UNIVERSITÀ DI PISA

Second hands-on: Depth of a node in a Random Search Tree

Algorithm Design (2021/2022)

Gabriele Pappalardo
Email: g.pappalardo4@studenti.unipi.it
Department of Computer Science

March 2022

1 Introduction

A **random binary search tree** for a set S can be defined as follows: if S is empty, then the null tree is a random search tree; otherwise, choose uniformly at random a key $k \in S$: the random search tree is obtained by picking k as *root*, and the random search trees on $L = \{x \in S : x < k\}$ and $R = \{x \in S : x > k\}$ become, respectively, the left and right subtrees of the root k.

Consider the *Randomized Quick Sort* (RQS) discussed in class and analyzed with indicator variables, and observe that the random selection of the pivots follows the above process, thus producing a random search tree of n nodes.

- Using a variation of the analysis with indicator variables X_{ij}, prove that the expected depth of a node (i.e. the random variable representing the distance of the node from the root) is nearly $2 \log n$.
- Prove that the expected size of its subtree is nearly $2 \log n$ too, observing that it is a simple variation of the previous analysis.
- Prove that the probability that the depth of a node exceeds $c2 \log n$ is small for any given constant $c > 2$.

2 Solution

2.1 Expected depth of a node

We define the elements of a random binary search tree with $\forall i \in [1, n].x_i \in S$. Let X_{ij} an indicator random variable defined as follows:

$$X_{ij} = \begin{cases} 1 & x_i \text{ is an ancestor of } x_j \\ 0 & \text{otherwise} \end{cases} \tag{1}$$

The depth of a node can be expressed as the sum of all the indicator random variables. Let D_i be the depth of a node x_i in a random binary search tree, defined as $D_i = \sum_{j=1}^{n} X_{ij}$. We can express the expected depth as:

$$E[D_i] = E\left[\sum_{j=1}^{n} X_{ij}\right] = \sum_{j=1}^{n} E[X_{ij}] = \sum_{j=1}^{n} \Pr(\{X_{ij} = 1\}) \tag{2}$$

The probability of $X_{ij} = 1$ can be computed using the following fact: if the node x_i is an ancestor of x_j it means that x_i is inserted before x_j (and $i > j$), therefore in the set $\{x_j, x_{j+1}, \ldots, x_i\}$ we have $i - j + 1$ elements to choice x_j, otherwise, if x_i is a descendant of x_j it means that x_i has been inserted after ($i < j$), therefore in the set $\{x_i, x_{i+1}, \ldots, x_j\}$ we have $j - i + 1$ elements to choice.

$$\Pr(\{X_{ij} = 1\}) = \begin{cases} \frac{1}{i-j+1} & i > j \\ \frac{1}{j-i+1} & i < j \end{cases} \tag{3}$$

The expected depth of a node can be computed using this fact:

$$\begin{aligned} E[D_i] = \sum_{j=0}^{n} \Pr(\{X_{ij} = 1\}) = \\ \sum_{j=0}^{i} \frac{1}{i-j+1} + \sum_{j=i+1}^{n} \frac{1}{j-i+1} = \\ \sum_{k=1}^{i+1} \frac{1}{k} + \sum_{k=2}^{n-i+1} \frac{1}{k} \leq \log n + \log n = 2\log n \end{aligned}$$

2.2 Expected size of a subtree

The analysis for the expected size of a subtree is similar to the previous one. We change the meaning of our indicator random variable as:

$$X_{ij} = \begin{cases} 1 & x_i \text{ is an descendant of } x_j \\ 0 & \text{otherwise} \end{cases} \tag{4}$$

2.3 Probability for the depth of a node

Knowing that the expected depth of a node $E[D_i] \leq 2\log n$, we can prove that probability that the depth of a node exceeds $c2\log n$ is small for any given constant $c \geq 2$.

$$\mathbf{Pr}(\{D_i \geq c2\log n\}) \leq e^{-\frac{\lambda^2}{2\mu+\lambda}} \tag{5}$$

Knowing that $\mu = E[D_i] \leq 2\log n$ we can derive λ:

$$\begin{aligned} \mu + \lambda = c2\log n \\ E[D_i] + \lambda = c2\log n \\ \lambda = c2\log n - E[D_i] \leq c2\log n - 2\log n = \log n(2c-2) \end{aligned}$$

Therefore, the Chernoff Bound becomes:

$$\begin{aligned} \mathbf{Pr}(\{D_i \geq c2\log n\}) &\leq e^{-\frac{(\log n(2c-2))^2}{2(2\log n)+\log n(2c-2)}} \\ &= e^{-\frac{(\log n)^2(2c-2)^2}{2(2\log n)+\log n(2c-2)}} \\ &= e^{-\frac{(\log n)^2(2c-2)^2}{\log n(2c+2)}} \\ &= e^{-\frac{\log n(2c-2)^2}{(2c+2)}} \\ &= n^{-\frac{(2c-2)^2}{(2c+2)}} \\ &= n^{-\frac{4c^2+4-8c}{2c+2}} \\ &= n^{-\frac{2c^2+2-4c}{c+1}} \\ &= \frac{1}{n^{\frac{2c^2+2-4c}{c+1}}} \end{aligned}$$

At the end, the probability is small for any constant $c > 2$.