

1. Universal Function

1.1 Definition

The universal function is defined as:

$$\begin{aligned}\psi_U: \mathbb{N}^2 &\rightarrow \mathbb{N} \\ \psi_U(x, y) &= \phi_x(y)\end{aligned}$$

General k-ary version:

$$\begin{aligned}\psi_U^{(k)}: \mathbb{N}^{(k+1)} &\rightarrow \mathbb{N} \\ \psi_U^{(k)}(e, \bar{x}) &= \phi_e^{(k)}(\bar{x})\end{aligned}$$

1.2 Key Properties

1. Computability predicates:

$$\begin{aligned}H(e, x, t): & \text{"}\phi_e(x)\downarrow \text{ in } t \text{ or fewer steps"} \\ S(e, x, y, t): & \text{"}\phi_e(x)\downarrow y \text{ in } t \text{ or fewer steps"}\end{aligned}$$

2. Kleene Normal Form:

$$\phi_e^{(k)}(x) = (\mu z. |\chi_{S^k}(e, \bar{x}, (z)_1, (z)_2) - 1|)_1$$

2. Second Recursion Theorem (SRT)

2.1 Statement

For any total computable function f , there exists e_0 such that:

$$\phi_{e_0} = \phi_f(e_0)$$

2.2 Application Template

1. Define suitable computable function f
2. Apply SRT to get e_0
3. Use the fact that $\phi_{e_0} = \phi_f(e_0)$ to reach conclusion

2.3 Common Applications

1. Proving K is not recursive:

```
Define  $f(x) = \begin{cases} e_0 & \text{if } x \in K \\ e_1 & \text{if } x \notin K \end{cases}$ 
```

2. Proving a set is not saturated:

```
Use SRT to construct  $e$  where:  
 $\phi_e(y) = \begin{cases} 0 & \text{if } y = e \\ \uparrow & \text{otherwise} \end{cases}$ 
```

3. Exercise Solving Strategies

3.1 For Universal Function Problems

1. Configuration tracking:

```
Consider:  
- Initial configuration  
- Step-by-step execution  
- Final configuration
```

2. Using predicates H and S:

```
To check if computation halts:  
use  $H(e, x, t)$ 
```

```
To check specific output:  
use  $S(e, x, y, t)$ 
```

3.2 For SRT Problems

1. Function construction:

```
- Define  $f$  based on desired property  
- Ensure  $f$  is total and computable
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- Apply SRT to get fixed point

2. Contradiction proofs:

- Assume property holds
- Define f to break property
- Use fixed point to reach contradiction

4. Common Problem Types

4.1 Fixed Point Problems

Goal: Find e where ϕ_e has special property P

1. Define f to enforce P
2. Apply SRT
3. Show fixed point has P

4.2 Impossibility Proofs

Goal: Show no program can have property P

1. Assume program exists
2. Define f to create contradiction
3. Use fixed point to show impossibility

4.3 Non-Computability Proofs

Goal: Show function g is not computable

1. Assume g computable
2. Define f using g
3. Use fixed point to reach contradiction

5. Exercise Examples

5.1 Self-Referential Programs

Goal: Find program that prints its own index

1. Define $f(x) = \text{"program that outputs } x\text{"}$

2. Apply SRT to get e_0
3. Show e_0 outputs e_0

5.2 Impossibility Results

Goal: Show no program can determine if another program outputs its own index

1. Assume such program exists
2. Define f to contradict it
3. Use fixed point to reach contradiction