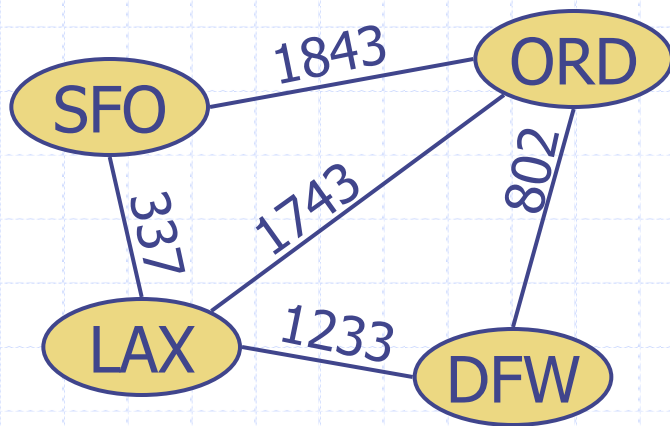


# Graphs: Recap



# Graphs: Basics

- ◆ Directed vs. undirected graph
  - e.g. edges have a direction associated with them
- ◆ (Non-uniformly) Weighted graph
  - e.g. edges have a weight associated with them
- ◆ Properties
  - $\sum_n \deg(v_n) = 2m$
  - $m \leq n(n-1)/2$
- ◆ Representation
  - Edge list structure,
  - Adjacency list structure, or
  - Adjacency matrix structure

# Graphs: Traversals

## ◆ Depth-first Search

- Traverse deeply first
- $O(n+m)$

## ◆ Breadth-first Search

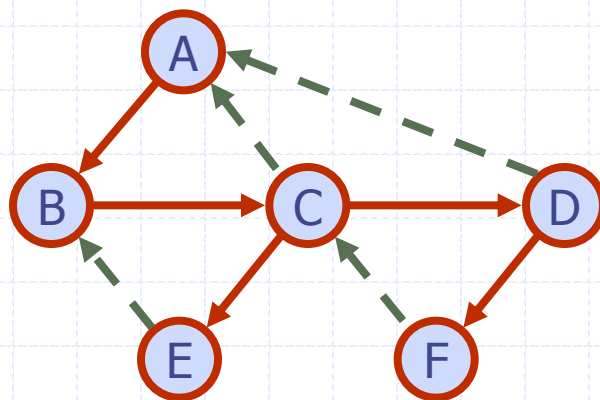
- Traverse broadly first (“breadth”)
- $O(n+m)$

# Graphs: Terminology

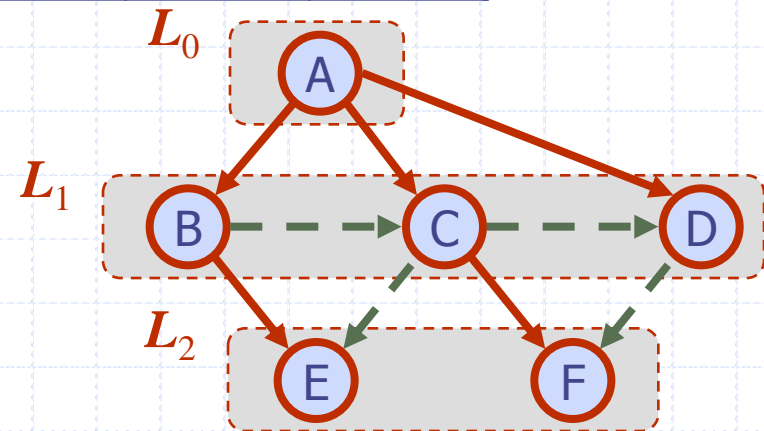
- ◆ Path
- ◆ Connected
- ◆ Subgraph
- ◆ Spanning
- ◆ Biconnected
  - e.g., separation edges or vertices
- ◆ What are these and how do you find them?
  - Connected component
  - Spanning Subgraph
  - Maximally-connected Subgraph
  - Spanning Tree
  - Spanning Forest
  - Biconnected Components

# Graphs: DFS vs. BFS

Applications	DFS	BFS
Spanning forest, connected components, paths, cycles	✓	✓
Shortest paths (for uniformly weighted graphs)		✓
Biconnected components	✓	



DFS



BFS

# Directed Graphs: Terminology

## ◆ Reachability

- Is a vertex  $u$  “reachable” from  $v$ 
  - ◆ e.g., can you get to MIA from HNL?

## ◆ Strongly-Connected Components

- Can you get to any city from any city

## ◆ Transitive Closure

- If I can get to MIA from HNL and to JFK from MIA, then I can get to JFK from HNL
- The transitive closure graph of a graph  $G$  extracts this information
- Algorithms:
  - ◆ Naïve:  $O(n(n+m))$  to  $O(n^3)$
  - ◆ Floyd-Warshall Algorithm:  $O(n^3)$  but low-cost

# Directed Graphs: Terminology

## ◆ Directed Acyclic Graph or DAG

- A directed graph with no cycles
- Permits a topological sorting
  - ◆ e.g., a sorting of the nodes from beginning to end
  - ◆ A topological sorting can be done using a modified DFS traversal

# Graphs: Shortest Path

- ◆ Given a weighted graph and two vertices  $u$  and  $v$ , we want to find a path of minimum total weight between  $u$  and  $v$ 
  - BFS gives us shortest paths for a uniformly weighted graph – this is the concept generalized to weighted graphs
- ◆ Note: related to Traveling Salesman problem which is finding the shortest path that visits all vertices (an NP-complete problem)



# Graphs: Shortest Path

- ◆ Algorithms for finding shortest paths from a start vertex
  - Dijkstra's
    - ◆ Naively grow a “cloud of connected vertices”
    - ◆ Assumes non-negative weights
    - ◆  $O(m \log n)$
  - Bellman-Ford's
    - ◆ Extends Dijkstra's by carrying along the total weight so far during the expansion
    - ◆ Supports negative weights
    - ◆  $O(nm)$
  - DAG-based
    - ◆ Assumes a DAG
    - ◆ Uses topological sorting
    - ◆  $O(n+m)$

# Graphs: Shortest Path

## ◆ All shortest path pairs

- Dijkstra's
  - ◆  $O(nm \log n)$
- Bellman Ford's
  - ◆  $O(n^2m)$
- Modified Floyd-Warshall
  - ◆  $O(n^3)$

# Graphs: Minimum Spanning Tree

- ◆ A spanning tree of a weighted graph with minimum total edge weight
  - e.g. the lowest cost network uniting all clients

# Graphs: Minimum Spanning Tree

## ◆ Algorithms

### ■ Prim-Jarnik's

- ◆ Similar to Dijkstra's: grows a cloud of connected vertices
- ◆  $O(m \log n)$

### ■ Kruskal's

- ◆ Maintains a forest of growing clouds of vertices
- ◆  $O((n+m) \log n)$

### ■ Baruvka's

- ◆ Similar to Kruskal's but at each iteration halves the number of connected components
- ◆  $O(m \log n)$

# Fun Fact: Fastest MST Algorithm

◆  $O(m\alpha)$

- $\alpha$  is function of  $(m, n)$  but in practice is  $\leq 4$

◆ Based on B. Chazelle's "Soft Heap"

- Amortized cost of operations is  $O(1)$  except insert, which is  $O(\log 1/e)$ , for  $e \in [0, 1/2]$
- At expense of  $e \cdot n$  of keys being "corrupted", faster heap is obtained