

Assignment 2

→ Problem Statement :

Write a program to find Maximum and Minimum elements in an array using Divide & Conquer strategy and verify the time complexity.

→ Theory :

A] Divide and Conquer Strategy :

Divide and Conquer is an algorithm used to divide a large set of problem statements into smaller sub-problems and then combining the solution of these sub-problems.

1. Divide Instance of problem into two or more smaller instances.
2. Solve smaller instances recursively
3. Obtain Solution to original instance by combining these solutions.

B] Finding Min-Max from an Array :

1. Conventional Way :

- Initialize values of min & max of the 1st 2 elements respectively. Starting from 3rd location, compare each element with max & min, and change max & min accordingly.

```

a. Algorithm MaxMin (a, n, max, min) {
b.   max := min := a[1]
c.   for i := 2 to n do {
d.     if (a[i] > max) then max := a[i]
e.     if (a[i] < min) then min := a[i]
f.   }
g. }

```

- This conventional approach requires $2(n-1)$ number of comparisons, which can be reduced to $n-1$ in best case by converting if into an if-else statement.
- If the elements of the arrays are polynomials, strings, vectors, etc; then the cost of element comparison is much higher.

2. Divide & Conquer Method :

```

a. Algorithm MaxMin (i, j, max, min) {
b.   if (i == j) then max := min := a[i]
c.   else if (i == j - 1) then {
d.     if (a[i] < a[j]) max := a[j], min := a[i]
e.     else max := a[i], min := a[j]
f.   }
g.   else {
h.     mid := (i + j) / 2
i.     MaxMin (i, mid, max, min)
j.     MaxMin (mid + 1, j, max, min)

```

k.
 l.
 m. } if $\begin{cases} \text{max} < \text{max1} \\ \text{min} > \text{min1} \end{cases}$ then $\text{max} := \text{max1}$
 n. } if $\text{min} > \text{min1}$ then $\text{min} := \text{min1}$

Evaluation

$$T(n) = \begin{cases} T(\lfloor n/2 \rfloor) + T(\lceil n/2 \rceil) + 2 & ; n > 2 \\ 1 & ; n = 2 \\ 0 & ; n = 1 \end{cases}$$

$$\Rightarrow T\left(\frac{n}{2}\right) = 2T\left(\frac{n/2}{2}\right) + 2 = 2T\left(\frac{n}{4}\right) + 2$$

$$\Rightarrow T(n) = 2(2T\left(\frac{n}{4}\right) + 2) + 2$$

$$T(n) = 4T\left(\frac{n}{4}\right) + 4 + 2 \quad \text{--- (1)}$$

$$\Rightarrow T\left(\frac{n}{4}\right) = 2T\left(\frac{n/4}{2}\right) + 2$$

$$= 2T\left(\frac{n}{8}\right) + 2$$

$$\Rightarrow T(n) = 8T\left(\frac{n}{8}\right) + 8 + 4 + 2$$

$$\therefore T(k-1) = 2^{k-1} \cdot T\left(\frac{n}{2^{k-1}}\right) + \sum_{i=1}^{k-1} 2^i$$

Put $n = 2^k$, therefore

$$T(n) = \frac{2^k}{2} \cdot T\left(\frac{2n}{2^k}\right) + \sum_{i=1}^{k-1} 2^i$$

$$= \frac{n}{2} \cdot T(2) + 2^k - 2$$

$$= \frac{n}{2} (1) + n - 2$$

$$= \boxed{\frac{3n - 2}{2}}$$

Total number of comparisons = $\frac{3n}{2} - 2$

Hence, Divide & Conquer Strategy is 25% faster than the conventional approach.

Program:

```
#include <iostream>
using namespace std;

typedef struct {
    int min;
    int max;
} Pair;

//find maximum between 2 numbers
int max(int x, int y) {
    if(x > y) {
        return x;
    }
    return y;
}

//find minimum between 2 numbers
int min(int x, int y) {
    if(x > y) {
        return y;
    }
    return x;
}

Pair minmax(int arr[], int left, int right) {
    Pair minMax, leftArr, rightArr;

    //one element
    if(left == right) {
        minMax.min = arr[left];
        minMax.max = arr[left];

        for(int i = left; i <= right; i++) {
            cout<< arr[i]<< " ";
        }
        cout<< endl<< minMax.min<< ", "<< minMax.max<< endl<<
endl;

        return minMax;
    }

    //2 elements
    if(left + 1 == right) {
        minMax.min = min(arr[left], arr[right]);
        minMax.max = max(arr[left], arr[right]);
    }
}
```

```

        for(int i = left; i <= right; i++) {
            cout<< arr[i]<< " ";
        }
        cout<< endl<< minMax.min<< ", "<< minMax.max<< endl<<
endl;

        return minMax;
    }

    //more than 2 elements
    int mid = (left + right)/2;

    //recursive calls to find min and max
    leftArr = minmax(arr, left, mid);
    rightArr = minmax(arr, mid + 1, right);

    //find the min and max from the divided arrays
    minMax.min = min(leftArr.min, rightArr.min);
    minMax.max = max(leftArr.max, rightArr.max);

    for(int i = left; i <= right; i++) {
        cout<< arr[i]<< " ";
    }
    cout<< endl<< minMax.min<< ", "<< minMax.max<< endl<< endl;

    return minMax;
}

int main() {

    int arr[10];
    cout<< endl<< "Enter 10 numbers: ";
    //accepting 10 elements
    for(int i = 0; i < 10; i++) {
        cin>>arr[i];
    }
    cout<< endl<< "*****"<< endl<< "Working:
"<< endl<< endl;
    Pair minMax = minmax(arr, 0, 9);
    cout<< "*****"<< endl<< "Final Output:
"<< minMax.min<< ", "<< minMax.max<< endl<<
"*****"<< endl<< endl;
    return 0;
}

```

Output:

```
Someshwars-MacBook-Pro:GroupB someshwargaikwad$ g++  
Assignment2.cpp  
Someshwars-MacBook-Pro:GroupB someshwargaikwad$ ./a.out
```

```
Enter 10 numbers: 32 15 46 78 12 90 21 98 23 55
```

```
*****
```

```
Working:
```

```
32 15  
15, 32
```

```
46  
46, 46
```

```
32 15 46  
15, 46
```

```
78 12  
12, 78
```

```
32 15 46 78 12  
12, 78
```

```
90 21  
21, 90
```

```
98  
98, 98
```

```
90 21 98  
21, 98
```

```
23 55  
23, 55
```

```
90 21 98 23 55  
21, 98
```

```
32 15 46 78 12 90 21 98 23 55  
12, 98
```

```
*****
```

```
Final Output: 12, 98
```

```
*****
```

```
Someshwars-MacBook-Pro:GroupB someshwargaikwad$
```


→ Conclusion :

Topics Covered :

1. Divide & Conquer
2. Divide & Conquer strategy to find minimum & maximum numbers in an array.