GCC Internals – Generic, Gimple, and RTL

J. Trávníček

Faculty of Information Technology Czech Technical University in Prague

GEN-4

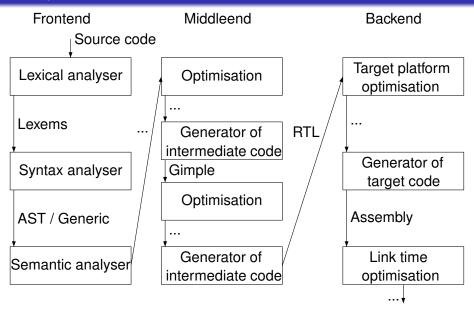
Outline

- Introduction
 - Compiler Design
- Generic
 - Generic
- Gimple
 - Gimple
- **RTL**
 - RTL
- **5** Live example
 - Live example

Outline

- Introduction
 - Compiler Design
- 2 Generic
 - Generic
- GimpleGimple
- 4 RTL
 - RTL
- 5 Live example
 - Live example

Compiler outline



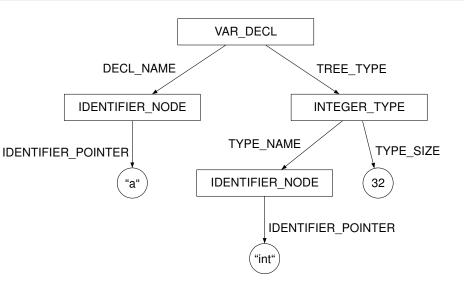
- Generic,
- Gimple,
- Register Transfer Language (RTL).

Outline

- Introduction
 - Compiler Design
 - 2 Generic
 - Generic
- GimpleGimple
- 4 RTL
 - RTL
- 5 Live example
 - Live example

- Tree representation.
- Internally union.
- Constructed by functions make node, build decl, build int cst, ..., build1, build2, ...
- Language independent.
- Represent entire functions, global definitions, ...
- Mostly close to C/C++ programming language.

Generic representation of variable declaration



Nodes of generic – binary

- EQ_EXPR, NE_EXPR
- LT_EXPR, LE_EXPR
- GT_EXPR, GE_EXPR
- TRUTH AND EXPR
- TRUTH_ANDIF_EXPR
- TRUTH_OR_EXPR
- TRUTH_ORIF_EXPR
- TRUTH_XOR_EXPR
- PLUS_EXPR, MINUS_EXPR
- MULT_EXPR, TRUNC_DIV_EXPR

build2(EQ_EXPR, integer_type_node, op1, op2)

Nodes of generic – unary

ABS_EXPR

Introduction

- NEGATE_EXPR
- TRUTH_NOT_EXPR

build1 (NEGATE_EXPR, integer_type_node, op1);

Nodes of generic – {pre, post}{inc, dec}

PREDECREMENT EXPR

Introduction

- POSTDECREMENT EXPR
- PREINCREMENT_EXPR
- POSTINCREMENT_EXPR

```
build2(PREDECREMENT_EXPR, integer_type_node, op1, by);
```

The second argument is how much to increment or decrement by. For a pointer, it would be the size of the object pointed to.

Nodes of generic – pointers

Nodes of generic – variables

```
tree declaration = build decl (UNKNOWN LOCATION, VAR DECL,
  get identifier(name), type);
TREE ADDRESSABLE(declaration) = true;
TREE USED(declaration) = true;
DECL INITIAL(declaration) = ... (expr);
tree declaration = build decl (UNKNOWN LOCATION, VAR DECL,
  get identifier(name), type);
TREE ADDRESSABLE(declaration) = true;
TREE USED(declaration) = true:
TREE STATIC(declaration) = true;
TREE PUBLIC(declaration) = true;
DECL INITIAL(declaration) = ... (expr);
```

Nodes of generic – function type

```
tree params = NULL_TREE;
params = chainon( params,
    tree_cons (NULL_TREE, integer_type_node, NULL_TREE));
tree resdecl = build_decl (BUILTINS_LOCATION, RESULT_DECL,
    NULL_TREE, integer_type_node);
tree fntype = build_function_type(TREE_TYPE(resdecl), params);
```

```
tree param decl = NULL TREE;
tree number = build decl (UNKNOWN LOCATION, PARM DECL,
  get identifier("number"), integer type node);
DECL ARG TYPE(number) = integer type node;
param decl = chainon( param decl. number );
tree fndecl = build decl( UNKNOWN LOCATION, FUNCTION DECL,
  get identifier("factorial"), ... (fntype) );
DECL ARGUMENTS(fndecl) = param decl;
DECL RESULT( fndecl ) = resdecl;
TREE STATIC(fndecl) = true;
TREE PUBLIC( fndecl ) = true;
DECL INITIAL(fndecl) = ... (block);
DECL SAVED TREE(fndecl) = ... (bind);
```

Nodes of generic – blocks

Generic

```
tree decls = NULL TREE;
tree variable i = build decl (...);
decls = chainon(decls, variable i);
tree block = build block(decls, NULL TREE,
  NULL TREE, NULL TREE);
TREE USED(block) = true;
tree stmts = alloc stmt list ();
append to statement list(... (expr), &stmts);
tree bind = build3 (BIND EXPR, void type node,
  BLOCK VARS(block), stmts, block);
TREE SIDE EFFECTS(bind) = true;
BLOCK SUPERCONTEXT(block) = fndecl;
```

```
tree decls = NULL_TREE;
tree variable_k = build_decl (...);
decls = chainon(decls, variable_k);

tree stmts = alloc_stmt_list ();
append_to_statement_list(... (expr), &stmts2);

tree bind = build3( BIND_EXPR, void_type_node,
    decls, stmts, NULL_TREE );
TREE SIDE EFFECTS(bind) = true;
```

- Intuitively no need for bind expression no variables to bind to statement list.
- moreover, not even block is needed.
- Conclusion, sequence of statements can be represented with statement list only.

Nodes of generic – modify and return expression

```
modify
build2(MODIFY_EXPR, TREE_TYPE(op1Generic),
   op1Generic, op2Generic);
return
build1(RETURN_EXPR, void_type_node, op1Generic);
```

Nodes of generic – call expression

```
tree * args_vec = XNEWVEC( tree, argumentsSize );
int i = 0;
for(argument : arguments) {
   args_vec[i++] = ... (expr);
}

tree fndecl = ...;

tree call = build_call_expr_loc_array( UNKNOWN_LOCATION,
   fndecl, argumentsSize, args_vec );
SET_EXPR_LOCATION(call, UNKNOWN_LOCATION);
TREE_USED(call) = true;
```

Nodes of generic – loops, if

```
loop
build1(LOOP_EXPR, void_type_node, ... (bind));
exit
build1(EXIT_EXPR, void_type_node, ... (expr));
if
build3(COND_EXPR, void_type_node, ... (expr), ... (bind), ... (bind));
```

- Nodes are defined in tree.def file.
- frontend can introduce new nodes,
- generic must be translated to gimple for middle-end,
- frontend must provide function to translate new nodes to gimple,
- generic can be dumped during compilation to file with fdump-tree-original-raw,
- c-like representation of generic can be dumped with fdump-tree-original.

- - Compiler Design
- Gimple Gimple
- - RTL
- - Live example

- Derived from Generic,
- three-address representation,
- temporaries hold intermediate values,
- labels and conditioned jumps represent all control structures

Temporaries

 Generic complex expressions are split to more expressions:

$$a = b + c + d$$

become

$$T1 = b + c$$

$$a = T1 + d$$

Control statements – if

 Control statements are converted to labels and (conditioned) gotos:

```
if ( ... ) goto <D.103>; else goto <D.104>;
<D.103>:
{
}
goto <D.105>;
<D.104>:
{
}
<D.105>:
```

Control statements – loops

 Control statements are converted to labels and (conditioned) gotos:

```
<D.99>:
{
    ...
    if ( ... ) goto <D.107>; else goto <D.109>;
    <D.109>:
}
goto <D.99>;
<D.107>:
```

Internal form dump

- gimple can be dumped during compilation to file with fdump-tree-gimple.
- all stages of gimple processing can be dumped during compilation to file with fdump-tree-all.

Outline

- Introduction
 - Compiler Design
- 2 Generic
 - Generic
- Gimple
 - Gimple
- 4 RTL
 - RTL
- 6 Live example
 - Live example

- Low level intermediate representation,
- algebraic representation of instructions,
- representation inspired by lisp.

- Constant: (const int i),
- Register: (reg:m n); n number of machine or pseudo register, m – mode (full word, half word, single byte, ...),
- Memory: (mem:m addr alias); m memory size, addr address, alias - reference alias,
- Addition: (plus:m x y); x, y arguments, m mode,
- Preincrement: (pre inc:m x), ...,
- Equal: (eq:m x y), ...,
- Call: (call (mem:fm addr) nbytes); nbytes number of bytes of arguments, fm - mode, addr - address of subroutine.

Examples

 all stages of rtl processing can be dumped with fdump-rtl-all.

- - Compiler Design
- - Generic
- Gimple
- - RTL
- 5 Live example
 - Live example

Live example

...