
CodeBruijn Programming, Categorically

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I'll write this later when I know what I'm saying.

CodeBruijn (“co de Bruijn”) programming is a methodology for representing and manipulating syntax in a nameless way, like de Bruijn representation [5], but making the opposite canonical choice for exposing the *non*-use of variables. The essence of the idea is to restrict the *scope* (which variables *may* occur free) of a term to its *support* (which variables *do* occur free). Variables are expelled from scope at the roots of maximal subtrees in which they are not used, where de Bruijn representation keeps variables in scope from their binding sites all the way to the leaves, i.e., the minimal subtrees in which they are not used. The codeBruijn approach thus relies on the maintenance of subtle invariants, reminiscent of Kennaway and Sleep’s ‘director strings’ representation [7]. Dependently typed programming languages such as Agda [10] readily taken on the task of minding that business for us. This paper shows how, and hopefully, why.

The key structure at work is the semi-simplicial category on scopes, i.e., the category of order-preserving embeddings (colloquially, ‘thinnings’) from one scope to another. From Bellegarde and Hook [3], via Bird and Paterson [4], and Altenkirch and Reus [2], it has become a commonplace to index types of terms by their scopes. Such types should really be thought of as ‘thinnables’ — *functors* from the thinnings — because thinnings act compositionally on terms. The operation of mapping a scope to its identity thinning induces a forgetful functor from thinnables to scope-indexed types. This forgetful functor has a celebrated right adjoint, amounting to abstraction over all scopes into which one’s own embeds, which is the basis of Altenkirch, Hofmann and Streicher’s

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Kripke-model construction which drives normalization by evaluation [1]. But, being a forgetful functor, one should ask after its *left* adjoint. That exists, of course, and is the basis of codeBruijn programming: we define *relevant* terms, indexed by *support*, then make them *freely* thinnable by attaching the thinning from support to scope at the root of the term. Further thinnings act by composition at the root, without traversing the term at all.

This paper is written as a literate Agda implementation of a codeBruijn toolkit, structured categorically. I formalise the active categorical abstractions, given provocation from the task at hand. I have adopted something of a tutorial style, partly because there is some novelty in teaching category theory to functional programmers with examples which are not sets-and-functions, but mostly because I am teaching myself. I hope it is also a useful engagement with dependently typed programming, for category theorists. I shall certainly draw the *diagrams* which drive the constructions. There will also be some transferable lessons about programming ‘craft’ which I shall seek to draw out.

1 de Bruijn Representation

In 1999, Altenkirch and Reus gave a de Bruijn treatment of simply typed λ -calculus, together with an implementation of simultaneous substitution [2]. Let us review how it goes.

The simple types are given inductively.

```
infixr 30 _▷_
data Ty : Set where
  base  : Ty
  _▷_   : Ty → Ty → Ty
```

In Agda, infix operators are named with / in their argument places. Types are arranged in backwards (‘snoc-’) lists to form *contexts*.

```
infixl 20 _→_,_
data Bwd (X : Set) : Set where
  []      : Bwd X
  _→_,_   : Bwd X → X → Bwd X
Cx = Bwd Ty
```

Craft 1 (Backwards Lists) *Forwards lists are much more commonplace in functional programming, but I have learned the hard way to use a separate type for lists which grow on the right. The cognitive cost of interpreting lists in reverse is higher, at times, than I can pay: I make mistakes. I also choose symbols for ‘snoc’ and ‘cons’ which avoid misleading reflectional symmetry and have modest pictographic value.*

Typed de Bruijn indices select *one* entry from a context.

```
infix 10 _←_
data _←_ (τ : Ty) : Cx → Set where
  ze :   τ ← Γ  →, τ
  su :   τ ← Γ
      → τ ← Γ  →, σ
```

Craft 2 (Parameters, Uniform Indices, Restrictable Indices) *In the **data** declaration of an Agda type former, some things are declared left of : and scope over the whole declaration. They must be used uniformly in the return types of value constructors. They are, however, free to vary in the types of recursive substructures. If they do so vary, we call them uniform indices. Only if they remain constant throughout should we refer to them as parameters. So, τ , above is a parameter, but Γ , below is a uniform index. The distinction impacts the category in which an initial object is being constructed. $(\tau \leftarrow _)$ is constructed in $Cx \rightarrow Set$, while $_ \vdash _$ is constructed in $Cx \rightarrow Ty \rightarrow Set$. Meanwhile, right of : come those things which may be restricted to particular patterns of value in the return types of value constructors, e.g., nonempty contexts above, and function types below.*

The type of terms reflect the typing rules, indexed by a context and the type being inhabited.

```
infix 10 _⊢_
data _⊢_ (Γ : Cx) : Ty → Set where
  va :   τ ← Γ
      → Γ ⊢ τ
  ap :   Γ ⊢ σ ⊃ τ
      → Γ ⊢ σ
      → Γ ⊢ τ
```

$$\begin{aligned} \text{la} : \quad & \Gamma \multimap, \sigma \vdash \tau \\ & \rightarrow \Gamma \vdash \sigma \supset \tau \end{aligned}$$

Observe that the context we are handed at the root of a term only ever gets larger, each time we use a **lambda**. Only when we reach a **variable** do we choose one thing from the context and disregard the rest.

Now, a *simultaneous substitution* is a type-respecting mapping from variables in some source context Γ to terms over target context Δ — from $(_ \leftarrow \Gamma)$ to $(\Delta \vdash _)$, if you will. When we push such a thing under **la**, we need instead a mapping from $(_ \leftarrow \Gamma \multimap, \sigma)$ to $(\Delta \multimap, \sigma \vdash _)$. We can map the newly bound **ze** to **va ze**, but the trouble is that all of Γ 's variables are mapped to terms over Δ , not $\Delta \multimap, \sigma$. It is thus necessary to traverse all those terms and adjust their leaves, because it is only at the leaves that we document *non*-usage of variables. Shift happens.

Worse, if we attempt to carry out the shift by simultaneous substitution, we leave the comfortable territory of structural recursion and have some explaining to do. It is useful to observe that shifts are merely simultaneous (order-preserving, injective) renumberings which may readily act on terms. Once we have simultaneous renumbering available, simultaneous substitution is easy. Moreover, they are very similar, so we may readily abstract the common structure, as I learned from Goguen and McKinna and demonstrated in my thesis [6, 8].

Shifts — simultaneous renumberings — are the problem, but they are also the key to the solution.

2 Thinnings

Definition 3 (Thinnings) We may define the thinnings, i.e., the order-preserving embeddings, our simultaneous renumberings, as follows:

```

module _ {  $X$  : Set } where    -- fix a set  $X$  of sorts, e.g., Ty

infix 10 _ $\leq$ _
infixl 20 _ $\multimap$ _

data _ $\leq$ _ : Bwd  $X$   $\rightarrow$  Bwd  $X$   $\rightarrow$  Set where

    _ $\multimap$ _ :                                $\gamma$            $\leq$   $\delta$ 
       $\rightarrow \forall x \rightarrow$   $\gamma$            $\leq$   $\delta \multimap, x$ 

```

$$\begin{aligned}
- \neg, - & : & \gamma & \leq \delta \\
& \rightarrow \forall x \rightarrow & \gamma \neg, x & \leq \delta \neg, x \\
[] & : & [] & \leq []
\end{aligned}$$

I am careful to speak of our backward lists as *scopes*, rather than *contexts*, as it is not necessary for them to document the *types* of the variables for this machinery to work.

We lift the constructors from lists to represent the situation where parts of the source scope are copied to the target, but we also introduce \neg to insert an extra element in the target.

Electronic engineers will notice that a thinning is more or less a vector of bits, with \neg for 0 and \neg for 1, but it is indexed by its *population* — the entries marked 1. Expect Boolean operations.

Definition 4 (identity thinnings, ι) *Let us observe that identity thinnings exist, copying their scope.*

$$\begin{aligned}
\iota & : \forall \{\gamma\} \rightarrow \gamma \leq \gamma \\
\iota \{[]\} & = [] \\
\iota \{\gamma \neg, x\} & = \iota \{\gamma\} \neg, x
\end{aligned}$$

Craft 5 (Implicits) *Agda uses curly braces to mark arguments which are normally suppressed. In general, it is sensible to adopt the suppression convention appropriate for the expected use sites. Here, the fact that \leq is a type constructor means that γ will be determined if the type is given in advance. Often, we then have to use an explicit override at definition sites. This sort of thing never happens in Hindley–Milner languages because any information for which inference is permitted is guaranteed to be operationally useless. The inference of operationally useful information represents progress.*

Further, thinnings compose. I write composition diagrammatically.

Definition 6 (thinning composition, \circ) *Thinnings fore and aft compose thus:*

$$\begin{aligned}
& \text{infixl } 20 \neg \circ \neg \\
\neg \circ \neg & : \gamma \leq \delta \rightarrow \delta \leq \zeta \rightarrow \gamma \leq \zeta \\
\theta & \circ (\phi \neg x) = \theta \circ \phi \neg x
\end{aligned}$$

$$\begin{aligned}
(\theta \multimap x) \mathbin{\text{\textcolor{teal}{;}}} (\phi \multimap, x) &= \theta \mathbin{\text{\textcolor{teal}{;}}} \phi \multimap x \\
(\theta \multimap, x) \mathbin{\text{\textcolor{teal}{;}}} (\phi \multimap, x) &= \theta \mathbin{\text{\textcolor{teal}{;}}} \phi \multimap, x \\
\boxed{} \mathbin{\text{\textcolor{teal}{;}}} \boxed{} &= \boxed{}
\end{aligned}$$

Craft 7 (Operator Priority and Association) *My habit is to arrange priority and association so that computation results in net decrease of parentheses.*

Craft 8 (Laziness and Definitional Equality) *Working in intensional type theories like Agda involves a certain amount of care with the equational properties of programs — the typechecker will run these programs, as defined, on open terms. So the order of the lines in the above program matters. If the aft-thinning inserts, there is no call to inspect the fore-thinning. Only if the aft thinning copies need we ask what the fore-thinning gives us. If the first line is moved later, the function becomes unnecessarily strict in the fore-thinning, and definitional equality loses power.*

The reader should note that I will shortly substitute a subtly different definition of composition for this one. Rest assured that its replacement will satisfy the above equations.

Before we move on, though, observe that our type of de Bruijn variables, $\tau \leftarrow \Gamma$ can readily be expressed as the *singleton thinning*, $\boxed{} \multimap, \tau \leq \Gamma$. Let us therefore refactor our previous definitions:

Definition 9 (Singleton Thinning) *Element selection is given by thinning from a singleton source*

$$\begin{aligned}
_ \leftarrow _ &: X \rightarrow \text{Bwd } X \rightarrow \text{Set} \\
x \leftarrow \gamma &= \boxed{} \multimap, x \leq \gamma
\end{aligned}$$

It is now immediately clear why thinnings have a sensible action on de Bruijn terms — thinnings act on de Bruijn variables *by composition*.

3 When 'Green Slime' is Bad, Avoid It

We are accustomed to reasoning by *equation*.

Definition 10 (Inductive Equality, \sim) *In Agda, we may give an inductive definition of equality, as follows:*

```

infix 5 _~_
data _~_ {l} {X : Set l} (x : X) : X → Set l where
  r~ : x ~ x

```

The l parameter is an arbitrary level in the Russell-style hierarchy: `Set` abbreviates `Set 0`; `Set 0 : Set 1`, and so on.

The above definition is *intensional*, in that we can give canonical evidence that $x \sim y$ only if x and y have the same *implementation*, up to definitional equality. We shall have trouble because of that, in what is to follow, but that is not the trouble I mean to discuss in this section.

Suppose we have an hypothesis $q : \theta \mathbin{\text{;}} \phi \sim \psi \neg, x$. We ought to be able to deduce that θ and ϕ are both made by $(_\neg, x)$. However, pattern matching on the equality proof q will fail, because it is not clear how to unify $\theta \mathbin{\text{;}} \phi$ with $\psi \neg, x$, unless we are gifted with the power to run functions backwards, which Agda is not. Our only option is to *remember* how $\mathbin{\text{;}}$ computes, then match on ϕ and θ , eliminating the impossible cases which arise by refuting q .

‘Green slime’ is a colloquialism for expressions involving recursively defined functions which do not compute to canonical form. It is toxic to unification, as it unifies only with variables (purple things). Fortunately, there are strategies to avoid it.

Craft 11 (Inductive Relations are Invertible) *It is often useful to replace equations $f\ s \sim t$ with relations $F\ s\ t$, where F is defined inductively to be the graph of f .*

Let us put this into practice at once!

Definition 12 (composition triangle) *Reading from the definition of $\mathbin{\text{;}}$, construct its graph, thus:*

```

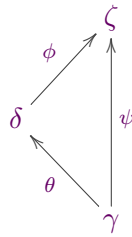
infixl 20 _⊢_,_
data ⊢_ ; _⊢~_ : (θ : γ ≤ δ) (φ : δ ≤ ζ) (ψ : γ ≤ ζ) → Set where
  _⊢_ :
    → ∀ x →
      ⊢ θ          ; φ          ⊢~ ψ
      ⊢ θ          ; φ ⊢ x ⊢~ ψ ⊢ x
  _⊢_,_ :
    → ∀ x →
      ⊢ θ          ; φ          ⊢~ ψ
      ⊢ θ ⊢ x ; φ ⊢, x ⊢~ ψ ⊢ x
  _⊢~_ :
    ⊢ θ          ; φ          ⊢~ ψ

```

$$\begin{array}{c} \rightarrow \forall x \rightarrow \\ \boxed{} : \end{array} \quad \begin{array}{c} \vdash \theta \multimap, x \mathbin{\textcolor{teal}{;}} \phi \multimap, x \multimap \psi \multimap, x \\ \vdash \boxed{} \quad \mathbin{\textcolor{teal}{;}} \boxed{} \quad \multimap \boxed{} \end{array}$$

And there was nothing green to be seen! The only green things on the right in the definition of $\mathbin{\textcolor{teal}{;}}$ were recursive calls to $\mathbin{\textcolor{teal}{;}}$, and these we have replaced by variables.

I call inhabitants of $\vdash \theta \mathbin{\textcolor{teal}{;}} \phi \multimap \psi$ *composition triangles* exactly because they witness the commutation of the diagram



Now let us get our hands on them. You might think that the thing to do is to prove

$$(\theta : \gamma \leq \delta) (\phi : \delta \leq \zeta) \rightarrow \vdash \theta \mathbin{\textcolor{teal}{;}} \phi \multimap \theta \mathbin{\textcolor{teal}{;}} \phi$$

but that amounts to implementing composition a *third* time, as well as being pragmatically suboptimal, as I shall explain later. A better move is to reimplement $\mathbin{\textcolor{teal}{;}}$ by proving, morally,

$$(\theta : \gamma \leq \delta) (\phi : \delta \leq \zeta) \rightarrow \text{‘}\exists \psi.\text{’} \vdash \theta \mathbin{\textcolor{teal}{;}} \phi \multimap \psi$$

We shall have need of existential quantification!

Definition 13 (Dependent Pair Types, Unit Type) *Dependent pair types (Σ -types, in the jargon) may be introduced as records, where the type of the second projection depends on the value of the first. The definition is polymorphic in the type theoretic hierarchy.*

```

infixr 5  $\multimap$ , $\multimap$ 

record  $\Sigma$  {  $k$   $l$  } (  $S$  : Set  $k$  ) (  $T$  :  $S$   $\rightarrow$  Set  $l$  ) : Set (  $l$   $\sqcup$   $k$  ) where
  constructor  $\multimap$ , $\multimap$ 
  field fst :  $S$ 
         snd :  $T$  fst
open  $\Sigma$  public    -- brings the projections into scope
syntax  $\Sigma$   $S$  (  $\lambda x \rightarrow T$  ) = (  $x$  :  $S$  )  $\times$   $T$ 

```

record One $\{l\} : \text{Set } l$ **where constructor** $\langle \rangle$

It is convenient to sugar dependent pair types as a binding form, after the fashion of dependent function types¹.

Three commonly occurring variations merit abbreviation.

infixr 4 $-\times-$
 $-\times- : \text{Set } k \rightarrow \text{Set } l \rightarrow \text{Set } (l \sqcup k)$
 $S \times T = (_ : S) \times T$
 $-\dot{\times}- : \{S : \text{Set } k\} (P\ Q : S \rightarrow \text{Set } l) \rightarrow (S \rightarrow \text{Set } l)$
 $(P \dot{\times} Q)\ s = P\ s \times Q\ s$
 $\langle \rangle : \{S : \text{Set } k\} (P : S \rightarrow \text{Set } l) \rightarrow \text{Set } (l \sqcup k)$
 $\langle P \rangle = (x : _) \times P\ x$

The first of these is ordinary non-dependent pairing. The second is its point-wise lifting to ‘predicates’: $P \dot{\times} Q$ holds whenever P holds and Q holds. The third, pronounced ‘possibly P ’ or ‘something is P ’, asserts that a ‘predicate’ (i.e., a function to $\text{Set } l$) is somehow satisfied: the domain of the predicate is elided. Now we can write $\langle P \dot{\times} Q \rangle$ for ‘something is both P and Q ’.

Craft 14 (Mixfix Operator Sections, Pattern Synonyms) The relation $\ulcorner _ \S _ \urcorner \sim _$ is given in mixfix form exactly because Agda supports operator sections — $(\ulcorner _ \S \phi \urcorner \sim \psi)$, $(\ulcorner \theta \S _ \urcorner \sim \psi)$ and $(\ulcorner \theta \S \phi \urcorner \sim _)$ — which are predicates designed to be used with $\langle \rangle$. We can write

$$\langle (\ulcorner \theta_0 \S \phi_0 \urcorner \sim _) \dot{\times} (\ulcorner \theta_1 \S \phi_1 \urcorner \sim _) \rangle$$

to assert the existence of a commuting square. Very often, the existential witness is not important. Agda provides the means to elide it:

infixr 6 $-\dot{;}-$
pattern $-\dot{;}-\ p\ q = _,\ p,\ q$

A pattern synonym is essentially a macro which can be used on either side of the $=$ sign: in this case, we ignore the existential witness on the left and demand its synthesis on the right.

¹Agda does not let you do precisely this, but L^AT_EX does the rest.

Lemma 15 (constructing composition triangles, $\langle \circ \rangle$) If θ and ϕ are thinnings which meet in the middle, it is possible to construct their composition triangle. Accordingly, we may redefine \circ to project the existential witness.

```

infix 10  $-\langle \circ \rangle-$ 
 $-\langle \circ \rangle-$  :  $(\theta : \gamma \leq \delta) (\phi : \delta \leq \zeta) \rightarrow \langle (\ulcorner \theta \circ \phi \urcorner \sim -) \rangle$ 
 $\theta \quad \langle \circ \rangle \phi \multimap x$  with  $-, v \leftarrow \theta \langle \circ \rangle \phi = -, v \multimap x$ 
 $\theta \multimap x \langle \circ \rangle \phi \neg, x$  with  $-, v \leftarrow \theta \langle \circ \rangle \phi = -, v \multimap x$ 
 $\theta \neg, x \langle \circ \rangle \phi \neg, x$  with  $-, v \leftarrow \theta \langle \circ \rangle \phi = -, v \neg, x$ 
 $\square \quad \langle \circ \rangle \square = -, \square$ 
 $-\circ-$  :  $\gamma \leq \delta \rightarrow \delta \leq \zeta \rightarrow \gamma \leq \zeta$ 
 $\theta \circ \phi = \text{fst } (\theta \langle \circ \rangle \phi)$ 

```

Craft 16 (with Programs) The **with** programming notation [9]

```

f  $\vec{p}$  with  $e$ 
f  $\vec{p}_1 \mid p_1 = e_1$ 
     $\vdots$ 
f  $\vec{p}_n \mid p_n = e_n$ 

```

allows us to extend the left-hand side of a program with an extra column for the value of e , so we may match patterns anywhere, refining the original patterns \vec{p} as well as asking about e . Usefully, e is abstracted from any types where it occurs, so matching on its value refines types, too.

By defining \circ as a projection of $\langle \circ \rangle$, we make **with** $\theta \langle \circ \rangle \phi$ abstract all occurrences of \circ at the same time as it yields up the composition triangle. The same is true of any program defined as the existential witness to the possibility of satisfying a specification.

The recently added notation, **with** $p \leftarrow e$, allows us to avoid introducing a new line to the pattern match if there is only one case, given by p .

Craft 17 (Invisible Programs) Agda allows the ‘don’t care’ / to be used both in patterns, where it neglects to ask a question, and in expressions, where it neglects to give an answer. In the latter case, the missing term must be inferable by unification. In the case of function graphs, the program already given in the relation determines by its construction the missing existential witnesses. The composition function has all but disappeared!

Now, composition triangles give the graph of a function, and functions are deterministic, accordingly, we should be able to recover the fact that the inputs determine both the output and the witness.

Lemma 18 (Uniqueness of Composition Triangles) *For any given inputs θ and ϕ , there is at most one ψ and at most one v such that $v : \lceil \theta \circ \phi \rceil \sim \psi$.*

```

infix 10 ~-~
~-~ : (v0 : ⌈θ∘ϕ⌋ ∼ ψ0) (v1 : ⌈θ∘ϕ⌋ ∼ ψ1)
      → ((★) : ψ0 ∼ ψ1) × v0 ∼ v1
v0 ~-~ x ~-~ v1 ~-~ x with ★ ← v0 ~-~ v1 = ★
v0 ~-~ x ~-~ v1 ~-~ x with ★ ← v0 ~-~ v1 = ★
v0 ~-~ x ~-~ v1 ~-~ x with ★ ← v0 ~-~ v1 = ★
[] ~-~ [] = ★

```

Craft 19 (Dependent Equations) *The equation $v_0 \sim v_1$ may look ill typed, but it is not. The inlined pattern match in the dependent pair type causes ψ_0 to unify with ψ_1 . This can be quite a convenient way to specify ‘telescopic’ equations.*

Craft 20 (Contraction Pattern, ★) *In typeset code, I write ★ instead of unique patterns which can be constructed by record expansion followed by at most one step of case analysis for each field. Matching with ★ collapses spaces to a point. Agda does not yet support this feature. Here, it abbreviates $r \sim$, $r \sim$, but it can scale to much larger uninteresting types.*

We have a notion of thinning, closed under identity and composition. I like to visualize thinnings as two horizontal sequences of dots. Each dot on the bottom is joined to a dot on top by a vertical chord, but there may be dots on top with no chord incident.



The identity thinning has all chords present. Composition is vertical pasting, followed by the contraction of chords which do not reach the bottom.



Spatial intuition makes it clear, informally, that identities are absorbed and that composition is associative. Let us make that intuition formal.

Lemma 21 (Degenerate Triangles) Every thinning θ induces two degenerate triangles, where θ and ι are composed, yielding θ .

$$\iota \circ \theta : (\theta : \gamma \leq \delta) \rightarrow \ulcorner \iota \circ \theta \urcorner \sim \theta$$

$$\iota \circ (\theta \multimap x) = \iota \circ \theta \multimap x$$

$$\iota \circ (\theta \multimap, x) = \iota \circ \theta \multimap, x$$

$$\iota \circ \square = \square$$

infixl 30 $\multimap \iota$

$$\multimap \iota : (\theta : \gamma \leq \delta) \rightarrow \ulcorner \theta \multimap \iota \urcorner \sim \theta$$

$$(\theta \multimap x) \multimap \iota = \theta \multimap \iota \multimap, x$$

$$(\theta \multimap, x) \multimap \iota = \theta \multimap \iota \multimap, x$$

$$\square \multimap \iota = \square$$

We may readily extract identity absorption laws in equational form.

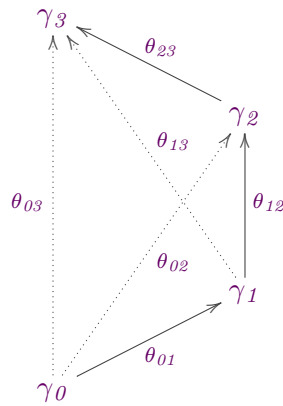
$$\iota \circ \sim : (\theta : \gamma \leq \delta) \rightarrow \iota \circ \theta \sim \theta$$

$$\iota \circ \sim \theta \text{ with } -, v \leftarrow \iota \langle \circ \rangle \theta \mid \star \leftarrow v \multimap \sim \iota \circ \theta = \star$$

$$\multimap \iota : (\theta : \gamma \leq \delta) \rightarrow \theta \multimap \iota \sim \theta$$

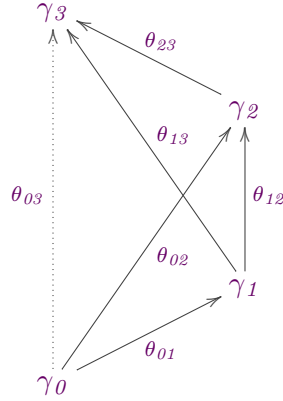
$$\theta \multimap \sim \iota \text{ with } -, v \leftarrow \theta \langle \circ \rangle \iota \mid \star \leftarrow v \multimap \sim \theta \multimap \iota = \star$$

Framing the associativity of composition in terms of triangles gives us a choice. When three arrows compose in sequence, they generate three more, together with four triangles.



Associativity amounts to the assertion that, given any two of the three composite arrows, with the two triangles they generate, the whole diagram can be recovered. All three results are useful, and they are interderivable, but one must be proven by induction — not on *thinnings* but on *triangles*.

Lemma 22 (Associativity (03))



$$\begin{aligned}
 \text{assoc03} &: \langle (\ulcorner \theta_{01} \circ \neg \sim \theta_{02} \urcorner \times \ulcorner \neg \circ \theta_{23} \neg \sim \theta_{13} \urcorner) \rangle \\
 &\rightarrow \langle (\ulcorner \theta_{01} \circ \theta_{13} \neg \sim \neg \urcorner \times \ulcorner \theta_{02} \circ \theta_{23} \neg \sim \neg \urcorner) \rangle \\
 \text{assoc03} (v \quad ; w \multimap x) \text{ with } v' ; w' \leftarrow &\text{assoc03} (v ; w) = v' \multimap x ; w' \multimap x \\
 \text{assoc03} (v \multimap x ; w \multimap x) \text{ with } v' ; w' \leftarrow &\text{assoc03} (v ; w) = v' \multimap x ; w' \multimap x \\
 \text{assoc03} (v \multimap x ; w \multimap x) \text{ with } v' ; w' \leftarrow &\text{assoc03} (v ; w) = v' \multimap x ; w' \multimap x \\
 \text{assoc03} (v \multimap x ; w \multimap x) \text{ with } v' ; w' \leftarrow &\text{assoc03} (v ; w) = v' \multimap x ; w' \multimap x \\
 \text{assoc03} (\Box \quad ; \Box) &= \Box \quad ; \Box
 \end{aligned}$$

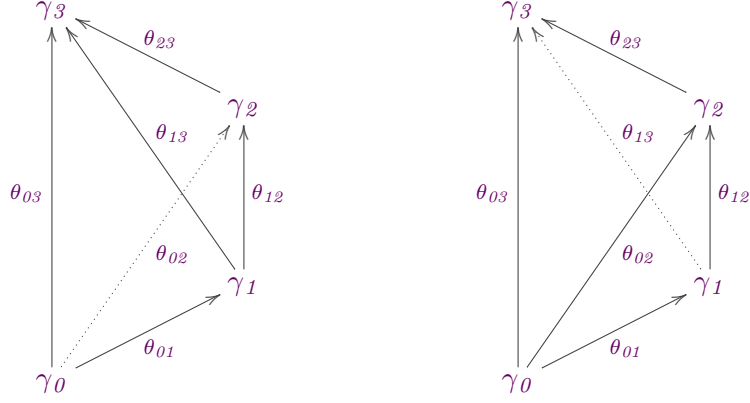
There are three step cases for an inserted x , covering the three stages in the sequence where it can have been inserted. There is one step case for a copied x , which must have been copied in all three stages.

The more familiar equational form of associativity follows by triangulation.

$$\begin{aligned}
 \text{assoc} &: (\theta_{01} : \gamma_0 \leq \gamma_1) (\theta_{12} : \gamma_1 \leq \gamma_2) (\theta_{23} : \gamma_2 \leq \gamma_3) \\
 &\rightarrow \theta_{01} \circ (\theta_{12} \circ \theta_{23}) \sim (\theta_{01} \circ \theta_{12}) \circ \theta_{23} \\
 \text{assoc } \theta_{01} \theta_{12} \theta_{23} \text{ with } \theta_{13} , v_{123} \leftarrow \theta_{12} \langle \circ \rangle \theta_{23} \mid \theta_{02} , v_{012} \leftarrow \theta_{01} \langle \circ \rangle \theta_{12} \mid \\
 &\quad - , v_{013} \leftarrow \theta_{01} \langle \circ \rangle \theta_{13} \mid - , v_{023} \leftarrow \theta_{02} \langle \circ \rangle \theta_{23} \mid \\
 &\quad v'_{013} ; v'_{023} \leftarrow \text{assoc03} (v_{012} ; v_{123}) \mid \star \leftarrow v_{013} \multimap v'_{013} , v_{023} \multimap v'_{023} = \star
 \end{aligned}$$

We may also derive the other two forms of diagrammatic composition.

Lemma 23 (Associativity (02, 13))



$$\begin{aligned}
\text{assoc02} &: \langle (\ulcorner \theta_{01} \urcorner \multimap \neg \sim \theta_{03}) \dot{\times} (\ulcorner \theta_{12} \urcorner \multimap \theta_{23} \neg \sim -) \rangle \\
&\rightarrow \langle (\ulcorner \theta_{01} \urcorner \multimap \theta_{12} \neg \sim -) \dot{\times} (\ulcorner - \urcorner \multimap \theta_{23} \neg \sim \theta_{03}) \rangle \\
\text{assoc02} \{ \theta_{01} = \theta_{01} \} \{ \theta_{12} = \theta_{12} \} (v_{013} ; v_{123}) &\textbf{with } \theta_{02}, v_{012} \leftarrow \theta_{01} \langle \multimap \rangle \theta_{12} \mid \\
v'_{013} ; v_{023} &\leftarrow \text{assoc03} (v_{012} ; v_{123}) \mid \star \leftarrow v_{013} \rightsquigarrow v'_{013} \\
&= v_{012} ; v_{023} \\
\text{assoc13} &: \langle (\ulcorner \theta_{01} \urcorner \multimap \theta_{12} \neg \sim -) \dot{\times} (\ulcorner - \urcorner \multimap \theta_{23} \neg \sim \theta_{03}) \rangle \\
&\rightarrow \langle (\ulcorner \theta_{01} \urcorner \multimap \neg \sim \theta_{03}) \dot{\times} (\ulcorner \theta_{12} \urcorner \multimap \theta_{23} \neg \sim -) \rangle \\
\text{assoc13} \{ \theta_{12} = \theta_{12} \} \{ \theta_{23} = \theta_{23} \} (v_{012} ; v_{023}) &\textbf{with } \theta_{13}, v_{123} \leftarrow \theta_{12} \langle \multimap \rangle \theta_{23} \mid \\
v_{013} ; v'_{023} &\leftarrow \text{assoc03} (v_{012} ; v_{123}) \mid \star \leftarrow v_{023} \rightsquigarrow v'_{023} \\
&= v_{013} ; v_{123}
\end{aligned}$$

That is, we construct the missing arrow and one of the triangles by composition. The `assoc03` lemma gives us the other triangle we want and a duplicate of one we already have. Contracting the duplication makes the triangles we want fit together.

We should now be able to construct the category of thinnings, if only we knew what it might mean to construct a category *in type theory*. That is when our troubles really begin.

4 Type Theorists Worry About Equality

An ingenue (or very sophisticated troll) once wrote to some mathematical mailing list asking whether category theory and type theory were the same. Some category theorists answered vaguely in the positive, at which point the type theorists accused them of insufficiently interrogating the meaning of ‘the

same'. The title of this section is tantamount to a definition of the discipline, especially if you come from the school which takes the classification of a thing to be the diagonal of the classified partial equivalence relation which says when two things are the same.

Informally, a category is given by

1. some notion of *objects*;
2. for every pair of objects, *source* and *target*, some notion of *arrows* from source to target;
3. for each object, an *identity* arrow from that object to itself;
4. for each pair of arrows which meet in the middle, a *composite* arrow from the source of the first to the target of the second;
5. ensuring that composition absorbs identity and associates, i.e., that some equations between arrows hold.

For thinnings, our objects are scopes and \leq tells us what the arrows are. We have candidates for identity and composition. We can take the same view of types and functions: `Set` gives our objects, \rightarrow our arrows, and we have the identity function and function composition.

But any type theorist will ask, or rather will *be asked* by their equipment, 'What is the status of equations between arrows?'. For thinnings, which are first order inductive data structures (indeed, bit vectors), the intensional \sim should suffice; for functions, \sim is dangerously restrictive, identifying only functions with the same *implementation*, up to definitional equality. Any workable notion of category within type theory has to negotiate this distinction, which is waved away in everyday mathematical practice.

We have three options:

1. *Worry About Equality*. Work to replace intensional \sim by something which better reflects mathematical intuition. That is the work of many lifetimes, mine included, and it is beginning to pay off. Observational Type Theory gave a good answer to when values are equal, but not such a good answer to when types are equal. (It was never the basis of a usable implementation, a fact for which I bear some blame.) Homotopy Type Theory gives a better answer, and in its Cubical variant, is beginning to materialise. This is the best option, if you have patience.

-
2. *Tell Lies.* Postulate that \sim has the properties we wish it had, e.g. that pointwise equal functions are equal. Get on with exploring the important ideas. Unfortunately, the computational properties of postulates frustrate the execution of actual, if sophisticated, programs when proofs of equations are used to transport actual values between merely provably equal types. None the less, this is the best option, if you have undergraduates.
 3. *Tell Weaker Truths.* Arrange to work up to \sim when you can (e.g., with thinnings) and to weaker notions (e.g., pointwise equality for functions) when you cannot. This is the best option, if you are in a hurry.

My head is with option 1, my heart is with option 2, but my entire digestive system is with option 3, so that is how I shall proceed in this paper.

The plan, which is far from original, is to work with setoids of arrows, carefully managing the appropriate notion of equivalence on a case-by-case basis. A *setoid* is a set equipped with an *arbitrary* equivalence relation.

Definition 24 (Setoid) *Every level of the hierarchy is equipped with a notion of Setoid.*

```
record Setoid l : Set (1+ l) where
  field El : Set l                -- Elements
      Eq : El → El → Set l      -- Equivalence
      Rf : (x : El) → Eq x x    -- Reflexivity
      Sy : (x y : El) → Eq x y → Eq y x -- Symmetry
      Tr : (x y z : El) → Eq x y → Eq y z → Eq x z -- Transitivity
open Setoid public
```

Let us now build some tools to enable the construction of some *Setoids* that we can use to specify arrows in categories and what we expect to be able to prove about them.

Definition 25 (Intensional Setoids) *Every Set gives rise to a Setoid whose equivalence is given by \sim .*

```
IN : Set l → Setoid l
El (IN X) = X
Eq (IN X) = _~_
Rf (IN X) x = r~
```

```

Sy (IN X) x y r~ = r~
Tr (IN X) x y z r~ r~ = r~

```

Definition 26 (Comprehension Setoids) Fix levels k and l .

```

module _ {k l : Level} where

```

From any **Setoid** k , we can construct a further **Setoid** by proof irrelevant comprehension on its elements with respect to a predicate in l .

```

_||_ : (T : Setoid k) (P : El T → Set l) → Setoid (k ⊔ l)
El (T || P) = (x : El T) × P x
Eq (T || P) (t₀, -) (t₁, -) = Eq T t₀ t₁ × One {l}
Rf (T || P) (t, -) = Rf T t, ⟨⟩
Sy (T || P) (t₀, -) (t₁, -) (t₀₁, ⟨⟩) = Sy T t₀ t₁ t₀₁, ⟨⟩
Tr (T || P) (t₀, -) (t₁, -) (t₂, -) (t₀₁, ⟨⟩) (t₁₂, ⟨⟩) = Tr T t₀ t₁ t₂ t₀₁ t₁₂, ⟨⟩

```

The **Eq** for comprehensions demands an element of the unit type instead of a proof that the proofs of P are equal: this is both a vestigial marker of some information that has been thrown away, and the means to bully Agda into accepting that the **Eq** type is at level $k \sqcup l$ rather than level k .

Definition 27 (Pointwise Setoids) Fix a **Set** S and a family of **Setoids** T .

```

module _ {k l} (S : Set k) (T : S → Setoid l) where

```

We may then construct **Setoids** which lift T pointwise by quantification (universal or existential over S).

```

PI : Setoid (k ⊔ l)
El PI = (s : S) → El (T s) -- explicit universal quantification
Eq PI f g = (s : S) → Eq (T s) (f s) (g s)
Rf PI f s = Rf (T s) (f s)
Sy PI f g q s = Sy (T s) (f s) (g s) (q s)
Tr PI f g h p q s = Tr (T s) (f s) (g s) (h s) (p s) (q s)
IM : Setoid (k ⊔ l)
El IM = {s : S} → El (T s) -- implicit universal quantification
Eq IM f g = (s : S) → Eq (T s) (f {s}) (g {s})

```

$$\begin{aligned}
\text{Rf IM } f \ s &= \text{Rf } (T \ s) \ f \\
\text{Sy IM } f \ g \ q \ s &= \text{Sy } (T \ s) \ f \ g \ (q \ s) \\
\text{Tr IM } f \ g \ h \ p \ q \ s &= \text{Tr } (T \ s) \ f \ g \ h \ (p \ s) \ (q \ s) \\
\text{SG} &: \text{Setoid } (k \sqcup l) \\
\text{EI SG} &= (s : S) \times \text{EI } (T \ s) \quad \text{-- existential quantification} \\
\text{Eq SG } (s_0, t_0) \ (s_1, t_1) &= \Sigma (s_0 \sim s_1) \lambda \{r \sim \rightarrow \text{Eq } (T \ s_0) \ t_0 \ t_1\} \\
\text{Rf SG } (s, t) &= r \sim, \text{Rf } (T \ s) \ t \\
\text{Sy SG } (s, t_0) \ (.s, t_1) \ (r \sim, t_{01}) &= r \sim, \text{Sy } (T \ s) \ t_0 \ t_1 \ t_{01} \\
\text{Tr SG } (s, t_0) \ (.s, t_1) \ (.s, t_2) \ (r \sim, t_{01}) \ (r \sim, t_{12}) &= r \sim, \text{Tr } (T \ s) \ t_0 \ t_1 \ t_2 \ t_{01} \ t_{12}
\end{aligned}$$

Definition 28 (Projection Setoids) Fix a Set X .

module $_ \{l\} \{X : \text{Set } l\}$ **where**

If we have a projection from X to the carrier of a Setoid, we can treat X as a quotient, working up to setoid equivalence of projections.

$$\begin{aligned}
\backslash\! &: (S : \text{Setoid } l) (p : X \rightarrow \text{EI } S) \rightarrow \text{Setoid } l \\
\text{EI } (S \backslash\! p) &= X \\
\text{Eq } (S \backslash\! p) \ x \ y &= \text{Eq } S \ (p \ x) \ (p \ y) \\
\text{Rf } (S \backslash\! p) \ x &= \text{Rf } S \ (p \ x) \\
\text{Sy } (S \backslash\! p) \ x \ y &= \text{Sy } S \ (p \ x) \ (p \ y) \\
\text{Tr } (S \backslash\! p) \ x \ y \ z &= \text{Tr } S \ (p \ x) \ (p \ y) \ (p \ z)
\end{aligned}$$

The notation necessarily identifies the target of the projection, but leaves the source implicit.

Craft 29 (Green Things in Blue Packaging) I anticipate that we shall need to construct explanations which look like equational proofs, but are constructed within the equivalence of a known Setoid. I therefore introduce a type constructor whose entire purpose is to fix the Setoid at work. There is no general way to infer the setoid X from a type which is known to be $\text{Eq } X \ x \ y$, so the craft lies in ensuring that we never forget which Setoid we work in.

```

record  $\_ \ni \_ \approx \_ \{l\} (X : \text{Setoid } l) (x \ y : \text{EI } X) : \text{Set } l$  where
  constructor eq
  field  $\text{qe} : \text{Eq } X \ x \ y$ 
  open  $\_ \ni \_ \approx \_$  public

```

When we formulate categorical laws, we shall use this wrapped version.

At last, we are ready to say what a category might be.

5 Something Not Entirely Unlike Categories in Type Theory

What follows is far from perfect. The best that can be said is that it is an effective pragmatic compromise. Neither is it an unusual recipe. I labour the point only to teach the craft of the cooking.

A category will have a **Set** of objects and, indexed by source and target objects, a **Setoid** of arrows.

But there's another catch: type theoretic level. There is no good reason to believe that the level objects live on is in any way related to the the level that arrows live on. Agda is particularly bad at supporting *cumulativity* — implicit upward flow between levels — and by 'bad', I mean it just does not. (Coq by contrast, is rather good at it.) Agda forces one to use level polymorphism instead of cumulativity. The two are poor stablemates, but they have backed the wrong horse. In the now, the pragmatic policy is to keep the levels of objects and arrows separate.

Definition 30 (Category) Fix k , the level of objects, and l , the level of arrows.

We may then define a notion of **Category**.

```
record Cat : Set (1+ (k ⊔ l)) where
  -- We have a Set of Objects, and a family of Setoids of Arrows.
  field Obj  : Set k
       Arr  : Obj → Obj → Setoid l
       -- Agda allows one to pause between fields to make definitions...
  ⊢▷ : Obj → Obj → Set l
  S ▷ T = El (Arr S T)
       -- ...and then resume requesting fields.
  -- We have identity and composition.
  field ι    : T ▷ T
       ⋄     : R ▷ S → S ▷ T → R ▷ T
       -- Locally define equality of arrows...
  _≈_ : {S T : Obj} (f g : S ▷ T) → Set l
  _≈_ {S} {T} f g = Arr S T ⊃ f ≈ g
       -- ...then require the laws.
  field coex : f ≈ f' → g ≈ g' → (f ⋄ g) ≈ (f' ⋄ g')
```

$$\begin{aligned}
\text{idco} &: (f : S \triangleright T) \rightarrow (\iota \circ f) \approx f \\
\text{coid} &: (f : S \triangleright T) \rightarrow (f \circ \iota) \approx f \\
\text{coco} &: (f : R \triangleright S) (g : S \triangleright T) (h : T \triangleright U) \rightarrow (f \circ (g \circ h)) \approx ((f \circ g) \circ h)
\end{aligned}$$

Note the inevitable necessity of `coex`, the explicit witness that composition respects the weak notion of equivalence given by \approx : let us ensure that this proof is always trivial.

As a warm-up, let us construct the category of sets and functions-up-to-pointwise-equality.

Definition 31 (Pointwise Set) *Every level l of the type theoretic hierarchy has a category of sets and functions, considered up to pointwise equality. The objects in the category are large, but the arrows are small.*

```

module _ (l : Level) where
  SET : Cat (1+ l) l
  Obj SET = Set l
  Arr SET S T = PI S λ _ → IN T
  ι SET x = x
  ∘ SET f g x = g (f x)
  qe (coex SET {f = f} (eq qf) (eq qg)) x with f x | qf x
  ... | _ | r~ = qg _
  qe (idco SET f) x = r~
  qe (coid SET f) x = r~
  qe (coco SET f g h) x = r~

```

When giving the extensionality witness for composition, we know only that its arguments agree pointwise. Fortunately for us, the definition of composition uses its arguments by invoking them at specific points.

Definition 32 (Discrete Category) *Every `Set` induces a discrete category with its elements for objects and only identity arrows, given by intensional equality.*

```

DISCRETE : (X : Set l) → Cat l l
Obj (DISCRETE X) = X
Arr (DISCRETE X) x y = IN (One {l}) || λ _ → x ~ y
ι (DISCRETE X) = ★
∘ (DISCRETE X) ★ ★ = ★

```

```

coex (DISCRETE X) ★★ = ★
idco (DISCRETE X) ★   = ★
coid (DISCRETE X) ★   = ★
coco (DISCRETE X) ★★★ = ★

```

I make the arrows carry trivial information, subject to the condition that source and target are equal. I am therefore not obliged to reason about equality between equality proofs.

Definition 33 (Opposite Category) Fix a category C . (The following declaration allows us to refer to its components as $C.\text{Obj}$, and so on.)

```

module _ {k l} (C : Cat k l) where
  private module C = Cat C

```

We construct OP , whose objects are those of C but whose arrows are reversed, as follows:

```

OP : Cat k l
Obj OP      = C.Obj
Arr OP S T  = C.Arr T S
ι    OP      = C.ι
∘OP f g    = g C.∘ f
coex OP qf qg = C.coex qg qf
idco OP      = C.coid
coid OP      = C.idco
coco OP f g h = eq (Sy (C.Arr _ _) _ (qe (C.coco h g f)))

```

Identities are preserved but composition is flipped. Correspondingly, the witness to extensionality of composition also flips and the identity absorption proofs swap over. Associativity of composition relies essentially on the symmetry of setoid equivalence for the arrows of C .

Outside our module, we may now write $\text{OP } C$ for the opposite category of C .

Definition 34 (Restricting Objects) One way to construct a subcategory is to restrict the objects to the image of a function.

```

module _ {k l} (C : Cat k l) {k'} {O : Set k'} where
  private module C = Cat C

```

```

OB : (O → C.Obj) → Cat k' l
Obj (OB f) = O
Arr (OB f) S T = C.Arr (f S) (f T)

```

I omit the remaining fields as they are copied from C .

Finally, in this section, let us assemble the jigsaw pieces which make up the category of thinnings.

Lemma 35 (Category of Thinnings) *Thinnings form the arrows of a category.*

```

THIN : Set → Cat 0 0
Obj (THIN X)      = Bwd X
Arr (THIN X) γ δ  = IN (γ ≤ δ)
ι (THIN X)        = ι
⊖ (THIN X)        = ⊖
coex (THIN X) ★ ★ = ★
idco (THIN X) θ   = eq (ι;~ θ)
coid (THIN X) θ   = eq (θ ~;ι)
coco (THIN X) θ φ ψ = eq (assoc θ φ ψ)

```

Now, the crucial structure that we exploit in codeBruijn programming arises from the relationship between the discrete category on scopes and the category of thinnings. That relationship requires us to consider ‘arrows between categories’, i.e., *functors*.

6 Functors as Arrows

When does one category tell you about another? When we have a structure-preserving translation — a *functor* — between them: maps of cities are useful in part because the lines on the map join up in the same pattern as the streets in the city. What is a functor, in our setting?

Definition 36 (Functor) *Fix a source category C and a target category D . Their relative sizes are unimportant.*

```

module _ {l k l' k'} (C : Cat l k) (D : Cat l' k') where
  private module C = Cat C; module D = Cat D

```

We may construct a **Setoid** of structure-preserving translations from C to D . Every C object must *Map* to a D object, and every C arrow must *map* to a D arrow, compatibly with *Map*.

```

Functor : Setoid (l ⊔ k ⊔ l' ⊔ k')
Functor =
  SG (C.Obj → D.Obj) λ Map →
  (IM C.Obj λ S → IM C.Obj λ T → PI (S C.▷ T) λ _ →
    D.Arr (Map S) (Map T))
  || λ map →

```

The use of a comprehension allows us to add the laws that *map* must obey: it must respect equivalence, and preserve identities and composition.

```

  ({ S T : C.Obj } { f g : S C.▷ T } →
    f C.≈ g → map f D.≈ map g)
× ({ X : C.Obj } →
  map (C.ι { X }) D.≈ D.ι { Map X })
× ({ R S T : C.Obj } (f : R C.▷ S) (g : S C.▷ T) →
  map (f C.⋄ g) D.≈ (map f D.⋄ map g))

```

Now, we should like Agda to recognize functors — elements of that **Setoid** — when she sees them and to know their source and target categories implicitly. This calls for ‘green things in blue packaging’.

```

record ⇒ : Set (l ⊔ k ⊔ l' ⊔ k') where
  constructor fun
  field nuf : El Functor
  Map      = fst nuf
  map      = fst (snd nuf)
  mapex    = fst (snd (snd nuf))
  mapid    = fst (snd (snd (snd nuf)))
  mapco    = snd (snd (snd (snd nuf)))
open ⇒ public

```

The \Rightarrow operator becomes infix when the module fixing C and D is opened.

I have made **Functor** a **Setoid** for a reason: they will shortly be the **Arrows** of a **Category**. However, before we can go to work with this definition, we should tool up for the reasoning within setoids that we shall inevitably face.

Craft 37 (Equivalence Reasoning for Setoids) Fix a **Setoid**, X , and name its laws.

```

module _ { $l$ } { $X$  : Setoid  $l$ } where
  private  $\text{RfX} = \text{Rf } X$ ;  $\text{SyX} = \text{Sy } X$ ;  $\text{TrX} = \text{Tr } X$ 

```

We may benefit from ‘green things in blue packaging’ by suppressing more detail when the setoid at work is fixed. E.g., we can drop the element from the reflexivity witness.

```

 $\text{r}\approx : X \ni x \approx x$ 
 $\text{r}\approx \{x\} = \text{eq } (\text{RfX } x)$ 

```

Moreover, we can build some combinators which facilitate readable chains of equational reasoning, showing the steps and the explanations which justify them.

```

infixr 5  $\_ \approx \_$ 
infixr 6  $\_ \square$ 

 $\_ \approx \_$  :  $\forall x \rightarrow X \ni x \approx y \rightarrow X \ni y \approx z \rightarrow X \ni x \approx z$ 
 $x \approx \text{eq } q \approx \text{eq } q' = \text{eq } (\text{TrX } \_ \_ \_ q \ q')$ 

 $\_ \approx \_$  :  $\forall x \rightarrow X \ni y \approx x \rightarrow X \ni y \approx z \rightarrow X \ni x \approx z$ 
 $x \approx \text{eq } q \approx \text{eq } q' = \text{eq } (\text{TrX } \_ \_ \_ (\text{SyX } \_ \_ q) \ q')$ 

 $\_ \square$  :  $(x : \text{El } X) \rightarrow X \ni x \approx x$ 
 $x \square = \text{eq } (\text{RfX } x)$ 

```

Lastly, we have a combinator which allows us to fix the setoid.

```

qprf :  $(X : \text{Setoid } l) \{x \ y : \text{El } X\} \rightarrow X \ni x \approx y \rightarrow \text{Eq } X \ x \ y$ 
qprf  $X = \text{qe}$ 

```

Let us have some examples.

Definition 38 (Forgetting Arrows) Fix a category

```

module _ { $k \ l$ } ( $C : \text{Cat } k \ l$ ) where
  private module  $C = \text{Cat } C$ 

```

The discrete category on the objects of C has a functor into C .

```

FORGET : DISCRETE  $C$ .Obj  $\Rightarrow C$ 
nuf FORGET

```

```

= (λ  $X \rightarrow X$ )      -- objects map to themselves
, (λ {★ →  $C.\iota$ })    -- arrows map to identities
, (λ {{-} {-} {★} {★} → eq (Rf (C.Arr - -) -)})
, r≈
, λ {★ ★ →
     $C.\iota$       ≈ (C.idco -) ≈
     $C.\iota C.;$   $C.\iota \square$ }

```

Definition 39 (Functions as Functors) A function $f : S \rightarrow T$ is trivially a functor between discrete categories.

```

FUN : ( $S \rightarrow T$ ) → DISCRETE  $S \Rightarrow$  DISCRETE  $T$ 
nuf (FUN  $f$ ) =  $f$ 
, (λ {★ → ★}) , (λ {{-} {-} {★} {★} → ★}) , ★ , λ {★ ★ → ★}

```

Definition 40 (Identity Functor) Fix a category C .

```

module _ { $k\ l : \text{Level}$ } { $C : \text{Cat } k\ l$ } where
  open Cat  $C$ 

```

There is a functor from C to itself.

```

I :  $C \Rightarrow C$ 
nuf I = (λ  $X \rightarrow X$ )  -- identity on objects
, (λ  $f \rightarrow f$ )         -- identity on morphisms
, (λ  $q \rightarrow q$ )         -- identity is extensional
, r≈                  -- identity preserves  $\iota$ 
, λ _ _ → r≈         -- identity preserves  $;$ 

```

Lemma 41 (Functors compose) Fix three categories.

```

module _ { $kC\ lC\ kD\ lD\ kE\ lE$ } { $C : \text{Cat } kC\ lC$ } { $D : \text{Cat } kD\ lD$ } { $E : \text{Cat } kE\ lE$ } where
  private module  $C = \text{Cat } C$ ; module  $D = \text{Cat } D$ ; module  $E = \text{Cat } E$ 

```

```

⋈ :  $C \Rightarrow D \rightarrow D \Rightarrow E \rightarrow C \Rightarrow E$ 
nuf (F ; G)
= (λ  $X \rightarrow \text{Map } G (\text{Map } F X)$ )  -- compose Maps

```

```

, (λ f → map G (map F f))    -- compose maps
, (λ q → mapex G (mapex F q)) -- compose extensionality witnesses
, ( map G (map F C.ι)        ≈[ mapex G (mapid F) ]≈
    map G D.ι                ≈[ mapid G ]≈
    E.ι                      □)
, λ f g →
    map G (map F (f C.⋈ g))   ≈[ mapex G (mapco F _ _) ]≈
    map G (map F f D.⋈ map F g) ≈[ mapco G _ _ ]≈
    (map G (map F f) E.⋈ map G (map F g)) □

```

We can acquire another example for the price of one useful data structure.

Definition 42 (Environment) *An environment is a scope-indexed sequence of sort-indexed values.*

```

module _ { X : Set } where
  data Env ( V : X → Set ) : Bwd X → Set where
    [] : Env V []
    --,-- : ∀ { xz x } → Env V xz → V x → Env V (xz --, x)

```

If we flip our thinking about thinnings and see them as the *selection* of a subscope from a scope, we should be able to make the corresponding selection of values from an environment. In particular, if a thinning tells us how the support of a term embeds in its scope, we can whittle an environment down to just those values that are pertinent to the term.

Lemma 43 (Selection) *Env V is the action on objects of a functor from reversed thinnings to sets-and-functions.*

```

module _ ( V : X → Set ) where
  SELECT : OP (THIN X) ⇒ SET 0
  fst (nuf SELECT) = Env V

```

To prove the lemma, we must construct the rest of the **Functor**. We should prepare the components separately, then slot them into place. Its action on morphisms is to copy where the thinning copies and discard where the thinning inserts.

```

module _ {  $V : X \rightarrow \text{Set}$  } where

  infixr 25  $\triangleleft$ 
   $\triangleleft : \forall \{ \gamma \delta \} \rightarrow \gamma \leq \delta \rightarrow \text{Env } V \delta \rightarrow \text{Env } V \gamma$ 
   $(\theta \hat{\rightarrow} x) \triangleleft (vz \neg, v) = \theta \triangleleft vz$ 
   $(\theta \neg, x) \triangleleft (vz \neg, v) = \theta \triangleleft vz \neg, v$ 
   $[] \triangleleft [] = []$ 

```

We may readily show that the identity thinning is the complete selection.

```

 $\iota \triangleleft : (vz : \text{Env } V \gamma) \rightarrow \iota \triangleleft vz \sim vz$ 
 $\iota \triangleleft [] = r \sim$ 
 $\iota \triangleleft (vz \neg, v) \text{ rewrite } \iota \triangleleft vz = r \sim$ 

```

The composition law is best shown by induction on the graph.

```

 $\Delta \triangleleft : (w : \ulcorner \theta \S \phi \urcorner \sim \psi) (vz : \text{Env } V \_) \rightarrow \psi \triangleleft vz \sim \theta \triangleleft \phi \triangleleft vz$ 
 $\Delta \triangleleft (w \hat{\rightarrow} x) (vz \neg, v) \text{ rewrite } \Delta \triangleleft w vz = r \sim$ 
 $\Delta \triangleleft (w \neg, x) (vz \neg, v) \text{ rewrite } \Delta \triangleleft w vz = r \sim$ 
 $\Delta \triangleleft (w \neg, x) (vz \neg, v) \text{ rewrite } \Delta \triangleleft w vz = r \sim$ 
 $\Delta \triangleleft [] [] = r \sim$ 

```

We can now complete the jigsaw.

```

snd (nuf SELECT) =  $\triangleleft$ , ( $\lambda \{ \star \rightarrow \text{eq } \lambda \_ \rightarrow r \sim \}$ )
               , eq  $\iota \triangleleft$ , ( $\lambda \phi \theta \rightarrow \text{eq } \lambda vz \rightarrow \Delta \triangleleft (\text{snd } (\theta \langle \S \rangle \phi)) vz$ )

```

Note, by the way, that a family of types induces a functor from a discrete category.

```

FAM : ( $X \rightarrow \text{Set } l$ )  $\rightarrow$  DISCRETE  $X \Rightarrow \text{SET } l$ 
FAM  $T = \text{FUN } T ; \text{FORGET } (\text{SET } \_)$ 

```

To finish this section, let us construct the category of categories, with functors for arrows. To start with, we should define the identity functor and functor composition.

Lemma 44 (Category of Categories) *Fix sizes for objects and arrows.*

```

module _ ( $k \ l : \text{Level}$ ) where

```

There is a category whose objects are **Categories** and whose arrows are **Functors**

```

CAT : Cat (1+ (k ⊔ l)) (k ⊔ l)
Obj CAT = Cat k l
Arr CAT C D = Functor C D \nuf {C = C} {D = D}

```

To prove this, we must complete the remaining fields.

1. identity and composition, as above

```

ι CAT = I
∘ CAT = ⋅

```

2. extensionality witness

```

qe (coex CAT {-} {-} {C} {F} {F'} {G} {G'})
  (eq (r~ , qf , ⟨⟩)) (eq (r~ , qg , ⟨⟩))
= r~
, (λ S T f → qprf (Arr C (Map G (Map F S)) (Map G (Map F T))) (
  map G (map F f) ≈ mapex G (eq (qf -- --)) ≈
  map G (map F' f) ≈ eq (qg -- --) ≈
  map G' (map F' f) □))
, ⟨⟩

```

3. absorption and associativity – these are trivial

```

idco CAT {-} {C} - = eq (r~ , (λ - - - → qprf (Arr C - -) r~) , ⟨⟩)
coid CAT {-} {C} - = eq (r~ , (λ - - - → qprf (Arr C - -) r~) , ⟨⟩)
coco CAT {-} {-} {-} {C} - - -
= eq (r~ , (λ - - - → qprf (Arr C - -) r~) , ⟨⟩)

```

7 Natural Transformations — Functors as Objects

We have seen functors as arrows in `CAT`, but they are also objects in *functor categories*, whose arrows are *natural transformations*: families of arrows in the target category which operate *uniformly*. Let us proceed directly.

Lemma 45 (Functor Category) *Fix source category C and target category D , irrespective of size.*

```

module _ { k l k' l' : Level } (C : Cat k l) (D : Cat k' l') where
  private module C = Cat C; module D = Cat D

```

Functors from *C* to *D* are the objects of a category with the following notion of arrow.

```

⇒_ : Cat (k ⊔ l ⊔ k' ⊔ l') (k ⊔ l ⊔ l')
Obj ⇒_ = C ⇒ D
Arr ⇒_ F G =
  (IM C.Obj λ X → D.Arr (Map F X) (Map G X))
  || λ n → { S T : C.Obj } (f : S C ▷ T) →
    (map F f D.; n { T }) D. ≈ (n { S } D.; map G f)

```

What are these arrows? They are families of arrows in *D* indexed by objects in *C*: for each *X*, we get a *D* arrow, *n* { *X* } from *F X* to *G X*. However, they are subject to a naturality condition.

$$\begin{array}{ccc}
 F X & \xrightarrow{n \{X\}} & G X \\
 Fm f \downarrow & & \downarrow Gm f \\
 F Y & \xrightarrow{n \{Y\}} & G Y
 \end{array}$$

That is, *n* turns *F*-shaped things into *G*-shaped things, without regard to contents.

Let us now construct the category.

1. identity, and the proof that it is natural

```

ι ⇒_ { F }
= D.ι , λ f →
  map F f D.; D.ι ≈ [ D.coid _ ] ≈
  map F f ≈ [ D.idco _ ] ≈
  (D.ι D.; map F f) □

```

2. composition, and the proof that it preserves naturality

```

∘ ⇒_ { F } { G } { H } (n , na) (p , pa) = n D.; p , λ f →
  map F f D.; (n D.; p) ≈ [ D.coco _ _ _ ] ≈
  (map F f D.; n) D.; p ≈ [ D.coex (na f) r≈ ] ≈
  (n D.; map G f) D.; p ≈ [ D.coco _ _ _ ] ≈
  n D.; (map G f D.; p) ≈ [ D.coex r≈ (pa f) ] ≈
  n D.; (p D.; map H f) ≈ [ D.coco _ _ _ ] ≈
  (n D.; p) D.; map H f □

```

3. laws — inherited pointwise from D

$$\begin{aligned}
\text{coex } \Rightarrow _ &= (\text{eq } (nq, \langle \rangle)) (\text{eq } (pq, \langle \rangle)) = \\
&= \text{eq } ((\lambda X \rightarrow \text{qe } (D.\text{coex } (\text{eq } (nq X)) (\text{eq } (pq X)))) , \langle \rangle) \\
\text{idco } \Rightarrow _ &= \text{eq } ((\lambda _ \rightarrow \text{qe } (D.\text{idco } _)) , \langle \rangle) \\
\text{coid } \Rightarrow _ &= \text{eq } ((\lambda _ \rightarrow \text{qe } (D.\text{coid } _)) , \langle \rangle) \\
\text{coco } \Rightarrow _ &= \text{eq } ((\lambda _ \rightarrow \text{qe } (D.\text{coco } _ _)) , \langle \rangle)
\end{aligned}$$

Natural transformations are indexed families of functions. It help to introduce notation for working with the latter. Let us do for functions what $\dot{\times}$ and $\langle \rangle$ do for pairs. Firstly, we lift the function type pointwise.

$$\begin{aligned}
_ \dot{\rightarrow} _ &: \{S : \text{Set } k\} (T \ U : S \rightarrow \text{Set } l) \rightarrow (S \rightarrow \text{Set } l) \\
(T \dot{\rightarrow} U) s &= T s \rightarrow U s
\end{aligned}$$

Secondly, we introduce the ‘necessarily’ operator, indicating that a predicate holds for its entire domain.

$$\begin{aligned}
[_] &: \{S : \text{Set } k\} (T : S \rightarrow \text{Set } l) \rightarrow \text{Set } (k \sqcup l) \\
[\ T] &= \forall \{s\} \rightarrow T s
\end{aligned}$$

Together, we obtain $[\ T \dot{\rightarrow} U]$ for an indexed family of functions.

Now, with this apparatus in place, let us revisit the \triangleleft operator, which is natural with respect to the type of values being stored in `Environments`.

Lemma 46 (natural selection) `Env` extends to a functor between functor categories as follows.

$$\begin{aligned}
\text{ENV} &: (\text{DISCRETE } X \Rightarrow \text{SET } 0) \Rightarrow (\text{OP } (\text{THIN } X) \Rightarrow \text{SET } 0) \\
\text{fst } (\text{nuf ENV}) \ V &= \text{SELECT } (\text{Map } V)
\end{aligned}$$

In order to prove this lemma, we must program to its specification. After unwrapping the setoid packaging, we find we must deliver an arrow action which commutes with selection.

$$\begin{aligned}
\text{env} &: [\ V \dot{\rightarrow} W] \rightarrow [\ \text{Env } V \dot{\rightarrow} \text{Env } W] \\
\text{env } f \ [\] &= [\] \\
\text{env } f \ (vz \dot{-}, v) &= \text{env } f \ vz \dot{-}, f v \\
\text{na} \triangleleft &: (f : [\ V \dot{\rightarrow} W]) (\theta : \gamma \leq \delta) (vz : \text{Env } V \delta) \rightarrow
\end{aligned}$$

$$\begin{aligned}
& \text{env } f (\theta \triangleleft vz) \sim \theta \triangleleft \text{env } f vz \\
& \text{na} \triangleleft f (\theta \hat{\rightarrow} x) (vz \rightarrow, v) \text{rewrite } \text{na} \triangleleft f \theta vz = r \sim \\
& \text{na} \triangleleft f (\theta \rightarrow, x) (vz \rightarrow, v) \text{rewrite } \text{na} \triangleleft f \theta vz = r \sim \\
& \text{na} \triangleleft f [] [] = r \sim
\end{aligned}$$

Of course, we must also ensure that `env` is extensional and respects identity and composition.

$$\begin{aligned}
\text{envex} & : ((x : X) (v : V x) \rightarrow f v \sim g v) \rightarrow (vz : \text{Env } V \gamma) \\
& \rightarrow \text{env } f vz \sim \text{env } g vz \\
\text{envid} & : (vz : \text{Env } V \gamma) \\
& \rightarrow \text{env } (\lambda x \rightarrow x) vz \sim vz \\
\text{envco} & : (f : [V_0 \rightarrow V_1]) (g : [V_1 \rightarrow V_2]) \{\gamma\} (vz : \text{Env } V_0 \gamma) \\
& \rightarrow \text{env } (\lambda x \rightarrow g (f x)) vz \sim \text{env } g (\text{env } f vz)
\end{aligned}$$

These are routine functional inductions. The whole jigsaw comes together, as follows:

$$\begin{aligned}
& \text{snd } (\text{nuf ENV}) \\
& = (\lambda \{ (f, -) \rightarrow \text{env } f, \lambda \theta \rightarrow \text{eq } (\text{na} \triangleleft f \theta) \}) \\
& , \quad (\lambda \{ (\text{eq } (q, \langle \rangle)) \rightarrow \text{eq } ((\lambda \gamma \rightarrow \text{envex } q), \langle \rangle) \}) \\
& , \quad \text{eq } ((\lambda \gamma \rightarrow \text{envid}), \langle \rangle) \\
& , \quad \lambda \{ (f, -) (g, -) \rightarrow \text{eq } ((\lambda \gamma \rightarrow \text{envco } f g), \langle \rangle) \}
\end{aligned}$$

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