IN 1400 - Fundamentals of Databases and Database Design

SQL: QUERIES 2

Week 12

Relational Database Schema

EMPLOYEE

FNAME MINIT LNAME SSN BDATE ADDRESS SEX	SALARY SUPERSSN DNO
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DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE	
				П

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME PN	<u>IUMBER</u>	PLOCATION	DNUM
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WORKS_ON

ESSN	PNO	HOURS

DEPENDENT

ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

EMPLOYEE	FNAME	MINIT	LNAME	<u>SSN</u>	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
	John	В	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	М	30000	333445555	5
	Franklin	Т	Wong	333445555	1955-12-08	638 Voss, Houston, TX	М	40000	888665555	5
	Alicia	J	Zelaya	999887777	1968-07-19	3321 Castle, Spring, TX	F	25000	987654321	4
	Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
	Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	М	38000	333445555	5
	Joyce	Α	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
	Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	М	25000	987654321	4
	James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	М	55000	null	1

Populated Database

DEPARTMENT	DEPARTMENT DNAME		MGRSSN	MGRSTARTDATE
	Research	5	333445555	1988-05-22
	Administration	4	987654321	1995-01-01
	Headquarters	1	888665555	1981-06-19

DNUMBER	DLOCATION
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

WORKS_ON ESSN PNO HOURS 123456789 1 32.5 123456789 2 7.5 666884444 3 40.0 453453453 1 20.0 333445555 2 10.0 333445555 3 10.0 333445555 10 10.0 333445555 20 10.0 999887777 30 30.0 999887777 10 10.0 987987987 10 35.0 987654321 30 20.0 987654321 20 15.0 888665555 20 null				
123456789 2 7.5 666884444 3 40.0 453453453 1 20.0 453453453 2 20.0 333445555 2 10.0 333445555 3 10.0 333445555 10 10.0 333445555 20 10.0 999887777 30 30.0 999887777 10 10.0 987987987 10 35.0 987987987 30 5.0 987654321 30 20.0 987654321 20 15.0	WORKS_ON	<u>ESSN</u>	<u>PNO</u>	HOURS
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987654321 30 20.0 987654321 20 15.0		987987987	10	35.0
987654321 20 15.0		987987987	30	5.0
33.33.132.		987654321	30	20.0
888665555 20 null		987654321	20	15.0
		888665555	20	null

PROJECT	PNAME	<u>PNUMBER</u>	PLOCATION	DNUM
	ProductX	1	Bellaire	5
	ProductY	2	Sugarland	5
	ProductZ	3	Houston	5
	Computerization	10	Stafford	4
	Reorganization	20	Houston	1
	Newbenefits	30	Stafford	4

DEPT_LOCATIONS

DEPENDENT	ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP
	333445555	Alice	F	1986-04-05	DAUGHTER
	333445555	Theodore	М	1983-10-25	SON
	333445555	Joy	F	1958-05-03	SPOUSE
	987654321	Abner	М	1942-02-28	SPOUSE
	123456789	Michael	М	1988-01-04	SON
	123456789	Alice	F	1988-12-30	DAUGHTER
	123456789	Elizabeth	F	1967-05-05	SPOUSE

SET OPERATIONS

- SQL has directly incorporated some set operations
- There is a union operation (UNION), and in some versions of SQL there are set difference (MINUS) and intersection (INTERSECT) operations
- The resulting relations of these set operations are sets of tuples; duplicate tuples are eliminated from the result
- The set operations apply only to union compatible relations; the two relations must have the <u>same</u> attributes and the <u>attributes must appear in the same</u> order

SET OPERATIONS (cont.)

 Query 9: Make a list of all project numbers for projects that involve an employee whose last name is 'Smith' as a worker or as a manager of the department that controls the project.

```
Q9: (SELECT PNAME
FROM PROJECT, DEPARTMENT, EMPLOYEE
WHERE DNUM=DNUMBER AND MGRSSN=SSN AND LNAME='Smith')
UNION
(SELECT PNAME
FROM PROJECT, WORKS_ON, EMPLOYEE
WHERE PNUMBER=PNO AND ESSN=SSN AND
LNAME='Smith')
```

NESTING OF QUERIES

- A complete SELECT query, called a nested query, can be specified within the WHERE-clause of another query, called the outer query
- Many of the previous queries can be specified in an alternative form using nesting
- Query 10: Retrieve the name and address of all employees who work for the 'Research' department.

```
Q10: SELECT FNAME, LNAME, ADDRESS
FROM EMPLOYEE
WHERE DNO IN (SELECT DNUMBER
FROM DEPARTMENT
WHERE DNAME='Research')
```

NESTING OF QUERIES (cont.)

- The nested query selects the number of the 'Research' department
- The outer query select an EMPLOYEE tuple if its DNO value is in the result of either nested query
- The comparison operator IN compares a value v with a set (or multi-set) of values V, and evaluates to TRUE if v is one of the elements in V
- In general, we can have several levels of nested queries
- A reference to an unqualified attribute refers to the relation declared in the innermost nested query
- In this example, the nested query is not correlated with the outer query

CORRELATED NESTED QUERIES

- If a condition in the WHERE-clause of a nested query references an attribute of a relation declared in the outer query, the two queries are said to be correlated
- The result of a correlated nested query is different for each tuple (or combination of tuples) of the relation(s) the outer query
- Query 11: Retrieve the name of each employee who has a dependent with the same first name as the employee.

```
Q11: SELECT E.FNAME, E.LNAME
FROM EMPLOYEE AS E
WHERE E.SSN IN (SELECT ESSN
FROM DEPENDENT
WHERE ESSN=E.SSN AND
E.FNAME=DEPENDENT_NAME)
```

CORRELATED NESTED QUERIES (cont.)

- In Q11, the nested query has a different result for each tuple in the outer query
- A query written with nested SELECT... FROM... WHERE... blocks and using the = or IN comparison operators can always be expressed as a single block query. For example, Q11 may be written as in Q11A

```
Q11A: SELECT E.FNAME, E.LNAME
FROM EMPLOYEE E, DEPENDENT D
WHERE E.SSN=D.ESSN AND E.FNAME=D.DEPENDENT_NAME
```

EXPLICIT SETS

- It is also possible to use an explicit (enumerated) set of values in the WHERE-clause rather than a nested query
- Query 14: Retrieve the social security numbers of all employees who work on project number 1, 2, or 3.

```
Q14: SELECT DISTINCT ESSN FROM WORKS_ON WHERE PNO IN (1, 2, 3)
```

NULLS IN SQL QUERIES

- SQL allows queries that check if a value is NULL (missing or undefined or not applicable)
- SQL uses IS or IS NOT to compare NULLs because it considers each NULL value distinct from other NULL values, so equality comparison is not appropriate.
- Query 15: Retrieve the names of all employees who do not have supervisors.

```
Q15: SELECT FNAME, LNAME
FROM EMPLOYEE
WHERE SUPERSSN IS NULL
```

 Note: If a join condition is specified, tuples with NULL values for the join attributes are not included in the result

AGGREGATE FUNCTIONS

- Include COUNT, SUM, MAX, MIN, and AVG
- Query 17: Find the maximum salary, the minimum salary, and the average salary among all employees.

```
Q17: SELECT MAX(SALARY),
MIN(SALARY), AVG(SALARY)
FROM EMPLOYEE
```

 Some SQL implementations may not allow more than one function in the SELECT-clause

AGGREGATE FUNCTIONS (cont.)

 Query 18: Find the maximum salary, the minimum salary, and the average salary among employees who work for the 'Research' department.

```
Q18: SELECT MAX(SALARY), MIN(SALARY),

AVG(SALARY)

FROM EMPLOYEE, DEPARTMENT

WHERE DNO=DNUMBER AND DNAME='Research'
```

AGGREGATE FUNCTIONS (cont.)

 Queries 19 and 20: Retrieve the total number of employees in the company (Q19), and the number of employees in the 'Research' department (Q20).

Q19: SELECT COUNT (*)

FROM EMPLOYEE

Q20: SELECT COUNT (*)

FROM EMPLOYEE, DEPARTMENT

WHERE DNO=DNUMBER AND DNAME='Research'

SUBSTRING COMPARISON

- The LIKE comparison operator is used to compare partial strings
- Two reserved characters are used: '%' (or '*' in some implementations) replaces an arbitrary number of characters,
- and '_' replaces a single arbitrary character

SUBSTRING COMPARISON (cont.)

 Query 24: Retrieve all employees whose address is in Houston, Texas. Here, the value of the ADDRESS attribute must contain the substring 'Houston,TX'.

Q24: SELECT FNAME, LNAME

FROM EMPLOYEE

WHERE ADDRESS LIKE '%Houston, TX%'

SUBSTRING COMPARISON (cont.)

Query 25: Retrieve all employees who were born during the 1950s. Here, '5' must be the 9th character of the string (e.g. 18-11-1950), so the BDATE value is '_____5_', with each underscore as a place holder for a single arbitrary character.

Q25: SELECT FNAME, LNAME FROM EMPLOYEE

WHERE BDATE LIKE '_____5_

 The LIKE operator allows us to get around the fact that each value is considered atomic and indivisible; hence, in SQL, character string attribute values are not atomic

ARITHMETIC OPERATIONS

- The standard arithmetic operators '+', '-'. '*', and '/' (for addition, subtraction, multiplication, and division, respectively) can be applied to numeric values in an SQL query result
- Query 26: Show the effect of giving all employees who work on the 'ProductX' project a 10% raise.

```
Q26: SELECT FNAME, LNAME, 1.1*SALARY
FROM EMPLOYEE, WORKS_ON, PROJECT
WHERE SSN=ESSN AND PNO=PNUMBER AND
PNAME='ProductX'
```

Joined Relations

- Can specify a "joined relation" in the FROM-clause
- Looks like any other relation but is the result of a join
- Allows the user to specify different types of joins (regular "theta" JOIN, NATURAL JOIN, LEFT OUTER JOIN, RIGHT OUTER JOIN, CROSS JOIN, etc)

Joined Relations (cont.)

Examples:

```
Q15: SELECT E.FNAME, E.LNAME, S.FNAME, S.LNAME
FROM EMPLOYEE E S
WHERE E.SUPERSSN=S.SSN
```

can be written as:

```
Q15A: SELECT E.FNAME, E.LNAME, S.FNAME, S.LNAME FROM (EMPLOYEE E LEFT OUTER JOIN EMPLOYEE S ON E.SUPERSSN=S.SSN)
```

Joined Relations (cont.)

Q16: SELECT FNAME, LNAME, ADDRESS

FROM EMPLOYEE, DEPARTMENT

WHERE DNAME='Research' AND DNUMBER=DNO

could be written as:

Q16A: SELECT FNAME, LNAME, ADDRESS

FROM (EMPLOYEE JOIN DEPARTMENT

ON DNUMBER=DNO)

WHERE DNAME='Research'

Do not use ON clause

or as:

Q16B: SELECT FNAME, LNAME, ADDRESS

FROM (EMPLOYEE NATURAL JOIN DEPARTMENT

WHERE DNAME='Research'

GROUPING

- In many cases, we want to apply the aggregate functions to subgroups of tuples in a relation
- Each subgroup of tuples consists of the set of tuples that have the same value for the grouping attribute(s)
- The function is applied to each subgroup independently
- SQL has a GROUP BY-clause for specifying the grouping attributes, which must also appear in the SELECT-clause

GROUPING (cont.)

 Query 21: For each department, retrieve the department number, the number of employees in the department, and their average salary.

```
Q21: SELECT DNO, COUNT (*), AVG (SALARY)
FROM EMPLOYEE
GROUP BY DNO
```

- In Q21, the EMPLOYEE tuples are divided into groups--each group having the same value for the grouping attribute DNO
- The COUNT and AVG functions are applied to each such group of tuples separately
- The SELECT-clause includes only the grouping attribute and the functions to be applied on each group of tuples
- A join condition can be used in conjunction with grouping

GROUPING (cont.)

 Query 22: For each project, retrieve the project number, project name, and the number of employees who work on that project.

Q22: SELECT PNUMBER, PNAME, COUNT (*)

FROM PROJECT, WORKS_ON

WHERE PNUMBER=PNO

GROUP BY PNUMBER, PNAME

 In this case, the grouping and functions are applied after the joining of the two relations

THE HAVING-CLAUSE

 Sometimes we want to retrieve the values of these functions for <u>only those groups that satisfy certain</u> <u>conditions</u>

 The HAVING-clause is used for specifying a selection condition on groups (rather than on individual tuples)

THE HAVING-CLAUSE (cont.)

 Query 23: For each project on which more than two employees work, retrieve the project number, project name, and the number of employees who work on that project.

```
Q23: SELECT PNUMBER, PNAME, COUNT(*)
FROM PROJECT, WORKS_ON
WHERE PNUMBER=PNO
GROUP BY PNUMBER, PNAME
HAVING COUNT (*) > 2
```

ORDER BY

- The **ORDER BY** clause is used to sort the tuples in a query result based on the values of some attribute(s).
- Query 27: Retrieve a list of employees and the projects each works in, ordered by the employee's department, and within each department ordered alphabetically by employee last name.
 - Q27: SELECT DNAME, LNAME, FNAME, PNAME

 FROM DEPARTMENT, EMPLOYEE, WORKS_ON, PROJECT

 WHERE DNUMBER=DNO AND SSN=ESSN AND PNO=PNUMBER

 ORDER BY DNAME, LNAME

ORDER BY (cont.)

- The default order is in ascending order of values
- We can specify the keyword **DESC** if we want a descending order; the keyword **ASC** can be used to explicitly specify ascending order, even though it is the default

Insert With Select

 To copy all columns from one table to another, existing table

```
INSERT INTO table2
SELECT * FROM table1;
```

INSERT WITH SELECT

 Example: Suppose we want to create a temporary table that has the name, number of employees, and total salaries for each department. A table DEPTS_INFO is created by U3A, and is loaded with the summary information retrieved from the database by the query in U3B.

```
CREATE TABLE DEPTS_INFO
U3A:
                (DEPT_NAME VARCHAR(10),
                NO_OF_EMPS INTEGER,
                 TOTAL_SAL INTEGER);
U3B:
         INSERT INTO DEPTS INFO (DEPT NAME,
               NO OF EMPS, TOTAL_SAL)
                      DNAME, COUNT (*), SUM (SALARY)
         SELECT
                      DEPARTMENT, EMPLOYEE
         FROM
                      DNUMBER=DNO
         WHERE
         GROUP BY
```

DNAME;

INSERT WITH SELECT

- Note: The DEPTS_INFO table may not be up-to-date if we change the tuples in either the DEPARTMENT or the EMPLOYEE relations after issuing U3B.
- We have to create a <u>view</u> to keep such a table up to date.

DELETE WITH SELECT

- Examples:
- Delete employees in Research department

```
U4C: DELETE FROM EMPLOYEE
WHERE DNO IN
(SELECT DNUMBER
FROM DEPARTMENT
WHERE DNAME='Research')
```

UPDATE WITH SELECT

<u>Example:</u> Give all employees in the 'Research' department a 10% raise in salary.

U6: UPDATE EMPLOYEE

SET SALARY = SALARY *1.1
WHERE DNO IN (SELECT DNUMBER

FROM DEPARTMENT

WHERE DNAME='Research')

- In this request, the modified SALARY value depends on the original SALARY value in each tuple
- The reference to the SALARY attribute on the right of = refers to the old SALARY value before modification
- The reference to the SALARY attribute on the left of = refers to the new SALARY value after modification

Summary of SQL Queries

 A query in SQL can consist of up to six clauses, but only the first two, SELECT and FROM, are mandatory. The clauses are specified in the following order:

Summary of SQL Queries (cont.)

- The SELECT-clause lists the attributes or functions to be retrieved
- The FROM-clause specifies all relations (or aliases) needed in the query but not those needed in nested queries
- The WHERE-clause specifies the conditions for selection and join of tuples from the relations specified in the FROM-clause
- GROUP BY specifies grouping attributes
- HAVING specifies a condition for selection of groups
- ORDER BY specifies an order for displaying the result of a query
- A query is evaluated by first applying the WHERE-clause, then GROUP BY and HAVING, and finally the SELECT-clause