# Computer Security: Principles and Practice

**Chapter 3 – User Authentication** 

#### **User Authentication**

- fundamental security building block
  - basis of access control & user accountability
- is the process of verifying an identity claimed by or for a system entity
- > has two steps:
  - identification specify identifier
  - verification bind entity (person) and identifier
- distinct from message authentication

#### **User Authentication**

- User authentication example:
  - User real name: Alice Toklas
  - User ID: ABTOKLAS
  - Password: A.df1618hJb
- These informations are stored in a system
  - Only Alice can access with this credential
  - But attackers can still do something ...

#### **Means of User Authentication**

- > four means of authenticating user's identity
- based on something the individual
  - knows e.g. password, PIN
  - possesses e.g. key, token, smartcard
  - is (static biometrics) e.g. fingerprint, retina
  - does (dynamic biometrics) e.g. voice, sign
- > can use alone or combined
- ➤ all can provide user authentication

#### **Password Authentication**

- > widely used user authentication method
  - user provides name/login and password
  - system compares password with that saved for specified login
- > authenticates ID of user logging and
  - that the user is authorized to access system
  - determines the user's privileges
  - is used in discretionary access control

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  - Attack: the attacker has the hash of the target password and he tries to break it
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    - Info related to the target
  - o Countermeasure:

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    - protect these informations

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  - o Countermeasure:
    - account lockout mechanisms (i.e., allow only few authentication attempts)

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  - Attack: the attacker tries popular password against a wide range of accounts
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  - o Countermeasure:
    - Policies that do not allow the use of simple and common passwords

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  - Attack: The attacker waits until a logged-in workstation is unattended
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  - Countermeasure:
    - Automatically logging-out mechanisms
    - Anomaly behaviour detection

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  - Attack: Users tend to write down passwords
    - E.g., post-it near the protected device
    - Devices with pre-configured passwords
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    - User training
    - Combined authentication mechanism
      - Password + token

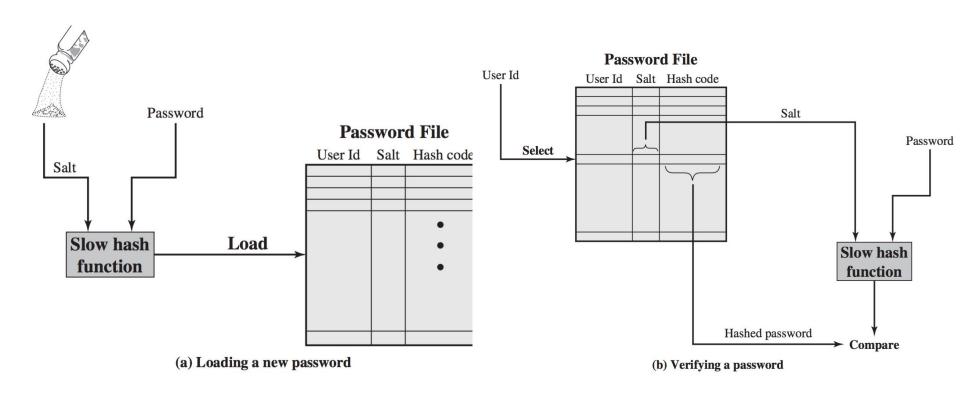
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  - o Countermeasure:
    - User training
    - Forbid the password-reuse in multiple systems
      - Feasible only on a specific network that we can control

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  - Attack: if the password is communicated through a network, an attacker can sniff these packets and steal the password
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  - Countermeasure:
    - Secure communication links

- Widely used security mechanism
- > Steps:
  - The user create a new password
  - The password is combined with a fixed length salt
    - The salt usually is pseudo-randomly generated
  - Hashed Password = Hash(password, salt)
  - ID, Hashed password and Salt are saved in a file
    - Password file
- These hashed functions are designed to be slow



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- We can identify three main reasons:
  - Password duplication prevention
    - If two users share the same password, the use of different salt produce different hashed passwords
  - Increase the difficultness of dictionary attacks
    - If the salt has b bits, the factor will be 2<sup>h</sup>b
  - Impossible to find out if a person uses the same password in different systems

# **Password Cracking**

#### ➤ dictionary attacks

- try each word then obvious variants in large dictionary against hash in password file
  - First try all common password
  - If there is no-matches, we try possible modifications (numbers, punctuation)
  - Computationally expensive

#### > rainbow table attacks

- oprecompute tables of hash values for all salts
- o a mammoth table of hash values
- e.g. 1.4GB table cracks 99.9% of alphanumeric
  Windows passwords in 13.8 secs
- not feasible if larger salt values used

#### **Password Choices**

- > users may pick short passwords
  - e.g. 3% were 3 chars or less, easily guessed
  - system can reject choices that are too short
- > users may pick guessable passwords
  - so crackers use lists of likely passwords
  - e.g. one study of 14000 encrypted passwords guessed nearly 1/4 of them
  - would take about 1 hour on fastest systems to compute all variants, and only need 1 break!

#### **Password File Access Control**

- can block offline guessing attacks by denying access to encrypted passwords
  - make available only to privileged users
  - often using a separate shadow password file
- > still have vulnerabilities
  - exploit O/S bug
  - accident with permissions making it readable
  - users with same password on other systems
  - access from unprotected backup media
  - sniff passwords in unprotected network traffic

#### Password Selection StrategieS

- > clearly have problems with passwords
- > goal to eliminate guessable passwords
- ➤ If we cannot trust users ... then we can guide them
- > techniques:
  - user education -> they can ignore it
  - computer-generated passwords -> difficult to remember
  - reactive password checking -> require resources
  - proactive password checking -> likely the best solution

# **Proactive Password Checking**

- > rule enforcement plus user advice, e.g.
  - 8+ chars, upper/lower/numeric/punctuation
  - may not suffice
- > password cracker
  - time and space issues
- ➤ Markov Model
  - generates guessable passwords
  - hence reject any password it might generate
- >> Bloom Filter
  - use to build table based on dictionary using hashes
  - check desired password against this table

# **Proactive Password Checking**

- Bloom Filter
  - use to build table based on dictionary using hashes
  - o check desired password against this table
- ➤ Bloom Filter mechanism
  - K order: K hash functions defined in a range [0, ..., N-1]
  - Given a password, it calculates k hashed passwords
  - Maintain a Table of Size N
    - Table[hash] = 1 over the words of the common words dictionary
  - Given a new password
    - It calculates the k-hashes
    - The password is rejected if all the corresponding bits of the hash table are equal to 1

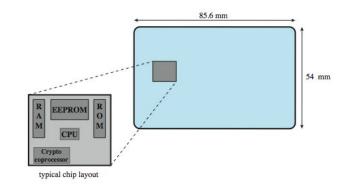
#### **Token Authentication**

- ➤ object user possesses to authenticate, e.g.
  - embossed card (e.g., old credit cards)
  - magnetic stripe card (e.g., hotel keys)
  - memory card (e.g., SIM)
  - Smartcard (e.g., Biometric ID card)

# **Memory Card**

- > store but do not process data
- magnetic stripe card, e.g. bank card
- > electronic memory card
- > used alone for physical access
- > with password/PIN for computer use
- > drawbacks of memory cards include:
  - need special reader (increase the cost of the security solution)
  - loss of token issues (we cannot trust users)
  - user dissatisfaction (not totally approved by users)

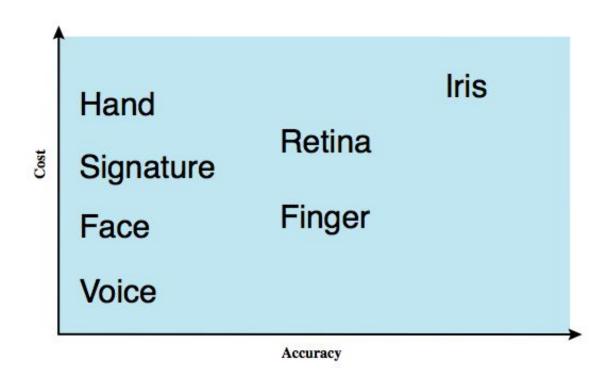
#### **Smartcard**

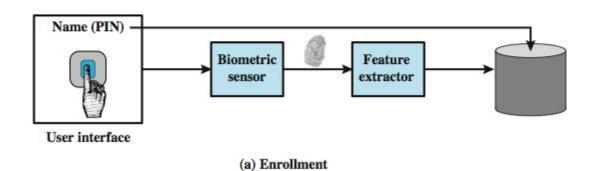


- > credit-card like
- ➤ has own processor, memory, I/O ports
  - wired or wireless access by reader
  - may have crypto co-processor
  - ROM, EEPROM, RAM memory
- executes protocol to authenticate with reader/computer
- ➤ Another example is the USB dongle

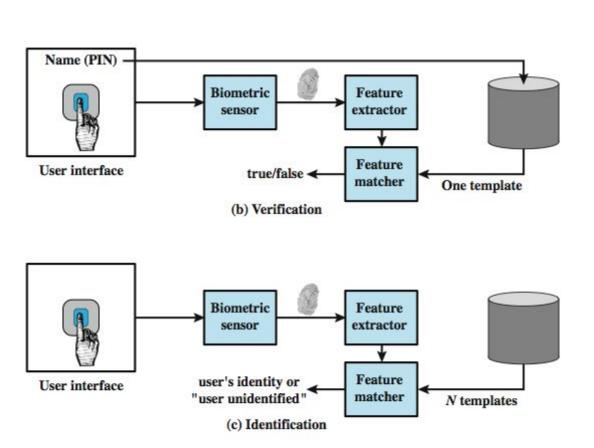
#### **Biometric Authentication**

authenticate user based on one of their physical characteristics



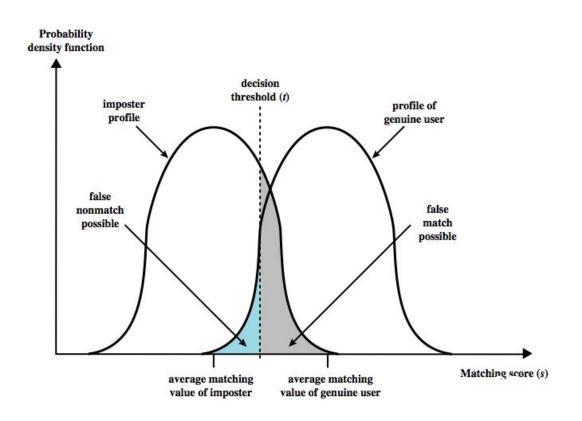


# Operation of a Biometric System



# **Biometric Accuracy**

- > never get identical templates
- > problems of false match / false non-match



#### Remote User Authentication

- > Two type of authentications:
  - Local and from remote
- > authentication over network more complex
  - o problems of eavesdropping, replay
- - ouser sends identity
  - o host responds with:
    - random number r (a.k.a. nonce)
    - An hash function h
    - A function *f*
  - ouser computes f(r,h(P)) and sends back
  - host compares value from user with own computed value, if match user authenticated
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#### **Authentication Security Issues**

- >client attacks
  - adversary attempts to masquerade as a legitimate user (e.g., exhaustive search)
- >host attacks
  - Attack to the host which contains hashed passwords
- > Eavesdropping
  - Learn the password (e.g., by observing the user)
- > Replay
  - Repetition of previously captured user response
- - Malicious application that register the password
- ➤ Denial-of-service
  - Disable the user authentication by flooding the service

### Summary

- > introduced user authentication
  - using passwords
  - using tokens
  - using biometrics
- > remote user authentication issues