

Theoretical Computer Science Exam, July 8, 2022

The exam consists of **4 exercises**. Available time: 1 hr 30 min.; you may write your answers in Italian, if you wish.

All exercises must be addressed for the written exam to be considered sufficient.

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Exercise 1 (8 pts)

Consider the language

$$L = \{ a^{2n+1} b^{3n-1} \mid n > 0 \}$$

For instance, strings $aaabb$, a^5b^5 , $a^7b^8 \in L$.

- Write a grammar, of the least general type, that generates language L . For this grammar, write a derivation of the string a^5b^5 .
- Design an automaton, of the least powerful type, that accepts language L . Write a computation of the automaton for the string $aaabb$ and one for the string a^5b^5 .

Exercise 2 (8 pts)

Consider a subprogram, called *ratios*(*in*, *nIn*, *out*), having three parameters: an array of integers *in*, its size *nIn*, an array of integers *out*. The input parameters are: *in* and *nIn*, whereas the array *out* is an output parameter. The two arrays can be assumed to have the same size (i.e., number of components), greater than 1; all the elements of array *in* are assumed to be strictly positive and strictly increasing; every element of the array *in*, except the first one, is a multiple of the preceding one. For instance, an acceptable input array *in* is the following [3, 9, 18, 54, 270].

After execution of the subprogram the first element of array *in* is also stored, in the first position, into the array *out*; furthermore, every element of array *out* in the other positions is equal to the value of the element having the same position in array *in* divided by the previous element in array *in*. For instance, the specification is satisfied by the following values: *nIn*=5, *in*=[3, 9, 18, 54, 270], *out*=[3, 3, 2, 3, 5].

Please answer the following questions:

- write pre- and post-conditions specifying the requirements of the above described procedure; assume that indices of the arrays start at 0;
- for each of the following input/output data say if the procedure *ratios* satisfies the specification if it computes the given value of the *out* array when it receives the given input parameter values; please provide suitable justifications for your answers.

a) *nIn*=5, *in*=[2, 7, 9, 11, 6], *out*=[2, 9, 7, 13, 8]

b) *nIn*=5, *in*=[2, 6, 12, 48, 96], *out*=[2, 3, 2, 3, 2]

c) *nIn*=5, *in*=[2, 6, 12, 48, 96], *out*=[2, 3, 2, 4, 2]

Exercise 3 (8 pts)

Let us consider a fixed finite set S of natural numbers.

1. Is the set S' , that includes exactly the natural numbers that are multiple of all elements of S , decidable?
2. Is the set S'' of Turing machines that, starting the computation with an empty input tape, *do not* print all the numbers that belong to S decidable?
3. Is the set S''' of Turing machines that, starting the computation with an empty input tape, *do not* print all the numbers of S semidecidable?

Please provide precise, detailed explanations for your answers.

Exercise 4 (8 pts)

The following C-like program reads from the input a positive integer value n and writes the first n powers of 2, that is, $2, 2^2, \dots, 2^n$.

```
int n, ev;
read(n);
ev = 1;
for ( i from 1 to n ) {
    ev = 2 * ev;
    write(ev);
}
```

Assuming that the algorithm, after a standard translation, is executed on a RAM machine, evaluate its time and space complexity class using as parameter the **value** of input n , adopting both the uniform and the logarithmic cost criteria.

In addition, evaluate the same complexity figures using as parameter the input **size**, that is, the length x of the string that encodes, in a standard way, the value of n .

Solutions

Exercise 1

Consider the language

$$L = \{ a^{2n+1} b^{3n-1} \mid n > 0 \}$$

For instance, strings $aaabb$, a^5b^5 , $a^7b^8 \in L$.

- a) Write a grammar, of the least general type, that generates language L . For this grammar, write a derivation of the string a^5b^5 .

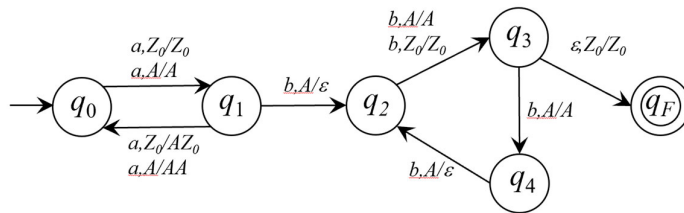
The language is clearly context-free.

$$S \rightarrow aa S bbb \mid aaabb$$

$$S \Rightarrow aa S bbb \Rightarrow aa aaabb bbb$$

- b) Design an automaton, of the least powerful type, that accepts language L . Write a computation of the automaton for the string $aaabb$ and one for the string a^5b^5 .

A deterministic pushdown automaton suffices



$\langle q_0, aaabb, Z_0 \rangle \vdash \langle q_1, aabb, Z_0 \rangle \vdash \langle q_0, abb, AZ_0 \rangle \vdash \langle q_1, bb, AZ_0 \rangle \vdash \langle q_2, b, Z_0 \rangle \vdash \langle q_3, \epsilon, Z_0 \rangle \vdash \langle q_F, \epsilon, Z_0 \rangle$

$\langle q_0, aaaaabbbbb, Z_0 \rangle \vdash \langle q_1, aaaabbbbb, Z_0 \rangle \vdash \langle q_0, aaabbbbb, AZ_0 \rangle \vdash \langle q_1, aabbbbb, AZ_0 \rangle \vdash \langle q_0, abbbbb, AAZ_0 \rangle \vdash \langle q_1, bbbbb, AAZ_0 \rangle \vdash \langle q_2, bbbb, AZ_0 \rangle \vdash \langle q_3, bbb, AZ_0 \rangle \vdash \langle q_4, bb, AZ_0 \rangle \vdash \langle q_2, b, Z_0 \rangle \vdash \langle q_3, \epsilon, Z_0 \rangle \vdash \langle q_F, \epsilon, Z_0 \rangle$

Exercise 2

1. write pre- and post-conditions specifying the requirements of the above described procedure; assume that indices of the arrays start at 0;

precondition

$$nIn > 1 \wedge a[0] > 0 \wedge \forall i (0 < i < nIn \rightarrow \exists k (k > 1 \wedge a[i] = k \cdot a[i-1]))$$

postcondition

$$out[0] = in[0] \wedge \forall i (0 < i < nIn \rightarrow out[i] = in[i] / in[i-1])$$

2. for each of the following input/output data say if the procedure *ratios* satisfies the specification if it computes the given value of the *out* array when it receives the given input parameter values; please provide suitable justifications for your answers.

a) $nIn=5, in=[2, 7, 9, 11, 6], out=[2, 9, 7, 13, 8]$ yes, precondition is not satisfied

b) $nIn=5, in=[2, 6, 12, 48, 96], out=[2, 3, 2, 3, 2]$ no, precondition is satisfied, postcondition is not, because $48 = 4 \cdot 12$

c) $nIn=5, in=[2, 6, 12, 48, 96], out=[2, 3, 2, 4, 2]$ yes, both precondition and postcondition are satisfied

Exercise 3

Let us consider a fixed finite set S of natural numbers.

1. Is the set S' , that includes exactly the natural numbers that are multiple of all elements of S , decidable?

Yes, because for any natural number it is possible to determine, by a simple procedure that always terminates, if it is multiple of all numbers in S .

2. Is the set S'' of Turing machines that, starting the computation with an empty input tape, *do not* print all the numbers that belong to S decidable?

No, as a consequence of the Rice theorem.

3. Is the set S'' of Turing machines that, starting the computation with an empty input tape, *do not* print all the numbers of S semidecidable?

First, notice that S'' is not decidable. Then notice that the set $\neg S''$, the complement of S'' , is the set of the TM that print all the elements of S (and possibly other ones). The set $\neg S''$ is semidecidable: any given TM can be executed with an empty input tape: if it eventually prints all elements of S then the given TM is established to belong to the set $\neg S''$. Otherwise the simulation never terminates and no answer is provided. Therefore, since the set $\neg S''$ is semidecidable, the set S'' cannot be semidecidable, otherwise it would also be decidable.

Exercise 4

Let us assume that the variables of this program are stored in the RAM memory as follows

$n \Rightarrow M[1]$
 $ev \Rightarrow M[2]$

With the constant cost criterion the time complexity class is $\Theta(n)$, because the loop is executed n times; the space complexity class is $\Theta(1)$, because a constant amount of memory is used.

With the logarithmic cost criterion, the space complexity class is $\Theta(\log 2^n) = \Theta(n)$ because the value 2^n must be stored and no value greater than 2^n is ever computed nor stored.

Concerning the time complexity class, the dominating factor in the complexity estimation derives from the execution of the code fragment $ev = 2 * ev$. This will be translated into a sequence of RAM instructions as shown here below, where the cost of each instruction, at the i -th iteration of the loop, is also reported.

LOAD 2 $\log 2 + \log 2^{i-1}$
 MUL =2 $\log 2 + \log 2 + \log 2^{i-1}$
 STORE 2 $\log 2 + \log 2^i$

The overall cost is therefore $T(n) \approx \sum_{i=1}^n (k + 3 \log 2^i) \approx k n + 3/2 n (n+1)$
 so that $T(n) \in \Theta(n^2)$.

If we adopt as parameter the input *size*, that is, the length x of the string that encodes, in a standard way, the value of n , then we have $x = \log n$, hence $n = 2^x$, and the following figures:

With the constant criterion, space is $\Theta(1)$ and time is $\Theta(2^x)$.

With the logarithmic criterion, space is $\Theta(2^x)$ and time is $\Theta(2^{2x})$.