# **Transport Layer**

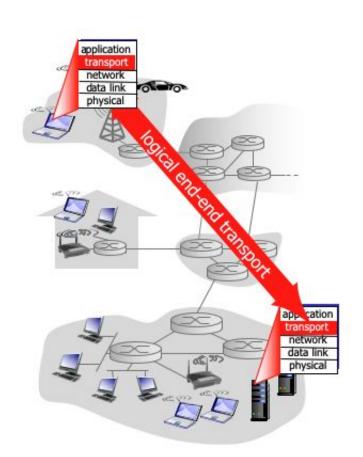
CS5700 Fall 2019

## Agenda

- Transport layer services
- UDP
- Reliable data transfer
- TCP
- Congestion control

## **Transport services and protocols**

- Provide logical communication between application processes
- Run in end systems (not the core)
- More than one transport protocol available to applications
  - TCP and UDP



## What does network layer do?

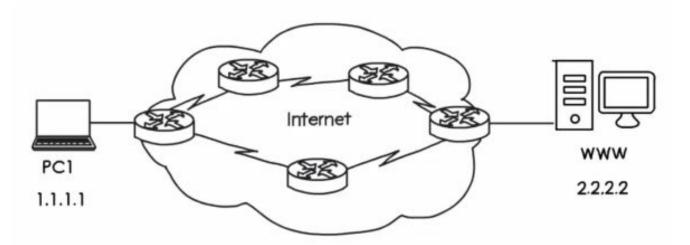
- What's the difference between transport layer and network layer?
- What services are provided by network layer?



application transport network link physical

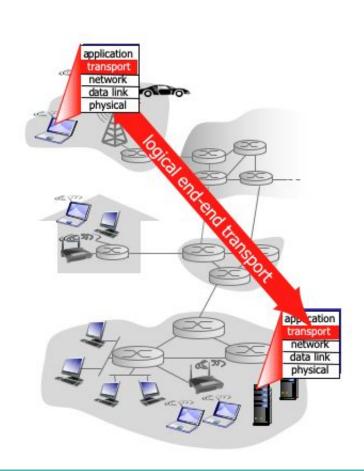
## **Network layer service model**

- Logical communication between hosts
- Every packet is treated individually and separately
- Best effort. No guarantee of delivery.



### **Transport layer protocols**

- TCP
  - Reliable in-order delivery
  - Connection oriented
  - Flow control
  - Congestion control
- UDP
- Services not available
  - Delay or bandwidth guarantee



## **UDP**

#### **UDP**

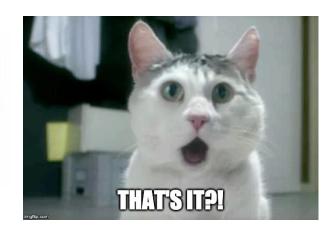
- User Datagram Protocol
  - Connection less
  - No guarantee of delivery
- Where do you see UDP used? Do you know why?



#### **UDP** header

Do you know what's each field for?

16 bit source port	16 bit destination por
16 bit UDP length	16 bit UDP checksum



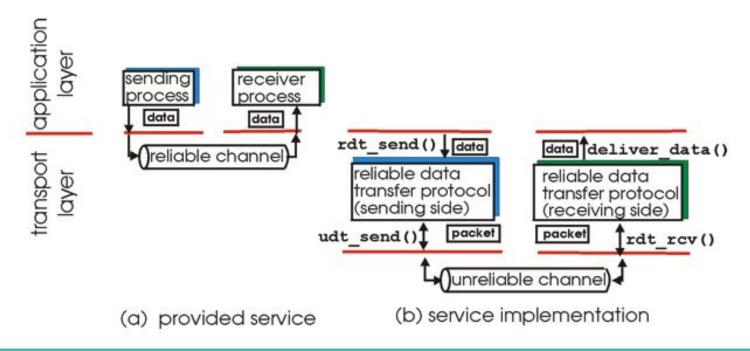
#### **UDP** checksum

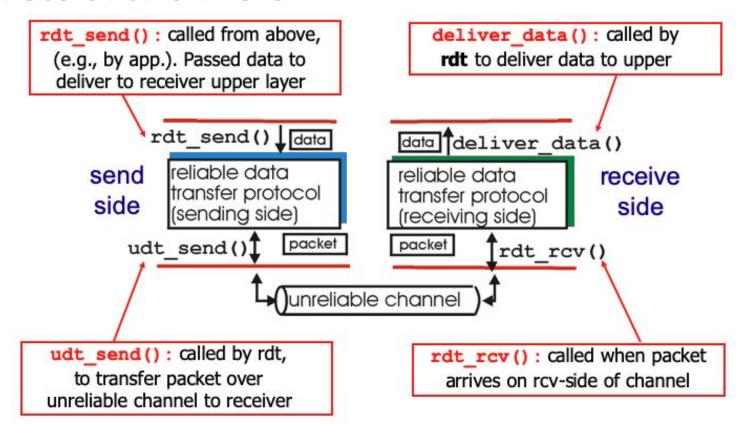
- Detect "errors" (e.g. flipped bits) in transmitted segment
- Sender
  - Treat data (include header) as seq of 16-bit integers
  - Add them up (1's complement), call it checksum
  - Put checksum into UDP header
- Receiver
  - Same algorithm, compute checksum and compare

### **UDP** socket

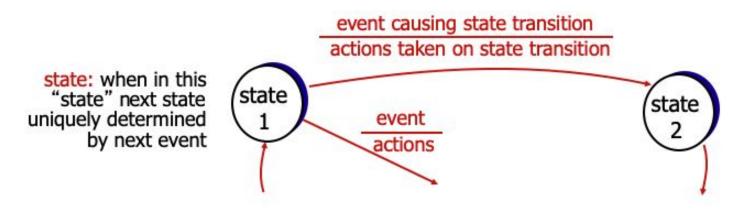


Important in application, transport, and link layers



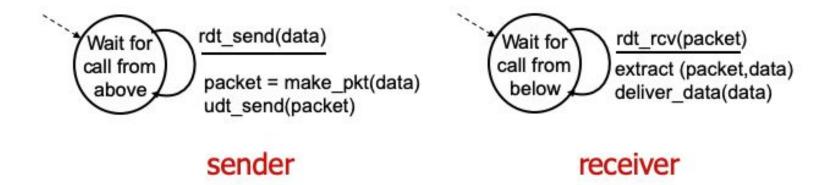


- Incrementally design sender/receiver of rdt
- Consider only unidirectional data transfer
  - But control info will flow on both directions
- Use FSM (finite state machines) to design algorithm



#### rdt1.0: over a reliable channel

- Underlying channel is perfectly reliable
  - No bit errors
  - No loss of packets



#### rdt2.0: channel with bit errors

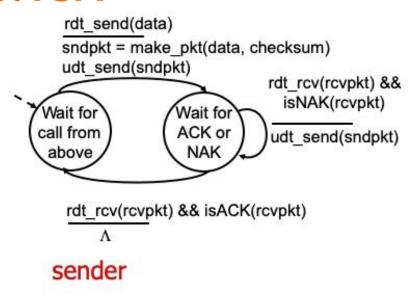
- Underlying channel may flip bits in packet
  - Checksum to detect bit errors
- How to recover from errors?



#### rdt2.0: channel with bit errors

- ACK (acknowledgement)
  - Receiver explicitly tells sender that pkt received OK
- NAK (negative acknowledgement)
  - Receiver explicitly tells sender that pkt had errors
- Sender needs to retransmit pkt on receipt of NAK
- Summary
  - Error detection
  - Feedback with control message ACK and NAK

#### rdt2.0: FSM



#### receiver

rdt\_rcv(rcvpkt) && corrupt(rcvpkt) udt send(NAK) Wait for call from below rdt rcv(rcvpkt) && notcorrupt(rcvpkt) extract(rcvpkt,data) deliver\_data(data) udt send(ACK)

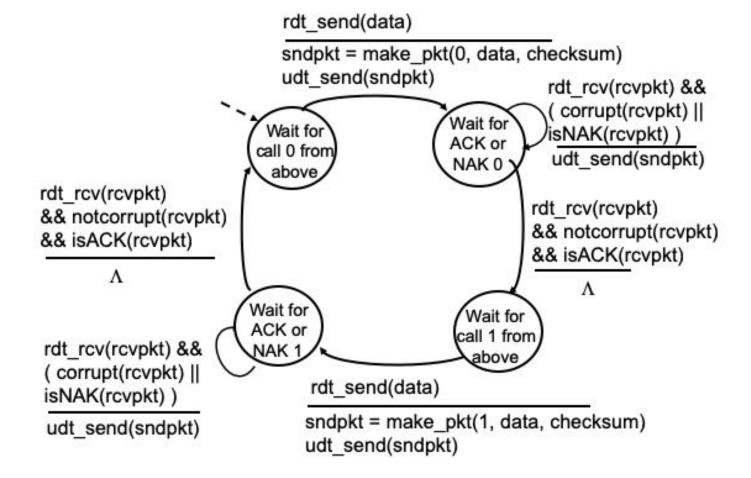
## rdt2.0: anything looks wrong?

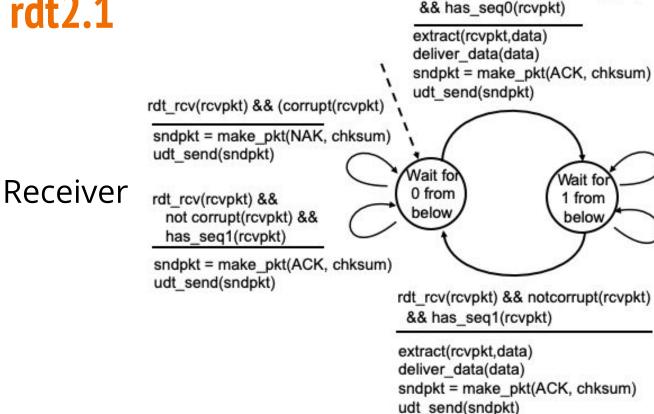


- What happens if ACK or NAK is corrupted?
  - Sender doesn't know what happened at receiver
- Can sender just retransmit?

- Receiver needs to handle duplicates when sender retransmit
- Need to use sequence number!
- Stop and wait algorithms
  - Sequence number either 0 or 1

Sender





rdt rcv(rcvpkt) && notcorrupt(rcvpkt)

1 from

below

rdt\_rcv(rcvpkt) && (corrupt(rcvpkt) sndpkt = make\_pkt(NAK, chksum) udt\_send(sndpkt)

rdt rcv(rcvpkt) && not corrupt(rcvpkt) && has seq0(rcvpkt)

sndpkt = make\_pkt(ACK, chksum) udt\_send(sndpkt)

## rdt2.1: summery

- Sender
  - Add sequence number to packets (either 0 or 1)
  - Retransmit if receives NAK
  - Retransmit if ACK/NAK is corrupted
- Receiver
  - Check if received packet is duplicate (use seq #)
  - Send ACK or NAK for each packet