



# Solving Most Systems of Random Quadratic Equations

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## Solving Quadratic Equations?

### Problem statement

Recover an unknown  $n$ -dimensional vector  $x$  from  $m$  phaseless quadratic equations

$$|\langle a_i, x \rangle| = \psi_i = \sqrt{y_i}, \quad 1 \leq i \leq m.$$

### Sufficient conditions for uniqueness

- Real case:  $m \geq 2n - 1$  ( $2n - 1$  necessary)
- Complex case:  $m \geq 4n - 4$

### Applications

- Phase retrieval in Xray crystallography, imaging, microscopy, ptychography, astronomy...
- Mixed linear regressions

## Prior Art

### Convex paradigm via matrix-lifting $X := xx^*$

$$y_i = a_i^* x x^* a_i = \text{Tr}(a_i a_i^* X), \quad 1 \leq i \leq m$$

[Candes'13, Waldspurger'15]

### Nonconvex paradigm (least-squares criterion)

$$\underset{z \in \mathbb{R}^n}{\text{minimize}} \quad f(z) = \frac{1}{2m} \sum_{i=1}^m (y_i - |a_i^* z|^2)^2$$

Algs: AltMinPhase [Netrapalli'15], truncated/projected/reshaped Wirtinger Flow (T/RWF) [Candes'15, Chen'15, Zhang'16, Soltanolkotabi'17], truncated amplitude flow [wang'16], composite optimization [Duchi'17]

### Two stages [Candes'15]

#### Spectral initialization

$$z_0 \leftarrow \text{leading eigenvector of } \frac{1}{m} \sum_{k=1}^m y_i a_i a_i^*$$

#### Iterative refinement

$$z_{t+1} = z_t - \mu_t \nabla f(z_t), \quad t = 0, 1, \dots$$

### Performance guarantees with $a_i \sim \mathcal{N}(0, I_n)$

Table 1: Comparisons of Different Algorithms

Algorithm	Sample complexity $m$	Computational complexity
PhaseLift	$\mathcal{O}(n)$	$\mathcal{O}(n^3/\epsilon^2)$
PhaseCut	$\mathcal{O}(n)$	$\mathcal{O}(n^3/\sqrt{\epsilon})$
AltMinPhase	$\mathcal{O}(n \log n (\log^2 n + \log(1/\epsilon)))$	$\mathcal{O}(n^2 \log n (\log^2 n + \log^2(1/\epsilon)))$
WF	$\mathcal{O}(n \log n)$	$\mathcal{O}(n^3 \log n \log(1/\epsilon))$
TWF	$\mathcal{O}(n)$	$\mathcal{O}(n^2 \log(1/\epsilon))$

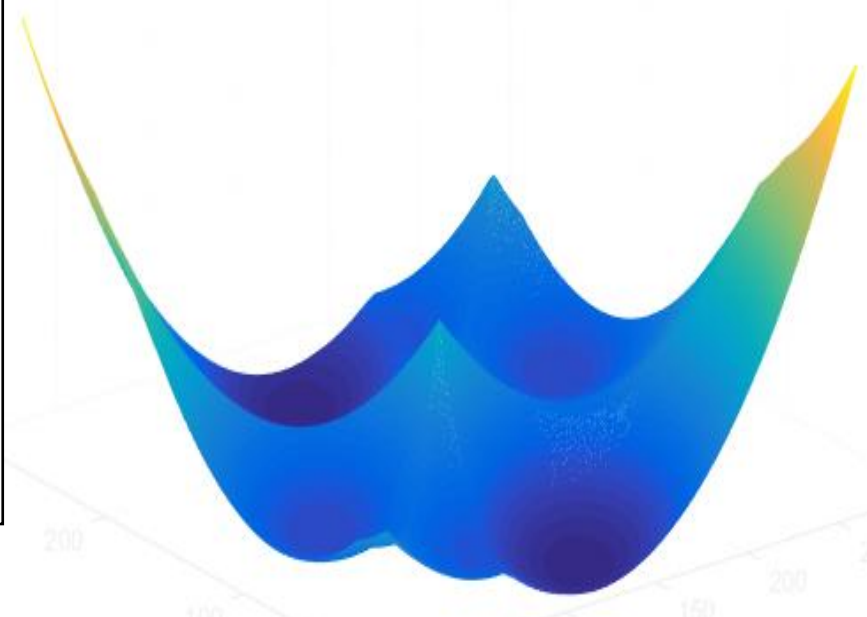
$\epsilon > 0$  given solution accuracy

## Problem Formulation and Algorithm

### The amplitude-based LS cost

$$\underset{z \in \mathbb{R}^n}{\text{minimize}} \quad \ell(z) = \frac{1}{2m} \sum_{i=1}^m (\psi_i - |a_i^* z|)^2$$

Nonsmooth, nonconvex;  
exponentially many  
stationary points;  
even NP-hard to find a  
local minimum



### Reweighted amplitude flow (RAF)

#### Reweighted spectral initialization

$$z_0 \leftarrow \text{1st eigenvector of } \bar{Y}_0 = \frac{1}{|\mathcal{I}_0|} \sum_{i \in \mathcal{I}_0} w_i a_i a_i^*$$

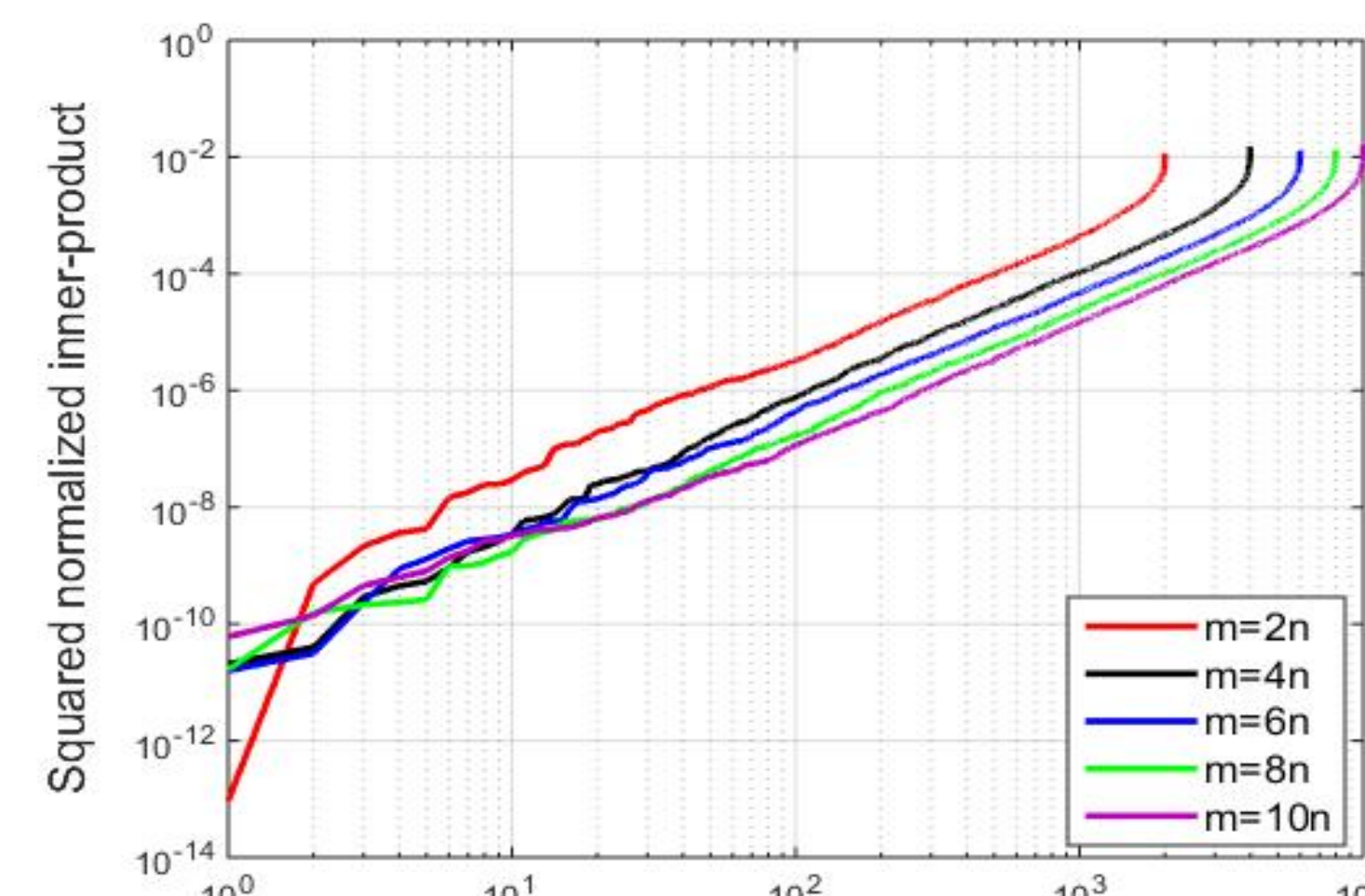
#### Reweighted gradient iterations

$$z_{t+1} = z_t - \mu \partial \ell_{\text{rw}}(z_t), \quad t = 0, 1, \dots$$

Reweighted gradients

## Reweighted Initialization

### An interesting experiment



$$x \sim \mathcal{N}(0, I)$$

$$a_i \stackrel{i.i.d.}{\sim} \mathcal{N}(0, I)$$

$$\frac{|\langle a_i, x \rangle|^2}{\|a_i\|^2 \|x\|^2}$$

Ordered squared normalized inner-products

High-dimensional random vectors are almost always nearly orthogonal to each other! [Cai'13]

**Key idea:** approximate  $x$  by a vector **maximally correlated** to a subset of vectors  $\{a_i\}_{i \in \mathcal{I}_0}$

**Find:** Evaluate and order the data

$$y_{[1]} \geq y_{[2]} \geq \dots \geq y_{[|\mathcal{I}_0|]} \geq y_{[|\mathcal{I}_0|+1]} \geq \dots \geq y_{[m]}$$

**Solve:** the largest eigenvalue problem

$$z_0 := \arg \max_{\|z\|=1} z^* \left( \frac{1}{|\mathcal{I}_0|} \sum_{i \in \mathcal{I}_0} \psi_i^2 a_i a_i^* \right) z$$

Solvable by power iterations in linear time

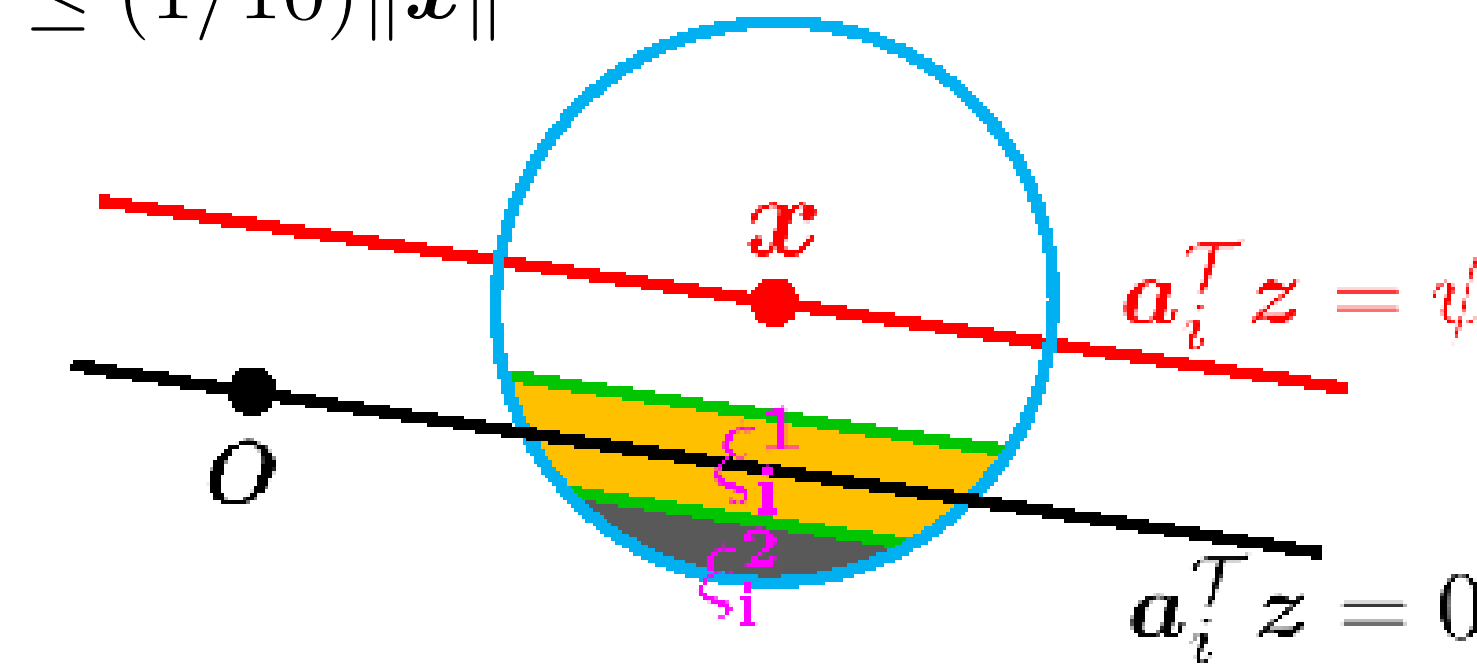
Scaled by norm estimate to have

$$\tilde{z}_0 = \sqrt{\frac{1}{m} \sum_{i=1}^m \psi_i^2} z_0$$

## Reweighted Gradients

### Iteratively reweighting?

$$\text{dist}(z_0, x) \leq (1/10) \|x\|$$



The larger  $\frac{|a_i^* z|}{|a_i^* x|}$ , the more likely  $\partial \ell_i(z)$  points to  $x$

Trick: Weight more the reliable gradients

### Iteratively reweighted gradient refinements

$$z_{t+1} = z_t - \frac{\mu}{m} \sum_{i=1}^m w_i^t (a_i^* z_t - \psi_i \frac{a_i^* z_t}{|a_i^* z_t|}) a_i$$

$$w_i := \frac{1}{1 + c_i / (|a_i^* z_t| / |a_i^* x|)}$$

Subsumes existing TAF/RWF as special cases  
Help eliminating non-smoothness

## Theoretical Guarantees

### Theorem If

- (a1) measurements are noiseless, i.i.d., Gaussian
- (a2) sample complexity satisfies  $m \geq c_1 |\mathcal{I}_0| \geq c_2 n$
- (a3) step size is chosen as  $\mu \lesssim \mu_0$

then w.h.p. for some  $0 < \nu < 1$ , RAF starts within basin of attraction and converges linearly to the global minimum

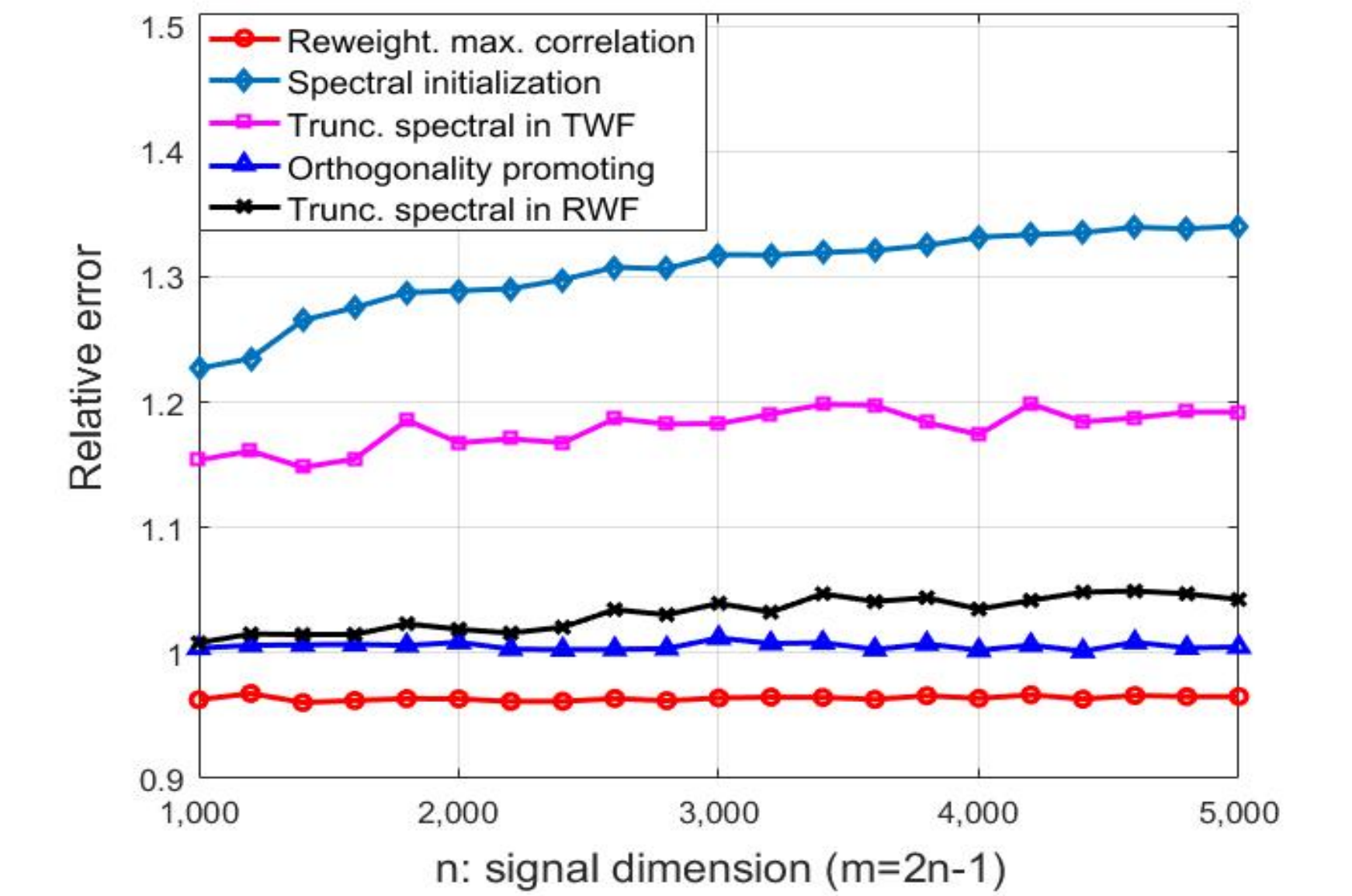
$$\text{dist}(z_t, x) = \min \|z_t \pm x\| \lesssim \frac{1}{10} (1 - \nu)^t \|x\|.$$

- Converges to global min of a nonconvex function
- Computational cost  $\mathcal{O}(mn \log(1/\epsilon))$
- Faster than state-of-the-art WF, TWF [Candes'15, Chen'15]

## Discussions

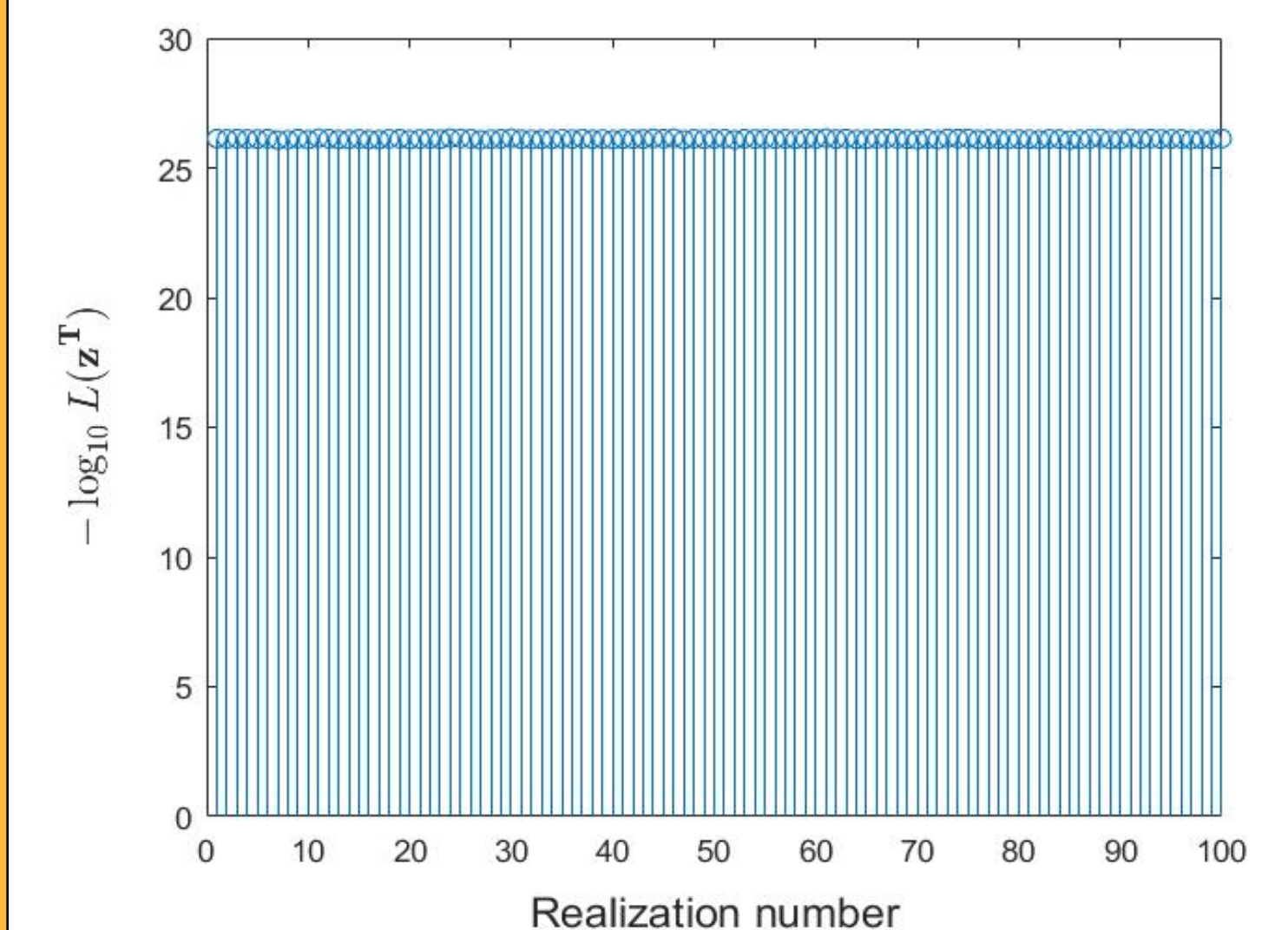
- Initializations for more applications?
- More practical Fourier phase retrieval
- Gradient regularization for other nonconvex opt. such as matrix/tensor completion, robust PCA, blind deconvolution, deep learning...

## Empirical Results



Noiseless Gaussian data

### Numerical surprise (objective in log scale)



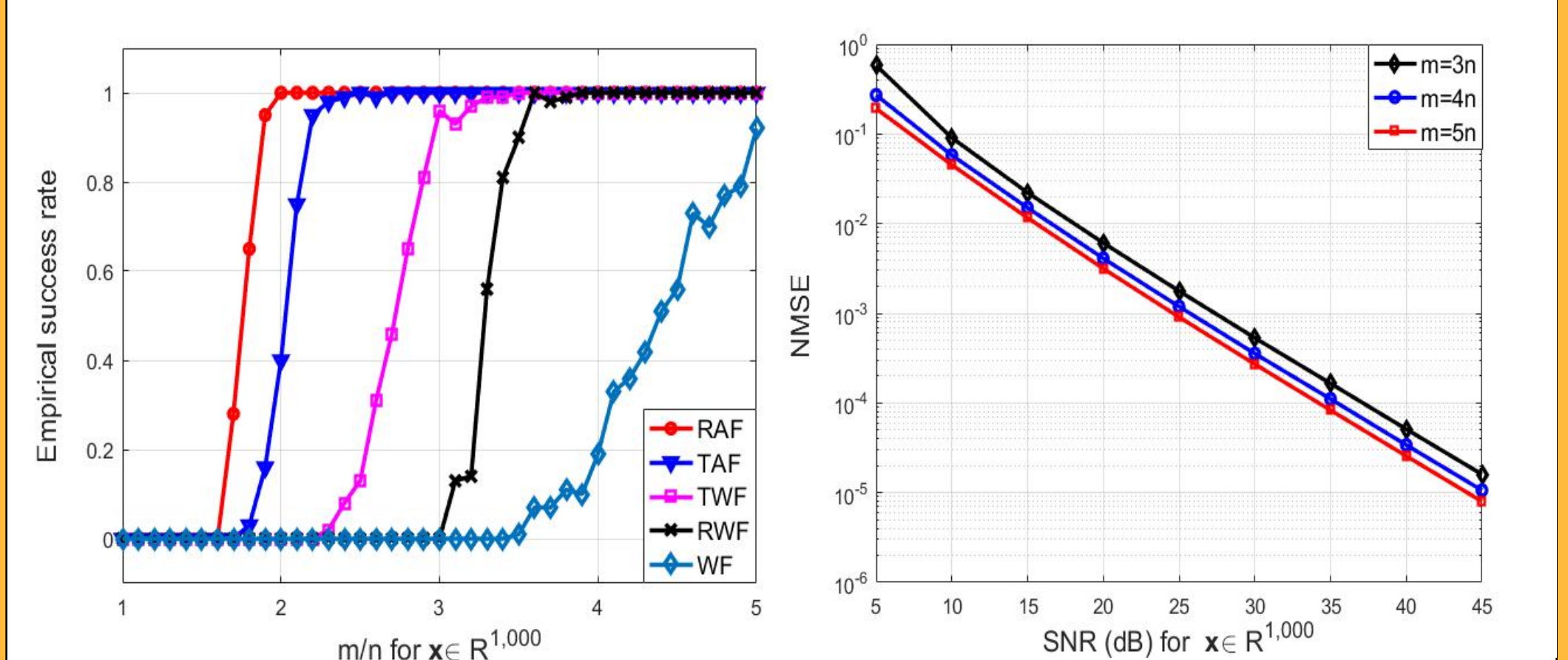
$$n = 1,000$$

$$a_i \sim \mathcal{N}(0, I)$$

$$T = 1,000$$

Under the information-limit condition  $m = 2n - 1$

### Successful recovery and mean-square error



### Real-data validation (CDP w/ 4 masks, $n \approx 10^6$ )

$$\psi^{(k)} = |FD^{(k)} x|, \quad 1 \leq k \leq K = 3$$

Top: recovered image by TAF; bottom: by RAF

