

CS 6110, Spring 2022, Assignment 1

Given 1/13/22 – Due 1/18/22 by 11:59 pm via your Github

NAME:

UNID:

(You **must** use this .tex as your answer template, inserting your answers after each question. You must retain the framebox so that I can easily visually locate your answers.)

(Make my grading easy by confining all your answers to a page or max 2 pages. Do not remove the clearpages inserted.)

Submission: Please submit via your Github. I'll pull and compile/make. Please also put your PDF there.

1. (10 points) Read the first 27 pages of Bradley's book. Also read the material around "CNF-conversion using gates" (AKA Tseitin transformation, CEATL, 18.3.4). (Tseitin or Tseytin is/was a Russian scientist; see https://en.wikipedia.org/wiki/Tseytin_transformation.) (This procedure is present in Bradley's book; locate it there.) Make a glossary of concepts covered such as satisfiability, validity, contradiction, equisat, etc. Study the truth-table based (1.3.1) and semantic argument-based (1.3.2) methods. Summarize all these in neat bullets in your PDF answer. Try to create your initial answer as a cheat-sheet of about a page. Later transplant it to your own place (say, a GDoc) and maintain these concepts. (We could merge them one day perhaps.)

Your Answer Here

2. (30 pts) The formula

$$a \cdot b + c$$

is given to you. (Note that in the Python BDD the syntax is different.) Convert this formula to an equisat formula using Tseitin's transformation, following the procedure in CEATL. Call the variable you introduce at the output of the "and gate" as p and call the final output z . Now, make a copy of `BDD.ipynb` found inside `pbl/` of Jove. (Jove is at <https://github.com/ganeshutah/Jove.git>.) Get rid of all the material in this file before submission (you can keep it to look at it while developing your solution). Then just have one title page "**Understanding Equisat Versus Equivalence.**" Then have just two code cells:

- A code cell

```
EquiSat = '''
Var_Order : a,b,c,p,z
fGiven = (c|(a&b))
fTseitin = ...your Tseitin-converted result...
Main_Exp : fGiven OP fTseitin # OP is -> and <- in turn
```

- The second code cell begins with `buildBDDmain` and is to draw the above.

Now answer these questions:

- (a) Which case (\rightarrow or \leftarrow) gave you a "1" node and why?
(Insert framebox here and answer.)
- (b) Within the case where you did not get a "0" node, insert that BDD which was non-0 here.
(framebox)
- (c) for all the paths to the "0" node, explain why that path exists (this is where equisat differed from equivalence). Write out your answer in neat bulleted steps per path.
(Insert framebox here and answer.)

3. (30 pts) First, begin reading Ben-Ari's SPIN book and get some practice following the commands there.

Run the program in CEATL, Exercises 21.4 (around page 400) on "Bubble sorting," and see how the bug is discovered (the "sortedness assertion" fails). Run this code under SPIN, listing the commands and flags you used. Insert your run-results in a framebox.

Fully explain all uses of nondeterminism in this Promela model. Fully explain also the use of data-abstraction (i.e., we got away with using "0" and "1" in sorting; is that representative-enough?) Write a good para arguing that (modulo modeling-errors which are assumed not to be there) this model-checking ended-up verifying the sorting for the size you considered. Can you extend your argument to say that this verification is good for arrays of any size? Write out the bullets of answers.

(Make my life easy; write clear concise bulleted arguments of 1-2 pages.)

4. (40 pts) Figure 21.2 gives you one version of how “system automata” and “property automata” are used—this matches my video-recording of Jan 11th. Run this example under SPIN, and produce a violation trace showing that the never automaton “accepts.” This reveals a liveness violation.

Diagramming, and Running SPIN on a terminal

You must attempt to draw message-sequence charts on a notepad or ipad to debug effectively. Please submit these with the assignment. The URL <http://spinroot.com/spin/Man/Spin.html> gives you info on SPIN's flags. There is more documentation on the **spinroot** page.

```
spin -a yourfile.pml
gcc -DMEMLIM=1024 -O2 -DXUSAFE -DSAFETY -DNOCLAIM -w -o pan pan.c
./pan -m10000
Pid: 66796
```

(Once the pan binary is generated),

GET pan HELP AS FOLLOWS:

```
[ganesh@thinmac Examples]$ ./pan --help
saw option --
Spin Version 6.4.5 -- 1 January 2016
Valid Options are:
-a,-l,-f -> are disabled by -DSAFETY
-A ignore assert() violations
-b consider it an error to exceed the depth-limit
-cN stop at Nth error (defaults to -c1)
-D print state tables in dot-format and stop
-d print state tables and stop
-e create trails for all errors
-E ignore invalid end states
-hN use different hash-seed N:0..499 (defaults to -h0)
-hash generate a random hash-polynomial for -h0 (see also -rhash)
    using a seed set with -RSn (default 12345)
-i search for shortest path to error
-I like -i, but approximate and faster
-J reverse eval order of nested unlessees
-mN max depth N steps (default=10k)
-n no listing of unreachable states
-QN set time-limit on execution of N minutes
-q require empty chans in valid end states
-r read and execute trail - can add -v,-n,-PN,-g,-C
-r filename read and execute trail in file
-rN read and execute N-th error trail
-C read and execute trail - columnated output (can add -v,-n)
-r -PN read and execute trail - restrict trail output to proc N
-g read and execute trail + msc gui support
-S silent replay: only user defined printf's show
-RSn use randomization seed n
-rhash use random hash-polynomial and randomly choose -p_rotateN, -p_permute, or p_reverse
-T create trail files in read-only mode
-t_reverse reverse order in which transitions are explored
-tsuf replace .trail with .suf on trailfiles
-V print SPIN version number
-v verbose -- filenames in unreachable state listing
-wN hashtable of 2^N entries (defaults to -w24)
-x do not overwrite an existing trail file
```

options -r, -C, -PN, -g, and -S can optionally be followed by
a filename argument, as in '-r filename', naming the trailfile
[ganesh@thinmac Examples]\$

== One tries to catch bugs at the most shallow depth ==
== SPIN also has a BFS mode and also a depth minimization mode ==

```
[ganesh@thinmac Examples]$ ./pan -m10000 <== search depth
```

```

=== LOOK AT HOW ERRORS ARE SUMMARIZED AND REPORTED ===
=== You MUST read these bugs and warnings carefully! ===
=== Also read the statistics carefully! ===

```

```

pan:1: invalid end state (at depth 20)
pan: wrote atomicphil.pml.trail

```

```

(Spin Version 6.4.5 -- 1 January 2016)
Warning: Search not completed
+ Partial Order Reduction

```

```

Full statespace search for:
never claim          - (not selected)
assertion violations +
cycle checks         - (disabled by -DSAFETY)
invalid end states +

```

```

State-vector 108 byte, depth reached 23, errors: 1      <==
    12 states, stored
    2 states, matched
    14 transitions (= stored+matched)
    5 atomic steps
hash conflicts:      0 (resolved)

```

```

Stats on memory usage (in Megabytes):
    0.002 equivalent memory usage for states (stored*(State-vector + overhead))
    0.291 actual memory usage for states
    128.000 memory used for hash table (-w24)
    0.534 memory used for DFS stack (-m10000)
    128.730 total actual memory usage

```

```

== THIS IS ERROR-TRAIL SIMULATON BELOW - understand all these flags! ==

```

```

[ganesh@thinmac Examples]$ spin -p -r -s -c -t atomicphil.pml <-- just an example
proc 0 = :init:
using statement merging
Starting phil with pid 1 <-- just an example showing what to expect, is below
proc 1 = phil
    1: proc 0 (:init::1) atomicphil.pml:28 (state 1) [(run phil(p0,v0,p2,v2))]
Starting phil with pid 2
proc 2 = phil
    2: proc 0 (:init::1) atomicphil.pml:32 (state 2) [(run phil(p1,v1,p0,v0))]
Starting phil with pid 3
proc 3 = phil
    3: proc 0 (:init::1) atomicphil.pml:36 (state 3) [(run phil(p2,v2,p1,v1))]
Starting fork with pid 4
proc 4 = fork
    4: proc 0 (:init::1) atomicphil.pml:39 (state 4) [(run fork(p0,v0))]
Starting fork with pid 5
proc 5 = fork
    5: proc 0 (:init::1) atomicphil.pml:41 (state 5) [(run fork(p1,v1))]
Starting fork with pid 6
proc 6 = fork
    6: proc 0 (:init::1) atomicphil.pml:43 (state 6) [(run fork(p2,v2))]
q\p  0  1  2  3  4  5  6
    5  .  .  .  lfp!0
    7: proc 3 (phil:1) atomicphil.pml:4 (state 1) [lfp!0]
    5  .  .  .  .  .  p?0
    8: proc 6 (fork:1) atomicphil.pml:10 (state 1) [p?0]
    3  .  .  .  rfp!0
    9: proc 3 (phil:1) atomicphil.pml:4 (state 2) [rfp!0]
    3  .  .  .  .  .  p?0
    10: proc 5 (fork:1) atomicphil.pml:10 (state 1) [p?0]
        Eating 11: proc 3 (phil:1) atomicphil.pml:4 (state 3) [printf('Eating')]
    6  .  .  .  lfv!0

```

```

12: proc 3 (phil:1) atomicphil.pml:4 (state 4) [lfv!0]
    6 . . . . . v?0
13: proc 6 (fork:1) atomicphil.pml:11 (state 2) [v?0]
    1 . lfp!0
14: proc 1 (phil:1) atomicphil.pml:4 (state 1) [lfp!0]
    1 . . . . . p?0
15: proc 4 (fork:1) atomicphil.pml:10 (state 1) [p?0]
    4 . . . rfv!0
16: proc 3 (phil:1) atomicphil.pml:4 (state 5) [rfv!0]
    4 . . . . . v?0
17: proc 5 (fork:1) atomicphil.pml:11 (state 2) [v?0]
    5 . . . lfp!0
18: proc 3 (phil:1) atomicphil.pml:4 (state 1) [lfp!0]
    5 . . . . . p?0
19: proc 6 (fork:1) atomicphil.pml:10 (state 1) [p?0]
    3 . . lfp!0
20: proc 2 (phil:1) atomicphil.pml:4 (state 1) [lfp!0]
    3 . . . . . p?0
21: proc 5 (fork:1) atomicphil.pml:10 (state 1) [p?0]
spin: trail ends after 21 steps
-----
final state:
-----
#processes: 7
21: proc 6 (fork:1) atomicphil.pml:11 (state 2)
21: proc 5 (fork:1) atomicphil.pml:11 (state 2)
21: proc 4 (fork:1) atomicphil.pml:11 (state 2)
21: proc 3 (phil:1) atomicphil.pml:4 (state 2)
21: proc 2 (phil:1) atomicphil.pml:4 (state 2)
21: proc 1 (phil:1) atomicphil.pml:4 (state 2)
21: proc 0 (:init::1) atomicphil.pml:46 (state 8) <valid end state>
7 processes created

=== NOW FIND OUT WHY THE ABOVE IS DESCRIBING A DEADLOCK ===
=== Drawing a diagram will help ===

```