# CS 564 Midterm Exam Fall 2018 Answers

## A: Let's write some queries! [30%]

We will use the following database schema:

Student (<u>stid</u>, firstname, lastname, major)
Class (<u>code</u>, title, semester, year, description)
Enrolled (<u>stid</u>, <u>code</u>, numcredits)

Furthermore, Enrolled.stid is a foreign key referring to Student.stid and Enrolled.code is a foreign key referring to Class.code

1. **[10%]** Express the following query in SQL: output the student IDs of the students that have taken at least 30 credits across all classes.

SELECT e.stid

FROM Enrolled e

GROUP BY e.stid

HAVING SUM(e.numcredits) > 29;

2. **[10%]** Express the following query in SQL: for each class offered in Spring 2018, output the class code and the number of CS majors that took that class.

SELECT c.code, COUNT(s.stid)

FROM Student s, Class c, Enrolled e

WHERE s.stid = e.stid AND c.code = e.code

AND s.major = 'CS' AND c.semester = 'Spring' AND c.year = 2018

GROUP BY c.code;

3. [10%] Express the following SQL query in Relational Algebra:

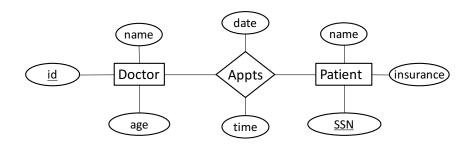
```
SELECT DISTINCT firstname, lastname FROM Student
WHERE stid NOT IN (
SELECT s.stid
```

```
FROM Student s, Enrolled e
WHERE s.stid = e.stid AND e.numcredits > 3);
```

## B: ER, RELATIONAL MODEL, AND SOME SQL [13%]

For the following questions, circle **exactly one** correct option.

1. [5%] Consider the following ER diagram:



If we translate the relationship **Appts** to a relation in the relational model, how many attributes will the relation have?

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2. [4%] Consider the schema from Part A. The following SQL queries will always give the same result:

```
SELECT * FROM Enrolled
LEFT OUTER JOIN Student on Student.stid = Enrolled.stid;
SELECT * FROM Enrolled
INNER JOIN Student on Student.stid = Enrolled.stid;
```

#### **TRUE**

3. [4%] Consider the following relational table for Student:

studentID	name	age
1234	George	NULL
1144	Anna	19
2214	Maria	NULL

How many tuples will the following SQL query return?

SELECT \* Student FROM WHERE name = 'George' OR age > 20;

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## C: NORMALIZATION AND DEPENDENCY THEORY [42%]

Consider the relation R(A, B, C, D, E) with the following set F of functional dependencies:

 $A \rightarrow B, C$   $D \rightarrow C$   $E \rightarrow D$   $B, E \rightarrow C$ 

[10%] Add a tuple to the instance of *R* below such that **all** fds in *F* are violated:

Α	В	C	D	E
a	b	С	d	e
a	b	С	ď	e'

ANSWER: a,b,c',d',e

For the following questions, circle the right option(s).

#1 There can be more than one correct options for every question!

#2 You can get partial credit by explaining how you came up with your answer.

1. **[12%]** The following attribute sets are **superkeys** but **not keys** in relation *R*:

(i) A, E

(ii) A, D, E (iii) A, B, C, D (iv) A, B, C, D, E

ANSWER: (ii) and (iv)

2. [8%] The following functional dependencies are **redundant** in *F* (in other words, if we remove the functional dependency, the fd closure remains the same):

(i)  $A \rightarrow B$ , C (ii)  $D \rightarrow C$  (iii)  $E \rightarrow D$  (iv) B,  $E \rightarrow C$ 

ANSWER: (iv)

3. [12%] The following hold for the decomposition of R into ABC, AE, DE. It is:

(i) lossless-join

(ii) dependency preserving

ANSWER: (i)

### D: Buffer Management [15%]

In this question, we consider a buffer pool with 5 frames, and two files: one with three pages  $A_1$ ,  $A_2$ ,  $A_3$ , and the other with four pages  $B_1$ ,  $B_2$ ,  $B_3$ ,  $B_4$ . We want to read these two files in a nested loop (as we will see later in class, this is one possible implementation of a join between the two files). The sequence of requests is as follows:

```
Request A_1, Request B_1, Release B_1, Request B_2, Release B_2, ..., Release A_1, Request A_2, Request B_1, Release B_1, Request B_2, Release B_2, ..., Release A_2, Request A_3, Request B_1, Release B_1, Request B_2, Release B_2, ..., Release A_3.
```

Initially, all buffer frames are free, and none of the pages to be accessed are in RAM. For the following questions, circle the right option.

#1 You can get partial credit by explaining how you came up with your answer.

1.	[10%] If the buffer manager uses the LRU replacement policy, how many total I/Os (both reads and writes) occur for the above sequence?					
	(i) 5	(ii) 7	(iii) 10	(iv) 15		
	ANSWER	R: 15				
2.	2. [5%] If the buffer manager knows the sequence a priori, how many total I/Os do the optimal replacement policy achieve?					
	(i) 0	(ii) 5	(iii) 7	(iv) 10		
	ANSWEE	₹• 7				

BONUS: write your own SQL joke!