Termination of original F5

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Main goal

Prove that F5 algorithm terminates on any input data represented by a set of homogeneous polynomials.

1. Definitions
   - Polynomial signature
   - Algorithm F5: original version

2. Termination proof
   - F5 terminates: scheme of proof
   - S-pair chains
   - Representations

Polynomial signature

- Fix ideal generated by "input" polynomials $I = (f_1, \ldots, f_n)$
- Fix order $\prec$ on pairs (m, f_i); m is monomial without coefficient

Definitions

Signature: $\max_{\prec}(m_k, f_{i_k})$ for some input-representation

$$g = \sum_k c_k m_k f_{i_k},\ c_k \neq 0, \text{ all pairs } (m_k, f_{i_k}) \text{ are distinct}$$

Minimal signature: $\mathcal{S}(g) \stackrel{\text{def}}{=} min_{\prec}$ (all signatures of g)
We say *mh sig-safe reduce g* if $\mathrm{LM}(mh) = \mathrm{LM}(g)$, $\mathcal{S}(mh) \prec \mathcal{S}(g)$

- F5 tracks $\mathcal{S}$ and performs only signature-safe reductions
- F5 use Position-over-Term order $\prec$ (compare i-s than m-s)
  - Order is compatible to multiplication by monomial
  - Allows to define *sig-lead ratio* with linear order:

$$\frac{\mathcal{S}(g)}{\mathrm{LM}(g)} \prec_{SL} \frac{\mathcal{S}(h)}{\mathrm{LM}(h)} \stackrel{\text{def}}{\Leftrightarrow} \mathcal{S}(g)\,\mathrm{LM}(h) \prec \mathcal{S}(h)\,\mathrm{LM}(g)$$

Algorithm F5: original version

Scheme: computing basis of $(G \cup \{f_i\})$ from basis G

1. $P = \{\text{S-pair}(p, f_i) | p \in G$, and S-pair parts pass 'F5' criterion$\}$
2. **while** $P \neq \varnothing$ **do:**
   - 2.1 $F =$ S-polys(take P all max-degree S-pairs passing 'Rewritten')
   - 2.2 **while** $\exists h \in F$, $\mathcal{S}(h)$ is minimal for non-zero F polynomials **do:**
     - 2.2.1 **if** $\exists p \in G$ – sig-safe h reductor, satisfying 'F5' and 'Rewritten'**:** $F = (F \setminus \{h\}) \cup$Reduce(h,p)
     - 2.2.2 **else:** $F = F \setminus \{h\}$; $R = R \cup \{h\}$
   - 2.3 **for** $r \in R$**:** $P = P \cup \{\text{S-pair}(p, r) | p \in G$,'F5' passed$\}$; $G = G \cup \{r\}$

Termination of original F5

- is topical: no algorithm was proved to be always faster than F5
- is not like Buchberger: G may be extended by polynomial h whose reductor $p \in G$ is not sig-safe or is not satisfying criteria
- proofed in original work only if $f_1, \ldots, f_n$ is regular sequence
- proof by reformulation as generic algorithm (TRB-F5; F5GEN)
  - equivalence proof has gaps

F5 terminates: scheme of proof

1* If the algorithm doesn't stop it fills G with polynomials containing infinite sequence with $\prec_{SL}$-increasing sig-lead ratio

2 This sequence contains p_1, p_2 with $\mathrm{LM}(p_2)|\mathrm{LM}(p_1)$

3 Polynomial p_2 is sig-safe reductor for p_1

4* Any polynomial p with signature smaller than signature of h on step 2.2 has representation of the form $p = \sum_k c_k m_k g_k$, $c_k \neq 0, g_k \in G$ where $\mathcal{S}(m_k g_k) \preccurlyeq \mathcal{S}(p), \mathrm{LM}(m_k g_k) \leqslant \mathrm{LM}(p)$ and all $m_k g_k$ satisfy F5 and Rewritten criteria

5 Applying this to p_2 on step with $h = p_1$ and selecting element with $\mathrm{LM}(m_k g_k) = \mathrm{LM}(p_2)$ as $m p_3$ gives polynomial p_3 in G.

6 p_3 is sig-safe reductor of p_1 and satisfies criteria

7 This leads to contradiction because p_1 was added to G without reduction by p_3

**studied below in more detail*

S-pair chains

Representations