

Fixed points are **fundamental** to computability theory because they provide the mathematical foundation for understanding **self-reference** and **recursive definitions**.

What Are Fixed Points in Computability?

Basic Concept

A **fixed point** of a functional Φ is a function f such that $\Phi(f) = f$.

Think of it like this: if you apply the functional to the function, you get back the same function unchanged.

Concrete Example: Factorial Function

Define factorial recursively:

- $\text{fact}(0) = 1$
- $\text{fact}(n+1) = (n+1) \times \text{fact}(n)$

This can be viewed as finding a fixed point of the functional:

- $\Phi(f)(0) = 1$
- $\Phi(f)(n+1) = (n+1) \times f(n)$

The factorial function is the unique fixed point: $\Phi(\text{factorial}) = \text{factorial}$

Why Fixed Points Are So Important

1. Mathematical Foundation for Recursion

Every recursive definition is really about finding a fixed point:

Without fixed point theory: "How do we know recursive definitions actually define anything?" **With fixed point theory:** "We prove the recursive functional has a computable fixed point"

From your knowledge:

"The First Recursion Theorem is used to give 'meaning' to programs, computing a recursive program, ensuring implementing the program will be defined rigorously over its inputs in a correct way."

2. Self-Referential Programs

The **Second Recursion Theorem** (the one that IS examined) uses fixed points to prove something amazing:

Given ANY program transformation f , there exists a program e_0 such that $\phi_{e_0} = \phi f(e_0)$

This means: *No matter how you try to transform programs, there's always some program that computes the same function before and after your transformation.*

3. Diagonalization and Undecidability

Fixed points are the mathematical machinery behind many impossibility results:

- **Rice's Theorem:** Uses Second Recursion Theorem (which uses fixed points)
- **Halting Problem:** Can be proven using fixed point arguments
- **Self-referential constructions:** "This program does X to itself"

Why Fixed Points Work for Self-Reference

The Deep Insight

Problem: How can a program refer to itself? **Solution:** Use fixed points to construct programs that "see" their own code.

From your project knowledge, the Second Recursion Theorem proof shows this construction:

1. Define $g(x,y) = \phi f(\phi x(x))(y)$ [program x applied to itself, then transformed by f]
2. Use smn-theorem to get $s(x)$ such that $\phi s(x)(y) = g(x,y)$
3. Since s is computable, $s = \phi m$ for some m
4. Take $e_0 = \phi m(m) = s(m)$
5. Then $\phi e_0 = \phi f(e_0)$ [the program e_0 is unchanged by transformation f]

This is **pure diagonalization** - the program applies itself to itself!

Practical Examples Where Fixed Points Matter

1. Ackermann Function

The Ackermann function is defined recursively and exists because the recursive functional has a computable fixed point.

2. μ -operator (Minimization)

The search operation $\mu y.f(x \square, y)$ can be viewed as a fixed point:

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 $\Phi(g)(x \sqcup, y) = \{$ 
   $y,$                 if  $f(x \sqcup, y) = 0$ 
   $g(x \sqcup, y+1),$     if  $f(x \sqcup, y) \neq 0$  and  $f(x \sqcup, y) \downarrow$ 
   $\uparrow,$           otherwise
 $\}$ 

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3. Virus Programs and Quines

Programs that copy themselves or modify themselves use fixed point constructions.

The Hierarchy of Fixed Point Theorems

First Recursion Theorem (Theoretical Foundation)

- **What:** Every recursive functional has a least computable fixed point
- **Why important:** Justifies that recursive definitions actually define computable functions
- **Exam relevance:** Rarely tested directly, but foundation for everything else

Second Recursion Theorem (Practical Power)

- **What:** Self-referential program construction
- **Why important:** Proves impossibility results, enables diagonalization arguments
- **Exam relevance:** **HEAVILY TESTED** - appears frequently in exercises

Myhill-Shepherdson Theorems (Bridge)

- **What:** Connect program transformations to functional transformations
- **Why important:** Allow us to work with programs as mathematical objects

Why You See Fixed Points Everywhere

The fundamental insight: Computability theory is full of self-reference:

- Programs that examine other programs
- Sets defined in terms of themselves
- Functions that compute their own properties

Fixed point theory provides the mathematical tools to handle this self-reference rigorously.

Bottom Line

Fixed points aren't just abstract mathematics - they're the **essential tool** for:

1. **Proving recursive definitions work** (First Recursion Theorem)

2. **Constructing self-referential programs** (Second Recursion Theorem)
3. **Proving impossibility results** (Rice's Theorem, Halting Problem)
4. **Understanding the limits of computation**