

1. Many-One Reducibility (\leq_m)

1.1 Definition

For sets $A, B \subseteq \mathbb{N}$, $A \leq_m B$ if there exists a total computable function f such that:

$$\forall x: x \in A \Leftrightarrow f(x) \in B$$

1.2 Properties

1. Transitivity: If $A \leq_m B$ and $B \leq_m C$ then $A \leq_m C$
2. Properties preserved:

- If B recursive $\Rightarrow A$ recursive
- If B r.e. $\Rightarrow A$ r.e.
- If A not recursive $\Rightarrow B$ not recursive
- If A not r.e. $\Rightarrow B$ not r.e.

1.3 Template for Reductions

To prove $A \leq_m B$:

1. Define $g(x,y)$ function:
$$g(x,y) = \begin{cases} \text{[something]} & \text{if } x \in A \\ \text{[something else]} & \text{otherwise} \end{cases}$$
2. Apply s-m-n theorem to get s where:
$$\phi_{s(x)}(y) = g(x,y)$$
3. Prove s is reduction function:
$$x \in A \Leftrightarrow s(x) \in B$$

2. Common Reduction Patterns

2.1 From K (Halting Set)

To reduce $K \leq_m A$:

1. Define $g(x,y) = \begin{cases} 1 & \text{if } x \in K \\ \uparrow & \text{otherwise} \end{cases}$
2. Get s where $\phi_s(x)(y) = g(x,y)$
3. Prove $x \in K \Leftrightarrow s(x) \in A$

2.2 From \bar{K} (Complement of Halting Set)

To reduce $\bar{K} \leq_m A$:

1. Define $g(x,y) = \begin{cases} y & \text{if } x \notin K \\ \uparrow & \text{if } x \in K \end{cases}$
2. Get s where $\phi_s(x)(y) = g(x,y)$
3. Prove $x \in \bar{K} \Leftrightarrow s(x) \in A$

3. Problem Types and Solutions

3.1 Non-Recursiveness Proofs

To prove A not recursive:

1. Show $K \leq_m A$ or $\bar{K} \leq_m A$
2. Since K/\bar{K} not recursive, A not recursive

3.2 Non-R.E. Proofs

To prove A not r.e.:

1. Show $\bar{K} \leq_m A$
2. Since \bar{K} not r.e., A not r.e.

3.3 Comparing Sets

To compare sets A and B:

1. Try to construct reduction $A \leq_m B$
2. If impossible, prove $B \leq_m A$
3. If neither possible, prove incomparability

4. Exercise Examples

4.1 Input Problem

Prove $A_n = \{x \mid \phi_x(n) \downarrow\}$ not recursive:

1. Define $g(x,y) = \begin{cases} 1 & \text{if } x \in K \\ \uparrow & \text{otherwise} \end{cases}$
2. Get s where $\phi_{s(x)}(y) = g(x,y)$
3. Show $x \in K \Leftrightarrow s(x) \in A_n$

4.2 Output Problem

Prove $B_n = \{x \mid n \in E_x\}$ not recursive:

1. Define $g(x,y) = \begin{cases} n & \text{if } x \in K \\ \uparrow & \text{otherwise} \end{cases}$
2. Get s where $\phi_{s(x)}(y) = g(x,y)$
3. Show $x \in K \Leftrightarrow s(x) \in B_n$

5. Common Pitfalls

1. Function not total:

Reduction function must be total computable

2. Wrong direction:

Must prove both directions:

$x \in A \Rightarrow f(x) \in B$

$f(x) \in B \Rightarrow x \in A$

3. Not computable:

Ensure reduction function is computable

6. Strategy for Finding Reductions

1. Look for patterns:

- How to encode membership in A into B?
- What structure in B can represent A?

2. Common transformations:

- Use input as parameter
- Encode multiple values
- Use s-m-n theorem

3. Check requirements:

- Total function
- Computable
- Preserves membership