

Efficient Concurrency-Bug Detection Across Inputs

State-of-art solution

Step1. Input design:产生test inputs来保证code coverage

Step2. Bug Detection: 使用dynamic bug-detection tool找到可能的buggy interleaving

Step3. Bug Validation: 对每个input, 执行多次。来排除suspicious interleaving

但这种方法会降低10-100倍的速度

Motivation

- 不同输入会导致相同的interleaving，现有工具会产生duplicate bug reports.

App.	# All Race Pairs			# All Atom. Vio.			# Buggy Race Pairs			# Buggy Atom. Vio.		
	Total	Unique	DupRate	Total	Unique	DupRate	Total	Unique	DupRate	Total	Unique	DupRate
Click	3114	848	3.6	6145	2298	2.7	6	1	6.0	6	1	6.0
FFT	300	66	4.5	1423	369	3.8	28	4	7.0	35	5	7.0
LU	238	58	4.1	874	163	5.4	28	4	7.0	21	3	7.0
Mozilla	1991	481	4.1	2459	723	3.4	42	6	7.0	42	7	6.0
PBZIP2	293	65	4.5	499	143	3.5	32	8	4.0	39	11	3.5

Table 2: Inefficiency in data-race and atomicity-violation detection across inputs (DupRate measures the average number of inputs that expose the same data race or atomicity violation)

- 减少冗余的工作，不产生duplicate bug reports.
- 使用CFP(concurrency function pair)解决

CFP指在某个Input下，可以并行执行的函数对

```
/*Thread 1*/
foo1(){
    lock(L);
    foo2();
    unlock(L);
    ...
}

/*Thread 2*/
lock(L);
bar();
unlock(L);
```

Figure 2: An example of concurrent functions (For illustration purpose, the vertical position of each code statement in the figure represents when the statement is executed)

Method

- Measure CFP for each input and generate aggregated CFP.
- Identify selected inputs and selected functions under these inputs. 用贪心算法
- Apply existing data-race and atomicity-violation detector to selected inputs functions under selected inputs.

Input1: $CFP_1 = \{\{f_1, f_2\}, \{f_2, f_3\}, \{f_2, f_4\}, \{f_4, f_5\}\}$

Input2: $CFP_2 = \{\{f_1, f_2\}, \{f_3, f_4\}, \{f_3, f_5\}\}$

Input3: $CFP_3 = \{\{f_2, f_3\}, \{f_3, f_4\}\}$

$CFP_{Aggregated} = \{\{f_1, f_2\}, \{f_2, f_3\}, \{f_2, f_4\}, \{f_4, f_5\}, \{f_3, f_4\}, \{f_3, f_5\}\}$

Step 1:

Selected input — Input1

Selected functions — $\{f_1, f_2, f_3, f_4, f_5\}$

$CFP_{Uncovered} = \{\{f_3, f_4\}, \{f_3, f_5\}\}$

Step 2:

Selected input — Input2

Selected functions — $\{f_3, f_4, f_5\}$

$CFP_{Uncovered} = \emptyset$

Figure 8: A toy example of input/function selection

如何产生CFP

- 根据同步信息来判断是否能并行执行
- 分析run-time trace of each thread
 - 计算locks和vector timestamp
 - 根据上一步的信息计算CFP

```
/* .ent, .exi: function entrance, exit;
   .ent_exi.lockset: lockset protecting the critical
   section that holds both entrance and exit;
   .vec_time: vector timestamps calculated using
   order—enforcing synchronization. */
Bool concurrentFunction(f1, f2)
{
    if(f1.thread_id == f2.thread_id) return false;

    /*Can f1 start between f2's entrance and exit?*/
    if ((f1.ent.lockset  $\cap$  f2.ent_exi.lockset) !=  $\emptyset$ ) return false;
    if (f1.ent.vec_time < f2.ent.vec_time) return false;
    if (f1.ent.vec_time > f2.exi.vec_time) return false;

    /*Can f2 start between f1's entrance and exit?*/
    if ((f2.ent.lockset  $\cap$  f1.ent_exi.lockset) !=  $\emptyset$ ) return false;
    if (f2.ent.vec_time < f1.ent.vec_time) return false;
    if (f2.ent.vec_time > f1.exi.vec_time) return false;

    return true; /*f1 and f2 are concurrent*/
}
```

Figure 3: Pseudo code that judges concurrent functions