

22-01

20-02-2023

→ PROVIDES SATURATED SET DEFINITION

$$A \subseteq \mathbb{N}$$

RECURSION  
THEOREMS  
(KLEEN)

$\Rightarrow$  RECURSIVENESS  
(R.E / R.E.C)

$$\exists m, n \in \mathbb{N} \Rightarrow \varphi_m = \varphi_n$$

NOT SATURATED  $\rightarrow$  2 R.E.C. THEOREM

② STATE THE 2. R.T

TOTAL  
COMPUTABLE  $f: \mathbb{N} \rightarrow \mathbb{N}$

$$\exists e \in \mathbb{N} \mid \varphi_e = \varphi_{f(e)}$$

③ USE 2. R.T TO PROVE

K IS NOT SATURATED

SMN-THEOREM

$$\varphi_{s(x)}(y) = g(x, y) = \dots$$



$$g(x, y) = \begin{cases} 1 & \text{if } x \in W_x \\ \uparrow & \text{otherwise} \end{cases} = \mathcal{C}_K(x) = \mu \mathbb{Z} \cdot |y - x|$$

SMN-THEOREM  $\quad \varphi_{s(x)}(y) = g(x, y)$

2.2.5  $\rightarrow \exists e \in \mathbb{N} \mid \varphi_R(e) = \varphi_e$

$$h(e, y) = \begin{cases} 1 & \text{if } e \in W_e \\ \uparrow & \text{otherwise} \end{cases} = \mu \mathbb{Z} \cdot |y - e|$$

NOT SATURATED  $\Rightarrow \varphi_m \neq \varphi_n$

$$e' \in \mathbb{N}, e' \neq e, \varphi_{e'} \neq \varphi_e \quad \begin{matrix} e' \notin K \\ e' \in K \end{matrix}$$

$[K \text{ is NOT SATURATED}]$

$\leftarrow$  TEMPLATE 2.2.5

SMN-THEOREM  $\rightarrow$  STATES IT

$\exists s: \mathbb{N} \rightarrow \mathbb{N}$  (TOTAL / COMPUTABLE)

$$\forall x \in \mathbb{N}, x > 0 \mid W_{s(x)} = \bigcup_A \{ \varphi_e(x) \mid \varphi_e(x) = 2^x \}$$

$$P_M \neq \mathbb{P}$$

SET OF  
GIVEN NUMBERS

- STATE IT

$$[S_m^n] \quad \exists m, n > 1$$

$$S_m^n : \mathbb{N}^{m+1} \rightarrow \mathbb{N} \text{ s.t. } \forall e \in \mathbb{N}, \vec{x} \in \mathbb{N}, \vec{y} \in \mathbb{N}^n$$

$$\varphi_e^{(m+n)}(\vec{x}, \vec{y}) = \varphi_{S_{m,n}(e, \vec{x})}^{(n)}(\vec{y})$$

$$[W_{SCX} = \mathbb{P}, \text{ if } \langle x \rangle = 2x] \quad g : \mathbb{N}^2 \rightarrow \mathbb{N}$$

$$g(x, y) = \begin{cases} q_A(2, y) \bmod 2x & y \in \mathbb{P}, x > 0 \\ \uparrow & \text{otherwise} \end{cases}$$

$$= \text{rem}(2x, q_A(2, y))$$

$$+ \mu \mathbb{Z}, (\text{rem}(2, y) + \overline{\Delta} g(x))$$

$$\overline{\Delta} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$|W_{SCX}| = \{x \in \mathbb{N} \mid g(x, y) \downarrow\}$$

$$= \{x \in \mathbb{P} \mid y \in \mathbb{P}\} = \{\mathbb{P}\}$$

$$|\mathcal{C}_{SCX}| = \{g(x, y) \downarrow \mid y < 2x\} = 2x$$


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SET = PROPERTY  $\{A \subseteq \mathbb{N}\}$

[RECURSIVE]  $\Rightarrow \mathcal{B}_A$

[NOT RECURSIVE]  $\begin{cases} K \in_m A \\ \text{RICE-THEOREM} \end{cases}$

[NOT R.E.]  $\begin{cases} \bar{K} \in_m A \\ \text{RICE-SHAPIRO} \end{cases}$

[NOT COMPUTABLE  $\subseteq$  NOT RECURSIVE]

$\chi_k$

DIAGONALIZATION  
 $\varphi_e \neq \varphi_{e+1}$

[SATURATED  $\Rightarrow$  RECURSIVENESS]  
[NOT SATURATED  $\Rightarrow$  2. R.E.]

$A = \{x \in \mathbb{N} \mid W_x \neq \emptyset \wedge W_x \subseteq \mathbb{N}_x\}$

R.E. / R.E. (?)

SATURATEDNESS  $\rightarrow \exists$  INPUTS  
FOR WHICH  
THE SET  
CONDITIONS  
|

HOLD IN  
A FINITE  
N. OF STEPS

$$A = \{x \in \mathbb{N} \mid W_x \neq \emptyset \wedge W_x \subseteq B_x\}$$

A IS SATURATED  $\Rightarrow A =$  SET OF COMPUTABLE FUNCTIONS

$$A = \{x \in \mathbb{N} \mid \varphi_x \in A\}$$

$$b.f.x / \phi x$$

$$\varphi_x(y) \Rightarrow f_x(y)$$

$$A = \{f \mid \text{dom}(f) \neq \emptyset \wedge \text{dom}(f) \subseteq \text{cod}(f)\}$$

$$\varphi_x(x) \Rightarrow f_x(x)$$

$$A = \{x \in \mathbb{N} \mid W_x \neq \emptyset \wedge W_x \subseteq B_x\}$$

A IS NOT R.E.  $\rightarrow$  RICE-SHAPIRO (?)

$$\rightarrow \left[ \begin{array}{l} \textcircled{1} \exists f \in A, \forall \theta \notin A \\ \textcircled{2} \exists f \notin A, \exists \theta \subseteq f \in A \end{array} \right]$$

$$ID = \mathbb{N} \quad ID \in A \mid W_x = \mathbb{N} \neq \emptyset$$

$$\wedge ID \in \mathbb{N}$$

$$\left[ \begin{array}{l} \text{NO FINITE SUBFUNCTION} \\ \text{DOES NOT RESPECT } A \end{array} \right]$$

$\Downarrow$   
 ALTERNATE  $\rightarrow$  USE  $\text{PROD}(x) = \begin{cases} x-1 \\ 1 \end{cases}$

$$\bar{A} = \{x \in \mathbb{N} \mid W_x = \emptyset \vee W_x \neq B_x\}$$

$\downarrow$   
 RSC (R.E. (NOT)?

$$\left[ \text{ID} \notin \bar{A} \mid W_x = \mathbb{N} \neq \emptyset \vee \right. \\ \left. \mathbb{N} \notin \mathbb{N} \right]$$

$\exists f \notin A, \exists \theta \subseteq f \in A \Leftarrow (2)$  RICE-  
 SHAPIRO  
 ...

$$\left[ \theta = \emptyset \in \bar{A} \mid W_x = \emptyset \vee \right. \\ \left. \emptyset \neq \emptyset \right]$$

$$\underbrace{x \neq x-1}_{\text{PROD}} \quad \left[ \begin{array}{l} x = \{0, 1, 2\} \\ x = \{0, 1\} \end{array} \right] \quad \text{RANGE - IDSA}$$

$$\text{dom}(\text{pred}) \subseteq \text{cod}(\text{pred})$$

$A/\bar{A}$  NOT R.E.  $\rightarrow$  NOT RECURSIVE



PIR  $\rightarrow$  STATE WHAT PIR MEANS

SHOW

①  $\text{isqrt} : \mathbb{N} \rightarrow \mathbb{N}$  s.t.  $\text{isqrt}(x) \geq \lfloor \sqrt{x} \rfloor$

②  $\text{lp} : \mathbb{N} \rightarrow \mathbb{N}$  (LARGEST PRIME DIVISOR OF  $x$ )

ASSUMES PIR FUNCTIONS  
SEEN IN THE COURSE

PIR  $\left( \begin{array}{l} \text{"NORMAL" DEFINITION (INDUCTION)} \\ \text{MINIMALIZATION (\mu - OPERATOR)} \end{array} \right.$

$\text{isqrt} : \mathbb{N} \rightarrow \mathbb{N}$  s.t.  $\text{isqrt}(x) \geq \lfloor \sqrt{x} \rfloor$

$$\sqrt{x} = y^2$$

LARGEST  $y$  s.t.  $y^2 \leq x$

SMALLEST  $y$  s.t.  $y^2 > x$

$$y^2 \leq x \Rightarrow y+1$$

$$\text{isqrt}(x) = \mu y < x+1. ((y+1)^2 > x)$$

$$= \mu y < x+1, \exists y ((y+1)^2 = x))$$

$$\textcircled{2} \quad l_p: \mathbb{N} \rightarrow \mathbb{N} \quad \left( \text{LARGEST PRIME DIVISOR OF } x \right)$$

BOUNDED SUM

$\in \mathbb{P} \mathbb{R}$

BOUNDED PRODUCT

$p_i \in \mathbb{P}$

COUNT:  $\mathbb{N} \rightarrow \mathbb{N}$  s.t.

$$\text{count}(x) = \sum_{i=1}^x \text{div}(p_i, x)$$

$$l_p(x) = \text{count}(x) + \overline{\exists} (x=1)$$

[EXAMPLES FROM LESSON  $\rightarrow$  lcd / gcd]

$$A = \{x \in \mathbb{N} \mid x+1 \in Bx\} \quad A/\bar{A} \text{ rsc. / r.b.}$$

$$K \leq_m A \quad [x \in K \mid f(x) \in A]$$

$$g(x, y) = \begin{cases} y & x \in K \\ 1 & \text{otherwise} \end{cases}$$

$\downarrow$   
SMN-THEOREM

$$= y \circ \underbrace{S_K(x)}$$

$$S_K = \begin{cases} 1 & x \in K \\ 1 & \text{otherwise} \end{cases}$$



$$\Downarrow$$

$$\varphi_{S(x)}(y) = g(x, y)$$

$$- x \in K, \varphi_{S(x)}(y) = g(x, y) = y \quad \forall y \in \mathbb{N}$$

$$y \equiv S(x) \mid S(x) + 1 \in B_{S(x)} = \mathbb{N}$$

$$\text{so } S(x) \in A$$

$$- x \notin K, \varphi_{S(x)}(y) = g(x, y) = \uparrow, \forall y \in \mathbb{N}$$

$$W_{S(x)} = B_{S(x)} = \emptyset \neq \mathbb{N}$$

A NOT RECURSIVE

( $\hookrightarrow$  ALTERNATIVELY  
(RICE'S THEOREM))

$$L_1, L_2 = \mathbb{N} / \emptyset$$

$$\Rightarrow A \neq \emptyset, A \neq \mathbb{N}$$

SATURATED

$\equiv$   
A NOT RECURSIVE

$$A = \{x \in \mathbb{N} \mid x+1 \in \underbrace{B_x}\}$$

$\downarrow$   
NOT REC. (R.O(?))

$$\downarrow \quad S(x, y, z, A)$$

$$W_x \equiv H(x, y, A)$$

$$K = H(x, x, y)$$

$$\# (\mu(x, y, A), S(x, y, x+1, A)) = \varphi_x$$

$$= \# (\mu(x, y, A), \chi_S(x, y, x+1, A) - 1)$$

$$\downarrow$$

$$(w)_1 = g, (w)_2 = A$$

$$= \mu w. (\exists x, (w)_1, x+1, (w)_2 - 1)$$

$$\Downarrow$$

$$1 \mu w. \exists (x, (w)_1, x+1, (w)_2)$$

$A \rightarrow \text{NOT RSC. / NOT R.E.}$

$$\bar{A} = \{x \in \mathbb{N} \mid x+1 \notin B_x\} \rightarrow \text{NOT R.E.}$$

$\bar{A}$  (NOT RECURSIVE  $\Rightarrow$   
OTHERWISE BOTH  
WOULD BE)