

Case study: Trading system



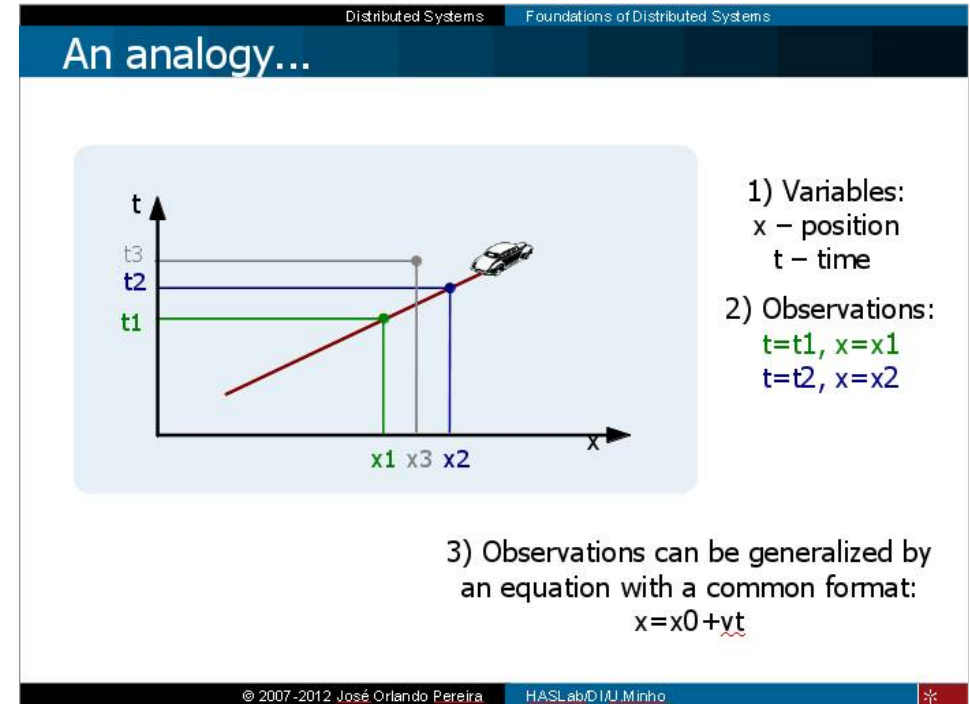
- Input: offer to buy or sell something
- No output until a suitable buyer/seller pair is found
 - The trading system never holds stock
- Output: confirmation

Case study: Trading system

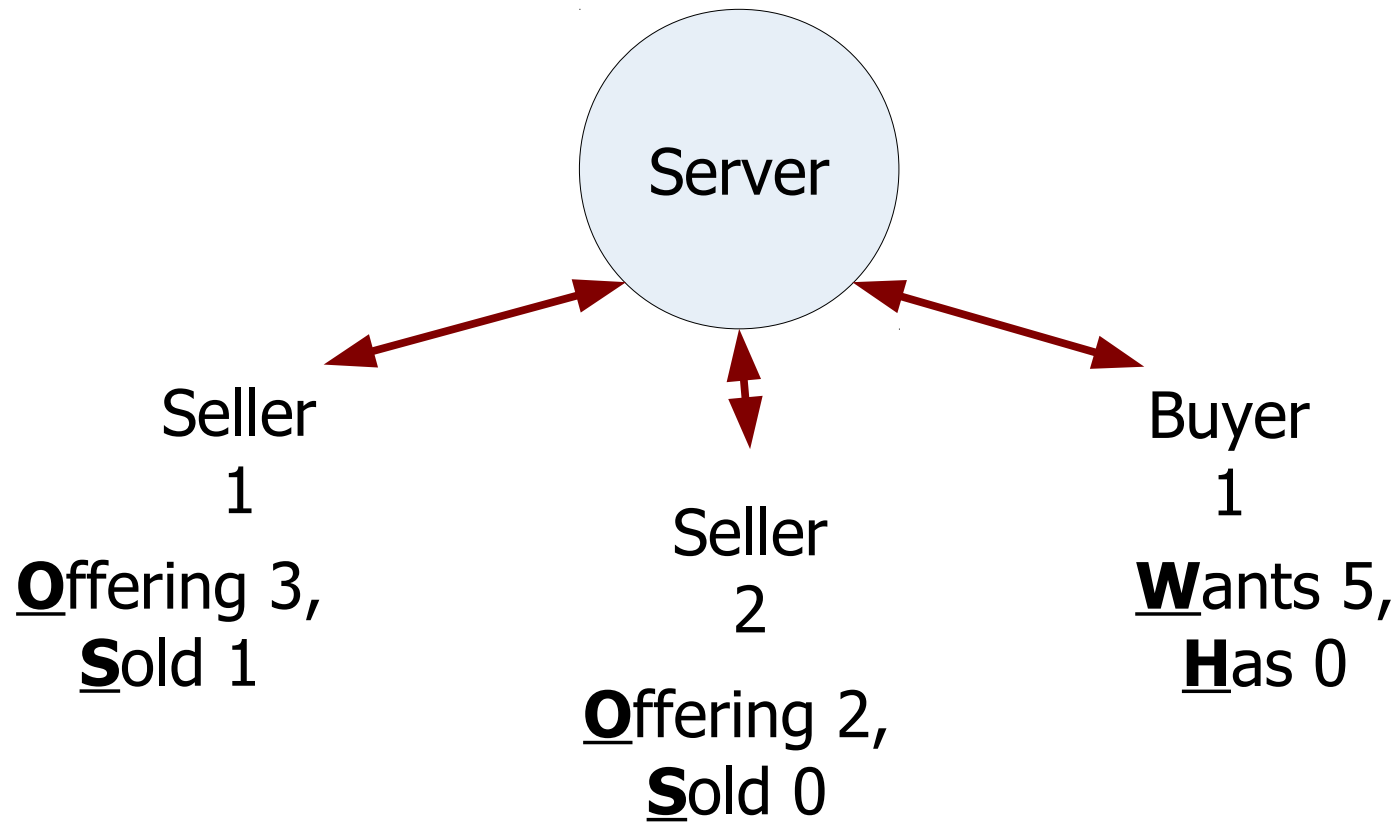
- What matters:
 - We don't buy/sell more than what has been offered/requested
 - If there are sellers and buyers for at least k items, eventually k items are sold and bought
 - If multiple buyers/sellers are competing, make sure no one is left behind
- What doesn't matter:
 - How the server is implemented
 - If there is a server at all...
 - ...

Step 1

- Select variables, observe and record...
- Major pitfall:
 - Programming variables vs. Model variables



Distributed state



Variable	O1	S1	O2	S2	W1	H1
Value
Value	3	1	2	0	5	0
Value

Distributed trace

- Sequence of states:
- Sequence of actions:

O1	S1	O2	S2	W1	H1
3	1	2	0	5	0
3	1	2	1	5	0
3	1	2	1	5	1
3	1	2	1	5	2
3	2	2	1	5	2
4	2	2	1	5	2
...



s2 sells 1
 b1 buys 1
 b1 buys 1
 s1 sells 1
 s1 offers 1
 ...

Distributed trace

- What actions are atomic?

O1	S1	O2	S2	W1	H1
3	1	2	0	5	0
3	1	2	1	5	0
4	1	2	1	5	1
4	1	2	1	5	2
...



Distributed trace

- How are alternative actions ordered?

O1	S1	O2	S2	W1	H1
3	1	2	0	5	0
3	1	2	1	5	0
3	1	2	1	5	1
3	1	2	1	5	2
3	2	2	1	5	2
4	2	2	1	5	2
...

O1	S1	O2	S2	W1	H1
3	1	2	0	5	0
3	1	2	1	5	0
3	1	2	1	5	1
3	2	2	1	5	1
3	2	2	1	5	2
4	2	2	1	5	2
...

Conclusion

- Our specifications are sets of possible traces
- Assuming a distributed system means that:
 - All possible orderings of events have to be considered
 - Large sets are required even for simple specifications
- Not necessarily a problem:
 - Recall: $x = x_0 + vt$ describes infinite “traces”

Step 2

- What can be observed?

An invalid trace


O1	S1	O2	S2	W1	H1
3	1	2	0	5	0
3	1	2	1	5	0
3	1	2	1	5	1
3	1	2	0	5	1
3	2	2	1	5	2
4	2	2	1	5	2
...



An invalid trace

- No action can explain this step:

O1	S1	O2	S2	W1	H1
3	1	2	0	5	0
3	1	2	1	5	0
3	1	2	1	5	1
3	1	2	0	5	1
3	2	2	1	5	2
4	2	2	1	5	2
...




An invalid trace

O1	S1	O2	S2	W1	H1
3	1	4	0	5	0
3	1	4	4	5	0
3	1	4	4	5	1
3	1	4	4	5	1
3	2	4	4	5	2
4	2	4	4	5	2
...

An invalid trace

- No previous sequence of actions can explain this state:

O1	S1	O2	S2	W1	H1
3	1	4	0	5	0
3	1	4	4	5	0
3	1	4	4	5	1
3	1	4	4	5	1
3	2	4	4	5	2
4	2	4	4	5	2
...



Safety property

- “Nothing bad ever happens...”
- Can be identified on a finite prefix of the trace


An invalid trace

O1	S1	O2	S2	W1	H1
1	0	1	0	1	0
1	1	1	0	1	0
1	1	1	0	1	1
1	1	1	0	2	1
1	1	1	0	2	1
1	1	1	0	2	1
...

An invalid trace

- Deadlock? Must see full trace...

O1	S1	O2	S2	W1	H1
1	0	1	0	1	0
1	1	1	0	1	0
1	1	1	0	1	1
1	1	1	0	2	1
1	1	1	0	2	1
1	1	1	0	2	1
...



An invalid trace

O1	S1	O2	S2	W1	H1
1	0	1	0	1	0
1	0	1	0	1	1
1	1	1	0	1	1
2	1	1	0	1	1
2	1	1	0	2	1
2	1	1	0	2	2
2	2	1	0	2	2
3	2	1	0	2	2
3	2	1	0	3	2
3	2	1	0	3	3
3	3	1	0	3	3
...

An invalid trace

- Does Seller 2 get a fair chance?
- Must see full trace..

O1	S1	O2	S2	W1	H1
1	0	1	0	1	0
1	0	1	0	1	1
1	1	1	0	1	1
2	1	1	0	1	1
2	1	1	0	2	1
2	1	1	0	2	2
2	2	1	0	2	2
3	2	1	0	2	2
3	2	1	0	3	2
3	2	1	0	3	3
3	3	1	0	3	3
...



Liveness

- “Something good eventually happens...”
- Cannot be identified on a finite prefix of the trace

Specification = Safety + Liveness

- Our specifications can always be decomposed in:
 - Safety
 - +
 - Liveness
- How to write them down?
- How to derive them?

Roadmap

- Recall:
 - A specification is a set of traces
- How to write a compact expression that generates such set?
- How to determine if such specification (in compact format) satisfies safety and liveness properties?

State machine

- **BuyOffer(i):**

- Pre-condition:

- True

- Effect:

- $O_i := O_i + 1$

- **Buy(i):**

- Pre-condition:

- $\Sigma H < \min(\Sigma O, \Sigma W)$

- $H_i < W_i$

- Effect:

- $H_i := H_i + 1$

- **SellOffer(i):**

- Pre-condition:

- True

- Effect:

- $W_i := W_i + 1$

- **Sell(i):**

- Pre-condition:

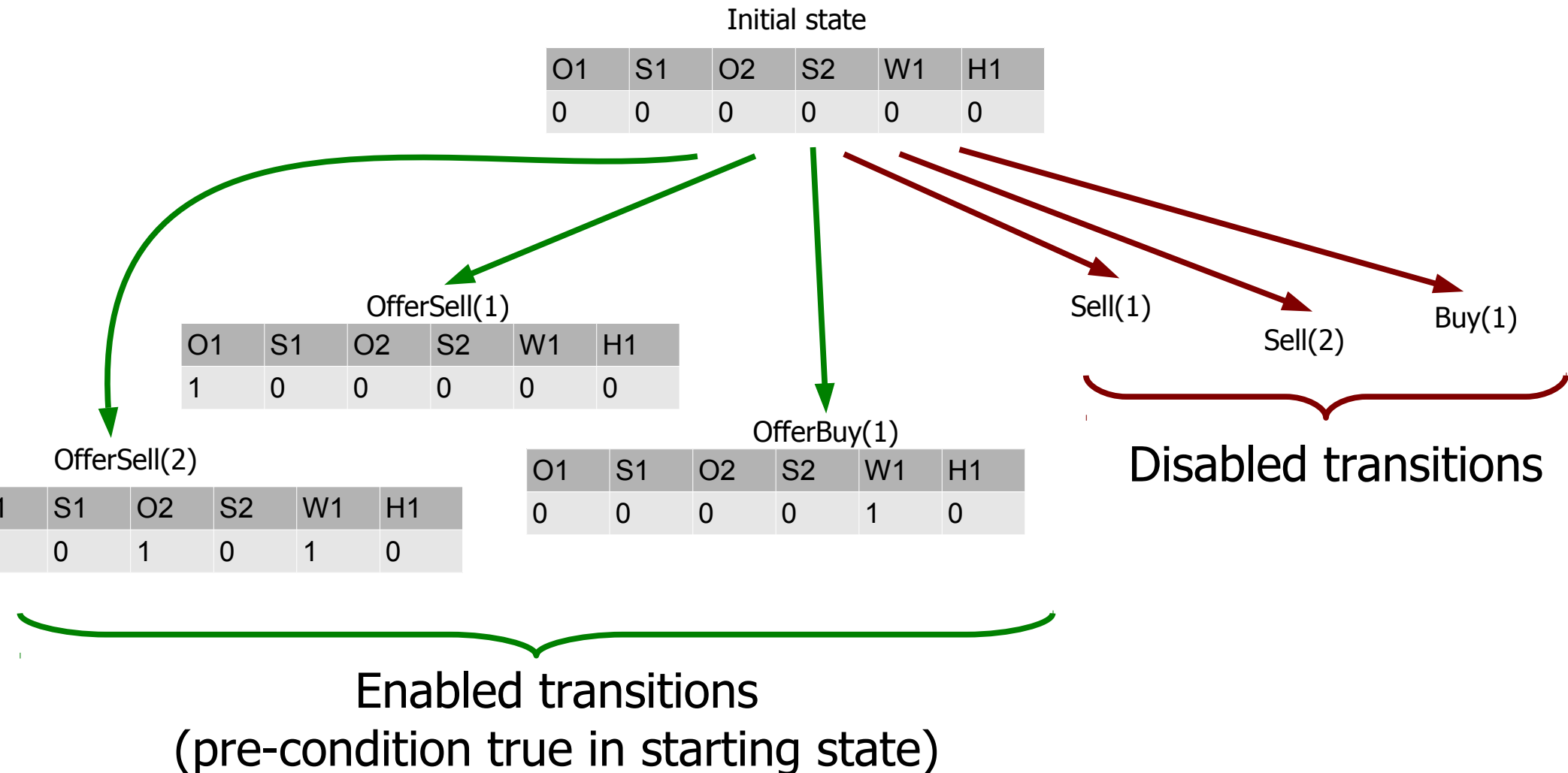
- $\Sigma S < \min(\Sigma O, \Sigma W)$

- $S_i < O_i$

- Effect:

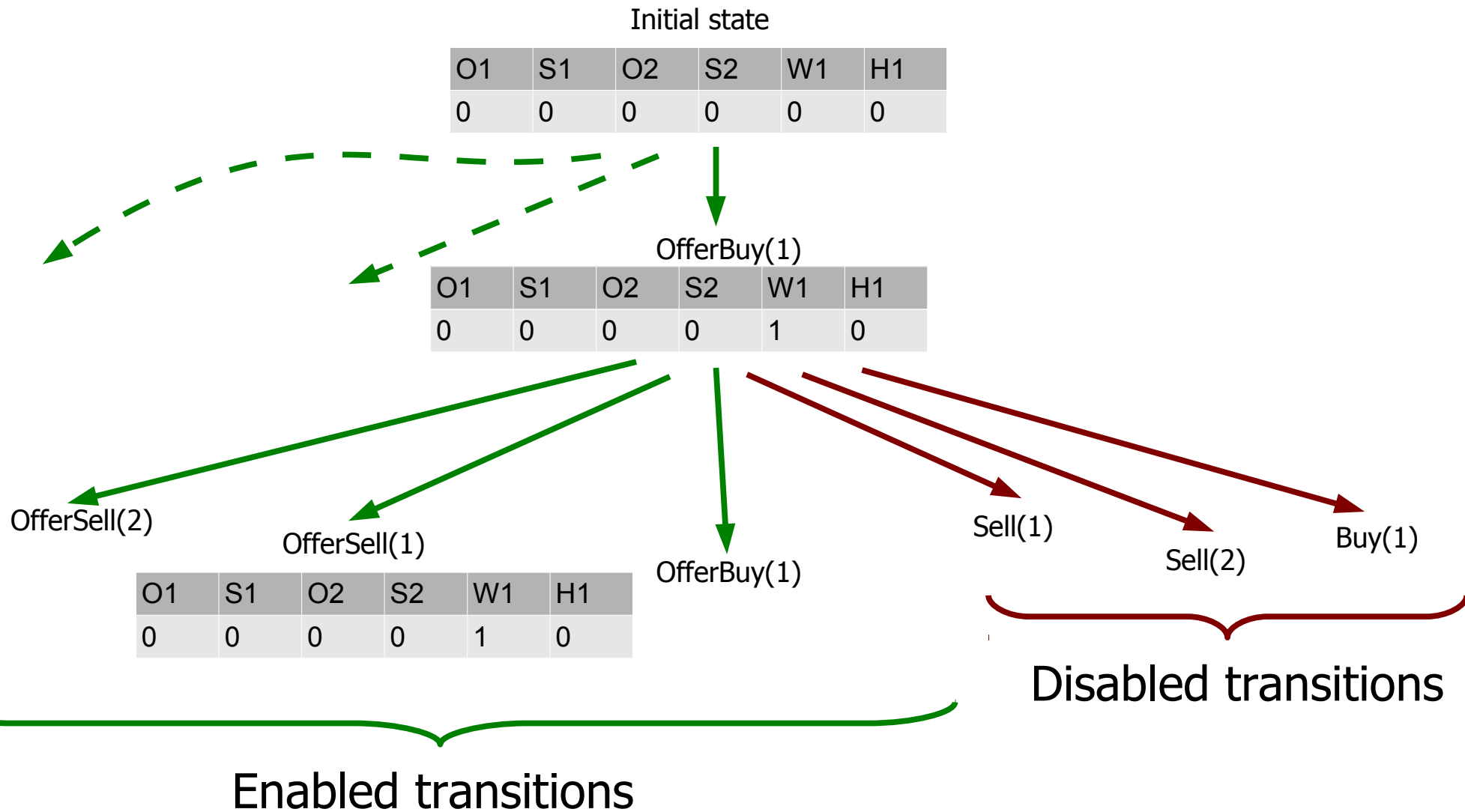
- $S_i := S_i + 1$

Safe traces

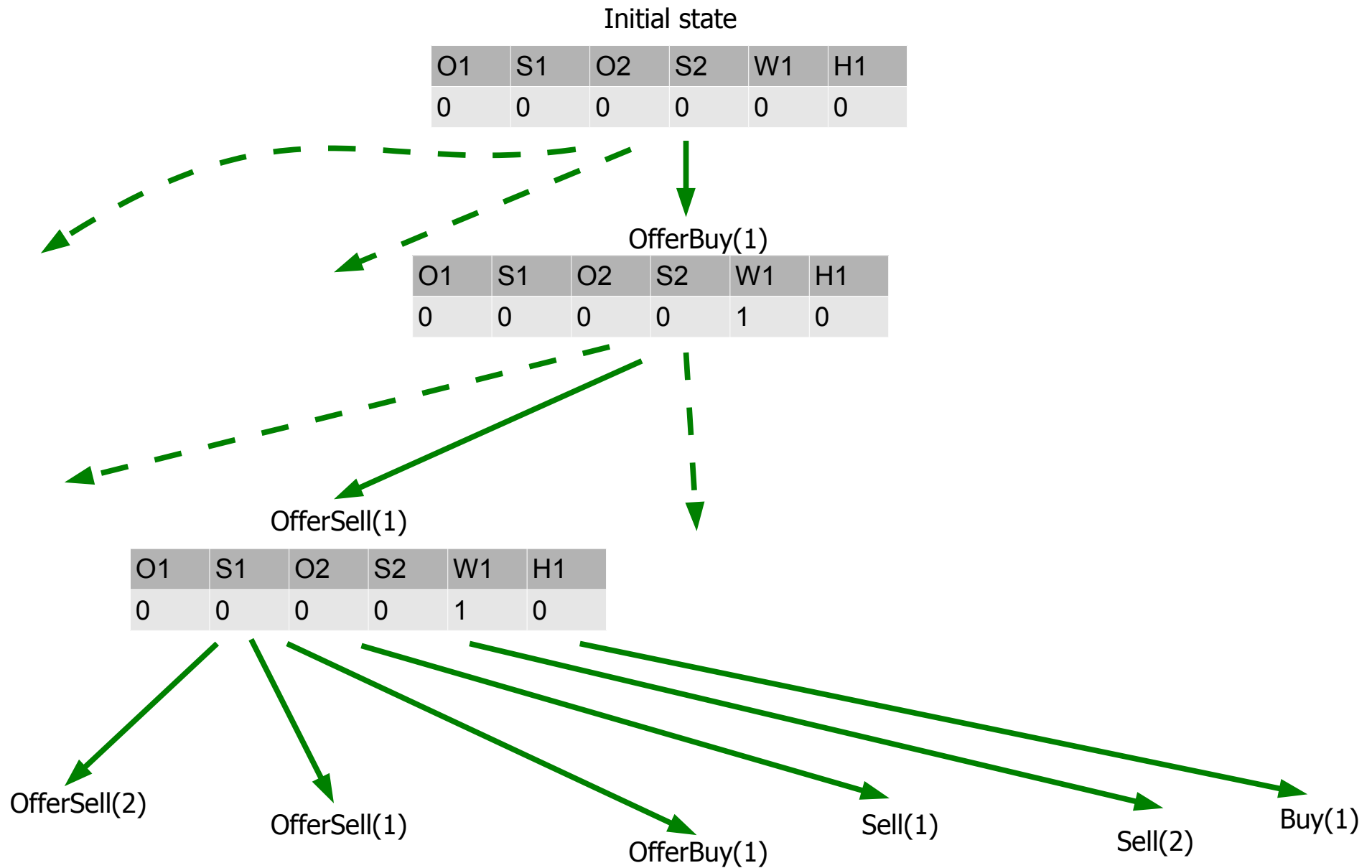


Three new safe traces found!

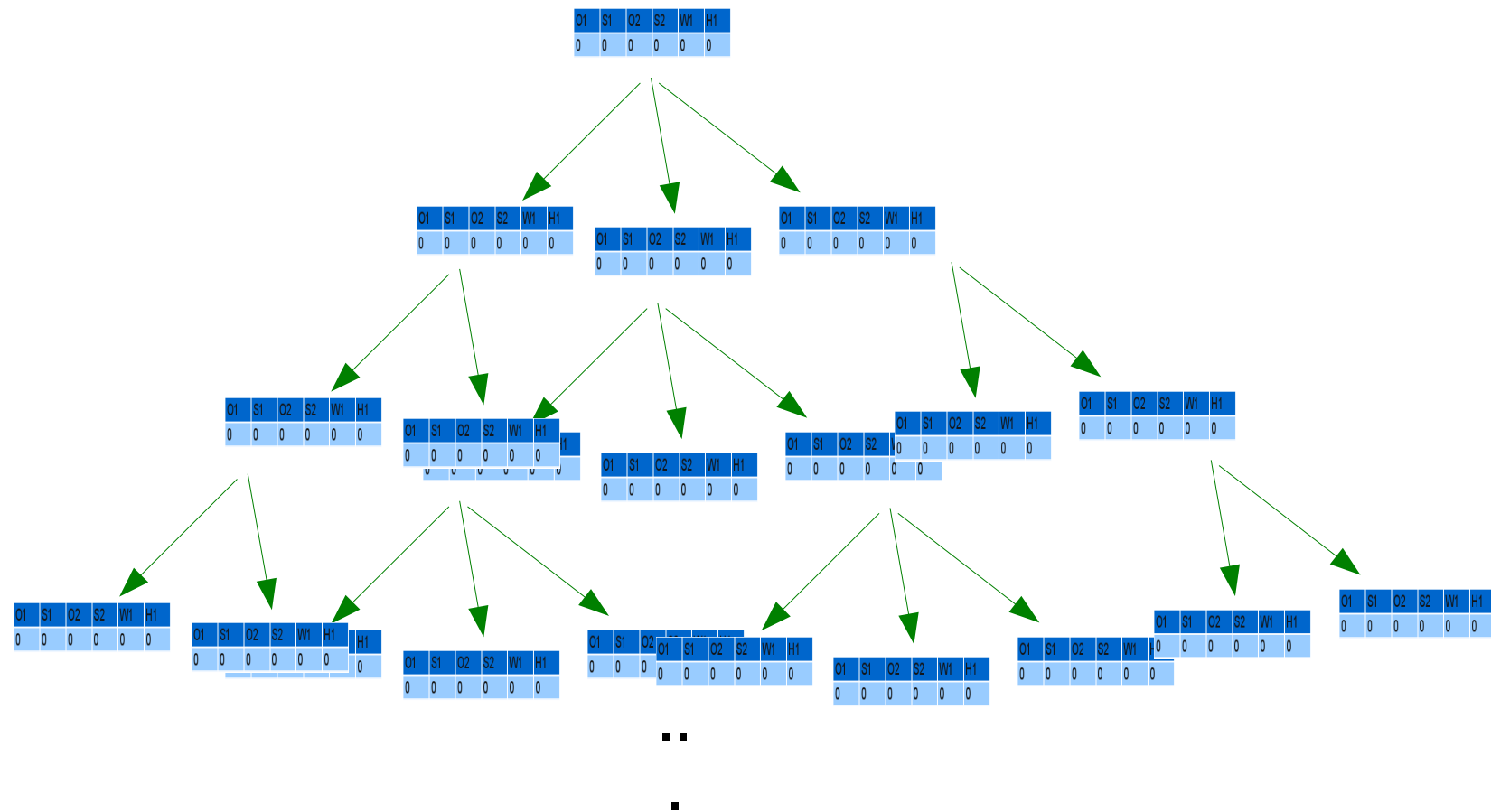
Safe traces



Safe traces



Safe traces



Safe traces

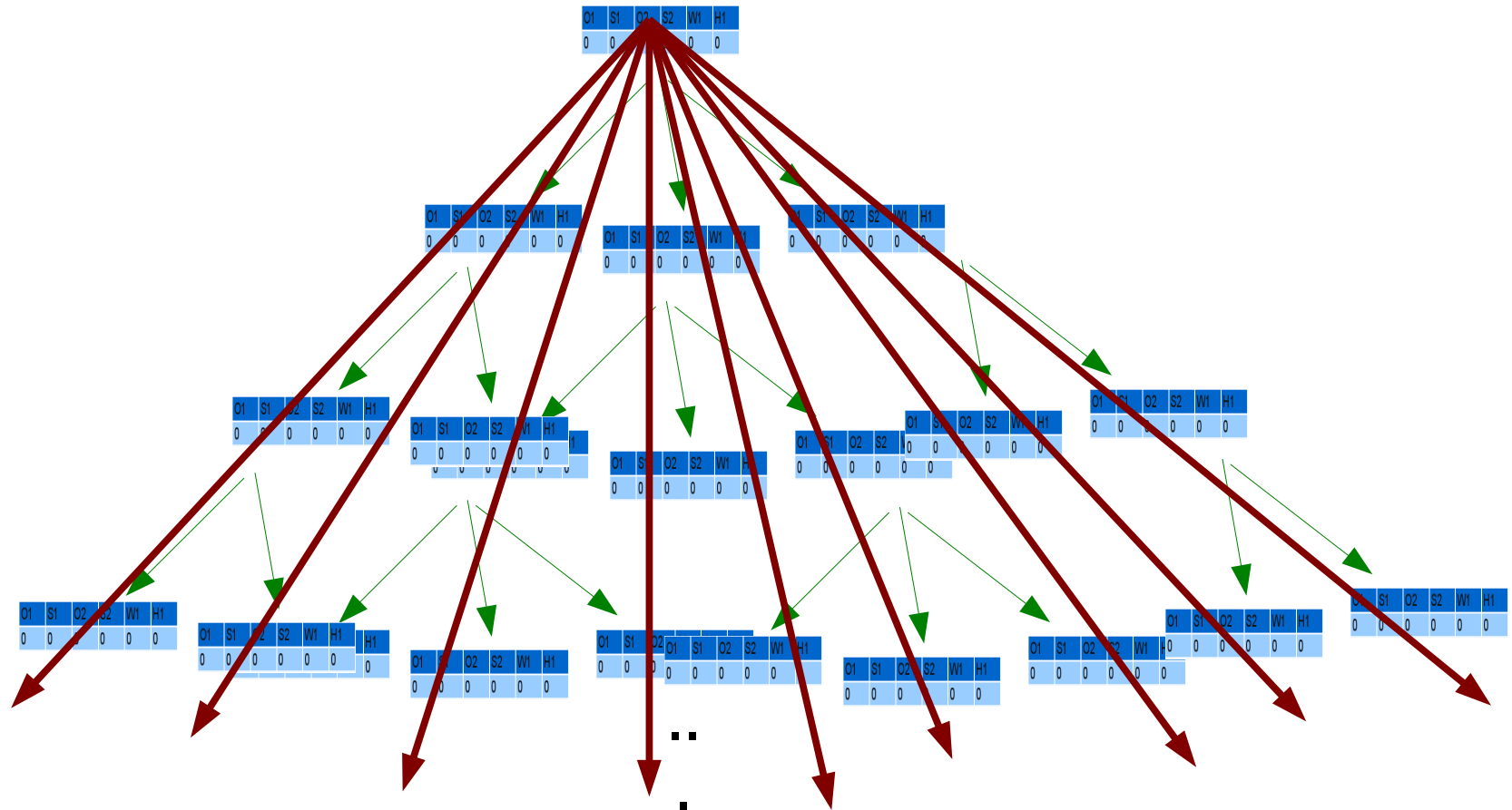
- Start with one trace for each initial state
- Repeat:
 - For any known safe trace
 - For any enabled transition
 - Append resulting state
 - Add new trace to set
- Most likely an infinite set...

Safety properties

- Does this specification meet the desired properties?
- Safety property:
 - We don't buy more than what has been offered (i.e. $\sum H \leq \sum O$)
- Observation:
 - This is a state invariant

Safety properties

!?



- Cannot evaluate all possible traces, but...

Base step

O1	S1	O2	S2	W1	H1
0	0	0	0	0	0

$\Sigma H \leq \Sigma O$? Yes!

Induction step

O1	S1	O2	S2	W1	H1
?	?	$\sum H \leq \sum O!$?	?	?



SellOffer(i)		Sell(i)		BuyOffer(i)	
O1	S1	O2	S2	W1	H1
?	?	?	?	?	?

$\sum H \leq \sum O?$ Yes!
 (all H unchanged, O never decreases)

Induction step

O1	S1	O2	S2	W1	H1
?	?	?	?	?	?

$$\Sigma H \leq \Sigma O!$$


Buy(i)

O1	S1	O2	S2	W1	H1
?	?	?	?	?	?+1

$$\Sigma H \leq \Sigma O?$$

Induction step

- We know:

Induction

- $\Sigma H \leq \Sigma O$

Pre-condition

- $\Sigma H < \min(\Sigma O, \Sigma W)$

- $\Sigma H < \Sigma O$

- $\Sigma H < \Sigma W$

- $H_i < W_i$

Effect

- $H_i := H_i + 1$

- Everything else unchanged

- We want to know if:

- $\Sigma H \leq \Sigma O$

Safety properties

O1	S1	O2	S2	W1	H1
0	0	0	0	0	0

O1	S1	O2	S2	W1	H1
?	?	?	?	?	?

$$\sum H \leq \sum O!$$



SellOffer(i)

Sell(i)

BuyOffer(i)

O1	S1	O2	S2	W1	H1
?	?	?	?	?	?

O1	S1	O2	S2	W1	H1
?	?	?	?	?	?

$$\sum H \leq \sum O!$$



Buy(i)

O1	S1	O2	S2	W1	H1
?	?	?	?	?	?+1



- Can easily inspect:
 - All initial states
 - All transitions (many are trivial)

Safety properties

- A state invariant can be proved by induction
- Other safety properties can be translated into state invariants by:
 - Strengthening the property
 - Recording past trace as state