

Effective C++ Programming

NIE-EPC (v. 2021):
 TEMPLATES, TEMPLATE INSTANTIATION, TEMPLATE
 ARGUMENT DEDUCTION, TEMPLATE SPECIALIZATION,
 VARIADIC TEMPLATES, TEMPLATE PARAMETER PACK
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Templates — templated entities

- *Templates* = parametrized *generators of source code entities*.
- Possible *template-generated entities*:
 - *classes*,
 - *functions* (both free and member),
 - *variables* (since C++17),
 - *types* (*type aliases*; since C++11),
 - *concepts* (since C++20).
- *Template definition* = prescription on how the entity should be generated (with respect to template arguments).
 - Its generation is called "*instantiation*".
 - Resulting entity is called "*instance*" of a template.
- *Example*:
 - `wrapper` = *class template*, `wrapper<int>` = its instance = *class*.

```
template <typename T>
struct wrapper {
    T obj;
};
```

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Templates — parameters and arguments

- Templates = parametrized generators of *source code entities*.
- *Implication* — *generation/instantiation* must happen *at compile time*.
 - It effectively generates source code that needs to be further compiled by a compiler.
 - *Note* — at *runtime*, templates no more exist (only their instances do).
- Generation of source code entities is parametrized by *template parameters*.
 - Arguments for these parameters — *template arguments* — need to be resolvable at compile time.
 - ⇒ Basically, there are two kinds of template parameters:
 - *type template parameters* — their arguments are types,
 - *non-type template parameters* — their arguments are compile-time constants; mostly *integral/Boolean*.
- *Example*:

```
template <typename T, size_t N> // T - type template param., N - non-type template param.
struct Array {
    T data[N];
};
```

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Templates — explicit instantiation

- Templates can be instantiated either *explicitly* or *implicitly*.
- *Explicit instantiation* forces a compiler to generate template instance without its usage.
- *Example* — source code...

```
template <typename T> bool nonzero(T t) { return t != 0; } // function template definition
template bool nonzero<int>(int); // explicit instantiation of function nonzero<int>
template bool nonzero<long>(long); // explicit instantiation of function nonzero<long>
```

- ...effectively equivalent source code after instantiation...

```
bool nonzero<int>(int t) { return t != 0; }
bool nonzero<long>(long t) { return t != 0; }
```

- ...resulting machine code:

```
bool nonzero<int>(int):
    test edi, edi
    setne al
    ret
bool nonzero<long>(long):
    test rdi, rdi
    setne al
    ret
```

- *Note* — explicit template instantiation is rarely used in practice.

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Templates — implicit instantiation

- *Implicit instantiation* forces a compiler to generate template instance by using the template.
- *Example* — source code...

```
template <typename T> bool nonzero(T t) { return t != 0; } // function template definition
int main() {
    std::cout << (void*)nonzero<int>; // print address of function nonzero<int>
}
```

- ...effectively equivalent source code after instantiation...

```
bool nonzero<int>(int t) { return t != 0; }
int main() {
    std::cout << (void*)nonzero<int>;
}
```

- ...resulting machine code:

```
bool nonzero<int>(int):
    test edi, edi
    setne al
    ret
main:
    ...
```

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Templates — header files

- *Implicit instantiation* is initiated by using the template.
- For instantiation, *template definition* needs to be available.
- *Implication* — if the template is not instantiated explicitly, its definition needs to be available in each translation unit where that template is used.
- *Typical solution of this problem* = putting definitions of templates into *header files*.
- *Note*:
 - In case of *function templates*, this can effectively result in the presence of a machine code of its instance — function — in multiple *translation units* = *object files*.
 - With templates, this does not break ODR (*one definition rule*).

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Template argument deduction

- In *all cases*, template arguments may be provided explicitly.
- In *some cases*, they may be implicitly *deduced* by a compiler.
- Such *deduction* is mostly based:
 - either on *function arguments* for function templates,
 - or on *initialization = constructor arguments* for class templates (since C++17).
- ⇒ For deduction, template parameter needs to be somehow involved in the form of a function (constructor) parameter.

Example:

```
template <typename T> bool nonzero(T t) { return t != 0; }
int main(int argc, char* argv[]) {
    std::cout << nonzero(argv) << '\n'; // explicitly-provided template argument
    std::cout << nonzero(argv) << '\n'; // template argument deduced by function argument
}
```

• ...effectively equivalent source code after instantiation...

```
bool nonzero(int& t) { return t != 0; }
int main(int argc, char* argv[]) {
    std::cout << nonzero(argv) << '\n';
    std::cout << nonzero(argv) << '\n';
}
```

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Template argument deduction (cont.)

- Deduction of template arguments must be *unambiguous*.
- Otherwise, its resolution cannot take place.
- Example:

```
template <typename T>
bool larger(T a, T b) { return a > b; }
int main() {
    std::cout << larger(1, 2.0);
    // error: no matching function to call 'larger(int, double)'
    // note: template argument deduction/substitution failed:
    // note: deduced conflicting types for parameter 'T' ('int' and 'double')
}
```

Resolution:

- According to the *first function argument* (literal of type `int`), the template argument is deduced as `int`.
- According to the *second function argument* (literal of type `double`), the template argument is deduced as `double`.
- None of these options has a higher priority; they are equivalent for the compiler.
- ⇒ The compiler cannot decide which one to use.

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Forwarding reference revisited

- Example — forwarding reference as a function template parameter:

```
template <typename T> void f(T&& param) { }
```

- It can bound both *lvalue* and *rvalue* arguments of *any type*...
- ...and it has special *template argument deduction rules* different for both these cases:
 - In case of *lvalue argument* of type `X`, `T` is deduced as `X&`.
 - In case of *rvalue argument* of type `X`, `T` is deduced as `X`.

```
int i = 1;
f(i); // i is lvalue of type int => T is deduced as int& => type of param is int&
f(true); // true is rvalue of type bool => T is deduced as bool => type of param is bool&&
```

- Effectively equivalent source code after instantiation:

```
void f(int&& param) { }
void f(bool&& param) { }
```

```
int i = 1;
f(i);
f(true);
```

- Note — once instantiation takes place, forwarding reference becomes either *lvalue* or *rvalue reference*.
- ⇒ There are no more forwarding references after template instantiation.

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Template argument substitution

- Once the template argument is resolved (explicitly provided, deduced,...) it is *substituted for the corresponding template parameter*.
- Substitution of template arguments effectively results in an instance of that template.

Example:

```
template <typename T> bool larger(const T& a, const T& b) { return a > b; }
int main() {
    std::cout << larger(1, 2);
}
```

- Here, template argument is deduced as `int` according to function arguments (both of type `int`).
- ⇒ `int` is substituted for `T`.
- ⇒ Effectively equivalent source code after instantiation:

```
bool larger(int& a, const int& b) { return a > b; }
int main(int argc, char* argv[]) {
    std::cout << larger(argv[1], argv[2]);
}
```

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Template specialization

- Template *specialization* allows having multiple definitions of the same template...
- ...while the one used for instantiation is selected according to the template arguments.
- One definition needs to be generic — it is called *primary template*.
- Other definitions are used for instantiation in special cases only ⇒ they are called "*specializations*".
- Instantiation:
 - When the template arguments match some specialization, it is used for instantiation.
 - Otherwise, primary template is used instead.
- Typically, only *class templates* are specialized in practice.
- With *function templates*, *overloading* is preferred.
- Mixing overloading with function template specialization is "*dangerous*".

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Template specialization — examples

- Example — class that *maps* type to a static Boolean constant which has value:
 - true* if the type is `bool`,
 - false* otherwise.

```
template <typename T>
struct is_bool { static const bool value = false; };
template <> // specialization for bool template argument
struct is_bool<bool> { static const bool value = true; };

int main() {
    std::cout << is_bool<int>::value; // (1) prints out "0"
    std::cout << is_bool<bool>::value; // (2) prints out "1"
}
```

Resolution:

- In (1), `int` does not match the specialization ⇒ class `is_bool<int>` is instantiated from the *primary template*.
- In (2), `bool` does match the specialization ⇒ class `is_bool<bool>` is instantiated from this *specialization*.

- ⇒ Effectively equivalent code after instantiation:

```
struct is_bool_int { static const bool value = false; };
struct is_bool_bool { static const bool value = true; };

int main() {
    std::cout << is_bool_int::value;
    std::cout << is_bool_bool::value;
}
```

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Template specialization — examples (cont.)

- Another example — factorial calculation.

- Primary template is based on recurrent factorial calculation:

```
template <int I>
struct factorial { static const int value = 1 * factorial<I-1>::value; }; // primary template

template <> struct factorial<0> { static const int value = 1; }; // specialization

int main() {
    std::cout << factorial<3>::value; // prints out 6
}
```

- Resolution:

- Class factorial<3> does not match the specialization \Rightarrow it is instantiated from the primary template as:

```
struct factorial<3> { static const int value = 3 * factorial<2>::value; }; ←
```

- This instance (class) contains the use of template factorial with template argument 2 \Rightarrow it needs to be instantiated as well.
- Class factorial<2> does not match the specialization \Rightarrow it is instantiated (again) from the primary template as:

```
struct factorial<2> { static const int value = 2 * factorial<1>::value; }; ←
```

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Template specialization — examples (cont.)

```
template <int I>
struct factorial { static const int value = 1 * factorial<I-1>::value; }; // primary template
template <> struct factorial<0> { static const int value = 1; }; // specialization

struct factorial<2> { static const int value = 2 * factorial<1>::value; }; ←
```

- Resolution (cont.):

- This instance (class) contains the use of template factorial with template argument 1 \Rightarrow it needs to be instantiated as well.
- Class factorial<1> does not match the specialization \Rightarrow it is instantiated (again) from the primary template as:

```
struct factorial<1> { static const int value = 1 * factorial<0>::value; }; ←
```

- This instance (class) contains the use of template factorial with template argument 0 \Rightarrow it needs to be instantiated as well.
- Class factorial<0> does match the specialization \Rightarrow it is instantiated from the specialization as:

```
struct factorial<0> { static const int value = 1; }; ←
```

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Template specialization — examples (cont.)

- Summary:

```
template <int I>
struct factorial { static const int value = 1 * factorial<I-1>::value; }; // primary template
template <> struct factorial<0> { static const int value = 1; }; // specialization

int main() {
    std::cout << factorial<3>::value; // prints out 6
}
```

- Effectively equivalent code after instantiation...

```
struct factorial<0> { static const int value = 1; };
struct factorial<1> { static const int value = 1 * factorial<0>::value; };
struct factorial<2> { static const int value = 2 * factorial<1>::value; };
struct factorial<3> { static const int value = 3 * factorial<2>::value; };

int main() {
    std::cout << factorial<3>::value;
}
```

- ...and after resolution of the constants:

```
struct factorial<0> { static const int value = 1; };
struct factorial<1> { static const int value = 1; };
struct factorial<2> { static const int value = 2; };
struct factorial<3> { static const int value = 6; };

int main() {
    std::cout << factorial<3>::value;
}
```

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Template metafunctions and type traits

- Recall — class templates that maps template arguments (types or compile-time constants) to other types or compile-time constants are generally called template metafunctions.
- Both introduced templates `is_bool` and `factorial` fall into this category.
 - `is_bool` maps type to compile-time Boolean constant,
 - `factorial` maps compile-time integral constant to other compile-time integral constant.
- Metafunctions that works with types are called "type traits".
- We have already seen some type traits from the C++ standard library, such as `std::is_reference` or `std::remove_reference`.
- Note — definition of metafunctions is mostly based on template specialization.

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Template metafunctions — examples

- Metafunction `std::is_reference` — maps type to a Boolean constant which is:

- true if type is a reference type,
- false otherwise.

- Solution:

- primary template for non-reference types,
- specialization for reference types.

```
template <typename T> struct is_reference // primary template
{ static const bool value = false; };
template <typename T> struct is_reference<T&> // specialization for lvalue references
{ static const bool value = true; };
template <typename T> struct is_reference<T&&> // specialization for rvalue references
{ static const bool value = true; };
```

- Partial vs full specialization:

- If a specialization is no more parametrized, it is called "full specialization".
- If a specialization is parametrized, it is called "partial specialization".

```
template <> struct factorial<0> { static const int value = 1; }; // (full) specialization ←
```

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Template metafunctions — examples (cont.)

- Example:

```
template <typename T> struct is_reference { static const bool value = false; };
template <typename T> struct is_reference<T&> { static const bool value = true; };
template <typename T> struct is_reference<T&&> { static const bool value = true; };

int main() {
    std::cout << is_reference<float>::value; // prints out "0"
    std::cout << is_reference<bool&>::value; // prints out "1"
    std::cout << is_reference<int&&>::value; // prints out "1"
}
```

- Resolution:

- `float` is non-reference type \Rightarrow it does not match any specialization \Rightarrow `is_reference<float>` is instantiated from primary template as:

```
struct is_reference<float> { static const bool value = false; }; ←
```

- `bool&` is reference type \Rightarrow it does match a specialization for lvalue references \Rightarrow `is_reference<bool&>` is instantiated from it as:

```
struct is_reference<bool&> { static const bool value = true; }; ←
```

- `int&&` is reference type \Rightarrow it does match a specialization for rvalue references \Rightarrow `is_reference<int&&>` is instantiated from it as:

```
struct is_reference<int&&> { static const bool value = true; }; ←
```

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Template metafunctions — examples (cont.)

• Summary:

```
template <typename T> struct is_reference { static const bool value = false; };
template <typename T> struct is_reference<T&> { static const bool value = true; };
template <typename T> struct is_reference<T&&> { static const bool value = true; };

int main() {
    std::cout << is_reference<float>::value; // prints out "0"
    std::cout << is_reference<bool&>::value; // prints out "1"
    std::cout << is_reference<int&&>::value; // prints out "1"
}
```

• Effectively equivalent code after instantiation:

```
struct is_reference<float> { static const bool value = false; };
struct is_reference<bool&> { static const bool value = true; };
struct is_reference<int&&> { static const bool value = true; };

int main() {
    std::cout << is_reference<float>::value;
    std::cout << is_reference<bool&>::value;
    std::cout << is_reference<int&&>::value;
}
```

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Template metafunctions — examples (cont.)

• Metafunction `std::remove_reference` — maps input type `T` to an output type which is:

- type referred to by `T` if `T` is a reference type,
- `T` otherwise (identity).

• Solution:

- primary template for non-reference types,
- specialization for reference types.

```
template <typename T> struct remove_reference { using type = T; };
template <typename T> struct remove_reference<T&> { using type = T; };
template <typename T> struct remove_reference<T&&> { using type = T; };
```

• Example:

```
int main() {
    remove_reference<bool>::type b = true; // type of b is bool
    remove_reference<int&>::type i = 1; // type of i is int
}
```

• Effectively equivalent code after instantiation:

```
struct remove_reference<bool> { using type = bool; };
struct remove_reference<int&> { using type = int; };

int main() {
    remove_reference<bool>::type b = true;
    remove_reference<int&>::type i = 1;
}
```

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Template metafunctions — “shortcuts”

• Recall — usage of metafunctions/type traits requires relatively long writing:

```
template <typename T>
typename remove_reference<T>::type && move(T&& param) {
    return static_cast<typename remove_reference<T>::type &&> param;
}
```

• For type traits which “produce” types, we can create a “shortcut” type (alias) template wrapper as follows:

```
template <typename T>
using remove_reference_t = typename remove_reference<T>::type;
```

• Consequently...

```
typename remove_reference<T>::type
```

• ...may be then replaced by:

```
remove_reference_t<T>
```

• For type traits in the C++ standard library, such “shortcut” type templates were introduced in C++14.

```
template <typename T> std::remove_reference_t<T> && move(T&& param)
{ return static_cast<std::remove_reference_t<T> &&> param; }
```

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Template metafunctions — “shortcuts” (cont.)

• Similarly, such “shortcut” wrappers may be created for type traits which “produce” values.

• However, these require variable templates, which were not available until C++17.

```
template <typename T>
inline bool is_reference_v = is_reference<T>::value;
```

• Consequently...

```
is_reference<T>::value
```

• ...may be then replaced by:

```
is_reference_v<T>
```

• For type traits in the C++ standard library, such “shortcut” variable templates were introduced in C++17.

```
std::is_reference_v<T> // since C++17
```

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Template specialization — example

• Template specialization has many uses other than just template metafunctions.

• C++ standard library examples:

- `std::atomic` — specializations for integers, pointers,...
- `std::hash` — specialization for different types;
- `std::vector<bool>` — possible bit-compressed implementation.

• Another example — some parallel computation based on MPI (see NI(E)-PDP):

- MPI library has a “C” API (no templates).
- MPI functions that send/receive data takes:
 - a non-type pointer to these data (of type `void*`),
 - special argument “type tag” that tells the library of what type these data are.
- Exemplary MPI type tags: `MPI_INT` for `int`, `MPI_DOUBLE` for `double`,...
- Exemplary function for sending data to another MPI process:

```
int MPI_Send(void* buf, int count, MPI_Datatype mdt, int dest, int tag, MPI_Comm comm);
```

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Template specialization — example (cont.)

• If we know the type of sent data, everything works fine:

```
const int n = 4;
double a[] = { 1.0, 2.0, 3.0, 4.0 };
MPI_Send(&a, 1, MPI_INT, 0, 0, MPI_COMM_WORLD);
MPI_Send(&a, n, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD);
```

• Problem — what if type of sent data is a template parameter?

• Example — parallel computation parametrized by floating-point data type, which defines its used precision:

```
template <typename T> // T allows to choose FP precision (expected float or double)
void some_parallel_computation() {
    std::vector<T> v; // result of local process
    ... // calculate local result
    // send it to root process (process number 0):
    MPI_Send(v.data(), v.size(), MPI_??? , 0, 0, MPI_COMM_WORLD);
}
```

• How to specify MPI type tag that specifies type of sent data?

• Required functionality:

- If `T` is `float`, `MPI_FLOAT` type tag should be used.
- If `T` is `double`, `MPI_DOUBLE` type tag should be used.

• Generally, we need to map types to corresponding MPI type tags.

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Template specialization — example (cont.)

- Possible solution — template specialization.
 - Note — MPI type tags may be in fact pointers into some type table ([link](#)).
 - ⇒ Generally, they may not be compile-time constants.
- Alternative — static member function that returns tag value.

```
template <typename T> struct GetMPIDataType; // primary template; declaration only
template< struct GetMPIDataType<float>
{ static MPI_Datatype get() { return MPI_FLOAT; } }; // specialization for float
template< struct GetMPIDataType<double>
{ static MPI_Datatype get() { return MPI_DOUBLE; } }; // specialization for double
... // similarly for other types supported by MPI library (MPI_INT, MPI_LONG,...)
```

- Application to our problem:

```
MPI_Send(v.data(), v.size(), GetMPIDataType<T>::get(), 0, 0, MPI_COMM_WORLD);
```

- Example — computation to be instantiated with *single precision*:

```
MPI_Send(v.data(), v.size(), GetMPIDataType<float>::get(), 0, 0, MPI_COMM_WORLD);
```

- This initiates instantiation of `GetMPIDataType<float>` from a corresponding specialization:

```
struct GetMPIDataType<float>
{ static MPI_Datatype get() { return MPI_FLOAT; } };
```

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Templates — complex example

- Consider our implementation of Vector class template...

```
template <typename T> class Vector {
    size_t capacity_, size_;
    T* data_;
    ...
};
```

- ...and its following exemplary use:

```
struct X {
    X(int); // converting constructor
    X(int, bool); // another converting constructor
    X(const X&); // copy constructor
    X(X&&); // move constructor
};

int main() {
    Vector<X> v;
}
```

- Vector class template is instantiated with X template argument, which results in `Vector<X>` class ⇒ X is substituted for T:

```
class Vector<X> {
    size_t capacity_, size_;
    X* data_;
    ...
};
```

- Let us now look at some Vector member functions.

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Templates — complex example (cont.)

- First, the reserve member function was defined as:

```
void reserve(size_t capacity) {
    if (capacity <= capacity_) return;
    T* data = (T*)operator new(capacity * sizeof(T));
    for (size_t i = 0; i < size_; i++) new (data + i) T(move( *(data_ + i) ));
    clear();
    ::operator delete(data_);
    data_ = data; capacity_ = capacity;
}
```

- After *instantiation/substitution*, this effectively results in the following definition:

```
void reserve(size_t capacity) {
    if (capacity <= capacity_) return;
    X* data = (X*)operator new(capacity * sizeof(X));
    for (size_t i = 0; i < size_; i++) new (data + i) X( move( *(data_ + i) ) );
    clear();
    ::operator delete(data_);
    data_ = data; capacity_ = capacity;
}
```

- This initiates the instantiation of move function template.
- Note — we deliberately used our custom version here instead of `std::move` to show what happens next.

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Templates — complex example (cont.)

- Our custom move function template definition:

```
template <typename T>
typename remove_reference<T>::type && move(T&& param) {
    return static_cast<typename remove_reference<T>::type &&> param;
}
```

- Note — similarly as above, we used our `remove_reference` trait.

- Resolution:

- Template argument of move call is not specified.
 - ⇒ It needs to be deduced according to the function argument.
- Function argument of move is lvalue expression of type X.
- param is a forwarding reference.
 - ⇒ T is deduced as X&.

- This effectively results in the following instantiation of move:

```
remove_reference<X&>::type && move<X&>(X&& param) {
    return static_cast<remove_reference<X&>::type &&> param;
}
```

- Now, this instantiation initiates instantiation of `remove_reference<X&>` class.

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Templates — complex example (cont.)

- `remove_reference<X&>` is instantiated as follows:

```
struct remove_reference<X&> { using type = X; }
```

- ⇒ For illustration, we can then effectively rewrite...

```
remove_reference<X&>::type && move<X&>(X&& param) {
    return static_cast<remove_reference<X&>::type &&> param;
}
```

- ...as...

```
X&& move<X&>(X&& param) { return static_cast<X&&> param; }
```

- ...and this instance of move is called in `Vector<X>::reserve`:

```
void reserve(size_t capacity) {
    ...
    for (size_t i = 0; i < size_; i++) new (data + i) X( move<X&>( *(data_ + i) ) );
    ...
}
```

- This is exactly what we expected — `move<X&>` call is an expression that:
 - represents the original vector element `*(data_ + i)` of type X,
 - and its category is *rvalue*.
- ⇒ New elements are initialized by *move constructor* of class X.

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Templates — complex example (cont.)

- Summary for `Vector<X>::reserve...`

```
template <typename T> class Vector {
    size_t capacity_, size_;
    T* data_;
public:
    void reserve(size_t capacity) {
        if (capacity <= capacity_) return;
        T* data = (T*)operator new(capacity * sizeof(T));
        for (size_t i = 0; i < size_; i++) new (data + i) T( move( *(data_ + i) ) );
        clear();
        ::operator delete(data_);
        data_ = data; capacity_ = capacity;
    }
};
```

- ...is effectively instantiated as follows:

```
struct remove_reference<X&> { using type = X; }
remove_reference<X&>::type && move<X&>(X&& param) {
    return static_cast<remove_reference<X&>::type &&> param;
}

class Vector<X> {
    size_t capacity_, size_;
    X* data_;
public:
    void reserve(size_t capacity) {
        if (capacity <= capacity_) return;
        X* data = (X*)operator new(capacity * sizeof(X));
        for (size_t i = 0; i < size_; i++) new (data + i) X( move<X&>( *(data_ + i) ) );
        clear();
        ::operator delete(data_);
        data_ = data; capacity_ = capacity;
    }
};
```

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Templates — complex example (cont.)

- Next, the `push_back` member functions were defined as:

```
void push_back(const T& param) { // overload for lvalues
    if (size_ == capacity_) reserve(capacity_ * 2 * capacity_ : 1);
    new (data_ + size_) T( param );
    size_++;
}

void push_back(T&& param) { // overload for rvalues
    if (size_ == capacity_) reserve(capacity_ * 2 * capacity_ : 1);
    new (data_ + size_) T( move( param ) );
    size_++;
}
```

- After *instantiation/substitution*, this effectively results in the following definitions:

```
void push_back(const X& param) {
    if (size_ == capacity_) reserve(capacity_ * 2 * capacity_ : 1);
    new (data_ + size_) X( param );
    size_++;
}

void push_back(X&& param) {
    if (size_ == capacity_) reserve(capacity_ * 2 * capacity_ : 1);
    new (data_ + size_) X( move(X&) param );
    size_++;
}
```

- Note — move call after deduction results in the call of the already instantiated function `move<X&>`.

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Variadic templates

- What about `emplace_back` member function?

```
template <typename... U> void emplace_back(US&... param) {
    if (size_ == capacity_) reserve(capacity_ * 2 * capacity_ : 1);
    new (data_ + size_) T( std::forward<U>(param)... );
    size_++;
}
```

- To understand it, we first need to explain *variadic templates*.
- Recall — *template parameter* starts with
 - typename (or class) keyword for *type template parameters*,
 - type name for *non-type template parameters*,
- and is then it is (optionally) followed by identifier (template parameter name).

```
template <typename T, class U, int I, size_t N> struct X { };
```

- When we write *ellipsis* (sequence of 3 dots) after the *typename/class/type name*, we create a so-called “*template parameter pack*” (instead of a *template parameter*).

```
template <typename T, typename... Ts> struct X { };
// T is template parameter, Ts is a template parameter pack
```

- Template with at least one template parameter pack is called **variadic template**.

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Variadic templates (cont.)

- Recall — when template is instantiated, *template parameters* are substituted by *template arguments*.
- Namely, a single template parameter is substituted by a single template argument.
- On the contrary, *template parameter pack* may be substituted by an arbitrary number of template arguments, including no one.

```
template <typename T, typename... Ts> struct X { };
X<int, bool, double> x1; // T is substituted by int, Ts is substituted by bool and double
X<int, bool> x2; // T is substituted by int, Ts is substituted by bool
X<int> x3; // T is substituted by int, Ts is empty
```

- Effectively, template parameter pack represents a list of virtual *unnamed* template parameters.
- It is not possible to refer to individual elements of template parameter pack.
- However, it is possible to so-call “*expand*” it.

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Variadic templates — expansion

- Template parameter pack expansion* happens when its *identifier* is followed by *ellipsis* in the template definition.
- Effect* — as if the list of virtual template parameters were written in the source code separated by commas.
- Within *instantiation*, this list is substituted by the list of comma-separated template arguments.

```
template <typename T, typename... Ts> struct X { Ts... };
X<int, bool, double> x; // Ts is substituted by bool and double
// => Ts... is literally expanded to “bool, double”
```

- ⇒ Instantiated class `X<int, bool, double>`:

```
struct X<int, bool, double> { bool, double }; // for illustration, does not compile
```

- Where such expansion can be used?
 - Wherever a *comma-separated list of template parameters/arguments* can be used.
 - In case of type template parameter list ⇒ wherever one can use a *comma-separated list of types*.

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Variadic templates — tuple

- For instance, a comma-separated list of types can be used as a list of template arguments.
- Example — *tuple-like class*.
- Tuple* = generalized pair.
- `std::pair` is a class template that accepts two type template arguments.
- `std::pair<T1, T2>` class is an “owner” that owns/manages two objects.
 - First object is of type `T1`, second object is of type `T2`.
 - Both objects are stored in the “included” storage of the pair owner.
- `std::tuple` class template generalizes this concept to any number of owned/managed objects.
 - ⇒ An arbitrary number of type template arguments needs to be accepted.
 - ⇒ `std::tuple` is a variadic template.

```
std::tuple
Defined in header <tuple>
template< class... Types >
class tuple;
```

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Variadic templates — tuple (cont.)

- How to define a tuple-like class template?

```
template <typename... Ts>
class Tuple { /* ??? */ };
```

- We would want to define member variables — one for each template argument (having its type).
- Unfortunately, this is not directly possible with the template parameter pack expansion mechanism.
- Possible solution — recurrent definition:
 - Tuple with some template arguments can be defined as a class with the following member variables:
 - First member variable has the type of the first template argument.
 - Second member variable is a tuple with the remaining template arguments.
 - ⇒ We need to “*recognize*” the first template argument.
 - ⇒ It cannot be “*mapped*” to the template parameter pack.

```
template <typename T, typename... Ts>
class Tuple {
    T t_;
    Tuple< Ts... > ts_;
    // ... other code (constructors, member functions, etc.)
}
```

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Variadic templates — tuple (cont.)

- As always with recurrence, we need some “base case”.
- In our case, it will be a *specialization for a single template argument*:

```
template <typename T, typename... Ts> // primary template
class Tuple {
    T t_;
    Tuple< Ts... > ts_;
    // no other code now
};

template <typename T> // specialization for a single template argument
class Tuple<T> {
    T t_;
};
```

- **Note** — instantiation of Tuple template with a single template argument in fact corresponds with both the primary template (with empty parameter pack) and the specialization.
- However, C++ standard rules gives higher “priority” to the specialization.

```
Tuple<int> t; // no ambiguity; will be instantiated from the specialization
```

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Variadic templates — tuple (cont.)

- **Example:**

```
template <typename T, typename... Ts> class Tuple {
    T t_;
    Tuple< Ts... > ts_;
};

Tuple< int, bool, double > t; // T substituted for int, Ts for bool and double
// Ts... is expanded to bool, double
```

- It does not match the specialization \Rightarrow it will be instantiated from the primary template as:

```
class Tuple< int, bool, double > {
    int t_;
    Tuple< bool, double > ts_;
};
```

- This instantiation initiates the instantiation of Tuple with bool and double template arguments.
- This is, again, instantiated from the primary template as:

```
class Tuple< bool, double > {
    bool t_;
    Tuple< double > ts_;
};
```

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Variadic templates — tuple (cont.)

```
class Tuple< bool, double > {
    bool t_;
    Tuple< double > ts_;
};
```

- This instantiation initiates the instantiation of Tuple with double template argument.
- This instantiation matches both primary template and specialization, but specialization gets priority.

```
class Tuple< double > {
    double t_;
};
```

- **Summary** of instantiated Tuple classes:

```
class Tuple< double > {
    double t_;
};

class Tuple< bool, double > {
    bool t_;
    Tuple< double > ts_;
};

class Tuple< int, bool, double > {
    int t_;
    Tuple< bool, double > ts_;
};
```

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Variadic templates — tuple (cont.)

- The initial instantiation is kind-of effectively equivalent to the following one:

```
class Tuple< int, bool, double > {
    int t_;
    class Tuple< bool, double > {
        bool t_;
        class Tuple< double > {
            double t_;
        } ts_;
    } ts_;
};
```

- \Rightarrow The instance effectively — but not directly — contains member variables of types of all template arguments.
- **Notes:**
 - This is a basic definition useless for practice.
 - However, it showed how variadic templates and templates parameter packs are commonly defined and resolved in practice.
 - For practical purposes, we would need to add constructors, functions for accessing owned objects, etc., which is beyond the scope of this lecture.

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