Machine Problem 3: Page Table Management

Gaurangi Sinha

UIN: 734004353

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The goal of this assignment is the implementation of a demand-paging based virtual memory system for a kernel, focusing on the paging mechanism of the x86 architecture. It aims for the setup and initialization of the paging system and page table infrastructure for managing memory, initially for a single address space, with the prospect of extending to multiple processes and address spaces. This involves changes in page table management, handling page faults, and understanding the memory layout and management within a kernel environment.

We work on handling modest amounts of memory, enabling us to maintain the page table directly in physical memory. As mentioned in the Machine problem handout, this setup serves as a precursor to more advanced exercises where the page table itself will be managed in virtual memory.

In this machine problem, we have created a manager for handling page frames in the memory system of an x86 architecture-based machine. This manager is built on a hierarchical structure with two levels: the "page directory" at the top level and the "page table" at the second level. Each of these tables contain entries that are 32 bits in length, with the first 10 bits dedicated to navigation within the page directory, the next 10 bits for navigation within the page table, and the final 12 bits designated for the offset inside the physical memory frame. The designed paging mechanism is prepared to allocate pages to processes and adeptly manage page faults should they occur

Assigned Task

Completed

Code Description

I made modifications to the files "page_table.c," "cont_frame_pool.c" and "cont_frame_pool.h". The codes for "cont_frame_pool.c" and "cont_frame_pool.h" were reused from the machine problem 2. To compile the modified code, the provided makefile should be executed. So I run the commands mentioned in compilation section of this document. For simulating the environment, the "copykernel.sh" script is used, followed by initiating the Bochs simulator with the command "bochs -f bochsrc.bxrc".

init paging: Initializes the static variables of the PageTable.

The PageTable constructor initially creates a page table tailored for directly mapped memory aligned with the kernel pool. Within this table, all pages are marked as "VALID" and granted "READ/WRITE" permissions. Conversely, the other frames are set to "INVALID" status, earmarked for user-level operations. For the page directory, only the first entry is marked as valid, with all subsequent entries being deemed invalid.

The "load" function initialises the current page table to reflect the current page table object and sets up the CR3 register with the active page directory.

Page fault handling falls under the purview of the "handle_fault" function. It initially looks up the error code in the REGS structure to see if the interruption it received is related to a page fault. The CR2 register contains the address that is generating the issue. The indices for the page directory and page table are calculated using this address. A new frame is allocated for the directory entry in the event that it is determined to be invalid, hence initialising a new page table. Next, the page table entry's validity is verified. In the event that it is invalid, a frame is also allocated for this entry, which is equivalent to assigning a page from the process pool that the page table entry would reference.

Compilation

To compile the code, we run the following commands:

make clean make ./copykernel.sh bochs –f bochsrc.bxrc

Testing

All the tests are passing, for given and additional scenarios. I have attached the screenshot if the output I got.











