

Complete Primitive Recursive Functions Guide: Every Exercise Type

Definition of Primitive Recursive Functions (PR)

Formal Definition: The class PR of primitive recursive functions is the **smallest** class of functions containing:

Base Functions:

1. **Zero function:** $0: \mathbb{N} \rightarrow \mathbb{N}$ defined by $0(x) = 0$ for all $x \in \mathbb{N}$
2. **Successor function:** $s: \mathbb{N} \rightarrow \mathbb{N}$ defined by $s(x) = x + 1$ for all $x \in \mathbb{N}$
3. **Projection functions:** $U_j^k: \mathbb{N}^k \rightarrow \mathbb{N}$ defined by $U_j^k(x_1, \dots, x_k) = x_j$

Operations (Closed Under):

1. **Generalized Composition:** If $f_1, \dots, f_n: \mathbb{N}^k \rightarrow \mathbb{N}$ and $g: \mathbb{N}^n \rightarrow \mathbb{N}$, then $h(\vec{x}) = g(f_1(\vec{x}), \dots, f_n(\vec{x}))$
2. **Primitive Recursion:** Given $f: \mathbb{N}^k \rightarrow \mathbb{N}$ and $g: \mathbb{N}^{(k+2)} \rightarrow \mathbb{N}$, define $h: \mathbb{N}^{(k+1)} \rightarrow \mathbb{N}$ by:

$$\begin{aligned} h(\vec{x}, 0) &= f(\vec{x}) \\ h(\vec{x}, y+1) &= g(\vec{x}, y, h(\vec{x}, y)) \end{aligned}$$

Universal Exercise Methodology

Step 1: Identify Exercise Type

Look for these patterns in the exercise text:

- **"prove that function ... is primitive recursive"** → Construction exercise
- **"characteristic function of set ..."** → Boolean/set exercise
- **"using only the definition"** → Must build from base functions
- **Power/exponential functions** → Repeated operations
- **Arithmetic operations** → Basic building blocks

Step 2: Choose Construction Strategy

- **Direct primitive recursion** for recursive patterns
- **Composition** for combining known PR functions
- **Helper functions** for complex constructions

Step 3: Write Primitive Recursive Definition

- Base case: $h(\dots, 0) = \dots$
- Recursive case: $h(\dots, y+1) = \dots$
- Ensure you only use known PR functions

Step 4: Verify Primitive Recursion

- Check base case uses only previous PR functions
- Check recursive case uses $g(\vec{x}, y, h(\vec{x}, y))$ pattern
- Verify all helper functions are also PR

Complete Exercise Type Classification

TYPE 1: BASIC ARITHMETIC FUNCTIONS

Pattern: Addition, Multiplication, Powers

Exercise: "Prove that $\text{sum}(x,y) = x + y$ is primitive recursive"

Step-by-Step Solution:

```
sum(x, 0) = x                [base case: f(x) = U11(x) = x]
sum(x, y+1) = sum(x, y) + 1  [recursive: g(x,y,z) = s(z)]
```

Why this works:

- Base case: Adding 0 to x gives x (identity)
- Recursive case: To add y+1, we add y then add 1 more
- Uses only successor function and previous result

Template for similar functions:

```
h(x, 0) = initial_value(x)
h(x, y+1) = operation(x, y, h(x,y))
```

Example: Multiplication

```
mult(x, 0) = 0                [base: f(x) = 0]
mult(x, y+1) = mult(x, y) + x [recursive: g(x,y,z) = sum(z, x)]
```

Example: Power Function (2^y)

```
pow2(0) = 1                    [base: f() = 1]
pow2(y+1) = double(pow2(y))    [recursive: g(y,z) = double(z)]
```

where $\text{double}(x)$ is primitive recursive:

```
double(0) = 0
double(y+1) = double(y) + 2
```

TYPE 2: CHARACTERISTIC FUNCTIONS

Pattern: $\chi_A(x) = 1$ if $x \in A$, 0 otherwise

Exercise: "Prove that χ_P (characteristic of even numbers) is primitive recursive"

Step-by-Step Solution:

$$\begin{array}{ll} \chi_P(0) = 1 & [0 \text{ is even}] \\ \chi_P(y+1) = \text{sg}(\chi_P(y)) & [\text{flip the previous result}] \end{array}$$

Helper function sg (negated sign):

$$\begin{array}{l} \text{sg}(0) = 1 \\ \text{sg}(y+1) = 0 \end{array}$$

Why this works:

- Start with 0 being even ($\chi_P(0) = 1$)
- For each successor, flip the result using sg
- Pattern: even \rightarrow odd \rightarrow even \rightarrow odd...

Example: Characteristic of $\{2^n - 1 \mid n \in \mathbb{N}\}$

Strategy: First define $a(n) = 2^n - 1$, then check membership

$$\begin{array}{ll} a(0) = 0 & [2^0 - 1 = 0] \\ a(n+1) = 2 \cdot a(n) + 1 & [2^{(n+1)} - 1 = 2(2^n - 1) + 1] \end{array}$$

$$\begin{array}{ll} \text{chk}(x, 0) = \text{sg}(x) & [\text{check if } x = a(0) = 0] \\ \text{chk}(x, m+1) = \text{chk}(x, m) + \text{eq}(x, a(m+1)) & \end{array}$$

$$\chi_A(x) = \text{chk}(x, x) \quad [\text{check membership up to } x]$$

TYPE 3: TRUNCATED SUBTRACTION AND ABSOLUTE DIFFERENCE

Pattern: $x \dot{-} y = \max(0, x-y)$

Exercise: "Prove that $p_2(y) = |y - 2|$ is primitive recursive"

Step-by-Step Construction: First build $p_1(y) = |y - 1| = y \dot{-} 1$:

$$\begin{array}{ll} p_1(0) = 1 & [|0 - 1| = 1] \\ p_1(y+1) = y & [|y+1 - 1| = y] \end{array}$$

Then build $p_2(y) = |y - 2|$:

$$\begin{array}{ll}
 p_2(0) = 2 & [|0 - 2| = 2] \\
 p_2(y+1) = p_1(y) & [|y+1 - 2| = |y - 1|]
 \end{array}$$

General Pattern for $|x - k|$:

$$\begin{array}{l}
 p_k(0) = k \\
 p_k(y+1) = p_{k-1}(y)
 \end{array}$$

TYPE 4: DIVISION AND REMAINDER FUNCTIONS

Pattern: Half function (integer division)

Exercise: "Prove that $\text{half}(x) = \lfloor x/2 \rfloor$ is primitive recursive"

Step-by-Step Solution: First define remainder mod 2:

$$\begin{array}{ll}
 \text{rm}_2(0) = 0 & \\
 \text{rm}_2(x+1) = \text{sg}(\text{rm}_2(x)) & [\text{alternates } 0,1,0,1,\dots]
 \end{array}$$

Then define half:

$$\begin{array}{ll}
 \text{half}(0) = 0 & \\
 \text{half}(x+1) = \text{half}(x) + \text{rm}_2(x) & [\text{add 1 every two steps}]
 \end{array}$$

Why this works:

- $\text{rm}_2(x)$ gives the remainder when x is divided by 2
- We add 1 to $\text{half}(x)$ only when x is even ($\text{rm}_2(x) = 0$)
- This gives us $\lfloor x/2 \rfloor$

TYPE 5: SUMMATION AND PRODUCT FUNCTIONS

Pattern: Σ and Π operations

Exercise: "Prove that $t(x) = \Sigma(y=0 \text{ to } x) y = 0+1+2+\dots+x$ is primitive recursive"

Step-by-Step Solution:

$$\begin{array}{ll}
 t(0) = 0 & [\text{sum from } 0 \text{ to } 0 \text{ is } 0] \\
 t(y+1) = t(y) + (y+1) & [\text{add next term}]
 \end{array}$$

This uses the sum function which we know is primitive recursive.

General Pattern for $\Sigma(i=0 \text{ to } x) f(i)$:

```
sum_f(0) = f(0)
sum_f(x+1) = sum_f(x) + f(x+1)
```

Example: Sum of k arguments

Exercise: "Prove that $\text{sum}_k(x_1, \dots, x_k) = \sum x_i$ is primitive recursive"

Recursive construction:

```
sum_2(x_1, x_2) = x_1 + x_2           [base case for k=2]
sum_{k+1}(x_1, \dots, x_{k+1}) = sum_k(x_1, \dots, x_k) + x_{k+1}
```

TYPE 6: BOUNDED SEARCH AND COUNTING

Pattern: Count elements satisfying a property

Exercise: "Prove that $\text{cpr}(x, y) = |\{p \mid x \leq p < y \wedge p \text{ prime}\}|$ is primitive recursive"

Step-by-Step Solution: Use helper function $\text{cpr}'(x, k) = |\{p \mid x \leq p < x+k \wedge p \text{ prime}\}|$:

```
cpr'(x, 0) = 0                        [no primes in empty interval]
cpr'(x, k+1) = cpr'(x, k) +  $\chi_{\text{Prime}}(x+k)$ 
```

Then: $\text{cpr}(x, y) = \text{cpr}'(x, y \div x)$

General Pattern for counting:

```
count(property, start, 0) = 0
count(property, start, k+1) = count(property, start, k) + property(start+k)
```

TYPE 7: COMPLEX RECURSIVE PATTERNS

Pattern: Nested recursions and multiple base cases

Exercise: "Prove that the largest prime $\leq x$ is primitive recursive"

Step-by-Step Construction:

```
count_primes(x) =  $\sum_{i=1}^x \chi_{\text{Prime}}(i)$    [count primes up to x]
largest_prime(x) =  $p_{\{\text{count\_primes}(x)\}}$    [use counting to index]
```

Complex helper construction:

```
lp(x) =  $\mu y \leq x. \text{sg}(\text{div}(p_{\{x-y\}}, x))$    [bounded search]
lp(x) =  $p_{\{x-\text{lp}(x)\}} \cdot \text{sg}(x-1) + \text{sg}(x-1)$    [handle edge cases]
```

PRACTICAL CONSTRUCTION STRATEGIES

Strategy 1: Build Helper Functions First

Pattern: Complex function requiring multiple steps

1. Identify needed operations (sg, $s\bar{g}$, rm, div, etc.)
2. Build these as primitive recursive functions
3. Combine using composition and recursion

Strategy 2: Use Characteristic Functions

Pattern: Functions involving set membership

1. Define characteristic function of relevant set
2. Use Boolean operations (sg, $s\bar{g}$, multiplication)
3. Combine with arithmetic operations

Strategy 3: Bounded Operations

Pattern: "For all $x \leq \text{bound}$ " or "There exists $x \leq \text{bound}$ "

1. Use bounded sum: $\sum_{i=0}^n f(i)$
2. Use bounded product: $\prod_{i=0}^n f(i)$
3. Use bounded minimization: $\mu_{y \leq x}. P(y)$

Strategy 4: Case Analysis

Pattern: Functions defined differently for different ranges

1. Use sg and $s\bar{g}$ to create "switches"
2. Multiply each case by its condition
3. Add all cases together

Template:

$$f(x) = \text{case1}(x) \cdot \text{condition1}(x) + \text{case2}(x) \cdot \text{condition2}(x) + \dots$$

VERIFICATION CHECKLIST

For every PR construction, verify:

✓ **Base case uses only:** Zero, successor, projections, or previously proven PR functions ✓ **Recursive case follows pattern:** $h(\vec{x}, y+1) = g(\vec{x}, y, h(\vec{x}, y))$ ✓ **All helper functions are PR:** Each must be built from base functions ✓ **Composition is valid:** Output types match input types ✓ **No unbounded operations:** No μ_y without bound

COMMON MISTAKE PATTERNS

Mistake 1: Using Non-PR Functions

Wrong: "Since division is obvious..." **Correct:** Build division from rm and bounded search

Mistake 2: Unbounded Operations

Wrong: $\mu y. P(y)$ (no bound) **Correct:** $\mu y \leq x. P(y)$ (bounded minimization)

Mistake 3: Wrong Base Cases

Wrong: Forgetting to handle $f(0)$ properly **Correct:** Always specify base case explicitly

Mistake 4: Non-Primitive Recursion Pattern

Wrong: $h(x, y+1) = h(x+1, y-1)$ (complex recursion) **Correct:** $h(x, y+1) = g(x, y, h(x, y))$ (simple pattern)

COMPLETE REAL EXAMPLES FROM YOUR MATERIALS

OFFICIAL EXERCISE 2.1: Power of 2 Function

Exact Exercise Text: "Prove that the function $\text{pow2} : \mathbb{N} \rightarrow \mathbb{N}$, defined by $\text{pow2}(y) = 2^y$, is primitive recursive."

Complete Step-by-Step Solution:

Step 1: Identify the pattern

- $\text{pow2}(0) = 2^0 = 1$
- $\text{pow2}(1) = 2^1 = 2$
- $\text{pow2}(2) = 2^2 = 4$
- Pattern: each step doubles the previous result

Step 2: Set up primitive recursion

$\text{pow2}(0) = 1$	[base case]
$\text{pow2}(y+1) = 2 * \text{pow2}(y)$	[recursive case]

Step 3: Make it formal using helper function We need $\text{double}(x) = 2 * x$ to be primitive recursive:

$\text{double}(0) = 0$	[base case]
$\text{double}(y+1) = \text{double}(y) + 2$	[recursive: add 2 each time]

Step 4: Complete definition

$\text{pow2}(0) = 1$	[$f() = 1$]
$\text{pow2}(y+1) = \text{double}(\text{pow2}(y))$	[$g(y, z) = \text{double}(z)$]

Step 5: Verify it's primitive recursive

- Base case uses constant 1 (composition of successor and zero)
 - Recursive case uses double (which we proved is PR) and previous result
 - Follows $h(y+1) = g(y, h(y))$ pattern ✓
-

OFFICIAL EXERCISE 2.2: Characteristic Function of $\{2^n - 1\}$

Exact Exercise Text: "Prove that the characteristic function χ_A of the set $A = \{2^n - 1 : n \in \mathbb{N}\}$ is primitive recursive."

Complete Step-by-Step Solution:

Step 1: Understand the set

- $A = \{0, 1, 3, 7, 15, 31, \dots\} = \{2^0 - 1, 2^1 - 1, 2^2 - 1, 2^3 - 1, \dots\}$

Step 2: Build helper function $a(n) = 2^n - 1$

$$\begin{array}{ll} a(0) = 0 & [2^0 - 1 = 0] \\ a(n+1) = 2 \cdot a(n) + 1 & [2^{(n+1)} - 1 = 2(2^n - 1) + 1] \end{array}$$

Step 3: Build checking function We need $\text{chk}(x, m) = 1$ if $\exists n \leq m$ such that $x = a(n)$, 0 otherwise:

$$\begin{array}{ll} \text{chk}(x, 0) = \text{sg}(x) & [\text{check if } x = a(0) = 0] \\ \text{chk}(x, m+1) = \text{chk}(x, m) + \text{eq}(x, a(m+1)) & [\text{accumulate matches}] \end{array}$$

Step 4: Final characteristic function

$$\chi_A(x) = \text{chk}(x, x) \quad [\text{check membership up to } x]$$

Step 5: Verify all helpers are PR

- $a(n)$ uses multiplication and addition (both PR)
 - $\text{eq}(x, y) = \text{sg}(|x - y|)$ uses truncated subtraction and sg
 - chk uses addition and eq
 - All compose correctly ✓
-

OFFICIAL EXERCISE 2.3: Even Numbers Characteristic Function

Exact Exercise Text: "Prove that χ_P , the characteristic function of the set of even numbers P is primitive recursive."

Complete Step-by-Step Solution:

Step 1: Understand the pattern

- $\chi_P(0) = 1$ (0 is even)
- $\chi_P(1) = 0$ (1 is odd)
- $\chi_P(2) = 1$ (2 is even)
- Pattern: alternates 1, 0, 1, 0, ...

Step 2: Use negated sign to flip First define $s\bar{g}$ (negated sign function):

$s\bar{g}(0) = 1$	[base case]
$s\bar{g}(y+1) = 0$	[recursive case]

Step 3: Build χ_P using alternation

$\chi_P(0) = 1$	[0 is even]
$\chi_P(y+1) = s\bar{g}(\chi_P(y))$	[flip previous result]

Step 4: Verify the pattern

- $\chi_P(0) = 1$ ✓
- $\chi_P(1) = s\bar{g}(\chi_P(0)) = s\bar{g}(1) = 0$ ✓
- $\chi_P(2) = s\bar{g}(\chi_P(1)) = s\bar{g}(0) = 1$ ✓
- Pattern continues correctly

Why this works: Each successor flips even/odd status, and $s\bar{g}$ flips $1 \leftrightarrow 0$.

OFFICIAL EXERCISE 2.4: Half Function (Integer Division)

Exact Exercise Text: "Prove the function $\text{half} : \mathbb{N} \rightarrow \mathbb{N}$, defined by $\text{half}(x) = \lfloor x/2 \rfloor$, is primitive recursive."

Complete Step-by-Step Solution:

Step 1: Build remainder function rm_2

$\text{rm}_2(0) = 0$	[$0 \bmod 2 = 0$]
$\text{rm}_2(x+1) = s\bar{g}(\text{rm}_2(x))$	[alternates 0,1,0,1,...]

Step 2: Build half function

$\text{half}(0) = 0$	[$\lfloor 0/2 \rfloor = 0$]
$\text{half}(x+1) = \text{half}(x) + \text{rm}_2(x)$	[add 1 every two steps]

Step 3: Verify the logic

- When x is even: $\text{rm}_2(x) = 0$, so $\text{half}(x+1) = \text{half}(x) + 0 = \text{half}(x)$
- When x is odd: $\text{rm}_2(x) = 1$, so $\text{half}(x+1) = \text{half}(x) + 1$
- This gives us: $\text{half}(0)=0$, $\text{half}(1)=0$, $\text{half}(2)=1$, $\text{half}(3)=1$, $\text{half}(4)=2$, ...
- Pattern: $\lfloor 0/2 \rfloor = 0$, $\lfloor 1/2 \rfloor = 0$, $\lfloor 2/2 \rfloor = 1$, $\lfloor 3/2 \rfloor = 1$, $\lfloor 4/2 \rfloor = 2$ ✓

Step 4: Verify primitive recursion

- Uses sg (which we built) and addition (basic PR function)
 - Follows proper recursive pattern ✓
-

OFFICIAL EXERCISE: Sum Function $t(x) = 0+1+2+\dots+x$

Exact Exercise Text: "Show that the function $t: \mathbb{N} \rightarrow \mathbb{N}$ defined by $t(x) = \sum_{y=0}^x y = 0 + 1 + 2 + \dots + x$ is in PR."

Complete Step-by-Step Solution:

Step 1: Set up primitive recursion

$$\begin{array}{ll} t(0) = 0 & [\text{sum from } 0 \text{ to } 0 \text{ is } 0] \\ t(y+1) = t(y) + (y+1) & [\text{add next term}] \end{array}$$

Step 2: Use known PR function sum Since $\text{sum}(x,y) = x + y$ is primitive recursive:

$$\begin{array}{ll} t(0) = 0 & [f() = 0] \\ t(y+1) = \text{sum}(t(y), y+1) & [g(y,z) = \text{sum}(z, s(y))] \end{array}$$

Step 3: Verify correctness

- $t(0) = 0$ ✓
 - $t(1) = t(0) + 1 = 0 + 1 = 1$ ✓
 - $t(2) = t(1) + 2 = 1 + 2 = 3$ ✓
 - $t(3) = t(2) + 3 = 3 + 3 = 6$ ✓
 - Formula: $t(n) = n(n+1)/2$ ✓
-

PRACTICAL TUTORING EXAMPLE: Counting Primes

Exercise: "Prove that $\text{cpr}(x,y) = |\{p \mid x \leq p < y \wedge p \text{ prime}\}|$ is primitive recursive"

Step-by-Step Construction:

Step 1: Use helper function $\text{cpr}'(x,k)$

$cpr'(x, 0) = 0$ [no primes in empty interval]
 $cpr'(x, k+1) = cpr'(x, k) + \chi_{\text{Prime}}(x+k)$ [add 1 if $x+k$ is prime]

Step 2: Relate to original function

$cpr(x, y) = cpr'(x, y \div x)$ [count primes in $[x, x+(y-x)) = [x, y]$]

Step 3: Verify with example

- $cpr(10, 15)$ counts primes in $[10, 15) = \{11, 13\}$
- $cpr'(10, 5)$ counts primes in $\{10, 11, 12, 13, 14\}$
- $\chi_{\text{Prime}}(10)=0, \chi_{\text{Prime}}(11)=1, \chi_{\text{Prime}}(12)=0, \chi_{\text{Prime}}(13)=1, \chi_{\text{Prime}}(14)=0$
- $\text{Sum} = 0+1+0+1+0 = 2 \checkmark$

Key insight: Use bounded counting with characteristic functions.

COMPLETE SOLUTION LIBRARY

Basic Building Blocks (Always Available):

$0(x) = 0$ [zero function]
 $s(x) = x + 1$ [successor]
 $U_j^k(x_1, \dots, x_k) = x_j$ [projections]

$sg(0) = 1, sg(x+1) = 0$ [negated sign]
 $sum(x, 0) = x, sum(x, y+1) = sum(x, y) + 1$
 $mult(x, 0) = 0, mult(x, y+1) = mult(x, y) + x$

Advanced Functions:

$eq(x, y) = sg(|x - y|)$ [equality test]
 $leq(x, y) = sg(x \div y)$ [less-or-equal test]
 $max(x, y) = x + (y \div x)$ [maximum]
 $min(x, y) = x \div (x \div y)$ [minimum]

Set Operations:

$\chi_{A \cup B}(x) = sg(\chi_A(x) + \chi_B(x))$ [union]
 $\chi_{A \cap B}(x) = \chi_A(x) \cdot \chi_B(x)$ [intersection]
 $\chi_{\bar{A}}(x) = sg(\chi_A(x))$ [complement]

This methodology covers EVERY type of primitive recursive exercise by providing:

1. **Pattern recognition** for exercise types
2. **Step-by-step construction** strategies
3. **Complete verification** methods
4. **Common mistake** avoidance
5. **Reusable building blocks** for complex constructions

You can now approach any PR exercise with confidence using these mechanical techniques!