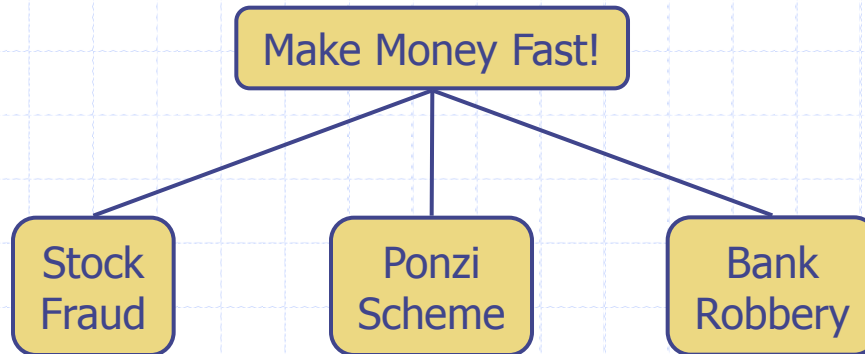


# Trees



# Outline and Reading

- ◆ Tree ADT
- ◆ Preorder and postorder traversals
- ◆ BinaryTree ADT
- ◆ Inorder traversal
- ◆ Euler Tour traversal
- ◆ Template method pattern
- ◆ Data structures for trees
- ◆ C++ implementation

# What is a Tree?

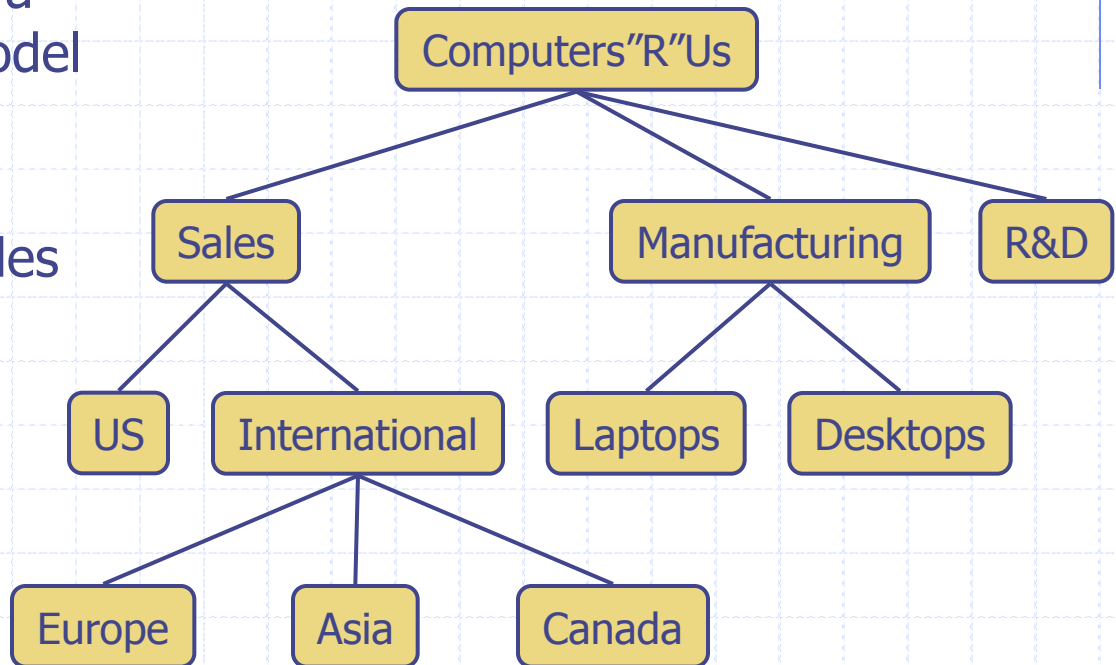


# What is a Tree?



# What is a Tree?

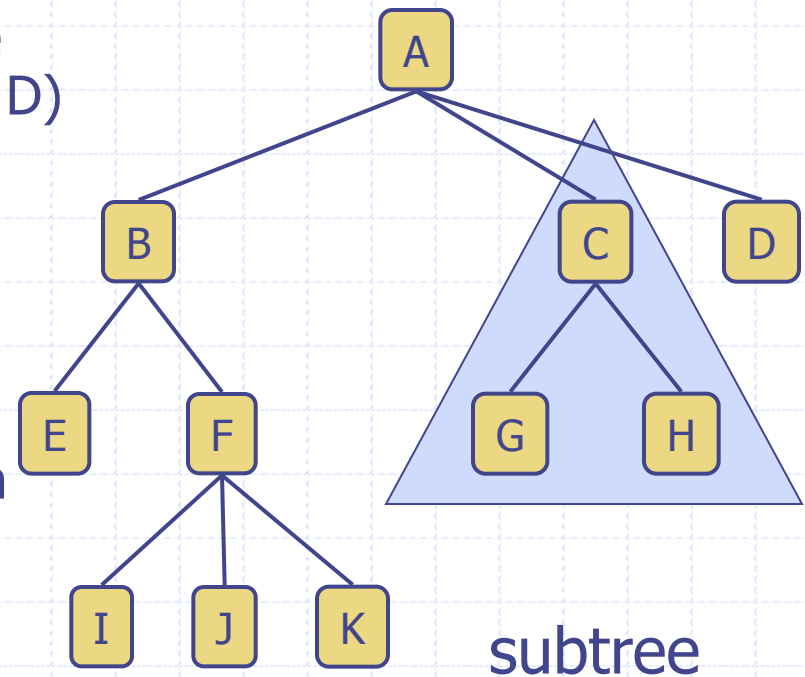
- ◆ In computer science, a tree is an abstract model of a hierarchical structure
- ◆ A tree consists of nodes with a parent-child relation
- ◆ Applications:
  - Organization charts
  - File systems
  - Programming environments



# Tree Terminology

- ◆ Root: node without parent (A)
- ◆ Internal node: node with at least one child (A, B, C, F)
- ◆ External node (a.k.a. leaf ): node without children (E, I, J, K, G, H, D)
- ◆ Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- ◆ Depth of a node: number of ancestors
- ◆ Height of a tree: maximum depth of any node (3)
- ◆ Descendant of a node: child, grandchild, grand-grandchild, etc.

- ◆ Subtree: tree consisting of a node and its descendants



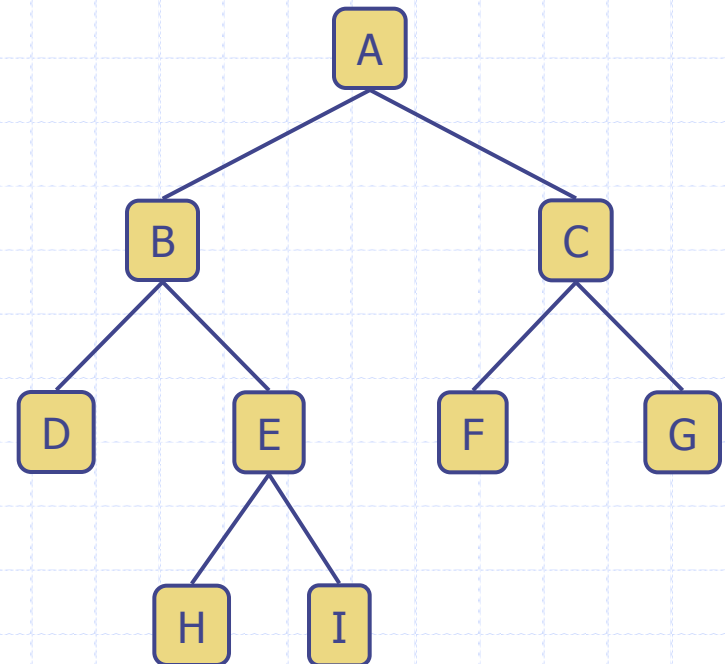
# Tree ADT (not a list/sequence!)

- ◆ We use positions to abstract nodes
- ◆ Generic methods:
  - integer **size()**
  - boolean **isEmpty()**
  - objectIterator **elements()**
  - positionIterator **positions()**
- ◆ Accessor methods:
  - position **root()**
  - position **parent(p)**
  - positionIterator **children(p)**
- ◆ Query methods:
  - boolean **isInternal(p)**
  - boolean **isExternal(p)**
  - boolean **isRoot(p)**
- ◆ Update methods:
  - **swapElements(p, q)**
  - object **replaceElement(p, o)**
- ◆ Additional update methods may be defined by data structures implementing the Tree ADT

# Binary Tree

- ◆ A binary tree is a tree with the following properties:
  - Each internal node has two children
  - The children of a node are an ordered pair
- ◆ We call the children of an internal node left child and right child
- ◆ Alternative recursive definition: a binary tree is either
  - a tree consisting of a single node, or
  - a tree whose root has an ordered pair of children, each of which is a binary tree

- ◆ Applications:
  - arithmetic expressions
  - decision processes
  - searching

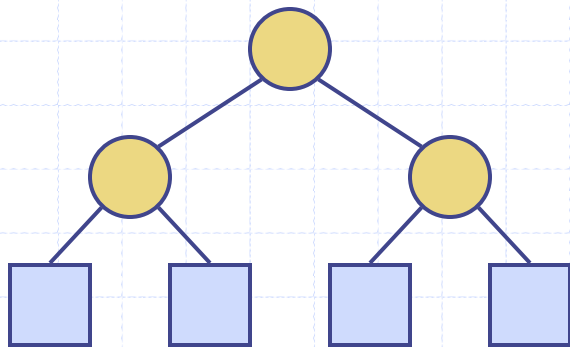




# Properties of Binary Trees

## ◆ Notation

- $n$  number of nodes
- $e$  number of external nodes
- $i$  number of internal nodes
- $h$  height



## ◆ Properties:

- $e =$

- ◆  $i + 1$

- ◆  $e \leq 2^h$

- $n =$

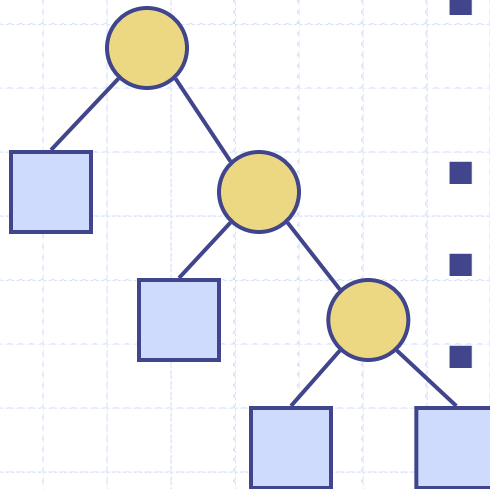
- ◆  $2e - 1$

- $h \leq i$

- $h \leq (n - 1)/2$

- $h \geq \log_2 e$

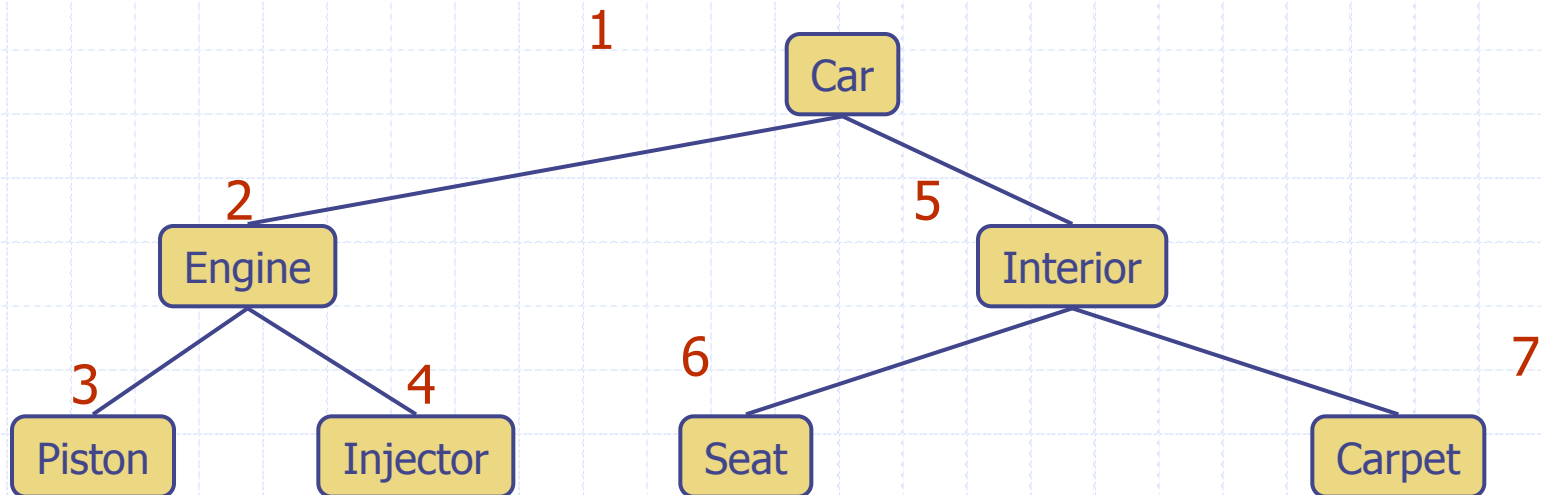
- $h \geq \log_2 (n + 1) - 1$



# Preorder Traversal

- ◆ A traversal visits the nodes of a tree in a systematic manner
- ◆ In a preorder traversal, a node is visited before its descendants
- ◆ Application: print a *structured document*

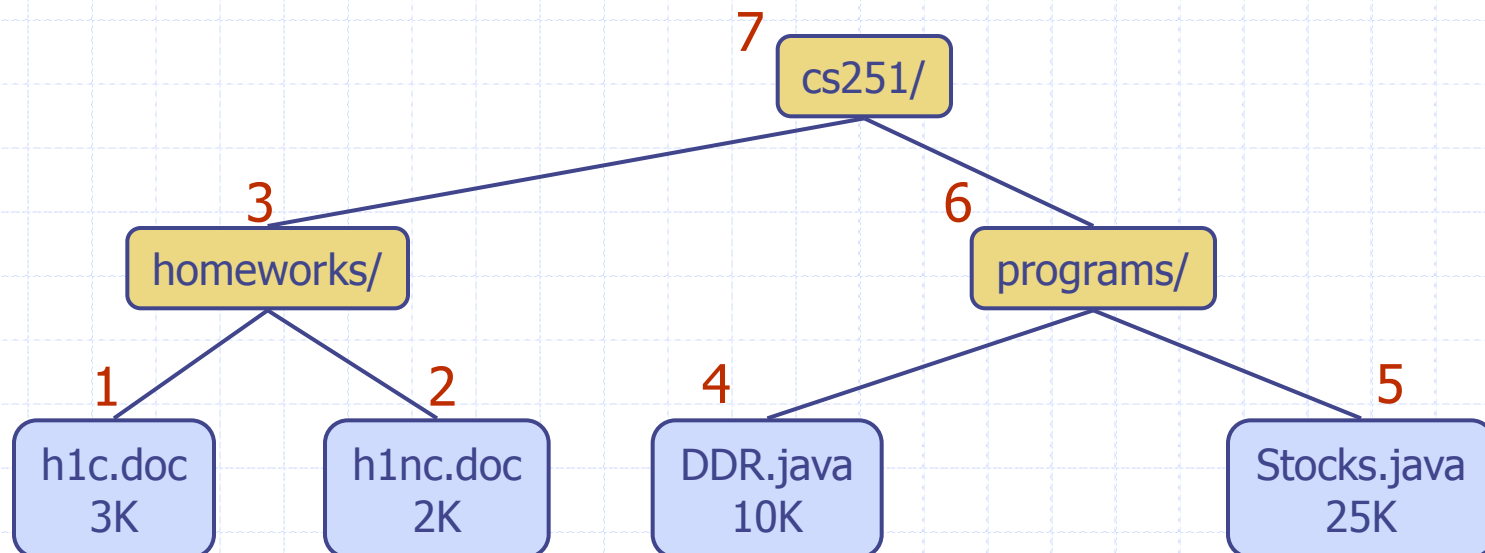
```
Algorithm preOrder(v)  
  visit(v)  
  for each child w of v  
    preorder (w)
```



# Postorder Traversal

- ◆ In a postorder traversal, a node is visited after its descendants
- ◆ Application: compute space used by files in a directory and its subdirectories

**Algorithm** *postOrder*( $v$ )  
  **for each** child  $w$  of  $v$   
    *postOrder* ( $w$ )  
  *visit*( $v$ )



# Inorder Traversal

- ◆ In an inorder traversal a node is visited after its left subtree and before its right subtree
- ◆ Application: draw a binary tree
  - $x(v)$  = inorder rank of  $v$
  - $y(v)$  = depth of  $v$

**Algorithm** *inOrder*( $v$ )

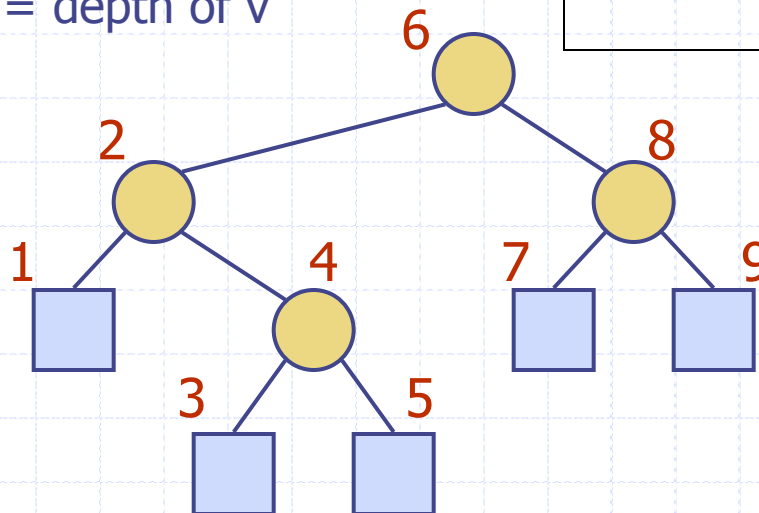
**if** *isInternal* ( $v$ )

*inOrder* (*leftChild* ( $v$ ))

*visit*( $v$ )

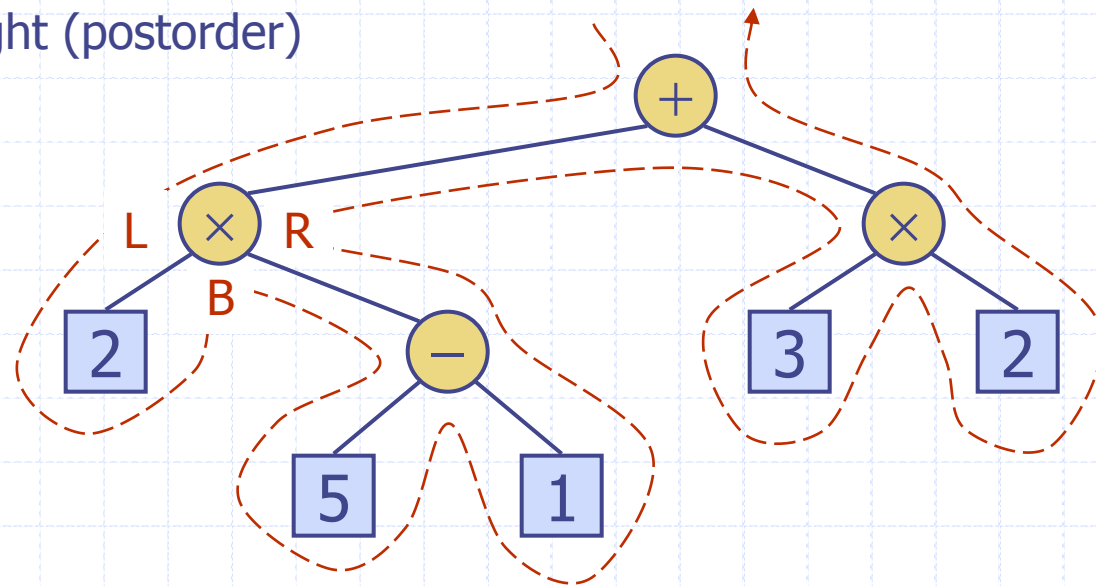
**if** *isInternal* ( $v$ )

*inOrder* (*rightChild* ( $v$ ))



# Euler Tour Traversal

- ◆ Generic traversal of a binary tree
- ◆ Includes a special cases the preorder, postorder and inorder traversals
- ◆ Walk around the tree and visit each node three times:
  - on the left (preorder)
  - from below (inorder)
  - on the right (postorder)



# Example

## ◆ 1. Build a Binary Tree for the following:

### City

- Residential Area
  - ◆ House1, House2
    - Each has Kitchen, Bedroom 1, Bedroom 2
- Commercial Area
  - ◆ Store 1, Store 2, Store 3
    - Each has Lobby, Storage
    - Each Lobby has a Cashier and Aisles

# Example



# Example

## ◆ 2. What traversal would print out the “structure”?)

e.g.: City

- Residential Area
  - ◆ House1
    - Kitchen
    - Bedroom 1, Bedroom 2
  - ◆ House2
    - Kitchen
    - Bedroom 1, Bedroom 2
- Commercial Area
  - ◆ Store 1
    - Lobby
      - Cashier, Aisles
  - ◆ Store 2
    - Lobby
      - Cashier, Aisles
  - ◆ Store 3
    - Lobby
      - Cashier, Aisles



# Example

## ◆ 3. What happens to the binary tree if: City

- Residential Area

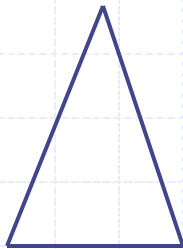
- ◆ House1, House2, House 3, House 4, House 5
  - Each has Kitchen, Bedroom 1, Bedroom 2

- Commercial Area

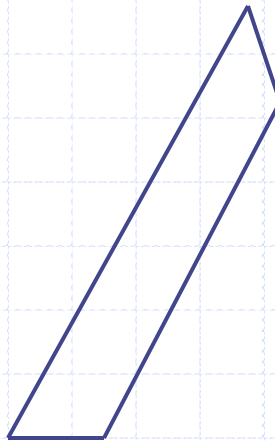
- ◆ Store 1, Store 2, Store 3, Store 4, Store 5, Store 6, Store 7
  - Each has Lobby, Storage
  - Each Lobby has a Cashier and Aisles

# Example

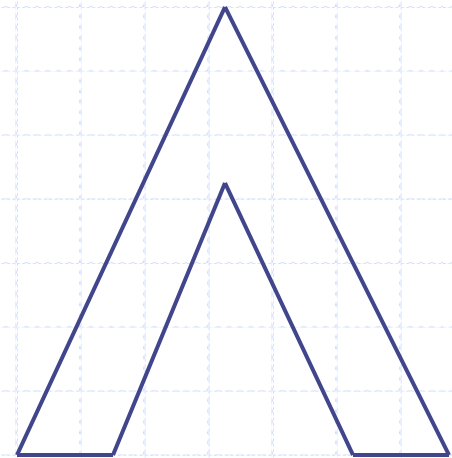
- ◆ An “unbalanced” tree is produced
- ◆ Conceptually:



balanced



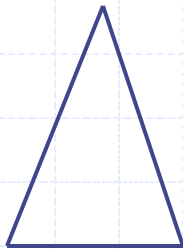
unbalanced



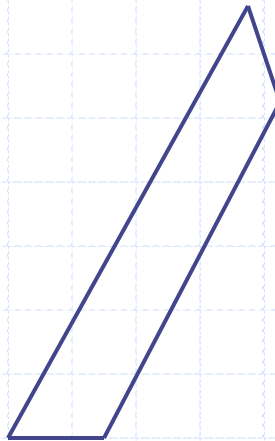
??

# Example

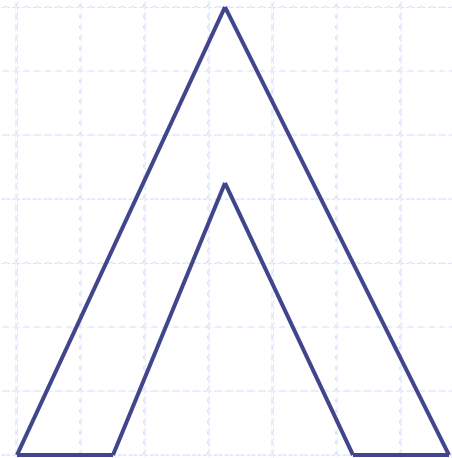
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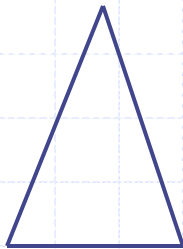
unbalanced



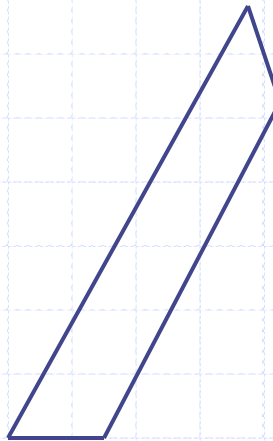
unbalanced

# Example

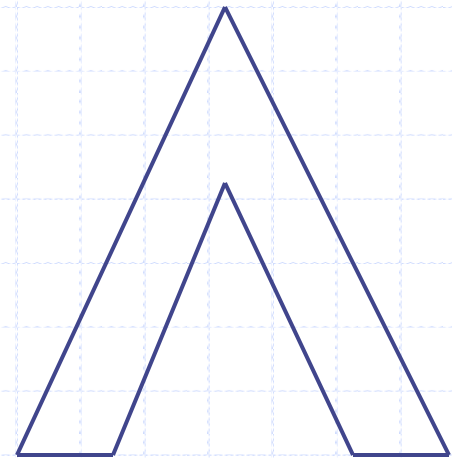
◆ What is the height of the tree of  $n$  nodes?



balanced



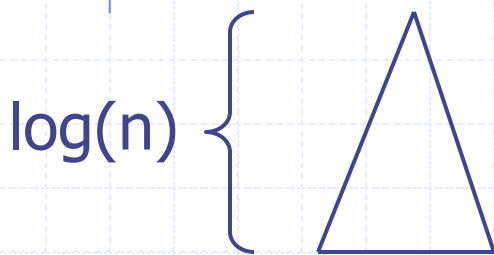
unbalanced



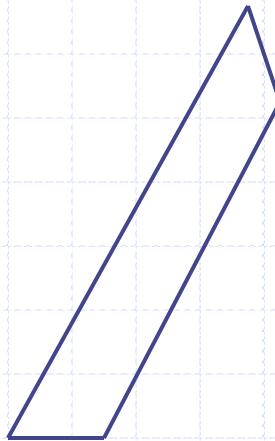
unbalanced

# Example

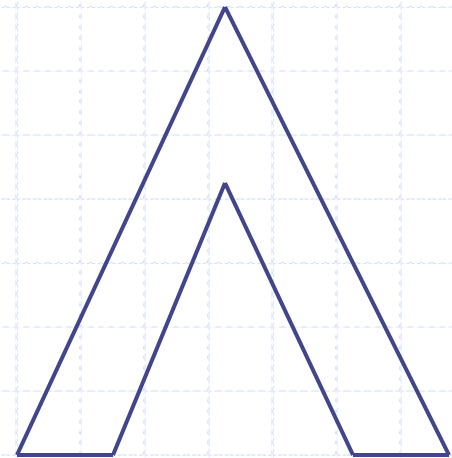
◆ What is the height of the tree of  $n$  nodes?



balanced



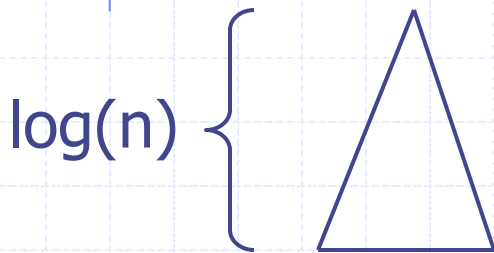
unbalanced



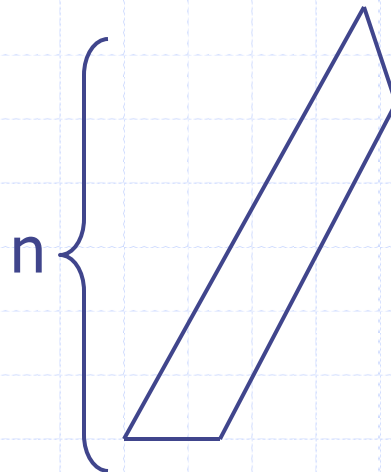
unbalanced

# Example

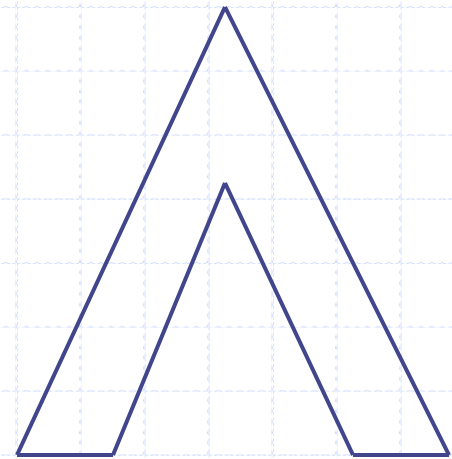
◆ What is the height of the tree of  $n$  nodes?



balanced



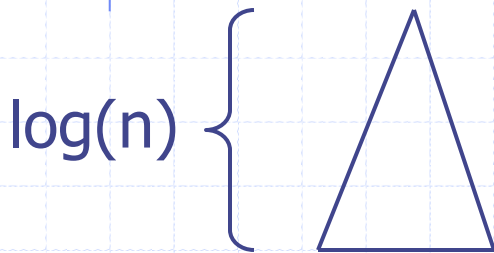
unbalanced



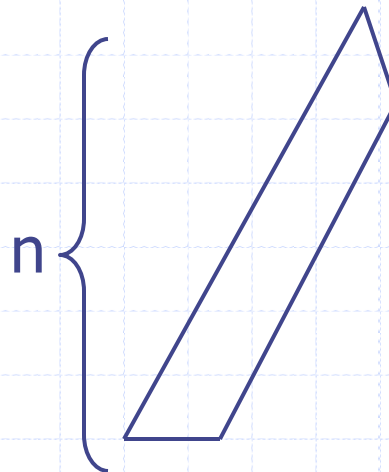
unbalanced

# Example

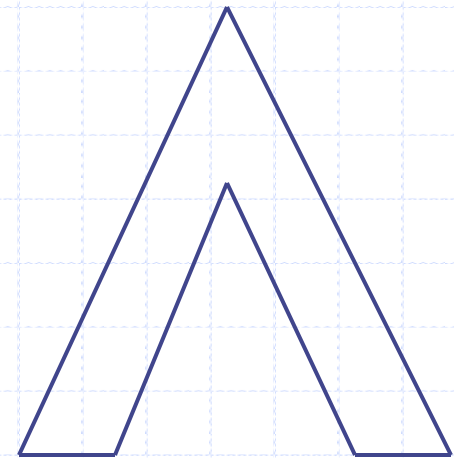
◆ Why do we care?



balanced



unbalanced

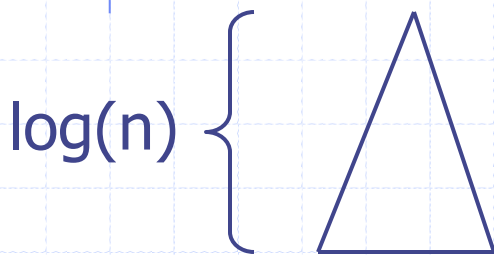


unbalanced

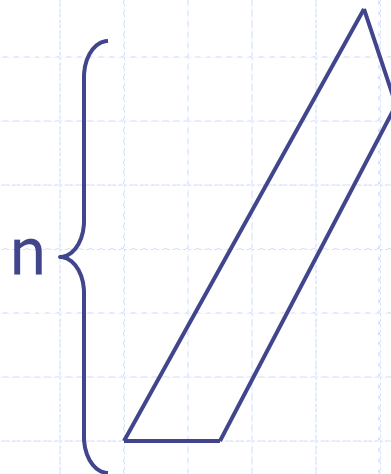
# Example

## ◆ Efficiency!

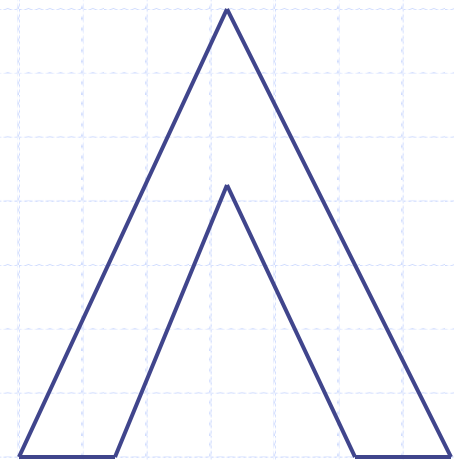
- $O(\log n)$  time to reach any node in a balanced tree and  $O(n)$  time to reach any node in an unbalanced tree



balanced



unbalanced

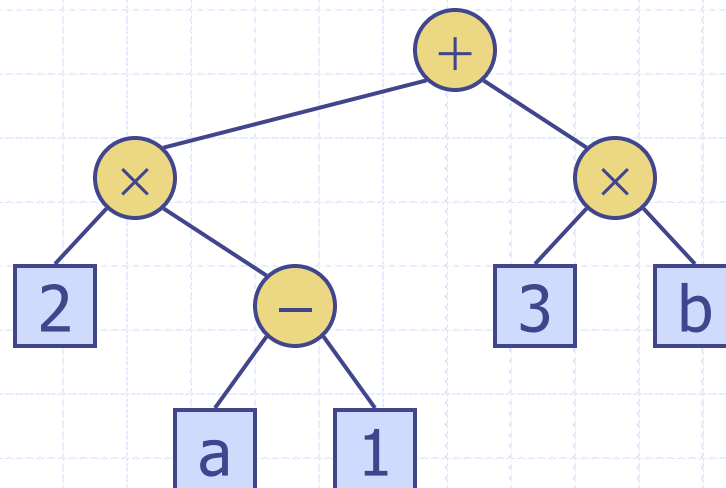


unbalanced



# Arithmetic Expression Trees

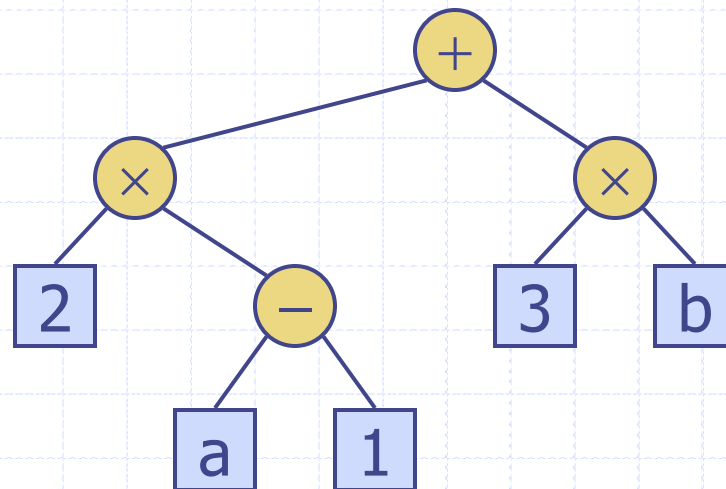
- ◆ Binary tree associated with an arithmetic expression
  - internal nodes: operators
  - external nodes: operands
- ◆ Example: arithmetic expression tree for the expression  $(2 \times (a - 1) + (3 \times b))$
- ◆ What traversal order is used here?



# Arithmetic Expression Trees

## ◆ Trick question!:

- "evaluation order"
- "print order"
- "notation/construction order"



# Evaluate Arithmetic Expressions

- ◆ Specialization of a postorder traversal
  - recursive method returning the value of a subtree
  - when visiting an internal node, combine the values of the subtrees

**Algorithm** *evalExpr(v)*

**if** *isExternal* (*v*)

**return** *v.element* ()

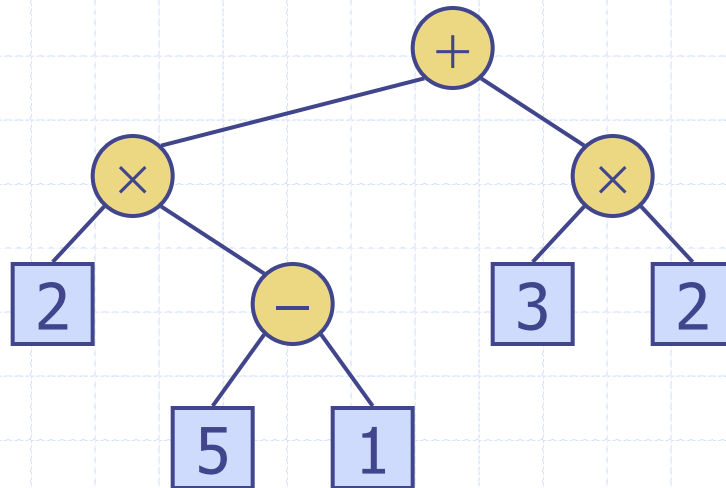
**else**

*x*  $\leftarrow$  *evalExpr*(*leftChild* (*v*))

*y*  $\leftarrow$  *evalExpr*(*rightChild* (*v*))

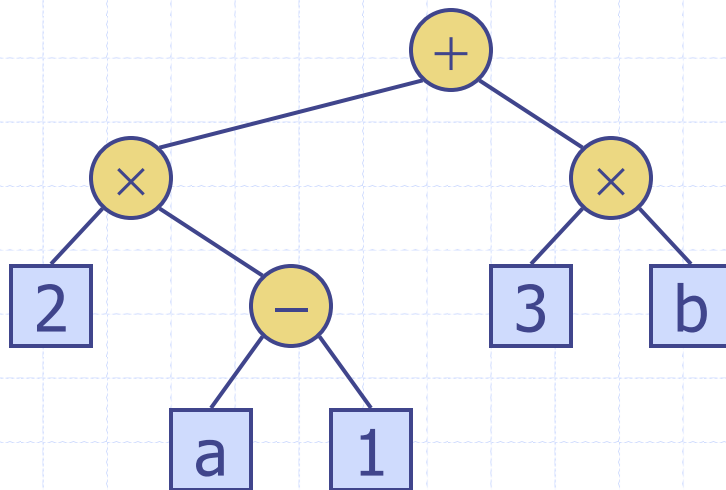
$\diamond$   $\leftarrow$  operator stored at *v*

**return** *x*  $\diamond$  *y*



# Print Arithmetic Expressions

- ◆ Specialization of an inorder traversal
  - print operand or operator when visiting node
  - print "(" before traversing left subtree
  - print ")" after traversing right subtree



```
Algorithm printExpression(v)  
  if isInternal (v)  
    print("(")  
    inOrder (leftChild (v))  
    print(v.element ())  
    if isInternal (v)  
      inOrder (rightChild (v))  
    print (")")
```

$((2 \times (a - 1)) + (3 \times b))$

# Notation/Construction Order

- ◆ Prefix notation

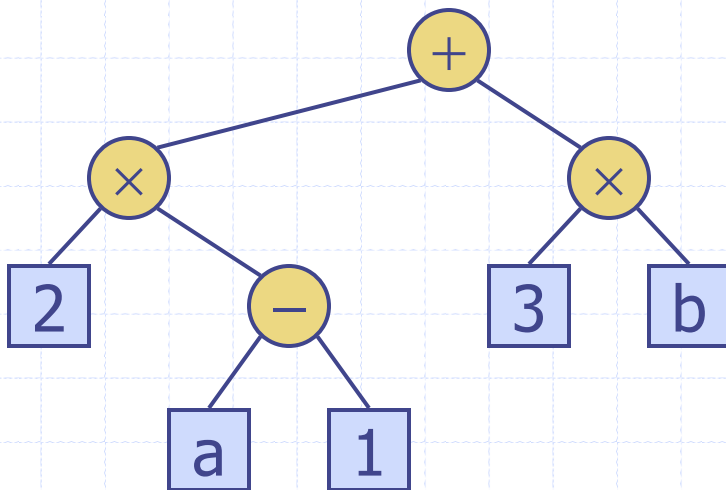
- $+x2-a1x3b$

- ◆ Inorder notation

- $2x(a-1)+(3xb)$

- ◆ Postfix notation

- $a1-2x3bx+$



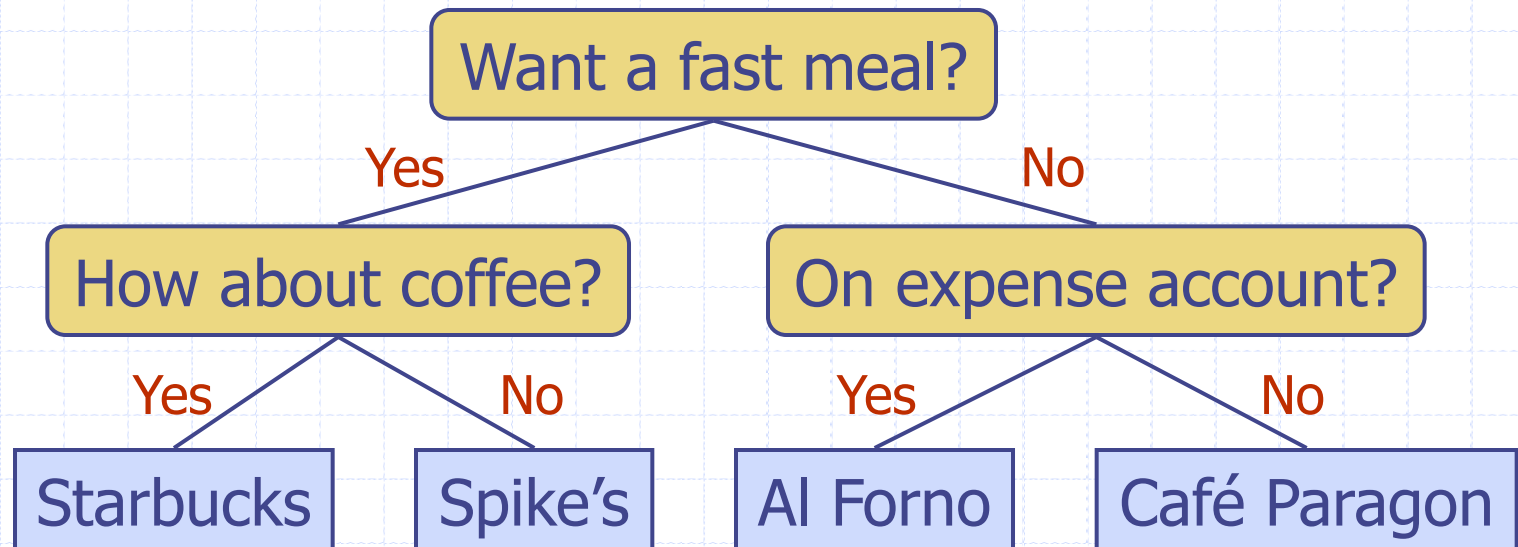
$$((2 \times (a - 1)) + (3 \times b))$$

# Example



# Decision Tree

- ◆ Binary tree associated with a decision process
  - internal nodes: questions with yes/no answer
  - external nodes: decisions
- ◆ Example: dining decision



# BinaryTree ADT

◆ The BinaryTree ADT extends the Tree ADT, i.e., it inherits all the methods of the Tree ADT

◆ Additional methods:

- position **leftChild**(p)
- position **rightChild**(p)
- position **sibling**(p)

◆ Update methods may be defined by data structures implementing the BinaryTree ADT



# Template Method Pattern

- ◆ Generic algorithm that can be specialized by redefining certain steps
- ◆ Implemented by means of an abstract C++ class
- ◆ Visit methods that can be redefined by subclasses
- ◆ Template method `eulerTour`
  - Recursively called on the left and right children
  - A `Result` object with fields `leftResult`, `rightResult` and `finalResult` keeps track of the output of the recursive calls to `eulerTour`

```
class EulerTour {
protected:
    BinaryTree* tree;
    virtual void visitExternal(Position p, Result r) {}
    virtual void visitLeft(Position p, Result r) {}
    virtual void visitBelow(Position p, Result r) {}
    virtual void visitRight(Position p, Result r) {}
    int eulerTour(Position p) {
        Result r = initResult();
        if (tree->isExternal(p)) { visitExternal(p, r); }
        else {
            visitLeft(p, r);
            r.leftResult = eulerTour(tree->leftChild(p));
            visitBelow(p, r);
            r.rightResult = eulerTour(tree->rightChild(p));
            visitRight(p, r);
            return r.finalResult;
        } // ... other details omitted
    }
};
```

# Specializations of EulerTour

◆ We show how to specialize class EulerTour to evaluate an arithmetic expression

◆ Assumptions

- External nodes support a function **value()**, which returns the value of this node.
- Internal nodes provide a function **operation(int, int)**, which returns the result of some binary operator on integers.

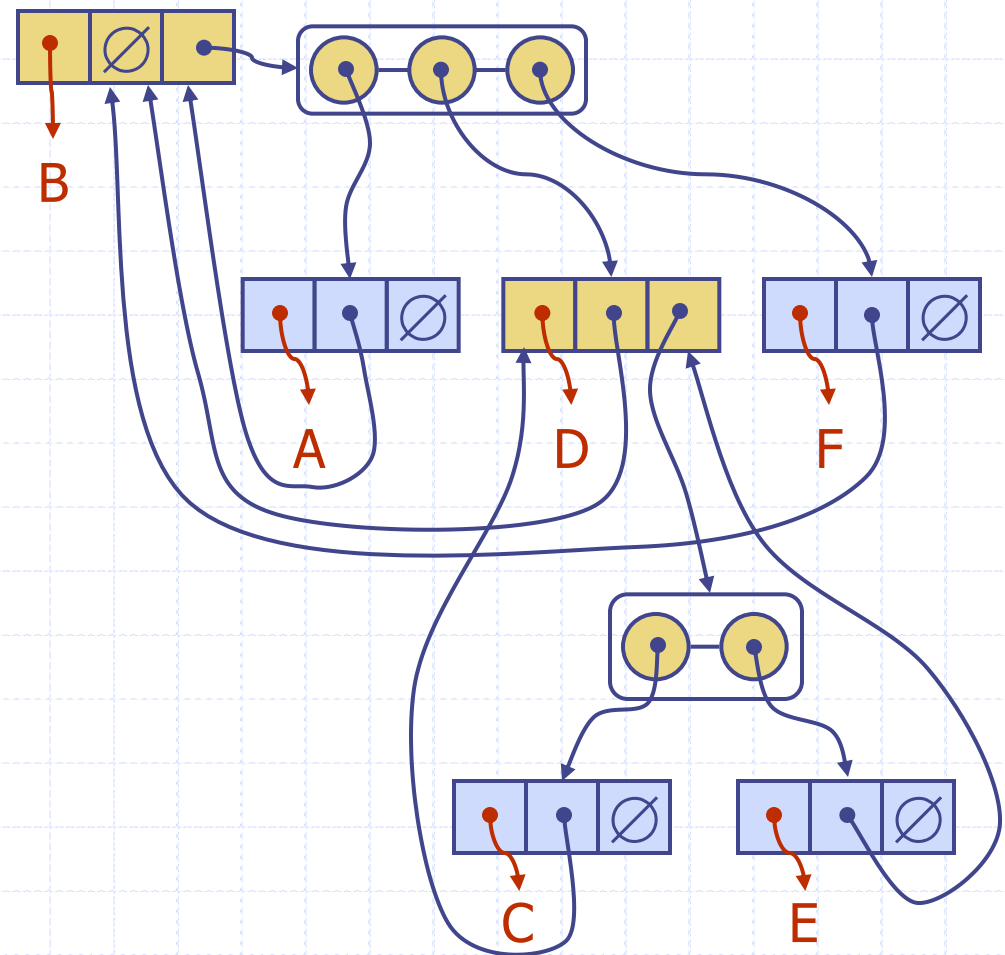
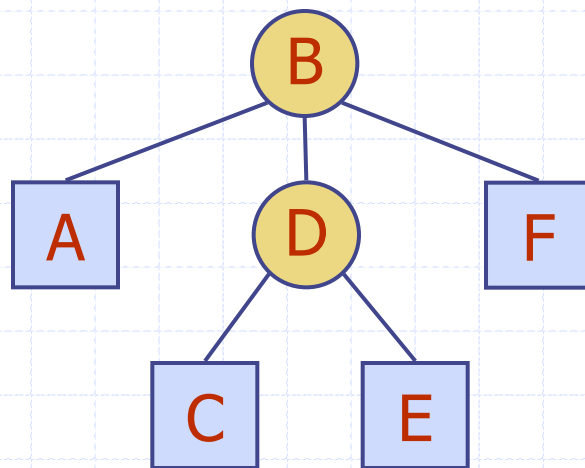
```
class EvaluateExpression
    : public EulerTour {
protected:
    void visitExternal(Position p, Result r) {
        r.finalResult = p.element().value();
    }

    void visitRight(Position p, Result r) {
        Operator op = p.element().operator();
        r.finalResult = p.element().operation(
            r.leftResult, r.rightResult);
    }

    // ... other details omitted
};
```

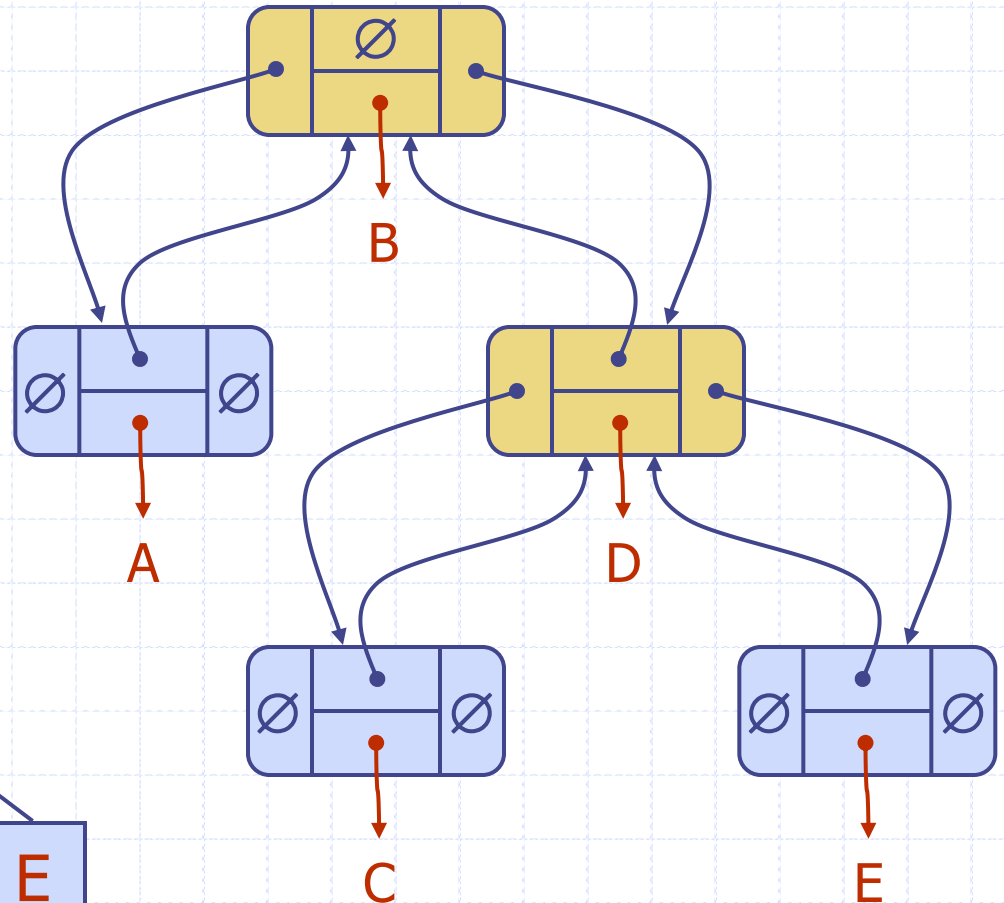
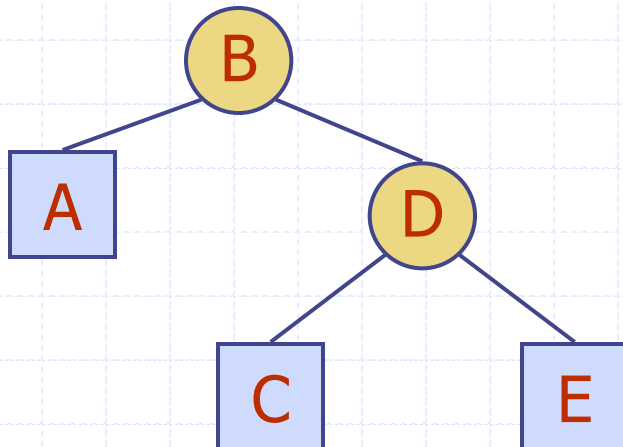
# Data Structure for Trees

- ◆ A node is represented by an object storing
  - Element
  - Parent node
  - Sequence of children nodes
- ◆ Node objects implement the Position ADT



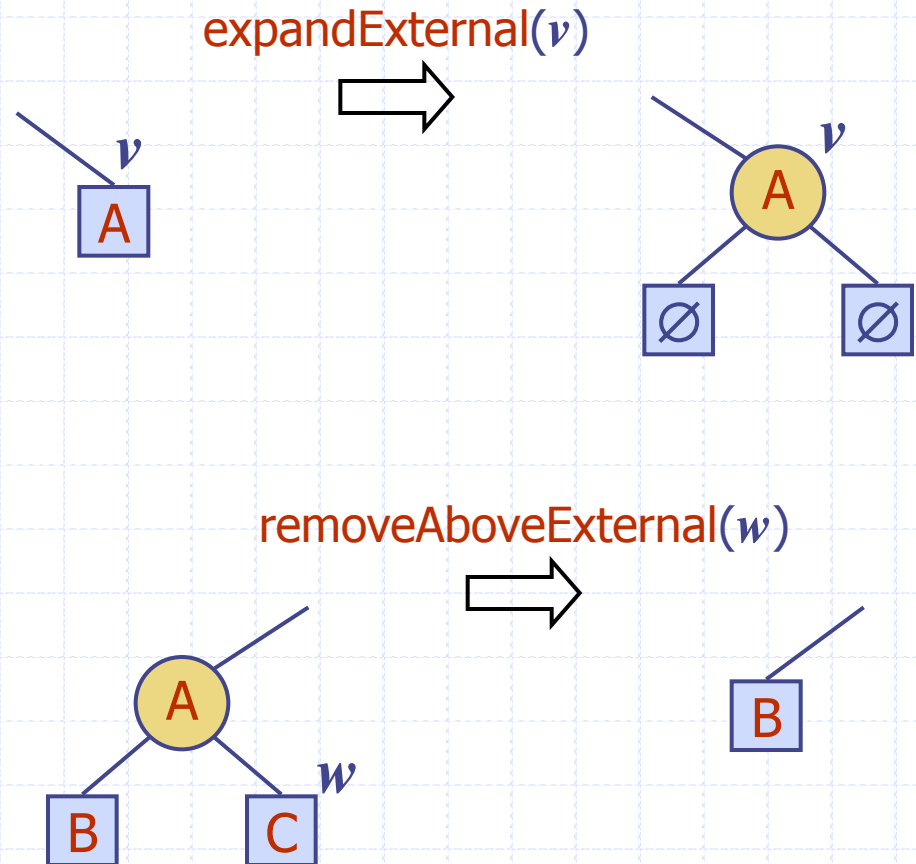
# Data Structure for Binary Trees

- ◆ A node is represented by an object storing
  - Element
  - Parent node
  - Left child node
  - Right child node
- ◆ Node objects implement the **Position** ADT



# C++ Implementation

- ◆ Tree interface
- ◆ BinaryTree interface extending Tree
- ◆ Classes implementing Tree and BinaryTree and providing
  - Constructors
  - Update methods
  - Print methods
- ◆ Examples of updates for binary trees
  - `expandExternal(v)`
  - `removeAboveExternal(w)`



# Let the Hunger Games begin!