

Università di Pisa

Tenth hands-on: Game Theory II

Algorithm Design (2021/2022)

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May 2022

1 Introduction

1.1 Problem 1: Crossing Streets

Two players drive up to the same intersection at the same time. If both attempt to cross, the result is a fatal traffic accident. The game can be modeled by a payoff matrix where crossing successfully has a payoff of 1, not crossing pays 0, while an accident costs 100.

- Build the payoff matrix
- Find the Nash equilibria
- Find a mixed strategy Nash equilibrium
- Compute the expected payoff for one player (the game is symmetric)

1.2 Problem 2: Mixed strategy for Back Stravinsky Game

Find the mixed strategy and expected payoff for the Back Stravinsky game.

1.3 Problem 3: Children to Kindergartens

The Municipality of your city wants to implement an algorithm for the assignment of children to kindergartens that, on the one hand, takes into account the desiderata of families and, on the other hand, reduces 'city traffic caused by taking children to school. Every school has a maximum capacity limit that cannot be exceeded under any circumstances. As a form of welfare the Municipality has established the following two rules:

- 1. in case of a child already attending a certain school, the sibling is granted the same school:
- 2. families with only one parent have priority for schools close to the workplace.

Model the situation as a stable matching problem and describe the payoff functions of the players. Question: what happens to twin siblings?

2 Solution

2.1 Solution for 'Crossing Streets'

The action set is build with two events $A = \{Cross, Stop\}$. With $D_i, i \in [1, 2]$ we define the i-th driver that is crossing the street. The payoff table is shown in Table 2.1.

As we can see from the table, we have two Nash equilibrium, which are the profiles: (Cross, Stop) and (Stop, Cross).

Table 1: Payoff Table of Crossing Streets

To find the mixed strategies for this problem, we solve the following equations:

$$\begin{cases} \mu_{D_1}(\mathit{Cross}) = \mu_{D_1}(\mathit{Stop}) \\ \mu_{D_2}(\mathit{Cross}) = \mu_{D_2}(\mathit{Stop}) \end{cases}$$
 \iff
$$\begin{cases} -100q + 1(1-q) = 0q + 0(1-q) \\ -100p + 1(1-p) = 0p + 0(1-p) \end{cases}$$
 \iff
$$\begin{cases} q = \frac{1}{101} \\ p = \frac{1}{101} \end{cases}$$

Since the game is symmetric, the expected payoff is equal for both players:

$$\begin{split} E[\![\mu_{D_1}]\!] &= E[\![\mu_{D_2}]\!] \\ &= pq(-100) + p(1-q)1 + (1-p)q0 + (1-p)(1-q)0 \\ &= -100pq + p - qp = -101pq + p = p(1-101q) = p(1-101p) = 0 \end{split}$$

2.2 Solution for 'Mixed Strategy for Back Stravinsky Game'

The payoff table for Bach-Stravinsky game is the following one:

		Wife	
		Bach	Stravinsky
Husband	Bach	(2,1)	(0,0)
	Stravinsky	(0,0)	(1, 2)

Table 2: Payoff Table for Bach-Stravinsky

With μ_H and μ_W we refer to the payoff function for the husband and wife. The set of strategies that the couple can choose is the following one:

$$S_H = S_W = \{ \text{Bach}, \text{Stravisnky} \}$$

To find a mixed strategy we need to convince the players to play a randomized strategy instead of a fixed one. We need to find a probability distribution (p, 1-p) for the wife that makes the husband indifferent. This means that we have to equal $\mu_W(\text{Bach}) = \mu_W(\text{Stravisnky})$, where:

$$\mu_W(\text{Bach}) = 1 \times p + 0 \times (1 - p)$$

$$\mu_W(\text{Stravinsky}) = 0 \times p + 2 \times (1 - p)$$

Computing the equation $\mu_W(Bach) = \mu_W(Stravisnky)$ we obtain:

$$\mu_W(\text{Bach}) = \mu_W(\text{Stravisnky})$$

$$\iff 1 \times p + 0 \times (1 - p) = 0 \times p + 2 \times (1 - p)$$

$$\iff p = 2 \times (1 - p)$$

$$\iff p = \frac{2}{3}$$

For the husband we have to act similar. Let us find a probability distribution (q, 1-q) for the husband, such that makes the wife indifferent regard to the selected choice. So find q by solving the equation $\mu_H(\text{Bach}) = \mu_H(\text{Stravinsky})$:

$$\mu_H(\text{Bach}) = \mu_H(\text{Stravisnky})$$

$$\iff 2 \times q + 0 \times (1 - q) = 0 \times q + 1 \times (1 - q)$$

$$\iff 3q = 1$$

$$\iff q = \frac{1}{3}$$

2.3 Solution for 'Children to Kindergartens'

Given two sets of equal size, the *stable matching problem* consists of matching each element of one set to one in the other. In this problem each element has some preferences.

The following solution gives the idea behind the real one. We can describe the stable matching problem creating two sets: the first one is the set containing the children will go at schools, and, since the number of schools of a city does not match the number of children, the second set is the number of "seats" available in each school. For example, given n schools indicated with S_i (the i-th school), each with its own capacity c_i , we define the slots in the following way:

$$s_{i_k} = k$$
-th slot available in the school i for $0 \le k \le c_i$

According to the game's description, we can model the preferences, for each child, as the distance between the child's location (indicated with H_i for the house and W_i for the parent's workplace) and the school (defining a function $d(H_i, S_i)$).

First, families with only parent have the higher priority overall, so we try to minimize the distance between the parent's workplace and the school as:

$$\min_{i} \{d(W_i, S_i)\}$$

Second, if a child (with a sibling) is already displaced in a school, then his/her sibling will be placed in the same one, setting the distance d equal to 0. The twins will happen to be in the same school, since they can be seen as a "unique" child with weight 2.