Processes and threads. Signals.

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Process

- Process is running program.
- Each process is identified by unique process identifier (PID) used by kernel.
- Each process knows its parent under the parent process identifier (PPID).
- New process is created by system call:
- fork()
 - function creates a new process
 - the address space of the new process (child process) is an exact copy of the address space of the calling process (parent process)
 - function returns
 - -1 (error) or child PID in parent process
 - 0 in child process
 - child process has some new properties (PID, PPID,...) and other properties are inherited (e.g. EUID, EGID,...) from parent process
- exec()
 - function replaces the current process image with a new process



Example

```
int main (void)
{ . . .
 pid = fork();
  switch (pid) {
  case -1: /* error */
   perror ("error in fork()");
    exit (1);
  case 0: /* child */
   printf ("PID of child: %d\n", (int) getpid ());
    execlp("sleep", "sleep", "30", (char *) NULL);
   perror ("error in execlp()");
    exit (1);
  default: /* parent */
   printf ("PID of parent : %d\n", (int) getpid ());
   wait(&status);
  };
  ...}
```



Thread

- Thread is running subprogram in process/kernel environment.
- Process containing n threads can be executed concurrently on n CPU's.
- Process with one thread (default) can use only one CPU at a given time.
- Thread creation is faster then process creation.
- New thread can be created by library function, e.g. pthread_create().

Example

```
void *thread_code(void *threadid)
{ printf("ID of thread: %ld\n", (long int) threadid);
  sleep(60);
  pthread_exit(NULL);
int main(void)
{ pthread t threads[NUM THREADS];
  int rc, i;
  for(i=0; i<NUM THREADS; i++) {</pre>
    rc = pthread_create(&threads[i], NULL, thread_code, (void *) i);
    if (rc) { perror("error in pthread_create()"); exit(1); }
  /* pthread exit(NULL); */
  /* pthread join(threads[i],&data); */
. . . }
```

Context switching

- Computing process of storing and restoring state (context) of a CPU so that execution can be resumed from the same point at a later time
- This enables multiple threads/processes to share available CPUs.
- Kernel determines when and who gets the CPU.
- Kernel uses several information to make decision
 - Process/thread priority (e.g. 0 169 in Solaris)
 - Priority class (e.g. TS, IA, FSS, FX, SYS, RT in Solaris)
 - Tread state, thread behavior in history, ...
- Thread/process priority can be fixed or dynamic.
- Thread can use CPU only during some time quantum.
- Size of time quantum can be different for different threads/processes and can vary in time (depends on Unix implementation).





CPU 0 - usage

Memory

Process B

Thread 1

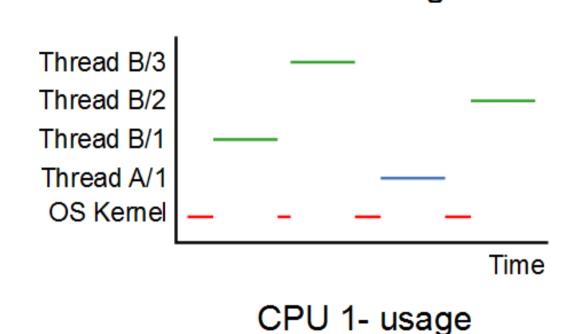
Thread 2

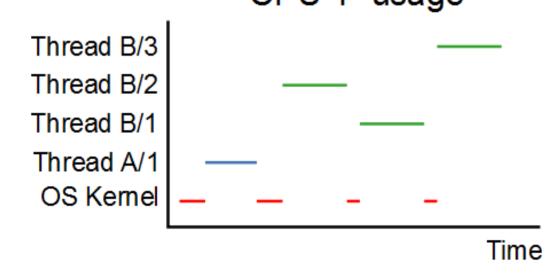
Thread 3

Process A

Thread 1

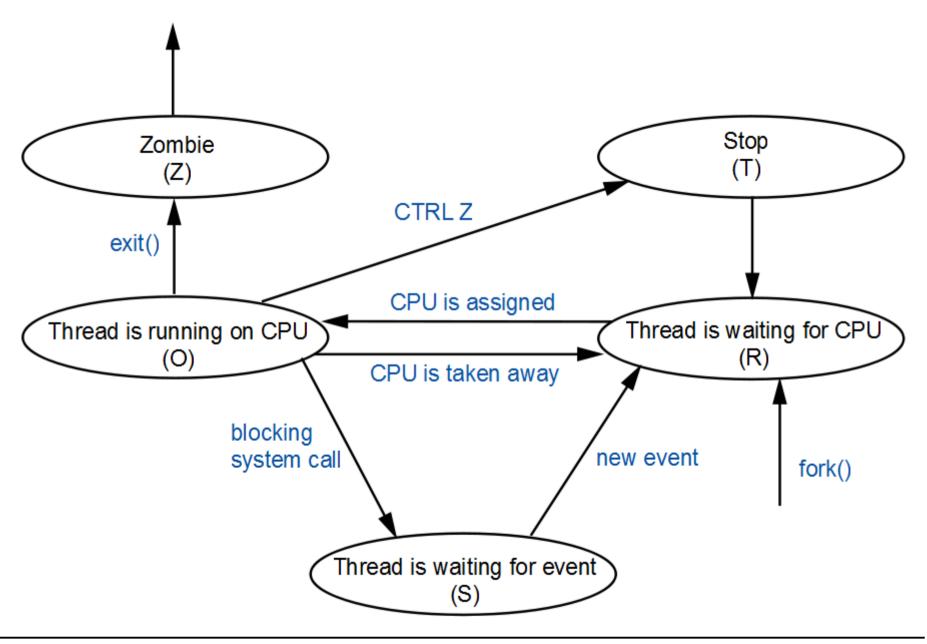
OS Kernel







Thread states







Process/thread listing

ps

 prints information about processes that have the same effective user ID and the same controlling terminal as the invoker

Lists information about every running process

$$ps-f$$
 or $ps-1$

- Prints more details about running processes :
 - S state(O, S, R, Z, T)
 - PID, PPID process ID, parent process ID
 - PRI priority
 - NI NICE value
 - STIME starting time
 - TIME cumulative execution time
 - TTY controlling terminal
 - CMD command name





Process/thread listing

ps -o format

- allows the output format to be specified under user control
- format specification must be a list of names:
 - user ruser group rgroup uid ruid gid rgid pid ppid pgid sid pri nice class time etime stime s c lwp ...

ps -Le

Prints information about each thread (light weight process)





pgrep [-lvx] [pattern] [-u users...]

- reports the process IDs of the processes whose attributes match the criteria specified on the command line
- -1 Prints the process name along with the process ID of each matching process.

-v Reverses the sense of the matching.





prstat or top

 iteratively examines all active processes on the system and reports statistics based on the selected output mode and sort order

```
ptree [-a] [ pid ] [ user ] in Solaris
pstree [-a] [ pid ] [ user ] in Linux
```

 prints the process trees containing the specified pids or users, with child processes indented from their respective parent processes





 Process priority can be decreased (root can also increase) by command:

```
nice -priority program
nice -n priority program
```

- where priority is integer number 1-19
- higher number = lower priority
- negative number = increasing of priority (only root)

 In Solaris: better command for priority modification is priocnt1





How to write a correct script?

What is the difference in the following loops?

 Which loop will determine earlier that the file is removed if the loop is executed concurrently 100x (every time with different file name \$1)?

Signals

- Limited form of inter-process communication used in Unix
- Kernel interrupts the process's normal flow of execution, when signal is sent to a process.
- Signal can be send by
 - kernel (e.g. arithmetic exception, segmentation fault...)
 - terminal driver (e.g. key sequence CTRL C, CTRL \ , ...)
 - other process (e.g. command kill or function kill())
- Signal is identified by name and number.
- Reaction to signal: none, exit, exit+core
- List of signals:
 - Command kill -1
 - In Unix manual man -s 3HEAD signal (Solaris)
- How to send signal:

```
kill -signal PID
pkill -signal [ -vx] [ pattern -u users ...]
```





Signals and terminal driver

- Some signal can be sent by sequence of keys:
- (see stty -a or man stty)

Key sequences	Meaning	Signal
CTRL C	Interrupt process	2 SIGINT
CRTL \	Quit process	3 SIGQUIT
CTRL Z	Suspend process (not in sh).	24 SIGTSTP
	Not process termination!!!	





Important signals

15 SIGTERM (TERMinate)

9 SIGKILL (KILL)

2 SIGINT (INTerrupt) and 3 SIGQUIT (QUIT)

1 SIGHUP (HangUP)

- Parent process sends this signal to its child processes during termination
- If the child process must continue after parent process termination you must start children process by command

nohup cmd &



Reaction to signal

- Default signal reaction are set during process startup.
- Process can modify signal reaction (except signals KILL and STOP).

- We can also modify signal reaction by command trap.
 - Signal reaction setup

```
trap 'commands' signals
```

Print definition of signal reaction:

```
trap
```

Setup ignoring:

```
trap ' signals
```

Setup default signal reaction:

```
trap - signals
```



Command execution

In foreground

\$ command

- Command (no built-in command) is executed like new process.
- Standard input and outputs are assigned to terminal...
- Shell waits for command termination.

In background

\$ command &

- Command (no built-in command) is executed like new process.
- Only standard outputs are assigned to terminal (if the command try to read from stdin the it is stopped)
- Shell doesn't wait for command termination.



Command execution

Immune to hang-ups

- \$ nohup command &
- Command is running after shell termination.
- Standard outputs are redirected to file nohup.out.

In given time

- Command can be started at a given time by commands at or crontab.
- System process cron executes the command under user identity at a given time.
- Standard outputs are sent by email.



Job control

Properties of shell (except /bin/sh)

 Every process which is running in given shell is assigned job identifier (JID) in this shell.

In the shell we can use the following commands

```
jobs list all running process inside this shell
```

fg [%JID] move process to foreground

bg [%JID] move process to foreground

kill -signál [%JID] send signal to process

• If JID is not specified, the last process is used.





Elapsed time report

time command

System calls used by command

truss command Solaris

strace příkaz Linux

