

BELLMAN FORD CODE:

```
class Graph
```

```
{
    map<int, map<int,int> > strg ;
    int V ;
    vector<vector<int> > edgeList ;

    public :

    Graph(int V)
    {
        this->V = V ;
    }

    void addEdge(int u, int v, int cost)
    {
        strg[u][v] = cost ;
        edgeList.push_back({u,v,cost}) ;
    }

    void display()
    {
        for(int i = 0 ; i < V ; i++)
        {
            cout << i << "\t" ;

            map<int, int>::iterator itr ;

            for(itr = strg[i].begin() ; itr != strg[i].end() ; itr++)
                cout << itr->first << "@" << itr->second << ", " ;
            cout << endl ;
        }
    }
}
```

$[0,1,2]$ $[0,2,20]$ $[1,2,10]$ $[2,3,5]$

```
void bellmanFord(int src)
```

```
{
    int cost[V] ;
    fill(cost, cost+V, 100000) ;
    cost[src] = 0 ;
```

$\rightarrow O(V)$

```
// V-1 times, relax every edge
```

```
for(int i = 1 ; i <= V ; i++)
```

$O(V)$

```
{
    for(auto edge : edgeList)
```

$[0,1,10]$ $O(E)$

```
{
    int u = edge[0] ;
    int v = edge[1] ;
    int c = edge[2] ;
```

```
// cost Relax
```

```
int oc = cost[v] ;
int nc = cost[u] + c ;
if(nc < oc)
```

```
{
    if(i <= V-1)
        cost[v] = nc ;
```

```
else
```

```
{
    cout << "-ve wt cycle present" ;
    return ;
}
```

```
}
```

```
}
```

```
for(int i = 0 ; i < V ; i++)
```

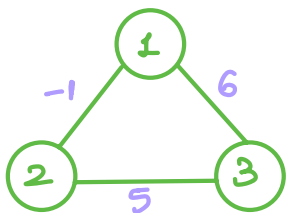
```
    cout << i << " -> " << cost[i] << endl ;
```

```
}
```

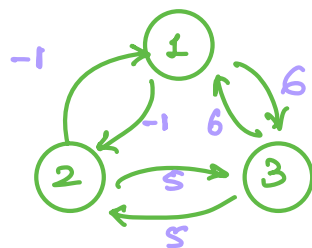
```
};
```

$V + VE = O(VE)$

$O(1)$



Undirected graph
with -ve edge



\Rightarrow Directed graph
with -ve wt cycle

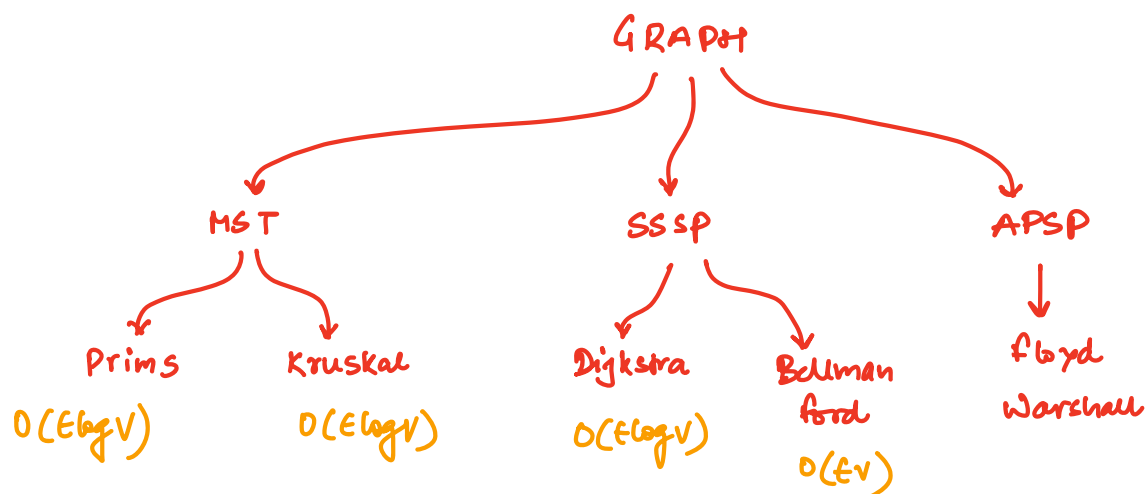
no SSSP algo that works with
-ve wt cycle.

DJKSTRA

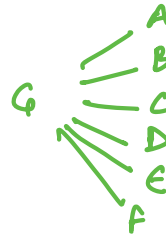
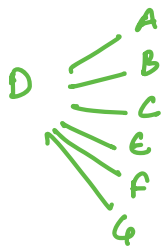
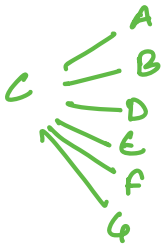
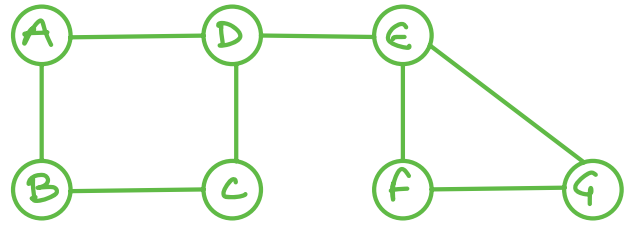
- undirected graph:
Works with all +ve wt edges
- directed graph:
Works with all +ve wt edges
- Time Complexity:
 $O(E \log V)$

BELLMAN FORD

- undirected graph:
Works with all +ve wt edges
- directed graph:
Works with +ve wt cycle and
-ve wt edge
Doesn't work with -ve wt
cycle.
- Time Complexity:
 $O(VE)$



All Pair Shortest Path



One option:

Dijkstra

Bellman ford

SSSP:

$E \log V$

$E \cdot V$

APSP

$V \cdot E \log V$

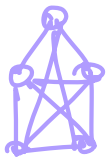
$V \cdot E \cdot V = V^2 \cdot E$

Complete Graphs $E = V^2$

\downarrow
 $V^3 \log V$

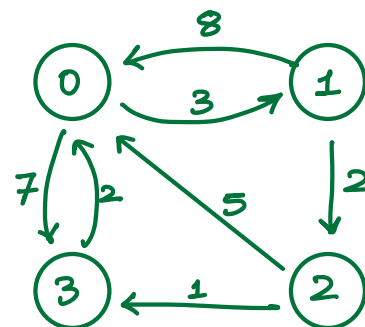
\downarrow
 V^4

Floyd
Warshall
 V^3



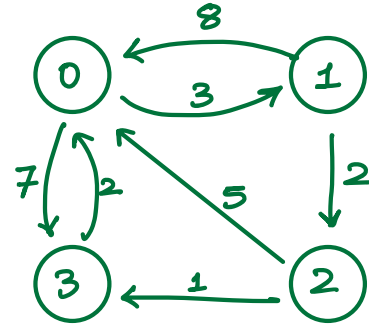
$$V C_2 = \frac{V!}{(V-2)! 2!} = \frac{(V-2)! (V-1)(V)}{(V-2)! 2!} = \frac{V(V-1)}{2} = V^2$$

	0	1	2	3
0	0	3	∞	7
1	8	0	2	∞
2	5	∞	0	1
3	2	∞	∞	0



Vertex 0

	0	1	2	3
0	0	3	∞	7
1	8	0	2	∞ 15
2	5	8 ∞	0	1
3	2	∞ 5	∞	0



$(1,1):0$

$$\underbrace{1 \rightarrow 0}_{8} \rightarrow \underbrace{0 \rightarrow 1}_{3} = 11 \quad \times$$

$(1,2):2$

$$\underbrace{1 \rightarrow 0}_{8} \rightarrow \underbrace{0 \rightarrow 2}_{\infty} \quad \times$$

$(1,3): \infty$

$$\underbrace{1 \rightarrow 0}_{8} \rightarrow \underbrace{0 \rightarrow 3}_{7} = 15 \quad \checkmark$$

$(2,1): \infty$

$$\underbrace{2 \rightarrow 0}_{5} \rightarrow \underbrace{0 \rightarrow 1}_{3} = 8 \quad \checkmark$$

$(2,3):1$

$$\underbrace{2 \rightarrow 0}_{5} \rightarrow \underbrace{0 \rightarrow 3}_{7} = 12 \quad \times$$

$(3,1): \infty$

$$\underbrace{3 \rightarrow 0}_{2} \rightarrow \underbrace{0 \rightarrow 1}_{3} = 5 \quad \checkmark$$

$(3,2): \infty$

$$\underbrace{3 \rightarrow 0}_{2} \rightarrow \underbrace{0 \rightarrow 2}_{\infty} \quad \times$$

Vertex 1

	0	1	2	3
0	0	3	5 ∞	7
1	8	0	2	15
2	5	8	0	1
3	2	5	∞ 7	0

$(0,2): \infty$

$$\underbrace{0 \rightarrow 1}_{3} \rightarrow \underbrace{1 \rightarrow 2}_{2} = 5 \quad \checkmark$$

$(0,3):7$

$$\underbrace{0 \rightarrow 1}_{3} \rightarrow \underbrace{1 \rightarrow 3}_{15} = 18 \quad \times$$

$(2,0):5$

$$\underbrace{2 \rightarrow 1}_{8} \rightarrow \underbrace{1 \rightarrow 0}_{8} = 16 \quad \times$$

$(2,3):1$

$$\underbrace{2 \rightarrow 1}_{8} \rightarrow \underbrace{1 \rightarrow 3}_{15} \quad \times$$

$(3,0):2$

$$\underbrace{3 \rightarrow 1}_{5} \rightarrow \underbrace{1 \rightarrow 0}_{8} \quad \times$$

$(3,2): \infty$

$$\underbrace{3 \rightarrow 1}_{5} \rightarrow \underbrace{1 \rightarrow 2}_{2} = 7 \quad \checkmark$$

Vertex 2

	0	1	2	3
0	0	3	5	7 6
1	8 7	0	2	15 3
2	5	8	0	1
3	2	5	7	0

(0,1):3

0→2→1
5 8 X

(0,3):7

0→2→3
5 1 = 6
✓

(1,0):8

1→2→0
2 5 = 7
✓

(1,3):15

1→2→3
2 1 = 3
✓

(3,0):2

3→2→0
7 5 X

(3,1):5

3→2→1
7 8 X

Vertex 3

	0	1	2	3
0	0	3	5	6
1	7 5	0	2	3
2	5 3	8 6	0	1
3	2	5	7	0

(0,1):3

0→3→1
6 5 X

(0,2):5

0→3→2
6 7 X

(1,0):7

1→3→0
3 2 = 5
✓

(1,2):2

1→3→2
3 7 X

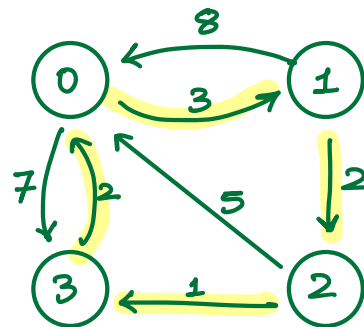
(2,0):5

2→3→0
1 2 = 3
✓

(2,1):8

2→3→1
1 5 = 6
✓

	0	1	2	3
0	0	3	5	6
1	5	0	2	3
2	3	6	0	1
3	2	5	7	0



FLOYD WARSHALL CODE:

```
void floydWarshall()
```

```
{
    int cost[V][V] ;

    for(int i= 0 ; i < V ; i++)
    {
        for(int j = 0 ; j < V ; j++)
        {
            if(i == j)
                cost[i][j] = 0 ;
            else
                cost[i][j] = 100000 ;
        }
    }

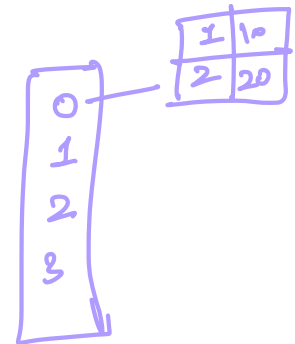
    for(int i = 0 ; i < V ; i++)
    {
        map<int, int>::iterator itr ;
        for(itr = strg[i].begin() ; itr != strg[i].end() ; itr++)
            cost[i][itr->first] = itr->second ;
    }

    for(int k = 0 ; k < V ; k++)
    {
        for(int i= 0 ; i < V ; i++)
        {
            for(int j = 0 ; j < V ; j++)
            {
                int oc = cost[i][j] ;
                int nc = cost[i][k] + cost[k][j] ;

                if(nc < oc)
                    cost[i][j] = nc ;
            }
        }
    }

    for(int i= 0 ; i < V ; i++)
    {
        for(int j = 0 ; j < V ; j++)
        {
            cout << cost[i][j] << " " ;
        }
        cout << endl ;
    }
}
```

∞



edges wt

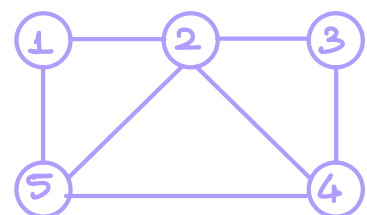
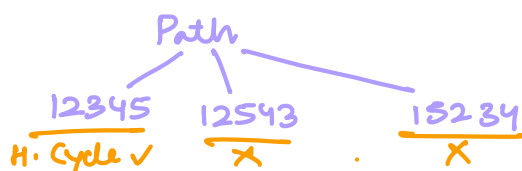
(i, j)
 $i \rightarrow k \rightarrow j$
 $O(V^3)$

TRAVELLING SALESMAN PROBLEM (TSP):

HAMILTONIAN GRAPH:

A **Hamiltonian Path** in an undirected graph is a path that visits every vertex exactly once.

grapher is a path that

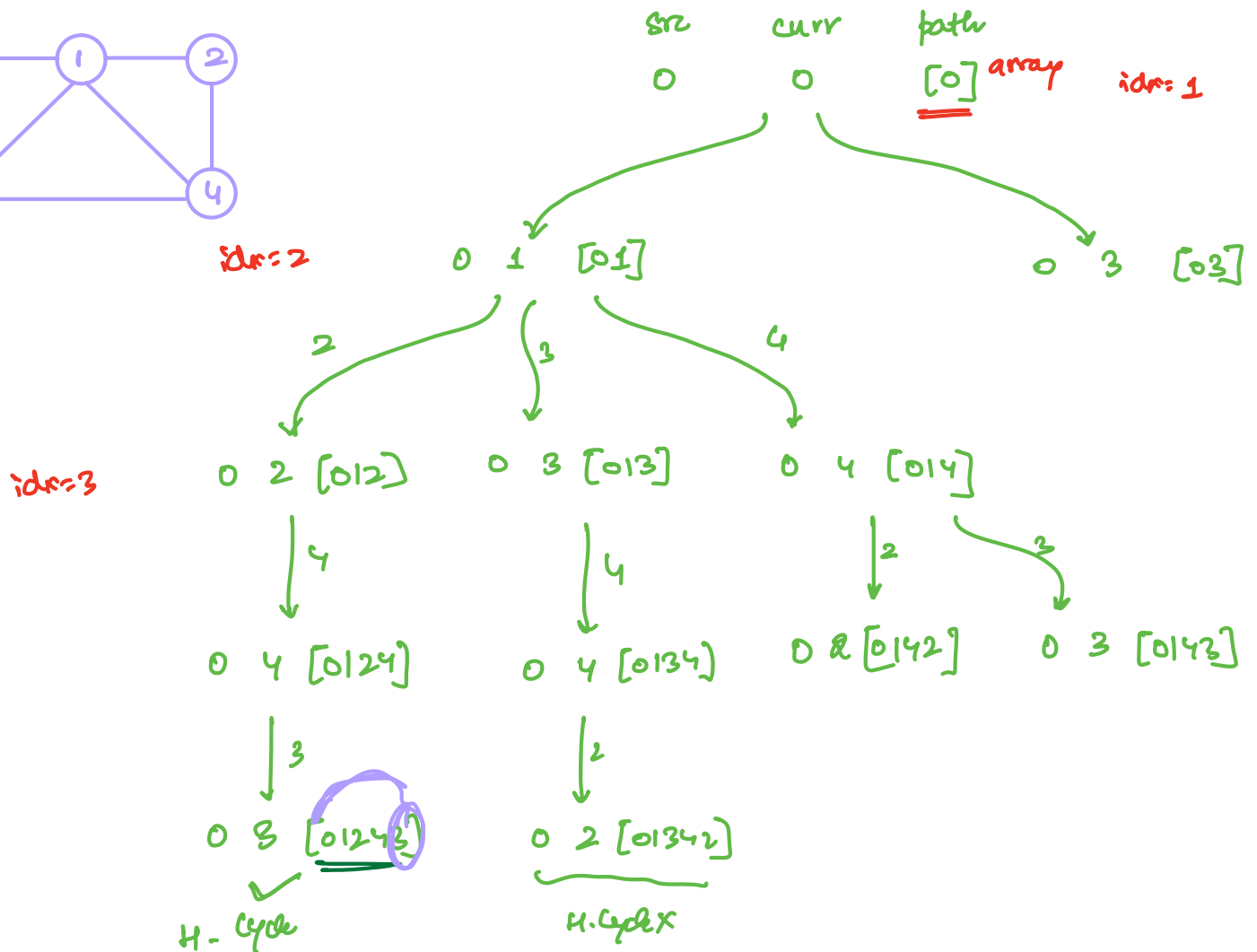
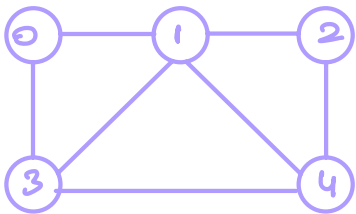


A **Hamiltonian Cycle (or Hamiltonian Circuit)** is a Hamiltonian Path such that there is an edge (in the graph) from the last vertex to the first vertex of the Hamiltonian Path.

If graph contains a Hamiltonian Cycle, it is called Hamiltonian graph otherwise it is non-Hamiltonian.

Task: Determine whether a given graph contains Hamiltonian Cycle or not. If it contains, then print the path.

Print all Hamiltonian Cycle:



PRINT ALL HAMILTONIAN CYCLE CODE

```
#include<iostream>
#include<map>
#include<queue>
#include<vector>

using namespace std ;

class Graph
{
    map<int, map<int,int> > strg ;
    int V ;

public :

    Graph(int V)
    {
        this->V = V ;
    }

    void addEdge(int u, int v, int cost)
    {
        strg[u][v] = cost ;
        strg[v][u] = cost ;
    }

    void display()
    {
        for(int i = 0 ; i < V ; i++)
        {
            cout << i << "\t" ;

            map<int, int>::iterator itr ;

            for(itr = strg[i].begin() ; itr != strg[i].end() ; itr++)
                cout << itr->first << "@" << itr->second << ", " ;

            cout << endl ;
        }
    }

    bool isItSafe(int *path, int nbr)
    {
        for(int i = 0 ; i < V ; i++)
        {
            if(path[i] == nbr)
                return false ;
        }

        return true ;
    }

    void hamiltonainCycle(int src, int curr, int *path, int idx)
    {
        if(idx == V)
        {
            if(strg[curr].count(src) != 0)
            {
                for(int i = 0 ; i < V ; i++)
                    cout << path[i] << " " ;

                cout << endl ;
            }

            return ;
        }

        map<int, int>::iterator itr ;
        for(itr = strg[curr].begin() ; itr != strg[curr].end() ; itr++)
        {
            int nbr = itr->first ;

            if(isItSafe(path,nbr))
            {
                path[idx] = nbr ;
                hamiltonainCycle(src, nbr, path, idx+1) ;
                path[idx] = -1 ;
            }
        }
    }
};
```



Cost(10) → 1
Cost(20) → 0


```

int main()
{
    int n = 5 ;
    Graph g(n) ;

    g.addEdge(0,1,3) ;
    g.addEdge(0,3,7) ;
    g.addEdge(1,2,2) ;
    g.addEdge(1,3,5) ;
    g.addEdge(1,4,1) ;
    g.addEdge(3,4,2) ;
    g.addEdge(2,4,2) ;

    g.display() ;
    int path[n] ;
    for(int i = 0 ; i < n ; i++)
        path[i] = -1 ;
    path[0] = 0 ;
    g.hamiltonainCycle(0,0,path,1) ;

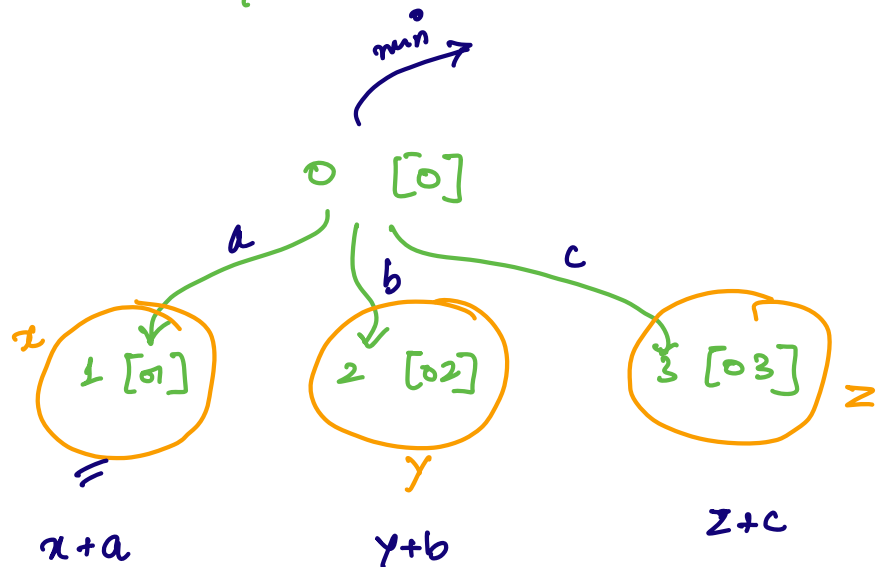
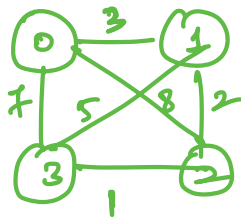
    return 0 ;
}

```

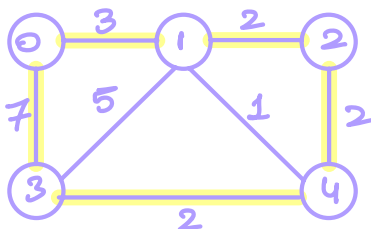
Travelling salesman Problem (TSP)

→ weighted graph

→ min wt hamiltonian cycle



$\min(x+a, y+b, z+c)$



TSP Cost: $7+2+2+2+3 = 16$

TSP CODE

```
int tsp(int src, int curr, int *path, int idx)
{
    if(idx == V)
    {
        if(strg[curr].count(src) != 0)
            return strg[curr][src] ;
        else
            return 100000;
    }

    int ans = 100000 ;
    map<int, int>::iterator itr ;
    for(itr = strg[curr].begin() ; itr != strg[curr].end() ; itr++)
    {
        int nbr = itr->first ;

        if(isItSafe(path,nbr))
        {
            path[idx] = nbr ;

            int rr = tsp(src, nbr, path, idx+1) ;
            ans = min(ans, rr+strg[curr][nbr]) ;

            path[idx] = -1 ;
        }
    }

    return ans ;
}
```