Effective C++ Programming

NIE-EPC (v. 2021):

TEMPLATES, TEMPLATE INSTANTIATION, TEMPLATE ARGUMENT DEDUCTION, TEMPLATE SPECIALIZATION, VARIADIC TEMPLATES, TEMPLATE PARAMETER PACK

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Templates — templated entities

- Templates = parametrized generators of source code entities.
- Possible template-generated entities:
 - classes,
 - functions (both free and member),
 - variables (since C++17),
 - types (type aliases; since C++11),
 - concepts (since C++20).
- Template definition = prescription on how the entity should be generated (with respect to template arguments).
 - Its generation is called "instantiation".
 - Resulting entity is called "instance" of a template.
- Example:
 - wrapper = class template, wrapper<int> = its instance = class.

```
template <typename T>
struct wrapper {
  T obj;
};
```

Templates — parameters and arguments

- Templates = parametrized generators of *source code* entities.
- Implication generation/instantiation must happen at compile time.
 - It effectively generates source code that needs to be further compiled by a compiler.
 - Note at runtime, templates no more exist (only their instances do).
- Generation of source code entities is parametrized by template parameters.
 - Arguments for these parameters template arguments need to be resolvable at compile time.
 - $\bullet \Rightarrow$ Basically, there are two kinds of template parameters:
 - type template parameters their arguments are types,
 - non-type template parameters their arguments are compile-time constants; mostly integral/Boolean).
- Example:

```
template <typename T, size_t N> // T - type template param., N - non-type template param.
struct Array {
   T data[N];
};
```

Templates — explicit instantiation

- Templates can be instantiated either explicitly or implicitly.
- Explicit instantiation forces a compiler to generate template instance without its usage.
 - Example source code...

```
template <typename T> bool nonzero(T t) { return t != 0; } // function template definition
template bool nonzero<int>(int); // explicit instantiation of function nonzero<int>
template bool nonzero<long>(long); // explicit instantiation of function nonzero<long>
```

• ...effectively equivalent source code after instantiation...

```
bool nonzero<int>(int t) { return t != 0; }
bool nonzero<long>(long t) { return t != 0; }
```

• ...resulting *machine code*:

```
bool nonzero<int>(int):
    test edi, edi
    setne al
    ret

bool nonzero<long>(long):
    test rdi, rdi
    setne al
    ret
```

Note — explicit template instantiation is rarely used in practice.

Templates — implicit instantiation

- *Implicit instantiation* forces a compiler to generate template instance by using the template.
 - Example source code...

```
template <typename T> bool nonzero(T t) { return t != 0; } // function template definition
int main() {
   std::cout << (void*)nonzero<int>; // print address of function nonzero<int>
}
```

• ...effectively equivalent source code after instantiation...

```
bool nonzero<int>(int t) { return t != 0; }
int main() {
  std::cout << (void*)nonzero<int>;
};
```

...resulting machine code:

```
bool nonzero<int>(int):
   test edi, edi
   setne al
   ret
main:
   ...
```

Templates — header files

- Implicit instantiation is initiated by using the template.
- For instantiation, template definition needs to be available.
- Implication if the template is not instantiated explicitly, its
 definition needs to be available in each translation unit where
 that template is used.
- Typical solution of this problem = putting definitions of templates into header files.
- Note:
 - In case of function templates, this can effectively result in the presence of a machine code of its instance function in multiple translation units = object files.
 - With templates, this does not break ODR (one definition rule).

Template argument deduction

- In *all cases*, template arguments may be provided explicitly.
- In some cases, they may be implicitly deduced by a compiler.
- Such deduction is mostly based:
 - either on function arguments for function templates,
 - or on initialization = constructor arguments for class templates (since C++17).
- → For deduction, template parameter needs to be somehow involved in the form of a function (constructor) parameter.
- Example:

```
template <typename T> bool nonzero(T t) { return t != 0; }
int main(int argc, char* argv[]) {
   std::cout << nonzero<int>(argc); // explicitly-provided template argument
   std::cout << nonzero(argc); // template argument deduced by function argument
}</pre>
```

...effectively equivalent source code after instantiation...

```
bool nonzero<int>(int t) { return t != 0; }
int main(int argc, char* argv[]) {
   std::cout << nonzero<int>(argc);
   std::cout << nonzero<int>(argc);
}
```

Template argument deduction (cont.)

- Deduction of template arguments must be unambiguous.
- Otherwise, its resolution cannot take place.
- Example:

```
template <typename T>
bool larger(T a, T b) { return a > b; }
int main() {
   std::cout << larger(1, 2.0);
   // error: no matching function to call 'larger(int, double)'
   // note: template argument deduction/substitution failed:
   // note: deduced conflicting types for parameter 'T' ('int' and 'double')
}</pre>
```

Resolution:

- According to the first function argument (literal of type int), the template argument is deduced as int.
- According to the second function argument (literal of type double), the template argument is deduced as double.
- None of these options has a higher priority; they are equivalent for the compiler.
- $\bullet \Rightarrow$ The compiler cannot decide which one to use.

Forwarding reference revisited

• Example — forwarding reference as a function template parameter:

```
template <typename T> void f(T&& param) { }
```

- It can bound both *lvalue* and *rvalue* arguments of *any type...*
- ...and it has special template argument deduction rules different for both these cases:
 - In case of *lvalue argument* of type X, T is deduced as X&.
 - In case of rvalue argument of type X, T is deduced as X.

```
int i = 1;
f(i);  // i  is Lvalue of type int => T is deduced as int& => type of param is int&
f(true);  // true is rvalue of type bool => T is deduced as bool => type of param is bool&&
```

Effectively equivalent source code after instantiation:

```
void f<int&>(int& param) { }
void f<bool>(bool&& param) { }

int i = 1;
f(i);
f(true);
```

- Note once instantiation takes place, forwarding reference becomes either lvalue or rvalue reference.
 - There are no more forwarding references after template instantiation.

Template argument substitution

- Once the template argument is resolved (explicitly provided, deduced,...) it is substituted for the corresponding template parameter.
- Substitution of template arguments effectively results in an instance of that template.
- Example:

```
template <typename T> bool larger(const T& a, const T& b) { return a > b; }
int main() {
  std::cout << larger(1, 2);
}</pre>
```

- Here, template argument is deduced as int according to function arguments (both of type int).
 - \Rightarrow int is substituted for T.
 - ⇒ Effectively equivalent source code after instantiation:

```
bool larger<int>(const int& a, const int& b) { return a > b; }
int main(int argc, char* argv[]) {
   std::cout << larger<int>(1, 2);
}
```

Template specialization

- Template specialization allows having multiple definitions of the same template...
- ...while the one used for instantiation is selected according to the template arguments.
- One definition needs to be generic it is called primary template.
- Other definitions are used for instantiation in special cases only ⇒ they are called "specializations".
- Instantiation:
 - When the template arguments match some specialization, it is used for instantiation.
 - Otherwise, primary template is used instead.
- Typically, only class templates are specialized in practice.
 - With function templates, overloading is preferred.
 - Mixing overloading with function template specialization is "dangerous".

Template specialization — examples

- Example class that maps type to a static Boolean constant which has value:
 - *true* if the type is bool,
 - *false* otherwise.

```
int main() {
  std::cout << is_bool<int>::value; // (1) prints out "0"
  std::cout << is_bool<bool>::value; // (2) prints out "1"
}
```

- Resolution:
 - In (1), int does not match the specialization ⇒ class is_bool<int> is instantiated from the primary template.
 - In (2), bool does match the specialization ⇒ class is_bool<bool> is instantiated from this specialization.
- ► ⇒ Effectively equivalent code after instantiation:

```
struct is_bool<int> { static const bool value = false; };
struct is_bool<bool> { static const bool value = true; };
int main() {
   std::cout << is_bool<int>::value;
   std::cout << is_bool<bool>::value;
}
```

Template specialization — examples (cont.)

• *Another example* — factorial calculation.

std::cout << factorial<3>::value; // prints out 6

• Primary template is based on recurrent factorial calculation:

```
template <int I>
struct factorial { static const int value = I * factorial<I-1>::value; };

• Each recurrence needs some base case = specialization in our case:

template <> struct factorial<0> { static const int value = 1; };  // specialization

int main() {
```

- Resolution:
 - Class factorial<3> does not match the specialization ⇒ it is instantiated from the primary template as:

```
struct factorial<3> { static const int value = 3 * factorial<2>::value; };
```

- This instance (class) contains the use of template factorial with template argument 2 ⇒ it needs to be instantiated as well.
- Class factorial
 does not match the specialization ⇒ it is instantiated (again) from the primary template as:

```
struct factorial<2> { static const int value = 2 * factorial<1>::value; };
```

Template specialization — examples (cont.)

```
template <int I>
struct factorial { static const int value = I * factorial<I-1>::value; };
template <> struct factorial<0> { static const int value = 1; };  // specialization

struct factorial<2> { static const int value = 2 * factorial<1>::value; };
```

- Resolution (cont.):
 - This instance (class) contains the use of template factorial with template argument 1 ⇒ it needs to be instantiated as well.
 - Class factorial<1> does not match the specialization ⇒ it is instantiated (again) from the *primary template* as:

```
struct factorial<1> { static const int value = 1 * factorial<0>::value; };
```

- This instance (class) contains the use of template factorial with template argument o ⇒ it needs to be instantiated as well.
- Class factorial<0> does match the specialization ⇒ it is instantiated from the specialization as:

```
struct factorial<0> { static const int value = 1; };
```

Template specialization — examples (cont.)

Summary:

```
template <int I>
struct factorial { static const int value = I * factorial<I-1>::value; };

template <> struct factorial<0> { static const int value = 1; };  // specialization

int main() {
   std::cout << factorial<3>::value; // prints out 6
}
```

Effectively equivalent code after instantiation...

```
struct factorial<0> { static const int value = 1; };
struct factorial<1> { static const int value = 1 * factorial<0>::value; };
struct factorial<2> { static const int value = 2 * factorial<1>::value; };
struct factorial<3> { static const int value = 3 * factorial<2>::value; };
int main() {
   std::cout << factorial<3>::value;
}
```

....and after resolution of the constants:

```
struct factorial<0> { static const int value = 1; };
struct factorial<1> { static const int value = 1; };
struct factorial<2> { static const int value = 2; };
struct factorial<3> { static const int value = 6; };
int main() {
   std::cout << factorial<3>::value;
}
```

Template metafunctions and type traits

- Recall class templates that maps template arguments (types or compile-time constants) to other types or compile-time constants are generally called template metafunctions.
- Both introduced templates is_bool and factorial fall into this category.
 - is_bool maps type to compile-time Boolean constant,
 - factorial maps compile-time integral constant to other compile-time integral constant.
- Metafunctions that works with types are called "type traits".
- We have already seen some type traits from the C++ standard library, such as std::is_reference or std::remove_reference.
- Note definition of metafunctions is mostly based on template specialization.

Template metafunctions — examples

- Metafunction std::is_reference maps type to a Boolean constant which is:
 - true if type is a reference type,
 - false otherwise.
- Solution:
 - primary template for non-reference types,
 - specialization for reference types.

- Partial vs full specialization:
 - If a specialization is no more parametrized, it is called "full specialization".......
 - If a specialization is parametrized, it is called "partial specialization".

```
template <> struct factorial<0> { static const int value = 1; }; // (full) specialization <==
```

Template metafunctions — examples (cont.)

• Example:

```
template <typename T> struct is_reference { static const bool value = false; };
template <typename T> struct is_reference<T&> { static const bool value = true; };
template <typename T> struct is_reference<T&&> { static const bool value = true; };
int main() {
   std::cout << is_reference< float >::value; // prints out "0"
   std::cout << is_reference< bool& >::value; // prints out "1"
   std::cout << is_reference< int&& >::value; // prints out "1"
}
```

Resolution:

float is non-reference type ⇒ it does not match any specialization ⇒ is_reference<float> is instantiated from primary template as:

```
struct is_reference<float> { static const bool value = false; };
```

boo1& is reference type ⇒ it does match a specialization for lvalue references ⇒ is_reference<boo1&> is instantiated from it as:

 int&& is reference type ⇒ it does match a specialization for rvalue references ⇒ is_reference<int&&> is instantiated from it as:

```
struct is_reference<int&&> { static const bool value = true; }; <-----
```

Template metafunctions — examples (cont.)

Summary:

```
template <typename T> struct is_reference { static const bool value = false; };
template <typename T> struct is_reference<T&> { static const bool value = true; };
template <typename T> struct is_reference<T&&> { static const bool value = true; };
int main() {
   std::cout << is_reference< float >::value; // prints out "0"
   std::cout << is_reference< bool& >::value; // prints out "1"
   std::cout << is_reference< int&& >::value; // prints out "1"
}
```

Effectively equivalent code after instantiation:

```
struct is_reference<float> { static const bool value = false; };
struct is_reference<bool&> { static const bool value = true; };
struct is_reference<int&&> { static const bool value = true; };
int main() {
   std::cout << is_reference< float >::value;
   std::cout << is_reference< bool& >::value;
   std::cout << is_reference< int&& >::value;
}
```

Template metafunctions — examples (cont.)

- Metafunction std::remove_reference maps input type T to an output type which is:
 - type referred to by T if T is a reference type,
 - T otherwise (identity).
- Solution:
 - primary template for non-reference types,
 - specialization for reference types.

```
template <typename T> struct remove_reference { using type = T; }
template <typename T> struct remove_reference<T&> { using type = T; }
template <typename T> struct remove_reference<T&&> { using type = T; }
```

• Example:

```
int main() {
  remove_reference< bool >::type b = true; // type of b is bool
  remove_reference< int& >::type i = 1; // type of i is int
}
```

Effectively equivalent code after instantiation:

```
struct remove_reference<bool> { using type = bool; }
struct remove_reference<int&> { using type = int; }
int main() {
  remove_reference< bool >::type b = true;
  remove_reference< int& >::type i = 1;
}
```

Template metafunctions — "shortcuts"

 Recall — usage of metafunctions/type traits requires relatively long writing:

```
template <typename T>
typename remove_reference<T>::type && move(T&& param) {
  return static_cast< typename remove_reference<T>::type && > param;
}
```

 For type traits which "produce" types, we can create a "shortcut" type (alias) template wrapper as follows:

```
template <typename T>
using remove_reference_t = typename remove_reference<T>::type;
```

Consequently...

```
typename remove_reference<T>::type
```

...may be then replaced by:

```
remove_reference_t<T>
```

For type traits in the C++ standard library, such "shortcut" type templates were introduced in C++14.

```
template <typename T> std::remove_reference_t<T> && move(T&& param)
{  return static_cast< std::remove_reference_t<T> && > param; }
```

Template metafuncts — "shortcuts" (cont.)

- Similarly, such "shortcut" wrappers may be created for type traits which "produce" values.
- However, these require variable templates, which were not available until C++17.

```
template <typename T>
inline bool is_reference_v = is_reference<T>::value;
```

Consequently...

```
is_reference<T>::value
```

...may be then replaced by:

```
is_reference_v<T>
```

 For type traits in the C++ standard library, such "shortcut" variable templates were introduced in C++17.

```
std::is_reference_v<T> // since C++17
```

Template specialization — example

- Template specialization has many uses other than just template metafunctions.
- C++ standard library examples:
 - std::atomic specializations for integers, pointers,...;
 - std::hash specialization for different types;
 - std::vector<bool> possible bit-compressed implementation.
- Another example some parallel computation based on MPI (see NI(E)-PDP):
 - MPI library has a "C" API (no templates).
 - MPI functions that send/receive data takes:
 - a non-type pointer to these data (of type void*),
 - special argument/"type tag" that tells the library of what type these data are.
 - Exemplary MPI type tags: MPI_INT for int, MPI_DOUBLE for double,...
 - Exemplary function for sending data to another MPI process:

```
int MPI_Send(void* buf, int count, MPI_Datatype mdt, int dest, int tag, MPI_Comm comm);
```

Template specialization — example (cont.)

• If we know the type of sent data, everything works fine:

```
const int n = 4;
double a[] = { 1.0, 2.0, 3.0, 4.0 };
MPI_Send(&n, 1, MPI_INT, 0, 0, MPI_COMM_WORLD);
MPI_Send(a, n, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD);
```

- Problem what if type of sent data is a template parameter?
- Example parallel computation parametrized by floating-point data type, which defines its used precision:

- How to specify MPI type tag that specifies type of sent data?
- Required functionality:
 - If T is float, MPI_FLOAT type tag should be used.
 - If T is double, MPI_DOUBLE type tag should be used.
- Generally, we need to map types to corresponding MPI type tags.

Template specialization — example (cont.)

- Possible solution template specialization.
 - Note MPI type tags may be in fact pointers into some type table (link).
 - ⇒ Generally, they may not be compile-time constants.
 - Alternative static member function that returns tag value.

```
template <typename T> struct GetMPIDatatype; // primary template; declaration only
template<> struct GetMPIDatatype<float>
{  static MPI_Datatype get() { return MPI_FLOAT; } }; // specialization for float
template<> struct GetMPIDatatype<double>
{  static MPI_Datatype get() { return MPI_DOUBLE; } }; // specialization for double
... // similarly for other types supported by MPI library (MPI_INT, MPI_LONG,...)
```

Application to our problem:

```
MPI_Send(v.data(), v.size(), GetMPIDatatype<T>::get(), 0, 0, MPI_COMM_WORLD);
```

Example — computation to be instantiated with single precision:

```
MPI_Send(v.data(), v.size(), GetMPIDatatype<float>::get(), 0, 0, MPI_COMM_WORLD);
```

This initiates instantiation of GetMPIDatatype<float> from a corresponding specialization:

```
struct GetMPIDatatype<float>
{ static MPI_Datatype get() { return MPI_FLOAT; } };
```

Templates — complex example

Consider our implementation of Vector class template...

```
template <typename T> class Vector {
  size_t capacity_, size_;
  T* data_;
  ...
```

...and its following exemplary use:

Vector class template is instantiated with X template argument,
 which results in Vector<X> class ⇒ X is substituted for T:

```
class Vector<X> {
  size_t capacity_, size_;
  X* data_;
  ...
```

Let us now look at some Vector member functions.

• First, the reserve member function was defined as:

```
void reserve(size_t capacity) {
  if (capacity <= capacity_) return;
  T* data = (T*)::operator new(capacity * sizeof(T));
  for (size_t i = 0; i < size_; i++) new (data + i) T( move( *(data_ + i) ) );
  clear();
  ::operator delete(data_);
  data_ = data; capacity_ = capacity;
}</pre>
```

 After instantiation/substitution, this effectively results in the following definition:

```
void reserve(size_t capacity) {
  if (capacity <= capacity_) return;
  X* data = (X*)::operator new(capacity * sizeof(X));
  for (size_t i = 0; i < size_; i++) new (data + i) X( move( *(data_ + i) ) );
  clear();
  ::operator delete(data_);
  data_ = data; capacity_ = capacity;
}</pre>
```

- This initiates the instantiation of move function template.
- Note we deliberately used our custom version here instead of std::move to show what happens next.

Our custom move function template definition:

```
template <typename T>
typename remove_reference<T>::type && move(T&& param) {
  return static_cast< typename remove_reference<T>::type && > param;
}
```

- Note similarly as above, we used our remove_reference trait.
- Resolution:
 - Template argument of move call is not specified.
 - \Rightarrow It needs to be deduced according to the function argument.
 - Function argument of move is lvalue expression of type X.
 - param is a forwarding reference.
 - \Rightarrow T is deduced as X&.
 - This effectively results in the following instantiation of move:

```
remove_reference<X&>::type && move<X&>(X& && param) {
  return static_cast< remove_reference<X&>::type && > param;
}
```

 Now, this instantiation initiates instantiation of remove_reference<X&> class.

remove_reference<X&> is instantiated as follows:

```
struct remove_reference<X&> { using type = X; }
```

• \Rightarrow For illustration, we can then effectively rewrite...

```
remove_reference<X&>::type && move<X&>(X& && param) {
  return static_cast< remove_reference<X&>::type && > param;
}
```

• ...as...

```
X&& move<X&>(X& param) { return static_cast< X&& > param; }
```

...and this instance of move is called in Vector<X>::reserve:

```
void reserve(size_t capacity) {
    ...
for (size_t i = 0; i < size_; i++) new (data + i) X( move<X&>( *(data_ + i) ) );
    ...
```

- This is exactly what we expected move<X&> call is an expression that:
 - represents the original vector element *(data_ + i) of type X,
 - and its category is *rvalue*.
- New elements are initialized by move constructor of class X.

Summary for Vector<X>::reserve...

```
template <typename T> class Vector {
    size_t capacity_, size_;
    T* data_;
public:
    void reserve(size_t capacity) {
        if (capacity <= capacity_) return;
        T* data = (T*)::operator new(capacity * sizeof(T));
        for (size_t i = 0; i < size_; i++) new (data + i) T( move( *(data_ + i) ) );
        clear();
        ::operator delete(data_);
        data_ = data; capacity_ = capacity;
}</pre>
```

...is effectively instantiated as follows:

```
struct remove_reference<X&> { using type = X; }

remove_reference<X&>::type && move<X&>(X& param) {
    return static_cast< remove_reference<X&>::type && > param;
}

class Vector<X> {
    size_t capacity_, size_;
    X* data_;

public:
    void reserve(size_t capacity) {
    if (capacity <= capacity_) return;
    X* data = (X*)::operator new(capacity * sizeof(X));
    for (size_t i = 0; i < size_; i++) new (data + i) X( move<X&>( *(data_ + i) ) );
    clear();
    ::operator delete(data_);
    data_ = data; capacity_ = capacity;
}
```

Next, the push_back member functions were defined as:

```
void push_back(const T& param) {
    if (size_ == capacity_) reserve(capacity_ ? 2 * capacity_ : 1);
    new (data_ + size_) T( param );
    size_++;
}

void push_back(T&& param) {
    if (size_ == capacity_) reserve(capacity_ ? 2 * capacity_ : 1);
    new (data_ + size_) T( move( param ) );
    size_++;
}

// overload for lvalues

// overload for rvalues
```

 After instantiation/substitution, this effectively results in the following definitions:

```
void push_back(const X& param) {
  if (size_ == capacity_) reserve(capacity_ ? 2 * capacity_ : 1);
  new (data_ + size_) X( param );
  size_++;
}

void push_back(X&& param) {
  if (size_ == capacity_) reserve(capacity_ ? 2 * capacity_ : 1);
  new (data_ + size_) X( move<X&>( param ) );
  size_++;
}
```

 Note — move call after deduction results in the call of the already instantiated function move < X&>.

Variadic templates

What about emplace_back member function?

```
template <typename... U> void emplace_back(U&&... param) {
  if (size_ == capacity_) reserve(capacity_ ? 2 * capacity_ : 1);
  new (data_ + size_) T( std::forward<U>(param)... );
  size_++;
}
```

- To understand it, we first need to explain variadic templates.
- Recall template parameter starts with
 - typename (or class) keyword for type template parameters,
 - type name for non-type template parameters,
- and is then it is (optionally) followed by identifier (template parameter name).

```
template <typename T, class U, int I, size_t N> struct X { };
```

• When we write ellipsis (sequence of 3 dots) after the typename/class/type name, we create a so-called "template parameter pack" (instead of a template parameter).

```
template <typename T, typename ... Ts> struct X { };
// T is template parameter, Ts is a template parameter pack
```

 Template with at least one template parameter pack is called variadic template.

Variadic templates (cont.)

- Recall when template is instantiated, template parameters are substituted by template arguments.
- Namely, a single template parameter is substituted by a single template argument.
- On the contrary, template parameter pack may be substituted by an arbitrary number of template arguments, including no one.

```
template <typename T, typename ... Ts> struct X { };

X<int, bool, double> x1; // T is substituted by int, Ts is substituted by bool and double
X<int, bool> x2; // T is substituted by int, Ts is substituted by bool
X<int> x3; // T is substituted by int, Ts is empty
```

- Effectively, template parameter pack represents a list of virtual unnamed template parameters.
- It is not possible to refer to individual elements of template parameter pack.
- However, it is possible to so-call "expand" it.

Variadic templates — expansion

- Template parameter pack expansion happens when its identifier is followed by ellipsis in the template definition.
- Effect as if the list of virtual template parameters were written in the source code separated by commas.
- Within instantiation, this list is substituted by the list of commaseparated template arguments.

Instantiated class X<int, bool, double>:

```
struct X<int, bool, double> { bool, double }; // for illustration, does not compile
```

- Where such expansion can be used?
 - Wherever a comma-separated list of template parameters/arguments can be used.
 - In case of type template parameter list ⇒ wherever one can use a commaseparated list of types.

Variadic templates — tuple

- For instance, a comma-separated list of types can be used as a list of template arguments.
- Example tuple-like class.
- Tuple = generalized pair.
- std::pair is a class template that accepts two type template arguments.
- std::pair<T1,T2> class is an "owner" that owns/manages two objects.
 - First object is of type T1, second object is of type T2.
 - Both objects are stored in the "included" storage of the pair owner.
- std::tuple class template generalizes this concept to any number of owned/managed objects.
 - → An arbitrary number of type template arguments needs to be accepted.
 - ⇒ std::tuple is a variadic template.

Defined in header <tuple>
template < class... Types >
class tuple;

std::tuple

How to define a tuple-like class template?

```
template <typename... Ts>
class Tuple { /* ??? */ };
```

- We would want to define member variables one for each template argument (having its type).
- Unfortunately, this is not directly possible with the template parameter pack expansion mechanism.
- Possible solution recurrent definition:
 - Tuple with some template arguments can be defined as a class with the following member variables:
 - First member variable has the type of the first template argument.
 - Second member variable is a tuple with the remaining template arguments.
 - $\bullet \Rightarrow$ We need to "*recognize"* the first template argument.
 - \Rightarrow It cannot be "mapped" to the template parameter pack.

```
template <typename T, typename... Ts>
class Tuple {
  T t_;
  Tuple< Ts... > ts_;
  // ... other code (constructors, member functions, etc.)
```

- As always with recurrence, we need some "base case".
- In our case, it will be a *specialization for a single template argument*:

```
template <typename T, typename... Ts> // primary template
class Tuple {
   T t_;
   Tuple< Ts... > ts_;
   // no other code now
};

template <typename T> // specialization for a single template argument
class Tuple<T> {
   T t_;
};
```

- Note instantiation of Tuple template with a single template argument in fact corresponds with both the primary template (with empty parameter pack) and the specialization.
- However, C++ standard rules gives higher "priority" to the specialization.

```
Tuple<int> t; // no ambiguity; will be instantiated from the specialization
```

• Example:

 It does not match the specialization ⇒ it will be instantiated from the primary template as:

```
class Tuple< int, bool, double > {
  int t_;
  Tuple< bool, double > ts_;
};
```

- This instantiation initiates the instantiation of Tuple with bool and double template arguments.
- This is, again, instantiated from the primary template as:

```
class Tuple< bool, double > {
  bool t_;
  Tuple< double > ts_;
};
```

```
class Tuple< bool, double > {
  bool t_;
  Tuple< double > ts_;
};
```

- This instantiation initiates the instantiation of Tuple with double template argument.
- This instantiation matches both primary template and specialization, but specialization gets priority.

```
class Tuple< double > {
  double t_;
};
```

Summary of instantiated Tuple classes:

```
class Tuple< double > {
   double t_;
};

class Tuple< bool, double > {
   bool t_;
   Tuple< double > ts_;
};

class Tuple< int, bool, double > {
   int t_;
   Tuple< bool, double > ts_;
};
```

• The initial instantiation is kind-of effectively equivalent to the following one:

```
class Tuple< int, bool, double > {
  int t_;
  class Tuple< bool, double > {
    bool t_;
    class Tuple< double > {
       double t_;
    } ts_;
  } ts_;
}
```

- → The instance effectively but not directly contains member variables of types of all template arguments.
- Notes:
 - This is a basic definition useless for practice.
 - However, it showed how variadic templates and templates parameter packs are commonly defined and resolved in practice.
 - For practical purposes, we would need to add constructors, functions for accessing owned objects, etc., which is beyond the scope of this lecture.