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fw.resize(ord.size() + 1);

```
coord.resize(fw.size());
    for(auto &a : pts) {
      swap(a.first, a.second);
    sort(pts.begin(), pts.end());
    for(auto &a : pts) {
      swap(a.first, a.second);
      for(int on = upper_bound(ord.begin(), ord.end(), a.first) - ord.begin
          (); on < fw.size(); on += on & -on) {
        if(coord[on].empty() || coord[on].back() != a.second) {
          coord[on].push_back(a.second);
    for(int i = 0; i < fw.size(); i++) {</pre>
      fw[i].assign(coord[i].size() + 1, 0);
  void upd(T x, T y, T v) {
    for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin(); xx <</pre>
         fw.size(); xx += xx & -xx) {
      for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y) -
          coord[xx].begin(); yy < fw[xx].size(); yy += yy & -yy) {</pre>
        fw[xx][yy] += v;
  \hat{T} qry (T x, T y) {
    T ans = 0;
    for(int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin(); xx >
         0; xx -= xx & -xx) {
      for(int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y) -
          coord[xx].begin(); yy > 0; yy -= yy & -yy) {
        ans += fw[xx][yy];
    return ans;
    qry(T x1, T y1, T x2, T y2) {
    return qry(x2, y2) - qry(x2, y1 - 1) - qry(x1 - 1, y2) + qry(x1 - 1, y1
         - 1);
  void upd(T x1, T y1, T x2, T y2, T v) {
    upd(x1, y1, v);
    upd(x1, y2 + 1, -v);
    upd (x2 + 1, y1, -v);
    upd(x2 + 1, y2 + 1, v);
private:
  vector<T> ord;
  vector<vector<T>> fw, coord;
```

1.2 Seg Tree

```
int n, t[2 * ms];

void build() {
    for(int i = n - 1; i > 0; --i) t[i] = t[i<<1] + t[i<<1|1]; // Merge
}

void update(int p, int value) { // set value at position p
    for(t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1]; // Merge
}

int query(int 1, int r) {
    int res = 0;
    for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1) {
        if(1&1) res += t[1++]; // Merge
        if(r&1) res += t[--r]; // Merge
    }
}
```

```
return res;
}

// If is non-commutative
S query(int 1, int r) {
    S resl, resr;
    for (1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1) {
        if (l&l) resl = combine(resl, t[l++]);
        if (r&l) resr = combine(t[--r], resr);
    }
    return combine(resl, resr);
}
```

1.3 Seg Tree cm update de range

```
int n, t[2 * ms];
void build() {
  for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1]; // Merge
void update(int v, int 1, int r) {
  for (\bar{1} + n, r + n+1; 1 < r; 1 >>= 1, r >>= 1) {
    if(1&1) t[1++] += v; // Merge
    if(r\&1) t[--r] += v; // Merge
int query(int p) {
  int res = 0;
  for(p += n; p > 0; p >>= 1) res += t[p]; // Merge
  return res;
void push() { // push modifications to leafs
  for (int i = 1; i < n; i++) {
    t[i<<1] += t[i]; // Merge
    t[i<<1|1] += t[i]; // Merge
    t[i] = 0;
```

1.4 Seg Tree cm Lazy Iterativa

```
struct LazyContext {
 LazyContext() { }
  void reset() { }
  void operator += (LazyContext o) { }
struct Node {
  Node() {
  Node() { }
  Node (Node 1, Node r) { }
  bool canBreak(LazyContext lazy) { } // false if non beats
  bool canApply(LazyContext lazy) { } // true if non beats
  void apply(LazyContext &lazy) { }
template <class i_t, class e_t, class lazy_cont>
class SegmentTree {
  void init(std::vector<e_t> base) {
    n = base.size();
    h = 0;
    while ((1 << h) < n) h++;
    tree.resize(2 * n);
    dirty.assign(n, false);
    lazy.resize(n);
```

```
for (int i = 0; i < n; i++) {
      tree[i + n] = i t(base[i]);
    for (int i = n - 1; i > 0; i--) {
      tree[i] = i_t(tree[i + i], tree[i + i + 1]);
      lazy[i].reset();
  i_t gry(int 1, int r) {
    if(l >= r) return i_t();
    1 += n, r += n;
    push(1);
    push(r - 1);
    i_t lp, rp;
    for (; 1 < r; 1 /= 2, r /= 2) {
     if(1 & 1) lp = i_t(lp, tree[l++]);
      if(r & 1) rp = i_t(tree[--r], rp);
    return i_t(lp, rp);
  void upd(int 1, int r, lazy_cont 1c) {
    if(1 >= r) return;
    1 += n, r += n;
    push(1);
    push(r - 1);
    int 10 = 1, r0 = r;
    for (; 1 < r; 1 /= 2, r /= 2) {
     if(1 & 1) downUpd(1++, 1c);
      if(r & 1) downUpd(--r, lc);
    build(10);
    build(r0 - 1);
  void upd(int pos, e_t v) {
    pos += n;
    push (pos);
    tree[pos] = i_t(v);
    build(pos);
private:
  int n, h;
  std::vector<bool> dirty;
  std::vector<i_t> tree;
  std::vector<lazy_cont> lazy;
  void apply(int p, lazy_cont lc) {
    tree[p].apply(lc);
    if(p < n) {
      dirty[p] = true;
      lazy[p] += lc;
  void pushSingle(int p) {
    if(dirty[p]) {
      downUpd(p + p, lazy[p]);
downUpd(p + p + 1, lazy[p]);
      lazy[p].reset();
      dirty[p] = false;
  void push(int p) {
    for (int s = h; s > 0; s--) {
      pushSingle(p >> s);
  void downUpd(int p, lazy_cont lc) {
    if(tree[p].canBreak(lc)) {
    } else if(tree[p].canApply(lc)) {
      apply(p, lc);
    } else {
      pushSingle(p);
      downUpd(p + p, lc);
      downUpd(p + p + 1, lc);
```

```
tree[p] = i_t(tree[p + p], tree[p + p + 1]);
}
void build(int p) {
  for(p /= 2; p > 0; p /= 2) {
    tree[p] = i_t(tree[p + p], tree[p + p + 1]);
    if(dirty[p]) {
        tree[p].apply(lazy[p]);
    }
};
```

1.5 Seg Tree cm Lazy

```
int arr[ms], seq[4 * ms], lazy[4 * ms], n;
void build (int idx = 0, int l = 0, int r = n-1) {
 int mid = (1+r)/2;
  lazy[idx] = 0;
  if(1 == r) {
    seg[idx] = arr[1];
    return;
  build(2*idx+1, 1, mid); build(2*idx+2, mid+1, r);
  seq[idx] = seq[2*idx+1] + seq[2*idx+2]; // Merge
void apply(int idx, int 1, int r) {
  if(lazy[idx] && !canBreak) { // if not beats canBreak = false
    if(1 < r) {
      lazy[2*idx+1] += lazy[idx]; // Merge de lazy
      lazy[2*idx+2] += lazy[idx]; // Merge de lazy
    if(canApply) { // if not beats canApply = true
      seg[idx] += lazy[idx] * (r - 1 + 1); // Aplicar lazy no seg
    } else {
      apply (2*idx+1, 1, mid); apply (2*idx+2, mid+1, r);
      seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
  lazy[idx] = 0; // Limpar a lazy
int query(int L, int R, int idx = 0, int l = 0, int r = n-1) {
 int mid = (1+r)/2;
 apply(idx, l, r);
  if(1 > R | | r < L) return 0; // Valor que nao atrapalhe</pre>
  if(L <= 1 && r <= R) return seg[idx];</pre>
  return query(L, R, 2*idx+1, 1, mid) + query(L, R, 2*idx+2, mid+1, r); //
void update(int L, int R, int V, int idx = 0, int 1 = 0, int r = n-1) {
  int mid = (1+r)/2;
  apply(idx, l, r);
  if(1 > R \mid \mid r < L) return;
  if(L <= 1 && r <= R)
    lazy[idx] = V;
    apply(idx, l, r);
   return:
  update(L, R, V, 2*idx+1, 1, mid); update(L, R, V, 2*idx+2, mid+1, r);
  seg[idx] = seg[2*idx+1] + seg[2*idx+2]; // Merge
```

1.6 Persistent Segment Tree

```
struct Node{
        int v = 0;
        Node *l = this, *r = this;
};
int CNT = 1:
Node buffer[ms * 20];
Node* update(Node *root, int 1, int r, int idx, int val) {
        Node *node = &buffer[CNT++];
        *node = *root;
        int mid = (1 + r) / 2;
        node->v += val;
        if(1+1 != r) {
                if(idx < mid) node->1 = update(root->1, 1, mid, idx, val);
                else node->r = update(root->r, mid, r, idx, val);
        return node;
int query(Node *node, int tl, int tr, int l, int r){
        if(l <= tl && tr <= r) return node->v;
        if(tr <= 1 || t1 >= r) return 0;
        int mid = (t1+tr) / 2;
        return query (node->1, tl, mid, l, r) + query (node->r, mid, tr, l, r
            );
```

1.7 Treap

```
mt19937 rng ((int) chrono::steady_clock::now().time_since_epoch().count());
typedef int Value;
typedef struct item * pitem;
struct item {
  item () {}
  item (Value v) { // add key if not implicit
    value = v:
    prio = uniform int distribution<int>() (rng);
   cnt = 1;
   rev = 0;
    1 = r = 0;
  pitem 1, r;
  Value value:
  int prio, cnt:
 bool rev;
int cnt (pitem it) {
  return it ? it->cnt : 0;
void fix (pitem it) {
  if (it)
    it->cnt = cnt(it->1) + cnt(it->r) + 1;
void pushLazy (pitem it) {
 if (it && it->rev) {
    it->rev = false:
    swap(it->l, it->r);
    if (it->1) it->1->rev ^= true;
    if (it->r) it->r->rev ^= true;
void insert (pitem & t, pitem it) {
 if (!t)
   t = it:
  else if (it->prio > t->prio)
    split (t, it->key, it->l, it->r), t = it;
  else
    insert (t->key <= it->key ? t->r : t->l, it);
void merge (pitem & t, pitem 1, pitem r) {
 pushLazy (1); pushLazy (r);
```

```
if (!1 || !r) t = 1 ? 1 : r;
  else if (1->prio > r->prio)
   merge (1->r, 1->r, r), t = 1;
  else
   merge (r->1, 1, r->1), t = r;
  fix (t);
void erase (pitem & t, int key) {
  if (t->key == key) {
    pitem th = t;
    merge (t, t->1, t->r);
    delete th:
    erase (key < t->key ? t->1 : t->r, key);
void split (pitem t, pitem & l, pitem & r, int key) {
  if (!t) return void(l = r = 0);
  pushLazy (t);
  int cur_key = cnt(t->1); // t->key if not implicit
  if (key <= cur_key)</pre>
    split (t->1, 1, t->1, key), r = t;
    split (t->r, t->r, r, key - (1 + cnt(t->1))), 1 = t;
  fix (t);
void reverse (pitem t, int 1, int r) {
  pitem t1, t2, t3;
  split (t, t1, t2, 1);
  split (t2, t2, t3, r-l+1);
  t2 \rightarrow rev = true:
  merge (t, t1, t2);
  merge (t, t, t3);
void unite (pitem & t, pitem l, pitem r) {
  if (!1 || !r) return void ( t = 1 ? 1 : r );
  if (l->prio < r->prio) swap (l, r);
  pitem lt, rt;
  split (r, lt, rt, l->key);
 unite (1->1, 1->1, 1t);
 unite (1-> r, 1->r, rt);
 t = 1;
pitem kth_element(pitem t, int k) {
        if(!t) return NULL;
        if(t->1) {
                if(t->l->size >= k) return kth element(t->l, k);
                else k -= t->1->cnt;
        return (k == 1) ? t : kth_element(t->r, k - 1);
int countLeft(pitem t, int key) {
        if(!t) {
                return 0;
        } else if(t->key < key) {</pre>
                return 1 + (t->1 ? t->1->size : 0) + countLeft(t->r, key);
        } else
                return countLeft(t->1, key);
```

1.8 Persistent Treap

```
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
typedef int Key;
struct Treap {
   Treap(){}
   Treap(char k) {
     key = 1;
     size = 1;
}
```

```
1 = r = NULL;
    val = k;
  Treap *1, *r;
  Key key;
  char val;
  int size;
typedef Treap * PTreap;
bool leftSide(PTreap 1, PTreap r) {
  return (int) (rng() % (l->size + r->size)) < l->size;
void fix(PTreap t) {
  if (t == NULL) {
    return;
  t->size = 1;
  t->key = 1;
  if (t->1) {
    t \rightarrow size += t \rightarrow l \rightarrow size;
    t\rightarrow key += t\rightarrow l\rightarrow size;
  if (t->r) {
    t->size += t->r->size;
void split(PTreap t, Key key, PTreap &1, PTreap &r) {
  if (t == NULL) {
    1 = r = NULL;
  } else if (t->key <= key) {</pre>
    1 = new Treap();
    *1 = *t;
    split(t->r, key - t->key, l->r, r);
    fix(1);
  } else
    r = new Treap();
    *r = *t;
    split(t->1, key, l, r->1);
    fix(r);
void merge(PTreap &t, PTreap 1, PTreap r) {
  if (!l || !r) {
    t = 1 ? 1 : r;
    return;
  t = new Treap();
  if (leftSide(l, r)) {
    *t = *1;
    merge (t->r, 1->r, r);
  } else {
    merge(t->1, 1, r->1);
  fix(t);
vector<PTreap> ver = {NULL};
PTreap build(int 1, int r, string& s) {
  if (1 >= r) return NULL;
  int mid = (1 + r) >> 1;
  auto ans = new Treap(s[mid]);
  ans->1 = build(1, mid, s);
  ans->r = build(mid + 1, r, s);
  fix(ans);
  return ans;
int last = 0;
void go(PTreap t, int f) {
  if (!t) return;
  go(t->1, f);
cout << t->val;
  last += (t->val == 'c') * f;
  go(t->r, f);
```

```
}
void insert(PTreap t, int pos, string& s) {
    PTreap 1, r;
    split(t, pos + 1, l, r);
    PTreap mid = build(0, s.size(), s);
    merge(mid, l, mid);
    merge(mid, mid, r);
    ver.push_back(mid);
}
void erase(PTreap t, int L, int R) {
    PTreap 1, mid, r;
    split(t, L, l, mid);
    split(mid, R - L + 1, mid, r);
    merge(l, l, r);
    ver.push_back(l);
}
```

1.9 KD-Tree

```
int d:
long long getValue(const PT &a) {return (d & 1) == 0 ? a.x : a.y; }
bool comp (const PT &a, const PT &b) {
 if((d & 1) == 0) { return a.x < b.x; }</pre>
  else { return a.y < b.y; }</pre>
long long sqrDist(PT a, PT b) { return (a - b) * (a - b); }
class KD_Tree {
public:
  struct Node {
    PT point;
    Node *left, *right;
  void init(std::vector<PT> pts)
    if(pts.size() == 0) {
      return;
    int n = 0;
    tree.resize(2 * pts.size());
    build(pts.begin(), pts.end(), n);
  long long nearestNeighbor(PT point) {
    long long ans = (long long) le18;
    nearestNeighbor(&tree[0], point, 0, ans);
    return ans;
private:
  std::vector<Node> tree;
  Node* build(std::vector<PT>::iterator 1, std::vector<PT>::iterator r, int
       &n, int h = 0) {
    int id = n++;
    if(r - 1 == 1) {
      tree[id].left = tree[id].right = NULL;
      tree[id].point = *1;
    else if(r - 1 > 1)
      std::vector\langle PT \rangle::iterator mid = 1 + ((r - 1) / 2);
      d = h;
      std::nth_element(l, mid - 1, r, comp);
      tree[id].point = *(mid - 1);
      // BE CAREFUL!
      // DO EVERYTHING BEFORE BUILDING THE LOWER PART!
      tree[id].left = build(1, mid, n, h^1);
      tree[id].right = build(mid, r, n, h^1);
    return &tree[id];
  void nearestNeighbor(Node* node, PT point, int h, long long &ans) {
    if(!node) {
      return;
```

```
if(point != node->point) {
    // THIS WAS FOR A PROBLEM
    // THAT YOU DON'T CONSIDER THE DISTANCE TO ITSELF!
    ans = std::min(ans, sqrDist(point, node->point));
}
d = h;
long long delta = getValue(point) - getValue(node->point);
if(delta <= 0) {
    nearestNeighbor(node->left, point, h^1, ans);
    if(ans > delta * delta) {
        nearestNeighbor(node->right, point, h^1, ans);
    }
} else {
    nearestNeighbor(node->right, point, h^1, ans);
    if(ans > delta * delta) {
        nearestNeighbor(node->left, point, h^1, ans);
    }
};
```

1.10 Sparse Table

```
vector<vector<int>> table;
vector<int> lg2;
void build(int n, vector<int> v) {
  lg2.resize(n + 1);
  lq2[1] = 0;
  for (int i = 2; i \le n; i++) {
    lg2[i] = lg2[i >> 1] + 1;
  table.resize(lg2[n] + 1);
  for (int i = 0; i < lg2[n] + 1; i++) {
    table[i].resize(n + 1);
  for (int i = 0; i < n; i++) {
    table[0][i] = v[i];
  for (int i = 0; i < lg2[n]; i++) {</pre>
    for (int j = 0; j < n; j++) {
      if (j + (1 << i) >= n) break;
      table[i + 1][j] = min(table[i][j], table[i][j + (1 << i)]);
int get(int 1, int r) {
 int k = lg2[r - l + 1];
 return min(table[k][1], table[k][r - (1 << k) + 1]);</pre>
```

1.11 Max Queue

```
template <class T, class C = less<T>>
struct MaxQueue {
   MaxQueue() { clear(); }
   void clear() {
      id = 0;
      q.clear();
   }
   void push(T x) {
      pair<int, T> nxt(1, x);
      while(q.size() > id && cmp(q.back().second, x)) {
         nxt.first += q.back().first;
        q.pop_back();
    }
   q.push_back(nxt);
}
```

```
T qry() { return q[id].second;}
void pop() {
    q[id].first--;
    if(q[id].first == 0) { id++; }
}
private:
    vector<std::pair<int, T>> q;
    int id;
    C cmp;
};
```

1.12 Policy Based Structures

1.13 Color Updates Structure

```
struct range {
  int 1, r;
  range(int l = 0, int r = 0, int v = 0) : l(l), r(r), v(v) {}
  bool operator < (const range &a) const {
    return 1 < a.1;</pre>
set<range> ranges;
vector<range> update(int 1, int r, int v) { // [1, r)
  vector<range> ans;
 if(l >= r) return ans;
  auto it = ranges.lower_bound(1);
  if(it != ranges.begin()) {
    it--;
    if(it->r>1) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, 1, cur.v));
      ranges.insert(range(l, cur.r, cur.v));
  it = ranges.lower_bound(r);
  if(it != ranges.begin()) {
    it--:
    if(it->r>r) {
      auto cur = *it;
      ranges.erase(it);
      ranges.insert(range(cur.1, r, cur.v));
      ranges.insert(range(r, cur.r, cur.v));
  for(it = ranges.lower_bound(l); it != ranges.end() && it->l < r; it++) {</pre>
    ans.push_back(*it);
  ranges.erase(ranges.lower_bound(1), ranges.lower_bound(r));
  ranges.insert(range(l, r, v));
  return ans;
int query(int v) { // Substituir -1 por flag para quando nao houver
```

```
auto it = ranges.upper_bound(v);
if(it == ranges.begin()) {
   return -1;
}
it--;
return it->r >= v ? it->v : -1;
```

2 Graph Algorithms

2.1 Simple Disjoint Set

```
struct dsu {
  vector<int> hist, par, sz;
  vector<ii> changes;
  int n:
  dsu (int n) : n(n) {
   hist.assign(n, 1e9);
    par.resize(n);
    iota(par.begin(), par.end(), 0);
    sz.assign(n, 1);
  int root (int x, int t) {
    if(hist[x] > t) return x;
    return root(par[x], t);
  void join (int a, int b, int t) {
    a = root(a, t);
    b = root(b, t);
    if (a == b) { changes.emplace_back(-1, -1); return; }
    if (sz[a] > sz[b]) swap(a, b);
    sz[b] += sz[a];
    hist[a] = t;
    changes.emplace_back(a, b);
 bool same (int a, int b, int t) {
    return root(a, t) == root(b, t);
  void undo () {
    int a, b;
    tie(a, b) = changes.back();
    changes.pop_back();
    if (a == -1) return;
    sz[b] -= sz[a];
    par[a] = a;
    hist[a] = 1e9;
    n++;
  int when (int a, int b) {
    while (1) {
      if (hist[a] > hist[b]) swap(a, b);
      if (par[a] == b) return hist[a];
      if (hist[a] == 1e9) return 1e9;
      a = par[a];
};
```

```
queue<int> q;
#define OUTER(x) (label[x] >= 0)
void L(int x, int y, int nxy) {
  int join, v, r = first[x], s = first[y];
 if (r == s) { return; }
  nxy += n + 1;
  label[r] = label[s] = -nxy;
  while (1) {
    if (s != 0) { swap(r, s); }
    r = first[label[mate[r]]];
    if (label[r] != -nxy) { label[r] = -nxy; }
    else {
      join = r;
      break;
  v = first[x];
  while (v != join) {
    if (!OUTER(v)) { q.push(v); }
    label[v] = nxy;
    first[v] = join;
    v = first[label[mate[v]]];
  v = first[y];
  while (v != join) {
    if (!OUTER(v)) { q.push(v); }
    label[v] = nxy;
    first[v] = join;
    v = first[label[mate[v]]];
  for (int i = 0; i <= n; i++)</pre>
    if (OUTER(i) && OUTER(first[i])) { first[i] = join; }
void R(int v, int w) {
  int t = mate[v];
  mate[v] = w;
  if (mate[t] != v) { return; ]
  if (label[v] >= 1 && label[v] <= n) {</pre>
    mate[t] = label[v];
    R(label[v], t);
   return;
  int x = from[label[v] - n - 1], y = to[label[v] - n - 1];
  R(x, y);
 R(y, x);
int E() {
 memset(mate, 0, sizeof(mate));
  int \mathbf{r} = 0:
  bool e7:
  for (int u = 1; u \le n; u++) {
    memset(label, -1, sizeof(label));
    while (!q.empty()) { q.pop(); }
    if (mate[u]) { continue; }
    label[u] = first[u] = 0;
    q.push(u);
    e7 = false;
    while (!q.empty() && !e7) {
      int x = q.front();
      q.pop();
      for (int i = 0; i < nadj[x]; i++) {</pre>
        int y = from[adj[x][i]];
        if (y == x) { y = to[adj[x][i]]; }
        if (!mate[y] && y != u) {
```

#define MAXN 110

int n, m;

#define MAXM MAXN * MAXN

int mate[MAXN], first[MAXN], label[MAXN];

int adj[MAXN] [MAXN], nadj[MAXN], from[MAXM], to[MAXM];

```
mate[y] = x;
          R(x, y);
          r++;
e7 = true;
          break;
        } else if (OUTER(y)) { L(x, y, adj[x][i]); }
        else {
          int v = mate[y];
          if (!OUTER(v)) {
           label[v] = x;
            first[v] = y;
            q.push(v);
        }
    label[0] = -1;
  return r;
/*Exemplo simples de uso*/
memset(nadj, 0, sizeof nadj);
for (int i = 0; i < m; ++i) { // arestas
 scanf("%d%d", &a, &b);
  a++, b++; // nao utilizar o vertice 0
  adj[a][nadj[a]++] = i;
  adi[b][nadi[b]++] = i;
  from[i] = a;
 to[i] = b;
printf("O emparelhamento tem tamanho %d\n", E());
for (int i = 1; i <= n; i++) {</pre>
 if (mate[i] > i) { printf("%d com %d\n", i - 1, mate[i] - 1); }
```

2.3 Boruvka

```
struct edge {
 int u, v;
 int w;
 int id;
 edge () {};
  edge (int u, int v, int w = 0, int id = 0) : u(u), v(v), w(w), id(id) {};
 bool operator < (edge &other) const { return w < other.w; };</pre>
vector<edge> boruvka (vector<edge> &edges, int n) {
  vector<edge> mst;
 vector<edge> best(n);
  initDSU(n):
 bool f = 1;
  while (f) {
   f = 0;
    for (int i = 0; i < n; i++) best[i] = edge(i, i, inf);</pre>
    for (auto e : edges) {
      int pu = root(e.u), pv = root(e.v);
      if (pu == pv) continue;
      if (e < best[pu]) best[pu] = e;</pre>
      if (e < best[pv]) best[pv] = e;</pre>
    for (int i = 0; i < n; i++) {
      edge e = best[root(i)];
      if (e.w != inf) {
        join(e.u, e.v);
        mst.push_back(e);
        f = 1;
  return mst;
```

2.4 Dinic Max Flow

```
const int ms = 1e3; // vertices
const int me = 1e5; // arestas
int adj[ms], to[me], ant[me], wt[me], z, n;
int copy adi[ms], fila[ms], level[ms];
void clear() { // Lembrar de chamar no main
 memset(adj, -1, sizeof adj);
void add(int u, int v, int k) {
 to[z] = v;
  ant[z] = adj[u];
 wt[z] = k;
  adj[u] = z++;
  swap(u, v);
  to[z] = v
 ant[z] = adj[u];
wt[z] = 0; // Lembrar de colocar = 0
 adj[u] = z++;
int bfs(int source, int sink) {
 memset(level, -1, sizeof level);
  level[source] = 0;
  int front = 0, size = 0, v;
  fila[size++] = source;
  while(front < size) {</pre>
    v = fila[front++];
    for(int i = adj[v]; i != -1; i = ant[i]) {
      if(wt[i] && level[to[i]] == -1) {
        level[to[i]] = level[v] + 1;
        fila[size++] = to[i];
  return level[sink] != -1;
int dfs(int v, int sink, int flow) {
  if(v == sink) return flow;
  for(int &i = copy_adj[v]; i != -1; i = ant[i]) {
    if(wt[i] && level[to[i]] == level[v] + 1 &&
      (f = dfs(to[i], sink, min(flow, wt[i])))) {
      wt[i] -= f;
      wt[i ^ 1] += f;
      return f:
  return 0;
int maxflow(int source, int sink) {
 int ret = 0, flow;
  while(bfs(source, sink)) {
    memcpy(copy_adj, adj, sizeof adj);
    while((flow = dfs(source, sink, 1 << 30))) {</pre>
     ret += flow:
  return ret;
```

2.5 Min Cost Max Flow

```
template <class T = int>
class MCMF {
public:
    struct Edge {
```

```
Edge(int a, T b, T c) : to(a), cap(b), cost(c) {}
   T cap, cost;
 MCMF(int size) {
   n = size;
    edges.resize(n);
    pot.assign(n, 0);
    dist.resize(n);
   visit.assign(n, false);
  pair<T, T> mcmf(int src, int sink) {
    pair<T, T > ans(0, 0);
    if(!SPFA(src, sink)) return ans;
    fixPot();
    // can use dijkstra to speed up depending on the graph
    while(SPFA(src, sink)) {
      auto flow = augment(src, sink);
      ans.first += flow.first;
     ans.second += flow.first * flow.second;
      fixPot();
    return ans;
 void addEdge(int from, int to, T cap, T cost) {
    edges[from].push_back(list.size());
    list.push_back(Edge(to, cap, cost));
    edges[to].push_back(list.size());
    list.push_back(Edge(from, 0, -cost));
private:
 int n;
 vector<vector<int>> edges;
 vector<Edge> list;
 vector<int> from;
  vector<T> dist, pot;
  vector<bool> visit;
  /*bool dij(int src, int sink) {
   T INF = std::numeric_limits<T>::max();
   dist.assign(n, INF);
    from.assign(n, -1);
    visit.assign(n, false);
    dist[src] = 0;
    for (int i = 0; i < n; i++) {
      int best = -1;
      for (int j = 0; j < n; j++) {
        if(visit[j]) continue;
        if(best == -1 \mid \mid dist[best] > dist[j]) best = j;
      if(dist[best] >= INF) break;
      visit[best] = true;
      for(auto e : edges[best]) {
        auto ed = list[e];
        if (ed.cap == 0) continue;
        T toDist = dist[best] + ed.cost + pot[best] - pot[ed.to];
        assert(toDist >= dist[best]);
        if(toDist < dist[ed.to]) {
         dist[ed.to] = toDist;
          from[ed.to] = e;
    return dist[sink] < INF;
  pair<T, T> augment(int src, int sink) {
    pair<T, T> flow = {list[from[sink]].cap, 0};
    for(int v = sink; v != src; v = list[from[v]^1].to) {
      flow.first = min(flow.first, list[from[v]].cap);
```

```
flow.second += list[from[v]].cost;
    for(int v = sink; v != src; v = list[from[v]^1].to) {
      list[from[v]].cap -= flow.first;
      list[from[v]^1].cap += flow.first;
    return flow;
  queue<int> q;
  bool SPFA(int src, int sink) {
   T INF = numeric_limits<T>::max();
    dist.assign(n, INF);
    from.assign(n, -1);
    q.push(src);
    dist[src] = 0;
    while(!q.empty())
      int on = q.front();
      q.pop();
      visit[on] = false;
      for(auto e : edges[on]) {
        auto ed = list[e];
        if(ed.cap == 0) continue;
        T toDist = dist[on] + ed.cost + pot[on] - pot[ed.to];
        if(toDist < dist[ed.to]) {</pre>
          dist[ed.to] = toDist;
          from[ed.to] = e;
          if(!visit[ed.to]) {
            visit[ed.to] = true;
            q.push(ed.to);
    return dist[sink] < INF;</pre>
  void fixPot() {
    T INF = numeric_limits<T>::max();
    for (int i = 0; i < n; i++) {
      if(dist[i] < INF) pot[i] += dist[i];</pre>
};
```

2.6 Euler Path and Circuit

```
int pathV[me], szV, del[me], pathE, szE;
int adj[ms], to[me], ant[me], wt[me], z, n;
// Funcao de add e clear no dinic
void eulerPath(int u) {
  for(int i = adj[u]; ^i; i = ant[u]) if(!del[i]) {
    del[i] = del[i^1] = 1;
    eulerPath(to[i]);
    pathE[szE++] = i;
  }
  pathV[szV++] = u;
}
```

2.7 Articulation Points/Bridges/Biconnected Components

```
int adj[ms], to[me], ant[me], z;
int num[ms], low[ms], timer;
bool art[ms], bridge[me], f[me];
int bc[ms], nbc;
stack<int> st, stk;
vector<vector<int>> comps;
```

```
void clear() { // Lembrar de chamar no main
 memset(adj, -1, sizeof adj);
 z = 0;
void add(int u, int v) {
 to[z] = v;
  ant[z] = adj[u];
 adj[u] = z++;
void generateBc (int v) {
  while (!st.emptv()) {
    int u = st.top();
    st.pop();
    bc[u] = nbc;
    if (v == u) break;
  ++nbc;
void dfs (int v, int p) {
  st.push(v), stk.push(v);
  low[v] = num[v] = ++timer;
  for (int i = adj[v]; i != -1; i = ant[i]) {
    if (f[i] || f[i^1]) continue;
    f[i] = 1;
    int u = to[i];
    if (num[u] == -1) {
      dfs(u, v);
      low[v] = min(low[v], low[u]);
      if (low[u] > num[v]) bridge[i] = bridge[i^1] = 1;
      if (low[u] >= num[v]) {
        art[v] = (num[v] > 1 || num[u] > 2);
        comps.push_back({v});
        while (comps.back().back() != u)
          comps.back().push_back(stk.top()), stk.pop();
      low[v] = min(low[v], num[u]);
  if (low[v] == num[v]) generateBc(v);
void biCon (int n) {
 nbc = 0, timer = 0;
 memset(num, -1, sizeof num);
 memset (bc, -1, sizeof bc);
 memset (bridge, 0, sizeof bridge);
 memset(art, 0, sizeof art);
  memset(f, 0, sizeof f);
  for (int i = 0; i < n; i++) {
   if (num[i] == -1) {
     timer = 0;
      dfs(i, 0);
vector<int> g[ms];
int id[ms];
void buildBlockCut (int n) {
  int z = 0;
  for (int u = 0; u < n; ++u) {
    if (art[u]) id[u] = z++;
  for (auto &comp : comps) {
    int node = z++;
    for (int u : comp) {
     if (!art[u]) id[u] = node;
```

```
g[node].push_back(id[u]);
    g[id[u]].push_back(node);
    }
}
```

2.8 SCC - Strongly Connected Components / 2SAT

```
const int ms = 212345;
vector<int> q[ms];
int idx[ms], low[ms], z, comp[ms], ncomp;
stack<int> st;
int dfs(int u) {
 if(~idx[u]) return idx[u] ? idx[u] : z;
  low[u] = idx[u] = z++;
  st.push(u);
  for(int v : g[u]) {
   low[u] = min(low[u], dfs(v));
  if(low[u] == idx[u]) {
    while(st.top() != u) {
      int v = st.top();
      idx[v] = 0;
      low[v] = low[u];
      comp[v] = ncomp;
      st.pop();
    idx[st.top()] = 0;
    st.pop();
    comp[u] = ncomp++;
  return low[u];
bool solveSat(int n) {
  memset(idx, -1, sizeof idx);
  z = 1; ncomp = 0;
  for (int i = 0; i < 2*n; i++) dfs(i);
  for(int i = 0; i < 2*n; i++) if(comp[i] == comp[i^1]) return false;
  return true;
int trad(int v) { return v < 0 ?(~v) *2^1 : v * 2; }</pre>
void add(int a, int b) { g[trad(a)].push_back(trad(b)); }
void addOr(int a, int b) { add(~a, b); add(~b, a); }
void addImp(int a, int b) { addOr(~a, b); }
void addEqual(int a, int b) { addOr(a, ~b); addOr(~a, b); }
void addDiff(int a, int b) { addEqual(a, ~b); }
// value[i] = comp[trad(i)] < comp[trad(~id)];</pre>
```

2.9 LCA - Lowest Common Ancestor

```
int par[ms] [mlg+1], lvl[ms];
void dfs(int v, int p, int l = 0) { // chamar como dfs(root, root)
    lvl[v] = 1;
    par[v][0] = p;
    for(int k = 1; k <= mlg; k++) {
        par[v][k] = par[par[v][k-1]][k-1];
    }
    for(int u : g[v]) {
        if(u != p) dfs(u, v, l + 1);
    }
} int lca(int a, int b) {
    if(lvl[b] > lvl[a]) swap(a, b);
    for(int i = mlg; i >= 0; i--) {
        if(lvl[a] - (1 << i) >= lvl[b]) a = par[a][i];
    }
```

```
if(a == b) return a;
for(int i = mlg; i >= 0; i--) {
   if(par[a][i] != par[b][i]) a = par[a][i], b = par[b][i];
}
return par[a][0];
}
```

2.10 LCA O(1)

```
template<class T>
struct RMQ {
  vector<vector<T>> jmp;
  RMQ(const vector<T>& V) : jmp(1, V) {
    for (int pw = 1, k = 1; pw * 2 <= (int) size(V); pw *= 2, ++k) {
      jmp.emplace_back(size(V) - pw * 2 + 1);
      for (int j = 0; j < (int) size(jmp[k]); ++j)
        jmp[k][j] = min(jmp[k - 1][j], jmp[k - 1][j + pw]);
  T query(int a, int b) {
    assert(a < b); // or return inf if a == b</pre>
    int dep = 31 - __builtin_clz(b - a);
    return min(jmp[dep][a], jmp[dep][b - (1 << dep)]);</pre>
};
struct LCA {
  int T = 0;
  vector<int> time, path, ret;
 RMO<int> rmg;
  LCA(vector<vector<int>>& C) : time(size(C)), rmq((dfs(C,0,-1), ret)) {}
  void dfs(vector<vector<int>>& C, int v, int par) {
    time[v] = T++;
    for (int y : C[v]) if (y != par) {
      path.push_back(v), ret.push_back(time[v]);
      dfs(C, y, v);
  int lca(int a, int b) {
    if (a == b) return a;
    tie(a, b) = minmax(time[a], time[b]);
    return path[rmq.query(a, b)];
};
```

2.11 Heavy Light Decomposition

```
class HLD {
public:
  void init(int n) { /* resize everything */ }
  void addEdge(int u, int v) {
    edges[u].push_back(v);
    edges[v].push_back(u);
  void setRoot(int r) {
    t = 0;
    p[r] = r;
    h[r] = 0;
    prep(r, r);
    nxt[r] = r;
    hld(r);
  int getLCA(int u, int v) {
    while(!inSubtree(nxt[u], v)) u = p[nxt[u]];
    while(!inSubtree(nxt[v], u)) v = p[nxt[v]];
```

```
return in[u] < in[v] ? u : v;</pre>
  // is v in the subtree of u?
  bool inSubtree(int u, int v) {
    return in[u] <= in[v] && in[v] < out[u];</pre>
  // returns ranges [1, r) that the path has
  vector<pair<int, int>> getPath(int u, int anc) {
    vector<std::pair<int, int>> ans;
    //assert(inSubtree(anc, u));
    while(nxt[u] != nxt[anc]) {
      ans.emplace_back(in[nxt[u]], in[u] + 1);
      u = p[nxt[u]];
    // this includes the ancestor! care
    ans.emplace_back(in[anc], in[u] + 1);
    return ans;
private:
  vector<int> in, out, p, rin, sz, nxt, h;
  vector<vector<int>> edges;
  void prep(int on, int par) {
    sz[on] = 1;
    p[on] = par;
    for(int i = 0; i < (int) edges[on].size(); i++) {</pre>
      int &u = edges[on][i];
      if(u == par) {
        swap(u, edges[on].back());
        edges[on].pop_back();
      } else {
        h[u] = 1 + h[on];
        prep(u, on);
        sz[on] += sz[u];
        if(sz[u] > sz[edges[on][0]]) {
          swap(edges[on][0], u);
  void hld(int on)
    in[on] = t++;
    rin[in[on]] = on;
    for(auto u : edges[on]) {
      nxt[u] = (u == edges[on][0] ? nxt[on] : u);
      hld(u);
    out[on] = t;
};
```

2.12 Centroid Decomposition

```
template<typename T>
struct CentroidDecomposition {
  vector<int> sz, h, dad;
  vector<vector<pair<int, T>>> adj;
  vector<vector<T>> dis;
  vector<bool> removed;
  CentroidDecomposition (int n) {
    sz.resize(n);
    h.resize(n);
    dis.resize(n);
    removed.resize(n);
    removed.resize(n, 0);
    dad.resize(n);
}
void add (int a, int b, T w = 1) {
    adj[a].push_back({b, w});
```

```
adj[b].push_back({a, w});
void dfsSize (int v, int par) {
  sz[v] = 1;
  for (auto u : adj[v]) {
    if (u.x == par || removed[u.x]) continue;
    dfsSize(u.x, v);
    sz[v] += sz[u.x];
int getCentroid (int v, int par, int tam) {
  for (auto u : adj[v]) {
    if (u.x == par || removed[u.x]) continue;
    if ((sz[u.x]<<1) > tam) return getCentroid(u.x, v, tam);
  return v;
void setDis (int v, int par, int nv) {
  for (auto u : adj[v]) {
    if (u.x == par || removed[u.x]) continue;
    dis[u.x][nv] = dis[v][nv]+u.y;
    setDis(u.x, v, nv);
void decompose (int v, int par = -1, int nv = 0) {
  dfsSize(v, par);
  int c = getCentroid(v, par, sz[v]);
dad[c] = par;
removed[c] = 1;
  h[c] = nv;
  setDis(c, par, nv);
  for (auto u : adj[c]) {
    if (!removed[u.x]){
      decompose(u.x, c, nv + 1);
int operator [] (const int idx) const {
  return dad[idx];
T dist (int u, int v) {
  if (h[u] < h[v]) swap(u, v);
  return dis[u][h[v]];
```

2.13 Sack

```
void dfs(int v, int par = -1, bool keep = 0) {
   int big = -1;
   for (int u : adj[v]) {
      if (u == par) continue;
      if (big == -1 || sz[u] > sz[big]) {
            big = u;
      }
   }
   for (int u : adj[v]) {
        if (u == par || u == big) {
            continue;
      }
      dfs(u, v, 0);
   }
   if (big != -1) {
        dfs(big, v, 1);
   }
   for (int u : adj[v]) {
      if (u == par || u == big) {
            continue;
      }
}
```

```
put(u, v);
}
if (!keep) {
}
```

2.14 Hungarian Algorithm - Maximum Cost Matching

```
int u[ms], v[ms], p[ms], way[ms], minv[ms];
bool used[ms];
pair<vector<int>, int> solve(const vector<vector<int>> &matrix) {
  int n = matrix.size();
  if(n == 0) return {vector<int>(), 0};
  int m = matrix[0].size();
  assert (n <= m);
  memset(u, 0, (n+1)*sizeof(int));
  memset(v, 0, (m+1)*sizeof(int));
  memset(p, 0, (m+1)*sizeof(int));
  for(int i = 1; i <= n; i++) {</pre>
    memset(minv, 0x3f, (m+1)*sizeof(int));
    memset (way, 0, (m+1) *sizeof(int));
    for (int j = 0; j <= m; j++) used[j] = 0;</pre>
    p[0] = i;
    int k0 = 0;
    do {
      used[k0] = 1;
      int i0 = p[k0], delta = inf, k1;
      for (int j = 1; j \le m; j++) {
        if(!used[j]) {
          int cur = matrix[i0-1][j-1] - u[i0] - v[j];
          if (cur < minv[i]) {</pre>
            minv[i] = cur;
            way[j] = k0;
          if(minv[j] < delta) {</pre>
            delta = minv[j];
            k1 = j;
      for (int j = 0; j <= m; j++) {
        if(used[j]) {
          u[p[j]] += delta;
          v[j] -= delta;
        } else {
          minv[j] -= delta;
      k0 = k1;
    } while(p[k0]);
      int k1 = way[k0];
      p[k0] = p[k1];
      k0 = k1;
    } while(k0);
  vector<int> ans(n, -1);
  for(int j = 1; j <= m; j++) {</pre>
    if(!p[i]) continue;
    ans[p[j] - 1] = j - 1;
  return {ans, -v[0]};
```

2.15 Burunduk

```
struct edge {
 int a, b, 1, r;
typedef vector <edge> List;
int cnt[N + 1], ans[N], u[N], color[N], deg[N];
void add (int a, int b) {
 g[a].pb(b), g[b].pb(a);
void dfs (int v, int value)
 u[v] = 1, color[v] = value;
  forn(i, sz(q[v]))
    if (!u[g[v][i]])
      dfs(g[v][i], value);
int compress (List &v1, int vn, int &add_vn) {
  int vn1 = 0;
  forn (i, vn) u[i] = 0;
  forn (i, vn) {
    if (!u[i]) deg[vn1] = 0, dfs(i, vn1++);
  forn (i, sz(v1)) {
    v1[i].a = color[v1[i].a];
    v1[i].b = color[v1[i].b];
    if (v1[i].a != v1[i].b)
     deg[v1[i].a]++, deg[v1[i].b]++;
  vn = vn1, vn1 = 0;
  forn (i, vn) {
    u[i] = vn1, vn1 += (deg[i] > 0), add_vn += !deg[i];
  forn (i, sz(v1)) {
   v1[i].a = u[v1[i].a];
    v1[i].b = u[v1[i].b];
  return vn1;
void go (int 1, int r, const List &v, int vn, int add_vn) {
 if (cnt[1] == cnt[r]) return;
  if (!sz(v)){
    while (1 < r)
     ans[1++] = vn + add vn;
    return;
  List v1;
  forn (i, vn) {
    g[i].clear();
  forn (i, sz(v)) {
    if (v[i].a != v[i].b) {
     if (v[i].1 <= 1 && v[i].r >= r)
        add(v[i].a, v[i].b);
      else if (1 < v[i].r \&\& r > v[i].1)
        v1.pb(v[i]);
  int vn1 = compress(v1, vn, add_vn);
  int m = (1 + r) / 2;
  go(1, m, v1, vn1, add_vn);
 go(m, r, v1, vn1, add_vn);
```

2.16 Minimum Arborescence

```
// uncommented O(V^2) arborescence
struct Edges {
   //set<pair<long long, int>> cost; O(Elog^2)
   long long cost[ms];
   // possible optimization, use vector of size n
```

```
// instead of ms
  long long sum = 0;
  Edges() {
    memset(cost, 0x3f, sizeof cost);
  void addEdge(int u, long long c) {
    // cost.insert({c - sum, u}); O(Elog^2)
    cost[u] = min(cost[u], c - sum);
  pair<long long, int> getMin() {
    //return *cost.begin(); O(E*log^2)
    pair<long long, int> ans(cost[0], 0);
    // in this loop can change ms to n to make it faster for many cases
    for (int i = 1; i < ms; i++) {
      if(cost[i] < ans.first) {</pre>
        ans = pair<long long, int>(cost[i], i);
    return ans;
  void unite(Edges &e) {
    O(E*log^2E)
    if(e.cost.size() > cost.size()) {
     cost.swap(e.cost);
      swap(sum, e.sum);
    for(auto i : e.cost) {
      addEdge(i.second, i.first + e.sum);
    e.cost.clear();
    // O(V^2)
    // can change ms to n
    for (int i = 0; i < ms; i++) {
      cost[i] = min(cost[i], e.cost[i] + e.sum - sum);
typedef vector<vector<pair<long long, int>>> Graph;
Edges ed[ms]:
int par[ms];
long long best[ms];
int col[ms];
int getPar(int x) { return par[x] < 0 ? x : par[x] = getPar(par[x]); }</pre>
void makeUnion(int a, int b) {
  a = getPar(a);
 b = getPar(b);
  if(a == b) return;
  ed[a].unite(ed[b]);
  par[b] = a;
long long arborescence(Graph edges) {
 // root is 0
  // edges has transposed adjacency list (cost, from)
  // edge from i to j cost c is
  // edge[j].emplace_back(c, i)
  int n = (int) edges.size();
  long long ans = 0;
  for (int i = 0; i < n; i++) {
   ed[i] = Edges();
    par[i] = -1;
    for(auto j : edges[i]) {
     ed[i].addEdge(j.second, j.first);
    col[i] = 0;
  // to change the root you can simply change this next line to
  // col[root] = 2;
  col[0] = 2;
  for (int i = 0; i < n; i++) {
    if(col[getPar(i)] == 2) {
```

```
continue;
  int on = getPar(i);
 vector<int> st;
  while (col[on] != 2) {
    assert(getPar(on) == on);
    if(col[on] == 1) {
      int v = on;
      vector<int> cycle;
//cout << "found cycle\n";</pre>
      while(st.back() != v) {
        //cout << st.back() << endl;</pre>
        cycle.push_back(st.back());
        st.pop_back();
      for(auto u : cycle) { // compress cycle
        makeUnion(v, u);
      v = getPar(v);
      col[v] = 0;
      on = v;
    } else {
      // still no cycle
      // while best is in compressed cycle, remove
      // THIS IS TO MAKE O(E*log^2) ALGORITHM!!
      // while(!ed[on].cost.empty() && getPar(on) == getPar(ed[on].getMin
           ().second)) {
           ed[on].cost.erase(ed[on].cost.begin());
      11 }
      // O(V^2)
      for (int x = 0; x < n; x++) {
        if(on == getPar(x)) {
          ed[on].cost[x] = 0x3f3f3f3f3f3f3f3f3fLL;
      // best edge
      auto e = ed[on].getMin();
      // O(E*log^2) assert(!ed[on].cost.empty()) if every vertex appears
           in the arborescence
      // O(V^2)
      assert (e.first < 0x3f3f3f3f3f3f3f3f3f3fLL);
      int v = getPar(e.second);
      //cout << "found not cycle to " << v << " of cost " << e.first + ed
           [on].sum << '\n';
      assert (v != on);
      best[on] = e.first + ed[on].sum;
      ans += best[on];
      // compress edges
      ed[on].sum = -(e.first);
      st.push_back(on);
      col[on] = 1;
      on = v;
  // make everything 2
  for (auto u : st)
   assert (getPar(u) == u);
    col[u] = 2;
return ans;
```

2.17 Dominator Tree

```
struct dominator_tree {
  vector<basic_string<int>> g, rg, bucket;
  vector<int> arr, par, rev, sdom, dom, dsu, label;
  int n, t;
  dominator_tree(int n) : g(n), rg(n), bucket(n), arr(n, -1),
      par(n), rev(n), sdom(n), dom(n), dsu(n), label(n), n(n), t(0) {}
```

```
void add_edge(int u, int v) { g[u] += v; }
  void dfs(int u) {
    arr[u] = t;
    rev[t] = u;
    label[t] = sdom[t] = dsu[t] = t;
    for (int w : g[u]) {
      if (arr[w] == -1) {
        dfs(w);
        par[arr[w]] = arr[u];
      rg[arr[w]] += arr[u];
  int find(int u, int x=0) {
    if (u == dsu[u])
      return x ? -1 : u;
    int v = find(dsu[u], x+1);
    if (\mathbf{v} < 0)
      return u;
    if (sdom[label[dsu[u]]] < sdom[label[u]])</pre>
      label[u] = label[dsu[u]];
    dsu[u] = v;
    return x ? v : label[u];
  vector<int> run(int root) {
    dfs(root);
    iota(dom.begin(), dom.end(), 0);
    for (int i=t-1; i>=0; i--) {
      for (int w : rq[i])
        sdom[i] = min(sdom[i], sdom[find(w)]);
      if (i)
        bucket[sdom[i]] += i;
      for (int w : bucket[i]) {
        int v = find(w);
        if (sdom[v] == sdom[w])
          dom[w] = sdom[w];
        else
          dom[w] = v;
      if (i > 1)
        dsu[i] = par[i];
    for (int i=1; i<t; i++) {</pre>
      if (dom[i] != sdom[i])
        dom[i] = dom[dom[i]];
    vector<int> outside_dom(n);
    iota(begin(outside_dom), end(outside_dom), 0);
    for (int i=0; i<n; i++)</pre>
      outside_dom[rev[i]] = rev[dom[i]];
    return outside_dom;
};
```

2.18 Kuhn

```
int n, m;
vector<vector<int>> g;
vector<int> mt;
vector<bool> used;

bool try_kuhn(int v) {
   if (used[v]) return false;
   used[v] = true;
   for (int to : g[v]) {
      if (mt[to] == -1 || try_kuhn(mt[to])) {
       mt[to] = v;
      return true;
   }
}
```

2.19 Bipartite Check

```
void make_set(int v) {
  parent[v] = make pair(v, 0);
  h[v] = 0;
pair<int, int> find_set(int v) {
 if (v != parent[v].first) {
    int parity = parent[v].second;
    parent[v] = find_set(parent[v].first);
    parent[v].second ^= parity;
  return parent[v];
bool add_edge(int a, int b) {
  int x, y;
  tie(a, x) = find_set(a);
  tie(b, y) = find_set(b);
  if (a == b) {
    if (x == y)
      return false;
  } else {
    if (h[a] < h[b])
      swap (a, b);
    else if (h[a] == h[b])
     ++h[a];
    parent[b] = make_pair(a, x^y^1);
  return true;
```

2.20 Link Cut Tree Padrao

```
node() \{ p = ch[0] = ch[1] = -1; \}
};
node t[MAX];
bool is_root(int x) {
        return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].
            ch[1] != x);
void rotate(int x) {
        int p = t[x].p, pp = t[p].p;
        if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
        bool d = t[p].ch[0] == x;
        t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
        if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
        t[x].p = pp, t[p].p = x;
void splav(int x) {
        while (!is_root(x)) {
                int p = t[x].p, pp = t[p].p;
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch
                     [0] == x)? x : p);
                rotate(x);
int access(int v) {
        int last = -1;
        for (int w = v; w+1; last = w, splay(v), w = t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
int find_root(int v) {
        access(v);
        while (t[v].ch[0]+1) v = t[v].ch[0];
        return splay(v), v;
void link(int v, int w) { // v deve ser raiz
        access(v):
        t[v].p = w;
void cut(int v) { // remove aresta de v pro pai
        access(v);
        t[v].ch[0] = t[t[v].ch[0]].p = -1;
int lca(int v, int w) {
        return access(v), access(w);
```

2.21 Link Cut Tree Aresta

```
// by brunomaletta
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
// Todas as operacoes sao O(log(n)) amortizado
// 9ce48f
namespace lct {
         struct node {
                 int p, ch[2];
                 11 val, sub;
                 bool rev;
                 int sz, ar;
                 ll lazy;
                 node() {}
                 node(int v, int ar_) :
                 p(-1), val(v), sub(v), rev(0), sz(ar_), ar(ar_), lazy(0) {
    ch[0] = ch[1] = -1;
```

```
node t[2*MAX]; // MAXN + MAXQ
map<pair<int, int>, int> aresta;
int sz;
void prop(int x) {
        if (t[x].lazy) {
                if (t[x].ar) t[x].val += t[x].lazy;
                t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
        if (t[x].rev) {
                swap(t[x].ch[0], t[x].ch[1]);
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
        t[x].lazy = 0, t[x].rev = 0;
void update(int x) {
        t[x].sz = t[x].ar, t[x].sub = t[x].val;
        for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
                prop(t[x].ch[i]);
                t[x].sz += t[t[x].ch[i]].sz;
                t[x].sub += t[t[x].ch[i]].sub;
bool is root(int x) {
        return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].
void rotate(int x) {
        int p = t[x].p, pp = t[p].p;
        if (!is\_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
        bool d = t[p].ch[0] == x;
        t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
        if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
        t[x].p = pp, t[p].p = x;
        update(p), update(x);
int splay(int x) {
        while (!is_root(x)) {
                int p = t[x].p, pp = t[p].p;
                if (!is_root(p)) prop(pp);
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch
                    [0] == x)? x : p);
                rotate(x);
        return prop(x), x;
int access(int v) {
        int last = -1;
        for (int w = v; w+1; update(last = w), splay(v), w = t[v].p
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
        return last;
void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar); }
int find_root(int v) {
        access(v), prop(v);
        while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
        return splay(v);
bool conn(int v, int w) {
        access(v), access(w);
        return v == w ? true : t[v].p != -1;
void rootify(int v) {
        access(v);
        t[v].rev ^= 1;
```

```
11 query(int v, int w) {
        rootify(w), access(v);
        return t[v].sub;
void update(int v, int w, int x) {
        rootify(w), access(v);
        t[v].lazy += x;
void link_(int v, int w) {
        rootify(w);
        t[w].p = v;
void link(int v, int w, int x) { // v--w com peso x
        int id = MAX + sz++;
        aresta[make_pair(v, w)] = id;
        make_tree(id, x, 1);
        link_(v, id), link_(id, w);
void cut_(int v, int w) {
        rootify(w), access(v);
t[v].ch[0] = t[t[v].ch[0]].p = -1;
void cut(int v, int w) {
        int id = aresta[make_pair(v, w)];
        cut_(v, id), cut_(id, w);
int lca(int v, int w) {
        access(v);
        return access(w);
```

2.22 Link Cut Tree Vertice

```
// by brunomaletta
// Valores nos vertices
// make_tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
\begin{subarray}{lll} \end{subarray} // \end{subarray} rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
// Todas as operacoes sao O(\log(n)) amortizado
namespace lct {
        struct node {
                 int p, ch[2];
                 11 val, sub;
                 bool rev:
                 int sz;
                 ll lazy;
                 node() {}
                 node(int v) : p(-1), val(v), sub(v), rev(0), sz(1), lazy(0)
                          ch[0] = ch[1] = -1;
        };
        node t[MAX];
        void prop(int x) {
                 if (t[x].lazy) {
                          t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].
                          if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
                          if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
                 if (t[x].rev) {
                          swap(t[x].ch[0], t[x].ch[1]);
                          if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
```

```
if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
        t[x].lazy = 0, t[x].rev = 0;
void update(int x) {
        t[x].sz = 1, t[x].sub = t[x].val;
        for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
                prop(t[x].ch[i]);
                t[x].sz += t[t[x].ch[i]].sz;
                t[x].sub += t[t[x].ch[i]].sub;
bool is_root(int x) {
        return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].
            ch[1] != x);
void rotate(int x) {
        int p = t[x].p, pp = t[p].p;
        if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
bool d = t[p].ch[0] == x;
        t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
        if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
        t[x].p = pp, t[p].p = x;
        update(p), update(x);
int splay(int x) {
        while (!is_root(x)) {
                int p = t[x].p, pp = t[p].p;
                if (!is_root(p)) prop(pp);
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch
                     [0] == x) ? x : p);
                rotate(x);
        return prop(x), x;
int access(int v) {
        int last = -1;
        for (int w = v; w+1; update(last = w), splay(v), w = t[v].p
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
        return last;
void make_tree(int v, int w) { t[v] = node(w); }
int find_root(int v) {
        access(v), prop(v);
        while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
        return splay(v);
bool connected(int v, int w) {
        access(v), access(w);
        return v == w ? true : t[v].p != -1;
void rootify(int v) {
        access(v);
        t[v].rev ^= 1;
11 query(int v, int w) {
        rootify(w), access(v);
        return t[v].sub;
void update(int v, int w, int x) {
        rootify(w), access(v);
        t[v].lazy += x;
void link(int v, int w) {
        rootify(w);
        t[w].p = v;
void cut(int v, int w) {
        rootify(w), access(v);
        t[v].ch[0] = t[t[v].ch[0]].p = -1;
int lca(int v, int w) {
```

access(v);
return access(w);
}

3 Dynamic Programming

3.1 Line Container

```
bool O;
struct Line {
  mutable 11 k, m, p;
  bool operator<(const Line& o) const {</pre>
    return Q ? p < o.p : k < o.k;
struct LineContainer : multiset<Line> {
 // (for doubles, use inf = 1/.0, div(a,b) = a/b)
  const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = inf; return false; }
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x -> p = div(v -> m - x -> m, x -> k - v -> k);
    return x->p >= y->p;
  void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p)
      isect(x, erase(y));
  11 query(11 x) {
    assert(!empty());
    Q = 1; auto 1 = *lower_bound(\{0, 0, x\}); Q = 0;
    return 1.k * x + 1.m;
};
```

3.2 Li Chao Tree

```
// by luucasv
typedef long long T;
const T INF = 1e18, EPS = 1;
const int BUFFER_SIZE = 1e4;
struct Line {
 T m, b;
  Line (T m = 0, T b = INF) : m(m), b(b) {}
 T apply(T x) { return x * m + b; }
struct Node {
 Node *left, *right;
  Line line;
 Node(): left(NULL), right(NULL) {}
struct LiChaoTree {
 Node *root, buffer[BUFFER_SIZE];
 T min_value, max_value;
 int buffer_pointer;
 LiChaoTree (T min_value, T max_value): min_value (min_value), max_value (
      max_value + 1) { clear(); }
  void clear() { buffer_pointer = 0; root = newNode(); }
```

```
void insert_line(T m, T b) { update(root, min_value, max_value, Line(m, b 3.4 Knuth Optimization
      )); }
  T eval(T x) { return query(root, min_value, max_value, x); }
  void update(Node *cur, T 1, T r, Line line) {
    T m = 1 + (r - 1) / 2;
    bool left = line.apply(1) < cur->line.apply(1);
    bool mid = line.apply(m) < cur->line.apply(m);
    bool right = line.apply(r) < cur->line.apply(r);
    if (mid) {
     swap(cur->line, line);
    if (r - 1 <= EPS) return;</pre>
    if (left == right) return;
    if (mid != left) {
     if (cur->left == NULL) cur->left = newNode();
     update(cur->left, 1, m, line);
      if (cur->right == NULL) cur->right = newNode();
      update(cur->right, m, r, line);
  T query (Node *cur, T l, T r, T x) {
    if (cur == NULL) return INF;
    if (r - 1 <= EPS) {
     return cur->line.apply(x);
    T m = 1 + (r - 1) / 2;
    T ans;
    if (x < m) {
      ans = query(cur->left, 1, m, x);
    } else {
     ans = query(cur->right, m, r, x);
    return min(ans, cur->line.apply(x));
  Node* newNode() {
     buffer[buffer_pointer] = Node();
      return &buffer[buffer_pointer++];
};
```

3.3 Divide and Conquer Optimization

```
int n, k;
11 dpold[ms], dp[ms], c[ms][ms]; // c(i, j) pode ser funcao
void compute(int 1, int r, int opt1, int optr) {
    if(l>r) return;
    int mid = (1+r)/2;
    pair<11, int> best = {inf, -1}; // long long inf
    for(int k = optl; k <= min(mid, optr); k++) {</pre>
        best = min(best, \{dpold[k-1] + c[k][mid], k\});
    dp[mid] = best.first;
    int opt = best.second;
    compute(l, mid-1, optl, opt);
    compute(mid+1, r, opt, optr);
il solve() {
    dp[0] = 0;
    for(int i = 1; i <= n; i++) dp[i] = inf; // initialize row 0 of the dp</pre>
    for (int i = 1; i \le k; i++) {
        swap (dpold, dp);
        compute(0, n, 0, n); // solve row i of the dp
    return dp[n]; // return dp[k][n]
```

```
int n, m, mid[ms][ms];
11 dp[ms][ms];
void knuth() {
  for (int i = n; i >= 0; i--) { // limites entre 0 \in n
    dp[i][i+1] = 0; mid[i][i+1] = i; // caso base
    for (int j = i+2; j \le n; j++) {
      dp[i][j] = inf; // long long inf
      for(int k = mid[i][j-1]; k <= mid[i+1][j]; k++) {</pre>
        if(dp[i][j] > dp[i][k] + dp[k][j]) {
          dp[i][j] = dp[i][k] + dp[k][j];
          mid[i][j] = k;
      dp[i][j] += c(i, j); // custo associado ao intervalo
```

4 Math

4.1 Chinese Remainder Theorem

```
long long modinverse (long long a, long long b, long long s0 = 1, long long
   s1 = 0) {
  if(!b) return s0;
  else return modinverse(b, a % b, s1, s0 - s1 * (a / b));
long long gcd(long long a, long long b) {
 if(!b) return a;
  else return gcd(b, a % b);
11 mul(11 a, 11 b, 11 m) {
  11 q = (long double) a * (long double) b / (long double) m;
  11 r = a * b - q * m;
  return (r + 5 * m) % m;
long long safemod(long long a, long long m) {
  return (a % m + m) % m;
struct equation{
 equation(long long a, long long m) {mod = m, ans = a, valid = true;}
  equation() {valid = false;}
  equation (equation a, equation b) {
    if(!a.valid || !b.valid) {
      valid = false;
      return;
    long long g = gcd(a.mod, b.mod);
    if((a.ans - b.ans) % g != 0) {
      valid = false;
      return;
    valid = true;
    mod = a.mod * (b.mod / q);
    ans = a.ans +
      mul(a.mod, modinverse(a.mod, b.mod), mod),
      (b.ans - a.ans) / g
      , mod);
    ans = safemod(ans, mod);
  long long mod, ans;
  bool valid;
```

```
void print()
{
    if(!valid)
        std::cout << "equation is not valid\n";
    else
        std::cout << "equation is " << ans << " mod " << mod << '\n';
    }
};</pre>
```

4.2 Diophantine Equations

```
int gcd_ext(int a, int b, int& x, int &y) {
  if (b == 0) {
    x = 1, y = 0;
    return a;
  int nx, ny;
  int gc = gcd_ext(b, a % b, nx, ny);
 x = ny;
  y = nx' - (a / b) * ny;
 return qc;
vector<int> diophantine(int D, vector<int> 1) {
 int n = l.size();
  vector<int> gc(n), ans(n);
  qc[n-1] = \hat{1}[n-1];
  for (int i = n - 2; i >= 0; i--) {
    qc[i] = qcd_ext(l[i], qc[i + 1], x, y);
  if (D % qc[0] != 0) {
    return vector<int>();
  for (int i = 0; i < n; i++) {
    if (i == n - 1) {
      ans[i] = D / l[i];
      D = l[i] * ans[i];
      continue;
    gcd_ext(l[i] / gc[i], gc[i + 1] / gc[i], x, y);
    ans[i] = (long long int) D / gc[i] * x % (gc[i + 1] / gc[i]);
    if (D < 0 \&\& ans[i] > 0) {
      ans[i] -= (gc[i + 1] / gc[i]);
    if (D > 0 \&\& ans[i] < 0) {
      ans[i] += (gc[i + 1] / gc[i]);
    D = l[i] * ans[i];
  return ans:
```

4.3 Discrete Logarithm

```
11 discreteLog (11 a, 11 b, 11 m) {
    a %= m; b %= m;
    l1 n = (11) sqrt (m + .0) + 1, an = 1;
    for (11 i = 0; i < n; i++) {
        an = (an * a) % m;
    }
    map<11, 11> vals;
    for (11 i = 1, cur = an; i <= n; i++) {
        if (!vals.count(cur)) vals[cur] = i;
        cur = (cur * an) % m;</pre>
```

```
}
ll ans = lel8; //inf
for (ll i = 0, cur = b; i <= n; i++) {
   if (vals.count(cur)) {
      ans = min(ans, vals[cur] * n - i);
   }
   cur = (cur * a) % m;
}
return ans;</pre>
```

4.4 Discrete Root

```
//x^k = a % mod

ll discreteRoot(ll k, ll a, ll mod) {
    ll g = primitiveRoot(mod);
    ll y = discreteLog(fexp(g, k, mod), a, mod);
    if (y == -1) {
        return y;
    }
    return fexp(g, y, mod);
}
```

4.5 Division Trick

```
for(int l = 1, r; l <= n; l = r + 1) {
    r = n / (n / 1);
    // n / i has the same value for l <= i <= r
}</pre>
```

4.6 Modular Sum

```
//calcula (sum(0 <= i <= n) P(i)) % mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
  ll calc(ll a, ll b, ll n, ll mod) {
    assert (a&&b);
    if(a >= b) {
      11 ret = ((n*(n+1)/2) *mod) *(a/b);
      if(a%b) ret = (ret + calc(a%b,b,n,mod))%mod;
      else ret = (ret+n+1)%mod;
      return ret;
    return ((n+1)*(((n*a)/b+1)*mod) - calc(b,a,(n*a)/b,mod) + mod + n/b +
        1)%mod;
//P(i) = a*i \mod m
  11 solve(ll a, ll n, ll m, ll mod) {
    a = (a%m + m)%m;
    if(!a) return 0;
    11 ret = (n*(n+1)/2)%mod;
    ret = (ret*a)%mod;
    11 g = __gcd(a,m);
ret -= m*(calc(a/g, m/g, n, mod)-n-1);
    return (ret%mod + mod)%mod;
//P(i) = a + r * i \mod m
  11 solve(ll a, ll r, ll n, ll m, ll mod) {
    a = (a\%m + m)\%m;
    r = (r%m + m)%m;
    if(!r) return (a*(n+1))%mod;
    if(!a) return solve(r, n, m, mod);
    11 g, x, y;
    g = gcdExtended(r, m, x, y);
    \dot{x} = (x\%m + m)\%m;
    11 d = a - (a/q) *q;
    a -= d;
    x = (x*(a/g)) %m;
```

4.7 Primitive Root

```
//is n primitive root of p ?
bool test(long long x, long long p) {
    long long m = p - 1;
    for(int i = 2; i * i <= m; ++i) if(!(m % i)) {
        if(fexp(x, i, p) == 1) return false;
        if(fexp(x, m / i, p) == 1) return false;
    }
    return true;
}
//find the smallest primitive root for p
int search(int p) {
    for(int i = 2; i < p; i++) if(test(i, p)) return i;
    return -1;
}</pre>
```

4.8 Linear Sieve

```
//check long long
vector <int> prime;
bool is_composite[MAXN];
int cnt[MAXN];
long long primePow[MAXN];
long long func[MAXN];
long long getFunction(int i, int p) {
 return cnt[i] + 1;
void sieve (int n) {
  fill(is_composite, is_composite + n, false);
  func[1] = 1;
  for (int i = 2; i < n; ++i) {
    if (!is_composite[i]) {
      prime.push_back (i);
      func[i] = 1; // base case
      cnt[i] = 1; primePow[i] = i;
    for (int j = 0; j < prime.size () && i * prime[j] < n; ++j) {
      is_composite[i * prime[j]] = true;
      if (i % prime[j] == 0) {
        func[i * prime[j]] = func[i / primePow[i]] * getFunction(i, prime[j
            ]); // f(ip) = f(i / primePow[i]) * f(primePow[i] * prime[j])
        cnt[i * prime[j]] = cnt[i] + 1;
        primePow[i * prime[j]] = primePow[i] * prime[j];
        break:
      } else {
        func[i * prime[j]] = func[i] * func[prime[j]]; // f(ip) = f(i) * f(ip)
        cnt[i * prime[j]] = 1;
        primePow[i * prime[j]] = prime[j];
```

```
// euclides estendido: acha u e v da equacao:
// u * x + v * y = gcd(x, y);
// u eh inverso modular de x no modulo y
// v eh inverso modular de y no modulo x

pair<ll, ll> euclides(ll a, ll b) {
    ll u = 0, oldu = 1, v = 1, oldv = 0;
    while(b) {
        ll q = a / b;
        oldv = oldv - v * q;
        oldu = oldu - u * q;
        a = a - b * q;
        swap(a, b);
        swap(v, oldu);
    }
    return make_pair(oldu, oldv);
}
```

4.10 Fast Exponentiation

```
const int mod = 1e9+7;
int fexp(int a, int b) {
  int ans = 1;
  while(b) {
    if(b & 1) ans = ans * a % mod;
    a = a * a % mod;
    b >>= 1;
  }
  return ans;
}
```

4.11 Matrix

4.12 FFT - Fast Fourier Transform

```
typedef double ld;
const ld PI = acos(-1);
struct Complex {
  ld real, imag;
  Complex conj() { return Complex(real, -imag); }
  Complex (id a = 0, ld b = 0) : real(a), imag(b) {}
  Complex operator + (const Complex &o) const { return Complex(real + o. real, imag + o.imag); }
  Complex operator - (const Complex &o) const { return Complex(real - o. real, imag - o.imag); }
  Complex operator * (const Complex &o) const { return Complex(real * o. real - imag * o.imag, real * o.imag + imag * o.real); }
  Complex operator / (ld o) const { return Complex(real / o, imag / o); }
```

```
void operator *= (Complex o) { *this = *this * o; }
  void operator /= (ld o) { real /= o, imag /= o; }
typedef std::vector<Complex> CVector;
const int ms = 1 \ll 22;
int bits[ms];
Complex root[ms];
void initFFT() {
 root[1] = Complex(1);
  for(int len = 2; len < ms; len += len) {</pre>
    Complex z(cos(PI / len), sin(PI / len));
    for(int i = len / 2; i < len; i++) {</pre>
      root[2 * i] = root[i];
      root[2 * i + 1] = root[i] * z;
void pre(int n) {
 int LOG = 0;
  while (1 << (LOG + 1) < n) {
   LOG++;
  for (int i = 1; i < n; i++) {
    bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << LOG);
CVector fft (CVector a, bool inv = false) {
  int n = a.size();
 pre(n);
  if(inv) {
    std::reverse(a.begin() + 1, a.end());
  for (int i = 0; i < n; i++) {
    int to = bits[i];
    if(to > i) {
      std::swap(a[to], a[i]);
  for(int len = 1; len < n; len *= 2) {</pre>
    for(int i = 0; i < n; i += 2 * len) {
      for (int j = 0; j < len; j++) {
        Complex u = a[i + j], v = a[i + j + len] * root[len + j];
        a[i + j] = u + v;
        a[i + j + len] = u - v;
  if(inv) {
    for (int i = 0; i < n; i++)
     a[i] /= n;
  return a;
void fft2in1(CVector &a, CVector &b) {
  int n = (int) a.size();
  for (int i = 0; i < n; i++) {
   a[i] = Complex(a[i].real, b[i].real);
  auto c = fft(a);
  for (int i = 0; i < n; i++) {
   a[i] = (c[i] + c[(n-i) % n].conj()) * Complex(0.5, 0);
    b[i] = (c[i] - c[(n-i) % n].conj()) * Complex(0, -0.5);
void ifft2in1(CVector &a, CVector &b) {
 int n = (int) a.size();
  for (int i = 0; i < n; i++) a[i] = a[i] + b[i] * Complex (0, 1);
  a = fft(a, true);
  for (int i = 0; i < n; i++) {
   b[i] = Complex(a[i].imag, 0);
   a[i] = Complex(a[i].real, 0);
```

```
std::vector<long long> mod_mul(const std::vector<long long> &a, const std::
    vector<long long> &b, long long cut = 1 << 15) {</pre>
  int n = (int) a.size();
  CVector C[4];
  for (int i = 0; i < 4; i++) C[i].resize(n);
  for (int i = 0; i < n; i++) {
    C[0][i] = a[i] % cut;
    C[1][i] = a[i] / cut;
    C[2][i] = b[i] % cut;
    C[3][i] = b[i] / cut;
  fft2in1(C[0], C[1]);
  fft2in1(C[2], C[3]);
  for (int i = 0; i < n; i++) {
    // 00, 01, 10, 11
    Complex cur[4];
    for (int j = 0; j < 4; j++) cur[j] = C[j/2+2][i] * C[j % 2][i];
    for(int j = 0; j < 4; j++) C[j][i] = cur[j];</pre>
  ifft2in1(C[0], C[1]);
  ifft2in1(C[2], C[3]);
  std::vector<long long> ans(n, 0);
  for (int i = 0; i < n; i++) {
    // if there are negative values, care with rounding
    ans[i] += (long long) (C[0][i].real + 0.5);
    ans[i] += (long long) (C[1][i].real + C[2][i].real + 0.5) * cut;
    ans[i] += (long long) (C[3][i].real + 0.5) * cut * cut;
  return ans:
std::vector<int> mul(const std::vector<int> &a, const std::vector<int> &b)
  int n = 1;
  while (n - 1 < (int) a.size() + (int) b.size() - 2) n += n;
  CVector poly(n);
  for (int i = 0; i < n; i++) {
    if(i < (int) a.size())</pre>
     polv[i].real = a[i];
    if(i < (int) b.size())</pre>
      poly[i].imag = b[i];
  poly = fft(poly);
  for (int i = 0; i < n; i++) {
   poly[i] *= poly[i];
  poly = fft(poly, true);
  std::vector<int> c(n, 0);
  for (int i = 0; i < n; i++) {
   c[i] = (int) (poly[i].imag / 2 + 0.5);
  while (c.size() > 0 && c.back() == 0) c.pop_back();
  return c;
```

4.13 NTT - Number Theoretic Transform

```
root[1] = 1;
  for(int len = 2; len < ms; len += len) {</pre>
    int z = fexp(gen, (MOD - 1) / len / 2);
    for(int i = len / 2; i < len; i++) {</pre>
     root[2 * i] = root[i];
     root[2 * i + 1] = (long long) root[i] * z % MOD;
void pre(int n) {
 int LOG = 0;
 while (1 << (LOG + 1) < n) {
   LOG++:
 for (int i = 1; i < n; i++) {
    bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << LOG);
vector<int> fft(vector<int> a, bool inv = false) {
 int n = (int) a.size();
 pre(n);
 if(inv) {
   reverse(a.begin() + 1, a.end());
 for (int i = 0; i < n; i++) {
    int to = bits[i];
    if(i < to)
      swap(a[i], a[to]);
 for(int len = 1; len < n; len *= 2) {</pre>
    for (int i = 0; i < n; i += len * 2) {
      for (int j = 0; j < len; j++) {
        int u = a[i + j], v = (l1) a[i + j + len] * root[len + j] % mod;
        a[i + j] = add(u, v);
        a[i + j + len] = add(u, mod - v);
  if(inv) {
    int rev = fexp(n, mod-2, mod);
    for (int i = 0; i < n; i++)
     a[i] = (ll) a[i] * rev % mod;
 return a:
std::vector<int> shift(const std::vector<int> &a, int s) {
 int n = std::max(0, s + (int) a.size());
 std::vector<int> b(n, 0);
  for(int i = std::max(-s, 0); i < (int) a.size(); i++) {</pre>
   b[i + s] = a[i];
 return b;
std::vector<int> cut(const std::vector<int> &a, int n) {
 std::vector<int> b(n, 0);
 for(int i = 0; i < (int) a.size() && i < n; i++) {</pre>
   b[i] = a[i];
 return b;
std::vector<int> operator +(std::vector<int> a, const std::vector<int> &b)
  int sz = (int) std::max(a.size(), b.size());
 a.resize(sz, 0);
  for(int i = 0; i < (int) b.size(); i++) {</pre>
    a[i] = add(a[i], b[i]);
 return a;
std::vector<int> operator -(std::vector<int> a, const std::vector<int> &b)
```

```
int sz = (int) std::max(a.size(), b.size());
  a.resize(sz, 0);
  for(int i = 0; i < (int) b.size(); i++) {</pre>
    a[i] = add(a[i], MOD - b[i]);
  return a;
std::vector<int> operator *(std::vector<int> a, std::vector<int> b) {
  while(!a.empty() && a.back() == 0) a.pop_back();
  while(!b.empty() && b.back() == 0) b.pop_back();
  if(a.empty() || b.empty()) return std::vector<int>(0, 0);
  int n = 1;
  while (n-1 < (int) \ a.size() + (int) \ b.size() - 2) \ n += n;
  a.resize(n, 0);
 b.resize(n, 0);
 a = fft(a, false);
  b = fft(b, false);
  for (int i = 0; i < n; i++) {
    a[i] = (int) ((long long) a[i] * b[i] % MOD);
 return fft(a, true);
std::vector<int> inverse(const std::vector<int> &a, int k) {
 assert(!a.empty() && a[0] != 0);
  if(k == 0) {
    return std::vector<int>(1, (int) fexp(a[0], MOD - 2));
  } else {
    int n = 1 \ll k;
    auto c = inverse(a, k-1);
    return cut(c * cut(std::vector<int>(1, 2) - cut(a, n) * c, n), n);
std::vector<int> log(const std::vector<int> &a, int k) {
  assert(!a.empty() && a[0] != 0);
  int n = 1 \ll k;
  std::vector<int> b(n, 0);
  for(int i = 0; i+1 < (int) a.size() && i < n; i++) {</pre>
   b[i] = (int)((i + 1LL) * a[i+1] % MOD);
  b = cut(b * inverse(a, k), n);
  assert((int) b.size() == n);
  for (int i = n - 1; i > 0; i--)
   b[i] = (int) (b[i-1] * fexp(i, MOD - 2) % MOD);
  b[0] = 0:
  return b;
std::vector<int> exp(const std::vector<int> &a, int k) {
  assert(!a.empty() && a[0] == 0);
  if(k == 0) {
    return std::vector<int>(1, 1);
  } else {
    auto b = \exp(a, k-1);
    int n = 1 \ll k;
    return cut (b * cut (std::vector<int>(1, 1) + cut (a, n) - log(b, k), n),
        n);
```

4.14 Fast Walsh Hadamard Transform

```
vector<11> FWHT(char oper, vector<11> a, const bool inv = false) {
  int n = (int) a.size();
  for(int len = 1; len < n; len += len) {
    for(int i = 0; i < n; i += 2 * len) {
      for(int j = 0; j < len; j++) {
        auto u = a[i + j] % mod, v = a[i + j + len] % mod;
        if(oper == '^') {
            a[i + j] = (u + v) % mod;
        }
}</pre>
```

```
a[i + j + len] = (u - v + mod) % mod;
        if(oper == '|') {
          if(!inv) {
            a[i + j + len] = (u + v) % mod;
          } else {
            a[i + j + len] = (v - u + mod) % mod;
        if(oper == '&') {
          if(!inv) {
            a[i + j] = (u + v) % mod;
            a[i + j] = (u - v + mod) % mod;
  if(oper == '^' && inv) {
    11 \text{ rev} = \text{fexp(n, mod - 2);}
    for (int i = 0; i < n; i++) {
      a[i] = a[i] * rev % mod;
  return a;
vector<ll> multiply(char oper, vector<ll> a, vector<ll> b) {
  int n = 1;
  while (n < (int) max(a.size(), b.size())) {</pre>
   n <<= 1;
 vector<ll> ans(n);
  while (a.size() < ans.size()) a.push_back(0);</pre>
  while (b.size() < ans.size()) b.push back(0);</pre>
  a = FWHT(oper, a);
  b = FWHT (oper, b);
  for (int i = 0; i < n; i++) {
    ans[i] = a[i] * b[i] % mod;
  ans = FWHT(oper, ans, true);
  return ans;
const int mxlog = 17;
vector<ll> subset_multiply(vector<ll> a, vector<ll> b) {
 int n = 1;
  while (n < (int) max(a.size(), b.size())) {</pre>
    n <<= 1;
  vector<ll> ans(n);
  while (a.size() < ans.size()) a.push_back(0);</pre>
  while (b.size() < ans.size()) b.push_back(0);</pre>
  vector<vector<ll>> A(mxlog + 1, vector<ll>(a.size())), B(mxlog + 1,
      vector<ll>(b.size()));
  for (int i = 0; i < n; i++) {
    A[__builtin_popcount(i)][i] = a[i];
    B[__builtin_popcount(i)][i] = b[i];
  for (int i = 0; i <= mxlog; i++) {</pre>
    A[i] = FWHT('|', A[i]);
    B[i] = FWHT('|', B[i]);
  for (int i = 0; i <= mxlog; i++) {</pre>
    vector<ll> C(n);
    for (int x = 0; x \le i; x++) {
      int y = i - x;
      for (int j = 0; j < n; j++) {
        C[j] = (C[j] + A[x][j] * B[y][j] % mod) % mod;
```

```
C = FWHT('|', C, true);
for (int j = 0; j < n; j++) {
   if (_builtin_popcount(j) == i) {
      ans[j] = (ans[j] + C[j]) % mod;
   }
}
return ans;
}</pre>
```

4.15 Miller and Rho

```
//miller rabin
typedef unsigned long long ull;
typedef long double 1d;
ull fmul(ull a, ull b, ull m) {
 ull q = (ld) a * (ld) b / (ld) m;
  ull r = a * b - q * m;
  return (r + m) % m;
bool miller (ull p, ull a) {
 ull s = p - 1;
  while (s % 2 == 0) s >>= 1;
  while (a >= p) a >>= 1;
 ull mod = fexp(a, s, p);
  while (s != p - 1 && mod != 1 && mod != p - 1) {
    mod = fmul(mod, mod, p);
    s <<= 1;
  if(mod != p - 1 && s % 2 == 0)return false;
  else return true;
bool prime (ull p) {
  if(p <= 3)
    return true;
  if(p % 2 == 0)
    return false;
  return miller(p, 2) && miller(p, 3)
    && miller(p, 5) && miller(p, 7)
    && miller(p, 11) && miller(p, 13)
    && miller(p, 17) && miller(p, 19)
    && miller(p, 23) && miller(p, 29)
    && miller(p, 31) && miller(p, 37);
//pollard_rho
ull func(ull x, ull c, ull n) {
  return (fmul(x, x, n) + c) % n;
ull gcd(ull a, ull b) {
  if(!b) return a;
  else return gcd(b, a % b);
ull rho(ull n) {
 if(n % 2 == 0) return 2;
  if(prime(n)) return n;
  while(1) {
    do {
     c = rand() % n;
    } while (c == 0 | | (c + 2) % n == 0);
    ull x = 2, y = 2, d = 1;
    ull pot = 1, lam = 1;
    do {
      if(pot == lam) {
        x = y;
pot <<= 1;
        lam = 0;
      \dot{y} = func(y, c, n);
      lam++;
      d = gcd(x >= y ? x - y : y - x, n);
```

```
} while(d == 1);
   if(d != n) return d;
vector<ull> factors(ull n) {
 vector<ull> ans, rest, times;
  if(n == 1) return ans;
  rest.push_back(n);
  times.push_back(1);
  while(!rest.empty()) {
    ull x = rho(rest.back());
    if(x == rest.back()) {
      int freq = 0;
      for(int i = 0; i < rest.size(); i++) {</pre>
        int cur_freq = 0;
        while (rest[i] % x == 0) {
          rest[i] /= x;
          cur_freq++;
        freq += cur_freq * times[i];
        if(rest[i] == 1) {
          swap(rest[i], rest.back());
          swap(times[i], times.back());
          rest.pop_back();
          times.pop_back();
      while(freq--) {
        ans.push_back(x);
      continue;
    ull e = 0;
    while(rest.back() % x == 0) {
     rest.back() /= x;
      e++;
    e *= times.back();
    if(rest.back() == 1) {
      rest.pop_back();
      times.pop_back();
    rest.push_back(x);
    times.push_back(e);
  return ans;
```

4.16 Determinant using Mod

```
swap (mat[i][k], mat[j][k]);

    res = -res;
    break;

}

if (!flag) {
    return 0;
}

for (int j = i + 1; j < n; j++) {
    while (mat[j][i]) {
        ll t = mat[i][i] / mat[j][i];
        for (int k = i; k < n; k++) {
            mat[i][k] = (mat[i][k] - t * mat[j][k]) % mod;
            swap (mat[i][k], mat[j][k]);
    }
    res = -res;
    }
}
res = (res * mat[i][i]) % mod;

return (res + mod) % mod;
</pre>
```

4.17 Gauss

```
const double eps = 1e-7;
int gauss (vector<vector<double>> a, vector<double> & ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
         int sel = row;
         for (int i=row; i<n; ++i) {</pre>
             if (abs (a[i][col]) > abs (a[sel][col]))
                 sel = i;
         if (abs (a[sel][col]) < eps) continue;</pre>
         for (int i=col; i<=m; ++i)</pre>
             swap (a[sel][i], a[row][i]);
         where [col] = row;
         for (int i=0; i<n; ++i) {</pre>
             if (i != row) {
                 double c = a[i][col] / a[row][col];
                 for (int j=col; j<=m; ++j)</pre>
                     a[i][j] -= a[row][j] * c;
         ++row;
    ans.assign (m, 0);
    for (int i=0; i<m; ++i) {</pre>
         if (where[i] != -1)
             ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {</pre>
         double sum = 0;
         for (int j=0; j<m; ++j)</pre>
             sum += ans[j] * a[i][j];
         if (abs (sum - a[i][m]) > eps)
             return 0;
    for (int i=0; i<m; ++i) {</pre>
        if (where[i] == -1)
             return INF;
    return 1;
```

```
// mod 2 (xor);
int gauss (vector <bitset<ms>> a, int m) {
    int n = (int) a.size();
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        for (int i=row; i<n; ++i) {</pre>
            if (a[i][col]) {
                 swap (a[i], a[row]);
                 break;
        if (!a[row][col]) continue;
        where[col] = row;
        for (int i=0; i<n; ++i) {</pre>
            if (i != row && a[i][col])
                a[i] ^= a[row];
        ++row;
    //same above
```

4.18 Lagrange Interpolation

```
class LagrangePoly {
public:
  LagrangePoly(vector<long long> _a) {
    //f(i) = \_a[i]
    //interpola o vetor em um polinomio de grau y.size() - 1
    y = _a;
    den.resize(y.size());
    int n = (int) y.size();
    for (int i = 0; i < n; i++) {</pre>
     y[i] = (y[i] % MOD + MOD) % MOD;
      den[i] = ifat[n - i - 1] * ifat[i] % MOD;
      if((n - i - 1) % 2 == 1) {
        den[i] = (MOD - den[i]) % MOD;
  long long getVal(long long x) {
    int n = (int) y.size();
    x %= MOD;
    if(x < n) {
      //return y[(int) x];
    vector<long long> 1, r;
    l.resize(n);
    1[0] = 1;
    for (int i = 1; i < n; i++) {
     l[i] = l[i - 1] * (x - (i - 1) + MOD) % MOD;
    r.resize(n);
    r[n - 1] = 1;
    for (int i = n - 2; i >= 0; i--) {
      r[i] = r[i + 1] * (x - (i + 1) + MOD) % MOD;
    long long ans = 0;
    for (int i = 0; i < n; i++) {
      long long coef = l[i] * r[i] % MOD;
      ans = (ans + coef * y[i] % MOD * den[i]) % MOD;
    return ans;
private:
  vector<long long> y, den;
```

```
int main(){
    fat[0] = ifat[0] = 1;
    for(int i = 1; i < ms; i++) {
        fat[i] = fat[i - 1] * i * MOD;
        ifat[i] = fexp(fat[i], MOD - 2);
}
// Codeforces 622F
int x, k;
cin >> x >> k;
vector<long long> a;
a.push_back(0);
for(long long i = 1; i <= k + 1; i++) {
        a.push_back((a.back() + fexp(i, k)) * MOD);
}
LagrangePoly f(a);
cout << f.getVal(x) << '\n';
}</pre>
```

4.19 Lagrange extracting polynomial

```
// O(n^2), receve v {x, y} e retorna o polinomio em fracao
vector<pair<int, int>> interpolate(vector<ii> v) {
  int n = v.size();
  vector<int> prod(n+1);
  prod[0] = 1;
  for (auto p : v) {
    for (int i = n; i > 0; i--) {
      prod[i] = prod[i-1] - p.first * prod[i];
    prod[0] = -p.first * prod[0];
  vector<pair<int, int>> ans(n+1);
  for(int i = 0; i <= n; i++) ans[i].second = 1;</pre>
  for (int i = 0; i < n; i++) {
    vector<int> pol(n+1); // (x - v[i].first)
    for (int j = n; j > 0; j--) {
      pol[j-1] = prod[j] + pol[j] * v[i].first;
    for (int j = 0; j < n; j++) {
      pol[j] *= v[i].second;
    int k = 1;
    for (int j = 0; j < n; j++) {
      if(i==j) continue;
      k *= v[i].first - v[j].first;
    if(k < 0) {
      k = -k
      for (auto &p : pol) p = -p;
    for (int i = 0; i < n; i++) {
      ans[i] = {ans[i].first*k + pol[i]*ans[i].second, k*ans[i].second};
      if(ans[i].first == 0) ans[i].second = 1;
      else {
        int gc = __gcd(abs(ans[i].first), ans[i].second);
        ans[i].first /= qc;
        ans[i].second /= qc;
  return ans;
```

1.20 Count integer points inside triangle

```
//gcd(p, q) == 1
ll get(ll p, ll q, ll n, bool floor = true) {
    if (n == 0) {
        return 0;
    }
    if (p % q == 0) {
        return n * (n + 1) / 2 * (p / q);
    }
    if (p > q) {
        return n * (n + 1) / 2 * (p / q) + get(p % q, q, n, floor);
    }
    ll new_n = p * n / q;
    ll ans = (n + 1) * new_n - get(q, p, new_n, false);
    if (!floor) {
        ans += n - n / q;
    }
    return ans;
}
```

4.21 Prime Counting

```
const int ms = 5001000, \lim n = 3e5, \lim p = 1e2;
std::vector<int> primes;
int id[ms];
int memo[lim_n][lim_p];
void pre() {
  std::vector<bool> isPrime(ms, true);
  for(int i = 2; i < ms; i++) {</pre>
    id[i] = (int) primes.size();
    if(!isPrime[i]) continue;
    id[i]++;
    primes.push_back(i);
    for(int j = i+i; j < ms; j += i) isPrime[j] = false;</pre>
  for(int i = 1; i < lim_n; i++) {</pre>
    memo[i][0] = i;
    for(int j = 1; j < lim_p; j++) memo[i][j] = memo[i][j-1] - memo[i/</pre>
        primes[j-1]][j-1];
int cbc(long long n) {
  int ans = std::max(0, (int) pow((double) n, 1.0 / 3) - 2);
  while ((11) ans * ans * ans < n) ans++;
  return ans;
long long dp(long long n, int i) {
  if(n == 0) return 0; if(i == 0) return n;
 if(primes[i-1] >= n) return 1;
  if((11) primes[i-1] * primes[i-1] > n && n < ms) return id[n] - (i-1);</pre>
  else if(n < lim_n && i < lim_p) return memo[n][i];</pre>
  else return dp(n, i-1) - dp(n / primes[i-1], i-1);
long long primeFunction(long long n) {
  if(n < ms) return id[(int)n];</pre>
  int i = id[cbc(n)];
  long long ans = dp(n, i) + i - 1;
  while((long long) primes[i] * primes[i] <= n) {</pre>
    ans -= primeFunction(n / primes[i]) - i;
    i++;
  return ans;
```

4.22 Berlekamp Massey

```
vector<int> berlekampMassey(const vector<int> &s) {
```

```
int n = (int) s.size(), l = 0, m = 1;
    vector<int> b(n), c(n);
    int 1d = b[0] = c[0] = 1;
    for (int i=0; i<n; i++, m++) {</pre>
         int d = s[i];
        for (int j=1; j<=1; j++)</pre>
             d = (d + c[j] * s[i-j]) % mod;
        if (d == 0)
             continue;
         vector<int> temp = c;
         int coef = d * fexp(ld, mod-2) % mod;
         for (int j=m; j<n; j++)</pre>
             (c[j] = (c[j] - coef * b[j-m]) % mod + mod) % mod;
         if (2 * 1 <= i) {
             1 = i + 1 - 1;
            b = temp;
             1d = d;
             \mathbf{m} = 0;
    c.resize(1 + 1):
    c.erase(c.begin());
    for (int &x : c)
       x = mod-x;
    return c;
// p = p*q % h
void mull(vector<int> &p,vector<int> &q, vector<int> &h, int m) {
         vector<int> t_(m+m);
         for (int i=0; i < m; ++i) if (p[i])</pre>
                 for(int j=0; j<m; ++ j)
                          t_{[i+j]} = (t_{[i+j]} + p[i] * q[j]) %mod;
         for(int i=m+m-1;i>=m;--i) if(t_[i])
                  //miuns t_{[i]}x^{i-m}(x^m-\sum_{j=0}^{m-1} x^{m-j-1}h_{j})
                 for (int j=m-1; ~ j; -- j)
                          t_{[i-j-1]} = (t_{[i-j-1]} + t_{[i]} *h[j]) %mod;
         for (int i=0;i<m;++i) p[i]=t_[i];</pre>
// a = caso base, h = recorrencia, m = tamanho da recorrencia
inline int calc(vector<int> &a, vector<int> &h, int K, int m) {
         vector<int> s(m), t(m);
         s[0]=1; if(m!=1) t[1]=1; else t[0]=h[0];
         //binary-exponentiation
         while(K) {
                 if(K&1) mull(s,t,h,m);
                 mull(t,t,h,m); K>>=1;
         for (int i=0; i < m; ++i) su=(su+s[i] *a[i]) %mod;</pre>
         return (su%mod+mod)%mod;
```

5 Geometry

5.1 Geometry

```
const double inf = 1e100, eps = 1e-12;
const double PI = acos(-1.0L);
int cmp (double a, double b = 0) {
  if (abs(a-b) < eps) return 0;
  return (a < b) ? -1 : +1;
}
struct PT {
  double x, y;</pre>
```

```
PT (double x = 0, double y = 0) : x(x), y(y) {}
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
  PT operator * (double c) const { return PT(x*c, y*c); }
  PT operator / (double c) const { return PT(x/c, y/c); }
  bool operator < (const PT &p) const {
    if(cmp(x, p.x) != 0) return x < p.x;
    return cmp(y, p.y) < 0;
  bool operator == (const PT &p) const {return !cmp(x, p.x) && !cmp(y, p.y)
      ; }
  bool operator != (const PT &p) const {return ! (p == *this);}
ostream &operator<<(ostream &os, const PT &p) {
  os << "(" << p.x << "," << p.y << ")";</pre>
double dot (PT p, PT q) { return p.x * q.x + p.y*q.y; }
double cross (PT p, PT q) { return p.x * q.y - p.y*q.x; }
double dist2 (PT p, PT q = PT(0, 0)) { return dot(p-q, p-q); }
double dist (PT p, PT q) { return hypot(p.x-q.x, p.y-q.y); }
double norm (PT p) { return hypot(p.x, p.y); }
PT normalize (PT p) { return p/hypot(p.x, p.y); }
double angle (PT p, PT q) { return atan2(cross(p, q), dot(p, q)); }
double angle (PT p) { return atan2(p.y, p.x); }
double polarAngle (PT p) {
  double a = atan2(p.y,p.x);
  return a < 0 ? a + 2*PI : a;
PT rotateCCW90 (PT p) { return PT(-p.y, p.x); }
PT rotateCW90 (PT p) { return PT(p.y, -p.x); }
PT rotateCCW (PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
typedef pair<PT, int> Line;
PT getDir (PT a, PT b) {
  if (a.x == b.x) return PT(0, 1);
  if (a.y == b.y) return PT(1, 0);
  int dx = b.x-a.x;
  int dy = b.y-a.y;
  int g = \underline{gcd(abs(dx), abs(dy))};
  if (dx < 0) q = -q;
  return PT(dx/g, dy/g);
Line getLine (PT a, PT b) {
  PT dir = getDir(a, b);
  return {dir, cross(dir, a)};
PT projPtLine (PT a, PT b, PT c) { // ponto c na linha a - b, a.b = |a|
    cost * Ibl
  return a + (b-a) * dot(b-a, c-a)/dot(b-a, b-a);
PT reflectPointLine (PT a, PT b, PT c) {
 PT p = projPtLine(a, b, c);
  return p*2 - c;
PT projPtSeg (PT a, PT b, PT c) { // c no segmento a - b double r = dot(b-a, b-a);
  if (cmp(r) == 0) return a;
  r = dot(b-a, c-a)/r;
  if (cmp(r, 0) < 0) return a;
  if (cmp(r, 1) > 0) return b;
  return a + (b - a) * r;
double distancePointSegment (PT a, PT b, PT c) { // ponto c e o segmento a
  return dist(c, projPtSeg(a, b, c));
bool ptInSegment (PT a, PT b, PT c) { // ponto c esta em um segmento a - b
  if (a == b) return a == c;
  a = a-c, b = b-c;
  return cmp(cross(a, b)) == 0 && cmp(dot(a, b)) <= 0;
bool parallel (PT a, PT b, PT c, PT d) {
  return cmp(cross(b - a, c - d)) == 0;
```

```
bool collinear (PT a, PT b, PT c, PT d) {
   return parallel(a, b, c, d) && cmp(cross(a - b, a - c)) == 0 && cmp(cross
             (c - d, c - a)) == 0;
 // Calcula distancia entre o ponto (x, y, z) e o plano ax + by + cz = d
double distPtPlane(double x, double y, double z, double a, double b, double
          c. double d) {
        return abs (a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c);
bool segInter (PT a, PT b, PT c, PT d) {
    if (collinear(a, b, c, d)) {
        if (cmp(dist(a, c)) == 0 \mid | cmp(dist(a, d)) == 0 \mid | cmp(dist(b, c)) == 0 \mid | cmp(dist(b, c
                0 || cmp(dist(b, d)) == 0) return true;
        if (cmp(dot(c - a, c - b)) > 0 \&\& cmp(dot(d - a, d - b)) > 0 \&\& cmp(dot
                 (c - b, d - b)) > 0) return false;
        return true;
    if (cmp(cross(d - a, b - a) * cross(c - a, b - a)) > 0) return false;
    if (cmp(cross(a - c, d - c) * cross(b - c, d - c)) > 0) return false;
    return true;
// Calcula a intersecao entre as retas a - b e c - d assumindo que uma
         unica intersecao existe
// Para intersecao de segmentos, cheque primeiro se os segmentos se
         intersectam e que nao sao paralelos
PT lineLine (PT a, PT b, PT c, PT d) {
   b = b - a; d = c - d; c = c - a;
    assert(cmp(cross(b, d)) != 0);
    return a + b * cross(c, d) / cross(b, d);
PT circleCenter (PT a, PT b, PT c) {
   b = (a + b) / 2; // bissector
    c = (a + c) / 2; // bissector
    return lineLine(b, b + rotateCW90(a - b), c, c + rotateCW90(a - c));
vector<PT> circle2PtsRad (PT p1, PT p2, double r) {
   vector<PT> ret;
    double d2 = dist2(p1, p2);
    double det = r * r / d2 - 0.25;
    if (det < 0.0) return ret;</pre>
    double h = sqrt(det);
    for (int i = 0; i < 2; i++) {
        double x = (p1.x + p2.x) * 0.5 + (p1.y - p2.y) * h;
        double y = (\bar{p}1.y + \bar{p}2.y) * 0.5 + (\bar{p}2.\bar{x} - \bar{p}1.\bar{x}) * h;
        ret.push_back(PT(x, y));
        swap(p1, p2);
   return ret:
bool circleLineIntersection(PT a, PT b, PT c, double r) {
        return cmp(dist(c, projPtLine(a, b, c)), r) <= 0;</pre>
vector<PT> circleLine (PT a, PT b, PT c, double r) {
   vector<PT> ret;
    PT p = projPtLine(a, b, c), p1;
    double h = norm(c-p);
    if (cmp(h,r) == 0) {
       ret.push_back(p);
    \} else if (cmp(h,r) < 0)
        double k = sqrt(r*r - h*h);
        p1 = p + (b-a) / (norm(b-a)) *k;
        ret.push_back(p1);
        p1 = p - (b-a) / (norm(b-a)) *k;
        ret.push_back(p1);
    return ret;
bool ptInsideTriangle(PT p, PT a, PT b, PT c) {
    if(cross(b-a, c-b) < 0) swap(a, b);
    if(ptInSegment(a,b,p)) return 1;
    if(ptInSegment(b,c,p)) return 1;
```

if(ptInSegment(c,a,p)) return 1;

```
bool x = cross(b-a, p-b) < 0;
 bool y = cross(c-b, p-c) < 0;
  bool z = cross(a-c, p-a) < 0;
  return x == y && y == z;
bool pointInConvexPolygon(const vector<PT> &p, PT q) {
  if (p.size() == 1) return p.front() == q;
 int 1 = 1, r = p.size()-1;
  while (abs(r-1) > 1) {
    int m = (r+1)/2;
    if(cross(p[m]-p[0], q-p[0]) < 0) r = m;
    else l = m;
  return ptInsideTriangle(q, p[0], p[1], p[r]);
// Determina se o ponto esta num poligono possivelmente nao-convexo
// Retorna 1 para pontos estritamente dentro, 0 para pontos estritamente
    fora do poligno
// e 0 ou 1 para os pontos restantes
bool pointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for (int i = 0; i < p.size(); i++) {
    int j = (i + 1) % p.size();
    if((p[i].y \le q.y \& q.y < p[j].y || p[j].y \le q.y \& q.y < p[i].y) \& 
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y)
      c = !c;
  return c;
// Determina se o ponto esta na borda do poligno
bool pointOnPolygon(const vector<PT> &p, PT q) {
  for(int i = 0; i < p.size(); i++)</pre>
    if(cmp(dist(projPtSeg(p[i], p[(i + 1) % p.size()], q), q)) == 0)
      return true:
    return false;
// area / semiperimeter
double rIncircle (PT a, PT b, PT c) {
  double ab = norm(a-b), bc = norm(b-c), ca = norm(c-a);
  return abs(cross(b-a, c-a)/(ab+bc+ca));
vector<PT> circleCircle (PT a, double r, PT b, double R) {
  vector<PT> ret;
  double d = norm(a-b);
  if (d > r + R \mid \mid d + min(r, R) < max(r, R)) return ret;
  double x = (d*d - R*R + r*r) / (2*d); // x = r*cos(R opposite angle)
  double y = sqrt(r*r - x*x);
  PT v = (b - a)/d;
  ret.push_back(a + v*x + rotateCCW90(v)*y);
  if (cmp(y) > 0)
    ret.push_back(a + v*x - rotateCCW90(v)*y);
  return ret;
double circularSegArea (double r, double R, double d)
  double ang = 2 * acos((d*d - R*R + r*r) / (2*d*r)); // cos(R opposite)
      angle) = x/r
  double tri = sin(ang) * r * r;
  double sector = ang * r * r;
  return (sector - tri) / 2;
double computeSignedArea (const vector<PT> &p) {
  double area = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area/2.0;
double computeArea(const vector<PT> &p) { return abs(computeSignedArea(p));
PT computeCentroid(const vector<PT> &p) {
  PT c(0,0);
```

```
double scale = 6.0 * computeSignedArea(p);
  for(int i = 0; i < p.size(); i++){</pre>
    int j = (i + 1) % p.size();
    c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
  return c / scale;
// Testa se o poligno listada em ordem CW ou CCW eh simples (nenhuma linha
    se intersecta)
bool isSimple(const vector<PT> &p) {
  for(int i = 0; i < p.size(); i++) {</pre>
    for (int k = i + 1; k < p.size(); k++) {
      int j = (i + 1) % p.size();
      int 1 = (k + 1) % p.size();
      if (i == 1 \mid | i == k) continue;
      if (segInter(p[i], p[j], p[k], p[l]))
        return false;
  return true;
vector< pair<PT, PT> > getTangentSegs (PT c1, double r1, PT c2, double r2)
  if (r1 < r2) swap(c1, c2), swap(r1, r2);
  vector<pair<PT, PT> > ans;
  double d = dist(c1, c2);
  if (cmp(d) <= 0) return ans;</pre>
  double dr = abs(r1 - r2), sr = r1 + r2;
  if (cmp(dr, d) >= 0) return ans;
  double u = acos(dr / d);
  PT dc1 = normalize(c2 - c1) *r1;
  PT dc2 = normalize(c2 - c1) \star r2;
  ans.push_back(make_pair(c1 + rotateCCW(dc1, +u), c2 + rotateCCW(dc2, +u))
      );
  ans.push_back(make_pair(c1 + rotateCCW(dc1, -u), c2 + rotateCCW(dc2, -u))
      );
  if (cmp(sr, d) >= 0) return ans;
  double v = acos(sr / d);
  dc2 = normalize(c1 - c2) *r2;
  ans.push_back({c1 + rotateCCW(dc1, +v), c2 + rotateCCW(dc2, +v)});
  ans.push_back(\{c1 + rotateCCW(dc1, -v), c2 + rotateCCW(dc2, -v)\});
  return ans:
```

5.2 Convex Hull

```
vector<PT> convexHull(vector<PT> p, bool needs = 1) {
  if(needs) sort(p.begin(), p.end());
  p.erase(unique(p.begin(), p.end()), p.end());
  int n = p.size(), k = 0;
  if(n <= 1) return p;</pre>
  vector<PT> h(n + 2); // se der wa bota n*2
  for (int i = 0; i < n; i++) {
    while (k \ge 2 \&\& cross(h[k-1] - h[k-2], p[i] - h[k-2]) \le 0) k--;
    h[k++] = p[i];
  for (int i = n - 2, t = k + 1; i >= 0; i--) {
    while (k \ge t \&\& cross(h[k-1] - h[k-2], p[i] - h[k-2]) \le 0) k--;
    h[k++] = p[i];
  h.resize(k); // n+1 points where the first is equal to the last
  return h;
vector<PT> splitHull(const vector<PT> &hull) {
  vector<PT> ans(hull.size());
  for (int i = 0, j = (int) \text{ hull.size}()-1, k = 0; k < (int) \text{ hull.size}(); k = 0
    if(hull[i] < hull[j]) {</pre>
      ans[k] = hull[i++];
```

```
} else {
      ans[k] = hull[i--];
 return ans;
// uniao de convex hulls
vector<PT> ConvexHull(const vector<PT> &a, const vector<PT> &b) {
 auto A = splitHull(a);
 auto B = splitHull(b);
 vector<PT> C(A.size() + B.size());
 merge(A.begin(), A.end(), B.begin(), B.end(), C.begin());
 return ConvexHull(C, false);
int maximizeScalarProduct(const vector<PT> &hull, PT vec) {
 // this code assumes that there are no 3 colinear points
 int ans = 0;
 int n = hull.size();
  if(n < 20) {
    for (int i = 0; i < n; i++) {
      if(dot(hull[i], vec) > dot(hull[ans], vec)) {
        ans = i:
  } else {
    if (dot(hull[1], vec) > dot(hull[ans], vec)) {
    for(int rep = 0; rep < 2; rep++) {</pre>
      int 1 = 2, r = n - 1;
      while (1 != r) {
        int mid = (1 + r + 1) / 2;
        bool flag = dot(hull[mid], vec) >= dot(hull[mid-1], vec);
        if(rep == 0) { flag = flag && dot(hull[mid], vec) >= dot(hull[0],
        else { flag = flag || dot(hull[mid-1], vec) < dot(hull[0], vec); }</pre>
        if(flag) {
          1 = mid;
        } else {
          r = mid - 1;
      if(dot(hull[ans], vec) < dot(hull[1], vec)) {</pre>
  return ans;
```

5.3 Cut Polygon

```
struct Segment {
    typedef long double T;
    PT p1, p2;
    T a, b, c;

Segment() {}

Segment(PT st, PT en) {
    p1 = st, p2 = en;
    a = -(st.y - en.y);
    b = st.x - en.x;
    c = a * en.x + b * en.y;
}

T plug(T x, T y) {
    // plug >= 0 is to the right
    return a * x + b * y - c;
}
```

```
T plug(PT p) {
    return plug(p.x, p.y);
  bool inLine(PT p) { return cross((p - p1), (p2 - p1)) == 0; }
  bool inSegment(PT p) {
    return inLine(p) && dot((p1 - p2), (p - p2)) >= 0 && dot((p2 - p1), (p
        - p1)) >= 0;
  PT lineIntersection(Segment s)
    long double A = a, B = b, C = c;
    long double D = s.a, E = s.b, F = s.c;
    long double x = (long double) C * E - (long double) B * F;
    long double y = (long double) A * F - (long double) C * D;
    long double tmp = (long double) A * E - (long double) B * D;
    x /= tmp;
    y /= tmp;
    return PT(x, y);
 bool polygonIntersection(const vector<PT> &poly) {
    long double 1 = -1e18, r = 1e18;
    for(auto p : poly) {
      long double z = plug(p);
      1 = \max(1, z);
      r = min(r, z);
    return 1 - r > eps;
};
vector<PT> cutPolygon(vector<PT> poly, Segment seg) {
  int n = (int) poly.size();
  vector<PT> ans;
  for (int i = 0; i < n; i++) {
    double z = seg.plug(poly[i]);
    if(z > -eps) {
      ans.push_back(poly[i]);
    double z2 = seg.plug(poly[(i + 1) % n]);
    if((z > eps \&\& z2 < -eps) || (z < -eps \&\& z2 > eps)) {
      ans.push_back(seg.lineIntersection(Segment(poly[i], poly[(i + 1) % n
          1)));
  return ans;
```

5.4 Smallest Enclosing Circle

```
for(int k = 0; k < j; k++) if (!inCircle(C, v[k])){
    PT o = circumcenter(v[i], v[j], v[k]);
    C = circle(o,dist(o,v[k]));
}
}
return C;
</pre>
```

5.5 Minkowski

```
bool comp (PT a, PT b) {
  int hp1 = (a.x < 0 \mid | (a.x==0 \&\& a.y<0));
  int hp2 = (b.x < 0 \mid | (b.x==0 \&\& b.y<0));
  if(hp1 != hp2) return hp1 < hp2;</pre>
  long long R = cross(a, b);
  if(R) return R > 0;
 return dot(a, a) < dot(b, b);
 // This code assumes points are ordered in ccw and the first points in
     both vectors is the min lexicographically
vector<PT> minkowskiSum(const vector<PT> &a, const vector<PT> &b) {
  if(a.empty() || b.empty()) return vector<PT>(0);
  vector<PT> ret;
  int n1 = a.size(), n2 = b.size();
  if(min(n1, n2) < 2){
    for (int i = 0; i < n1; i++) {
      for (int j = 0; j < n2; j++) {
        ret.push_back(a[i]+b[j]);
    return ret;
  PT v1, v2, p = a[0]+b[0];
  ret.push_back(p);
  for (int i = 0, j = 0; i + j + 1 < n1+n2; ) {
    v1 = a[(i+1)%n1]-a[i];
    v2 = b[(j+1)%n2]-b[j];
    if(j == n2 \mid | (i < n1 && comp(v1, v2))) p = p + v1, i++;
    else p = p + v2, j++;
    while(ret.size() >= 2 && cmp(cross(p-ret.back(), p-ret[(int)ret.size()
         -2])) == 0) {
      // removing colinear points
      // needs the scalar product stuff it the result is a line
      ret.pop_back();
    ret.push_back(p);
  return ret;
```

5.6 Half Plane Intersection

```
struct L {
    PT a, b;
    L(){}
    L()T a, PT b) : a(a), b(b) {}
};
double angle (L la) { return atan2(-(la.a.y - la.b.y), la.b.x - la.a.x); }
bool comp (L la, L lb) {
    if (cmp(angle(la), angle(lb)) == 0) return cross((lb.b - lb.a), (la.b - lb.a)) > eps;
    return cmp(angle(la), angle(lb)) < 0;
}
PT computeLineIntersection (L la, L lb) {
    return computeLineIntersection(la.a, la.b, lb.a, lb.b);
}
bool check (L la, L lb, L lc) {</pre>
```

```
PT p = computeLineIntersection(lb, lc);
            double det = cross((la.b - la.a), (p - la.a));
            return cmp(det) < 0;
\label{eq:convergence} \mbox{vector}\mbox{\ensuremath{\mbox{PT}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}\mbox{\ensuremath{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{\mbox{}}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}}\mbox{\ensuremath{}
            sort(line.begin(), line.end(), comp);
            vector<L> pl(1, line[0]);
            for (int i = 0; i < (int)line.size(); ++i) if (cmp(angle(line[i]),</pre>
                         angle(pl.back())) != 0) pl.push_back(line[i]);
             deque<int> dq:
            dq.push_back(0);
            dq.push_back(1);
            for (int i = 2; i < (int)pl.size(); ++i) {</pre>
                         while ((int)dq.size() > 1 && check(pl[i], pl[dq.back()], pl[dq[dq.
                                     size() - 2]])) dq.pop_back();
                         while ((int)dq.size() > 1 && check(pl[i], pl[dq[0]], pl[dq[1]])) dq
                                      .pop_front();
                         dq.push_back(i);
            while ((int)dq.size() > 1 \&\& check(pl[dq[0]], pl[dq.back()], pl[dq[dq.
                         size() - 2]])) dq.pop_back();
            while ((int)dq.size() > 1 \&\& check(pl[dq.back()], pl[dq[0]], pl[dq[1]])
                         ) dq.pop_front();
             vector<PT> res;
            for (int i = 0; i < (int) dq.size(); ++i) {
                         res.emplace_back(computeLineIntersection(pl[dq[i]], pl[dq[(i + 1) %
                                         dq.size()]]));
            return res;
```

5.7 Closest Pair

```
double closestPair(vector<PT> p) {
  int n = p.size(), k = 0;
  sort(p.begin(), p.end());
  double d = inf;
  set<PT> ptsInv;
  for(int i = 0; i < n; i++) {
    while(k < i && p[k].x < p[i].x - d) {
        ptsInv.erase(swapCoord(p[k++]));
    }
  for(auto it = ptsInv.lower_bound(PT(p[i].y - d, p[i].x - d));
        it != ptsInv.end() && it->x <= p[i].y + d; it++) {
        d = min(d, dist(p[i] - swapCoord(*it), PT(0, 0)));
    }
    ptsInv.insert(swapCoord(p[i]));
}
return d;
}</pre>
```

5.8 Voronoi

```
Segment getBisector(PT a, PT b) {
    Segment ans(a, b);
    swap(ans.a, ans.b);
    ans.b *= -1;
    ans.c = ans.a * (a.x + b.x) * 0.5 + ans.b * (a.y + b.y) * 0.5;
    return ans;
}

// BE CAREFUL!
// the first point may be any point
// O(N^3)
vector<PT> getCell(vector<PT> pts, int i) {
    vector<PT> ans;
    ans.emplace_back(0, 0);
    ans.emplace_back(le6, 0);
```

```
ans.emplace_back(1e6, 1e6);
  ans.emplace_back(0, 1e6);
  for(int j = 0; j < (int) pts.size(); <math>j++) {
   if(j != i) {
      ans = cutPolygon(ans, getBisector(pts[i], pts[j]));
 return ans;
// O(N^2) expected time
vector<vector<PT>> getVoronoi(vector<PT> pts) {
  // assert (pts.size() > 0);
  int n = (int) pts.size();
  vector<int> p(n, 0);
  for (int i = 0; i < n; i++) {
   p[i] = i;
  shuffle(p.begin(), p.end(), rng);
  vector<vector<PT>> ans(n);
  ans[0].emplace_back(0, 0);
  ans[0].emplace_back(w, 0);
  ans[0].emplace_back(w, h);
  ans[0].emplace_back(0, h);
  for(int i = 1; i < n; i++) {</pre>
   ans[i] = ans[0];
  for(auto i : p) {
    for(auto j : p) {
      if(j == i) break;
      auto bi = getBisector(pts[j], pts[i]);
      if(!bi.polygonIntersection(ans[j])) continue;
      ans[j] = cutPolygon(ans[j], getBisector(pts[j], pts[i]));
      ans[i] = cutPolygon(ans[i], getBisector(pts[i], pts[j]));
  return ans;
```

6 String Algorithms

6.1 KMP

```
vector<int> getBorder(string str) {
  int n = str.size();
  vector<int> border(n, -1);
  for (int i = 1, j = -1; i < n; i++) {
    while(j \ge 0 \&\& str[i] != str[j + 1]) {
      j = border[j];
    if(str[i] == str[j + 1]) {
      j++;
    border[i] = j;
  return border;
int matchPattern(const string &txt, const string &pat, const vector<int> &
    border) {
  int freq = 0;
  for (int i = 0, j = -1; i < txt.size(); i++) {
    while(j \ge 0 \&\& txt[i] != pat[j + 1]) {
      j = border[j];
    if(pat[j + 1] == txt[i]) {
      j++;
    if(j + 1 == (int) pat.size()) {
```

```
//found occurence
freq++;
j = border[j];
}
return freq;
```

6.2 KMP Automaton

```
// trad converts a char to its index
int trad(char ch) { return (int) ch; }
// sigma should be greater then the greatest value returned by trad
vector<vector<int>> buildAutomaton(string p, int sigma=300) {
  vector<vector<int>> A(p.size() + 1, vector<int>(sigma));
  int brd = 0;
  for (int i = 0; i < sigma; i++) A[0][i] = 0;
  A[0][trad(p[0])] = 1;
  for (int i = 1; i <= p.size(); i++) {
    for (int ch = 0; ch < sigma; ch++) {
        A[i][ch] = A[brd][ch];
    }
    if (i < p.size()) {
        A[i][trad(p[i])] = i + 1;
        brd = A[brd][trad(p[i])];
    }
  }
  return A;
}</pre>
```

6.3 Aho-Corasick

```
// quantidade de caracteres
const int ms = 1e6;
const int sigma = 26; // tamanho do alfabeto
int trie[ms][sigma], fail[ms], superfail[ms], terminal[ms], qtd;
void init() {
  qtd = 1;
  memset(trie[0], -1, sizeof trie[0]);
void add(string &s) {
  int node = 0;
  for (char ch : s) {
    int pos = val(ch); // no caso de alfabeto a-z: val(ch) = ch - 'a'
    if (trie[node][pos] == -1) {
     memset(trie[qtd], -1, sizeof trie[qtd]);
     terminal[qtd] = 0;
     trie[node][pos] = qtd++;
    node = trie[node][pos];
  ++terminal[node]; // trocar pela info que quiser
void buildFailure() {
  memset(fail, 0, sizeof(int) * qtd), memset(superfail, 0, sizeof(int) *
      qtd);
  queue<int> Q;
  Q.push(0);
  while (Q.size()) {
    int node = 0.front();
    () qoq.();
    for (int pos = 0; pos < sigma; ++pos) {</pre>
     int &v = trie[node][pos];
      int f = node == 0 ? 0 : trie[fail[node]][pos];
      // int sf = present[f] ? f : superfail[f];
      // present marks if that vertex is a terminal node or not
      // if summing up on terminal, doesn't work
      if (v == -1) {
```

```
v = f;
} else {
    fail[v] = f;

    // superfail[v] = sf;
    Q.push(v);
    // dar merge nas infos (por ex: terminal[v] += terminal[f])
}

void search(string &s) {
    int node = 0;
    for (char ch : s) {
        int pos = val(ch);
        node = trie[node][pos];
        // processar infos no no atual (por ex: ocorrencias += terminal[node])
}

// se tiver usando super fail, cuidado com o estado que voce ta, antes de subir pro sf, porque pode ser que o estado que ta nao seja no terminal
```

6.4 Algoritmo de Z

```
template <class T>
vector<int> ZFunc(const vector<T> &v) {
  vector<int> z(v.size(), 0);
  int n = (int) v.size(), a = 0, b = 0;
  if (!z.empty()) z[0] = n;
  for (int i = 1; i < n; i++) {
    int end = i; if (i < b) end = min(i + z[i - a], b);
    while(end < n && v[end] == v[end - i]) ++end;
    z[i] = end - i; if (end > b) a = i, b = end;
  }
  return z;
}
```

6.5 Suffix Array

```
vector<int> buildSa(const string& in) {
  int n = in.size(), c = 0;
  vector<int> temp(n), posBucket(n), bucket(n), bpos(n), out(n);
  for (int i = 0; i < n; i++) out [i] = i;
  sort(out.begin(), out.end(), [&](int a, int b) { return in[a] < in[b]; })</pre>
  for (int i = 0; i < n; i++) {</pre>
   bucket[i] = c;
    if (i + 1 == n || in[out[i]] != in[out[i + 1]]) c++;
  for (int h = 1; h < n && c < n; h <<= 1) {
    for (int i = 0; i < n; i++) posBucket[out[i]] = bucket[i];</pre>
    for (int i = n - 1; i >= 0; i--) bpos[bucket[i]] = i;
    for (int i = 0; i < n; i++) {
     if (out[i] >= n - h) temp[bpos[bucket[i]]++] = out[i];
    for (int i = 0; i < n; i++)
      if (out[i] >= h) temp[bpos[posBucket[out[i] - h]]++] = out[i] - h;
    for (int i = 0; i + 1 < n; i++) {
        int a = (bucket[i] != bucket[i + 1]) || (temp[i] >= n - h)
            || (posBucket[temp[i + 1] + h] != posBucket[temp[i] + h]);
        bucket[i] = c;
        c += a;
   bucket[n - 1] = c++;
    temp.swap(out);
```

```
return out:
vector<int> buildLcp(string s, vector<int> sa) {
 int n = (int) s.size();
 vector < int > pos(n), lcp(n, 0);
  for (int i = 0; i < n; i++) {
    pos[sa[i]] = i;
  int k = 0;
  for (int i = 0; i < n; i++) {
    if (pos[i] + 1 == n) {
      \mathbf{k} = 0;
      continue;
    int j = sa[pos[i] + 1];
    while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
    lcp[pos[i]] = k;
    k = \max(k - 1, 0);
  return 1cp;
```

6.6 Suffix Tree

```
#define fpos adla
const int inf = 1e9;
const int maxn = 1e4;
char s[maxn];
map<int, int> to[maxn];
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
int sz = 1, n = 0;
int make_node(int _pos, int _len)
    fpos[sz] = _pos;
    len [sz] = len;
    return sz++;
void go_edge()
    while(pos > len[to[node][s[n - pos]]])
        node = to[node][s[n - pos]];
        pos -= len[node];
void add_letter(int c)
    s[n++] = c;
    pos++;
    int last = 0;
    while (pos > 0)
        go edge();
        int edge = s[n - pos];
        int &v = to[node][edge];
        int t = s[fpos[v] + pos - 1];
        if(v == 0)
            v = make_node(n - pos, inf);
            link[last] = node;
            last = 0;
        else if(t == c)
            link[last] = node;
            return;
        else
```

```
int u = make_node(fpos[v], pos - 1);
    to[u][c] = make_node(n - 1, inf);
    to[u][t] = v;
    fpos[v] += pos - 1;
    len [v] -= pos - 1;
    v = u;
    link[last] = u;
    last = u;
}
if(node == 0)
    pos--;
else
    node = link[node];
}
//len[0] = inf
```

6.7 Suffix Automaton

```
int len[ms*2], link[ms*2], nxt[ms*2][sigma];
int sz, last;
void build(string &s) {
  len[0] = 0; link[0] = -1;
  sz = 1; last = 0;
  memset(nxt[0], -1, sizeof nxt[0]);
  for(char ch : s) {
    int c = ch-'a', cur = sz++;
len[cur] = len[last]+1;
    memset(nxt[cur], -1, sizeof nxt[cur]);
    int p = last;
    while (p != -1 && nxt[p][c] == -1) {
      nxt[p][c] = cur; p = link[p];
    if(p == -1) {
      link[cur] = 0;
    } else {
      int q = nxt[p][c];
      if(len[p] + 1 == len[q]) {
        link[cur] = q;
        len[sz] = len[p]+1; link[sz] = link[q];
        memcpy(nxt[sz], nxt[q], sizeof nxt[q]);
        while (p != -1 && nxt[p][c] == q) {
          nxt[p][c] = sz; p = link[p];
        link[q] = link[cur] = sz++;
    last = cur:
```

6.8 Manacher

```
return p;
} // pra cada centro o tamanho max do palindromo centrado ali, qualquer
    coisa printa a saida pra abacabaab
```

6.9 Polish Notation

```
inline bool isOp(char c) {
        return C=='+' || C=='-' || C=='*' || C=='/' || C=='^';
inline bool isCarac(char c) {
        return (c>='a' && c<='z') || (c>='A' && c<='Z') || (c>='0' && c<='9
int paren2polish(char* paren, char* polish) {
        map<char, int> prec;
        prec['('] = 0;
prec['+'] = prec['-'] = 1;
        prec['*'] = prec['/'] = 2;
        prec['^'] = 3;
        int len = 0;
        stack<char> op;
        for (int i = 0; paren[i]; i++) {
                if (isOp(paren[i])) {
                         while (!op.empty() && prec[op.top()] >= prec[paren[
                             i]]) {
                                 polish[len++] = op.top(); op.pop();
                         op.push(paren[i]);
                else if (paren[i] == '(') op.push('(');
                else if (paren[i]==')')
                         for (; op.top()!='('; op.pop())
                                 polish[len++] = op.top();
                         op.pop();
                else if (isCarac(paren[i]))
                         polish[len++] = paren[i];
        for(; !op.empty(); op.pop())
                polish[len++] = op.top();
        polish[len] = 0;
        return len;
```

6.10 String Hash

```
struct StringHashing {
  const uint64_t MOD = (1LL << 61) - 1;</pre>
  const int base = 31;
 vector<uint64_t> h, p;
 uint64_t modMul(uint64_t a, uint64_t b) {
    uint64_t 11 = (uint32_t)a, h1 = a >> 32, 12 = (uint32_t)b, h2 = b >>
        32;
    uint64 t 1 = 11 \star 12, m = 11 \star h2 + 12 \star h1, h = h1 \star h2;
    uint64_t ret = (1 \& MOD) + (1 >> 61) + (h << 3) + (m >> 29) + ((m <<
        35) >> 3) + 1;
    ret = (ret & MOD) + (ret >> 61);
    ret = (ret & MOD) + (ret >> 61);
    return ret - 1;
 uint64_t getKey(int 1, int r) { // [1, r]
    uint64 t res = h[r];
    if(1 > 0) res = (res + MOD - modMul(p[r - 1 + 1], h[1 - 1])) % MOD;
    return res;
```

```
wint64_t getInt(char c) {
    return c - 'a' + 1;
}

StringHashing(string &s) {
    int n = s.size();
    h.resize(n);
    p.resize(n);
    p[0] = 1;
    h[0] = getInt(s[0]);
    for(int i = 1; i < n; ++i) {
        p[i] = modMul(p[i - 1], base);
        h[i] = (modMul(h[i - 1], base) + getInt(s[i])) % MOD;
    }
};
</pre>
```

7 Miscellaneous

7.1 Ternary Search

```
// R
for(int i = 0; i < LOG; i++) {</pre>
  long double m1 = (A \times 2 + B) / 3.0;
  long double m2 = (A + 2 * B) / 3.0;
  if(f(m1) > f(m2))
    A = m1;
  else
    B = m2:
ans = f(A);
//Z
while (B - A > 4) {
  int m1 = (A + B) / 2;
  int m2 = (A + B) / 2 + 1;
 if(f(m1) > f(m2))
   A = m1;
  else
    B = m2:
ans = inf;
for (int i = A; i \le B; i++) ans = min(ans, f(i));
```

7.2 Count Sort

```
int H[(1<<15)+1], to[mx], b[mx];
void sort(int m, int a[]) {
    memset(H, 0, sizeof H);
    for (int i = 1; i <= m; i++) H[a[i] % (1<<15)]++;
    for (int i = 1; i < 1<<15; i++) H[i] += H[i-1];
    for (int i = m; i; i--) to[i] = H[a[i] % (1 << 15)]--;
    for (int i = 1; i <= m; i++) b[to[i]] = a[i];
    memset(H, 0, sizeof H);
    for (int i = 1; i <= m; i++) H[b[i]>>15]++;
    for (int i = 1; i < 1<<15; i++) H[i] += H[i-1];
    for (int i = m; i; i--) to[i] = H[b[i]>>15]--;
    for (int i = 1; i <= m; i++) a[to[i]] = b[i];
}</pre>
```

7.3 Random Number Generator

```
// mt19937_64 se LL
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
// Random_Shuffle
shuffle(v.begin(), v.end(), rng);
// Random number in interval
int randomInt = uniform_int_distribution(0, i)(rng);
double randomDouble = uniform_real_distribution(0, 1)(rng);
// bernoulli_distribution, binomial_distribution, geometric_distribution
// normal_distribution, poisson_distribution, exponential_distribution
```

7.4 Rectangle Hash

```
namespace {
  struct safe_hash {
    static uint64_t splitmix64(uint64_t x) {
      // http://xorshift.di.unimi.it/splitmix64.c
     x += 0x9e3779b97f4a7c15;
     x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
     x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
     return x ^ (x >> 31);
    size_t operator()(uint64_t x) const {
      static const uint64_t FIXED_RANDOM = std::chrono::steady_clock::now()
          .time_since_epoch().count();
      return splitmix64(x + FIXED RANDOM);
 };
struct rect {
  int x1, y1, x2, y2; // x1 < x2, y1 < y2
  rect () {};
  rect (int x1, int y1, int x2, int y2) : x1(x1), x2(x2), y1(y1), y2(y2)
  rect inter (rect other) {
    int x3 = max(x1, other.x1);
    int y3 = max(y1, other.y1);
    int x4 = min(x2, other.x2);
    int y4 = min(y2, other.y2);
    return rect(x3, y3, x4, y4);
  uint64 t get hash() {
    safe_hash sh;
    uint64_t ret = sh(x1);
    ret ^= sh(ret ^ y1);
    ret ^= sh(ret ^ x2);
    ret ^= sh(ret ^ y2);
    return ret;
};
```

7.5 Safe Hash

```
namespace {
    struct safe_hash {
        static uint64_t splitmix64(uint64_t x) {
            // http://xorshift.di.unimi.it/splitmix64.c
            x += 0x9e3779b97f4a7c15;
            x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
            x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
            return x ^ (x >> 31);
```

7.6 Unordered Map Tricks

```
// pair<int, int> hash function
struct HASH{
    size_t operator() (const pair<int,int>&x) const{
        return (size_t) x.first * 37U + (size_t) x.second;
    }
};

unordered_map<int,int>mp;
mp.reserve(1024);
mp.max_load_factor(0.25);
```

7.7 Iterate masks in bitcount order

```
for(int k = n-1; k >= 0; k--) {
  unsigned int i = (1 << k) -1;
  while(i < (1 << n)) {
    // do what you want
    unsigned int t = (i | (i - 1)) + 1;
    if(i == 0) break;
    i = t | ((((t & -t) / (i & -i)) >> 1) - 1);
  }
}
```

7.8 Submask Enumeration

```
for (int s=m; ; s=(s-1)&m) {
   ... you can use s ...
if (s==0) break;
}
```

7.9 Sum over Subsets DP

```
// F[i] = Sum of all A[j] where j is a submask of i
for(int i = 0; i < (1 < N); ++i)
   F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1 < N); ++mask) {
   if(mask & (1 < i))
        F[mask] += F[mask^(1 < i)];
}</pre>
```

7.10 Dates

```
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y){
```

```
1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
 n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x = 1461 * i / 4 - 31;
  j = 80 * x / 2447;
 d = x - 2447 * j / 80;
 x = \frac{1}{2} / 11;
 m = j + 2 - 12 * x;
 y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
```

7.11 Regular Expressions

```
import java.util.*;
import java.util.regex.*;

public class Main {
   public static String BuildRegex () {
      return "^" + sentence + "$";
   }

   public static void main (String args[]) {
      String regex = BuildRegex();
      // check pattern documentation
      Pattern pattern = Pattern.compile (regex);
      Scanner s = new Scanner(System.in);
      String sentence = s.nextLine().trim();
      boolean found = pattern.matcher(sentence).find()
}
```

7.12 Lat Long

```
/*
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
*/
struct l1
{
   double r, lat, lon;
};

struct rect
{
   double x, y, z;
};

ll convert(rect& P)
{
   ll Q;
   Q.r = sqrt(P.x*P.x*P.y*P.y*P.z*P.z);
   Q.lat = 180/M_PI*asin(P.z/Q.r);
```

```
Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
return Q;
}
rect convert(ll& Q)
{
    rect P;
    P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
    P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
    P.z = Q.r*sin(Q.lat*M_PI/180);
    return P;
}
```

7.13 Stable Marriage

```
std::vector<std::vector<int>> stableMarriage(std::vector<std::vector<int>>
    first, std::vector<std::vector<int>> second, std::vector<int> cap) {
        assert(cap.size() == second.size());
        int n = (int) first.size(), m = (int) second.size();
        // if O(N * M) first in memory, use table
        std::map<std::pair<int, int>, int> prio;
        std::vector<std::set<std::pair<int, int>>> current(m);
        for (int i = 0; i < n; i++) {
                std::reverse(first[i].begin(), first[i].end());
        for (int i = 0; i < m; i++) {
                for (int j = 0; j < (int) second[i].size(); <math>j++) {
                        prio[{second[i][j], i}] = j;
        for (int i = 0; i < n; i++) {
                int on = i;
                while(!first[on].empty()) {
                        int to = first[on].back();
                        first[on].pop_back();
                        if(cap[to]) {
                                cap[to]--;
                                assert(prio.count({on, to}));
                                current[to].insert({prio[{on, to}], on});
                                break;
                        assert(!current[to].empty());
                        auto it = current[to].end();
                        if(it->first > prio[{on, to}]) {
                                int nxt = it->second;
                                current[to].erase(it);
                                current[to].insert({prio[{on, to}], on});
                                on = nxt:
        std::vector<std::vector<int>> ans(m);
        for (int i = 0; i < m; i++) {
                for(auto it : current[i]) {
                        ans[i].push_back(it.second);
        return ans;
```

7.14 Mo

```
const int blk_sz = 170;
```

```
struct Query {
  int 1, r, idx;
  bool operator < (Query a) {</pre>
    if (1 / blk_sz == a.1 / blk_sz) {
      return r < a.r;
    return (1 / blk_sz) < (a.1 / blk_sz);
};
vector<Query> queries;
int a[MAXN], ans[MAXN], qnt[1000010];
int diff = 0;
void add(int x) {
 x = a[x];
  if (qnt[x] == 0) {
    diff++:
  qnt[x]++;
void remove(int x) {
 x = a[x];
  qnt[x]--;
  if (qnt[x] == 0) {
    diff--;
void mos() {
  int curr_1 = 0, curr_r = -1;
  sort(queries.begin(), queries.end());
  for (Query q : queries) {
    while (curr_1 > q.1) {
      curr_1--;
      add(curr_l);
    while (curr_r < q.r) {</pre>
      curr_r++;
      add(curr_r);
    while (curr_l < q.l) {</pre>
      remove(curr_1);
      curr_l++;
    while (curr_r > q.r) {
      remove (curr_r);
      curr_r--;
    ans[q.idx] = diff;
```

8 Teoremas e formulas uteis

8.1 Grafos

```
Formula de Euler: V - E + F = 2 (para grafo planar)
Handshaking: Numero par de vertices tem grau impar
Kirchhoff's Theorem: Monta matriz onde Mi,i = Grau[i] e Mi,j = -1 se houver
aresta i-j ou 0 caso contrario, remove uma linha e uma coluna
qualquer e o numero de spanning trees nesse grafo eh o det da matriz
```

Grafo contem caminho hamiltoniano se:
Dirac's theorem: Se o grau de cada vertice for pelo menos n/2
Ore's theorem: Se a soma dos graus que cada par nao-adjacente de vertices for pelo menos n
Boruvka's: enquanto grafo nao conexo, para cada componente conexa use a aresta que sai de menor custo.

Trees:

Tem Catalan(N) Binary trees de N vertices Tem Catalan (N-1) Arvores enraizadas com N vertices Caley Formula: n^(n-2) arvores em N vertices com label Prufer code: Cada etapa voce remove a folha com menor label e o label do vizinho eh adicionado ao codigo ate ter 2 vertices

Recuperar min cut eh ver se level[u] != -1 ai eh do lado do source Max Edge-disjoint paths: Max flow com arestas com peso 1

Max Node-disjoint paths: Faz a mesma coisa mas separa cada vertice em um com as arestas de chegadas e um com as arestas de saida e uma aresta de peso 1 conectando o vertice com aresta de chegada com ele mesmo com arestas de saida

Koniq's Theorem: minimum node cover = maximum matching se o grafo for bipartido, complemento eh o maximum independent set

Min vertex cover sao os vertices da particao do source que nao tao do lado do source do cut e os do sink que tao do lado do source, independent set o contrario

Min edge cover eh N - match, pega as arestas do match mais qualquer aresta dos outros vertices

Min Node disjoint path cover: formar grafo bipartido de vertices duplicados , onde aresta sai do vertice tipo A e chega em tipo B, entao o path cover eh N - matching

Min General path cover: Mesma coisa mas colocando arestas de A pra B sempre que houver caminho de A pra B

Dilworth's Theorem: Min General Path cover = Max Antichain (set de vertices tal que nao existe caminho no grafo entre vertices desse set)

Hall's marriage: um grafo tem um matching completo do lado X se para cada subconjunto W de X,

 $|W| \leftarrow |vizinhosW|$ onde |W| eh quantos vertices tem em W

feasible flow in a network with both upper and lower capacity constraints, no source or sink: capacities are changed to upper bound - lower bound . Add a new source and a sink. let M[v] = (sum of lower bounds of lower bouningoing edges to v) - (sum of lower bounds of outgoing edges from v). For all v, if M[v] > 0 then add edge (S, v) with capcity M, otherwise add (v,T) with capacity -M. If all outgoing edges from S are full, then a feasible flow exists, it is the flow plus the original lower_bounds

8.2 Math

Goldbach's: todo numero par n > 2 pode ser representado com n = a + b onde a e b sao primos

Twin prime: existem infinitos pares p, p + 2 onde ambos sao primos Legendre's: sempre tem um primo entre n^2 e (n+1)^2

Lagrange's: todo numero inteiro pode ser inscrito como a soma de 4

Zeckendorf's: todo numero pode ser representado pela soma de dois numeros de fibonnacis diferentes e nao consecutivos

Euclid's: toda tripla de pitagoras primitiva pode ser gerada com

 $(n^2 - m^2, 2nm, n^2+m^2)$ onde n, m sao coprimos e um deles eh par Wilson's: $n \in P$ primo quando $(n-1)! \mod n = n - 1$

Mcnugget: Para dois coprimos \mathbf{x} , \mathbf{y} o maior inteiro que nao pode ser escrito como ax + by eh (x-1)(y-1)/2

Fermat: Se p eh primo entao a(p-1) % p = 1

Se x e m tambem forem coprimos entao $x^k % m = x^k (k \mod (m-1)) % m$ Euler's theorem: $x^{(phi(m))} \mod m = 1$ onde phi(m) eh o totiente de euler

Chinese remainder theorem:

Para equacoes no formato $x = a1 \mod m1$, ..., $x = an \mod mn$ onde todos os pares m1, ..., mn sao coprimos

Deixe Xk = m1*m2*..*mn/mk e Xk^-1 mod mk = inverso de Xk mod mk, entao $x = somatorio com k de 1 ate n de ak*Xk*(Xk,mk^-1 mod mk)$

Para achar outra solucao so somar m1*m2*..*mn a solucao existente

```
Catalan number: exemplo expressoes de parenteses bem formadas
C0 = 1, Cn = somatorio de <math>i=0 \rightarrow n-1 de Ci*C(n-1+1)
outra forma: Cn = (2n \text{ escolhe } n)/(n+1)
Bertrand's ballot theorem: p votos tipo A e q votos tipo B com p>q, prob de
      em todo ponto ter mais As do que Bs antes dele = (p-q)/(p+q)
Se puder empates entao prob = (p+1-q)/(p+1), para achar quantidade de
    possibilidades nos dois casos basta multiplicar por (p + q escolhe q)
Propriedades de Coeficientes Binomiais:
Somatorio de k = 0 \rightarrow m de (-1)^k * (n escolhe k) = (-1)^m * (n -1 escolhe m)
(N \text{ escolhe } K) = (N \text{ escolhe } N-K)
(N \text{ escolhe } K) = N/K * (n-1 \text{ escolhe } k-1)
Somatorio de k = 0 \rightarrow n de (n escolhe k) = 2^n
Somatorio de m = 0 \rightarrow n de (m = 0) = (n+1) escolhe k + 1
Somatorio de k = 0 \rightarrow m de (n+k) escolhe k) = (n+m+1) escolhe m)
Somatorio de k = 0 \rightarrow n de (n \text{ escolhe } k)^2 = (2n \text{ escolhe } n)
Somatorio de k = 0 ou 1 \rightarrow n de k*(n escolhe k) = n * 2^(n-1)
Somatorio de k = 0 \rightarrow n de (n-k \text{ escolhe } k) = \text{Fib}(n+1)
Hockey-stick: Somatorio de i = r \rightarrow n de (i escolhe r) = (n + 1) escolhe r +
Vandermonde: (m+n \ escolhe \ r) = somatorio \ de \ k = 0 -> r \ de \ (m \ escolhe \ k) * (
    n = scolher - k
Burnside lemma: colares diferentes nao contando rotacoes quando m = cores e
     n = comprimento
(m^n + somatorio i = 1 - > n-1 de m^gcd(i, n))/n
Distribuicao uniforme a,a+1, ..., b Expected[X] = (a+b)/2
Distribuicao binomial com n tentativas de probabilidade p, X = sucessos:
    P(X = x) = p^x * (1-p)^(n-x) * (n escolhe x) e E[X] = p*n
Distribuicao geometrica onde continuamos ate ter sucesso, X = tentativas:
    P(X = x) = (1-p)^(x-1) * p \in E[X] = 1/p
Linearity of expectation: Tendo duas variaveis X e Y e constantes a e b, o
    valor esperado de aX + bY = a*E[X] + b*E[X]
V(X) = E((X-u)^2)
V(X) = E(X^2) - E(X^2)
PG: a1 * (q^n - 1)/(q - 1)
Mobius Inverse: Sum(d|n): mobius(d) = [n = 1] (expressao booleana)
Soma dos cubos de 1 a N = a^2 onde a = somatorio de 1 a N
Lindstrom-Gessel-Viennot: quantidade de caminhos disjuntos nas linhas do
    grid eh o determinante da matriz de gnts caminhos
```

8.3 Geometry

s = (a+b+c+d)/2

```
Formula de Euler: V - E + F = 2
Pick Theorem: Para achar pontos em coords inteiras num poligono Area = i +
    b/2 - 1 onde i eh o o numero de pontos dentro do poligono e b de
    pontos no perimetro do poligono
Two ears theorem: Todo poligono simples com mais de 3 vertices tem pelo
    menos 2 orelhas, vertices que podem ser removidos sem criar um
    crossing, remover orelhas repetidamente triangula o poligono
Incentro triangulo: (a(Xa, Ya) + b(Xb, Yb) + c(Xc, Yc))/(a+b+c) onde a =
    lado oposto ao vertice a, incentro eh onde cruzam as bissetrizes, eh o
     centro da circunferencia inscrita e eh equidistante aos lados
Delaunay Triangulation: Triangulacao onde nenhum ponto esta dentro de
    nenhum circulo circunscrito nos triangulos
Eh uma triangulacao que maximiza o menor angulo e a MST euclidiana de um
    conjunto de pontos eh um subconjunto da triangulacao
```

Brahmagupta's formula: Area cyclic quadrilateral

area = sqrt((s-a)*(s-b)*(s-c)*(s-d))

 $d = 0 \Rightarrow area = sqrt((s-a)*(s-b)*(s-c)*s)$

8.4 Dynamic Programming

```
Divide and conquer - dp[i][j] = mink < j{dp[i - 1][k] + C[k][j]} dividir o subsegmento ate j em i segmentos com custo, valido se A[i][j] <= A[i][j+1]

Knuth - p[i][j] = mini < k < j{dp[i][k] + dp[k][j]} + C[i][j], valido se A[i, j - 1] <= A[i][j] <= A[i+1, j]

onde A[i][j] eh o menor k que da a resposta otima slope trick - funcao piecewise linear convexa, descrita pelos pontos de mudanca de slope (multiset/heap)

lembre que existe fft, cht, aliens trick e bitset
```

8.5 Mersenne's Primes

Primos de Mersenne 2^n - 1
Lista de Ns que resultam nos primeiros 41 primos de Mersenne:
2; 3; 5; 7; 13; 17; 19; 31; 61; 89; 107; 127; 521; 607; 1.279; 2.203;
2.281; 3.217; 4.253; 4.423; 9.689; 9.941; 11.213; 19.937; 21.701;
23.209; 44.497; 86.243; 110.503; 132.049; 216.091; 756.839; 859.433;
1.257.787; 1.398.269; 2.976.221; 3.021.377; 6.972.593; 13.466.917;
20.996.011; 24.036.583;