

CS 506 → Network Analysis (3/24)

- Feel free to reach out for extensions — good to identify when you need more time
- Last classes: Logistic/Linear Regression
- Network Analysis: tools we may encounter, very present in daily lives

• Example Networks — social networks, epigenetic networks, networks between brain regions, internet, wifi (obviously very DIVERSE)

~~Networks~~

— want to formalize the network to ask specific Q's

⇒ Internet: — what is the internet traffic @ a location?

— Anomalous traffic patterns?

— Model of the internet?

⇒ Biology: — which regions of brain communicate for a specific task

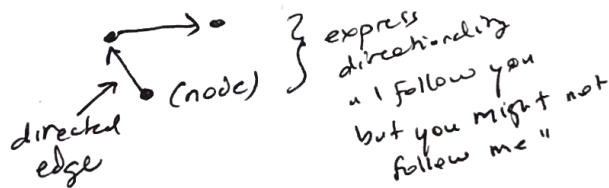
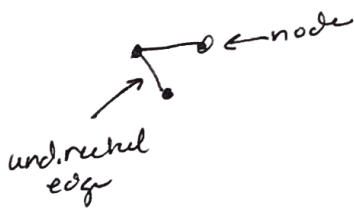
— patterns of interactions between genes

⇒ Social: — who is friends with whom? Who are influencers?

— what social groups are present? — Is there a clear representation of this?

Graphs (Also see: Lecture 4) → 2/8

— use graph/Graph Theory to model/represent and analyze Network



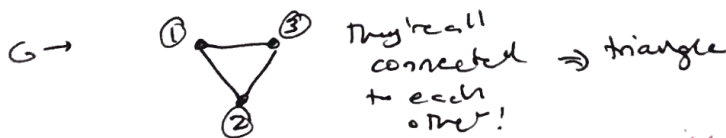
} express directionality
"I follow you
but you might not
follow me"

Formally, graph G is ordered pair of sets (V, E)

Set of all nodes → V
Set of all edges → E

— Let $G = (V, E)$

$V = \{1, 2, 3\}$ $E = \{(1, 2), (1, 3), (2, 3)\}$ ← undirected



— could be non-symmetric matrix when directed

— can also assign weights to matrix (indicates strength of connection)

computational perspective

— best way to store this info, so certain queries are more efficient?

ex. adjacency matrix

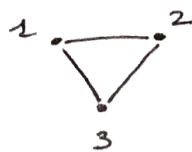
$$\begin{matrix} & \begin{matrix} ① & ② & ③ \end{matrix} \\ \begin{matrix} ① \\ ② \\ ③ \end{matrix} & \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix} \end{matrix}$$

— entry represents whether there's an edge between nodes

note: typically no edge between node with itself (zero diagonal)

Adjacency matrix (further illustrated)

		nodes		
		1	2	3
nodes	1	0	1	1
	2	1	0	1
	3	1	1	0



- walk through columns and rows and draw it out

- note: maximum possible # of edges

↳ when there are only 3 nodes, that means it's a triangle

Adjacency List → lists ^{all} the nodes something is connected to

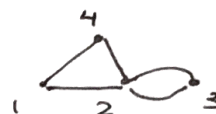
1: { 2, 3 }

2: { 1, 3 }

3: { 1, 2 }

} more space efficient (could use a dict)
BUT makes certain processes harder to solve

→ See Slides: Question from Goad will thinking



- ~~that's a graph~~

"complex" (nah)

Graphs Characteristics

degree of a node → # edges connected to it

↳ average # connected on network → minimum? maximum?

↳ ~~least~~ want to make

sure min degree is low esp. for machines w/ sensitive workload

path between two nodes → sequence of edges that join them

complete when there is an edge between ~~every~~ every pair of nodes

Question: ① sum of degrees of all nodes in graph as a function of N_E (# edges)

→ see triangle



each node degree = 2
answer = 6

② # edges in a complete graph as function of the # nodes (N_V)

↳ 2 endpoints to each edge
↳ so count the edges twice

note: must be undirected

(everyone meets everyone ... how many handshakes?)

$$= \frac{N_V (N_V - 1)}{2}$$

Graph Problems ① Clique Problem (largest complete subgraph)

② coloring problem (no two vertices that share an edge share a color) → coloring a graph w/ certain # colors

③ Traveling Salesman (most efficient route through all cities + return home)

④ shortest path (given edges have weights, path w/ shortest sum of weights)

⑤ vertex cover (min # of nodes to remove + get rid of all the edges)

↳ actually really easy

- Dijkstra algorithm (CS330)

Network characteristics

- Distribution of edges / node degrees:
 - anomaly detection
 - describe flow through the Network
 - ranking / recommendation
 - Centrality of a node:
 - identify influencers (internet → server that projects a lot of info)
 - Discover groups / clusterings
 - How nodes affect connectivity / flow
- ↳ find network metrics to characterize things in a mathematical way

Network Analysis = find a model of the network

- adding / removing nodes & edges = generation
- The state @ given point is the stochastic result of these processes
- model characteristics from prev. ↑ by modeling state of Graph
 - ↳ can we model its creation?
- ↳ find dist. of nodes @ any given point

Random Graph Model

- 1) N nodes M edges $G(N, M) = \{G = (V, E) \mid |V| = N, |E| = M\}$
 - ↳ all graphs w/ set $\# N$ and E
 - ↳ pick a graph randomly from set of graphs

→ distribute M edges uniformly

(Ex) $G(3, 2) = \{ \text{graph 1}, \text{graph 2}, \text{graph 3} \} \rightarrow$ pick a graph w/ prob. $1/3$

In general, probability w/ which you pick graph from $G(N, M)$?

aka $\frac{1}{C(\frac{N(N-1)}{2}, M)}$ ← $p = \left(\frac{C(N, 2)}{M} \right)^{-1}$

→ # max possible edges?
→ how many ways are there to choose M from all possible N

inverse of N "choose" 2

- 2) Let $G(N, p)$ be generated ... randomly connect nodes w/ prob. p , independently
- What is $G(N, p)$ as a function of M ? Note: $\frac{N(N-1)}{2}$ Bernoulli trials inserting edges w/ prob. p

generate graphs randomly

$$f_{G(N, M)} = p^M (1-p)^{\frac{N(N-1)}{2} - M}$$

N "choose" 2

↳ both methods related...

$$G(N, p) \text{ (constant)} \text{ given } M \text{ edges} = G(N, M)$$

- see slides for PROOF - what's the

$$P(G(N, p) \cong E_{G(N, p)} | | = M)$$

- what's the distribution of degree of Nodes?

$$P(\deg(v) = k) = \binom{N-1}{k} p^k (1-p)^{N-1-k} \leftarrow \text{Binomial distribution}$$

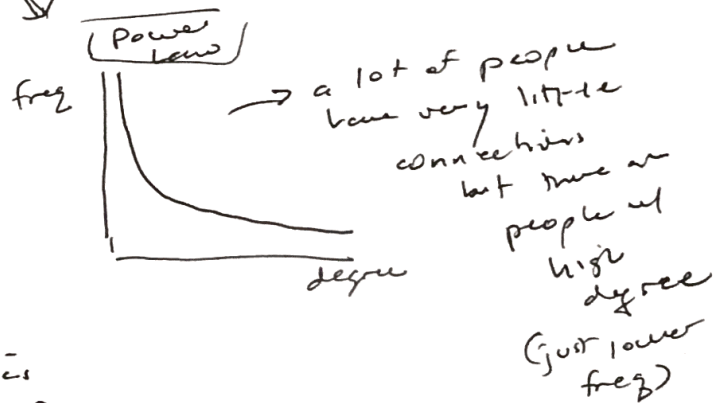
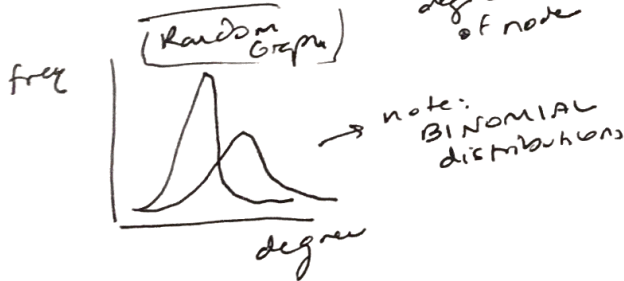
As $N \rightarrow \infty$ (Np constant), converges to Poisson Dist.

Q] - Expected # of connections / degree for a node in this Graph?

$\sim Np$ — does this make sense?
 $N = \# \text{ nodes}$
 think: social networks → what does Bin ^{distr} _{like} look like?
 (means on avg. everyone has similar # connections)
 - NOT so accurate re: social nets

Social Networks = power law distribution

$$P(k) = C k^{-\alpha} \rightarrow \text{true constants } C \text{ and } \alpha$$



- Always think → what are characteristics graph based on model?

After that,
 - compute things on the graph/model
 to get intuition about the network

Does it accurately describe the situation?

Describing/Comparing Graphs

- define metrics that represent characteristics & ^{compare} ~~describe~~ these
- metrics \rightarrow describe whole graphs
 \rightarrow describe individual nodes

① Diameter = ~~the~~, maximum shortest path between nodes

$$d_{ij} = \text{shortest path between } i \text{ and } j \quad \text{Diam}(G) = \max_{i,j} d_{ij}$$

"Small world phenomenon" (for example)

\hookrightarrow only takes few jumps to know someone famous

Q]



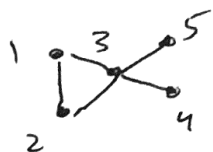
max diameter = ~~2~~ 2

② Clustering coefficient

$$C = \frac{\# \text{ triangles}}{\# \text{ triplets}}$$

\nearrow closed triplet
 \nearrow 3 nodes connected by 2 edges

defined as being centered on a node



$$C = \frac{\begin{matrix} (1) & (2) & (3) & (4) & (5) \\ 1 & 1 & 1 & 0 & 0 \end{matrix}}{\begin{matrix} 1 & 1 & 6 & 0 & 0 \end{matrix}} = \frac{3}{8}$$

\rightarrow captures # times you don't close the triangle

- doesn't matter path, just what node you start on

- start @ each node, count # triangles, add \Rightarrow numerator

- start @ each node, count # triplets, add \Rightarrow denominator

③ Density $N = \# \text{ nodes}$; $M = \# \text{ edges}$ $\text{density} = \frac{2M}{N(N-1)} \rightarrow \text{max possible edges}$

Q] density of complete graph? 1

$$2M = N(N-1)$$

metrics on Nodes...

Degree centrality → more central a Node is, higher # connections

$$C_{deg}(v) = \text{Deg}(v) \quad \leftarrow \text{the degree}$$

Closeness Centrality → more central = closer to all other nodes

$$C_{close}(v) = \frac{1}{\sum_u d(u,v)}$$

↑ ↑
node #1 node #2

→ Look @ slides & 2. exercises

Harmonic Centrality → $C_h(v) = \sum_{u \neq v} \frac{1}{d(u,v)}$ → need to set $d^{-1}(u,v) = 0$ if no path between u and v
- sum of inverse distances

Betweenness Centrality
→ # times a node acts as a bridge along shortest path to other nodes

$$C_b(v) = \sum_{s \neq v \neq t \in V} \frac{\sigma_{st}(v)}{\sigma_{st}}$$

σ_{st} → total # shortest paths between s and t

$\sigma_{st}(v)$ → # of shortest paths between s and t that go through v

(see code demo)
↳ networkx library in python