



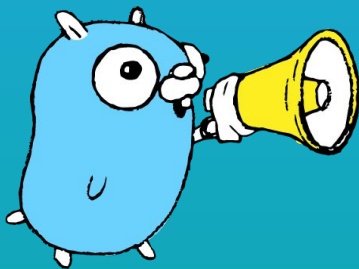
Taiwan, Taipei 2021

# The life of heap memory



---

Gaston Qiu



## Today's Agenda

Background Knowledge

---

TCMalloc

---

Garbage collection

---

---

---

---

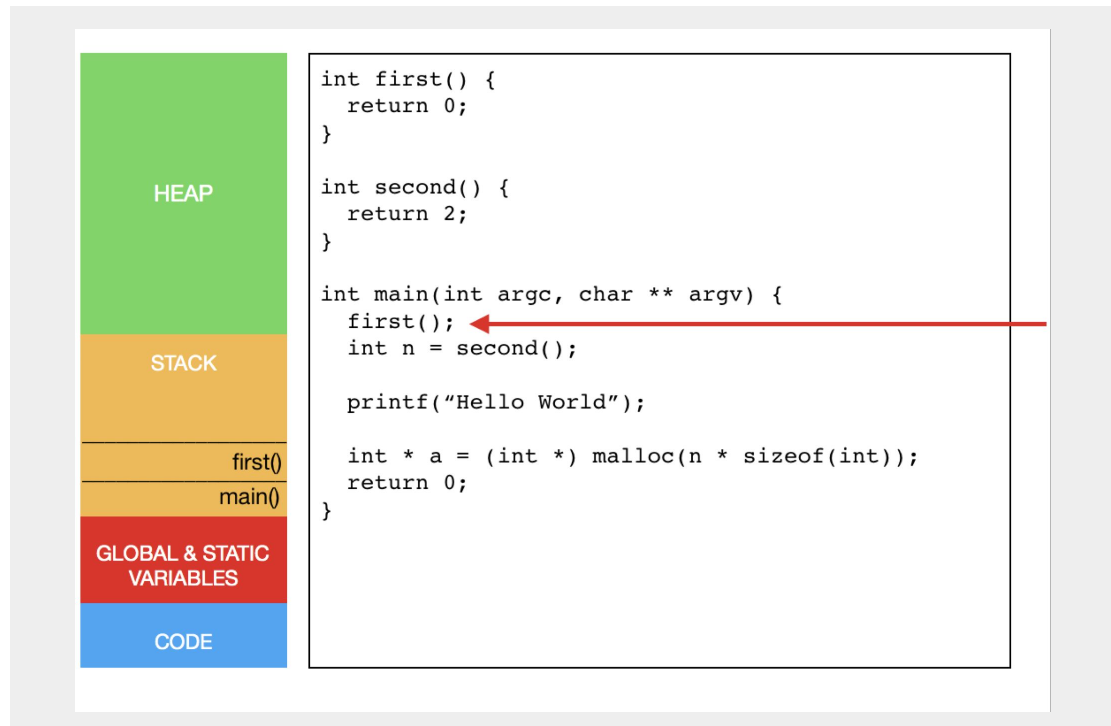
---

## SECTION ONE

---

# Background Knowledge

# Heap and Stack Example



Stack is fixed size linear data struct and only handle the memory will be release after the function finished.

# Heap and Stack Example



```
int first() {  
    return 0;  
}  
  
int second() {  
    return 2;  
}  
  
int main(int argc, char ** argv) {  
    first();  
    int n = second();  
  
    printf("Hello World");  
  
    int * a = (int *) malloc(n * sizeof(int));  
    return 0;  
}
```

Stack is fixed size linear data struct and only handle the memory will be release after the function finished.

# Heap and Stack Example



`malloc(n*sizeof(int))`

HEAP

STACK

`int * a`

`int n`

`main()`

GLOBAL & STATIC  
VARIABLES

CODE

```
int first() {  
    return 0;  
}  
  
int second() {  
    return 2;  
}  
  
int main(int argc, char ** argv) {  
    first();  
    int n = second();  
  
    printf("Hello World");  
  
    int * a = (int *) malloc(n * sizeof(int));  
    return 0;  
}
```



Heap is the runtime determinate hierarchical data structure which we need to manage its life cycle.

# Heap and Stack Example



```
int first() {  
    return 0;  
}  
  
int second() {  
    return 2;  
}  
  
int main(int argc, char ** argv) {  
    first();  
    int n = second();  
  
    printf("Hello World");  
  
    int * a = (int *) malloc(n * sizeof(int));  
    return 0;  
}
```

The memory allocation on heap is slow and handle its life cycle is difficult and tricky. So there is a lots of things we can optimize.

# Escape Analysis



## New an object on heap

```
func main() {  
    num := getRandom()  
    println(*num)  
}  
  
//go:noinline  
func getRandom() *int {  
    tmp := rand.Intn(100)  
    return &tmp  
}
```

Go decide the object should go to heap or stack through the escape analysis.



# Garbage collection (Mark and Sweep)



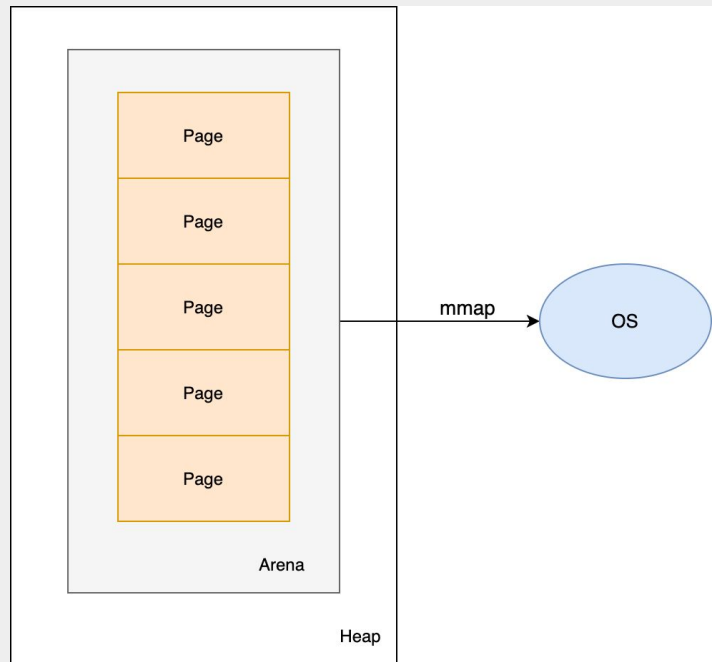
## Free an allocated object

```
func main() {  
    num := getRandom()  
    println(*num)  
    runtime.GC() // trigger gc  
}  
  
//go:noinline  
func getRandom() *int {  
    tmp := rand.Intn(100)  
    return &tmp  
}
```

The mechanism that decide the object should be released is called garbage collection. (In really life it will trigger automatically)

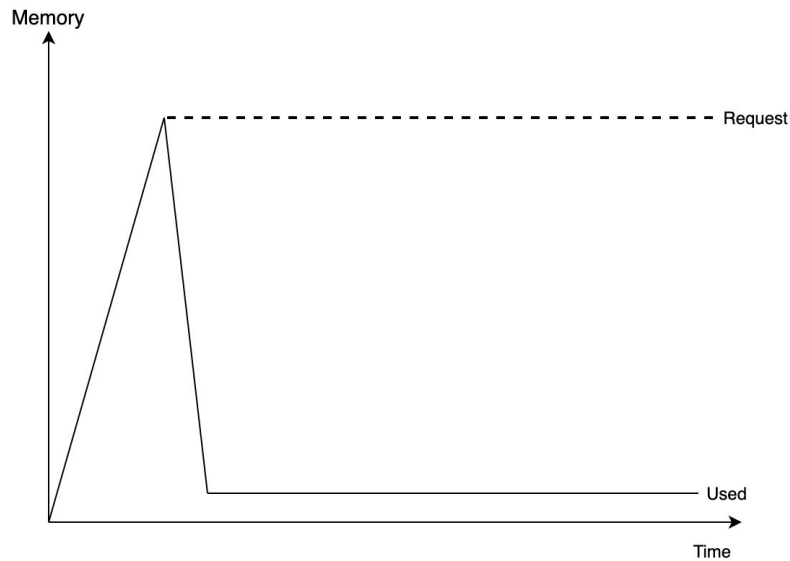
# Heap struct in go

## New an object on heap



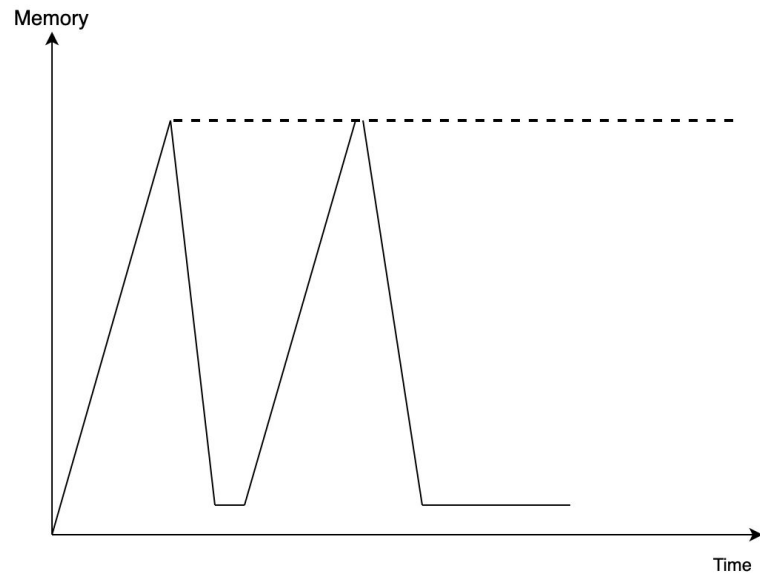
System calls (mmap) are very expensive. We pre-allocate an arena of memory.

# Since we hold the memory, we need to release it.



# Since we hold the memory, we need to release it.

---



# Garbage collection (Scavenging)



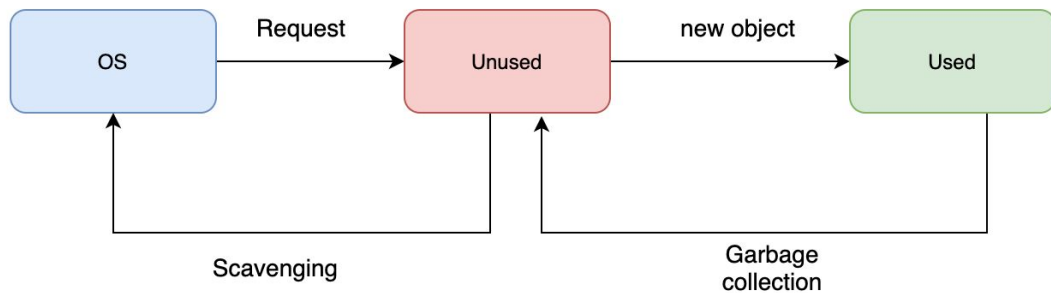
Release the memory to OS

```
func bgscavenge(c chan int)

func (p *pageAlloc) scavenge(
    nbytes uintptr,
    mayUnlock bool
) uintptr
```

After a memory is marked as unused. We should release it OS.

# Memory life cycle

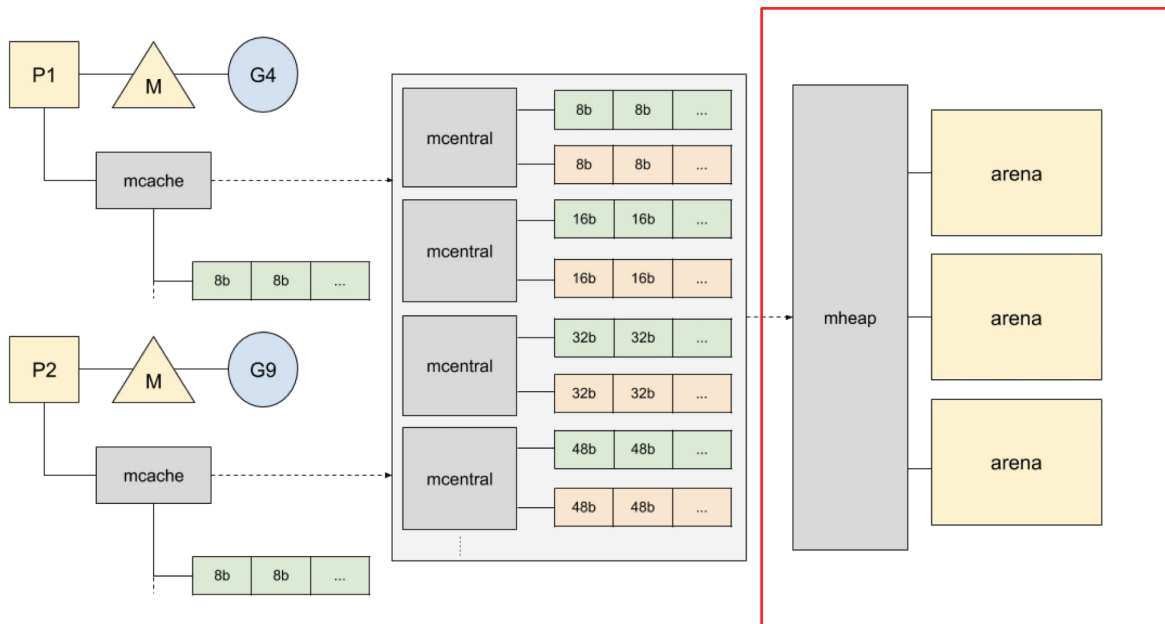


## SECTION TWO

---

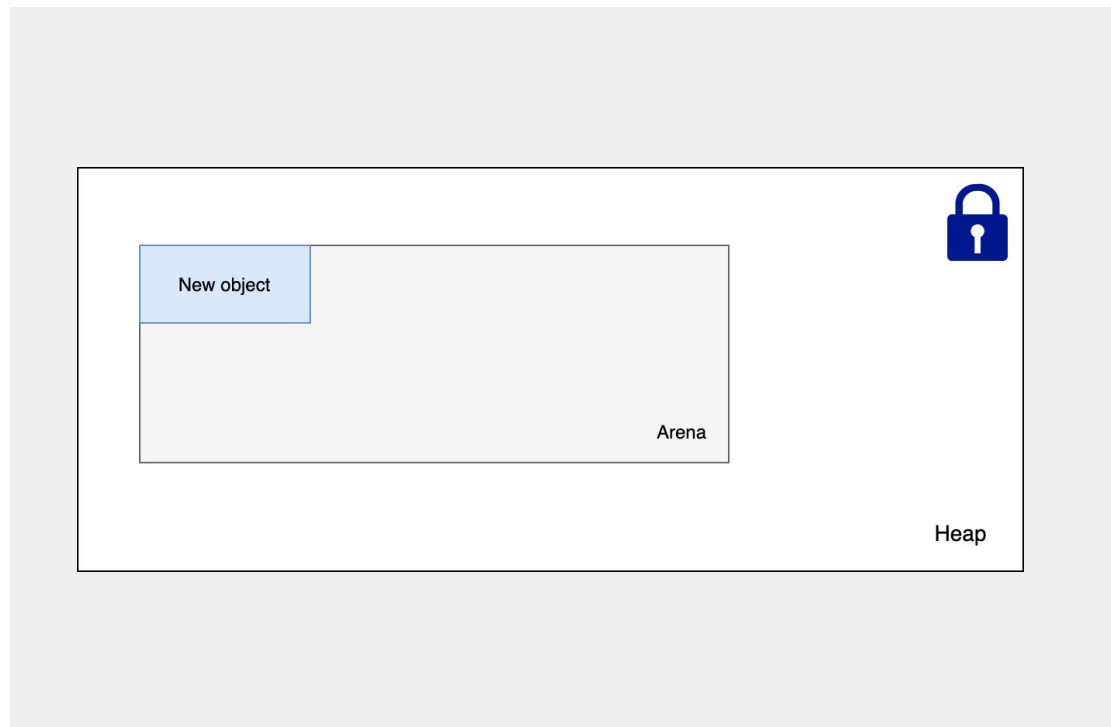
# Thread-cache Malloc (TCMalloc)

# Large allocation



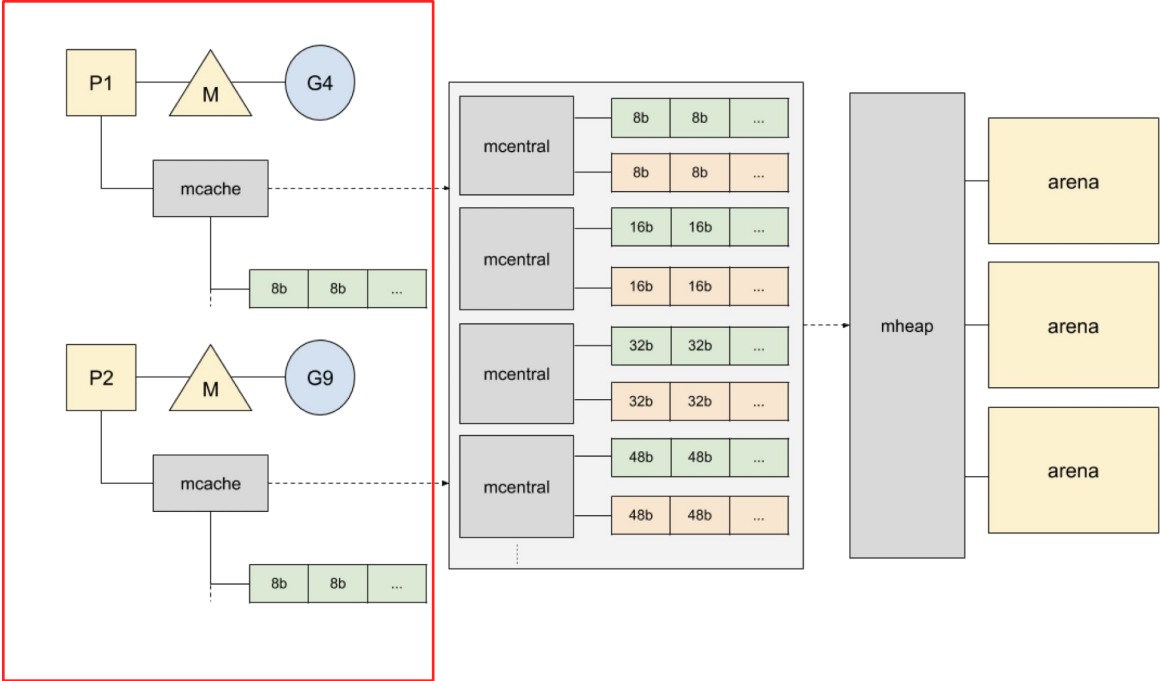


# Large allocation

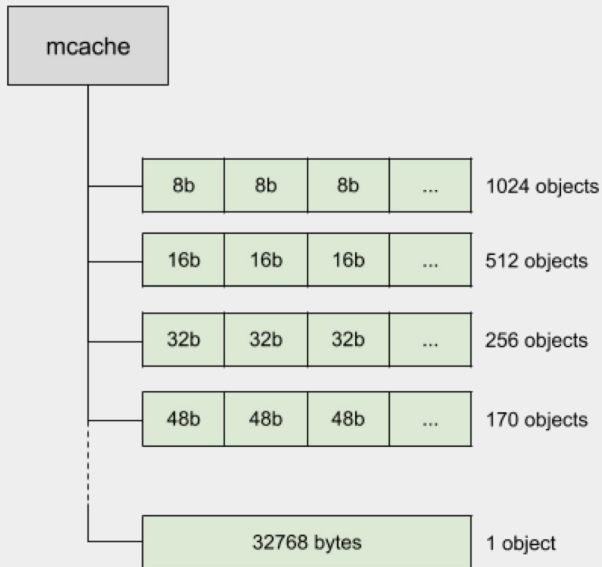


- Object size > 32KB
- Directly allocate on arena
- Lock required

# Small allocation



# Small allocation



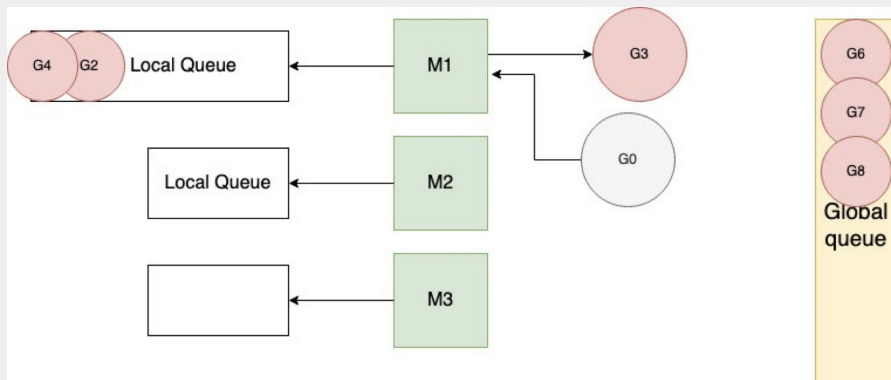
- Smallest object allocation (  $\leq$  32KB)

- There is about 70 different classes of span from 8 bytes to 32KB.

- Round up to the class size

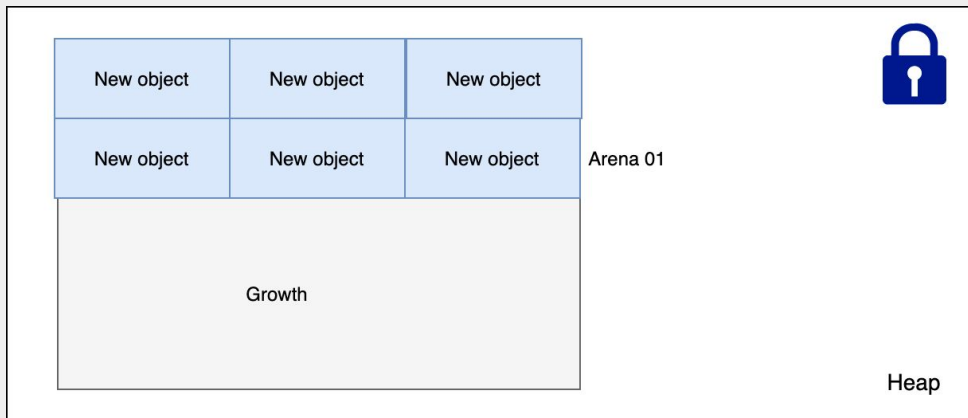
- No lock required in the small allocation.

# Mcache benefit



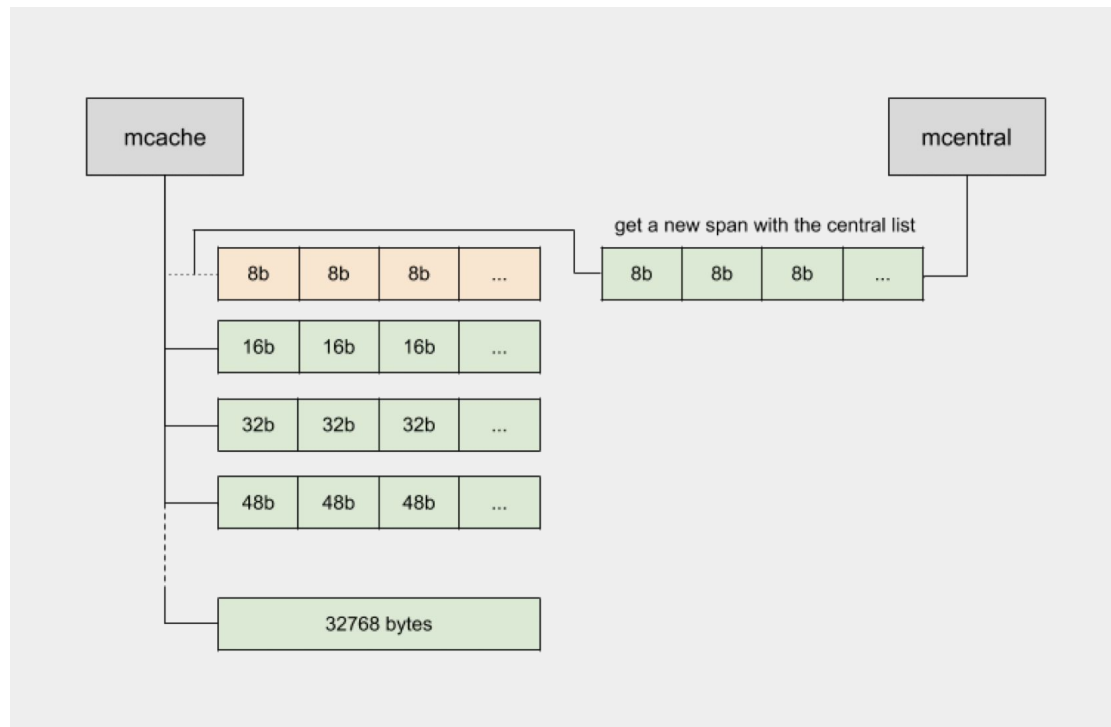
- Assuming most allocation is small object
- The mcache design align with golang m:n scheduler
- Better locality
- Less fragmentation

# The arena has no space (large allocation)

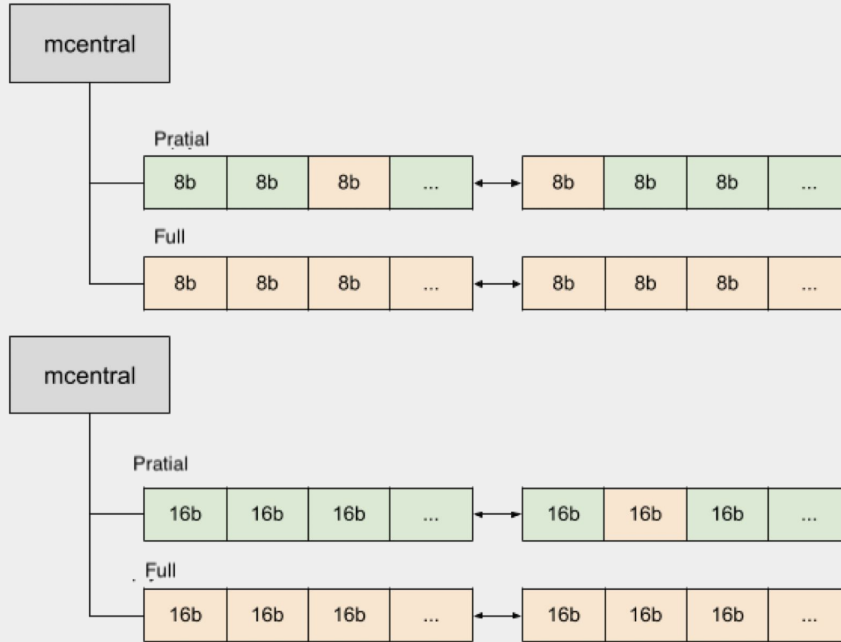


- Ask more pages from OS.

# The mcache has no space (small allocation)



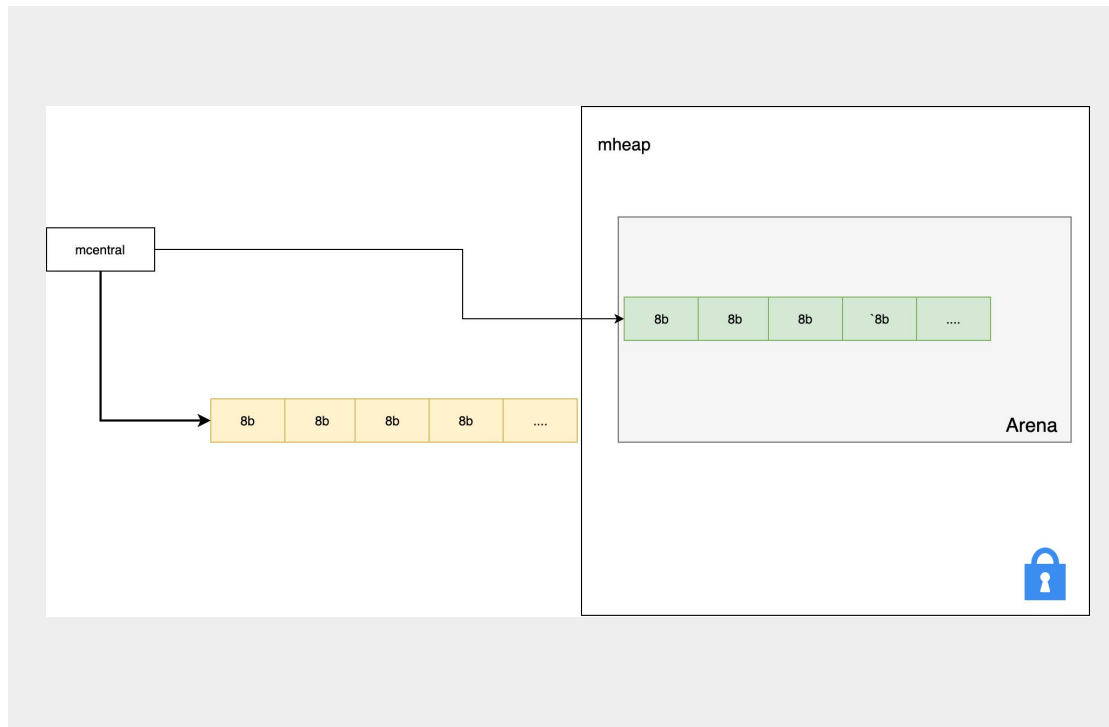
- Ask a new span from mcentral.
- All the mcache's share the same mcentral.
- The lock is required when ask a new span from mcentral.



- Partial: Partial of the span is in-used.

- Full: The span is full.

# The central list has no space (small allocation)



- When mcentral has no space, we can allocate a new span on heap.



# Request more memory

---



## SECTION Three

---

# Garbage collection

# Mark and Sweep

## Trigger garbage collection

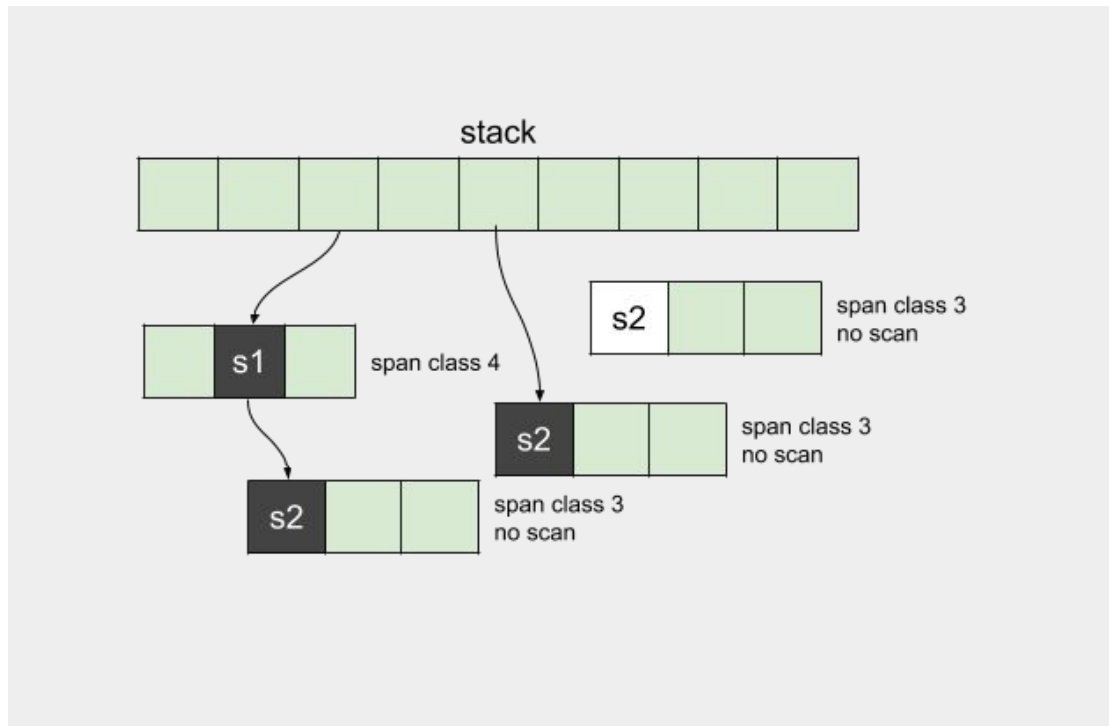


-  $2 * \text{GC percentage} / 100 * \text{current memory size}$  (In default GC percentage = 100)

- Every few minutes

# Mark and Sweep

## Mark phase (tri-color mark)

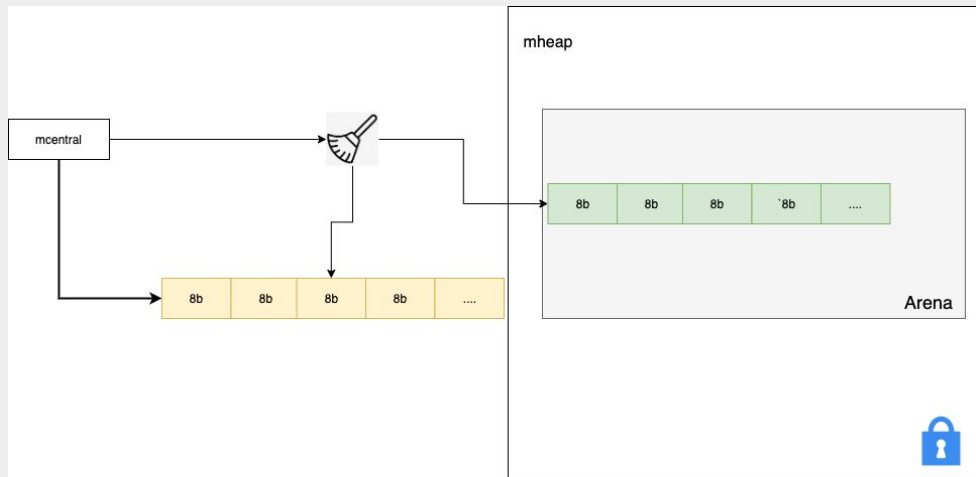


- Scan the heap and check the object is reference or not.

- Black: In-used, White: Unused

# Mark and Sweep

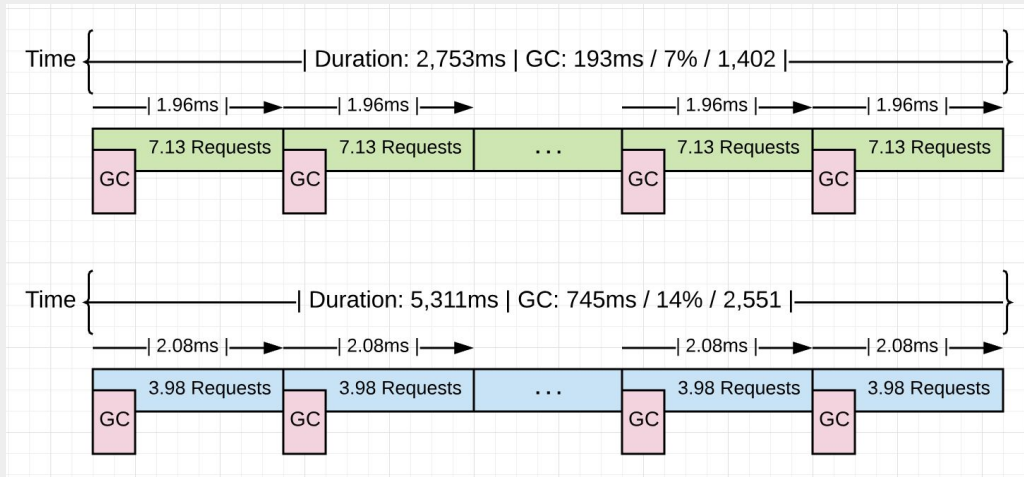
## Sweep phase



- Sweep concurrently in a background goroutine.
- When the goroutine needs another space on heap, it first attempts to reclaim that much memory by sweeping. (Before ask more from mheap or mcentral)

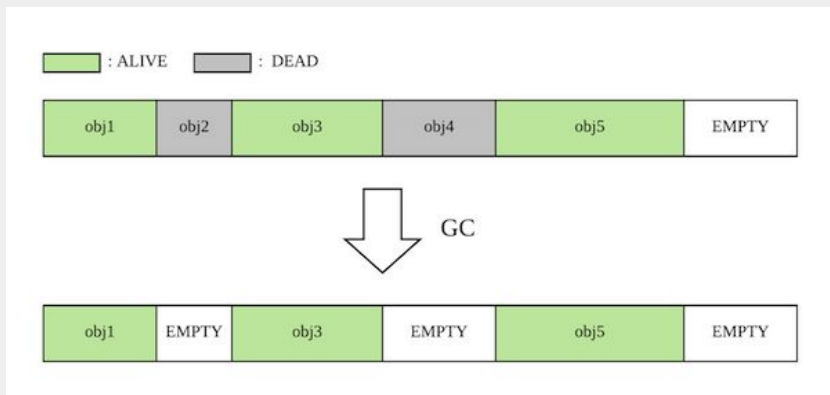
# Mark and Sweep

## Pacing



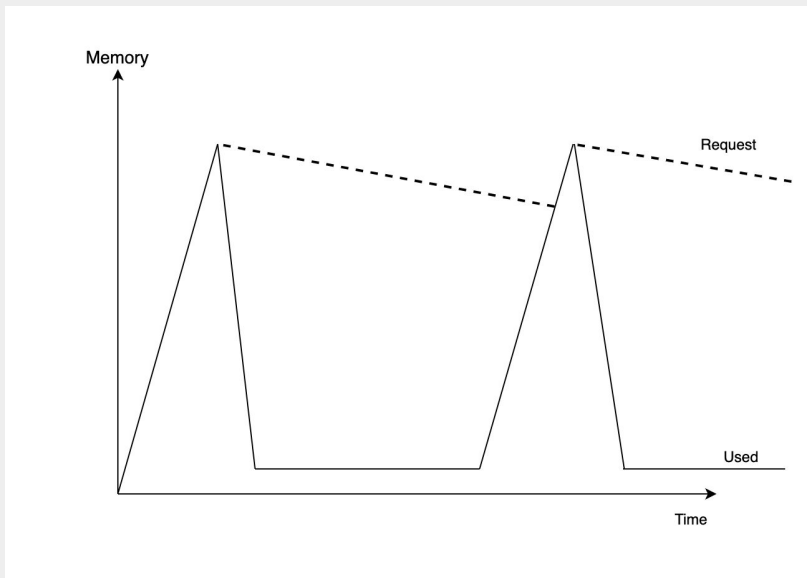
- Pacing calculate how fast the next garbage collection trigger.

# Heap fragmentation



- The heap memory will be fragment because of GC.

- The fragmentation reduce the efficiency.



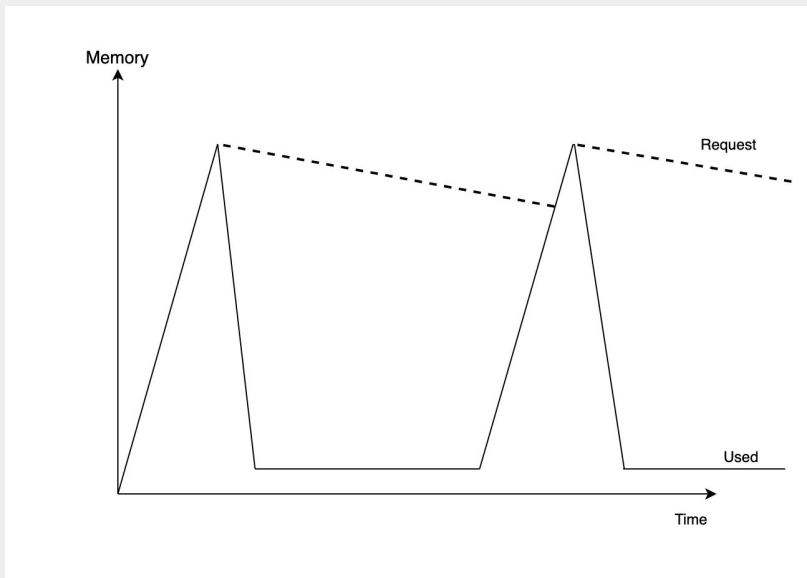
- Free and release the physical page backing mapping memory to OS.
- Reduce the page level fragmentation and Resident set size.



# Scavenging



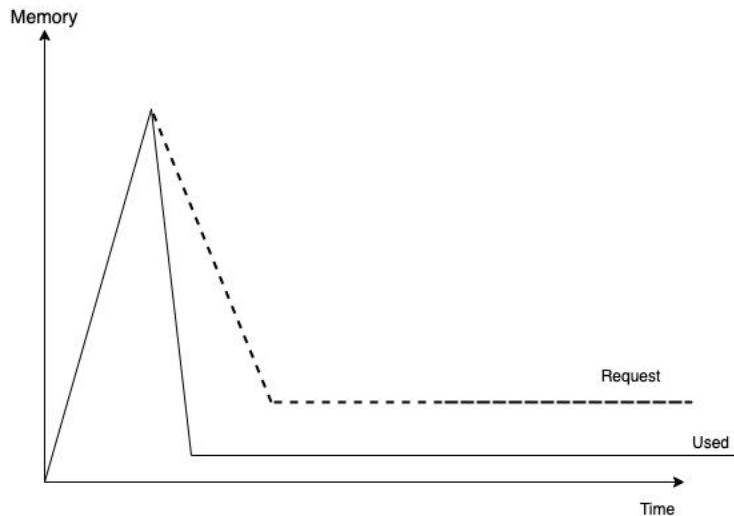
## Trigger scavenging



- Background scavenging
- Heap growth

# Scavenging

## How many memory we should retain



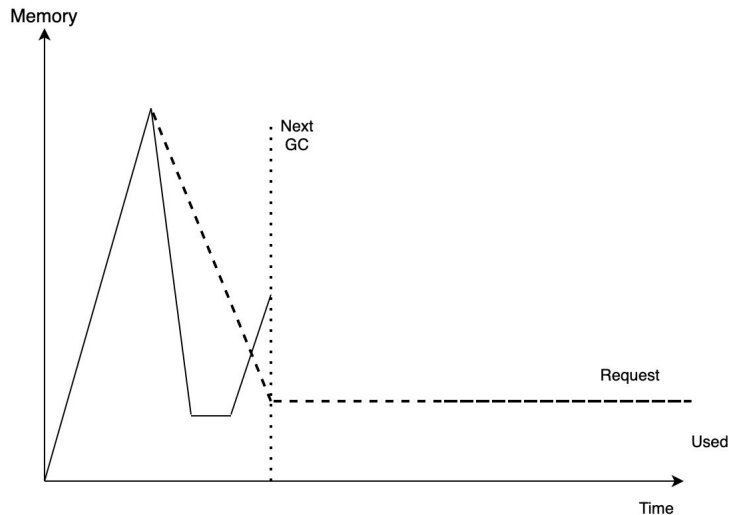
$C * (\text{next\_gc\_goal} / \text{last\_next\_gc\_goal}) * \text{last\_heap\_inuse}$

(C = 1.1 in 1.16 version)

# Scavenging



At what rate is memory scavenged

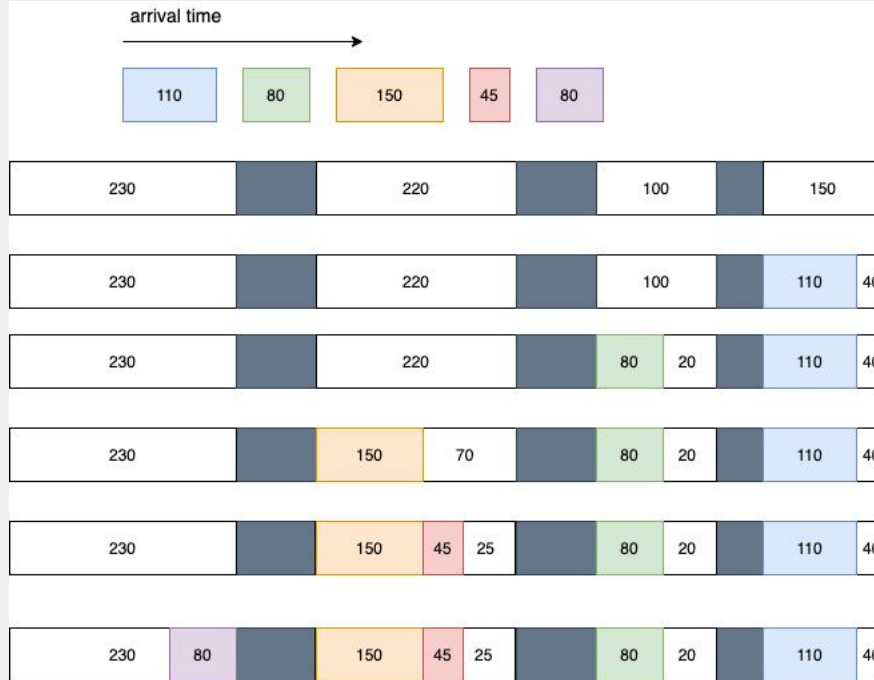


The goal is achieving the goal before the next collection trigger. (We record the interval between each trigger in pacing)

# Scavenging



## How to handle new allocation on heap



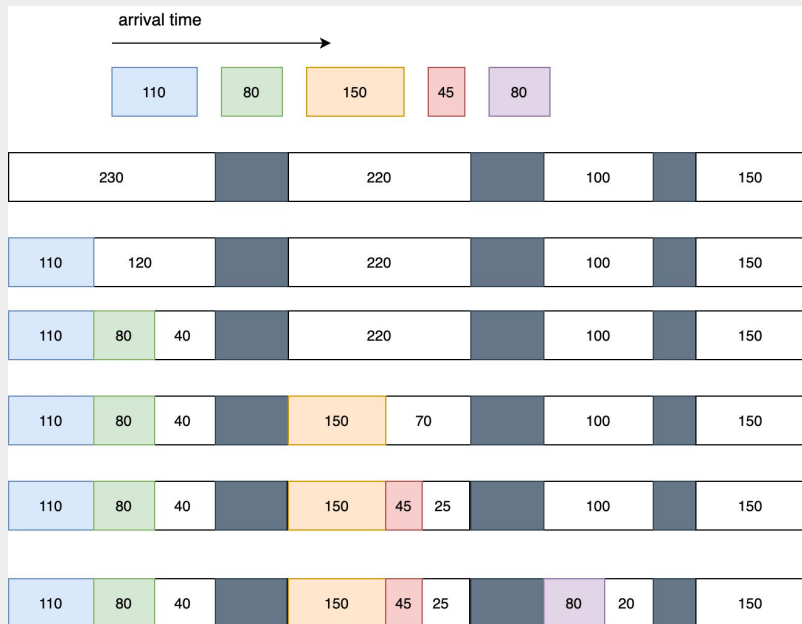
- Best-fit: the memory allocated at the closest fragmentation.

- This should provide the less fragmentation.

# Scavenging



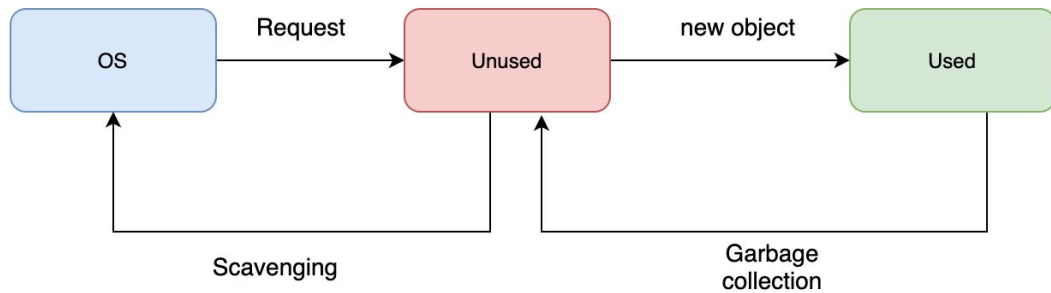
## How heap handle new allocation



- First-fit: the partition is allocated which is first sufficient.

- Scavenging from the highest base addresses first.

# Recap



## Reference:

- Smarter Scavenging:  
<https://go.googlesource.com/proposal/+/master/design/30333-smarter-scavenging.md>
- Deep understanding go memory allocation:  
<https://programmer.group/deep-understanding-go-memory-allocation.html>
- Evolving the Go Memory Manager's RAM and CPU Efficiency: [https://youtu.be/S\\_1YfTfuWmo](https://youtu.be/S_1YfTfuWmo)
- Tcmalloc: <https://google.github.io/tcmalloc/design.html>
- GC in go: <https://slides.com/jalex-chang/gc-in-go>



Thank you