

### THESIS DISSERTATION

# Cybersecurity Evaluation Methodology based on Metrics for Industrial Embedded Systems

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November 24, 2022

# **Declaration of Authorship**

I, Ángel LONGUEIRA-ROMERO, declare that this thesis titled, "Cybersecurity Evaluation Methodology based on Metrics for Industrial Embedded Systems" and the work presented in it are my own. I confirm that:

- This work was done wholly or mainly while in candidature for a research degree at this University.
- Where any part of this thesis has previously been submitted for a degree or any other qualification at this University or any other institution, this has been clearly stated.
- Where I have consulted the published work of others, this is always clearly attributed.
- Where I have quoted from the work of others, the source is always given. With the exception of such quotations, this thesis is entirely my own work.
- I have acknowledged all main sources of help.
- Where the thesis is based on work done by myself jointly with others, I have made clear exactly what was done by others and what I have contributed myself.

Signed:			
Date:			

"Aurrera begiratzen ez duena, atzean dago."

Those who don't look forward, stay behind.

Traditional basque proverb

# **Abstract**

#### MONDRAGON UNIBERTSITATEA

Faculty of Engineering
Electronics and Computer Department

Doctor of Philosophy

# Cybersecurity Evaluation Methodology based on Metrics for Industrial Embedded Systems

by Ángel LONGUEIRA-ROMERO

During the last decades, Embedded Systems (ESs) have evolved from isolated systems into a fully connected devices. Today, due to the development of processing capabilities, ESs have the ability to control or manage a wide range of systems or applications, including the critical infrastructures.

Historically, the development of ESs has been focused on functionality rather than security, and today this still applies in many sectors and applications. Moreover, there is an increasing number of security threats over ESs, and a successful attack could have severe economical or environmental consequences, including the loss of human lives.

A standardized and general accepted security testing framework is needed to provide guidance, common reporting forms and the possibility to compare the results along the time. This can be achieved by introducing security metrics into the evaluation or assessment process. If carefully designed and chosen, metrics could provide a quantitative, repeatable and reproducible value that would reflect the security level of the ES.

This dissertation is divided into three main topics, each one specialized in a different task of security evaluation: (1) Analysis of existing security metrics, (2) Vulnerability analysis, and (3) Aggregation of security metrics.

In the first section of the thesis, more than 500 metrics were collected and analyzed. The current issues of existing metrics were also extracted. As expected, the 77.5% of them is related exclusively to software, and only the 0.6% of them addresses exclusively hardware security. We also surveyed the literature to find the which are the desirable properties of a comprehensive metric. Based on these data, we proposed a new taxonomy to classify security metrics based on the properties of ESs.

In the second section, we present a model to analyze the known vulnerabilities of industrial components over time. The proposed Extended Dependency Graph (EDG) model is based on two main elements: a directed graph-based representation of the internal structure of the component, and a set of quantitative metrics based on the Common Vulnerability Scoring System (CVSS). The EDG model can be applied

throughout the entire lifespan of a device to track vulnerabilities, identify new requirements, root causes, and test cases. It also helps prioritize patching activities. The model was validated by application to the OpenPLC project.

Finally, in the third section we propose a CVSS aggregation algorithm that integrates information about the functionality disruption of the SUT, exploitation difficulty, existence of exploits, and the context where the SUT operates. The aggregation algorithm was applied to OpenPLC V3, showing that it is capable of filtering out vulnerabilities that cannot be exploited in the real conditions of deployment of the particular system.

This work aims to lay the foundations for constructing a security evaluation methodology that uses standardized metrics to quantify the security level of an ES.

# Laburpena

# MONDRAGON UNIBERTSITATEA Ingeniaritza Goi Eskola Politeknikoa Elektronika eta Informatikako Saila

Doktorego Programa

## Sistema Txertatu Industrialetarako Metriketan oinarritutako Zibersegurtasuna Ebaluatzeko Metodologia

by Ángel LONGUEIRA-ROMERO

Azken hamarkadetan, Sistema Txertatuak (ES) sistema isolatuak izatetik erabat konektatutako gailuak izatera igaro dira. Gainera, prozesatzeko gaitasuna handitzeari esker, ESek sistema eta aplikazio ugari kontrolatzen eta kudeatzen dituzte, kritikoenak barne.

Historikoki, ESen garapena funtzionaltasunean zentratu da segurtasunean baino gehiago, eta gaur egun ere horrela izaten jarraitzen du sektore eta aplikazio askotan. Horrek, ESei eragiten dieten mehatxuen kopurua handitzearekin batera, eraso arrakastatsu batek ondorio ekonomiko larriak izatea errazten du, baita giza bizitzak galtzea ere.

Orientazioa, txosten estandarizatuak eta emaitzak denboran zehar alderatzeko aukera emateko, beharrezkoa da alderdi horiek guztiak bateratuko dituen eta ebaluazio-prozesuan segurtasun-metrikak integratuko dituen metodologia bat. Kontu handiz diseinatzen eta aukeratzen badira, metrikek balio kuantitatiboa, errepikagarria eta erreproduzigarria eman lezakete, ESen segurtasun-maila islatuko lukeena.

Tesi hau hiru bloke nagusitan banatzen da, bakoitza zibersegurtasuna ebaluatzeko prozesuko zeregin batean espezializatua: (1) dauden segurtasun-metriken azterketa, (2) kalteberatasunen analisia, eta (3) Segurtasun-metriken agregazioa.

Tesi honen lehen atalean, 500 metrika baino gehiago bildu eta aztertu ziren, eta gaur egun metrikek dituzten arazoak atera ziren. Espero zen bezala, horien 77,5% softwarearekin soilik lotuta daude, eta 0,6% baino ez da hardware segurtasunari buruzkoa. Eskura dagoen literatura ere aztertu zen, metrika sakonean zer propietate nahi diren hautemateko. Datu horietatik guztietatik abiatuta, taxonomia berri bat garatu zen, segurtasun-metrikak ESen barne-egitura kontuan hartuta sailkatzen dituena.

Bigarren atalean, industria-osagaiek denboran zehar izan dituzten kalteberatasunak aztertzeko eredu bat aurkezten dugu. Proposatutako Extended Dependency Graph (EDG) eredua bi elementu nagusitan oinarritzen da: (2) System Under Test (SUT) delakoaren mendekotasunen barne-egituraren irudikapena, zuzendutako grafoetan oinarrituta; eta Common Vulnerability Scoring System (CVSS) delakoan oinarritutako metrika kuantitatiboen multzo bat. EDG eredua gailu baten balio-bizitza osoan aplika

daiteke, ahultasunen jarraipena egiteko, baldintza berriak identifikatzeko, urrakortasunen funtsezko kausak identifikatzeko eta test-kasu berriak sortzeko. Adabakiak aplikatzeko jarduerak lehenesten ere laguntzen du. Eredua OpenPLC proiektuan aplikatuta baliozkotu zen.

Azkenik, hirugarren atalean, GZEKren balioak gehitzeko algoritmo bat proposatzen da. Algoritmo horrek GLSren testuinguruko hainbat faktoreri buruzko informazioa biltzen du, hala nola funtzionaltasuna eteteari, ustiatzeko zailtasunari, esploiten existentziari eta SUTak jarduten duen testuinguruari buruzkoa. Agregazio-algoritmoa OpenPLC V3-ri aplikatu zitzaion, sistemaren hedapen-baldintza errealetan ustiatu ezin diren kalteberatasunak iragazteko gai dela frogatuz.

Lan honen bidez, IH baten segurtasun-maila kuantifikatzeko metrika estandarizatuak integratuko dituen zibersegurtasuna ebaluatzeko metodologia garatzeko oinarriak ezarri nahi dira.

# Resumen

### MONDRAGON UNIBERTSITATEA

Escuela Politécnica Superior de Ingeniería Departamento de Electrónica e Informática

Programa de Doctorado

## Metodología de Evaluación de la Ciberseguridad basada en Métricas para Sistemas Embebidos Industriales

by Ángel LONGUEIRA-ROMERO

En las últimas décadas, los Sistemas Empotrados (ES) han pasado de ser sistemas aislados a ser dispositivos totalmente conectados. Además, gracias al aumento en su capacidad de procesamiento, los ESs controlan y gestionan una amplia gama de sistemas y aplicaciones, incluidas aquellas más críticas.

Históricamente, el desarrollo de los ESs se ha centrado en la funcionalidad más que en la seguridad, y a día de hoy, esto sigue siendo así en muchos sectores y aplicaciones. Esto, sumado al aumento del número de amenazas que afectan a los ESs, facilita que un ataque exitoso pudiera tener graves consecuencias económicas, e incluso la pérdida de vidas humanas.

Para proporcionar orientación, informes estandarizados, y la posibilidad de comparar los resultados a lo largo del tiempo, se hace necesario una metodología que unifique todos estos aspectos, y que integre métricas de seguridad en el proceso de evaluación. Si se diseñan y eligen cuidadosamente, las métricas podrían proporcionar un valor cuantitativo, repetible y reproducible que reflejaría el nivel de seguridad de los ESs.

Esta tesis se divide en tres bloques principales, cada uno de ellos especializado en una tarea distinta del proceso de evaluación de la ciberseguridad: (1) Análisis de las métricas de seguridad existentes, (2) Análisis de vulnerabilidades, y (3) Agregación de las métricas de seguridad.

En la primera sección de esta tesis, se recopilaron y analizaron más de 500 métricas, y se extrajeron los problemas actuales a los que se enfrentan las métricas actualmente. Como era de esperar, el 77,5% de ellas están relacionadas exclusivamente con el software, y sólo el 0,6% de ellas aborda exclusivamente la seguridad hardware. También se analizó la literatura disponible para detectar cuáles son las propiedades deseables en métrica exhaustiva. A partir de todos estos datos, se desarrolló una nueva taxonomía que clasifica las métricas de seguridad teniendo en cuenta la estructura interna de los ESs.

En la segunda sección, presentamos un modelo para analizar las vulnerabilidades conocidas de los componentes industriales a lo largo del tiempo. El modelo Extended Dependency Graph (EDG) propuesto se basa en dos elementos principales: (2) La representación de la estructura interna de dependencias del System Under Test (SUT) basada en grafos dirigidos; y un conjunto de métricas cuantitativas basadas en el

Common Vulnerability Scoring System (CVSS). El modelo EDG puede aplicarse a lo largo de toda la vida útil de un dispositivo para hacer un seguimiento de las vulnerabilidades, identificar nuevos requisitos, identificar las causas fundamentales de las vulnerabilidades, y generar nuevos casos de test. También ayuda a priorizar las actividades de aplicación de parches. El modelo se validó mediante su aplicación al proyecto OpenPLC.

Por último, en la tercera sección se propone un algoritmo de agregación de valores CVSS que integra información sobre varios factores de contexto del SUT, como la interrupción de la funcionalidad, la dificultad de explotación, la existencia de exploits y el contexto en el que opera el SUT. El algoritmo de agregación se aplicó a OpenPLC V3, demostrando que es capaz de filtrar las vulnerabilidades que no pueden ser explotadas en las condiciones reales de despliegue del sistema.

Este trabajo pretende sentar las bases para el desarrollo de una metodología de evaluación de la ciberseguridad que integre métricas estandarizadas para cuantificar el nivel de seguridad de un ES.

# Acknowledgements

This achievement would not be possible without the support, patience, and pieces of advice of many people. People that might not be directly involved in this work at first, but without them, all this work would not have turned into reality.

I want to start thanking my advisors, Iñaki Garitano and Rosa Iglesias. I am grateful for their support, encouragement, time, and knowledge. They managed to guide me during the difficult times, and celebrate each little achievement that was made. Moreover, they taught me interpersonal skills that I would not be able to learn in any other situation. A PhD. is not only about acquiring technical knowledge, but also about learning how to deal with people, frustration, and uncertainty. This is about knowing yourself, and improving every day.

As it could not be otherwise, I also want to thank Jose Luis Flores for giving me his most valuable resources: his time, and his experience. The quality of this work would not be the same without his advice.

I cannot forget my parents, that pushed me forward and always encouraged me to do my best. They mean a world to me. Thanks to them, I am able to stand where I am now. Leaving home to follow your own way is not an easy decision, but their support was and still is invaluable. Thank you both.

As important and my family, my chosen family is a core piece in my life. My friends, my colleagues at Ikerlan, my friends that were also working on their thesis. They were with me to help me to deal with every problem, failure, or issue. They were also always there to celebrate with me every achievement. And of course, I cannot forget my partner, the biggest serendipity I have ever lived. He had to see bits of me that I did not know even existed. Nevertheless, he managed to guide me, and turned bad times into good moments and memories. Every weekend was a new adventure to disconnect from the hard work of the week.

And finally, I want to thank Euskal Herria, because since I came here, everyone was really charming. I have never felt so accepted and integrated. I learned a new language, I was immersed in a new culture, and I also met new people. Mila esker bihotz bihotzez denagatik.

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# List of Abbreviations

CAPEC Common Attack Pattern Enumeration and Classification

CC Common Criteria
CPS Cyber-Physical System

CVSS Common-Vulnerability Scoring System
CVE Common-Vulnerabilities and Exposures

CWE Common-Weakness Enumeration

ES Embedded System SM Security Metric

**ECSO** European Cyber Security Organisation

ICS Industrial Control System

**IoT** Internet of Things

NIST National Institute of Standards and Technology
OSSTMM Open Source Security Testing Methodology Manual

OWASP Open Web Application Security Project

PC Personal Computer SUT System Under Test

# Chapter 1

# Introduction

Since the First Industrial Revolution in the 17th century, the society we live in has evolved along with the advances brought about by industry. Today, we live in the fourth industrial revolution, or Industry 4.0, which is characterized by trends such as automation and data exchange, including technologies like Cyber-Physical Systems (CPS), Internet of Things (IoT), and edge computing. Unfortunately, quite often, technological advances of each industrial revolution bring some drawbacks which need to be addressed. All the advances introduced so far have one issue in common that was not a key aspect until now: cybersecurity.

In this chapter, the critical role of current CPS is highlighted. Second, we describe the features that are more critical to the security of CPS to set the motivation of this thesis. Then, we put forward the hypotheses that this work uses as a basis. Finally, we review the publication associated with this research and the structure of this dissertation.

# 1.1 Motivation and Objective of the Research

Embedded Systems (ES) evolved from fully analog devices to digital devices with the Third Industrial Revolution in the late 20th century, when the industry started to adopt digital computers. With this digital revolution, and the proliferation of digital record-keeping, data became an even more critical asset, and so did security. Nevertheless, the breakthrough for ESs came later, in the 21st century. The so-called Fourth Industrial Revolution (or Industry 4.0) introduced interconnectivity and smart automation to the landscape of the existing industries. With these new trends, ESs evolved again to integrate all these features, serving as a starting point for developing other technologies (*e.g.*, the Internet of Things). The new connectivity capabilities of ESs increase their exposures, and therefore, their attack surface, making them a promising target for attackers.

#### 1.1.1 Motivation

ESs, and more broadly industrial components, are the driving force of almost every industrial field, such as automotive, energy production, and transportation, including critical infrastructures [1–7]. A successful malicious attack could have severe consequences, turning security into a key issue for ESs. These types of components are rapidly evolving [8,9] and increasing in number [10]. This fact is related to several factors:

- The reuse of hardware and software: Open-source hardware and software, and Commercial Off-The-Shelf (COTS) components are being integrated to speed up the development of industrial components [11–13]. COTS are easy to use, but they can introduce vulnerabilities, creating potential entry points for attackers [14,15].
- New connectivity features: Industrial components are providing more advanced connectivity features, enabling new automation applications, services, and data exchange. This new connectivity, boost by the fifth generation (5G) of wireless technology for cellular networks, is further opening the window of exposure to any threat [6, 10, 16, 17].
- More complex systems: The complexity of industrial components is also increasing with the integration of new trends, such as the Internet of Things (IoT) [17–20], cloud computing, Artificial Intelligence (AI) [20,21], and big data. The extensive use of these technologies further opens the windows for attackers [22–27]. Complexity is a critical aspect of industrial components design, because it is closely related to the number of vulnerabilities [28,29].

This scenario points security is a key aspect of industrial components. Moreover, numerous attacks have been reported targeting industrial enterprises across the globe since 2010 [30]. An exponential rise in such attacks is predicted for future years [31,32].

In summary, the rapid evolution of ESs, their connectivity, and the integration of more and more features, increase their attack surface. This makes it essential to protect their use in environments such as critical infrastructures [33,34]. The sophistication of attacks, a larger attack surface, and the ease of attacks thanks to exploits and tools that decrease the necessary knowledge of the attackers, highlights the need to invest more in cybersecurity. The numerous attacks targeting industrial enterprises across the globe since 2010 reinforce this fact [30], and an exponential rise in such attacks is predicted for the upcoming years [31,32]. As a consequence, security is turning into a critical issue for ESs [35]. However, security by itself is not enough, and the degree of coverage of the implemented countermeasures also has to be evaluated to know whether they are sufficient. Tracking the security status of an ES [36,37] and considering both software and hardware in the evaluation would also be desirable [38–41].

# 1.2 Research Hypothesis

In this section, the hypotheses of the research are defined. To overcome the process of definition of hypotheses, it has been necessary to analyze the research gaps having as basis the context of the area described in Chapter 2 and related scientific contributions analyzed in Chapter 3.

The hypotheses to which this research seeks to provide an answer are set out and explained below:

H1. The intrinsic characteristics of embedded systems make it necessary to create a tailored security evaluation methodology.

ES have two main factors that define them: (1) They are resource-constrained, so power consumption, memory, or functionality are usually limited as a requirement; (2) Their security level is also related to the location and environment where they are working. These differences have to be taken into account when considering existing metrics or creating new ones for ESs, when classifying security metrics, and when assessing the level of security of an ES.

# H2. The result of the evaluation process should be a combination of a numeric value and a detailed description.

The evaluation result is not enough if the conditions under which the test was carried out are not set. Suppose the report describing the procedure followed during the evaluation is attached to the very first result of the evaluation. In that case, it is possible to re-evaluate the device under the same conditions, obtaining a similar and comparable result over time.

# H3. Not all existing security metrics provide a quantitative, repeatable, and reproducible value that reflects the security level of the ES.

A metric that returns a similar result under the same circumstances is needed as a solid foundation for an evaluation methodology. Moreover, this stability allows extracting conclusions, and make decisions.

# H4. Breaking down the ES into its assets makes the security evaluation more repeatable:

This break down gives the evaluator the ability to understand the inner of the System Under Test (SUT). Moreover, this break down let us track possible issues that can appear during the life cycle of ES, effectively implementing a constant security evaluation. It is also a powerful tool to develop more security metrics.

# H5. The aggregation of metrics can serve to identify the more critical elements in an ES, while helping to make decisions about the security status of the device.

An aggregation mechanism can be applied partially to know which parts are the more vulnerable, and therefore, the ones that need to be addressed first (decision-making tool). Moreover, they can give also an overview of the security status of an ES.

### 1.2.1 Objective of the Research

Measuring security is not an easy task, but it can be a useful tool to make decisions. Nevertheless, the result of a security evaluation process does not have to be extremely accurate, but meaningful and useful, as well as representative to serves as a decision making tool. This idea is the foundation of this research that can be summarized in the following main, high-level objective:

## **MAIN OBJECTIVE**

Design, implement, and validate a security assessment methodology tailored to embedded systems.

This research is based on two main subjects: (1) evaluation methodologies, and (2) security metrics. Independent sub-objectives are set for each:

### 1.2.1.1 Evaluation Methodology objectives

- EM1. Identify the main differences between evaluating an ES and a traditional IT system: Knowing which key aspects are different in each device type will serve as the basis for building the security evaluation methodology oriented to ES.
- EM2. **Identify and define repeatable, reproducible and quantitative security metrics, considering dimensions such monetary cost and time:** These properties ensure that the security evaluation methodology is objective, so evaluations carried out under the same conditions and using the same devices give the same result. Moreover, the methodology will track the evolution of the security state of the device through multiple evaluations over time. The cost in money and time of the evaluation process will also be considered to make the methodology applicable.
- EM3. **Identify and define how the security evaluation process should be documented:** The evaluation result will be accompanied by a document detailing the conditions under which it was carried out. This ensures the assessment result is repeatable, and reproducible.
- EM4. Break down the ES into its minimal assets for their joint and individual evaluation: Breaking down an ES into its most basic assets can give us a more detailed picture of the security status. Moreover, we can use this information to extract conclusions about the relationships between assets.
- EM5. **Re-evaluate ESs over time:** Tracking the evolution of the security status of ESs is a crucial to monitor potential vulnerabilities, prioritize patching activities, and track the development cycle.

### 1.2.1.2 Security Metric objectives

- SM1. **Analyze the applicability of the existing security metrics and the proposed taxonomies for classifying them:** Limitations in resources (*e.g.*, computation, and power consumption) mean that not every metric is suitable for evaluating the security of ESs.
- SM2. **Develop a meaningful method for metric aggregation:** Metrics carry a wealth of information, but a summarized value can be useful for non-technical users.

1.3. Scope 5

To better understand the relationships among hypothesis and objectives, in Table 1.1 a mapping of the objectives and the hypothesis is presented.

TABLE 1.1: Mapping of the initial objectives of the thesis to the pro-				
posed hypothesis.				

Objectives	Hypothesis
EM1	H1
EM2	H3, H4
EM3	H2
EM4	H3, H4
EM5	H2, H4
SM1	H1, H3
SM2	H2, H3, H5

## 1.3 Scope

Developing a comprehensive security evaluation methodology for ES is an ambitious goal. Therefore, in this section we explain which aspects of our objectives we have been able to achieve, and which have remained outside the scope of this research work. The latter will be considered as future work.

## 1.3.1 Evaluation Methodologies Objectives

- EM1. **Identify the main differences between evaluating an ES and a traditional IT system:** We found that there are differences when evaluating ES and IT systems, being the most important one that ESs usually have restrictive safety requirements.
- EM2. Identify and define repeatable, reproducible and quantitative security metrics, considering dimensions such a monetary cost and time: We were able to identify existing security metrics, but we were not able to test them for repeatability and reproducibility. Nevertheless, we proposed a set of security metrics oriented to vulnerabilities that are indeed repeatable and reproducible.
- EM3. **Identify and define how the security evaluation process will be documented:** Although we explored how each one of the analyzed standards document their findings, we have not proposed a novel documentation flow.
- EM4. Break down the ES into its minimal assets for their joined and individual evaluation: In our use case, we were able to break down the SUT into its minimal assets. Moreover, we used a directed graph-based representation to also take their relationship into account.
- EM5. The developed evaluation should be carried out over time: Our vulnerability analysis model makes it possible to evaluate an ES during its whole life spam in a continuous manner.

## 1.3.2 Security Metrics Objectives

- SM1. Check the applicability of the existing security metrics and the proposed taxonomies for classifying them: Although we could obtain a list with the existing metrics, and we analyzed their possible applicability to ESs, we did not test them. Nevertheless, we proposed a new taxonomy to classify them according to the structure and properties of ESs.
- SM2. Develop a meaningful method for metric aggregation. Validate the developed method on a real use case: We proposed a method to aggregate CVSS scores specifically.

### 1.4 Contributions

The contributions of this thesis are described below:

- 1. Systematic review of existing security metrics.
- 2. Proposed ES-oriented taxonomy to classify security metrics.
- 3. Directed graph-based model to track vulnerabilities over time in ESs.
- 4. Quantitative metrics to measure the vulnerability status of ES over time.
- 5. CVSS scores aggregation method.

## 1.5 Publications

Table 1.2 shows the publications of this thesis, and Table 1.4 shows the correspondence between the publications and the state of the art.

TABLE 1.2: List of publications achieved in the thesis.

Type of Publication	Name	Status
INDIN 2020 Conference	How to Quantify the Security Level of Em-	Published
	bedded Systems? A Taxonomy of Security	
	Metrics	
MDPI Sensors	A Novel Model for Vulnerability Analysis	Published
	through Enhanced Directed Graphs and	
	Quantitative Metrics	
Conference	Gotta Catch 'em All: Aggregating CVSS	Published
	Scores	

TABLE 1.4: Mapping of the topics reviewed in the state of the art to the publications of the thesis.

Topic	Publication
1) Evaluation methodologies	How to Quantify the Security Level of Embedded Systems? A Taxonomy of Security Metrics
2) Fundamental features of a comprehensive methodology	How to Quantify the Security Level of Embedded Systems? A Taxonomy of Security Metrics
3) Testing types and techniques	Thesis Research Plan
4) Existing standards	Thesis Research Plan
5) Security metrics	How to Quantify the Security Level of Embedded Systems? A Taxonomy of Security Metrics
6) Directed Graphs	A Novel Model for Vulnerability Analysis through Enhanced Directed Graphs and Quantitative Metrics
7) Vulnerability Analysis	A Novel Model for Vulnerability Analysis through Enhanced Directed Graphs and Quantitative Metrics
8) Metric Aggregation	Are you Gonna Buggy?: Aggregating CVSS Scores in EDG

## 1.6 Document Structure

This document is structured as follows:

- 1. **Introduction:** In this chapter, the motivations of this research work is described, summarizing the main reasons why ES security is important, and describing the problem to solve in a simplified way.
- 2. **Background:** In order to understand every aspect of this thesis, the most important concepts and key terms are described and explained, such as ES, security evaluation, and security metrics.
- 3. **Literature Review:** To set the foundations, and better understand the context, and the gaps in ESs security, a systematic literature review is presented in this chapter. In this section, the context of this research work is also introduced.
- 4. How to Quantify the Security Level of Embedded Systems? A Taxonomy of Security Metrics: In this chapter, we systematically review the literature to gather all existing security metrics to analyze them. With our conclusions, we

- proposed a new taxonomy to classify them, according to the structure of ES. This is the first contribution of this thesis.
- 5. A Novel Model for Vulnerability Analysis through Enhanced Directed Graphs and Quantitative Metrics: In this chapter, we introduce a new model based on directed graphs to track known vulnerabilities of ESs over time. To reinforce the proposed model, we also developed an initial set of quantitative metrics. This is the second contribution of the thesis.
- 6. **Gotta Catch 'em All: Aggregating CVSS Scores:** In the last contribution of this thesis, we proposed a novel method of aggregation of CVSS values.
- 7. **Discussion, Conclusions and Future Work:** Finally, in this last chapter, we extract the conclusions from the contributions that we have proposed during this document, and we discuss them. The future work is also explained.

# Chapter 2

# Background

Cybersecurity is a very complex subject, and so are CPSs. Commonly, CPSs are designed and developed to solve a particular problem. The interactions between the different components and systems have to be taken into account to really make an CPS secure, including its connection to the Internet, if any.

In this chapter, the basic concepts related to embedded systems, embedded security, and security evaluation are presented.

## 2.1 Embedded System

An ES is an applied computer system. Although the definition of ES is not standardized, and it is difficult to pin down, there are some common properties (summarized in Figure 2.1) that characterize them when compared with a Personal Computer (PC) [42–46]:

- **ESs objective:** most embedded devices are primarily designed to accomplish a particular task or a group of specific functions (e.g., an automatic dishwasher, or a digital camera).
- **ESs hardware constrains:** contains in processing performance, power consumption, memory, hardware functionality, and so forth.
- ESs software limitations: fewer applications, scaled-down applications, no Operating System (OS) or a limited OS, or less abstraction-level code.
- **ESs quality and reliability requirements:** some families of ESs have a very high threshold of quality, reliability and safety requirements. However, there are also embedded devices in which a malfunction is an inconvenience but not usually a life-threatening situation (e.g., a vending machine).

From this point, any definition of ES must be augmented with examples to better understand the concept. And one of the most significant examples is an automobile, more precisely, all the components and systems that integrate it. An automobile itself is not an ES, but its infotainment head-unit, anti-lock breaking system, powertrain engine control unit, digital instrument cluster, and a plethora of other electronic subsystems - dozens in the typical modern card - are all examples of embedded systems (see Figure 2.2).

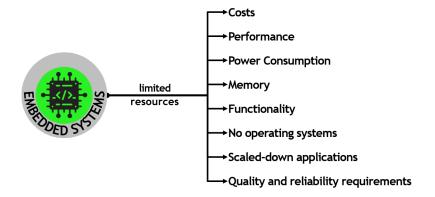


FIGURE 2.1: Most common limited resources in an embedded system.

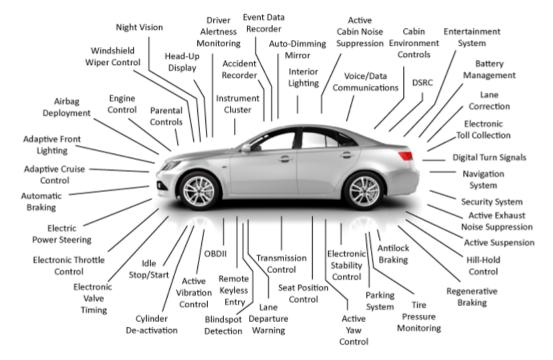


FIGURE 2.2: Summary of some ESs that can be found inside a car.

Nowadays, ESs are present in almost every electronic device: automotive, railway, consumer electronic (microwaves, ovens, washing machines...), vending machines, and manufacturing, to name a few. Even though this definition seems clear, there is some controversy that arises when classifying smartphones. When they were invented, they were conceived to perform just one task - i.e., calling to another device. They were simple devices whose functions were very limited. Nowadays, smartphones have way more functionalities that those initial mobile phones, and calling to another mobile phone is no longer the priority tasks. Some say that smartphones are not ES, but just a miniature desktop computer because of this. Nevertheless, there is little debate that individual components within the phone are ESs.

Finally, it is worth mentioning that ESs are also present in critical infrastructures, such as nuclear power plants, or the power grid. Because of their key role in the modern society, these kinds of infrastructures require a special protection and care, as

a malicious attack could have severe consequences.

## 2.1.1 Cyber-Physycal System

CPS [7] is a term equivalent to ES, but it makes emphases in the interaction with the physical world. They integrate computation, networking, and physical processes. Embedded computers and networks monitor and control the physical processes, with feedback loops where physical processes affect computations and vice versa. CPS integrates the dynamics of the physical processes with those of the software and networking, providing abstractions and modelling, design, and analysis techniques for the integrated whole.

## 2.1.2 Industrial Control System (ICS)

Industrial Control System (ICS) [7] is a general term that encompasses several types of control systems and associated instrumentation used for industrial process control. They are usually considered a special kind of ES. Such systems can range in size from a few modular panel-mounted controllers to large interconnected and interactive distributed control systems with many thousands of field connections.

## 2.1.3 Internet of Things

The IoT can be defined as an ES with internet connection [47]. Initially, this referred to the trend of connecting everyday objects to the internet, connecting to it more objects than people. Nowadays, the definition of the Internet of things has evolved due to the convergence of multiple technologies, real-time analytics, machine learning, commodity sensors, and embedded systems. Traditional fields of embedded systems, wireless sensor networks, control systems, automation (including home and building automation), and others all contribute to enabling the Internet of Things. Nowadays, they can be seen as an open and comprehensive network of intelligent objects that have the capacity to auto-organize, share information, data and resources, reacting and acting in face of situations and changes in the environment.

### 2.1.4 Industrial Automation and Control Systems (IACS)

According to ISA/IEC 62443 [48] Industrial Automation and Control System (IACS) are the collection of personnel, hardware, software, procedures, and policies involved in the operation of the industrial process and that can affect or influence its safe, secure and reliable operation.

### 2.1.5 Industrial Component

According to the ISA/IEC 62443-4-2 standard, an Industrial Component (IC) is one of the parts that make up an industrial product or system, including hardware, software, or other components [48]. The ISA/IEC 62443-4-2 document defines four types of components [49]:

• **Software Application**: One or more software programs and their dependencies<sup>1</sup> that are used to interface with the process or the control system itself (*e.g.*,

<sup>&</sup>lt;sup>1</sup>Any software program that is necessary for the software application to function, such as database packages, reporting tools, or any third party or open source software.

configuration software and historian). Software applications typically execute on host devices or embedded devices.

- Embedded Device: Special purpose device designed to directly monitor or control an industrial process (e.g., PLCs, wired or wireless field sensor devices, wired or wireless field actuator devices, Safety Instrumented System (SIS) controllers, distributed control system (DCS) controllers). According to the ISA/IEC 62443, typical attributes of these devices are limited storage, limited number of exposed services, programmed through an external interface, embedded operating systems (OSs), or firmware equivalent, real-time scheduler, may have an attached control panel, and may have a communications interface.
- Host Devices: General purpose device running an operating system capable
  of hosting one or more software applications, data stores or functions from
  one or more suppliers. This kind of devices typically include attributes such
  as filesystem, programmable services, no real-time scheduler, and full HMI
  (keyboard, mouse, etc.).
- **Network Device**: Device that facilitates data flow between devices, or restricts the flow of data, but may not directly interact with a control process. According to the standard, these devices typically include embedded OS or firmware, no HMI, no real-time scheduler, and configured through an external interface.

## 2.2 System Under Test (SUT)

The System Under Test (SUT) is defined as a system such as a Control System (CS), SCADA system or Monitoring System that is available from a single system supplier [50]. A SUT may be comprised of hardware and software components from several manufacturers but must be integrated into a single system and supported, as a whole, by a single supplier. The SUT may be an IC, a part of an IC, a set of ICs, a unique technology that may never be made into a product, or a combination of these.

### 2.3 Weakness

Weaknesses are flaws, faults, bugs, and other errors in software and hardware design, architecture, code, or implementation that, if left unaddressed, could result in systems, networks, and hardware being vulnerable to attacks [51] (e.g., buffer overflow).

# 2.4 Vulnerability

A vulnerability is a flaw in a software, firmware, hardware, or service component resulting from a weakness that can be exploited, causing a negative impact to the confidentiality, integrity, or availability of an impacted component or components. It can be used to gain remote or physical access to a system [52]. To put it in other words: All vulnerabilities rely on weaknesses, but not all weaknesses entail vulnerabilities.

Vulnerabilities can be classified as both known and unknown [53]. Known vulnerabilities are those that have been published in publicly available sources such as the National Vulnerability Database (NVD) [54]. On the other hand, unknown vulnerabilities are dormant vulnerabilities that have not been publicly exposed or exploited.

2.5. Attack Pattern 13

These potential vulnerabilities have to be discovered by using other methods such as penetration testing [55]. In some cases, unknown vulnerabilities might be known for a group of attackers that do not want to disclose their knowledge to take malicious advantages of it (zero-day vulnerabilities) [56].

### 2.5 Attack Pattern

An attack pattern is a description of the common attributes and approaches employed by adversaries to exploit known weaknesses in cyber-enabled capabilities [57]. Attack patterns define the challenges that an adversary may face and how they go about solving it.

# 2.6 Common Platform Enumeration (CPE) Scheme

CPE is an abstract structured naming scheme for describing and identifying applications, operating systems, software, and hardware, including industrial control systems, such as Supervisory Control And Data Acquisition (SCADA) [58]. The logical construct of a CPE is called "Well-Formed CPE Name (WFN)" [59]. The CPE scheme, however, cannot describe and identify specific instances of products (e.g., using serial numbers, particular licenses, or physically discernible products) [60]. CPE is operated by the NIST [61]. The latest version at the time this paper was written is version 2.3.

Version 8.0.6001 of Internet Explorer for its *beta* update can be represented using version 2.3 of the CPE naming specification<sup>2</sup>. The "a" after the CPE version indicates that the following element is a software application.

# 2.7 Common Weakness Enumeration (CWE)

CWE is a community-developed list of common software and hardware weakness types, each one associated with some CVEs (explained in the next subsection) [51,62]. CWE is operated by the MITRE Corporation [63]. The latest version at the time this paper was written is version 4.3.

CWE-119 is an example of a weakness present in all updates of version 8 of Internet Explorer. It is related to improper restriction of operations within the bounds of a memory buffer.

# 2.8 Common Vulnerabilities and Exposures (CVE)

CVE is a list of common identifiers for publicly known cybersecurity vulnerabilities [51,64,65] operated by the MITRE Corporation [63]. Each CVE includes a unique identification number, a description, one public reference, a reference to the software or hardware that is affected (using CPE), a reference to a weakness, and its severity [66]. As Dimitriadis *et al.* state in [66], "it can be assumed that CVEs are "known knowns" (things we are aware of and understand) or specific vulnerabilities, while

CWEs are "unknown knowns" (things we understand but are not aware of) or generic vulnerability types." The latest CVE version is always available in its official site<sup>3</sup>.

CVE-2015-6154 is an example of a vulnerability that exploits the CWE-119 weakness present in Internet Explorer 8, allowing remote attackers to execute arbitrary code or cause a denial of service (memory corruption) via a crafted website.

# 2.9 Common Vulnerability Scoring System (CVSS)

CVSS is a public framework that provides a standardized method for assigning quantitative values (scores) to security vulnerabilities (CVE) [65] according to their severity [67]. A CVSS score is a decimal number in the range [0.0, 10.0]. The latest version at the time this paper was written is version 3.1.

The CVSS value for vulnerability CVE-2015-6154 is 9.3 out of 10. This value is classified as critical.

# 2.10 Common Attack Pattern Enumeration and Classification (CAPEC)

CAPEC is a comprehensive dictionary that provides a publicly available classification taxonomy of known attack patterns (security threats) [68]. In [66], Dimitriadis *et al.* explain that each CAPEC describes the common characteristics of a cyber threat (aka attack-pattern), helping to explain how applications and other cyber-enabled capabilities are likely to be attacked. CAPEC also provides the likelihood that the attack pattern can occur and its impact. CAPEC utilizes a qualitative approach, rating both likelihood and impact in a five-step value scale ranging from very low to very high. Finally, each CAPEC records the weaknesses (CWEs) that the attack pattern can exploit. The latest version at the time this article was written is version 3.4. Figure 2.3 shows the relationship between the different standards [66].

# 2.11 White-box Testing

White-Box Testing is the detailed investigation of internal logic and structure of the code. In white box testing, it is necessary for a tester to have full knowledge of source code [69].

# 2.12 Black-box Testing

Black-Box Testing is a technique of testing without having any knowledge of the internal working of the application. It only examines the fundamental aspects of the system and has no or little relevance with the internal logical structure of the system [69].

<sup>&</sup>lt;sup>3</sup>https://cve.mitre.org/

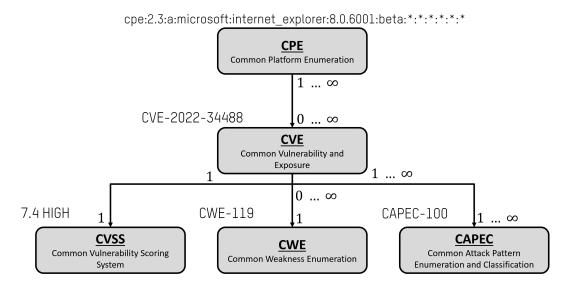


FIGURE 2.3: Relationship between the security standards defined by MITRE and NIST. Taken from [66].

# 2.13 Grey-box Testing

Grey-Box Testing is a technique to test the application with limited knowledge of the internal working of an application and also has the knowledge of fundamental aspects of the system [69]. It is an intermediate point between white-box and blackbox testing.

# Chapter 3

# Literature Review: State of the Art

Cybersecurity assessment involves not only cybersecurity assessment methodologies, but cybersecurity metrics as well. This chapter describes and discusses relevant scientific works performed within those two cybersecurity assessment categories. Bellow, we list the points that we have researched in our state of the art:

- 1. **Security metrics:** Proposed security metrics in scientific research papers and standards.
- 2. **Security Evaluation methodologies:** Existing security evaluation methodologies, in both literature and standards.
- 3. Fundamental features of a comprehensive security assessment methodology: Analysis of the features that an evaluation methodology has to meet to be comprehensive.
- 4. **Testing types and techniques:** Existing testing techniques in the literature and standards.
- 5. **Existing cybersecurity evaluation standards:** Both national and international.
- 6. **Directed Graphs:** Contemporary applications of directed graph in vulnerability analysis.
- 7. **Vulnerability Analysis:** Existing techniques o perform vulnerability analysis.
- 8. **Metric Aggregation:** Existing techniques to aggregate security metrics.

In the following sections, we will review each one in more detail.

# 3.1 Security Metrics

Previously, we described evaluation methodologies as the instructions and steps that evaluators must follow, and testing techniques as the tools that they have to perform each step. Following this example, metrics would be the way for the evaluators to interpret the results obtained.

First, to understand what a security metric is, it is important to distinguish between "metric" and "measurement".

#### **MEASUREMENT**

A measurement is a concrete and objective attribute that provides a single-point-in-time view of a specific and discrete factor [70].

An example of measurement would be the number of spams detected.

#### **METRIC**

A metric is generated from the analysis of the raw data provided by measurements [70].

A possible corresponding metric for the previous measurement could be the number of spams detected last month compared to the number of spam detected during this month.

Measuring security, both qualitatively and quantitatively, is not an easy task [71–73]. It is a long-standing open problem to the research community, and it is of practical importance to software industry today. Suitable metrics are needed to achieve a foundational science to guide system design and development, and to reveal the safety, security, and possible fragility of complex systems. Using metrics allows systems to be compared and evaluated, so we can increase accountability, demonstrate compliance and determine how much our investments in products and processes are making our systems more secure [74,75]. There is clearly a growing interest in adopting and using security metrics, yet a lack of widely accepted solutions.

# 3.1.1 Metrics in the literature

Security cannot be measured as a universal concept due to the complexity, uncertainty, non-stationary, limited observability of operational systems, and malice of attackers [73,76]. For this reason, there is a growing interest in adopting and using security metrics to measure concrete properties of systems [77]. The most exhaustive surveys about security metrics are commented below:

Atzeni *et al.* discussed the importance of security metrics, and summarized security system assessment methods and metrics in [71]. They concluded that measurements and metrics are necessary to justify investments in security and to manage security.

In [78], Pendleton *et al.* provided a comprehensive survey on quantitative system security metrics. In their work, they analyzed software, network, economic, and effectiveness security metrics. To classify the surveyed metrics, a hierarchical ontology was proposed, and then, each metric was briefly described.

In [79], Ramos *et al.* presented a survey that focuses on model-based quantitative security metrics. The authors review the state of the art of network security metrics in depth, more specifically, in the realm of model-based quantitative metrics. They present a complete and thorough review of the main metrics proposals, remarking the pros and the cons of each metric. Finally, the authors review the open issues and suggestions for future research directions.

#### 3.1.2 Metrics in the standards

We reviewed the industrial standards that are more cited in different articles to find whether they encourage the use of metrics during the evaluation process [80].

OSSTMM [81] defines the *rav*, a scale measurement of the attack surface. In this scale, 100 *rav* means perfect balance and a lower value indicates insufficient controls and therefore a greater attack surface. A value greater than 100 *rav* shows more controls than necessary, which might be inefficient and sometimes a problem, since excessive controls often increase complexity and maintenance issues. The *rav* does not measure risk for an attack surface, rather it enables the measurement of it.

The Common Criteria (CC) [50] defines the Evaluation Assurance Levels (EALs) as an increasing scale to describe the rigor and depth of evaluation. CC list seven levels, from EAL1 (the most basic one) to EAL7 (the most stringent security level). It is important to notice that the EAL levels do not measure security itself. Instead, the emphasis is given to "functional testing", confirming the overall security architecture and design, and performing some testing depending on the EAL to be achieved. In other words, the highest level means that more testing was done, but not that the product is necessarily more secure.

The CC suggests that the output documentation should include the numeric values of the metrics used, as well as the actions taken as a result of the measurements, but there is no reference about how to compute these metrics [82].

The National Institute of Standards and Technology (NIST) has done a lot of research on security metrics, and has proposed nine security metrics for three different aspects: (1) implementation, (2) effectiveness/efficiency, and (3) impact [83]. Moreover, NIST proposes a comprehensive taxonomy to classify the metrics presenting three security categories [84,85]: (1) management, (2) technical, and (3) operational; and 17 subcategories. According to NIST, metrics should be used from the moment systems are conceived, from their design, creation to operation. Such metrics help identify security problems early and assist in faster and more efficient management and governance [86,87].

The IEC 62443-1-3 standard [88] includes cybersecurity conformance metrics for managing a security program throughout the life cycle of an industrial automation control system. These metrics provide evidence of conformance for verifiability, completeness, and accuracy. They can also be of assistance to others for assessing the cyber robustness of their control system solutions. This standard also provides a list with the characteristics of a good metric, a metrics development and implementation process, and the steps to be followed for creating metrics. The 62443-4-1 [89] standard also includes some examples of metrics that can be used at the component level, that is, for ES, network components, host software and applications, e.g., functional security, deployment security and current backlog. In Annex A of [88], possible metrics are described to be used in the case of highest maturity level for an organization. Some of these metrics are: results of static code analysis, analysis of attack surface and deviations from Secure Coding Guidelines. They show the effectiveness of the development process with measurable improvements.

# 3.1.3 Current Issues with Security Metrics

There are many recurrent problems associated with security metrics that have been identified in the literature [90]:

- There are hardly any empirical results confirming the predictive power of security indicators. Hardly any security indicator has a solid theoretical foundation. Many indicators seem to correlate only weakly to actual security properties. A proper method for evaluating the correlation between the degree of process compliance and the actual security level is missing in the literature.
- It is essential to explain how the measure relates to the attribute it captures and to define how it is calculated.
- Empirical evidence in support of the claimed correlation is missing.
- Many security metrics lack an adequate description of the measure's scale, unit
  and reference values for comparison and interpreting the results. The applicability criteria for the method as well as its assumptions, possible advantages
  and disadvantages, eventual openly available tool support, and comparison to
  other properties is not always explained.
- Some indicators implicitly or explicitly rely on extrapolation.
- Non-existing mapping of security metrics and measurements to concrete security goals.

Sentilles *et al.* [91] carried out a literature survey using a systematic mapping methodology to analyze the security metrics proposed in the publications until that moment. The main research question was "What percentage of evaluation method descriptions contains pertinent information to facilitate its use?". They wanted to know (1) what proportion of properties were explicitly defined, (2) what proportion of evaluation methods provided explicitly key information to enable their use (*e.g.*, formulas) and (3) what proportion of other supporting elements of the evaluation methods was explicit (*e.g.*, available implementations). The authors identified the possible gaps concerning security metrics. Only half of the papers clearly specify the assumptions and hypotheses that are assumed to hold when using the method. Despite being

often mentioned in the papers as part of the solution being implemented, in practice, only a few implementations or programs are available to directly support the method evaluation. Sentilles *et al.* could find mentions of available tool support in only 12.5% papers, and only one performed some kind of benchmarking or comparison to similar properties or evaluation methods. This corresponds to only 6 % of all the papers. The authors concluded to a large extent that the information provided in the papers is insufficient to directly apply the method, especially by non-experts.

Even though metrics and methodologies have been proposed in the literature, it is difficult to know about them and understand whether they are applicable in a given context and if they are, how to use them. As Sentilles *et al.* concluded in [91]: "Despite the abundance of metrics found in the literature, those available give us an incomplete, disjointed, hazy view of software security." According with that, it seem reasonable that the future research should be focused on the development of a convincing theoretical foundation, empirical evaluation, and systematic improvement of existing approaches. It should not be focused on the development of more security metrics or approaches.

# 3.1.4 Granularity vs. Overall Security

As Pendleton *et al.* commented in [78], there is an issue that is far from being trivial when determining whether to aggregate multiple metrics into one or not. If multiple metrics are aggregated into a single value, granularity is lost. On the other hand, if every metric is kept as a single value, reporting details of attack, impact or countermeasures would be endless. This is known as the challenge in "measurability." The authors identify the advantages and the disadvantages of both options, and different proposed alternatives for each one.

#### 3.1.5 Scales of Metrics

Attributes of an object should be measured by certain scales [92]. The following classification has been used in various previous works [78,93]:

- **Nominal:** also known as *categorical scale*, they are represented by a discrete unordered set. Measurement in a nominal scale can be compared using standard set operations. Nominal scales can be transformed between each other via some methods for one-to-one mapping.
- Ordinal: this scale is a discrete ordered set, in which we can use operators such as "greater than" to compare two measurements on the same ordinal scale. Ordinal scales can be transformed between each other based on monotonic (i.e., order-preserving) one-to-one mapping methods. Ordinal scales are one step higher (i.e., more accurate in the measurement) than nominal scales in the hierarchy of scales because it can downgrade to nominal scale at the cost of information loss.
- Interval: this scale extends the ordinal scale such that the distance between
  adjacent points on the scale is constant. This equal-distance property allows
  the difference operator by which one can compare the difference between two
  points and the difference between another two points. One can define linear
  transformations between two interval scales, which preserve constant distance

between two points on the same scale. Interval scales can be analyzed via statistical analysis, such as mean, variance, and even higher moments. Interval scales can be downgraded to ordinal (and therefore nominal) scales by relaxing the equal-distance condition on the scale.

- Ratio: this scale extends the interval scale by defining the origin (i.e., value of 0) in a natural fashion. Ratio scales can be operated with multiplication, logarithm, and finding roots. Transformations can be defined between ratio scales.
- **Absolute:** this measures things measurable only in a single method. Counting is a simple example of absolute scale. Any kind of statistical analysis can be applied to measurements in the absolute scale.
- **Distribution:** this scale is used to measure an attribute with inherent uncertainty.

As it is said in [78], "It is worth noting that using a different scale gives a different level of accuracy, leading to a different degree of understanding in a given value as a metric." A higher level of scale in the hierarchy gives higher accuracy that can be easily transformed to a low level of scale. Measurements using a higher level of scale permits richer statistical inference.

# 3.1.6 Features of a Comprehensive Metric

Metrics are important for quantifying the security level achieved; however, computing metrics can be exhausting and time-consuming resulting in not being practical. Furthermore, it is not resource-effective to measure everything, thus using the appropriate metrics is critical.

Some authors have tried to characterize which criteria should meet a comprehensive metric. Among their efforts, three main ideas can be found in the literature: (1) the SMART criterion [94], (2) the PRAGMATIC criterion [95], and (3) the characteristics identified in the work carried out by Savola [77].

Since Doran proposed the "SMART" (Specific, Measurable, Attainable, Repeatable, Time-dependent) criteria in [94] (see Table 3.1), it has been widely used, developed and extended, and it has been also adapted for security metrics development [96]. Although several interpretations have arisen about what the letters in SMART actually mean, it is important to note that, while the exact definition of each letter is not critical, it should be meaningful and not redundant in a larger context. On the other hand, Krag *et al.* [95] propose a set of nine criteria for assessing and selecting metrics, using the acronym "PRAGMATIC" (Predictive, Relevant, Actionable, Genuine, Meaningful, Accurate, Timely, Independent, Cheap) (see Table 3.2).

According to Jaquith [97], a good metric should be consistently measured, without subjective criteria; cheap to gather, preferably in an automated way; expressed as a cardinal number or percentage and not with quality labels; expressed using at least one unit of measure; and should be contextually specific, so it is relevant enough to take action based on their values.

In their survey [78], Pendleton *et al.* discuss three main perspectives towards a good metric: (1) the conceptual perspective, that says that a good metric should be intuitively easy to understand and with high usability; (2) the measurement

SMART	Definition
Specific	Well-defined, unambiguous wording
Measurable	Quantitative when feasible
Attainable	Within budgetary and technical limitations
Repeatable	Do not vary depending on the person taking them
Time-dependent	Multiple time slices

TABLE 3.1: Definition for each letter of the SMART criterion.

TABLE 3.2: Definition for each letter of the PRAGMATIC criterion.

PRAGMATIC	Definition
Predictive	Predicts security outcomes
Relevant	Useful information concerning aspects that we truly care about
Actionable	Less lag between measurement and reaction
Genuine	Unambiguous, straightforward, credible, real, and objective
Meaningful	Meaningful to the intended recipients, audiences, or users of the metrics
Accurate	Accuracy and precision
Timely	Minimising the delay between collecting, analysing, and presenting the data
Independent	Truthful, honest, credible, objective representation, and resistant to manipulation
Cheap	Cost of collecting, analysing, and using metrics

perspective (as done in [97]); and (3) the utility perspective, that claims that a metric should allow both horizontal and temporal comparison between enterprise systems.

In [77,98], Savola analyses the literature to identify the quality criteria and dimensions of security metrics and measurement process. In [98], the pre-existing goodness criteria are listed. Afterwards, in [77], Savola takes research one step further by conducting expert surveys and interviews to 141 security experts from 21 different countries, and extracting 19 quality criteria (see Table 3.3). According to his results, correctness, measurability, and meaningfulness are the main quality criteria along with usability.

The quality criteria that Savola proposes are intended to support human decision-making in every step of the life cycle, e.g., security engineering activities at design and development practices or in security management activities at run time. In his work, Savola concludes that there is no security approach that would allow the measurement of security as a universal property, since there is no security metric that could fulfill all the quality criteria, so security metrics cannot be used to measure security as a whole. He finally points out that there is a need for widely-accepted approaches for security measuring.

#### 3.1.7 Taxonomies for security metrics

In [93,99], Savola reviews the most common taxonomies for security metrics, and also provides a new model to taxonomize security metrics for technical systems to systematize and organize their development activities. The most common classification for metrics divides them as organizational, technical and operational [93,100].

 Organizational security metrics: They evaluate the security level of a whole organization. They describe and track how effectively organizational programs

Dimension	Explanation
Correctness	The SM are correctly implemented and error-free.
Cwanularitu	The SM allow the measurement results that differ from each other
Granularity	to be distinguished at an adequate level.
Objectivity	Results are not influenced by the measurer's will, beliefs, or feelings.
Unbiasedness	Results are not influenced by any bias.
Controllability	Results should be kept within the defined limits or the measurement
Controllability	window.
Time dependability	The time-dependent behavior of SM can be leading, coincident, or lagging.
-	The time dependability of the SM should be part of their specification.
Comparability	Support comparison of the targets that they represent.
Measurability	Capable of having dimensions, quantity or capacity ascertained in the SuI.
Attainability	Measurement results can be achieved from the SuI
Availability	Results are generally available
Easiness	How easy it is to achieve the measurement results.
Reproducibility,	The same results are achieved if a measurement is repeated in the same
repeatability	context, with exactly the same conditions.
Scale reliability	Reproducibility of the measurement by different persons.
Cost effectiveness	The measurement gathering and measurement approaches should be cost
	effective.
Scalability	The SM should be applicable to SuIs of different sizes
Portability	Applicability of the SM to various target systems.
Non-intrusiveness	The measurements should not harm the normal operation of the SuI, require only
	minimum changes to the SuI, and not affect the measurement results.
Meaningfulness	The SM should be relevant and respond to the needs.
Effectiveness	The expectations of the adequacy of the SM sufficiency in the final use environment
	are satisfied.
Efficiency	The adequate requirements of the SM are achieved with only minimal undesired use
	of effort and time.
Representativeness,	The SM correspond to the actual system characteristics in
contextual specificity	the SuI in a contextually focused way.
Clarity	The SM are clearly formulized.
Succinctness	Only important parameters are considered.
Ability to	The SM should be able to show progression on the
show progression	dimension it addresses.
Completeness	The collection of SM should be complete from a
Completeness	measurement objectives perspective.

TABLE 3.3: Desired quality criteria of Security Metrics (SM) identified in [77,98].

and processes achieve cybersecurity [74,100–103]. They do not just include all the systems, but could also take into account all the people involved for the enterprise to work. Organizations can use measures and metrics to set goals (or benchmarks), and determine success or failure against the benchmarks. Some more references on this topic can be found in [95,104–106].

- **Technical security metrics:** They indicate the level of security a specific system exhibits. They describe and compare technical objects, *e.g.*, algorithms, specifications and architectures, and to indicate the level of security a specific system exhibits. Since the definition of what a "system" is covers a broad spectrum, technical security metrics also cover a wide range of systems, going from network security metrics to embedded devices security metrics.
- Operational security metrics: They are describe and manage the risk to an operational environment, including as-used systems and operating practices [100]. They include measures of operational readiness or security postures, measures used in risk management, metrics describing threat environment, and metrics

supporting incident response and vulnerability analysis. They could also include metrics produced as part of normal operations that can serve as input to other security metrics. The estimation of operational metrics generally requires experimental or empirical measurements.

# 3.1.8 Gaps Detected and Critics to the Current Status

Security metrics are a key piece to obtain an objective result in every evaluation process. Metrics are developed to measure concrete properties in very specific contexts. This means that they are not usually adaptable to a different target. As happens with security evaluation methodologies and security standards, there is also a bunch of security metrics to choose from. They can be used to evaluate almost every single aspect of a system or an organization, and they can be found in many forms. Under this scenario, it is not difficult to realize that choosing the right metrics is also an overwhelming task.

On the other hand, when talking about metrics, generalization is not always possible. This means that each dimension of the ES has associated a set of metrics that is the most suitable for evaluating that concrete aspect of the device. So having a vast amount of metrics is useless if there is no method to choose among them each time a security evaluation has to be made. It is as important having metrics to choose from, as it is to have a method to choose the right metric for each case, and how to apply each one.

# 3.2 Security Evaluation Methodologies and Standards

There are many security evaluation methodologies available in the literature. In this section, we review the most cited security evaluation methodologies surveyed by ENISA [107]. Table 3.4 shows them in alphabetical order.

Methods	Last version	Year	Status	Standard Type	Metric proposal	Scopeof the evaluation	Documentation accesibility
			Up to date     Deprecated	1. International		1. Organization 2. Network 3. Device 4. Hardware 5. Software 6. Web 7. User 8. Data	1. Free 2. Payment 3. Unaccesible
CC [108]	3.1	2017					
CPA [109,110]	1.4	2018					
CSPN [111]	3.0	2021					
ICCF [112]	1.0	2018					
ISA/IEC 62443 [48]	1.0	2018					
ISSAF [113,114]	0.2.1	2016					
NESCOR [115]	3	2015					
NITES [116]	2	2009					
NIST [117]	2	2008					
OSSTMM [118]	3	2010					
OWASP [119]	4.2	2014					
PTES [120, 121]	1	2014			□■		

TABLE 3.4: Comparison of security evaluation methodologies.

#### 3.2.1 Common Criteria (CC)

The Common Criteria for Information Technology Security Evaluation (Common Criteria or CC) is an international standard (ISO/IEC 15408) for computer security certification [108,122]. CC is a framework which provides assurance that the process of specification, implementation, and evaluation of a computer security product has been conducted in a rigorous, standard and repeatable manner [50,82]. CC defines seven levels of security or Evaluation Assurance Level (EAL) of a product or system, from EAL 1 (the lowest one) to EAL 7 (the highest one). EAL 7 indicates the product satisfies the most demanding development process. The CC Certification is recognized by 30 countries around the world that have signed the Common Criteria Recognition Agreement (CCRA).

CC has become so popular that it has served as the foundation of a vast amount of other methodologies, and standards all over the world. Nevertheless, CC exhibits some drawback [123,124]:

- The evaluation is a costly process and the vendor's return on that investment is not necessarily a more secure product.
- The evaluation focuses primarily on assessing the evaluation documentation, not on the actual security, technical correctness or merits of the product itself.

- The effort and time necessary to prepare evaluation evidence and other evaluationrelated documentation is so cumbersome that by the time the work is completed, the product in evaluation is generally obsolete.
- CC recognizes the need to limit the scope of evaluation, therefore, evaluation activities are only performed to a certain depth, use of time, and resources.
- Certification schemes are different depending on the country, and they can lead to different assessment results.

#### 3.2.2 CPA

The Commercial Product Assurance (CPA) scheme is a UK scheme to assess commercial products [109,110]. The certification scheme is governed by the National Cyber Security Center (NCSC). The CPA is based on the Common Criteria.

The evaluation is mainly black box because it does not require confidential information. The products are tested against the so-called security features defined by the CPA, and the security features depend on the product categories (e.g., encryption of data, software execution control, remote desktop, firewall, real-time secure communication, virtualization platforms, or smart meters).

Nowadays, the CPA only certifies smart meters or smart metering products. Certification is usually valid for six years.

#### 3.2.3 CSPN

The Certification Sécuritaire de Premier Niveau (CSPN) is a French IT security certification scheme established in 2008 [111]. Created by ANSSI in 2008, the CSPN's main objective is to offer a faster and cheaper alternative for IT security certification compared to Common Criteria. CSPN assessments follow a black box testing approach within a fixed and limited time frame. Therefore, it is less rigorous and targeted to low threat or risk environments. The CSPN covers mostly IT security products such as antivirus, firewall access control mechanisms and secure storage devices. It can also be applied to hardware.

# 3.2.4 IACS Cybersecurity Certification Framework (ICCF)

The IACS components Certification Framework (ICCF) is a certification scheme developed by the European Reference Network for Critical Infrastructure Protection (ERNCIP) [112]. ICCF focuses on the security of IACS components, rather than on subsystems or even entire IACS.

The ICCF scheme takes particular account of the work done on the ISA/EC 62443 standards. For example, IEC 62443-4-2 was adopted as the basic set of cybersecurity requirements. In addition, the proposed assessment scheme is heavily influenced by the Common Criteria.

# 3.2.5 IEC 62443-4-1

Document IEC 62443-4-1 serves as a reference for the definition of a life cycle for the development of industrial cybersecure components. These components must comply with a series of technical requirements for cybersecurity that are defined in document IEC 62443-4-2 and which, after implementation, would result in product instances configured to be integrated within a complete industrial solution (system). System-level integration should comply with IEC 62443-2-4 and IEC 62443-3-2, which should also consider the technical requirements at the system level defined in IEC 62443-3-3.

IEC 62443-4-1 standard describes the requirements that apply in the life cycle of a product or industrial component. IEC 62443-4-1 classifies the requirements into the following practices:

- 1. **Practice 1:** Security management.
- 2. **Practice 2:** Specification of security requirements.
- 3. Practice 3: Secure by design.
- 4. Practice 4: Secure implementation.
- 5. **Practice 5:** Security verification and validation testing.
- 6. Practice 6: Management of security-related issues.
- 7. **Practice 7:** Security update management.
- 8. **Practice 8:** Security guidelines.

The standard also tries to establish how to perform the concrete tests and how to document them, however, this process is not clear. For example, the following table shows the different tests to be performed in an IACS, however, these are defined in an ambiguous way and there is no relationship between testing and the required Security Level (SL).

Security Levels (SL) are described in IEC 62443-4-2, and for each component must decide what level of protection is desired for each of the fundamental protection requirements. The following table shows the relationship between the SL and the type of attacker, with SL 1 being classified as protection against accidental errors and SL 4 as protection against intentional attacks with sophisticated and highly motivated means.

# 3.2.6 Information Systems Security Assessment Framework (ISSAF)

It is a well-established penetration testing methodology developed by the Open Information Systems Security Group [125]. ISSAF provides comprehensive step-by-step penetration testing technical guidance. It is designed to evaluate the security of networks, systems and application controls.

The methodology outlines three well-defined action areas, and details the nine steps composing the main one (see Figure 3.1) as follows [113,114]:

1. **Planning and Preparation:** The first phase encompasses the steps needed to set the testing environment up, such as: planning and preparing test tools, contracts and legal protection, definition of the engagement team, deadlines, requirements and structure of the final reports.

- 2. **Assessment:** This phase is the core of the methodology, where the real penetration tests are carried out. The assessment phase is articulated in the following activities: (1) Information Gathering, (2) Network Mapping, (3) Vulnerability Identification, (4) Penetration, (5) Gaining Access and Privilege Escalation, (6) Enumerating Further, (7) Compromise Remote Users Sites, (8) Maintaining Access, (9) Covering Tracks.
- 3. **Reporting, Clean-up and Destroy Artifacts:** During this phase, at the very end of the active parts of the methodology, testers have to write a complete report and to destroy artifacts built during the Assessment phase.

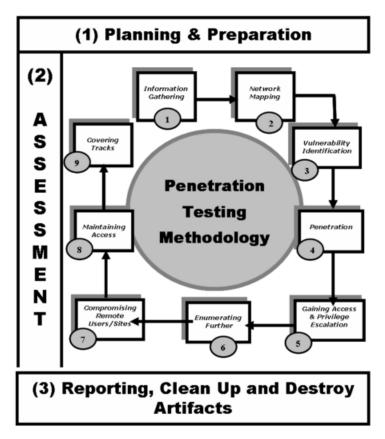


FIGURE 3.1: ISSAF workflow [113].

ISSAF has a clear and very intuitive structure, which guides the tester through the complicated assessment steps. The order in which the methodology describes the penetration testing process is optimized to help the tester perform a complete and correct penetration testing, avoiding the mistakes commonly associated with randomly selected attack strategies.

On the negative side, the reporting section is poorly implemented. There is not a well-defined and accurate guideline to develop final reports, and some suggestions are outdated. For example, the methodology suggests destroying and cleaning up the developed artifacts, while the most current practice is to leave the developed artifacts on the tested system, to ease further and deeper analysis [126].

# 3.2.7 NESCOR Guide to Penetration Testing for Electric Utilities

The National Electric Sector Cybersecurity Organization Resource (NESCOR) [127] has created a test template to provide network operators with a guide to help them perform penetration testing of smart grids [115]. It can be used for advanced metering infrastructures, distributed energy resources, and Electric Transportation [128].

The methodology provides for a linear sequence of tests and is divided into six phases with nine work packages, as Figure 3.2 shows.

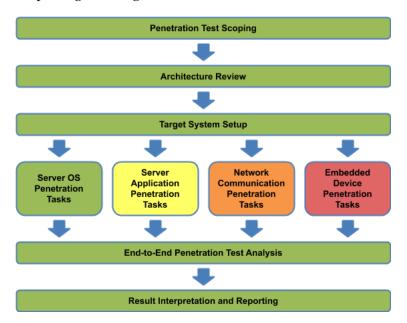


FIGURE 3.2: NESCOR workflow.

This methodology is based on the assumption that testers will initially gain a comprehensive overview of the network to be investigated by means of expert interviews, among other things.

#### **3.2.8 NITES**

The Singaporean National IT Evaluation Scheme (NITES) was established by the Security Accreditation Committee (SAC) [116]. NITES is an adaptation of Common Criteria v3.1, equivalent to the EAL4+ package.

There are four categories of products that can be assessed under NITES:

- 1. Secure portable storage.
- 2. Network related devices (i.e. VPN).
- 3. File/folder encryption.
- 4. Key management solutions (i.e. HSM).

As the SAC was accredited by Common Criteria as an issuing certification body in January 2019, the NITES scheme has lost relevance. The Common Criteria certification scheme is called the Singapore Common Criteria Scheme (SCCS).

# 3.2.9 NIST Guidelines on Network Security Testing (GNST)

The Guideline on Network Security Testing (GNST) [117] was the first methodology that introduced a formal process for reporting and to take advantage of inducted hypotheses [126]. GNST was developed by National Institute of Standards and Technology (NIST) [61].

GNST follows four main steps (Figure 3.3):

- **Planning:** The system is analyzed to find out the most interesting test targets.
- **Discovery:** The tester searches the system, looking for vulnerabilities.
- Attack: The tester verifies whether the found vulnerabilities can be exploited.
- **Reporting:** In the last step, every result is reported.

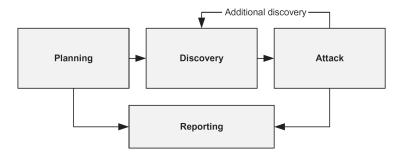


FIGURE 3.3: GNST workflow.

GNST was updated by Technical Guide to Information Security Testing and Assessment in 2008 [129].

#### 3.2.10 Open Source Security Testing Methodology Manual (OSSTMM)

The Open Source Security Testing Methodology Manual (OSSTMM) was developed by the Institute for Security and Open Methodologies (ISECOM) in January 2001 [130]. It was the first methodology to include human factors in the tests, taking into account the established fact that humans may be very dangerous for the system. Capturing the human factors such as insider attacks and social engineering attacks is a valuable feature, which no one else have never tried to capture.

OSSTMM is the *de facto* standard for security testers [126] to find vulnerabilities. It describes a complete testing methodology, offering comprehensive tools to report the result set. It is designed to test the operational security of physical locations, workflows, human security testing, physical security testing, wireless security testing, telecommunication security testing, data networks security testing and compliance [118].

Although it is a great methodology, OSSTMM fails to deliver some key concepts, namely [126]:

Control flow analysis: It is very biased towards communication analysis, forcing the tester to put a lot of effort on data flow, while control flow and application analysis is, comparatively, neglected.

- **Readable diagrams:** OSSTMM does not provide intuitive diagrams, mainly because of the high number of items and of many messy loops inside the flow.
- data loss process on writing reports and traceable reports: It offers nice templates to fill up the reports, but unfortunately these templates follow a strictly linear reading process. A good practice would be writing reports after each action; otherwise, many details easily fall into oblivion.

# 3.2.11 Open Web Application Security Project Testing Guide (OWASP)

The OWASP Testing Project [119] is a methodology developed by the Open Web Application Security Project (OWASP) to help people understand the what, why, when, where, and how of testing web applications [131,132]. Figure 3.4 shows the workflow of the OWASP Testing Guide.

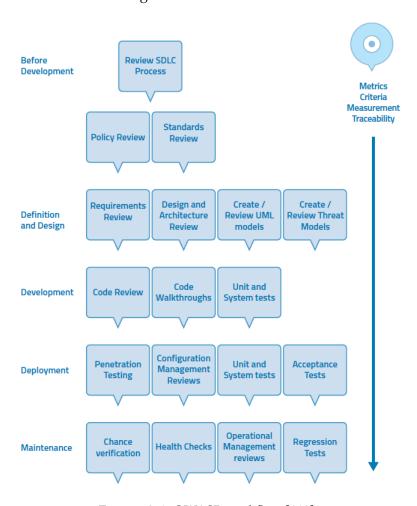


FIGURE 3.4: OWASP workflow [119].

# 3.2.12 Penetration Testing Execution Standard (PTES)

The Penetration Testing Execution Standard (PTES) consists of seven main sections [120]. These cover everything related to a penetration test. The main sections (Figure 3.5) defined by the standard as the basis for penetration testing execution are:

(1) Pre-engagement Interactions, (2) Intelligence Gathering, (3) Threat Modelling, (4) Vulnerability Analysis, (5) Exploitation, (6) Post Exploitation, (7) Reporting.

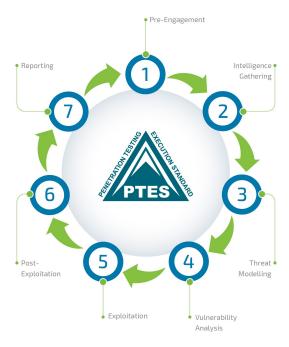


FIGURE 3.5: PTES workflow [133].

As the standard does not provide any technical guidelines as far as how to execute an actual pentest, a technical guide to accompany the standard itself was proposed [121].

# 3.2.13 Fundamental Features of a Comprehensive Evaluation Methodology

We have already reviewed the most relevant security evaluation methodologies for the industry in the previous section, but we lack a means of comparing them with each other. For that reason, we also reviewed the literature to find which are the fundamental features that a comprehensive evaluation methodology should have. In this section, we analyze them, and apply them to compare the security evaluation methodologies described in Section 3.2.

In [126], the fundamental features of a comprehensive methodology were identified:

- Modelling: The methodology should explicitly define the key concepts in order
  to facilitate the tester in modelling both the system and the evaluation process,
  by removing potential ambiguities and leading the tester to the kind of model
  that better suits the subsequent activities.
- 2. **Planning:** The methodology should support the tester in laying detailed test plans out, *i.e.*, definition of phases, prerequisites for each phase, tools to use in each phase, expected outcomes. During planning, the methodology should help the tester in selecting the most appropriate and useful activities for the specific system.
- 3. **Flexibility:** While statically defining a test plan is an important step, an even more powerful feature would be providing structured means of dynamically

integrating additions (deriving from the results that are acquired at each step) in the initially defined plan, leading to richer, or more specific, new plans. Without this possibility, the tester would be strongly exposed to the risk of defining either too general or too specific targets; in the former case, the procedure to achieve them would lack the right amount of detail to ensure its correctness; in the latter, significant alternatives would possibly be excluded a priori.

- 4. **Adaptation:** The concepts and models defined within a methodology should certainly be unambiguous, but this quality should not hinder the possibility to adapt them to many variations of the real systems to be tested.
- 5. Guidance: Given the huge amount of different aspects involved in security evaluation, the methodology should offer practical guidance about what activities compose a testing session, and which tasks are needed before, during, and after each activity, for example by means of up-to-date checklists related to the vulnerabilities and the most effective testing procedures for different testing contexts.
- 6. Reporting: Guidance should not be limited to the active phases of testing, but also extend to the documentation of every useful information related to the test setup, environment, progress and results. Supporting the tester in the reporting activity means not only helping him not to omit important details, but also letting him format the information in one or more ways that are suitable for different kinds of readers (technicians, policy-makers, managers, etc.).
- 7. Granularity: Finding good guidance and reporting features in a methodology, commonly means having lots of highly-detailed information at hand. However, capturing the details only where needed, while not uselessly encumbering the testing and reporting activities, is equally important. This criterion applies both to data collection and to task planning. With regard to the former, the methodology should not force filling out detailed reports about low-severity, low-priority or low-probability scenarios. With regard to the latter, the methodology should cater for the easy selection of sensible steps and the provision for skipping the useless ones, possibly foreseeing nested levels of planned tasks.

Table 3.5 shows an at-a-glance comparison of the methodologies previously discussed with respect to the properties defined above.

TABLE 3.5: Feature map of the security evaluation methodologies (+ good coverage, = average coverage, - limited or no coverage) [126,128, 134].

Features	OSSTMM	ISSAF	CEM	PTES	OWASP	NESCOR	GNST
Modeling	=	+	-	=	+	+	=
Planning	+	+	=	+	+	+	=
Flexibility	-	-	=	+	-	-	+
Adaptation	+	=	-	+	=	-	=
Guidance	=	=	=	=		+	+
Reporting	-	-	-	=	=	=	-
Granularity	=	+	-	=		=	=

The main issue with this scheme is that it is developed from a theoretical point of view, and it does not consider the cost of the evaluation process (both time and money). This is a key aspect, because it is the one that best reflects the reality of the evaluation process: a security evaluation methodology could theoretically be perfect in all aspects, but if it takes too long, or if it is too expensive to obtain a result, the methodology will be useless. This is why a last factor called "applicability", or "cost" should be introduced to balance the mix, and to adapt it to the industry.

# 3.2.14 Gaps Detected and Critics to the Current Status

Each analyzed methodology was devised with a different goal, and a different environment in mind. There are so many standards and methodologies that choosing one for any application can be overwhelming. So it is very common to find a bunch of standards for the same application, with different requirements. Moreover, they differ from country to country, and international standards are not always used. This scenario shows that it is not an easy task to develop a security evaluation methodology that can be applied to every single device, case and situation.

Taking a step forward in the analysis, the majority of the analyzed methodologies are conceived for a high-level system, namely network or organization level. This is because of the high demand of the sector, where every organization wants to be sure that they are protected both from competitors, and possible attackers. This encourages most efforts to focus on developing robust and comprehensive methodologies to protect business. As a consequence, other scenarios, which also need to be evaluated to know their security level, are neglected, and the methodologies available are more scarce, as is the case with ESs.

The security evaluation methodologies shown here can be classified according to their drawbacks when trying to apply them to an ES. Three scenarios can be differentiated:

- 1. **Application is possible with modifications:** This means that the general structure and content of the methodology have potential to assess the security level of ESs, but because of the nature of ESs, the methodology has to be adapted in order to be used. This is the case of OSSTMM, that was conceived for organizations, but it is so systematic that can be extrapolated to ESs.
- 2. Application is possible, but the evaluation takes a long time (with or without modification): This means that the evaluation process takes too long, and depending on the product, it could be obsolete by the time the evaluation is done. Sometimes the evaluation process can be expensive not only in time but also in money. This is the case of CEM of Common Criteria, where the evaluations time can take somewhat between six months up to one year.
- 3. Impossibility to apply the methodology: This means that there is no possible way to alter or modify the methodology to apply it to an ES. This is usually cased by very specialized methodologies developed for very concrete scenarios. This is the case of the OWASP testing methodology that is oriented to web applications.

As can be seen, a lot of work has been done at the organization level, but to the authors' knowledge, there is no security evaluation methodology that is specifically design to assess the security level of ESs.

Finally, it is very common to find recommendations of using metrics in the security evaluation process, and the advantages that they introduce, but they are not usually defined in any way. Analyzed methodologies let the evaluators choose their own set of metrics for each target of evaluation, increasing the subjectivity of the process.

# 3.3 Testing Types and Techniques in the Literature

Security evaluation methodologies, and in general, any methodology or standard, set the foundations for a given process. They standardize and list the steps to follow, and how to report the results. They are developed in this way in order to cover the largest number of cases in a specific domain. But it is the case that most of the time, standards specify what has to be done, but not how it should be done. If standards are like the instructions that the evaluators have to follow, testing techniques are the tools that the evaluator has to carry out the tasks described in the standards.

In this section, we define what is security testing, and we reviewed the literature to find the most relevant testing types and techniques used today. We classify them according to when they can be applied in the development life cycle.

# 3.3.1 Security Testing

Security testing is one type of assessment method used for determining whether the information system being assessed meets the security objectives (confidentiality, integrity, and availability). Security testing, as defined by the National Institute of Standards and Technology (NIST), is the "process of exercising one or more assessment objects under specified conditions to compare actual with expected behavior, the results of which are used to support the determination of security control existence, functionality, correctness, completeness, and potential for improvement over time" [135] (see Figure 3.6). Security testing can be used to determine if technical mechanisms or organizational activities meet a set of predefined specifications, or it can be used to evaluate whether security controls can be circumvented.

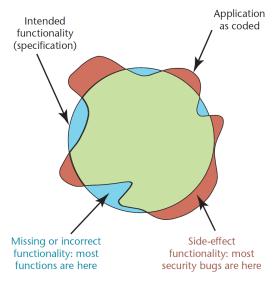


FIGURE 3.6: Intended versus implemented software behavior in applications. The circle represents the software's intended behavior as defined by the specification. The amorphous shape superimposed on it represents the application's true behavior. Most security bugs lay in the areas of the figure beyond the circle, as side effects of normal application functionality [136].

Security testing can be conducted using manual or automated techniques. Both testing techniques have advantages and disadvantages. Manual security testing is not only time-consuming; it can also be more expensive than automated testing, and qualified testers are required. Automated testing uses tools to execute tests and automate tasks, thereby improving the efficiency of security testing. Automated tools can also improve the time to complete security testing, and can be more frequently updated based on new types of vulnerabilities and exploits. However, automated security testing reduces the types of testing techniques (e.g., social engineering) that can be performed based on the limitation of automated tools to mimic human behavior. Therefore, both manually and automated testing should be considered against the depth and coverage required for the security testing [137].

It is essential to take testing into account in all phases of the secure software development life cycle [138]. Thus, security testing must be holistic, covering the whole secure software development life cycle. In concrete terms, Fig. 3.7 shows a recommended distribution of static and dynamic testing efforts among the phases of secure software development life cycle [139].

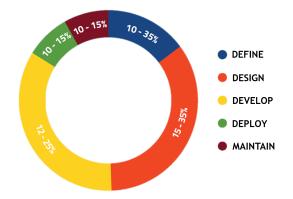


FIGURE 3.7: Proportion of test effort in secure software development life cycle [139].

### 3.3.2 Security Testing Techniques

In [140], a classification according to the test basis within the secure software development life cycle (Fig.3.8) is presented:

- Model-based security testing (MBST): It is a Model-based testing (MBT) approach that validates software system requirements related to security properties. It combines security properties like confidentiality, integrity, availability, authentication, authorization, and non-repudiation with a model of the System Under Test (SUT) and identifies whether the specified or intended security features hold in the model.
- 2. **Code-based testing and static analysis:** Many vulnerabilities can only be detected by looking at the code. Static analysis of the program code is an important part of any security development process as it allows detecting vulnerabilities at early stages of the development life cycle, where fixing vulnerabilities is comparatively cheap. This is also known as white-box testing.

- 3. **Penetration testing and dynamic analysis:** These testing techniques do not require access to the source code or other development artifacts of the application under test. Instead, the security test is conducted via interaction with the running software. This is also known as black-box testing.
- 4. **Security regression testing:** It is a combination of regression and security testing which ensures that changes made to a system do not harm its security and do not cause unintended effects, as unchanged parts and changed parts of the software behave as intended.

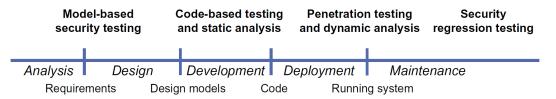


FIGURE 3.8: Security testing techniques in the secure software development life cycle [140].

Table 3.6 describes the testing techniques that are associated to each step in the development life cycle [81,117,140–142].

TABLE 3.6: Testing techniques according to the step in the development life cycle step.

Software Development Life-Cycle Step	lesting lechnique	Description
Model-Based Security Testing		Selected algorithms generate test cases automatically from a set of models of the system under test or of its environment.
Code-Based Testing and Static Analysis	Domain Testing	Only valid input.
	Error-Handling Testing	Properly process erroneous transactions.
	Log Review	Operate according to policies.
	Loop Testing	Exercise program loops.
	Mutation Testing	Using artificial defects, keeping them mini-
		mal.
	Path Testing	Satisfy coverage criteria for each logical path
		through the program.
	Recovery Testing	How well a system recovers from crashes,
		hardware failures, or other catastrophic prob-
		lems.
	Statement Testing	Each statement in a program is executed at
		least once during program testing.
	Static Analysis	Sanity of the code, algorithm, or document.
		The software is not actually executed.
	Storage Testing	Store data files in the correct directories. Re-
		serve sufficient space to prevent unexpected
		termination resulting from lack of space.
Penetrations Testing and Dynamic Analysis	Combinatorial Software Testing	Identify and test all unique combinations of
		software inputs.
	Denial of Service Testing (DoS)	Make the system unavailable by overloading
		it, or exploiting an existing vulnerability.

TABLE 3.6: Testing techniques according to the step in the development life cycle step.

Software Development Life-Cycle Step	Testing Technique	Description
	Docture Hother	The desired of CITT's characterial accurations
	Destructive resumb	Onderstand a 501 s structural periorniance
		or material behaviour under ditterent loads.
		Carried out to the SUT's failure.
	Endurance Testing	Check for memory leaks or other problems
		that may occur with prolonged execution.
	Fault Injection	Physically exploit an existing vulnerability
		through voltage, or temperature variations or
		laser light exposure.
	Functional Testing	Check test cases based on the specifications
		of the software SUT.
	Fuzzing	Sumit random, malformed data as inputs into
		software programs to determine if they will
		crash.
	Invasive Attack	The aim is to capture information stored in
		memory areas, or data flowing through the
		data bus, registers. Such techniques are also
		used to disconnect circuits, to override sen-
		sors, or to defeat blown fuse links. The physi-
		cal properties of the SUT are irreversibly mod-
		ified.
	Network Scanning	Identify all host potentially connected to an
		organization's network services operating on
		those hosts, and the specific application run-
		ning the identified service.
	Password Cracking	Identify weak passwords.

TABLE 3.6: Testing techniques according to the step in the development life cycle step.

Software Development Life-Cycle Step	Testing Technique	Description
	Penetration Testing	Test from the outside in a setup that is com-
		parable to an actual attack from a malicious
		third party. Limited information.
	Security Scanning	Identify network and system weaknesses, and
		later provide solutions for reducing these
		risks.
	Side-channel Attack	Use existence of physically observable phe-
		nomenons caused by the execution of com-
		puting tasks in present microelectronic de-
		vices (power consumption, electromagnetic
		emision or time) to extract information.
	Test Vector Leakage Assessment (TVLA)	Know whether the SUT leakages informa-
		tion through any side channel. TVLA can-
		not determine the magnitude of the leakage o
		whether it is exploitable.
	Social Engineering Testing	Querying personnel.
	Vulnerability Scanning / Testing	Identify hosts and host attributes (operat-
		ing systems, applications, open ports), but it
		also attentifies to identify vunterabilities rather than relying on human interpretation of the
		scanning results.
Security Regression Testing	Age testing	Evaluate the SUT'sability to perform in the future.
	Comparison Testing	compares the product strengths and weak-
		nesses with previous versions or other similar products.

TABLE 3.6: Testing techniques according to the step in the development life cycle step.

		•
Software Development Life-Cycle Step	Testing Technique	Description
	Parallel Testing	Check that the new application which has replaced its older version has been installed and is running correctly.
Others	A/B Testing	Determine whether a proposed change is more effective than the current approach.
	Compatibility Testing Compliance Testing	Check for backward compatibility.  Compliance with internal or external stan-
	E : : : : : : : : : : : : : : : : : : :	dards
	Comiguration resumb	hardware and software, and the effect of
		adding or modifying resources such as memory, disk drives and CPU.
	Conformance Testing	SUT conforms to the specification on which it
		is based.
	Dependency Testing	Examine an application's requirements for
		pre-existing software, initial states and configuration in order to maintain proper function-
		ality.
	File Integrity Checkers	Check whether the SUT computes and stores
		a checksum for every guarded file and estab-
		lishes a database of file checksums.
	Formal verification Testing	Proving or disproving the correctness of in-
		tended algorithms underlying a system with
		respect to a certain formal specification or property using formal mathods of mathemat-
		property, using rothinal menious of maniemat-

TABLE 3.6: Testing techniques according to the step in the development life cycle step.

Software Development Life-Cycle Step	Testing Technique	Description
	Posture Assessment	Security scanning, Ethical Hacking and Risk
		ture of an organization.
	Services Probing	Active examination of the application listen-
		ing behind the service.
	Smoke Testing	Preliminary testing to reveal simple fail-
		ures severe enough to, for example, reject a
		prospective software release.
	SQL Injection	Database errors
	Stability Testing	Check the quality and how the SUT behaves
		in different environmental parameters like
		temperature, voltage etc.
	Stress Testing	Evaluate the SUT at or beyond the limits of
		its specified requirements.

#### 3.3.3 Test Case Prioritization

With the increasing size of software systems and the continuous changes that are committed to the software's codebase, regression testing has become very expensive for real-world software applications. Test case prioritization is a classic solution in this context. Test case prioritization is the process of ranking existing test cases for execution with the goal of finding defects sooner. It is useful when the testing budget is limited and one needs to limit their test execution cost, by only running top n test cases, according to the testing budget. There are many heuristics and algorithms to rank test cases [143].

# 3.3.4 Gaps Detected and Critics to the Current Status

Even though some described test techniques can be applied to ES, most of them were developed for high-level systems. This means that those that cannot be applied directly —penetration testing, fuzzing, or A/B testing— should be adapted [144].

# 3.4 Vulnerability Analysis

This section presents the current status of vulnerability analysis. This review aims to find similar approaches in the literature, including the current standard.

# 3.4.1 Vulnerability Analysis in Security Standards

This review is focused on how standards conduct vulnerability analysis, the use of metrics, their management of the life cycle of the device, the techniques that they propose, and the security evaluation of both software and hardware.

#### 3.4.1.1 ISA/IEC 62443

The ISA/IEC 62443-4-1 technical document is divided into eight practices, which specify the secure product development life cycle requirements for both the development and the maintenance phases [89]. "Practice 5 - Security verification and validation testing" (SVV) section of this document specifies that a process shall be employed to identify and characterize potential security vulnerabilities in the product, including known and unknown vulnerabilities [53,56]. Two requirements in Practice 5 are in charge of the task of analyzing vulnerabilities, as follows [89]:

- Requirement SVV-3. Vulnerability Testing. This requirement states that a process shall be employed to perform tests that focus on identifying and characterizing potential and known security vulnerabilities in the product (*i.e.*, fuzz testing, attack surface analysis, black box known vulnerability scanning, software composition analysis, and dynamic runtime resource management testing).
- **Requirement SVV-4. Penetration Testing**. This requirement states that a process shall be employed to identify and characterize security-related issues via tests that focus on discovering and exploiting security vulnerabilities in the product (*i.e.*, penetration testing).

Although the ISA/IEC 62443-4-1 document considers the possibility of analyzing and characterizing the vulnerabilities of an industrial component, it does not propose a technique to perform this task, but instead refers to other standards for vulnerability handling processes [145]. In addition, it does not indicate how the data obtained from the analysis should be interpreted, and it does not define metrics or reference values for the current state of compliance with the requirement. Finally, it does not take into account neither the dependencies among the assets of the industrial component (dependency trees), nor their evolution of the number of vulnerabilities over time.

#### 3.4.1.2 Common Criteria

The CC defines five tasks in the Vulnerability Assessment class, which manage the deepness of the vulnerability assessment. The higher the EAL to be achieved, the greater the number of tasks in the list to be performed [55]:

- 1. Vulnerability survey,
- 2. Vulnerability analysis,
- 3. Focused vulnerability analysis,

- 4. Methodical vulnerability analysis, and
- 5. Advanced methodical vulnerability analysis.

Every task checks for the presence of publicly known vulnerabilities. Penetration testing is also performed. The main difference among the five levels of vulnerability analysis described here is the deepness of the analysis of known vulnerabilities and the penetration testing.

The CC scheme defines the general activities, but it does not specify how to perform them, therefore no technique for analyzing vulnerabilities is proposed. The evaluator decides the most appropriated techniques for each test in each scenario and for each device, which adds a large degree of subjectivity to the evaluation. Furthermore, dependencies among vulnerabilities and assets are not considered in the analysis. Moreover, the CC does not define a procedure to manage the life cycle of the device. In other words, when updated, the whole device has to be reevaluated [124,146–148]. Finally, although the usage of metrics is encouraged by the CC, it does not propose any explicitly defined metric to be used during the evaluation.

# 3.4.2 Vulnerability Analysis in the Literature: Directed Graphs

Directed graphs, and structures derived from them (*e.g.*, attack graphs), are graphs in which the edges have a defined direction. This direction is usually indicated with an arrow on the edge. They are mainly used for vulnerability analysis in computer networks. Vulnerability analysis is a key step towards the security evaluation of a device. Consequently, many research efforts have been focused on solving this issue. In this section, the most relevant works related to vulnerability analysis are reviewed.

Homer *et al.* [149] present a quantitative model for computer networks that objectively measures the likelihood of a vulnerability. Attack graphs and individual vulnerability metrics, such as CVSS, and probabilistic reasoning are applied to produce a sound risk measurement. However, the main drawback is that their work is only applicable to computer networks. Although they propose new metrics based on the CVSS for probabilistic calculations, they do not integrate standards such as CAPEC to enhance their approach centered on possible attacks and privilege escalation. They also fail to establish a relationship among existing vulnerabilities, and they fail to obtain the source problem causing each vulnerability.

Zhang *et al.* [150,151] developed a quantitative model that can be used to aggregate vulnerability metrics in an enterprise network based on attack graphs. Their model measures the likelihood that breaches can occur within a given network configuration, taking into consideration the effects of all possible interplays between vulnerabilities. This research is centered on computer networks, using attack graphs. Although the proposed model is capable of managing shared dependencies and cycles, only CVSS-related metrics are used. Moreover, this model assumes that the attacker knows all the information in the generated attack graphs. Finally, the method that they proposed for the aggregation of metrics is only valid for attack graphs, and is not valid for vulnerability analysis.

George *et al.* [152] propose a graph-based model to address the security issues in Industrial IoT (IIoT) networks. Their model is useful because it represents the relationships among entities and their vulnerabilities, serving as a security framework for the

risk assessment of the network. Risk mitigation strategies are also proposed. Finally, the authors discuss a method to identify the strongly connected vulnerabilities. However, the main drawback of this work is that each node of the generated attack graph represents a vulnerability instead of representing a device or an asset of that device. This leads to a loss of information in the analysis, because there is no way to know which vulnerability belongs to which device. Moreover, these methods need to know the relationships among present vulnerabilities in the devices. This information is not trivially obtained, and a human in the loop is needed. The proposals of [153] and [154] follow a similar graph-based approach to study the effects of cascade failures in the power grid, and a subway network.

Poolsappasit *et al.* [155] propose a risk management framework using Bayesian networks that enables a system administrator to quantify the chances of network compromise at various levels. The authors are able to model attacks on the network, and also to integrate standardized information of the vulnerabilities involved, such as their CVSS score. Although their proposed model lends itself to dynamic analysis during the deployed phase of the network, these results can only be applied to computer networks. Meanwhile, the prior probabilities that are used in the model are assigned by network administrators, and hence are subjective. The proposed model also has some issues related to scalability.

Muñoz-González *et al.* [156] propose the use of efficient algorithms to make an exact inference in Bayesian attack graphs, which enables static and dynamic network risk assessments. This model is able to compute the likelihood of a vulnerability, and can be extended to include zero-day vulnerabilities, attacker's capabilities, or dependencies between vulnerability types. Although this model is centered on studying possible attacks, it fails to integrate standards (such as CAPEC) that are related to attack patterns. Moreover, the generated graphs are focused on privilege escalation, trust, and users, rather than including information about vulnerabilities and the analyzed device.

Liu *et al.* [157] carry out a detailed assessment of vulnerabilities in IoT-based critical infrastructures from the perspectives of applications, networking, operating systems, software, firmware, and hardware. They highlight the three key critical infrastructure IoT-based cyber-physical systems (*i.e.*, smart transportation, smart manufacturing, and smart grid). They also provide a broad collection of attack examples upon each of the key applications. Finally, the authors provide a set of best practices and address the necessary steps to enact countermeasures for any generic IoT-based critical infrastructure system. Nevertheless, their proposal is focused on attacks and countermeasures, and it leaves aside the inner analysis of the targets. Continuous evaluation over time is not considered in this proposal, and no enhancements of the development process are generated.

Hu *et al.* [158] Hu et al. propose a network security risk assessment method that is based on the improved hidden Markov model (I-HMM). The proposed model reflects the security risk status in a timely and intuitive manner, and it detects the degree of risk that different hosts pose to the network. Although this is a promising approach, it is centered on computer networks and is at a higher abstraction level. No countermeasure or enhancement in the development process is proposed or generated.

Zografopoulos et al. [159] provide a comprehensive overview of the Cyber-Physical System (CPS) security landscape, with an emphasis on Cyber-Physical Energy Systems (CPES). Specifically, they demonstrate a threat modeling methodology to accurately represent the CPS elements, their interdependencies, as well as the possible attack entry points and system vulnerabilities. They present a CPS framework that is designed to delineate the hardware, software, and modeling resources that are required to simulate the CPS. They also construct high-fidelity models that can be used to evaluate the system's performance under adverse scenarios. The performance of the system is assessed using scenario-specific metrics. Meanwhile, risk assessment enables system vulnerability prioritization, while factoring the impact on the system's operation. Although this research work is comprehensive, it is focused on enhancing the existing adversary and attack modeling techniques of CPSs of the energy industry. Moreover, their model does not integrate the internal structure of the target of evaluation, and it does not take both software and hardware into account for the evaluation. Continuous evaluation over time is not considered. Finally, they do not propose countermeasures, or any kind of mechanism to enhance the security or the development of the CPSs.

#### 3.4.3 Gaps Detected and Critics to the Current Status

Most of the works reviewed here are more focused on modeling threats and attacks. They do not propose solutions to protect CPSs, enhancing their development, or manage their update throughout their whole life cycle. It is worth noting that they are still more focused on software evaluation, while hardware is usually neglected in their proposals.

As shown in this review, most of the research has adopted dependency trees, attack graphs, or directed graphs as the main tool to manage and assess vulnerabilities in computer networks. Graphs are an efficient technique to represent the relationships between entities, and they can also effectively encode the vulnerability relations in the network. Furthermore, the analysis of the graph can reveal the security-relevant properties of the network. For fixed infrastructure networks, graphical representations, such as attack graphs, are developed to represent the possible attack paths by exploiting the vulnerability relationships. For these reasons, vulnerability analysis techniques based on directed graphs are frequently found in the literature [160]. However, despite their potential, these analysis techniques have been relegated to vulnerability analysis in computer networks. Graph-based analysis has rarely been applied to industrial components.

#### 3.5 Metric Aggregation

System security quantification is not an easy task [75]. There exist both a lack of consensus and standardization around security metrics [71–74,90,91,161]. For this reason, research efforts keep aiming to unify this field [162].

Among these efforts, the Common Vulnerability Scoring System (CVSS) is a widely extended standard for vulnerability quantification [67]. CVSS is a public framework that provides a standardized method for assigning quantitative values to security vulnerabilities according to their severity. A CVSS score is a decimal number in the range [0, 10]<sup>1</sup> [65].

The CVSS is aimed to quantify the severity of vulnerabilities in individual and specific software items, however the majority of systems are actually a composition of simpler isolated items with different interdependencies. This situation highlights one of the biggest problems related to security quantification [163], the difficulty to really measure the global security state of a composite system. To do so, it would be necessary to aggregate each individual CVSS value into a global one in a consistent and coherent way.

The official CVSS documentation does not propose any kind of aggregation mechanism, and nowadays, there is no standardized method [164]. In addition to this, previous research works do not usually integrate contextual or interdependency information about the vulnerabilities to update the CVSS. This means that aspects such as whether affected functionalities, the environment of deployment, or the existence of exploits are usually neglected.

Context is a critical aspect to integrate in the aggregation process. This can be illustrated using a device implementing multiple functionalities as an example. To perform those functionalities, usually it will contain assets that implement those functionalities. But depending on the context where the device is deployed, some of its functionalities might not be needed. So the assets implementing unused functionalities would be disabled, and therefore, their vulnerabilities could not be exploited. It can also be the case that the asset implementing a functionality is simply inaccessible, so it could not also be exploited.

Nowadays, there is no widely-accepted method to aggregate CVSS values for soft-ware composition. All of them can be classified into one of the following categories [165, 166]: (1) Arithmetic Aggregation, (2) Attack Graph-based Aggregation, and (3) Bayesian Network-based Aggregation.

#### 3.5.1 Arithmetic Aggregation

This method uses arithmetic operations to aggregate the values [167–170]. Common examples of this approach are taking the maximum of the CVSS values, their arithmetic mean, or a combination of them. For example, Heyman *et al.* [167], proposed an algorithm to aggregate CVSS values in dependency graph that is based on taking the maximum value in each case, according to certain conditions.

Although their simplicity makes them suitable for initial approximations, their results can be biased in two ways:

<sup>&</sup>lt;sup>1</sup>The latest version at the time this paper was written is version 3.1.

- 1. **Exploitable by quantity:** When a system poses several vulnerabilities that by their own are not critical and cannot be exploited, they can sum up to an aggregated value of a high impact vulnerability (overfitting). This can happen when multiple simple mechanisms are combined as the aggregation algorithm.
- 2. **Exploitable by criticality:** When there exist a critical vulnerability, the whole system will be usually classified as critical. Nevertheless, that vulnerability might not be exploitable, nor being affecting the functionality of the system. This is specially common when using the maximum as the aggregation algorithm.

Since these metrics are defined over a set, they do not depend on where vulnerabilities are located, or how they are related to each other. These naive approaches lead to misleading results because they ignore causal relationships between vulnerabilities.

#### 3.5.2 Attack Graph-based Aggregation

This approach models the relationships between vulnerabilities using attack graphs, converting CVSS scores into probabilities [171–176]. In this way, both the CVSS value and the place of the vulnerability in the whole graph are taken into account.

Cheng *et al.* in [165] proposed a graph-based aggregation method that uses the underlying metrics of CVSS, where the dependency relationships between vulnerabilities are usually visible. As the center of the aggregation algorithm, they use the product of the CVSS used as probabilities, also known as the join probability of both vulnerability.

The main drawback with these approaches is that the relationship between individual vulnerabilities cannot be obtained straightforwardly from existing databases. This means that establishing a relation between two vulnerabilities implies that they can be chained during an attack, which is not always obvious. Moreover, factors such as exploitability of the vulnerabilities, or existing exploits are not taken into account.

#### 3.5.3 Bayesian Network-based Aggregation

Going a step further, these methods integrate the conditional relationship between vulnerabilities, modeling them using Bayesian networks [177–179]. Poolsappasit *et al.* [178] proposed a CVSS aggregation framework using Bayesian networks. They used the Bayesian probability factorization formula as the aggregation mechanism:

$$p(x) = \prod_{i=0} p(x_v | x_{pa(v)})$$

Bayesian network-based approaches have to deal with establishing the relationships between the vulnerabilities, but also with the calculation of conditional probabilities, that have to be usually estimated. As the previous ones, these techniques do not integrate information about how functionality if affected by existing vulnerabilities, or the possibility to actually exploit them.

#### 3.5.4 Gaps Detected and Critics to the Current Status

The aggregation methods found in the literature have evolved to address for an increasing complexity in the environment. So they evolved from simple averages,

to more complex models based on Bayesian networks. Nevertheless, their main drawback is not related to how they work, but to the data they need to work: CVSS scores where design, so most of the vulnerabilities have a value higher than 6.0. In fact, the current weighted average CVSS score is 6.5. This, together with the fact that any given device will have a high amount of vulnerabilities, makes the aggregated result to tend easily to 10. This is not desirable as the results are no longer meaningful.

Finally, these methods do not take the context or their functionality into account. In other words, a device could carry out several functionalities while only one of those functionalities could depend on a library such as OpenSSL. If there would be an OpenSSL vulnerability with a CVSS value of 10, and that particular device would be used to accomplish other functionalities not depending on OpenSSL, that would not mean that the CVSS value of the whole device would be 10. The resulting CVSS value for the whole device should be adjusted according to the used or accessible functionalities. But it can be the case that that vulnerability only affects the specific functions of OpenSSL that the devices does not use. Furthermore, it can happen that the device is indeed affected by the vulnerability, but its presence does not have any repercussion in its functionality.

## Chapter 4

# How to Quantify the Security Level of Embedded Systems? A Taxonomy of Security Metrics

This chapter covers the existing evaluation methodologies and security metrics in order to classify the existing security metrics based on a new taxonomy. In order to propose a new taxonomy, existing evaluation methodologies and standards have been analyzed to check whether they propose or define any kind of security metric. Then, we perform a systematic review of the literature, gathering existing security metrics. We highlight the importance of security metrics, describe the existing metrics proposed in international standards, and analyzes the proposed taxonomies. Finally, we classify them by developing a new taxonomy.

# 4.1 Features of Comprehensive Security Metrics for Embedded System

Some authors have tried to characterize which criteria should meet a comprehensive metric. Among their efforts, three main ideas can be found in the literature: (1) the SMART criterion [94], (2) the PRAGMATIC criterion [95] and (3) the characteristics identified in the work carried out by Savola [77].

Figure 4.1 shows the mapping among the three different concepts explained by each criterion. This mapping was carried out using the definition and the data given by the authors, to find the common points among them<sup>1</sup>.

The approach proposed by Savola classifies properties into 19 items; whereas the other two criteria have nine and five criteria, resulting in a higher mapping ratio. With this figure it is not possible to deduce which property is more important in any context, but it is possible to conclude that cost, accuracy, genuineness and repeatability are the most common properties in all three criteria.

It should be noted that none of these criteria explicitly mentions that metrics should be quantitative. This may be because, depending on the context, a nominal scale could be used instead. Some authors explicitly avoid this approach and prefer metrics to be quantitative [97]. In this mapping, many criteria have associated the cheap

<sup>&</sup>lt;sup>1</sup>This figure can be better visualized at https://embeddedsecuritymetrics.github.io/.

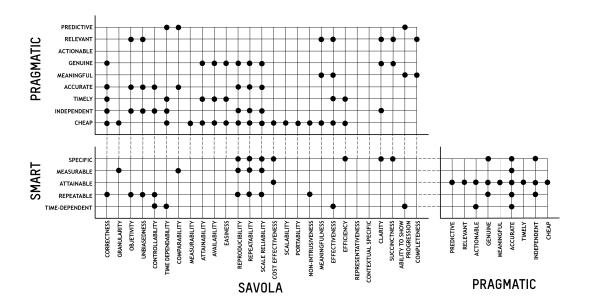


FIGURE 4.1: Mapping of criteria for identifying a good metric, according to the SMART, and PRAGMATIC criteria; and the survey carried out by Savola. Each black circle represents a common point between each one of the properties.

feature. This is because there are properties that directly affect the cost of computing a metric (the more properties it meets, the more expensive it is).

On the other hand, there are two properties identified by Savola (*i.e.*, contextual specific and representativeness) that do not have any correspondence with the rest of the criteria.

In the context of ES, we propose that the quality criteria that Security Metrics (SM) should address are (in no particular order):

- Comparability: SM should support comparison of the targets that they represent.
- **Cost effectiveness:** measurement gathering and approaches should be cost-effective.
- **Measurability:** SM are capable of having dimensions, quantity or capacity ascertained in the ES.
- **Repeatability:** The same results are achieved if a measurement is repeated in the same context, with exactly the same conditions.
- **Reproducibility:** The same results are achieved if a measurement is repeated in the same context, with exactly the same conditions, and different persons.

### 4.2 Selecting Metrics for Embedded Systems

There is neither a single system metric nor "one-perfect" set of metrics that could be suitable for every single case. The set of metrics that will be most suitable depends on multiple factors. For that reason, it is convenient to analyze whether the existing

metrics can be applied to ES. We carry out a systematic review of the state of the art to search for the currently available security metrics. The objective is to analyze to what extent the proposed metrics can be applicable to the security evaluation of ES. The details of this analysis are explained in the following points.

#### 4.2.1 Search and selection strategy

The data sources used are online databases, conference proceedings and academic journals of IEEE Xplore, Elsevier, AMC Digital Library, Springer, along with Google Scholar search engine, Web Of Science and Scopus. The search terms included the keywords: "security", "metric", "measure", "evaluation" and "assessment", considering another synonyms. The inclusion criteria included:

- Security measurements or metrics are the main subjects.
- Surveys collecting security metrics were preferred.
- The paper is primarily related to measuring security.

The initial number of metrics analyzed was 531, obtained mainly from [78,180–182]. For each one, the following data was collected:

- **Definition:** Definition of the metric given by the author.
- Scale: Nominal, ordinal, interval, ratio, absolute and distribution [78].
- Scope: User, software, hardware, device or organization.
- Automation: Automatic, semi-automatic or manual.
- **Measurement:** Static or dynamic measurement.

#### 4.2.2 Filtering and exclusion criteria

In order to identify the subset of metrics that can be also applied to ES from that initial set, the collected metrics were filtered according to the following criteria:

- Organizational specific metrics.
- Network specific metrics.
- Metrics without a clear definition.
- Repeated metrics.

When the concept of a metric was found to be applicable with some modification, it was considered to be eligible.

#### 4.2.3 Data extraction and interpretation

After filtering the data, the final count of metrics was 169. An interesting, albeit not very surprising fact, is that most of the metrics (77.5% of them) were related exclusively to software (*e.g.*, lines of code, number of functions and so on). On the other hand, only 0.6% of them were related exclusively to hardware (*e.g.*, sidechannel vulnerability factor metric); and finally, 14.8% of them could be applied to both software and hardware (*e.g.*, the historically exploited vulnerability metric that measures the number of vulnerabilities exploited in the past). The remaining 7.1%

is focused on other aspects, such as user usability. This shows that there is a clear lack of hardware security metrics in the literature, and main efforts are centered in developing software-related metrics. It is worth commenting these metrics are not only valid for ES, but could also be applied to other systems.

All the data extracted from the articles, including the 169 security metrics, as well as the filtering and selection processes, are available at GitHub.

#### 4.3 A New Taxonomy for Security Metrics

Although we discussed the most commonly used taxonomy for classifying security metrics in Chapter 3, to the best of our knowledge, there is no taxonomy that focuses on the specific assessment needs and limitations of ES. On the other hand, the taxonomies proposed in the literature are usually at a very high level, and therefore not applicable to ES. Thus, we propose a taxonomy that is based on which assets have to be protected in ES. Figure 4.2 shows the structure of the proposed taxonomy.

According to the research work in [183], embedded security cannot be addressed at a single level of abstraction; instead, it must be addressed at all abstraction levels. The main structure of the proposed taxonomy is built on three main blocks that aim to address all those abstraction levels: (1) Assets: elements of the ES that can be evaluated, as hardware, software, data, services, crypto keys, or communications; (2) Security Dimension: which attribute is to be protected (foundational requirements from the IEC 62443); and (3) Metric Properties: which are the characteristic parameters of the metric (such as automation, or computing cost). It is important to notice that a metric might cover more than one asset, and/or more than one security dimension, but at least, one value for each one have to be assigned (dashed line in Figure 4.2). Nevertheless, all values in metric properties are needed for each metric, as they represent the way the metric behaves, the cost associated to a metric and the information it returns (solid line in Figure 4.2).

The use of this taxonomy can be further explained using the Side-Channel Leakage (SCL) metric as an example. This metric uses statistical tests to measure if a device is prone to a side channel attack by comparing two sets of data; a random one and a fixed one. The SCL metric can be classified as follows:

• Assets: system hardware.

• Security dimension: data confidentiality.

• Metric properties:

- Automation: semi-automation.

- Scale: nominal (yes/no).

- Measurement: dynamic.

- Computing cost: constant.

As can be seen, the main feature of the proposed taxonomy is that considers all assets in ES. Unlike other proposals in the literature that are at a high level (*e.g.*, enterprise level) this one is specifically designed to evaluate all aspects that comprise ES. Furthermore, it includes not only the CIA triad (*i.e.*, Confidentiality, Integrity

#### SECURITY METRICS TAXONOMY METRIC PROPERTIES **ASSETS** SECURITY DIMENSION DATA / **SERVICES ACCESS RESOURCE** AUTOMATION INFORMATION CONTROL **AVAILABILITY** (automatic, semi-automatic, manual) HARDWARE SOFTWARE USE RESTRICT DATA CONTROL FL0W (nominal, ordinal, interval, ratio, INFORMATION AUXILIARY SUPPORTS EQUIPMENT DATA DATA MEASUREMENT CONFIDENTIALITY INTEGRITY CRYPTOGRAPHIC KEYS (static, dynamic) COMPUTING COST TIMELY RESPONSE TO EVENTS (constant, linear, exponential, COMMUNICATION NETWORKS logarithmic)

FIGURE 4.2: Taxonomy structure for security evaluation metrics for ES. Dashed line indicates AND/OR condition, and solid line AND condition.

and Availability), but also other dimensions, such as access control, use control, restrict data flow and timely-response to events. These dimensions are foundational requirements in IEC 62443 that need to be evaluated in ES.

#### 4.4 Conclusions and Future Work

Security metrics are not standardized, and they would help in quantifying the security level of ES, that is, in giving the degree of confidence in the effectiveness of protections and countermeasures. This research work proposes a new taxonomy for security metrics, which takes into account the different assets of ES. Since it is critical to identify proper and non-exhausting security metrics for this security assurance, a conceptual mapping of criteria for choosing the most appropriate metrics is given. On the other hand, the literature has been analyzed, and 169 security metrics have been identified. 77.5% of them are dedicated exclusively to evaluating software security, while only 0.6% of them are oriented to hardware evaluation. This difference clearly shows the wider adoption by the community of software-related security metrics.

As a future work, we propose to select and refine a reduced set of metrics from those that have been analyzed, based on the 5 criteria that we have proposed, and applied to a real ES. This will be the foundation for proposing a security assessment methodology specifically designed for ES, which will provide qualitative and quantitative results. These results will give trust to end-users, make it possible to monitor the evolution of the security level of a device over time, and to compare different versions or devices.

## **Chapter 5**

# A Novel Model for Vulnerability Analysis through Enhanced Directed Graphs and Quantitative Metrics

Chapter 4 presented an analysis of the current status of security metrics, both in the literature and in the standards. The analysis showed that even though there is a plethora of available metrics to use, existing standards and methodologies do not integrate any of them in their workflows. Nevertheless, they encourage their integration in the evaluation process. With this analysis, we have covered the current status of security metrics and evaluation methodologies.

In this chapter, we focus on other aspects of ESs: A methodology that enables the monitoring of the security of ESs during their whole life cycle. To achieve this, we developed a model for continuous vulnerability assessment to manage the security of existing industrial components during their whole lifespan to deal with emerging threats.

#### 5.1 Proposed Approach

In this contribution, we propose an Extended Dependency Graph (EDG) model that performs continuous vulnerability assessment to determine the source and nature of vulnerabilities, and enhance security throughout the entire life cycle of industrial components. The proposed model is built on a directed graph-based structure, and a set of metrics based on globally accepted security standards. The proposed EDG model is intended to:

- Identify the root causes and nature of vulnerabilities.
- Extract new requirements and test cases.
- Support the prioritization of patching.
- Track vulnerabilities during the whole lifespan of industrial components.
- Support the development and maintenance of industrial components.

To accomplish this task, the proposed model comprises two basic elements: the model itself, which is capable of representing the internal structure of the system under test; and a set of metrics, which allow conclusions to be drawn about the origin, distribution, and severity of vulnerabilities. Both the model and metrics are very flexible and exhibit some properties that make them suitable for industrial components, and can also be applied to enhance the ISA/IEC 62443 standard.

The content in this section is distributed into four sections, namely:

- 1. **Model:** The proposed model is explained, together with the systems in which it can be applied and the algorithms that are used to built it.
- 2. **Metrics:** Metrics are a great tool to measure the state of the system and to track its evolution. The proposed metrics and their usage are described in this section.
- 3. **Properties:** The main features of the proposed model and metrics (*e.g.*, granularity of the analysis, analysis over time, and patching policy prioritization support) are described in detail.
- 4. **Applicability:** Even though the reviewed standards exhibit some gaps, the proposed model aims to serve as the first step towards generating a set of tools to perform a vulnerability analysis in a reliable and continuous way. This last section will discuss the requirements of the ISA/IEC 62443-4-1 that can be enhanced using our model.

#### 5.1.1 Description of the Model

The proposed model in this research work is based on directed graphs, and requires knowledge of the internal structure of the device to be evaluated (i.e., the assets, both hardware and software, that comprise it and the relationships between them). This section defines the most basic elements that make up the model, the algorithms to build it for any give system, and its graphical representation.

**Definition 5.1.1.** A System Under Test (SUT<sup>1</sup>) is now represented by an Extended Dependency Graph (EDG) model  $G = (\langle A, V \rangle, E)$  that is based on directed graphs, where A and V represent the nodes of the graphs, and E represents its edges or dependencies:

- $A = \{a_1, ..., a_n\}$  represents the set of assets in which the SUT can be broken down, where n is the total number of obtained assets. An asset a is any component of the SUT that supports information-related activities, and includes both hardware and software [107, 184, 185]. Each asset is characterized by its corresponding CPE identifier, while its weaknesses are characterized by the corresponding CWE identifier. In the EDG model, the assets are represented by three types of nodes in the directed graphs (*i.e.*, root nodes, asset nodes and cluster).
- $V = \{v_1, ..., v_q\}$  represents the set of known vulnerabilities that are present in each asset of A, where q is the total number of vulnerabilities. They are characterized by the corresponding CVE and CVSS values. In the EDG model, vulnerabilities are represented using two types of nodes in the directed graphs (*i.e.*, known vulnerability nodes and clusters).
- $E = \{e_{ij} | \forall i, j \in \{1, ..., n+q\}$  such that  $i \neq j\}$  represents the set of edges or dependencies among the assets, and between assets and vulnerabilities.  $e_{ij}$  indicates that a dependency relation is established from asset  $a_i$  to asset  $a_j$ . Dependencies are represented using two different types of edges in the EDG (*i.e.*, normal dependency and deprecated asset/updated vulnerability edges).

In other words, the EDG model can represent a system, from its assets to its vulner-abilities, and its dependencies as a directed graph. Assets and vulnerabilities are represented as nodes, whose dependencies are represented as arcs in the graph. The information in the EDG is further enhanced by introducing metrics.

The EDG model of a given SUT will include four types of node and two types of dependency. The graphical representation for each element is shown in Table 5.1. Figure 5.1 shows an example of a simple EDG and its basic elements. All of the elements that make up an EDG will be explained in more detail below:

#### **5.1.1.1 Types of Node**

The EDG model uses four types of node:

Root nodes represent the SUT,

<sup>&</sup>lt;sup>1</sup>Following the denomination in the ISA/IEC 62443 standard [48], the SUT may be an industrial component, a part of an industrial component, a set of industrial components, a unique technology that may never be made into a product, or a combination of these.

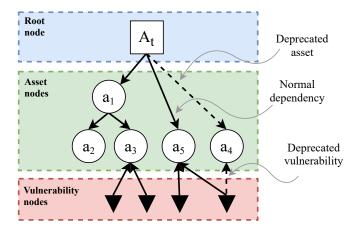


FIGURE 5.1: Basic elements of an EDG. Note that clusters are not displayed in this figure. For clusters, see Figure 5.3. For metrics definition, see Section 5.1.2.

TABLE 5.1: Overview of the information that is necessary to define each of the EDG elements.

SYMBOL	NOTATION	MEANING	VALUES
	A(t)	Root Node / Device Node	CPE <sub>current</sub>
$\bigcirc$	a(t)	Asset Node	$CPE_{previous}, CPE_{current}, CWE_{a_i}(t)$
$\circ$	$\underline{a}(t)$	cluster	$\{CPE_{previous}, CPE_{current}, CWE_{a_i}(t)\}, \{CVE_{a_i}(t), CVSS_{v_i}(t), CAPEC_{w_i}(t)\}, \{Dependencies\}$
lacktriangle	v(t)	Known Vulnerability Node	$CVE_{a_i}(t), CVSS_{v_i}(t), CAPEC_{w_i}(t)$
$\longrightarrow$	e(t)	Dependency Relation	_
>	e(t)	Updated Asset / Patched Vulnerability	_

- **Asset nodes** represent each one of the assets of the SUT,
- Known vulnerability nodes represent the vulnerabilities in the SUT, and
- Clusters summarize the information in a subgraph.

**Root nodes** (collectively, set  $G_R$ ) are a special type of node that represent the whole SUT. Any EDG starts in a root node and each EDG will only have one single root node, with an associated timestamp (t) that indicates when the last check for changes was done. This timestamp is formatted following the structure defined in the ISO 8601 standard for date and time [186].

**Asset nodes** (collectively, set  $G_A$ ) represent the assets that comprise the SUT. The EDG model does not impose any restrictions on the minimum number of assets that the graph must have. However, the SUT can be better monitored over time when there is a higher number of assets. Moreover, the results and conclusions obtained will be much more accurate. Nevertheless, each EDG will have as many asset nodes as necessary, and the break-down of assets can go as far and to as low-level as needed.

Each known vulnerability node will be characterized by the following set of values:

- *CPE*<sub>current</sub>: Current value for the CPE. This points to the current version of the asset it refers to.
- *CPE*<sub>previous</sub>: Value of the CPE that identifies the previous version of this asset. This will be used by the model to trace back all the versions of the same asset over time, from the current version, to the very first version.
- $CWE_{a_i}(t)$ : Set of all the weaknesses that are related to the vulnerabilities present in the asset. The content of this list can vary depending on the version of the asset.

Figure 5.2 illustrates how the tracking of the versions of an asset using CPE works. On the one hand, version  $a_i$  is the current version of asset a. It contains its current CPE value, and the CPE of its previous version. On the other hand,  $a_2$  and  $a_1$  are previous versions of asset a. The last value of  $a_1$  points to a null value. This indicates that it is the last value in the chain, and therefore the very first version of the asset a.

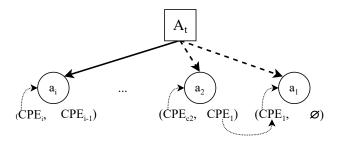


FIGURE 5.2: Tracking dependencies between the previous and current CPE values for asset *a*.

Known vulnerability nodes (collectively, set  $G_V$ ) represent a known vulnerability present in the asset that it relates to. Each asset will have a known vulnerability node for each known vulnerability belonging to that asset. Assets alone cannot tell how severe or dangerous the vulnerabilities might be, so unique characterization of vulnerabilities is crucial [152].

To identify each known vulnerability node, each will be characterized by the following set of features (formally defined in 5.1.2:

- $CVE_{a_i}(t)$ : This serves as the identifier of a vulnerability of asset  $a_i$ .
- $CVSS_{v_i}(t)$ : This metric assigns a numeric value to the severity of vulnerability  $v_i$ . Each CVE has a corresponding CVSS value.
- $CAPEC_{w_i}(t)$ : Each vulnerability (CVE) is a materialization of a weakness (CWE)  $w_i$  that can be exploited using a concrete attack pattern (CAPEC). In many cases, each CWE has more than one CAPEC associated. Consequently, this field is a set that contains all the possible attack patterns that can exploit the vulnerability that is being analyzed.

**Clusters** (collectively, set  $G_S$ ) are a special type of node that summarizes and simplifies the information contained in a subgraph in an EDG. Figure 5.3 shows how the clusters work.

To identify each cluster, and to be able to recover the information that they summarize, each is characterized by the data that define each of the elements that they contain:

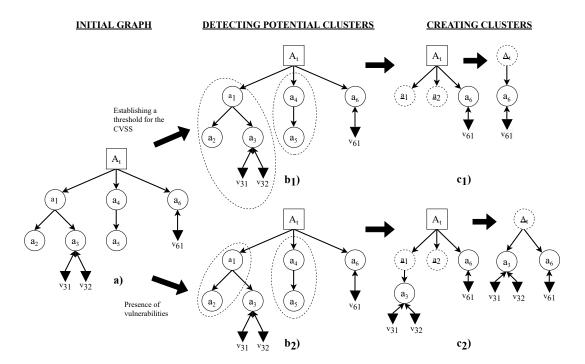


FIGURE 5.3: Creating clusters. Application of the two proposed criteria to the creation of clusters to simplify the graph: (1) Establishing a threshold to select which vulnerability stays outside the cluster (upper side). (2) Choosing the absence of vulnerability as the criterion to create clusters (lower side). The severity value (CVSS) for  $v_{211}$  and  $v_{212}$  is supposed to be lower than the establish threshold.

 $\{CPE_{previous}, CPE_{current}, CWE_{a_i}(t)\}, (CVE_{a_i}(t), CVSS_{v_i}(t), \{CAPEC_{w_i}(t))\},$ and their dependencies.

Two types of criteria can be used to create clusters and to simplify the obtained graph (Figure 5.3):

- 1. **Absence of vulnerabilities**: Using this criterion, clusters will group all nodes that contain no associated vulnerabilities.
- 2. **CVSS** score below a certain threshold: With this criterion, a threshold for the CVSS scores will be chosen. Nodes whose CVSS score is less than the defined threshold will be grouped into a cluster.

#### 5.1.1.2 Types of Edge

In the EDG model, edges plays a key role representing dependencies. Two types of edge can be identified:

- Normal dependencies relate two assets, or an asset and a vulnerability. They represent that the destination element depends on the source element. Collectively, they are known as set  $G_D$ .
- Deprecated asset or patched vulnerability dependencies indicate when an asset or a vulnerability is updated or patched. They represent that the destination

element used to depend on the source element. Collectively, they are known as set  $G_U$ .

The possibility of representing old dependencies brings the opportunity to reflect the evolution of the SUT over time. When a new version of an asset is released, or a vulnerability is patched, the model will be updated. Their dependencies will change then from a normal dependency to a deprecated asset or vulnerability dependency to reflect that change.

#### 5.1.1.3 Conditions of Application of EDGs

The EDG model is applicable to SUTs that meet the following set of criteria:

- Software and hardware composition: In our approach, the model is created by means of a white-box analysis. The absence or impossibility to perform a white-box analysis limits the ability to create an accurate model. Some knowledge about the internal structure and code is expected. This information is usually only known by the manufacturer of the component, unless the component is publicly available or open-source. It should be also possible to break down the SUT into simpler assets to generate a relevant EDG.
- Existence of publicly known vulnerabilities: The EDG model focuses on known vulnerabilities. This is not critical because many industrial components use commercial or open-source elements. The SUT must be composed of assets for which public information is available. If the majority of SUT assets are proprietary, or the SUT is an *ad hoc* development that is never exposed, then the generated EDG will not evolve. Therefore, the analysis will not be relevant.

#### 5.1.1.4 Steps to Build the Model

This section explains the process and algorithms that were used to build the corresponding EDG of a given SUT. The main scenarios that can be found are also described.

Before extracting useful information about the SUT, the directed graph associated with the SUT has to be built. This comprises several steps, which are described in the following paragraphs (see the flowchart in Fig. 5.4, and Fig. 5.5):

**Step 1** — **Break down the SUT into assets**. For the model to work properly, it relies on the SUT being able to be broken down into assets. With this in mind, the first step involves obtaining the assets of the SUT, either software or hardware. In the CC, this process is called modular decomposition of the SUT [50]. Ideally, every asset should be represented in the decomposition process, but this is not compulsory for the model to work properly. Each one of the assets obtained in this step will be represented as an asset node. In this step, the dependencies among the obtained assets are also added as normal dependencies.

**Step 2** — **Assign a CPE to each asset**. Once the assets and their dependencies have been identified, the next task is to assign the corresponding CPE identifier to each asset. If there is no publicly available information of a certain asset, and therefore, it does not have a CPE identifier, then it is always possible to generate one using the fields described in the CPE naming specification documents [59] for internal use in the model.

Step 3 — Add known vulnerabilities to the assets. In this step, the vulnerabilities  $(CVE_{a_i}(t))$  of each asset are set. This is done by consulting public databases of known vulnerabilities [54, 64] looking for existing vulnerabilities for each asset. When a vulnerability is found, it is added to the model of the SUT, including its dependencies. If there were no known vulnerabilities in an asset, then the asset would become the last leaf of its branch. In this step, the corresponding value of the CVSS of each vulnerability is also added to the model.

Step 4 — Assign to each asset its weaknesses and possible CAPECs. After the vulnerabilities, the corresponding weaknesses to each vulnerability ( $CWE_{a_i}(t)$ ) are added, along with the corresponding attack patterns ( $CAPEC_{w_i}(t)$ ) for each weakness. If there is no known vulnerability in an asset, then there will be no weaknesses. Meanwhile, it would be possible to have a known vulnerability in an asset, but no known weakness or attack pattern for that vulnerability. Finally, more than one CAPEC can be assigned to the same weakness. Consequently, it would be common to have a set of possible CAPECs that can be used to exploit the same weakness. It is worth noting that not all of them could be applied in every scenario.

Step 5 — Computing Metrics and tracking the SUT. At this point, the EDG of the SUT is completed with all the public information that can be gathered. This last step is to calculate the metrics defined (for further information, see Section 5.1.2.), generating the corresponding reports, and tracking the state of the SUT for possible updates in the information of the model. This step is always triggered when the SUT is updated. This can imply that a new asset can appear, an old asset can disappear, an old vulnerability can be patched, or a new one can appear in the SUT. All of these scenarios will be reflected in the model as they arise during its life cycle.

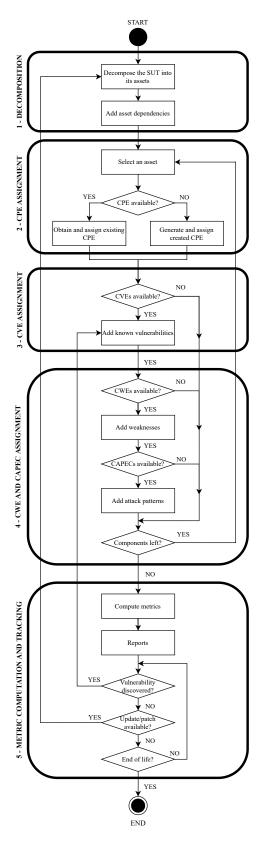
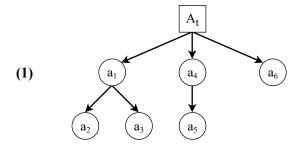
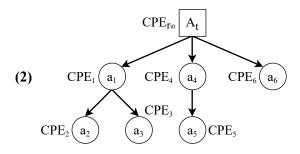
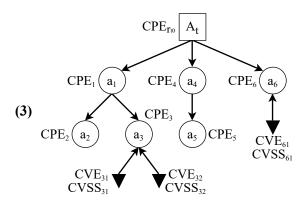


FIGURE 5.4: Algorithm to generate the initial EDG of a give SUT.







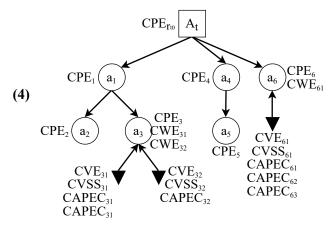


FIGURE 5.5: Example of the process of building the EDG model of a given SUT A.

#### 5.1.2 Security Metrics

The EDG model that was proposed in the previous sections is by itself capable of representing the internal structure of the SUT, and it can display it graphically for the user. This representation not only includes the internal assets of the SUT, but it also captures their relationships, existing vulnerabilities, and weaknesses. Moreover, assets, vulnerabilities, and weaknesses are easily identified using their corresponding CPE, CVE, and CWE values, respectively. All together, this constitutes a plethora of information that the model can use to improve the development and maintenance steps of the SUT, enhance its security, and track its status during its whole life cycle. Metrics are a great tool to integrate these features into the model.

Metrics can serve as a tool to manage security, make decisions, and compare results over time. They can also be used to systematically improve the security level of an industrial component or to predict its security level in a future point in time.

In this section, the basic definitions that serve as the foundation of the metrics are described. Then, the proposed metrics are introduced to complement the functionality of the EDG model. The main feature of these metrics is that they all depend on time as a variable, so it is possible to capture the actual state of the SUT, track its evolution over time, and compare the results.

#### 5.1.2.1 Basic Definitions

In this section, the basic concepts on which the definitions of the metrics will be based are formalized.

**Definition 5.1.2.** The set of all possible weaknesses at a time t is represented as CWE(t), where

$$CWE(t) = \{cwe_1, ..., cwe_m\}$$
(5.1)

and *m* is the total number of weaknesses at time t. This set contains the whole CWE database defined by MITRE [62].

**Definition 5.1.3.** The set of all of the possible vulnerabilities at a time t is represented as CVE(t) where

$$CVE(t) = \{cve_1, ..., cve_v\}$$

$$(5.2)$$

and p is the total number of vulnerabilities. This set contains the whole CVE database defined by MITRE [64].

**Definition 5.1.4.** The set of all possible attack patterns at a time t is represented as CAPEC(t), where

$$CAPEC(t) = \{capec_1, ..., capec_a\}$$
 (5.3)

and q is the total number of attack patterns at time t. This set contains the whole CAPEC database defined by MITRE [68].

**Definition 5.1.5.** The set of weaknesses of an asset  $a_i$  at a time t is defined as

$$CWE_{a_i}(t) = \{cwe_j | cwe_j \text{ is in the asset } a_i \text{ at time } t \land cwe_j \in CWE(t) \\ \land \forall k \neq j, cwe_j \neq cwe_k \}$$
(5.4)

From this expression, the set of all the weaknesses of a particular asset throughout its life cycle is defined as

$$CWE_{a_i} = \bigcup_{t=1}^{T} CWE_{a_i}(t)$$
 (5.5)

where  $|CWE_{a_i}|$  is the total number of non-repeated weaknesses in its entire life cycle.

**Definition 5.1.6.** The set of vulnerabilities of an asset  $a_i$  at a time t is defined as

$$CVE_{a_i}(t) = \{cve_i | cve_i \text{ is in the asset } a_i \text{ at time } t \land cve_i \in CVE(t)\}$$
 (5.6)

From this expression, the set of vulnerabilities of an asset throughout its entire life cycle is defined as

$$CVE_{a_i} = \bigcup_{t=1}^{T} CVE_{a_i}(t)$$
 (5.7)

where  $|CVE_{a_i}|$  is the total number of vulnerabilities in its entire life cycle.

**Definition 5.1.7.** The set of weaknesses of a SUT *A* with *n* assets at a time *t* is defined as:

$$CWE_A(t) = \bigcup_{i=1}^{n} CWE_{a_i}(t)$$
(5.8)

**Definition 5.1.8.** The set of vulnerabilities of a SUT A with n assets at a time t is defined as:

$$CVE_A(t) = \bigcup_{i=1}^n CVE_{a_i}(t)$$
 (5.9)

**Definition 5.1.9.** The set of vulnerabilities associated to the weakness  $cwe_j$  and to the asset  $a_i$  at a time t is defined as:

$$CVE_{a_i|cwe_j}(t) = \{cve_k|cve_k \text{ associated to weakness } cwe_j \text{ and to asset } a_i \text{ at time } t\}$$
(5.10)

It is worth noting that CWE is used as a classification mechanism that differentiates CVEs by the type of vulnerability that they represent. A vulnerability will usually have only one associated weakness, and weaknesses can have one or more associated vulnerabilities [66].

**Definition 5.1.10.** The partition j of an asset  $a_i$  at time t conditioned by a weakness  $cwe_k$  is defined as

$$CVE_{a_i|cwe_l}(t) = \{cwe_l|cwe_l = cwe_k \land cwe_l \in CVE_{a_i}(t)\}$$
(5.11)

**Definition 5.1.11.** The partition j of the SUT A at time t conditioned by a weakness  $cwe_k$  is defined as

$$CVE_{A|cwe_l}(t) = \{cwe_l | cwe_l = cwe_k \land cwe_l \in CVE_A(t)\}$$
(5.12)

**Definition 5.1.12.** The set of attack patterns associated to a weakness  $w_i$  at a time t is defined as

$$CAPEC_{w_i}(t) = \{capec_j | capec_j \text{ can exploit weakness } w_i \text{ at time } t \land capec_j \in CAPEC(t)\}$$
 (5.13)

**Definition 5.1.13.** The set of metrics that are defined in this research work based on the EDG model is defined as

$$M = \{m_1, ..., m_r\} \tag{5.14}$$

where r is the total number of metrics. This set can be extended, defining more metrics according to the nature of the SUT.

#### **5.1.2.2** Metrics

This section will describe the metrics that were defined based on the EDG model and the previous definitions. Although it might seem trivial, the most interesting feature of these metrics is that they all depend on time. Using time as an input variable for the computation of the metrics opens the opportunity to track results over time, compare them, and analyze the evolution of the status of the SUT. Furthermore, some metrics take advantage of time to generate an accumulated value, giving information about the life cycle of the SUT. Table 5.2 shows all of the proposed metrics, their definition, and their reference values.

In addition to the metrics in Table 5.2, the model allows the definition of other types of metrics according to the analysis to be performed, and the nature of the SUT (e.g., the vulnerability evolution function for SUT A up to time t for all vulnerabilities could be defined as the linear regression of the total number of vulnerabilities in each time t for SUT A, or using any other statistical model).

TABLE 5.2: Proposed metrics for the model.

N	METRIC	DEFINITION	REFERENCE VALUE				
	$M_0(A) = rac{ CVE_A(t) }{n(t)}$	Arithmetic mean of vulnerabilities in the SUT $A$ , where $n(t)$ is the number of assets in a SUT at a time $t$ . $M_0$ shows how many vulnerabilities would be present in each asset if they were evenly distributed among the assets of the SUT. The result of $M_0$ can serve as a preliminary analysis of the SUT, related to the criticality of its state. From Equation 8.	$M_0 < 1$ : The number of vulnerabilities is lower than the number of assets. $M_0 \ge 1$ : Every asset has at least one vulnerability.				
THES	$M_1(A,t) =  CVE_A(t) $	Number of vulnerabilities in a SUT $A$ at time $t$ . From Equation 8.	Ideally, the values of $M_1$ should be zero (no vulnerability in $A$ ), but the lower the value of $M_1$ , the better.				
VULNERABILITIES	$M_2(A) = \sum_{t=1}^{T}  CVE_A(t)  = \sum_{t=1}^{T} M_1(A, t)$	Number of vulnerabilities in a SUT $A$ throughout its entire life cycle $T$ . This metric computes the accumulated value of the number of vulnerabilities of a SUT throughout its entire life cycle. From Equation 8.	The lower the value of $M_2$ , the better.				
	$M_3(a_i,t) =  CVE_{a_i}(t) $	Number of vulnerabilities in an asset $a_k$ at time $t$ The values of $M_3$ can be useful during a vulnerability analysis, or when performing a penetration test, to identify the asset with more vulnerabilities. From Equation 6.	Ideally, the value of $M_3$ should be zero.				
	$M_4(a_k,t)=rac{ CVE_{a_k}(t) }{\sum_{i=1}^{n} CVE_{a_i}(t) }$	Relative frequency of vulnerabilities of the asset $a_k$ at a time $t$ . From Equation 6.	Ideally, the value of $M_4$ should be zero, or at least $M_4 \leq \frac{1}{n(t)}$ , being $n(t)$ the number of assets in the SUT. This value can also be expressed as the percentage of vulnerabilities of asset $a_i$ respect to the total number of vulnerabilities in the SUT, $M_4$ $(a_k,t) = \frac{ CVE_{a_k}(t) }{\sum_{i=1}^n  CVE_{a_i}(t) } \cdot 100$				
	$M_5(a_i, cwe_j, t) =  CVE_{a_i cwe_j}(t) $	Multiplicity of weakness $\mathit{cwe}_i$ of the asset $a_i$ at a time $t$ . This metric represents the number of times a weakness is present among the vulnerabilities of the asset $a_i$ . This is possible because a vulnerability can have associated the same weakness as other vulnerabilities. From Equation From Equation 9.	Ideally, the value of $M_5$ should be zero, or at least, $M_5 \leq \frac{ CVE_{A cuc_j}(t) }{n(t)}$ , being $n(t)$ the number of assets in the SUT. The value of the metric could be further narrowed by assuming that $cwe_j$ will be present in all but one asset, so $M_5 \leq \frac{ CVE_{A cuc_j}(t) }{n(t)-1}$ to be in acceptable values.				
	$M_6(A, cwe_j, t) =  CVE_{A cwe_j}(t) $	Multiplicity of weakness $cwe_j$ of the SUT $A$ at a time $t$ . This metric represents the number of times a weakness is present among the vulnerabilities of the SUT $A$ . From Equation 11.	Ideally, the value of $M_6$ should be zero.				
<b>VEAKNESSES</b>	$M_7(A,t) =  CWE_A(t) $	Number of weaknesses in a SUT $A$ at time $t$ . From Equation 7.	Ideally, the value of $M_7$ should be zero (no weakness in $A$ ), but the lower the value of $M_7$ , the better.				
WEA	$M_8(A) = \sum_{t=1}^{T}  CWE_A(t)  = \sum_{t=1}^{T} M_7(A, t)$	Number of weaknesses in a SUT <i>A</i> throughout its entire life cycle <i>T</i> . This metric computes the accumulated value of weaknesses of a SUT throughout its entire life cycle. From Equation 7	The lower the value of $M_8$ , the better.				

#### 5.1.3 Properties

Together, the EDG model and the defined metrics exhibit a series of characteristics that make them suitable for vulnerability assessment. These properties represent an advantage over the techniques reviewed in the state of the art, including automatic inference of root causes, spatial and temporal distribution of vulnerabilities, and prioritization of patching, which will be described in the following subsections.

#### 5.1.3.1 Automatic Inference of Root Causes

Each CWE natively contains information that is directly related to the root cause of a vulnerability. From this information, new requirements and test cases can be proposed.

#### 5.1.3.2 Spatial and Temporal Distribution of Vulnerabilities

The key feature of the proposed model is the addition of the temporal dimension in the analysis of vulnerabilities. This makes it possible to analyze the location of the vulnerabilities both in space (in which asset) and time (their recurrence), which allows us to track the state of the device throughout the whole life cycle. This approach also enables a further analysis of the SUT, by updating data in the model, such as new vulnerabilities that are found or new patches that are released.

Each time that a new vulnerability is found, or an asset is patched (*i.e.*, via an update), the initial EDG is updated to reflect those changes. An example of this process can be seen in Fig. 5.6.

At time  $t_0$ , the initial graph of the SUT A is depicted in Fig. 5.6. Because there is no vulnerability at that time, this graph can be simplified using the cluster notation, with just a cluster containing all assets. At time  $t_1$ , a new vulnerability that affects the asset  $a_2$  is discovered. At time  $t_2$ , the asset  $a_2$  is updated. This action creates a new version of asset  $a_2$ , asset  $a_3$ . Because the vulnerability was not corrected in the new update, both versions contain the vulnerability that was initially presented in asset  $a_2$ . Finally, at time  $t_3$ , the asset  $a_3$  is updated to its new version  $a_4$ , and the vulnerability is corrected.

This approach enables a further analysis of the SUT, including updated data, according to new vulnerabilities that are found or new patches that are released.

#### 5.1.3.3 Patching Policies Prioritization Support

The proposed model is not only able to include known vulnerabilities associated with an asset, but it also provides a relative importance sorting of vulnerabilities by CVSS. Relying on the resulting value, it is possible to assist in the vulnerability patching prioritization process. Furthermore, the presence of an existing exploit for a known vulnerability can be also be taken into account, when deciding which vulnerabilities need to be patched first. A high CVSS value combined with an available exploit for a given vulnerability is a priority when patching.

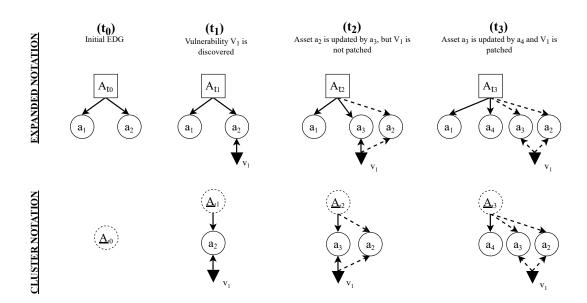


FIGURE 5.6: Representation of the temporal behavior in the graphical model using the two kinds of dependencies of the model. It is worth mentioning that these graphs could be further simplified by taking advantage of the cluster notation, as shown at the bottom of this figure.

#### 5.1.4 Applicability in the Context of ISA/IEC 62443

In this section, the potential application of the proposed EDG model to the existing security standards is described. The proposed EDG model can be used isolated by itself, or in combination with other techniques that complement the analysis. In this sense, the EDG model can be used to enhance some task in the security evolution processes defined by security standards.

The ISA/IEC 62443-4-1 standard specifies 47 process requirements for the secure development of products used in industrial automation and control systems [89]. Thus, the EDG model was developed to enhance the execution of one of those requirements defined by the standard: the "SVV-3: Vulnerability testing" requirement, serving as a support for the execution of Practice 5 — Security Verification and Validation testing. According to the SVV-3 requirement, both known and unknown vulnerability analysis has to be performed. The EDG model proposed in this research work is intended to support the identification of known vulnerabilities, their dependencies, and the possible consequences of their propagation, yielding the opportunity to analyze them systematically. Nevertheless, more requirements of the ISA/IEC 62443 can be mapped to one or more of the metrics defined in this research work. Using this relationship, it is possible to apply the EDG model to enhance the analysis and review of the following requirements:

#### 5.1.4.1 Security Requirements - 2: Threat Model (SR-2)

"A process shall be employed to ensure that all products have a threat model specific to the current development scope of the product. The threat model shall be reviewed and verified periodically" [89]. The proposed EDG model can serve as an abstraction of the threat model that has to be obtained. Moreover, the standard states that this threat model has to be reviewed periodically for updates. Given that the EDG of a given SUT evolves with every update, the threat model would be always up-to-date. Potential threats and their severity using the CVSS can also be analyzed with this proposal. Finally, these results can be used to enhance the risk assessment of the SUT.

#### 5.1.4.2 Security Management - 13: Continuous Improvement (SM-13)

"A process shall be employed for continuously improving the secure development life cycle" [89]. The EDG model can be used to identify recurrent issues in the development of an industrial component, due to its ability to track the state of a SUT over time. Consider the scenario where a piece of code contains an unknown vulnerability. For example, this code can implement a communication protocol, or the generation of a cryptographic key. If this piece of code is recurrently integrated in many type of devices, then when they are released to the market, the end users can identify that vulnerability and report it to the product supplier. The EDG can reflect the presence of that vulnerability. If an EDG is done for each type of device, then this problem can be detected beforehand. Using the CWE, the root problem can be detected. With this information, new training and corrective actions can be proposed to avoid this issue.

## 5.1.4.3 Specification of Security Requirements - 5: Security Requirements Review (SR-5)

"A process shall be employed to ensure that security requirements are reviewed, updated, and approved" [89]. As before, taking advantage of the previous scenario, the information extracted from the generated EDG model can be used to propose new requirements or to update the existing requirements.

# 5.1.4.4 Security Verification and Validation Testing - 4: Penetration Testing (SVV-4)

"A process shall be employed to identify and characterize security-related issues via tests that focus on discovering and exploiting security vulnerabilities in the product" [89]. The EDG model facilitates the identification of possible entry points to the SUT when carrying out a penetration test. In addition, existing attack patterns (CAPEC) and weaknesses (CWE) can serve as a starting point to discover unknown vulnerabilities and exploits.

# 5.1.4.5 Management of Security-related Issues - 3: Assessing Security-related issues (DM-3)

"A process shall be employed for analyzing security-related issues in the product" [89]. When a new vulnerability is detected, end users will report it to the product suppliers. Then, the corresponding EDG model of that SUT will be updated to reflect that change. This information, in addition to that previously contained in the EDG, can be used to obtain the severity value of the discovered vulnerability using the CVSS. This also facilitates the identification of root causes, related security issues, or the impact.

Finally, the ISA/IEC 62443-4-2 document defines four types of components of an IACS (*i.e.*, software applications, embedded devices, host devices, network devices) [49]. The proposed model is capable of representing the inherent complexity of each of them.

TABLE 5.3: Mapping between the developed metrics and the requirements they refer in the ISA/IEC 62443.

SR (Security Requirements), SM (Security Management), SVV (Security Validation and Verification), DM (Management of Security-Related Issues).

METRIC	SR-2	SR-5	SM-13	SVV-4	DM-3
$M_0(A) = \frac{ CVE_A(t) }{n(t)}$	•		•	•	•
$M_1(A,t) =  CVE_A(t) $		•			
$M_2(A) = \sum_{t=1}^{T}  CVE_A(t)  = \sum_{t=1}^{T} M_1(A, t)$			•		
$M_3(A,t) =  CVE_{a_i}(t) $		•	•	•	
$M_4(a_k,t) = rac{  extit{CVE}_{a_k}(t) }{\sum_{i=1}^n   extit{CVE}_{a_i}(t) }$		-	•		
$M_5(a_i, cwe_j, t) =  CVE_{a_i cwe_j}(t) $			•	•	
$M_6(A, cwe_j, t) =  CVE_{A cwe_j}(t) $				•	
$M_7(A,t) =  CWE_A(t) $	•			•	
$M_8(A) = \bigcup_{t=1}^T  CWE_A(t)  = \bigcup_{t=1}^T M_7(A, t)$		•	•		

#### 5.2 Real Use Case: General Analysis of OpenPLC

In this section, the EDG model and the proposed metrics will be applied to perform a vulnerability assessment of the OpenPLC project, which will be the SUT. In the subsections, we will assess the three available versions of the OpenPLC project. For each, the EDG model will be obtained, and the proposed metrics will be applied to draw conclusions about the vulnerability status of each version.

OpenPLC is the first functional standardized open source Programmable Logic Controller (PLC), both in software and hardware [187]. It was mainly created for research purposes in the areas of industrial and home automation, Internet of Things (IoT), and SCADA. Given that it is the only controller that provides its entire source code, it represents an engaging low-cost industrial solution — not only for academic research but also for real-world automation [188,189].

#### 5.2.1 Structure of OpenPLC

The OpenPLC project consists of three parts:

- 1. **Runtime:** It is the software that plays the same role as the firmware in a traditional PLC. It executes the control program. The runtime can be installed in a variety of embedded platforms, such as the Raspberry Pi, and in Operating Systems (OSs) such as Windows or Linux. Industrial Modbus slave devices can be attached to expand the number of inputs and outputs.
- 2. **Editor:** An application that runs on a Windows or Linux OS that is used to write and compile the control programs that will be later executed by the runtime.
- 3. **HMI Builder:** This software is to create web-based animations that will reflect the state of the process, in the same manner as a traditional HMI.

When installed, the OpenPLC runtime executes a built-in webserver that allows OpenPLC to be configured and new programs for it to run to be uploaded.

#### 5.2.2 Initial Scenario

The OpenPLC project is constituted by three different versions [190–192] as can be seen in Table 5.4. For this research work, Ubuntu Linux was selected as the platform to install the OpenPLC runtime. Ubuntu Linux provides comprehensive documentation, previous versions are accessible, and software dependencies can easily be obtained.

To make the analysis of OpenPLC fair, a contemporary operating system was selected for each of the OpenPLC versions, according to the version of Ubuntu that was available at the release time of each OpenPLC version (see Table 5.4). The Long Term Support (LTS) version was chosen, given that the industry tends to work with the most stable version available of any software and security updates are provided for a longer time.

The scenario used for the analysis consists of OpenPLC installed on Ubuntu Linux in a virtual machine, following the OSs shown in Table 5.4.

VERSION	RELEASE	OS	OS RELEASE
OpenPLC V1	2016/02/05	Ubuntu 14.04 LTS	2014/04/17
OpenPLC V2	2016/05/13	Ubuntu 16.04 LTS	2016/04/21
OpenPLC V3	2018/06/14	Ubuntu 18.04 LTS	2018/04/26

TABLE 5.4: Versions and release dates of OpenPLC and the available Ubuntu Linux LTS at that time for each date.

#### 5.2.3 Steps of the Analysis

The EDG model of OpenPLC was built following the steps described in Section 5.1.1.4 (see flowchart in Fig. 5.4). It is worth noting that these steps are followed for each one of the three versions available of OpenPLC, so an EDG is generated for each one. Obtaining the EDG for each version of OpenPLC will give information about the evolution of vulnerabilities over time.

For the sake of clarity and ease of analysis, the three obtained dependency graphs for each OpenPLC version were not merged into a single diagram. Fig. 5.7, Fig. 5.8, and Fig. 5.9, show the obtained EDG for versions V1, V2, and V3, respectively. In reality, these three diagrams would be the result of applying this method over time, updating the graph each time that a new vulnerability is discovered or a new patch/update is issued.

#### 5.2.4 Analysis

In this stage, the analysis of the SUT is performed based on the generated EDG and the value of the computed metrics. This process can be structured into three main steps, as follows:

- 1. **Analysis of the induced EDG model:** This step involves the analysis of the obtained directed-graph model. The structure, assets and dependencies are the focus of this first step.
- 2. **Vulnerability analysis:** Vulnerability number, distribution, and severity are analyzed in this step, which is supported by metrics. A proposal for vulnerability prioritization is also proposed.
- 3. **Root causes analysis (weaknesses):** Finally, the root cause of each vulnerability if found (related to the associated weakness). In this step, new requirements, test cases, and training activities are proposed based on the results of the analysis.

For each of the previously described steps, the analysis is done in both the spatial and temporal dimensions:

- 1. **Spatial Dimension:** This kind of analysis focuses on the distribution of vulnerabilities among the assets at a time *t*. In this example, this corresponds to independently analyzing each obtained EDG, because each one represents an instant in time (a different version of OpenPLC).
- 2. **Temporal Dimension:** This kind of analysis focuses on the evolution and distribution of vulnerabilities over time. In this example, this corresponds to

analyzing the changes in the number of assets and vulnerabilities, and their distribution over all three EDGs (over all three versions of OpenPLC).

#### 5.2.4.1 Analysis of the Induced EDG Model

OpenPLC V1 is analyzed in this subsection, focusing on the internal structure and dependencies among the assets.

The first result of the EDG model is thegraph obtained for OpenPLC V1 (Fig. 5.7). From the **spatial dimension** point of view, assets depend on a main service, server . js, based in Java. This web server offers a web GUI for the user to configure, start, and stop the PLC execution. Below this level, the main components of OpenPLC can be seen: OPLC Starter (responsible to start the OpenPLC and constantly monitor if it is running or not), OPLC Compiler (compiler from ladder logic to ANSI C code), and openplc (initialization procedures for the hardware, network and the main loop). The other assets are dynamic libraries of the system, such as libstdc++, libm, textttlibc (for C programming), and textttlibssl (C library for SSL and TLS).

Moreover, the libc library is a wrapper around the system calls of the Linux kernel, which provides and defines system calls and other basic functions. Thus, it is expected that all of the identified assets depend on this library in every version of OpenPLC. Fig. 5.7 shows that this is indeed the case.

From the **temporal dimension** point of view, OpenPLC V2 has to be analyzed in comparison to the previous version of OpenPLC to find changes in the structure and the number of assets. Fig. 5.8 shows the EDG for OpenPLC V2. Comparing both OpenPLC V1 and OpenPLC v2, it can be noted that their structure is very similar. In OpenPLC V2, new assets were introduced to provide new features to the project: Matiec compiler (which is an open-source compiler for programming languages defined in the IEC 61131-3 standard), ST optimizer (which is responsible for the optimization process after the initial compilation from OpenPLC Editor), Glue Generator (which is responsible for gluing the variables from the IEC program to the OpenPLC memory pointers), and OpenDNP3 (which is an implementation of the DNP3 protocol stack written in C++11).

In OpenPLC V2, and in comparison with OpenPLC V1, the OPLC compiler has disappeared and has been substituted by the Matiec Compiler, which supports all programming languages defined in the IEC 61131-3 standard. Moreover, the ST Optimizer and the Glue Generator were added to support the compilation process. Finally, the OpenDNP3 library has been added for high-performance applications, such as many concurrent TCP sessions.

Finally, and keeping the analysis in the temporal dimension, OpenPLC V3 can be compared to its ancestor, OpenPLC V2. Fig. 5.9 shows the EDG for OpenPLC V3. This last version of the project is the simplest. Now, the Java-based web server has been replaced by a Python-based web server, *webserver.py*. The main components of this version are: Matiec compiler, ST optimizer, Glue Generator, OpenDNP3, and LibModbus.

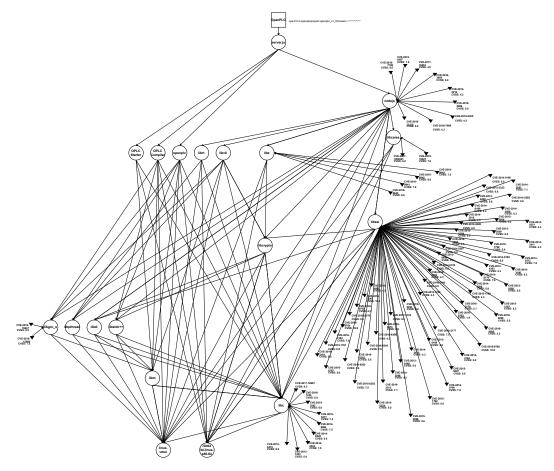


FIGURE 5.7: EDG for OpenPLC V1. Notice that, for simplicity, CWE, and CAPEC values are omitted, and only the CPE identifier of the SUT is shown.

#### 5.2.4.2 Vulnerability Analysis

In this step, all three available versions of the OpenPLC project are analyzed from a vulnerability perspective, both in the spatial and temporal dimensions. The goal here is to use the defined metrics as a support of the analysis, obtaining the number of vulnerabilities, their distribution, and their severity score. Table 5.5 shows the values of the proposed metrics for all three versions of OpenPLC. Finally, a ranked prioritization by severity for the vulnerabilities of each version is proposed.

From the **spatial dimension** perspective, just by observing the generated EDG for OpenPLC V1 (Fig. 5.7), it can be stated that:

• Most of the vulnerabilities are present in third-party open-source components, such as libssl ( $M_3 = 65$ ), and node-js ( $M_3 = 9$ ), having libssl the  $M_4 = 71.4\%$  of the vulnerabilities in this version. A *priori*, just by looking the number of vulnerabilities in each asset, it can be said that libssl is the most vulnerable asset in this version. Nevertheless, it is never enough to analyze the number of vulnerabilities ( $M_3$ ), and the CVSS score for each has to be taken into account, as well as the existence of known exploits (*e.g.*, updating the value of the CVSS using the temporal score).

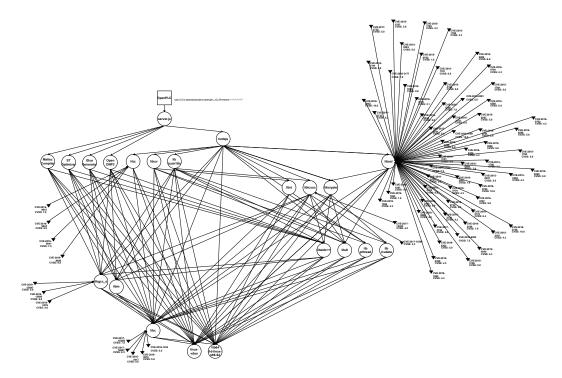


FIGURE 5.8: EDG for OpenPLC V2. Note that for the sake of simplicity, CWE, and CAPEC values are omitted and only the CPE identifier of the SUT is shown.

- The other vulnerabilities in OpenPLC V1 affect the assets of the operating system, such as libc ( $M_3 = 9$ ).
- The *ad hoc* assets developed for this project have no known vulnerabilities. This does not mean that they are secure, but rather that there are no known vulnerabilities available for them. Zero-day vulnerabilities could be present in the SUT.

Values of  $M_1 = 91$  suggests a high number of vulnerabilities in this project. This is expected given that OpenPLC V1 was the first version of OpenPLC. In the following versions, the value of  $M_1$  is expected to decrease.

The most striking fact when observing the EDG for OpenPLC V1 is the large number of vulnerabilities that are present in *libssl*. From this perspective, if a pen-tester were to evaluate OpenPLC V1, *libssl* would be a promising target. When the values of the CVSS are checked, three vulnerabilities<sup>2</sup> have a score of 10.0 out of 10.0. As can be seen, EDGs are a powerful tool for inspecting the structure of the SUT, and they can be used to analyze how the exploitation of a vulnerability can affect the rest of the SUT.

All of the assets point to *libc*, which is a wrapper around the system calls of the Linux kernel. An evaluator, even without this information, can understand the importance of this dependency with *libc* using the EDG. On closer inspection, the highest score of CVSS for the vulnerabilities of *libc* is 9.3 out of 10.0<sup>3</sup>. If the exploitation of this

<sup>&</sup>lt;sup>2</sup>CVE-2016-2842, CVE-2016-0799, and CVE-2016-0705.

<sup>&</sup>lt;sup>3</sup>CVE-2017-16997

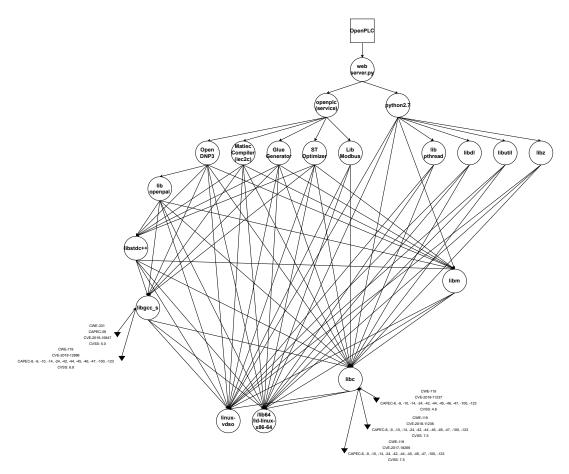


FIGURE 5.9: EDG for OpenPLC V3. Note that only the CPE of SUT is shown for the sake of simplicity.

vulnerability is possible, then it would allow a local user to gain privileges, exposing other assets of the SUT.

Using the value of severity of each vulnerability, it is possible to generate a list of vulnerabilities to be patched. This list can either be ordered by the global CVSS or by asset. Table 5.6 shows all the vulnerabilities whose CVSS value is between 6.0 and 10.0 ordered by asset and by descending CVSS. This can be used to decide in which order the vulnerabilities have to be patched, both at asset level or at SUT level.

From the information in Table 5.6, it is clear that all the three vulnerabilities with a CVSS value of 10.0 are in libssl, which makes this asset the most vulnerable by number of vulnerabilities and by CVSS. These vulnerabilities should be a priority during the patching stages. It is worth noting that libc, which is an important asset in Ubuntu Linux, has a vulnerability whose score is 9.3 (CVE-2017-16997). This should also be a priority when patching.

Moving now to OpenPLC V2, and comparing it with OpenPLC V1 (temporal dimension), similar conclusions can be drawn:

• Most of the vulnerabilities are in third-party open-source components. Moreover, libssl ( $M_3 = 63$ ) remains as the asset with the majority of vulnerabilities (with  $M_4 = 81.8\%$  of them), followed by libz ( $M_3 = 4$ ) from node-js.

METRIC																	
		libgcc_s	libc	libz		nodejs	libssl	others	libgcc_s	libc			libssl	others	libgcc_s		others
n(t)																	
$M_0(A) = \frac{ CVE_A(t) }{n(t)}$																	
$M_1(A,t) =  CVE_A(t) $					91							77				5	
$M_2(A) = \sum_{t=1}^{I}  CVE_A(t)  = \sum_{t=1}^{I}$	$_{1}M_{1}(A,t)$																
$M_3(A) =  CVE_{a_i}(t) $																	0
$M_4(a_k, t) = \frac{ CV L_{a_k}(t) }{\sum_{i=1}^{n}  CV L_{a_i}(t) }$		0.02	0.10	0.04	0.02	0.10	0.71	0.00	0.04	0.06	0.05	0.03	0.82	0.00	0.40	0.60	0.00
•	CWE-17	-	1	-	-	-	2	-	-	-	-	-	3	-	-	-	-
		-	-	-	-		-	-	-	-	-	-	-	-	-	-	-
		-	-	-	-	3	5	-	-	-	-	-	3	-	-	-	-
		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
		-	1	-	-	1	-	-	-	-	-	-	-	-	-	-	-
		1	5	-		1			1	3	-	1	6	-	1	3	-
		-	-	_	_	_		-	_	-	-	-		_	-	-	_
		-	_	4	_	_	2	-	-	-	4	_	4	-	_	-	_
	CWE-190	-	-	-	-	-	1	-	-	-	-	1	1	-	-	-	-
$M_{\tau}(q, cros, t) =  CVE  \cdot \cdot \cdot (t)$	CWE-200	-	-	-	1	3	5	-	1	-	-	-	12	-	-	-	-
$W15(u_i, cwe_j, t) =  C V L_{a_i} _{cwe_j}(t)$	CWE-295	-	-	-	-	-	-	-	-	-	-	-	1	-	-	-	-
		-	-	-	-	-		-	-	-	-	-		-	-	-	-
		-	-	-	-	-	-	-	-	-	-	-		-	-	-	-
		1	-	-	-	-	-	-	1	-	-	-		-	1	-	-
		1	-	-	-	-	4	-	1	-	-	-		-	1	-	-
				-				-		1	-					-	
		_	_	_	_	_		_	_	-	_	_		_	_	_	_
		-	1	_	-	_	-	-	-	1	_	_	-	_	-	-	_
	CWE-787	-	-	-	1	1	2	-	-	-	-	-	2	-	-	-	-
	CWE-NULL	-	-	-	-	-	12	-	-	-	-	-	10	-	-	-	-
												3				-	
												-				-	
												3				-	
												-				-	
												_				_	
									11						4		
	CWE-125				2							3				-	
	CWE-189				6							8				-	
	CWE-190				1							2				-	
$M_{\epsilon}(A.cme_i, t) \equiv  CVE_{Al}(t) $					9									-			
																-	
					12									-			
					-											-	
					1											1	
					-							-				-	
					8							7				-	
	CWE-400				1							1				_	
	CWE-426				1							1				-	
	CIVIE DOD	4					2					-					
	CWE-787															-	
$M_7(A,t) =  CWE_A(t) $	CWE-787 CWE-NULL				12 19							10 18				- 2	
	$M_1(A,t) =  CVE_A(t) $	$\begin{array}{c} n(t) \\ M_0(A) = \frac{ \nabla VE_A(t) }{n(t)} \\ M_1(A,t) =  \nabla VE_A(t)  \\ M_2(A) = \frac{ \nabla VE_A(t) }{n(t)} \\ M_2(A) = \frac{ \nabla VE_A(t) }{n(t)} \\ M_3(A) =  \nabla VE_{a_1}(t)  \\ M_4(a_k,t) = \frac{ \nabla VE_{a_1}(t) }{ \nabla VE_{a_1}(t) } \\ (CWE-19) \\ (CWE-12) \\ (CWE-13) \\ (CWE-13) \\ (CWE-14) \\ (CWE-14) \\ (CWE-15) \\ (CWE-15) \\ (CWE-16) \\ (CWE-16) \\ (CWE-17) \\ (CWE-19) \\ (CWE-19) \\ (CWE-19) \\ (CWE-19) \\ (CWE-19) \\ (CWE-19) \\ (CWE-10) \\ (CWE-10) \\ (CWE-10) \\ (CWE-20) \\ (CWE-20) \\ (CWE-30) \\ (CWE-30) \\ (CWE-30) \\ (CWE-30) \\ (CWE-10) \\ (CWE-11) \\ (CWE-12) \\ (CWE-12) \\ (CWE-13) \\ (CWE-12) \\ (CWE-14) \\ (CWE-12) \\ (CWE-12) \\ (CWE-13) \\ (CWE-13) \\ (CWE-14) \\ (CWE-12) \\ (CWE-12) \\ (CWE-13) \\ (CWE-13) \\ (CWE-14) \\ (CWE-12) \\ (CWE-12) \\ (CWE-13) \\ (CWE-13) \\ (CWE-13) \\ (CWE-13) \\ (CWE-13) \\ (CWE-14) \\ (CWE-15) \\ (CWE-16) \\ (CWE-16) \\ (CWE-16) \\ (CWE-17) \\ (CWE-17) \\ (CWE-18) \\ (CWE-113) \\ (CWE-130) \\ (CWE-30) \\ (CWE-30) \\ (CWE-30) \\ (CWE-30) \\ (CWE-30) \\ (CWE-311) \\ (CWE-320) \\ (CWE-331) \\ (CWE-332) \\ (CWE-331) \\ (CWE-332) \\ (CWE-331) \\ (CWE-332) \\ (CWE-3331) \\ (CWE-332) \\ (CWE-332) \\ (CWE-3331) \\ (CWE-332) \\ (CWE-332) \\ (CWE-3331) \\ (CWE-332) \\ (CWE-332) \\ (CWE-3331) \\ (CWE-332) \\ $	$\begin{array}{c c} & & \text{libgec_s} \\ \hline m(t) & & & \\ \hline M_0(A) = \frac{ CVE_A(t) }{n(t)} & & \\ \hline M_1(A,t) =  CVE_A(t)  & & \\ \hline M_2(A) = \sum_{i=1}^{ CVE_A(t) } & & \\ \hline M_2(A) = \sum_{i=1}^{ CVE_A(t) } & & \\ \hline M_2(A) = \sum_{i=1}^{ CVE_A(t) } & & \\ \hline M_2(A) = \frac{ CVE_A(t) }{\sum_{i=1}^{ CVE_A(t) }} & & \\ \hline CWE-17 & & \\ \hline CWE-19 & - \\ \hline CWE-20 & - \\ \hline CWE-22 & - \\ \hline CWE-21 & & \\ \hline CWE-113 & - \\ \hline CWE-113 & - \\ \hline CWE-119 & 1 \\ \hline CWE-119 & 1 \\ \hline CWE-125 & - \\ \hline CWE-190 & - \\ \hline CWE-190 & - \\ \hline CWE-200 & - \\ \hline CWE-200 & - \\ \hline CWE-200 & - \\ \hline CWE-310 & - \\ \hline CWE-311 & - \\ \hline CWE-320 & - \\ \hline CWE-331 & - \\ \hline CWE-320 & - \\ \hline CWE-331 & - \\ \hline CWE-320 & - \\ \hline CWE-331 & - \\ \hline CWE-320 & - \\ \hline CWE-300 & - \\ \hline CWE-200 & - \\ \hline CWE-125 & - \\ \hline CWE-113 & - \\ \hline CWE-190 & - \\ \hline CWE-190 & - \\ \hline CWE-190 & - \\ \hline CWE-200 & - \\ \hline CWE-200 & - \\ \hline CWE-200 & - \\ \hline CWE-301 & - \\ \hline CWE-302 & - \\ \hline CWE-301 & - \\ \hline CWE-302 & - \\ \hline CWE-302 & - \\ \hline CWE-3031 & - \\ \hline CWE-302 & - \\ \hline CWE-302 & - \\ \hline CWE-3031 & - \\ \hline CWE-302 & -$	$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	$ \begin{array}{c c c c c c c c c c c c c c c c c c c $	$\begin{array}{c c c c c c c c c c c c c c c c c c c $	$ \begin{array}{c c c c c c c c c c c c c c c c c c c $					H				

TABLE 5.5: Metric values for each asset and version of OpenPLC. Notice that "CWE-NULL" refers to a void value of CWE for a certain CVE value.

• The remaining vulnerabilities are related to the operating system, just as in OpenPLC V1, libc ( $M_3 = 5$ ).

For this version,  $M_1 = 77$ , less than  $M_1 = 98$  for the previous version. As was expected, the total number of vulnerabilities has decreased. From the **spatial dimension**, it is worth noting that libssl is still the asset with the highest number of vulnerabilities and is still a promising target for a pen-tester or an attacker. Nevertheless, the number of vulnerabilities for libssl ( $M_3$ ) has decreased from OpenPLC V1 to OpenPLC V2, as it was expected.

As was done earlier, a list of vulnerabilities can be generated that is ordered by asset and by descending CVSS. Table 5.7 shows all of the vulnerabilities for OpenPLC V2 whose CVSS value is between 6.0 and 10.0.

Table 5.7 shows that libssl now has four instead of three vulnerabilities whose CVSS value is 10.0 (CVE-2016-0799, CVE-2016-2108, CVE-2016-0705, CVE-2016-2842). In OpenPLC V2, libc is still affected by the same vulnerability with a score of 9.3 (CVE-2017-16997). All of these vulnerabilities should be a priority during patching activities.

TABLE 5.6: Vulnerability prioritization by asset and by CVSS for Open-PLC V1. Only the vulnerabilities whose value is between 6.0 and 10.0 are shown.

CVE	CVSS	ASSET
CVE-2019-15847	7.5	libcares
CVE-2018-12886	7.5	libgcc
CVE-2016-9843	7.5	libz
CVE-2016-9841	7.5	libz
CVE-2016-9840	6.8	libz
CVE-2016-9842	6.8	libz
CVE-2017-16997	9.3	libc
CVE-2014-9984	7.5	libc
CVE-2014-4043	7.5	libc
CVE-2015-5277	7.2	libc
CVE-2015-7547	6.8	libc
CVE-2014-0475	6.8	libc
CVE-2016-2842	10.0	libssl
CVE-2016-0705	10.0	libssl
CVE-2016-0799	10.0	libssl
CVE-2016-6304	7.8	libssl
CVE-2016-0798	7.8	libssl
CVE-2014-8176	7.5	libssl
CVE-2016-2182	7.5	libssl
CVE-2014-3512	7.5	libssl
CVE-2016-6303	7.5	libssl
CVE-2015-0292	7.5	libssl
CVE-2016-2177	7.5	libssl
CVE-2014-3567	7.1	libssl
CVE-2014-3513	7.1	libssl
CVE-2015-1791	6.8	libssl
CVE-2012-2333	6.8	libssl
CVE-2015-0209	6.8	libssl
CVE-2014-3509	6.8	libssl
CVE-2014-0195	6.8	libssl
CVE-2014-3505	5.0	libssl

Finally, OpenPLC V3, which is the last available version, is analyzed. The most striking fact when comparing OpenPLC V3 with the other versions (**temporal dimension**) is the significant reduction in the global number of vulnerabilities. This effect is caused by the absence of libssl in this version.

Focusing on the **spatial dimension** for this version, vulnerabilities in OpenPLC V3 are only related to assets in the operating system,  $libgcc\_s$  ( $M_3 = 2$ ), and libc ( $M_3 = 3$ ). In OpenPLC V3, neither third-party open-source assets nor *ad hoc* assets have any known vulnerability. In fact, OpenPLC V3 has the lowest values of  $M_1 = 5$  for the whole life cycle of the project, and therefore has the lowest number of vulnerabilities.

Table 5.8 shows an ordered list of all the vulnerabilities in OpenPLC V3 whose CVSS value is between 6.0 and 10.0.

For OpenPLC V3, the most vulnerable asset is 1ibc, according to the CVSS score shown in Table 5.8.

CVE	CVSS	ASSET
CVE-2018-12886	6.8	libgcc_s
CVE-2017-16997	9.3	libc
CVE-2017-18269	7.5	libc
CVE-2016-9843	7.5	libz
CVE-2016-9841	7.5	libz
CVE-2016-9842	6.8	libz
CVE-2016-9840	6.8	libz
CVE-2016-0799	10.0	libssl
CVE-2016-2108	10.0	libssl
CVE-2016-0705	10.0	libssl
CVE-2016-2842	10.0	libssl
CVE-2016-0798	7.8	libssl
CVE-2016-6304	7.8	libssl
CVE-2016-2109	7.8	libssl
CVE-2016-2177	7.5	libssl
CVE-2016-2182	7.5	libssl
CVE-2016-6303	7.5	libssl
CVE-2015-0209	6.8	libssl
CVE-2015-1791	6.8	libssl
CVE-2016-2106	6.5	libssl
CVE-2016-2176	6.4	libssl

TABLE 5.7: Vulnerability prioritization by asset and by CVSS for Open-PLC V2. Only the vulnerabilities whose value is between 6.0 and 10.0 are shown.

TABLE 5.8: Vulnerability prioritization by asset and by CVSS for Open-PLC V3. Only the vulnerabilities whose value is between 6.0 and 10.0 are shown.

CVE	CVSS	ASSET
CVE-2018-12886	6.8	libgcc_s
CVE-2018-11236	7.5	libc
CVE-2017-18269	7.5	libc

# 5.2.4.3 Root Causes Analysis (Weaknesses)

In this last step of the analysis, the information about the root causes of the vulnerabilities is extracted. This can be achieved through weaknesses. Each CWE contains data about the main issue to be solved for any given vulnerability. This information is useful to propose new requirements, test cases, and training. The analysis of the root causes will also be carried out in both the spatial and temporal dimensions.

The analysis of the root causes starts by extracting the weaknesses for each vulnerability for all three versions of OpenPLC. All of the extracted weaknesses for all three available versions of OpenPLC are shown in Fig. 5.10 and in Table 5.9.

Focusing on the spatial dimension, OpenPLC V1 has the widest range of weaknesses ( $M_7 = 19$ ), and therefore has the widest range of root causes. From the total number of vulnerabilities in OpenPLC V1 ( $M_3 = 91$ ),  $M_6 = 15$  of them are related to weaknesses CWE-119<sup>4</sup>: "The software performs operations on a memory buffer, but it can read

<sup>&</sup>lt;sup>4</sup>Description for CWE-119: https://cwe.mitre.org/data/definitions/119.html

	METRIC	OpenPLC V1	OpenPLC V2	OpenPLC V3	$\sum_{i=1}^{T} M_8(A, cwe_j, t)$
	CWE-17	3	3	-	6
	CWE-19	1	-	-	1
	CWE-20	8	3	-	11
	CWE-22	1	-	-	1
	CWE-94	1	-	-	1
$\overline{}$	CWE-113	1	-	-	1
$_{j}$ $(t$	CWE-119	15	11	4	30
cwe	CWE-125	2	3	-	5
$E_A$	CWE-189	6	8	-	14
$\Lambda$	CWE-190	1	2	-	3
<u> </u>	CWE-200	9	13	-	22
· ·	CWE-295	-	1	-	1
2 j, t	CWE-310	12	5	-	17
$M_8(A, cwe_j, t) =  CVE_{A cwe_j}(t) $	CWE-311	-	2	-	2
A,	CWE-320	-	3	-	3
$I_8$	CWE-331	1	1	1	3
~	CWE-362	4	1	-	5
	CWE-399	8	7	-	15
	CWE-400	1	1	-	2
	CWE-426	1	1	-	2
	CWE-787	4	2	-	6
	CWE-NULL	12	10	-	22
$M_7(A)$	$A,t) =  CWE_A(t) $	19	18	2	22

TABLE 5.9: Weakness analysis for all versions of OpenPLC.

from or write to a memory location that is outside the intended boundary of the buffer."  $M_6=12$  of them are related to CWE-310  $^5$ : "Weaknesses in this category are related to the design and implementation of data confidentiality and integrity. Frequently, these deal with the use of encoding techniques, encryption libraries, and hashing algorithms. The weaknesses in this category could lead to a degradation of the quality of data if they are not addressed."

Repeating this same analysis with OpenPLC V2, we can see that the number of different weaknesses ( $M_7 = 18$ ) is almost the same as before. But in this version, the most repeated weakness is CWE-200<sup>6</sup>: "The product exposes sensitive information to an actor that is not explicitly authorized to have access to that information." Followed by CWE-119.

Comparing the previous versions with OpenPLC V3, we find now that the total number of weaknesses ( $M_7 = 2$ ) has been reduced drastically. Nevertheless, weakness CWE-119 is still present in this version, and is the most repeated.

This analysis can also be done by combining the information of all three versions, drawing conclusions for the whole life cycle of OpenPLC. Requirements and test cases can be extracted from the most recurrent weaknesses, and training can be proposed to avoid weaknesses that repeat over time. In our example, CWE-119 has the greatest

<sup>&</sup>lt;sup>5</sup>Description for CWE-310: https://cwe.mitre.org/data/definitions/310.html.

<sup>&</sup>lt;sup>6</sup>Description for CWE-200: https://cwe.mitre.org/data/definitions/200.html

frequency over time ( $\sum_{i=1}^T M_6 = 30$ ), followed by CWE-200 ( $\sum_{i=1}^T M_6 = 22$ ), and CWE-310 ( $\sum_{i=1}^T M_6 = 17$ ). Table 5.10 shows the generated requirements, Table 5.11 shows the proposed training, and Table 5.12 shows the proposed test cases for OpenPLC.

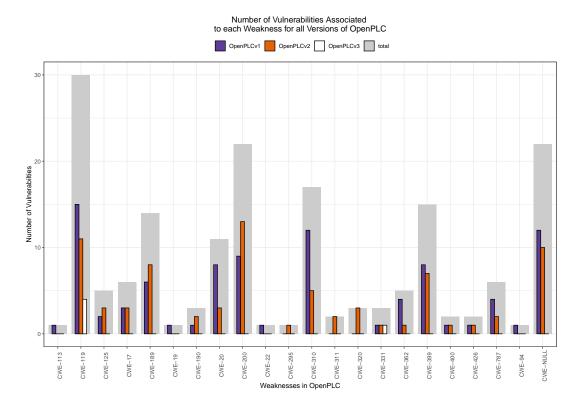


FIGURE 5.10: Evolution of the number of vulnerabilities over consecutive versions of OpenPLC.

Table 5.10: An example of generated requirements for OpenPLC.

CWE ID	REQUIREMENTS
CWE-19, CWE-20, CWE-119, CWE-125, CWE-189, CWE-190, CWE-362, CWE-399, CWE-400, CWE-787	Use languages that perform their own memory management.
CWE-19, CWE-20, CWE-119, CWE-125, CWE-189, CWE-190, CWE-787	Use libraries or frameworks that make it easier to handle numbers without unexpected consequences. Examples include safe integer handling packages such as SafeInt (C++) or IntegerLib (C or C++).
CWE-19, CWE-20, CWE-119, CWE-125, CWE-189, CWE-190, CWE-200, CWE-362, CWE-787	Use a CPU and operating system that offers Data Execution Protection (NX) or its equivalent.
CWE-19, CWE-20, CWE-125, CWE-189, CWE-190, CWE-200, CWE-362	Ensure that all protocols are strictly defined, such that all out- of-bounds behaviors can be identified simply, and require strict conformance to the protocol.
CWE-19, CWE-20, CWE-399, CWE-400	Run or compile the software using features or extensions that randomly arrange the positions of a program's executable and libraries in memory ( <i>e.g.</i> , Address Space Layout Randomization (ASLR), or Position-Independent Executables (PIE)).
CWE-19, CWE-295, CWE-310, CWE-311, CWE-320, CWE-331	Clearly specify which data or resources are valuable enough that they should be protected by encryption. Require that any transmission or storage of this data/resource should use well-vetted encryption algorithms. Up-to-date algorithms must be used, and the entropy of the keys must be sufficient for the application.
CWE-22, CWE-94, CWE-113, CWE-399, CWE-400, CWE-426	Hard-code the search path to a set of known-safe values (such as system directories), or only allow them to be specified by the administrator in a configuration file. Do not allow these settings to be modified by an external party.
CWE-20, CWE-22, CWE-94, CWE-113, CWE-399, CWE-400, CWE-426	Use an input validation framework such as Struts or the OWASP ESAPI Validation API.
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-295, CWE-426	Assume all input is malicious. Use an "accept known good" input validation strategy, i.e., use a list of acceptable inputs that strictly conform to specifications. Reject any input that does not strictly conform to specifications, or transform it into something that does.
CWE-19, CWE-20, CWE-119, CWE-125, CWE-399, CWE-400, CWE-787	Run or compile the software using features or extensions that automatically provide a protection mechanism that mitigates or eliminates buffer overflows.
CWE-19, CWE-22, CWE-119, CWE-125, CWE-399, CWE-400, CWE-787	Replace unbounded copy functions with analogous functions that support length arguments, such as strcpy with strncpy. Create these if they are not available.

TABLE 5.11: Example of proposed training for OpenPLC.

CWE ID	TRAINING
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-119, CWE-125, CWE-426	Identification of all potentially relevant properties of an input (length, type of input, the full range of acceptable values, missing or extra inputs, syntax, consistency across related fields).
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-119, CWE-125, CWE-426	Input validation strategies.
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-119, CWE-125, CWE-200, CWE-426	Allowlists and Denylists.
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-119, CWE-125, CWE-189, CWE-190, CWE-426	Character encoding compatibility.
CWE-19, CWE-20, CWE-94, CWE-113, CWE-119, CWE-125, CWE-426	Buffer overflow detection during compilation ( <i>e.g.</i> , Microsoft Visual Studio /GS flag, Fedora/Red Hat FORTIFY_SOURCE GCC flag, StackGuard, and ProPolice).
CWE-19, CWE-20, CWE-94, CWE-113, CWE-119, CWE-125, CWE-200, CWE-426	Secure functions, such as <i>strcpy</i> with <i>strncpy</i> . Create these if they are not available.
CWE-19, CWE-20, CWE-94, CWE-113, CWE-119, CWE-125, CWE-189, CWE-190, CWE-426,CWE-787	Secure programming: memory management.
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-119, CWE-125, CWE-189, CWE-190, CWE-787	Understand the programming language's underlying representation and how it interacts with numeric calculation.
CWE-19, CWE-20, CWE-22, CWE-94, CWE-113, CWE-119, CWE-125	System compartmentalization.
CWE-200, CWE-295, CWE-310, CWE-311, CWE-320, CWE-331	Certificate management.
CWE-200, CWE-295, CWE-310, CWE-311, CWE-320, CWE-331	Certificate pinning.
CWE-310, CWE-311, CWE-320, CWE-331	Encryption integration (Do not develop custom or private cryptographic algorithms). $ \\$
CWE-310, CWE-311, CWE-320, CWE-331	Secure up-to-date cryptographic algorithms.
CWE-189, CWE-200, CWE-362, CWE-399, CWE-400, CWE-426	Shared resources management.
CWE-200, CWE-362, CWE-399, CWE-400, CWE-426	Thread-safe functions.

TABLE 5.12: Example of generated test cases for OpenPLC.

CAPEC ID	TEST CASES
CAPEC-10	Check for buffer overflows through manipulation of environment variables. This test leverages implicit trust often placed in environment variables.
CAPEC-14	Static analysis of the code: secure functions and buffer overflow.
CAPEC-24	Feed overly long input strings to the program in an attempt to overwhelm the filter (by causing a buffer overflow) and hoping that the filter does not fail securely (i.e. the user input is let into the system unfiltered)
CAPEC-45	This test uses symbolic links to cause buffer overflows. The evaluator can try to create or manipulate a symbolic link file such that its contents result in out of bounds data. When the target software processes the symbolic link file, it could potentially overflow internal buffers with insufficient bounds checking.
CAPEC-46	Static analysis of the code: secure functions and buffer overflow.
CAPEC-47	In this test, the target software is given input that the evaluator knows will be modified and expanded in size during processing. This test relies on the target software failing to anticipate that the expanded data may exceed some internal limit, thereby creating a buffer overflow.
CAPEC-59	This test targets predictable session ID in order to gain privileges. The evaluator can try to predict the session ID used during a transaction to perform spoofing and session hijacking.
CAPEC-97	Automated measurement of the entropy of the keys generated.
CAPEC-475	Fuzz testing to externally provided data, such as directories and filenames.

#### 5.2.5 Second Scenario

For this use case, the setup consisted of OpenPLC installed on 14.04 LTS Ubuntu Linux in a virtual machine. All configuration options were by default.

# 5.2.6 Building the EDG

Using the generated EDG for OpenPLC V1 shown in Figure 5.7, we extracted the information about security updates (discarding updates that introduced more functionalities), for both libssl and nodejs. Table 5.13 shows the security updates and their date of availability for both libssl [193] and nodejs [194] for Ubuntu 14.04 LTS. There were two security updates available for the amd64 architecture for each asset.

ASSET	$1_{st}UPDATE$	SOLVED VULNERABILITIES (CVSS)	$2_{nd}UPDATE$	SOLVED VULNERABILITIES (CVSS)
libssl	2014/04/07	CVE-2014-0076 (1.9) CVE-2014-0160 (5.0)	2018/12/06	CVE-2018-5407 (1.9) CVE-2018-0734 (4.3)
nodeis	2014/03/27	_	2018/08/10	CVE-2016-5325 (4.3)

TABLE 5.13: Update information of both libssl and nodejs.

From this data, we can extract that:

- Updates for nodejs were released before the updates for libssl.
- libssl shows more vulnerabilities than node js.
- The highest CVSS score in the period of this analysis is 5.

Then, the EDG for these two assets and their updates were built. Figure 5.11 shows the updates over time of the EDG, whereas Figure 5.12 shows the final EDG with all the information included.

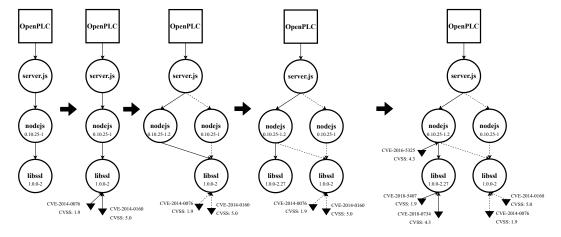


FIGURE 5.11: Temporal evolution of the EDG for OpenPLC V1 for both libss and nodejs.

# 5.2.7 Analysis of the EDG

Using Figure 5.12, and Table 5.14, we can analyze the obtained EDG:

1. **Analysis of the induced EDG model:** The structure, assets, and dependencies are the focus of this first step.

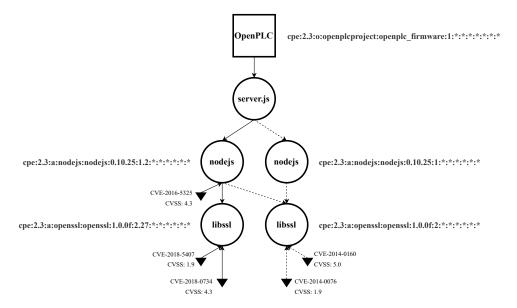


FIGURE 5.12: Final EDG for libssl and nodejs integrating all the updates for Ubuntu Linux 14.04 for amd64 architecture.

We can observe that libssl is used by nodejs, and they are not at the same level of the hierarchy. So vulnerabilities could propagate upwards and downwards through the EDG.

2. **Vulnerability analysis:** Vulnerability number, distribution, and severity are analyzed in this step. A proposal for vulnerability prioritization is also generated.

We can highlight that nodejs had one vulnerability discovered after its first update, whereas libssl had vulnerabilities in both periods of time. We could argue that, as nodejs is the most accessible asset from the exterior, its vulnerabilities should be first addressed, even though the associated CVSS is not the highest one.

3. **Weaknesses analysis:** Finally, the root cause of each vulnerability is found. In this step, new requirements, test cases, and training activities are proposed based on the results of the analysis.

Table 5.14 shows the root cause for each vulnerability. Using this data, new requirements (Table 5.15), test cases (Table 5.16) and training activities (Table 5.17) were proposed.

It is worth noticing that this use case is focused on reflecting the temporal evolution of the EDG. For this reason, metrics cannot be computed here, because of the low number of vulnerabilities available.

Table 5.14: Relationship between vulnerabilities and weaknesses for both libssl and nodejs.

CVE	CVSS	CWE	DESCRIPTION
CVE-2014-0076	1.9	CWE-310	Cryptographic Issues
CVE-2014-0160	7.5	CWE-119	Improper Restriction of Operations within the Bounds of a Memory Buffer
CVE-2016-5325	6.1	CWE-113	Improper Neutralization of CRLF Sequences in HTTP Headers ('HTTP Response Splitting')
CVE-2018-0734	5.9	CWE-327	Use of a Broken or Risky Cryptographic Algorithm
CVE-2018-5407	4.7	CWE-203 CWE-200	Observable Discrepancy Exposure of Sensitive Information to an Unauthorized Actor

Table 5.15: An example of generated requirements for OpenPLC V1.

CWE ID	REQUIREMENTS
CWE-119	Use languages that perform their own memory management.
CWE-119	Use libraries or frameworks that make it easier to handle numbers without unexpected consequences. Examples include safe integer handling packages such as SafeInt (C++) or IntegerLib (C or C++).
CWE-119, CWE-200	Use a CPU and operating system that offers Data Execution Protection (NX) or its equivalent.
CWE-190, CWE-200	Ensure that all protocols are strictly defined, such that all out-of-bounds behaviors can be identified simply, and require strict conformance to the protocol.
CWE-310	Clearly specify which data or resources are valuable enough that they should be protected by encryption. Require that any transmission or storage of this data/resource should use well-vetted encryption algorithms. Up-to-date algorithms must be used, and the entropy of the keys must be sufficient for the application.
CWE-113	Use an input validation framework such as Struts or the OWASP ESAPI Validation API.
CWE-113	Assume all input is malicious. Use an "accept known good" input validation strategy, i.e., use a list of acceptable inputs that strictly conform to specifications. Reject any input that does not strictly conform to specifications, or transform it into something that does.
CWE-113	Hard-code the search path to a set of known-safe values (such as system directories), or only allow them to be specified by the administrator in a configuration file. Do not allow these settings to be modified by an external party.
CWE-119	Run or compile the software using features or extensions that automatically provide a protection mechanism that mitigates or eliminates buffer overflows.
CWE-119	Replace unbounded copy functions with analogous functions that support length arguments, such as strcpy with strncpy. Create these if they are not available.

TABLE 5.16: Example of generated test cases for OpenPLC V1.

CAPEC ID	TEST CASES
CAPEC-119	Check for buffer overflows through manipulation of environment variables. This test leverages implicit trust often placed in environment variables.
CAPEC-119	Static analysis of the code: secure functions and buffer overflow.
CAPEC-119	Feed overly long input strings to the program in an attempt to overwhelm the filter (by causing a buffer overflow) and hoping that the filter does not fail securely (i.e. the user input is let into the system unfiltered)
CAPEC-119	This test uses symbolic links to cause buffer overflows. The evaluator can try to create or manipulate a symbolic link file such that its contents result in out of bounds data. When the target software processes the symbolic link file, it could potentially overflow internal buffers with insufficient bounds checking.
CAPEC-119	Static analysis of the code: secure functions and buffer overflow.

TABLE 5.17: Example of proposed training for OpenPLC V1.

CWE ID	TRAINING
CWE-113, CWE-119	Identification of all potentially relevant properties of an input (length, type of input, the full range of acceptable values, missing or extra inputs, syntax, consistency across related fields).
CWE-113, CWE-119	Input validation strategies.
CWE-113, CWE-119, CWE-200	Allowlists and Denylists.
CWE-113, CWE-119	Character encoding compatibility.
CWE-113, CWE-119	Buffer overflow detection during compilation ( <i>e.g.</i> , Microsoft Visual Studio /GS flag, Fedora/Red Hat FORTIFY_SOURCE GCC flag, StackGuard, and ProPolice).
CWE-113, CWE-119CWE-200	Secure functions, such as <i>strcpy</i> with <i>strncpy</i> . Create these if they are not available.
CWE-113, CWE-119CWE-190	Secure programming: memory management.
CWE-113, CWE-119	Understand the programming language's underlying representation and how it interacts with numeric calculation.
CWE-113, CWE-119	System compartmentalization.
CWE-200, CWE-310	Certificate management.
CWE-200, CWE-310	Certificate pinning.
CWE-310	Encryption integration (Do not develop custom or private cryptographic algorithms).
CWE-310	Secure up-to-date cryptographic algorithms.
CWE-200	Shared resource management.
CWE-200	Thread-safe functions.

# 5.3 Conclusions and Future Work

Vulnerability analysis is a critical task which ensures the security of industrial components. The EDG model that we propose performs continuous vulnerability assessment throughout the entire life cycle of industrial components. The model is built on a directed graph-based structure and a set of metrics based on globally accepted security standards. Metrics can be used by the model to improve the development process of the SUT, enhance its security, and track its status. The key feature of the proposed model is the addition of the temporal dimension in the analysis of vulnerabilities. The location of vulnerabilities can be analyzed in both space (in which asset) and time (their recurrence), which allows the state of the device to be tracked throughout the whole life cycle.

The model was successfully applied to the OpenPLC use case, which demonstrated its advantages, applicability, and potential. The use case showed that the model can assist in updating management activities, applying patching policies, launching training activities, and generating new test cases, and requirements. This has significant implications for cybersecurity evaluators, as it can serve as a starting point for identifying vulnerabilities, weaknesses, and attack patterns.

Further research will enhance the EDG by adding a mathematical model to aggregate the values of the CVSS metric for each asset, and a value for the whole SUT. This will enable the comparison of different SUTs over time. More improvements will be made in the prioritization of patching, taking into account the context and the functionalities of the SUT. Finally, historical information about the developers can be integrated into the EDG model to predict future vulnerabilities.

# Chapter 6

# Gotta Catch 'em All: Aggregating CVSS Scores

In this chapter, we propose a novel aggregation algorithm for a set of CVSS values which takes into account context-related properties of the SUT.<sup>1</sup>. This approach is based on the EDGs proposed in Chapter 5. Because EDGs are capable of modeling dependencies, this algorithm can also be applied to computer networks. Our proposal is capable of selecting the most relevant CVSS to be aggregated, taking into account four different context-related properties of the SUT:

- 1. Functionality disruption.
- 2. Exploitation difficulty.
- 3. Existence of exploits, and their development state.
- 4. Context of deployment.

This approach increases the granularity of the CVSS base, environment and temporal metrics, where not every possible value in the scale [0,10] is achievable, or the result of changing the value of a submetric has almost no effect on the final CVSS [164,195]. Moreover, our proposal is capable of detecting which branch in the EDG, (in the case of an ES the most critical set of assets), is contributing the most (more critical) to the final score.

<sup>&</sup>lt;sup>1</sup>The Python code implementing the aggregation algorithm is available at GitHub https://github.com/aaalongueira/CVSS\_Aggregation.

# 6.1 Proposed Approach for Metric Aggregation

In this section, we propose a CVSS aggregation algorithm inspired by the risk propagation formula [196] described in MAGERIT [197,198]. First, we describe the corrections factors involved in our proposal. Then, the aggregation formula is introduced. Finally, the algorithm and the interpretation of the results is explained in detail.

#### 6.1.1 Correction Factors

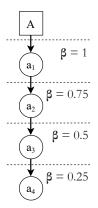
The proposed aggregation algorithm integrates correction factors to adapt the formula described in MAGERIT. These correction factors apply individually for each CVSS, except for the average and summarized factors. Correction factors are summarized in Table 6.1.

TABLE 6.1: Correction factors proposed for adapting the Bayesian sum proposed in MAGERIT.

CORRECTION FACTOR	DESCRIPTION			
Functionality factor $(\rho)$	Binary value indicating whether a vulnerability affects			
	or not the functionality of the SUT.			
Deepness factor ( $\beta$ )	Value between [0, 1] proportional to the position of the			
	affected asset in the EDG of the SUT.			
Context factor $(\gamma)$	Binary value indicating whether the vulnerability can			
	be exploited in the real and particular conditions of the			
	SUT.			
Exploit factor ( $\mu$ )	This factor accounts for the existence of a public exploit			
	for a given vulnerability, being proportional to its state			
	of development: Not defined ( $\mu = 0.5$ ), Theoretical			
	( $\mu = 1.25$ ), Proof-Of-Concept ( $\mu = 1.5$ ), Functional			
	$(\mu = 1.75)$ , and Automated $(\mu = 2)$ .			
Summarized factor ( $\lambda$ )	This factor summarizes the effect of all the above ones,			
	$\lambda = \rho \beta \gamma \mu \sigma.$			
Average factor $(\sigma)$	Function that adjust the value of the sum to avoid its			
	rapid evolution to 10.			

- 1. Functionality factor ( $\rho$ ): This correction factor represents whether any functionality of the systems is affected by its vulnerabilities.
  - It is represented by a binary value, being 0 when no functionality is affected, and 1 when any of them is affected.
  - For example, a cryptographic library with a vulnerability in SHA1. If the SUT does not make use of SHA1 in any way, the vulnerability would not be exploitable, and could be removed from the analysis ( $\rho = 0$ ).
- 2. Deepness Factor ( $\beta$ ): This factor represents the difficulty of chained exploitation of each vulnerability. It is represented by a value between [0, 1] inversely proportional to the amount of assets to compromise in order to exploit vulnerability. Vulnerabilities close to the entry point will account more for the final aggregation, whereas those that are far away will account less. In this approach, linear

interpolation is proposed to calculate the weight of each layer, because of its simplicity. Nevertheless, different interpolations could be used according to the criticality of the system. Figure 6.1 shows the corresponding  $\beta$  for a four-layer system.



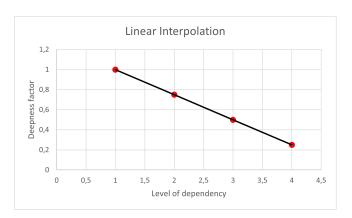


FIGURE 6.1: Calculation of the deepness factor for a four-layer of dependency example.

- 3. Context factor ( $\gamma$ ): This factor considers whether the exploitation of a vulnerability is actually possible in the real scenario where the system is deployed. It is represented by a binary value, where 0 indicated that it is not possible, and 1, that it is possible. It is calculated comparing the attack vector of the CVSS with the real conditions where the device is deployed. For example, this can happen when a vulnerability with a high CVSS score needs physical access to be exploited, but in reality the device is physically isolated. To reflect this, the CVSS should be updated, lowering the resulting value [165]. This factor aims to complement the existing submetrics in the temporal and the environment metrics of the CVSS. Both the temporal and the environmental scores lack of an "isolated" value for the attack vector.
- 4. *Exploit factor* ( $\mu$ ): This factor accounts for the existence of a public exploit for a given vulnerability, being proportional to its state of development. The temporal score of the CVSS already implements this feature, but the CVSS values are not updated in practice [195]. Moreover, taking into account the temporal score has almost no effect as opposed to using the raw initial base score. This means that a CVSS just considering the base score is higher than a CVSS considering an exploit code maturity of "functional exploit exists". To solve this issue, we introduce the following values for the exploit factor: Not defined ( $\mu$  = 0.5), Theoretical ( $\mu$  = 1.25), Proof-Of-Concept ( $\mu$  = 1.5), Functional ( $\mu$  = 1.75), and Automated ( $\mu$  = 2). These values are equivalent to the scale defined in the CVSS Specification Document [67].
- 5. *Summarized factor* ( $\lambda$ ): The  $\lambda$  factor accounts for the effect of all the factors above:

$$\lambda = \rho \beta \gamma \mu \tag{6.1}$$

6. Average factor ( $\sigma$ ): This factor defines the behavior of the aggregation function. It can be chosen as needed (e.g., the arithmetic or harmonic mean), but taking into account all the values to be added.

# 6.1.2 Aggregation Formula

The aggregation function is defined as:

$$\Gamma(\overrightarrow{V}) = 10 - \frac{1}{\sigma} f(\overrightarrow{V}) \tag{6.2}$$

Where  $\overrightarrow{V}$  is a vector  $(cvss_0, cvss_1, \dots, cvss_n)$  with all the corrected CVSS values to be added, cvss, being n the last value to be added.  $f(\overrightarrow{V}) = a_n$  is defined as the following recursive function:

$$a_n = 10 \left[ 1 - \left( 1 - \frac{\lambda_{a_{n-1}}}{10} a_{n-1} \right) \cdot \left( 1 - \frac{\lambda_{cvss_n}}{10} cvss_n \right) \right]$$

$$(6.3)$$

Where the base case is defined as:

$$a_0 = \lambda_{cvss_0} cvss_0 \tag{6.4}$$

## 6.1.3 Algorithm

The proposed aggregation algorithm is divided into the following steps (see Figure 6.2):

- 1. Calculation of the correction factors for each CVSS,
- 2. Calculation of the summarized factor for each CVSS,
- 3. Calculation of the corrected CVSS values,
- 4. Calculation of the average correction function, and
- 5. Aggregation.

Notice that the dependency graph of the SUT, the vulnerabilities associated to each element of the dependency graph, and their CVSS value are needed.

#### Correction factors for each CVSS

The first step obtains the values of each correction factor for each CVSS:

- 1. **Functionality factor** ( $\rho$ ): This factor is obtained using the description provided in the corresponding CVE of each CVSS. The description provides enough information to decide whether the functionality of the system is affected.
- 2. **Context factor** ( $\gamma$ ): This factor is obtained by comparing the value of the Attack Vector (AV) submetric of the CVSS, with the real environment of deployment of the SUT.
- 3. **Exploit factor** ( $\mu$ ): To obtain this factor, public databases have to be queried to find any potential exploit for each vulnerability.
- 4. **Deepness factor** ( $\beta$ ): For any given CVSS, its value is obtained according to the deepness in the exploit chain for the SUT [199].

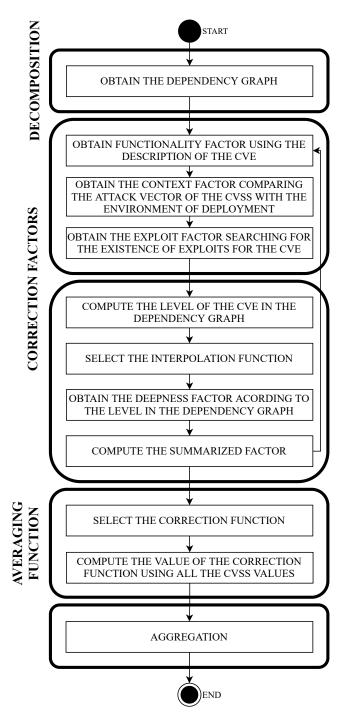


FIGURE 6.2: Flowchart showing the main steps of the aggregation algorithm for each CVSS.

#### Summarized factor for each CVSS

The summarized factor,  $\lambda$ , is obtained by multiplying all the corrections factors obtained in the previous step, following Equation 6.1.

#### **Corrected CVSS values**

The corrected CVSS values are obtained by multiplying each CVSS by its corresponding summarized factor,  $\lambda$ . At this point, it is necessary to check for overflows, because the exploitation factor generated corrected CVSS values higher than 10. Values higher than 10 are set to 10 at this stage.

#### **Correction function**

At this point, it is necessary to choose an averaging function. Choosing one function over the other will cause the aggregation result to grow slower or faster toward 10 in each addition. In this case, and for the sake of clarity, we chose the arithmetic mean, but any other kind of mean (*e.g.*, *harmonic mean*) could be used according to each scenario.

## Aggregation

Finally, the aggregated value is computed using Equation 6.2.

## 6.1.4 Interpretation of the result

The advantage of this method is that the result can be interpreted in the same way that a normal CVSS would be interpreted. This is because of the correction factors in Equation 6.2, that only let the algorithm return high values when vulnerabilities with high CVSS values are exploitable in reality ( $\lambda$  is close to 1). This mechanism ensures that multiple aggregated low CVSS values do not result in a critical score just because there are a large number of them.

6.2. Use Case 105

# 6.2 Use Case

To test the potential of our proposal, we analyzed Version 3 of OpenPLC project, obtaining a CVSS aggregated value for its vulnerabilities using the proposed algorithm.

OpenPLC is the first functional open source Programmable Logic Controller (PLC), both in software and hardware [187]. It was mainly created for research purposes, because it provides its entire source code [188,189]. The current version of the project is OpenPLC V3 [192].

#### 6.2.1 Use Case Scenario

For this use case, we are going to make the next assumptions:

- The system executing OpenPLC V3 is deployed in an isolated network.
- The system running OpenPLC V3 is physically isolated.
- The attacker is an insider without access to the systems.
- The reference point for the deepness factor will be the webserver.py in Figure 6.3.

# 6.2.2 Structure of OpenPLC

The first step was to obtain the inner structure of OpenPLC V3 using the Extended Dependency Graph (EDG) proposed in [199]. To simplify the obtained graph, we only represented the shortest path to each node, so the worst case scenario (more accessible from the outside) is considered. The result is shown in Figure 6.3.

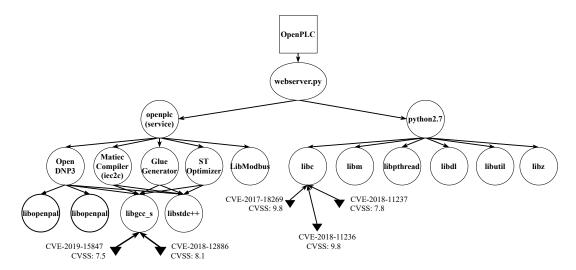


FIGURE 6.3: Extended Dependency Graph of OpenPLC V3. Circles represent individual assets, black triangles are the vulnerabilities associated to each asset, and the square represent the entry point to the system, or root node of dependency.

TABLE 6.3: Vulnerabilities present in OpenPLC V3. For each one, the CVSS is shown, together with their associated Attack Vector (AV), and their correction factors.

CVE	CVSS	Attack Vector	Functionality (ρ)	Deepness (β)	Context $(\gamma)$	Exploit (µ)	Summarized ( $\lambda$ )	Corrected CVSS
CVE-2017-18269	9.8	Network	1	0.5	1	1.25	0.625	6.125
CVE-2018-11236	9.8	Network	0	0.5	1	0	0	0
CVE-2018-11237	7.8	Local	1	0.5	0	1.25	0	0
CVE-2018-12886	8.1	Network	1	0.25	1	1.25	0.313	2.530
CVE-2019-15847	7.5	Network	1	0.25	1	1.25	0.313	2.344

# 6.2.3 Calculation of the Correcting Factors

OpenPLC V3 has five vulnerabilities: two vulnerabilities affecting libgcc\_s, and three vulnerabilities affecting libc. Table 6.3 shows each vulnerability in more detail.

From these data, it is possible to obtain all corrections factors for each vulnerability, as follows (Table 6.3 summarizes the results):

# Functionality Factor $(\rho)$

This factor is obtained from the analysis of the description of each CVE. From these data, we have to decide whether the functionality of OpenPLC V3 is affected ("1") or not ("0").

#### **Deepness Factor** ( $\beta$ )

By taking a look at Figure 6.3, it can be seen that the maximum deepness level is four. So the possible values for the deepness factor are the ones shown in Figure 6.1. More precisely, vulnerabilities CVE-2019-15847 and CVE-2018-12886 have a deepness factor of 0.25, because they are at level four. By contrast, vulnerabilities CVE-2017-18269, CVE-2018-11236, and CVE-2018-11237 have a deepness factor of 0.5, because they are at level three.

#### Context Factor $(\gamma)$

From the initial assumptions, insiders can only exploit the existing vulnerabilities from the local network. This means that every vulnerability that has an attack vector of "network" (N) can be exploited, thus CVE-2017-18269, CVE-2018-11236, CVE-2018-12886, and CVE-2019-15847 are exploitable by the attacker. Vulnerabilities whose attack vector is "local" (L) cannot be exploited, because physical access is needed. Therefore, CVE-2018-11237 cannot be exploited.

# **Exploit Factor** $(\mu)$

Public databases have to be queried to find existing exploits for each vulnerability. According to their state of development, a different value is assigned.

#### **Summarized Factor** $(\lambda)$

The summarized factor for each vulnerability is obtained as the product of the previous factors, as shown in Equation 6.1. At this step, by taking a look at the resulting values of  $\lambda$ , it is possible to know which CVSS will contribute to the final aggregation and in which percentage ( $\lambda > 0$ ), and which ones will not contribute at all ( $\lambda = 0$ ).

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# Average Factor $(\sigma)$

Finally, we obtained the average factor by calculating the arithmetic mean of all the initial CVSS values:  $\sigma = 8.6$ .

# 6.2.4 Aggregation

The previous step before the aggregation is obtaining the corrected CVSS value for each initial CVSS. This is done by multiplying each CVSS by their corresponding summarized value ( $\lambda$ ). The corrected values are shown in Table 6.3.

Finally, the aggregation is performed using the corrected CVSS values. The aggregation is an iterative process that takes the first two values to be added, and adds them using Equation 6.2. Then, this result is added to the third value to be added, and so on, until there are no more values.

For OpenPLC V3, this process returns a final aggregated value of 9.1. Without the correction factors, the result would be 10. Nevertheless, taking into account features such as the exploitability of the vulnerabilities, the context of the SUT, or its functionalities, we can select the most important CVSS values to be aggregated. With such process, the total amount of CVSS values to be added is simplified. This also helps to simplify potential attack paths.

This result was obtained aggregating three of the five CVSS values present in Open-PLC V3. The associated CVSS for CVE-2018-11236 and CVE-2018-11237 were not taking into account for the aggregation, because they do not affect to any functionality of the system, Moreover, CVE-2018-11237 cannot be exploited in the conditions described in the use case.

CVE-2017-18269 (with an associated CVSS of 9.8) is the vulnerability with the highest value for  $\lambda$ . Therefore, it is going to contribute the most to the final aggregated value. CVE-2018-12886 and CVE-2019-15847 follow with a CVSS of 8.1 and 7.5 respectively. As it is shown, the selected vulnerabilities have a high CVSS, so it is expected that the aggregated value would be also high. This is reflected in the obtained result of 9.1.

Finally, it is worth highlighting that the final result is lower than the highest CVSS value present in OpenPLC V3. This difference is due to the effect of the correction factors: as the CVE-2017-18269 is further away from the entry point of the system (in layer 3), its real CVSS value in lower.

# 6.3 Conclusions and Future Work

In this research work, we proposed a new aggregation algorithm for CVSS values. The proposed approach integrates correction factors to select the most relevant CVSS values to be added based on contextual information. For each vulnerability, we check for:

- 1. Functionality disruption.
- 2. Exploitation difficulty.
- 3. Existence of exploits, and their development state.
- 4. Context of deployment.

We assigned a different correction factor to each one of the previous properties to further ponder the initial CVSS value and adjust it to the real context where the system is operating.

The proposed aggregation algorithm was applied to OpenPLC V3 in a use case. Two of the existing vulnerabilities were filtered out by the algorithm, as they cannot be exploited in the described context of OpenPLC V3. The rest of the vulnerabilities were aggregated, and the result (9.1) was indeed lower than the highest CVSS present in the system (9.8). This shows that the CVSS for each vulnerability was correctly adjusted to the real context of deployment of OpenPLC V3.

As future work, we plan to perform the aggregation at the submetric level of the CVSS, instead of using the base metric value, giving more granular values for each factor.

# Chapter 7

# **Conclusions and Future Work**

ESs constitute a critical piece for the industry, where they are present in almost every field, including critical infrastructures. Moreover, they are also rapidly evolving, and increasing in number. Under this scenario, a successful attack could have disastrous consequences, both economic and human losses. These characteristics together highlight the necessity of a methodology to assess the level of security of ESs.

This thesis focuses on quantifying the security level of ESs. For this reason, we reviewed both industrial standards and scientific literature to collect existing metrics. The properties of a comprehensive metrics were also found during this analysis. After analyzing the data, we found that most of the collected metrics are design to assess the level of security of an organization, and cannot be applied to ESs. Those that can be applied to ES are focused on software-related aspect of the system, while hardware-related ones are often neglected. Because existing taxonomies are defined to be used at a very high level (organization level), we proposed a new ES-centered taxonomy to classify security metrics. This taxonomy focuses on the specific assessment needs and constrains of ESs.

On the other hand, we found that graph-based techniques that are typically used in computer networks analysis can be applied to ES. Representing both the inner dependencies of an ES and its vulnerabilities helps to find root causes, make a decision about patching prioritization, and to generate new requirements and test cases. With this idea in mind, we proposed the Extended Dependency Graphs (EDGs), a directed graph-based structure that is capable of representing the inner structure of an ES. Moreover, we integrate security standards such as the CVSS, and the CWE to quantify, and find potential root cause for each vulnerability. In addition, we were also able to generate new requirements, new test cases, and propose training.

Finally, to complement the previously proposed model, we developed an aggregation algorithm for CVSS values. The proposed approach integrates correction factors to select the most relevant CVSS values to be added based on contextual information. For each vulnerability, we check for functionality disruption, exploitation difficulty, existence of exploits, and context of deployment.

As future work, analyzing the obtained metrics in more deepness could be promising. Choosing a smaller set of metrics, and test them empirically to check their real applicability to ESs. This would allow a more detailed knowledge of existing metrics, their applicability, and the relation between the property they are measuring and the

returned value. Moreover, it would be interesting to know which metrics are indeed more used in the industry.

With regard to the EDGs, we plan to generate a tool that automates the generation of EDGs. This tool could also integrate automatic updating of the current state of the SUT. So with every change (discovered vulnerability, update release) the corresponding EDG would be updated. A system to include warnings, available exploits, or patching prioritization suggestions would be also interesting. Historical information about the developers can also be integrated into the EDG model to predict future vulnerabilities, and tendencies among the developers. Additionally, according to the properties of a comprehensive methodology described in Section 3.2.13 (Table 3.5), we should enhance the guidance and reporting of the EDGs.

Finally, we plan to explore the advantages of performing the aggregation at the submetric level of the CVSS, instead of using the base metric value, giving more granular values for each factor. Integrating this aggregation algorithm to the EDG tool, could lead to better analysis of each SUT.

# **Chapter 8**

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## Appendix A

## **Collected Metrics**

# METRIC	DEFINITION	SCALE	SCOPE AU	AUTOMATION	MEASUREMENT	USABLE	REASON
1 False positives (FP)	Percentage of flagging genuine email as phishing email	RATIO	USER		DYNAMIC	ON	Phishing affects only people. It exploits a user vulnerability. This evaluates the user, not the device
2 False negatives (FN)	Percentage of detecting a phishing email as a genuine email	RATIO	USER		DYNAMIC	NO	Phishing affects only people. It exploits a user vulnerability. This evaluates the user, not the device
3 User's online behavior		КАПО	USER		DYNAMIC	ON	Phishing affects only people. It exploits a user vulnerability. This evaluates the user, not the device. This is only a suggestion. Not defined
4 Users' susceptibility to attacks		RATIO	USER		UNKNOWN	NO	This evaluates the user, not the device. This is only a suggestion. Not defined
5 Entropy			USER AUT	AUTOMATIC	STATIC	YES	This could be used with keys instead of passwords to measure randomness
6 Password guessability (statistical / parameterized)	This metric measures the guessability of a set of passwords (rather than for individual passwords) by an "idealized" attacker with the assumption of a perfect guess / This metric measures the number of guesses an attacker needs to make via a particular cracking algorithm (i.e., a particular threat model) before recovering a particular password and training data. Different password cracking algorithms leads to different results	RATIO ABSOLUTE	USER AUT	АUТОМАПС	STATIC	NO	Passwords are for people. Devices use keys (that are much more larger)
7 Password strength meter	This metric measures the password strength via an estimated effort to guess the password, while considering factors such as the character set (e.g., special symbols are required or not), password length, whitespace permissibility, password entropy, and blacklisted passwords being prevented or not. One variant metric is adaptive password strength, which is used to improve the accuracy of strength estimation	ORDINAL	USER AUT	аυтоматіс	STATIC	ON	Passwords are for people. Devices use keys (that are much more larger)
8 Attack surface	It aim to measure the ways by which an attacker can compromise a targeted software $% \left( 1\right) =\left\{ 1\right\} $	ABSOLUTE	DEVICE		STATIC	YES	
9 Historical vulnerability metric	It measures the degree that a system is vulnerable (i.e., frequency of vulnerabili- RATIO ties) in the past DISTRI	- RATIO ABSOLUTE DISTRIBUTION	DEVICE		STATIC	YES	This could be part of a checklist of task, to check if there is any vulnerability that affects the device
10 Historically exploited vulnerability metric	It measures the number of vulnerabilities exploited in the past	RATIO ABSOLUTE DISTRIBUTION	DEVICE		STATIC	YES	This could be part of a checklist of task, to check if there is any vulnerability that affects the device
11 Future vulnerability metric	It measures the number of vulnerabilities that will be discovered during a future period of time	RATIO ABSOLUTE DISTRIBUTION	DEVICE		STATIC	ON	If this has to do with statistic and probabilities, I think it is better not to use it
12 Future exploited vulnerability metric	It measures the number of vulnerabilities that will be exploited during a future period of time	RATIO ABSOLUTE DISTRIBUTION	DEVICE		STATIC	ON	If this has to do with statistic and probabilities, I think it is better not to use it
13 Tendency-to-be-exploited metric	It measures the tendency that a vulnerability may be exploited, where the "tendency" may be derived from information sources such as posts on Twitter before vulnerability disclosures	RATIO ABSOLUTE DISTRIBUTION	DEVICE		STATIC	??	This research work might not be part of a security evaluation, but could be applied if so
14 Vulnerability lifetime	Client-end vulnerabilities Server-end vulnerabilities Cloud-end vulnerabilities	ABSOLUTE	NETWORK		STATIC	NO	Clearly related to nerworks
15 Attack vector	Describes whether a vulnerability can be exploited remotely, adjacently, locally, or physically (i.e., attacker must have physical access to the computer)	NOMINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
16 Attack complexity	Describes the conditions that must hold before an attacker can exploit the vulnerability, such as low or high	- ORDINA	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
17 Privilege required	Describes the level of privileges that an attacker must have in order to exploit a vulnerability, such as none, low, or high	ORDINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
18 User interaction	Describes whether the exploitation of a vulnerability requires the participation of a user (other than the attacker), such as none or required	NOMINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
19 Authorization scope	Describes whether or not a vulnerability has an impact on resources beyond its privileges (e.g., sandbox or virtual machine), such as unchanged or changed	NOMINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
20 Confidentiality impact	The impact of a successful exploitation of a vulnerability in terms of confidentiality such as none, low, high	ORDINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
21 Integrity impact	The impact of a successful exploitation of a vulnerability in terms of integrity such as none, low, high	ORDINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
22 Availability impact	The impact of a successful exploitation of a vulnerability in terms of availability, such as none, low, high	ORDINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
23 Exploit code maturity	The likelihood a vulnerability can be attacked based on the current exploitation techniques, such as undefined, unproven, proof-of-concept, functional, or high	ORDINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
24 Remediation level	Describes whether or not a remediation method is available for a given vulnerability, such as undefined, unavailable, workaround, temporary fix, or official fix	- NOMINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
25 Report confidence	Measures the level of confidence for a given vulnerability as well as known technical details, such as unknown, reasonable, or confirmed	NOMINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
26 Security requirement	Describes environment-dependent security requirements in terms of confidentiality, integrity, and availability, such as not defined, low, medium, or high	NOMINAL	SOFTWARE		STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
27 Depth metric	Ratio of the diameter of a domain-level attack graph over the diameter in the most secure case, implying that the larger the diameter, the more secure the network		NETWORK		STATIC	ON	It is related to the topological properties of attack graphs
28 Existence, number, and lengths of attack paths metrics	It uses the attributes of attack paths from an initial state to the goal state. These metrics can be used to compare two attacks	1	NETWORK		STATIC	ON	It is related to the topological properties of attack graphs

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29 Necessary defense metric	it estimates a minimal set of defense countermeasures necessary for thwarting a certain attack	DEVICE	SIAIIC	YES II	This is related with the detenses that are necessary to implement so the system is secure for a certain vulnerability
30 Effort-to-security-failure metric	It measures an attacker's effort to reach its goal state	DEVICE	DYNAMIC*	YES A	Although it might be difficult to estimate, this could be also included in the evaluation process
31 Weakest adversary metric	It estimates minimum adversary capabilities required to achieve an attack goal	DEVICE	DYNAMIC*	H ??	If the effort can be calculated in an objective manner, this could be used in embedded systems
32 k-zero-day-safety metric	It measures a number of zero-day vulnerabilities for an attacker to compromise a target	DEVICE	STATIC	II ON	This has to do with probabilities. Can be applied, but I won't
33 Likelihood of exploitation	It measures the probability that an exploit will be executed by an attacker with certain capabilities	SOFTWARE	STATIC	П ??	This is related with statistics and risk assessment. It could be applied
34 Effective base metric / scope	Given an attack graph and a subset of exploits, the effective base metric aims to adjust the CVSS base metric by taking the highest value of the ancestors of an exploit to reflect the worst-case scenario, while the effective base score metric is calculated based on the effective base metric	SOFTWARE	STATIC		
35 Reaction time metric	It captures the delay between the observation of the malicious entity at time $t$ ABSOLUTE and the blacklisting of the malicious entity at time $t(ic,,t-t)$	DEVICE	DYNAMIC	YES If ar	If the device has any kind of detection mechanism, this could be used. Detection and reaction might have different times
36 Coverage metric	It estimates the portion of blacklisted malicious players ORDINAL	DEVICE	STATIC	YES T	This could be used, but the device has to have a blacklist
37 NO METRIC IS DEFINED	No metrics have been defined to measure the effectiveness of DEP. The effectiveness of DEP can be measured based on the probability of being compromised by a certain attack A(t) over all possible classes of attacks	DEVICE		N -	No metric is defined
38 Average indirect target reduction	It counts the overall reduction of targets for any indirect control-flow transfer. It ABSOLUTE measures the overall reduction in terms of the number of targets exploitable by the attacker where smaller targets are more secure	UNKNOWN	UNKNOWN	UNKNOWN U	UNKNOWN
39 Average target size	Ratio between the size of the largest target and the number of targets	UNKNOWN	UNKNOWN	UNKNOWN UNKNOWN	NKNOWN
40 Evasion resistance	It is measured against control flow bending attacks, reflecting the effort (or premises) that an attacker must make (or satisfy) for evading the CFI scheme	UNKNOWN	UNKNOWN	UNKNOWN UNKNOWN	NKNOWN
41 Coverage metric	It measures the fraction of events detectable by a specific sensor deployment, reflecting a defender's need in monitoring events	DEVICE	DYNAMIC	UNKNOWN U	UNKNOWN
42 Redundancy metric	It estimates the amount of evidence provided by a specific sensor deployment to detect an event $\  $	DEVICE	DYNAMIC	YES TI	This can be applied to devices as a whole, to check any mechanism of intrusion detection (mostly, tanti-tampering and software)
43 Confidence metric	It measures how well-deployed sensors detect an event in the presence of compromised sensors	DEVICE	DYNAMIC	YES TI	This can be applied to devices as a whole, to check any mechanism of intrusion detection (mostly, tanti-tampering and software)
44 Cost metric	It measures the amount of resources consumed by deploying sensors including - the cost for operating and maintaining sensors	ORGANISATION	STATIC	NO Ti	This is related with the cost of a specific countermeasure. This is more likely to be applied in early steps of the development
45 Detection time	For instrument-based attack detection, this metric is used to measure the delay between the time t0 at which a compromised computer sends its first scan packet and the time t that a scan packet is observed by the instrument	UNKNOWN	UNKNOWN	UNKNOWN U	UNKNOWN
46 True-positive rate	It is the probability that an intrusion is detected as an attack	NETWORK	DYNAMIC	II ON	This is clearly a network-oriented metric
47 False-negative rate	It is the probability that an intrusion is not detected as an attack	NETWORK	DYNAMIC	NO TI	This is clearly a network-oriented metric
48 True-negative rate	It is the probability that a non-intrusion is not detected as an attack	NETWORK	DYNAMIC	II ON	This is clearly a network-oriented metric
49 False-positive rate	Also known as false alarm rate, it is the probability that a nonintrusion is detected as an attack	NETWORK	DYNAMIC	II ON	This is clearly a network-oriented metric
50 Intrusion detection capability metric	It is the normalized metric of $I(I,O)$ with respect to $H(I)$ based on the base rate where I is the input to the IDS as a stream of $0/1$ random variables (0 for benign/normal and I for malicious/abnormal), O is the output of the IDS as a stream of $0/1$ random variables (0 for no alert or normal; 1 for alert / abnormal), $H(I)$ and $H(O)$ , respectively, denote the entropy of I and O, and $I(I,O)$ is the mutual information between I and O	NETWORK	UNKNOWN	ON T	This is clearly a network-oriented metric
51 Receiver operating characteristic (ROC)	It reflects the dependence of the true-positive rate on the false-positive rate, reflecting a tradeoff between the true-positive and the false-positive rates	NETWORK	STATIC	II ON	This is clearly a network-oriented metric
52 Intrusion detection operating characteristic (IDOC)	It describes the dependence of the true positive rate $\Pr(A \mid I)$ on the Bayesian detection rate $\Pr(I \mid A)$ , while accommodating the base rate $\Pr(I)$	NETWORK	STATIC	II ON	This is clearly a network-oriented metric
53 Cost metric (damage / response / operational)	It includes the damage cost incurred by undetected attacks, the response cost spent on the reaction to detected attacks including both true and false alarms, and the operational cost for running an IDS	NETWORK	STATIC	TI ON	This is clearly a network-oriented metric, that aims to translate the impact into moneu
54 Relative effectiviness	This metric reflects the strength of a defense tool when employed in addition to RATIO other defense tools. A defense tool does not offer any extra strength if it cannot detect any attack undetected by other defense tools in place	NETWORK	DYNAMIC	TI ON	This is clearly a network-oriented metric
55 Collective effectiviness	This metric measures the collective strength of IDSs and anti-malware programs RATIO	NETWORK	DYNAMIC	II ON	This is clearly a network-oriented metric
				Ţ	This is clearly a network-oriented metric
57 ASLR-induced Effective entropy metric	Measure of entropy in user space memory which quantitatively considers an adversary's ability to leverage low entropy regions of memory via absolute and dynamic inter-section connections	SOFTWARE	DYNAMIC	YES If	If the device runs an OS, this could be applied
58 Moving Target Defense (MTD)	It measures the degree that an enterprise system can tolerate some undesirable security configurations in order to direct the global security state S(t) towards a desired stable system state	NETWORK	DYNAMIC	II ON	This is only applicable to an enterprise system
59 Penetration resistance (PR)	Running a penetration test to estimate the level of effort (e.g., person-day or cost) required for a red team to penetrate into a system	DEVICE NETWORK	DYNAMIC	YES Pe	Penetration testing is a really common way to test this kind of device

60 Network di	Network diversity (ND)	It measures the least or average effort an attacker must make to compromise - a target entity based on the causal relationships between resource types to be	NETWORK	DYNAMIC	ON	This is clearly a network-oriented metric
61 Lifetime of	Lifetime of zero-day attacks	It measures the period of time between when an attack was launched and when - the corresponding vulnerability is disclosed to the public	SOFTWARE HARDWARE	DYNAMIC	ON	This can be computed, but it is not feasible for a security evaluation
62 Victims by	/ zero-day attacks	It measures the number of computers compromised by zero-day attacks	ORGANISATION	DYNAMIC*	NO	This is usefull for networks or an organisation. It can be computed for a network of devices, but it is not useful for a security evaluation
	Susceptibility of a device to zero-day attacks	If we can capture the degree of the susceptibility of a device to zero-day attacks, then we can prioritize resource allocation to the ones requiring higher attention for mitigating the damage	DEVICE			This metric is proposed in [2] as a future concept, but it is not developed
1	Targeted threat index	Level of targeted malware attacks	NETWORK ORGANISATION	DYNAMIC	ON	This is usefull for networks or an organisation. It can be computed for a network of devices, but it is not useful for a security evaluation
65 Botnet size	9	Number of bots, x, that can be instructed to launch attacks (e.g., distributed denial-of-service attacks) at time t, denoted by y(t). Due to time zone difference, y(t) is often much smaller than the actual x as some of x is turned off during night time at time zones	NETWORK	STATIC	ON	This is clearly a network-oriented metric
66 Network bandwidth	oandwidth (1900)	Network bandwidth that a botnet can use to launch denial-of-service attacks	NETWORK	STATIC DYNAMIC	ON	This is clearly a network-oriented metric
67 Botnet efficiency	ciency	Network diameter of the botnet network topology. It measures a botnet's capability in communicating command-and-control messages and updating bot programs	NETWORK	?3	ON	This is clearly a network-oriented metric
68 Botnet robustness	ustness	Robustness of botnets under random or intelligent disruptions. The idea was derived from the concept of complex network robustness, characterized by the percolation threshold under which a network is disrupted into small components	NETWORK	<i>ξ?</i>	ON	This is clearly a network-oriented metric
69 Infection rate	ate	Average number of vulnerable computers that are infected by a compromised computer (per time unit) at the early stage of spreading	NETWORK	DYNAMIC	ON	This is usefull for networks or an organisation. It can be computed for a network of devices, but it is not useful for a security evaluation
70 Increased f	Increased false-positive rate	The strength of attacks as a consequence of applying a certain evasion method	DEVICE NETWORK	DYNAMIC	ON	This can only be calculated if the embedded system has implemented an evasion mechanism against machine learning attacks
71 Increased t	Increased false-negative rate	The strength of attacks as a consequence of applying a certain evasion method	DEVICE NETWORK	DYNAMIC	ON	This can only be calculated if the embedded system has implemented an evasion mechanism against machine learning attacks
72 Obfuscatio	Obfuscation prevalence metric	Occurrence of obfuscation in malware samples	MALWARE	STATIC	NO	Malware-related
73 Structural	Structural complexity metric	Runtime complexity of packers in terms of their layers or granularity	MALWARE	;?	ON	Malware-related
74 Evasion ca	Evasion capability of attacks	Allows comparing the evasion power of two attacks, but also allows computing $$ - the damage caused by evasion attacks $$	MALWARE	DYNAMIC	NO	Proposed, but not defined
75 Obfuscatio	Obfuscation sophistication	In terms of the amount of effort required for unpacking packed malware	MALWARE	STATIC	ON	Proposed, but not defined. Malware related
76 Network n	Network maliciousness metric	It estimates the fraction of blacklisted IP addresses in a network	NETWORK	STATIC DYNAMIC	NO	This metric works at network level, for systems in a network. Cannot be applied to embedded systems
77 Rogue netx	Rogue networkmetric	Population of networks used to launch drive-by download or phishing attacks	NETWORK	DYNAMIC	ON	This metric works at network level, for systems in a network. Phishing cannot be applied to embedded systems
78 ISP badness metric	ss metric	Quantifies the effect of spam from one ISP or Autonomous System (AS) on the rest of the Internet	NETWORK	DYNAMIC	ON	This is clearly a network-oriented metric
79 Control-pla	Control-plane reputation metric	Calibrates the maliciousness of attacker-owned (i.e., rather than legitimate but mismanaged/abused) ASs based on their control plane information (e.g., routing behavior), which can achieve an early-detection time of 50–60 days (before these malicious ASs are noticed by other defense means)	ATTACKER	?5	ON	This is related to external systems that could be potential attackers
80 Cybersecui	Cybersecurity posturemetric	Measures the dynamic threat imposed by the attacking computers	ATTACKER	DYNAMIC	NO	This is related to attacks themselves
81 Fraction of	Fraction of compromised computers at time t		NETWORK	DYNAMIC	NO	Directly related to the systems in a network
82 Probability	Probability a computer is compromised at time t	It quantifies the degree of security in enterprise systems by using modeling - techniques based on a holistic perspective	DEVICE NETWORK	53	53	Although this is a network-oriented metric, this might be adapted to be used with an embedded system
83 Encounter rate	rate	Measures the fraction of computers that encountered some malware or attack $$ - during a period of time	NETWORK	STATIC	ON	Directly related to the systems in a network
84 Incident rate	ate	Measures the fraction of computers successfully infected or attacked at least once during a period of time	NETWORK	STATIC	NO	Directly related to the systems in a network
85 Blocking rate	ate	The rate an encountered attack is successfully defended by a deployed defense	NETWORK	DYNAMIC	<i>??</i>	Although this is a network-oriented metric, this might be adapted to be used with an embedded system (maybe for a penetration test)
86 Breach freq	Breach frequencymetric	How often breaches occur	NETWORK	STATIC	ON	Directly related to the systems in a network
87 Breach size metric	e metric	Number of records breached in individual breaches	NETWORK	STATIC	ON	Directly related to the systems in a network
88 Time-betwo	Time-between-incidents metric	Period of time between two incidents	NETWORK	STATIC	ON	Directly related to the systems in a network
89 Time-to-fir	Time-to-first-compromise metric	Estimates the duration of time between when a computer starts to run and the first malware alarm is triggered on the computer where the alarm indicates detection rather than infection	DEVICE	DYNAMIC	YES	This can be used in embedded systems, but it might not be suitable for a security evaluation (maybe for a penetration test)
90 Delay in in	Delay in incident detection	Time between the occurrence and detection	ORGANISATION	DYNAMIC	ON	This metric could be used to transmit the importance of security in the organisation, but cannot be used in embedded systems
91 Cost of incidents	zidents		ORGANISATION	STATIC	ON	Similar to 56
92 Security spending	pending	Percentage of IT budget	ORGANISATION	STATIC	ON	This metric shows the investing in security in the organization. It Cannot be applied to embedded systems

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55 Security budget allocation	it estimates now the security budget is allocated to various security activities and - resources	OKGANISALION		SIAIIC	NO	Ins metric shows the investing in security in the organization. It cannot be applied to embedded systems	
94 Return on security investment (ROSI)	It is a variant of the classic return on investment (ROI) metric, measuring the financial net gain of an investment based on the gain from investment minus the cost of investment	ORGANISATION		STATIC	NO	This metric shows the investing in security in the organization. It Cannot be applied to embedded systems	
95 Net present value	It measures the difference between the present economic value of future inflows and the present economic value of outflows with respect to an investment	ORGANISATION		STATIC	ON	This metric shows the investing in security in the organization. It Cannot be applied to embedded systems	
96 Attack Impact	It is the quantitative measure of the potential harm caused by an attacker to exploit a vulnerability	NETWORK		STATIC			
97 Attack Cost	It is the cost spent by an attacker to successfully exploit a vulnerability (i.e., security weakness) on a host	NETWORK		STATIC DYNAMIC	YES	This is related to network, but this could be applicable to devices too	
98 Structural Important Measured	It it used to qualitatively determine the most critical event (attack, detection or mitigation) in a graphical attack model	NETWORK		STATIC	ON	Related to networks	
99 Mincut Analysis		NETWORK					
100 Mean-time-to-compromise (MTTC)	It is used to measure how quickly a network can be penetrated	NETWORK		DYNAMIC	۶۶	This is no directly applicable, but its equivalent would be a penetration attack in an embedded system	
101 Mean-time-to-recovery (MTTR)	It is used to assess the effectiveness of a network to recovery from an attack - incidents. It is defined as the average amount of time required to restore a system out of attack state	NETWORK		DYNAMIC	?خ	Although this is developed for networks, the same concept can be translated into embeded systems	
102 Mean-time-to-First-Failure (MTFF)	23	NETWORK					
103 Mean-Time-To-Breach (MTTB)		NETWORK					
104 Return on Investment	It measures the financial net gain of an investment based on the gain from investment minus the $\cos t$ of investment	NETWORK		STATIC	ON	Network-related and cost related. Useful for organisation managing	
105 Return on Attack	The gain the attacker expects from successful attack over the losses he sustains due to the countermeasure deployed by his target	NETWORK		;?	?خ	Network-related, but it scent can be translated into embedded systems	
106 Common Vulnerability Scoring System (CVSS) - Overall value	I The overall score is determined by generating a base score and modifying it through the temporal and environmental formulas	DEVICE		STATIC	YES	This could be used tohether with other metics to obtain a numeric value	
107 Probability of Vulnerability Exploited	It is used to assess the likelihood of an attacker exploiting a specific vulnerability on a host	SOFTWARE HARDWARE NETWORK		STATIC	¿?	Network-related metric. This is also related to risk assessment. Can be traslated into embedded systems	
108 Probability of attack detection	It is used to assess the likelihood of a countermeasure to successfully identify the event of an attack on a target.	SOFTWARE HARDWARE NETWORK		DYNAMIC	¿?	Network-related metric. This is also related to risk assessment. Can be traslated into embedded systems	
109 Probability of host compromised	It is used to assess the likelihood of an attacker to successfully compromise $a$ -target	HARDWARE SOFTWARE NETWORK		STATIC DYNAMIC	¿?	Network-related metric. This is also related to risk assessment. Can be traslated into embedded systems	
110 Network Compromise Percentage	It quantifies the percentage of hosts on the network on which an attacker canobtain an user or administration level privilege	NETWORK		STATIC DYNAMIC	ON	It's about networks, and we can't predict where the device will be used, so this metric is not representative	
111 Waekest Adversary	It is used to assess the security strength of a network in terms of the weakest part of the network that an attacker can successfully penetrate	SOFTWARE HARDWARE NETWORK		STATIC DYNAMIC	ON	Similar to 31. It's about networks, and we can't predict where the device will be used, so this metric is not representative	
112 Vulnerable Host Percentage	This metric quantifies the percentage of hosts with vulnerability on a network. It is used to assess the overall security of a network	NETWORK		DYNAMIC	ON	It's about networks, and we can't predict where the device will be used, so this metric is not representative	
113 Attack Shortest Path	It is the smallest distance from the attacker to the target. This metric represents the minimum number of hosts an attacker will use to compromise the target host	NETWORK	AUTOMATIC*	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
114 Number of Attack Paths	It is the total number of ways an attacker can compromise the target	NETWORK	AUTOMATIC*	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
115 Mean of Attack Path Lengths	It is the average of all path lengths. It gives the expected effort that an attacker may use to breach a network policy	NETWORK	АUТОМАПС⁴	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
116 Normalised Mean of Path Lengths	Expected number of exploits an attacker should execute in order to reach the target	NETWORK	AUTOMATIC*	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
117 Standand Deviation of Path Lengths	It is used to determine the attack paths of interest	NETWORK	АUТОМАПС⁴	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
118 Mode of Path Lengths	It is the attack path length that occurs most frequently	NETWORK	AUTOMATIC*	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
119 Median of Path Lengths	It is used by network administrator to determine how close is an attack path length to the value of the median path length (i.e. path length that is at the middle of all the path length values)	NETWORK	AUTOMATIC*	DYNAMIC	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
120 Attack Resistance Metric	It is use to assess the resistance of a network configuration based on the composition of measures of individual exploits. It is also use for assessing and comparing the security of different network configurations	NETWORK		53	ON	If it has to do with paths, it has to do with networks, and we can't assume where is this going to be used	
121 Mean Time To Failure (MTTF) - Time to compromise	Mean time for a potential intruder to reach the specified target	NETWORK	AUTOMATIC	DYNAMIC	NO		
122 Mean Effort To Failure (METF) - Effort to compromise	Mean effort for a potential attacker to reach the specified target	NETWORK	SEMI-AUTOMATIC	C DYNAMIC	ON		
123 Mean Time to Security Failure (MTTSF) - Time to compromise	<ul> <li>Expected amount of time attackers need to successfully complete an attack - process</li> </ul>	NETWORK	SEMI-AUTOMATIC	C STATIC	ON		
124 Mean Time to First Failure (MTFF) - Time to compromise	e Mean time interval from system initiation to the first system failure	NETWORK	SEMI-AUTOMATIC	C STATIC	?خ	Similar to 102. Depending on the type of test, the time to compromise can be measured (penetration test, fuzzing) but the result could not be representative	

125 Mean Time To Compromise (MTTC) - Time to con	Mean Time To Compromise (MTTC) - Time to compro-Mean time needed for an attacker to gain some level of privilege on some system	NETWORK	SEMI-AUTOMATIC	STATIC	ON	Similar to 100. How can the effort be calculated in an objective manner? This
126 Shortest path	component	NETWORK	AUTOMATIC	DYNAMIC	ON	Similar to 113
127 Number of paths		NETWORK	AUTOMATIC	DYNAMIC	NO	Similar to 114
128 Mean of Path lengths		NETWORK	AUTOMATIC	DYNAMIC	ON	Similar to 115
129 Weakest adversary		NETWORK	SEMI-AUTOMATIC		ON	Similar to 31
130 Network compromise percentage		NETWORK	AUTOMATIC	DYNAMIC	ON	If we are evaluating a device, it has no sense talking about networks
131 Compromise probability		NETWORK	SEMI-AUTOMATIC	DYNAMIC	NO	
132 Attack Resistance	It is use to assess the resistance of a network configuration based on the composition of measures of individual exploits. It is also use for assessing and comparing the security of different network configurations	NETWORK	SEMI-AUTOMATIC	STATIC	ON	Similar to 120
133 Expected difficulty		NETWORK	AUTOMATIC	STATIC	ON	
134 Percentage of distinct hosts (d1-Diversity)		NETWORK	SEMI-AUTOMATIC	STATIC	ON	This is related with servers in a network, cannot be applied to a single device
135 Number of bits leaked (mean privacy)		NETWORK	SEMI-AUTOMATIC	STATIC	;?	If adapted, this metric can be used to test memory related vulnerabilities
136 K-zero day safety	Instead of attempting to rank unknown vulnerabilities, the metric counts how many such vulnerabilities would be required for compromising network assets; a larger count implies more security because the likelihood of having more unknown vulnerabilities available, applicable, and exploitable all at the same time will be significantly lower	NETWORK	SEMI-AUTOMATIC*	* STATIC*	ON	Similar to 32. This has to do with probabilities. Can be applied, but I won't
137 rav	It is a scale measurement of the attack surface, the amount of uncontrolled interactions with a target, which is calculated by the quantitative balance between operations, limitations, and controls					
138 30-Day Chum		SOFTWARE				
139 Churn	Count of Source Lines of Code that has been added or changed in a component since the previous revision of the software	SOFTWARE			1	
140 Frequency	Number of times that a binary was edited during its development cycle	SOFTWARE	AUTOMATIC*	STATIC*	ON	This is more related to the development of the software
141 Lines Changed	The cumulated number of code lines changed since the creation of a file	SOFTWARE	AUTOMATIC*	STATIC*	NO	This is more related to the development of the software
142 Lines new	The cumulated number of new code lines since the creation of a file	SOFTWARE	AUTOMATIC*	STATIC*	ON	This is more related to the development of the software
143 Number of lines deleted between revisions		SOFTWARE	AUTOMATIC*		ON	This is more related to the development of the software
144 Number of commits	Counting the commits a given developer made	SOFTWARE	AUTOMATIC*	STATIC*	ON	This is more related to the development of the software
145 Percentage of interactive churn (PIC)						
146 Relative churn	Normalized by parameters such as lines of code, files counts	SOFTWARE	AUTOMATIC*	STATIC*		
147 Repeat frequency	The number of consecutive edits that are performed on a binary	SOFIWARE	AUTOMATIC*	STATIC*		
148 Total chum	The total added, modified, and deleted lines of code of a binary during the - development	SOFTWARE	AUTOMATIC*	STATIC*		
149 Average service depth (ASD)						
150 Code complexity		SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software
151 Complexity		SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software
152 Count path complexity		SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software
153 Cyclomatic number	M = $E$ - $N$ + $P$ , where $E$ = the number of edges of the graph. $N$ = the number of nodes of the graph. $P$ = the number of connected components.	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
154 Dependency network complexity						
155 Execution complexity						
156 Fan in (FI)	Number of inputs a function uses. Inputs include parameters and global variables that are used (read) in the function / Number of functions calling a function $$	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
157 Fan out (FO)	Number of outputs that are set. The outputs can be parameters or global variables (modified) / Number of functions called by a function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
158 Henry Kafura: SLOC *( (FI*FO)2)						
159 Lack of cohesion of methods (LCOM)						
160 Max fan in	The largest number of inputs to a function such as parameters and global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
161 Max fan out	The largest number of assignment to the parameters to call a function or global - variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
162 MaxMaxNesting	The maximum of MaxNesting in a file	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
163 MaxNesting	Maximum nesting level of control constructs such as if or while statements in a function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
164 Nesting complexity	Maximum nesting level of control constructs (if, while, for, switch, etc.) in the function	SOFTWARE	АUТОМАПС⁴	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software

165 Nimbor of children (MOC)	Mumbar of immadista enhalacae of a place or the count of daring of accees 16 -	SOFTWARE	AITTOMATIC*	STATIC*	VFC	If the coturns is plaine at near the property of the contract
to intitude of chiracet (NOC)	Numer or numerators so that are of the control of control o	SOTIWAKE			1153	n the souware is available, it can be appired. It can be also appired to enfocued software
6 SumCyclomaticStrict	The sum of the strict cyclomatic complexity, where strict cyclomatic complexity is defined as the number of conditional statements in a function ${\bf n}$	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
77 SumEssential	The sum of essential complexity, where essential complexity is defined as the number of branches after reducing all the programming primitives such as a for loop in a function's control flow graph into a node iteratively until the graph cannot be reduced any further. Completely well-structured code has essential complexity 1	SOFTWARE	АUТОМАПС⁴	SТАТІС⁴	YES	If the software is available, it can be applied. It can be also applied to embedded software
8 SumFanIn	The sum of Fanin, where Fanin is defined as the number of inputs to a function such as parameters and global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
9 SumFanOut	The sum of FanOut, where FanOut is defined as the number of assignment to the parameters to call a function or global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
0 SumMaxNesting	The sum of the MaxNesting in a file, where MaxNesting is defined as the maximum nesting level of control constructs such as if or while statements in a function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
1 McCabe Cyclomatic Complexity	$M=E\cdot N+2P$ where $E=$ the number of edges of the graph. $N=$ the number of nodes of the graph. $P=$ the number of connected components	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
2 Weighted methods per class (WMC)	Number of local methods defined in the class	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
3 Annual Loss Expectancy (ALE)						
4 Cost of security breach					ON	
5 Remediation Impact (RI)	The remediation impact provides an overview of how much impact the remediation could have on the affected system				ON	
6 Risk Potential						
7 Threat-to-impact transitions						
3 Access accuracy	Number of correctly configured user accounts, against the overall number of accounts created, including badly configured accounts and hanging accounts	ORGANISATION	AUTOMATIC*	DYNAMIC*	ON	This is not related to embedded systems
Administrator and operator logs	Total number corrective actions taken after specific event is recorded in log $/$ -total number of events recorded in $\log$ * $100$					
180 Anomalous session count	The first phase derives a SessionTableAccessProfile by correlating application server user log entries for a user session with accessed database tables; the resulting value of SessionTableAccessProfile is represented as a user ID followed by a set of ordered pairs with a table name and a count. The second phase derives the AnomalousSessionCount by counting how many SessionTableAccessCounts don't fit a predefined user profile. If AnamalousSessionCount is greater than one for any user, especially a privileged user, it could indicate the need for significant refactoring and redesign of the Web application's persistence layer. This is a clear case in which detection at design time is preferable.	SOFTWARE			5.5	If adapted, it could be used
. Approval accuracy	Number of approved provisioning activities, against the overall provisioning activities, including the unauthorized ones	ORGANISATION	SEMI-AUTOMATIC*	DYNAMIC*	NO	This is not related to embedded systems
2 Arc coverage						
3 Audit logging	Total number of log files monitored in specific time interval / number of available log files * 100	SOFTWARE	SEMI-AUTOMATIC*	DYNAMIC*	ON	Although this has to do with larger systems, embedded systems could not be able to produce log files
l Block coverage						
Glassified Accessor Attribute Interactions (CAAI)	The ratio of the sum of all interactions between accessors and classified attributes - to the possible maximum number of interactions between accessors and classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	:5	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Attributes Interaction Weight (CAIW)	The ratio of the number of all interactions with classified attributes to the total $$ number of all interactions with all attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Class Data Accessibility (CCDA)	The ratio of the number of non-private classified class attributes to the number of classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Instance Data Accessibility (CIDA)	The ratio of the number of non-private classified instance attributes to the number of classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Method Extensibility (CME)	The ratio of the number of the non-finalised classified methods in program to the total number of classified methods in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Methods Weight (CMW)	The ratio of the number of classified methods to the total number of methods in the program $% \left( \frac{1}{2}\right) =\frac{1}{2}\left( \frac{1}{2}\right) =\frac{1}$	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Mutator Attribute Interactions (CMAI)	The ratio of the sum of all interactions between mutators and classified attributes - to the possible maximum number of interactions between mutators and classified attributes in the program.	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	¿?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Operation Accessibility (COA)	The ratio of the number of non-private classified methods to the number of $$ -classified methods in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	¿?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classified Writing Methods Proportion (CWMP)	The ratio of the number of methods which write classified attributes to the total number of classified methods in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
: Composite-Part Critical Classes (CPCC)	The ratio of the number of critical composed-part classes to the total number of critical classes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Glasses Design Proportion (CDP)	Ratio of the number of critical classes to the total number of classes, from the group of Design Size Metrics.	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	:?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software

196 Controls against malicious code	Number of incidents as results of malware (malicious software) outbreaks on system / number of detected and blocked malware occurrences * 100	SOFTWARE	SEMI-AUTOMATIC* DYNAMIC*	ON	This metric could be adapted to embedded systems, but it might not make sense in the testing phase
197 Countermeasure-effectiveness					
198 Coverage					
199 Coverage of Hazard Analysis					
200 Critical Class Extensibility (CCE)	The ratio of the number of the non-finalised critical classes in program to the total number of critical classes in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC*	:?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
201 Critical Design Proportion (CDP)	The ratio of the number of critical classes to the total number of classes in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
202 Critical Element Ratio	A process may not require all aspects of an object to be instantiated. Critical Elements nan Object) / (Total number of elements in the object)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
203 Critical Serialized Classes Proportion (CSCP)	The ratio of the number of critical serialized classes to the total number of critical classes in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
204 Critical Superclasses Inheritance (CSI)	The ratio of the sum of classes which may inherit from each critical superclass to the number of possible inheritances from all critical classes in the program's inheritance hierarchy	SOFTWARE		5?	
205 Critical Superclasses Proportion (CSP)	The ratio of the number of critical superclasses to the total number of critical classes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC* STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
206 Depth of inspection	It calculates the depth of inspection by collecting the statistical data of inspection and testing approaches and finds out the defect capturing capability as a ratio. DI can be measured phase-wise or as a whole before the deployment of the product. He defects captured by inspection process / Hefects captured by both inspection and testing approaches	SOFTWARE	SEMI-AUTOMATIC* STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
207 EIR - ratio between external and internal data flow					
208 Fail-Safe Defaults (PFSD)					
210 Hazard Analysis Achieved					
211 Isolation (PI)	This assesses the level of security isolation between system components. This means getting privileges to a component does not imply accessibility of other co-located components. This metric can be assessed using system architecture and deployment models. Components marked as confidential should not be hosted with nonconfidential (public) components. Methods that are not marked as confidential should not have access to confidential attributes or methods	SOFTWARE	SEMI-AUTOMATIC* STATIC*	÷?	I am not sure if this metric can be applied to embedded software
212 Least Common Mechanism (PLCM)					
213 Number of Catch Blocks Per Class (NCBC)	It measures the exception handing factor (EHF) in some way, it is defined as the percentage of catch blocks in a class over the theoretical maximum number of catch blocks in the class	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
214 Percentage High-, medium- and moderate-Risk Software Hazards with Safety Requirements	It reveals whether all high-, medium- and moderate-risk software hazards have resulted in applicable safety requirements through hazard analysis	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	This is related to software and the design
215 Percentage of Software Safety Requirements (PSSR)	How sufficient hazard analysis has been performed. (# Software Safety Requirenents) * 100 ments / # SoftwareRequirements) * 100	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	This is related to software and the design
216 Percentage of Software Safety Requirements Traceable to Hazards	Ensuring traceability to system hazards or System of Systems hazards increases - the validation case. Percentage indicator of traceability of requirements. All derived software safety requirements must be traceable to system hazards or Systems hazards eff Traceable Software Safety Requirements) * 100	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	This is related to software and the design
217 Percentage of identified corrective action that has not been implemented		ORGANISATION		5?	
218 Percentage of incidents that are a recurrence of previous incidents		ORGANISATION	MANUAL* DYNAMIC*	ON	This a high level metric
219 Percentage of increases in reported security incidents		ORGANISATION	SEMI-AUTOMATIC* DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
220 Percentage of IT assets for which fault tolerance mechanisms have been implemented		ORGANISATION	AUTOMATIC* DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
221 Percentage of IT assets for which recovery procedures have been defined and implemented		ORGANISATION	AUTOMATIC* DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
222 Percentage of IT assets for which redundancy mechanisms have been implemented		ORGANISATION	AUTOMATIC* DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
223 Percentage of new components that were deployed in the system all at once		ORGANISATION			
224 Percentage of organization attended security training		ORGANISATION			
225 Percentage of Organization contributing to development		ORGANISATION			
226 Percentage of reported incidents that have been followed up and mitigated		ORGANISATION	SEMI-AUTOMATIC* DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
227 Percentage of reported security incidents where the cause of the incident was identified		ORGANISATION			

228 Percentage of security incidents related to incorrect, incomplete, missing or compromised audit data		ORGANISATION	MANUAL*	DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
229 Percentage of security incidents related to lack of an audit capability		ORGANISATION	MANUAL*	DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
230 Percentage of security incidents related to the ability to bypass an audit function		ORGANISATION	MANUAL*	DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
231 Percentage of security incidents that exploited existing vulnerabilities with known solutions		ORGANISATION	SEMI-AUTOMATIC* DYNAMIC	DYNAMIC*	ON	This is a high level metric, and it is related with attacks on a organisation's network
232 Percentage of servers with installed and active automatic hard-disk encryption		NETWORK			ON	Although this cannot be applied directly, this can be something to check for in the checklist
233 Percentage of system architecture for which failure modes have been thoroughly identified		ORGANISATION	SEMI-AUTOMATIC* DYNAMIC	DYNAMIC*	;?	This is a high level metric, but the concept might be applied to embedded systems
234 Percentage of system changes that were reviewed for security impacts before implementation		ORGANISATION				
235 Percentage of system for which approved configuration settings has been implemented		ORGANISATION	AUTOMATIC*	DYNAMIC*	ON	This is a high level metric. Different configurations of the device can be tested, but this might not be suitable for assessment
236 Percentage of system that depends on external components that the organization lacks control over		ORGANISATION	SEMI-AUTOMATIC*	DYNAMIC*	<i>i?</i>	This can be applied to components of thirds parties, as libraries that have not been tested
237 Percentage of system that has been subject to risk analysis		ORGANISATION	AUTOMATIC*	STATIC*	;?	If the risk analysis can be done just to a percentage of the device, this can be applied. Although it might not make sense to do that
238 Percentage of system that is continuously monitored for configuration policy compliance		ORGANISATION	AUTOMATIC*	DYNAMIC*	NO	This is a high level metric
239 Percentage of system where permissions to install non- standard software is limited or prohibited		ORGANISATION	AUTOMATIC*	DYNAMIC*	ON	Software cannot be installed on a embedded system
240 Percentage of system where the authority to make configuration changes is limited in accordance to policy		ORGANISATION	AUTOMATIC*	DYNAMIC*	ON	This is a high level metric
241 Percentage of systems with the latest approved patches installed		ORGANISATION	AUTOMATIC*	DYNAMIC*	ON	This is related to an organisation security
242 Percentage of sytems exposed at time of malware						
PercentValidatedInput	To compute this metric, let T equal the count of the amount of input forms or interfaces the application exposes (the number of HTML form POSTs, GETs, and so on) and let V equal the number of these interfaces that use input validation mechanisms. The ratio V/T makes a strong statement about the Web application's vulnerability to exploits from invalid input	SOFTWARE			YES	This could be adapted to other pieces of code, rather than just HTML
244 Ratio of Design Decisions (RDD)	Ratio of the number of design decisions related to security to the total number - of design decisions of the entire system. The objective is to assess the portion of design declicated to security	SOFIWARE	SEMI-AUTOMATIC*	STATIC*	¿?	This is related to design phase
245 Ratio of Design Flaws Related to Security (RDF)	Ratio of design flaws related to security to the total number of design flaws applicable to the whole system	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	<i>??</i>	This is related to design phase
246 Ratio of extension misuse cases extended once to the total number of extension misuse cases	Are the functional decompositions between misuse cases correctly handled? 1 - (# Extension Misuse Cases included once / # Extension Misuse Cases)	SOFTWARE	MANUAL*	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
247 Ratio of Implementation Errors That Have an Impact on Security (RSERR)	Ratio of the number of errors related to security to the total number of errors in the implementation of the system (i.e. NERR)	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	<i>i</i> ?	This is related to design phase
248 Ratio of inclusion misuse cases included once to the total number of inclusion misuse cases	I	SOFTWARE	MANUAL*	STATIC*	NO	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
249 Ratio of misuse cases used as pre/post conditions of other misuse cases to the total number of misuse cases	Are the functional decompositions between misuse cases correctly handled? (# - Misuse Cases used as pre/post conditions) / (# Misuse Cases)	SOFTWARE	MANUAL*	STATIC*	NO	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
250 Ratio of Patches Issued to Address Security Vulnerabilities (RP)		SOFTWARE	SEMI-AUTOMATIC*	STATIC*	;?	If the design is available, it can be applied
251 Ratio of Security Requirements (RSR)	Ratio of the number of requirements that have direct impact on security to the total number of requirements of the system	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	:?	If the design is available, it can be applied
252 Ratio of Security Test Cases (RTC)	Number of test cases designed to detect security issues to the total number of $\;$ - test cases of the entire system	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	¿?	If the design is available, it can be applied
253 Ratio of Security Test Cases that Fail (RTCP)	Number of test cases that detect implementation errors (i.e. the ones that fail) to the total number of test cases, designed specifically to target security issues	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	<i>i?</i>	If the design is available, it can be applied
254 Ratio of Shared Resources (RSR)						
Ratio of Software Changes Due to Security Considerations (RSC)	Number of changes in the system triggered by a new set of security requirements. Software changes due to security considerations include patches that are released after a system is delivered, or any other security updates	SOFTWARE	SEMI-AUTOMATIC* STATIC	STATIC*	¿?	If the design is available, it can be applied
256 Ratio of the number of included security requirements to the total number of stated security requirements	How vulnerable is the application based on the stated security requirements? 1 ( (# Stated Security Requirements - # Excluded Stated Security Requirements ) / # Stated Security Requirements )	SOFTWARE	MANUAL*	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
257 Ratio of the number of misuse cases that do not threaten the application to the total number of misuse cases	Do the misuse cases correctly represent the application vulnerabilities and are they consistent with application security use cases? (# Non-Threatening Misuse Cases) / (# Misuse Cases)	SOFTWARE	MANUAL*	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
258 Ratio of the Number of Omitted Exceptions (ROEX)	Ratio of the number of omitted exceptions (i.e. Noex) to the total number of exceptions that are related to security	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	23	If the software is available, it can be applied. It can be also applied to embedded software

259 Ratio of the Number of Omitted Security Requirements (ROSR)	Ratio of the Number of Omitted Security Requirements Ratio of the number of security requirements that have not been considered dur- (ROSR) to the total number of security requirements ing the analysis phase (i.e. NOSR) to the total number of security requirements	SOFTWARE	SEMI-AUTOMATIC* STATIC*	;?	If the design is available, it can be applied
260 Ratio of the number of the base misuse cases associated	identified during the analysis phase (NSK) Are the misusers presented and handled correctly in the misuse case model? 1	SOFTWARE	MANUAL* STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a
261 Ratio of the number of unmitigated misuse cases that threaten the application to the total number of misuse cases		SOFTWARE	MANUAL* STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
262 Readability of Classified Attributes (RCA)					
263 Readability via Classified Methods (RCM)					
264 Readability via Critical Classes (RCC)					
265 Readability via Critical Classes (RCC)					
266 Software Hazard Analysis Depth	Hazardous software, or safety-critical software, allocated as high- or mediumrisk, according to the Software Hazard Criticality Matrix, requires analysis of second- and third-order causal factors (should they exist)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	Hazard is similar to risk.
267 Static analysis alert density		SOFTWARE			
268 Static analysis vulnerability density	Vulnerability metrics are typically normalized by code size to create vulnerability density metrics. Normalizing by size lets us compare software components of varying size. Number of vulnerabilities a static-analysis tool finds per thousand LOC				
269 Target Distribution (TD)					
270 Unaccessed Assigned Classified Attribute (UACA)	The ratio of the number of classified attributes that are assigned but never used to the total number of classified attributes in the program	SOFTWARE	-		Object-oriented programs
271 Uncalled Classified Accessor Method (UCAM)	The ratio of the number of classified methods that access a classified attribute but are never called by other methods to the total number of classified methods in the program	SOFTWARE			Object-oriented programs
272 Unused Critical Accessor Class (UCAC)	The ratio of the number of classes which contain classified methods that access classified attributes but are never used by other classes to the total number of critical classes in the program	SOFTWARE		1	Object-oriented programs
273 Window of exposure					
274 Writability of Classified Attributes (WCA)					
275 Writability via Classified Methods (WCM)					
276 Writability via Critical Classes (WCC)					
277 Average internal data flow (AIDF)					
278 Average external data flow (EDF)					
279 Average incoming data flow (EDFIN)					
280 Average outgoing data flow (EDFOUT)					
281 Classified Attributes Inheritance (CAI)	The ratio of the number of classified attributes which can be inherited in a hierarchy to the total number of classified attributes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC* STATIC*	:?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
282 Classified Methods Inheritance (CMI)	The ratio of the number of classified methods which can be inherited in a hierarachy to the total number of classified methods in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>؟؟</i>	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
283 Coupling	Measures the following three coupling dimensions between modules: referential dependency, structural dependency, and data integrity dependency				
284 Coupling between components					
285 Coupling Between Object classes (CBOC)	Number of other classes coupled to a class C	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
286 Coupling Between Objects (CBO)					
287 Coupling Corruption Propagation	Coupling corruption propagation is meant to measure the total number of methods that could be affected by erroneous originating method. Coupling Corruption Propagation = Number of child methods invoked with the parameter(s) based on the parameter(s) of the original invocation	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
288 Critical Classes Coupling (CCC)	The ratio of the number of all classes' links with classified attributes to the total number of possible links with classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
289 Depth of Inheritance Tree (DIT)	Maximum depth of the class in the inheritance tree	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
290 Eigenvector Centrality (EvCent)					
291 Flow_Betweenness					
292 Incoming closure					
293 Incoming data flow of (EDFIN)					
294 Incoming direct					
295 InDegree					
296 InDegree_w					

18   Statistical Control Con	298 Lack of cohesion of methods (LCOM)	The LCOM value for a class C is defined as $LCOM(C) = \frac{(1- E(C) )}{ V(C) M(C) }$ , where V(C) is the set of instance variables, M(C) is the set of instance methods, and E(C) is the set of pairs (v,m) for each instance variable v in V(C) that is used by method m in M(C)	SOFTWARE	А∪ТОМАПС*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
The number of other to be service without into meeting.  A Machine and the regions of the first interest interest of the first interest interest interest of the first interest i	99 Layer information						
The number of calls to the functions delands in an entry profit of the functions delands in an entry of the function of the functions delands in an entry of the functions delands in an entry of the function	0 Number of edges between two arbitrary nodes (IE)						
Processing data Bear.  A benderming by the particular of the particular of particular	1 NumCalls	The number of calls to the functions defined in an entity	SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embed software
Notice that the control of the contr	2 OutDegree						
Aboolean value representing whether the Java program impact the reduction . SOLTWAKE	3 OutDegree_w						
A booken when the imposenting whether the love program imposs the reflection - SOTTWARE - STATIC - STA	1 Outgoing closure						
A backetor value impresenting whether the Java program impacts the reduction - SOPTWARE	Outgoing data flow of (EDFOUT)						
Publication to the expressibility in Properties of Part (Control of Part) Proc. Learner (Contr	Outgoing direct						
Packetan value representing whether the Jean program imports the reflection Aboolean value representing whether the Jean program imports the reflection and the Aboolean value representing whether the Jean program imports the reflection and the Aboolean value representing the Aboolean value and the	Paths						
Abdough when the representing a before the jane program imports the reflection - SOTTWARE ALTDAATIC STATIC TO THIS class of the control of the processing of	Ratio between average outgoing and incoming data (OIR)	flow					
the control of the properties of the control of the	Reflection Package Boolean (RPB)	A boolean value representing whether the Java program imports the reflection package (1) or not (0)	SOFTWARE				Object-oriented programs
When who have greater that the control of the contr	RW_Betweenness						
When chairs by propositor by Carbon and Carb	SP_Betweenness						
Also, there are more than one that one than one that one than one than one than one than one than one than one that one than one than one than one than one that one than one than one than one that of the than one that the that one that one that one that the that one that one that the that one that one that the that o	Vulnerability Propagation (VP)	Vulnerability Propagation (VP) of a class C, denoted by VP(C), is the set containing classes in the hierarchy which directly or indirectly inherit class C. Cardinality of set VP(C) represents the number of classes which have become vulnerable due to class C. Alternatively, Vulnerability Propagation of a class C is the total number of classes which directly or indirectly inherit class C.	SOFTWARE	AUTOMATIC	STATIC*	YES	This metric can be applied to embedded software
Neces Completely (AC)  This metric insurance to completely to a vibrate set obtaining the venders delity A vibrate in seeds  Neces Ventro(AV)  This metric insurance to completely on the explorite at wallier setting to a vibrate and a vibrate a vibrate and a vibrate an	Vulnerability Propagation Factor (VPF)	An Inheritance hierarchy may contain no class or one or more vulnerable classes.  Also, there are more than one Inheritance hierarchies present in a design. So, calculation of Vulnerability Propagation Factor (VPF) of a design requires calculation of Vulnerability Propagation due to each vulnerable class in Inheritance hierarchies present in the design. Then, union of Vulnerability Propagation due to each vulnerable class will give overall Vulnerability Propagation in design.	SOFTWARE	А∪ТОМАПС	STATIC*	YES	This metric can be applied to embedded software
Necess (Verbor (AV)  Adversary Work Eactor  Adversary Work Eactor  Adversary Eactor  Adversa	1	This metric measures the complexity of attacks exploiting the vulnerability. A vulnerability with low complexity can be for example a buffer overflow in a web server, the vulnerability can be exploited at will					
Makeslay Work Earther and Adversary work facts it an information system. Black a comparison between the season of the continuous comparisons of modern information systems. However, this nettice can be difficult to measure and configuration systems. However, this nettic can be defined to the season of the seas		This metric indicates from where an attacker can exploit the vulnerability					
Attendability         Description         Description of the standard palls and on the change of the	1	Adversary work factor is an informative metric for gauging the relative strengths and weaknesses of modern information systems. However, this metric can be difficult to measure and evaluate		MANUAL*	DYNAMIC*	ON	This metric has nothing to do with embedded systems and it is directly re to networks
Authentication (AU)  In this metric counts know often an attacker must authenticate before the valnera- Condictantiality Requirement (CR)  Damage potential-effort ratio  In indicates the level of damage an attacker can potentially cause to the system - Damage potential-effort ratio  Depth  Exclusive Exclusive Inner for the set of functions, S, defined in an entity evaluding the execution time for the set of functions called by the functions in S defined in an entity including all the execution time for the set of functions S, defined in an entity including all the execution time for the set of functions S, defined in an entity including all the execution time spent by the functions alled directly or indirectly by the functions alled the functions functions alled the functions	Attackability	The number of inbound edges of a node in the hostbased attack graph represents all the direct attack paths that other hosts in the considered network can compromise to that node. This number can be used to compute the attackability metric of a host in the context of the system under study. The metric is based on intuitive properties derived from common sense. For example, the metric will indicate reduced confidence when more inbound edges exist. (1 - # inbound edges of the node/ total inbound edges)	NETWORK	А∪ТОМАПС*	STATIC*	ON	This metric has nothing to do with embedded systems and it is directly retaled to networks
Damage potential (CR)  It indicates the level of damage an attacker can potentially cause to the system - DEVICE SEMI-AUTOMATIC* DYNAMIC*  Damage potential effort ratio and the effort required for the attacker to cause such damage  Depth  Execution time for the set of functions. S, defined in an entity excluding the execution time spent by the functions called by the functions in S  Execution time for the set of functions. S, defined in an entity including all the execution time spent by the functions alled directly or indirectly by the functions  Execution time for the set of functions. S, defined in an entity including all the execution time spent by the functions alled directly or indirectly by the functions  Execution time spent by the functions alled directly or indirectly by the functions  Execution time spent by the functions alled directly or indirectly by the functions  Execution time spent by the functions alled directly or indirectly by the functions  Execution time spent by the functions alled directly or indirectly by the functions  Execution time spent by the functions alled directly or indirectly by the functions  Execution of an attack copy asystem  An intuition behind this metric is the following: the less steps an attacker has - NETWORK	Authentication (AU)	This metric counts how often an attacker must authenticate before the vulnerability can be exploited					
Damage potential effort ratio and the effort required for the attacker to cause such damage  Depth  Excution time for the set of functions, S, defined in an entity excluding the excution time spent by the functions called by the functions in S  Exposted Time to Completion  Exposted Time to Completion  Excution time spent by the functions called by the functions in S  Exposted Time to Completion  Exposted Time to Completion  Execution time spent by the functions called directly or indirectly by the functions  Expostated Time to Completion  Exposted Time to Completion  Expostrated Time to Completion  Expostrate Time to Completion  Expostrated Time to Completion  Expostrate Time Time to Completion  Expostrate Time Time to Completion  Expostrate Time Time Time Time Time Time Time Tim	Confidentiality Requirement (CR)						
Execution time for the set of functions, S, defined in an entity excluding the execution time spent by the functions called by the functions in S. Execution time spent by the functions called by the functions in S. Execution time for the set of functions, S, defined in an entity including all the execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions in S. Execution time spent by the functions called directly or indirectly by the functions called directly or indirectly by the functions in S. Execution of an attack cone a system.  Minimal cost of attack represents the minimal cost that the attacker has to pay an attacker has to pay an attacker has to pay an attacker has a secure the stands are consistilly, and the less secure the spent by the system is system in the system in the spent by the spent by the system is spiken in the spent by the spent by the spent by the spiken in the spent by the spent by the spent by the spiken in the spiken	Damage potential-effort ratio	It indicates the level of damage an attacker can potentially cause to the system and the effort required for the attacker to cause such damage	DEVICE	SEMI-AUTOMATIC		YES	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase
Execution time spent by the functions, S, defined in an entity excluding the execution time spent by the functions called by the functions in S. Execution time spent by the functions called by the functions are spent by the functions called by the functions in S. Execution time for the set of functions called directly or indirectly by the functions in S. Indirectly by the functions called directly or indirectly by the functions called directly called a function called directly or indirectly the simple called directly called a function behind this metric is the following: the less steps an attacker has called a function is system is system is system in called a function of an attacker has called a function behind this metric is the following: the less secure the system is system is system in called a function of an attacker has called a function of an entity the function of an entity the function of an entity the simple called a function of an entity the function of an entity in the function of an entity in the function of an entity of the fun	Depth						
Expected Time to Completion  Exploitability (TE)  Inclusive Execution time for the set of functions, S, defined in an entity including all the execution time spent by the functions called directly or indirectly by the functions  Inclusive Exercise  Inclusive Exercise time execution time spent by the functions called directly or indirectly by the functions in S  Inclusive Exercise time execution time spent by the functions called directly or indirectly by the functions in S  Inclusive Exercise time execution of an attack represents the minimal cost that the attacker has to pay - Includent to security Failure  Minimal cost of attack  An intuition behind this metric is the following: the less steps an attacker has - Includent to an attack the simpler is to execute the attack successfully, and the less secure the system is	ExclusiveExeTime	Execution time for the set of functions, S, defined in an entity excluding the execution time spent by the functions called by the functions in S					
Execution time for the set of functions, S, defined in an entity including all the execution time spent by the functions called directly or indirectly by the functions all the execution time spent by the functions called directly or indirectly by the functions all the execution time spent by the functions called directly or indirectly by the function of an attack on a system  Minimal length of attack  Mi	Expected Time to Completion						
Inclusive ExeTime  Execution time spent by the functions, S, defined in an entity including all the execution time spent by the functions called directly or indirectly by the functions in S  Integrity Requirement (IR)  Mean Effort to security Failure  Minimal cost of attack represents the minimal cost that the attacker has to pay -  Minimal cost of attack  Minimal cost of attack on a system  An intuition behind this metric is the following: the less steps an attacker has -  Minimal length of attacks  An intuition behind this metric is the following: the less secure the system is system is system is a system is a system is system	Exploitability (TE)						
Mean Effort to security Failure       Minimal cost of attack represents the minimal cost that the attacker has to pay - for the execution of an attack on a system       NETWORK*       -       NO         Minimal length of attack       An intuition behind this metric is the following: the less steps an attacker has to make, the simpler is to execute the attack successfully, and the less secure the system is system is       NETWORK*       -       -       NO	InclusiveExeTime	Execution time for the set of functions, 5, defined in an entity including all the execution time spent by the functions called directly or indirectly by the functions in S					
Mean Effort to security Failure       Minimal cost of attack represents the minimal cost that the attacker has to pay -       NETWORK*       -       NO         Minimal cost of attack       An intuition behind this metric is the following: the less steps an attacker has to make, the simpler is to execute the attack successfully, and the less secure the system is       NETWORK*       -       NO	Integrity Requirement (IR)						
Minimal cost of attack Minimal cost of attack represents the minimal cost that the attacker has to pay - NETWORK* NO NO NETWORK*	Mean Effort to security Failure						
Minimal length of attacks An intuition behind this metric is the following: the less steps an attacker has - NETWORK* NO NO to make, the simpler is to execute the attack successfully, and the less secure the system is	Minimal cost of attack	Minimal cost of attack represents the minimal cost that the attacker has to pay for the execution of an attack on a system	NETWORK*	-	-	ON	Formal model. Formal verification. Too much notation. I think this is n related to networks
	Minimal length of attacks	An intuition behind this metric is the following: the less steps an attacker has to make, the simpler is to execute the attack successfully, and the less secure the system is	NETWORK*	1		ON	Formal model. Formal verification. Too much notation. I think this is more related to networks

330 Protection Rate (PR)	It expresses the efficiency of security mechanisms. Percentage of the known vulnerabilities protected by a given security mechanism. It is based on the Attack Surface metric. PR = (1 - AttackSurfaceEvaluatedSystem) $^\prime$ 100 enceSystem) $^\ast$ 100	SOFTWARE	SEMI-AUTOMATIC* DYNAMIC*	¿?	This metric was developed specificaly for OSGi platform, but I think it can be adapted to be applied to embedded systems
331 Rigor					
332 Side-channel Vulnerability Factor (SVF)	SVF quantifies patterns in attackers' observations and measures their correlation to the victim's actual execution patterns and in doing so captures systems' vulnerability to side-channel attacks, we can measure information leakage by computing the correlation between ground-truth patterns and attacker observed patterns. We call this correlation between ground-truth patterns and attacker observed patterns. We call this correlation Side-channel Vulnerability Factor (SVF). SVF measures the signal-to-noise ratio in an attacker's observations. While any amount of leakage could compromise a system, a low signal-to-noise ratio means that the attacker must either make do with inaccurate results (and thus make many observations to create an accurate result) or become much more intelligent about recovering the original signal. We use phase detection techniques to find patterns in both sets of data then compute the correlation between actual patterns in the victim and observed patterns.	HARDWARE	SEMI-AUTOMATIC* DYNAMIC*	YES	Embedded devices can be vulnerable to side-channel attacks and they can be tested
333 Social Engineering Resistance (SER)	$SER = \langle SERlow, SERhigh>; SERlow = (1 - (P + N - A)/N) * 100; SERhigh = (1 - P/N) * 100; where N is the number of individuals selected or invited to take part in the experiment, A is the number of active participants; P is the total number of valid password/username pairs obtained during the experiment$	USER	SEMI-AUTOMATIC* DYNAMIC*	ON	Social engineering cannot be applied to a device
334 Structural severity	Measure that uses software attributes to evaluate the risk of an attacker reaching a vulnerability location from attack surface entry points. It is measured based on three values: high (reachable from an entry point), medium (reachable from an entry point with no dangerous system calls), low (not reachable from any entry points)	SOFTWARE	SEMI-AUTOMATIC* DYNAMIC*	53	I am not sure if this metric can be applied to embedded software
335 Vulnerability exploitability					
336 Weakest adversary	It measures the security strength of a network in terms of the strength of the weakest adversary (i.e., requiring least amount of effort) who can successfully penetrate the network. We define a network configuration that is vulnerable to a weaker set of initial attacker attributes as less secure than a network configuration vulnerable to a stronger set of attributes	NETWORK	AUTOMATIC* DYNAMIC*	ON	Similar to 31. This is for networks, not for embedded systems
337 Vulnerability Index (VI)	Probability of a component's vulnerability being exposed in a single execution	DEVICE	MANUAL* DYNAMIC*	5?	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase
338 Effective Vulnerability Index (EVI)	Relative measure of the impact of a component on the system's insecurity	DEVICE	SEMI-AUTOMATIC* DYNAMIC*	:?	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase
339 CNBetweenness				ON	
340 CNCloseness				ON	
341 Depth of Master Ownership	This metric determines the level of ownership of the binary depending on the number of edits done. The organization level of the person whose reporting engineers perform more than 75% of the rolled up edits is considered as the DMO. This metric determines the binary owner based on activity on that binary. Our choice of 75% is based on prior historical information on Windows to quantify ownership.			ON	
342 DNAvgBetweenness				ON	
343 DNAvgCloseness				ON	
344 DNAvgEdgeBetweenness				ON	
345 DNMaxCloseness				ON	
346 DNMaxDegree				ON	
347 DNMaxEdgeBetweenness				ON	
346 D.N.Minbetweenness				ON OX	
350 Edit Frequency	Total number of times the source code that makes up the binary was edited. An edit is when an engineer checks out code from the version control system, alters it, and checks it in again. This is independent of the number of lines of code altered during the edit			ON	
351 Level of Organizational Code Ownership	The percent of edits from the organization that contains the binary owner (or if there is no owner the percent of edits from the organization that made the majority of the edits to that binary)	ORGANISATION		ON	It is at higher level. It is related to the organisation's security
352 New Effective Author (NEA)				ON	
353 Number of Developers		ORGANISATION		ON	It is at higher level. It is related to the organisation's security
354 Number of Ex-Engineers	Total number of unique engineers who have touched a binary and have left the company as of the release date of the software system $ \\$	ORGANISATION		ON	It is at higher level. It is related to the organisation's security
355 Organization Intersection Factor	The number of different organizations that contribute more than $10\%$ of edits, as measured at the level of the overall org owners			ON	
356 Overall Organization Ownership	This is the ratio of the people at the DMO level making edits to a binary relative $$ to total engineers editing the binary	ORGANISATION		ON	It is at higher level. It is related to the organisation's security
357 Arithmetic Expression					
358 Array index					

SOFTWARE	359 Attack Surface Metric	This metric has been proposed by Howard and Manadhata and Wing. Here we consider one of the latest versions of attack surface metric presented by Manadhata et al.	DEVICE*			•	Similar to 8. Formal model. Formal verification. Too much notation. I think this is more related to networks
Contact Action of Contact Ac	1		SOFTWARE				
Control Methods (100 CAC)  Control Methods (100		The total number of classified attributes in the program	SOFTWARE				Object-oriented programs
Continue bacteries   Continu	1	The total number of classified methods in the program	SOFTWARE		,	,	Object-oriented programs
Control Appropriate Contro			SOFTWARE				
Centrel Specified         SCHTPAME         CENTRAME         CENTRAME <td></td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td>							
Control Clause National Control Clause and							
Decision of Statistics of St		Number of base classes	SOFTWARE	-	-		
Company and Auto-Auto-Auto-Auto-Auto-Auto-Auto-Auto-		The total number of critical classes in the program	SOFTWARE				Object-oriented programs
Execute Section   Execute Se							
Commission requested betalety							
Interface of color (1992)							
Name of a planting   Name of							
Description							
Number of Attacks  Number of Design Decisione Related in Security Number of Attacks  Number of Design Decisione Related in Security Number of Attacks  Number of Design Decisione Related in Security Number of Maria Shi Promoting States (States States)  Number of Design Decisione Related in Security Number of Design Decisioner Related Security Number of Security Number			SOFTWARE				
Number of Abesides Microbial State of Abesides State S			SOFTWARE				
Name of A Markes    A Mission should show any state of the Appear of The Age A Markes    A Mission should show any state of the Age A Markes    A Mission should show and the Age A Markes    A Mission should show any state of the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should show and the Age A Markes    A Mission should s			NETWORK				
Number of beging connection to accountly requirements that remarks and objects that the second of th		How many attacks on a system exist. The idea behind this metric is that the more attacks on a system exist the less secure the system is. This metric is applied for the simplest analysis of attack graphs. Number of attacks also can be used for the analysis of results of the penetration testing	NETWORK*	1	1	ON	Formal model. Formal verification. Too much notation. I think this is more related to networks
Number of Deept Decision Related to Security (NDD)  Number of Decision Related to Security (NDD)  Number of Security requirements that recurs and security and security or security delicities to those armseg alteriates are problem, showner and security or security and security and security or security and security or security and security and security or security and security requirements security and secur							
Number of breslopes  Number of blooksesses, but the street and of several by the street and by the			SOFTWARE	SEMI-AUTOMATIC*	l e	<i>3</i>	If the design is available, it can be applied
Number of elicited security requirements that tensure is session identifier created on server saled. Is more saled in the control of the control of server saled. Is more saled in the control of server saled. It may be control or server saled in the control or server saled. It may be control or server saled in the control or server saled. It may be control or server saled in the control or server sale in the control or server saled in the control or saled in the control or server saled in the control or server saled in the control or server saled in the control or saled in the control			ORGANISATION			ON	It has to do with the organisation and the development of the product. It is applied at the beginning of the life-cycle
Number of sociative generations that ensure because the tensor of admitted can administrations be associated bandle of a control of the control of		What is the number of elicited security use cases? # Elicited Security Use Cases $$ -	SOFTWARE	MANUAL	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
Number of seclaridy requirements that put the Summation of the excluded security requirements that put the System at risk of possible attacks.  Number of solida variables		SOFTWARE	MANUAL	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment	
Number of global variables*  Number of identified misuse cases found? # Misuse Cases found? # Misuse Cases found? # Misuse Cases  Number of incoming connections (RAN)  Number of mountain method signature  Number of parameters in the method signature  Number of parameters in the method signature  Number of security Algorithms (NSA)  Number of security algorithms that are supported by the application  Number of Security Algorithms (NSA)  Number of Security Requirements (NSB)  Number			SOFTWARE	MANUAL	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
Number of identified misuse cases to what is the number of misuse cases found? # Misuse Cases  Number of incenting connections (RIN)  Number of member and security Algorithms (NSA)  Number of Security Algorithms (NSA)  Number of Security Requirements (ISR)*  Number of security Requirements (ISR)*  Number of subclasses  Number of			SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software
Number of accounting connections (RIN)  Number of burnished voluenting connections (RIN)  Number of burnished vulnerabilities  Number of Security Algorithms (NSR)  Number of Security Incidents Reported (NSR)  Number of Security Requirements (NSR)  Number of Security Requirements (NSR)  Number of Security Incidents Reported (NSR)  Numb		What is the number of misuse cases found? # Misuse Cases	SOFTWARE	MANUAL	STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
Number of ougoing connections from (ROUT)  Number of ougoing connections from (ROUT)  Number of parameters in the method signature  Number of published vulnerabilities  Number of security Algorithms (NSA)  Number of Security Algorithms (NSA)  Number of Security Incidents Reported (NSR)  Number of Security Incidents Reported (NSR)  Number of Security Requirements (NSR)  Number of security requirements identified during the analysis phase of the -  SOFTWARE  SOFTWARE  SOFTWARE  SOFTWARE  SOFTWARE  SOFTWARE  Number of Security Requirements (NSR)  Number of security requirements identified during the analysis phase of the -  SOFTWARE  Number of security Requirements (NSR)  Number of security (NSR)  Number of security (NSR)  Number of security (NSR)  Num							
Number of outgoing connections from (ROUT)  Number of parameters in the method signature  Number of published vulnerabilities  Number of return points in the method  Number of security Algorithms (NSA)  Number of Security Algorithms (NSA)  Number of Security Requirements (NSR)*  Number							
Number of parameters in the method signature Number of published vulnerabilities  Number of published vulnerabilities  Number of security Algorithms that are supported by the application  Number of Security Requirements (NSR)  Number of sub classes  Number of sub classes  Number of sub classes  Number of sub classes  Number of Number of sub classes  Numb							
Number of published vulnerabilities  Number of return points in the method  Number of security Algorithms that are supported by the application  Number of Security Algorithms (NSA)  Number of Security Incidents Reported (NSR)  Number of Security Requirements (NSR)  Number							
Number of return points in the method  Number of security algorithms that are supported by the application  Number of Security Algorithms (NSA)  Number of Security Algorithms (NSA)  Number of Security Incidents Reported (NSR)  Number of Security Requirements (NSR)*  Number of Security			DEVICE				
Number of Security Algorithms (NSA)  Number of security algorithms that are supported by the application  Number of Security Incidents Reported (NSR)  System  Number of Security Incidents Reported (NSR)*  Number of Security Requirements (NSR)*  Number of sub classes  Number of sub classes  Number of sub classes  Number of Security Requirements (NSR)*  Number of sub classes  Number of Security Requirements (NSR)*  Number of sub classes  Number of Security Requirements (NSR)*  Number of sub classes  Number of Security Requirements (NSR)*  Number of Security Requ			SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software
Number of Security Incidents Reported (NSR)       Number of incidents related to security that are reported by the users of the system       SOFTWARE       SEMI-AUTOMATIC*       DYNAMIC*       NO         Number of Security Requirements (NSR)*       Number of security requirements identified during the analysis phase of the spplication       SOFTWARE       SOFTWARE       STATIC*       \$7         Number of sub classes       Number of sub classes       SOFTWARE       SOFTWARE       YES		Number of security algorithms that are supported by the application	SOFTWARE	SEMI-AUTOMATIC*		<i>??</i>	If the design is available, it can be applied, but this metric can be used in finals stages of the development life-cycle
Number of Security Requirements (NSR) *       Number of Security requirements identified during the analysis phase of the -       SOFTWARE       STATIC*       2?         Number of sub classes       SOFTWARE       YES         NumFunctions       SOFTWARE       YES		Number of incidents related to security that are reported by the users of the system	SOFTWARE	SEMI-AUTOMATIC*		ON	By definition, the values of this metric are collectec by the user. Son it can't be used in the valuation
Number of sub classes         SOFTWARE         YES           NumFunctions         SOFTWARE         YES		Number of security requirements identified during the analysis phase of the application	SOFTWARE	MANUAL	STATIC*	<i>??</i>	If the design is available, it can be applied
NumFunctions       SOFTWARE   YES			SOFTWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software
			SOFIWARE			YES	If the software is available, it can be applied. It can be also applied to embedded software

397 Paths		SOFTWARE			
398 Reduce Attack Surface (PRAS)					
399 Resource allocation					
400 Response set for a Class (RFC)	Set of methods that can potentially be executed in response to a message received by an object of that class. RFC is simply the number of methods in the set, including inherited methods	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
401 Security Absolute Measurements (SAM)					
402 Stall Ratio	Stall ratio is a measure of how much a program's progress is impeded by frivolous activities, stall ratio = (lines of non-progressive statements in a loop) / (total lines in the loop)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
403 Total Security Index (TSI)					
404 User authority		USER			
405 Volume of email correspondence with vulnerability handling team	han-				
407 Availability Requirement (AR)					
408 CMMI Level					
409 Comment ratio	Ratio of number of comment lines to number of code lines	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
410 CommentDensity	The ratio of lines of comments to lines of code				
411 Compartmentalization	Compartmentalization means that systems and their components run in different compartments, isolated from each other. Thus a compromise of any of them does not impact the others. This metric can be measured as the number of independent components that do not trust each other (performs authentication and authorization for requests/ calls coming from other system components) that the system is based on to deliver its function. The higher the compartmentalization value, the more secure the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	53	If the software is available, it can be applied. It can be also applied to embedded software
412 Completeness of fix					
413 Confidentiality					
414 Confidentiality Impact (C)					Similar to 22
415 Defense-In-Depth	This metric verifies that security controls are used at different points in the system chain including network security, host security, and application security. Components that have critical data should employ security controls in the network, host, and component layers. To assess this metric we need to capture system architecture and deployment models as well as the security architecture model. Then we can calculate the rath of components with critical data that apply the layered security principle compared to number of critical components	SOFTWARE	SEMI-AUTOMATIC* STATIC*	5.5	I am not sure if this metric can be applied to embedded software, or at verification phases
416 Expected Reliability					
417 Fail Securely	The system does not disclose any data that should not be disclosed ordinarily at system failure. This includes system data as well as data about the system in case of exceptions. This metric can be evaluated from the security control responses – i.e. how the control behaves in case it failed to operate. From the system architecture perspective, we can assess it as the number of critical attributes and methods that can be accessed in a given component. The smaller the metric value, the likely more secure the system in case of failure	SOFIWARE	SEMI-AUTOMATIC* STATIC*	YES	This could be used in evaluation
418 Independence of Verification					
419 Inspection Performance Metric (IPM)	It measures the performance of the people during the inspection process using the aforementioned parameters. # defects captured by inspection process / inspection effort	TESTER	SEMI-AUTOMATIC* STATIC*	ON	This metric is related to the tester
420 Integrity					
421 Isolation (P1)	This assesses the level of security isolation between system components. This means getting privileges to a component does not imply accessibility of other co-located components. This metric can be assessed using system architecture and deployment models. Components marked as confidential should not be hosted with nonconfidential (public) components. Methods that are not marked as confidential should not have access to confidential attributes or methods	SOFTWARE	SEMI-AUTOMATIC* STATIC*	55	I am not sure if this metric can be applied to embedded software
422 K-zero day safety	Instead of attempting to rank unknown vulnerabilities, the metric counts how many such vulnerabilities would be required for compromising network assets; a larger count implies more security because the likelihood of having more unknown vulnerabilities available, applicable, and exploitable all at the same time will be significantly lower	NETWORK	SEMI-AUTOMATIC* STATIC*	ON	Similar to 32. This has to do with probabilities. Can be applied, but I won't
423 Least Privilege	This metric states that each component and user should be granted the minimal privileges required to complete their tasks. This metric can be assessed from two perspectives: from the security controls perspective we can review users' granted privileges. From the architectural analysis perspective this can be assessed as how the system is broken down to minimal possible actions i.e. the number of components that can access critical data. The smaller the value, the more secure the system	SOFIWARE	SEMI-AUTOMATIC⁴ STATIC⁴	53	If the software is available, it can be applied. It can be also applied to embedded software
424 Mean Time To Failure					
425 Report Confidence (RC)					

Security resource indicator  Service Mechanism Strength Software project management experience  Trustworthiness  Trustworthiness  Use of (automated) tools Variable Security Vulnerability  Vulnerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average time from incident detection until incident has been reported  BrokenAccountCount	Trustworthiness of software means its worthy of being trusted to fulfill require- ments which may be needed for a particular software component, application, system, or network. It involves attributes of stability, data security, quality, pri-	ORGANISATION				
Service Mechanism Strength  Software project management experience  Trustworthiness  Use of (automated) tools  Variable Security Vulnerability  Vulnerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average inne from incident detection until incident has been reported  BrokenAccountCount	tware means its worthy of being trusted to fulfill require- reeded for a particular software component, application, involvos attributes of stability data security, quality, pri-	ORGANISATION				
Software project management experience  Trustworthiness  Use of (automated) tools Variable Security Vulnerability  Valuerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average inne from incident detection until incident has been reported  BrokenAccountCount	tware means its worthy of being trusted to fulfill require- reeded for a particular software component, application, involvos attibutes of stability data security, quality pri-	ORGANISATION				
Trustworthiness Use of (automated) tools Variable Security Vulnerability  Vulnerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average time from incident detection until incident has been reported  BrokenAccountCount	tware means its worthy of being trusted to fulfill require- needed for a particular software component, application, involvos atributes of sability data security, cutilty, pri-				NO	It has to do with the organisation and the development of the product. It is applied at the beginning of the life-cycle
Use of (automated) tools Variable Security Vulnerability  Vulnerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average time from incident detection until incident has been reported  BrokenAccountCount	yactive of the work. The covers attractive of statistics, was accurately years, party, years,	SOFIWARE				
Variable Security Vulnerability  Vulnerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average time from incident detection until incident has been reported  BrokenAccountCount						
Vulnerability Confidence Coefficient (VCC)  Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average time from incident detection until incident has been reported  BrokenAccountCount	The security vulnerability of a variable v@l is determined by the combined security damage it might cause to the overall system security when v@l is attacked. The more security properties that can be left infact, the less vulnerable the variable is; on the other hand, the more security properties are violated, the more criticall the variable is.	SOFTWARE	AUTOMATIC* STA	STATIC⁴	¿?	This is only applicable if the source code is available
Vulnerability Free Days (VFD)  Attack Graph Probability  Average Active Vulnerabilities per day (AAV)  Average time from incident detection until incident has been reported  BrokenAccountCount						
Attack Graph Probability Average Active Vulnerabilities per day (AAV) Average time from incident detection until incident has been reported BrokenAccountCount	Percent of days in which the vendor's queue of reported vulnerabilities for the product is empty	SOFIWARE	SEMI-AUTOMATIC* DY	DYNAMIC*	NO	This metric is not suitable for the evaluation phase
Average Active Vulnerabilities per day (AAV)  Average time from incident detection until incident has been reported  BrokenAccountCount						
Average time from incident detection until incident has been reported BrokenAccountCount	Median number of software vulnerabilities which are known to the vendor of a particular piece of software but for which patch has been publidy released by the vendor	SOFTWARE	SEMI-AUTOMATIC* DYI	DYNAMIC*	ON	This metric is not suitable for the evaluation phase
BrokenAccountCount		ORGANISATION	SEMI-AUTOMATIC* DYI	DYNAMIC*	;?	This a high level metric, but it might be used during a penetration testing if the device has countermeasures implemented
ing illegal access.	Number of accounts that have no activity for more than 90 days and will never - expire. Such accounts represent a clear risk of password compromise and resulting illegal access.	SOFTWARE	,		¿?	
440 Collateral Damage Potential (CDP)						
441 CVSS Base Score This group represents time, such as access concompromises the conficent to auther requirement for auther	This group represents the properties of a vulnerability that don't change over time, such as access complexity, access vector, degree to which the vulnerability compromises the confidentiality, integrity, and availability of the system, and requirement for authentication to the system	DEVICE				Similar to 15-26
42 Developer Risk						
43 Development Risk						
44 Environment Risk						
45 Excess Privilege						
446 Expected threat frequency						
47 Expected Vulnerability						
448 ExploitedFlawCount						
449 Faults found during manual inspection		SOFTWARE			YES	It is related to software and it can be applied to embedded software
450 InjectionFlawCount						
451 Integrity Impact (I)						
Maximal probability of successful attack	The probability to accomplish an attack successfully is a well-known metric. The metric describes the most probable way to compromise the system	NETWORK*			ON	Formal model. Formal verification. Too much notation. I think this is more related to networks
453 Mean Failure Cost Mean Failure Cost MFC within the system may lost data, resulting in: variable, and define Mean variable, and define Mean further, this quantity is	Mean Failure Cost MFC) used in the operational sense because the lack of security within the system may cause damage, in terms of lost productivity, lost business, lost data, resulting in security violations. We represent this loss by a random variable, and define MFC as the mean of this random variable. As discussed further, this quantity is not intrinsic to the system, but varies by stakeholder	ORGANISATION	,		ON	This metric is not related to embeded systems, but to organisations
454 Mean time from vendor patch availability to patch approval to patch installation		ORGANISATION	SEMI-AUTOMATIC* DYI	DYNAMIC*	NO	This a high level metric, and it is related with a network inside an organisation
455 Mean time to incident detection		ORGANISATION				
Monitoring system use	Number of unauthorized attempts to access file, folders, device attachment, attempts to change security settings	ORGANISATION	AUTOMATIC* DYI	DYNAMIC*	ON	This can be adapted to be used in embedded systems, but this metric makes no sense in the evaluation phase
457 Non-security failure reports						
Number of Design Flaws Related to Security (NSDF)	Security-related design flaws occur when software is planned and specified with- out proper consideration of security requirements and principles. For instance, clear-lext passwords are considered as design flaws. Design flaws can be detected using design inspection techniques (e.g., design reviews). Identifying the number of design flaws related to security can help detect security issues earlier in the design process	SOFTWARE	SEMI-AUTOMATIC* STA	STATIC*	¿?	If the design is available, it can be applied

459 Number of Exceptions That Have Been Implemented to Handle Execution Failures Related to Security (NEX)	to Number of exceptions that have been included in the code to handle possible - failures of the system due to an error in a code segment that has an impact on security	SOFIWARE	SEMI-AUTOMATIC* STATIC*	?5	If the design is available, it can be applied
460 Number of excluded security requirements that ensure input/output handling	re Is a specific encoding scheme defined for all inputs? Is a process of canonicalization applied to all inputs? Is an appropriate validation defined and applied to all inputs in terms of type, length, format/syntax and range? Is a whitelist Filtering approach is applied to all inputs? Are all the validations performed on the client and server side? Is all unsuccessful input handling rejected with an error message? Is all unsuccessful input handling logged? Is output data to the client filtered and encoded? Is output encoding performed on server side?	SOFTWARE	MANUAL STATIC*	ON	This metric is intended to be use in the design phase. This cannot be applied in a security assessment
461 Number of function calls that don't check return values					
462 Number of Implementation Errors Found in the System (NERR)	m Number of implementation errors of the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>?</i> ?	Related to the implementation phase
463 Number of Implementation Errors Related to Security (NSERR)	y Number of implementation errors of the system that have a direct impact on security. One of the most common security related implementation errors is the buffer overflow	SOFTWARE	SEMI-AUTOMATIC* STATIC*	;?	Related to the implementation phase
464 Number of Omitted Exceptions for Handling Execution Failures Related to Security (NOEX)	n Number of missing exceptions that have been omitted by software engineers - when implementing the system. These exceptions can easily be determined through testing techniques.	SOFTWARE	SEMI-AUTOMATIC* STATIC*	??	Related to the implementation phase
465 Number of Omitted Security Requirements (NOSR)	Number of requirements that should have been considered when building the application	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	If the design is available, it can be applied
466 Number of open security bugs					
467 Number of reported security incidents		ORGANISATION	SEMI-AUTOMATIC* DYNAMIC*	ON *C	This is high level metric
468 Number of security bulletins issues per year		DEVICE			
469 Number of service accounts with weak or default passwords	-53	SOFTWARE			
470 Number of successful attempts to execute recovery this period	sī	ORGANISATION	SEMI-AUTOMATIC* DYNAMIC	ON *	This is a high level metric, and it is related with attacks on a organisation's network
471 Number of violations of the LP principle	Number of violations of the Least Privilege principle. A component does not adhere to LP if it, based upon the permissions attributed as described before, is capable of executing tasks it is not responsible for	SOFTWARE	AUTOCATIC* STATIC*	YES	This could be used in embedded systems
472 OverflowVulnCount					
473 Percentage Software Hazards (PSH)	Comparing the number of software hazards identified against historical data, it indicates the sufficiency of the software hazard identification based on the identified hazards. (# software hazards / # System hazards ) * 100	SOFTWARE	SEMI-AUTOMATIC* STATIC*	SÍ	Hazard is similar to risk.
474 Project Management Risk					
475 Remediation Level (RL)				ON	Similar to 24
476 Remediation Potency (RP)	Remediation potency refers to how good the remediation is against the vulnerability			NO	
477 Remediation Scheme (RS)	It is used to describe the remediation process			ON	
478 Requirements Risk					
479 Security of system documentation	Total number of unauthorized accesses to system documentation / total number of access to system documentation $^\star100$	ORGANISATION	SEMI-AUTOMATIC* DYNAMIC	ON *D	This might be useful to measure anything related to social engineering and information gathering and OSINT
480 Static analysis alert count (number of)		SOFTWARE			
481 Temporal Score	Temporal metrics represent the time dependent features of the vulnerabilitie	SOFTWARE	SEMI-AUTOMATIC* DYNAMIC	C* YES	Similar to 23, 24 and 25. This is the final result that is obtained with those partial metrics. CVSS
482 Time to close bug/vulnerability					
483 Time to Problem Correction					
484   Time to Problem Report					
485 User Risk		USER			
486 Vulnerabilities found during requirements, design and coding	id	DEVICE			
487 Vulnerabilities found post-development		DEVICE			
488 Vulnerability Density	Number of vulnerabilities in the unit size of a code. VD = size of the software / $\#$ -vulnerabilities in the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase
490 XsiteVulnCount	It is obtained via a penetration-testing tools	SOFTWARE		ON	It is not clearly defined
491 Operational Environment Security Measurement					
492 Program Implementation Security Measurement					
493 Security Indictor of Software Systems (sic)					
494 Security metrics of Arithmetic Expression	The security flaws of arithmetic expression include divide by zero, overflow of arithmetic operation, misused data type of arithmetic expression, and truncation error of arithmetic operation	SOFTWARE	AUTOMATIC* STATIC*	?خ	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
495 Security Metrics of Array Index	The security flaws of array index include index out of range, incorrect index variable and uncontrollable array index	SOFTWARE	AUTOMATIC* STATIC*	<i>ز؟</i>	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist

496 Security Metrics of Component Interfaces	The security flaws of component interface include inconsistent parameter num-	SOFTWARE	AUTOMATIC*	STATIC*	3.5	This is not a metric, but a group of them. Individually, they could be used if they
	bers, inconsistent data types, and inconsistent size of data types				,	are specified. Part of a checklist
497 Security Metrics of Control Operation	The security flaws of control operations include infinite loop, deadlock situations and boundary defects	SOFIWARE	AUTOMATIC*	STATIC*	53	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
498 Security Metrics of I/O Management	File is an important media which can store critical information and record log data. However, file privilege definition, file access management, file contents, file organization and file size are major factors to affect the security software operation	SOFTWARE	AUTOMATIC*	STATIC*	:?	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
499 Security Metrics of Input Format	The security flaws of data format include mismatched data contents, mismatched adata type, mismatch data volume	SOFTWARE	AUTOMATIC*	STATIC*	??	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
500 Security Metrics of Kernel Operation	A released software system must be adapted to a suitable operating system and environment. The kernel of operating system and environment become a critical factor to affect the operating security of software system. Resource management, execution file protection, main memory control and process scheduling are major tasks of operating system. The security flaws, which embedded in the tasks of operating system, become high security risk for software system operation	SOFTWARE	А∪ТОМАПС	STATIC*	¿?	This is not a metric, but a group of them. Individually, they could be used if they are specified. Checklist. Only when there is an OS
501 Security Metrics of Network Environment	Network environment is a critical tool in the e-commerce years. However, data transmission management, remote access control, and transmission contents always become the security holes of software system operation	SOFTWARE	AUTOMATIC⁴	STATIC*	;?	This is not a metric, but a group of them. Individually, they could be used if they are specified
502 Security Metrics of Resources Allocation	The security flaws of resource allocation include exhausted memory, unreleased - memory exhausted resources, and unreleased resources	SOFTWARE	AUTOMATIC*	STATIC*	;?	This is not a metric, but a group of them. Individually, they could be used if they are specified
503 Security Metrics of User Authority	Many users may use the released system to do specific jobs. However, user privilege definition, user password management and user detailed information maintenance are important tasks to control and manage the user authority. Without user authority control and management, the operation environment of software system will be on high security risk	SOFTWARE	А∪ТОМАПС	STATIC*	:?	This is not a metric, but a group of them. Individually, they could be used if they are specified
504 Instruction Count						
505 Attack-reward (URL Jumping)						
506 Classified Accessor Attribute Interactions (CAAI)	The ratio of the sum of all interactions between accessors and classified attributes to the possible maximum number of interactions between accessors and classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 187. Object-oriented programs
507 Classified Attributes Inheritance (CAI)	The ratio of the number of classified attributes which can be inherited in a hierarchy to the total number of classified attributes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC* STATIC*	STATIC*	ON	Similar to 284. Object-oriented programs
508 Classified Attributes Interaction Weight (CAIW)	The ratio of the number of all interactions with classified attributes to the total number of all interactions with all attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 188. Object-oriented programs
509 Classified Class Data Accessibility (CCDA)	The ratio of the number of non-private classified class attributes to the number of classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 189. Object-oriented programs
510 Classified Instance Data Accessibility (CIDA)	The ratio of the number of non-private classified instance attributes to the number of classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	NO	Similar to 190. Object-oriented programs
511 Classified Method Extensibility (CME)	The ratio of the number of the non-finalised classified methods in program to the total number of classified methods in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 191. Object-oriented programs
512 Classified Methods Inheritance (CMI)	The ratio of the number of classified methods which can be inherited in a hierarchy to the total number of classified methods in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 258. Object-oriented programs
513 Classified Methods Weight (CMW)	The ratio of the number of classified methods to the total number of methods in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	NO	Similar to 192. Object-oriented programs
514 Classified Mutator Attribute Interactions (CMAI)	The ratio of the sum of all interactions between mutators and classified attributes - to the possible maximum number of interactions between mutators and classified attributes in the program.	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 193. Object-oriented programs
515 Classified Operation Accessibility (COA)	The ratio of the number of non-private classified methods to the number of $$ -classified methods in the program	SOFIWARE	SEMI-AUTOMATIC*	STATIC*	NO	Similar to 194. Object-oriented programs
516 CNBetweenness						
517 Composite-Part Critical Classes (CPCC)	The ratio of the number of critical composed-part classes to the total number of critical classes in the program	SOFTWARE				Similar to 196
518 Confidentiality						Similar to 413
519 Critical Class Extensibility (CCE)	The ratio of the number of the non-finalised critical classes in program to the total number of critical classes in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC	STATIC*	ON	Similar to 201. Object-oriented programs
520 Critical Design Proportion (CDP)	The ratio of the number of critical classes to the total number of classes in the program	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 202. Object-oriented programs
521 Critical Superclasses Inheritance (CSI)	The ratio of the sum of classes which may inherit from each critical superclass - to the number of possible inheritances from all critical classes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC*	· STATIC*	ON	Similar to 205. Object-oriented programs
522 Critical Superclasses Proportion (CSP)	The ratio of the number of critical superclasses to the total number of critical $$ -classes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC*	STATIC*	ON	Similar to 206. Object-oriented programs
523 Design Size		SOFTWARE HARDWARE				
524 DNMaxEdgeBetweenness						
525 Knot Count		HIG THE AND A				
526 Kolmogorov Complexity		SOFTWARE				

527 Measurement of Cost		
528 Number of Developers	ORGANISATION	NO It is no related directly with the device
529 Potency		
530 Resilience		
531 Shortest Path		

#	METRIC	DEFINITION SCALE	SCOPE AUTO	TOMATION MEASUREMENT	USABLE	REASON
r.	Епtropy		USER AUTO	AUTOMATIC STATIC	YES	This could be used with keys instead of passwords to measure randomness
<u>«</u>	Attack surface	It aim to measure the ways by which an attacker can compromise a targeted software ABSOLUTE	DEVICE	STATIC	YES	
6	Historical vulnerability metric	It measures the degree that a system is vulnerable (i.e., frequency of vulnerabilities) in RATIO the past DISTRIBUTION	DEVICE	STATIC	YES	This could be part of a checklist of task, to check if there is any vulnerability that affects the device
10	Historically exploited vulnerability metric	It measures the number of vulnerabilities exploited in the past ABSOLUTE ABSOLUTE DISTRIBUTION	DEVICE	STATIC	YES	This could be part of a checklist of task, to check if there is any vulnerability that affects the device
13	Tendency-to-be-exploited metric	It measures the tendency that a vulnerability may be exploited, where the "tendency" RATIO may be derived from information sources such as posts on Twitter before vulnerability ABSOLUTE disclosures	DEVICE	STATIC	?خ	This research work might not be part of a security evaluation, but could be applied if so
15	Attack vector	Describes whether a vulnerability can be exploited remotely, adjacently, locally, or NOMINAL physically (i.e., attacker must have physical access to the computer)	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
16	Attack complexity	Describes the conditions that must hold before an attacker can exploit the vulnerability, ORDINA such as low or high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
17	Privilege required	Describes the level of privileges that an attacker must have in order to exploit a vulnera- ORDINAL bility, such as none, low, or high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
18	User interaction	Describes whether the exploitation of a vulnerability requires the participation of a user NOMINAL (other than the attacker), such as none or required	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
19	Authorization scope	Describes whether or not a vulnerability has an impact on resources beyond its privileges NOMINAL (e.g., sandbox or virtual machine), such as unchanged or changed	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
20	Confidentiality impact	The impact of a successful exploitation of a vulnerability in terms of confidentiality such ORDINAL as none, low, high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
21	Integrity impact	The impact of a successful exploitation of a vulnerability in terms of integrity such as ORDINAL none, low, high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
22	Availability impact	The impact of a successful exploitation of a vulnerability in terms of availability, such as ORDINAL none, low, high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
23	Exploit code maturity	The likelihood a vulnerability can be attacked based on the current exploitation tech- ORDINAL niques, such as undefined, unproven, proof-of-concept, functional, or high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
24	Remediation level	Describes whether or not a remediation method is available for a given vulnerability, NOMINAL such as undefined, unavailable, workaround, temporary fix, or official fix	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
25	Report confidence	Measures the level of confidence for a given vulnerability as well as known technical NOMINAL details, such as unknown, reasonable, or confirmed	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
26	Security requirement	Describes environment-dependent security requirements in terms of confidentiality, NOMINAL integrity, and availability, such as not defined, low, medium, or high	SOFTWARE	STATIC	YES	For calculating new vulnerabilities, or to know each dimension of known vulnerabilities
29	Necessary defense metric	It estimates a minimal set of defense countermeasures necessary for thwarting a certain - attack	DEVICE	STATIC	YES	This is related with the defenses that are necesary to implement so the system is secure for a certain vulnerability
30	Effort-to-security-failure metric	It measures an attacker's effort to reach its goal state	DEVICE	DYNAMIC*	YES	Although it might be difficult to estimate, this could be also included in the evaluation process
31	Weakest adversary metric	It estimates minimum adversary capabilities required to achieve an attack goal	DEVICE	DYNAMIC*	53	If the effort can be calculated in an objective manner, this could be used in embedded systems
33	Likelihood of exploitation	It measures the probability that an exploit will be executed by an attacker with certain capabilities	SOFTWARE	STATIC	:?	This is related with statistics and risk assessment. It could be applied
34	Effective base metric / scope	Given an attack graph and a subset of exploits, the effective base metric aims to adjust the CVSS base metric by taking the highest value of the ancestors of an exploit to reflect the worst-case scenario, while the effective base score metric is calculated based on the effective base metric	SOFTWARE	STATIC		
35	Reaction time metric	It captures the delay between the observation of the malicious entity at time $t$ and the ABSOLUTE blacklisting of the malicious entity at time $t$ (i.e., $t$ - $t$ )	DEVICE	DYNAMIC	YES	If the device has any kind of detection mechanism, this could be used. Detection and reaction might have different times
36	Coverage metric	It estimates the portion of blacklisted malicious players	DEVICE	STATIC	YES	This could be used, but the device has to have a black-list
38	Average indirect target reduction	It counts the overall reduction of targets for any indirect control-flow transfer. It measures ABSOLUTE the overall reduction in terms of the number of targets exploitable by the attacker where smaller targets are more secure	UNKNOWN	UNKNOWN	UNKNOWN	n unknown
39	Average target size	Ratio between the size of the largest target and the number of targets	UNKNOWN	UNKNOWN	UNKNOWN	N UNKNOWN
40	Evasion resistance	It is measured against control flow bending attacks, reflecting the effort (or premises) that an attacker must make (or satisfy) for evading the CFI scheme	UNKNOWN	UNKNOWN	UNKNOW	UNKNOWN UNKNOWN
41	Coverage metric	It measures the fraction of events detectable by a specific sensor deployment, reflecting - a defender's need in monitoring events	DEVICE	DYNAMIC	UNKNOWN	N UNKNOWN
45	Redundancy metric	It estimates the amount of evidence provided by a specific sensor deployment to detect an event	DEVICE	DYNAMIC	YES	This can be applied to devices as a whole, to check any mechanism of intrusion detection (mostly, lantitampering and software)

43 Confidence metric	It measures how well-deployed sensors detect an event in the presence of compromised - sensors	DEVICE		DYNAMIC	YES	This can be applied to devices as a whole, to check any mechanism of intrusion detection (mostly, tanti-tampering and software)
45 Detection time	For instrument-based attack detection, this metric is used to measure the delay between the time to at which a compromised computer sends its first scan packet and the time that a scan packet is observed by the instrument	UNKNOWN		UNKNOWN	UNKNOWN	UNKNOWN UNKNOWN
57 ASLR-induced Effective entropy metric	Measure of entropy in user space memory which quantitatively considers an adversary's ability to leverage low entropy regions of memory via absolute and dynamic intersection connections	SOFTWARE		DYNAMIC	YES	If the device runs an OS, this could be applied
59 Penetration resistance (PR)	Running a penetration test to estimate the level of effort (e.g., person-day or cost) required for a red team to penetrate into a system	DEVICE NETWORK		DYNAMIC	YES	Penetration testing is a really common way to test this kind of device
82 Probability a computer is compromised at time t	It quantifies the degree of security in enterprise systems by using modeling techniques - based on a holistic perspective	DEVICE NETWORK		;?	;?	Although this is a network-oriented metric, this might be adapted to be used with an embedded system
85 Blocking rate	The rate an encountered attack is successfully defended by a deployed defense	NETWORK		DYNAMIC	?;	Although this is a network-oriented metric, this might be adapted to be used with an embedded system (maybe for a penetration test)
89 Time-to-first-compromise metric	Estimates the duration of time between when a computer starts to run and the first malware alarm is triggered on the computer where the alarm indicates detection rather than infection	DEVICE		DYNAMIC	YES	This can be used in embedded systems, but it might not be suitable for a security evaluation (maybe for a penetration test)
96 Attack Impact	It is the quantitative measure of the potential harm caused by an attacker to exploit a vulnerability $% \left( \frac{1}{2}\right) =\frac{1}{2}\left( \frac{1}{$	NETWORK		STATIC		
97 Attack Cost	It is the cost spent by an attacker to successfully exploit a vulnerability (i.e., security weakness) on a host	NETWORK		STATIC DYNAMIC	YES	This is related to network, but this could be applicable to devices too
100 Mean-time-to-compromise (MTTC)	It is used to measure how quickly a network can be penetrated	NETWORK		DYNAMIC	;?	This is no directly applicable, but its equivalent would be a penetration attack in an embedded system
101 Mean-time-to-recovery (MTTR)	It is used to assess the effectiveness of a network to recovery from an attack incidents. It is defined as the average amount of time required to restore a system out of attack state	NETWORK		DYNAMIC	¿?	Although this is developed for networks, the same concept can be translated into embeded systems
105 Return on Attack	The gain the attacker expects from successful attack over the losses he sustains due to the countermeasure deployed by his target	NETWORK		:?	;?	Network-related, but it scent can be translated into embedded systems
106 Common Vulnerability Scoring System (CVSS) - Overall value	The overall score is determined by generating a base score and modifying it through the temporal and environmental formulas	DEVICE		STATIC	YES	This could be used tohether with other metics to obtain a numeric value
107 Probability of Vulnerability Exploited	It is used to assess the likelihood of an attacker exploiting a specific vulnerability on a host	DEVICE NETWORK		STATIC	;?	Network-related metric. This is also related to risk assessment. Can be traslated into embedded systems
108 Probability of attack detection	It is used to assess the likelihood of a countermeasure to successfully identify the event of an attack on a target	DEVICE NETWORK		DYNAMIC	¿?	Network-related metric. This is also related to risk assessment. Can be traslated into embedded systems
109 Probability of host compromised	It is used to assess the likelihood of an attacker to successfully compromise a target	DEVICE NETWORK		STATIC DYNAMIC	۶۶	Network-related metric. This is also related to risk assessment. Can be traslated into embedded systems
137 Number of bits leaked (mean privacy)		NETWORK	SEMI-AUTOMATIC	STATIC	;?	If adapted, this metric can be used to test memory related vulnerabilities
139 rav	It is a scale measurement of the attack surface, the amount of uncontrolled interactions with a target, which is calculated by the quantitative balance between operations, limitations, and controls	DEVICE				
141 Churn	Count of Source Lines of Code that has been added or changed in a component since the previous revision of the software	SOFTWARE			,	
148 Relative churn	Normalized by parameters such as lines of code, files counts	SOFTWARE	AUTOMATIC*	STATIC*		
149 Repeat frequency	The number of consecutive edits that are performed on a binary	SOFTWARE	AUTOMATIC*	STATIC*		
150 Total churn	The total added, modified, and deleted lines of code of a binary during the development	SOFTWARE	AUTOMATIC*	STATIC*		
155 Cyclomatic number	$M=E\cdot N+P, where \ E=the \ number \ of \ edges \ of \ the \ graph. \ N=the \ number \ of \ nodes \ of \ the \ graph. \ P=the \ number \ of \ connected \ components.$	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
158 Fan in (FI)	Number of inputs a function uses. Inputs include parameters and global variables that are used (read) in the function / Number of functions calling a function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
159 Fan out (FO)	Number of outputs that are set. The outputs can be parameters or global variables (modified) / Number of functions called by a function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
162 Max fan in	The largest number of inputs to a function such as parameters and global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
163 Max fan out	The largest number of assignment to the parameters to call a function or global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
164 MaxMaxNesting	The maximum of MaxNesting in a file	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
165 MaxNesting	Maximum nesting level of control constructs such as if or while statements in a function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
166 Nesting complexity	Maximum nesting level of control constructs (if, while, for, switch, etc.) in the function	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
167 Number of children (NOC)	Number of immediate sub-classes of a class or the count of derived classes. If class CA inherits class CB, then CB is the base class and CA is the derived class. In other words, CA is the children of class CB, and CB is the parent of class CB.	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
168 SumCyclomaticStrict	The sum of the strict cyclomatic complexity, where strict cyclomatic complexity is defined as the number of conditional statements in a function $% \left( \frac{1}{2}\right) =\frac{1}{2}\left( \frac{1}$	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software

169 SumEssential	The sum of essential complexity, where essential complexity is defined as the number of branches after reducing all the programming primitives such as a for loop in a function's control flow graph into a node iteratively until the graph cannot be reduced any further. Completely well-structured code has essential complexity 1	SOFTWARE	AUTOMATIC⁴	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
170 SumFanIn	The sum of FanIn, where FanIn is defined as the number of inputs to a function such as parameters and global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
171 SumFanOut	The sum of FanOut, where FanOut is defined as the number of assignment to the parameters to call a function or global variables	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
172 SumMaxNesting	The sum of the MaxNesting in a file, where MaxNesting is defined as the maximum nesting level of control constructs such as if or while statements in a function	SOFTWARE	AUTOMATIC*	ЅТАПС*	YES	If the software is available, it can be applied. It can be also applied to embedded software
173 McCabe Cyclomatic Complexity	$M=E\cdot N+2P$ where $E=$ the number of edges of the graph. $N=$ the number of nodes $$ of the graph. $P=$ the number of connected components	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
174 Weighted methods per class (WMC)	Number of local methods defined in the class	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
187 Classified Accessor Attribute Interactions (CAAI)	The ratio of the sum of all interactions between accessors and classified attributes to the possible maximum number of interactions between accessors and classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	:3	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
188 Classified Attributes Interaction Weight (CAIW)	The ratio of the number of all interactions with classified attributes to the total number of all interactions with all attributes in the program	SOFTWARE	SEMI-AUTOMATIC	* STATIC*	:?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
189 Classified Class Data Accessibility (CCDA)	The ratio of the number of non-private classified class attributes to the number of - classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	53	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
190 Classified Instance Data Accessibility (CIDA)	The ratio of the number of non-private classified instance attributes to the number of - classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC	* STATIC*	5?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
191 Classified Method Extensibility (CME)	The ratio of the number of the non-finalised classified methods in program to the total number of classified methods in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	:?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
192 Classified Methods Weight (CMW)	The ratio of the number of classified methods to the total number of methods in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	53	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
193 Classified Mutator Attribute Interactions (CMAI)	The ratio of the sum of all interactions between mutators and classified attributes to the possible maximum number of interactions between mutators and classified attributes in the program.	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	۶۶:	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
194 Classified Operation Accessibility (COA)	The ratio of the number of non-private classified methods to the number of classified - methods in the program	SOFTWARE	SEMI-AUTOMATIC	* STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
195 Classified Writing Methods Proportion (CWMP)	The ratio of the number of methods which write classified attributes to the total number of classified methods in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	:?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
196 Composite-Part Critical Classes (CPCC)	The ratio of the number of critical composed-part classes to the total number of critical classes in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	?ئ	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
Classes Design Proportion (CDP)	Ratio of the number of critical classes to the total number of classes, from the group of Design Size Metrics.	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	5?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
201 Critical Class Extensibility (CCE)	The ratio of the number of the non-finalised critical classes in program to the total number of critical classes in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
202 Critical Design Proportion (CDP)	The ratio of the number of critical classes to the total number of classes in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	;?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
203 Critical Element Ratio	A process may not require all aspects of an object to be instantiated. Critical Elements - Ratio = (Critical Data Elements in an Object) $/$ (Total number of elements in the object)	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
204 Critical Senalized Classes Proportion (CSCP)	The ratio of the number of critical serialized classes to the total number of critical classes $$ in the program	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	55	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
205 Critical Superclasses Inheritance (CSI)	The ratio of the sum of classes which may inherit from each critical superclass to the number of possible inheritances from all critical classes in the program's inheritance hierarchy	SOFTWARE			:?	
206 Critical Superclasses Proportion (CSP)	The ratio of the number of critical superclasses to the total number of critical classes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC*	* STATIC*	55	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
212 Isolation (PI)	This assesses the level of security isolation between system components. This means getting privileges to a component does not imply accessibility of other co-located components. This netric can be assessed using system architecture and deployment models. Components marked as confidential should not be hosted with nonconfidential (public) components. Methods that are not marked as confidential should not have access to confidential attributes or methods	SOFTWARE	SEMI-AUTOMATIC* STATIC	* STATIC*	?5	I am not sure if this metric can be applied to embedded software
214 Number of Catch Blocks Per Class (NCBC)	It measures the exception handing factor (EHF) in some way. It is defined as the percentage of catch blocks in a class over the theoretical maximum number of catch blocks in the class	SOFTWARE	AUTOMATIC*	STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software

ments	ments				
218 Percentage of Software Safety Requirements (PSSR)	How sufficient hazard analysis has been performed. (# Software Safety Requirements / - # Software Requirements) * 100	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	This is related to software and the design
219 Percentage of Software Safety Requirements Traceable to Hazards	Ensuring traceability to system hazards or System of Systems hazards increases the validation case. Percentage indicator of traceability of requirements. All derived software safety requirements must be traceable to system hazards or System of Systems hazards (# Traceable Software Safety Requirements / # Software Safety Requirements) * 100	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	This is related to software and the design
246 PercentValidatedInput	To compute this metric, let T equal the count of the amount of input forms or interfaces the application exposes (the number of HTML form POSTs, GETs, and so on) and let V equal the number of these interfaces that use input validation mechanisms. The ratio V/T makes a strong statement about the Web application's vulnerability to exploits from invalid input	SOFTWARE		YES	This could be adapted to other pieces of code, rather than just HTML
247 Ratio of Design Decisions (RDD)	Ratio of the number of design decisions related to security to the total number of design - decisions of the entire system. The objective is to assess the portion of design dedicated to security	SOFTWARE	SEMI-AUTOMATIC* STATIC*	?خ	This is related to design phase
248 Ratio of Design Flaws Related to Security (RDF)	Ratio of design flaws related to security to the total number of design flaws applicable to the whole system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	This is related to design phase
250 Ratio of Implementation Errors That Have an Impact on Security (RSERR)	Ratio of the number of errors related to security to the total number of errors in the implementation of the system (i.e. NERR)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	:?	This is related to design phase
253 Ratio of Patches Issued to Address Security Vulnerabilities (RP)	Ratio of the number of patches that are issued to address security vulnerabilities to the otal number of patches of the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	If the design is available, it can be applied
254 Ratio of Security Requirements (RSR)	Ratio of the number of requirements that have direct impact on security to the total number of requirements of the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	?خ	If the design is available, it can be applied
255 Ratio of Security Test Cases (RTC)	Number of test cases designed to detect security issues to the total number of test cases of the entire system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	۶۶	If the design is available, it can be applied
256 Ratio of Security Test Cases that Fail (RTCP)	Number of test cases that detect implementation errors (i.e. the ones that fail) to the total number of test cases, designed specifically to target security issues	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	If the design is available, it can be applied
258 Ratio of Software Changes Due to Security Considerations (RSC)	Number of changes in the system triggered by a new set of security requirements Software changes due to security considerations include patches that are released after a system is delivered, or any other security updates	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	If the design is available, it can be applied
261 Ratio of the Number of Omitted Exceptions (ROEX)	Ratio of the number of omitted exceptions (i.e. Noex) to the total number of exceptions - that are related to security	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>ز</i> ؟	If the software is available, it can be applied. It can be also applied to embedded software
262 Ratio of the Number of Omitted Security Requirements (ROSR)	Ratio of the number of security requirements that have not been considered during the analysis phase (i.e. NOSR) to the total number of security requirements identified during the analysis phase (NSR)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	If the design is available, it can be applied
269 Software Hazard Analysis Depth	Hazardous software, or safety-critical software, allocated as high- or medium-risk, according to the Software Hazard Criticality Matrix, requires analysis of second- and third-order causal factors (should they exist)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	Hazard is similar to risk.
273 Unaccessed Assigned Classified Attribute (UACA)	The ratio of the number of classified attributes that are assigned but never used to the otolal number of classified attributes in the program	SOFTWARE			Object-oriented programs
274 Uncalled Classified Accessor Method (UCAM)	The ratio of the number of classified methods that access a classified attribute but are never called by other methods to the total number of classified methods in the program	SOFTWARE			Object-oriented programs
275 Unused Critical Accessor Class (UCAC)	The ratio of the number of classes which contain classified methods that access classified attributes but are never used by other classes to the total number of critical classes in the program	SOFTWARE			Object-oriented programs
284 Classified Attributes Inheritance (CAI)	The ratio of the number of classified attributes which can be inherited in a hierarchy to the total number of classified attributes in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>ξ?</i>	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
285 Classified Methods Inheritance (CMI)	The ratio of the number of classified methods which can be inherited in a hierarchy to the total number of classified methods in the program's inheritance hierarchy	SOFTWARE	SEMI-AUTOMATIC* STATIC*	5?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
286 Coupling	Measures the following three coupling dimensions between modules: referential dependency, structural dependency, and data integrity dependency	SOFTWARE			
288 Coupling Between Object classes (CBOC)	Number of other classes coupled to a class C	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
290 Coupling Corruption Propagation	Coupling corruption propagation is meant to measure the total number of methods that could be affected by erroneous originating method. Coupling Corruption Propagation = Number of child methods invoked with the parameter(s) based on the parameter(s) of the original invocation	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
291 Critical Classes Coupling (CCC)	The ratio of the number of all classes' links with classified attributes to the total number of possible links with classified attributes in the program	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	Object-oriented programs. If the software is available, it can be applied. It can be also applied to embedded software
292 Depth of Inheritance Tree (DIT)	Maximum depth of the class in the inheritance tree	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
301 Lack of cohesion of methods (LCOM)	The LCOM value for a class C is defined as $LCOM(C) = \frac{(1- E(C) )}{ V(C) M(C) } \cdot 100\%$ , where · V(C) is the set of instance variables, M(C) is the set of instance methods, and E(C) is the set of pairs (y,m) for each instance variable v in V(C) that is used by method m in M(C)	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
304 NumCalls	The number of calls to the functions defined in an entity	SOFTWARE		YES	If the software is available, it can be applied. It can be also applied to embedded software

312 Reflection Package Boolean (RPB)	A boolean value representing whether the Java program imports the reflection package $$ - $$ (1) or not (0)	SOFTWARE			Object-oriented programs
Vulnerability Propagation (VP)	Vulnerability Propagation (VP) of a class C, denoted by VP(C), is the set containing classes in the bierarchy which directly or indirectly inherit class C. Cardinality of set VP(C) represents the number of classes which have become vulnerable due to class C. Alternatively, Vulnerability Propagation of a class C is the total number of classes which directly or indirectly inherit class C.	SOFTWARE	AUTOMATIC* STATIC*	YES	This metric can be applied to embedded software
315 Vulnerability Propagation Factor (VPF)	An Inheritance hierarchy may contain no class or one or more vulnerable classes. Also, there are more than one Inheritance hierarchies present in a design. So, calculation of Vulnerability Propagation Factor (VPF) of a design requires calculation of Vulnerability Propagation due to each vulnerable class in Inheritance hierarchies present in the design. Then, union of Vulnerability Propagation due to each vulnerable class will give overall Vulnerability Propagation due to each vulnerable class will give overall	SOFTWARE	AUTOMATIC* STATIC*	YES	This metric can be applied to embedded software
316 Access Complexity (AC)	This metric measures the complexity of attacks exploiting the vulnerability. A vulnerability with low complexity can be for example a buffer overflow in a web server, the vulnerability can be exploited at will	DEVICE			
317 Access Vector (AV)	This metric indicates from where an attacker can exploit the vulnerability	DEVICE			
320 Authentication (AU)	This metric counts how often an attacker must authenticate before the vulnerability can be exploited	SOFTWARE			
322 Damage potential-effort ratio	It indicates the level of damage an attacker can potentially cause to the system and the effort required for the attacker to cause such damage	DEVICE	SEMI-AUTOMATIC* DYNAMIC*	YES	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase
332 Protection Rate (PR)	It expresses the efficiency of security mechanisms. Percentage of the known vulnerabilities protected by a given security mechanism. It is based on the Attack Surface metric. PR = (1 - AttackSurfaceEvaluatedSystem / AttackSurfaceReferenceSystem) * 100	SOFTWARE	SEMI-AUTOMATIC* DYNAMIC*	53	This metric was developed specificaly for OSGi platform, but I think it can be adapted to be applied to embedded systems
334 Side-channel Vulnerability Factor (SVF)	SVF quantifies patterns in attackers' observations and measures their correlation to the victim's actual execution patterns and in doing so captures systems' vulnerability to side-channel attacks. we can measure information leakage by computing the correlation between ground-truth patterns and attacker observed patterns. We call this correlation Side-channel Vulnerability Factor (SVF). SVF measures the signal-to-noise ratio in an attacker's observations. While any amount of leakage could compromise a system, a low signal-to-noise ratio means that the attacker must either make do with inaccurate results (and thus make many observations to create an accurate result) or become much more intelligent about recovering the original signal. We use phase detection techniques to find patterns in both sates of data then compute the correlation between actual patterns in the victim and observed patterns.	HARDWARE	SEMI-AUTOMATIC* DYNAMIC*	YES	Embedded devices can be vulnerable to side-channel attacks and they can be tested
336 Structural severity	Measure that uses software attributes to evaluate the risk of an attacker reaching a vulnerability location from attack surface entry points. It is measured based on three values: high (reachable from an entry point), medium (reachable from an entry point with no dangerous system calls), low (not reachable from any entry points)	SOFTWARE	SEMI-AUTOMATIC* DYNAMIC*	5?	I am not sure if this metric can be applied to embedded software
Vulnerability Index (VI)	Probability of a component's vulnerability being exposed in a single execution	DEVICE	MANUAL* DYNAMIC*	<i>ڊ</i> ج	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase
Effective Vulnerability Index (EVI)	Relative measure of the impact of a component on the system's insecurity	DEVICE	SEMI-AUTOMATIC* DYNAMIC*	55	It seems that it can be applied to embedded systems, but I am not sure if it is applicable in the evaluation phase
361 Classified Attributes Total (CAT)	The total number of classified attributes in the program	SOFTWARE	1	1	Object-oriented programs
362 Classified Methods Total (CMT)	The total number of classified methods in the program	SOFTWARE			Object-oriented programs
366 Count of Base Classes (CBC)	Number of base classes	SOFTWARE			
		SOFTWARE		· .	Object-oriented programs  If the decime is available it can be awalted
3/8 Number of Design Decisions Kelated to Secu- rity (NDD)	In proposed metric arms at assessing the number of design decisions that address the security requirements of the system. During the design phase, it is common to end up with multiple solutions to the same problem. Software engineers make many design decisions in order to choose among alternative design solutions.	SOFIWAKE		<i>}?</i>	If the design is available, it can be applied
383 Number of global variables *		SOFTWARE		YES	If the software is available, it can be applied. It can be also applied to embedded software
390 Number of return points in the method		SOFTWARE		YES	If the software is available, it can be applied. It can be also applied to embedded software
391 Number of Security Algorithms (NSA)	Number of security algorithms that are supported by the application	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>؟</i> غ	If the design is available, it can be applied, but this metric can be used in finals stages of the development life-cycle
393 Number of Security Requirements (NSR) *	Number of security requirements identified during the analysis phase of the application	SOFTWARE	MANUAL STATIC*	<i>ڊ</i> ?	If the design is available, it can be applied
394 Number of sub classes		SOFTWARE		YES	If the software is available, it can be applied. It can be also applied to embedded software
395 NumFunctions		SOFTWARE		YES	If the software is available, it can be applied. It can be also applied to embedded software
400 Response set for a Class (RFC)	Set of methods that can potentially be executed in response to a message received by an object of that class. RFC is simply the number of methods in the set, including inherited methods	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software
402 Stall Ratio	Stall ratio is a measure of how much a program's progress is impeded by frivolous activities. stall ratio = (lines of non-progressive statements in a loop) $/$ (total lines in the loop)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be also applied to embedded software

409 Comment ratio	Ratio of number of comment lines to number of code lines	SOFTWARE	AUTOMATIC* STATIC*	YES	If the software is available, it can be applied. It can be
					also applied to embedded software
410 CommentDensity	The ratio of lines of comments to lines of code	SOFTWARE			
411 Compartmentalization	Compartmentalization means that systems and their components run in different compartments, isolated from each other. Thus a compromise of any of them does not impact the others. This metric can be measured as the number of independent components that do not trust each other (performs authentication and authorization for requests/calls coming from other system components) that the system is based on to deliver its function. The higher the compartmentalization value, the more secure the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>??</i>	If the software is available, it can be applied. It can be also applied to embedded software
415 Defense-In-Depth	This metric verifies that security controls are used at different points in the system chain including network security, host security, and application security. Components that have critical data should employ security controls in the network, host, and component layers. To assess this metric we need to capture system architecture and deployment models as well as the security architecture model. Then we can calculate the ratio of components with critical data that apply the layered security principle compared to number of critical components	SOFTWARE	SEMI-AUTOMATIC* STATIC*	¿?	I am not sure if this metric can be applied to embedded software, or at verification phases
417 Fail Securely	The system does not disclose any data that should not be disclosed ordinarily at system failure. This includes system data as well as data about the system in case of exceptions. This metric can be evaluated from the security control responses – i.e. how the control behaves in case it failed to operate. From the system architecture perspective, we can assess it as the number of critical attributes and methods that can be accessed in a given component. The smaller the metric value, the likely more secure the system in case of failure	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	This could be used in evaluation
421 Isolation (PI)	This assesses the level of security isolation between system components. This means getting privileges to a component does not imply accessibility of other co-located components. This metric can be assessed using system architecture and deployment models. Components marked as confidential should not be hosted with nonconfidential (public) components. Methods that are not marked as confidential should not have access to confidential attributes or methods	SOFTWARE	SEMI-AUTOMATIC* STATIC*	÷?	I am not sure if this metric can be applied to embedded software
423 Least Privilege	This metric states that each component and user should be granted the minimal privi- leges required to complete their tasks. This metric can be assessed from two perspectives: from the security controls perspective we can review users' granted privileges. From the architectural analysis perspective this can be assessed as how the system is broken down to minimal possible actions i.e. the number of components that can access critical data. The smaller the value, the more secure the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	÷?	If the software is available, it can be applied. It can be also applied to embedded software
431 Trustworthiness	Trustworthiness of software means its worthy of being trusted to fulfill requirements which may be needed for a particular software component, application, system, or network. It involves attributes of stability, data security, quality, privacy, safety and so on. Software trustworthiness is interrelated with not only risk control in the software process, but also the quality management of the software development process. Furthermore, vision is needed to avoid excessive costs and schedule delays in development and risks management costs in operation; to improve development efforts; and above all	SOFTWARE			
433 Variable Security Vuherability	The security vulnerability of a variable v@l is determined by the combined security damage it might cause to the overall system security when v@l is attacked. The more security properties that can be left intact, the less vulnerable the variable is; on the other hand, the more security properties are violated, the more criticall the variable is.	SOFTWARE	АUТОМАТІС⁺ STATIC⁺	:?	This is only applicable if the source code is available
449 Faults found during manual inspection		SOFTWARE		YES	It is related to software and it can be applied to embedded software
458 Number of Design Flaws Related to Security (NSDF)	Secunity-related design flaws occur when software is planned and specified without proper consideration of security requirements and principles. For instance, clear-text passwords are considered as design flaws. Design flaws can be detected using design inspection techniques (e.g., design reviews). Identifying the number of design flaws related to security can help detect security issues earlier in the design process	SOFTWARE	SEMI-AUTOMATIC* STATIC*	<i>??</i>	If the design is available, it can be applied
459 Number of Exceptions That Have Been Implemented to Handle Execution Failures Related to Security (NEX)	Number of exceptions that have been included in the code to handle possible failures of the system due to an error in a code segment that has an impact on security	SOFTWARE	SEMI-AUTOMATIC* STATIC*	53	If the design is available, it can be applied
462 Number of Implementation Errors Found in the System (NERR)	Number of Implementation Errors Found in Number of implementation errors of the system the System (NERR)	SOFTWARE	SEMI-AUTOMATIC* STATIC*	:3	Related to the implementation phase
463 Number of Implementation Errors Related to Security (NSERR)	Number of implementation errors of the system that have a direct impact on security. One of the most common security related implementation errors is the buffer overflow	SOFTWARE	SEMI-AUTOMATIC* STATIC*	:?	Related to the implementation phase
464 Number of Omitted Exceptions for Handling Execution Failures Related to Security (NOEX)	Number of missing exceptions that have been omitted by software engineers when implementing the system. These exceptions can easily be determined through testing techniques.	SOFTWARE	SEMI-AUTOMATIC* STATIC*	:?	Related to the implementation phase
465 Number of Omitted Security Requirements (NOSR)	Number of requirements that should have been considered when building the application	SOFTWARE	SEMI-AUTOMATIC* STATIC*	53	If the design is available, it can be applied
471 Number of violations of the LP principle	Number of violations of the Least Privilege principle. A component does not adhere to LP if it, based upon the permissions attributed as described before, is capable of executing tasks it is not responsible for	SOFTWARE	АUТОСАПС⁴ STATIС⁴	YES	This could be used in embedded systems
473 Percentage Software Hazards (PSH)	Comparing the number of software hazards identified against historical data, it indicates the sufficiency of the software hazard identification based on the identified hazards. (# software hazards / # System hazards ) * 100	SOFTWARE	SEMI-AUTOMATIC* STATIC*	SÍ	Hazard is similar to risk.
480 Static analysis alert count (number of)		SOFTWARE			
488 Vulnerability Density	Number of vulnerabilities in the unit size of a code. VD = size of the software / $\#$ -vulnerabilities in the system	SOFTWARE	SEMI-AUTOMATIC* STATIC*	YES	It seems that it can be applied to embeded systems, but I am not sure if it is applicable in the evaluation phase

495 Security Metrics of Array Index	The security flaws of array index include index out of range, incorrect index variable -	SOFTWARE	AUTOMATIC*	STATIC*	53	This is not a metric, but a group of them. Individually,
	and uncontrollable array index					they could be used if they are specified. Part of a checklist
496 Security Metrics of Component Interfaces	The security flaws of component interface include inconsistent parameter numbers, inconsistent data types, and inconsistent size of data types	SOFTWARE	AUTOMATIC*	STATIC*	<i>?</i> غ	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
497 Security Metrics of Control Operation	The security flaws of control operations include infinite loop, deadlock situations and boundary defects	SOFTWARE	AUTOMATIC*	STATIC*	?5	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
498 Security Metrics of I/O Management	File is an important media which can store critical information and record log data However, file privilege definition, file access management, file contents, file organization and file size are major factors to affect the security software operation	SOFTWARE	AUTOMATIC*	STATIC*	<i>?</i> غ	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
499 Security Metrics of Input Format	The security flaws of data format include mismatched data contents, mismatched data type, mismatch data volume	SOFTWARE	AUTOMATIC*	STATIC*	<i>ڊ</i> غ	This is not a metric, but a group of them. Individually, they could be used if they are specified. Part of a checklist
500 Security Metrics of Kernel Operation	A released software system must be adapted to a suitable operating system and environment. The kernel of operating system and environment become a critical factor to affect the operating security of software system. Resource management, execution file protection, main memory control and process scheduling are major tasks of operating system. The security flaws, which embedded in the tasks of operating system, become high security risk for software system operation	SOFTWARE	А <b>∪ТОМАТІС</b> *	STATIC*	53	This is not a metric, but a group of them. Individually, they could be used if they are specified. Checklist. Only when there is an OS
501 Security Metrics of Network Environment	Network environment is a critical tool in the e-commerce years. However, data transmission management, remote access control, and transmission contents always become the security holes of software system operation	SOFTWARE	AUTOMATIC*	STATIC*	??	This is not a metric, but a group of them. Individually, they could be used if they are specified
502 Security Metrics of Resources Allocation	The security flaws of resource allocation include exhausted memory, unreleased memory - exhausted resources, and unreleased resources	SOFTWARE	AUTOMATIC*	STATIC*	;?	This is not a metric, but a group of them. Individually, they could be used if they are specified
503 Security Metrics of User Authority	Many users may use the released system to do specific jobs. However, user privilege - definition, user password management and user detailed information maintenance are important tasks to control and manage the user authority. Without user authority control and management, the operation environment of software system will be on high security risk	SOFTWARE	AUTOMATIC*	STATIC*	;?	This is not a metric, but a group of them. Individually, they could be used if they are specified
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## Appendix B

## **Curriculum Vitae**

# Ángel

# Longueira-Romero

#### Education

- 2019–2022 PhD in Industrial Cybersecurity, Faculty of Engineering, Mondragon University, Metric-Based Cybersecurity Evaluation Methodology for Industrial Embedded Systems. Specialization in:
  - Cybersecurity Evaluation Standards.
  - Vulnerability analysis.
  - Cybersecurity measurement.
- 2017–2019 Master's degree in Automation Engineering and Industrial Informatics, Gijón Polytechnic School of Engineering, University of Oviedo, Specialised in Optical Techniques for Industrial Inspection.

Other Activities:

- Class President.
- 2012–2017 **Bachelor's degree in Industrial Electronics and Automation Engineering**, *Gijón Polytechnic School of Engineering*, University of Oviedo, Specialised in Industrial Robotics. Other Activities:
  - MENTOR Programme. Providing help to new Erasmus students in the campus.
  - Member of the self-evaluation committee of the Industrial Electronics and Automation Engineering degree.
  - Tandem language exchange programme of the University of Oviedo.

#### Academic Dissertations

- 2019–2022 **Ph.D. Thesis**, *Mondragon University*, Arrasate, Metric-Based Cybersecurity Evaluation Methodology for Industrial Embedded Systems.
- 2018–2019 **Master's Thesis**, *IKERLAN*, Arrasate, Design and Implementation of a PKI Infrastucture with Automatic Reenrolment of Certificates for Industrial Embedded Systems. With Honors
  - 2017 **Undergraduate Thesis**, Thermal machinery and engines team. Department of Energy, Gijón, Design of a temperature regulation system for a Stirling engine heater. With Honors

#### Training

2018 **IEC 61131 Edu Net Certificate**, Automation technology in theory and practice according to IEC 61131.

#### Experience

- 2022-Present Industrial Cybersecurity Researcher, IKERLAN, Arrasate (Spain).
  - 2019–2022 **PhD Candidate**, *IKERLAN*, Arrasate (Spain).
  - 2018–2019 **Industrial Cybersecurity Intern**, *IKERLAN*, Arrasate (Spain), Worked at the Industrial Cybersecurity Team deploying a PKI with automatic reenrolment using SCEP, EST and OCSP for embedded devices with two HSMs.

- 2017–2018 Amplification Equipment Repair Apprentice, Ortodoxo Amplification (Now: Bell Tone Vintage), Gijón, Worked with Marco Fernández designing, building and repairing vacuum tube and solid state musical equipment, and developing ad hoc tools.
- 2016–2017 Mentoring and Tutoring, Private lessons of industrial technology, mathematics, physics, chemistry, technical drawing and music to high school students.

#### Technical Skills

OSs Windows, GNU/Linux.

Programming C, C++, Python, assembly.

Databases PostgreSQL, MySQL.

Software MatLab, R, Altium.

Hardware FPGAs, Microcontrollers, PCB design.

Standards ISO 9001, ISO 27005, ISA/IEC 62443, Common Criteria, SAE J3061.

Methodologies NIST 800 Serie Reports, OWASP, OSSTMM.

Data Analysis Machine Learning, Computer vision, Data Processing, Data Visualisation.

Cybersecurity Cryptography, Public Key Infrastructure (PKI), Risk Analysis (MAGERIT), Threat Mod-

elling (STRIDE).

Protocolos SCEP, EST, OCSP.

Other Microsoft Office, LATEX.

#### Languages

Spanish Native

English Fluent

Basque Intermediate Conversationally fluent

#### Transferable Skills

- Teamwork - Self management

- Written and verbal communication - Research and analytical skills

- Fast learner and proactive problem solving

#### Interests

- 3D design with Fusion 360 (Autodesk) - 3D printing (Simplify3D and Cura)

- Arduino projects - Raspberry projects - Language learning - CrossFit and Bike

- Travelling - Music (Electric Bass and Oboe)

#### Other Information

- Willing to travel - Interested in a wide variety of topics

- Driving license