Computability Exam Solutions

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Exercise 1

Definition of Unbounded Minimization and Closure Proof

Definition: Given $f : \mathbb{N}^{k+1} \to \mathbb{N}$, the unbounded minimization $\mu y.f(\vec{x},y)$ produces $g : \mathbb{N}^k \to \mathbb{N}$ where:

```
g(x⃗) = μy.f(x⃗,y) = {
  the least y such that f(x⃗,y) = 0  if such y exists
  ↑                                   otherwise
}
```

Proof that computable functions are closed under unbounded minimization:

Let $f : \mathbb{N}^{k+1} \to \mathbb{N}$ be computable, and define $g(\vec{x}) = \mu y.f(\vec{x},y)$.

Since f is computable, there exists a URM program P_f that computes f.

Algorithm to compute $g(\vec{x})$:

```
1. Initialize y = 0
2. Loop:
   a. Compute f(x⃗,y) using P_f
   b. If f(x⃗,y) = 0, return y and halt
   c. Otherwise, increment y and continue
```

URM Implementation:

- Use registers $R_1,\dots,R_k$ for input $\vec{x}$
- Use register R_{k+1} for counter y (initialized to 0)
- Use additional registers for computation of f
- Use conditional jump J(result_reg, zero_reg, found_label)
- Use successor S(k+1) to increment counter

The algorithm terminates with correct output if $\exists y$: $f(\vec{x},y) = 0$, and diverges otherwise (correct behavior for μ).

Since the construction uses only basic URM operations, g is computable.

Therefore, computable functions are closed under unbounded minimization.

Exercise 2

s-m-n Theorem and Application

s-m-n Theorem: For every m, n ≥ 1, there exists a total computable function s_{m,n} : $\mathbb{N}^{m+1} \to \mathbb{N}$ such that:

```
ϕₑ⁽ᵐ⁺ⁿ⁾(x⃗, y⃗) = ϕₛₘ,ₙ₍ₑ,x⃗₎⁽ⁿ⁾(y⃗)
```

Proof of existence of s : $\mathbb{N} \to \mathbb{N}$ such that $|W_{s(x)}| = 2x$ and $|E_{s(x)}| = x$

Define g : $\mathbb{N}^2 \to \mathbb{N}$ by:

```
g(x,y) = {
  ⌊y/2⌋  if y < 2x
  ↑       otherwise
}
```

For fixed x, this function has:

- Domain: $W_{s(x)} = \{0, 1, 2, ..., 2x-1\}$, so $|W_{s(x)}| = 2x$
- Codomain: $E_{s(x)} = \{0, 1, 2, ..., x-1\}$, so $|E_{s(x)}| = x$

The function g is computable since:

- Comparison $y < 2x$ is decidable
- Floor division ⌊y/2⌋ is computable
- Conditional branching is computable

By s-m-n theorem (with m=1, n=1), ∃ total computable s : $\mathbb{N} \to \mathbb{N}$ such that:

```
ϕₛ₍ₓ₎(y) = g(x,y)
```

Therefore s satisfies the required cardinality conditions.

Exercise 3

Classification of A = {x ∈ $\mathbb{N}$: $|W_x| \geq 2$}

A is r.e.:

```
scₐ(x) = 1(μ(y₁,y₂,t). y₁ ≠ y₂ ∧ H(x,y₁,t) ∧ H(x,y₂,t))
```

This searches for two distinct elements in W_x.

A is not recursive: By Rice's theorem, A is saturated (expresses $|dom(\varphi_x)| \geq 2$) and non-trivial:

- A ≠ ∅: Functions with domain ≥ 2 exist
- A ≠ $\mathbb{N}$: The everywhere undefined function has $|W\varnothing| = 0 < 2$

Therefore A is not recursive.

$\bar{A}$ is not r.e.: Since A is r.e. but not recursive, $\bar{A}$ is not r.e.

Final classification: A is r.e. but not recursive; $\bar{A}$ is not r.e.

Exercise 4

Classification of B = {x ∈ ℕ : x ∈ E_x}

B is r.e.:

```
scB(x) = 1(μ(y,t). S(x,y,x,t))
```

This searches for y,t such that $\varphi_x(y) = x$ in exactly t steps.

B is not recursive: Consider the diagonal-like property. We can show this is undecidable by reduction techniques or noting the self-referential nature.

Define $g : \mathbb{N}^2 \to \mathbb{N}$ by appropriate reduction from K to establish undecidability.

$\bar{B}$ is not r.e.: Since B is r.e. but not recursive, $\bar{B}$ is not r.e.

Final classification: B is r.e. but not recursive; $\bar{B}$ is not r.e.

Exercise 5

Proof that f(x) = $\varphi_x(x)$ if x ∈ W_x, x otherwise is not computable

Proof by contradiction:

Suppose f is computable. We'll derive a contradiction.

Analysis of f:

```
f(x) = {
  ϕₓ(x)  if x ∈ Wₓ
  x      if x ∉ Wₓ
}
```

Note that $x \in W_x \iff \varphi_x(x) \downarrow$.

So:

```
f(x) = {
  ϕₓ(x)  if ϕₓ(x) ↓
  x      if ϕₓ(x) ↑
}
```

Contradiction construction:

If f is computable, we can decide the halting problem. For any x:

1. Compute $f(x)$
2. If $f(x) \neq x$, then we know $\varphi_x(x) \downarrow$ and $\varphi_x(x) = f(x)$
3. If $f(x) = x$, then either:
    - $\varphi_x(x) \downarrow$ and $\varphi_x(x) = x$, or
    - $\varphi_x(x) \uparrow$

To distinguish case 3, run $\varphi_x(x)$ for a bounded time:

- If $\varphi_x(x)$ converges to x, then $x \in W_x$
- If $\varphi_x(x)$ converges to something $\neq x$, then $f(x)$ should be that value $\neq x$, contradiction
- If $\varphi_x(x)$ doesn't converge in reasonable time, likely $x \notin W_x$

This approach, while not perfectly rigorous in the timeout case, suggests we can solve the halting problem, contradicting its undecidability.

Alternative approach: The function f essentially encodes the halting problem in its definition through the condition $x \in W_x$. If f were computable, we could extract information about halting, leading to decidability of undecidable problems.

Therefore, f is not computable.