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Theory of Computations

Assignment #4

Question #1:

Let $a^n b^{3n} \in L$, where $n \geq 0$. Let $S(n)$ be a string such that $S(n) = a^n b^{3n}$

Choosing $n = 5$, then: $S(5) = a^5 b^{15} = \text{aaaaabbbbbbbbbbbbbb}$

By choosing $x = a$, $y = \text{aaaa}$, $z = \text{bbbbbbbbbbbbbb}$, then we have $S(5) = xyz$

Now suppose that $xy^i z$ is regular, then $xy^{i+1} z$ must be regular.

So if xyz is regular, then $xy^2 z$ must also be regular.

$$\therefore xy^2 z = \text{aaaaaaaaabbbbbbbbbbbbbb} = a^9 b^{15}$$

Since $15 < 3 * 9$, then there is a contradiction.

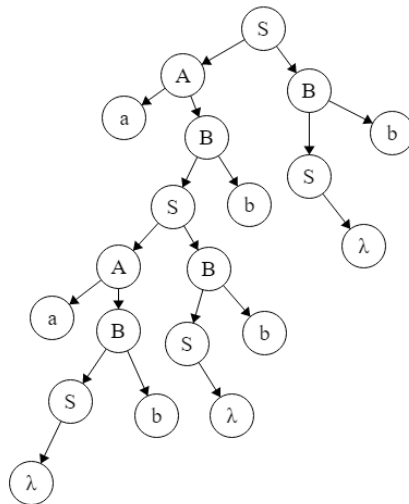
$\therefore L$ is not a regular language.

Question #2:

I. Rightmost derivation:

$$S \Rightarrow AB \Rightarrow ASb \Rightarrow Ab \Rightarrow aBb \Rightarrow aSbb \Rightarrow aABbb \Rightarrow aASbbb \Rightarrow aAbbbb \\ \Rightarrow aaBbbbb \Rightarrow aaSbbbbb \Rightarrow aabbbbb$$

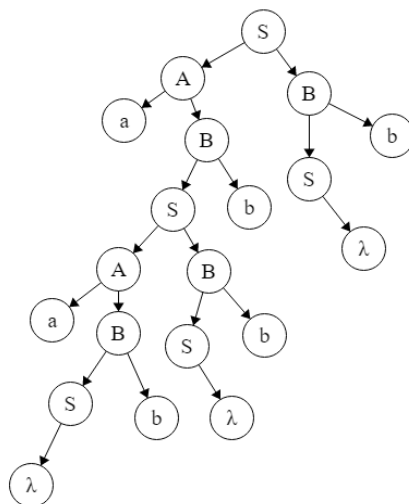
Parse tree:



II. Leftmost derivation:

$$S \Rightarrow AB \Rightarrow aBB \Rightarrow aSbB \Rightarrow aABbB \Rightarrow aaBBbB \Rightarrow aaSbBbB \Rightarrow aabBbbB \\ \Rightarrow aabSbbbB \Rightarrow aabbbbSb \Rightarrow aabbbbb$$

Parse tree:



Question #3:

a) $S \Rightarrow aaB \Rightarrow aaAa \Rightarrow aabBba \Rightarrow aabbBbba \Rightarrow aabbAabba \Rightarrow aabbabba$

b) The string is not accepted since the grammar forces us to begin with “aa”

Question #4:

a)

$$S \rightarrow aA$$

$$A \rightarrow aBC$$

$$B \rightarrow b \mid aBb$$

b)

$$S \rightarrow aA$$

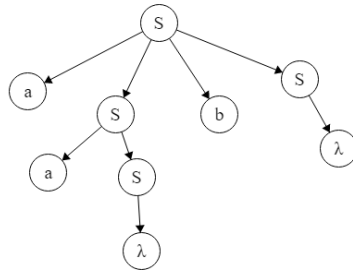
$$A \rightarrow bbb \mid aAB$$

$$B \rightarrow bb$$

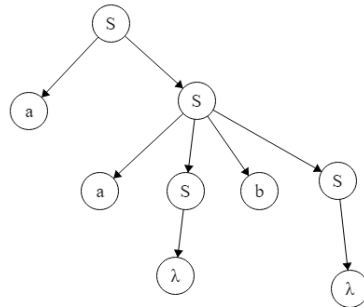
Question #5:

a)

One way to draw a parse tree for the string “aab”



Another way to draw a parse tree for the string “aab”



b)

$$S \Rightarrow aSbS \Rightarrow aaSbS \Rightarrow aabS \Rightarrow aab$$

$$S \Rightarrow aS \Rightarrow aaSbS \Rightarrow aabS \Rightarrow aab$$

Since there were two separate derivations and parse trees for the same string, then the language is ambiguous.

Question #6:

a)

$$S \rightarrow ASBb \mid SBb \mid aA \mid a$$

$$A \rightarrow aAS \mid aS \mid a$$

$$B \rightarrow SbS \mid A \mid bb \mid CD$$

$$C \rightarrow AC \mid C \mid CC$$

$$D \rightarrow DA \mid D \mid aA$$

b)

$$S \rightarrow ASBb \mid SBb \mid aA \mid a$$

$$A \rightarrow aAS \mid aS \mid a$$

$$B \rightarrow SbS \mid a \mid bb \mid CD$$

$$C \rightarrow AC \mid C \mid CC$$

$$D \rightarrow DA \mid a \mid b$$

c)

Part I: Find all Non-Terminal symbols that derive the elements of Σ and put them in a set. The process is done repeatedly until the last two sets are equal.

$$S_1 = \{S, A, B, D\}, S_2 = \{S, A, B, D\}$$

Since S_1 and S_2 are equal, then part I is complete.

The new grammar:

$$S \rightarrow ASBb \mid SBb \mid aA \mid a$$

$$A \rightarrow aAS \mid aS \mid a$$

$$B \rightarrow SbS \mid a \mid bb$$

$$D \rightarrow DA \mid a \mid b$$

Part II: Have the start symbol in a set, and append all the non-Terminal symbols that are derived from it until a dead-end is reached.

$$S1 = \{S\}, S2 = \{S, A, B\}, S3 = \{S, A, B\}$$

Since S2 and S3 are equal, then part II is complete.

The new grammar:

$$S \rightarrow ASBb \mid SBb \mid aA \mid a$$

$$A \rightarrow aAS \mid aS \mid a$$

$$B \rightarrow SbS \mid a \mid bb$$

d)

$$S \rightarrow S_1B_2 \mid SB_2 \mid A_1A \mid a$$

$$A \rightarrow S_1S \mid A_1S \mid a$$

$$B \rightarrow SS_2 \mid B_1B_1 \mid a$$

$$B_1 \rightarrow b$$

$$B_2 \rightarrow B B_1$$

$$A_1 \rightarrow a$$

$$S_1 \rightarrow AS$$

$$S_2 \rightarrow B_1S$$

Question #7:

Let $a^n b^n c^{n+1} \in L$, where $n \geq 0$. Let $S(n) = a^n b^n c^{n+1}$

By choosing $n = 4$, then $S(4) = aaaaabbbbccccc$

By choosing $u = aa$, $v = aa$, $x = bbbbc$, $y = c$, $z = ccc$

Then $uvxyz = S(4)$

Since $uv^i xy^i z$ is context free, then $uv^{i+1} xy^{i+1} z$ must also be context free.

\therefore By applying the property on $uvxyz$, $uv^2 xy^2 z$ must also be context free

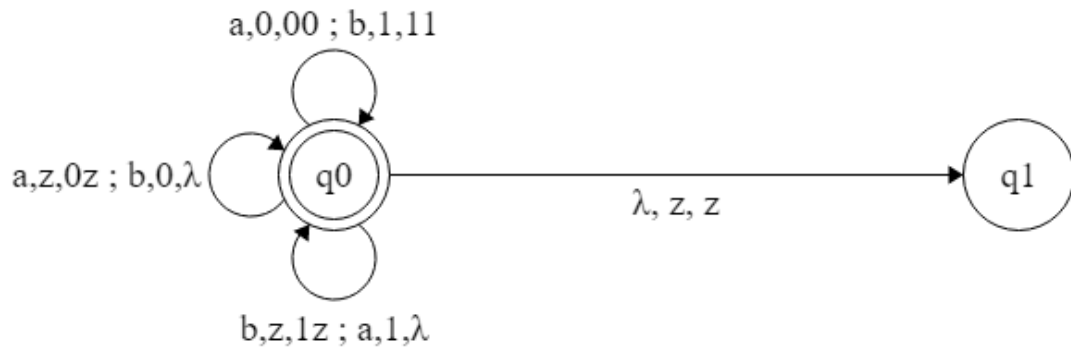
$\therefore uv^2 xy^2 z = aaaaaabbbbccccc = a^6 b^4 c^6$

Since the language L was defined as $L = \{a^n b^n c^m \mid n \geq 0, m > n\}$

$|a|$ is greater than $|b|$ and $|a|$ is equal to $|c|$, which contradicts the definition.

$\therefore L$ is not context free

Question 8:



Initial State is q0

$$\delta(q0, a, z) = (q0, 0z),$$

$$\delta(q0, b, z) = (q0, 1z),$$

$$\delta(q0, b, 0) = (q0, \lambda)$$

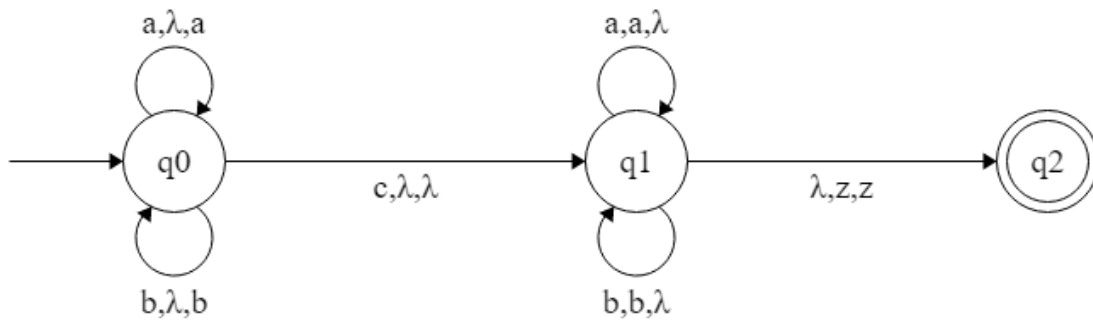
$$\delta(q0, a, 1) = (q0, \lambda),$$

$$\delta(q0, a, 0) = (q0, 00),$$

$$\delta(q0, b, 1) = (q0, 11)$$

$$\delta(q0, \lambda, z) = (q1, z)$$

Question 9:



$$\delta(q_0, a, \lambda) = (q_0, a),$$

$$\delta(q_0, b, \lambda) = (q_0, b),$$

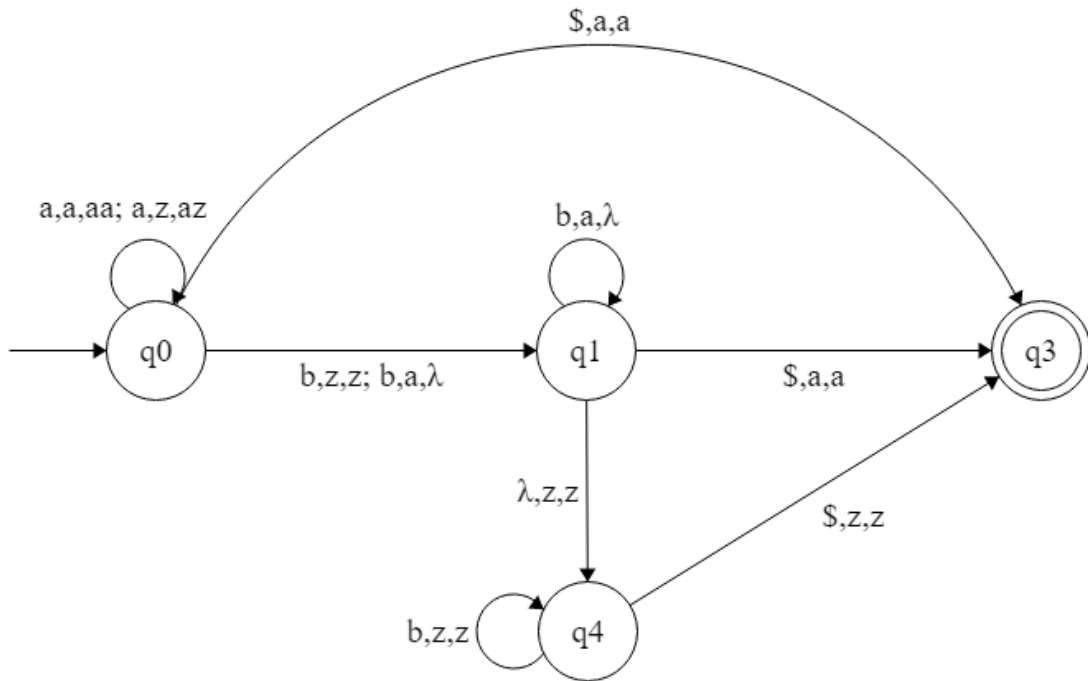
$$\delta(q_0, c, \lambda) = (q_1, \lambda)$$

$$\delta(q_1, a, a) = (q_1, \lambda),$$

$$\delta(q_1, b, b) = (q_1, \lambda),$$

$$\delta(q_1, \lambda, z) = (q_2, z)$$

Question 10:



$$\delta(q0, a, a) = (q0, aa),$$

$$\delta(q0, a, z) = (q0, az),$$

$$\delta(q0, \$, a) = (q3, a)$$

$$\delta(q0, b, z) = (q1, z),$$

$$\delta(q0, b, a) = (q1, \lambda),$$

$$\delta(q1, b, a) = (q1, \lambda)$$

$$\delta(q1, \lambda, z) = (q4, z),$$

$$\delta(q1, \$, a) = (q3, a),$$

$$\delta(q4, b, z) = (q4, z)$$

$$\delta(q4, \$, z) = (q3, z)$$