

1. Definition by Cases

Core Concept

A function can be defined through cases using decidable predicates and known computable functions:

$$f(\bar{x}) = \{ \begin{array}{l} f_1(\bar{x}) \text{ if } Q_1(\bar{x}) \\ f_2(\bar{x}) \text{ if } Q_2(\bar{x}) \\ \dots \\ f_n(\bar{x}) \text{ if } Q_n(\bar{x}) \end{array} \}$$

where:

- Each Q_i is a decidable predicate
- Predicates are mutually exclusive
- Each f_i is computable

Characteristic Functions

For a set $A \subseteq \mathbb{N}$, its characteristic function is:

$$\chi_A(x) = \{ \begin{array}{l} 1 \text{ if } x \in A \\ 0 \text{ if } x \notin A \end{array} \}$$

2. Algebra of Decidability

Basic Operations

For decidable predicates Q_1, Q_2 :

1. Negation ($\neg Q$):

$$\chi_{\neg Q}(x) = \text{sg}(\chi_Q(x))$$

2. Conjunction ($Q_1 \wedge Q_2$):

$$X_{(Q_1 \wedge Q_2)}(x) = X_{Q_1}(x) \cdot X_{Q_2}(x)$$

3. Disjunction ($Q_1 \vee Q_2$):

$$X_{(Q_1 \vee Q_2)}(x) = \max\{X_{Q_1}(x), X_{Q_2}(x)\}$$

Closure Properties

If Q_1, \dots, Q_n are decidable predicates and $f: \{0,1\}^n \rightarrow \{0,1\}$, then:

- The predicate Q corresponding to $f(X_{Q_1}, \dots, X_{Q_n})$ is decidable

3. Bounded Operations

Bounded Sum

For $f: \mathbb{N}^{k+1} \rightarrow \mathbb{N}$ computable:

$$g(\bar{x}, y) = \sum_{z < y} f(\bar{x}, z)$$

Defined by primitive recursion:

$$g(\bar{x}, 0) = 0$$

$$g(\bar{x}, y+1) = g(\bar{x}, y) + f(\bar{x}, y)$$

Bounded Product

Similarly:

$$h(\bar{x}, y) = \prod_{z < y} f(\bar{x}, z)$$

Defined by:

$$h(\bar{x}, 0) = 1$$

$$h(\bar{x}, y+1) = h(\bar{x}, y) \cdot f(\bar{x}, y)$$

Bounded Quantification

For decidable predicate Q :

1. Universal:

$$Q_1(\bar{x}, y) = \forall z < y. Q(\bar{x}, z)$$
$$\neg Q_1(\bar{x}, y) = \exists z < y. \neg Q(\bar{x}, z)$$

2. Existential:

$$Q_2(\bar{x}, y) = \exists z < y. Q(\bar{x}, z)$$
$$\neg Q_2(\bar{x}, y) = \forall z < y. \neg Q(\bar{x}, z)$$

4. Bounded Minimalization

Definition

For $f: \mathbb{N}^{k+1} \rightarrow \mathbb{N}$ total computable:

$$h(\bar{x}, y) = \mu z < y. f(\bar{x}, z) = 0 = \begin{cases} \min\{z < y : f(\bar{x}, z) = 0\} & \text{if such } z \text{ exists} \\ y & \text{otherwise} \end{cases}$$

Key Properties

1. Always terminates (bounded search)
2. Total when f is total
3. Computable through bounded operations:

$$h(\bar{x}, y) = \sum_{z < y} \prod_{w \leq z} \text{sg}(f(\bar{x}, w))$$

Examples

1. Integer square root:

$$\lfloor \sqrt{x} \rfloor = \max\{y \leq x : y^2 \leq x\}$$
$$= \min\{y \leq x : (y+1)^2 > x\}$$

2. Greatest common divisor:

$$\text{gcd}(x,y) = \min\{z \leq \min(x,y) : z|x \wedge z|y\}$$

5. Applications to Arithmetic Functions

The following functions can be proven computable using bounded operations:

1. Number of divisors $D(x)$
2. Primality test $\text{Pr}(x)$
3. nth prime number p_n
4. Prime power decomposition $(x)_i$

These concepts form the foundation for understanding more complex computability results and provide essential tools for proving functions computable through bounded operations.