# CheckMPI-Documentation

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# Contents

1	Intr	oduction	4
2	$\mathbf{Cod}$	e structure	5
3		Case Validation	<b>5</b>
4	$\operatorname{Cod}$	e versions	7
	4.1	Step0	8
	4.2	Step1	9
	4.3	Step2	10
	4.4	Step3	11
	4.5	Step4	13
	4.6	Step5	14
	4.7	Step6	15
	4.8	Step7	16
	4.9	Step8	17
	4.10	Step9	18
5	Figu	ires	19
	5.1	Compilers options	20
	5.2	TEST Case Options	20
	5.3	Licence	20
6	Rep	ository Structure	20
7	Hov	to compile	21
	7.1	Further Pre-processing Flags	22
	7.2		22
8	Con	apiler options	22
	8.1		22
	8.2	1 1	22

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### 1 Introduction

This code was developed to exploit MPI comms.

It implements a simple 3D domain decomposition for a periodic box. It was intended to check different communication patterns and their issues. In details we will explore

- sendrecv
- mpi datatype
- handmade packing/unpaking data
- OpenACC directives
- Cuda Aware mpi
- $\bullet\,$  Non-blocking comms
- ullet Communication-Computation overlap

#### 2 Code structure

#### 3 Test Case

The test case is a *sinthetic* test case. It is a:

- 3D flow
- with periodic boundary condition
- 3 fields to propagate, initialized with a complete sinusoidal for each task
- All the tests are done using single precision.
- The propagation is
  - from left to right
  - from rear to front
  - from bottom to right

Note that no flops are performed, only "rigid" data movement.

```
do k = n, 1, -1
  do j = m, 1, -1
      do i = l, 1, -1
      field1(i,j,k) = field1(i-1,j,k)
      field2(i,j,k) = field2(i,j-1,k)
      field3(i,j,k) = field3(i,j,k-1)
      end do
  end do
end do
```

The input file used for the test is the following:

```
%parameters
lx = 600
ly = 600
lz = 600
proc_x = 2
proc_y = 2
proc_z = 2
itfin=15000
icheck=1510 /
```

#### 3.1 Validation

The initial condition, a sinusoidal wave, it the same for each task, so the profile are the *the same* independently for all the tests: this means that the UNIX md5sum command should return the same hash.

```
md5sum RUN_STEP*/prof_i*0.dat
3cdb4eef4d04a99e786046dead098237 RUN_STEP0/prof_i.000000.dat
3cdb4eef4d04a99e786046dead098237 RUN_STEP1/prof_i.000000.dat
3cdb4eef4d04a99e786046dead098237 RUN_STEP2/prof_i.000000.dat
3cdb4eef4d04a99e786046dead098237 RUN_STEP3/prof_i.000000.dat
```

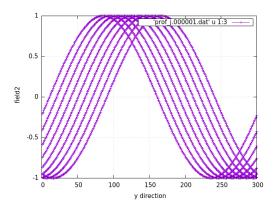


Figure 1: Validation for variable field2.

3cdb4eef4d04a99e786046dead098237 RUN\_STEP4/prof\_i.000000.dat 3cdb4eef4d04a99e786046dead098237 RUN\_STEP5/prof\_i.000000.dat

Table 1: GPU/CPU support for different versions

version	CPU	GPU	develop ver.
Step0	yes	no	-
Step1	yes	no	yes
Step2	yes	no	-
Step3	yes*	yes	-
Step4	yes*	yes	-
Step5	yes*	yes	yes
Step6	yes*	yes	yes
Step7	yes*	yes	yes
Step8	yes*	yes	-
Step9	yes*	yes	-

### 4 Code versions

At the moment 9 different versions are available. They can be selected at compile time via pre-processing flag syntax.

The syntax, e.g. to choose Step6 version, is:

#### make STEP6=1

In detail, these are the different options to choose from:

- ullet Step0:  $3 \times 2$  sendrecv per task/timestep and mpi datatype
- Step1:  $3 \times 2$  sendrecv and explicit pack-unpack (only for x direction)
- Step2: 2 sendrecv and explicit pack-unpack (only all directions)
- Step3: 2 sendrecv and explicit pack-unpack (only all directions) with OpenACC
- Step4: 2 sendrecv and explicit pack-unpack (only all directions) with OpenACC and CUDAWARE MPI call
- Step6: like Step5 but with isend and recv for all the three directions (development step)
- Step7: like Step6 but with isend and irecv for all the three directions (development step).
- $\bullet$  Step8: communication-computation overlap.
- Step9: communication-computation overlap with async clause.

#### 4.1 Step0

It is the starting point version: it uses one different single sendrecv call for each population to propagate information. It calls bcond\_comm\_step0 subroutine. It means 3 calls to sendrecv for each direction (+/-) and each decomposition along x, y, z.

```
tag = 04
call mpi_sendrecv(field1(0,0,n), 1, xyplane, up(2), tag,
                                                              &
                  field1(0,0,0), 1, xyplane, down(2), tag,
                                                              &
                  lbecomm, status,ierr)
tag = 06
call mpi_sendrecv(field2(0,0,n), 1, xyplane, up(2), tag,
                                                              &
                  field2(0,0,0), 1, xyplane, down(2), tag,
                                                              &
                  lbecomm, status,ierr)
tag = 07
call mpi_sendrecv(field3(0,0,n), 1, xyplane, up(2), tag,
                                                              &
                  field3(0,0,0), 1, xyplane, down(2), tag,
                  lbecomm, status,ierr)
```

It use explicity mpi datatype

These are the options activated with the command make

```
# default (NVIDIAcompiler)
CC = mpicc
FC = mpifort
FIX = -DPGI
FOPT = -fast -Mcontiguous -Mnodepchk
OPT = -fast -Mcontiguous -Mnodepchk
VER = step0
```

This version is a CPU-only.

### 4.2 Step1

In this version, the data for the x direction is explicitly packed/unpacked using a do loop. It calls bcond\_comm\_step1 subroutine.

These are the options activated with the command make STEP1=1

```
# default (NVIDIAcompiler)
CC = mpicc
FC = mpifort
FIX = -DPGI
FOPT = -fast -Mcontiguous -Mnodepchk
COPT =
OPT = -fast -Mcontiguous -Mnodepchk
ifdef STEP1
FIX=-DSTEP1
VER=step1
endif
```

This version is a CPU-only.

#### 4.3 Step2

In this version, the data for the three directions is explicitly packed/unpacked using a do loop and a single sendrecv is used for all the 3 fields to be propagated It calls bcond\_comm\_step2 subroutine.

```
do k = 0,n+1
        do i = 0,1+1
           bufferYIN(i,k,1)=field1(i,1,k)
           bufferYIN(i,k,2)=field2(i,1,k)
           bufferYIN(i,k,3)=field3(i,1,k)
        enddo
     enddo
!
    call mpi_sendrecv(bufferYIN(0,0,1),msgsizeY,MYMPIREAL,left(2),tag,&
                       bufferYOUT(0,0,1),msgsizeY,MYMPIREAL,right(2),tag,&
                       lbecomm,status,ierr)
     do k = 0,n+1
        do i = 0,1+1
          field1(i,m+1,k)=bufferYOUT(i,k,1)
          field2(i,m+1,k)=bufferYOUT(i,k,2)
           field3(i,m+1,k)=bufferYOUT(i,k,3)
        enddo
     enddo
```

These are the options activated with the command make STEP2=1

```
# default (NVIDIAcompiler)
CC = mpicc
FC = mpifort
FIX = -DPGI
FOPT = -fast -Mcontiguous -Mnodepchk
COPT =
OPT = -fast -Mcontiguous -Mnodepchk
ifdef STEP2
   FIX=-DSTEP2
VER=step2
endif
```

This version is a CPU-only.

### 4.4 Step3

In this version OpenACC directives are use for

• loop inside do\_somethingGPU subroutine pack/unpack fields to propagate

It calls bcond\_comm\_step3 subroutine. For the loop inside do\_somethingGPU subroutine a copy has to be done to exploit GPU parallelism, activated by FAST preprocessing flags.

```
ifdef FAST
!$acc kernels
        do k = 0, n+1
           do j = 0, m+1
              do i = 0, 1+1
                 temp1(i,j,k) = field1(i,j,k)
                 temp2(i,j,k) = field2(i,j,k)
                 temp3(i,j,k) = field3(i,j,k)
              end do
           end do
        end do
!$acc end kernels
!$acc kernels
        do k = 1, n
           do j = 1, m
              do i = 1, 1
                 field1(i,j,k) = temp1(i-1,j,k)
                 field2(i,j,k) = temp2(i,j-1,k)
                 field3(i,j,k) = temp3(i,j,k-1)
              end do
           end do
        end do
!$acc end kernels
#else
```

For the pack/unpack section these are the directive used

```
do k = 0,n+1
    do j = 0,m+1
    field1(0,j,k) = bufferXOUT(j,k,1)
    field2(0,j,k) = bufferXOUT(j,k,2)
    field3(0,j,k) = bufferXOUT(j,k,3)
    enddo
enddo
!$acc end kernels
```

These are the options activated with the command  ${\tt make}$  STEP3=1

```
# default (NVIDIAcompiler)
CC = mpicc
FC = mpifort
FIX = -DPGI
FOPT = -fast -Mcontiguous -Mnodepchk
COPT =
OPT = -fast -Mcontiguous -Mnodepchk
ifdef STEP3
FIX=-DSTEP3 -acc -Minfo=acc -DOPENACC -DFAST
VER=step3
endif
```

To activate the GPU version the compile line is:

```
make STEP3=1 GPUENABLE=1
```

#### 4.5 Step4

In this version OpenACC directives are use for exploit Cuda-aware MPI comms. It calls bcond\_comm\_step4 subroutine.

These are the options activated with the command make STEP4=1

```
# default (NVIDIAcompiler)
CC = mpicc
FC = mpifort
FIX = -DPGI
FOPT = -fast -Mcontiguous -Mnodepchk
COPT =
OPT = -fast -Mcontiguous -Mnodepchk
ifdef STEP4
FIX=-DSTEP4 -acc -Minfo=acc -DOPENACC -DFAST
VER=step4
endif
```

To activate the GPU version the compile line is:

```
make STEP4=1 GPUENABLE=1
```

# 4.6 Step5

This is a development step.

It is like Step4 but using isend and recv subroutines only for the x direction. To activate the GPU version the compile line is:

make STEP5=1 GPUENABLE=1

# 4.7 Step6

This is a development step.

It is like Step6 but with isend and recv subroutines for all the three directions.

To activate the GPU version the compile line is:

make STEP6=1 GPUENABLE=1

# 4.8 Step7

This is a development step.

It is like Step7 but with isend and irecv for all three directions and restructured to allow computation and communication overlap.

It works correctly only if there is more than one task for each direction (e.b. total tasks  $\geq 8$ ).

To activate the GPU version the compile line is:

make STEP7=1 GPUENABLE=1

# 4.9 Step8

In this version, the operations of do\_something is splitted to allow the overlap. It works correctly only if there is more than one task for each direction (e.b. total tasks  $\geq 8$ ).

To activate the GPU version the compile line is:

make STEP8=1 GPUENABLE=1

### 4.10 Step9

In this version, the operations of do\_something is splitted to allow the overlap. It works correctly only if there is more than one task for each direction (e.b. total tasks  $\geq 8$ ). Respect step8 the async clause is used for the kernel loops. This implementation shows his efficiency for more than 64 tasks.

To activate the GPU version the compile line is:

make STEP9=1 GPUENABLE=1

Table 2: Figures for different versions (using single DGX node, with 8 tasks

version	tot	%COLL	%MPI	х	У	z	type
Step0	661"	53%	47%	268"	15.4"	10.6"	CPU
Step1	484"	72%	28%	101"	15.4"	11.0"	CPU
Step2	488"	72%	28%	74"	19.1"	13.9"	CPU
Step7	680"	77%	23%	26.3"	23.8"	23.8"	CPU
Step8	1650"	77%	23%	26.3"	23.8"	23.8"	CPU
Step3	97.5"	16%	82%	26.3"	23.8"	23.8"	GPU
Step4	27.7"	59%	35%	3.52"	1.74"	3.62"	GPU
Step5	26.5"	61%	32%	2.51"	2.04"	3.25"	GPU
Step6	25.4"	64%	29%	2.46"	1.34"	3.02"	GPU
Step7	24.6"	-%	-%	-"	-"	-"	GPU
Step8	24.0"	-%	-%	-"	_'''	_''	GPU
Step9	21.5"	-%	-%	-"	_'''	-"	GPU

Table 3: Figures for scale-up (using Leonardo, with 4 tasks per node

proc_x	proc_y	proc_x	STEP4	STEP8	STEP9
2	2	2	185"	159"	143"
3	3	3	267"	230"	196"
4	4	4	290"	280"	218"
5	5	5	404"	285"	224"
6	6	6	568"	291"	233"
7	7	7	621"	268"	232"
8	8	8	683"	285"	241"

### 5 Figures

All the figures reported in the tables are obtained from file task.000000.log. In tab 2 are reported some performance figures for a 8-task run on DGX. Each task uses a  $300^3$  gridpoint. About 100'000 iterations were performed. In tab. 3 a scale up using version STEP4 and STEP8 are reported with respect different number of tasks. Each task uses a  $300^3$  gridpoint. About 100'000 iterations were performed

#### 5.1 Compilers options

### 5.2 TEST Case Options

#### 5.3 Licence

This code is released under the MIT license. It can be downloaded from https://github.com/gamati01/Check\_MPI.

### 6 Repository Structure

The code is developed to be self-consistent. No external libraries are needed to compile/run the code. The directory structure is the following

- DOC: In this directory some documentation
- RUN: In this directory the executable file will be created (i.e., bgk2d.\*.x) after the compilation step.
- SRC: In this directory all the source files, with the Makefile, are present.
- UTIL: In this directory some utility scripts and some input files are present.
- TEST: In this directory the four test cases, with some scripts, used for the code validation are present, each in a different directory.
- BENCH: In this directory the three benchmark scripts (for single core, single compute node, and GPU) are present, each in a different directory.
- CI: In this directory some script to perform a *fast* check for compilation and execution of all possible configurations.

# 7 How to compile

These are the steps to compile the code:

- 1. Go in the SRC directory.
- Compile the desired configuration: e.g., make serial DOUBLE=1 FUSED=1 LDC=1 GNU=1 for a lid-driven cavity test in double precision, with the fused implementation and GNU compiler.
- 3. If the compilation is successful, the exe file, i.e., bgk2d.serial.x, will be copied in the ../RUN/ directory.
- 4. Go in the RUN directory.
- 5. Copy from the UTIL directory a bgk.input file. For the single core run: ../UTIL/bgk.core.input bgk.input.
- 6. Run the code: ./bgk2d.serial.x.

The code uses dynamic allocation, so, once fixed, the test case does not need to be recompiled if you need to vary the size of the simulation box.

### 7.1 Further Pre-processing Flags

### 7.2 Input/Output

The code will produce the following file, together with some standard output:

### 8 Compiler options

### 8.1 Compiler options for CPU version

- GNU: -Ofast
- INTEL: -03 -xCORE-AVX512 -mtune=skylake-avx512 -assume contiguous\_pointer
- NVIDIA: -02 -Mnodepchk -Mcontiguous
- AMD: -03
- CRAY: -03

### 8.2 Compiler options for GPU version

- GNU: -fopenmp -ftree-parallelize-loops=4
- INTEL: -qopenmp
- NVIDIA:-stdpar=multicore
- AMD: -h thread\_do\_concurrent