## Computer Science Department, Kuwait University

Date: 21/4/2022 **Due: 7/5/2022** 

Course: 0418211/01 Homework 4

- Q1. Prove using pumping theorem for regular languages that  $L = \{ a^n b^m , m \ge 3n \}$  is not regular.
- Q2. Find a <u>rightmost</u> derivation and a <u>leftmost</u> derivation for the string *aabbbb* using the grammar

$$S \rightarrow AB | \lambda$$

$$A \rightarrow aB$$

$$B \rightarrow Sb$$

Draw the parsing tree of each derivation.

Q3. Consider the grammar with productions

$$S \rightarrow aaB$$

$$A \rightarrow bBb|\lambda$$

$$B \rightarrow Aa \mid bBb$$

For each of the following strings decide whether or the string is in the language generated by this grammar. Justify your answer.

- a) aabbabba
- b) abbababa
- **Q4.** Find an *s-grammar* for each of the following languages:

a- 
$$L=\{a^nb^n : n \ge 2\}$$

b- L=
$$\{a^nb^m : n \ge 1, m = 2n+1\}$$

Q5. Show that the following grammar G is ambiguous

$$S \rightarrow aSbS|aS|\lambda$$

(Consider the string aab. Show that it has two:

- a) Parse trees.
- b) Leftmost derivations.
- **Q6**. (*Grammar Simplification*) Begin with the grammar:

$$S \rightarrow ASBb \mid aA$$

$$A \rightarrow aAS \mid a \mid \lambda$$

$$B \rightarrow SbS \mid A \mid bb \mid CD$$

$$C \rightarrow AC \mid CC$$

$$D \rightarrow DA \mid a \mid b$$

- a) Remove  $\lambda$ -productions.
- b) Remove any unit-productions in the resulting grammar.
- c) Remove any useless symbols in the resulting grammar.
- d) Put the resulting grammar into Chomsky Normal Form.
- e) Put the resulting grammar into Greibach Normal Form.
- Q7. Prove using pumping theorem for CFL that  $L = \{ a^n b^n c^m, n \ge 0, m > n \}$  is not context free.
- **Q8.** Design an NPDA for the language L={  $w \in \{a,b\}^+$ ,  $n_a(w) \neq n_b(w)$  }.

For example: aabb is not accepted. aabab is accepted.

- **Q9.** Design a DPDA for the language  $L=\{ w c w^R, w \in \{a,b\}^+ \}$ .
- **Q10.** Design a *DPDA* that accepts the following language and trace its configuration changes given the input *aaabb*\$: (\$ is a marker for the end of input)

$$L = \{ a^n b^m \$, n \neq m \}$$