

Graphical Ingnuity: Spectral Solutions to Combinatorial Conondrums

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Table of Contents

1 Preface

- SGT Talk Thursday 4/25

2 Background

3 Structure and Spectra

4 Problem I

5 Problem II

6 Problem 3: Open

Preface

SGT Talk Thursday 4/25

Reimagining Spectral Graph Theory



Dr. Stephen Young

Algorithms, Combinatorics, and Optimization Team Lead
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April 25, 2024 at 4:20
Spark 335

Zoom ID: 992 1304 8021 Passcode: 483214

Refreshments at 3:30pm
Neill 216 (Hacker Lounge)

Background

Basics

Definition

The adjacency matrix of a graph G , $A(G) = A = [a_{ij}]$, is the $n \times n$ matrix such that

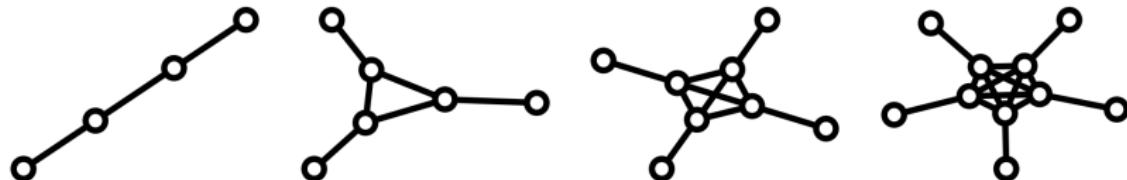
$$a_{ij} = \begin{cases} 1 & \text{if } i \text{ adjacent to } j \text{ in } G \\ 0 & \text{otherwise} \end{cases}$$

for $i, j = 1, 2, \dots, n$

Definition

Throughout, the spectrum of a graph G will refer to the collection of eigenvalues of A , $\sigma(A)$, with the ordering $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$

Examples I



Adjacency spectra starting from the left where $\phi = \frac{1+\sqrt{5}}{2}$:

- $\sigma(A) = \{-\phi, -\frac{1}{\phi}, \frac{1}{\phi}, \phi\}$
- $\sigma(A) = \{-\phi, -\phi, 1 - \sqrt{2}, \frac{1}{\phi}, \frac{1}{\phi}, 1 + \sqrt{2}\}$
- $\sigma(A) = \{-\phi, -\phi, -\phi, -.302 \dots, \frac{1}{\phi}, \frac{1}{\phi}, \frac{1}{\phi}, 3.302 \dots\}$
- $\sigma(A) = \{-\phi, -\phi, -\phi, -\phi, .236 \dots, \frac{1}{\phi}, \frac{1}{\phi}, \frac{1}{\phi}, \frac{1}{\phi}, 4.236 \dots\}$

Examples II

- Let $[n] = \{1, 2, \dots, n\}$.
- Let nodes of graph G correspond to distinct k -element subsets of $[n]$.
- If $|i \cap j| = 0$ for $i, j \in V(G)$, then let i be adjacent to j i.e. $(i, j) \in E(G)$.

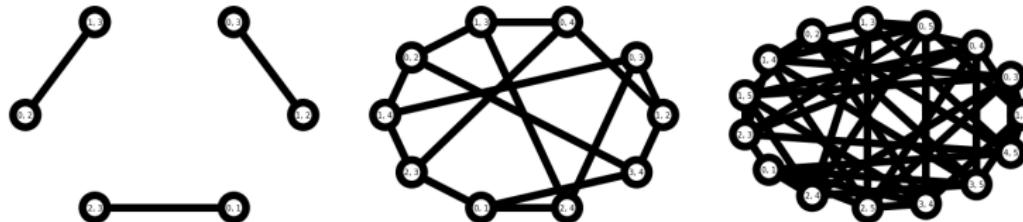


Figure: The Kneser Graphs $K(n, k)$ with
 $\sigma(A) = \{(-1)^j \binom{n-k-j}{k-j} \mid j = 0, 1, 2, \dots, k\}$

Structure and Spectra

Basic Examples

Let G be a simple, undirected graph on n nodes. Then,

- $\sum_{i=1}^n \lambda_i = 0$
- $\text{tr}(A^2) = \sum_{i=1}^n \lambda_i^2 = 2m$ where m is the number of edges in the graph
- $\Delta(G) \geq \lambda_1 \geq \bar{d}(G)$ where $\bar{d}(G)$ is the average degree in G and $\Delta(G)$ is the max. Also, $\lambda_1 = \bar{d}(G)$ for k -regular graphs where $\bar{d} = k$.

Hoffman's Bound

Theorem

Let $G = (V, E)$ be finite, k -regular graph on n nodes with adjacency matrix A . Let $S \subset V(G)$ be an independent set of vertices in G . Then

$$\frac{|S|}{n} \leq \frac{-\lambda_n}{k - \lambda_n}$$

Proof

- Let \mathbb{R}^n be equipped with inner product

$$\langle u, v \rangle = \frac{1}{n} \sum_{x \in V(G)} u(x)v(x)$$

- Let $S \subset V(G)$ be an independent set with indicator vector 1_S
- Using the orthonormal basis of eigenvectors for A , $\{v_1, v_2, \dots, v_n\}$, we can rewrite 1_S as: $1_S = \sum_{i=1}^n \alpha_i v_i$.
- Let \mathbf{j} be the all ones vector. Note that
$$\alpha_1 = \langle 1_S, \mathbf{j} \rangle = \frac{|S|}{n} = \langle 1_S, 1_S \rangle$$
- Since all v_i are normalized, $\langle 1_S, 1_S \rangle = \sum_{i=1}^n \alpha_i^2$

Proof

- Lastly, since $0 = \langle A\mathbf{1}_S, \mathbf{1}_S \rangle = \sum_{i=1}^n \lambda_i \alpha_i^2 \geq d\alpha_1^2 + \lambda_n \sum_{i=2}^n \alpha_i^2 = d\alpha_1^2 + \lambda_n(\alpha_1 - \alpha_1^2)$
- Thus, $\frac{|S|}{n} = \alpha_1 \leq \frac{-\lambda_n}{d-\lambda_n}$

Problem I

Erdős-Ko-Rado

Definition

Let \mathcal{F} be a family of sets. If $|X \cap Y| \neq 0 \forall X, Y \subset \mathcal{F}$, then \mathcal{F} is an *intersecting family*.

Question

For $n, k \in \mathbb{N}$, what is the maximum size of an intersecting family of k -element subsets of an n -element set?

Proof

- Consider the Kneser Graph $K(n, k)$.
- Note that the Kneser Graph is $\binom{n-k}{k}$ -regular with spectrum $\{(-1)^j \binom{n-k-j}{k-j} \mid j = 0, 1, 2, \dots, k\}$ making $\lambda_n = -\binom{n-k-1}{k-1}$.
- So, using Hoffman's Bound,

$$\begin{aligned}\frac{|S|}{n} &\leq \frac{-\lambda_n}{k - \lambda_n} \longrightarrow \frac{\alpha}{\binom{n-k}{k}} \leq \frac{\binom{n-k-1}{k-1}}{\binom{n-k}{k} + \binom{n-k-1}{k-1}} \\ \iff \alpha &\leq n \frac{\binom{n-k-1}{k-1}}{\binom{n-k}{k} + \binom{n-k-1}{k-1}} = \binom{n-1}{k-1}\end{aligned}$$

Problem II

Dimer Problem

- The ‘dimer’ problem named after the two dimensional lattices of ‘diatomic’ molecules that form on surface of crystals.
- Dimers are forced to occupy two neighboring points on the lattice.

Question

How many ways can dimers cover the surface of the crystal?

Question

How many perfect matchings, p , are there for an $m \times n$ grid?

Short answer: take the permanent! ($\text{per}(A) = p^2$)

Count

- Let the undirected graph G be an $m \times n$ grid with adjacency matrix A and H be the graph obtained by replacing each edge in G with oppositely oriented edges.
- Note $\text{perm}(A) = \sum_{\sigma \in S_n} \prod_{i=1}^n a_{i,\sigma(i)} = p^2 = \sum_{L \subset H} 1$ where L is a spanning subgraph of H .
- Likewise,
$$\det(A) = \sum_{\sigma \in S_n} (\text{sgn}(\sigma) \prod_{i=1}^n a_{i,\sigma(i)}) = \sum_{L \subset H} (-1)^{c(L)}$$
where $c(L)$ is the number of components in L .
- For every $L \subset H$, $(-1)^{c(L)} = i^{h(L)}$ where $h(L)$ is the number of horizontal edges in L .
- A can be factored as $A_m \otimes \mathbb{I}_n + A_n \otimes \mathbb{I}_m$.

Count

- Take $A^* = A_m \otimes \mathbb{I}_n + iA_n \otimes \mathbb{I}_m$ (multiplying each horizontal edge by i).
- Then, since $(-1)^{c(L)} = i^{h(L)}$, $\text{per}(A) = \det(A^*)$ and so $\det(A^*) = p^2$
- Thus, $p^2 = \prod_{i=1}^n \lambda_i$ for $\lambda_i \in \sigma(A^*)$ i.e.

$$p^2 = \prod_{j=1}^m \prod_{k=1}^n \left(2 \cos\left(\frac{\pi j}{m+1}\right) + 2i \cos\left(\frac{\pi k}{n+1} I\right) \right)$$

Problem 3: Open

Bounded Permutations

- Consider permutations π of $[n]$ where $|\pi(i) - i| \leq k$ for $i = 1, 2, \dots, n$
- Denote the number of restrictions of this type by $N(\mathbf{R}^{(k)}, n)$
- Let $k < h < n$ with $k \ll n$ and consider a partially constructed permutation of the type $\mathbf{R}^{(k)}$:

$$\begin{bmatrix} 1 & 2 & \cdots & h-1 \\ \pi(1) & \pi(2) & \cdots & \pi(h-1) \end{bmatrix}$$

- The next step in our permutation is to select a value for $\pi(h)$
- At this point there are a certain number of ‘states of incompletion’ of the permutation
- Assign each node i in a digraph G a distinct state such that i is adjacent to j if and only if there is a value of $\pi(h)$ which will take state i and produce state j .

Bounded Permutations

Conjecture

The limit

$$\lim_{n \rightarrow \infty} N(\mathbf{R}^{(k)}, n)^{1/n}$$

exists and is equal to the largest eigenvalue of G .

Question

What about for permutations π such that $j \leq \pi(i) - i \leq k$?