Homework Assignment 3

Matthew Tiger

September 19, 2015

Problem 1. Let A and B be square $n \times n$ matrices and $x \in \mathbb{R}^{n \times 1}$ be a column vector. Count the number of multiplications needed to compute (AB)x versus A(Bx). Which one is better for large values of n?

Solution. To determine the number of multiplications necessary to compute the product (AB)x we must find out how many multiplications it takes to compute AB = C and how many multiplications it takes to compute Cx. Since A and B are square $n \times n$ matrices, each entry in their product will require n multiplications. Since AB has n^2 entries, it will take n^3 multiplications to compute AB = C. Now, C is a square $n \times n$ matrix and x is a $n \times 1$ matrix, so each entry in the product Cx will require n multiplications. Since there are n entries in Cx, the product will require n^2 multiplications. Therefore, $n^3 + n^2$ multiplications are necessary to compute (AB)x.

Similarly the number of multiplications necessary to compute A(Bx) is determined by the number of multiplications necessary to compute Bx = D and AD. Since Bx has n entries and it takes n multiplications to compute each entry, it takes n^2 multiplications to compute Bx. For similar reasons it takes n^2 multiplications to compute AD. Therefore, it takes $2n^2$ multiplications to compute A(Bx).

Clearly, it is better to compute A(Bx) for large values of n if the goal is to reduce the number of multiplications necessary to compute the product.

Problem 2. Let A and B be square $n \times n$ upper triangular matrices. Show that C = AB is also upper triangular. How many multiplications are needed to compute C?

Solution. Note that a matrix $A = (a_{ij})$ is upper triangular if every entry of the matrix below the main diagonal is 0, i.e. if $a_{ij} = 0$ when i > j.

If C = AB, then it is clear that by definition $c_{ij} = \sum_{k=1}^{n} a_{ik} b_{kj}$. Thus, the matrix product C is upper triangular if $c_{ij} = 0$ when i > j. So, let's consider c_{ij} for which $1 \le j < i \le n$. Then

$$c_{ij} = \sum_{k=1}^{n} a_{ik} b_{kj} = \sum_{k=1}^{j} a_{ik} b_{kj} + \sum_{k=j+1}^{n} a_{ik} b_{kj}$$
(1)

Now for $1 \le k \le j < i$, the entry $a_{ik} = 0$ since A is upper triangular showing that the left sum in (1) is 0 and for $j < j + 1 \le k \le n$, the entry $b_{kj} = 0$ since B is upper triangular showing that the right sum in (1) is 0. Therefore, $c_{ij} = 0$ for i > j and C is an upper triangular matrix.

To compute the number of multiplications necessary to compute this product, we must determine what the other entries c_{ij} are when $i \leq j$.

When $1 \le i \le j \le n$,

$$c_{ij} = \sum_{k=1}^{n} a_{ik} b_{kj} = \sum_{k=1}^{i-1} a_{ik} b_{kj} + \sum_{k=i}^{j} a_{ik} b_{kj} + \sum_{k=j+1}^{n} a_{ik} b_{kj} = \sum_{k=i}^{j} a_{ik} b_{kj}$$

since $a_{ik} = 0$ for $1 \le k \le i - 1$ due to the fact that A is upper triangular and $b_{kj} = 0$ for $j + 1 \le k \le n$ due to the fact that B is upper triangular. Thus, to compute the entry c_{ij} when $i \le j$, there are j - i + 1 multiplications necessary to compute the entry and 0 multiplications are necessary when i > j.

If $x(c_{ij})$ is the number of multiplications necessary to calculate the entry c_{ij} , then X, the number of computations necessary to calculate the product C, using the above, is given by

$$X = \sum_{j=1}^{n} \sum_{i=1}^{n} x(c_{ij}) = \sum_{j=1}^{n} \sum_{i=1}^{j} x(c_{ij})$$

$$= \sum_{j=1}^{n} \sum_{i=1}^{j} j - i + 1$$

$$= \frac{1}{2} \sum_{j=1}^{n} 2j^{2} - j(j+1) + 2j$$

$$= \frac{1}{2} \sum_{i=1}^{n} j^{2} + j = \frac{n(n+1)(2n+1)}{12} + \frac{n(n+1)}{4} = \frac{n(n+1)(n+2)}{6}$$

Therefore, the number of multiplications necessary to compute the product of two $n \times n$ upper triangular matrices is, in general, n(n+1)(n+2)/6.

Problem 3. Let A be a square $n \times n$ upper triangular matrix. Show that A^{-1} is also upper triangular.

Solution. Suppose a matrix A is an upper triangular matrix. Then

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ 0 & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & a_{nn} \end{bmatrix} = \begin{bmatrix} a_{11} & 0 & \dots & 0 \\ 0 & a_{22} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & a_{nn} \end{bmatrix} + \begin{bmatrix} 0 & a_{12} & \dots & a_{1n} \\ 0 & 0 & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 0 \end{bmatrix}$$

$$= \begin{bmatrix} a_{11} & 0 & \dots & 0 \\ 0 & a_{22} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & a_{nn} \end{bmatrix} \begin{pmatrix} \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{bmatrix} + \begin{bmatrix} 0 & a_{12} & \dots & a_{1n} \\ 0 & 0 & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 0 \end{bmatrix}$$

$$= \Lambda(I + U)$$

where Λ is a diagonal matrix consisting of the main diagonal of A, I is the identity matrix, and U is an upper triangular nilpotent matrix.

Using this definition of A, we can see easily that $A^{-1} = (I+U)^{-1}\Lambda^{-1}$. Let $B = I - U + U^2 + \cdots + (-1)^n U^n$. Then it is easy to see that B commutes with I + U and

$$(I+U)B = (I+U)(I-U+U^{2}+\cdots+(-1)^{n}U^{n})$$

$$= (I+U)\sum_{i=0}^{n}(-1)^{i}U^{i}$$

$$= \sum_{i=0}^{n}(-1)^{i}U^{i} + \sum_{i=0}^{n}(-1)^{i}U^{i+1}$$

$$= I + \sum_{i=1}^{n}(-1)^{i}U^{i} + \sum_{i=1}^{n}(-1)^{i-1}U^{i} + U^{n+1}$$

$$= I + \sum_{i=1}^{n}((-1)^{i}+1^{i})U^{i} + U^{n+1} = I + U^{n+1}.$$
(2)

Now $(I+U)B=I+U^{n+1}=I$ since U is a nilpotent matrix. Thus, $(I+U)^{-1}=B=I-U+U^2+\cdots+(-1)^nU^n$. Note that since the sum of two upper triangular matrices is an upper triangular matrix and the product of two upper triangular matrices is an upper triangular matrix, $(I+U)^{-1}=I-U+U^2+\cdots+(-1)^nU^n$ is an upper triangular matrix since I and U are upper triangular matrices.

Therefore, $A^{-1}=(I+U)^{-1}\Lambda^{-1}$ must be upper triangular since Λ^{-1} is an upper triangular matrix.