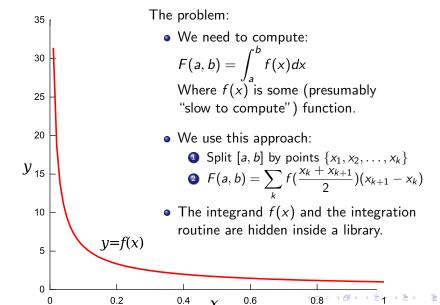
MPI profiling with Allinea MAP

Alexander Gaenko

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Problem: Calculation of a definite integral.



Sequential program code

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```
#include <stdio.h>
     #include <math.h>
3
     // Declare integrand() and integral() from "mymath" library
     #include "mylib/mymath.hpp"
     int main(int argc, char** argv)
9
       unsigned long int n;
       if (argc!=2 || sscanf(argv[1],"%lu",&n)!=1) {
10
11
         fprintf(stderr, "Usage:\n%s integration_steps\n\n\n", argv[0]);
         return 1;
13
      // Integration limits.
      const double global a=1E-5:
       const double global b=1:
19
      // Perform integration
       const double v=integral(integrand, n, global a, global b);
       const double y_exact=4*(pow(global_b,0.25)-pow(global_a,0.25));
       printf("Result=%lf Exact=%lf Difference=%lf\n", y, y_exact, y-y_exact);
       return 0:
```

Problem size behavior: how to measure.

Do we even need to parallelize?

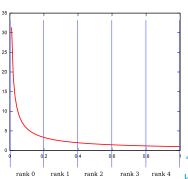
```
1 | # Compilation:
   $ gcc -03 -o integral_seq.x integral_seq.cxx\
         -L./mylib -lmymath
4 | # Timed runs:
   $ time -p ./integral_seq.x 1000000
   Result=3.775045 Exact=3.775063 Difference=-0.000019
   real 2.12
   $ time -p ./integral_seq.x 2000000
   Result=3.775058 Exact=3.775063 Difference=-0.000005
   real 4.28
10
   $ time -p ./integral_seq.x 8000000
11
   Result=3.775062 Exact=3.775063 Difference=-0.000001
12
   real 17.61
13
```

The time grows linearly with the problems size. Acceptable accuracy at 8M points. Can we speed it up?

Parallellization: domain decomposition.

Approach:

- Split [a, b] into several domains;
- Compute integrals independently.



$$F(a,b) = \int_a^b f(x) dx$$

- Assign a process to each domain $[x_k, x_{k+1}]$
- 2 Let each process compute $F(x_k, x_{k+1})$

3
$$F(a,b) = \sum_{k} F(x_k, x_{k+1})$$

"Embarassingly parallel" problem, high speedup is expected.

A sketch of the parallel code

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```
MPI_Init(&argc, &argv);
int rank, nprocs;
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &nprocs);
// ...Get total number of steps, broadcast it...
// Each rank figures out its integration limits and number of steps
unsigned long my_stepbase, my_nsteps;
get_steps(nsteps_all, nprocs, rank, &my_stepbase, &my_nsteps);
const double per step=(global b-global a)/nsteps all:
const double x1=global_a + my_stepbase*per_step;
const double x2=x1 + mv nsteps*per step;
// Compute my own part of the integral
double my_y=integral(integrand, my_nsteps, x1, x2);
// Sum all numbers on master
double y=0;
MPI Reduce(&mv v. &v. 1. MPI DOUBLE, MPI SUM, O. MPI COMM WORLD):
// ... print results ...
MPI_Barrier(MPI_COMM_WORLD);
// Here we could start another computation.
MPI Finalize():
```

Parallel performance: how to measure.

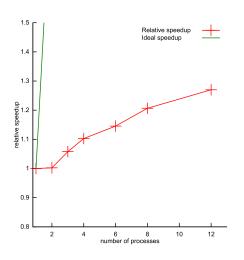
Now let's see how much we achieved...

- Strong scaling: as we add processes, how do we fare?
- Weak scaling: as we add both processes and work?

```
$ mpicc -03 -o integral_par.x integral_par.cxx \
           -L./mylib -lmymath
2
   $ time -p mpirun -np 1 ./integral_par.x 8000000
   Result=3.775062 Exact=3.775063 Difference=-0.000001
   real 17.23
   user 17.08
   sys 0.02
   $ time -p mpirun -np 2 ./integral_par.x 8000000
   Result=3.775062 Exact=3.775063 Difference=-0.000001
   real 17.24
10
   user 31.98
11
   sys 0.05
12
```

Parallel performance: results

Is there a performance problem?



- Relative speedup: $s(p) = \frac{\text{(time with 1 process)}}{\text{(time with } p \text{ processes)}}$
- Ideal relative speedup: $s_{\text{ideal}}(p) = p$
- Our speedup is
 25% on 12 nodes!
- I'd call it "dismal". We **do** have a problem!
- Why? How to figure it out?

How does performance analysis work?

How to collect data?

- Instrumentation:
 - Insert timers & counters in the code
 - Requires source or binary processing
- Sampling:
 - Interrupt & check the program at regular intervals
 - Introduces statistical error

What kind of data?

- Profile:
 - Summary information only
 - Relatively small file
- Trace:
 - Detailed recording during the run
 - Potentially huge file
 - Profile can be restored

Allinea MAP does tracing by sampling.

Prepare for profiling

2

To prepare for profiling/tracing, one needs to:

- Compile with full optimization
- Generate debugging symbols
- Link with system libs dynamically
 - Usually the default
 - Notable exception: Cray
- On Flux: load ddt module

Running Map: simple way (demo)

- Get interactive access to a compute node
- Change to your working directory
- Optionally, set sampling interval
- Run as you would, prefixed by map

Caution:

- Too small interval: large overhead!
- Too large interval: not enough samples!
- Allinea recommends at least 1000 samples/process



Running Map: other options

What if you can not or would not run a GUI?

- Have slow or non-existing X connection to compute nodes.
- Do not want to wait for interactive session.

Use -profile option.

```
#PBS -V
#PBS -q flux -l qos=flux -A account_flux
#PBS -l nproc=12,walltime=10:0:0

cd $PBS_0_WORKDIR
export ALLINEA_SAMPLER_INTERVAL=5
map -profile mpirun -np 12 ./integral_par.x 8000000
```

This will create a *.map file. Then run from the login node:

```
1 $ map integral_par_even_12p_*.map
```

If you are submitting to a Flux queue...

- Run map from the login node:
- \$ map ./integral_par.x 8000000
- Set number of processes
- Check Submit to queue
- Olick Configure...
- Solution
 Load a proper submission
 template file (see next page)
- Click OK
- Click Run



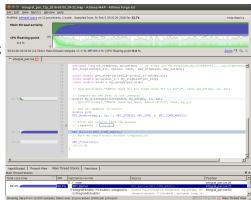
Submission template for Flux

```
#PBS -V
#PBS -1 walltime=WALL_CLOCK_LIMIT_TAG
#PBS -1 nodes=NUM_NODES_TAG:ppn=PROCS_PER_NODE_TAG
#PBS -q QUEUE_TAG -1 qos=flux -A account_flux
#PBS -o PROGRAM_TAG-allinea.stdout
#PBS -e PROGRAM_TAG-allinea.stderr

cd $PBS_O_WORKDIR
AUTO_LAUNCH_TAG
```

Time for a live demo!

- Most of the time is spent in MPI;
- As run progresses, *even more* time is spent in MPI;
- Problem: some processes spend more time computing f(x) then others!
- It's called "Load Imbalance".



Possible solutions:

- Distribute work unevenly (but how?)
- Implement Dynamic Load Balancing.



Dynamic load balancing

Idea: If a process has nothing to do, make it to do something.

Manager-Workers approach:

Manager

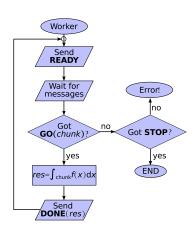
- Listen to all workers
- Worker sent READY ?
 - send GO with a job chunk
- Worker sent DONE?
 - add result to the sum
- No more job chunks?
 - send STOP to the worker
- No more workers?
 - we are done!
 - Otherwise, go to (1)

Worker

- Send READY to the Manager
- 2 Listen to the Manager
- Manager sent GO ?
 - Get job chunk
 - Do the calculation
 - Send DONE with result to the Manager
 - Go to (1)
- Manager sent STOP?
 - exit.

Dynamic load balancing





Dynamic Load Balancing: Worker code

```
1
     void do worker(int manager, int rank)
2
       for (;;) {
4
        // Inform manager that i am ready:
         MPI_Send(NULL,0, MPI_DOUBLE, manager, TAG_READY, MPI_COMM_WORLD);
6
         // ...and wait for the task.
7
         double range data[3]: // expect x1, x2 and the number of steps (as double!)
         MPI Status stat:
         MPI_Recv(range_data,3,MPI_DOUBLE, manager, MPI_ANY_TAG, MPI_COMM_WORLD, &stat);
10
         switch (stat.MPI TAG) {
11
         case TAG GO:
12
           break; // do normal work
13
         case TAG_STOP:
14
           return:
15
         default:
16
           fprintf(stderr, "Rank %d: Got unexpected tag=%d, aborting. \n", rank, stat.MPI_TAG);
17
           MPI Abort (MPI COMM WORLD, 2):
18
19
         // we are here on TAG GO.
         const double x1=range data[0]:
20
21
         const double x2=range data[1]:
22
         const unsigned long my_nsteps=range_data[2];
23
         // Compute my own part of the integral
24
         double mv_v=integral(integrand, my_nsteps, x1, x2);
25
26
         // Send the result to manager
27
         MPI Send(&mv v.1.MPI DOUBLE, manager, TAG DONE, MPI COMM WORLD):
28
29
```

Dynamic Load Balancing: Manager code

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```
double do_manager(const double global_a, const double global_b,
                 const unsigned long natepa all, const unsigned long points per block.
                 const int nprocs, const int rank)
 const double per step=(global b-global a)/nsteps all:
  int nworkers_left=nprocs-1;
 unsigned long ipoint=0: // next point to be processed
 double v=0:
 for (;;) {
   // Get a tagged message and possibly a result from any worker
    double v worker=0:
   MPI_Status stat;
   MPI_Recv(&y_worker,1,MPI_DOUBLE, MPI_ANY_SOURCE,MPI_ANY_TAG, MPI_COMM_WORLD, &stat);
    const int rank worker=stat.MPI SOURCE:
    switch (stat.MPI TAG) {
    case TAG_READY:
     // Do we have any work for this worker?
     if (ipoint>=nsteps all) {
       // if not, stop the worker
       MPI_Send(NULL,0,MPI_DOUBLE, rank_worker, TAG_STOP, MPI_COMM_WORLD);
        --nworkers left:
       break;
```

Dynamic Load Balancing: Manager code (cont.)

```
1
           { // Prepare chunk of work for the worker
 2
             const unsigned long ns_worker=
                (ipoint+points_per_block > nsteps_all)? (nsteps_all - ipoint) : points_per_block;
             const double x1=global_a+ipoint*per_step;
             const double x2=x1+ns worker*per step:
6
             double range_data[3];
7
             range_data[0]=x1;
8
             range data[1]=x2:
g
             range data[2]=ns worker:
10
             // Send the chunk
11
             MPI Send(range data.3.MPI DOUBLE, rank worker, TAG GO, MPI COMM WORLD):
12
             // adjust the amount of points
13
             ipoint += ns_worker;
14
15
           break:
16
         case TAG_DONE:
           y += y_worker;
17
18
           break:
19
         default:
20
           fprintf(stderr, "Rank %d (manager): Got unexpected tag=%d from %d,"
21
                          " aborting.\n".rank.rank worker.stat.MPI TAG):
22
           MPI Abort (MPI COMM WORLD.1):
23
24
         if (nworkers left<=0) { // Anv active workers left?
25
           return y;
26
27
28
```

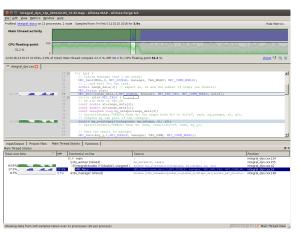
Dynamic Load Balancing: main

```
1
     int main(int argc, char** argv)
2
3
       MPI_Init(&argc, &argv);
4
5
       int rank, nprocs;
       MPI_Comm_rank(MPI_COMM_WORLD, &rank);
7
       MPI_Comm_size(MPI_COMM_WORLD, &nprocs);
8
9
       // Get command line arguments, broadcast
10
       unsigned long int nsteps all, points per block:
11
12
13
      // Global integration limits.
14
       const double global a=1E-5:
15
       const double global_b=1;
16
17
       // Split into workers and manager:
18
       if (rank==0) { // Run as the manager and get the result:
19
         double y=do_manager(global_a,global_b,nsteps_all,points_per_block,nprocs,rank);
20
21
         const double y_exact=4*(pow(global_b,0.25)-pow(global_a,0.25));
22
         printf("Result=%lf Exact=%lf Difference=%lf\n", y, y_exact, y-y_exact);
23
       } else { // Run as a worker
24
         do worker(0, rank):
25
       7
26
27
       MPI Barrier (MPI COMM WORLD):
28
       // Here we could start another computation.
29
       MPI_Finalize();
30
       return 0:
31
```

Dynamic Load Balancing: large block size (2500)

Run with:

\$ map mpirun -np 12 ./dyn_integral.x 8000000 2500



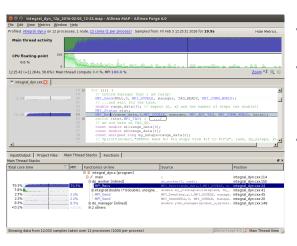
- CPU occupied till half-way
- Spikes in MPI use: Manager receiving data
- Looks like the last worker was holding everyone
- Other workers:

worker starvation
(no work to do)

Dynamic Load Balancing: small block size (2)

Run with:

\$ map mpirun -np 12 ./dyn_integral.x 8000000 2

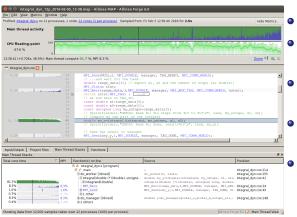


- Brief useful CPU work, then all time is in MPI
- Just moving data around:
 - Workers receiving data
 - 8% (1/12) of time: Manager sending data.
- Very low computation/ communication ratio.

Dynamic Load Balancing: good block size (100)

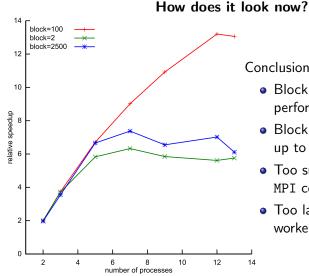
Run with:

\$ map mpirun -np 12 ./dyn_integral.x 8000000 100



- 80% time CPU is busy!
- 8% (1/12) of time: Manager work.
- Mostly MPI by the end of the run.
- OK computation/ communication ratio.
- Still room for improvement!

Dynamic Load Balancing: Strong scaling graph.



Conclusions:

- Block size does affect. performance.
- Block size 100 grows up to node size (12).
- Too small block: MPI communication overhead.
- Too large block: workers starvation.

Concluding slide

Take-home message

- Once you made your program parallel, do simple scaling experiments.
- If scaling is bad, use profiling tools to understand why.
- Allinea Map is available for all Flux users.
- Map can analyze not only parallel, but single-node and OpenMP performance. (Can show something if time permits!)
- If you need any advise and/or help with your parallel programming,
 ARC provides free consulting service
 - Just send a mail to HPC support...
 - ...or directly to: "Alexander Gaenko" <galexv@umich.edu>

Thank you for your attention!