Metaheuristics for the Vehicle Routing Problem: Fifteen Years of Research

by

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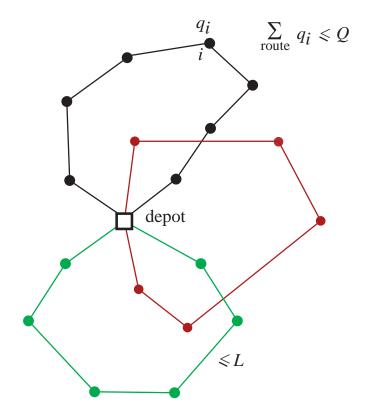
Web site: www.hec.ca/chairedistributique/metaheuristics.pdf

Vehicle Routing Problem

- Depot
- \bullet m (or at most m) identical vehicles based at the depot
- *n* customers
- Distance (cost, travel time) matrix (c_{ij})
- q_i : demand of customer i
- Q: vehicle capacity
- *L*: maximal route length (duration)

VRP: determine a set of m or at most m vehicle routes

- 1. Starting and ending at the depot
- 2. Visiting each customer exactly once
- 3. Satisfying the capacity constraint
- 4. Satisfying the maximal length constraint
- 5. Of minimal total cost



- Problem introduced by Dantzig and Ramser (Management Science, 1959)
- NP-hard
- Has multiple applications
- Exact algorithms: relatively small instances
- In practice heuristics are used
- Several variants e.g.
 - heterogeneous vehicle fleet (Gendreau et al., 1999)
 - time windows (Cordeau et al., VRP book, 2002)
 - pickup and deliveries (Desaulniers et al., VRP book, 2002)
 - periodic visits (Cordeau et al., Networks, 1997)
 - etc.
- Recommended book:
 - P. Toth and D. Vigo, The Vehicle Routing Problem, SIAM Monographs on Discrete Mathematics and Applications, Philadelphia, 2002.

Exact Algorithms

1981 Dynamic Programming with State Space Relaxation (Christofides, Mingozzi, Toth, *Networks*)

 $(10 \le n \le 25)$

Branch-and-bound (k-shortest spanning trees, q-paths) (Christofides, Mingozzi, Toth, *Mathematical Programming* ($10 \le n \le 25$)

Branch-and-cut (Laporte, Nobert, Desrochers, *Operations Research*) $(n \le 60)$

Branch-and-cut (for a restricted version of the VRP) (Fisher, *Operations Research*) $(n \le 135)$

Branch-and-cut (Ralphs et al., website) $(n \le 101)$

Branch-and-cut (Augerat et al., working paper) $(n \le 135)$

2000	Branch-and-cut (Blasum and Hochstättler, working paper) $(n \leq 76)$
2002	Branch-and-cut (Naddef and Rinaldi, <i>VRP Book</i>) (survey)
2003	Branch-and-cut-and-price (Fukasawa et al., <i>Relato-rios de Pesquisa en Engenharia de Produção</i>)
	Branch-and-cut (Wenger, Ph.D. dissertation, University of Heidelberg)
2004	Two-commodity network flow formulation (Baldacci, Hadjiconstantinou, Mingozzi, Operations Research) $(n \leq 135)$

Heuristic Algorithms

- Classical algorithms (Laporte, Semet, VRP Book, 2002)
 - savings (Clarke, Wright, Operations Research, 1965)
 - sweep (Gillett, Miller, Operations Research, 1974)
 - cluster first, route second (Fisher, Jaikumas, Networks, 1981)
 - intra-route improvement methods (TSP heuristics)
 - inter-route improvement methods (λ -interchanges, Osman, 1993; cyclic exchanges, Thompson and Psaraffis, 1993; edge exchange schemes, Kindervater and Savelsbergh, 1997; ejection chains (Xu and Kelly, 1996; Rego and Roucairol, 1996; Rego, 1998); very large neighbourhood search (Ergun et al., 2003)
 - SERR (De Franceschi, Fischetti, Toth, working paper, 2004)

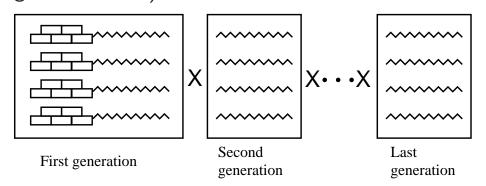
- Metaheuristics (Gendreau, Laporte, Potvin, VRP Book, 2002)
 - local search (simulated annealing, deterministic annealing, tabu search)

Single constructionimprovement thread

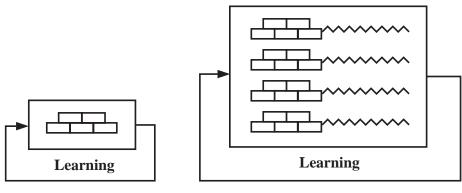
Constructive phase
followed by improvement
in several ways (may be
executed in parallel)

Several constructionimprovement threads (may be executed in parallel)

population search (adaptive memory procedures, genetic search)



learning mechanisms (neural networks, ant colony systems)



a) Neural networks

b) Ant algorithms

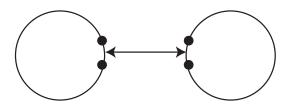
15 years of metaheuristics

1989	First tabu search implementation (Willard, M.Sc. thesis, Imperial College)
1991	First version of Taburoute (Gendreau, Hertz, Laporte, Tristan I Conference)
1993	Tabu search (Taillard, Networks)
1993	Simulated Annealing and tabu search (Osman, Annals of Operations Research)
1994	Taburoute (Gendreau, Hertz, Laporte, <i>Management Science</i>)
1995	Adaptive memory (Rochat, Taillard, <i>Journal of Heuristics</i>)
1996	Ejection chains (Rego, Roucairol, <i>Meta-Heuristics:</i> Theory and Applications)

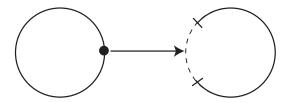
2001	Unified tabu search algorithm (Cordeau, Laporte, Mercier, <i>Journal of the Operational Research Society</i>)
2002	Adaptive memory (Tarantilis, Kiranoudis, <i>Annals of Operations Research</i>)
2003	Granular tabu search (Toth, Vigo, INFORMS Journal on Computing)
2003	Very large neighbourhood search (Ergun, Orlin, Steele-Feldman, working paper, MIT)
2004	Deterministic annealing (Li, Golden, Wasil, Computers & Operations Research)
2004	Population search (Prins, Computers & Operations Research; Mester and Bräysy, Computers & Operations Research)
2004	Ant systems optimization (Reinmann, Doerner, Hartl, Computers & Operations Research)

Algorithmic ideas

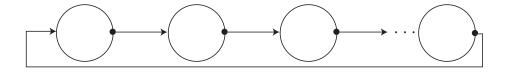
- Neighbourhood structures
 - 2-interchanges (Taillard, 1993)



simple vertex moves combined with local reoptimization (GENI) (Taburoute and UTSA)



composite moves (ejection chains, very large neighbourhood search)



- attribute sets $B(x) = \{(i,k): i \text{ is visited by vehicle } k \\ \text{in solution } x\}$ Remove (i,k) from B(x) and replace it with (i,k') $(k' \neq k)$

Tabu tags

- Assign tabu tag to an attribute (instead of maintaining an actual tabu list)
- Tabu duration: variable in Taillard (1993) and in Taburoute, fixed but size-dependent in UTSA
- Aspiration criteria (overriding tabu status)
 - Attribute related in UTSA
- Intermediate infeasible solutions (Taburoute, UTSA)

$$F'(x) = F(x) + \alpha Q(x) + \beta D(x)$$

where α and β are periodically updated (almost necessary if simple vertex moves are used).

Continuous diversification (Taillard)

Penalize cost of worsening candidate solutions by adding to their cost a penalty proportional to the frequency of move:

$$F(x) := F(x) + \gamma \sqrt{mn} f_{ik}$$

- Periodic route reoptimization
- False starts

Used in Taburoute but not in UTSA: better perform 10^5 iterations on one solution than 10^4 iterations on each of 10 solutions.

Intensification

Used in Taburoute but not in UTSA.

 Data perturbation (Codenetti et al., INFORMS Journal on Computing, 1996)

Used in Latest version of UTSA (0.69% \rightarrow 0.56%): temporarily relocate the depot to next vertex of a route.

- Granularity (Toth, Vigo)
 - Remove long edges from data to obtain a sparse distance matrix.

granularity threshold: $\nu=\beta\bar{c}$, where \bar{c} is the average edge cost in a good feasible solution sparsification parameter $\beta\in[1.0,2.0]$ keep edges incident to the depot and those for which $c_{ij}\leq\nu$

- Applied by Toth and Vigo: 4 times faster than Taburoute, also better.
- Applied by Li, Golden, Wasil in conjunction with record-to-record principle (Dueck, 1993): accept candidate neighbour if cost does not exceed 1.01 times cost of best known solution.

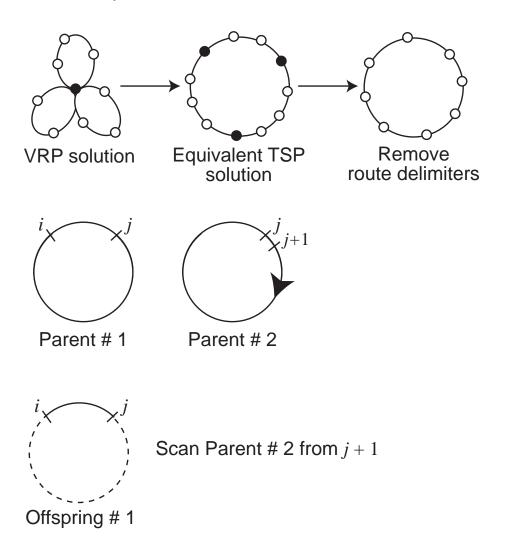
Adaptive memory (Rochat, Taillard)

Keep a pool of good solutions, combine them and reoptimize.

- Rochat, Taillard: select a route from each of several solutions until this cannot be done without overlaps (→ several routes + loose vertices). Reoptimize.
- BoneRoute (Tarantilis, Kiranoudis): extract segments (bones) from good quality routes.

• Solution recombination (used in genetic search, Prins, 2004)

Solution representation:



For offspring # 2 reverse the role of the two parents

 Guided evolution (AGES: active guided evolution strategy, Mester, Bräysy)

Create each offspring from a single parent: apply local search, penalize some solution features (e.g. very long edges), use continuous diversification, 2-opt moves, 2-interchanges, very large neighbourhoods, restarts from best known solution.

Memetic search (Moscatto and Cotta, 2003)

Combines genetic search with local search. Improve offspring by local search.

Advantage: provides width and depth. Applied by Prins and Mester and Bräysy.

 Learning (D-ants savings algorithm of Reimann, Doerner and Hartl)

Generate a pool of good solution by Clarke and Wright savings algorithm and improve them. Replace saving criterion $s_{ij}=c_{i0}+c_{0j}-c_{ij}$ with $t_{ij}^{\alpha}s_{ij}^{\beta}$ where

 t_{ij}^{α} contains information on how good combining i and j turned out to be in previous solutions and α , β are user-controlled parameters.

Apply saving s_{ij} with probability p_{ij} .

Computational results

• 14 Christofides, Mingozzi, Toth (1979) instances

$$(50 \le n \le 199)$$

20 Golden et al. (1998) instances

$$(200 \le n \le 480)$$

- Comparisons with best known results
- Two average statistics are reported:

% over best: average deviation over best

minutes: CPU time in minutes

 Machine is reported (for comparison of machine sppeds, see J.J. Dongarra, Performance of various computers using standard linear equations software. Technical report CS-89-85, Computer Science Department, University of Tennessee).

Christofides, Mingozzi, Toth instances $(50 \le n \le 199)$

%					
Authors	over best	Minutes	Machine		
Taillard (1993) (tabu search parallel)	0.02	21.30 (1%)	Silicon Graphics (4D135)		
Gendreau, Hertz, Laporte (1994) (Taburoute standard version)	0.82	55.59	Silicon Graphics (36 MHz)		
Rochat, Taillard (1995) (adaptive memory)	0.00	16.19 (1%)	Silicon Graphics Indigo (100 MHz)		
Rego, Roucairol (1996) (Tabuchain sequential) (Tabuchain parallel)	0.58 0.37	29.06 16.27	Sun Sparc IPC Sun Sparc IPC		
Rego (1998) (Flower)	1.68	2.51	HP9000/712		
Tarantilis, Kiranoudis (2002) (BoneRoute)	0.23	5.22	Pentium II (400 MHz)		
Ergun et al. (2003) (very large neighbourhoods)	0.23 (best of 5 runs)	28.91 (best value first time)	Pentium III (733 MHz)		
Toth, Vigo (2003) (granular tabu search)	0.64	3.84	Pentium (200 MHz)		
Cordeau, Laporte, Mercier (2001) (UTSA, 2004 version)	0.56	24.62	Pentium IV (2 GHz)		

	%		
Authors	over best	Minutes	Machine
Prins (2004) (Memetic algorithm)	0.24	5.19	GHz PC
Mester, Bräysy (2004) (AGES best)	0.03	7.72	Pentium IV (2 GHz)
(AGES fast)	0.07	0.27	Pentium IV (2 GHz)
Berger, Berkaoui (2004) (genetic search)	0.49	21.25	Pentium (400 MHz)
Li, Golden, Wasil (2004) (record-to-record)	0.41 (best variant)	_	_

Golden, Wasil, Kelly, Chao instances $(200 \le n \le 480)$

Authors	% over best	Minutes	Machine
Ergun et al. (2003) (very large neighbourhoods)	3.76 (best of 5 runs)	137.95 (best value first time)	Pentium III (733 MHz)
Toth, Vigo (2003) (granular tabu search)	2.87	17.55	Pentium (200 MHz)
Cordeau, Laporte, Mercier (2001) (UTSA, 2004 version)	1.45	56.11	Pentium IV (2 GHz)
Prins (2004) (Memetic algorithm)	0.91	66.90	GHz PC
Tarantilis, Kiranoudis (2002) (BoneRoute)	0.74 (8 C instances)	42.05	Pentium II (400 MHz)
Mester, Bräysy (2004) (AGES best) (AGES fast)	0.00 0.93	72.94 0.63	Pentium IV (2 GHz)
Li, Golden, Wasil (2004) (record-to-record)	1.05	_	
Reimann, Doerner, Hartl (2004) (D-ants)	0.60	49.33	Pentium (900 MHz)

Assessment

Heuristics should be assessed on

- accuracy
- speed
- simplicity
- flexibility

(Cordeau, J.-F., Gendreau, M., Laporte, G., Potvin, J.-Y., Semet, F., "A guide to vehicle routing heuristics", *Journal of the Operational Research Society*, 53, 512–522, 2002)

- Overall winners: Mester, Bräysy (AGES), Tarantilis, Kiranoudis (BoneRoute), Prins: these algorithms combine population search and local search
- Winner (accuracy and speed): Mester, Bräysy (AGES)
- Winner (simplicity): Li, Golden, Wasil (record-to-record)
- Runners up (simplicity): Prins; Toth, Vigo (granular tabu search); Cordeau, Laporte, Mercier (UTSA)
- Mentions (flexibility): Toth, Vigo (granular tabu search);
 Rochat, Taillard (adaptive memory), Ergun et al. (very large neighbourhood search); Cordeau, Laporte, Mercier (UTSA: JORS 2001, 2004)

Reference

Cordeau, J.-F., Gendreau, M., Hertz, A., Laporte, G., Sormany, J.S., "New heuristics for the vehicle routing problem", Les Cahiers du GERAD, G-2004-33, HEC Montréal, 2004.

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