

1. Main Techniques

1.1 Diagonalization

The most fundamental technique for proving functions non-computable. Key steps:

1. Assume the function is computable
2. Use it to construct a function that differs from every computable function
3. Reach a contradiction

Example:

```
To prove  $f(x) = \begin{cases} \phi_x(x) + 1 & \text{if } \phi_x(x) \downarrow \\ 0 & \text{otherwise} \end{cases}$  is not computable:
```

1. Assume f is computable
2. Then $\exists e$ such that $f = \phi_e$
3. Consider $f(e)$:
 - If $\phi_e(e) \downarrow$ then $f(e) = \phi_e(e) + 1 \neq \phi_e(e)$
 - If $\phi_e(e) \uparrow$ then $f(e) = 0 \neq \phi_e(e)$
4. Contradiction: $f(e) \neq \phi_e(e) = f(e)$

1.2 Reduction from Known Non-Computable Functions

Common non-computable functions to reduce from:

- Halting function $h(x,y) = "\phi_x(y) \downarrow"$
- Diagonal function $K(x) = "x \in W_x"$
- Total function $t(x) = "\phi_x \text{ is total}"$

Example:

```
To prove  $g(x) = "\phi_x \text{ is total}"$  is not computable:
```

1. Assume g is computable
2. Define $h(x) = \begin{cases} \phi_x(x) + 1 & \text{if } g(x) = 1 \\ 0 & \text{otherwise} \end{cases}$
3. Show h would be computable (contradiction)

2. Special Cases

2.1 Total Non-Computable Functions

To prove a total function is not computable:

1. Show it's total
2. Prove non-computability by diagonalization or reduction

Example:

```
f(n) = {  
     $\chi_K(n)$    if  $n \in K$   
    0           otherwise  
}
```

1. Total because output always defined
2. Not computable as it would solve halting problem

2.2 Functions with Special Properties

Some functions might be non-computable due to special properties:

- Growth rate exceeding all computable functions
- Having undecidable properties in their domain/range
- Encoding solutions to undecidable problems

3. Common Proof Patterns

3.1 Contradiction via Halting Problem

1. Assume f is computable
2. Use f to decide membership in K
3. Reach contradiction since K is not decidable

3.2 Construction of Uncomputable Function

1. Start with a computable function g
2. Modify g to create f that differs from every ϕ_e
3. Show f cannot be computable

3.3 Rice's Theorem Application

1. Show function defines a non-trivial property of computable functions
2. Apply Rice's Theorem to prove non-computability

4. Common Mistakes to Avoid

1. Not proving totality when claiming a function is total
2. Confusing partial and total functions
3. Incorrect application of diagonalization
4. Assuming composition preserves non-computability

5. Exercise Strategy

1. Identify which technique might work:
 - Does it look like diagonalization would work?
 - Can you reduce from a known non-computable function?
 - Does it involve properties of computable functions?
2. Set up the proof structure:
 - Assume computability
 - Plan contradiction
 - Build helper functions if needed
3. Execute the proof:
 - Be precise about function definitions
 - Track where computability is used
 - Verify contradiction is reached
4. Verify:
 - Check all cases are covered
 - Verify all functions used are computable
 - Confirm contradiction is valid