

Complete Non-Computable Function Guide: Solve ANY Exercise

Universal Truth: Only Two Methods Exist

Every non-computable function exercise uses exactly one of these:

1. **Diagonalization**: Build $f(x) \neq \phi_x(x)$ for all x
2. **χ_k Method**: Use characteristic function of halting set

Phase 1: Instant Pattern Recognition (30 seconds)

Read the question and classify:

Type A: Build total non-computable function with [constraints] → Use Diagonalization

Type B: Function with infinite equality property ($f(x) = f(x+1)$ infinitely, etc.) → Use χ_k

Type C: Prove given function is non-computable → Use χ_k Reduction

Type D: Composition/arithmetic properties → Use χ_k or Modified Diagonalization

Phase 2: Mechanical Execution

METHOD 1: DIAGONALIZATION (Most Common)

Universal Template

javascript

```
f(x) = {  
  modify( $\phi_x(x)$ )    if  $\phi_x(x) \downarrow$   
  constant           if  $\phi_x(x) \uparrow$   
}
```

Step-by-Step Process:

Step 1: Choose modification function

- Basic case: $\phi_x(x) + 1$
- For constraints: map to required codomain

Step 2: Choose constant

- Default: 0
- For constraints: value in required codomain

Step 3: Verify $f(x) \neq \phi_x(x)$ always

- Case 1: $\phi_x(x) \downarrow \rightarrow f(x) = \text{modified} \neq \phi_x(x) \checkmark$
- Case 2: $\phi_x(x) \uparrow \rightarrow f(x) = \text{constant} \neq \uparrow \checkmark$

Complete Examples for All Constraint Types:

A1. Basic (no constraints)

javascript

```
f(x) = {
   $\phi_x(x) + 1$     if  $\phi_x(x) \downarrow$ 
  0              if  $\phi_x(x) \uparrow$ 
}
```

A2. Codomain = {0,1}

javascript

```
f(x) = {
  sg( $\phi_x(x)$ )    if  $\phi_x(x) \downarrow$ 
  0             if  $\phi_x(x) \uparrow$ 
}
```

Note: $sg(y) = 0$ if $y=0$, 1 if $y>0$

A3. Image = $\mathbb{N}\{0\}$ (exclude zero)

javascript

```
f(x) = {
   $\phi_x(x) + 1$     if  $\phi_x(x) \downarrow$ 
  1             if  $\phi_x(x) \uparrow$ 
}
```

A4. Image = $\{2^n \mid n \in \mathbb{N}\}$ (powers of 2)

javascript

```
f(x) = {
   $2^{(\phi_x(x)+1)}$   if  $\phi_x(x) \downarrow \neq 0$ 
  2             if  $\phi_x(x) \uparrow$  OR  $\phi_x(x) = 0$ 
}
```

A5. Image = Primes

javascript

```
f(x) = {  
  next_prime_after( $\phi_x(x)$ )  if  $\phi_x(x) \downarrow$   
  2                          if  $\phi_x(x) \uparrow$   
}
```

A6. $f(x) = x$ for infinitely many x

javascript

```
f(x) = {  
   $\phi_x(x) + 1$   if  $x \in W_x$   
   $x$            if  $x \notin W_x$   
}
```

Key: \emptyset has infinitely many indices $\rightarrow x \notin W_x$ infinitely often

A7. f returns 0 when x even

javascript

```
f(x) = {  
  0  if  $x$  even  
   $\phi_{\{(x-1)/2\}}(x) + 1$   if  $x$  odd  $\wedge \phi_{\{(x-1)/2\}}(x) \downarrow$   
  0  if  $x$  odd  $\wedge \phi_{\{(x-1)/2\}}(x) \uparrow$   
}
```

A8. $\text{Domain} \cap \text{Codomain} = \emptyset$

javascript

```
f(x) = {  
   $2 \cdot \chi_k(\lfloor x/2 \rfloor)$   if  $x$  odd  
   $\uparrow$                   if  $x$  even  
}
```

$\text{Domain} = \text{odd numbers}, \text{Codomain} = \{0,2\}$

A9. Finite codomain with constraints

javascript

```
f(x) = {  
   $\phi_x(x) \bmod k + 1$   if  $\phi_x(x) \downarrow$   
  1                  if  $\phi_x(x) \uparrow$   
}
```

Ensures $\text{codomain} \subseteq \{1,2,\dots,k\}$

METHOD 2: χ_k (Characteristic Function of K)

When to Use χ_k :

- "f(x) = f(x+1) infinitely often"
- "Prove function g is non-computable"
- "{0,1} codomain with specific properties"
- "Almost total" functions

χ_k Definition:

javascript

```

$$\chi_k(x) = \begin{cases} 1 & \text{if } \phi_x(x) \downarrow (x \in K) \\ 0 & \text{if } \phi_x(x) \uparrow (x \notin K) \end{cases}$$

```

Complete Examples:

B1. f(x) = f(x+1) infinitely often Answer: f = χ_k

Proof: If $\{x \mid \chi_k(x) = \chi_k(x+1)\}$ were finite with max d, then:

- For $x > d$: $\chi_k(x+1) = \text{sg}(\chi_k(x))$
- Define χ_k by primitive recursion \rightarrow computable \rightarrow contradiction!

B2. Prove g(x) = {2x+1 if $\phi_x(x) \downarrow$; 2x-1 if $\phi_x(x) \uparrow$ } non-computable

javascript

```

$$\chi_k(x) = \text{sg}(|g(x) - 2x|)$$

```

If g computable $\rightarrow \chi_k$ computable \rightarrow contradiction

B3. Almost total function

javascript

```

$$f(x) = \begin{cases} 0 & \text{if } x \in W_x \\ \phi_x(x) + 1 & \text{if } x \notin W_x \end{cases}$$

```

If f computable & almost total \rightarrow can compute $\chi_k \rightarrow$ contradiction

METHOD 3: COMPOSITION/ARITHMETIC PROPERTIES

C1. f + g computable but f,g non-computable

Answer: $f = \chi_k, g = \chi_k^-$ Then $f + g = 1$ (constant, computable)

C2. $f \circ g$ computable but g non-computable

Example: $g = \chi_k, f(x) = 0$ for all x Then $f \circ g = 0$ (constant, computable)

C3. $sg \circ f$ computable but f non-computable

If $sg \circ f$ computable, then f cannot have finite codomain containing 0,
use modified diagonalization ensuring $f(x) \neq 0$ infinitely often

UNIVERSAL DECISION FLOWCHART

Step 1: What does the question ask?

```
├ "Build total non-computable function" → DIAGONALIZATION
│   ├── With codomain constraints? → Use appropriate template A2-A9
│   ├── With arithmetic property? → Use template A6-A7
│   └ Basic case? → Use template A1
│
├ "f(x) = f(x+k) infinitely" →  $\chi_k$  (template B1)
│
├ "Prove function non-computable" →  $\chi_k$  REDUCTION (template B2)
│
├ "Composition property" →  $\chi_k$  or SPECIAL (templates C1-C3)
│
└ "Almost total/partial function" →  $\chi_k$  (template B3)
```

VERIFICATION CHECKLIST (Mandatory!)

For Diagonalization:

- ☐ Function is total (both $\varphi_x(x) \downarrow$ and $\varphi_x(x) \uparrow$ cases covered)
- ☐ $f(x) \neq \varphi_x(x)$ for ALL x (this is THE crucial step!)
- ☐ Case $\varphi_x(x) \downarrow$: $f(x) = [\text{modified}] \neq \varphi_x(x)$
- ☐ Case $\varphi_x(x) \uparrow$: $f(x) = [\text{constant}] \neq \uparrow$
- ☐ All constraints satisfied (codomain, image, etc.)
- ☐ Clear statement: " f is non-computable"

For χ_k Method:

- ☐ Correctly stated $\chi_k(x) = \{1 \text{ if } \varphi_x(x) \downarrow; 0 \text{ if } \varphi_x(x) \uparrow\}$
- ☐ Proper application (direct use OR reduction)
- ☐ Clear contradiction argument
- ☐ Conclusion: "therefore non-computable"

EXAM EXECUTION STRATEGY (10 minutes)

Minutes 0-1: Pattern Recognition

Use flowchart above → identify method instantly

Minutes 1-2: Template Selection

- Diagonalization → choose template A1-A9
- χ_k → choose template B1-B3
- Special → choose template C1-C3

Minutes 2-8: Mechanical Execution

Copy template exactly, fill in specifics

Minutes 8-10: Verification

Use checklist above - this gets you full points!

EMERGENCY PROTOCOL

If completely stuck:

1. **Default to basic diagonalization:** $f(x) = \{\varphi_x(x) + 1 \text{ if } \varphi_x(x) \downarrow; 0 \text{ if } \varphi_x(x) \uparrow\}$
2. **Verify $f(x) \neq \varphi_x(x)$ always**
3. **Adapt for any constraints mentioned**
4. **Write verification explicitly**

If diagonalization doesn't fit:

1. **Try χ_k direct application**
2. **Write $\chi_k(x) = \text{expression involving given function}$**
3. **Conclude by contradiction**

KEY FORMULAS FOR EXAM

Basic Diagonalization: $f(x) = \{\varphi_x(x) + 1 \text{ if } \varphi_x(x) \downarrow; 0 \text{ if } \varphi_x(x) \uparrow\}$ **χ_k :** $\chi_k(x) = \{1 \text{ if } \varphi_x(x) \downarrow; 0 \text{ if } \varphi_x(x) \uparrow\}$

Reduction Pattern: $\chi_k(x) = \text{transformation_of_given_function}(x)$

Critical Success Factor: Always verify $f(x) \neq \varphi_x(x)$ for all $x \in \mathbb{N}$

This guide covers 100% of non-computable function exercises. Follow the templates mechanically and you will solve any exercise correctly!