LexBFS and its applications

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Overview

Graph Searches

What's in a graph?

Graph

We consider non-oriented, simple and connected graphs

INCLUDE GRAPHS EXAMPLES AND COUNTEREXAMPLES

Generic Search

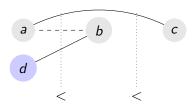
```
for i in [1, ..., n]:
if i == 1:
    u = any vertex
else:
    u = any unvisited marked vertex
visit(u)
for v in neighbours(u):
    mark(v)
```

Let's number vertices in the order they are visited.

Theorem

An order σ corresponds to a Generic Search if and only if

 $\forall a <_{\sigma} b <_{\sigma} c, ac \in E \text{ and } ab \notin E, \exists d <_{\sigma} b \text{ st } db \in E$



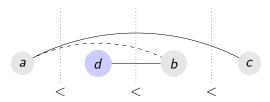
INCLUDE GRAPH EXAMPLE

```
for i in [1, ..., n]:
if i == 1:
    u = any vertex
else:
    u = any unvisited vertex w/ max label
visit(u)
for v in neighbours(u):
    label[v] = i
```

Theorem

An order σ corresponds to a DFS if and only if

 $\forall a <_{\sigma} b <_{\sigma} c, ac \in E \text{ and } ab \notin E, \exists a <_{\sigma} d <_{\sigma} b \text{ st } db \in E$



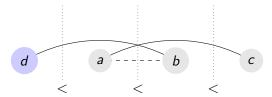
```
for i in [n, ..., 1]:
if i == n:
    u = any vertex
else:
    u = any unvisited vertex w/ max label
visit(u)
for v in neighbours(u):
    if v has no label:
        label[v] = i
```

INCLUDE GRAPH EXAMPLE

Theorem

An order σ corresponds to a BFS if and only if

 $\forall a <_{\sigma} b <_{\sigma} c, ac \in E \text{ and } ab \notin E, \exists d <_{\sigma} a \text{ st } db \in E$



Let's rewrite BFS

```
for i in [n, ..., 1]:
if i == n:
    u = any vertex
else:
    u = any unvisited vertex
        w/ max first element of label
visit(u)
for v in neighbours(u):
    label[v].append(i)
```

Here is LexBFS

```
for i in [n, ..., 1]:
if i == n:
    u = any vertex
else:
    u = any unvisited vertex
    w/ max lexicographical label
visit(u)
for v in neighbours(u):
    label[v].append(i)
```

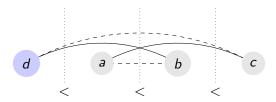
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INCLUDE GRAPH EXAMPLE

Theorem

An order σ corresponds to a LexBFS if and only if

 $\forall a <_{\sigma} b <_{\sigma} c, ac \in E \text{ and } ab \notin E, \exists d <_{\sigma} a \text{ st } db \in E \text{ and } dc \notin E$

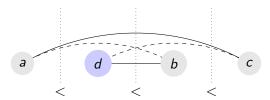


Iterating on [1, ..., n] and prepending, we obtain LexDFS

$\mathsf{Theorem}$

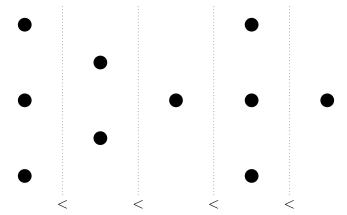
An order σ corresponds to a LexDFS if and only if

 $\forall a <_{\sigma} b <_{\sigma} c, ac \in E, ab \notin E, \exists a <_{\sigma} d <_{\sigma} b \text{ st } db \in E, dc \notin E$



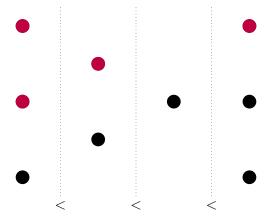
How to implement LexBFS: Pivoting

Let's order vertices by label:



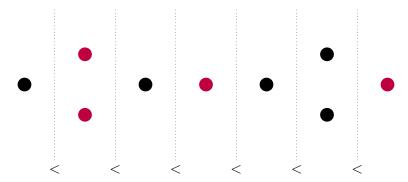
How to implement LexBFS: Pivoting

We remove one of the maximum vertex :



How to implement LexBFS: Pivoting

Now, we split each bucket



Data structure for pivoting

- Vertices in a doubly linked list
- Pointer from each vertex to its class
- Pointers from each class to its first and last vertex

