

# CS 354

# Machine Organization and Programming

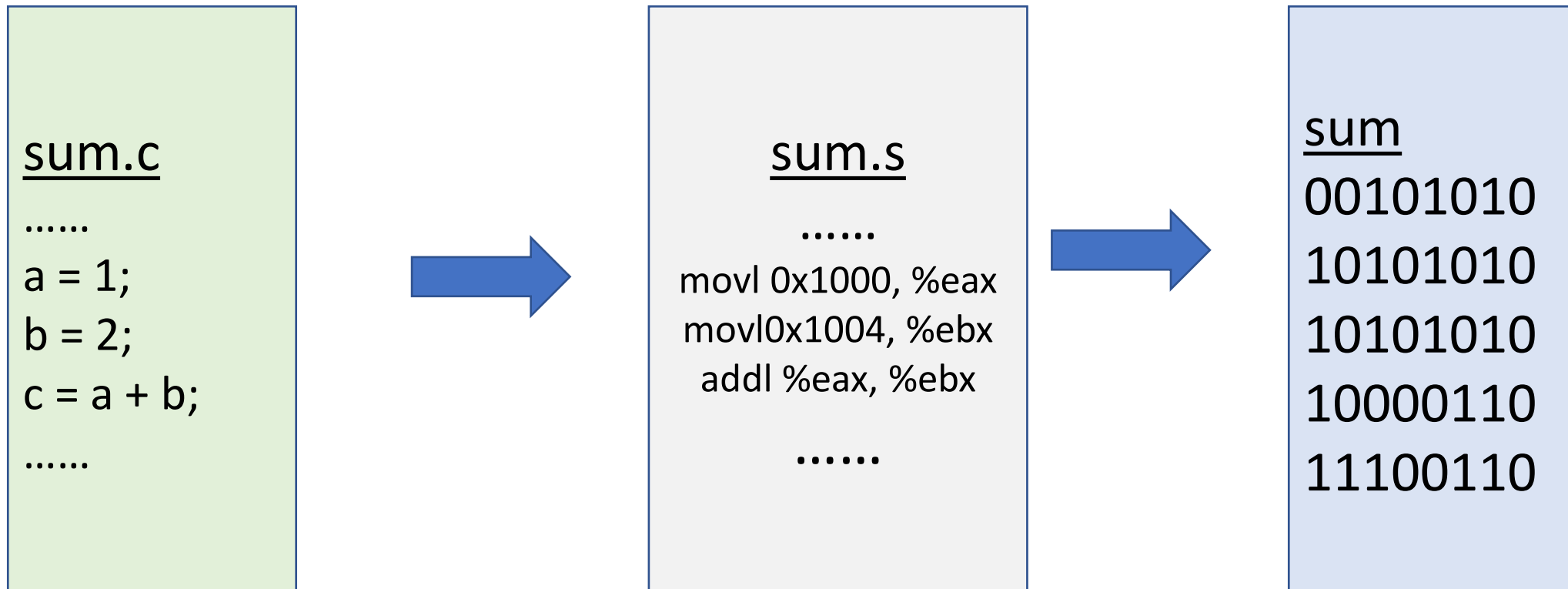
## Lecture 17

Michael Doescher  
Summer 2020

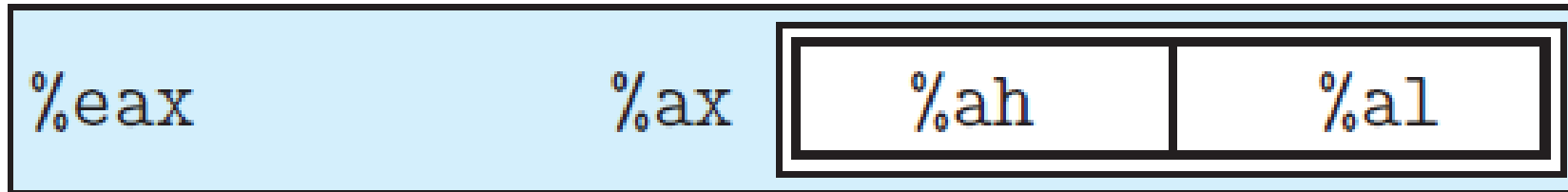
Assembly Languages  
Control Flow  
Conditionals

# What happens when we run a program?

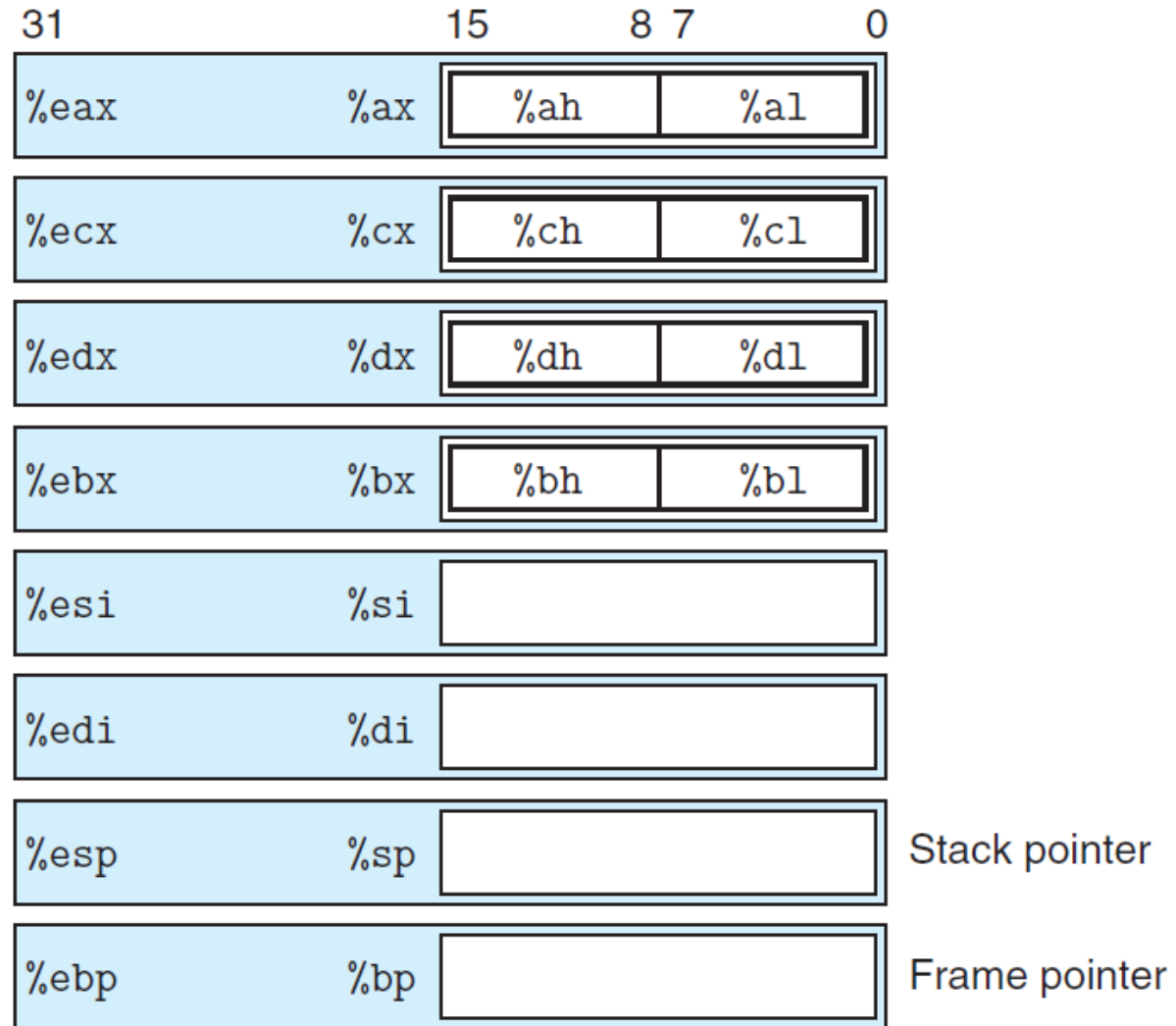
## Compiling



# Registers



# Registers



# Assembly Instructions : CS:APP 3.4

AT&T Syntax of x86

mov (movl, movb, movw)

Copies data from one location to another (source remains unchanged)

# Assembly Instructions : CS:APP 3.4

AT&T Syntax of x86

mov (movl, movb, movw)

Copies data from one location to another (source remains unchanged)

movl imm, register

movl imm, memory

mov register, register

mov register, memory

mov memory, register

# Assembly Instructions : CS:APP 3.4

AT&T Syntax of x86

mov (movl, movb, movw)

Copies data from one location to another (source remains unchanged)

movl imm, register

movl imm, memory

mov register, register

mov register, memory

mov memory, register

~~mov anything, immediate~~

~~mov memory, memory ::: movl 8(%eax), 0x7008~~

# Assembly Instructions : CS:APP 3.4

AT&T Syntax of x86

mov (movl, movb, movw)

Copies data from one location to another (source remains unchanged)

movl imm, register

movl imm, memory

mov register, register

mov register, memory

mov memory, register

~~mov anything, immediate~~

~~mov memory, memory ::: movl 8(%eax), 0x7008~~

General Form

Imm(%R1, %R2, Scale) ::: computes address as  $\text{Imm} + \%R1 + \%R2 * \text{scale}$



# Assembly Instructions : CS:APP 3.5

AT&T Syntax of x86

mov (movl, movb, movw)

lea : load effective address (& operator in C)

# Assembly Instructions : CS:APP 3.5

AT&T Syntax of x86

mov (movl, movb, movw)

lea : load effective address (& operator in C)

Arithmetic

add, sub, imul, idiv

(addl, addb, addw)

destination  $\leftarrow$  source – destination

integer division: dividend stored in %edx:%eax

remainder in %edx and quotient in %eax

# Assembly Instructions : CS:APP 3.5

AT&T Syntax of x86

mov (movl, movb, movw)

lea : load effective address (& operator in C)

Arithmetic

add, sub, imul, idiv

(addl, addb, addw)

destination  $\leftarrow$  source – destination

integer division: dividend stored in %edx:%eax

remainder in %edx and quotient in %eax

inc, dec

# Assembly Instructions : CS:APP 3.5

## AT&T Syntax of x86

mov (movl, movb, movw)

lea : load effective address (& operator in C)

## Arithmetic

add, sub, imul, idiv

(addl, addb, addw)

destination  $\leftarrow$  source – destination

integer division: dividend stored in %edx:%eax

remainder in %edx and quotient in %eax

inc, dec

## Bitwise operations

and, or, not, shifting

# Control Flow : CS:APP 3.6

- Sequential

```
Stmt 1;
```

```
Stmt 2;
```

```
Stmt 3;
```

```
Stmt 4;
```

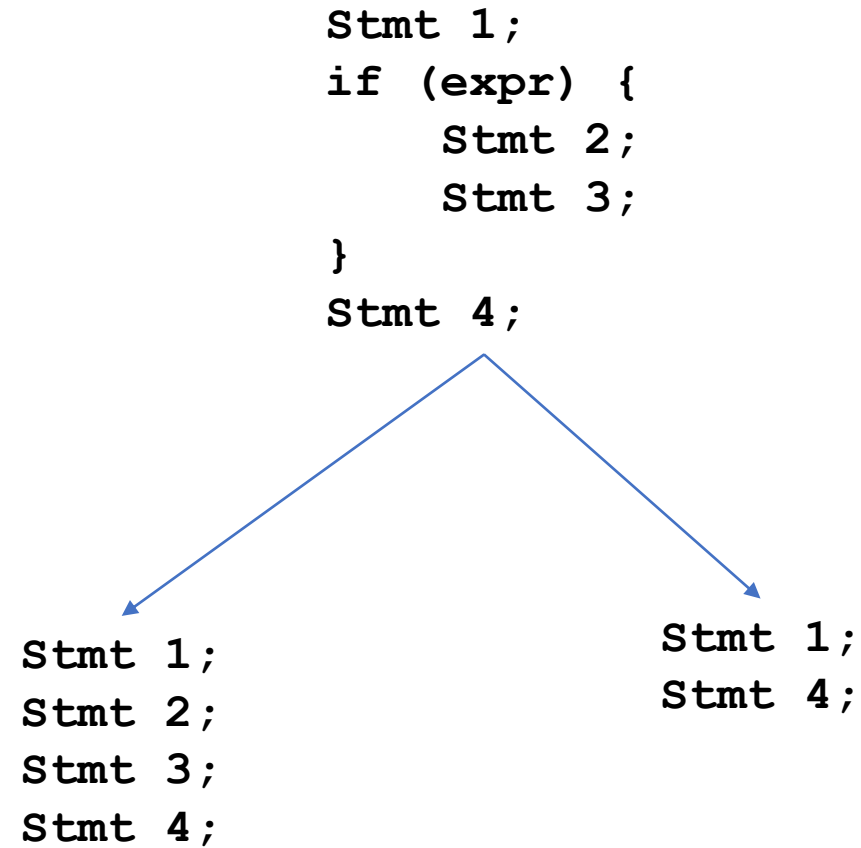
# Control Flow : CS:APP 3.6

- Sequential
- Conditional

```
Stmt 1;  
if (expr) {  
    Stmt 2;  
    Stmt 3;  
}  
Stmt 4;
```

# Control Flow : CS:APP 3.6

- Sequential
- Conditional



# Control Flow : CS:APP 3.6

- Sequential
- Conditional

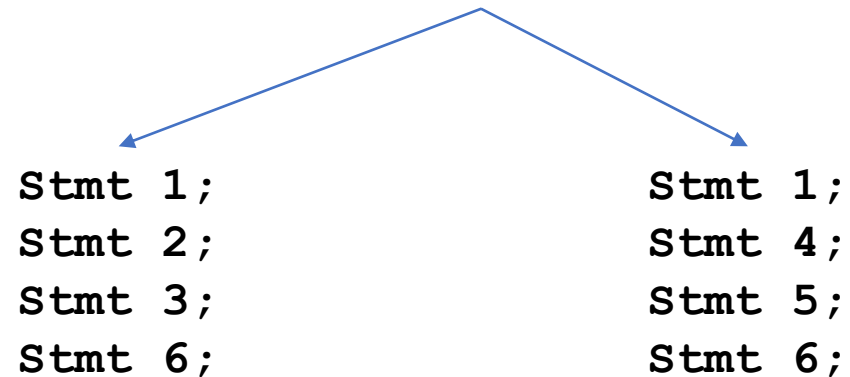
```
Stmt 1;  
if (expr) {  
    Stmt 2;  
    Stmt 3;  
}  
else {  
    Stmt 4;  
    Stmt 5;  
}  
Stmt 6;
```



# Control Flow : CS:APP 3.6

- Sequential
- Conditional

```
Stmt 1;  
if (expr) {  
    Stmt 2;  
    Stmt 3;  
}  
else {  
    Stmt 4;  
    Stmt 5;  
}  
Stmt 6;
```



# Control Flow : CS:APP 3.6

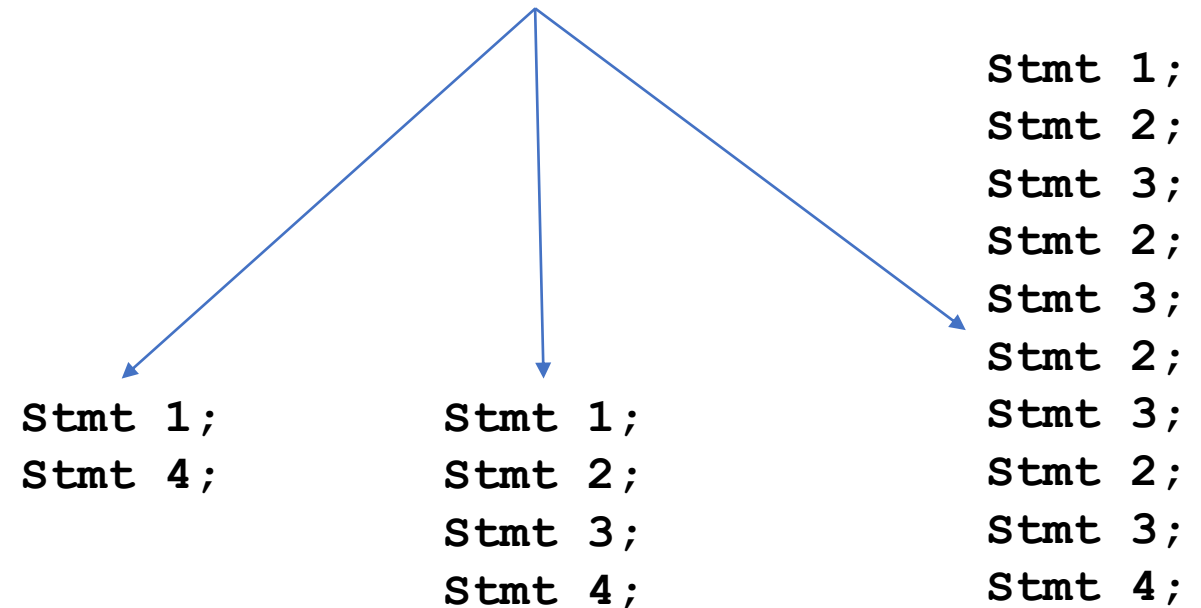
- Sequential
- Conditional
- Iteration

```
Stmt 1;  
while (expr) {  
    Stmt 2;  
    Stmt 3;  
}  
Stmt 4;
```

# Control Flow : CS:APP 3.6

- Sequential
- Conditional
- Iteration

```
Stmt 1;  
while (expr) {  
    Stmt 2;  
    Stmt 3;  
}  
Stmt 4;
```



# Control Flow : CS:APP 3.6

- Sequential
- Conditional
- Iteration
- Functions

# Control Flow : CS:APP 3.6

- Sequential
- Conditional
- Iteration
- Functions

What do we need at the assembly level to implement these?

# Control Flow : CS:APP 3.6

- Sequential
- Conditional
- Iteration
- Functions

What do we need at the assembly level to implement these?

Boolean Operators

<, <=, >, >=, !=, ==

# Control Flow : CS:APP 3.6

- Sequential
- Conditional
- Iteration
- Functions

What do we need at the assembly level to implement these?

Remember we have a bunch of registers, most importantly

- Instruction Pointer %eip
- Condition Code Register, %eflags

# Condition Codes: Addition

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly after every operation



# Condition Codes: Addition: ZF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly after every operation

```
T = a + b
```

```
T = 3 + (-3)
```

```
if (T == 0)
```

```
    ZF -> 1 (set)
```

```
else
```

```
    ZF -> 0 (not set)  
            (unset)
```

# Condition Codes: Addition: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} 3 \quad 011 \\ +4 \quad 100 \\ \hline 7 \quad 111 \end{array}$$

No carry out

Required

$$CF = 0$$

# Condition Codes: Addition : CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

**T = a + b**

3	011
+4	100
<hr/>	
7	111

**T = a + b**

3	011
+5	101
<hr/>	

**No carry out  
Required  
CF = 0**

# Condition Codes: Addition : CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

**T = a + b**

3	011
+4	100
<hr/>	
7	111

**T = a + b**

	1
3	011
+5	101
<hr/>	
	0

**No carry out  
Required  
CF = 0**

# Condition Codes: Addition : CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

**T = a + b**

3	011
+4	100
<hr/>	
7	111

**T = a + b**

	11
3	011
+5	101
<hr/>	
	00

**No carry out  
Required  
CF = 0**

# Condition Codes: Addition : CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

**T = a + b**

3	011
+4	100
<hr/>	
7	111

**T = a + b**

	111
3	011
+5	101
<hr/>	
	000

**No carry out  
Required  
CF = 0**

# Condition Codes: Addition : CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

**T = a + b**

$$\begin{array}{r} 3 \quad 011 \\ +4 \quad 100 \\ \hline 7 \quad 111 \end{array}$$

**No carry out  
Required  
CF = 0**

**T = a + b**

$$\begin{array}{r} \quad \quad 111 \\ 3 \quad \quad 011 \\ +5 \quad \quad 101 \\ \hline (1)000 \end{array}$$

**Carry out  
CF = 1**

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

$$T = a + b$$

$$\begin{array}{r} 1 \quad 001 \\ -2 \quad 010 \\ \hline \end{array}$$

Updated implicitly  
after every operation

Unsigned Numbers



# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} 1 \quad 001 \\ -2 \quad 010 \\ \hline 1 \end{array}$$

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} 1 \\ 1 \quad 001 \\ -2 \quad 010 \\ \hline 1 \end{array}$$

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} 02 \\ 1 \quad 001 \\ -2 \quad 010 \\ \hline 1 \end{array}$$

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} \phantom{0}012 \\ 1 \phantom{0}001 \\ -2 \phantom{0}010 \\ \hline \phantom{0}1 \end{array}$$

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} \phantom{0}012 \\ 1 \phantom{0}001 \\ -2 \phantom{0}010 \\ \hline \phantom{0}11 \end{array}$$

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} \phantom{0}012 \\ 1 \phantom{0}001 \\ -2 \phantom{0}010 \\ \hline \phantom{0}111 \end{array}$$

# Condition Codes: Subtraction: CF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Unsigned Numbers

$$T = a + b$$

$$\begin{array}{r} \phantom{0}012 \\ 1 \phantom{0}001 \\ -2 \phantom{0}010 \\ \hline 7 \phantom{0}111 \end{array}$$

**Carry Flag is set when we  
Borrow during subtraction  
Also.**

$$CF = 1$$

# Condition Codes: Addition: SF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

**T = a + b**

**if T < 0 SF = 1**  
**else SF = 0**

3	011
+ (-4)	100
<hr/>	

Signed Numbers



# Condition Codes: Addition: SF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

**T = a + b**

**if T < 0 SF = 1**  
**else SF = 0**

3	011
+ (-4)	100
<hr/>	
-1	111

Signed Numbers

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

**T = a + b**

$$\begin{array}{r} 2 \quad 010 \\ +1 \quad 001 \\ \hline \end{array}$$

**T = a + b**

$$\begin{array}{r} 3 \quad 011 \\ +1 \quad 001 \\ \hline \end{array}$$

Updated implicitly  
after every operation

Signed Numbers

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

**T = a + b**

$$\begin{array}{r} 2 \quad 010 \\ +1 \quad 001 \\ \hline 3 \quad 011 \end{array}$$

**T = a + b**

$$\begin{array}{r} 3 \quad 011 \\ +1 \quad 001 \\ \hline \end{array}$$

Updated implicitly  
after every operation

Signed Numbers

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

**T = a + b**

$$\begin{array}{r} 2 \quad 010 \\ +1 \quad 001 \\ \hline 3 \quad 011 \end{array}$$

**T = a + b**

$$\begin{array}{r} 3 \quad 011 \\ +1 \quad 001 \\ \hline -4 \quad 100 \end{array}$$

Updated implicitly  
after every operation

Signed Numbers

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

$$T = a + b$$

$$\begin{array}{r} 2 \quad 010 \\ +1 \quad 001 \\ \hline 3 \quad 011 \end{array}$$

$$OF = 0$$

$$T = a + b$$

$$\begin{array}{r} 3 \quad 011 \\ +1 \quad 001 \\ \hline -4 \quad 100 \end{array}$$

$$OF = 1$$

Signed Numbers

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Signed Numbers

$$T = a + b$$

$$\begin{array}{r} 2 \quad 010 \\ +1 \quad 001 \\ \hline 3 \quad 011 \end{array}$$

$$OF = 0$$

$$T = a + b$$

$$\begin{array}{r} 3 \quad 011 \\ +1 \quad 001 \\ \hline -4 \quad 100 \end{array}$$

$$OF = 1$$

-4, -3, -2, -1, 0, 1, 2, 3

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Signed Numbers

$$T = a + b$$

$$\begin{array}{r} 2 \quad 010 \\ +1 \quad 001 \\ \hline 3 \quad 011 \end{array}$$

$$OF = 0$$

$$T = a + b$$

$$\begin{array}{r} 3 \quad 011 \\ +1 \quad 001 \\ \hline -4 \quad 100 \end{array}$$

$$OF = 1$$

-4, -3, -2, -1, 0, 1, 2, 3

Any positive number + any negative  
number yield a number from this set

# Condition Codes: Addition: OF

%eflags

- ZF = Zero Flag
- CF = Carry Flag
- SF = Sign Flag
- OF = Overflow Flag

Updated implicitly  
after every operation

Signed Numbers

`T = a + b`

```
  2  010
+1  001
-----
  3  011
```

`OF = 0`

`T = a + b`

```
  3  011
+1  001
-----
 -4  100
```

`OF = 1`

`-4, -3, -2, -1, 0, 1, 2, 3`

```
if (a>0 && b>0 && t<0) OF = 1
or (a<0 && b<0 && t>0) OF = 1
else OF = 0
```



# COMPARE INSTRUCTION

`cmpl b,a`

Evaluates  $a-b$   
But does not store  
the result

%eflags

- ZF =
- CF =
- SF =
- OF =

# COMPARE INSTRUCTION

`cmpl b,a`

Evaluates  $a-b$   
But does not store  
the result

%eflags

- $ZF = a-b == 0$
- $CF =$
- $SF =$
- $OF =$

# COMPARE INSTRUCTION

`cmpl b,a`

Evaluates  $a-b$   
But does not store  
the result

%eflags

- $ZF = a-b == 0$
- $CF = a < b$
- $SF =$
- $OF =$

- CF is set when we borrow in during subtraction of unsigned numbers.
- Borrowing is required when we subtract a bigger number from a smaller number

# COMPARE INSTRUCTION

`cmpl b,a`

Evaluates  $a-b$   
But does not store  
the result

%eflags

- $ZF = a-b == 0$
- $CF = a-b < 0$  or  $a < b$
- $SF = a-b < 0$  or  $a < b$
- $OF =$

- SF similar to CF but for signed numbers
- Subtracting a bigger number from a smaller number results in a negative number
- Same as saying  $a < b$

# COMPARE INSTRUCTION

`cmpl b,a`

Evaluates  $a-b$   
But does not store  
the result

%eflags

- $ZF = a-b == 0$
- $CF = a-b < 0$  or  $a < b$
- $SF = a-b < 0$  or  $a < b$
- $OF =$   
 $(a > 0 \ \&\& \ b < 0 \ \&\& \ (a-b) < 0)$   
or  $(a < 0 \ \&\& \ b > 0 \ \&\& \ (a-b) > 0)$

# TEST INSTRUCTION

`testl b,a`

Evaluates `a&b`  
But does not store  
the result

Almost always used with  
the same operands

`testl a,a`

`%eflags`

- `ZF =`
- `CF =`
- `SF =`
- `OF =`

# TEST INSTRUCTION

`testl b,a`

Evaluates `a&b`  
But does not store  
the result

Almost always used with  
the same operands

`testl a,a`

`%eflags`

- `ZF = a&b == 0` (`a&a == 0 -> a == 0`)
- `CF =`
- `SF =`
- `OF =`

# TEST INSTRUCTION

`testl b,a`

Evaluates `a&b`  
But does not store  
the result

Almost always used with  
the same operands

`testl a,a`

`%eflags`

- `ZF = a&b == 0` (`a&a == 0 -> a == 0`)
- `CF =`
- `SF = a&b < 0` (`a&a -> a < 0`)
- `OF =`



# TEST INSTRUCTION

`testl b,a`

Evaluates `a&b`  
But does not store  
the result

`%eflags`

- `ZF = a&b == 0` (`a&a == 0 -> a == 0`)
- `CF =`
- `SF = a&b < 0` (`a&a -> a < 0`)
- `OF =`

<code>If %eax == 0</code>	<code> </code>	<code>testl %eax, %eax</code>	<code> </code>	<code>ZF = 1, SF = 0</code>
<code>If %eax &lt; 0</code>	<code> </code>	<code>testl %eax, %eax</code>	<code> </code>	<code>ZF = 0, SF = 1</code>
<code>If %eax &gt; 0</code>	<code> </code>	<code>testl %eax, %eax</code>	<code> </code>	<code>ZF = 0, SF = 0</code>

# JUMP INSTRUCTION

```
Instr 1
jmp Target      // unconditional jump
Instr 2
Instr 3
Target:         // a label
Instr 4
Instr 5
```

# Conditional JUMP INSTRUCTIONS

```
Instr 1          // set the flags
jz Target        // jump if ZF
Instr 2          // if block
Instr 3
Target:          // label no else block
Instr 4
Instr 5
```

# Conditional JUMP INSTRUCTIONS

```
Instr 1          // set the flags
jz Target        // jump if ZF
Instr 2          // if block
Instr 3
Target:          // label no else block
Instr 4
Instr 5
```

We also have jump instructions

jz, jnz (jump if zero, jump if not zero)

j1, jle, jg, jge, je, jne (<, <=, >, >=, ==, !=)

jb, jbe, ja, jae (above, below - for unsigned numbers)

# JUMP CONDITIONS

jmp No Flag Requirements

je ZF

jne  $\sim$ ZF

j1 SF^OF

jle SF^OF | ZF

jg  $\sim$  (SF^OF)

jge  $\sim$  (SF^OF) | ZF

jb CF

jbe CF | ZF

ja  $\sim$ CF &  $\sim$ ZF

jae  $\sim$ CF

# Worksheet Problem 6

Assume the value of a is in %eax, and the value of b is in %ebx

Write x86 assembly code for:

```
if (a>b) {  
    a++;  
}
```

# Worksheet Problem 6

Assume the value of a is in %eax, and the value of b is in %ebx

Write x86 assembly code for:

```
if (a>b) {  
    a++;  
}
```

Condition

Jump

DO:

DONT:

# Worksheet Problem 6

Assume the value of a is in %eax, and the value of b is in %ebx

Write x86 assembly code for:

```
if (a>b) {  
    a++;  
}
```

```
    cmpl %ebx, %eax  
    jle DONT  
DO:  
    addl $1, %eax  
DONT:
```



# Worksheet Problem 6

Assume the value of a is in %eax, and the value of b is in %ebx

Write x86 assembly code for:

```
if (a>b) {  
    a++;  
}
```

```
    cmpl %ebx, %eax  
    jle DONT  
    addl $1, %eax  
DONT:
```

# Worksheet Problem 7

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
if (a>b) {  
    a++;  
} else {  
    b = a;  
}
```

# Worksheet Problem 7

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
if (a>b) {  
    a++;  
} else {  
    b = a;  
}
```

```
    condition  
    jump  
DO:
```

```
DONT:
```

```
END:
```

# Worksheet Problem 7

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
if (a>b) {  
    a++;  
} else {  
    b = a;  
}  
  
    cmpl %ebx, %eax  
    jle DONT  
DO:  
    addl $1, %eax  
    jmp END  
DONT:  
    movl %eax, %ebx  
END:
```

# Worksheet Problem 7

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
if (a>b) {  
    a++;  
} else {  
    b = a;  
}  
  
    cmpl %ebx, %eax  
    jle DONT  
    addl $1, %eax  
    jmp END  
DONT:  
    movl %eax, %ebx  
END:
```

# Worksheet Problem 8

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
while (b > 0) {  
    a++;  
    b--;  
}
```

# Worksheet Problem 8

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
while (b > 0) {  
    a++;  
    b--;  
}
```

```
TOP:  
    condition  
    jump BOTTOM  
    statements  
    jump TOP  
BOTTOM:
```

# Worksheet Problem 8

Assume the value of a is in %eax, and the value of b is in %ebx  
Write x86 assembly code for:

```
while (b > 0) {  
    a++;  
    b--;  
}
```

```
TOP:  
    testl %ebx, %ebx  
    jle BOTTOM  
    incl %eax  
    decl %ebx  
    jmp TOP  
BOTTOM:
```