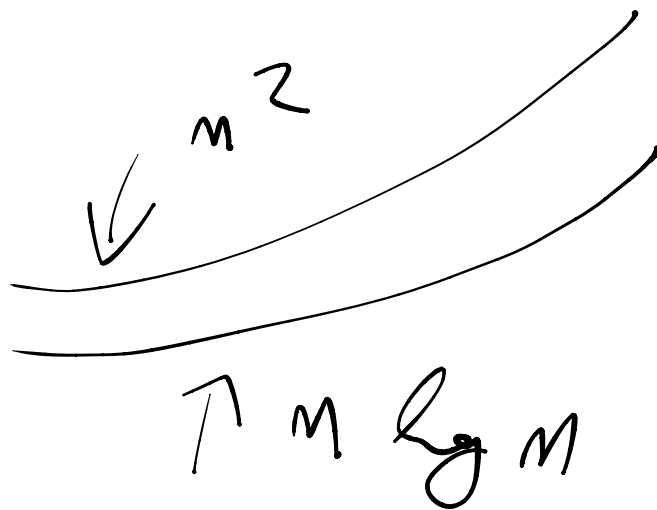


→ INSERTION SORT (n^2)

→ MERGE SORT ($n \lg n$)

→ HEAP SORT ($n \lg n$)

$$\lim_{n \rightarrow \infty} \frac{f(n)}{n \lg n}$$



$$O(n^2) > O(n)$$

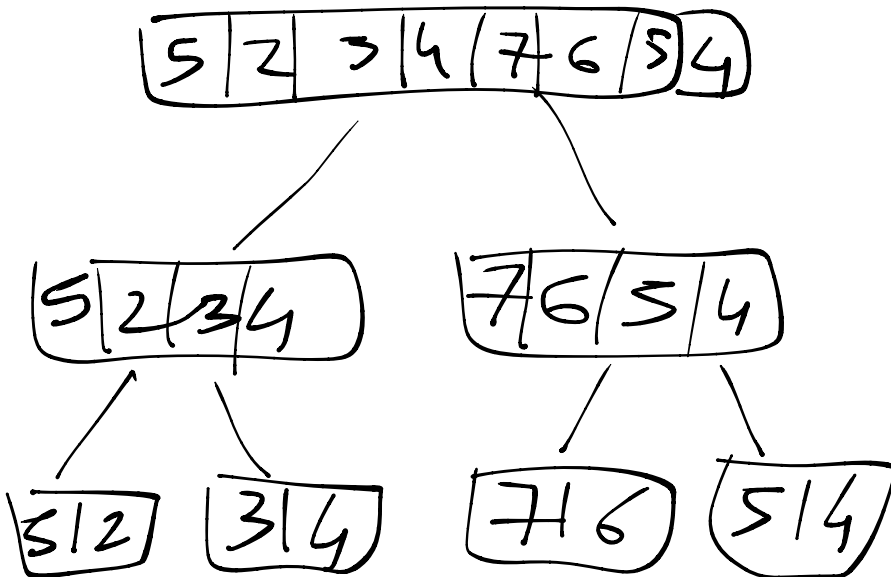
$$> O(n \lg n) \dots > O(1)$$

INSERTION SORT

→ 2 Cikli

MERGE SORT

(DIVIDES IT IMPROV)



$$n \log(m) < m^2$$

↑

SPACE AND

LE SOLUTION

IN IF

RICORRENZA \Rightarrow

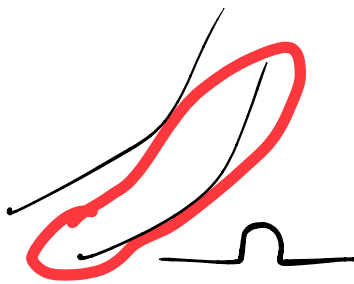
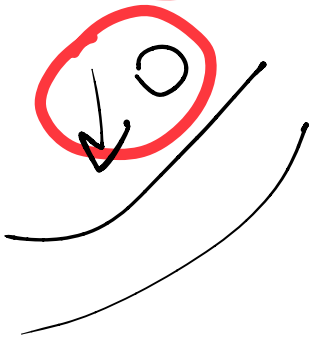
CALCOLO

COMPLESSITÀ

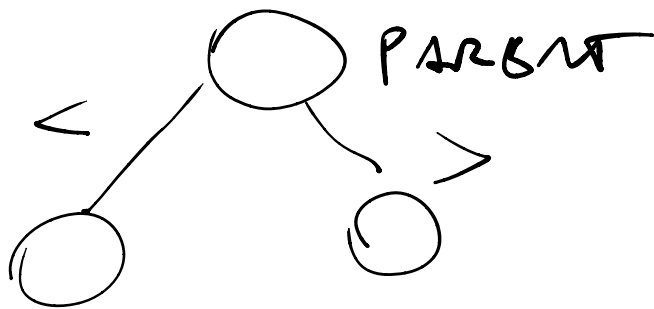
ALGORITMI AD ∞

$\bigcirc \leftarrow$ **OGRADE**

$\bigcirc(\infty)$ \rightarrow **PICCOLO**



HSAP



$A[1]$ $A[2 \times 1]$

LEFT

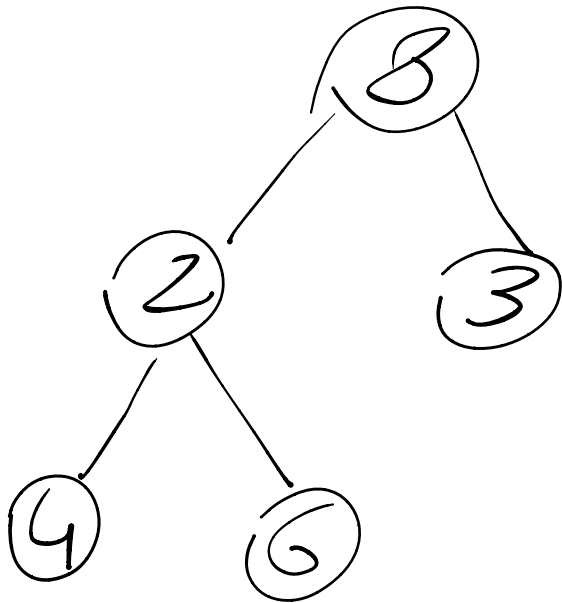
HSAPIFY

1, 2, 3, 4, 5, 6, 7

①

$$A[1/2] \leq A[P]$$

8, 2, 3, 4, 6



8, 10, 12, 14



(8)

(10)

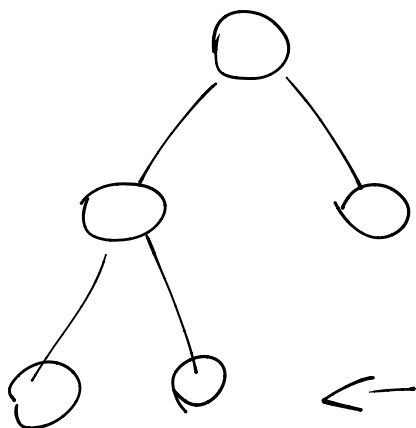
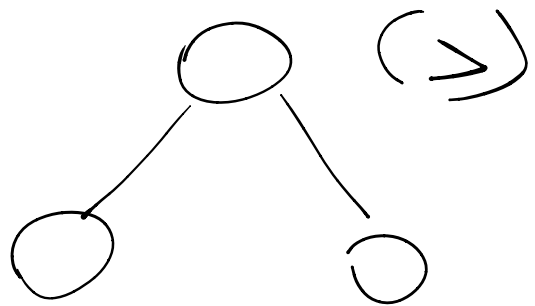
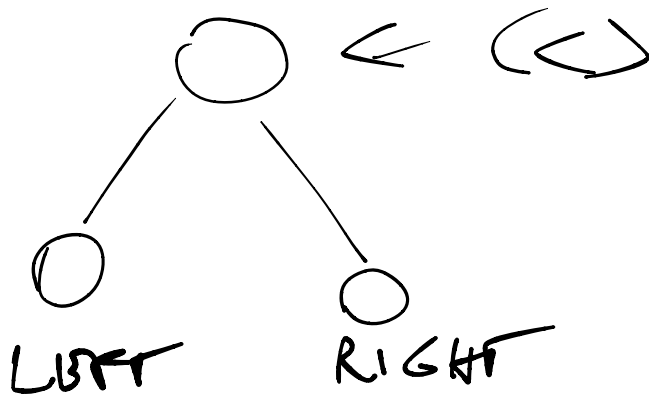
(8)

PARENT > RIGHT & LEFT

→ VOGLIARO

RUSHPIUS 1 1 1 1 1 1

MIN-HEAP / MAX-HEAP



$O(h)$

$= O(\lg(n) + 1)$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{n \lg b(a)} \begin{cases} 1 & (O) \\ 2 & (\Theta) \\ 3 & (\Omega) \end{cases}$$

$$T(n) = a T\left(\frac{n}{b}\right) + f(n)$$

STRUTTURA

FISSA

RECURSIVE!

TRAD.

$O \rightarrow$ SOPRA

$\Theta \rightarrow$ METÀ

$\Omega \rightarrow$ SOTTO

$$T(n) = \overset{a=1}{1} T(\overset{\downarrow M}{n-2}) + \underbrace{2M}_{f(n)}$$

\uparrow
 $T\left(\frac{n}{b}\right)$

$$\lim_{n \rightarrow \infty} \frac{2(n) \rightarrow f(n)}{n \lg b(a) \rightarrow n \lg 1} = \frac{2}{1} = 2$$

CASO 2



COSTANTE

$$\lim_{n \rightarrow \infty} (\dots) \rightarrow 0 \rightarrow \frac{1}{0} = \infty$$

$$\rightarrow K = \text{constant}(\theta)$$

$$\rightarrow n = \frac{1}{\infty} = 0$$

$$\begin{aligned} \text{O} &\rightarrow T(n) \leq C(n) \\ \Omega &\rightarrow T(n) \geq d(n) \end{aligned}$$

$$\text{A. } \sim \nearrow$$

$$\text{O} \rightarrow T(n) \leq C(n)$$

$$\rightarrow T(n) \geq d(n)$$

$$n=2 \rightarrow n^k$$

$$\rightarrow T(n) \leq C(n^2)$$

$$\rightarrow T(n) \geq o(n^2)$$

$$\underline{T(n) \leq cn^2}$$

↑

SOLUZIONI

→ M

→ C

→ VALORI ∃ !

VALORI

DI C e di M

$$C \geq \frac{4}{3}$$

0

$$C \geq 2 \quad \forall \quad n \geq 1$$

$$T(n) \leq c(n^2)$$

$$\leq (T(n-2))^2 + 2n$$

$$\Rightarrow c(n-2)^2 + 2n \leq cn^2$$

$$c(n^2 + 4 - 4n) + 2n \leq cn^2$$

$$\cancel{cn^2} + 4c - 4cn + 2n \leq \cancel{cn^2}$$

$$\underline{4C - 4Cn + 2n} \leq 0$$

$$\left[C \geq \frac{1}{2}, n \geq 1 \right]$$

$$\begin{aligned} T(n) &\geq d(n^2) \\ &\geq d(n-2)^2 + 2n \end{aligned}$$

$$d \geq 1, \forall n \in \mathbb{N}$$

→ ① RICORRENZA
EQUAZIONE

DVZ

→ ② DIVISIONE CASI

$$T(n) = \underline{O(n \lg n)}$$

CASE 3

↳

CONDIZIONE n

RE LARGITA'

CASO 3 \rightarrow

$$\left| f\left(\frac{m}{b}\right) - f(m) \right| \leq \frac{1}{b}$$

PER RISOLVERE UN CASO 3
DEVO USARE QUESTA!

SCQ QUIZZATO

