



# Model Predictive Control

## 3. Fundamentals of Optimization

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## Gradient, Hessian, and Jacobian

### Tutorial

**Definition 3.1** The **gradient** of a function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is defined as  $\nabla f(x_1, \dots, x_n) = \left( \frac{\partial f}{\partial x_1} \quad \dots \quad \frac{\partial f}{\partial x_n} \right)^T$ .

**Definition 3.2** The **Hessian** of a function  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is defined as  $\mathbf{H}_f(x_1, \dots, x_n) = \begin{pmatrix} \frac{\partial^2 f}{\partial x_1^2} & \dots & \frac{\partial^2 f}{\partial x_1 \partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial^2 f}{\partial x_n \partial x_1} & \dots & \frac{\partial^2 f}{\partial x_n^2} \end{pmatrix}$ .

**Definition 3.3** The **Jacobian** of a function  $\mathbf{f}: \mathbb{R}^n \rightarrow \mathbb{R}^m$  is defined as  $\mathbf{J}_f(x_1, \dots, x_n) = \begin{pmatrix} \frac{\partial f_1}{\partial x_1} & \dots & \frac{\partial f_1}{\partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial f_m}{\partial x_1} & \dots & \frac{\partial f_m}{\partial x_n} \end{pmatrix}$ .

## Nonlinear Optimization Problem

**Problem 3.1** A **nonlinear optimization problem** is defined in **standard form** as

$$\min_x f(\mathbf{x}) \quad \text{with } f: \mathbb{R}^n \rightarrow \mathbb{R} \quad \text{cost function or objective function} \quad (3.1)$$

$$\text{subject to } \begin{cases} \mathbf{h}(\mathbf{x}) = \mathbf{0} & \text{with } \mathbf{h}: \mathbb{R}^n \rightarrow \mathbb{R}^m & \text{equality constraints} & (3.2) \\ \mathbf{g}(\mathbf{x}) \leq \mathbf{0} & \text{with } \mathbf{g}: \mathbb{R}^n \rightarrow \mathbb{R}^p & \text{inequality constraints} & (3.3) \end{cases}$$

- **Symbols**

- The vector  $\mathbf{x} = (x_1 \ x_2 \ \cdots \ x_n)^T \in \mathbb{R}^n$  is denoted as **decision variable** or **optimization variable**
- The solution  $\mathbf{x}^* \in \mathbb{R}^n$  of Problem 3.1 is denoted as **minimizer**

- **Remark**

- For  $m < n$  the equality constraints (3.2) are **underdetermined** ✓
- For  $m = n$  the equality constraints (3.2) are **determined** for  $h_i, i \in \{1, \dots, m\}$  independent ✗
- For  $m > n$  the equality constraints (3.2) are **overdetermined** ✗

## Nonlinear Optimization Problem

- **Assumption**

- Cost function  $f \in \mathcal{C}^2$ , functions  $h_i \in \mathcal{C}^1, i \in \{1, \dots, m\}$  and  $g_j \in \mathcal{C}^1, j \in \{1, \dots, p\}$  where  $\mathcal{C}^j$  is the set of  $j$  times continuously differentiable functions

- **Remarks**

- **Nonsmooth optimization** if assumption not fulfilled (not considered in this lecture)
- **Integer optimization** if  $x \in \mathbb{Z}^n$  (not considered in this lecture)

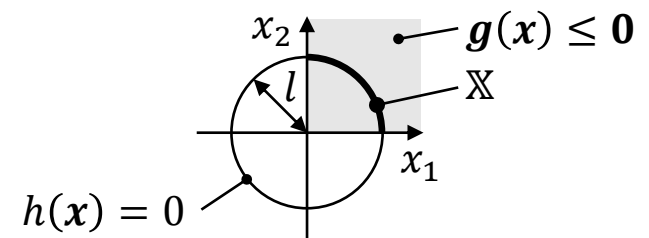
- **Example**

- **Maximization** of the **area of a right triangle** with legs  $x_1$  and  $x_2$  and a given hypotenuse  $l$

- Cost function  $f(x) = -\frac{1}{2}x_1x_2$

- Equality constraint  $h(x) = x_1^2 + x_2^2 - l^2 = 0$

- Inequality constraints  $g_1(x) = -x_1 \leq 0, g_2(x) = -x_2 \leq 0$



## Nonlinear Optimization Problem

**Problem 3.2** A **nonlinear optimization problem** is defined as

$$\min_x f(\mathbf{x}) \quad \text{with } f: \mathbb{R}^n \rightarrow \mathbb{R} \quad \text{cost function} \quad (3.4)$$

$$\text{subject to } \mathbf{x} \in \mathbb{X} \quad \text{with } \mathbb{X} = \{\mathbf{x} \in \mathbb{R}^n \mid \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\} \quad \text{feasible set} \quad (3.5)$$

- **Symbols**

- A vector  $\mathbf{x} \in \mathbb{X}$  is denoted as **feasible point**

- **Remarks**

- Problem 3.2 is an alternative formulation of Problem 3.1
- Problem 3.2 can be written even more briefly as  $\min_{\mathbf{x} \in \mathbb{X}} f(\mathbf{x})$
- Note that considering a minimization problem is not restrictive since a maximization problem can be transformed into a minimization problem using  $\max_{\mathbf{x} \in \mathbb{X}} f(\mathbf{x}) = \min_{\mathbf{x} \in \mathbb{X}} -f(\mathbf{x})$

## Local Minimum and Global Minimum

**Definition 3.4** The cost function  $f(x)$  has a **local minimum** at the point  $x^* \in \mathbb{X}$  if there exists an  $\varepsilon > 0$  such that  $f(x^*) \leq f(x)$  for all  $x \in \mathbb{X} \setminus \{x^*\}$  and  $\|x - x^*\| < \varepsilon$ . If  $\leq$  is replaced by  $<$ , then the local minimum is a **strict local minimum**.

**Definition 3.5** The cost function  $f(x)$  has a **global minimum** at the point  $x^* \in \mathbb{X}$  if  $f(x^*) \leq f(x)$  for all  $x \in \mathbb{X} \setminus \{x^*\}$ . If  $\leq$  is replaced by  $<$ , then the global minimum is a **unique** or **strict global minimum**.

**Theorem 3.1** A global minimum exists if

- (1) the feasible set  $\mathbb{X}$  is bounded, i.e.  $\exists \alpha \in \mathbb{R}: \|x\| \leq \alpha \forall x \in \mathbb{X}$ ,
- (2) the feasible set is not empty, i.e.  $\mathbb{X} \neq \emptyset$ .

- **Remark**

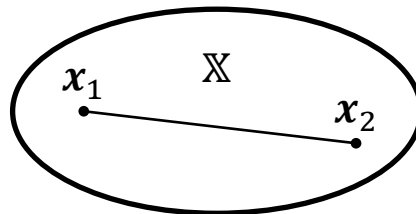
- Note the Theorem 3.1 is only sufficient

## Convex Sets

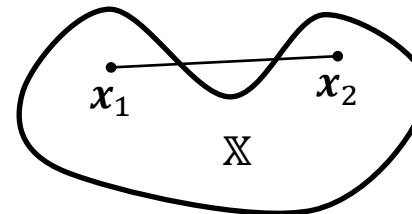
**Definition 3.6** A set  $\mathbb{X}$  is **convex** if  $\alpha x_1 + (1 - \alpha)x_2 \in \mathbb{X}$  for any  $x_1, x_2 \in \mathbb{X}$  and  $\alpha \in [0,1]$ .

- Interpretation**

- Note that  $\alpha x_1 + (1 - \alpha)x_2$  with  $\alpha \in [0,1]$  represents the line segment between the points  $x_1$  and  $x_2$
- A set is thus convex if the line segment connecting two arbitrary points  $x_1$  and  $x_2$  is also in the set



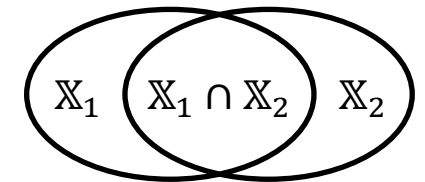
convex



not convex

- Properties**

- (1)  $\mathbb{X}$  convex,  $\beta \in \mathbb{R} \Rightarrow \beta\mathbb{X} = \{x | x = \beta v, v \in \mathbb{X}\}$  convex
- (2)  $\mathbb{X}_1, \mathbb{X}_2$  convex  $\Rightarrow \mathbb{X}_1 + \mathbb{X}_2 = \{x | x = v_1 + v_2, v_1 \in \mathbb{X}_1, v_2 \in \mathbb{X}_2\}$  convex
- (3)  $\mathbb{X}_1, \mathbb{X}_2$  convex  $\Rightarrow \mathbb{X}_1 \cap \mathbb{X}_2$  convex



## Convex Functions

**Definition 3.7** A function  $f: \mathbb{X} \rightarrow \mathbb{R}$  is **convex** on a convex set  $\mathbb{X}$  if

$$f(\alpha x_1 + (1 - \alpha)x_2) \leq \alpha f(x_1) + (1 - \alpha)f(x_2) \quad \forall x_1, x_2 \in \mathbb{X} \quad \forall \alpha \in [0,1].$$

**Definition 3.8** A function  $f: \mathbb{X} \rightarrow \mathbb{R}$  is **strictly convex** on a convex set  $\mathbb{X}$  if

$$f(\alpha x_1 + (1 - \alpha)x_2) < \alpha f(x_1) + (1 - \alpha)f(x_2) \quad \forall x_1, x_2 \in \mathbb{X}, x_1 \neq x_2 \quad \forall \alpha \in (0,1).$$

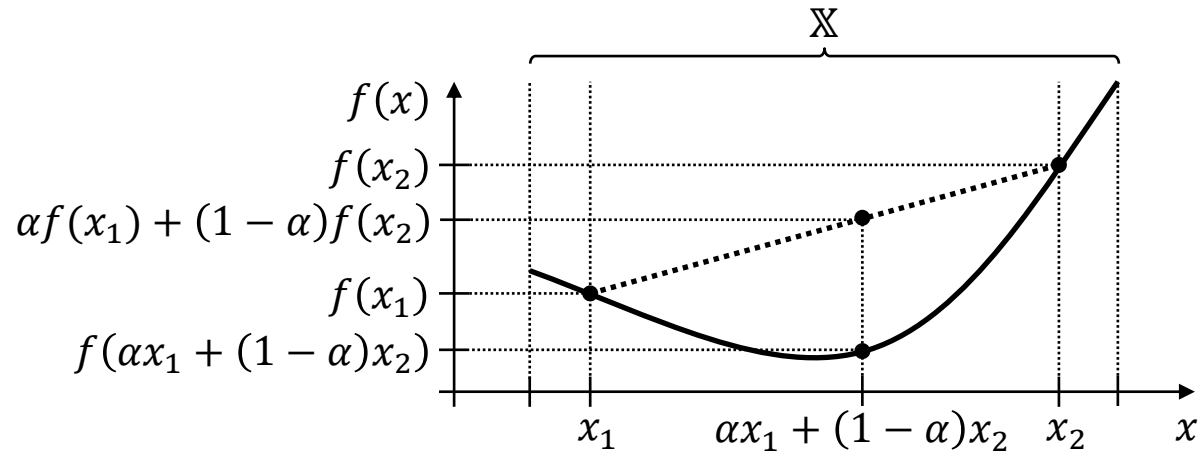
**Definition 3.9** A function  $f: \mathbb{X} \rightarrow \mathbb{R}$  is **(strictly) concave** on a convex set  $\mathbb{X}$  if  $-f$  is (strictly) convex.

- **Interpretation**

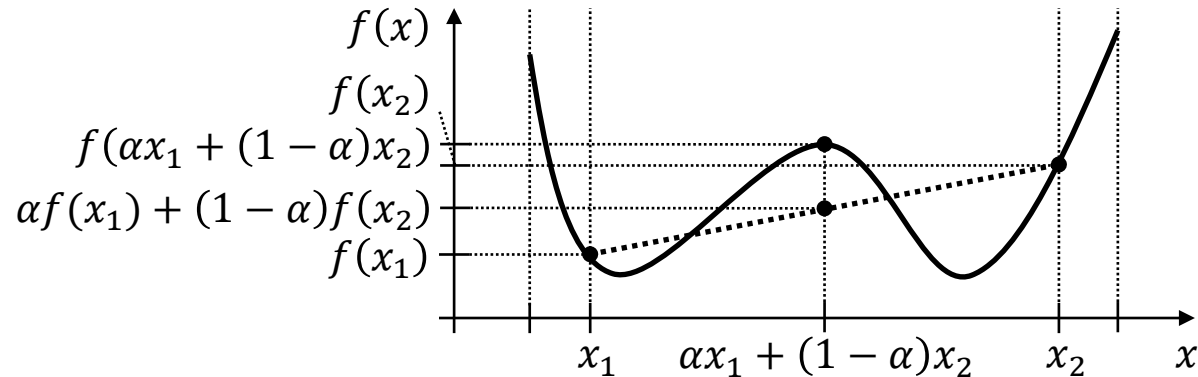
- A function  $f$  is convex if the secant connecting two arbitrary points  $(x_1, f(x_1))$  and  $(x_2, f(x_2))$  lies on or above the graph of  $f$



## Convex Functions



convex



not convex

## Convex Functions

- **Example**

- Is the function  $f(\mathbf{x}) = x_1 x_2$  convex on  $\mathbb{X} = \{\mathbf{x} | x_1 \geq 0, x_2 \geq 0\}$ ?

- Consider the points  $\mathbf{x}_1 = (1 \ 2)^T \in \mathbb{X}$  and  $\mathbf{x}_2 = (2 \ 1)^T \in \mathbb{X}$ , then

$$\alpha \mathbf{x}_1 + (1 - \alpha) \mathbf{x}_2 = \alpha \begin{pmatrix} 1 \\ 2 \end{pmatrix} + (1 - \alpha) \begin{pmatrix} 2 \\ 1 \end{pmatrix} = \begin{pmatrix} \alpha + 2 - 2\alpha \\ 2\alpha + 1 - \alpha \end{pmatrix} = \begin{pmatrix} 2 - \alpha \\ 1 + \alpha \end{pmatrix}$$

$$f(\alpha \mathbf{x}_1 + (1 - \alpha) \mathbf{x}_2) = (2 - \alpha)(1 + \alpha) = 2 + \alpha - \alpha^2$$

$$\alpha f(\mathbf{x}_1) + (1 - \alpha) f(\mathbf{x}_2) = 2\alpha + 2(1 - \alpha) = 2$$

- Consider e.g.  $\alpha = \frac{1}{2}$ , then

$$f\left(\frac{1}{2} \mathbf{x}_1 + \frac{1}{2} \mathbf{x}_2\right) = 2 + \frac{1}{2} - \frac{1}{4} = \frac{9}{4} > \frac{1}{2} f(\mathbf{x}_1) + \frac{1}{2} f(\mathbf{x}_2) = 2$$

- The function  $f(\mathbf{x})$  is not convex on  $\mathbb{X}$

## Convex Functions

- **Properties**

- (1)  $f_i(x)$  convex on  $\mathbb{X}$ ,  $\alpha_i \geq 0$ ,  $i \in \{1, \dots, N\} \Rightarrow f(x) = \sum_{i=1}^N \alpha_i f_i(x)$  convex on  $\mathbb{X}$
- (2)  $f(x)$  convex on  $\mathbb{X}$ ,  $x_1, x_2 \in \mathbb{X} \Rightarrow f(\alpha x_1 + (1 - \alpha)x_2)$  convex on  $\mathbb{X}$  for  $\alpha \in [0, 1]$
- (3)  $f(x)$  convex on  $\mathbb{X} \Rightarrow \{x \in \mathbb{X} | f(x) \leq 0\}$  convex
- (4)  $\{x \in \mathbb{X} | f(x) \leq 0\}$  convex  $\nRightarrow f(x)$  convex on  $\mathbb{X}$
- (5)  $f(x) \in \mathcal{C}^1$  convex on  $\mathbb{X} \Leftrightarrow f(x_2) \geq f(x_1) + (x_2 - x_1)^T \nabla f(x_1) \quad \forall x_1, x_2 \in \mathbb{X}$
- (6)  $f(x) \in \mathcal{C}^1$  strictly convex on  $\mathbb{X} \Leftrightarrow f(x_2) > f(x_1) + (x_2 - x_1)^T \nabla f(x_1) \quad \forall x_1, x_2 \in \mathbb{X}, x_1 \neq x_2$
- (7)  $f(x) \in \mathcal{C}^2$  convex on  $\mathbb{X} \Leftrightarrow H_f(x) \succcurlyeq 0 \quad \forall x \in \mathbb{X}$
- (8)  $f(x) \in \mathcal{C}^2$  strictly convex on  $\mathbb{X} \Leftrightarrow H_f(x) \succ 0 \quad \forall x \in \mathbb{X}$

- **Example**

- When is the **quadratic form**  $f(x) = x^T P x$  with  $P = P^T$  **convex** and **strictly convex** on  $\mathbb{R}^n$ ?
- It is  $H_f(x) = P$ . Thus,  $f(x)$  is convex on  $\mathbb{R}^n$  iff  $P \succcurlyeq 0$  and strictly convex on  $\mathbb{R}^n$  iff  $P \succ 0$  (!)

## Convex Optimization Problem

**Problem 3.3** Consider the **nonlinear optimization problem**

$$\min_x f(x) \quad \text{with } f: \mathbb{R}^n \rightarrow \mathbb{R} \quad \text{cost function} \quad (3.6)$$

$$\text{subject to } x \in \mathbb{X} \quad \text{with } \mathbb{X} = \{x \in \mathbb{R}^n \mid h(x) = 0, g(x) \leq 0\} \quad \text{feasible set} \quad (3.7)$$

The problem is **convex** if the **feasible set**  $\mathbb{X}$  is **convex** and the **cost function**  $f$  is **convex** on the feasible set  $\mathbb{X}$ . It is furthermore **strictly convex** if the **cost function**  $f$  is also **strictly convex** on the feasible set  $\mathbb{X}$ .

- **Remark**

- Proving convexity of the feasible set  $\mathbb{X}$  is very involved except in special cases
- For example, if the functions  $h_i(x), i \in \{1, \dots, m\}$  are linear and the functions  $g_j(x), j \in \{1, \dots, p\}$  are convex on  $\mathbb{X}$ , then the feasible set  $\mathbb{X}$  is an intersection of convex sets and therefore convex

**Theorem 3.2** Let  $f: \mathbb{X} \rightarrow \mathbb{R}$  be a convex function defined on the convex set  $\mathbb{X}$ . Then **each local minimum** of  $f$  on  $\mathbb{X}$  is **also a global minimum** of  $f$  on  $\mathbb{X}$  and the **set of global minima** of  $f$  on  $\mathbb{X}$  is **convex**.

## Definitions

**Definition 3.10** An **inequality constraint**  $g_j(x) \leq 0$  is denoted as **active** at a feasible point  $x \in \mathbb{X}$  if  $g_j(x) = 0$  and as **inactive** at a feasible point  $x \in \mathbb{X}$  if  $g_j(x) < 0$ .

- **Remark**

- Active inequality constraints will be denoted in the following by  $\mathbf{g}^a: \mathbb{R}^n \rightarrow \mathbb{R}^{p^a}, \mathbf{g}^a(x) = \mathbf{0}$
- Inactive inequality constraints will be denoted in the following by  $\mathbf{g}^i: \mathbb{R}^n \rightarrow \mathbb{R}^{p^i}, \mathbf{g}^i(x) < \mathbf{0}$
- Note that  $p^a + p^i = p$

**Definition 3.11** The feasible point  $x \in \mathbb{X}$  is denoted as **regular point** if the vectors

$$\nabla h_i(x), i \in \{1, \dots, m\} \text{ and } \nabla g_j^a(x), j \in \{1, \dots, p^a\}$$

are linearly independent.

## Karush-Kuhn-Tucker (KKT) Conditions

**Theorem 3.3** Let  $\mathbf{x}^* \in \mathbb{R}^n$  be a regular point and a local minimizer to Problem 3.1 and introduce the function  $L(\mathbf{x}, \boldsymbol{\lambda}, \boldsymbol{\mu}) = f(\mathbf{x}) + \boldsymbol{\lambda}^T \mathbf{h}(\mathbf{x}) + \boldsymbol{\mu}^T \mathbf{g}(\mathbf{x})$ . Then there exist  $\boldsymbol{\lambda}^* \in \mathbb{R}^m$  and  $\boldsymbol{\mu}^* \in \mathbb{R}^p$  such that

- (1)  $\nabla_{\mathbf{x}} L(\mathbf{x}^*, \boldsymbol{\lambda}^*, \boldsymbol{\mu}^*) = \nabla f(\mathbf{x}^*) + \mathbf{J}_h^T(\mathbf{x}^*) \boldsymbol{\lambda}^* + \mathbf{J}_g^T(\mathbf{x}^*) \boldsymbol{\mu}^* = \nabla f(\mathbf{x}^*) + \sum_{i=1}^m \nabla h_i(\mathbf{x}^*) \lambda_i^* + \sum_{j=1}^p \nabla g_j(\mathbf{x}^*) \mu_j^* = \mathbf{0}$
- (2)  $\nabla_{\boldsymbol{\lambda}} L(\mathbf{x}^*, \boldsymbol{\lambda}^*, \boldsymbol{\mu}^*) = \mathbf{h}(\mathbf{x}^*) = \mathbf{0}$
- (3)  $\mathbf{g}(\mathbf{x}^*) \leq \mathbf{0}$
- (4)  $\mathbf{g}^T(\mathbf{x}^*) \boldsymbol{\mu}^* = 0$
- (5)  $\boldsymbol{\mu}^* \geq \mathbf{0}$ .

- **Remarks**

- **No constraints?** Only the green term is relevant.
- **Only equality constraints?** Only the green term and blue terms are relevant.
- Condition (4) can also be written as  $g_j(\mathbf{x}^*) \mu_j^* = 0, j \in \{1, \dots, p\}$

## Karush-Kuhn-Tucker (KKT) Conditions

- Symbols

- The function  $L(x, \lambda, \mu) = f(x) + \lambda^T h(x) + \mu^T g(x)$  is called **Lagrangian**
- The vector  $\lambda$  is called **Lagrange multiplier**
- The vector  $\mu$  is called **Karush-Kuhn-Tucker multiplier**

- Properties

- $\mu_j^* = 0$  if  $g_j(x^*) < 0$  (i.e. if the inequality constraint is **inactive**) due to conditions (3) to (5)
- $\mu_j^* \geq 0$  if  $g_j(x^*) = 0$  (i.e. if the inequality constraint is **active**) due to conditions (3) to (5)
- $\mu_j < 0$  and  $g_j(x) = 0$  (i.e. the inequality constraint is **active**) while (1) to (4) fulfilled indicates that the cost  $f(x)$  can be reduced by setting  $g_j(x) < 0$  (i.e. by setting the inequality constraint **inactive**)

- Remarks

- The KKT conditions presume **constraint qualification**. Constraint qualification is ensured in most optimization problems, e.g. if  $h$  and  $g^a$  are linear, see [PLB12, p. 78] for details.

## Karush-Kuhn-Tucker (KKT) Conditions

- **Remarks**

- The KKT conditions are **only necessary** for general **nonlinear optimization problems** (Problem 3.1)
- The KKT conditions are **necessary** and **sufficient** for **convex optimization problems** (Problem 3.3)
- The KKT conditions can usually be evaluated **analytically** for **simple optimization problems**
- The KKT conditions must generally be evaluated **numerically** for **complex optimization problems**

- **Example**

- **Maximization** of the **area of a right triangle** with legs  $x_1$  and  $x_2$  and a given hypotenuse  $l$  (Slide 3-4)
- Cost function  $f(\mathbf{x}) = -\frac{1}{2}x_1x_2$
- Constraints  $h(\mathbf{x}) = x_1^2 + x_2^2 - l^2 = 0$ ,  $g_1(\mathbf{x}) = -x_1 \leq 0$ ,  $g_2(\mathbf{x}) = -x_2 \leq 0$
- Lagrangian  $L(\mathbf{x}, \lambda, \boldsymbol{\mu}) = -\frac{1}{2}x_1x_2 + \lambda(x_1^2 + x_2^2 - l^2) - \mu_1x_1 - \mu_2x_2$
- An **analytical solution** can be obtained by analyzing all **combinations** of **active** and **inactive inequality constraints** to determine **candidate solutions** and then comparing the candidate solutions w.r.t. cost



## Karush-Kuhn-Tucker (KKT) Conditions

- **Example**

- **Case 1**  $g_1(\mathbf{x}^*) < 0$  (**inactive**),  $g_2(\mathbf{x}^*) < 0$  (**inactive**), then  $\mu_1^* = \mu_2^* = 0$

$$\left. \begin{aligned} \frac{\partial}{\partial x_1} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) &= -\frac{1}{2}x_2^* + 2\lambda^*x_1^* = 0 \\ \frac{\partial}{\partial x_2} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) &= -\frac{1}{2}x_1^* + 2\lambda^*x_2^* = 0 \\ \frac{\partial}{\partial \lambda} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) &= x_1^{*2} + x_2^{*2} - l^2 = 0 \end{aligned} \right\} \begin{aligned} &x_1^* = x_2^* \\ &x_1^* = x_2^* = \frac{l}{\sqrt{2}}, \lambda^* = \frac{1}{4} \quad \checkmark \end{aligned}$$

- **Case 2**  $g_1(\mathbf{x}^*) = 0$  (**active**),  $g_2(\mathbf{x}^*) < 0$  (**inactive**), then  $\mu_1^* \geq 0, \mu_2^* = 0$

$$\left. \begin{aligned} \frac{\partial}{\partial x_1} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) &= -\frac{1}{2}x_2^* + 2\lambda^*x_1^* - \mu_1^* = 0 \\ \frac{\partial}{\partial x_2} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) &= -\frac{1}{2}x_1^* + 2\lambda^*x_2^* = 0 \\ \frac{\partial}{\partial \lambda} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) &= x_1^{*2} + x_2^{*2} - l^2 = 0 \\ -x_1^*\mu_1^* &= 0 \end{aligned} \right\} \begin{aligned} &x_1^* = 0, x_2^* = \pm l, \mu_1^* = \mp \frac{1}{2}l \quad \times \\ &\text{or} \\ &\mu_1^* = 0, x_1^* = x_2^* = \frac{l}{\sqrt{2}}, \lambda^* = \frac{1}{4} \quad \checkmark \\ &x_1^* = 0 \text{ or } \mu_1^* = 0 \end{aligned}$$

## Karush-Kuhn-Tucker (KKT) Conditions

- Example

- **Case 3**  $g_1(\mathbf{x}^*) < 0$  (inactive),  $g_2(\mathbf{x}^*) = 0$  (active), then  $\mu_1^* = 0, \mu_2^* \geq 0$

Analogous to Case 2

- **Case 4**  $g_1(\mathbf{x}^*) = 0$  (active),  $g_2(\mathbf{x}^*) = 0$  (active), then  $\mu_1^* \geq 0, \mu_2^* \geq 0$

$$\frac{\partial}{\partial x_1} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) = -\frac{1}{2}x_2^* + 2\lambda^*x_1^* - \mu_1^* = 0$$

$$\frac{\partial}{\partial x_2} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) = -\frac{1}{2}x_1^* + 2\lambda^*x_2^* - \mu_2^* = 0$$

$$\frac{\partial}{\partial \lambda} L(\mathbf{x}^*, \lambda^*, \boldsymbol{\mu}^*) = x_1^{*2} + x_2^{*2} - l^2 = 0$$

$$-x_1^*\mu_1^* = 0$$

$$-x_2^*\mu_2^* = 0$$

$$\left. \begin{array}{l} x_1^* = 0 \text{ or } x_2^* = 0 \text{ ✗} \\ \mu_1^* = 0 \text{ and } \mu_2^* = 0 \end{array} \right\}$$

$$\left. \begin{array}{l} \mu_1^* = 0, \mu_2^* = 0, \\ x_1^* = x_2^* = \frac{l}{\sqrt{2}}, \lambda^* = \frac{1}{4} \checkmark \end{array} \right\}$$

- The maximum area is obtained for the legs  $x_1^* = x_2^* = \frac{l}{\sqrt{2}}$  and has the value  $\frac{1}{2}x_1^*x_2^* = \frac{l^2}{4}$

## Hyperplanes and Half-Spaces

**Definition 3.12** The set  $\{\mathbf{x} \in \mathbb{R}^n | \mathbf{a}^T \mathbf{x} = b\}$  with  $\mathbf{a} = (a_1 \ a_2 \ \dots \ a_n)^T \in \mathbb{R}^n \setminus \{\mathbf{0}\}$ ,  $b \in \mathbb{R}$  is called **hyperplane**.

- **Remarks**

- The vector  $\mathbf{a}$  is orthogonal to the hyperplane and therefore called normal
- For  $b = 0$  the hyperplane contains the origin and thus is a subspace of  $\mathbb{R}^n$
- For  $n = 2$  the hyperplane becomes  $a_1 x_1 + a_2 x_2 = b$  and thus describes a line in  $\mathbb{R}^2$
- For  $n = 3$  the hyperplane becomes  $a_1 x_1 + a_2 x_2 + a_3 x_3 = b$  and thus describes a plane in  $\mathbb{R}^3$
- A hyperplane is a **convex set**

**Definition 3.13** The set  $\{\mathbf{x} \in \mathbb{R}^n | \mathbf{a}^T \mathbf{x} \leq b\}$  with  $\mathbf{a} = (a_1 \ a_2 \ \dots \ a_n)^T \in \mathbb{R}^n \setminus \{\mathbf{0}\}$ ,  $b \in \mathbb{R}$  is called **half-space**.

- **Remarks**

- Partly  $\{\mathbf{x} \in \mathbb{R}^n | \mathbf{a}^T \mathbf{x} \geq b\}$  is called positive half-space and  $\{\mathbf{x} \in \mathbb{R}^n | \mathbf{a}^T \mathbf{x} \leq b\}$  negative half-space
- A half-space is a **convex set**

## Linear Varieties

### Tutorial

**Definition 3.14** The set  $\{\mathbf{x} \in \mathbb{R}^n | \mathbf{A}\mathbf{x} = \mathbf{b}\}$  with  $\mathbf{A} \in \mathbb{R}^{m \times n}$ ,  $\mathbf{b} \in \mathbb{R}^m$  is called **linear variety** or **flat**.

- **Remarks**

- A linear variety can also be written as  $\mathbf{a}_i^T \mathbf{x} = b_i, i \in \{1, \dots, m\}$  ( $\mathbf{a}_i^T$  rows of  $\mathbf{A}$ ,  $b_i$  components of  $\mathbf{b}$ )
- A linear variety is therefore the **intersection of  $m$  hyperplanes**
- A linear variety is therefore a **convex set** (intersection of convex sets, cf. Slide 3-7, Property (3))

- **Examples**

$$(a_1 \quad a_2 \quad a_3) \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = b \quad \Leftrightarrow \quad a_1 x_1 + a_2 x_2 + a_3 x_3 = b \quad \text{describes a plane in } \mathbb{R}^3$$

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix} \quad \Leftrightarrow \quad \begin{aligned} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 &= b_1 \\ a_{21}x_1 + a_{22}x_2 + a_{23}x_3 &= b_2 \end{aligned} \quad \text{describes a line in } \mathbb{R}^{3*}$$

\* if  $\mathbf{a}_1$  and  $\mathbf{a}_2$  are linearly independent

## Polyhedra and Polytopes

Tutorial

**Definition 3.15** The set  $\{x \in \mathbb{R}^n | Ax \leq b\}$  with  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$  is called **polyhedron**.

- **Remarks**

- A polyhedron can also be written as  $a_i^T x \leq b_i, i \in \{1, \dots, m\}$  ( $a_i^T$  rows of  $A$ ,  $b_i$  components of  $b$ )
- A polyhedron is therefore the **intersection of  $m$  half-spaces**
- A polyhedron is therefore a **convex set** (intersection of convex sets, cf. Slide 3-7, Property (3))
- The  $0, 1, \dots, (k - 1)$ -dim. polyhedra forming the boundary of a  $k$ -dim. polyhedron are called **faces**
- The faces of dimension  $0, 1, (k - 2)$ , and  $(k - 1)$  are called **vertices, edges, ridges, and facets**

**Definition 3.16** A **polytope** is a **bounded polyhedron** (i.e.  $\exists \alpha \in \mathbb{R}: \|y\| \leq \alpha \forall y \in \{x \in \mathbb{R}^n | Ax \leq b\}$ ).

- **Remark**

- Note that the definition of a polytope is not unique in the literature
- Definition 3.16 is based on [BBM15, Section 3.1]

## Polyhedra and Polytopes

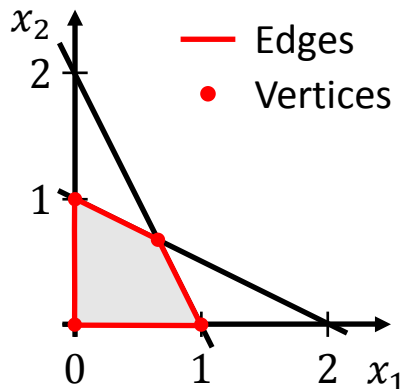
Tutorial

**Definition 3.17** The set  $\{x \in \mathbb{R}^n | x = \sum_{i=1}^V \alpha_i V_i, 0 \leq \alpha_i \leq 1, \sum_{i=1}^V \alpha_i = 1\}$  is called **polytope** where  $V_i \in \mathbb{R}^n$  are the **vertices** and  $V$  is the number of vertices.

- **Remark**

- The representation according to Definition 3.15 is called **half-space representation** (H-representation)
- The representation according to Definition 3.17 is called **vertex representation** (V-representation)

- **Example**



$$\begin{pmatrix} -1 & 0 \\ 0 & -1 \\ 2 & 1 \\ 0.5 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} \leq \begin{pmatrix} 0 \\ 0 \\ 2 \\ 1 \end{pmatrix} \Leftrightarrow \begin{cases} x_1 \geq 0 \\ x_2 \geq 0 \\ x_2 \leq -2x_1 + 2 \\ x_2 \leq -0.5x_1 + 1 \end{cases}$$

The polyhedron is bounded and therefore a polytope

The polyhedron is unbounded if the first or second row are removed

## Linear Programming Problem

**Problem 3.4** The **linear programming problem** is defined as

$$\min_x \mathbf{c}^T \mathbf{x} \quad \text{with } \mathbf{c}, \mathbf{x} \in \mathbb{R}^n \quad \text{linear cost function} \quad (3.8)$$

$$\text{subject to } \begin{cases} \mathbf{A}_{\text{eq}} \mathbf{x} = \mathbf{b}_{\text{eq}} & \text{with } \mathbf{A}_{\text{eq}} \in \mathbb{R}^{m \times n}, \mathbf{b}_{\text{eq}} \in \mathbb{R}^m \\ \mathbf{A}_{\text{ieq}} \mathbf{x} \leq \mathbf{b}_{\text{ieq}} & \text{with } \mathbf{A}_{\text{ieq}} \in \mathbb{R}^{p \times n}, \mathbf{b}_{\text{ieq}} \in \mathbb{R}^p \end{cases} \quad \begin{array}{l} \text{linear equality constraints} \\ \text{linear inequality constraints} \end{array} \quad \begin{array}{l} (3.9) \\ (3.10) \end{array}$$

- **Remarks**

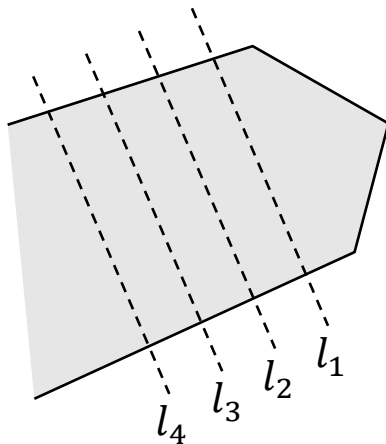
- The linear cost function is a convex function. The linear equality constraints (linear variety) and the linear inequality constraints (polyhedron) are convex sets and thus the feasible set is a convex set.
- The linear programming problem is therefore **convex**.
- Several methods exist for solving the linear programming problem. The most important are the **simplex method** (exponential complexity) and **Karmarkar's method** (polynomial complexity)
- The linear programming problem can be solved in **MATLAB/Optimization Toolbox** with `linprog`

## Characterization of the Solution

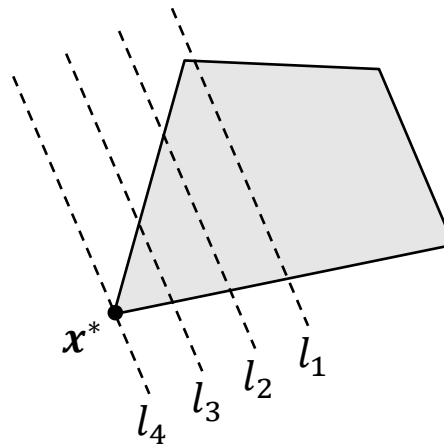
- Cases

- (1) The cost is unbounded, i.e.  $\mathbf{c}^T \mathbf{x}^* = -\infty$
- (2) The cost is bounded, i.e.  $\mathbf{c}^T \mathbf{x}^* > -\infty$ , the minimizer  $\mathbf{x}^*$  unique (vertex of the feasible set for  $\mathbb{R}^2$ )
- (3) The cost is bounded, i.e.  $\mathbf{c}^T \mathbf{x}^* > -\infty$ , the minimizer  $\mathbf{x}^*$  not unique (edge of the feasible set for  $\mathbb{R}^2$ )

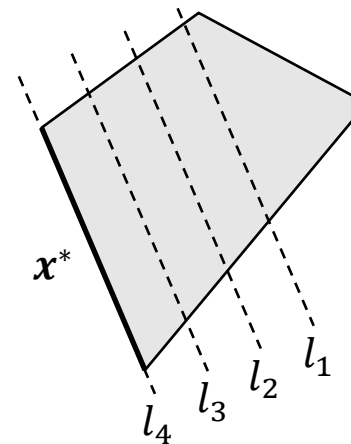
- Graphical Interpretation in  $\mathbb{R}^2$



Case (1)



Case (2)



Case (3)

--- Level curves  
 $\mathbf{c}^T \mathbf{x} = l_i$   
 $l_i > l_{i+1}$   
 $i \in \mathbb{N}$   
 (parallel lines)



## Quadratic Programming Problem

**Problem 3.5** The **quadratic programming problem** is defined as

$$\min_x \frac{1}{2} \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{f}^T \mathbf{x} \quad \text{with } \mathbf{H} \in \mathbb{R}^{n \times n}, \mathbf{H} = \mathbf{H}^T \succcurlyeq \mathbf{0}, \mathbf{f} \in \mathbb{R}^n \text{ quadratic cost function} \quad (3.11)$$

$$\text{subject to } \begin{cases} \mathbf{A}_{\text{eq}} \mathbf{x} = \mathbf{b}_{\text{eq}} & \text{with } \mathbf{A}_{\text{eq}} \in \mathbb{R}^{m \times n}, \mathbf{b}_{\text{eq}} \in \mathbb{R}^m & \text{linear equality constr.} \\ \mathbf{A}_{\text{ieq}} \mathbf{x} \leq \mathbf{b}_{\text{ieq}} & \text{with } \mathbf{A}_{\text{ieq}} \in \mathbb{R}^{p \times n}, \mathbf{b}_{\text{ieq}} \in \mathbb{R}^p & \text{linear inequality constr.} \end{cases} \quad (3.12)$$

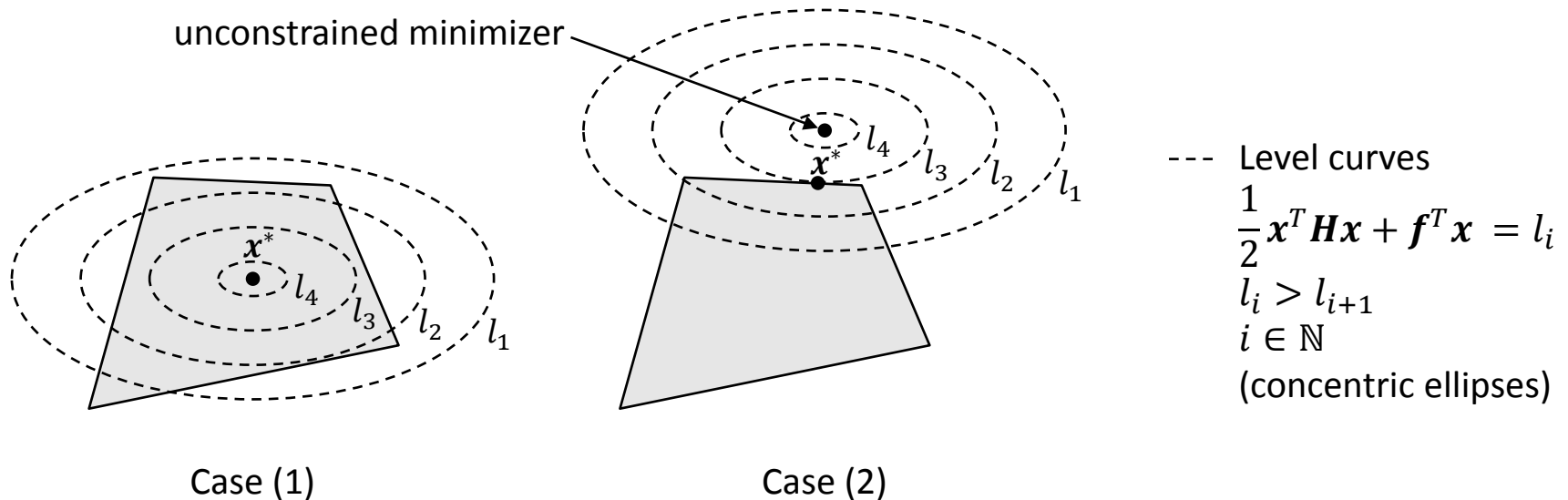
$$(3.13)$$

- **Remarks**

- The quadratic cost function is a convex function for  $\mathbf{H} \succcurlyeq \mathbf{0}$  and a strictly convex function for  $\mathbf{H} \succ \mathbf{0}$ . The linear equality constraints (linear variety) and the linear inequality constraints (polyhedron) are convex sets and thus the feasible set is a convex set.
- The quadratic programming problem is therefore **convex for  $\mathbf{H} \succcurlyeq \mathbf{0}$**  and **strictly convex for  $\mathbf{H} \succ \mathbf{0}$**
- The quadratic programming problem can be solved in **MATLAB/Optimization Toolbox** with `quadprog`

## Characterization of the Solution

- Cases
  - (1) The cost is bounded and the minimizer  $\mathbf{x}^*$  lies strictly inside the feasible set
  - (2) The cost is bounded and the minimizer  $\mathbf{x}^*$  lies on the boundary of the feasible set
- Graphical Interpretation in  $\mathbb{R}^2$



## Solution based on the Active Set Method

- Approach

- Consider that a **feasible point**  $\mathbf{x}^{(i)}$  and related **active inequality constraints**  $\mathbf{A}_{\text{ieq}}^{\text{a}} \mathbf{x}^{(i)} = \mathbf{b}_{\text{ieq}}^{\text{a}}$  are known
- Find an **improved point**  $\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}$  considering only  $\mathbf{A}_{\text{eq}} \Delta \mathbf{x}^{(i)} = \mathbf{0}$  and  $\mathbf{A}_{\text{ieq}}^{\text{a}} \Delta \mathbf{x}^{(i)} = \mathbf{0}$

For the improved point  $\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}$  the cost function becomes

$$\begin{aligned} f(\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}) &= \frac{1}{2} (\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)})^T \mathbf{H} (\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}) + \mathbf{f}^T (\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}) \\ &= f(\mathbf{x}^{(i)}) + \frac{1}{2} \Delta \mathbf{x}^{(i)T} \mathbf{H} \Delta \mathbf{x}^{(i)} + \underbrace{(\mathbf{f}^T + \mathbf{x}^{(i)T} \mathbf{H})}_{\mathbf{f}^{(i)T}} \Delta \mathbf{x}^{(i)} \\ &= f(\mathbf{x}^{(i)}) + \frac{1}{2} \Delta \mathbf{x}^{(i)T} \mathbf{H} \Delta \mathbf{x}^{(i)} + \mathbf{f}^{(i)T} \Delta \mathbf{x}^{(i)} \end{aligned}$$

The improved point thus results from the **optimization problem**

$$\begin{aligned} \min_{\Delta \mathbf{x}^{(i)}} & \frac{1}{2} \Delta \mathbf{x}^{(i)T} \mathbf{H} \Delta \mathbf{x}^{(i)} + \mathbf{f}^{(i)T} \Delta \mathbf{x}^{(i)} \\ \text{subject to } & \mathbf{A}_{\text{eq}} \Delta \mathbf{x}^{(i)} = \mathbf{0}, \mathbf{A}_{\text{ieq}}^{\text{a}} \Delta \mathbf{x}^{(i)} = \mathbf{0} \end{aligned} \tag{3.14}$$

## Solution based on the Active Set Method

- Approach

The **Lagrangian** to the optimization problem (3.14) obeys

$$L(\Delta \mathbf{x}^{(i)}, \boldsymbol{\lambda}^{(i+1)}, \boldsymbol{\mu}^{(i+1)}) = \frac{1}{2} \Delta \mathbf{x}^{(i)T} \mathbf{H} \Delta \mathbf{x}^{(i)} + \mathbf{f}^{(i)T} \Delta \mathbf{x}^{(i)} + \boldsymbol{\lambda}^{(i+1)T} \mathbf{A}_{\text{eq}} \Delta \mathbf{x}^{(i)} + \boldsymbol{\mu}^{(i+1)T} \mathbf{A}_{\text{ieq}}^a \Delta \mathbf{x}^{(i)}$$

The **KKT conditions** (only (1) and (2) relevant) to optimization problem (3.14) are then given by

$$\nabla_{\Delta \mathbf{x}^{(i)}} L(\Delta \mathbf{x}^{(i)}, \boldsymbol{\lambda}^{(i+1)}, \boldsymbol{\mu}^{(i+1)}) = \mathbf{H} \Delta \mathbf{x}^{(i)} + \mathbf{f}^{(i)} + \mathbf{A}_{\text{eq}}^T \boldsymbol{\lambda}^{(i+1)} + \mathbf{A}_{\text{ieq}}^{aT} \boldsymbol{\mu}^{(i+1)} = \mathbf{0}$$

$$\nabla_{\boldsymbol{\lambda}^{(i+1)}} L(\Delta \mathbf{x}^{(i)}, \boldsymbol{\lambda}^{(i+1)}, \boldsymbol{\mu}^{(i+1)}) = \mathbf{A}_{\text{eq}} \Delta \mathbf{x}^{(i)} = \mathbf{0}$$

$$\nabla_{\boldsymbol{\mu}^{(i+1)}} L(\Delta \mathbf{x}^{(i)}, \boldsymbol{\lambda}^{(i+1)}, \boldsymbol{\mu}^{(i+1)}) = \mathbf{A}_{\text{ieq}}^a \Delta \mathbf{x}^{(i)} = \mathbf{0}$$

which can be written as a **system of linear equations (SLE)**

$$\begin{pmatrix} \mathbf{H} & \mathbf{A}_{\text{eq}}^T & \mathbf{A}_{\text{ieq}}^{aT} \\ \mathbf{A}_{\text{eq}} & \mathbf{0} & \mathbf{0} \\ \mathbf{A}_{\text{ieq}}^a & \mathbf{0} & \mathbf{0} \end{pmatrix} \begin{pmatrix} \Delta \mathbf{x}^{(i)} \\ \boldsymbol{\lambda}^{(i+1)} \\ \boldsymbol{\mu}^{(i+1)} \end{pmatrix} = \begin{pmatrix} -\mathbf{f}^{(i)} \\ \mathbf{0} \\ \mathbf{0} \end{pmatrix} \quad (3.15)$$

## Solution based on the Active Set Method

- **Approach**

The **solution** of the optimization problem (3.14) finally follows by **solving the SLE** (3.15), e.g. based on the inverse (slow) or QR/LU decomposition (fast), cf. [Mac02, Section 3.3], [PLB12, Section 5.4.3]

- Check if the improved point  $\mathbf{x}^{(i)} + \Delta\mathbf{x}^{(i)}$  is a minimizer of the original quadratic programming problem (Problem 3.5) by **evaluating the KKT conditions** (1) to (5)
- If not, then consider another improved point

- **Remarks**

- Solving a **quadratic programming problem** with **only equality constraints** is obviously quite **easy**
- The **active set method** is based on solving quadratic programming problems with equality constraints iteratively for different combinations of active inequality constraints (active sets)
- This can be formalized as an **algorithm**

## Solution based on the Active Set Method

- Algorithm

1. Determine initial feasible point  $\mathbf{x}^{(0)}$  and active inequality constraints  $\mathbf{A}_{\text{ieq}}^a \mathbf{x}^{(0)} = \mathbf{b}_{\text{ieq}}^a$  (active set)
2. Set  $i := 0$
3. Determine  $\Delta \mathbf{x}^{(i)}$ ,  $\boldsymbol{\lambda}^{(i+1)}$ , and  $\boldsymbol{\mu}^{(i+1)}$  by solving the SLE (3.15)
4. Evaluate the KKT conditions (1) to (5) for Problem 3.5
  - a. If  $\Delta \mathbf{x}^{(i)} = \mathbf{0}$  and  $\boldsymbol{\mu}^{(i+1)} \geq \mathbf{0}$ , then stop since  $\mathbf{x}^{(i)}$  is a feasible global minimizer for Problem 3.5
  - b. If  $\Delta \mathbf{x}^{(i)} = \mathbf{0}$  and at least one  $\boldsymbol{\mu}^{(i+1)} < \mathbf{0}$ , then set  $\mathbf{x}^{(i+1)} := \mathbf{x}^{(i)}$  and remove the active inequality constraint with the smallest  $\boldsymbol{\mu}^{(i+1)}$  from the active set
  - c. If  $\Delta \mathbf{x}^{(i)} \neq \mathbf{0}$  and  $\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}$  feasible, then set  $\mathbf{x}^{(i+1)} := \mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}$  and retain the active set
  - d. If  $\Delta \mathbf{x}^{(i)} \neq \mathbf{0}$  and  $\mathbf{x}^{(i)} + \Delta \mathbf{x}^{(i)}$  infeasible, then find the largest  $\alpha^{(i)} > 0$  for which  $\mathbf{x}^{(i+1)} := \mathbf{x}^{(i)} + \alpha^{(i)} \Delta \mathbf{x}^{(i)}$  is feasible and add resulting active inequality constraint to active set
5. Set  $i := i + 1$  and go to 3.

## Solution based on the Active Set Method

- Remarks
  - An **initial feasible point** can be determined from a **linear** or **quadratic programming problem**, see [Mac02, Section 3.3] for details
  - The **variables  $x^{(i+1)}$**  resulting after each iteration  $i$  are **feasible points** of the original quadratic programming problem (Problem 3.5), allowing an **early termination** (relevant for MPC)
  - A **warm start**, i.e. an initialization of the iteration with a point which is known to be close to the minimizer (initial guess) for reducing the number of iterations, is **straightforward** (relevant for MPC)
  - The active set method has **exponential complexity**

## Solution based on the Interior Point Method

- Approach

- Transform the **constrained optimization problem** to an **unconstrained optimization problem**, i.e.

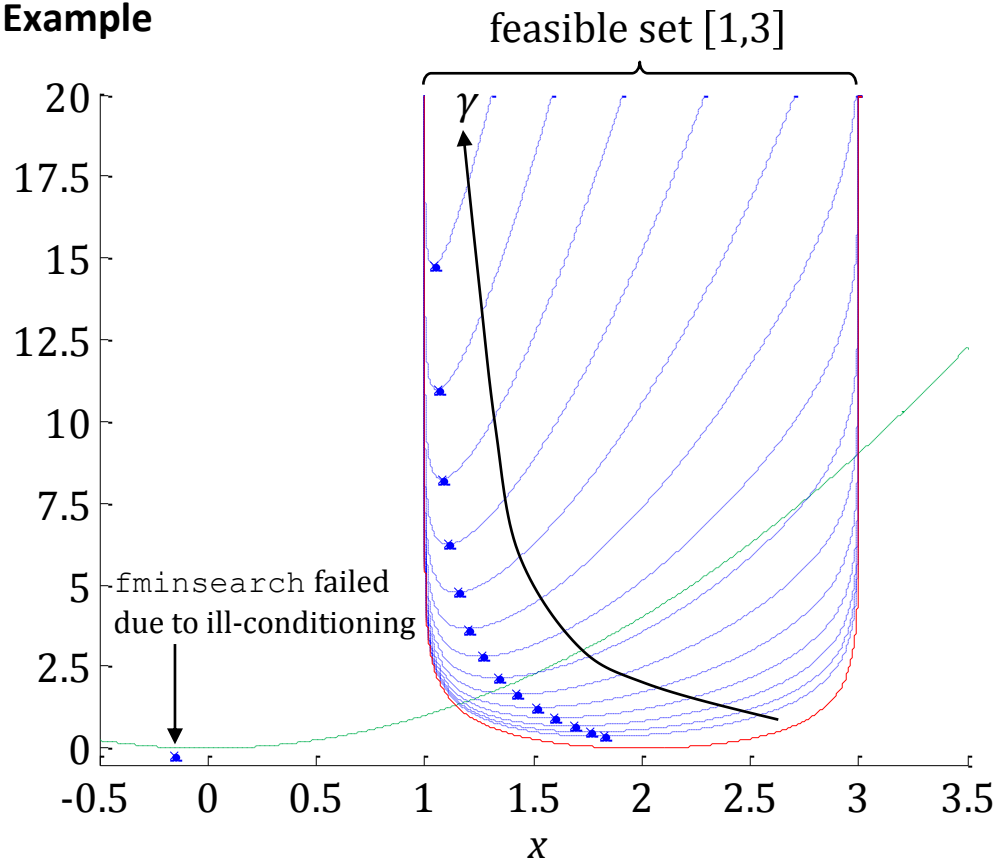
$$\begin{aligned} \min_{\mathbf{x}} \frac{1}{2} \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{f}^T \mathbf{x} \quad & \rightarrow \quad \min_{\mathbf{x}} \gamma \left( \frac{1}{2} \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{f}^T \mathbf{x} \right) - \underbrace{\sum_{j=1}^p \ln(b_{\text{ieq},j} - \mathbf{a}_{\text{ieq},j}^T \mathbf{x})}_{\text{barrier function}} \\ \text{subject to } \mathbf{A}_{\text{ieq}} \mathbf{x} &\leq \mathbf{b}_{\text{ieq}} \end{aligned}$$

- The **barrier function** is finite in the interior but infinite on the boundary of the feasible set
- Let  $\mathbf{x}_{\gamma}^*$  be the minimizer of the unconstrained optimization problem for some  $\gamma > 0$  and  $\mathbf{x}^*$  be the minimizer of the constrained optimization problem. It can be shown that  $\mathbf{x}_{\gamma}^* \rightarrow \mathbf{x}^*$  as  $\gamma \rightarrow \infty$ . However, the unconstrained optimization problem becomes ill-conditioned as  $\gamma \rightarrow \infty$ .
- The **interior point method** is based on solving the unconstrained optimization problem iteratively for an increasing  $\gamma$  until  $\mathbf{x}_{\gamma}^*$  does not change significantly anymore
- The path followed by  $\mathbf{x}_{\gamma}^*$  is denoted as **central path**



## Solution based on the Interior Point Method

- Example



$$\min_{x \in [1,3]} x^2$$

$$\min_x \gamma x^2 - \ln(-1+x) - \ln(3-x)$$

—  $x^2$

—  $-\ln(-1+x) - \ln(3-x)$

—  $\gamma x^2 - \ln(-1+x) - \ln(3-x)$   
(for  $\gamma = \text{logspace}(-1, 2, 20)$ )

×  $x_\gamma^*$  (obtained with fminsearch)

## Solution based on the Interior Point Method

- Remarks

- The **equality constraint**  $A_{\text{eq}}x = b_{\text{eq}}$  can be regarded in the interior point method by reformulation into two inequality constraints  $A_{\text{eq}}x \leq b_{\text{eq}}$  and  $-A_{\text{eq}}x \leq -b_{\text{eq}}$
- The **minimizers**  $x_{\gamma}^*$  of the unconstrained optimization problem are **feasible points** of the constrained optimization problem, allowing an **early termination** of the iterations (relevant for MPC)
- The interior point method requires **modifications** to address **ill-conditioning**. The minimizers  $x_{\gamma}^*$  are usually no feasible points under the these modifications, not allowing an early termination (MPC)
- A **warm start**, i.e. an initialization of the iteration with a point which is known to be close to the minimizer (initial guess) for reducing the number of iterations, is usually **difficult** (relevant for MPC)
- The interior point method has **polynomial complexity**

## Remarks on Optimization Software

- **Overviews**

- [plato.asu.edu/guide.html](http://plato.asu.edu/guide.html)
- [yalmip.github.io/allsolvers/](http://yalmip.github.io/allsolvers/)
- [neos-guide.org/optimization-tree](http://neos-guide.org/optimization-tree)
- <https://www.coin-or.org/>

- **Modeling Languages and Solvers**

- YALMIP ([yalmip.github.io/](http://yalmip.github.io/))
- CVX ([cvxr.com/cvx/](http://cvxr.com/cvx/))
- CVXGEN ([cvxgen.com](http://cvxgen.com))
- FORCES ([forces.ethz.ch](http://forces.ethz.ch))
- $\mu$ AO-MPC ([ifatwww.et.uni-magdeburg.de/syst/muAO-MPC/](http://ifatwww.et.uni-magdeburg.de/syst/muAO-MPC/))