CS 354
Machine Organization and Programming
Lecture 23

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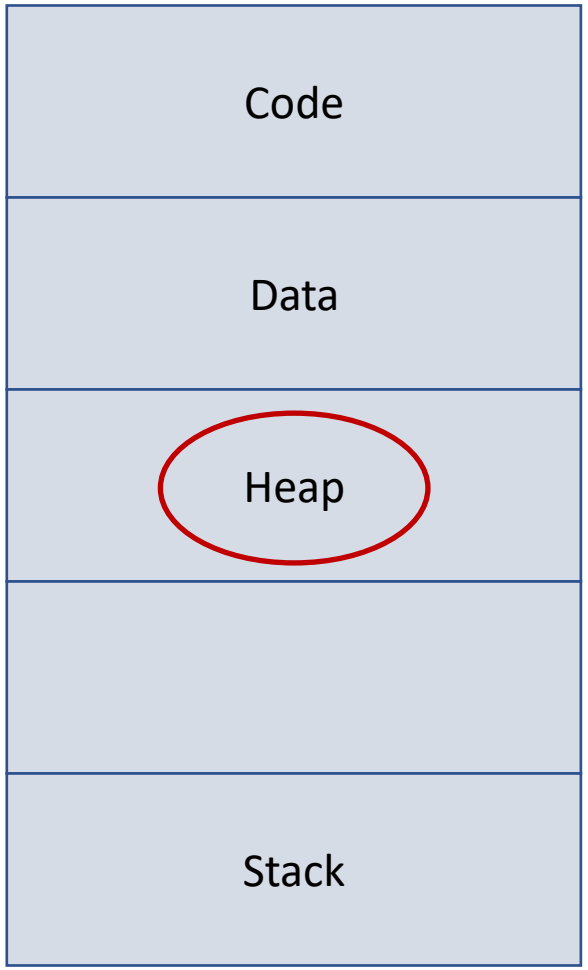
Summer 2020

Dynamic Memory Allocation
Part 2

Address Space

Memory Allocator Requirements

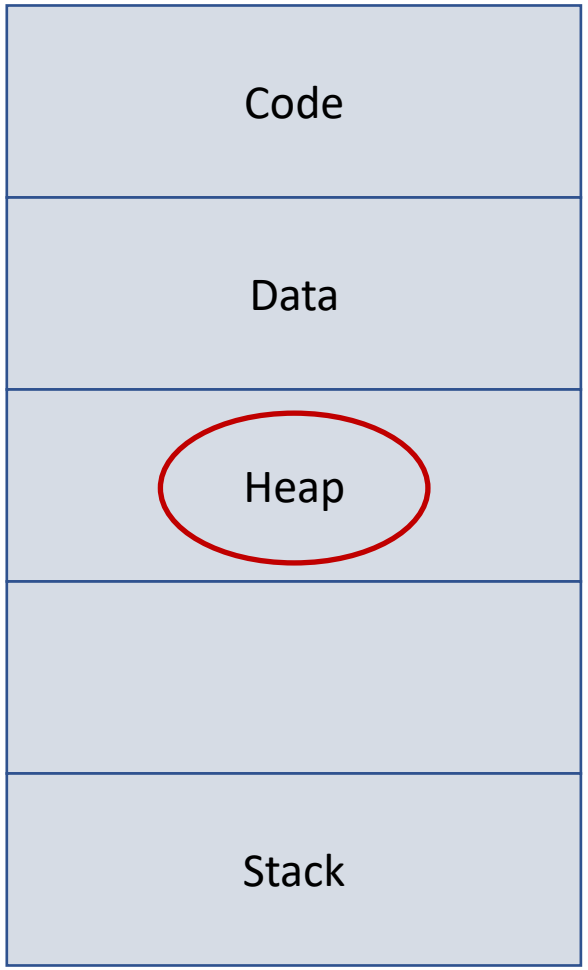
1. Keep track of block sizes for freeing
2. Arbitrary Request Sequences
3. Immediate Response
    1. no buffering
    2. or reordering
4. Use only the heap
5. Block Alignment
6. No modification of allocated blocks



Address Space

Implementation Issues

1. Tracking Free Blocks – Free List
2. Placement Policy
3. Splitting
4. Coalescing



Address Space

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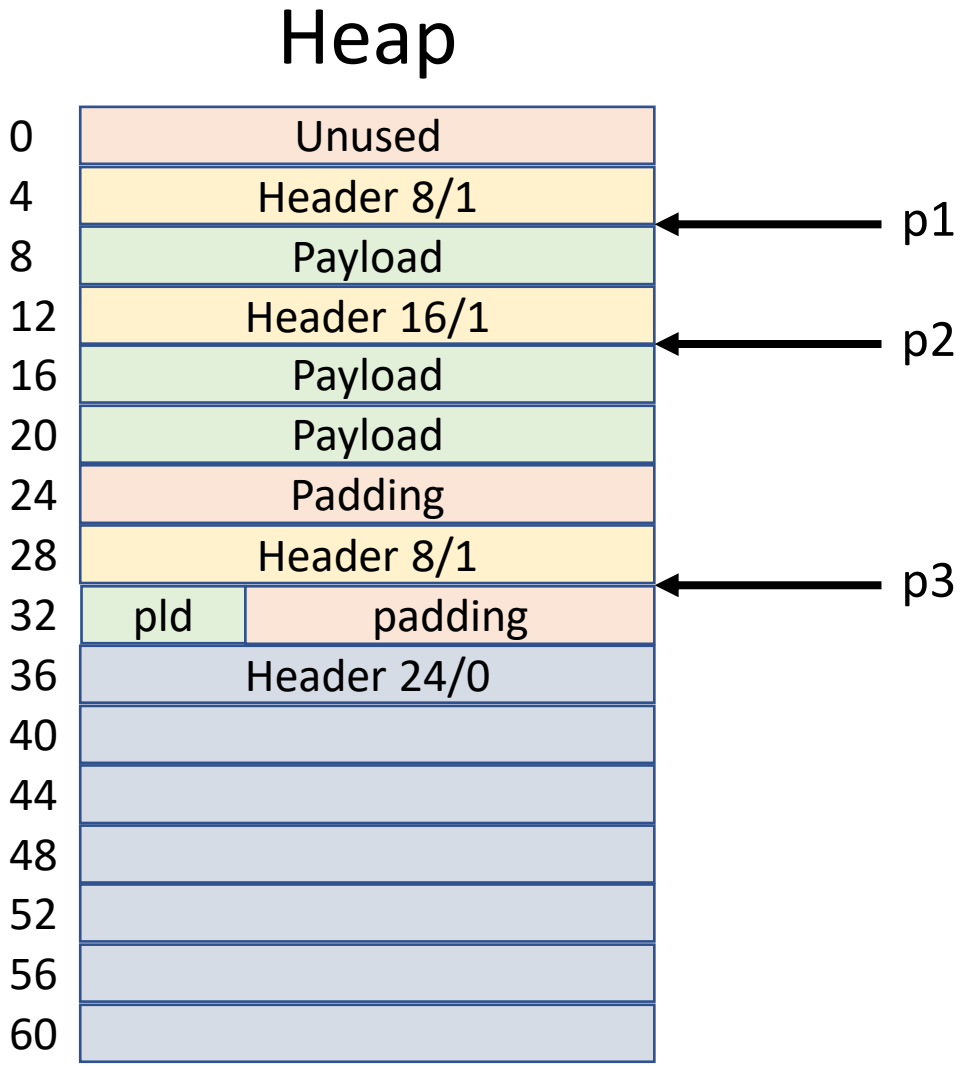
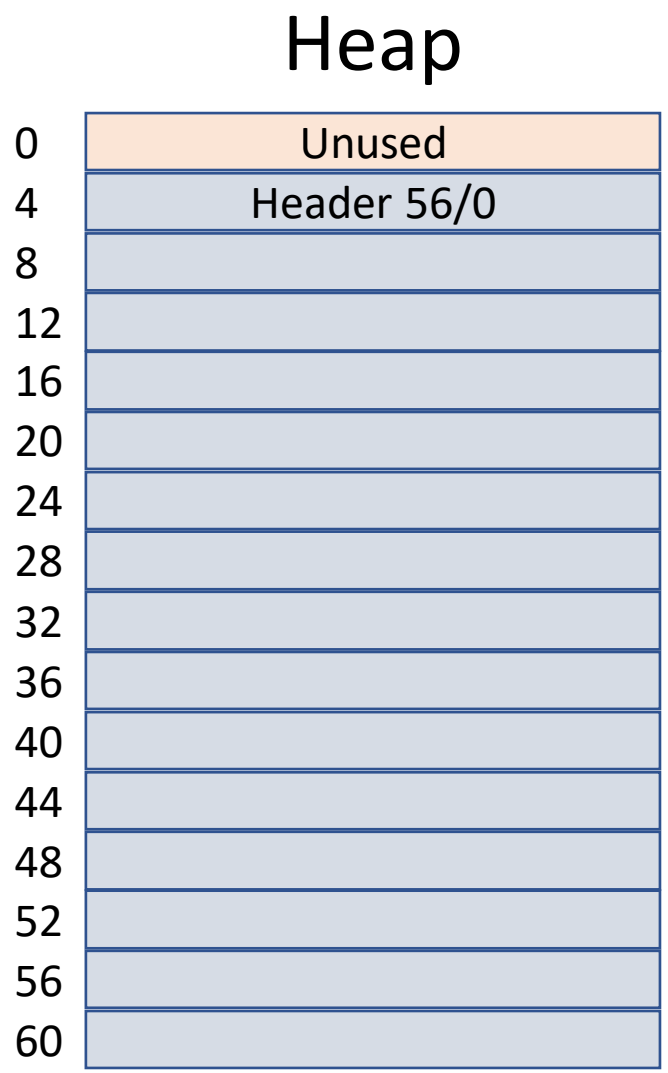


Implicit Free List

```
int *p1 = malloc(4);
int *p2 = malloc(8);
int *p3 = malloc(1);

free (p2)
```



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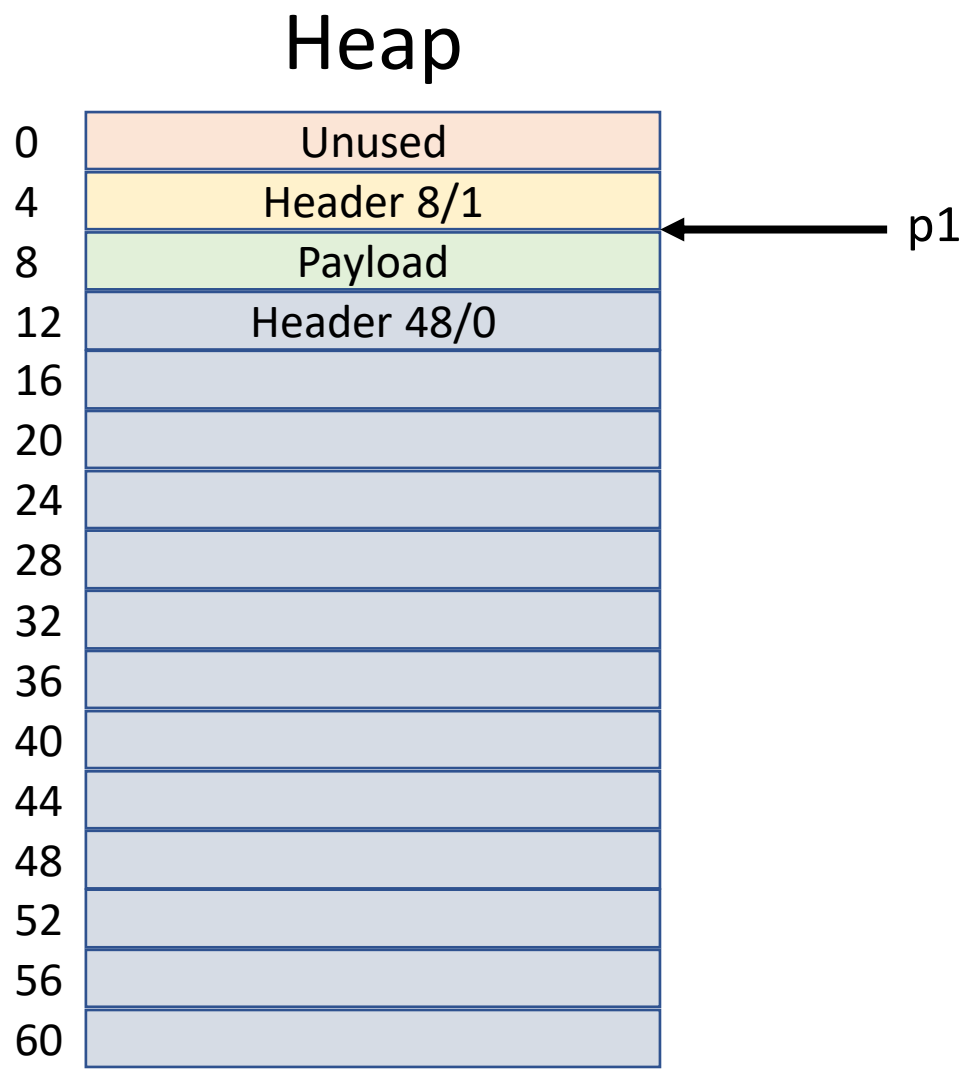

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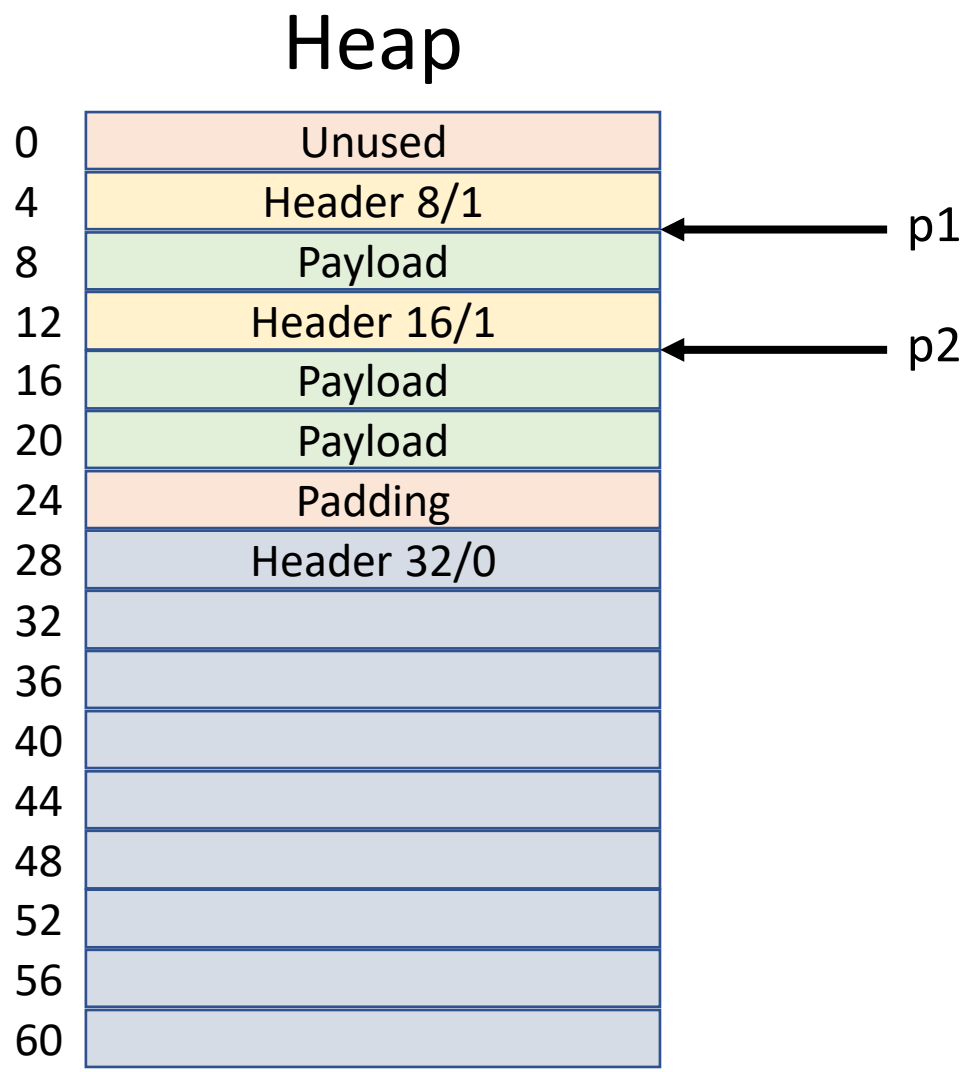

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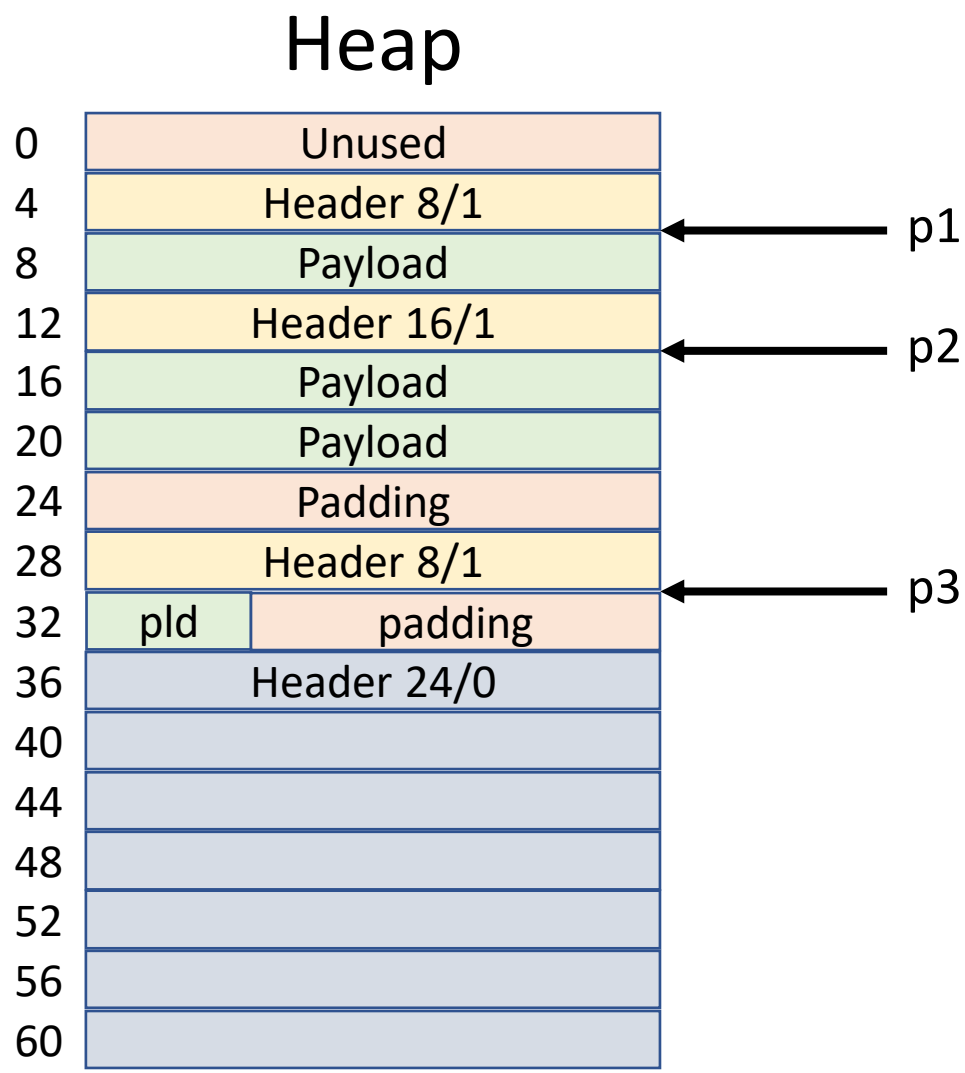

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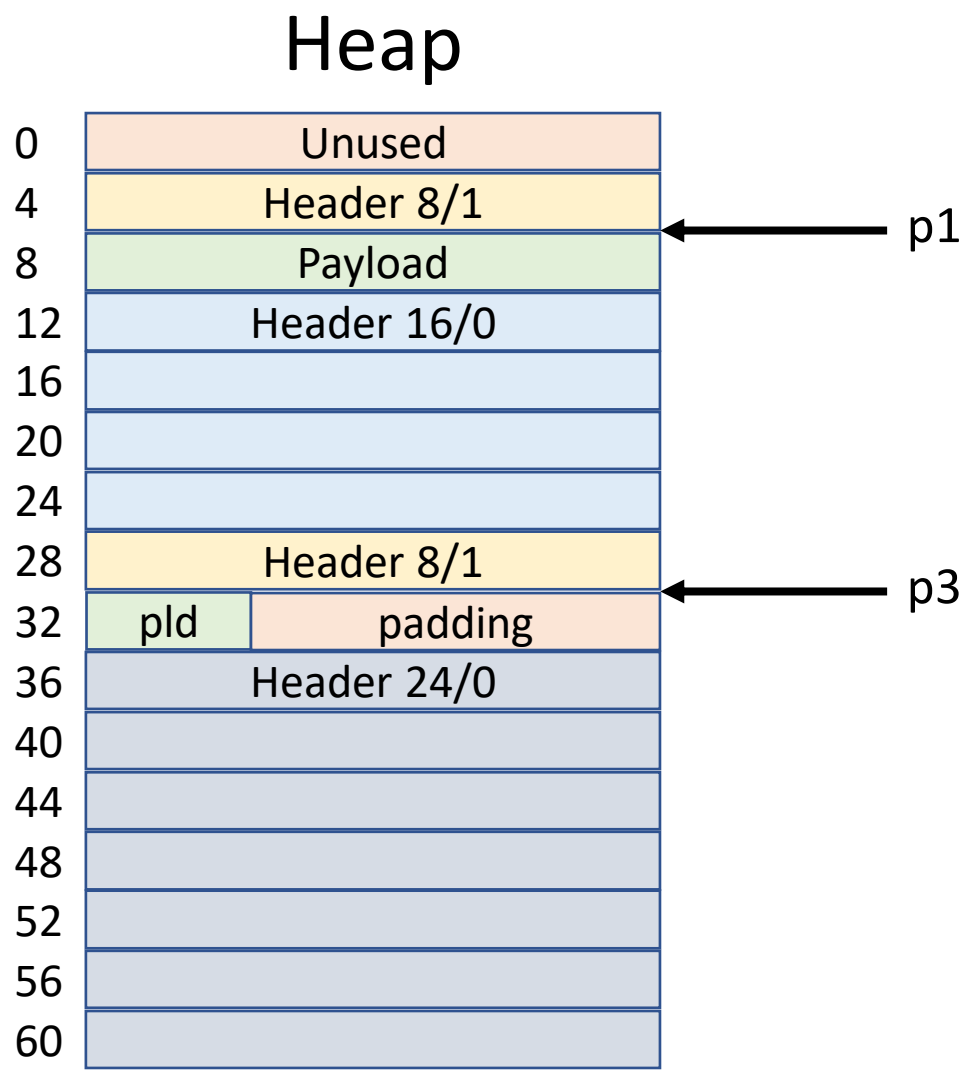

Implicit Free List

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free (p2)
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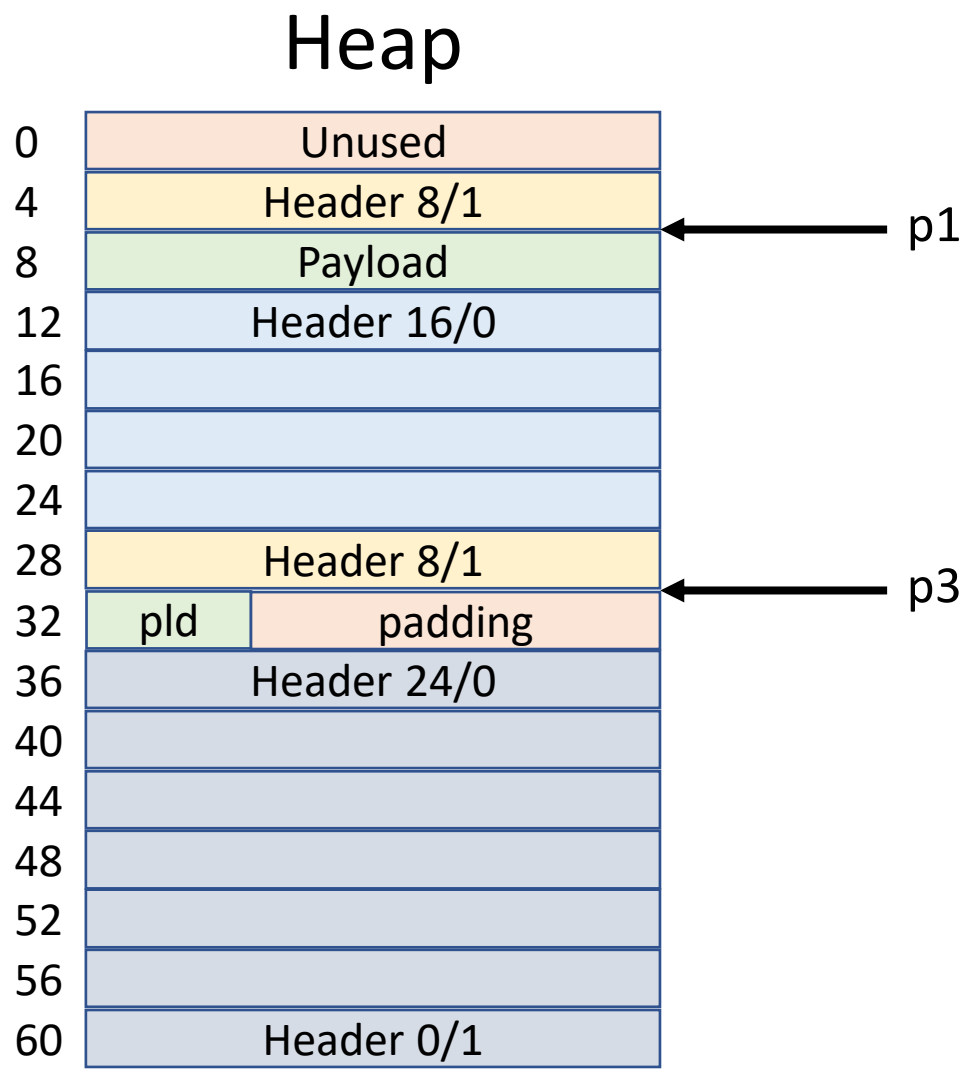
End of the Stack Marker



Implicit Free List

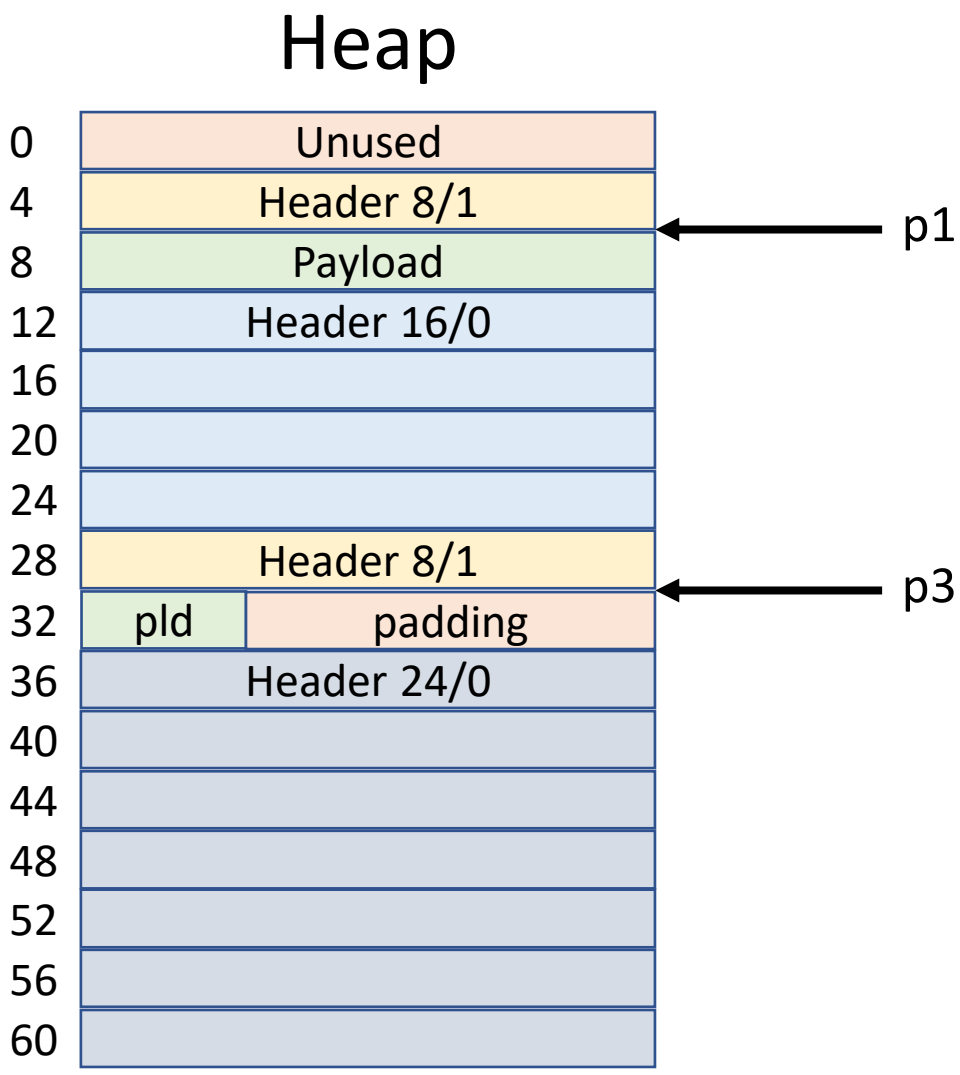
How to pack the size of the allocated block and the allocated / free flag into 4 bytes?

The size must be a multiple of 8

0x 0008 : 8 : 0000 1000
0x 0010 : 16 : 0001 0000
0x 0018 : 24 : 0001 1000
0x 0020 : 32 : 0010 0000

Just use the b0 bit for the flag

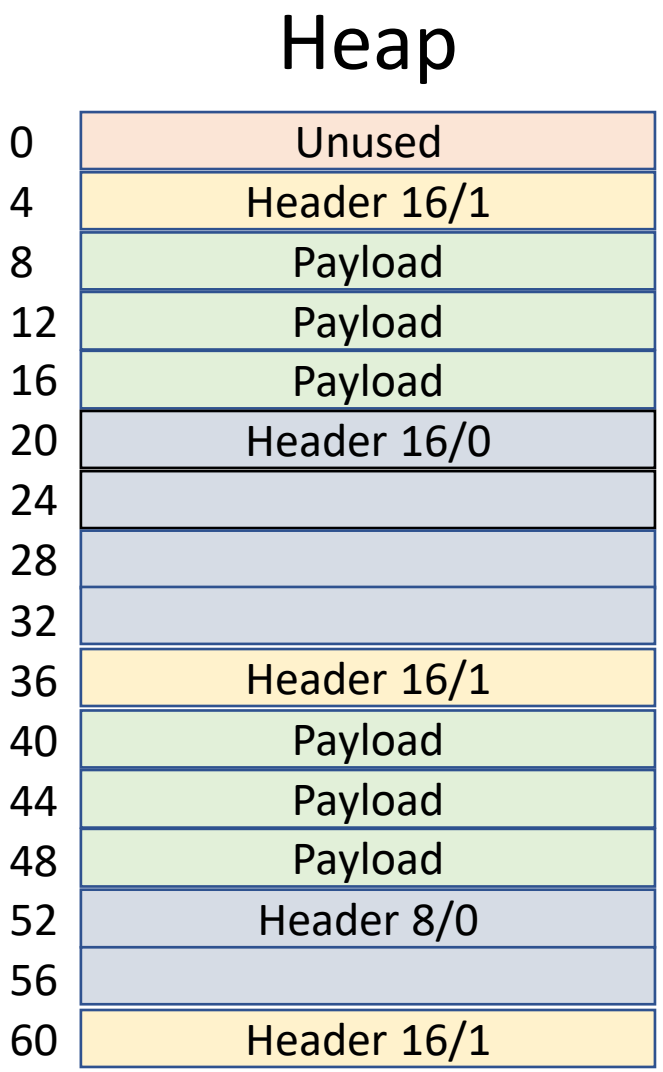
8/0 -> 0000 1000 = 8
8/1 -> 0000 1001 = 9
16/1 -> 0001 0001 = 17



Implicit Free List

1. Tracking Free Blocks – Free List
2. Placement Policy
3. Splitting
4. Coalescing

malloc(4)
malloc(12)



Implicit Free List

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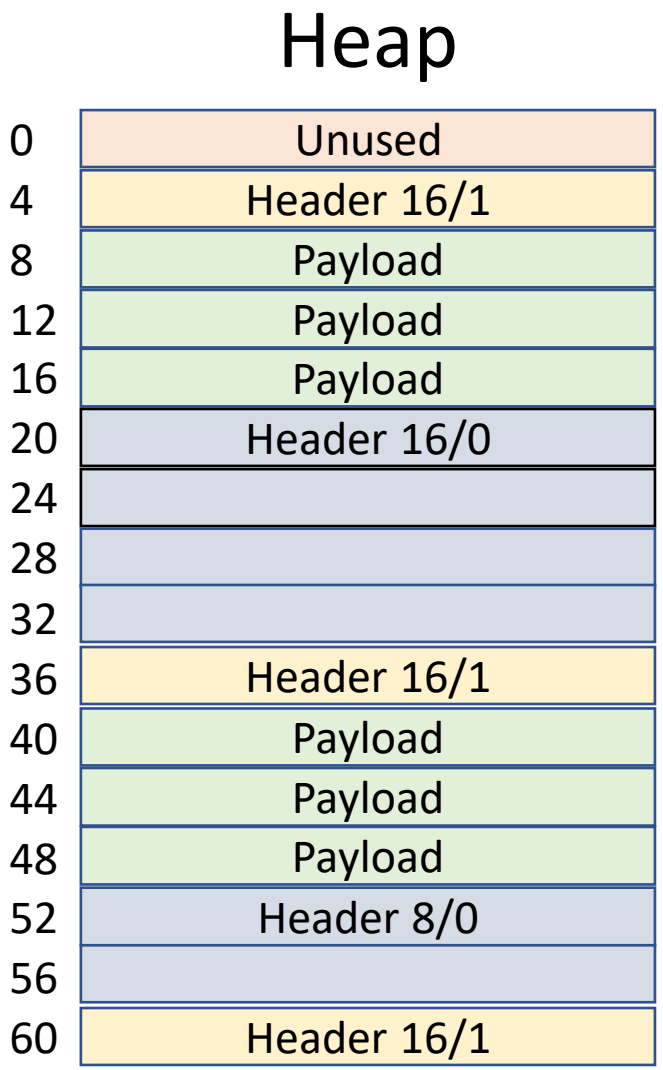
Linked List



Implicit Free List

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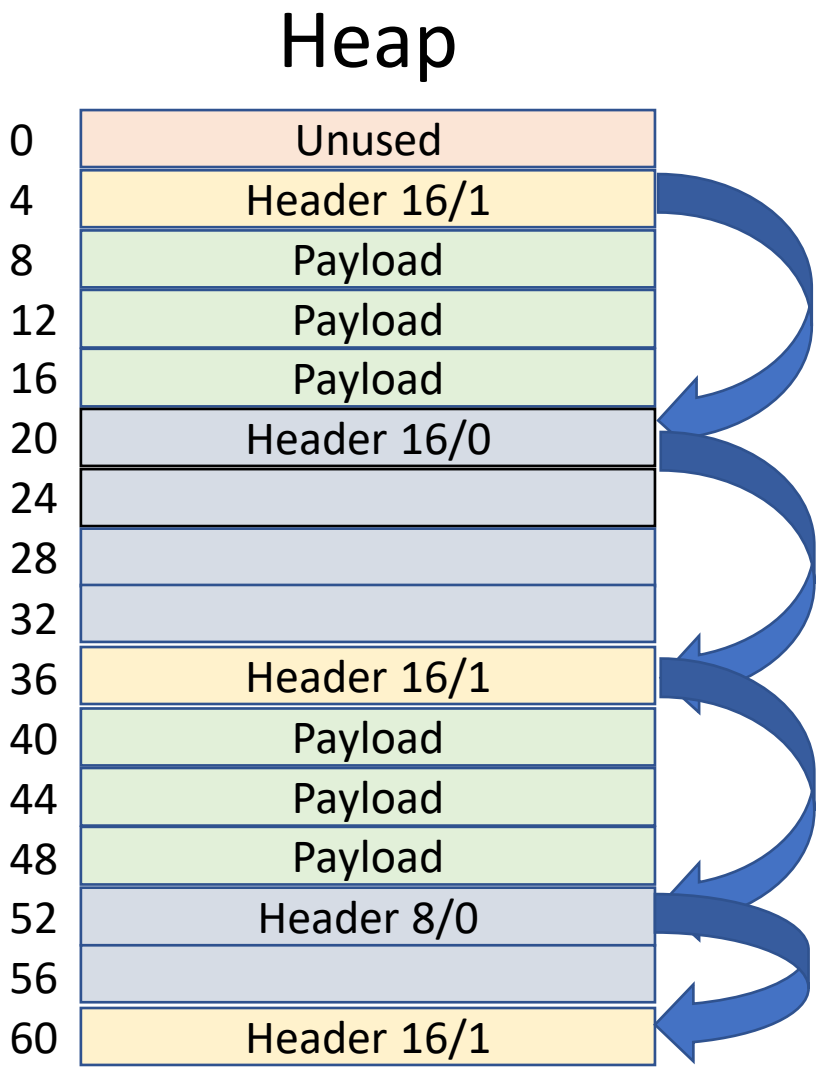
Linked List

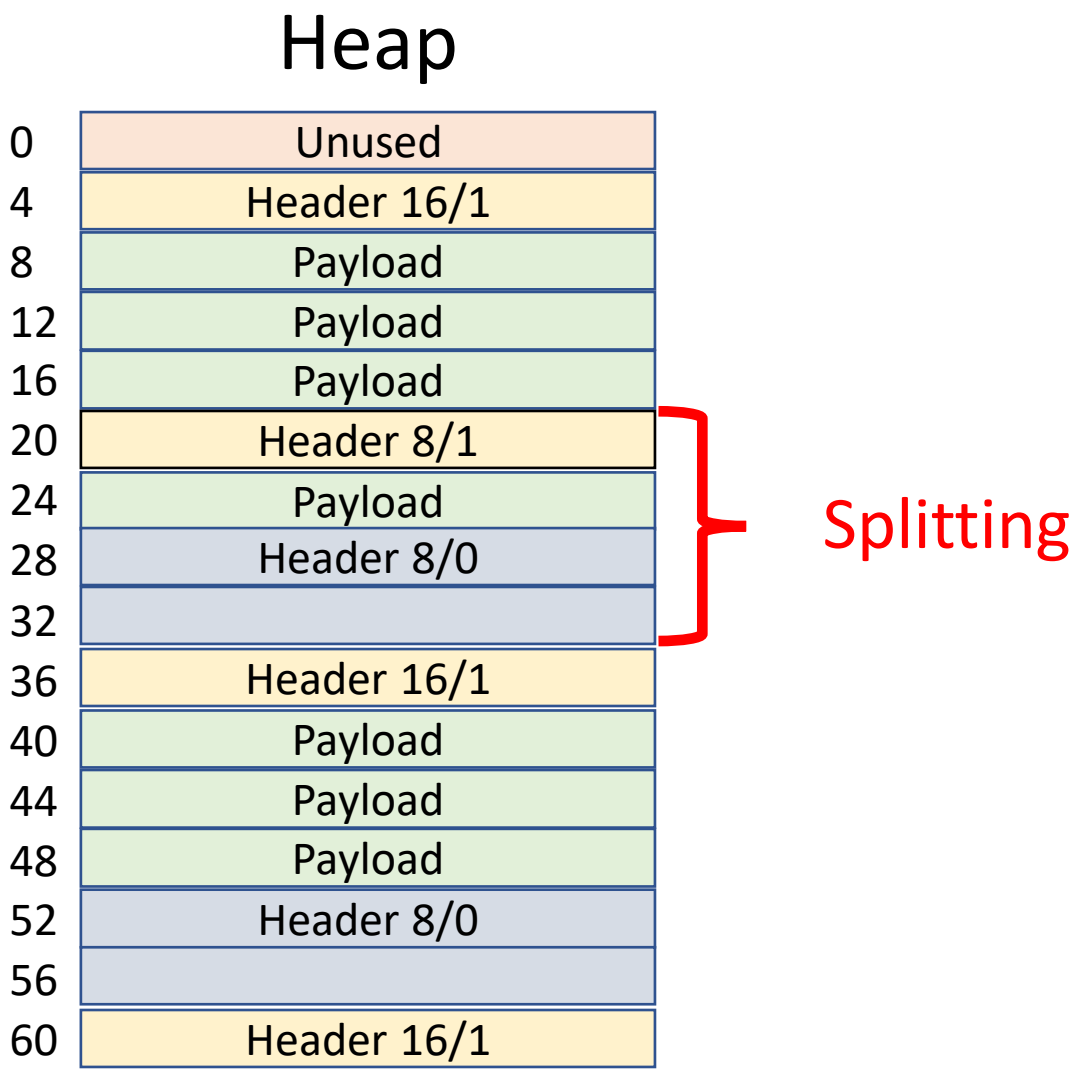


Placement Policies

1. Tracking Free Blocks – Free List
2. Placement Policy
3. Splitting
4. Coalescing

First Fit

- malloc(4) -> success
- malloc(12)



Placement Policies

1. Tracking Free Blocks – Free List
2. Placement Policy
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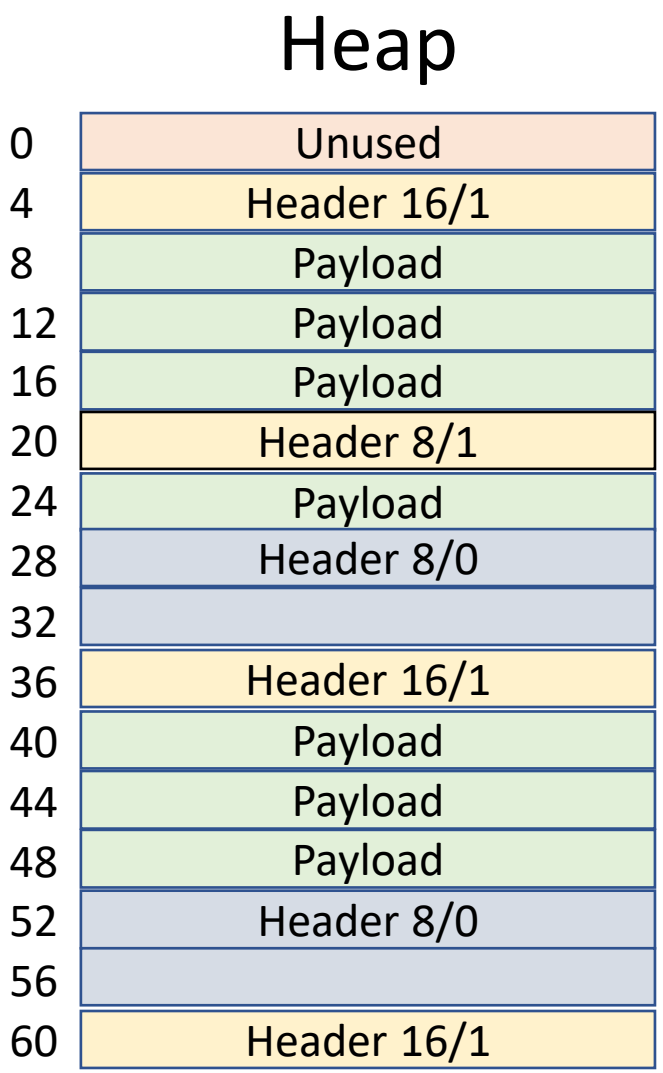
4. Coalescing

First Fit

- malloc(4) -> success
- malloc(12) -> fails



Placement Policies

1. Tracking Free Blocks – Free List
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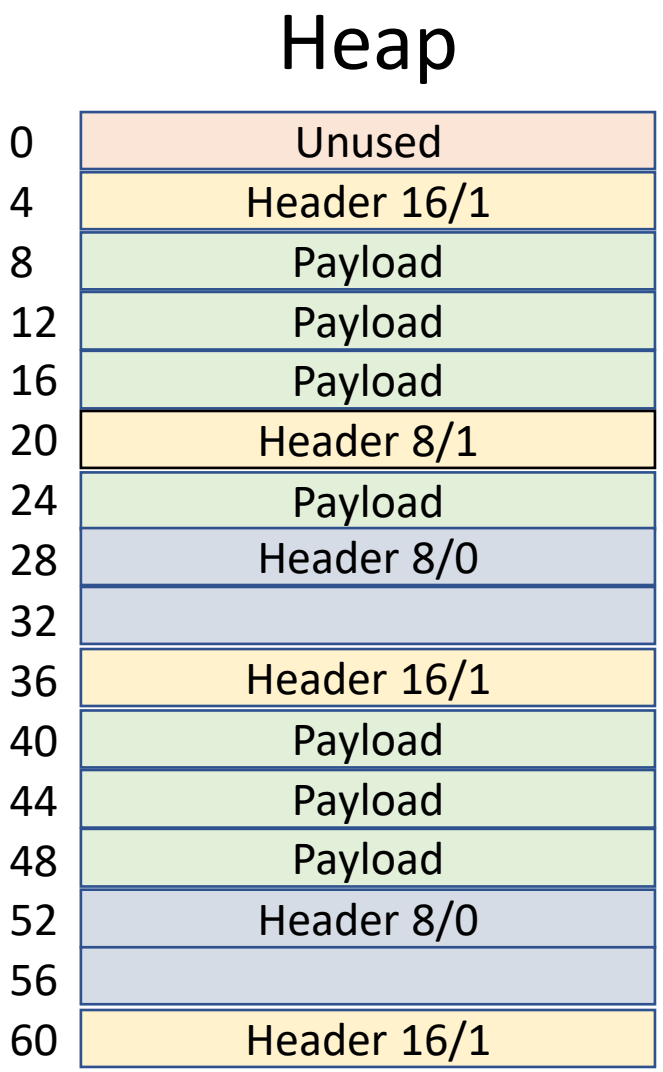
First Fit

- retains large blocks at the end of the list
- but fragments the beginning of the list
- increases search time



Placement Policies

1. Tracking Free Blocks – Free List
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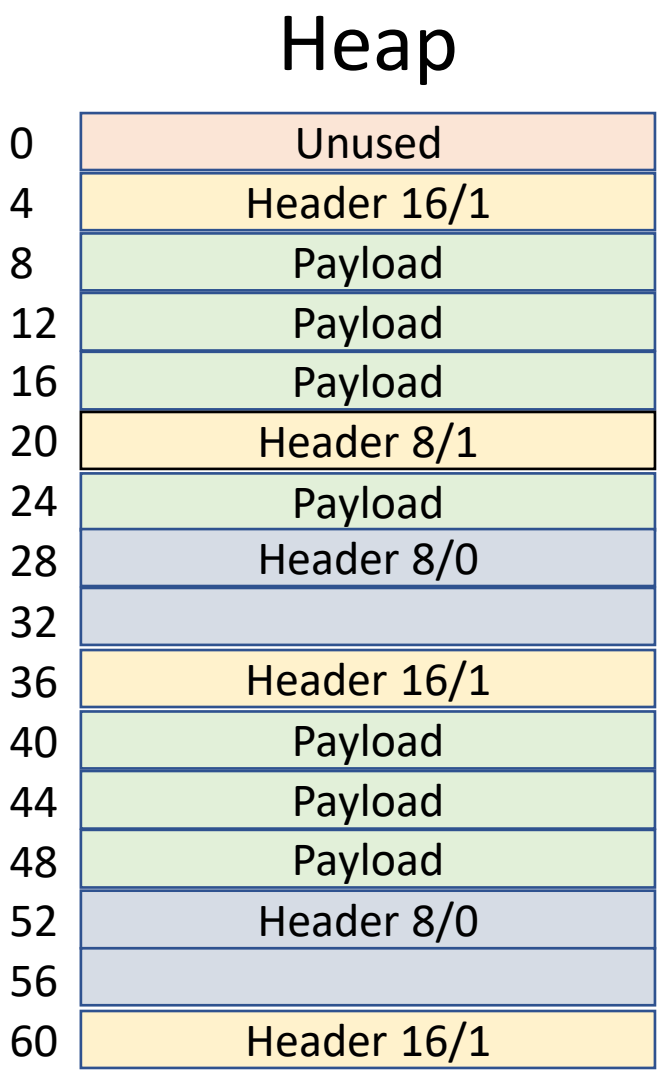
First Fit

Next Fit

- begin searching for the next available block from the last found.



Placement Policies

1. Tracking Free Blocks – Free List
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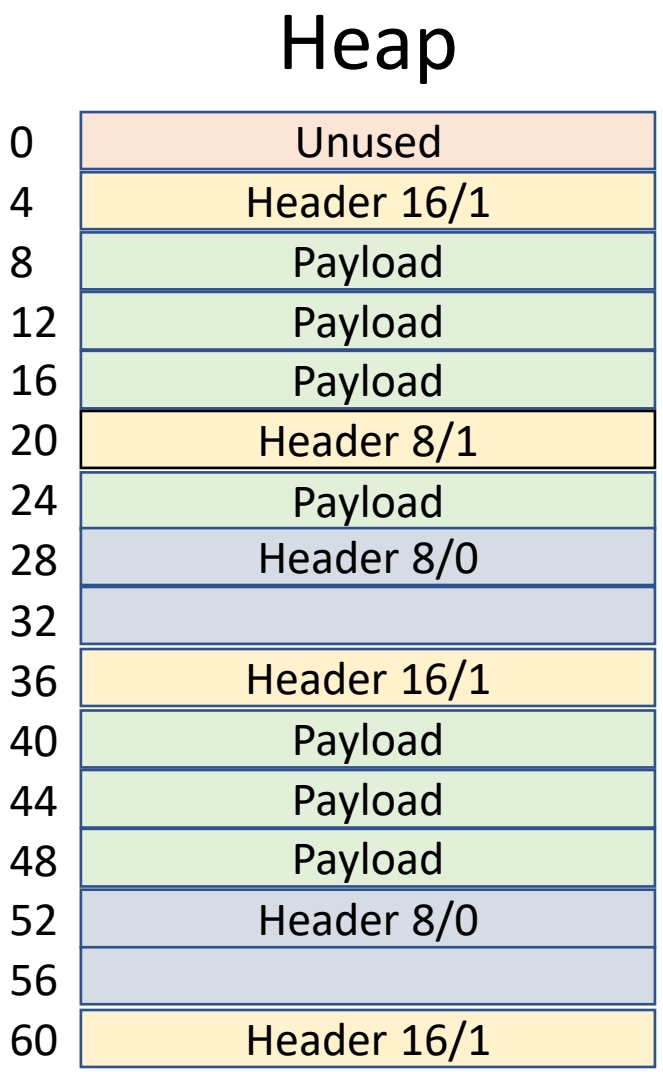
First Fit

Next Fit

- begin searching for the next available block from the last found.
- Faster than First Fit
- Worse Memory Utilization



Placement Policies

1. Tracking Free Blocks – Free List
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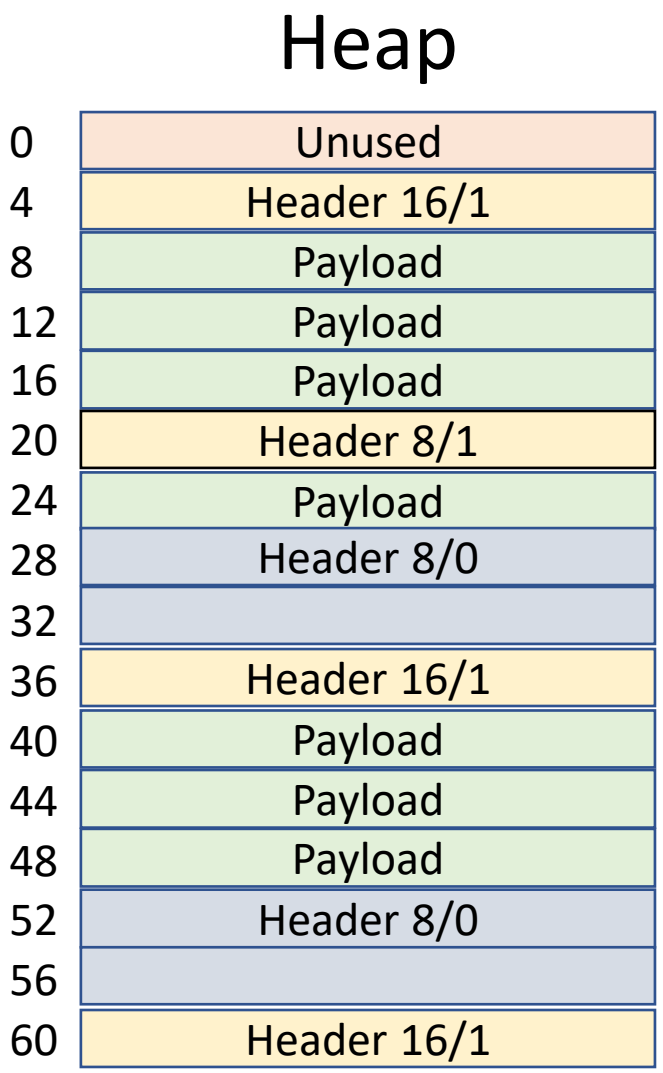
3. Splitting
4. Coalescing

First Fit

Next Fit

Best Fit

- Search the entire memory for the smallest block that will work
- Slow
- Good memory utilization
- Small Fragments



Placement Policies

1. Tracking Free Blocks – Free List
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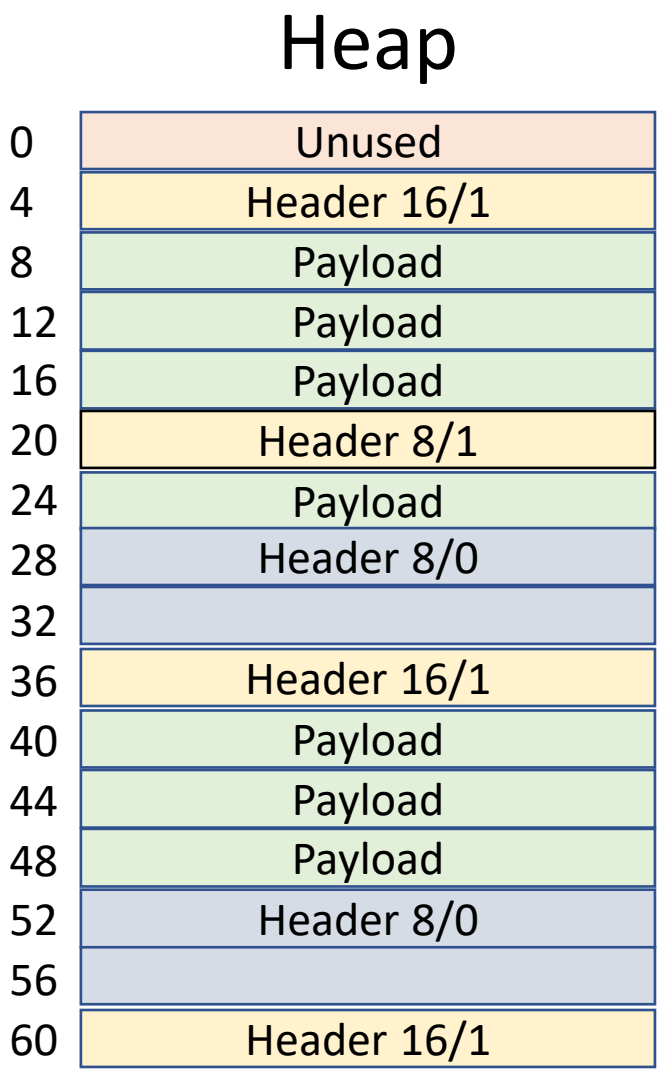
First Fit

Next Fit

Best Fit

Worst Fit

- Allocate from the largest block
- Doesn't leave small fragments behind
- "Maximizes the chance that the next allocation will fit"



Splitting

1. Tracking Free Blocks – Free List
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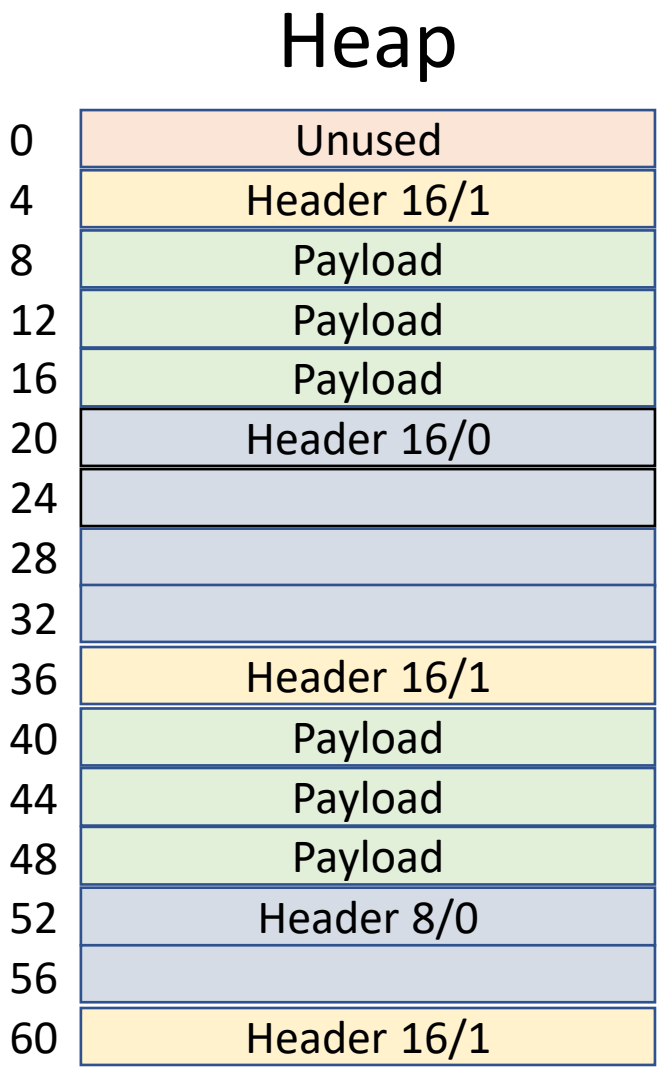
Allocate the entire block
Split the block



Coalescing

1. Tracking Free Blocks – Free List
2. Placement Policy
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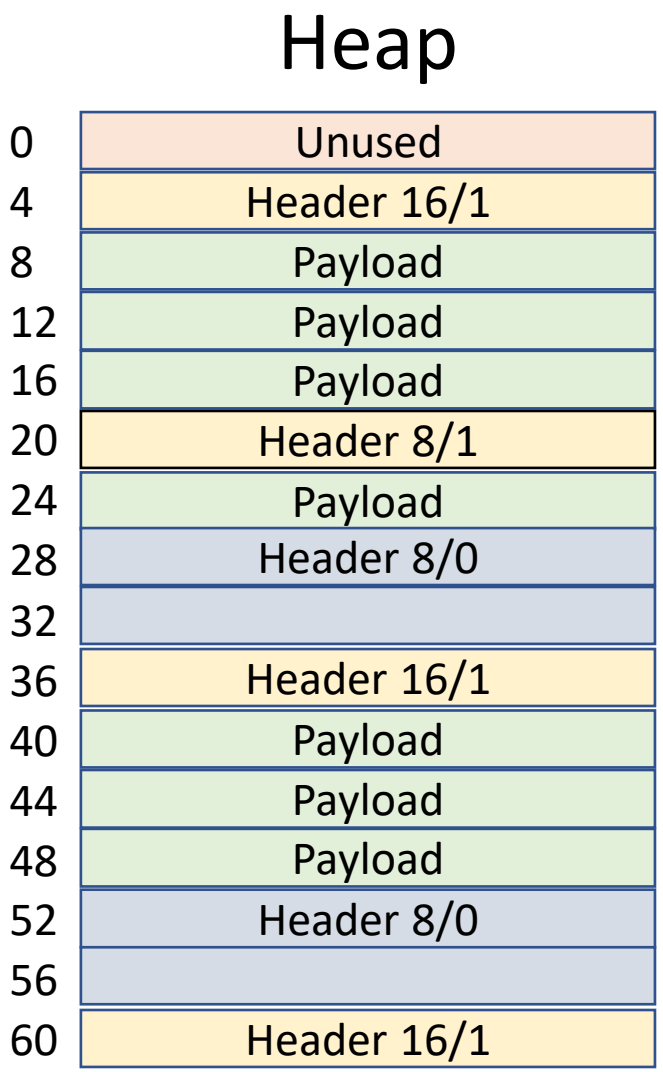
Immediate
Deferred



Coalescing

1. Tracking Free Blocks – Free List
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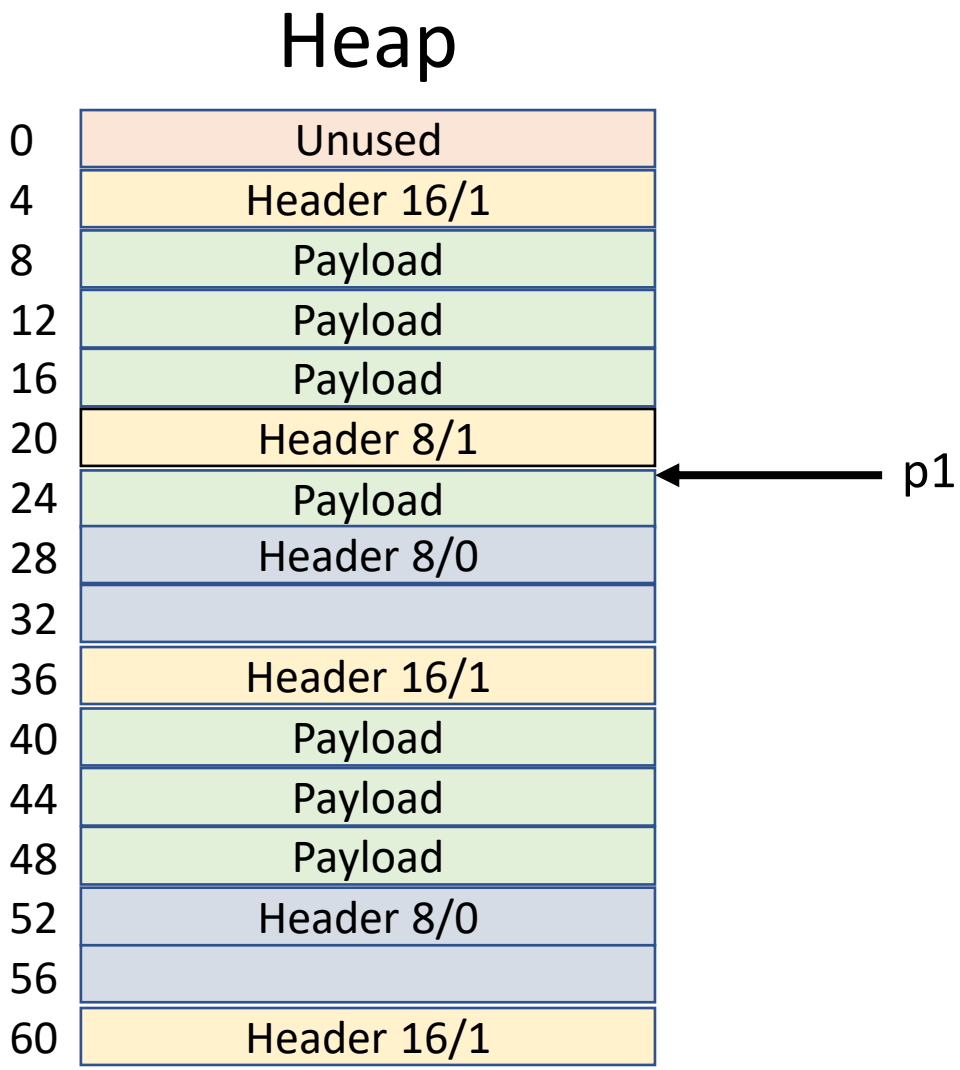
4. Coalescing

Example 1: Free block at 24



Coalescing

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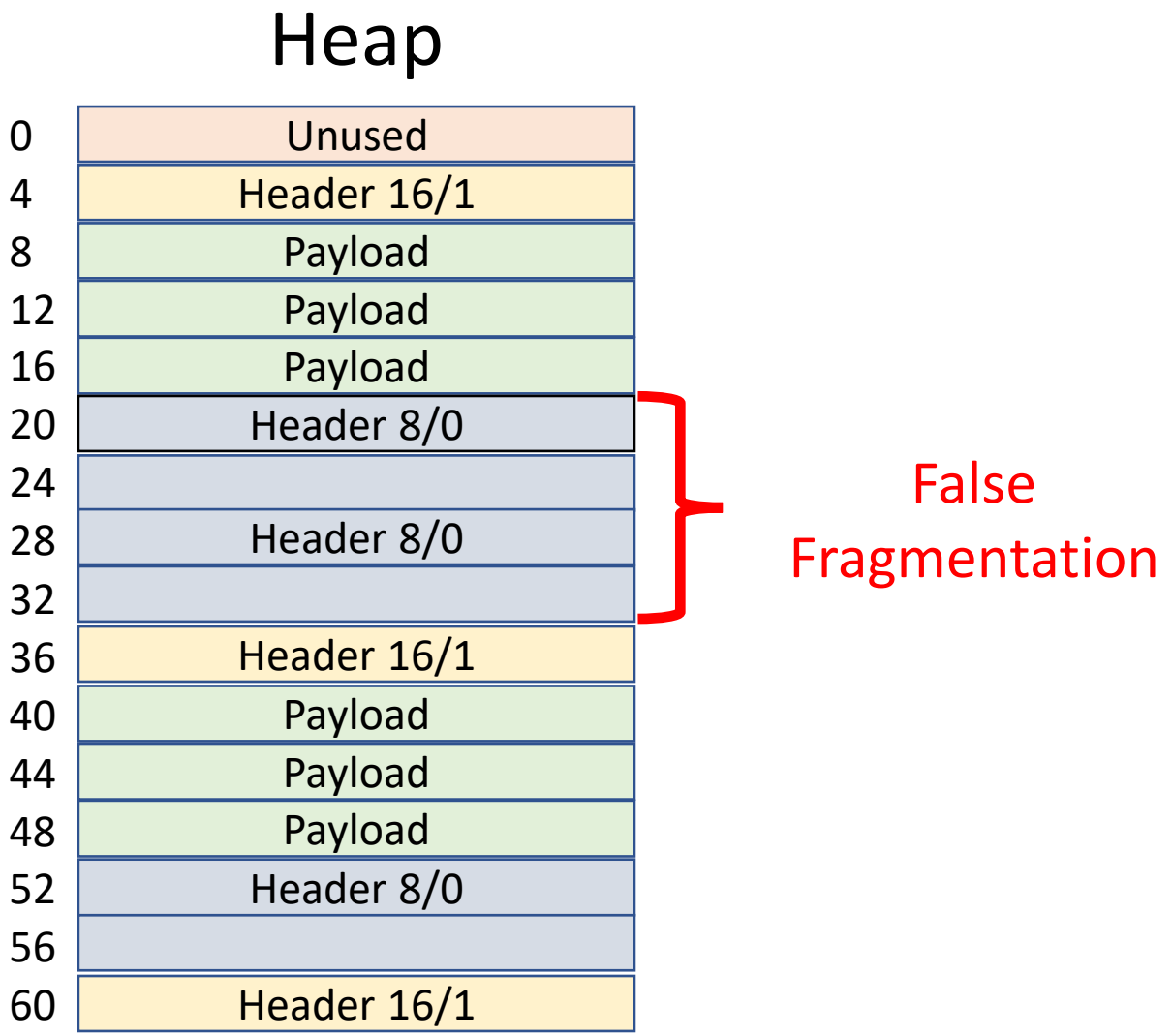

Coalescing

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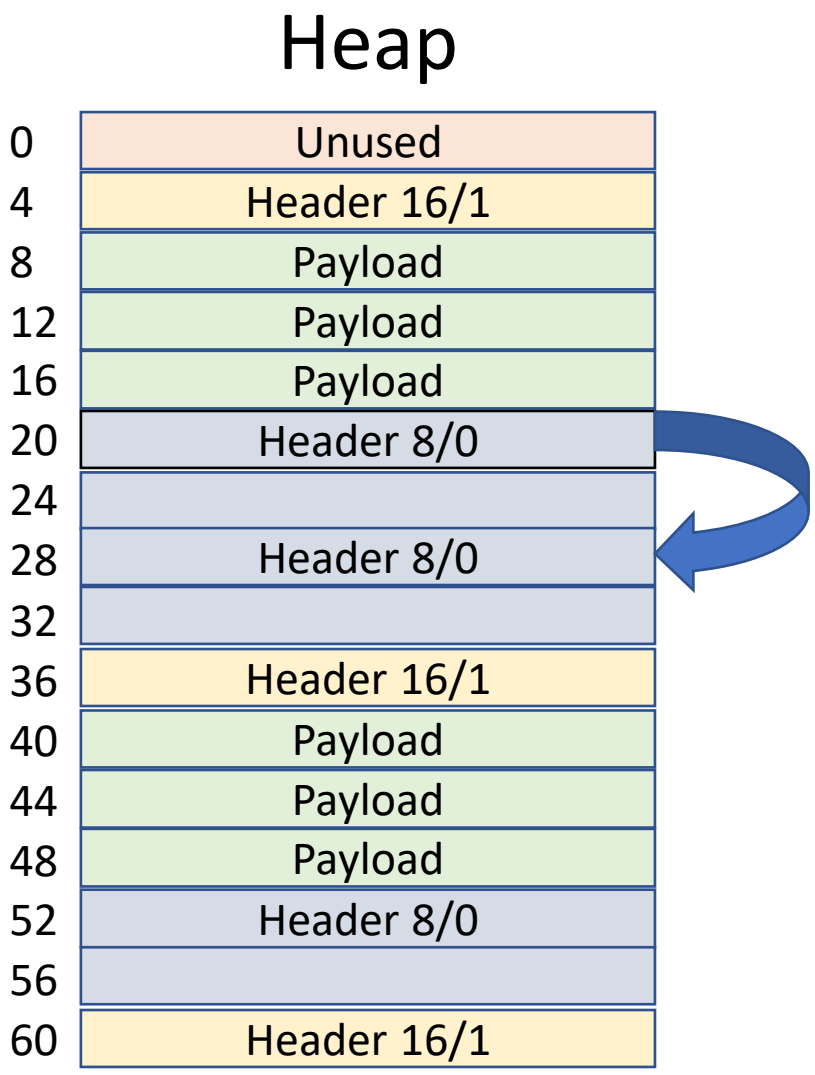

Coalescing

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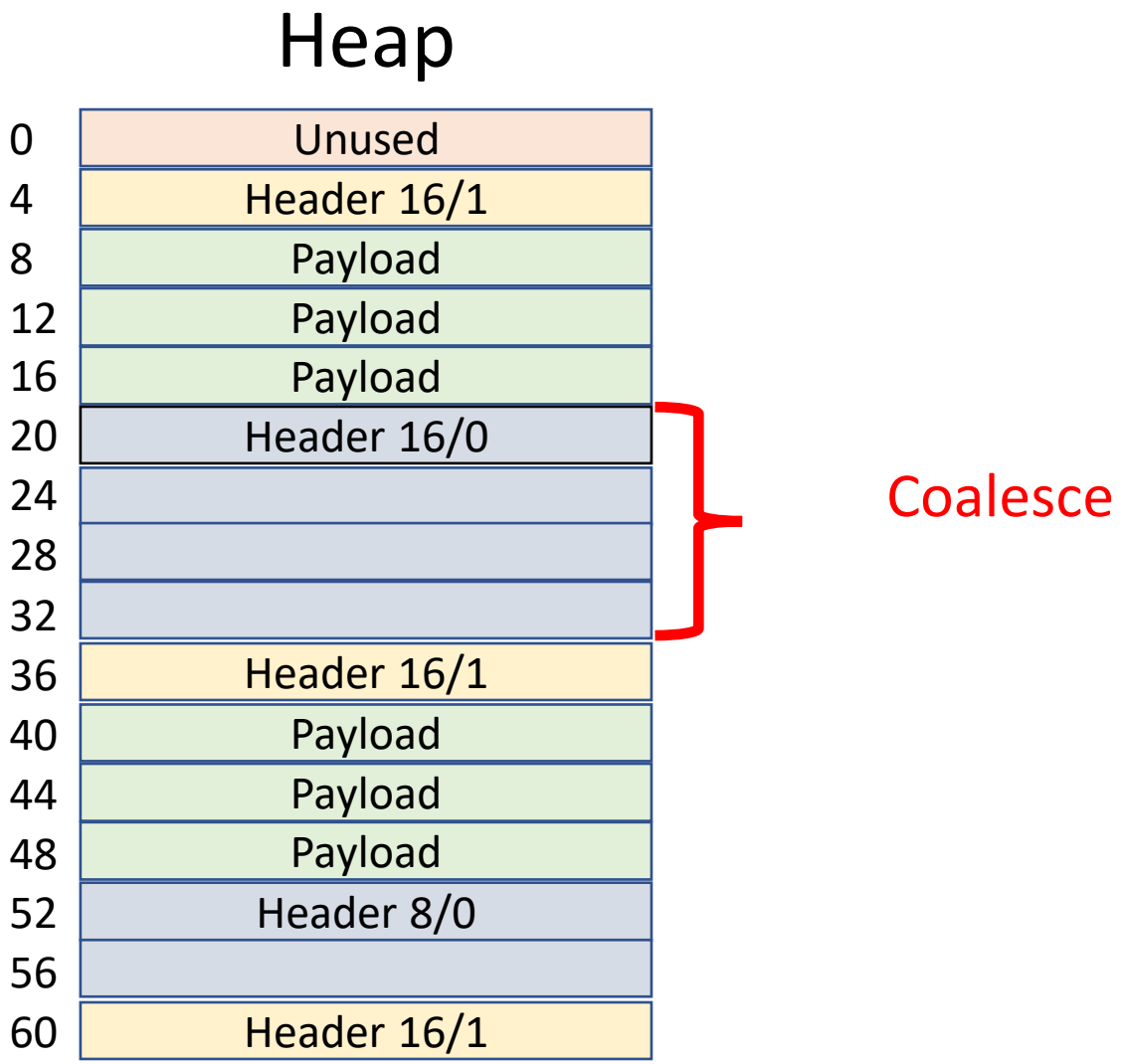
Example 1: Free block at 24

Immediate



Coalescing

1. Tracking Free Blocks – Free List
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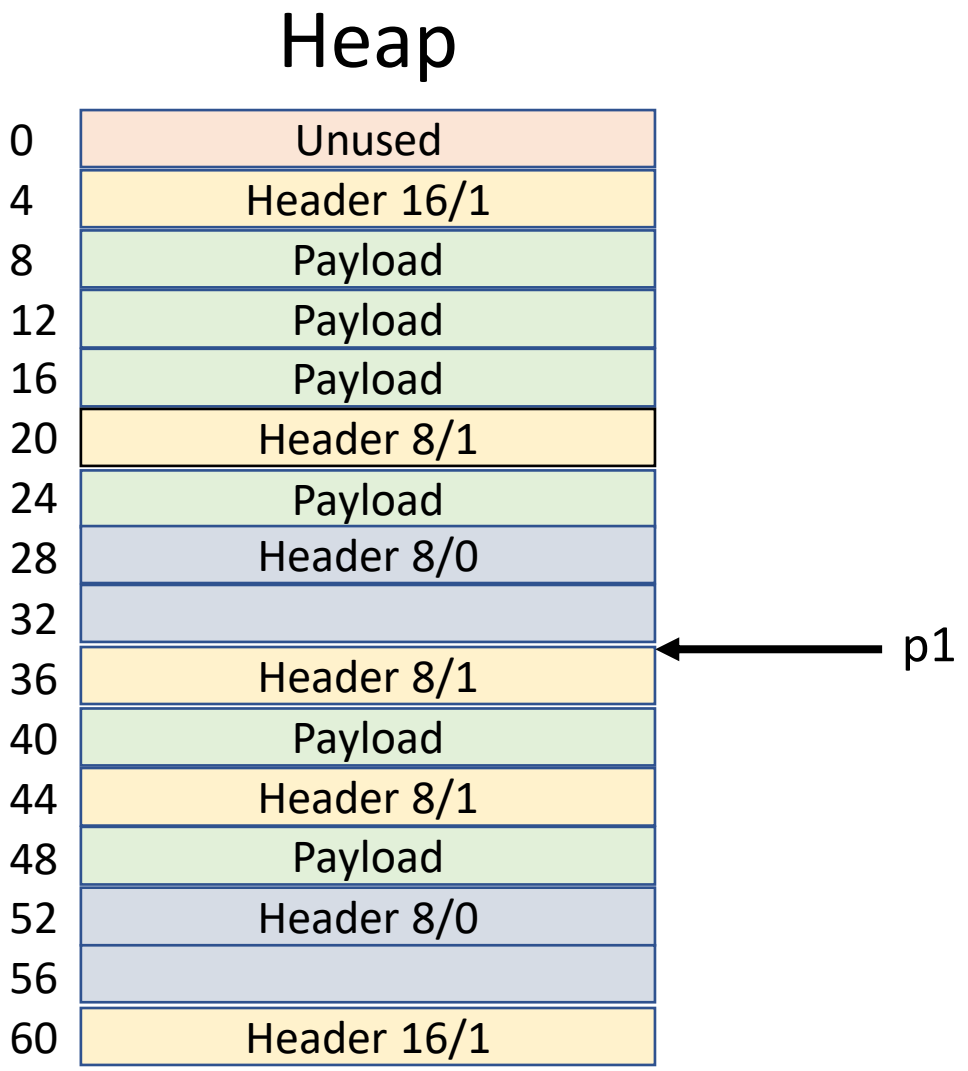
Example 2: Free block at 36



Coalescing

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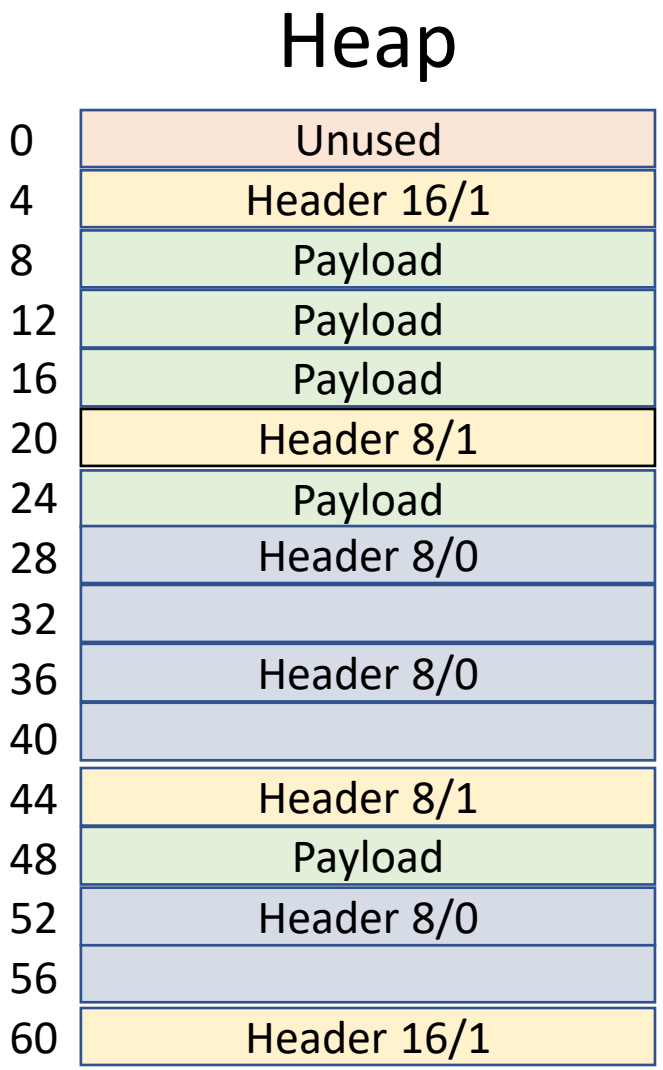
Example 2: Free block at 36



Coalescing

1. Tracking Free Blocks – Free List
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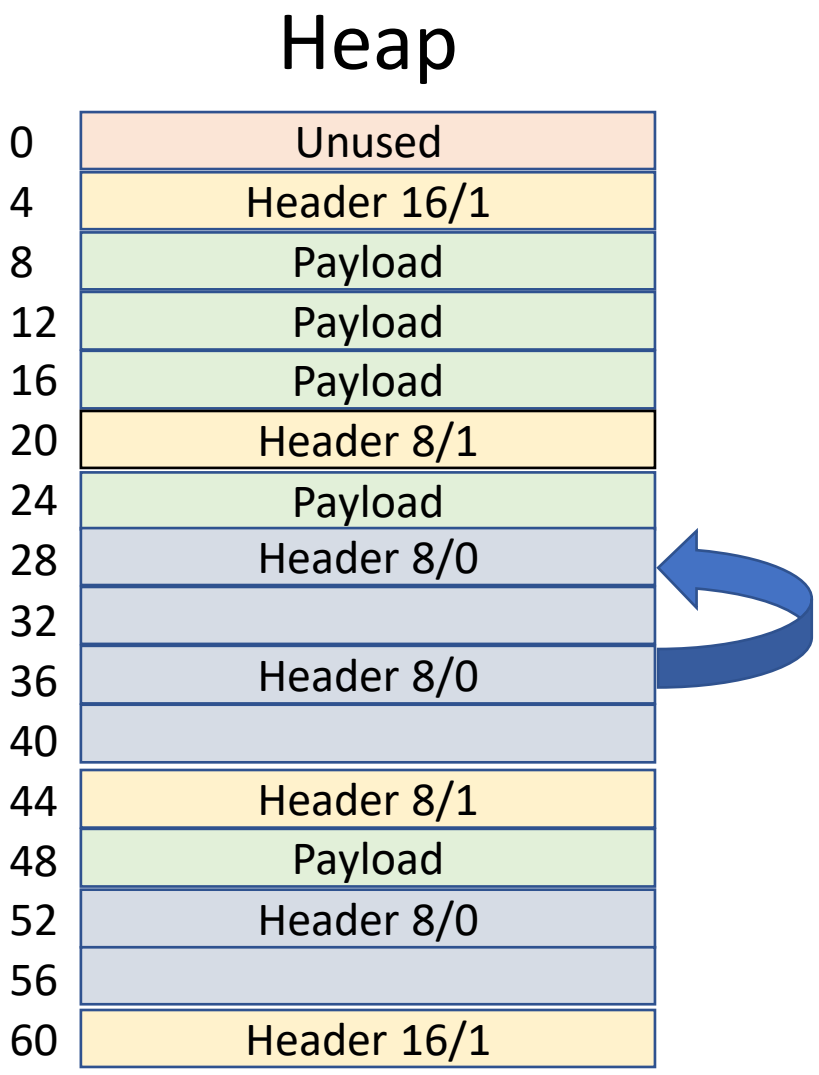
Example 2: Free block at 36
No pointer backward



Coalescing

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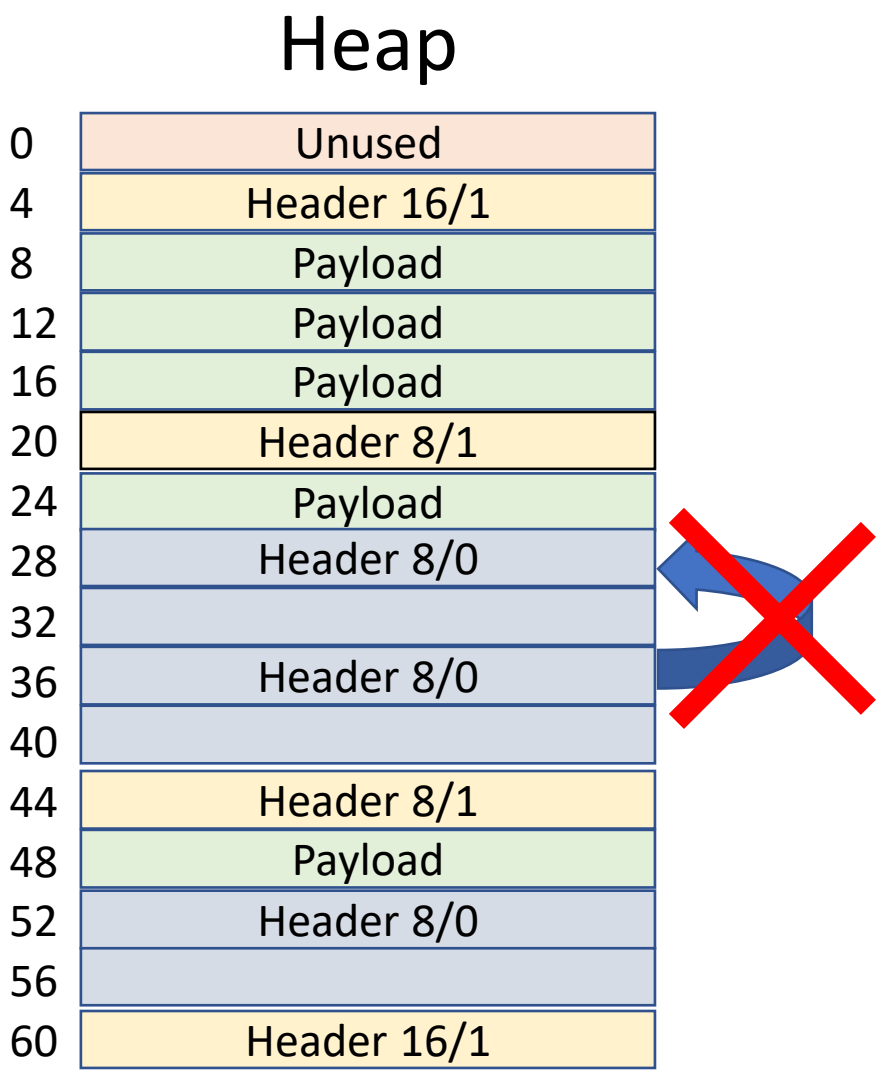
Example 2: Free block at 36

Need to traverse the entire list

Deferred



Coalescing

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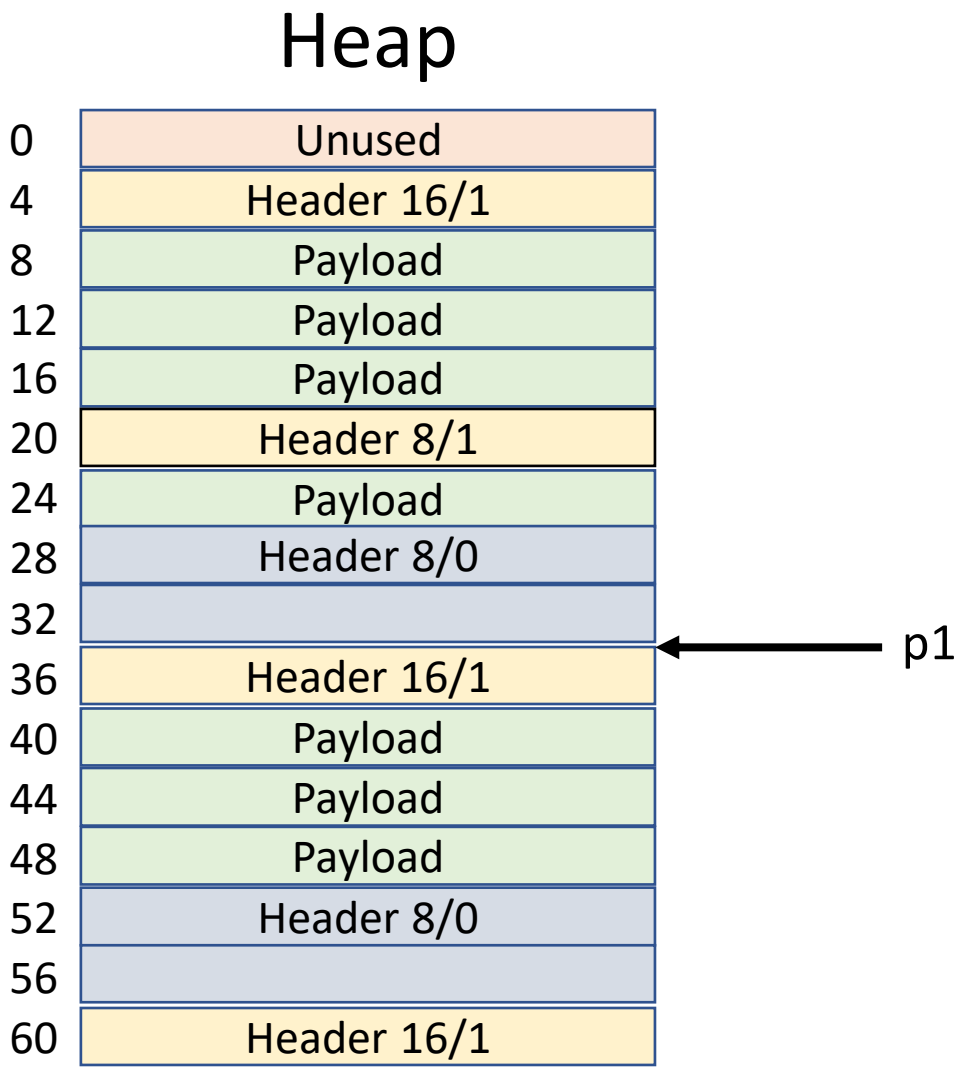
Example 3: Free block at 36



Coalescing

1. Tracking Free Blocks – Free List
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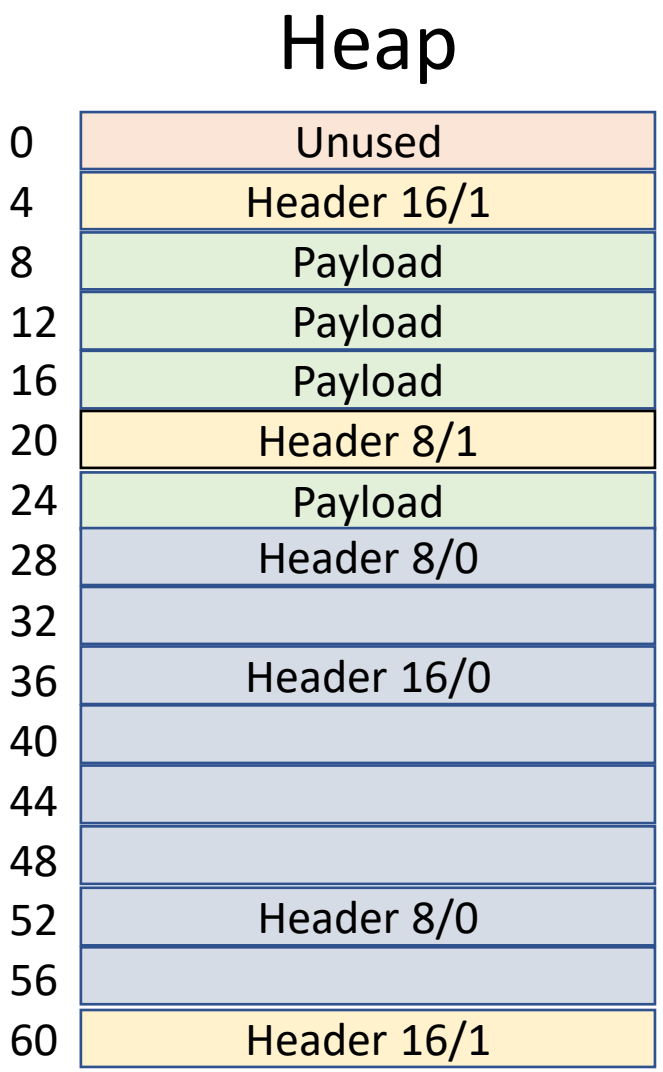
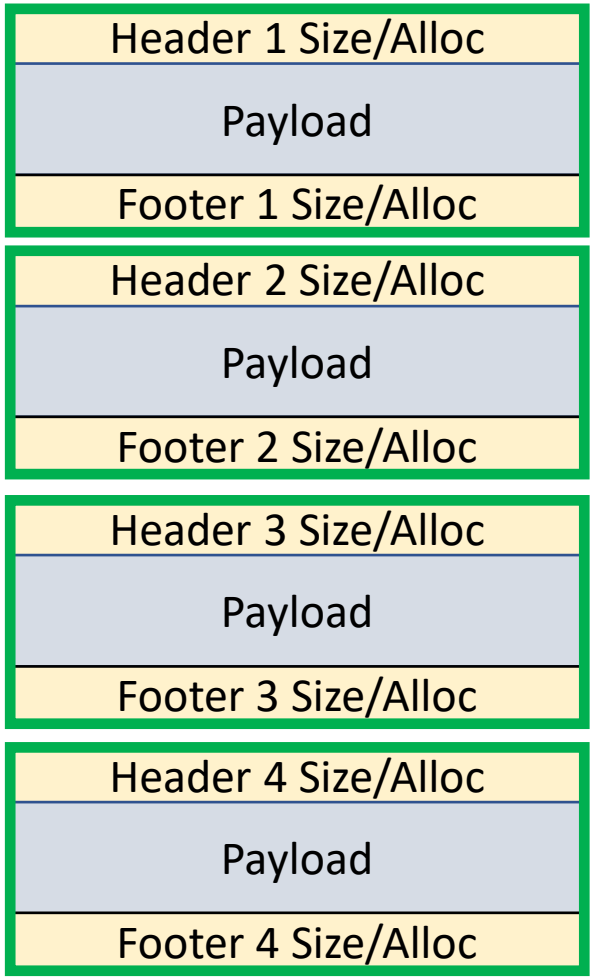
Example 3: Free block at 36
One free block ahead
One free block behind



Coalescing Strategies – Boundary Tags

Strategy: Duplicate the header as the last block



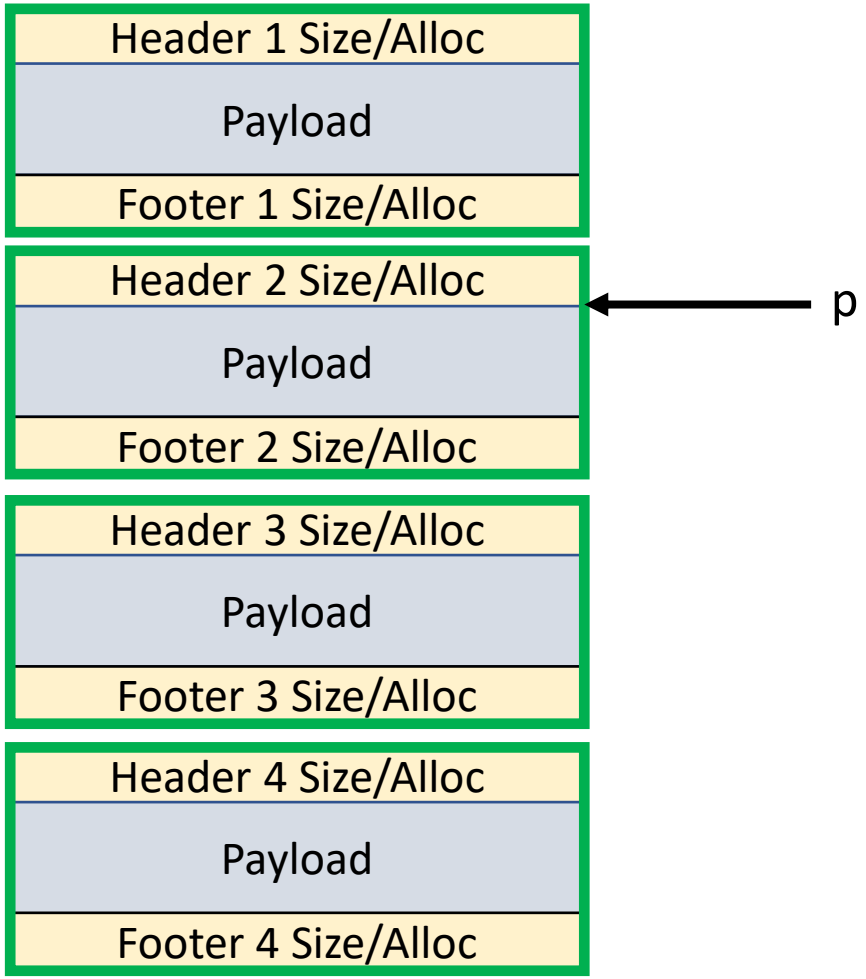
Coalescing Strategies – Boundary Tags

Strategy: Duplicate the header as the last block

free (p)

- byte before p holds the size of this block
- 2 bytes before holds the size/alloc of the previous block

- Decreases memory utilization – especially for applications that use many small blocks (linked lists)
- Fast coalescing of adjacent blocks



Coalescing Strategies – Improved Boundary Tags

- Only free blocks need footers
- Headers are multiple of 8 so we get three bits that must be 0 for the size
- We're already using b0 for the **alloc** bit
- Use b1 for the **previous alloc** bit

- Slightly more complicated
- Still constant time coalescing
- Better memory utilization