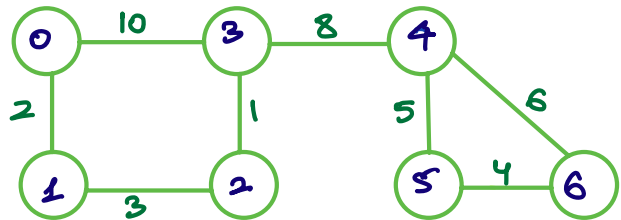


Dijkstra:

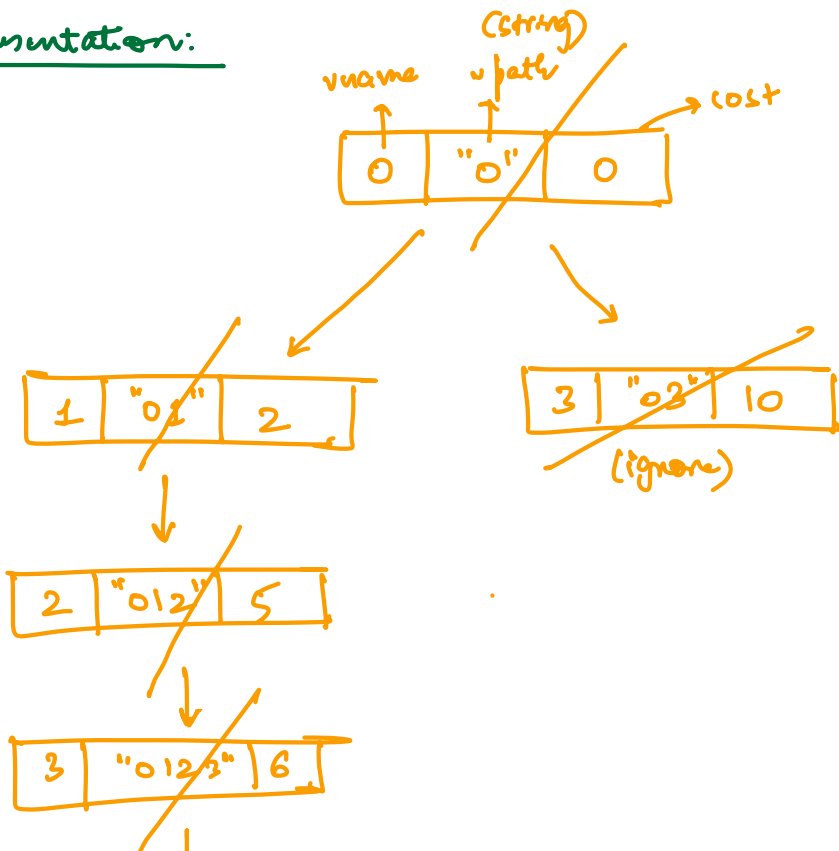
- Greedy Algo (Similar to Prim's)
- Priority Queue
- Source Node

src node = 0

0 → 1 : 2
 0 → 2 : 5
 0 → 3
 0 → 4
 0 → 5
 0 → 6



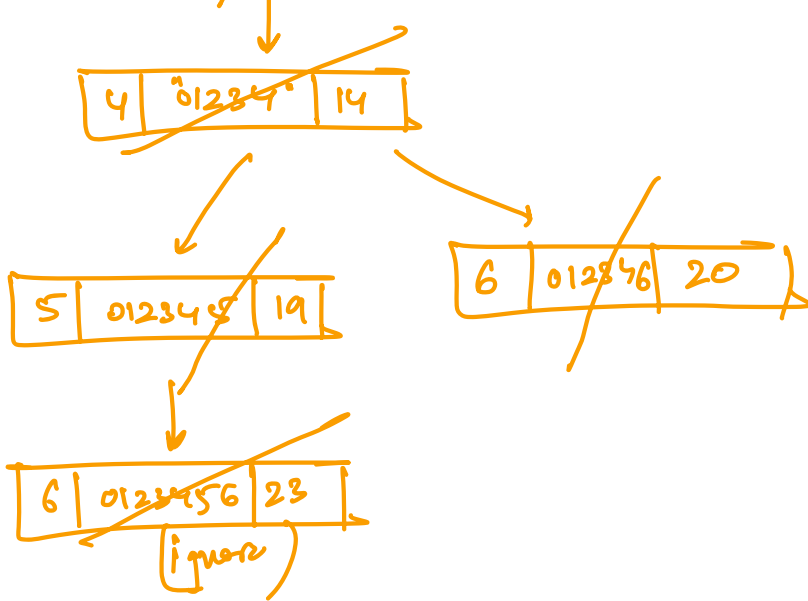
Implementation:



1. Remove
 2. visited
 2. Print
 4. nodes unvisited
- 0 ✓
 1 ✓
 2 ✓
 3 ✓
 4 ✓
 5 ✓
 6 ✓

0 via 0 @ 0
 1 via 01 @ 2
 2 via 012 @ 5
 3 via 0123 @ 6
 4 via 01234 @ 14
 5 via 012345 @ 19

6 via 012346 @ 20



$$\epsilon(1 + \log \epsilon + 1 + 1) + 2\epsilon$$

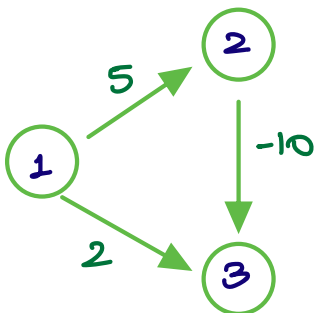
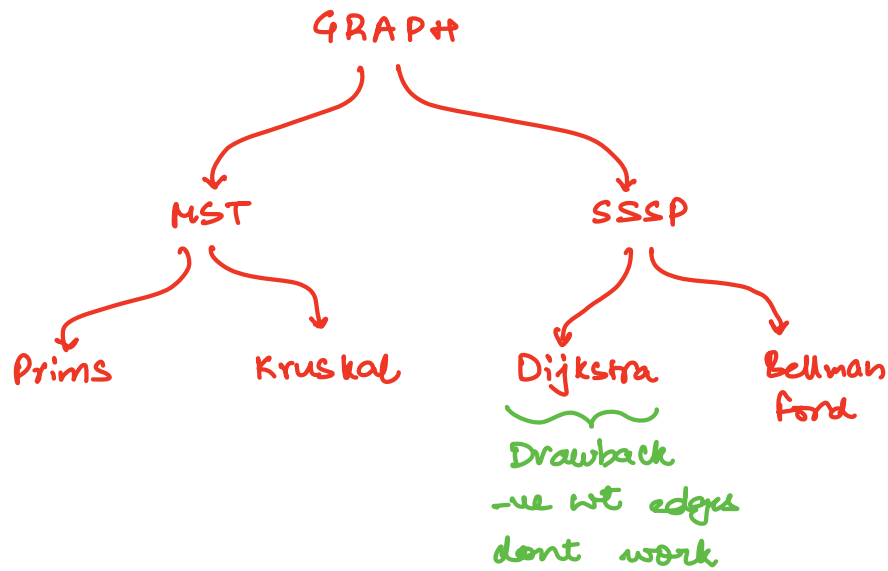
$$\epsilon \log \epsilon + 2\epsilon$$

$$\epsilon \log \epsilon$$

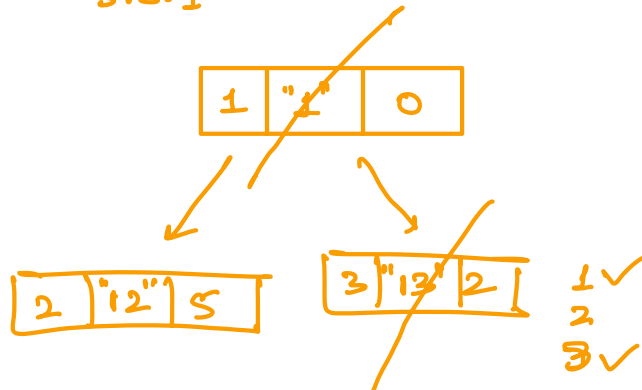
$$\epsilon \log V^2$$

$$O(\epsilon \log V)$$

$$\epsilon = V^2$$



src: 1

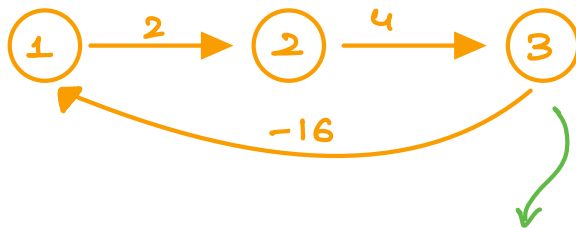


- remove
- visited
- print
- colors

1 via "1" @ 0
 3 via "12" @ 2
wrong answer

Bellman Ford:

- -ve wt edge ✓
- -ve wt cycle X
SSSP X



$$122 \ 122 \ 122 : -4 - 10 = -14$$

$$123 \ 123 : 6 - 16 + 6 = -4$$

$$123 : 6$$

SSSP \rightarrow no -ve wt cycle X

\rightarrow -ve wt edge but no -ve wt cycle : bellman ford

\rightarrow +ve wt edges: dijkstra, bf

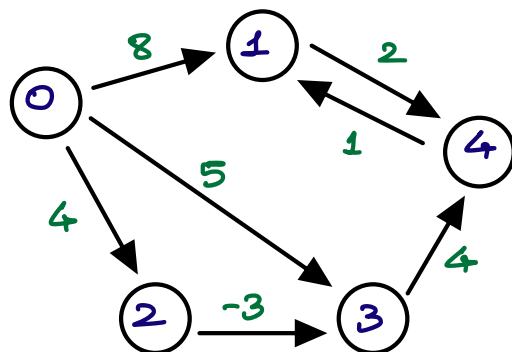
Example:

Edges

$0 \rightarrow 2 : 4$
 $0 \rightarrow 1 : 8$
 $0 \rightarrow 3 : 5$
 $2 \rightarrow 3 : -5$
 $1 \rightarrow 4 : 2$
 $4 \rightarrow 1 : 1$
 $3 \rightarrow 4 : 4$

Cost

$0 \rightarrow 0$
 $1 \rightarrow \infty$
 $2 \rightarrow \infty$
 $3 \rightarrow \infty$
 $4 \rightarrow \infty$



Relax every edge

$V-1$ times



edge $(u \rightarrow v)$

$\text{if } (\underbrace{\text{cost}[v]}_{\text{old cost}} > \underbrace{\text{cost}[u] + \text{cost edge } u \rightarrow v}_{\text{new cost}})$

$\text{cost}[v] = \text{cost}[u] + \text{cost edge } u \rightarrow v$

Edges

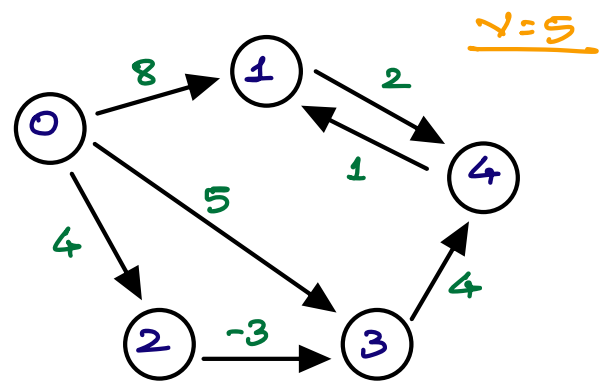
0 → 2 : 4
0 → 1 : 8
0 → 3 : 5
2 → 3 : -3
1 → 4 : 2
4 → 1 : 1
4 → 3 : 1
3 → 4 : 4

Cost

Src = 0

Update → Cost

0 → 0
1 → ~~∞~~ 2 6
2 → ~~∞~~ 4
3 → ~~∞~~ 8 1
4 → ~~∞~~ 10 5



1st time

0 → 2
∞ > 0 + 4
oc nc
✓

0 → 1
∞ > 0 + 8
nc
✓

0 → 3
∞ > 0 + 5
✓

2 → 2
5 4 - 2
5 > +1
✓

1 → 4
∞ > 8 + 2
✓

4 → 1
8 10 + 1
8 > 11
X

3 → 4
10 + 1 + 4
10 > 5

2nd time

0 → 2
4 0 + 4
X

0 → 1
8 0 + 8
X

0 → 3
1 0 + 5
X

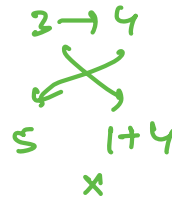
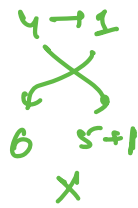
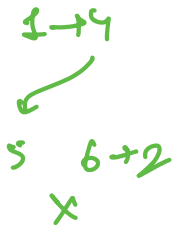
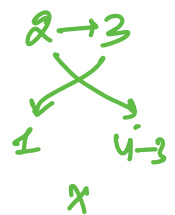
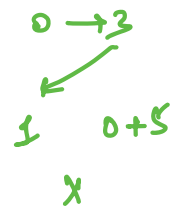
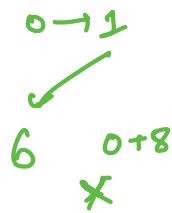
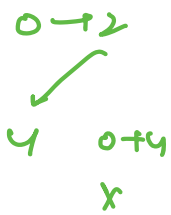
2 → 3
1 4 - 3
X

1 → 4
5 8 + 2
X

4 → 1
8 5 + 1
8 > 6
✓

3 → 4
5 1 + 4
X

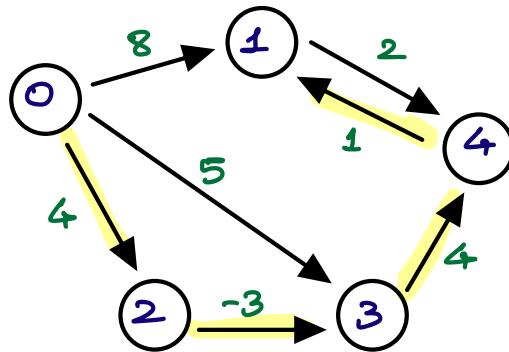
3rd time



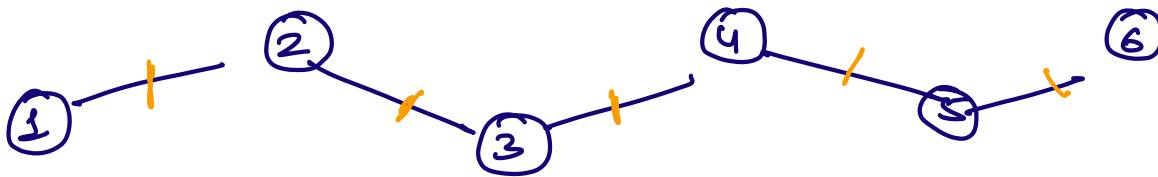
stop no changes

4th time

Vertex → Cost
0 → 0
1 → ~~2~~ 6
2 → ~~4~~
3 → ~~8~~ 1
4 → ~~10~~ 5

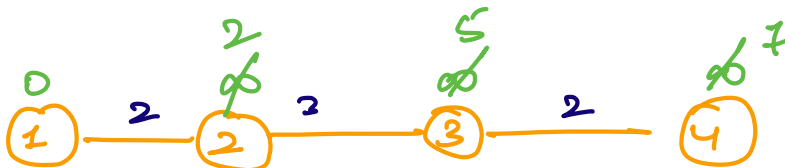


why V-1?



V=6
V-1 is the max distance possible b/w any 2 nodes

Eg:

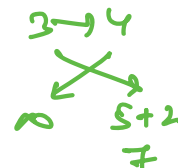
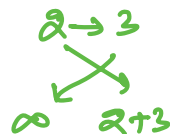
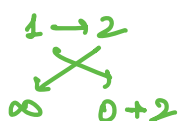


V=4

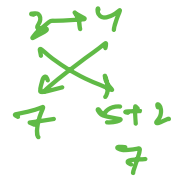
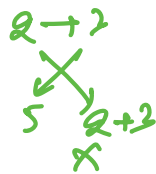
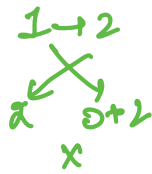
Edges:

1 → 2 : 2
2 → 3 : 3
3 → 4 : 2

1st time:



2nd time:



3rd time:

no change



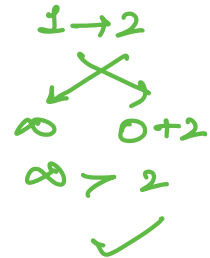
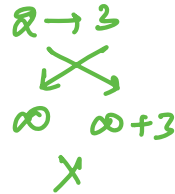
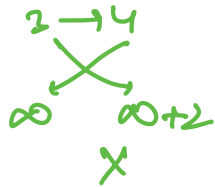
Edges:

3 → 4 : 2

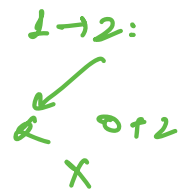
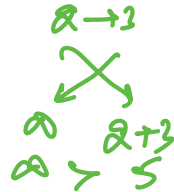
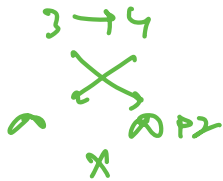
2 → 3 : 2

1 → 2 : 2

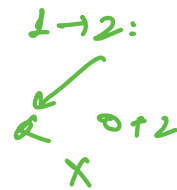
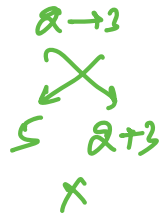
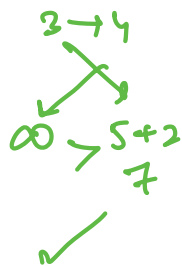
1st time



2nd time



3rd time



Bellman Ford → -ve wt edge ✓

-ve wt cycle

V-1 times relax

i th time : vertex wt update \Rightarrow -ve wt cycle



src=1

$N=4$

1st relax ✓
2nd relax ✓
2nd relax ✓
4th relax ↓
 claps

