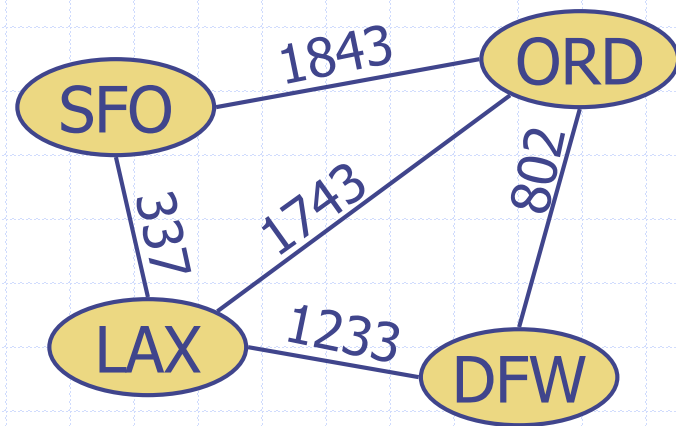


Graphs: Recap



Graphs: Basics

- ◆ Directed vs. undirected graph
 - e.g. edges have a direction associated with them
- ◆ (Non-uniformly) Weighted graph
 - e.g. edges have a weight associated with them
- ◆ Properties
 - $\sum_n \deg(v_n) = 2m$
 - $m \leq n(n-1)/2$
- ◆ Representation
 - Edge list structure,
 - Adjacency list structure, or
 - Adjacency matrix structure

Graphs: Traversals

◆ Depth-first Search

- Traverse deeply first
- $O(n+m)$

◆ Breadth-first Search

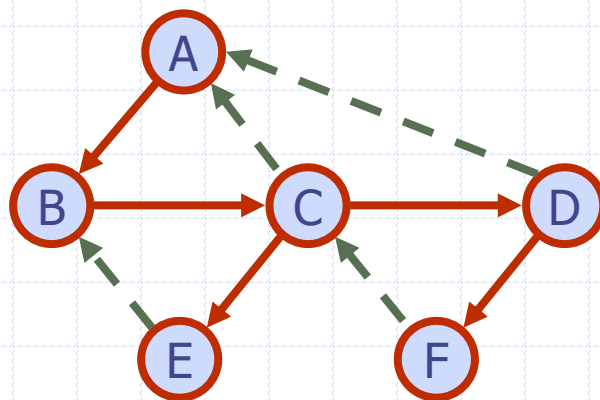
- Traverse broadly first (“breadth”)
- $O(n+m)$

Graphs: Terminology

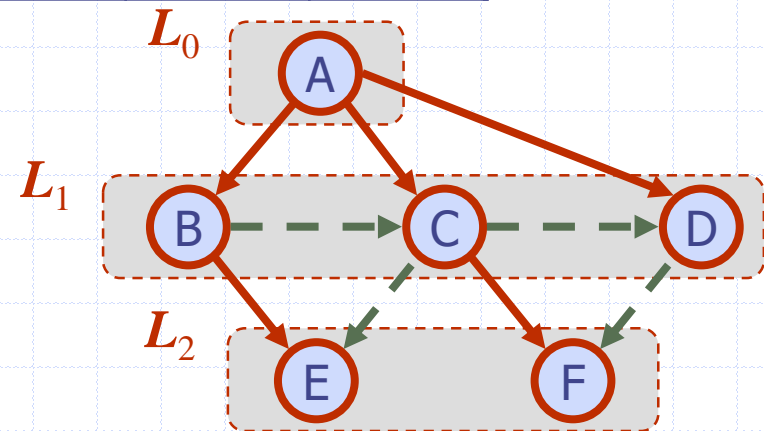
- ◆ Path
- ◆ Connected
- ◆ Subgraph
- ◆ Spanning
- ◆ Biconnected
 - e.g., separation edges or vertices
- ◆ What are these and how do you find them?
 - Connected component
 - Spanning Subgraph
 - Maximally-connected Subgraph
 - Spanning Tree
 - Spanning Forest
 - Biconnected Components

Graphs: DFS vs. BFS

| Applications | DFS | BFS |
|--|-----|-----|
| Spanning forest, connected components, paths, cycles | ✓ | ✓ |
| Shortest paths (for uniformly weighted graphs) | | ✓ |
| Biconnected components | ✓ | |



DFS



BFS

Directed Graphs: Terminology

◆ Reachability

- Is a vertex u “reachable” from v
 - ◆ e.g., can you get to MIA from HNL?

◆ Strongly-Connected Components

- Can you get to any city from any city

◆ Transitive Closure

- If I can get to MIA from HNL and to JFK from MIA, then I can get to JFK from HNL
- The transitive closure graph of a graph G extracts this information
- Algorithms:
 - ◆ Naïve: $O(n(n+m))$ to $O(n^3)$
 - ◆ Floyd-Warshall Algorithm: $O(n^3)$ but low-cost

Directed Graphs: Terminology

◆ Directed Acyclic Graph or DAG

- A directed graph with no cycles
- Permits a topological sorting
 - ◆ e.g., a sorting of the nodes from beginning to end
 - ◆ A topological sorting can be done using a modified DFS traversal

Graphs: Shortest Path

- ◆ Given a weighted graph and two vertices u and v , we want to find a path of minimum total weight between u and v
 - BFS gives us shortest paths for a uniformly weighted graph – this is the concept generalized to weighted graphs
- ◆ Note: related to Traveling Salesman problem which is finding the shortest path that visits all vertices (an NP-complete problem)

Graphs: Shortest Path

◆ Algorithms for finding shortest paths from a start vertex

- Dijkstra's
 - ◆ Naively grow a “cloud of connected vertices”
 - ◆ Assumes non-negative weights
 - ◆ $O(m \log n)$
- Bellman-Ford's
 - ◆ Extends Dijkstra's by carrying along the total weight so far during the expansion
 - ◆ Supports negative weights
 - ◆ $O(nm)$
- DAG-based
 - ◆ Assumes a DAG
 - ◆ Uses topological sorting
 - ◆ $O(n+m)$

Graphs: Shortest Path

◆ All shortest path pairs

- Dijkstra's
 - ◆ $O(nm \log n)$
- Bellman Ford's
 - ◆ $O(n^2m)$
- Modified Floyd-Warshall
 - ◆ $O(n^3)$

Graphs: Minimum Spanning Tree

- ◆ A spanning tree of a weighted graph with minimum total edge weight
 - e.g. the lowest cost network uniting all clients

Graphs: Minimum Spanning Tree

◆ Algorithms

■ Prim-Jarnik's

- ◆ Similar to Dijkstra's: grows a cloud of connected vertices
- ◆ $O(m \log n)$

■ Kruskal's

- ◆ Maintains a forest of growing clouds of vertices
- ◆ $O((n+m) \log n)$

■ Baruvka's

- ◆ Similar to Kruskal's but at each iteration halves the number of connected components
- ◆ $O(m \log n)$

Fun Fact: Fastest MST Algorithm

◆ $O(m\alpha)$

- α is function of (m, n) but in practice is ≤ 4

◆ Based on B. Chazelle's "Soft Heap"

- Amortized cost of operations is $O(1)$ except insert, which is $O(\log 1/e)$, for $e \in [0, 1/2]$
- At expense of $e \cdot n$ of keys being "corrupted", faster heap is obtained