

Next Greater Element.

4 5 2 25

$4 \rightarrow 5$

$5 \rightarrow 25$

$2 \rightarrow 25$

$25 \rightarrow -1$

Brute force Approach.

T.C $\rightarrow O(n^2)$

4 5 2 25

5	25	25	-1
---	----	----	----

Worst case

1000 999 998 997 - - - 1 2000
 $i=2$
 $j=2$ - - -

NHΣ

2000	2000	2000	2000	2000	-
------	------	------	------	------	---

1001th times 1000th 999th times

$$\text{arr} = n \quad \begin{matrix} n-1 \\ i=1 \\ j=1 \end{matrix} \quad n-2 \quad - \quad - \quad - \quad - \quad 1 \quad n+1$$

$$\text{NGF} \quad \boxed{n^{\text{th}} \mid (n-1)^{\text{th}} \mid (n-2)^{\text{th}} \mid \dots \mid 1}$$

$$\begin{aligned} & n + (n-1) + (n-2) + \dots + 1 \\ &= \frac{n(n+1)}{2} = O(n^2). \end{aligned}$$

Stack.

4	5	2	25
---	---	---	----

0
↑

1
↑

2
↑

3
~~↑~~

5
2
25

5	25	25	-1
---	----	----	----

NGE()

$O(n)$

↓

Stack s;

next h [];

for(i = arr.length - 1 ; i >= 0 ; i--)

{

while(!s.empty() && s.peek() <= arr[i])

{

s.pop();

}

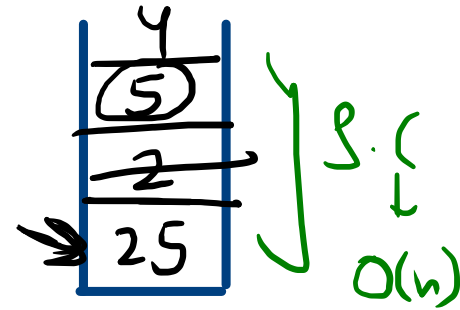
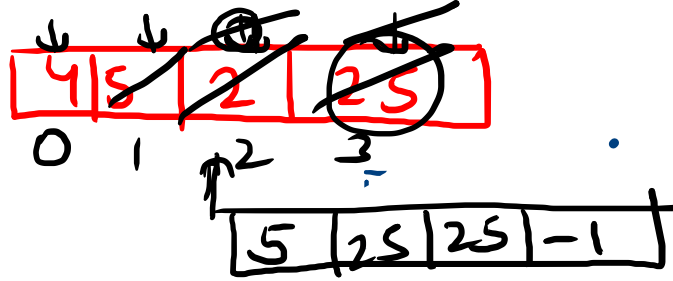
}

↓

```

NGE() {
    Stack s;
    nge[];

```



```

for (int i = arr.length - 1; i >= 0; i--)

```

```

{
    while (!s.empty() && s.peek() <= arr[i])
        s.pop();

```

out

```

}

```

```

if (s.empty()) nge[i] = -1;
else nge[i] = s.peek();

```

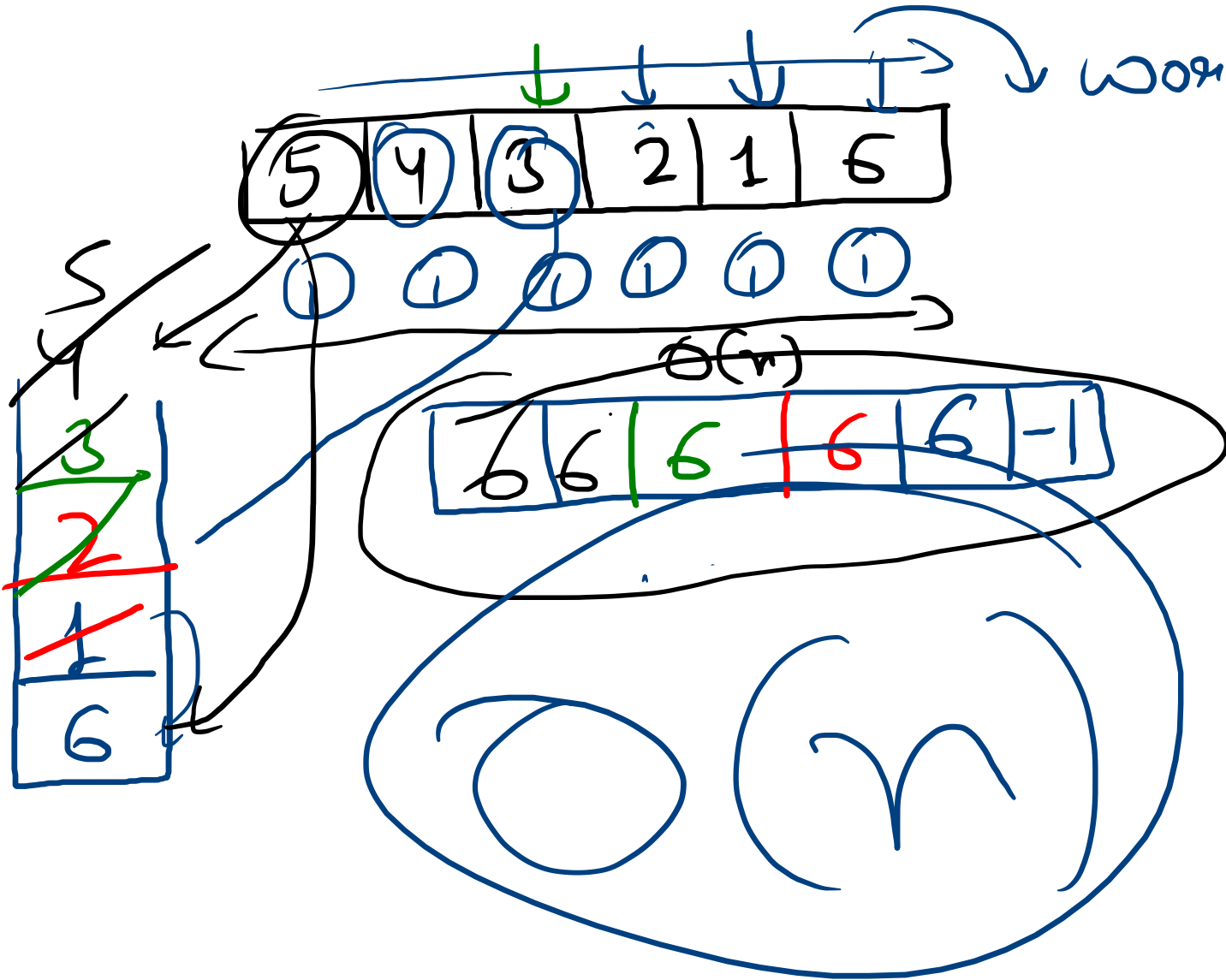
```

s.push(arr[i]);

```

y

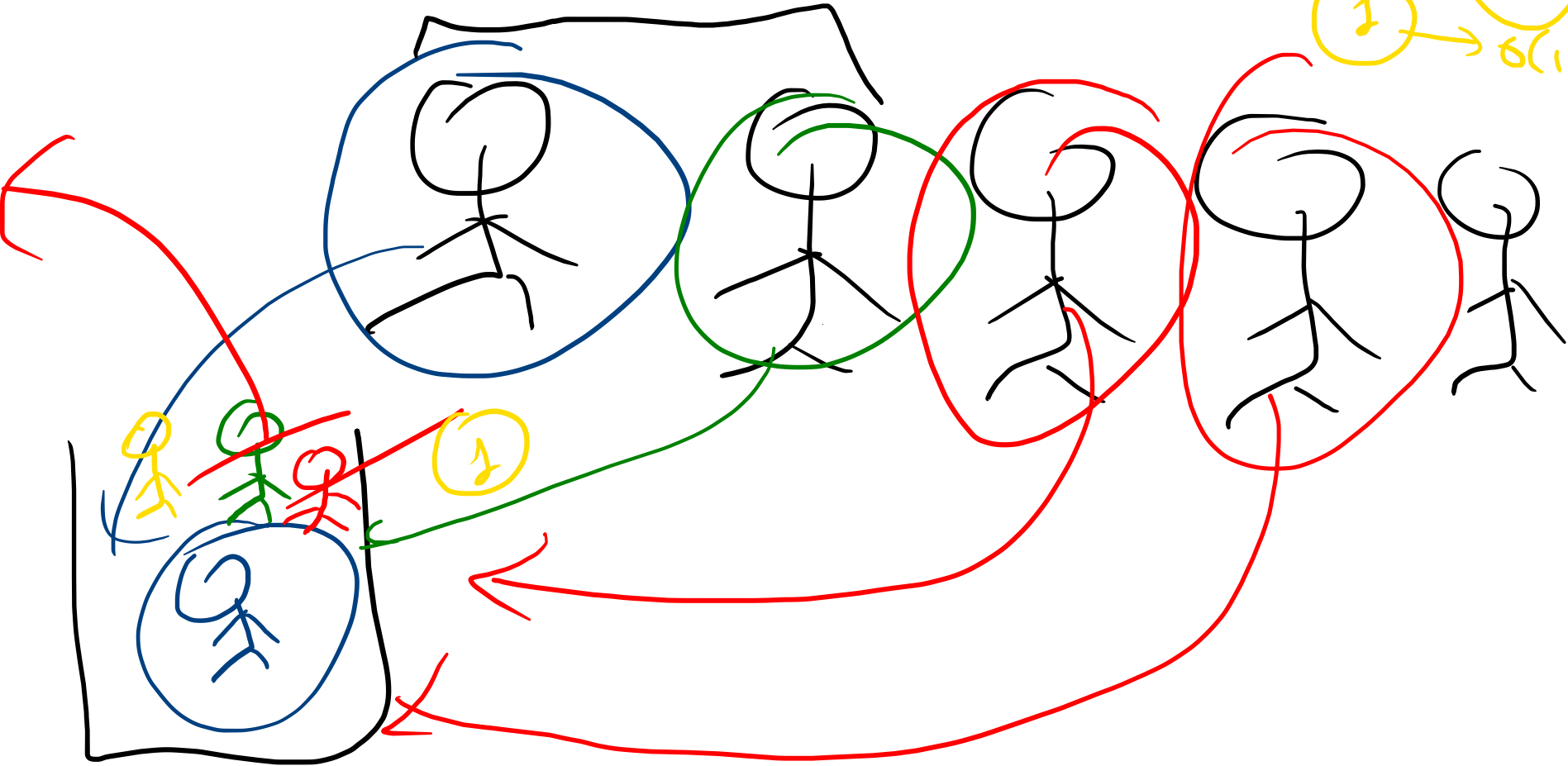
While (true)



Worst case

Dadebhaiya

while loop \Rightarrow $\textcircled{1}$
 $\textcircled{1} \rightarrow \delta(1)$



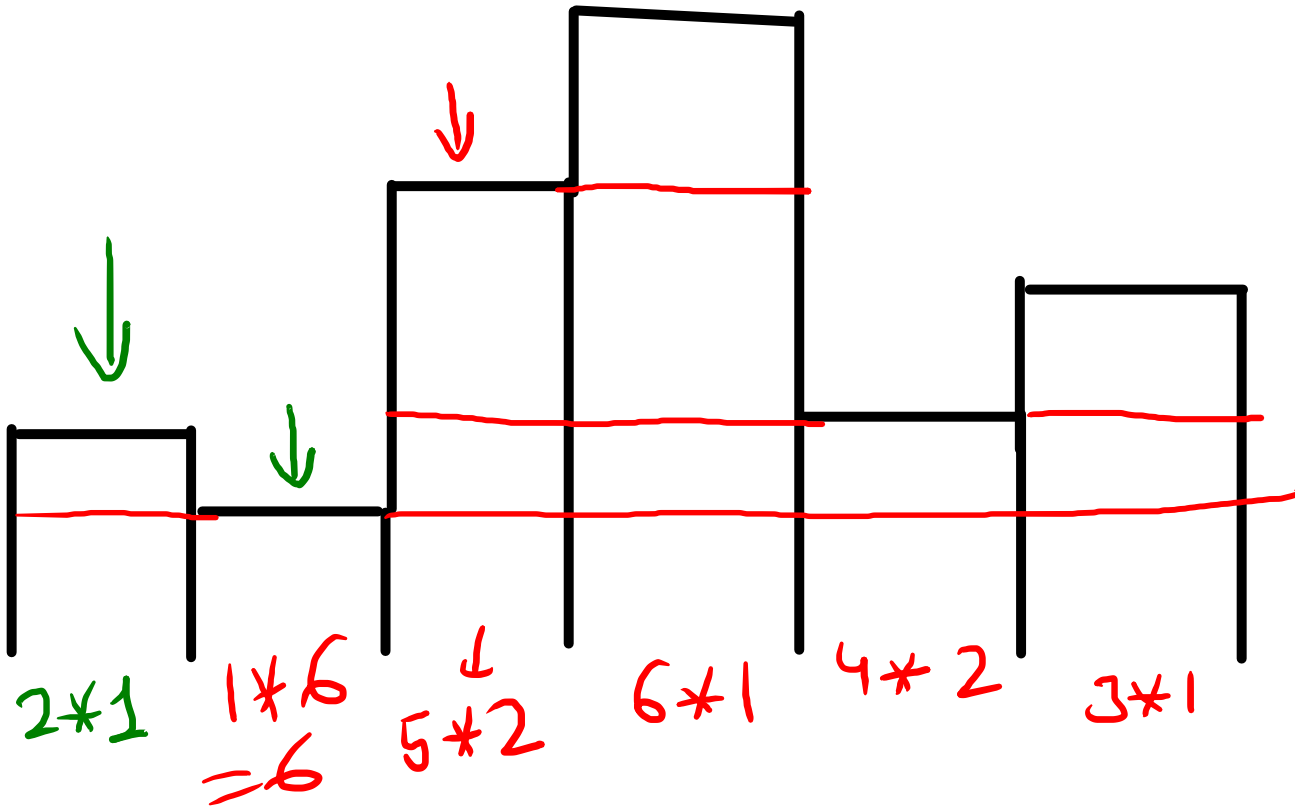
Maximum Rectangle area in histogram.

2 1 5 6 2 3



Maximum Rectangle area in histogram.

2 1 5 6 2 3



Maximum Rectangle area in histogram.

$(r-l) - 1$ ✓
 $(6+1) - 1 = 7 - 1 = 6$

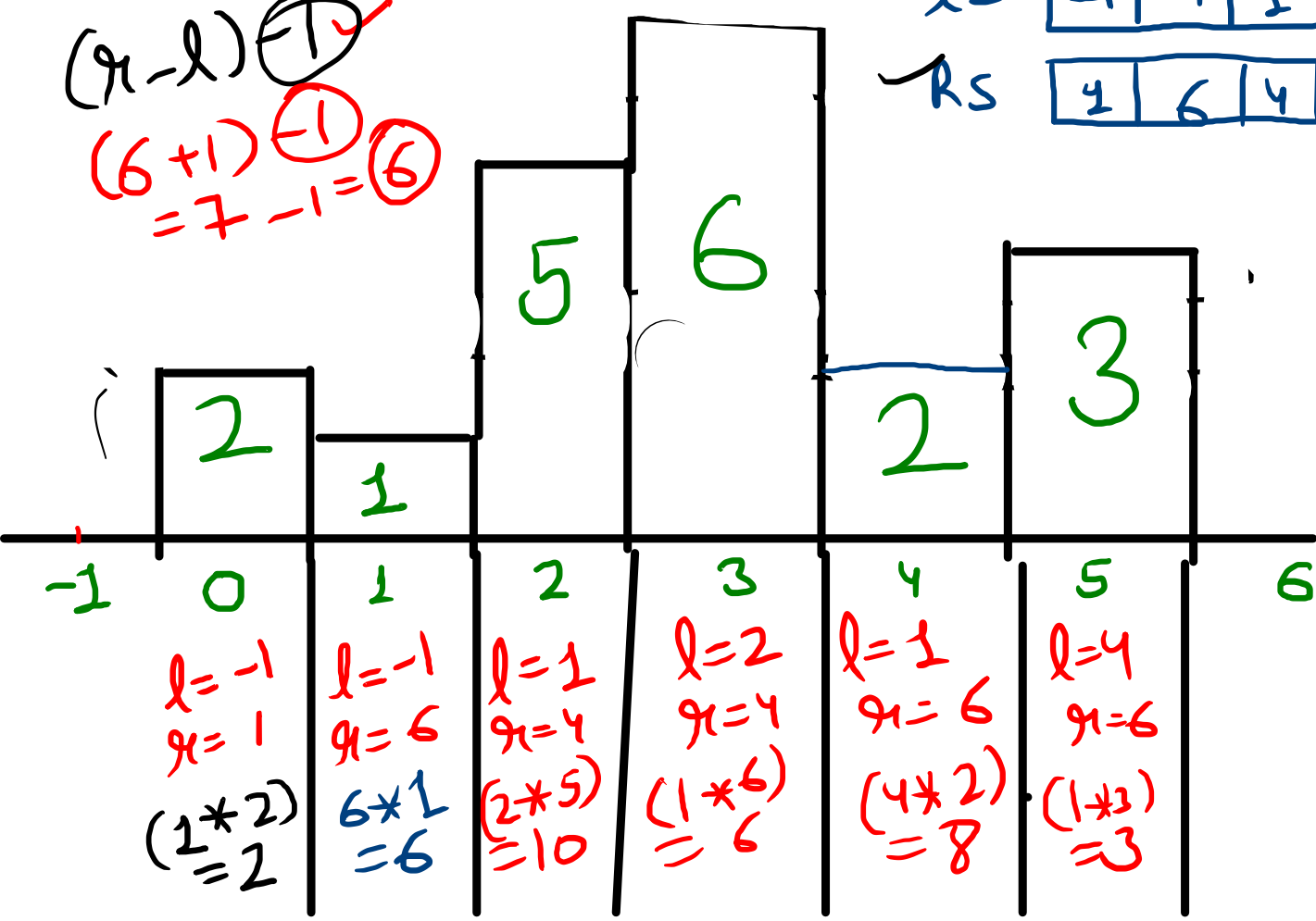
l_s

-1	-1	1	2	2	4
----	----	---	---	---	---

 r_s

1	6	4	4	6	6
---	---	---	---	---	---

indexes.



Left^{Next} Smallest
Right Next Smallest



Smallest to left.

Smallest \rightarrow chota
bhai

↓	↓	↓	↓	↓	↓
2	1	5	6	2	3
0	1	2	3	4	5

pos {

4
3
2
1
0

$\rightarrow O(n)$

qs =

-1	-1	1	2	1	4
----	----	---	---	---	---

as =

1	6	4	4	6	6
---	---	---	---	---	---

$\rightarrow O(n)$

	0	1	2	3	4	5
ls =	-1	-1	1	2	1	4

	0	1	2	3	4	5
rs =	1	6	4	4	6	6

	0	1	2	3	4	5
ans =	2	1	5	6	2	3

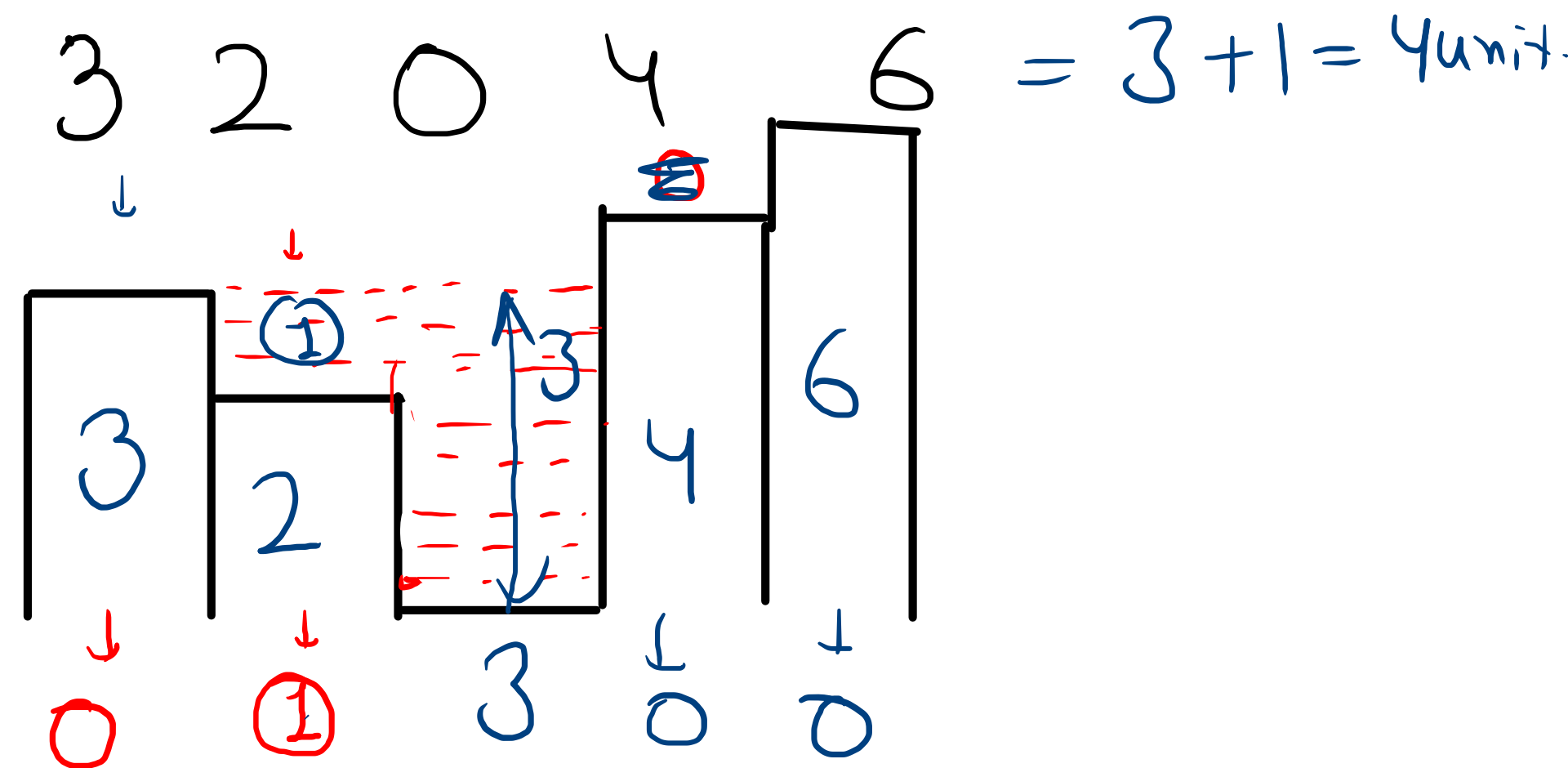
$\max (1 \times 2, 6 \times 1, 2 \times 5, 1 \times 6, 4 \times 2, 1 \times 3)$
 \Downarrow
 $\textcircled{10}$

$O(n)$

$$((rs[i] - ls[i] - 1) * ans[i]) \leftarrow \max$$

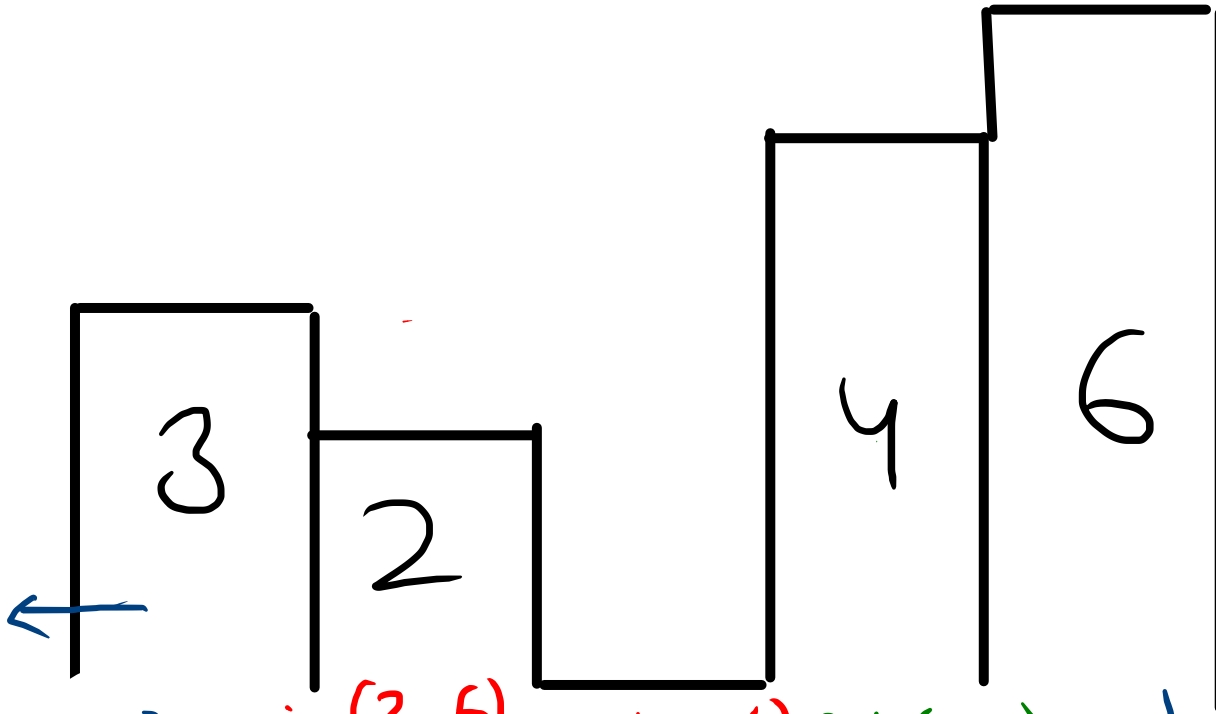
$$O(n) + O(n) + O(n) \\ = O(n)$$

Rain water problem.



3 2 0 4 6

$1+3=4$ unit



$$\begin{aligned} \min(\underline{3}, 6) \\ &= 3 - 3 \\ &= 0 \end{aligned}$$

$$\begin{aligned} \min(\underline{3}, 6) \\ &= 3 \\ (3-2) &= 1 \end{aligned}$$

$$\begin{aligned} \min(\underline{3}, 6) \\ &= 3 \\ (3-0) &= 3 \end{aligned}$$

$$\begin{aligned} \min(\underline{4}, 6) \\ &= 4 \\ (4-4) &= 0 \end{aligned}$$

$$\begin{aligned} \downarrow \min(\underline{6}, 6) \\ (6-6) &= 0 \end{aligned}$$

Max left
Max Right

3	2	0	4	6
3	3	3	4	6
6	6	6	6	6

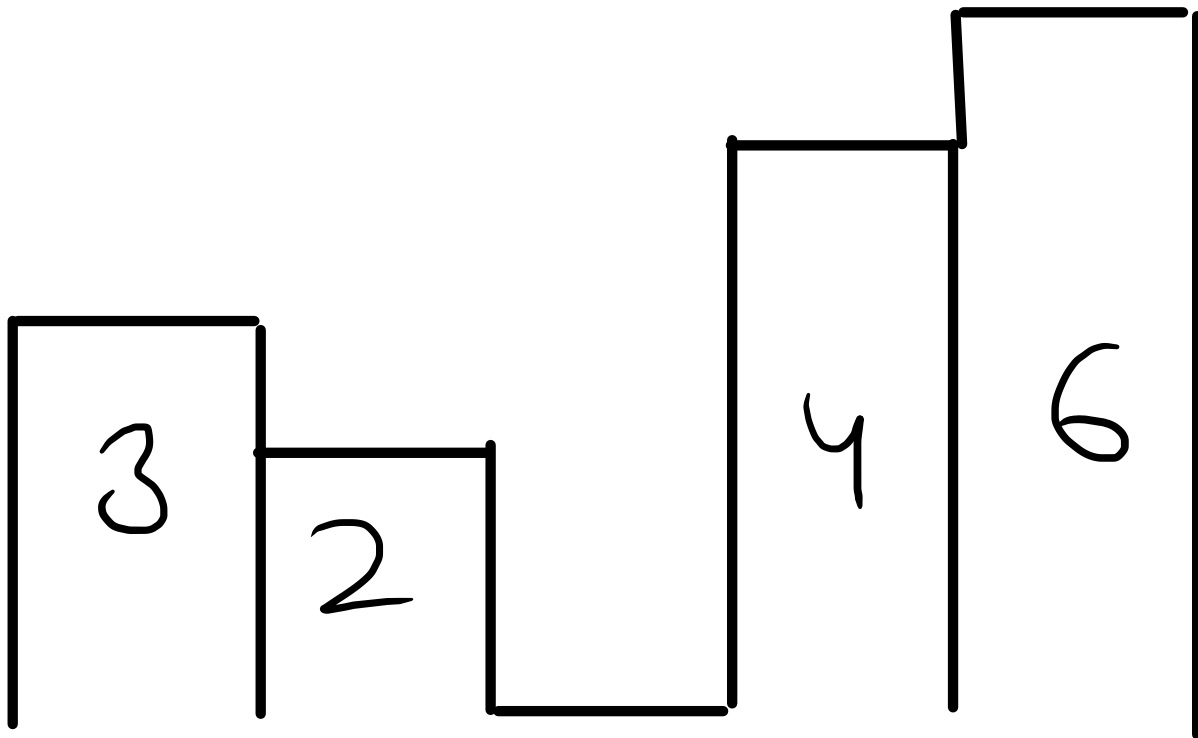
$$0 + 1 + 3 + 0 + 0 = \underline{4 \text{ unit}}$$

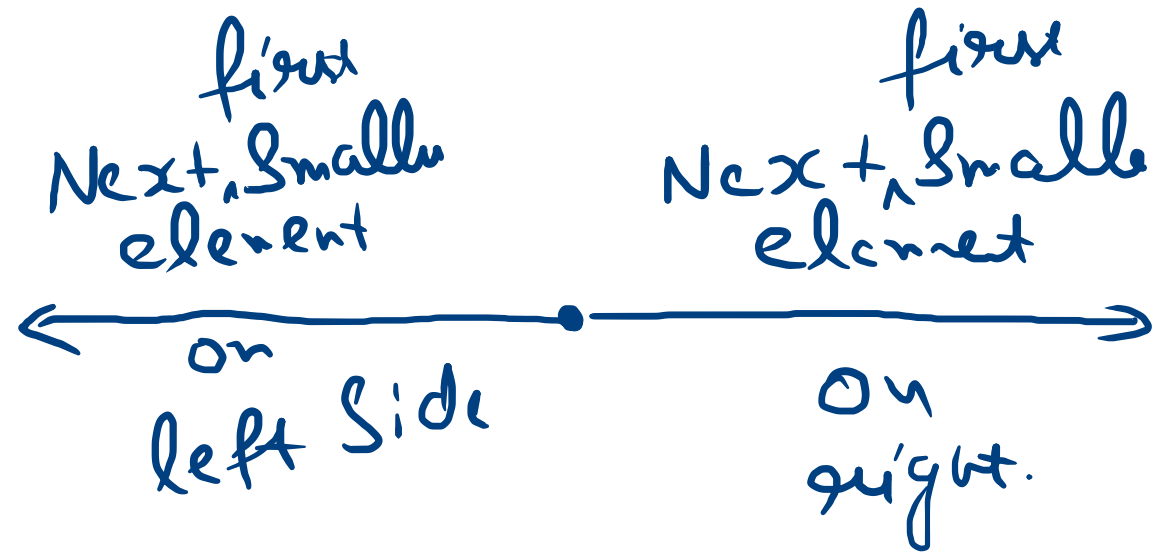
for (i=0; i<n; i++)
{

water += min(left[i], right[i]) - arr[i];

}

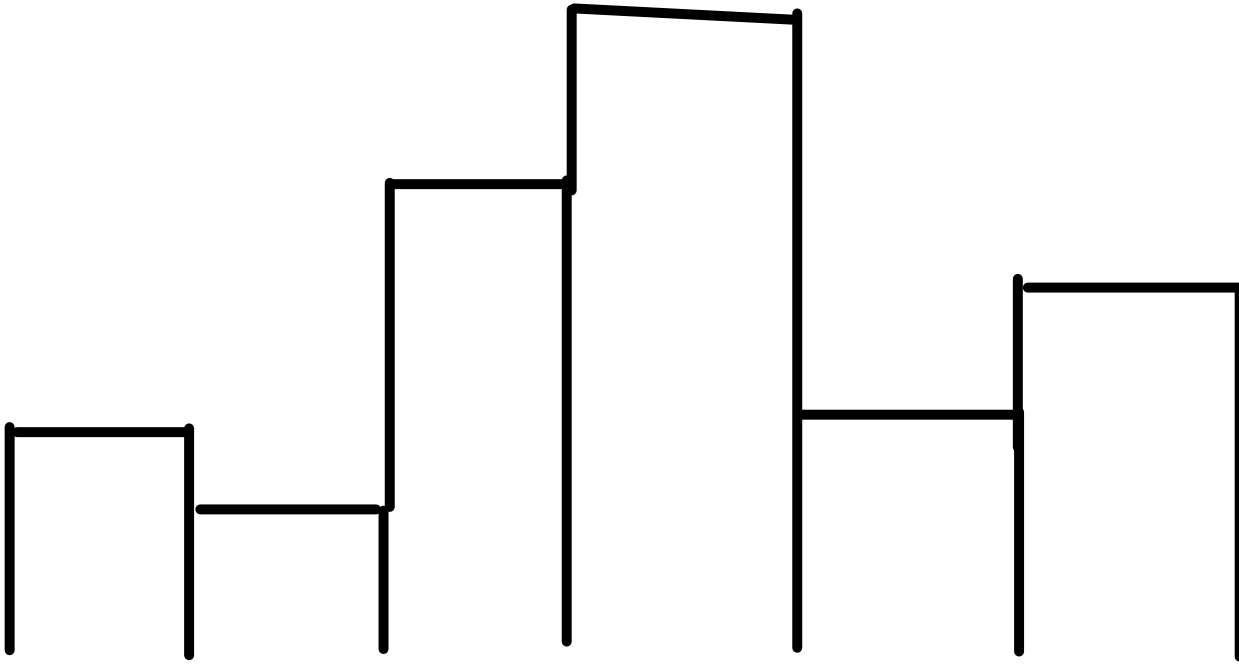
3 2 0 4 6



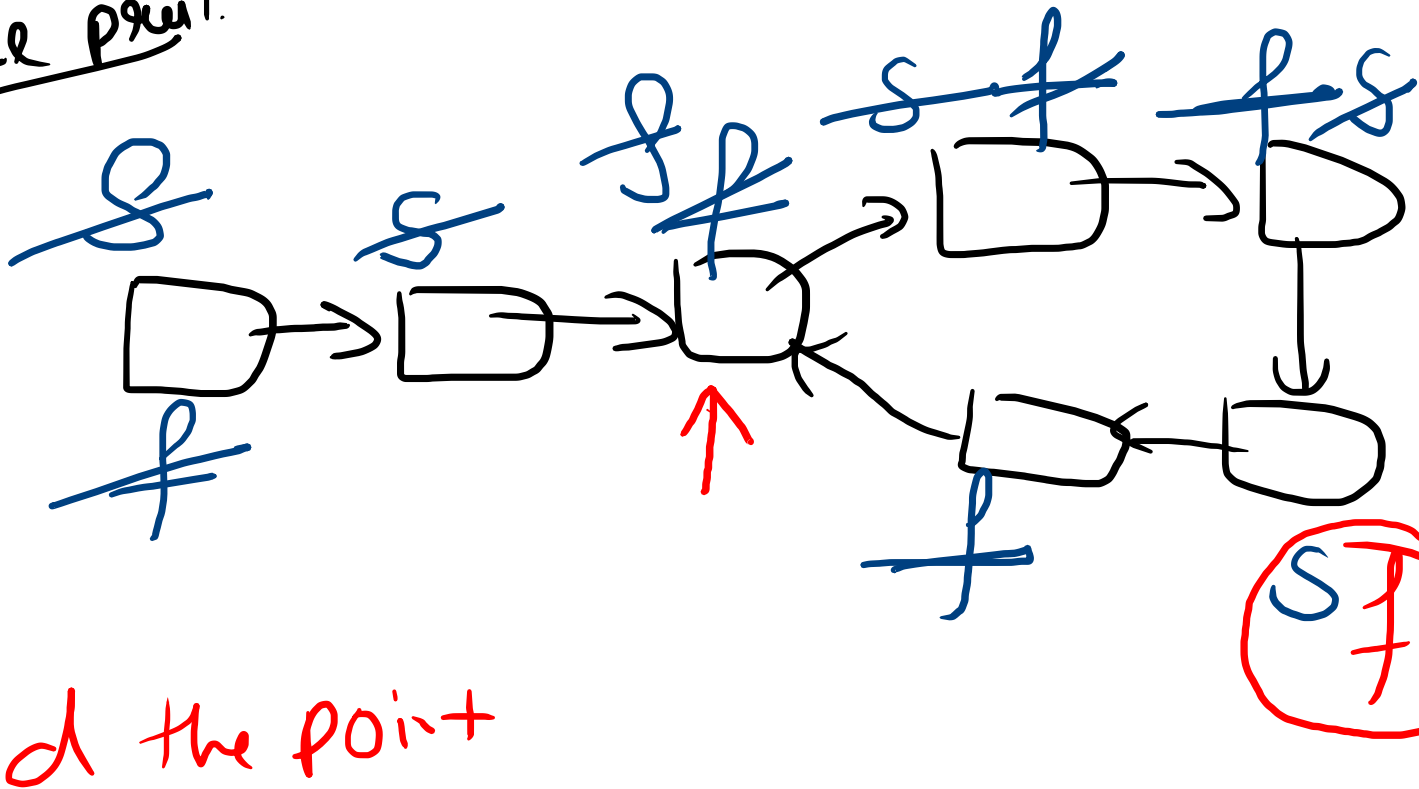


Maximum Rectangle area in histogram.

2 1 5 6 2 3



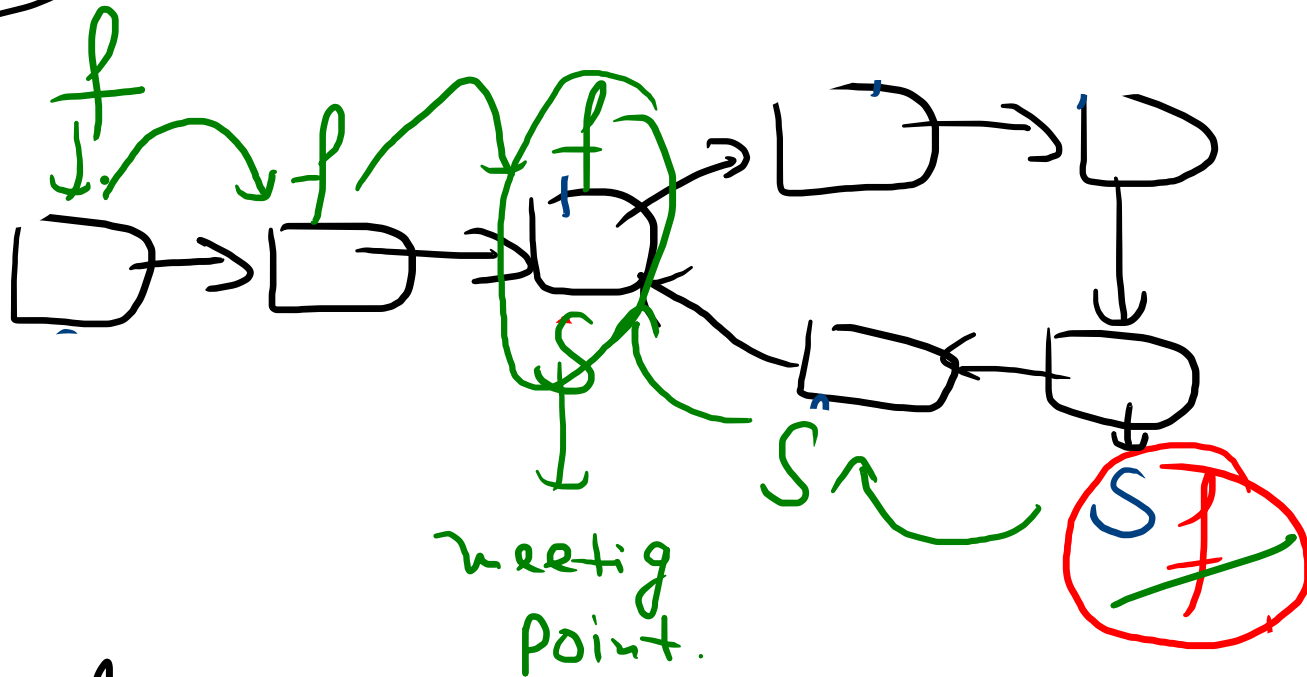
Cycle present.



find the point
where cycle is
starting.

Cycle
present.

Cycle print.



$slow = slow.next$

$fast = fast.next.next$

Constraint = 100 nodes

4-5

Insert
Delete.
cycle.
Reverse.
Merge

