int &v = trie[node][pos];
      int f = node == 0 ? 0 : trie[fail[node]][pos];
      // int sf = present[f] ? f : superfail[f];
      // present marks if that vertex is a terminal node or not
      // if summing up on terminal, doesn't work
      if (!v) {
      } else {
        fail[v] = f;
      // superfail[v] = sf;
        Q.push(v);
        // dar merge nas infos (por ex: terminal[v] += terminal[f])
void search(string &s) {
  int node = 0;
  for (char ch : s)
   int pos = val(ch);
   node = trie[node][pos];
    // processar infos no no atual (por ex: ocorrencias += terminal[node])
// se tiver usando super fail, cuidado com o estado que voce ta, antes de subir pro sf
     , porque pode ser que o estado que ta nao seja no terminal
```

6.3 Algoritmo de Z

```
template <class T>
vector<int> ZFunc(const vector<T> &v) {
  vector<int> z(v.size(), 0);
  int n = (int) v.size(), a = 0, b = 0;
  if (!z.empty()) z[0] = n;
  for (int i = 1; i < n; i++) {
    int end = i; if (i < b) end = min(i + z[i - a], b);
    while(end < n && v[end] == v[end - i]) ++end;
    z[i] = end - i; if (end > b) a = i, b = end;
}
return z;
}
```

6.4 Suffix Array

```
vector<int> buildSa(const string& in) {
   int n = in.size(), c = 0;
   vector<int> temp(n), posBucket(n), bucket(n), bpos(n), out(n);
   for (int i = 0; i < n; i++) out[i] = i;
   sort(out.begin(), out.end(), [&](int a, int b) { return in[a] < in[b]; });
   for (int i = 0; i < n; i++) {
      bucket[i] = c;
      if (i + 1 == n || in[out[i]] != in[out[i + 1]]) c++;
   }
   for (int h = 1; h < n && c < n; h <<= 1) {
      for (int i = 0; i < n; i++) posBucket[out[i]] = bucket[i];
      for (int i = n - 1; i >= 0; i--) bpos[bucket[i]] = i;
      for (int i = 0; i < n; i++) {
        if (out[i] >= n - h) temp[bpos[bucket[i]]++] = out[i];
    }
   for (int i = 0; i < n; i++) {
      if (out[i] >= h) temp[bpos[posBucket[out[i] - h]]++] = out[i] - h;
   }
```

```
c = 0;
   for (int i = 0; i + 1 < n; i++) {
        int a = (bucket[i] != bucket[i + 1]) || (temp[i] >= n - h)
           || (posBucket[temp[i + 1] + h] != posBucket[temp[i] + h]);
        bucket[i] = c;
       c += a;
   bucket[n - 1] = c++;
   temp.swap(out);
 return out;
vector<int> buildLcp(string s, vector<int> sa) {
 int n = (int) s.size();
 vector<int> pos(n), lcp(n, 0);
 for (int i = 0; i < n; i++) {
   pos[sa[i]] = i;
 int k = 0;
 for(int i = 0; i < n; i++) {</pre>
   if (pos[i] + 1 == n) {
     k = 0:
      continue;
   int j = sa[pos[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
   lcp[pos[i]] = k;
   k = \max(k - 1, 0);
 return lcp;
```

6.5 Suffix Automaton

```
int len[ms*2], link[ms*2], nxt[ms*2][sigma];
int sz, last;
void build(string &s) {
  len[0] = 0; link[0] = -1;
  sz = 1; last = 0;
  memset(nxt[0], -1, sizeof nxt[0]);
  for(char ch : s) {
    int c = ch-'a', cur = sz++;
    len[cur] = len[last]+1;
    memset(nxt[cur], -1, sizeof nxt[cur]);
    int p = last;
    while (p !=-1 \&\& nxt[p][c] == -1) {
      nxt[p][c] = cur; p = link[p];
    if(p == -1) {
      link[cur] = 0;
    } else {
      int q = nxt[p][c];
      if(len[p] + 1 == len[q]) {
        link[cur] = q;
      } else {
        len[sz] = len[p]+1; link[sz] = link[q];
        \label{eq:memcpy} \texttt{(nxt[sz], nxt[q], sizeof } \texttt{nxt[q]);}
        while (p != -1 && nxt[p][c] == q) {
          nxt[p][c] = sz; p = link[p];
        link[q] = link[cur] = sz++;
    last = cur;
```

6.6 Manacher

```
std::array<std::vector<int>, 2> manacher(const std::string& s) {
   int n = (int) s.size();
   std::array<std::vector<int>, 2> p = {std::vector<int>(n+1), std::vector<int>(n
   )};
   for(int z = 0; z < 2; z++) for (int i = 0, l = 0, r = 0; i < n; i++) {
        int t = r-i+!z;
        if (i<r) p[z][i] = std::min(t, p[z][l+t]);</pre>
```

6.7 Polish Notation

```
inline bool isOp(char c) {
        return c=='+' || c=='-' || c=='*' || c=='/' || c=='^';
inline bool isCarac(char c) {
        return (c>='a' && c<='z') || (c>='A' && c<='Z') || (c>='0' && c<='9');
int paren2polish(char* paren, char* polish) {
        map<char, int> prec;
prec['('] = 0;
prec['+'] = prec['-'] = 1;
        prec['*'] = prec['/'] = 2;
        prec['^'] = 3;
        int len = 0;
        stack<char> op;
        for (int i = 0; paren[i]; i++) {
                if (isOp(paren[i])) {
                         while (!op.empty() && prec[op.top()] >= prec[paren[i]]) {
                                  polish[len++] = op.top(); op.pop();
                         op.push(paren[i]);
                 else if (paren[i] == '(') op.push('(');
                 else if (paren[i]==')')
                         for (; op.top()!='('; op.pop())
                                  polish[len++] = op.top();
                         op.pop();
                 else if (isCarac(paren[i]))
                         polish[len++] = paren[i];
        for(; !op.empty(); op.pop())
                polish[len++] = op.top();
        polish[len] = 0;
        return len;
```

6.8 String Hash

```
struct StringHashing {
  const uint64_t MOD = (1LL << 61) - 1;</pre>
  const int base = 31;
  vector<uint64 t> h, p;
  uint64_t modMul(uint64_t a, uint64_t b) {
   uint64_t 11 = (uint32_t)a, h1 = a >> 32, 12 = (uint32_t)b, h2 = b >> 32;
   uint64_t 1 = 11 * 12, m = 11 * h2 + 12 * h1, h = h1 * h2;
    uint64 t ret = (1 \& MOD) + (1 >> 61) + (h << 3) + (m >> 29) + ((m << 35) >> 3) +
    ret = (ret & MOD) + (ret >> 61);
    ret = (ret & MOD) + (ret >> 61);
    return ret - 1;
  uint64_t getKey(int 1, int r) { // [1, r]
   uint64_t res = h[r];
    if(1 > 0) res = (res + MOD - modMul(p[r - 1 + 1], h[1 - 1])) % MOD;
    return res;
 uint64_t getInt(char c) {
  return c - 'a' + 1;
  StringHashing(string &s) {
    int n = s.size();
```

```
h.resize(n);
p.resize(n);
p[0] = 1;
h[0] = getInt(s[0]);
for(int i = 1; i < n; ++i) {
    p[i] = modMul(p[i - 1], base);
    h[i] = (modMul(h[i - 1], base) + getInt(s[i])) % MOD;
}
};
</pre>
```

7 Miscellaneous

7.1 Random Number Generator

```
// mt19937_64 se LL
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
// Random_Shuffle
shuffle(v.begin(), v.end(), rng);
// Random number in interval
int randomInt = uniform_int_distribution(0, i)(rng);
double randomDouble = uniform_real_distribution(0, 1)(rng);
// bernoulli_distribution, binomial_distribution, geometric_distribution
// normal_distribution, poisson_distribution, exponential_distribution
```

7.2 Safe Hash

7.3 Unordered Map Tricks

```
// pair<int, int> hash function
struct HASH(
    size_t operator()(const pair<int,int>&x)const(
    return (size_t) x.first * 37U + (size_t) x.second;
};

unordered_map<int,int>mp;
mp.reserve(1024);
mp.max_load_factor(0.25);
```

7.4 Iterate masks in bitcount order

```
for(int k = n-1; k >= 0; k--) {
  unsigned int i = (1 << k) -1;
  while(i < (1 << n)) {
    // do what you want
    unsigned int t = (i | (i - 1)) + 1;</pre>
```

```
if(i == 0) break;
i = t | ((((t & -t) / (i & -i)) >> 1) - 1);
}
```

7.5 Submask Enumeration

```
for (int s=m; ; s=(s-1)&m) {
    ... you can use s ...
if (s==0) break;
}
```

7.6 Sum over Subsets DP

```
// F[i] = Sum of all A[j] where j is a submask of i
for(int i = 0; i<(1<<N); ++i)
   F[i] = A[i];
for(int i = 0;i < N; ++i) for(int mask = 0; mask < (1<<N); ++mask){
   if(mask & (1<<i))
      F[mask] += F[mask^(1<<i)];
}</pre>
```

7.7 Subset Sum

```
* Given N non-negative integer weights w and a non-negative target t,
 * computes the maximum S \le t such that S is the sum of some subset of the weights.
 * Time: O(N \max(w_i))
int knapsack(vector<int> w, int t) {
  int a = 0, b = 0;
  while (b < w.size() && a + w[b] <= t) a += w[b++];</pre>
  if (b == w.size()) return a;
  int m = *max_element(w.begin(), w.end());
  vector < int > u, v(2*m, -1);
  v[a+m-t] = b;
  for(int i = b; i < w.size(); i++) {</pre>
    for (int x = 0; x < m; x++) v[x+w[i]] = max(v[x+w[i]], u[x]);
    for (int x = 2*m; --x > m;)
      for (int j = max(011, u[x]); j < v[x]; j++)
        v[x-w[j]] = max(v[x-w[j]], j);
  for (a = t; v[a+m-t] < 0; a--);
  return a;
```

7.8 Regular Expressions

```
import java.util.*;
import java.util.regex.*;

public class Main {
    public static String BuildRegex () {
        return "^" + sentence + "$";
    }

    public static void main (String args[]) {
        String regex = BuildRegex ();
        // check pattern documentation
        Pattern pattern = Pattern.compile (regex);
        Scanner s = new Scanner(System.in);
        String sentence = s.nextLine().trim();
        boolean found = pattern.matcher(sentence).find()
    }
}
```

7.9 Lat Long

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
struct 11
  double r, lat, lon;
struct rect
  double x, y, z;
11 convert (rect& P)
 11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
 Q.lat = 180/M_PI*asin(P.z/Q.r);
Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
  return ():
rect convert(11& 0)
  P.x = O.r*cos(O.lon*M PI/180)*cos(O.lat*M PI/180);
 P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
 P.z = O.r*sin(O.lat*M PI/180);
  return P:
```

7.10 Stable Marriage

```
std::vector<std::vector<int>> stableMarriage(std::vector<std::vector<int>> first, std
     ::vector<std::vector<int>> second, std::vector<int> cap) {
        assert(cap.size() == second.size());
        int n = (int) first.size(), m = (int) second.size();
        // if O(N \star M) first in memory, use table
        std::map<std::pair<int, int>, int> prio;
        std::vector<std::set<std::pair<int, int>>> current(m);
        for (int i = 0; i < n; i++) {
                std::reverse(first[i].begin(), first[i].end());
        for (int i = 0; i < m; i++)
                for(int j = 0; j < (int) second[i].size(); j++) {</pre>
                        prio[{second[i][j], i}] = j;
        for(int i = 0; i < n; i++) {</pre>
                int on = i;
                while(!first[on].empty()) {
                        int to = first[on].back();
                         first[on].pop_back();
                        if(cap[to]) {
                                cap[to]--:
                                assert(prio.count({on, to}));
                                current[to].insert({prio[{on, to}], on});
                        assert(!current[to].empty());
                        auto it = current[to].end();
                        if(it->first > prio[{on, to}]) {
                                int nxt = it->second;
                                current[to].erase(it);
                                current[to].insert({prio[{on, to}], on});
                                on = nxt;
        std::vector<std::vector<int>> ans(m);
        for(int i = 0; i < m; i++) {</pre>
                for(auto it : current[i]) {
```

```
ans[i].push_back(it.second);
}
return ans;
```

8 Teoremas e formulas uteis

8.1 Grafos

```
Formula de Euler: V - E + F = 2 (para grafo planar)
Handshaking: Numero par de vertices tem grau impar
Kirchhoff's Theorem: Monta matriz onde Mi,i = Grau[i] e Mi,j = -1 se houver aresta i-j
      ou 0 caso contrario, remove uma linha e uma coluna qualquer e o numero de
     spanning trees nesse grafo eh o det da matriz
Grafo contem caminho hamiltoniano se:
Dirac's theorem: Se o grau de cada vertice for pelo menos n/2
Ore's theorem: Se a soma dos graus que cada par nao-adjacente de vertices for pelo
Boruvka's: enquanto grafo nao conexo, para cada componente conexa use a aresta que sai
      de menor custo.
Tem Catalan(N) Binary trees de N vertices
Tem Catalan (N-1) Arvores enraizadas com N vertices
Caley Formula: n^(n-2) arvores em N vertices com label
Prufer code: Cada etapa voce remove a folha com menor label e o label do vizinho eh
     adicionado ao codigo ate ter 2 vertices
Recuperar min cut eh ver se level[u] != -1 ai eh do lado do source
Max Edge-disjoint paths: Max flow com arestas com peso 1
Max Node-disjoint paths: Faz a mesma coisa mas separa cada vertice em um com as
     arestas de chegadas e um com as arestas de saida e uma aresta de peso 1
     conectando o vertice com aresta de chegada com ele mesmo com arestas de saida
Konig's Theorem: minimum node cover = maximum matching se o grafo for bipartido,
     complemento eh o maximum independent set
Min vertex cover sao os vertices da particao do source que nao tao do lado do source
     do cut e os do sink que tao do lado do source, independent set o contrario
Min edge cover eh N - match, pega as arestas do match mais qualquer aresta dos outros
     vertices
Min Node disjoint path cover: formar grafo bipartido de vertices duplicados, onde
     aresta sai do vertice tipo A e chega em tipo B, entao o path cover eh N -
Min General path cover: Mesma coisa mas colocando arestas de A pra B sempre que houver
      caminho de A pra B
Dilworth's Theorem: Min General Path cover = Max Antichain (set de vertices tal que
     nao existe caminho no grafo entre vertices desse set)
Hall's marriage: um grafo tem um matching completo do lado X se para cada subconjunto
    |W| \le |vizinhosW| onde |W| eh quantos vertices tem em W
feasible flow in a network with both upper and lower capacity constraints, no source
     or sink: capacities are changed to upper bound - lower bound. Add a new source
     and a sink. let M[v] = (sum \ of \ lower \ bounds \ of \ ingoing \ edges \ to \ v) - (sum \ of \ lower \ bounds \ of \ ingoing \ edges \ to \ v)
     lower bounds of outgoing edges from v). For all v, if M[v] > 0 then add edge (S, v
     ) with capcity M, otherwise add (v,T) with capacity -M. If all outgoing edges
     from S are full, then a feasible flow exists, it is the flow plus the original
     lower bounds
```

8.2 Math

```
Goldbach's: todo numero par n > 2 pode ser representado com n = a + b onde a e b sao primos

Twin prime: existem infinitos pares p, p + 2 onde ambos sao primos

Legendre's: sempre tem um primo entre n^2 e (n+1)^2

Lagrange's: todo numero inteiro pode ser inscrito como a soma de 4 quadrados

Zeckendorf's: todo numero pode ser representado pela soma de dois numeros de fibonnacis diferentes e nao consecutivos

Euclid's: toda tripla de pitagoras primitiva pode ser gerada com (n^2 - m^2, 2nm, n^2+m^2) onde n, m sao coprimos e um deles eh par

Wilson's: n eh primo quando (n-1)! mod n = n - 1
```

```
Mcnugget: Para dois coprimos x, y a quantidade de inteiros que nao pode ser escrito como ax + by eh (x-1)(y-1)/2,
     o maior inteiro que nao conseque eh x*y-x-y
Fermat: Se p eh primo entao a(p-1) % p = 1
Se x e m tambem forem coprimos entao x^k % m = x^k (k \mod (m-1)) % m
Euler's theorem: x^{(phi(m))} \mod m = 1 onde phi(m) eh o totiente de euler
Chinese remainder theorem:
Para equacoes no formato x = al \mod ml, ..., x = an \mod mn onde todos os pares ml,
      ..., mn sao coprimos
Deixe Xk = m1 * m2 * ... * mn/mk e Xk^-1 mod mk = inverso de Xk mod mk, entao
x = somatorio com k de 1 ate n de ak*Xk*(Xk,mk^-1 mod mk)
Para achar outra solucao so somar ml*m2*..*mn a solucao existente
Catalan number: exemplo expressoes de parenteses bem formadas
C0 = 1, Cn = somatorio de <math>i=0 \rightarrow n-1 de Ci*C(n-1+1)
outra forma: Cn = (2n \text{ escolhe } n)/(n+1)
Bertrand's ballot theorem: p votos tipo A e q votos tipo B com p>q, prob de em todo
     ponto ter mais As do que Bs antes dele = (p-q)/(p+q)
Se puder empates entao prob = (p+1-q)/(p+1), para achar quantidade de possibilidades
     nos dois casos basta multiplicar por (p + g escolhe g)
Propriedades de Coeficientes Binomiais:
Somatorio de k = 0 \rightarrow m de (-1)^k \star (n \text{ escolhe } k) = (-1)^m \star (n - 1 \text{ escolhe } m)
(N escolhe K) = (N escolhe N-K)
(N \text{ escolhe } K) = N/K * (n-1 \text{ escolhe } k-1)
Somatorio de k = 0 \rightarrow n de (n escolhe k) = 2^n
Somatorio de m = 0 \rightarrow n de (m \ escolhe \ k) = (n+1 \ escolhe \ k+1)
Somatorio de k = 0 \rightarrow m de (n+k \text{ escolhe } k) = (n+m+1 \text{ escolhe } m)
Somatorio de k = 0 \rightarrow n de (n escolhe k)^2 = (2n escolhe n)
Somatorio de k = 0 ou 1 \rightarrow n de k*(n escolhe k) = n * 2^(n-1)
Somatorio de k = 0 \rightarrow n de (n-k \text{ escolhe } k) = \text{Fib}(n+1)
Hockey-stick: Somatorio de i = r \rightarrow n de (i = scolhe r) = (n + 1 = scolhe r + 1)
Vandermonde: (m+n \text{ escolhe } r) = \text{somatorio de } k = 0 \rightarrow r \text{ de } (m \text{ escolhe } k) \star (n \text{ escolhe } r)
Burnside lemma: colares diferentes nao contando rotacoes quando m = cores e n =
     comprimento
(m^n + somatorio i = 1 - > n-1 de m^gcd(i, n))/n
Distribuicao uniforme a, a+1, ..., b Expected[X] = (a+b)/2
Distribuicao binomial com n tentativas de probabilidade p, X = sucessos:
   P(X = x) = p^x * (1-p)^(n-x) * (n escolhe x) e E[X] = p*n
Distribuicao geometrica onde continuamos ate ter sucesso, X = tentativas:
    P(X = x) = (1-p)^(x-1) * p e E[X] = 1/p
Linearity of expectation: Tendo duas variaveis X e Y e constantes a e b, o valor
     esperado de aX + bY = a*E[X] + b*E[X]
V(X) = E((X-u)^2)
V(X) = E(X^2) - E(X^2)
```

```
PG: al * (q^n - 1)/(q - 1)
Mobius Inverse: Sum(d|n): mobius(d) = [n = 1] (expressao booleana)
Soma dos cubos de 1 a N = a^2 onde a = somatorio de 1 a N
Lindstrom-Gessel-Viennot: quantidade de caminhos disjuntos nas linhas do grid eh o determinante da matriz de qnts caminhos
```

8.3 Geometry

Formula de Euler: V - E + F = 2

```
Pick Theorem: Para achar pontos em coords inteiras num poligono Area = i + b/2 - 1
     onde i eh o o numero de pontos dentro do poligono e b de pontos no perimetro do
     poligono
Two ears theorem: Todo poligono simples com mais de 3 vertices tem pelo menos 2
     orelhas, vertices que podem ser removidos sem criar um crossing, remover orelhas
     repetidamente triangula o poligono
Incentro triangulo: (a(Xa, Ya) + b(Xb, Yb) + c(Xc, Yc))/(a+b+c) onde a = lado oposto
     ao vertice a, incentro eh onde cruzam as bissetrizes, eh o centro da
     circunferencia inscrita e eh equidistante aos lados
Delaunay Triangulation: Triangulacao onde nenhum ponto esta dentro de nenhum circulo
    circunscrito nos triangulos
Eh uma trianqulacao que maximiza o menor anqulo e a MST euclidiana de um conjunto de
    pontos eh um subconjunto da triangulação
Brahmagupta's formula: Area cyclic quadrilateral
s = (a+b+c+d)/2
area = sqrt((s-a)*(s-b)*(s-c)*(s-d))
d = 0 \Rightarrow area = sqrt((s-a)*(s-b)*(s-c)*s)
```

8.4 Dynamic Programming

```
Divide and conquer - dp[i][j] = mink < j{dp[i - 1][k] + C[k][j]} dividir o subsegmento ate j em i segmentos com custo, valido se A[i][j] <= A[i][j+1] Knuth - p[i][j] = mini < k < j{dp[i][k] + dp[k][j]} + C[i][j], valido se A[i, j - 1] <= A[i][j] <= A[i+1, j] onde A[i][j] eh o menor k que da a resposta otima slope trick - funcao piecewise linear convexa, descrita pelos pontos de mudanca de slope (multiset/heap) lembre que existe fft, cht, aliens trick e bitset
```