

BCH Codes

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Abstract—This manual provides an introduction to BCH codes.

1 GENERATOR POLYNOMIAL

1.1 For a BCH code, the minimal polynomials are given by

$$g_1(x) = 1 + x + x^3 + x^5 + x^{14} \quad (1.1)$$

$$g_2(x) = 1 + x^6 + x^8 + x^{11} + x^{14} \quad (1.2)$$

$$g_3(x) = 1 + x + x^2 + x^6 + x^9 + x^{10} + x^{14} \quad (1.3)$$

$$g_4(x) = 1 + x^4 + x^7 + x^8 + x^{10} + x^{12} + x^{14} \quad (1.4)$$

$$g_5(x) = 1 + x^2 + x^4 + x^6 + x^8 + x^9 + x^{11} + x^{13} + x^{14} \quad (1.5)$$

$$g_6(x) = 1 + x^3 + x^7 + x^8 + x^9 + x^{13} + x^{14} \quad (1.6)$$

$$g_7(x) = 1 + x^2 + x^5 + x^6 + x^7 + x^{10} + x^{11} + x^{13} + x^{14} \quad (1.7)$$

$$g_8(x) = 1 + x^5 + x^8 + x^9 + x^{10} + x^{11} + x^{14} \quad (1.8)$$

$$g_9(x) = 1 + x + x^2 + x^3 + x^9 + x^{10} + x^{14} \quad (1.9)$$

$$g_{10}(x) = 1 + x^3 + x^6 + x^9 + x^{11} + x^{12} + x^{14} \quad (1.10)$$

$$g_{11}(x) = 1 + x^4 + x^{11} + x^{12} + x^{14} \quad (1.11)$$

$$g_{12}(x) = 1 + x + x^2 + x^3 + x^5 + x^6 + x^7 + x^8 + x^{10} + x^{13} + x^{14} \quad (1.12)$$

Obtain the minimal polynomial matrix.

Solution:

https://raw.githubusercontent.com/gadepall/EE6317/master/BCH/codes/min_poly_mat.py

1.2 Obtain the generator polynomial vector.

Solution: The generator polynomial is obtained as

$$g(x) = \prod_{i=1}^m g_i(x) \quad (1.13)$$

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The following code computes \mathbf{g} .

```
https://raw.githubusercontent.com/gadepall/
EE6317/master/BCH/codes/gen_poly.py
```

2 ENCODING

2.1 Let \mathbf{m} be a $k \times 1$ message vector and

$$m(x) = m_{k-1}x^{k-1} + m_{k-2}x^{k-2} + \dots + m_1x + m_0 \quad (2.1)$$

be the corresponding Message polynomial.

2.2 Let

$$m(x)x^{n-k} = q(x)g(x) + d(x) \quad (2.2)$$

and

$$c(x) = m(x)x^{n-k} + d(x) \quad (2.3)$$

Write a program to compute the corresponding coefficient vector \mathbf{c} . This is the output of the BCH encoder.

3 DECODING

3.1 Let the Received polynomial be $r(x)$ i.e which contains both transmitted codeword polynomial $c(x)$ and the error polynomial $e(x)$

$$e(x) = e_0 + e_1x^1 + \dots + e_{n-1}x^{n-1} \quad (3.1)$$

Where e_i represents the value of the error at the location. For binary BCH codes e_i is either 0 or 1.

$$r(x) = c(x) + e(x) = r_{n-1}x^{n-1} + r_{n-2}x^{n-2} + \dots + r_1x + r_0 \quad (3.2)$$

Define, Syndrome

$$S_i = r(\alpha^i) = c(\alpha^i) + e(\alpha^i) = e(\alpha^i) \quad (3.3)$$

Where α^i is a root of the codeword.

Suppose that v errors occurred, and $0 \leq v \leq t$.

Let the error occurs at i_1, i_2, \dots, i_v .

The Decoding Process, for a t -error correcting code will follow the basic steps,

3.2 Compute the Syndrome $S = (S_1, S_2, \dots, S_{2t})$ from the received polynomial $r(x)$

3.3 Determine the error-location polynomial $\sigma(x)$ from the syndrome components S_1, S_2, \dots, S_{2t} using the Berlekamp's Algorithm.

3.4 Using the Chain searching algorithm, determine the error-locations by finding the roots of $\sigma(x)$, the flip the positions in $r(x)$. Which is the estimated message vector polynomial $\hat{c}(x)$.

3.5 where \mathbf{w} is $p \times 1$ and \mathbf{X} is $N \times p$. Show that

$$E(\hat{\mathbf{w}}) = \mathbf{w} \quad (3.4)$$

3.6 If the covariance matrix of \mathbf{y} is

$$\mathbf{C}_y = \sigma^2 \mathbf{I} \quad (3.5)$$

show that

$$\mathbf{C}_w = \sigma^2 (\mathbf{X}^T \mathbf{X})^{-1} \quad (3.6)$$

3.7 Let

$$\hat{\mathbf{y}} = \mathbf{X}\hat{\mathbf{w}} \quad (3.7)$$

$$\hat{\sigma}^2 = \frac{1}{N-p} \|\mathbf{y} - \hat{\mathbf{y}}\|^2 \quad (3.8)$$

$$\mathbf{y} - \hat{\mathbf{y}} \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}) \quad (3.9)$$

Show that

$$(N-p)\hat{\sigma}^2 \sim \sigma^2 \chi_{N-p}^2 \quad (3.10)$$

3.8 Let

$$\hat{z} = \frac{w_j}{\hat{\sigma} \sqrt{v_j}} \quad (3.11)$$

where v_j is the diagonal element of $(\mathbf{X}^T \mathbf{X})^{-1}$. If $w_j = 0$, show that z_j has a t_{N-p} distribution.

3.9 Plot $\Pr(|Z| > z)$ for t_{30}, t_{100} and the standard normal distribution.

4 APPLICATIONS

4.1 Explain how (??) can be used to obtain the Nearest Neighbour approximation.

4.2 Repeat the exercise for the least squares method.