

LinkedList - 1

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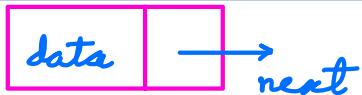
Issues with Array

Continuous memory allocation.

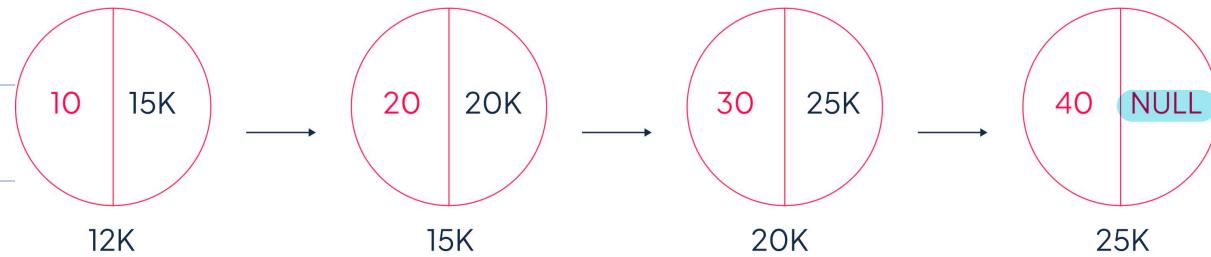


linked list → Linear Data Structure

Non-continuous memory allocation.



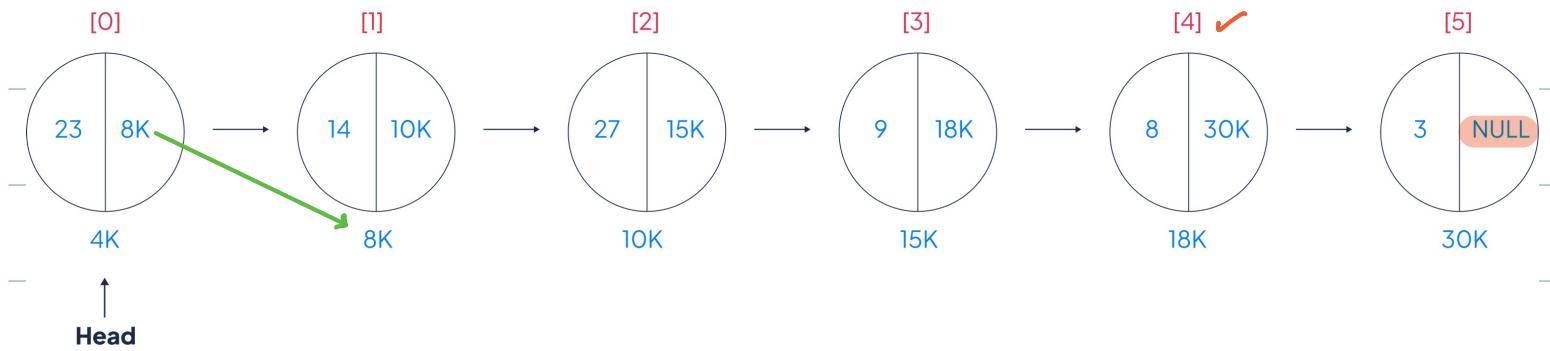
LinkedList



Head

```
class Node {  
    int val;  
    Node next;  
    public Node(int v);  
        this.val = v;  next = null;  
    }  
}
```

How to access the element at k th.idx in a linked-list



$idx = 4 \rightarrow \underline{8}$

if (Head == null) return -1

temp = Head

for i → 1 to K {

 temp = temp.next

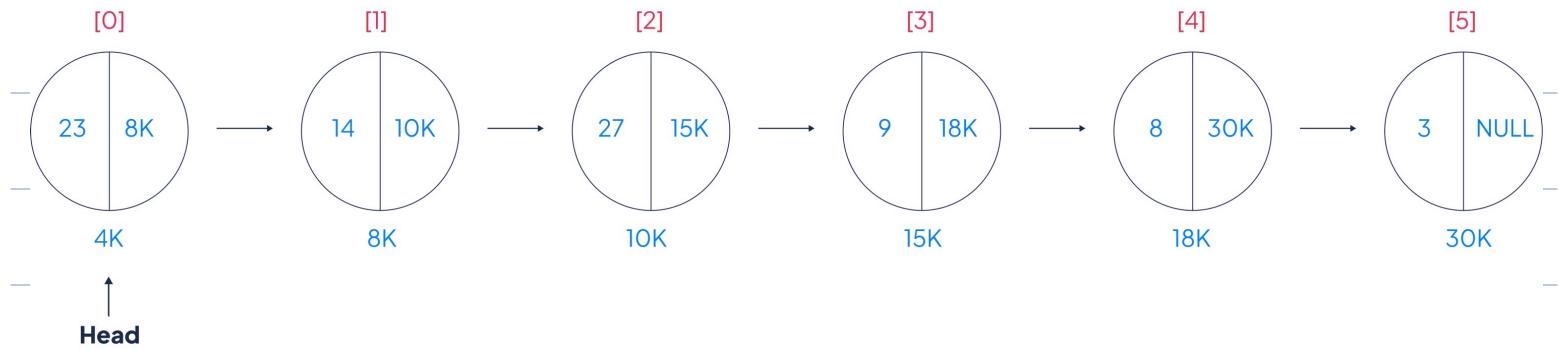
 if (temp == null) return -1

}

return temp.val

TC = $O(K)$

Search in linked-list



search(3) → true
X

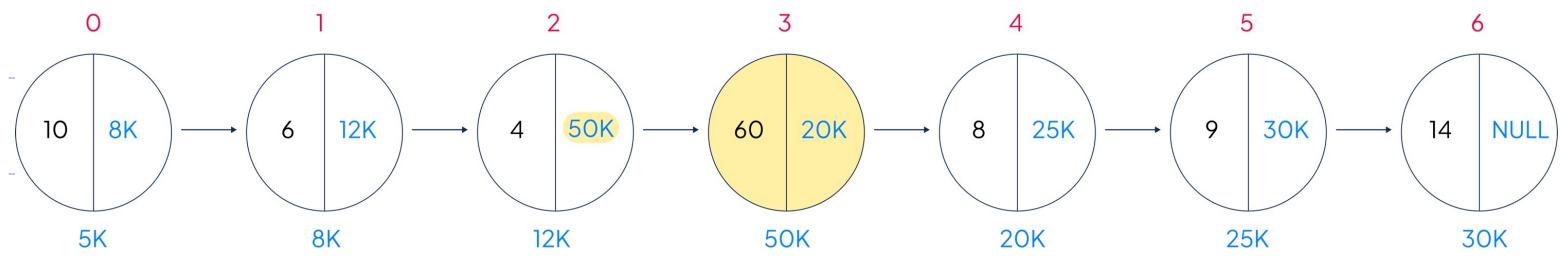
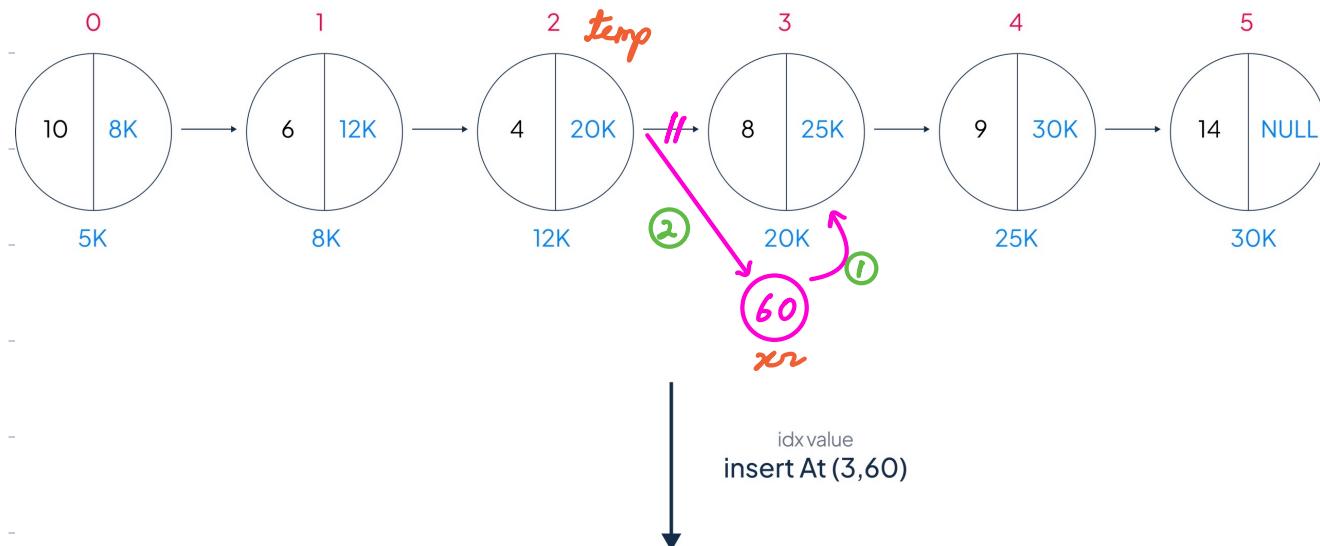
temp = Head

```
while (temp != null) {  
    if (temp.val == X) return true  
    temp = temp.next  
}
```

return false $T.C = O(N)$ // Linear Search

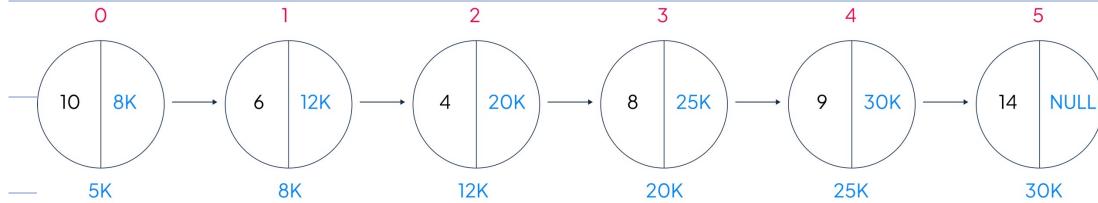
Binary Search → Not possible \because we cannot jump to middle element.

Insertion in linked-list at Kth index



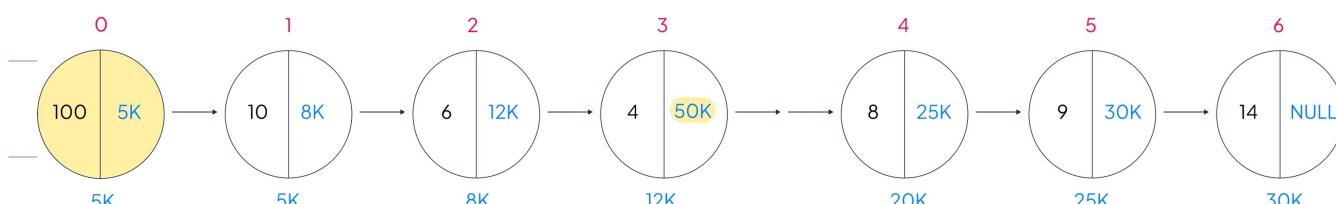
Edge-cases

If $K = 0$



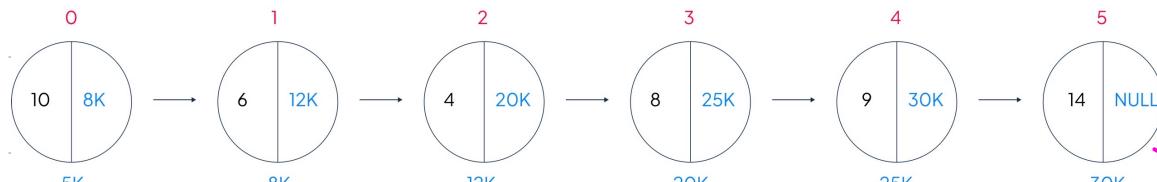
Head

insert (100,0)



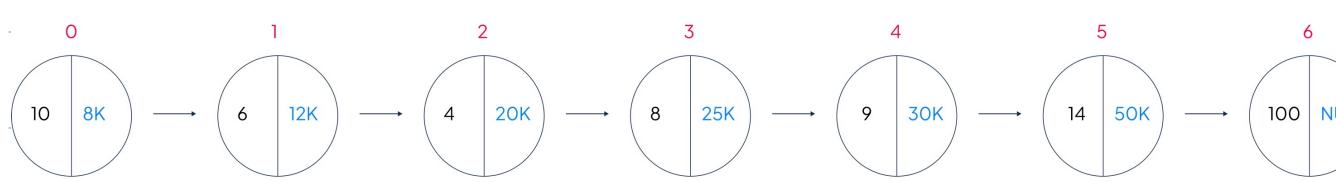
Head

If $K = N$



Head

insert (100,6)



Head



|| insert (x, k)

Node xr = new Node (x)

if (k == 0) {

 xr.next = Head

 Head = xr

 return Head

}

temp = Head

for i → 1 to (k-1) {

 temp = temp.next

}

xr.next = temp.next

temp.next = xr

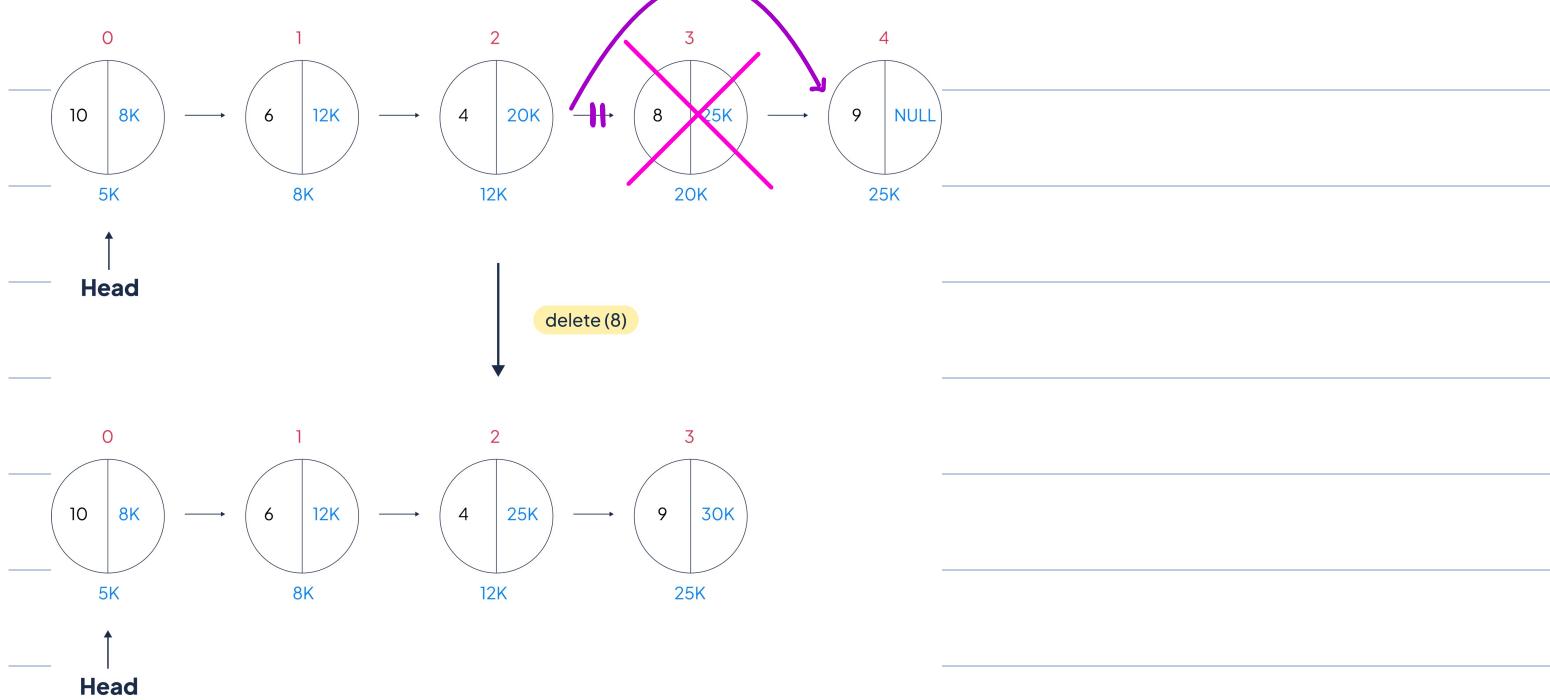
return Head

TC = O (K)

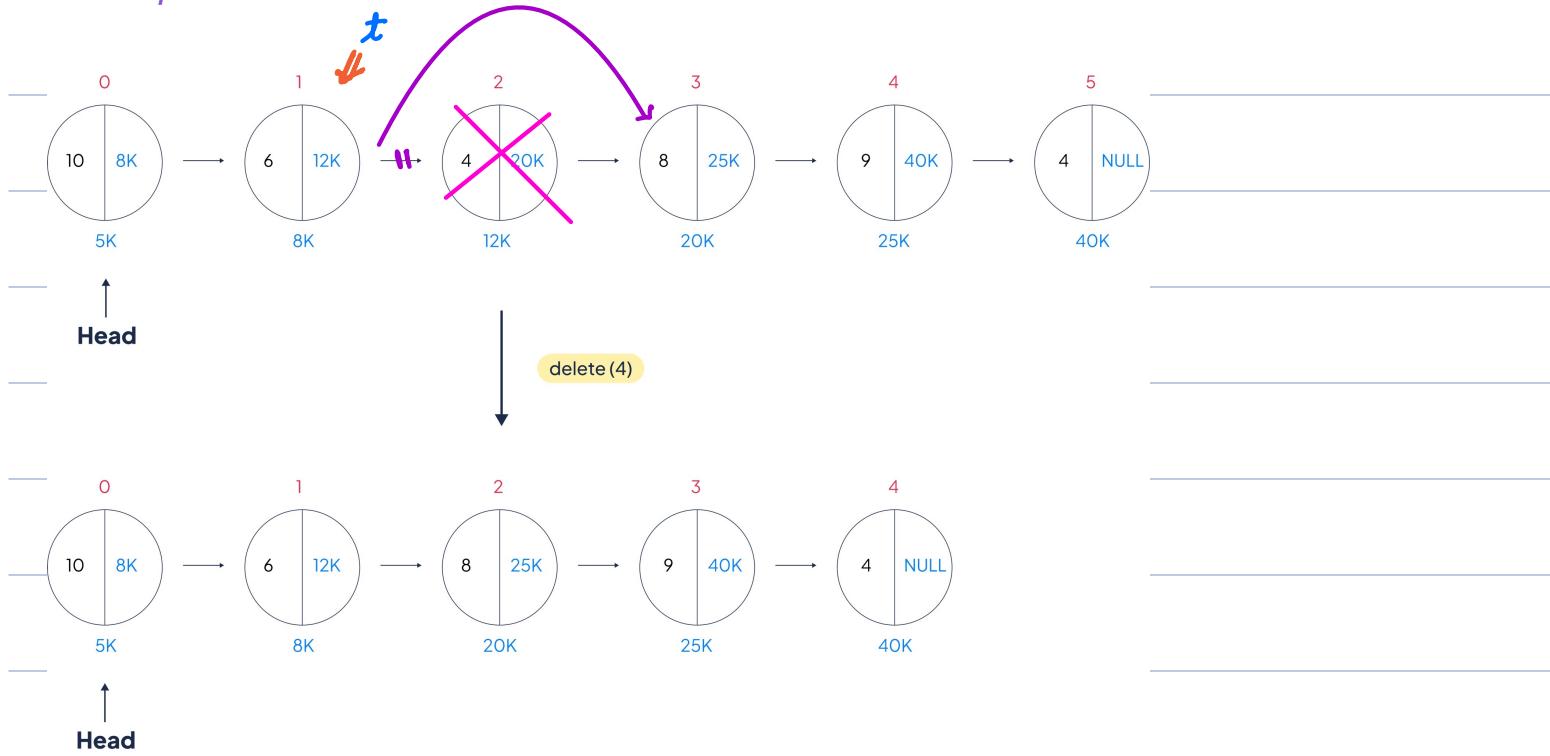
Deletion in linked-list

Delete the first occurrence of value X in the given linked-list. If element is not present, leave as is.

Example 1



Example 2

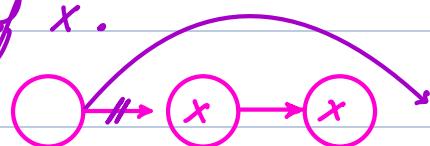




```
if ( Head == null ) return null  
if ( Head.val == x ) { Head = Head.next  
    return Head }  
  
temp = Head  
  
while ( temp.next != null ) {  
    if ( temp.next.val == x ) {  
        temp.next = temp.next.next  
        break  
    }  
  
    temp = temp.next  
}  
  
return Head
```

$TC = \underline{O(N)}$

H.W \rightarrow Delete all occurrences of x .





Scenario

Comment

OnePlus has a lineup of N mobile phones ready in their manufacturing line.

It has detected a defect in one of their phone models during production.

They have decided to recall all phones of the defective model from their manufacturing line. Your task is to help OnePlus remove all defective phones from their production lineup efficiently

Problem

You are given a linked list A of N nodes where each node represents a specific model type of a OnePlus mobile phone in the manufacturing line.

Each node contains an integer representing the model number of the phone. You will also be given an integer B which represents the model number of the defective phone that needs to be removed.

Your goal is to remove all nodes (phones) from the linked list that have the model number B and return the modified linked list representing the updated manufacturing line.

Delete all occurrences of B from list A .

if (Head == null) return null

while (Head != null && Head.val == B)

Head = Head.next

// B = 10

temp = Head



while (temp.next != null) {

if (temp.next.val == B)

temp.next = temp.next.next

else

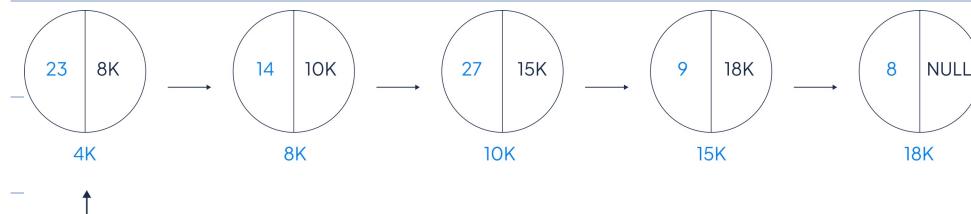
temp = temp.next

}

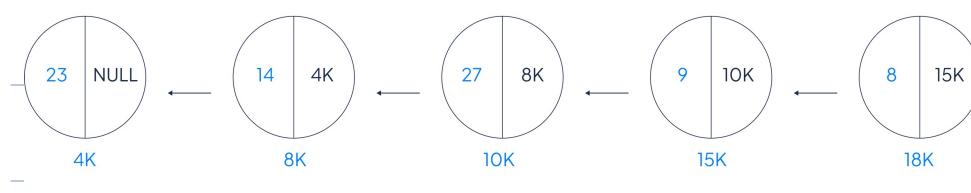
return Head

TC = $O(N)$

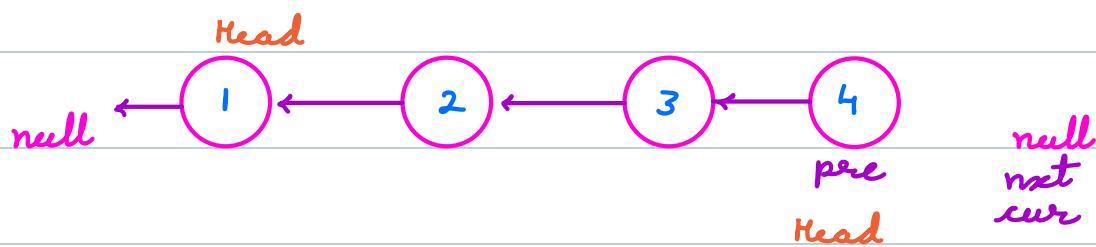
Reverse a linked-list



Reverse



Head



$nxt = cur.next$

$cur.next = pre \checkmark$

$pre = cur$

$cur = nxt$





cur = Head

pre = null

while (cur != null) {

 nxt = cur.next

 cur.next = pre

 pre = cur

 cur = nxt

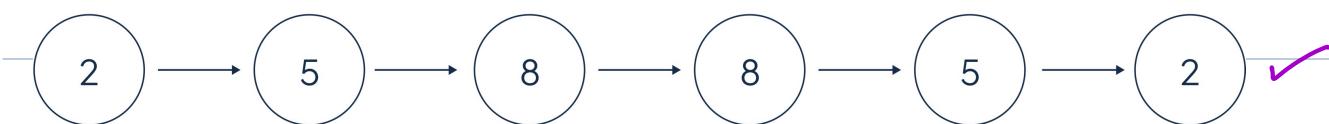
}

TC = $O(N)$ SC = $O(1)$

return pre // head

- Check if a linked-list is a palindrome or not

$\xrightarrow{\quad} = \xleftarrow{\quad}$
"racecar"



1 \rightarrow null ✓

Sol 1 → 1) create copy of list

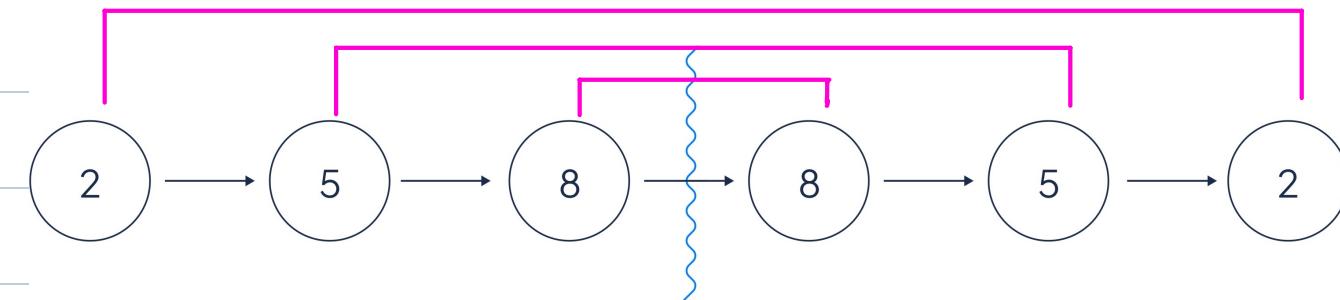
2) Reverse the copy

3) Check if reverse of copy & original list is same.

TC = O(N)

SC = O(N) → copy





Reverse 2nd half ✓



- 1) Find middle element & divide the list. ✓
Find size & travel half. ✓
- 2) Reverse second half. ✓ $TC = O(N)$
- 3) Compare first half with reversed second half. ✓

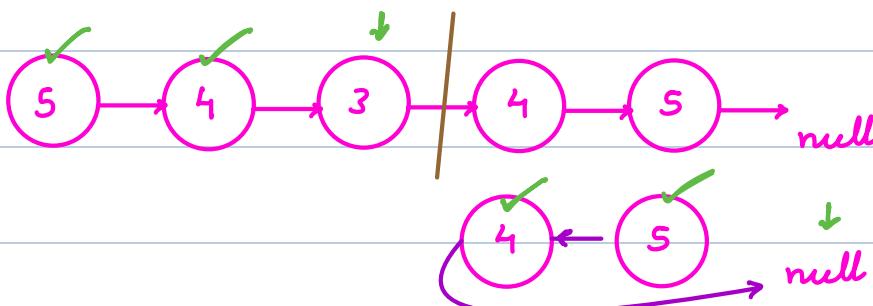
$TC = O(N)$ $SC = O(1)$

If first half and reversed second half are equal

→ return true;

otherwise

→ return false;



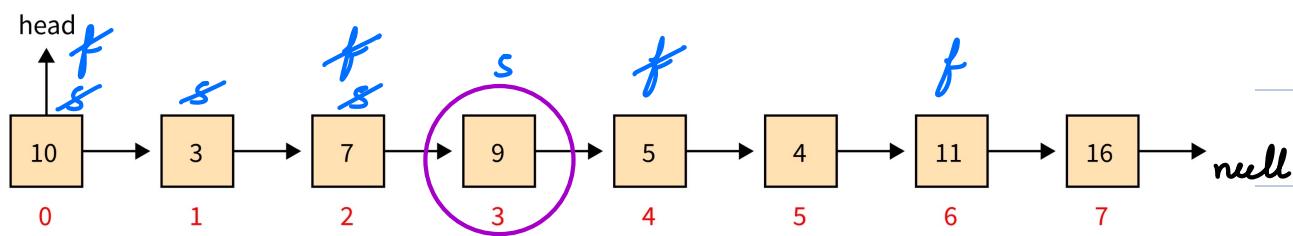
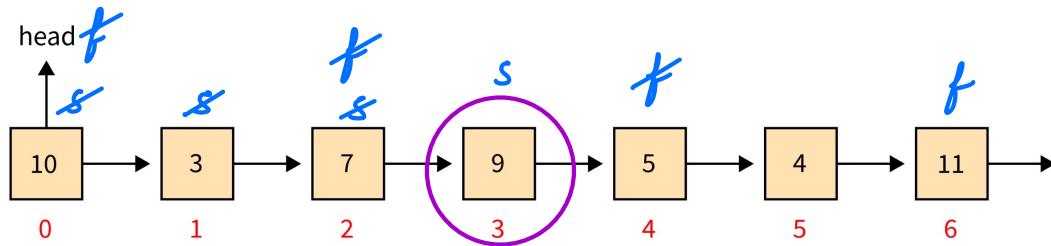
LinkedList - 2

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2. Middle point in Linked-list
3. Merge two sorted Linked-list
4. Merge sort in Linked-list
5. Check if there is a loop in Linked-list
6. Find the start point of the Loop



Middle of a Linked-list



Sol 1 → \Rightarrow Find length of linked list (N). $\rightarrow \underline{O(N)}$

\Rightarrow travel half length ($N/2$). $\rightarrow \underline{O(N/2)}$

$TC = \underline{O(N)}$

$SC = \underline{O(1)}$

Sachin (200 Km/h)



if ($Head == null$) return null

$s = f = Head$

while ($f.next != null \& \& f.next.next != null$) {

$s = s.next$

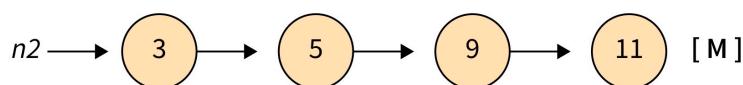
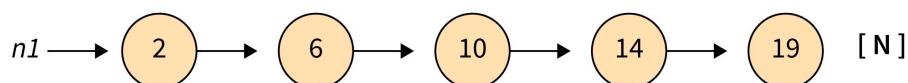
$f = f.next.next$

}

$TC = \underline{O(N/2)} = \underline{O(N)}$ $SC = \underline{O(1)}$



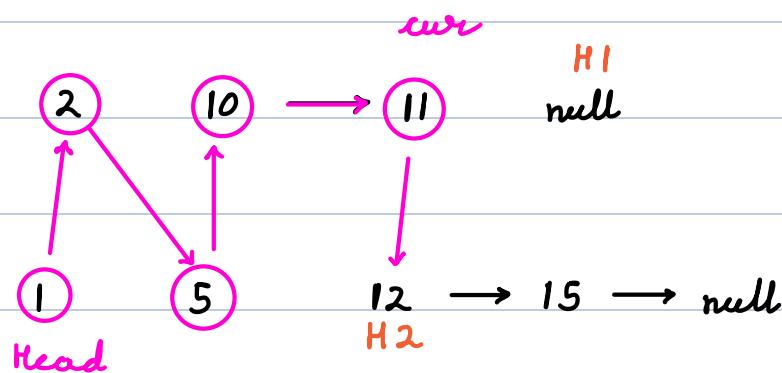
Merge Two Sorted Linked-list



Head $2 \rightarrow 3 \rightarrow 5 \rightarrow 6 \rightarrow 9 \rightarrow 10 \rightarrow 11 \rightarrow 14 \rightarrow 19 \rightarrow \text{null}$

H1 $2 \rightarrow 10 \rightarrow 11 \rightarrow \text{null}$

H2 $1 \rightarrow 5 \rightarrow 12 \rightarrow 15 \rightarrow \text{null}$



if (H1 == null) return H2

if (H2 == null) return H1

Head = null



if ($H1.data \leq H2.data$) {

 Head = $H1$

$H1 = H1.next$

} else {

 Head = $H2$

$H2 = H2.next$

}

$cur = Head$

while ($H1 \neq \text{null}$ & & $H2 \neq \text{null}$) {

 if ($H1.data \leq H2.data$) {

$cur.next = H1$

$H1 = H1.next$

 } else {

$cur.next = H2$

$H2 = H2.next$

 } $cur = cur.next$

}

if ($H1 == \text{null}$) $cur.next = H2$

if ($H2 == \text{null}$) $cur.next = H1$

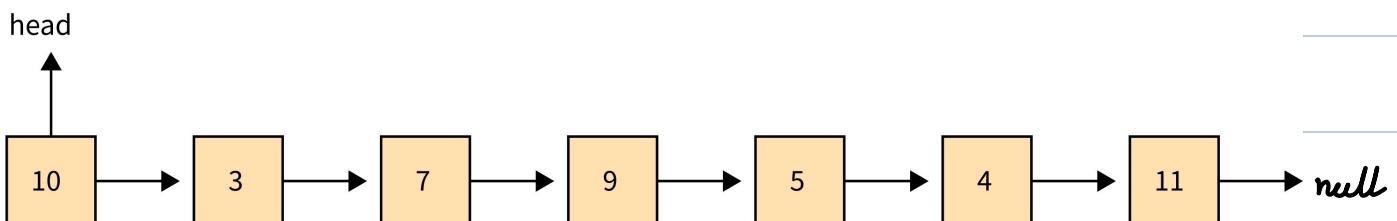
return Head

$TC = \underline{O(N+M)}$

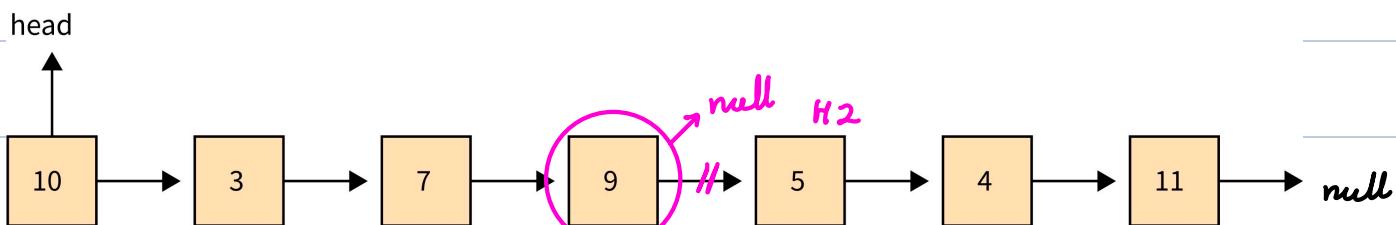
$SC = \underline{O(1)}$



Merge Sort a Linked-list



Head 3 → 4 → 5 → 7 → 9 → 10 → 11 → null



H1 10 → 3 → null H2 7 → 9 → null 5 → 4 → null 11 → null

10 → null 3 → null 7 → null 9 → null 5 → null 4 → null 11 → null

10 → null 3 → null 7 → null 9 → null 5 → null 4 → null

3 → 10 → null

7 → 9 → null

4 → 5 → null

3 → 7 → 9 → 10 → null

4 → 5 → 11 → null

3 → 4 → 5 → 7 → 9 → 10 → 11 → null



Idea

1. Find the middle node
2. Make recursive calls to sort 1st half and 2nd half
3. Finally, merge two sorted linked-list



Node sort (Head) {

if (Head == null || Head.next == null)
return Head

mid = findMiddle (Head) $\rightarrow TC = O(N)$

H2 = mid.next

mid.next = null

H1 = sort (Head)

H2 = sort (H2)

return mergeSortedLists (H1, H2) $\rightarrow TC = O(N)$

}

$TC = O(N \log(N))$

$SC = O(\log(N))$

Scenerio

You are using **Google Maps** to help you find your way around a new place. But, you get lost and end up walking in a circle. **Google Maps** has a way to keep track of where you've been with the help of special **sensors**.

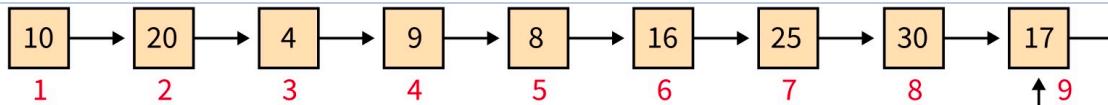
These sensors notice that you're **walking in a loop**, and now, **Google** wants to create a new feature to help you figure out exactly where you started going in circles.

Problem

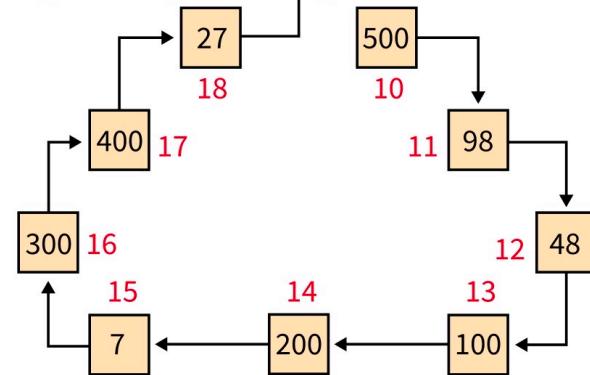
You have a **linked list** that shows each **step** of your **journey**, like a chain of events. Some of these steps have accidentally led you back to a place you've already been, making you **walk in a loop**. The goal is to find out the exact point where you first started walking in this loop.

Check if there is a loop

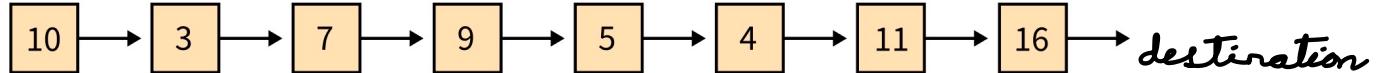
1



Ans = true



2



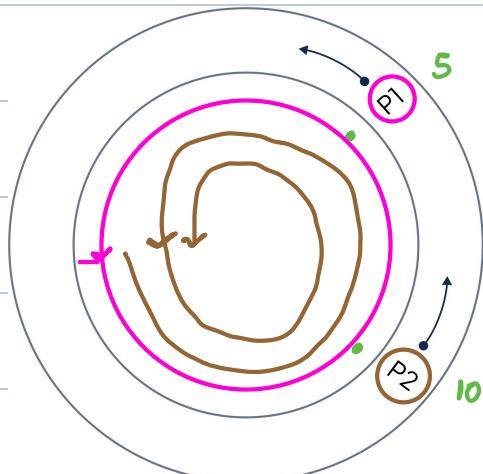
Ans = false

sol 1 → Use hashset to check already visited node. ✓

SC = O(1)

O(1)

If two people are running with different speeds on a circular track, they will 100% meet at some point.





if (Head == null) return false

s = f = Head

while (f != null && f.next != null) {

s = s.next

f = f.next.next

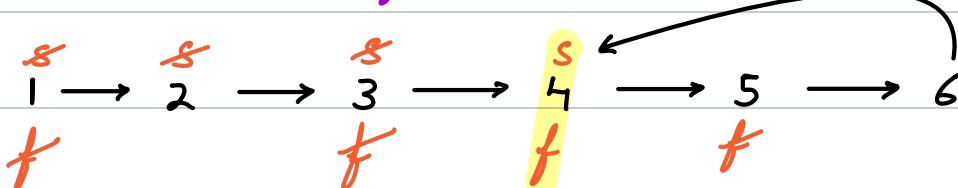
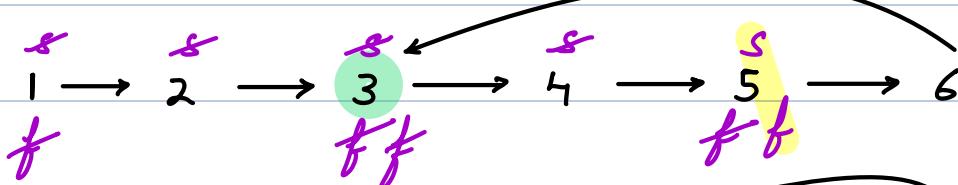
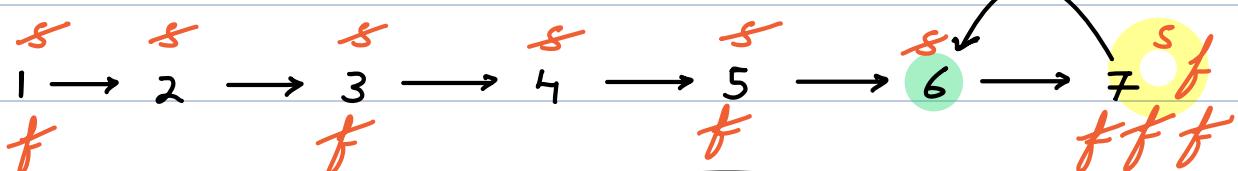
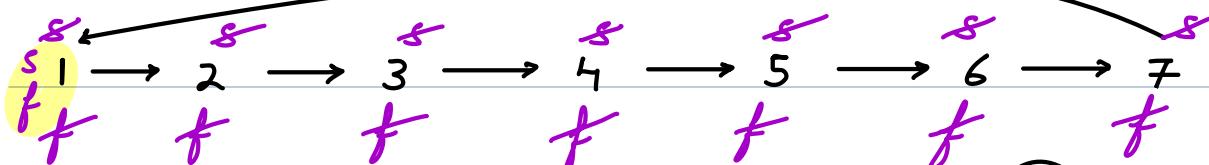
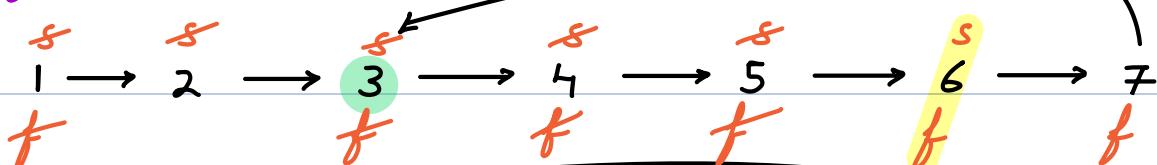
if (s == f) return true

}

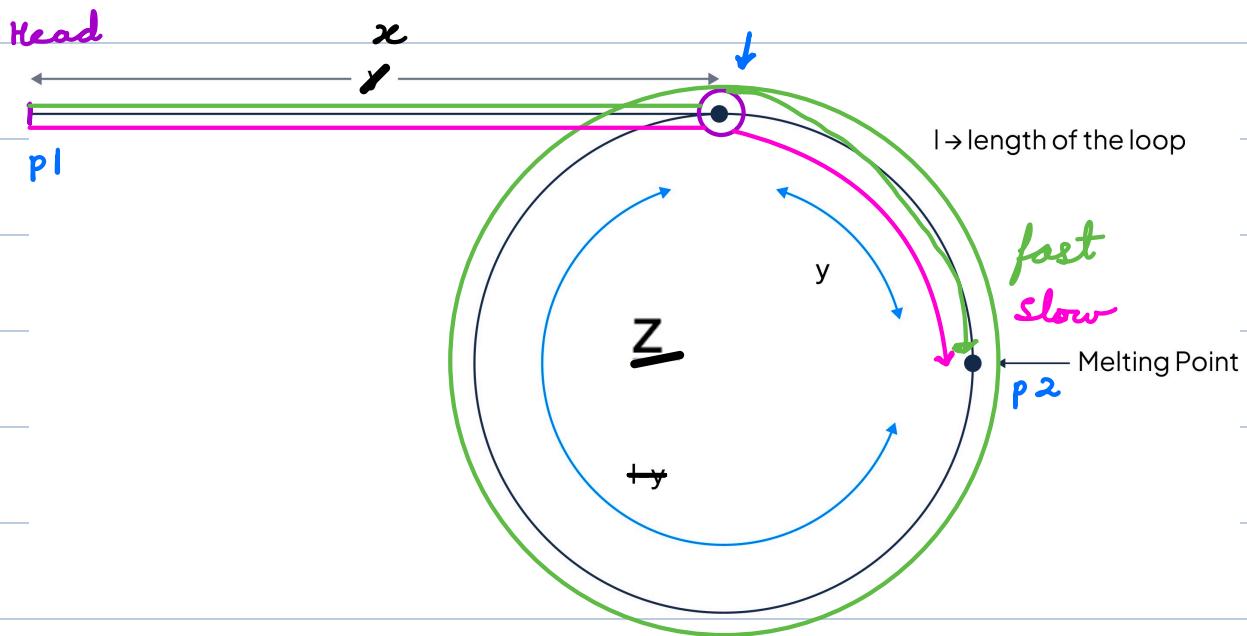
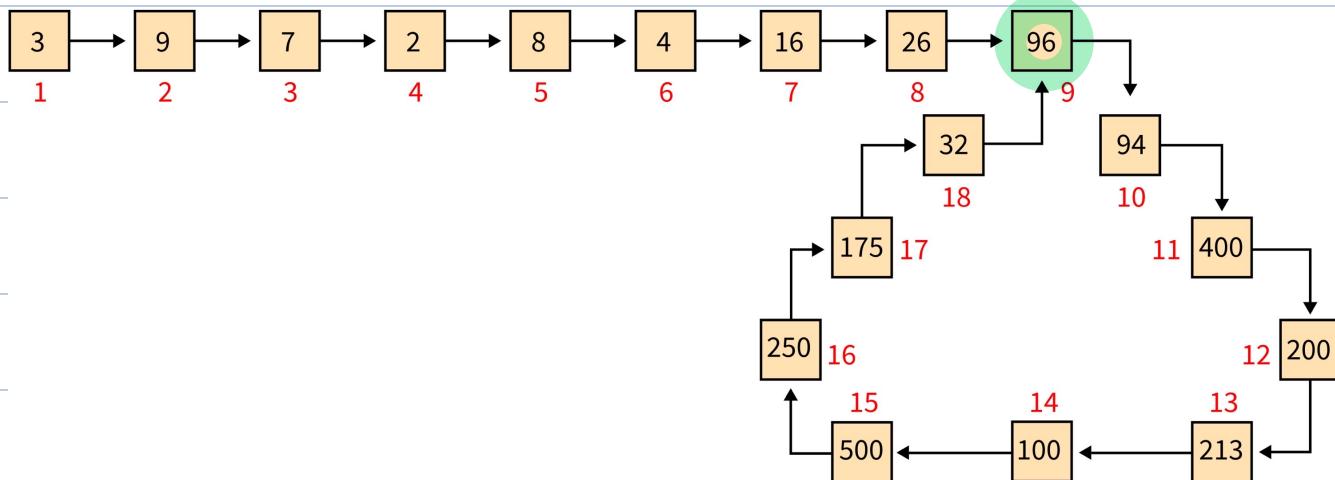
return false

TC = O(N) SC = O(1)

Head



★ Find the start point of the loop



Distance

$$\text{slow} = x + y$$

$$\text{fast} = x + y + z + y$$

Fast is all double speed \Rightarrow

$$x + y + z + y = 2 * (x + y)$$

$$x + y + z + y = x + y + x + y$$

$$\Rightarrow z = x$$



x

if (Head == null) return false

s = f = Head

while (f != null && f.next != null) {

s = s.next

f = f.next.next

if (s == f) break

{

p1 = Head

p2 = s

while (p1 != p2) {

p1 = p1.next

p2 = p2.next

{

return p1

TC = O(N) SC = O(1)

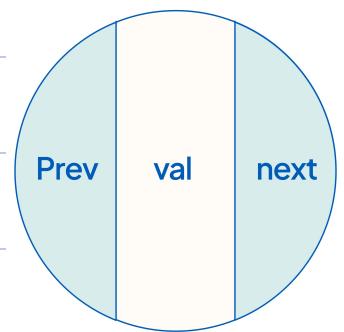
Agenda:

1. What is doubly linked list?
 2. LRU Cache
 3. Check if LL is palindrome



Doubly Linked List

```
class Node {  
    int val;  
    Node next;  
    Node prev;  
    public Node (int v) {  
        this.val = v;  
        this.next = null;  
        this.prev = null;  
    }  
}
```





Scenario

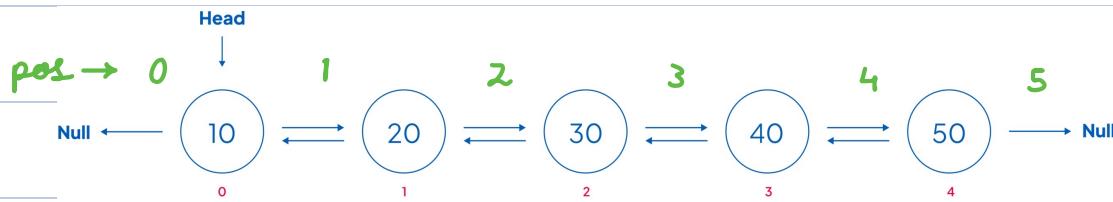
Spotify wants to enhance its user experience by allowing users to navigate through their music playlist seamlessly using "next" and "previous" song functionalities.

Problem

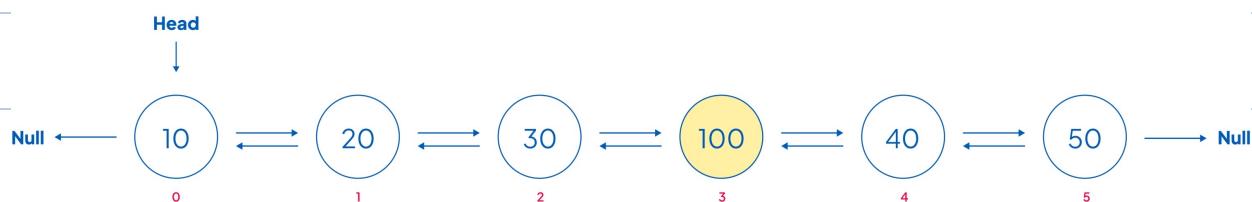
You are tasked to implement this feature using a doubly linked list where each node represents a song in the playlist. The system should support the following operations:

- **Add Song :** Insert a new song into the playlist. If the playlist is currently empty, this song becomes the "Current song".
- **Play Next Song :** Move to the next song in the playlist and display its details.
- **Play Previous Song :** Move to the previous song in the playlist and display its details.
- **Current Song :** Display the details of the current song being played.

Insertion in Doubly Linked List

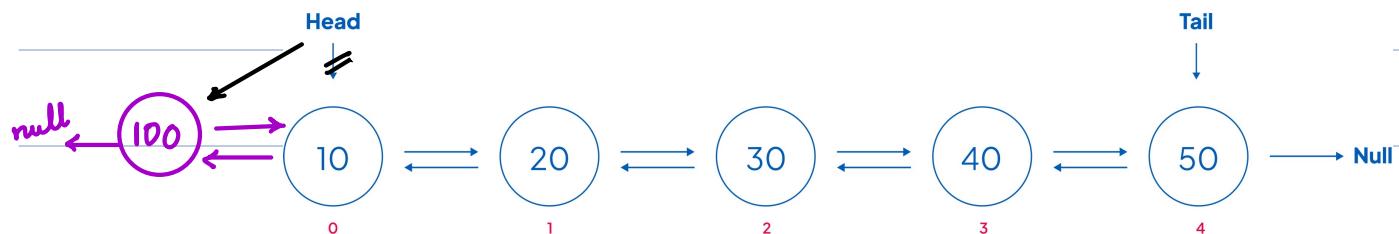


↓
insert(100,3)



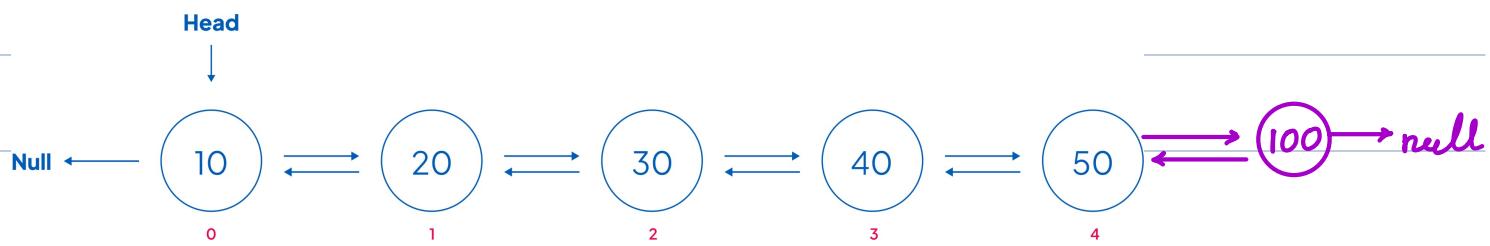
Edge - Case

- At index 0





- At index 5



// Head, data, pos

$0 \leq pos \leq \text{length of LL}$

$xr = \text{new Node (data)}$

$\text{if } (pos == 0) \{$

$xr.\text{next} = \text{Head}$

$\text{if } (\text{Head} \neq \text{null}) \quad \text{Head}.pre = xr$

$\text{return } xr \quad // \text{new Head}$

}

$\text{temp} = \text{Head}$

$\text{for } i \rightarrow 1 \text{ to } (pos-1) \{$

$\text{temp} = \text{temp}.next$

}

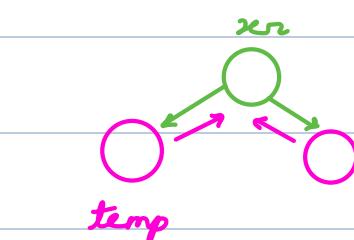
$xr.pre = \text{temp}$

$xr.next = \text{temp}.next$

$\text{if } (\text{temp}.next \neq \text{null}) \quad \text{temp}.next.pre = xr$

$\text{temp}.next = xr$

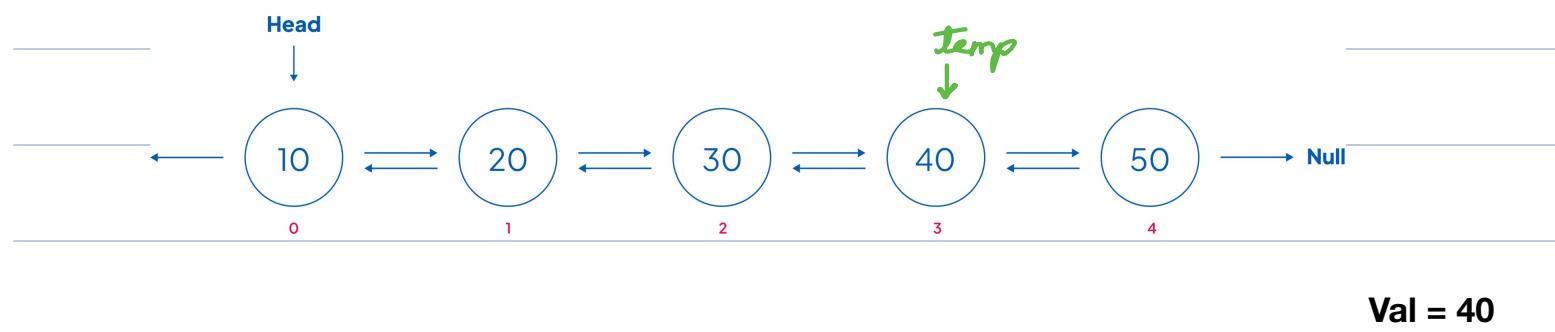
return Head



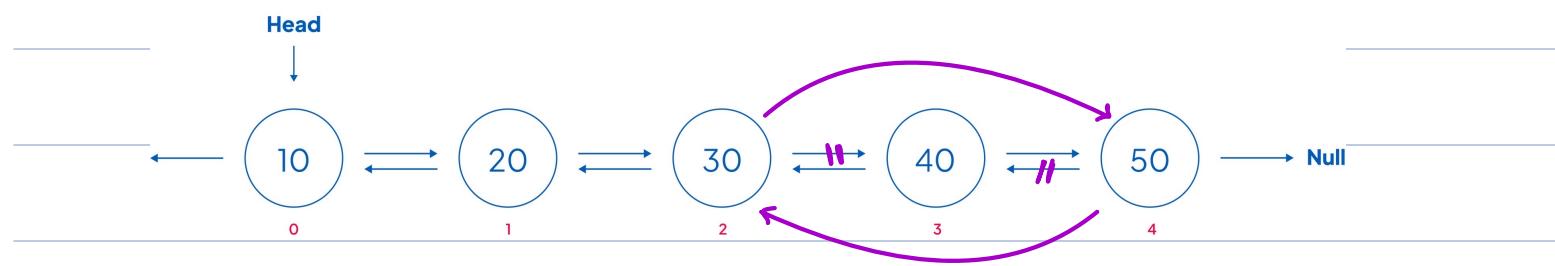
$TC = \underline{O(N)}$

$SC = \underline{O(1)}$

Delete a node from D.L.L



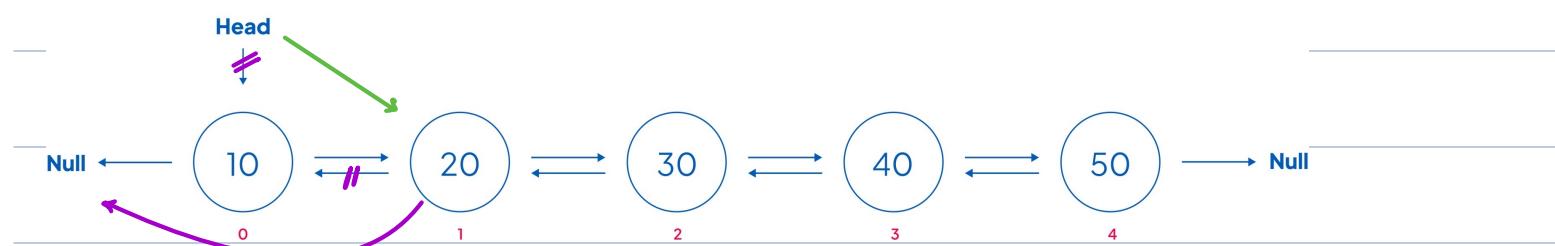
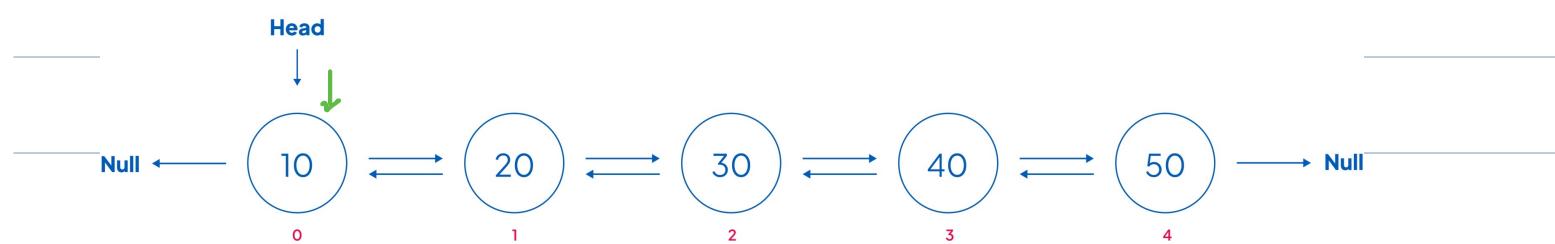
Delete the node with val = 40 →



Edge - Case

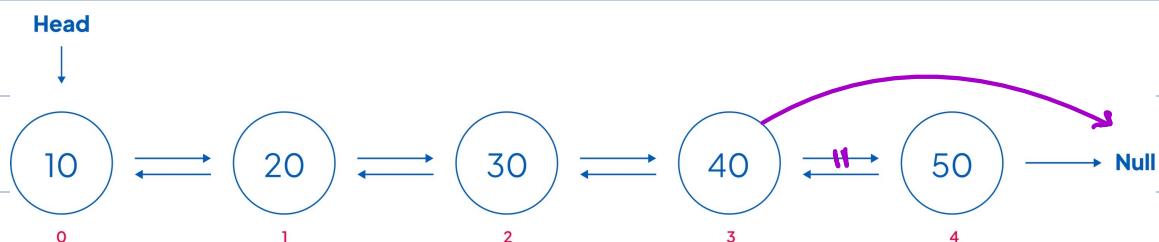
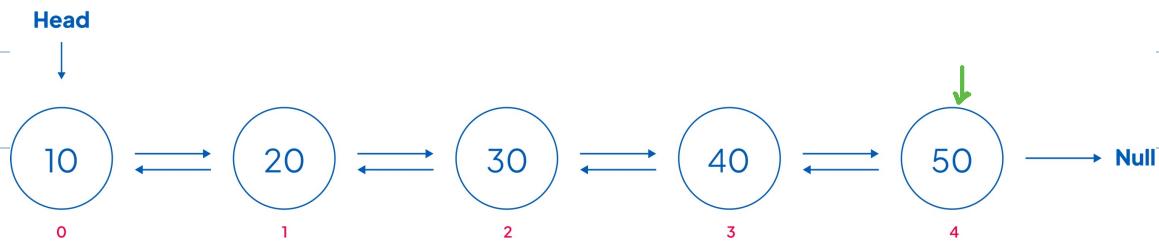
- idx → 0

val → 10





• last - node

val \rightarrow 50

temp = Head

while (temp != null) {

 if (temp.val == X) break

 temp = temp.next

}

if (temp == null) return Head

{ if (temp.pre == null)

 Head = Head.next

 if (temp.next != null)

 temp.next.pre = temp.pre

 if (temp.pre != null)

 temp.pre.next = temp.next

return Head



TC = O(N) // Searching

SC = O(1)

L.R.U Cache

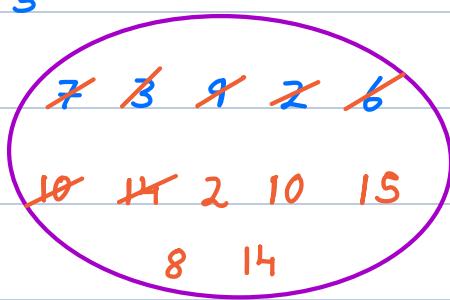
(Google)

↓
Least Recently Used

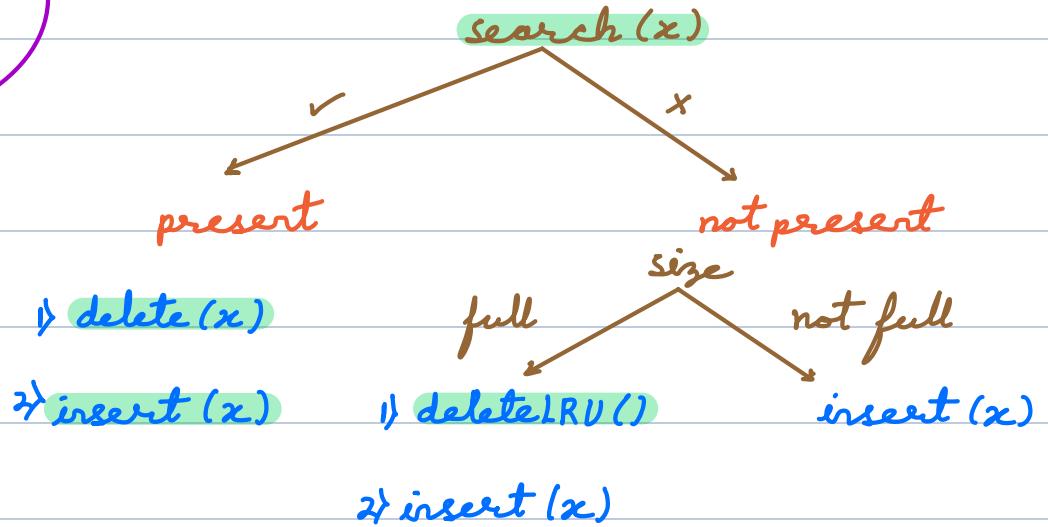
Q → Given a running stream of integers & a fixed memory M ($SC = O(M)$). Maintain most recent M items i.e. if the memory is full, delete least recent element.

i/p → 7 3 9 2 6 10 14 2 10 15 8 14

M = 5



new i/p → x



searching(x) → HashSet / HashMap < data, node > X ✓

deleteLRU() → Array (static / dynamic) X

(deleteHead()) LinkedList (single / Doubly) X ✓



$\text{delete}(x) \rightarrow$ data $x \rightarrow x_n = \text{findNode}(x) \checkmark$

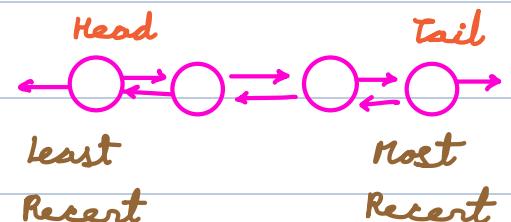
$\text{deleteNode}(x_n)$

$\text{insert}(x) \rightarrow \checkmark$

1) $\text{search}(x) \rightarrow \text{hm. contains}(x)$

2) $\text{deleteLRU}() \rightarrow x = \text{deleteHead}()$

$\text{hm. remove}(x)$



3) $\text{delete}(x) \rightarrow x_n = \text{hm. get}(x)$

$\text{deleteNode}(x_n)$

4) $\text{insert}(x) \rightarrow x_n = \text{new Node}(x)$

$\text{insertTail}(x_n)$

$\text{hm. put}(x, x_n)$

$TC = \underline{O(1)}$

$SC = \underline{O(M)}$



Q → Check if the given linked list is palindrome.

SC = O(1)

$\rightarrow = \leftarrow$

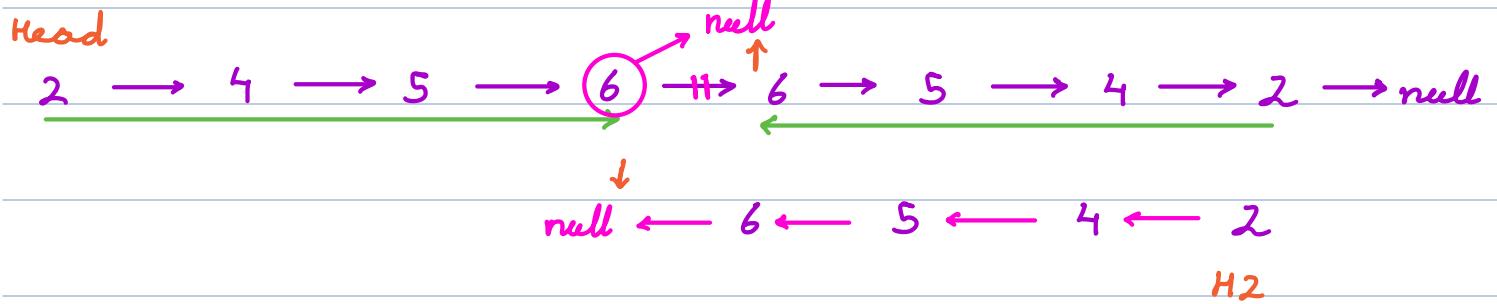
$x = \text{rev}(x)$

Head

2 → 4 → 5 → 4 → 2 → null Ans = true

Head

2 → 4 → 5 → 2 → 2 → null Ans = false



Sol → 1) Find middle element.

2) Reverse second half of LL.

3) Compare first & second half.

TC = O(N)

SC = O(1)