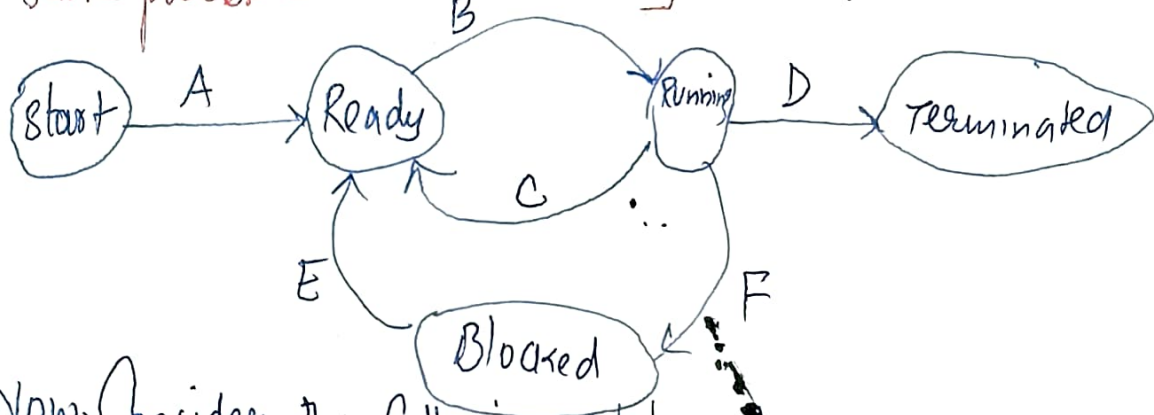


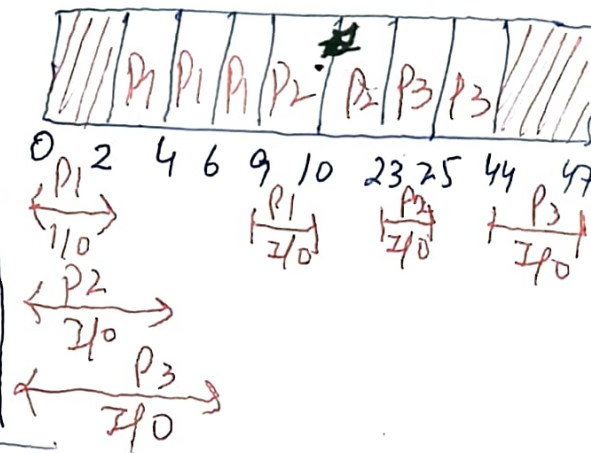
Q₁ - In the following process state transition diagram for a Uniprocessor system, assume that there are always some processes in the Ready state:



Now, Consider the following statements:

Q₂ - Consider 3 processes, all arriving at time 0, with BT of 10, 20, and 30 units respectively. Each process spends the first 20% of exec. time doing I/O, the next 70% of time doing Computation, and the last 10% of time doing I/O again. The OS uses SRTF scheduling. Assume that all I/O operations can be overlapped as much as possible. For what %age of time does the CPU remain idle?

PNO	AT	BT	I/O	CPU	I/O
P ₁	0	10	2	7	1
P ₂	0	20	4	14	2
P ₃	0	30	6	21	3

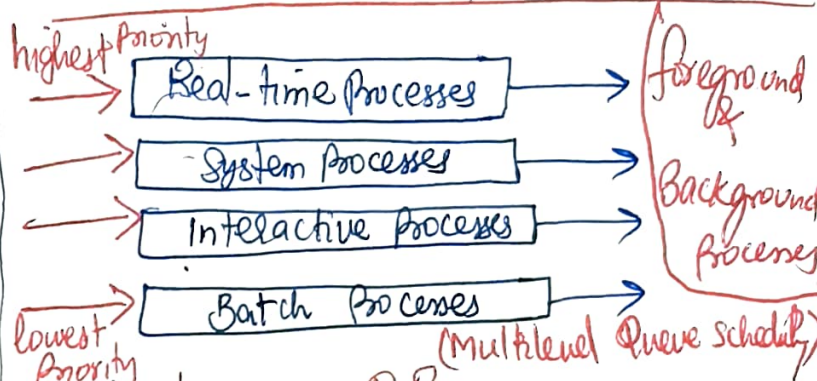
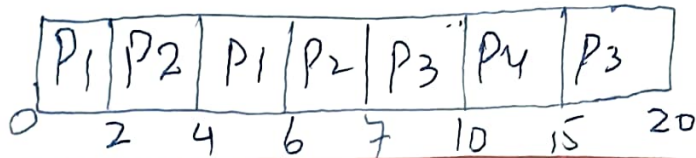


CPU	I/O	0	2	4	6	8	10	23	25	44	47
		///	①	①	①	②	②	③	③	///	///
		②③	②③	③		①		②		③	

$$\begin{aligned} \text{CPU idle time} &= 2 + 2 + 3 \\ &= 10.638\% \end{aligned}$$

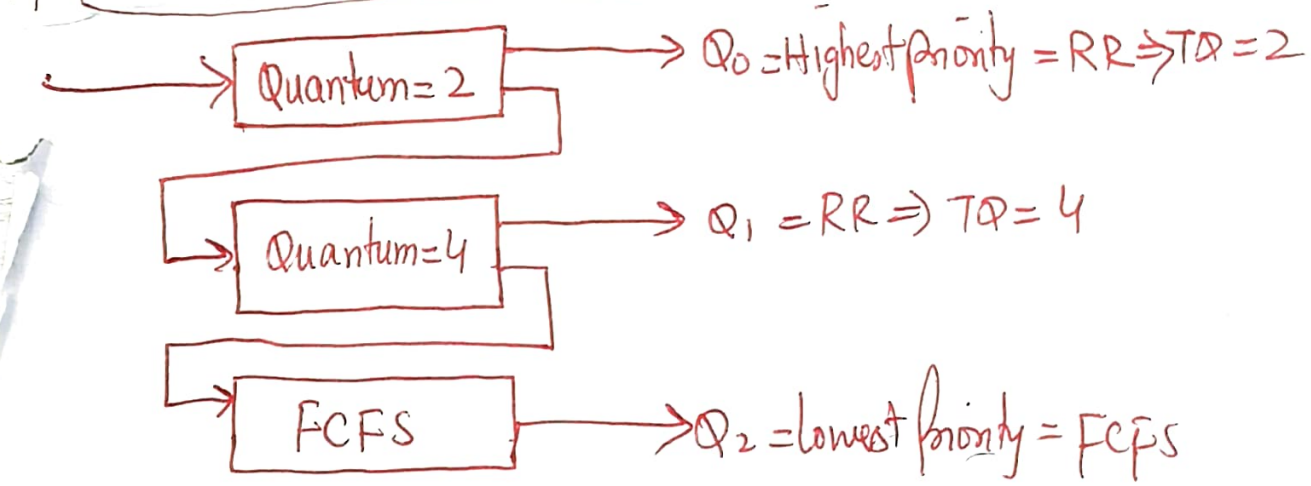
Multilevel Queue

PNO	AT	BT	Queue No
P1	0	4	1
P2	0	3	1
P3	0	8	2
P4	10	5	1

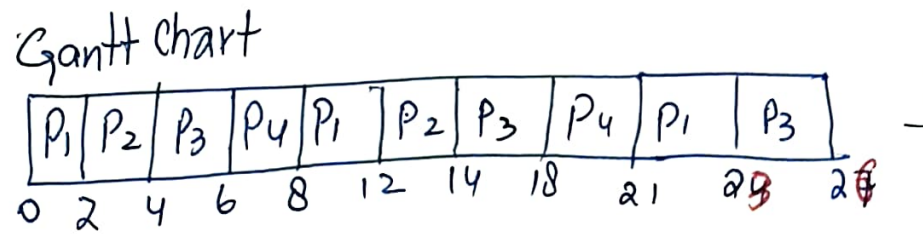


NOTE: Queue 1 has higher priority than Queue 2. RR is used in Queue 1 (TQ=2), while FCFS is used in Queue 2.

Multilevel Feedback Scheduling (Aging) prevents Starvation



PNO	AT	BT
P1	0	8
P2	1	4
P3	2	9
P4	3	5



Operation on Processes

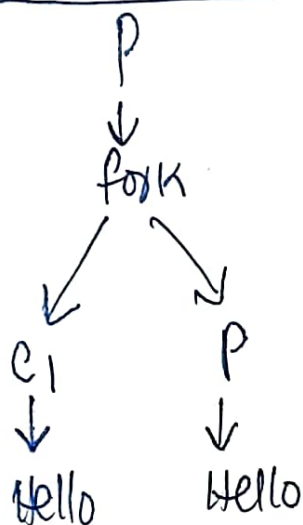
↳ Creation

↳ Termination

$\text{fork()} = \begin{cases} 0 & \text{child} \\ +1 > 0 & \text{parent} \\ -1 < 0 & \text{error} \end{cases}$

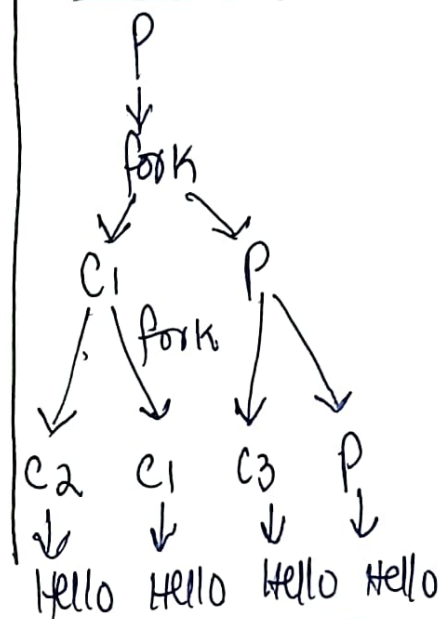
Ex1: `main() {
fork();
printf("Hello");
}`

Process Tree



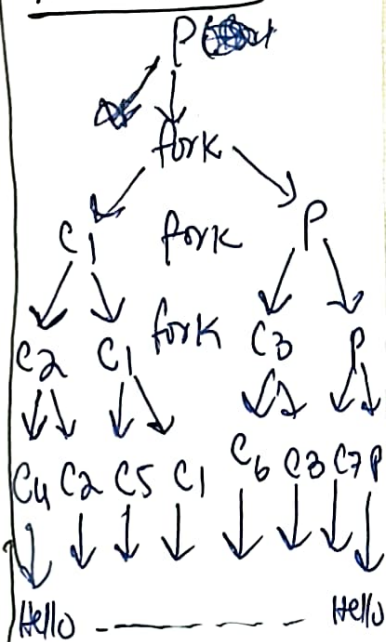
Ex2: `main() {
fork();
fork();
printf("Hello");
}`

Process Tree



Ex3 `main() {
fork();
fork();
fork();
printf("Hello");
}`

Process Tree



NOTE ① No. of Hello statements printed = 2^n

② No. of child processes created is $2^n - 1$

Zombie process

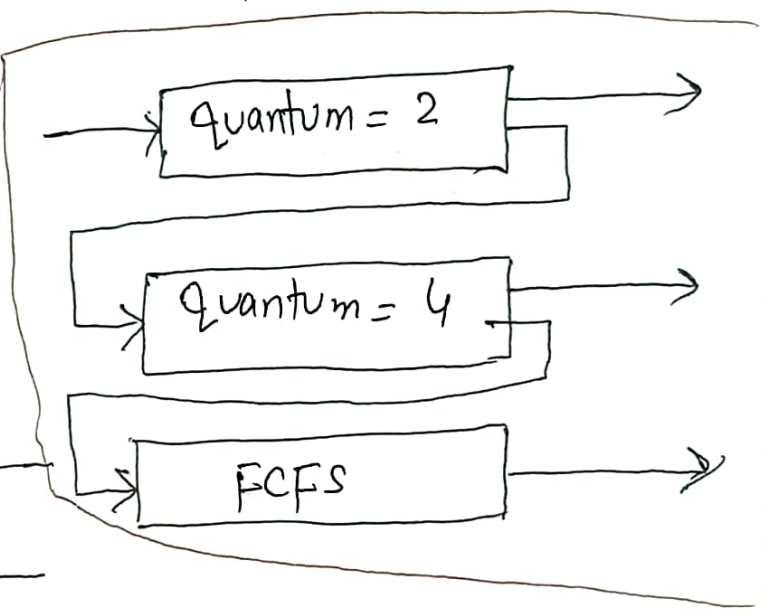
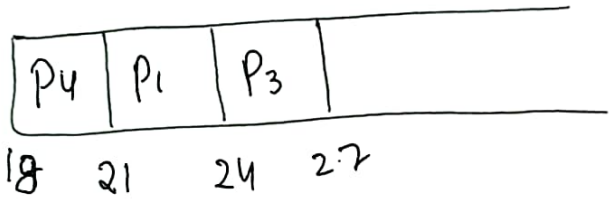
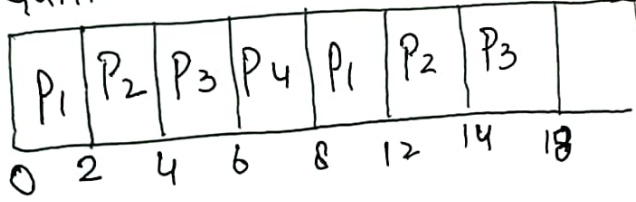
Orphan process

Process Information

Process	AT	BT	
P1	0	8	③
P2	1	4	x ①
P3	2	9	④
P4	3	5	x ②

Q0 (Highest priority) = RR TD=2
 Q1 = RR TD=4
 Q2 Lowest p = FCFS

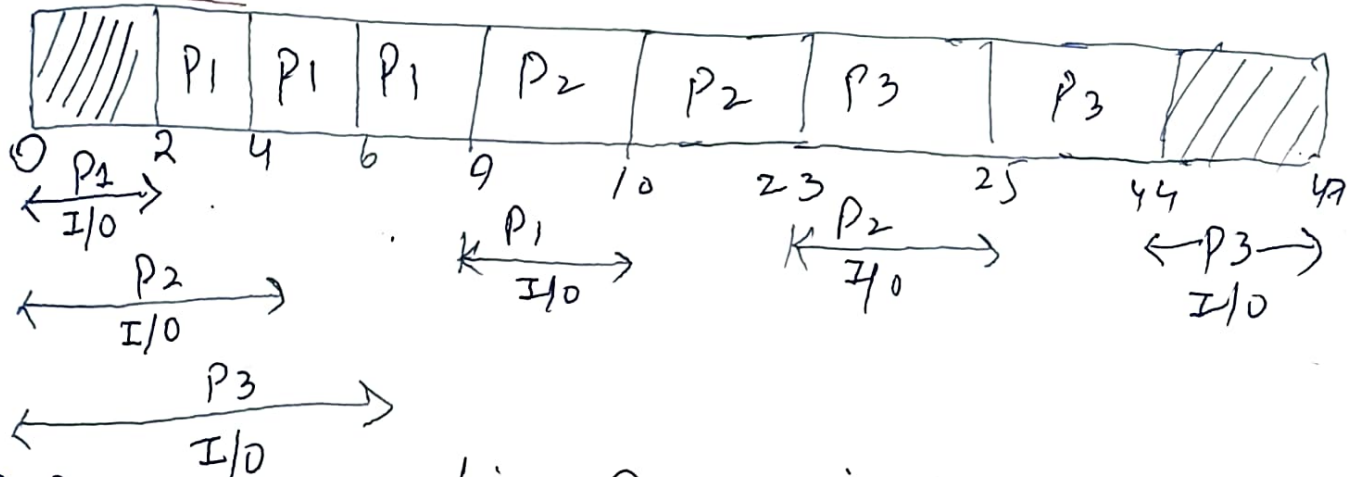
Gantt Chart



Q. -

PNO	AT	BT	I/O	CPU Burst	I/O Burst
P1	0	10	2	7	1
P2	0	20	4	14	2
P3	0	30	6	21	3

Gantt Chart



∴ percentage of time CPU remains idle = $\frac{5}{47} \times 100$
 $= 10.638\%$

Round Robin :- Criteria: "Time Quantum" ; Mode: "Preemptive"
 TQ = 2

Process No	AT	BT	CT	TAT	WT	RT
P1	0	5	12	12	7	0
P2	1	4	11	10	6	1
P3	2	2	6	4	2	2
P4	4	1	9	5	4	4

Ready Queue: P1 P2 P3 P1 P4 P2 P1

Running Queue (Gantt chart): P1 P2 P3 P1 P4 P2 P1
 0 2 4 6 8 9 11 12

Pre-emptive Priority Scheduling Criteria: "Priority" Mode: "Preemptive"
 NOTE: Higher the no. Higher the priority

Priority	Process No	AT	BT	CT	TAT	WT
10	P1	0	5	12	12	7
20	P2	1	4	8	7	3
30	P3	2	2	4	2	0
40	P4	4	1	5	1	0

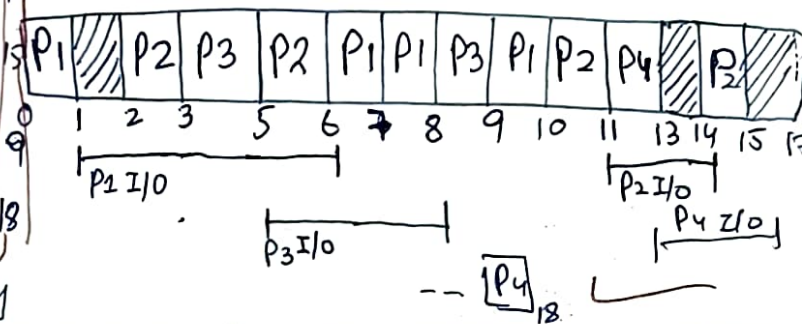
Gantt chart

P1	P2	P3	P3	P4	P2	P1	
0	1	2	3	4	5	8	12

Non-preemptive Priority Scheduling

Mix BT (CPU & I/O Both) Mode: Pre-emptive Criteria: Priority based
 Find CT of P1, P2, P3, P4

Process	AT	Priority	CPU	I/O	CPU	CT
P1	0	2	1	5	3	10
P2	2	3	3	3	1	10
P3	3	1H	2	3	1	9
P4	3	4	2	4	1	10



NOTE: Lower the no. higher the priority

Categories of Scheduling Algos.

Based on the environments viz

1. Batch
2. Interactive
3. Real Time

↳ Scheduling in Batch Systems: FCFS, SJF, SRTF

↳ Scheduling in Interactive Systems: Priority, RR, Multiple Queues, SRTF, Guaranteed scheduling, Lottery scheduling, Fair-share Scheduling.

↳ Scheduling in Real-time Systems:
 Hard Real Time
 → Soft Real Time (video whose audio isn't properly synced)
 Rate monotonic, Earliest Deadline First.

↳ Static Scheduling

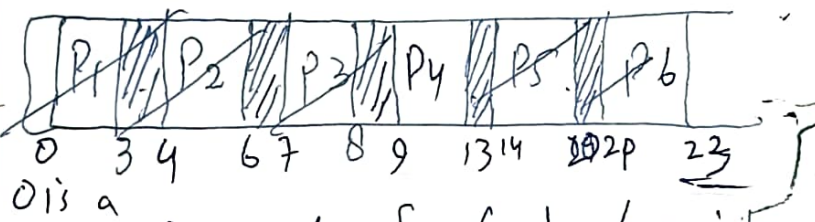
↳ Dynamic Scheduling

02/Sep/2025

Q. Consider the set of 6 processes whose AT and BT are given below. If the CPU Scheduling Policy is FCFS & there is 1 unit of overhead in scheduling the processes, find the efficiency of the algorithm?

PNO	AT	BT	CT	TAT	WT
P1	0	3	-		
P2	1	2	-		
P3	2	1	-		
P4	3	4	-		
P5	4	5	-		
P6	5	2	-		

Gantt chart



Now, Useless Time or wasted Time = $6 \times 1 = 6 \times 1 = 6$ unit
 Total Time = 23 unit
 Useful time = $23 - 6 = 17$ unit
 \therefore Efficiency (η) = $\frac{\text{Useful Time}}{\text{Total Time}} = \frac{17}{23} = 0.7391 = 73.91\%$

- NOTE:
- RR Scheduling is FCFS Scheduling with preemptive mode.
 - When Time Quantum tends to infinity, RR Scheduling becomes FCFS Scheduling
 - The Waiting time for the process having the highest priority will always be zero in preemptive mode

NOTE: → Convoy effect
 → Starvation
 → Aging

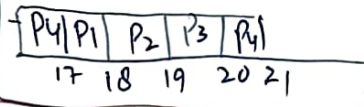
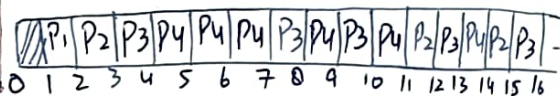
02 Sep 2025

Longest Job First Algorithm

Process NO	AT	BT
1	0	3
2	1	2
3	2	4
4	3	5
5	4	6

Longest remaining Time first Algo

PNO	AT	BT
1	1	2
2	2	4
3	3	6
4	4	8



Highest Response Ratio Next (HRRN)

Criteria: Response Ratio (RR) = $\frac{W+S}{S}$

W = Waiting time for a process so far

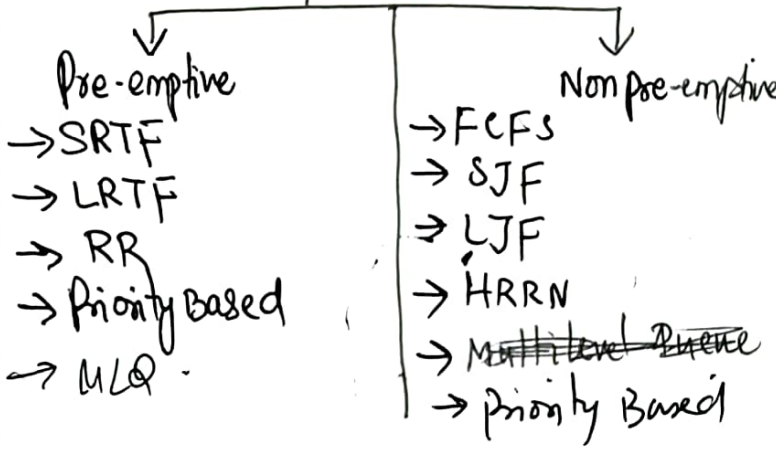
S = Service time or BT

→ HRRN not only favours shorter jobs but also limits the waiting time of longer jobs.

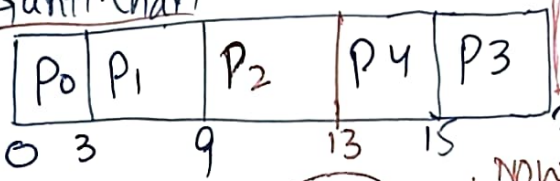
Mode: non-preemptive

Scheduling Algorithms

PNO	AT	BT	CT	TAT	WT	RT
0	0	3				
1	2	6				
2	4	4				
3	6	5				
4	8	2				



Gantt Chart



$RR_2 = \frac{S_2 + 4}{4} = 2.25$ ✓

$RR_3 = \frac{3+5}{5} = 1.6$

$RR_4 = \frac{1+2}{2} = 1.5$

Now, again

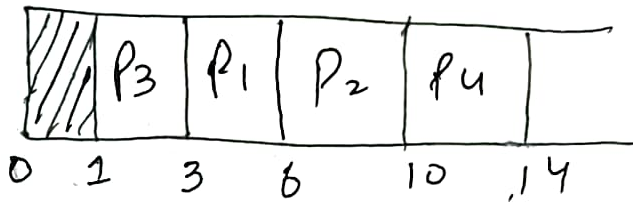
$RR_3 = \frac{7+5}{5} = \frac{12}{5} = 2.4$

$RR_4 = \frac{5+2}{2} = \frac{7}{2} = 3.5$

② SJF Criteria: Burst Time | Mode: Non-Preemptive

Process No	AT	BT	CT	TAT	WT	RT
P ₁	1	3	6	5	2	2
P ₂	2	4	10	8	4	4
P ₃	1	2	3	2	0	0
P ₄	4	4	14	10	6	6

Gantt Chart



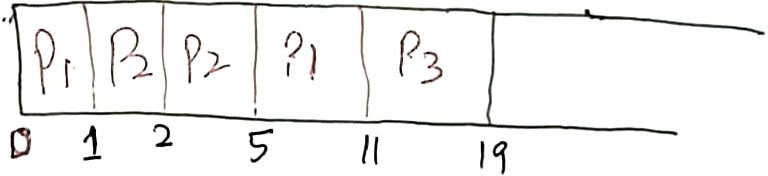
SRTF

④

Process	AT	BT
P ₁	0	7
P ₂	1	4
P ₃	2	8

The Gantt chart for Preemptive SJF scheduling algo is ?

Gantt chart



③ SRTF Criteria: "BT" mode: "Preemptive"

Process No	AT	BT	CT	TAT	WT	RT
P1	0	5 ⁴	9	9	4	0
P2	1	3 ²	4	3	0	0
P3	2	4	13	11	7	7
P4	4	5	5	6	0	0

Ganttchart (SRTF)

P1	P2	P2	P2	P4	P1	P3	
----	----	----	----	----	----	----	--

$$RT = \text{cpu first Time} - \text{AT}$$

① FCFS (Criteria: ~~BT~~ Arrival Time) mode: Non-preemptive

Process No	AT	BT or Exec. Time	CT	TAT	WT	RT
P1	0	2	2	2	0	0
P2	1	2	4	3	1	1
P3	5	3	8	3	0	0
P4	6	4	12	6	2	2

②

P1	P2	 	P3	P4	
----	----	--------------	----	----	--

0 2 4 5 7 12