Iterative deepening Search

Introduction

- The iterative deepening algorithm is a combination of DFS and BFS algorithms. This search algorithm finds out the best depth limit and does it by gradually increasing the limit until a goal is found.
- This algorithm performs depth-first search up to a certain "depth limit", and it keeps increasing the depth limit after each iteration until the goal node is found.

IDS

- This Search algorithm combines the benefits of Breadth-first search's fast search and depth-first search's memory efficiency.
- The iterative search algorithm is useful uninformed search when search space is large, and depth of goal node is unknown.

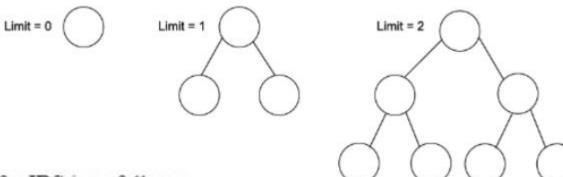
Advantages:

• It combines the benefits of BFS and DFS search algorithm in terms of fast search and memory efficiency.

Disadvantages:

 The main drawback of IDDFS is that it repeats all the work of the previous phase.

IDS



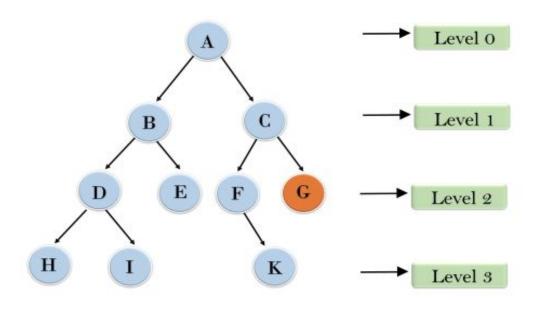
The algorithm for IDS is as follows:

- 1. Set depth-limit ← 0
- Do Solution = DLS(depth-limit, initial node)

If(solution= goal state) then return else Depth-limit = depth-limit + 1 continue

Example:

Iterative deepening depth first search



1'st Iteration----> A
2'nd Iteration----> A, B, C
3'rd Iteration-----> A, B, D, E, C, F, G
4'th Iteration-----> A, B, D, H, I, E, C, F, K, G
In the fourth iteration, the algorithm will
find the goal node.

Completeness:

This algorithm is complete is if the branching factor is finite.

Time Complexity:

 Let's suppose b is the branching factor and depth is d then the worst-case time complexity is O(b^d).

Space Complexity:

The space complexity of IDDFS will be O(bd).

Optimal:

 IDDFS algorithm is optimal if path cost is a nondecreasing function of the depth of the node.

Bidirectional Search Algorithm

- Bidirectional search algorithm runs two simultaneous searches, one form initial state called as forward-search and other from goal node called as backward-search, to find the goal node.
- Bidirectional search replaces one single search graph with two small subgraphs in which one starts the search from an initial vertex and other starts from goal vertex.

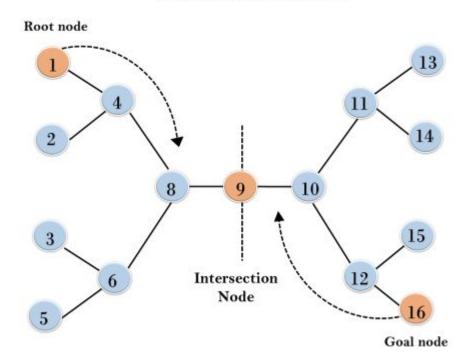
Bidirectional Search Algorithm

- The search stops when these two graphs intersect each other.
- Bidirectional search can use search techniques such as BFS, DFS, DLS, etc.
- Advantages:
- Bidirectional search is fast.
- Bidirectional search requires less memory
- Disadvantages:
- Implementation of the bidirectional search tree is difficult.
- In bidirectional search, one should know the goal state in advance.

Example:

- In the below search tree, bidirectional search algorithm is applied. This algorithm divides
 one graph/tree into two sub-graphs. It starts traversing from node 1 in the forward
 direction and starts from goal node 16 in the backward direction.
- The algorithm terminates at node 9 where two searches meet.

Bidirectional Search



Completeness: Bidirectional Search is complete if we use BFS in both searches.

Time Complexity: O(b^d). Space Complexity: O(b^d).

Optimal: Bidirectional search is

Optimal.