IMMEDIATELY RECURSIVE - "Safe Patterns"

Pattern 1: Pure Arithmetic/Finite Conditions

```
✓ A = \{x \in \mathbb{N} \mid x \text{ is even}\}
✓ A = \{x \in \mathbb{N} \mid x < 100\}
✓ A = \{x \in \mathbb{N} \mid x \equiv 3 \pmod{7}\}
✓ A = \{x \in \mathbb{N} \mid x = 2^k \text{ for some } k \le 10\}
✓ A = \{0, 1, 4, 9, 16, 25, ...\} (perfect squares)
```

Why: Pure number theory, no halting involved.

Pattern 2: Finite Sets

```
✓ A = {2, 5, 17, 42}

✓ A = {x | x ≤ 1000 and x is prime}
```

Why: Finite = always recursive.

Pattern 3: Co-finite Sets

```
✓ A = \mathbb{N} \setminus \{17\} (all numbers except 17)
✓ A = \{x \mid x \neq 0 \text{ and } x \neq 1\}
```

Why: Complement is finite.

IMMEDIATELY NON-RECURSIVE - "Danger Patterns"

Pattern 1: The Self-Reference Triangle

```
\times A = {x | x \in W_x} (halting on own index)

\times A = {x | \phi_x(x) \downarrow} (K - the halting set)

\times A = {x | \phi_x(x) = x} (outputs own index)

\times A = {x | \phi_x(x) > x} (outputs larger than index)

\times A = {x | x \in E_x} (index in own range)
```

Pattern 2: Domain/Range Equality Patterns

```
\times A = {x | W_x = E_x} (domain equals range)

\times A = {x | W_x = N} (total functions)

\times A = {x | E_x = {0}} (specific range)

\times A = {x | |W_x| = 5} (specific domain size)

\times A = {x | W_x \subseteq E_x} (domain subset of range)
```

Why: These are **saturated sets** \rightarrow Rice's Theorem applies.

Pattern 3: The "Intersection/Union" Danger Zone

```
\times A = {x | W_x n E_x \neq \emptyset} (domain meets range)

\times A = {x | W_x U E_x = N} (domain union range = all)

\times A = {x | x \in W_x n E_x} (self-reference + intersection)
```

Why: Combines halting with complex relationships.

Pattern 4: Function Properties

```
\times A = {x | \phi_{-}x is total} (totality is undecidable)

\times A = {x | \phi_{-}x is injective} (injectivity is undecidable)

\times A = {x | \phi_{-}x is increasing} (function properties)

\times A = {x | \phi_{-}x(y) = y² for infinitely many y}
```

Why: All non-trivial function properties are undecidable.

USUALLY NON-RECURSIVE - "Suspicious Patterns"

Pattern 1: Quantified Halting

```
A = \{x \mid \exists y > x. \ y \in W_x\} (probably not recursive)

A = \{x \mid \forall y \in W_x. \ \phi_x(y) > 0\} (probably not recursive)

A = \{x \mid \exists k. \ k \cdot x \in W_x\} (probably not recursive)
```

Strategy: Try reduction from K.

Pattern 2: Bounded but Complex

```
O A = \{x \mid \exists y \le x. \ \phi_x(y) = x^2\} (might be r.e. but not recursive)
O A = \{x \mid |W_x \cap [0,x]| \ge 2\} (finite intersection, but...)
```

Strategy: Check if r.e., then try $K \leq_m A$.

DECISION FLOWCHART

6 INSTANT RECOGNITION EXAMPLES

Immediately Recursive:

```
A = \{x \mid x < 2025\} \leftarrow Finite

A = \{x \mid x \text{ is odd}\} \leftarrow Simple arithmetic
```

Immediately Non-Recursive:

A =
$$\{x \mid x \in W_x\}$$

A = $\{x \mid W_x = \mathbb{N}\}$
A = $\{x \mid \phi_x \text{ is total}\}$
A = $\{x \mid E_x = \{0,1,2\}\}$

- ← Self-reference danger!
- ← Rice's theorem!
- ← Function property!
- ← Specific range!

Need Investigation:

```
A = \{x \mid \exists y \in W_x. \ y > x\}
A = \{x \mid |W_x| \le 5\}
A = \{x \mid x^2 \in E_x\}
```

- ← Try K ≤_m A
- ← Check if saturated
- ← Try reduction approach

🙀 KEY WARNING SIGNALS

- 1. " $x \in W_x$ " or " $\phi_x(x)$ " \rightarrow Immediate red flag
- 2. "W_x = [something]" → Usually Rice's Theorem
- 3. "E_x = [something]" → Usually Rice's Theorem
- 4. "φ_x is [property]" → Usually Rice's Theorem
- 5. " $\forall y \in W_x$ " or " $\exists y \in W_x$ " \rightarrow Try reduction
- 6. No W_x, E_x, ϕ _x \rightarrow Probably recursive

PRACTICAL RULE OF THUMB

"If it talks about what programs DO (halting, outputs, domains), it's probably not recursive.

If it talks about what numbers ARE (even, prime, bounded), it's probably recursive."

This pattern recognition works for ~90% of exam problems. For the remaining 10%, apply formal reduction techniques or Rice's theorem systematically.