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- Like stacks, queues are an Abstract Data Type with more than one implementation.

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 - integer size(): returns the number of elements stored.
 - boolean isEmpty(): indicates whether no elements are stored.

Example

	Lxample	
Operation	Return value	Queue
enqueue(5)	-	(5)
enqueue(3)	-	(5, 3)
dequeue()	5	(3)
enqueue(7)	_	(3, 7)
dequeue()	3	(7)
first()	7	(7)
dequeue()	7	()
dequeue()	null	()
<pre>isEmpty()</pre>	true	()
enqueue(9)	-	(9)
enqueue(7)	-	(9, 7)
size()	2	(9, 7)
enqueue(3)	_	(9, 7, 3)
enqueue(5)	_	(9, 7, 3, 5)
dequeue()	9	(7, 3, 5)

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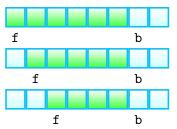
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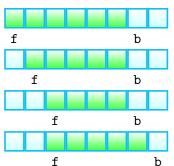
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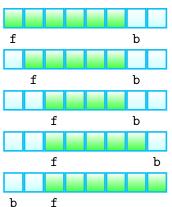
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- Two variables keep track of the front and size:
 - front: index of the front element
 - size: number of stored elements
- When the queue has fewer than N elements, array location back = (front + size) mod N is the first empty slot past the rear of the queue.

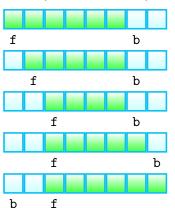


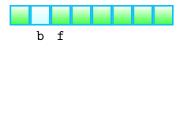


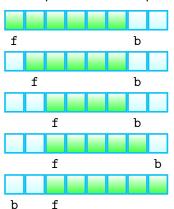


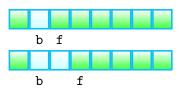


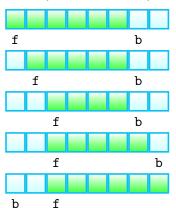


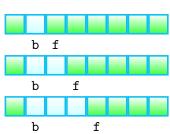


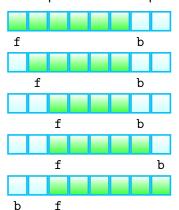


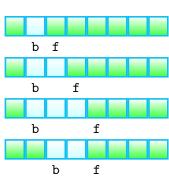












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 \begin{array}{l} \textbf{procedure} \ \mathtt{ENQUEUE}(o) \\  back \leftarrow (\mathit{front} + \mathit{size}) \ \mathtt{mod} \ \mathit{N} \\  arr[\mathit{back}] \leftarrow o \\  \mathit{size} \leftarrow \mathit{size} + 1 \\ \textbf{end} \ \mathbf{procedure} \end{array}
```

Keep in mind that you need to consider the case when the queue is full and how to handle it.

Notice that the cell in the array, arr[f] is not set to null after it is dequeued. Consider the case, when the queue is empty and you attempt to dequeue.

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- Elements would be added and removed to opposite ends of the list. The most efficient way to do this would depend on the list.

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 - Each operations runs in time O(1).

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 - During the process of resizing, the queue would be "unwound", so that the front of the queue would be in index 0 of the resized array.
- Linked list-backed queue do not have either of the above two limitations

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- We can implement a round robin scheduler using a queue processes by repeatedly performing the following steps:
 - 1. process = processes.dequeue()
 - 2. Do work on process
 - 3. processes.enqueue(process)